

# Buggy Robot

SER 2016

(Problem also available on Kattis)

For spoilers, see:

<https://github.fit.edu/ballard2014/BuggyRobot>

# Problem Statement

- You Control a robot on a grid
- Goal: reach the goal coordinate
- Some squares are blocked off by obstacles
- You are given a sequence of commands (U,D,L,R) which are not guaranteed to bring your robot to the goal
- What's the minimum number of changes to this sequence which brings the robot to the goal?

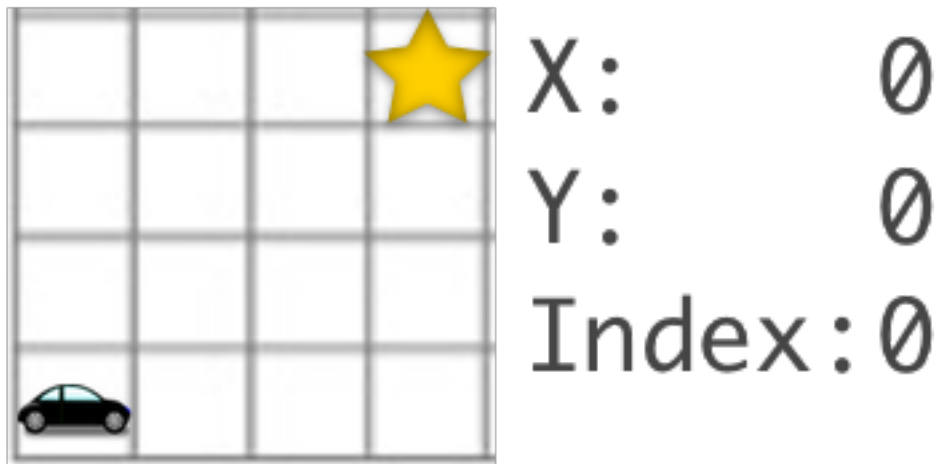
# Input:

- A rectangular map containing obstacles
- A start coordinate
- A goal coordinate
- A string of commands
- Constraints:
  - Map is  $\leq 50 \times 50$
  - Command length  $\leq 50$

# Rules

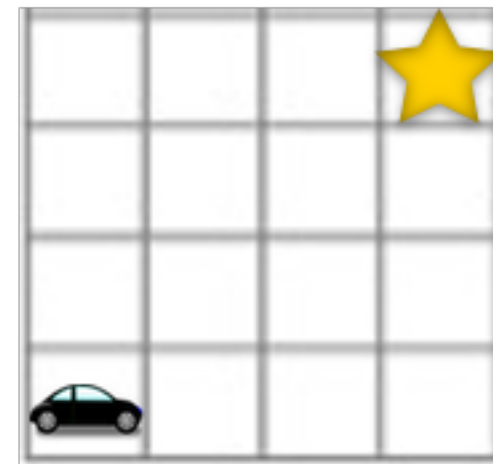
- A single change consists of:
  - Inserting a command (U,D,L,R)
  - Deleting a command
- The robot will not obey commands asking it to:
  - Run into an obstacle
  - Run off the map's boundary

