Relational Algebra (and SQL)

Annoucements

Homework 1: Due beginning of class Thursday

Project 1 Part 1: End of class Thursday

Project 1 Part 2:

Available Thursday; due in 2 weeks

Comment: lecture timing is bad

Yes, sorry. Will attempt to avoid.

This week: Relational Algebra. Next week: SQL.

-> 1 week to do the project part 2

Comment: Live SQL example from ER

Goal: full example next week, with queries

Q: Still confused about aggregates

Example in the middle of class

Reading

Ramakrishnan Sections 4.1 and 4.2

Helpful Reference

https://en.wikipedia.org/wiki/Relational_algebra

Queries: ask your data a question

Declarative language: ask *what* you want, now *how*

Meaning: Queries are high level, “readable”

Not Turing complete (can’t execute any algorithm)

Supports easy, efficient access to large databases

Domain specific language for data access

SQL: Data Definition Language (DDL)

Data Manipulation Language (DML)

Basic Single Table SELECT

SELECT *(output)* FROM *(input)*

WHERE *(condition)*

```
SELECT * FROM Students
SELECT name FROM Students
SELECT name FROM Students WHERE age < 21
SELECT name, login FROM Students WHERE gpa >= 3
```

sid	name	login	age	gpa
1	eugene	ewu	20	2.5
2	luis	gravano	25	3.5
3	martha	martha	32	3.9

Single Table Semantics

conceptual evaluation method:

1. FROM clause: retrieve Students relation
2. WHERE clause: Check conditions, discard tuples that fail
3. SELECT clause: Delete unwanted fields

Real evaluation is *much* more efficient, but must produce the same answers.

Ambiguous names

E.g.   Students: (sid, name, …)

        Enrolled: (sid, cid, grade)

Qualified names: Use table name: Students.age

Rename: AS (optional): shortcuts, ambiguity, clarity

SELECT S.sid, S.name FROM Students AS S

SELECT S.sid, S.name FROM Students S

Related data: Multiple tables

What does this return?

```
SELECT S.name, E.cid
FROM Students S, Enrolled E
WHERE S.sid = E.sid AND
      E.grade = 'A'
```

Enrolled

sid	cid	grade
1	2	A
1	3	B
2	2	A+

SELECT ~~S.name, E.cid~~
~~FROM Students S~~, Enrolled E
WHERE ~~S.sid = E.sid AND~~
      E.grade = 'A'

Students

sid	name
1	eugene
2	luis

Result

name	cid
eugene	2

Multi-Table Semantics

- Modify the FROM clause evaluation
  - 1. FROM clause: compute *cross-product* of Students and Enrolled

Enrolled

sid	cid	grade
1	2	A
1	3	B
2	2	A+

Students

sid	name
1	eugene
2	luis

Cross-product

sid	cid	grade	sid	name
1	2	A	1	eugene
1	3	B	1	eugene
2	2	A+	1	eugene
1	2	A	2	luis
1	3	B	2	luis
2	2	A+	2	luis

Multi-Table Semantics

Modify the FROM clause evaluation

1. FROM clause: compute *cross-product* of Students, Enrolled
2. WHERE clause: Check conditions, discard tuples that fail
3. SELECT clause: Delete unwanted fields

Formal Relational Query Languages

Formal basis for real languages e.g., SQL

Relational Algebra

- Function of operations applied to relations
- Operational: step-by-step execution plans

Relational Calculus

- Logical, describes what data users want
- (not operational, fully declarative)

Definitions

Relation (for this lecture)

- Instance is a set of tuples
- Schema defines field names and types (domains)
- Students(<u>sid</u> int, name text, major text, gpa int)

- Fields: Reference by name: (e.g. major)
  - Reference by position starting at 1: (e.g. 3)

Names are for humans; positions for “intermediate” results

Definitions

Query is a function over relation instances

$$Q(R_1,\ldots,Rn) = R_{result}$$

Schemas of output relations are well defined by query Q.

