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MARIE SKODOWSKA-CURIE ACTIONS

**Individual Fellowships (IF)
Call: H2020-MSCA-IF-2015**

PART B

“SATIR”

“STATIC ANALYSIS VIA TYPE INFERENCE”

THIS PROPOSAL IS TO BE EVALUATED AS:

[GLOBAL FELLOWSHIP (GF)]

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List of Participants

Participants	Legal Entity Short Name	Academic (tick)	non-Academic (tick)	Country	Dept./ Division/ Laboratory	Supervisor	Role of partner organization
<u>Outgoing host</u>							
University of Pennsylvania	UPenn	✓		USA	Department of Computer and Information Science	Steve Zdancewic	
<u>Return host</u>							
University of Birmingham	Bham	✓		UK	School of Computer Science	Dan R. Ghica	

Sections 1,2,3 and 4 are limited to a total of 10 pages.

1 Summary

Software is becoming increasingly complex and critical (autonomous vehicles, the internet of things, financial systems, etc.). Bugs, security vulnerabilities, or inefficient resource management are costly. *Type systems* are a scalable way to analyze software but is hindered by a certain lack of expressiveness. *Graded type systems* can improve the state-of-the-art by including more information in the types and by incorporating some inspiration from *abstract interpretation*, another *static analysis* method which is extremely expressive, but as opposed to types lacks compositionality. *Compositionality* is important because it gives great support for *scalability*.

Functional programming languages often use complex type systems describing the behavior of the programs. These types are statically inferred before the compilation, acting as a first test for the correctness of the program. The advantage of type systems is their inherent compositionality. This means that libraries come with pre-computed types which are easily checkable. However, the verification power is low, restricted to the fact that composition of programs never fails due to incompatibility of format.

More refined static, whole-program analysis can generally be performed by abstract interpretation (AI). Those techniques are very efficient but fail to compose. This means that any new analysis has to be performed through the whole program, including libraries. Such a global analysis has two flaws: it is often too resource consuming to be used on a regular basis, and the behaviors of common libraries are approximated while it could be interesting to manually optimize their interpretation once and for all.

In parallel, several functional languages such as Coq, Agda, F* or Idris offer to associate to programs certain proofs of correctness encoded in a dependent type system. These dependent type systems are extremely expressive: not only can they catch up to any abstract interpretation analyses but they can describe many more refined behaviors. Moreover, these systems are still compositional and can produce certificates. However, the automatic inference is hardly ever automatisable, so the programmer has to produce the proofs mostly manually, even though the system can provide some help.

Notice, also, that performing statical analysis under richly typed languages should be advantageous: advanced type systems offer several ways to encode semantical information inside the type (for readability or for safety). However, actual tools are basically forgetting the type of a program before running the analysis, losing all this semantical information given by the programmer.

In this project, we propose to investigate a way to recover as much information as possible from type inference which is one of the most complex kinds of compositional and static analysis. We do not intend to stamp out global analyses that are proven more efficient in the general case. The expected outcome would be a static analyzer for functional languages that do not fall too far behind existing ones, while being more scalable and able to interact with other type constraints.

For this purpose, we are focusing on emerging parameterized type systems. The parameters represent either qualitative statements or quantitative information over resources. Such type systems will be called *graded types systems*.

A good point of comparison would be Hoare triples with preconditions and postconditions that need to match during composition, but some differences must be emphasized:

- A program is not only associated with a precondition and a post-condition corresponding to the parameterization of the input and output types, but also with **higher order conditions** parameterizing every subtype (especially arrow types).
- All of these conditions are dependent over “resource variables” that are instantiated by unification during the composition (adding some new constraints). In particular, this allows any program to have a **most general type** that characterizes it; this most general type is computed once and can be reused each time the program is used in a larger project.
- Constraints are not logical propositions but element of an **algebraic structure**. These structures take shape of ordered monoids or ordered semirings. They are reminiscent of abstract interpretation domains.

The SATIR project has three objectives. The **first** (and main) **objective** is to develop the theory of graded type systems in order to be as expressive as possible, to adapt to existing functional languages and to recover AI’s constructions (Galois relations and widening). The **second objective** is to track existing instances of graded type systems and to integrate them in our general framework. The **third objective** is to select a fragment of the system for the purpose of implementation: this fragment need to be inferenceable up-to the use of some widening and calls over an SMT solver (to resolve algebraic constraints).

