

Data.Map.Strict

An efficient implementation of ordered maps from keys to values (dictionaries).

API of this module is strict in both the keys and the values. If you need value-lazy maps, use [Data.Map.Lazy](#) instead. The [Map](#) type is shared between the lazy and strict modules, meaning that the same [Map](#) value can be passed to functions in both modules (although that is rarely needed).

These modules are intended to be imported qualified, to avoid name clashes with Prelude functions, e.g.

```
import qualified Data.Map.Strict as Map
```

The implementation of [Map](#) is based on *size balanced* binary trees (or trees of *bounded balance*) as described by:

- Stephen Adams, "*Efficient sets: a balancing act*", Journal of Functional Programming 3(4):553-562, October 1993, <http://www.swiss.ai.mit.edu/~adams/BB/>.
- J. Nievergelt and E.M. Reingold, "*Binary search trees of bounded balance*", SIAM journal of computing 2(1), March 1973.

Bounds for [union](#), [intersection](#), and [difference](#) are as given by

- Guy Blelloch, Daniel Ferizovic, and Yihan Sun, "*Just Join for Parallel Ordered Sets*", <https://arxiv.org/abs/1602.02120v3>.

Note that the implementation is *left-biased* -- the elements of a first argument are always preferred to the second, for example in [union](#) or [insert](#).

Warning: The size of the map must not exceed `maxBound :: Int`. Violation of this condition is not detected and if the size limit is exceeded, its behaviour is undefined.

Operation comments contain the operation time complexity in the Big-O notation (http://en.wikipedia.org/wiki/Big_O_notation).

Be aware that the `Functor`, `Traversable` and `Data` instances are the same as for the [Data.Map.Lazy](#) module, so if they are used on strict maps, the resulting maps will be lazy.

Strictness properties

This module satisfies the following strictness properties:

1. Key arguments are evaluated to WHNF;
2. Keys and values are evaluated to WHNF before they are stored in the map.

Here's an example illustrating the first property:

```
delete undefined m == undefined
```

Here are some examples that illustrate the second property:

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License	BSD-style
Maintainer	libraries@haskell.org
Stability	provisional
Portability	portable
Safe	Safe
Haskell	
Language	Haskell98

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```
map (\ v -> undefined) m == undefined    -- m is not empty
mapKeys (\ k -> undefined) m == undefined -- m is not empty
```

Map type

```
data Map k a
```

Source

A Map from keys `k` to values `a`.

Instances

Functor (Map k)	# Source
Foldable (Map k)	# Source
Traversable (Map k)	# Source
Ord k => IsList (Map k v)	# Source
(Eq k, Eq a) => Eq (Map k a)	# Source
(Data k, Data a, Ord k) => Data (Map k a)	# Source
(Ord k, Ord v) => Ord (Map k v)	# Source
(Ord k, Read k, Read e) => Read (Map k e)	# Source
(Show k, Show a) => Show (Map k a)	# Source
Ord k => Semigroup (Map k v)	# Source
Ord k => Monoid (Map k v)	# Source
(NFData k, NFData a) => NFData (Map k a)	# Source
type Item (Map k v)	# Source

Operators

```
(!) :: Ord k => Map k a -> k -> a
```

infixl 9

Source

$O(\log n)$. Find the value at a key. Calls **error** when the element can not be found.

```
fromList [(5,'a'), (3,'b')] ! 1    Error: element not in the map
fromList [(5,'a'), (3,'b')] ! 5 == 'a'
```

```
(\\) :: Ord k => Map k a -> Map k b -> Map k a
```

infixl 9

Source

Same as **difference**.

Query

```
null :: Map k a -> Bool
```

Source

$O(1)$. Is the map empty?

```
Data.Map.null (empty)      == True
Data.Map.null (singleton 1 'a') == False
```

```
size :: Map k a -> Int
```

Source

$O(1)$. The number of elements in the map.

```
size empty == 0
size (singleton 1 'a') == 1
size (fromList [(1,'a'), (2,'c'), (3,'b')]) == 3
```

```
member :: Ord k => k -> Map k a -> Bool
```

[# Source](#)

$O(\log n)$. Is the key a member of the map? See also `notMember`.