Use positional or named field notation

Names in results inherited from input (unless renamed)

Relational Algebra Overview

Core 5 operations

- PROJECT (π)
- SELECT (σ)
- UNION (U)
- SET DIFFERENCE (-)
- CROSSPRODUCT (x)

Additional operations

- RENAME (p)
- INTERSECT (∩)
- JOIN (⋈)
- DIVIDE (/)

Instances Used Today: Library

Students, Reservations

R1

sid	rid	day
1	101	10/10
2	102	11/11

S1

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

S2

sid	name	gpa	age
4	aziz	3.2	21
2	barak	3	21
3	trump	2	88
5	rusty	3.5	21

Project

$$\pi_{<attr1,\ldots>}(A) = R_{result}$$

Extract desired fields (subset of columns)

Schema is subset of input schema in the projection list

- $\pi_{<a,b,c>}(A)$ has output schema (a,b,c) w/ types carried over

Project

S2

sid	name	gpa	age
4	aziz	3.2	21
2	barak	3	21
3	trump	2	88
5	rusty	3.5	21

$$\pi_{name,age}(S2) =$$

name	age
aziz	21
barak	21
trump	88
rusty	21

Project

S2

sid	name	gpa	age
4	aziz	3.2	21
2	barak	3	21
3	trump	2	88
5	rusty	3.5	21

$$\pi_{age}(S2) =$$

age
21
88

Where did all the rows go?

Real systems typically don't remove duplicates. Why?

Select

$$\sigma_{<p>}(A) = R_{result}$$

Select subset of rows that satisfy condition *p*

Won't have duplicates in result. Why?

Result schema same as input

Select

S1

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

$\sigma_{age<30}$ (S1) =

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21

$\pi_{name}(\sigma_{age<30}$ (S1)) =

name
eugene
barak

Commutativity

$$A + B = B + A$$

$$A * B = B * A$$

$$A + (B * C) = (B * C) + A$$

Associativity:

$$A + (B + C) = (A + B) + C$$

$$A + (B * C) = (A + B) * C$$

Commutativity

$$A + B = B + A$$

$$A * B = B * A$$

$$A + (B * C) = (B * C) + A$$

Associativity:

$$A + (B + C) = (A + B) + C$$

$$\cancel{A + (B * C) = (A + B) * C}$$

Commutativity

$\pi_{age}(\sigma_{age<30}(S1))$

$\sigma_{age<30}$ (

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

)

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21

Commutativity

$$\pi_{age}(\sigma_{age<30}(S1))$$

$\sigma_{age<30}$

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

π_{age}

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21

=

age
20
21

Commutativity

$$\sigma_{age<30}(\pi_{age}(S1))$$

$$\underbrace{\pi_{age}\left(\begin{array}{|c|c|c|c|}\hline \text{sid} & \text{name} & \text{gpa} & \text{age} \\ \hline 1 & \text{eugene} & 4 & 20 \\ \hline 2 & \text{barak} & 3 & 21 \\ \hline 3 & \text{trump} & 2 & 88 \\ \hline \end{array}\right)}$$

age
20
21
88

Commutativity

$$\sigma_{age<30}(\pi_{age}(S1))$$

π_{age} (

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

)

$\sigma_{age<30}$ (

age
20
21
88

) =

age
20
21

Commutativity

Does Project and Select commute?

$$\pi_{age}(\sigma_{age<30}\ (S1)) = \sigma_{age<30}(\pi_{age}(S1))$$

What about

$$\pi_{name}(\sigma_{age<30}\ (S1))?$$

Commutativity

Does Project and Select commute?

$$\pi_{age}(\sigma_{age<30}\ (S1)) = \sigma_{age<30}(\pi_{age}(S1))$$

What about

$$\pi_{name}(\sigma_{age<30}\ (S1)) \mathrel{!=} \sigma_{age<30}(\pi_{name}(S1))$$

