

# Frequent Values

<http://www.spoj.com/problems/FREQUENT/>

# Problem Statement

- You are given a sequence of  **$n$  integers  $a_1, a_2, \dots, a_n$**  in non-decreasing order. In addition to that, you are given several **queries consisting of indices  $i$  and  $j$  ( $1 \leq i \leq j \leq n$ )**. For each query, determine the most frequent value among the integers  $a_i, \dots, a_j$ .
- That's what it says, at least. The example question shows that what they really want is **the number of occurrences of the most frequent value in the query range**

# Input

- The input consists of several test cases
  - Each test case starts with a line containing two integers  $n$  and  $q$  ( $1 \leq n, q \leq 100000$ )
  - The next line contains  $n$  integers  $a_1, \dots, a_n$  ( $-100000 \leq a_i \leq 100000$ , for each  $i \in \{1, \dots, n\}$ ) separated by spaces
  - You can assume that for each  $i \in \{1, \dots, n-1\}$ :  $a_i \leq a_{i+1}$
  - The following  $q$  lines contain one **query** each, consisting of two integers  $i$  and  $j$  ( $1 \leq i \leq j \leq n$ ), which indicate the boundary indices for the query.
  - The final test case is followed by a zero
- 
- Upper limit for  $n, q$  is about 100,000. Rules out solutions of  $n^2$  or worse

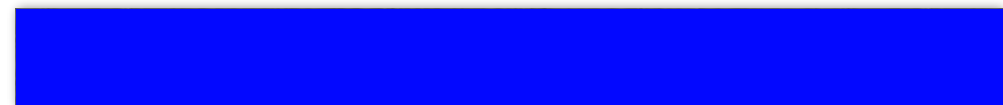
# Example Problem

10 3  
-1 -1 1 1 1 1 3 10 10 10  
2 3  
1 10  
5 10  
0



1  
4  
3

Index 1 2 3 4 5 6 7 8 9 10



Value -1 -1 1 1 1 1 3 10 10 10

# Analysis

- There is an obvious  $n^2$  solution, but doesn't work well on large input sizes
- As seen in the previous graphic, queries can be intuitively modeled as intervals over a number line

- So can repeated values!

-1 -1 1 1 1 1 3 10 10 10      →      [1,2] [3,6] [7,7] [8,10]

- Solutions are the maximum length of an intersection between a query interval and the repeated value intervals

# Interval Tree

- Interval trees are optimized to answer the question “Which intervals, from a set, intersect a given interval?”
- Balanced binary tree
- In addition to children, each node contains:
  1. A point on the number line
  2. A List of Intervals in the set which intersect with that point
- If a node has point ‘p’, its left child must have a point less than p, and its right child must have a point greater than p

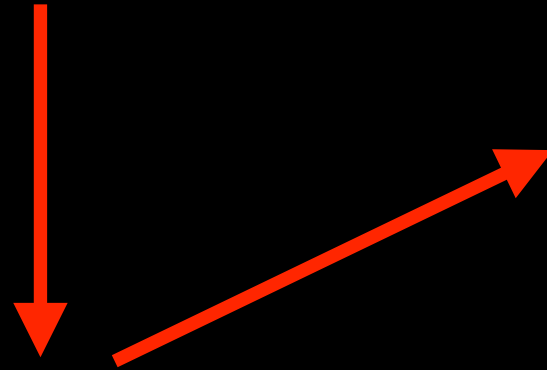
# Interval Tree - Construction

- Interval trees act like quick sort, partitioning space by repeatedly picking a point, and sorting all intervals into three categories:
  1. Completely left of the point
  2. Completely right of the point
  3. Intersecting that point
- Eventually, all intervals intersect one of the points in the tree
- Selecting these points effectively is important
  - one strategy: average of all endpoints

```

Node root;
IntervalTree(LinkedList<Interval> intervals) {
    if (!intervals.isEmpty()) // avoid div by 0 error when getting average
        root = new Node(intervals); //Node constructor handles heavy lifting
}