2 Excellence

2.1 Quality, innovative aspects and credibility of the research (including inter/multidisciplinary aspects)

The SATIR project intends to specify a particular subclass of abstract interpretation that we call graded type systems. Another way to state it is that SATIR project intends to reinject the whole power of abstract interpretation inside

a type system by substituting the choice of the domain by the choice of some more refine algebraic structures called grading structures.

Objective 1: Theory Our main objective is to build up a comprehensive theory of graded types in its most general version. This objective can be split in three successive goals. **Goal 1.1** consists in identifying the exact nature of grading structures. **Goal 1.2** aims at transporting abstract interpretation principles to graded type systems. Finally, **Goal 1.3** extends the result in several directions (generalizations to a full fledged language and to intersection type systems).

Objective 2: Applications The second objective pursues real applications. Indeed, we intend to spend an important time looking for existing related work and possible applications. The intent behind this objective is both to get some inspiration for the general case and also to disseminate our results. There are three phases, or goals, regarding the applications. **Goal 2.1** aims at studying and internalizes actual fully formed instances of graded type systems such as those of BLL, D_ℓ PCF or DFuzz.¹ **Goal 2.2** focuses on qualitative examples and **Goal 2.3** focuses on quantitative ones.

Objective 3: Implementation The third objective is the delivery of an implementation for the end of the three years project. This line of work also takes into account the studies relative to the whole type inference process which form **Goal 3.1**. The implementation itself is the **Goal 3.2**, we are aiming for a simplified functional language where the user only provides the grading structure and heuristics for the widening. The objective also contains **Goal 3.3** which is pure learning over later-phase and future complements such as the use of SMT-solvers.

State of the art. In 2001, Patrick Cousot was writing that “the most severe restrictions [of type inference] are on the considered properties”². The situation has now changed and type systems are able to treat a large variety of problems.

Our first example was developed in late 90’s in the security community: the use of types for information-flow analyses.³ The idea consists in automatically inferring the security level of a program. For example, a program $\lambda xy.x : l.int \rightarrow l.int \rightarrow l.int$ use low security assertion and only take lowly secured input resulting in a lowly secured output. The apparent weakness of this type system is that a program may have a lot of different types, in particular $\lambda xy.x$ can also have the types $h.int \rightarrow l.int \rightarrow h.int$ or $l.int \rightarrow h.int \rightarrow l.int$. To resolve this weakness, we remark that a lowly secured assumption can always be assumed highly secured so that $l.int \rightarrow \tau$ is less precise than $h.int \rightarrow \tau$; moreover, one will never creat highly secured objects, so that we always assume that the result is lowly secured by default. With these assumptions, the main type⁴ of $\lambda xy.x$ is $l.int \rightarrow h.int \rightarrow int$.

The above is (one of) the most basic instance of BLL-like type systems that we will develop later. Nowadays, information-flow analysis has been developed further along four axes:⁵ the expressivity, the concurrency, the covert channel (modularity regarding the observation function of the attacker) and the security policy (allowing restricted declassifications). A special considetation will be applied to strenghten graded type system towerd each of these directions (WP 1.3.1, WP 2.2.2 and WP 2.3.2).

Flavien to Steve: I kinda remember that you had several papers in this area, maybe can you add a few sentences/references???

Linear types were introduced⁶ in different areas, and in particular for finer information-flow annalysis. It consists at insuring that some argument is used exactly once (or at most once), which is usefull for a more efficient compilation or for security. A more refined version of a linear types system is a multilinear type system where arguments are tagged with an integer bounding there number of use. Even more interesting, this tag can be infered automatically via type inference. Bellow is an example of a type derivation in call by name:

$$\frac{x : 1.int \vdash \lambda f.f (f x) : 2.(1.int \rightarrow int) \rightarrow int \quad x : 1.int \vdash \lambda y.y + x : 1.int \rightarrow int}{x : 3.int \vdash (\lambda f.f (f x)) (\lambda y.y + x) : int} \quad 2*1+1=3$$

The function $\lambda f.f (f x)$ has type $2.(1.int \rightarrow int) \rightarrow int$ meaning that it uses twice its argument f which, itself is using its argument once. In the end, we know that the argument x is used **3** times: once called by the function $\lambda f.f (f x)$ and twice called by the function $(\lambda y.y + x)$ (that is linear but used twice).