```
member 5 (fromList [(5,'a'), (3,'b')]) == True
member 1 (fromList [(5,'a'), (3,'b')]) == False
```

```
notMember :: Ord k => k -> Map k a -> Bool
```

[# Source](#)

$O(\log n)$. Is the key not a member of the map? See also `member`.

```
notMember 5 (fromList [(5,'a'), (3,'b')]) == False
notMember 1 (fromList [(5,'a'), (3,'b')]) == True
```

```
lookup :: Ord k => k -> Map k a -> Maybe a
```

[# Source](#)

$O(\log n)$. Lookup the value at a key in the map.

The function will return the corresponding value as (`Just` value), or `Nothing` if the key isn't in the map.

An example of using lookup:

```
import Prelude hiding (lookup)
import Data.Map

employeeDept = fromList [("John","Sales"), ("Bob","IT")]
deptCountry = fromList [("IT","USA"), ("Sales","France")]
countryCurrency = fromList [("USA", "Dollar"), ("France", "Euro")]

employeeCurrency :: String -> Maybe String
employeeCurrency name = do
    dept <- lookup name employeeDept
    country <- lookup dept deptCountry
    lookup country countryCurrency

main = do
    putStrLn $ "John's currency: " ++ (show (employeeCurrency "John"))
    putStrLn $ "Pete's currency: " ++ (show (employeeCurrency "Pete"))
```

The output of this program:

```
John's currency: Just "Euro"
Pete's currency: Nothing
```

```
findWithDefault :: Ord k => a -> k -> Map k a -> a
```

[# Source](#)

$O(\log n)$. The expression (`findWithDefault` def k map) returns the value at key k or returns default value def when the key is not in the map.

```
findWithDefault 'x' 1 (fromList [(5,'a'), (3,'b')]) == 'x'
findWithDefault 'x' 5 (fromList [(5,'a'), (3,'b')]) == 'a'
```

```
lookupLT :: Ord k => k -> Map k v -> Maybe (k, v)
```

[# Source](#)

$O(\log n)$. Find largest key smaller than the given one and return the corresponding (key, value) pair.

```
lookupLT 3 (fromList [(3,'a'), (5,'b')]) == Nothing
lookupLT 4 (fromList [(3,'a'), (5,'b')]) == Just (3, 'a')
```

lookupGT :: Ord k => k -> Map k v -> Maybe (k, v) # Source

$O(\log n)$. Find smallest key greater than the given one and return the corresponding (key, value) pair.

```
lookupGT 4 (fromList [(3,'a'), (5,'b')]) == Just (5, 'b')
lookupGT 5 (fromList [(3,'a'), (5,'b')]) == Nothing
```

lookupLE :: Ord k => k -> Map k v -> Maybe (k, v) # Source

$O(\log n)$. Find largest key smaller or equal to the given one and return the corresponding (key, value) pair.

```
lookupLE 2 (fromList [(3,'a'), (5,'b')]) == Nothing
lookupLE 4 (fromList [(3,'a'), (5,'b')]) == Just (3, 'a')
lookupLE 5 (fromList [(3,'a'), (5,'b')]) == Just (5, 'b')
```

lookupGE :: Ord k => k -> Map k v -> Maybe (k, v) # Source

$O(\log n)$. Find smallest key greater or equal to the given one and return the corresponding (key, value) pair.

```
lookupGE 3 (fromList [(3,'a'), (5,'b')]) == Just (3, 'a')
lookupGE 4 (fromList [(3,'a'), (5,'b')]) == Just (5, 'b')
lookupGE 6 (fromList [(3,'a'), (5,'b')]) == Nothing
```

Construction

empty :: Map k a # Source

$O(1)$. The empty map.

```
empty == fromList []
size empty == 0
```

singleton :: k -> a -> Map k a # Source

$O(1)$. A map with a single element.

```
singleton 1 'a' == fromList [(1, 'a')]
size (singleton 1 'a') == 1
```

Insertion

insert :: Ord k => k -> a -> Map k a -> Map k a # Source

$O(\log n)$. Insert a new key and value in the map. If the key is already present in the map, the associated value is replaced with the supplied value. **insert** is equivalent to **insertWith const**.