Commutativity

Does Project and Select commute?

$$\pi_{age}(\sigma_{age<30}\ (S1)) = \sigma_{age<30}(\pi_{age}(S1))$$

What about

$$\pi_{name}(\sigma_{age<30}\ (S1))\ != \sigma_{age<30}(\pi_{name,\ age}(S1))$$

Commutativity

Does Project and Select commute?

$$\pi_{age}(\sigma_{age<30}(S1)) = \sigma_{age<30}(\pi_{age}(S1))$$

What about

$$\pi_{name}(\sigma_{age<30}(S1)) = \pi_{name}(\sigma_{age<30}(\pi_{name,\, age}(S1)))$$

OK!

Union, Set-Difference

A op B = R_{result}

A, B must be *union-compatible*

- Same number of fields
- Field i in each schema have same type

Result Schema borrowed from first arg (A)

- Student(sid int, age int) U Class(cid int, max int) = ?

Union, Set-Difference

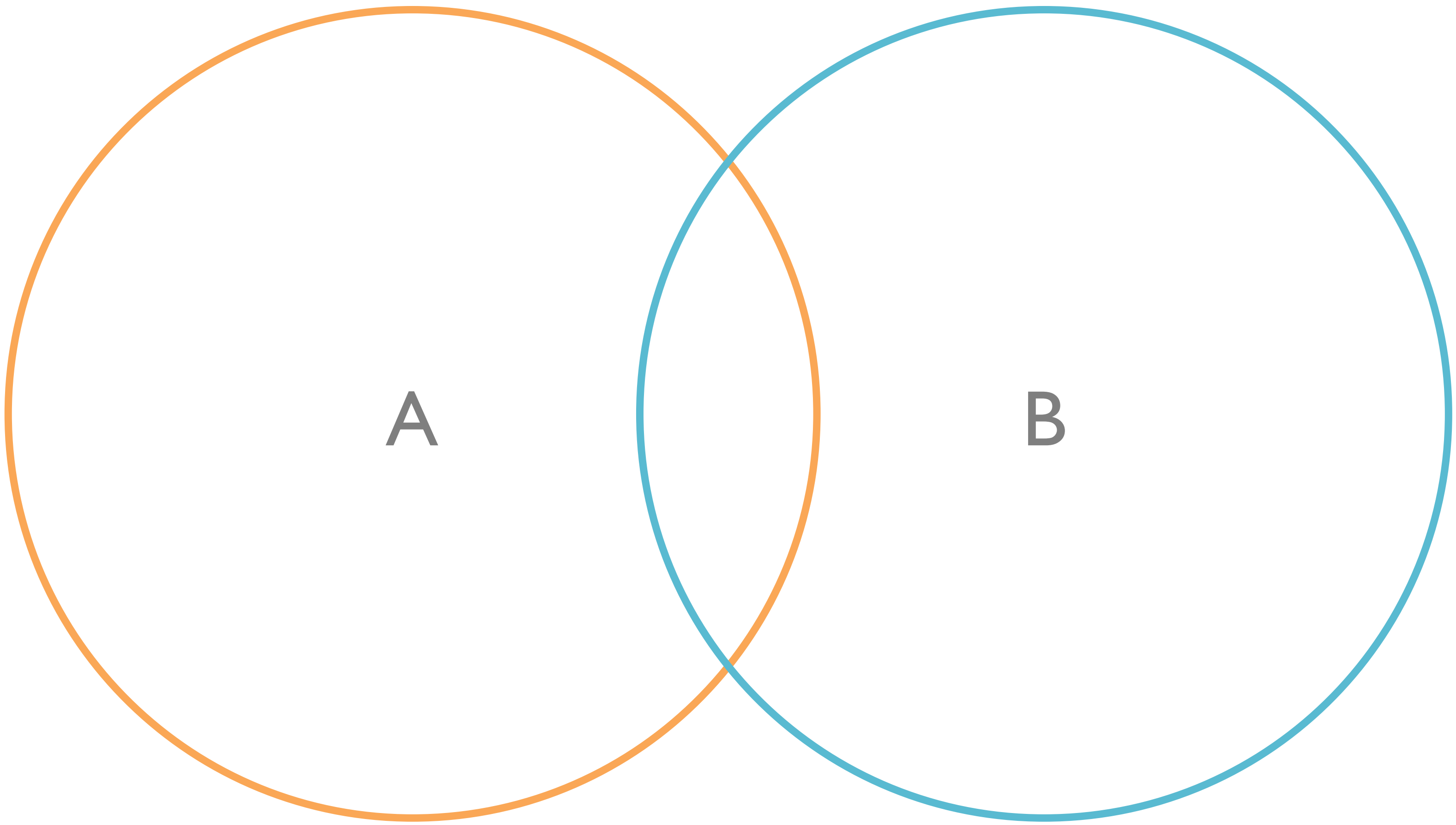
$$A \text{ op } B = R_{result}$$

A, B must be *union-compatible*

- Same number of fields
- Field i in each schema have same type

Result Schema borrowed from first arg (A)

- Student(sid int, age int) U Class(cid int, max int) = R_{result}(sid int, age int)


A
B

Union, Intersect, Set-Difference

S1

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

S2

sid	name	gpa	age
4	aziz	3.2	21