I did not find certain references, for me it seems folklor but I am not sure...

One of the culminative point of information-flow security area is the language DFuzz⁷. DFuzz was developed by the host team conjointly with the security team and a former Marie SKODOWSKA-CURIE fellow. This language extends over the covert channel direction by providing only differentially secured programs. Extending the idea of linearity to probabilities, programs of DFuzz are of type $r.\sigma \rightarrow \tau$, where $r \in \mathbb{R}^{\geq 0}$ is the expected value over the number of use of the argument. A novelty, here, is that the resource r can also depend over some “resource variable” corresponding, for example, to the number of loop in a recursion.

¹See paragraph “State of the art”

²Patrick Cousot. ”Abstract interpretation based formal methods and future challenges (Electronique version)”. In LNCS 2001.

³F. Nielson, H.R. Nielson and C. Hankin. “Principles of program analysis”, 2015, In Springer

⁴assuming that y have to be an integer

⁵A. Sabelfeld and A.C. Myers, “Language-based information-flow security”, 2003, In IEEE journal on selected areas in communications

⁶independantly?

⁷M. Gaboardi, A. Haeberlen, J. Hsu, A. Narayan and B.C. Pierce

- BLL and DIPCF
- BLL-based systems and Ghica/Smith
- (symbolic) refinement types
- Schopenhauer types ⁸
- petricek et al
- graded monads
- HO model checking, mezzo and shape annalysis

Methodology.

- Extension of BLL-like languages with monads and resource dependances,
- logics- and semantics-oriented,
- classification of graded algebras.
- more general types,
- operational semantics,
- categorical axiomatisation,
- concrete semantics,
- soundness/completeness

Targeted practical examples:

- information/control-flow
- HO model checking ⁹
- Sequentiality ¹⁰
- Probabilities ¹¹
- (complexity analysis) ¹²

Novelty, timeliness and quality.

2.2 Clarity and quality of transfer of knowledge/training for the development of the researcher in light of the research objectives

2.3 Quality of the supervision and the hosting arrangements

Outgoing Host Institution The University of Pennsylvania, founded in 1740 by Benjamin Franklin, is a premier institution providing doctoral education in the United States, and has offered Ph.D. programs for over one hundred years. The University has pre-eminent scholars in all of its disciplines. 21 Nobel Prizes have been awarded to University of Pennsylvania faculty and alumni, and 6 Engineering faculty are in the National Academy of Engineering.

Founded in 1972, the Department of Computer and Information Science of the University of Pennsylvania is part of the school of School of Engineering and Applied Science and is one of the birthplace of the modern computer. It was here that the ENIAC, the world's first electronic, large-scale, general-purpose digital computer, was developed in 1946. Nowadays, the Department of Computer and Information Science combines energies of professors, researchers and doctoral students resulting from different research areas: Artificial Intelligence, Graphics, Information Management, Software Principles, Systems, Theory. The faculty of the Department of Computer and Information Science counts 30 primary faculty members as well as 500 undergradates, and 300 graduate students (Masters and PhD).

Outside from the hosting group (Programming Language group), collaborations are expected with the "Logic and Computation group" (objective 1.1) and the "Security and Privacy group" (objectives 2.1 and 2.2) . Notice that these groups already share numeral active collaborations with the host group.

⁸K. Hammond, M. Hofmann, S. Jost, H.-W. Loidl. "Static determination of quantitative resource usage for higher-order programs". In ACM SIGPLAN-SIGACT POPL10.

⁹Ong, Kobayashi, Mellies, Grellois

¹⁰Ghica, Smith

¹¹DFuzz

¹²BLL, DIPCF

Outgoing Phase Group Among the different research groups, the Programming Language group (Penn PL club) works on several specific topics in programming language theory and practice related to the research topic of this project. In particular, it studies advanced type systems including subtyping, functional-oriented languages, ad hoc polymorphism, dependant types, polytypic programming, substructural type systems, language design and implementation, language-based security, metaprogramming, as well as foundations for concurrent, distributed, and mobile computing.