```



```

static class Node {

    LinkedList<Interval> ltor = new LinkedList<>();
    LinkedList<Interval> rtol = new LinkedList<>();
    Node r;
    Node l;
    int point;
}

```

```

public Node(LinkedList<Interval> intervals) {
    //get the average of the endpoints
    this.point = averageOfEndpoints(intervals);
    //sort intervals into three categories:
    LinkedList<Interval> left = new LinkedList<>();
    LinkedList<Interval> right = new LinkedList<>();
    while (!intervals.isEmpty()) {
        Interval i = intervals.remove();
        if (i.containsPoint(point)) {
            ltor.add(i);
            rtol.add(i);
        }
        else if (i.r < point) {
            left.add(i);
        }
        else /*if (i.l > point)*/ {
            right.add(i);
        }
    }
    //sort the intersecting list by
    // increasing left endpoint
    // decreasing right endpoint
    ltor.sort((Interval a, Interval b) ->
        Integer.compare(a.l, b.l));
    rtol.sort((Interval a, Interval b) ->
        -Integer.compare(a.r, b.r));

    //create children
    if (!right.isEmpty()) r = new Node(right);
    if (!left.isEmpty()) l = new Node(left);
}

```

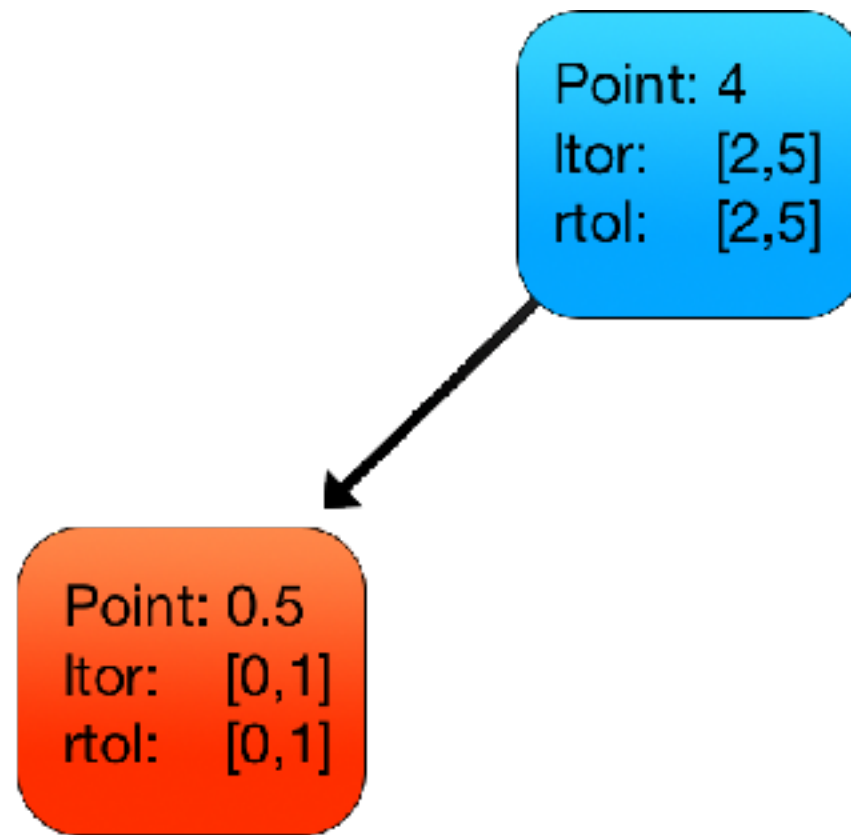


# Interval Tree - Construction

Point: 4  
ltor: [2,5]  
rtol: [2,5]