# Executing Commands



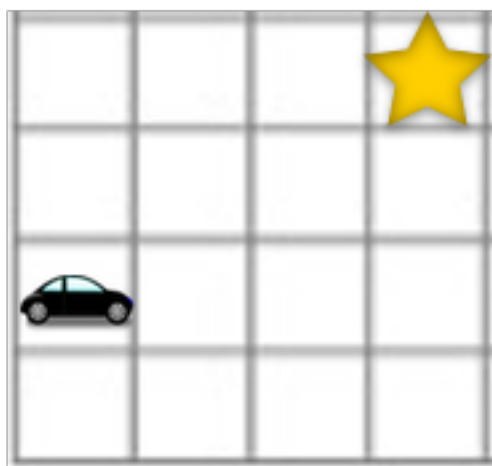
X: 0  
Y: 0  
Index: 0

Commands: DUDLRR



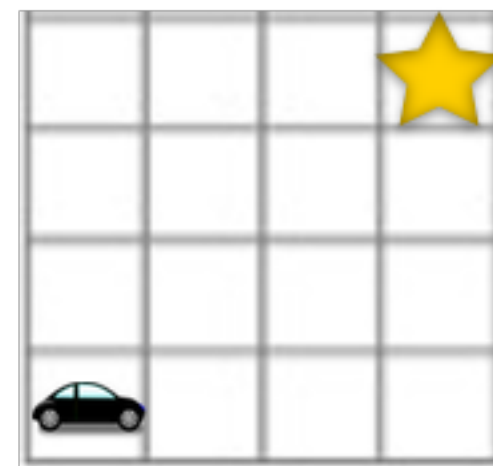
X: 0  
Y: 0  
Index: 1

Commands: DUDLRR



X: 0  
Y: 1  
Index: 2

Commands: DUDLRR

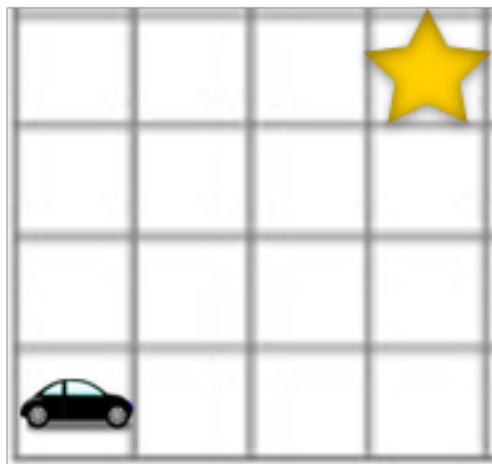


X: 0  
Y: 0  
Index: 3

Commands: DUDLRR

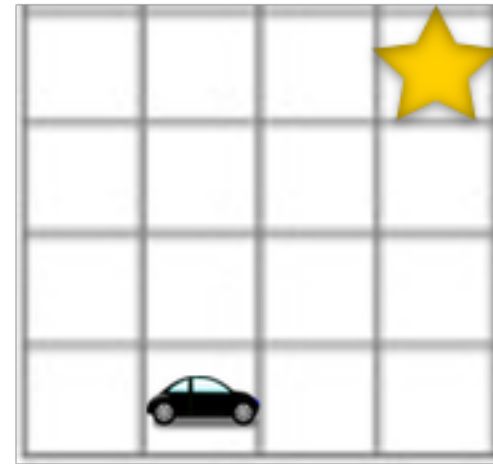


# Executing Commands



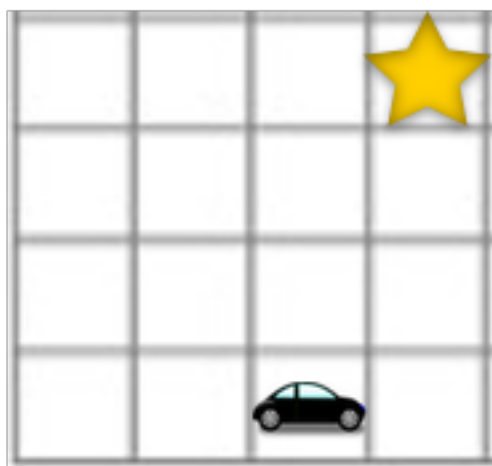
X: 0  
Y: 0  
Index: 4

Commands: DUDLRR



X: 1  
Y: 0  
Index: 5

Commands: DUDLRR



X: 2  
Y: 0  
Index: 6

Commands: DUDLRR



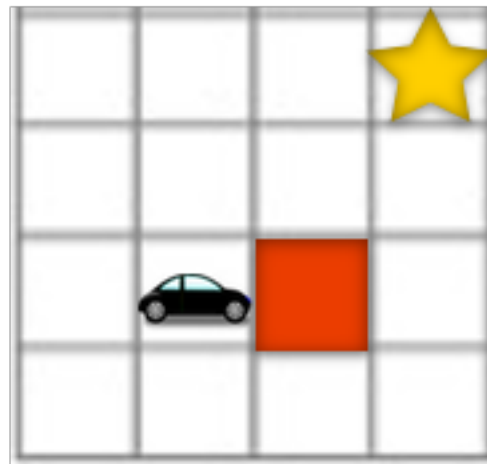
**Done!**

# Observations

- The previous example has a special property:
  1. There's no benefit to appending commands anywhere but the end of the string
- Does this always hold?