The PL club actually counts 3 professors (the supervisor, Benjamin C. Pierce and Stephanie Weirich) and more than 10 doctoral students and post-doc researchers. The group stand as one of the worldwide top theoretical programming language groups with yearly publications in ACM POPL.

Outgoing Phase Supervisor Steve Zdancewic is reconised as an international expert of functional programming languages and security. His research interests lie in theoretical computer science, programming language design and implementation (Coq verification of LLVM program transformations and randomized algorithms, type-directed program synthesis, linear types and GUI programming), language enforcing verification (language-based enforcement of information-flow policies, low-level code memory safety, understanding dynamic security policies, and authorization logic) as well as secure concurrent and distributed computing, functional programming languages, type theory, linear and modal logics, theorem proving and mechanized metatheory.

Steve Zdancewic received the Ph.D. in Computer Science from Cornell University in 2002 and imediatly joint the University of Pennsylvania as Assistant Professor. Hi is now full profesor in the University of Pennsylvania since 2014. He is the author of more than fifty publications in international journals and conferences.

Return Host Institution The School of Computer Science of the University of Birmingham is internationally leading in several areas, including foundations of computer science and security. The foundations group includes 9 Academic Staff and the security group 7. In REF 2014, 29% of its research output was judged in the 4* (world-leading) category, and 57% was judged 3* (internationally leading). Overall, 100% of its output is of international standard. It has a thriving research environment following its rapid research-led expansion in the last ten years, having c.60 research students, c.30 research fellows, and 36 academic staff, six regular seminar series.

Return Host Group In addition to the host supervisor, the Theory group includ numerous strong figures in fields that are particularly relevent for us: Krishnaswami, Escardo are reconised specialists of linear and dependent type systems which graded type system is a refinement. Levy, Reddy are specialists of categorical semantics of programming languages and monads in particular. Jung, Vickers (and Escardo) are specialist of domain theory and can help with the formalisation of graded algebras. Finally, Ritter is a former specialist in linear logic and is now working on security issues, as such he is a potential collaborator for Goals 2.2 and 2.3.