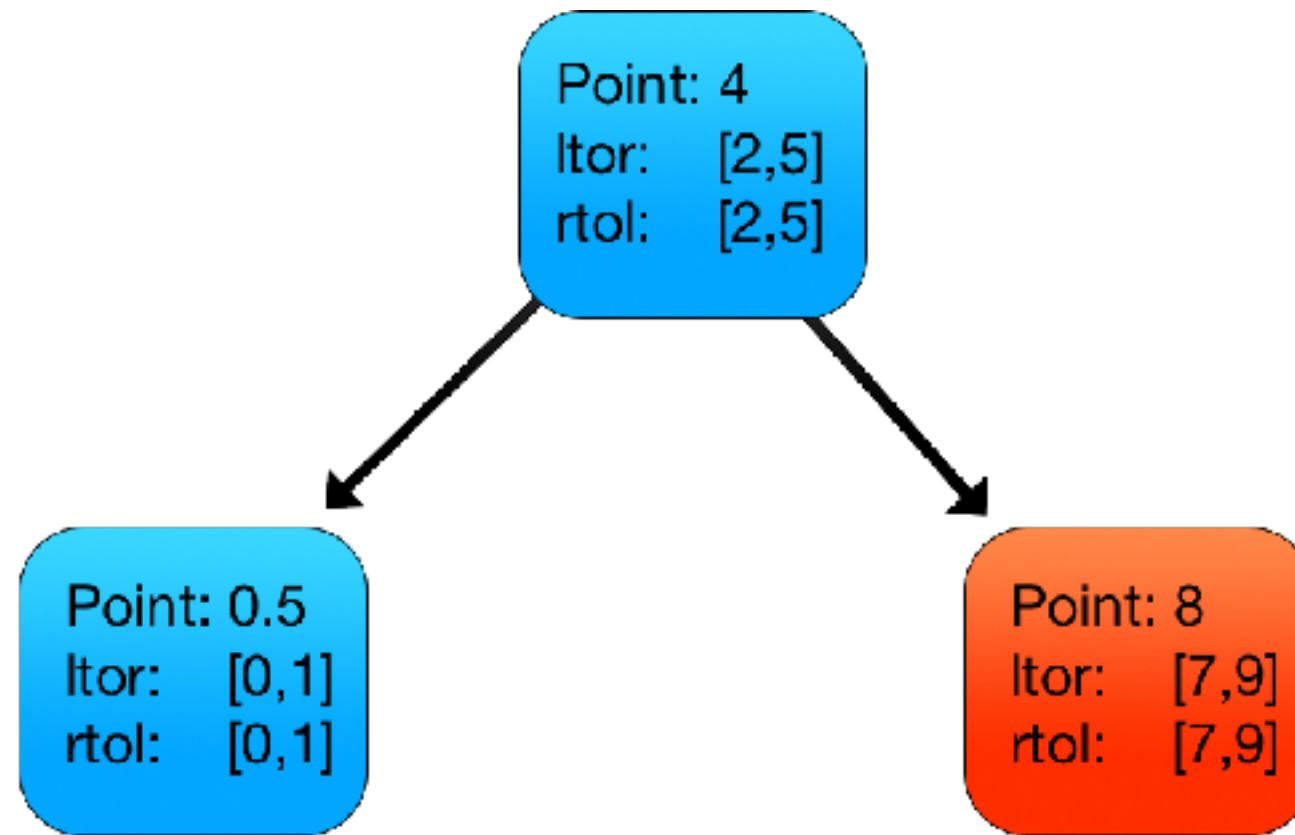
[0,1] [2,5] [7,9]

# Interval Tree - Construction



[0,1]