```
insert 5 'x' (fromList [(5,'a'), (3,'b')]) == fromList [(3, 'b'), (5, 'x')]
insert 7 'x' (fromList [(5,'a'), (3,'b')]) == fromList [(3, 'b'), (5, 'a'), (7, 'x')]
insert 5 'x' empty == singleton 5 'x'
```

```
insertWith :: Ord k => (a -> a -> a) -> k -> a -> Map k a -> Map k a
```

```
| # Source
```

$O(\log n)$. Insert with a function, combining new value and old value. `insertWith f key value mp` will insert the pair (key, value) into mp if key does not exist in the map. If the key does exist, the function will insert the pair (key, `f new_value old_value`).

```
insertWith (++) 5 "xxx" (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "xxx")]
insertWith (++) 7 "xxx" (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "xxx")]
insertWith (++) 5 "xxx" empty == singleton 5 "xxx"
```

```
insertWithKey :: Ord k => (k -> a -> a -> a) -> k -> a -> Map k a -> Map k a
```

```
Source
```

```
| #
```

$O(\log n)$. Insert with a function, combining key, new value and old value. `insertWithKey f key value mp` will insert the pair (key, value) into mp if key does not exist in the map. If the key does exist, the function will insert the pair (key, `f key new_value old_value`). Note that the key passed to `f` is the same key passed to `insertWithKey`.

```
let f key new_value old_value = (show key) ++ ":" ++ new_value ++ "|" ++ old_value
insertWithKey f 5 "xxx" (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "xxx|a")]
insertWithKey f 7 "xxx" (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "xxx|a")]
insertWithKey f 5 "xxx" empty == singleton 5 "xxx|a"
```

```
insertLookupWithKey :: Ord k => (k -> a -> a -> a) -> k -> a -> Map k a -> (Maybe a, Map k a)
```

```
| # Source
```

$O(\log n)$. Combines insert operation with old value retrieval. The expression `(insertLookupWithKey f k x map)` is a pair where the first element is equal to `(lookup k map)` and the second element equal to `(insertWithKey f k x map)`.

```
let f key new_value old_value = (show key) ++ ":" ++ new_value ++ "|" ++ old_value
insertLookupWithKey f 5 "xxx" (fromList [(5,"a"), (3,"b")]) == (Just "a", fromList [(3, "b"), (5, "xxx|a")])
insertLookupWithKey f 7 "xxx" (fromList [(5,"a"), (3,"b")]) == (Nothing, fromList [(3, "b"), (5, "xxx|a")])
insertLookupWithKey f 5 "xxx" empty == (Nothing, singleton 5 "xxx|a")
```

This is how to define `insertLookup` using `insertLookupWithKey`:

```
let insertLookup kx x t = insertLookupWithKey (\_ a _ -> a) kx x t
insertLookup 5 "x" (fromList [(5,"a"), (3,"b")]) == (Just "a", fromList [(3, "b"), (5, "x")])
insertLookup 7 "x" (fromList [(5,"a"), (3,"b")]) == (Nothing, fromList [(3, "b"), (5, "x")])
```

Delete/Update

```
delete :: Ord k => k -> Map k a -> Map k a
```

```
| # Source
```

$O(\log n)$. Delete a key and its value from the map. When the key is not a member of the map, the original map is returned.

```
delete 5 (fromList [(5,"a"), (3,"b")]) == singleton 3 "b"
delete 7 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "a")]
delete 5 empty == empty
```

```
adjust :: Ord k => (a -> a) -> k -> Map k a -> Map k a
```

```
| # Source
```

$O(\log n)$. Update a value at a specific key with the result of the provided function. When the key is not a member of the map, the original map is returned.

```
adjust ("new " ++) 5 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "new
adjust ("new " ++) 7 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "a"
adjust ("new " ++) 7 empty == empty
```

adjustWithKey :: Ord k => (k -> a -> a) -> k -> Map k a -> Map k a | # Source

$O(\log n)$. Adjust a value at a specific key. When the key is not a member of the map, the original map is returned.

```
let f key x = (show key) ++ ":new " ++ x
adjustWithKey f 5 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "5:new
adjustWithKey f 7 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "a")]
adjustWithKey f 7 empty == empty
```

update :: Ord k => (a -> Maybe a) -> k -> Map k a -> Map k a | # Source

$O(\log n)$. The expression (`update f k map`) updates the value `x` at `k` (if it is in the map). If (`f x`) is `Nothing`, the element is deleted. If it is (`Just y`), the key `k` is bound to the new value `y`.