2	barak	3	21
3	trump	2	88
5	rusty	3.5	21

S1US2 =

sid	name	gpa	age
1	eugene	4	20
4	aziz	3.2	21
5	rusty	3.5	21
3	trump	2	88
2	barak	3	21

Union, Intersect, Set-Difference

S1

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

S2

sid	name	gpa	age
4	aziz	3.2	21
2	barak	3	21
3	trump	2	88
5	rusty	3.5	21

S1 - S2 =

sid	name	gpa	age
1	eugene	4	20

Note on Set Difference & Performance

Most operators are monotonic

increasing size of inputs → outputs grow

can compute *incrementally*

Set Difference is *not monotonic:*

compute A – B

add data in X to B: B2 = B ∪ X

compute A – B2

could be smaller than A – B

Set difference is *blocking*:

For T – S, must wait for all S tuples before any results

Intersect

$$A \cap B = R_{result}$$

A, B must be *union-compatible*

Intersect

S1

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

S2

sid	name	gpa	age
4	aziz	3.2	21
2	barak	3	21
3	trump	2	88
5	rusty	3.5	21

$S1 \cap S2$ =

sid	name	gpa	age
2	barak	3	21
3	trump	2	88

Intersect

$$A \cap B = R_{result}$$

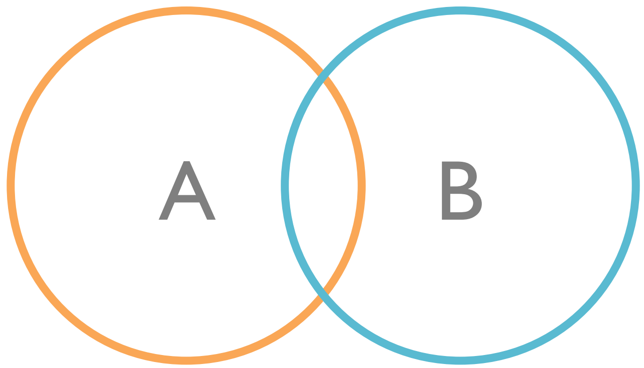
A, B must be *union-compatible*

Can we express using core operators?

