

QCQP: Progress Report

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The QCQP

We consider the QCQP

$$\begin{aligned} (\mathbf{QCQP}) \quad & \max \quad x^T Q x + q^T x \\ & \text{s.t.} \quad x^T A_i x + a_i^T x (\leq, =, \geq) b_i \end{aligned} \tag{1}$$

Rank- r Indefiniteness I

Formally, a quadratic inequality induced by a symmetric matrix A can be expressed as the following,

$$x^T \left(\sum_{j \in J_+} \lambda_j v_j v_j^T \right) x + a^T x \leq b + \left(\sum_{j \in J_-} \lambda_j v_j v_j^T \right) \quad (2)$$

- ▶ For clarity, we say a quadratic inequality $x^T A x + a^T x \leq b$ is *rank- r indefinite* if first r eigenvalues are negative.
- ▶ In comparison, a maximization problem with matrix Q is rank- r indefinite if last r eigenvalues are nonnegative.

Rank-1 Indefinite Quadratic Inequality I

Suppose $Q = RR^T - aa^T$

For quadratic inequality, it can be written as,

$$x^T RR^T x \leq (a^T x)^2$$

Then we naturally solve the problem in two disjoint subregions, by $a^T x \geq 0 \vee a^T x \leq 0$,

$$(P1) \quad \|R^T x\| \leq -a^T x$$

$$(P2) \quad \|R^T x\| \leq a^T x$$

Implementation,

- ▶ Create offline subregions at depth 0 (multiple root nodes)
- ▶ **Seems** impossible to do explicit disjunctions for inhomogeneous case.
- ▶ Apply branch and bound.

Rank-1 Indefinite Maximization I

Consider maximization

$$\begin{aligned} \max \quad & x^T(aa^T - RR^T)x + q^Tx \\ \Rightarrow \max \quad & z \\ \text{s.t.} \quad & z + x^TRR^Tx - q^Tx \leq (a^Tx)^2 \end{aligned}$$

We analyze,

$$S = \{(z, x) : x^TRR^Tx + z - q^Tx \leq (a^Tx)^2\}$$

1. Simply branch on $\text{sign}(a^Tx)$ will prune one region very fast by bound.
2. **Seems** impossible to do explicit disjunctions as homogeneous inequality.
3. If using **best bound** rule in node selection, it has no obvious advantage.
4. Apply branch and bound.

Extension to rank- r I

Consider quadratic maximization,

$$\begin{aligned} \max \quad & z \\ \text{s.t.} \quad & \sum_{j \in J_-} \lambda_j z_j - q^T x + z \leq \sum_{j \in J_+} \lambda_j z_j \\ & y_j \geq (x_j^T v_j)^2, j = 1, \dots, n \end{aligned} \tag{3}$$

With box or ball as regularity (else (3) will be unbounded.)

- ▶ ball: $x \in B(0, \delta)$, $x \in [0, 1]^n$
- ▶ ellipsoid: $x \in \{x : x^T A x + a^T x \leq \delta^2\}$
- ▶ regularity implies, (explains why ball constraint is easier.)

$$y^T e = \|x\|^2 \tag{4}$$

Computational Results I

- ▶ We test on “unconstrained” (with only a box or ball on x)
- ▶ MSC is very effective in the test problems. as $r \nearrow$, the problem becomes hard.
- ▶ MSC is tight with only ball constraints? (TRS)

Future Work I

- ▶ Extension to an indefinite quadratic inequalities
- ▶ Plug in ADMM as primal feasible solution.

Future Work: Multiple Quadratic Constraints I

Our old MSC.

$$\begin{aligned} \text{Maximize} \quad & y_0^T \lambda_0 + q^T x \\ \text{s.t.} \quad & V_i z_i = x & i = 0, \dots, m & (5) \\ & y_i^T \lambda_i + a_i^T x \leq b_i & i = 1, \dots, m & (6) \\ & y_i = z_i \circ z_i & i = 0, \dots, m & (7) \end{aligned}$$

- ▶ In this formulation, we need $m \times n$ auxillary pairs (z, y) by allowing different bases V_i
- ▶ Actually this may not be necessary...