Return Host Supervisor (Scientist in charge) Dan Ghica

CV

2.4 Capacity of the researcher to reach and re-enforce a position of professional maturity in research

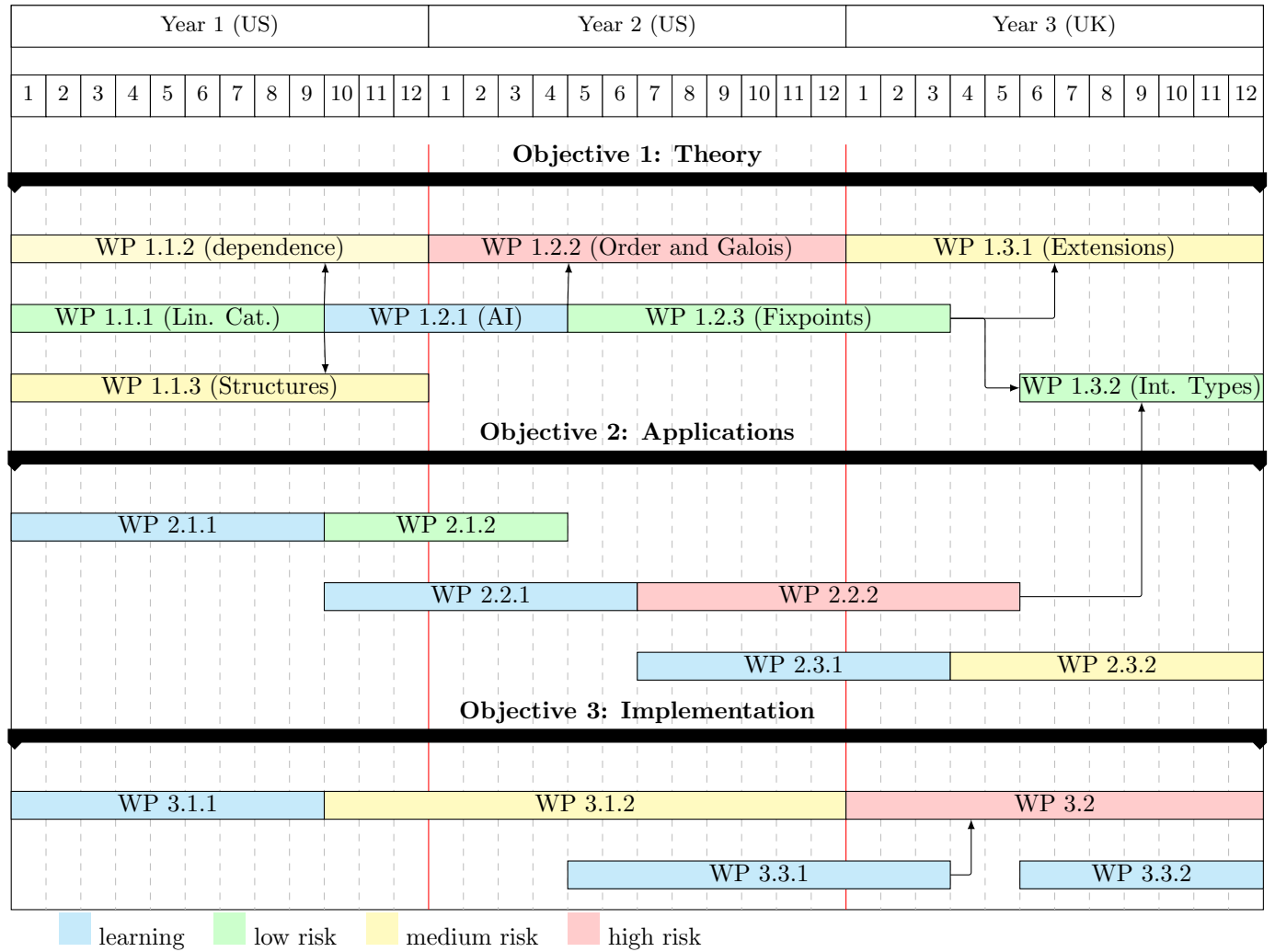
3 Impact

3.1 Enhancing research- and innovation-related human resources, skills, and working conditions to realize the potential of individuals and to provide new career perspectives

3.2 Effectiveness of the proposed measures for communication and results dissemination

4 Implementation

4.1 Overall coherence and effectiveness of the work plan, including appropriateness of the allocation of tasks and resources



The project is divided in three parallel lines. The theory line focuses on developing the main theory behind our project. The Application line has for goal to understand related fields and to apply our result to those fields. The Implementation line is targeting an implementation of our language.

Theory: The main objective follows a sequential pattern.

Goal 1.1 is supposed to be the basis for all the remaining work, it is thus scheduled the first year. Notice, moreover, that no other research-focused work package has been planned in parallel (except for the three last months that are assumed to be in final phase). Goal 1.1 has three different work packages. Work package 1.1.1 is a package that focuses on grading non conventional models of linear logics in order to investigate natural extensions of graded type systems. It is a support package for packages 1.1.2 and 1.1.3, and no important direct result are expected here. Work package 1.1.2 has the objective to specify the notion of dependence over grading structures. Work package 1.1.3 has the objective to distinguish a set of independent and interacting layers in the grading structure such as the monad-comonad separation. Early results on these two directions are quite promising and several researchers are active in

this area; thus expecting a reasonable formalisation before the two year deadline can be considered medium risk.

Goal 1.2 is at the heart of this project and investigates the AI-related features of graded type systems. It is planned the second year, but is extended with preliminaries and some remaining work. Work package 1.2.1 is a learning-focused package that aims at expanding the knowledge of the applicant regarding abstract interpretation. Work package 1.2.2 explores in more detail the order that naturally arises from grading structures in order to link it with domain theory. We expect this package to end with a notion of Galois relation between grading structures. Work package 1.2.3 is focusing on fixpoints, and in particular on the corresponding notion of widening. This package is considered low risk once the Galois connection has been properly set.

Goal 1.3 consists in optional extensions of the theoretical work. In particular, Work package 1.3.1 is aimed at extending the theory for a λ -calculus to a full fledged language and Work package 1.3.2 is investigating the extension to intersection types.