$$A \cap B = ?$$

Intersect

$$A \cap B = R_{result}$$

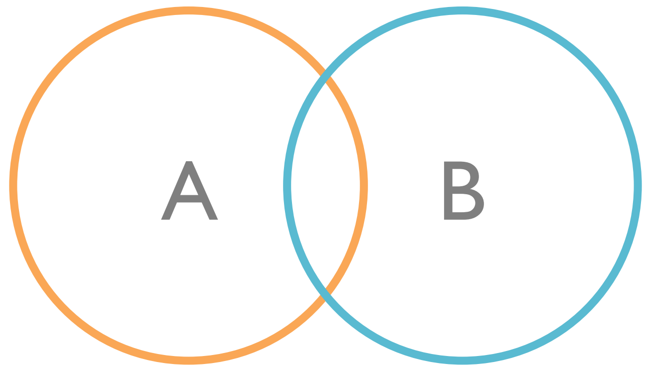


Can we express using core operators?

$A \cap B = A - ?$ (think venn diagram)

Intersect

$$A \cap B = R_{result}$$



Can we express using core operators?

$$A \cap B = A - (A - B)$$

Cross-Product

$$A(a_1,\ldots,a_n) \times B(a_{n+1},\ldots,a_m) = R_{result}(a_1,\ldots,a_m)$$

Each row of A paired with each row of B

- Result schema: Combine A and B's fields, inherit if possible
- Conflict: students and reservations have *sid* field: need positions

Cross-Product

S1

sid	name	gpa	age
1	eugene	4	20
2	barak	3	21
3	trump	2	88

R1

sid	rid	day
1	101	10/10
2	102	11/11

S1 x R1 =

(sid)	name	gpa	age	(sid)	rid	day
1	eugene	4	20	1	101	10/10
2	barak	3	21	1	101	10/10
3	trump	2	88	1	101	10/10
1	eugene	4	20	2	102	11/11
2	barak	3	21	2	102	11/11
3	trump	2	88	2	102	11/11

Rename

ρ(<new_name>(<mappings>), Q)

Explicitly defines/changes field names of schema

ρ(C(1 → sid1, 5 → sid2), S1 x R1)

C =

sid1	name	gpa	age	**sid2**	rid	day
1	eugene	4	20	1	101	10/10
2	barak	3	21	1	101	10/10
3	trump	2	88	1	101	10/10
1	eugene	4	20	2	102	11/11
2	barak	3	21	2	102	11/11
3	trump	2	88	2	102	11/11