```
let f x = if x == "a" then Just "new a" else Nothing
update f 5 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "new a")]
update f 7 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "a")]
update f 3 (fromList [(5,"a"), (3,"b")]) == singleton 5 "a"
```

updateWithKey :: Ord k => (k -> a -> Maybe a) -> k -> Map k a -> Map k a | # Source

$O(\log n)$. The expression (`updateWithKey f k map`) updates the value `x` at `k` (if it is in the map). If (`f k x`) is `Nothing`, the element is deleted. If it is (`Just y`), the key `k` is bound to the new value `y`.

```
let f k x = if x == "a" then Just ((show k) ++ ":new a") else Nothing
updateWithKey f 5 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "5:new
updateWithKey f 7 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "a")]
updateWithKey f 3 (fromList [(5,"a"), (3,"b")]) == singleton 5 "a"
```

updateLookupWithKey :: Ord k => (k -> a -> Maybe a) -> k -> Map k a -> (Maybe a, Map k a) | # Source

$O(\log n)$. Lookup and update. See also `updateWithKey`. The function returns changed value, if it is updated. Returns the original key value if the map entry is deleted.

```
let f k x = if x == "a" then Just ((show k) ++ ":new a") else Nothing
updateLookupWithKey f 5 (fromList [(5,"a"), (3,"b")]) == (Just "5:new a", fromList [(3, "b"), (5, "5:new a")])
updateLookupWithKey f 7 (fromList [(5,"a"), (3,"b")]) == (Nothing, fromList [(3, "b"), (5, "a")])
updateLookupWithKey f 3 (fromList [(5,"a"), (3,"b")]) == (Just "b", singleton 5 "a")
```

alter :: Ord k => (Maybe a -> Maybe a) -> k -> Map k a -> Map k a | # Source

$O(\log n)$. The expression (`alter f k map`) alters the value `x` at `k`, or absence thereof. `alter` can be used to insert, delete, or update a value in a `Map`. In short: `lookup k (alter f k m) = f (lookup k m)`.

```
let f _ = Nothing
alter f 7 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "a")]
alter f 5 (fromList [(5,"a"), (3,"b")]) == singleton 3 "b"

let f _ = Just "c"
```

```
alter f 7 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "a"), (7, "c")]
alter f 5 (fromList [(5,"a"), (3,"b")]) == fromList [(3, "b"), (5, "c")]
```

```
alterF :: (Functor f, Ord k) => (Maybe a -> f (Maybe a)) -> k -> Map k a -> f (Map k a) | # Source
```

$O(\log n)$. The expression (`alterF f k map`) alters the value `x` at `k`, or absence thereof. `alterF` can be used to inspect, insert, delete, or update a value in a `Map`. In short: `lookup k <$> alterF f k m = f (lookup k m)`.

Example:

```
interactiveAlter :: Int -> Map Int String -> IO (Map Int String)
interactiveAlter k m = alterF f k m where
  f Nothing -> do
    putStrLn $ show k ++
      " was not found in the map. Would you like to add it?"
    getUserResponse1 :: IO (Maybe String)
  f (Just old) -> do
    putStrLn "The key is currently bound to " ++ show old ++
      ". Would you like to change or delete it?"
    getUserResponse2 :: IO (Maybe String)
```

`alterF` is the most general operation for working with an individual key that may or may not be in a given map. When used with trivial functors like `Identity` and `Const`, it is often slightly slower than more specialized combinators like `lookup` and `insert`. However, when the functor is non-trivial and key comparison is not particularly cheap, it is the fastest way.

Note on rewrite rules:

This module includes GHC rewrite rules to optimize `alterF` for the `Const` and `Identity` functors. In general, these rules improve performance. The sole exception is that when using `Identity`, deleting a key that is already absent takes longer than it would without the rules. If you expect this to occur a very large fraction of the time, you might consider using a private copy of the `Identity` type.

Note: `alterF` is a flipped version of the `at` combinator from `At`.

Combine

Union

```
union :: Ord k => Map k a -> Map k a -> Map k a | # Source
```

$O(m \cdot \log(n/m + 1))$, $m \leq n$. The expression (`union t1 t2`) takes the left-biased union of `t1` and `t2`. It prefers `t1` when duplicate keys are encountered, i.e. (`union == unionWith const`).

