A QCQP Solver Strengthen SOCP relaxation

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QCQP

Recall QCQP:

Maximize
$$x^{T}Qx + q^{T}x$$

s.t. $x^{T}A_{i}x + a_{i}^{T}x (\leq, =, \geq) b_{i}$ (1)
 $0 \leq x \leq e$

 $ightharpoonup Q, A_i$ maybe indefinite

Related Literature

- **S-free** sets $P = \{x : x \in R int(S), S \text{ convex}\}$
 - ✗ Gonzalo Muñoz, '20, '21, IPCO, Maximal quadratic free sets by LP (extreme rays)
 - Bienstock, '14, SIOPT, Cutting planes for optimization over convex function over nonconvex sets. also lifting

SOCPr and aggregation (SOCP-representable) For convex inequalities $\{f_i \leq 0\}$, consider aggregation $\{\sum_i \lambda_i f_i \leq 0\}$

- ▶ Burer '17, Modaresi, Vielma' 17 two quadratic functions
- Dey' 20, quadratic and polytope. app. for bilinear: Dey et al' 19, New SOCP relaxation and branching rule for bipartite bilinear programs
- ▶ Dey et al' 21, 3 quadratic intersection.

Linear disjunction

Burer '17

For intersection with a paraboloid and quadratic function (indefinite)

$$\mathcal{Y} = \begin{cases} y = \begin{pmatrix} \tilde{y} \\ y_n \end{pmatrix} \in R^n : & \tilde{y}^T \tilde{Q} \tilde{y} + 2g^T y + f \leq 0 \end{cases}$$
Homogenize: $x = [\tilde{y}, y_n, x_{n+1}]$

$$A_0 := \begin{bmatrix} I & 0 & 0 \\ 0^T & 0 & -\frac{1}{2} \\ 0^T & -\frac{1}{2} & 0 \end{bmatrix}, \quad A_1 := \begin{bmatrix} \tilde{Q} & 0 & \tilde{g} \\ 0^T & 0 & g_n \\ \tilde{g}^T & g_n & f \end{bmatrix}, \quad H^1 := \{x : x_{n+1} = 1\}$$

$$Def \quad A_t := \begin{bmatrix} (1-t)I + t\tilde{Q} & 0 & \\ 0^T & 0 & (1-t)(-\frac{1}{2}) + tg_n \\ t\tilde{g}^T & (1-t)(-\frac{1}{2}) + tg_n & tf \end{bmatrix}$$

$$t = 1/(1-\lambda), \lambda = \min(\lambda(Q)) < 0$$

$$\mathcal{Y}_t = \begin{cases} y : \tilde{y}^T \tilde{y} \leq y_n, x^T A_t x \leq 0 \\ \end{cases}$$

$$\Rightarrow \mathcal{Y}_t = \text{cl.convex.hull}(\mathcal{Y})$$

Cutting planes for QCQP

- ▶ use SOCPr to improve QCQP relaxations
- consider cases for box constraint (so as to incorporate Branch and Cut)