# Interval Tree - Construction

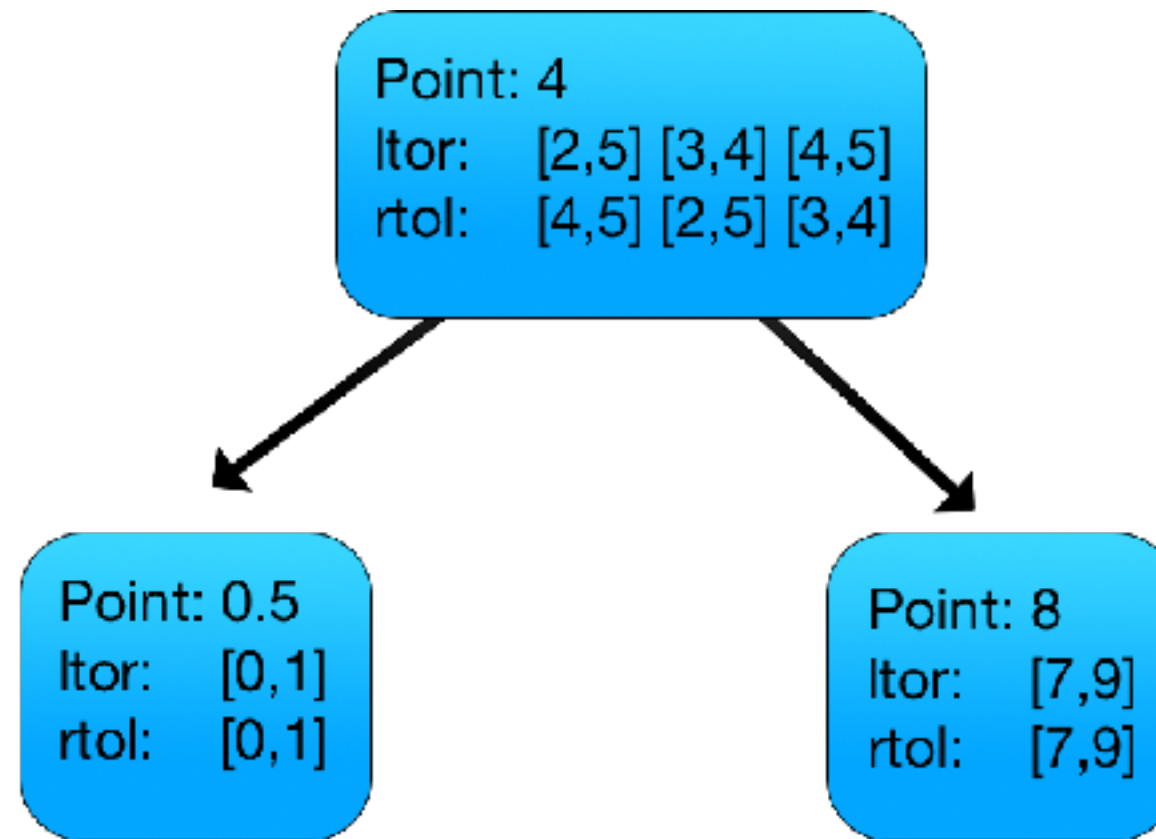


[7,9]

# Interval Tree - Construction

- Why is it necessary to have **l<sub>tor</sub>** and **r<sub>tol</sub>**?
- The following example and the querying code will make it clear

# Interval Tree - Construction 2



[0,1] [2,5] [3,4] [4,5] [7,9]

# Interval Tree - Querying

- Given an interval, find all interval in the tree which intersect it
  1. Get intervals at a node
    - If you intersect a node's point, you can add all its intervals to your set
    - If not, peel away from ltor or rtol, depending on which side of the point the interval is on
  2. Recurse through child..
    - R if the interval extends to the right of point
    - L if the interval extends to the left of point
- Efficiency is  $O(\log(n) + m)$ 
  - $n$ : intervals in tree
  - $m$ : intervals in tree that intersect your query