```
union (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (7, "C")]) == fromList
```

```
unionWith :: Ord k => (a -> a -> a) -> Map k a -> Map k a -> Map k a | # Source
```

$O(m \cdot \log(n/m + 1))$, $m \leq n$. Union with a combining function.

```
unionWith (++) (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (7, "C")]) ==
```

```
unionWithKey :: Ord k => (k -> a -> a -> a) -> Map k a -> Map k a -> Map k a | # Source
```

$O(m \cdot \log(n/m + 1))$, $m \leq n$. Union with a combining function.

```
let f key left_value right_value = (show key) ++ ":" ++ left_value ++ "|" ++ right_value
unionWithKey f (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (7, "C")]) ==
```

```
unions :: Ord k => [Map k a] -> Map k a | # Source
```

The union of a list of maps: (`unions == foldl union empty`).

```
unions [(fromList [(5, "a"), (3, "b")]), (fromList [(5, "A"), (7, "C")]), (fromList [(3, "b"), (5, "a"), (7, "C")])]
unions [(fromList [(5, "A3"), (3, "B3")]), (fromList [(5, "A"), (7, "C")]), (fromList [(3, "B3"), (5, "A3"), (7, "C")])]
```

```
unionsWith :: Ord k => (a -> a -> a) -> [Map k a] -> Map k a | # Source
```

The union of a list of maps, with a combining operation: (`unionsWith f == foldl (unionWith f) empty`).

```
unionsWith (++) [(fromList [(5, "a"), (3, "b")]), (fromList [(5, "A"), (7, "C")])]
== fromList [(3, "bB3"), (5, "aAA3"), (7, "C")]
```

Difference

```
difference :: Ord k => Map k a -> Map k b -> Map k a | # Source
```

$O(m \log(n/m + 1))$, $m \leq n$. Difference of two maps. Return elements of the first map not existing in the second map.

```
difference (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (7, "C")]) == singleton 3 "b"
```

```
differenceWith :: Ord k => (a -> b -> Maybe a) -> Map k a -> Map k b -> Map k a | # Source
```

$O(n+m)$. Difference with a combining function. When two equal keys are encountered, the combining function is applied to the values of these keys. If it returns `Nothing`, the element is discarded (proper set difference). If it returns `(Just y)`, the element is updated with a new value `y`.

```
let f al ar = if al == "b" then Just (al ++ ":" ++ ar) else Nothing
differenceWith f (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (3, "B"), (3, "b")])
== singleton 3 "b:B"
```

```
differenceWithKey :: Ord k => (k -> a -> b -> Maybe a) -> Map k a -> Map k b -> Map k a | # Source
```

$O(n+m)$. Difference with a combining function. When two equal keys are encountered, the combining function is applied to the key and both values. If it returns `Nothing`, the element is discarded (proper set difference). If it returns `(Just y)`, the element is updated with a new value `y`.

```
let f k al ar = if al == "b" then Just ((show k) ++ ":" ++ al ++ "|" ++ ar) else Nothing
differenceWithKey f (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (3, "B")])
== singleton 3 "3:b|B"
```


Intersection

```
intersection :: Ord k => Map k a -> Map k b -> Map k a
```

Source

$O(m \log(n/m + 1))$, $m \leq n$. Intersection of two maps. Return data in the first map for the keys existing in both maps. (`intersection m1 m2 == intersectionWith const m1 m2`).

```
intersection (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (7, "C")]) == s
```

```
intersectionWith :: Ord k => (a -> b -> c) -> Map k a -> Map k b -> Map k c
```

Source
#

$O(m \log(n/m + 1))$, $m \leq n$. Intersection with a combining function.

```
intersectionWith (++) (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (7, "C")])
```

```
intersectionWithKey :: Ord k => (k -> a -> b -> c) -> Map k a -> Map k b -> Map k c
```

Source

$O(m \log(n/m + 1))$, $m \leq n$. Intersection with a combining function.