Applications: The applications line follows a classic 2-2-2 pattern corresponding to the three applicative objectives, each separated into a learning phase (9 months) and an active phase (duration varying with the difficulty level). Notice that this objective is composed only of collaborative material.

Goal 2.1 (work packages 2.1.1 and 2.2) targets the embedding of BLL, D_ℓ PCF and DFuzz in our system. It is considered low risk as these languages are well established so that we roughly have to verify that we are taking into consideration each of its components. This objective is mainly aiming at supporting Goal 1.1 with real examples. Moreover, this goal has been scheduled just after one year spent in Focus team (Bologna) that created D_ℓ PCF and in the UPen's programming language team that created DFuzz.

Goal 2.2 (work packages 2.2.1 and 2.2.2) targets the embedding of qualitative examples such as the control flow analyzes and¹³ the higher order model checking. This objective is classed high risk for several reasons: this will be the first new application and the relation between graded type systems and higher order model checking is not yet fully established. Its scheduling at middle position among applicative goals permit to manage the high risk by potentially extending the deadline. Moreover, the parallelizing of Goal 2.2 and the work package 1.2.3 is important as the fixpoint construction is fundamental in higher order model checking application. Finally, its scheduling permit a potential additional investigation over intersection type systems at the end of the project.

Goal 2.3 (work packages 2.3.1 and 2.3.2) targets the embedding of quantitative examples such as probability or scheduling. This goal is classified medium risk despite their intrinsic difficulties since the applicant has some experience with these cases. Remark that this goal is planed in Birmingham with Dan Ghica that is an expert of the scheduling example (which we believe critical).

Implementation:

Not yet fixed
 WP 3.1.1: learning type inference
 WP 3.1.2: working on type inference of graded type systems
 WP 3.2: concrete implementation
 WP 3.3.1: Learning over SMT solvers
 WP 3.3.2: optional learning

Milestones and deliverables: There are three main milestones at the end of every year and three secondary milestones around midyear.

The first secondary milestone (around month 9) does not have a deliverable but it marks the end of the adjustment and warm-up period. In particular, WP 2.1 and WP 3.1 will switch from learning phase to active file while WP 1.1 will go to a writing phase. The end of this writing phase is marked by the **Milestone 1** with 2 deliverable articles corresponding to WP 1.1.2 and 1.1.3; at least one of these article should be a collaboration with the host team, corresponding with the ongoing WP 2.1.2.

The second secondary milestone (from month 16 to month 18) is a smooth transition with an collaborative article as delivery focused on Objective 2.1 and with early results of WP 3.1.2. **Milestone 2** corresponds to the delivery of two articles. The first is a collaborative article treating WP 1.2.2 and the first advancements over WP 2.2.2. The second is a single-authored article over WP 3.1.2.

The last secondary milestone (from month 27 to month 29) is an other fuzzy transition with a collaborative article as delivery around work packages 1.2.3 and 2.2.2; this milestone will also mark the point where an engineer will be hired to work on the implementation. **Milestone 3** marks the end of the project with the full implementation of WP 3.2 as delivery as well as a collaborative article over WP 2.3.2 together with results from WP 1.3.1.

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¹³the intersection-free fragment of

4.2 Appropriateness of the management structure and procedures, including quality management and risk management

4.3 Appropriateness of the institutional environment (infrastructure)

UPen and Bhm both have much experience in helping international employees, many of whom supported by Marie Curie grants. Services are devoted to help foreign employees in finding living accommodation as well as dealing with administrative tasks due to mobility (contracts, taxes, etc.). The fellow will therefore receive help not only on his arrival, but for the whole time of his stay, whenever necessary. This will allow him to attend to all administrative aspects of Marie Curie grants, such as preparation of deliverables in a best possible form.

The administrative management will further be supported by the EU Office of Bhm which has been involved in the management of several EU projects which handles the financial and personnel aspects of national and international research projects including EU projects. The Faculty of Science was the recipient of ???Marie Curie Individual Fellowships and has great experience of project administration, leaving the fellow free to focus on scientific aspects of the project by lightening the administrative task load.