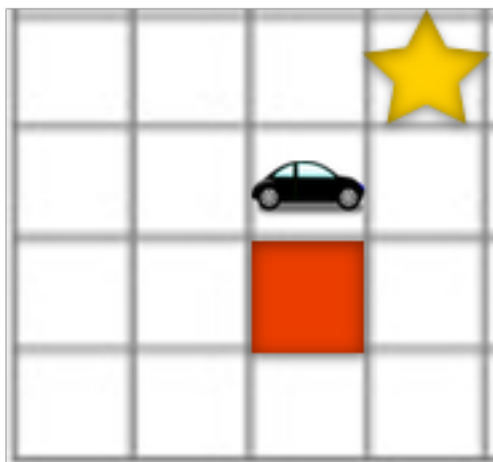
# Counter Example

Starting Point —>



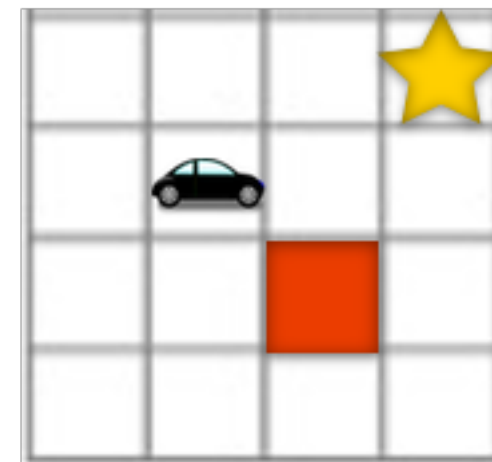
X: 1  
Y: 1  
Index: 0

Commands: D  
↑



X: 2  
Y: 2  
Index: 3

Commands: URD



X: 1  
Y: 2  
Index: 3

Commands: DUU





# More Thoughts

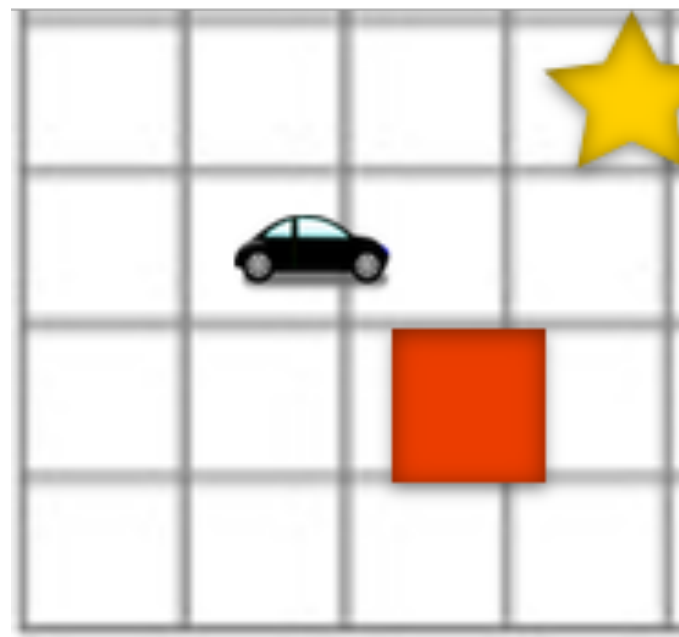
- For each command we can add, we have to consider what would happen if we add it anywhere in the current string
  - Large search space
- This isn't a textbook pathfinding problem, but we can borrow some ideas from  $A^*$ , Dijkstra, etc.

# More Thoughts

- Cost Function: The fewer changes we make to the command sequence, the better
- Secondary Cost Function: Manhattan Distance
- Greedy Strategy: Consider 'Universes' in the order of increasing changes (priority queue?)
- Terminate upon reaching the goal

# A Universe

- It will be useful to define 'Universe' objects, which store the robot's state as we follow different possible sequences of altering the command string
- Example:



X: 1  
Y: 2  
Index: 3  
Changes: 0

# Brute Force

Note: priority = # of changes

```
PriorityQueue<Universe> q = empty;
Command[] commands = given;
Commands[] commandSet = {up, down, left, right};

q.add(new Universe(startCoordinate, 0, 0))
while (true) {
    Universe u = q.remove();
    if atGoal(u) return u.changes;
    if (u.commandIndex < commands.length) {
        Coordinate coordAfter = coordinateAfter(u.coord, commands[u.index]);
        q.add(new Universe( coordAfter, u.index+1, u.changes ));
        q.add(new Universe( u.coord, u.index+1, u.changes+1));
    }
    for (Command command : commandSet) {
        Coordinate coordAfter = coordinateAfter(u.coord, command);
        q.add(new Universe( coordAfter, u.index, u.changes+1 ));
    }
}
```

# Analysis

- Complexity:  $L * 5^n$ 
  - L: length of command string
  - n: Solution (min # of changes to reach goal)
- This solution introduces a lot of redundancy, doubling the remaining problem size with every redundant Universe creation

# Eliminating Redundancy

- LPA\*: Lifelong Planning Algorithm
  - behavior like  $A^*$
  - records the cost of the best method for reaching any given point in map
- We can avoid redundancy by tracking the lowest cost of any (x position, y position, command index) tuple

```
Map<Tuple<Coordinate, Integer>> lowestCosts = *default all values to infinity*;
lowestCosts(startCoordinate, 0) = 0;
q.add(new Universe(startCoordinate, 0, 0))
while (true) {
    Universe u = q.remove();
    if atGoal(u) return u.changes;
    if u.changes > lowestCosts(u.coordinate, u.index)
        continue;
    if (u.commandIndex < commands.length) {
        Coordinate coordAfter = coordinateAfter(u.coord, commands[u.index]);
        if (lowestCosts(coordAfter, u.index+1) > u.changes) {
            q.add(new Universe( coordAfter, u.index+1, u.changes ));
            lowestCosts(u.coordinate, u.index+1) = u.changes;
        }
        if (lowestCosts(u.coordinate, u.index+1) > u.changes+1) {
            q.add(new Universe( u.coord, u.index+1, u.changes+1));
            lowestCosts(u.coordinate, u.index+1) = u.changes+1;
        }
    }
}
for (Command command : commandSet) {
    Coordinate coordAfter = coordinateAfter(u.coord, command);
    if (lowestCosts(coordAfter, u.index) > u.changes+1) {
        q.add(new Universe( coordAfter, u.index, u.changes+1 ));
        lowestCosts(coordAfter, u.index) = u.changes+1;
    }
}
}
```

# Analysis

- Complexity (time *and* space):  $L * W * H$ 
  - L: length of command string
  - W: Width of map
  - H: Height of map



Questions, Comments?