```
let f k al ar = (show k) ++ ":" ++ al ++ "|" ++ ar
intersectionWithKey f (fromList [(5, "a"), (3, "b")]) (fromList [(5, "A"), (7, "C")])
```

General combining functions

See [Data.Map.Strict.Merge](#)

Deprecated general combining function

```
mergeWithKey :: Ord k => (k -> a -> b -> Maybe c) -> (Map k a -> Map k c) -> (Map k b -> Map k c) -> Map k a -> Map k b -> Map k c
```

Source

$O(n+m)$. An unsafe universal combining function.

WARNING: This function can produce corrupt maps and its results may depend on the internal structures of its inputs. Users should prefer `merge` or `mergeA`.

When `mergeWithKey` is given three arguments, it is inlined to the call site. You should therefore use `mergeWithKey` only to define custom combining functions. For example, you could define `unionWithKey`, `differenceWithKey` and `intersectionWithKey` as

```
myUnionWithKey f m1 m2 = mergeWithKey (\k x1 x2 -> Just (f k x1 x2)) id id m1 m2
myDifferenceWithKey f m1 m2 = mergeWithKey f id (const empty) m1 m2
myIntersectionWithKey f m1 m2 = mergeWithKey (\k x1 x2 -> Just (f k x1 x2)) (const empty) id m1 m2
```

When calling `mergeWithKey combine only1 only2`, a function combining two `Maps` is created, such that

- if a key is present in both maps, it is passed with both corresponding values to the combine function. Depending on the result, the key is either present in the result with specified value, or is left out;
- a nonempty subtree present only in the first map is passed to `only1` and the output is added to the result;
- a nonempty subtree present only in the second map is passed to `only2` and the output is added to the result.

The `only1` and `only2` methods *must return a map with a subset (possibly empty) of the keys of the given map*. The values can be modified arbitrarily. Most common variants of `only1` and `only2` are `id` and `const empty`, but for example `map f` or `filterWithKey f` could be used for any `f`.

Traversal

Map

```
map :: (a -> b) -> Map k a -> Map k b
```

Source

O(n). Map a function over all values in the map.

```
map (++ "x") (fromList [(5,"a"), (3,"b")]) == fromList [(3, "bx"), (5, "ax")]
```

```
mapWithKey :: (k -> a -> b) -> Map k a -> Map k b
```

Source

O(n). Map a function over all values in the map.

```
let f key x = (show key) ++ ":" ++ x
mapWithKey f (fromList [(5,"a"), (3,"b")]) == fromList [(3, "3:b"), (5, "5:a")]
```

```
traverseWithKey :: Applicative t => (k -> a -> t b) -> Map k a -> t (Map k b)
```

Source
#

O(n). `traverseWithKey f m == fromList $ traverse ((k, v) -> (v' -> v' seq (k,v'))) $ f k v` (`toList m`) That is, it behaves much like a regular `traverse` except that the traversing function also has access to the key associated with a value and the values are forced before they are installed in the result map.

```
traverseWithKey (\k v -> if odd k then Just (succ v) else Nothing) (fromList [(1,
traverseWithKey (\k v -> if odd k then Just (succ v) else Nothing) (fromList [(2,
```

```
traverseMaybeWithKey :: Applicative f => (k -> a -> f (Maybe b)) -> Map k a -> f (Map k b)
```

Source

O(n). Traverse keys/values and collect the `Just` results.

```
mapAccum :: (a -> b -> (a, c)) -> a -> Map k b -> (a, Map k c)
```

Source

O(n). The function `mapAccum` threads an accumulating argument through the map in ascending order of keys.

```
let f a b = (a ++ b, b ++ "X")
mapAccum f "Everything: " (fromList [(5,"a"), (3,"b")]) == ("Everything: ba", from
```

```
mapAccumWithKey :: (a -> k -> b -> (a, c)) -> a -> Map k b -> (a, Map k c)
```

Source

O(n). The function `mapAccumWithKey` threads an accumulating argument through the map in ascending order of keys.

```
let f a k b = (a ++ " " ++ (show k) ++ "-" ++ b, b ++ "X")
mapAccumWithKey f "Everything:" (fromList [(5,"a"), (3,"b")]) == ("Everything: 3-l
```

```
mapAccumRWithKey :: (a -> k -> b -> (a, c)) -> a -> Map k b -> (a, Map k c)
```

Source
#

O(n). The function `mapAccumR` threads an accumulating argument through the map in descending order of keys.