Furthermore, office space and supplies are provided by UPen's department of computer sciences, while research supplies are available through research grants. Each student is guaranteed access to a computer and the Internet from laboratories and from his or her office as well as to systems administration support for computing and data storage. The University Library System holds more than 5,000,000 books and serial volumes in hard copy and digitalized format. Penn strongly encourages students to make use of available university equipment and resources. Although space permits only this brief survey of SEAS and CIS equipment available to postdoctoral fellows, Penn has equipment and student equipment access policies which have been judged exemplary by its peers.

Similarly, Bhm will furnish office space, access to scientific papers from all major publishers, books, and computing facilities.

4.4 Competences, experience and complementarity of the participating organisations and institutional commitment

5 CV of the Experienced Researcher

Personal Informations

Name Flavien BREUVART

Birthday April 6th, 1988

Family Married

Email breuvart@pps.univ-paris-diderot.fr

Website <http://www.pps.univ-paris-diderot.fr/breuvart/>

(Future) PhD Thesis

title Could the syntax be the chicken and the semantics the egg?

subtitle From the old style H to the newborn quantitative coefficients.

supervisors Antonio Bucciarelli and Michele Pagani

defense Octobre 23th, 2015

jury Antonio Bucciarelli (Paris Diderot), Pierre-Louis Curien (CNRS), Dan Ghica (University of Birmingham), Giulio Manzonetto (University Paris 13), Guy McCusker (rapporteur, University of Bath), Michele Pagani (Paris Diderot), Simona Ronchi della Rocca (president, University of Turin), Thomas Streicher (rapporteur, Technische Universitat Darmstadt)

Position

2015-2017: **Iria post-doctoral positions**, in *Focus Team*, Bologna, Italy

2012-2015: **PhD candidate**, with Antonio Bucciarelli and Michele Pagani, at *PPS, University Paris Diderot*, Paris, France

2012-2015: **Teaching Assistant**, at *UFR Informatique, University Paris Diderot*, Paris, France

2008-2012: **Fonctionnaire stagiaire**, at *ENS Cachan*, Cachan, France

Other research experience

true?

correct name?

find the number

- 2011-2012: **12 months internship**, with Antonio Bucciarelli, *at PPS, University Paris Diderot*, Paris, France
- 2011: **5 months internship**, with Thomas Ehrhard and Antonio Bucciarelli, *at PPS, University Paris Diderot*, Paris, France
- 2010: **5 months internship**, with Peter Selinger, *at Dalhousie University*, Halifax, Canada
- 2009: **2 months internship**, with Achim Jung, *at University of Birmingham*, Birmingham, UK

Grants and Awards

- 2015: **Iria post-doctoral positions**, *in Focus Team*, Bologna, Italy
- 2014: **Kleene Award**, for best student paper, *in CSL-LICS*, Vienna, Austria
- 2012: **Three-Years PhD fellowship**, *from the french "Ministry of Higher Education and Research"*, France

Publications

Conference papers

- CSL'15: **Breuvart Flavien and Pagani Michele**, *Berlin, Germany*
"Modelling Coeffects in the Relational Semantics of Linear Logic"
- CSL-LICS'14 **Breuvart Flavien** *Vienna, Austria*
"On the characterization of models of \mathcal{H}^* "
- TLCA'13 **Breuvart Flavien** *Eindhoven, Netherlands*
"The resource lambda calculus is short-sighted in its relational model"

Submitted

- 2016 **Breuvart Flavien, Marco Gaboardi, Shin-ya Katsumata and Dominic Orchard**,
"Combining effects and coeffects"
- 2016 **Breuvart Flavien**, long version
"On the characterization of models of \mathcal{H}^* "

non peer-reviewed

- 2012 **Breuvart Flavien**, arXiv preprint arXiv:1205.4691,
"On the discriminating power of tests in resource lambda-calculus"

Responsibilities

- 2014: **workshop organization**, Participation at the organization of the international workshop Domain XI, *at PPS, University Paris Diderot*, Paris, France
- 2013: **seminar organization**, Organization of the temporary working group "Sémantique des calculs classiques", *at PPS, University Paris Diderot*, Paris, France
- 2012-2015: **Representative of PhD students**, in the Board of directors, *of the laboratory PPS*, Paris, France