Future Work: Multiple Quadratic Constraints II

Consider one indefinite A with a rank- r indefinite Q , we can convexify A by...

$$Q = \sum_{J_+} \lambda_j v_j v_j^T - \sum_{J_-} \lambda_j v_j v_j^T$$

$$x^T A_+ x - x^T A_- x + a^T x \leq b \quad (8)$$

$$\Rightarrow x^T (A + V_- \Gamma V_-^T) x + a^T x \leq b + \underbrace{x^T (V_- \Gamma V_-^T) x}_{\text{diag}(\Gamma)^T y} \quad (9)$$

- ▶ Then we do not need to add more $\{y\}$ for this constraint.
- ▶ The only question left is whether we can find Γ such that,

$$V_- \Gamma V_-^T - A_- \succeq 0 \quad (10)$$

- ▶ This is weaker than simultaneous diagonalization via congruence.

Future Work: ADMM in the original x I

Notice,

$$\|x\|^2 = \max_{\|\xi\| \leq \sqrt{s}} \xi^T x \quad (11)$$

So we add slack variable s, t, ξ and bilinear constraint.

$$(\mathbf{MSC}) \quad \text{Maximize : } y_0^T \lambda_0 \quad (12)$$

$$\text{s.t. } (y, z, x) \in \Omega \quad (13)$$

$$y_i^T e \leq t \quad i = 0, \dots, m \quad (14)$$

$$(\kappa) \quad t = s \quad i = 0, \dots, m \quad (15)$$

$$(\mu) \quad \xi^T x = t \quad (16)$$

$$\xi^T \xi \leq s \quad (17)$$

If s, t, ξ, y, z, x is the solution, then y, z, x is the optimal solution for MSC.

This allows the augmented Lagrangian function,

$$\mathcal{L}(x, y, z, \xi, s, \kappa, \mu) = -y_0^T \lambda_0 + \kappa(t - s) + \mu(\xi^T x - t) + \frac{\rho}{2}(t - s)^2 + \frac{\rho}{2}(\xi^T x - s)^2$$

Future Work: ADMM in the original x II

The ADMM iteration,

$$\begin{aligned}(x, y, z, t)^{k+1} &= \arg \min_{(x, y, z) \in \Omega, t \geq 0} L(x, y, z, \xi^k, s^k, \kappa^k, \mu^k) \\(s, \xi)^{k+1} &= \arg \min_{(s, \xi) \in \mathcal{Q}} L((x, y, z, t)^{k+1}, \xi, s, \kappa^k, \mu^k) \\ \kappa^{k+1} &= \kappa^k + \rho(t^{k+1} - s^{k+1}) \\ \mu^{k+1} &= \mu^k + \rho(\langle \xi^{k+1}, x^{k+1} \rangle - s^{k+1})\end{aligned}$$

where $\mathcal{Q}(\cdot)$ forms a simple SOCP for s, ξ ,

$$\mathcal{Q}(x) = \{(s, \xi) : \|\xi\|^2 \leq s\} \quad (18)$$

Future Work: ADMM in the original x III

The size of auxillary variable ξ equals to n , if r is small,
We can actually shrink the size of above problem.

$$(\mathbf{MSC}) \quad \text{Maximize : } y_0^T \lambda_0 \quad (19)$$

$$\text{s.t. } (y, z, x) \in \Omega \quad (20)$$

$$y_i^T e = s \quad i = 0, \dots, m \quad (21)$$

$$(\mu_i, i \in I_+) \quad \xi_i \cdot (v_i^T x) = y_i \quad i \in I_+ \quad (22)$$

$$\xi_i^2 \leq y_i, (v_i^T x)^2 \leq y_i \quad (23)$$

The ADMM,

$$\mathcal{L} = y^T \lambda + \sum_{i \in I_+} \mu_i [\xi_i \cdot (v_i^T x) - y_i] + \frac{\rho}{2} \sum_{i \in I_+} (\xi_i \cdot (v_i^T x) - y_i)^2 \quad (24)$$

Which is easier than old ADMM for general case.