```
mapKeys :: Ord k2 => (k1 -> k2) -> Map k1 a -> Map k2 a
```

Source

$O(n \log n)$. `mapKeys f s` is the map obtained by applying `f` to each key of `s`.

The size of the result may be smaller if `f` maps two or more distinct keys to the same new key. In this case the value at the greatest of the original keys is retained.

```
mapKeys (+ 1) (fromList [(5,"a"), (3,"b")]) == fromList [(6,"a"), (4,"b")]
mapKeys (\ _ -> 1) (fromList [(1,"b"), (2,"a"), (3,"d"), (4,"c")]) == singleton 1
mapKeys (\ _ -> 3) (fromList [(1,"b"), (2,"a"), (3,"d"), (4,"c")]) == singleton 3
```

mapKeysWith :: Ord k2 => (a -> a -> a) -> (k1 -> k2) -> Map k1 a -> Map k2 a # Source

$O(n \log n)$. `mapKeysWith c f s` is the map obtained by applying `f` to each key of `s`.

The size of the result may be smaller if `f` maps two or more distinct keys to the same new key. In this case the associated values will be combined using `c`.

```
mapKeysWith (++) (\ _ -> 1) (fromList [(1,"b"), (2,"a"), (3,"d"), (4,"c")]) == singleton 1
mapKeysWith (++) (\ _ -> 3) (fromList [(1,"b"), (2,"a"), (3,"d"), (4,"c")]) == singleton 3
```

mapKeysMonotonic :: (k1 -> k2) -> Map k1 a -> Map k2 a # Source

$O(n)$. `mapKeysMonotonic f s` == `mapKeys f s`, but works only when `f` is strictly monotonic. That is, for any values `x` and `y`, if `x < y` then `f x < f y`. *The precondition is not checked.* Semi-formally, we have:

```
and [x < y ==> f x < f y | x <- ls, y <- ls]
    ==> mapKeysMonotonic f s == mapKeys f s
    where ls = keys s
```

This means that `f` maps distinct original keys to distinct resulting keys. This function has better performance than `mapKeys`.

```
mapKeysMonotonic (\ k -> k * 2) (fromList [(5,"a"), (3,"b")]) == fromList [(6,"a"), (4,"b")]
valid (mapKeysMonotonic (\ k -> k * 2) (fromList [(5,"a"), (3,"b")])) == True
valid (mapKeysMonotonic (\ _ -> 1) (fromList [(5,"a"), (3,"b")])) == False
```

Folds

foldr :: (a -> b -> b) -> b -> Map k a -> b # Source

$O(n)$. Fold the values in the map using the given right-associative binary operator, such that `foldr f z == foldr f z . elems`.

For example,

```
elems map = foldr (:) [] map
```

```
let f a len = len + (length a)
foldr f 0 (fromList [(5,"a"), (3,"bbb")]) == 4
```

foldl :: (a -> b -> a) -> a -> Map k b -> a # Source

$O(n)$. Fold the values in the map using the given left-associative binary operator, such that `foldl f z == foldl f z . elems`.

For example,

```
elems = reverse . foldl (flip (:)) []
```

```
let f len a = len + (length a)
foldl f 0 (fromList [(5,"a"), (3,"bbb")]) == 4
```

```
foldrWithKey :: (k -> a -> b -> b) -> b -> Map k a -> b
```

Source

$O(n)$. Fold the keys and values in the map using the given right-associative binary operator, such that `foldrWithKey f z == foldr (uncurry f) z . toAscList`.

For example,

```
keys map = foldrWithKey (\k x ks -> k:ks) [] map

let f k a result = result ++ "(" ++ (show k) ++ ":" ++ a ++ ")"
foldrWithKey f "Map: " (fromList [(5,"a"), (3,"b")]) == "Map: (5:a)(3:b)"
```

```
foldlWithKey :: (a -> k -> b -> a) -> a -> Map k b -> a
```

Source

$O(n)$. Fold the keys and values in the map using the given left-associative binary operator, such that `foldlWithKey f z == foldl (\z' (kx, x) -> f z' kx x) z . toAscList`.

For example,

```
keys = reverse . foldlWithKey (\ks k x -> k:ks) []

let f result k a = result ++ "(" ++ (show k) ++ ":" ++ a ++ ")"
foldlWithKey f "Map: " (fromList [(5,"a"), (3,"b")]) == "Map: (3:b)(5:a)"
```

```
foldMapWithKey :: Monoid m => (k -> a -> m) -> Map k a -> m
```

Source

$O(n)$. Fold the keys and values in the map using the given monoid, such that

```
foldMapWithKey f = fold . mapWithKey f
```

This can be an asymptotically faster than `foldrWithKey` or `foldlWithKey` for some monoids.

Strict folds

```
foldr' :: (a -> b -> b) -> b -> Map k a -> b
```

Source

$O(n)$. A strict version of `foldr`. Each application of the operator is evaluated before using the result in the next application. This function is strict in the starting value.