Teaching

- 2015: **advanced Java**, TA for L3 course named POOIG, *at university Paris Diderot*, Paris, France
- 2014: **advanced Java**, TA for M3 course for linguists named POO3, *at university Paris Diderot*, Paris, France
- 2014 & 2015: **architecture**, TA for M1 course, *at university Paris Diderot*, Paris, France
- 2014: **"UE libre" vulgarization**, In charge of the "UE libre" (optional short course) on scientific vulgarization in M1, *at university Paris Diderot*, Paris, France
- 2013 & 2014: **mid-beginner Java**, TA for L1 course named TO2, *at university Paris Diderot*, Paris, France
- 2012: **Java projects**, Supervising Java projects in L2 (more exactly for the "classe préparatoire intégrée de l'EIDD"), *at university Paris Diderot*, Paris, France
- 2008-2012: **vulgarization**, occasional (5-10 times per years) 1 hour of vulgarization on mathematics and computer sciences for middle school students, *with the association PRACCIS*, Cachan, France
- 2008-2009: **vulgarization of game theory**, 1h/week of vulgarization of game theory to high school students, *with the association Sciences Ouvertes*, Drancy, France

Communications

- 10/09/2015 **CSL'2015**, Berlin, Germany
"Modelling Coeffects in Rel"
- 03/09/2014 **seminar of logic team**, of LAMA, *University of Savoie*, Chambéry, France
"On the characterization of models of \mathcal{H}^* "
- 17/07/2014 **CSL-LICS**, Vienna, Austria
"On the characterization of models of \mathcal{H}^* "
- 04/04/2014 **seminar of team LCR**, in LIPN, *Univerisity Paris Nort*, Paris, France
"De la caractérisation des modèles de \mathcal{H}^* "
- 09/04/2014 **LIAFA-PPS' PHD seminar**, *University Paris Diderot*, Paris, France
"Basic notions on weighted models"
- 14/03/2014 **"Theory seminar"**, School of Computer Science, *University of Birmingham*, Birmingham, UK
"A bridge between semirings"
- 03/12/2013 **BLL'workshop**, Fontainebleau, France
"A bridge between semirings"
- 27/06/2013 **TLCA**, Eindhoven, Netherlands
"The resource lambda calculus is short-sighted in its relational model"
- 15/02/2013 **GEOCAL Meeting**, Lyon, France
"Une bien étrange caractérisation de l'adéquation complète dans les domaines de Scott"
- 24/10/2012 **LIAFA-PPS' PHD seminar**, *University Paris Diderot*, Paris, France
"Petites histoires entre modèles et syntaxes"
- 12/09/2012 **PPS' anual meating**, *University Paris Diderot*, Trouville, France
"Petites histoires entre modèles et syntaxes"
- 02/09/2011 **LIAFA-PPS' PHD seminar**, *University Paris Diderot*, Paris, France
"Et si les PPSiens se mettaient à faire de l'algèbre?!!"

Education

- 2010-2011: **PHD thesis in Paris Diderot**, *Paris, France*
- 2010-2011: **M2 in MPRI**, *Cachan, France*
- 2009-2010: **M1 in MPRI**, *Cachan, France*
- 2008-2009: **L3 in computer sciences and in mathematics**, *at ENS Cachan, Cachan, France*
- 2008: **admission at ENS Cachan**, *Cachan, France*
- 2006-2008: **CPGE MPSI and MP***, *at Lycée Fénélon, Paris 5^{iem}, France*
- 2006: **French baccalaureate**, "mention bien", *at Lycée W.A.Mozart,, Le Blanc Mesnil, France*

Associative activities

- 2010-2011: **president of PRACCIS project**, in *SONO*, *Cachan, France*
- 2009-2010: **active member**, of *BDE ENS Cachan*, *with high implication in the clubs BDA, Mediatek, PRACCIS Cachan, France*
- 2007-2008: **active member** of *Sciences Ouvertes* *Drancy, France*
- 2005-2007: **board of directors** of the association *MJC Daniel André* *Drancy, France*

6 Capacities of the Participating Organisations

7 Ethical Aspects

There are no ethical sensitive issues associated with the subject of this proposal.

8 Letters of Commitment of Partner Organisations