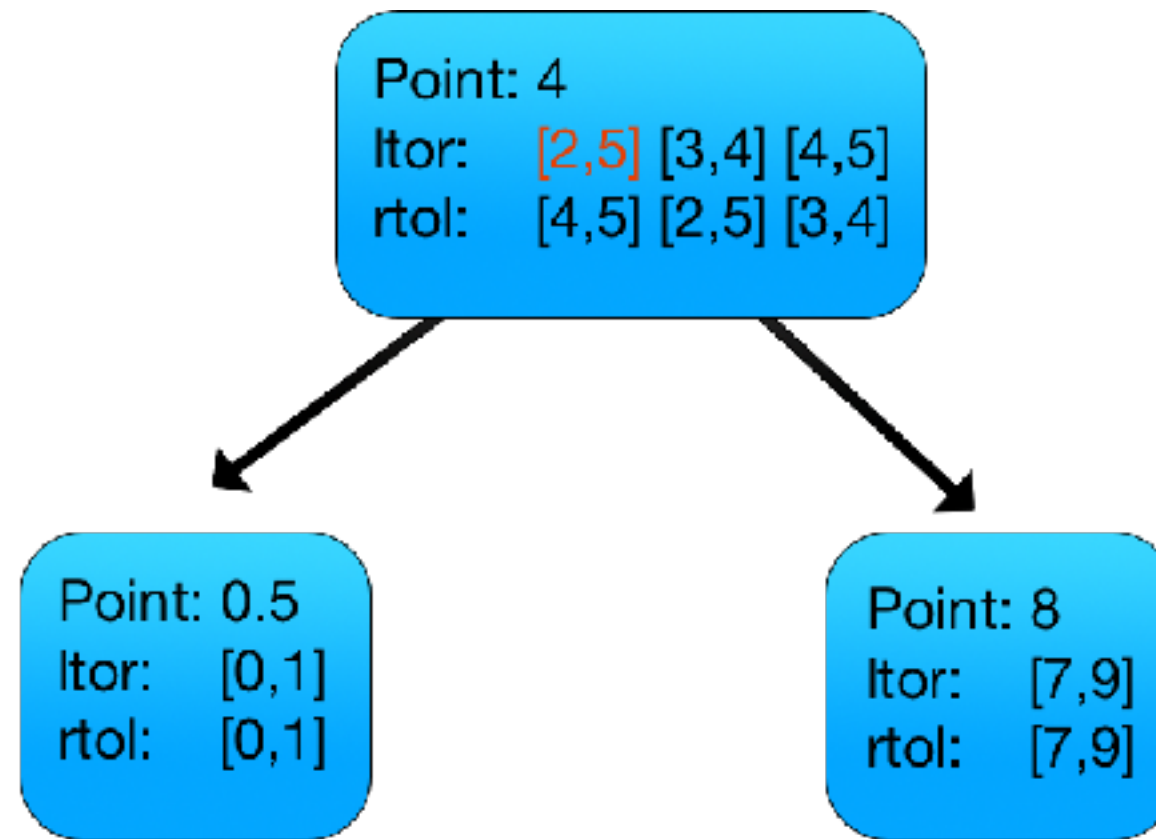
```

LinkedList<Interval> intersectingIntervals(Interval i) {
    LinkedList<Interval> set = new LinkedList<>();
    addIntersectingIntervals(i, set, root);
    return set;
}

void addIntersectingIntervals(Interval i, LinkedList<Interval> set, Node current) {
    if (current != null) {
        if (i.containsPoint(current.point)) {
            set.addAll(current.ltor); //same content as rtol, just a different ordering
        }
        else {
            ListIterator<Interval> it;
            if (current.point < i.l) //interval is to the right
                it = current.rtol.listIterator();
            else //interval is to the left
                it = current.ltor.listIterator();
            while (it.hasNext()) {
                Interval temp = it.next();
                if (i.intersects(temp)) set.add(temp);
                else break; //nothing after this one will intersect
            }
        }
        if (i.l < current.point)
            addIntersectingIntervals(i, set, current.l);
        if (i.r > current.point)
            addIntersectingIntervals(i, set, current.r);
    }
}