Example: Aggregate

Three entities: Company, Product, Service Rep

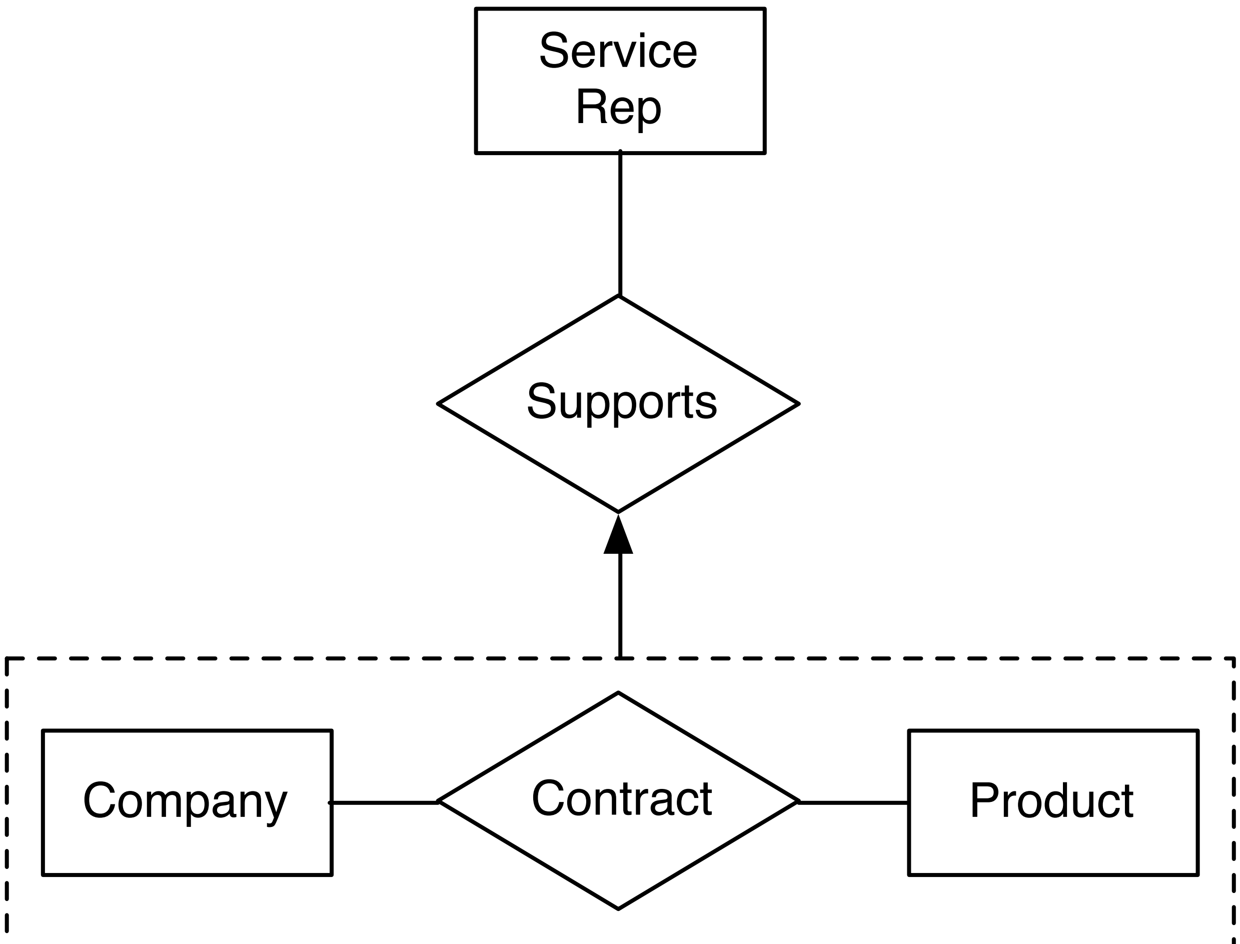
Companies purchase products with contracts

Big contracts have one dedicated service rep

Q1: Draw the E-R diagram (hint: use aggregate)

Q2: Create a schema

Diagram



SQL Schema: Entities

```
CREATE TABLE Company(
    cid int PRIMARY KEY,
    email text);
CREATE TABLE Product(
    pid int PRIMARY KEY,
    description text);
CREATE TABLE ServiceRep(
    sid int PRIMARY KEY,
    name text);
```

Relationships (no constraints)

```
CREATE TABLE Contract(
    cid int REFERENCES Company,
    pid int REFERENCES Product,
    PRIMARY KEY (cid, pid));
CREATE TABLE Supports(
    cid int,
    pid int,
    sid int REFERENCES SupportRep,
    PRIMARY KEY (cid, pid, sid),
    FOREIGN KEY (cid, pid) REFERENCES Contract));
```

At most one: Make unique

```
CREATE TABLE Supports(
    cid int,
    pid int,
    sid int REFERENCES SupportRep,
    PRIMARY KEY (cid, pid, sid),
    FOREIGN KEY (cid, pid) REFERENCES Contract),
    UNIQUE (cid, pid));
```

DUPLICATION: UNIQUE + PRIMARY KEY

Reduce duplication

```
CREATE TABLE Supports(
    cid int,
    pid int,
    sid int REFERENCES SupportRep NOT NULL,
    PRIMARY KEY (cid, pid),
    FOREIGN KEY (cid, pid) REFERENCES
Contract));
```

DUPLICATION: Same primary key as Contract

Combine tables

```
CREATE TABLE Contract(
    cid int REFERENCES Company,
    pid int REFERENCES Product,
    sid int REFERENCES SupportRep,
    PRIMARY KEY (cid, pid));
```