```

# Interval Tree - Querying

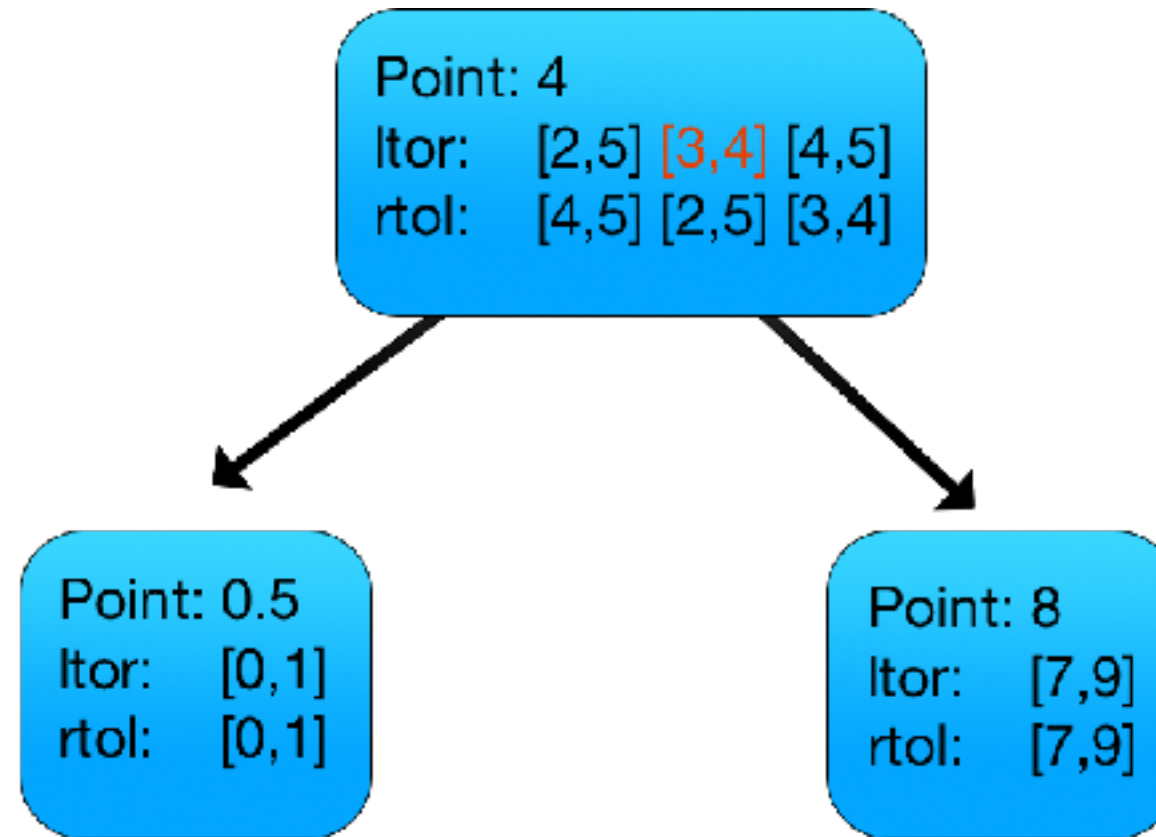


Query: [1,2]  
Solution Set: [2,5]



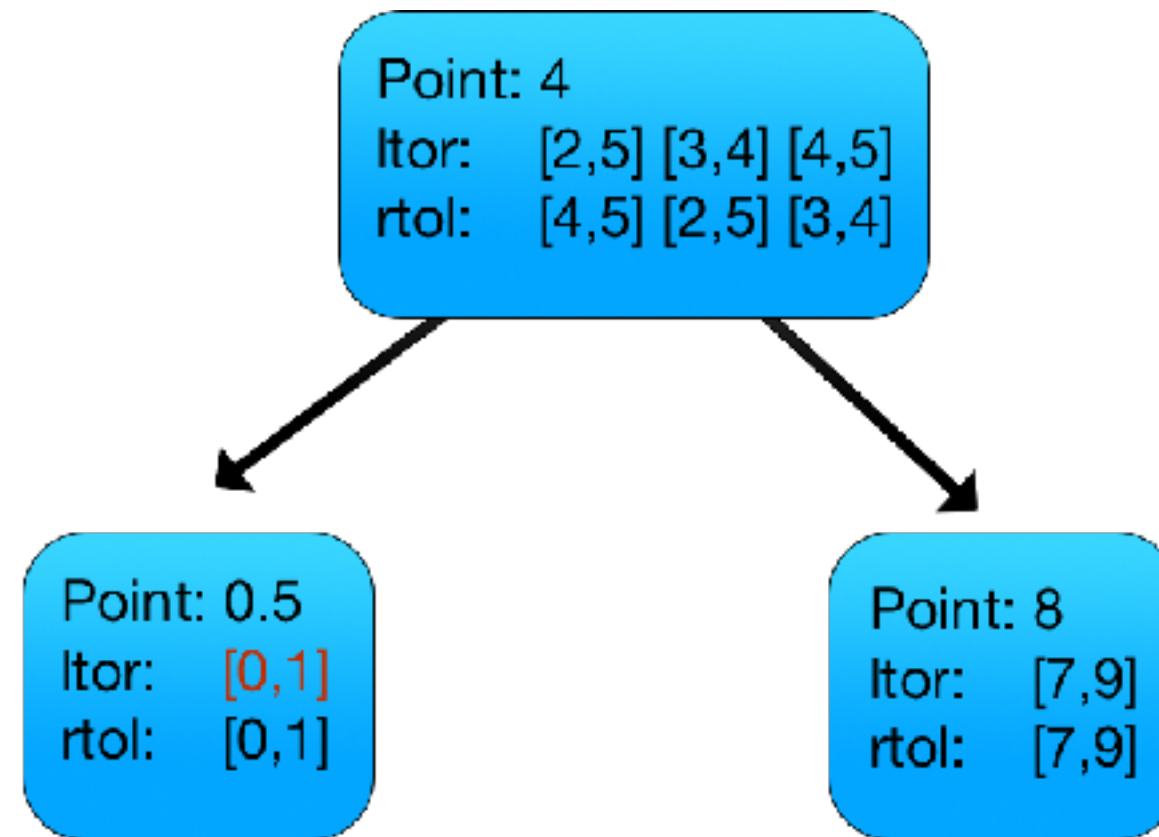
# Interval Tree - Querying

Because  $[1,2]$  and  $[3,4]$  don't intersect, we don't need to check  $[4,5]$ !



Query:  $[1,2]$   
Solution Set:  $[2,5]$

# Interval Tree - Querying



Query: [1,2]

Solution Set: [2,5] [0,1]

# Problem-Specific Notes

- No two intervals that go in the tree overlap
  - makes `rtol` and `ltor` unnecessary
- Once you get the set of intervals that overlap your query, you need to find the size of their intersection
  - example: '5' repeats through interval  $[4,9]$ , but my query only considers  $[3,7]$ . The correct answer is 4, not 6
- You know that all solutions are  $\geq 1$ . You can halve the time of building the tree by only constructing it after throwing out all intervals of length one

# Main() Code

```
public static void main (final String[] args) throws IOException {
    BufferedReader in = new BufferedReader(new InputStreamReader(System.in));
    BufferedWriter out = new BufferedWriter(new OutputStreamWriter(System.out));
    while (true) {
        //get the data
        StringTokenizer st = new StringTokenizer(in.readLine());
        int size = Integer.parseInt(st.nextToken());
        if (size == 0) break; //end of tests
        int queries = Integer.parseInt(st.nextToken());
        int[] values = new int[size];
        st = new StringTokenizer(in.readLine());
        for (int i = 0; i < size; i++) {
            values[i] = Integer.parseInt(st.nextToken());
        }
        //get all nontrivial (size > 1) intervals of consecutive values
        LinkedList<Interval> intervals = new LinkedList<>();
        int base = 0;
        while (base < size) {
            int current = base + 1;
            while (current < size) {
                if (values[current] == values[base]) current++;
                else break;
            }
            if (base - current > 1)
                intervals.add(new Interval(base, right: current-1));
            base = current;
        }
    }
}
```

# More Main() Code

```
//create the interval tree
IntervalTree tree = new IntervalTree(intervals);
//process the queries
for (int q = 0; q < queries; q++) {
    st = new StringTokenizer(in.readLine());
    int ql = Integer.parseInt(st.nextToken()) - 1;
    int qr = Integer.parseInt(st.nextToken()) - 1;
    Interval query = new Interval(ql, qr);
    LinkedList<Interval> intersecting = tree.intersectingIntervals(query);
    int bestLength = -1;
    for (Interval i : intersecting) {
        Interval intersection = query.intersection(i);
        if (intersection.length() > bestLength) {
            bestLength = intersection.length();
        }
    }
    out.write(Integer.toString(bestLength));
    out.write("str: "\n");
}
}
```