```
foldl' :: (a -> b -> a) -> a -> Map k b -> a
```

Source

$O(n)$. A strict version of `foldl`. Each application of the operator is evaluated before using the result in the next application. This function is strict in the starting value.

```
foldrWithKey' :: (k -> a -> b -> b) -> b -> Map k a -> b
```

Source

$O(n)$. A strict version of `foldrWithKey`. Each application of the operator is evaluated before using the result in the next application. This function is strict in the starting value.

```
foldlWithKey' :: (a -> k -> b -> a) -> a -> Map k b -> a
```

Source

$O(n)$. A strict version of `foldlWithKey`. Each application of the operator is evaluated before using the result in the next application. This function is strict in the starting value.

Conversion

elems :: Map k a -> [a]

Source

$O(n)$. Return all elements of the map in the ascending order of their keys. Subject to list fusion.

```
elems (fromList [(5,"a"), (3,"b")]) == ["b","a"]
elems empty == []
```

keys :: Map k a -> [k]

Source

$O(n)$. Return all keys of the map in ascending order. Subject to list fusion.

```
keys (fromList [(5,"a"), (3,"b")]) == [3,5]
keys empty == []
```

assocs :: Map k a -> [(k, a)]

Source

$O(n)$. An alias for `toAscList`. Return all key/value pairs in the map in ascending key order. Subject to list fusion.

```
assocs (fromList [(5,"a"), (3,"b")]) == [(3,"b"), (5,"a")]
assocs empty == []
```

keysSet :: Map k a -> Set k

Source

$O(n)$. The set of all keys of the map.

```
keysSet (fromList [(5,"a"), (3,"b")]) == Data.Set.fromList [3,5]
keysSet empty == Data.Set.empty
```

fromSet :: (k -> a) -> Set k -> Map k a

Source

$O(n)$. Build a map from a set of keys and a function which for each key computes its value.

```
fromSet (\k -> replicate k 'a') (Data.Set.fromList [3, 5]) == fromList [(5,"aaaaa'
fromSet undefined Data.Set.empty == empty
```

Lists

toList :: Map k a -> [(k, a)]

Source

$O(n)$. Convert the map to a list of key/value pairs. Subject to list fusion.

```
toList (fromList [(5,"a"), (3,"b")]) == [(3,"b"), (5,"a")]
toList empty == []
```

fromList :: Ord k => [(k, a)] -> Map k a

Source

$O(n \log n)$. Build a map from a list of key/value pairs. See also `fromAscList`. If the list contains more than one value for the same key, the last value for the key is retained.

If the keys of the list are ordered, linear-time implementation is used, with the performance equal to `fromDistinctAscList`.

```
fromList [] == empty
fromList [(5,"a"), (3,"b"), (5, "c")] == fromList [(5,"c"), (3,"b")]
fromList [(5,"c"), (3,"b"), (5, "a")] == fromList [(5,"a"), (3,"b")]
```

```
fromListWith :: Ord k => (a -> a -> a) -> [(k, a)] -> Map k a
```

Source

$O(n \log n)$. Build a map from a list of key/value pairs with a combining function. See also [fromAscListWith](#).

```
fromListWith (++) [(5,"a"), (5,"b"), (3,"b"), (3,"a"), (5,"a")] == fromList [(3,
fromListWith (++) [] == empty
```

```
fromListWithKey :: Ord k => (k -> a -> a -> a) -> [(k, a)] -> Map k a
```

Source

$O(n \log n)$. Build a map from a list of key/value pairs with a combining function. See also [fromAscListWithKey](#).

```
let f k a1 a2 = (show k) ++ a1 ++ a2
fromListWithKey f [(5,"a"), (5,"b"), (3,"b"), (3,"a"), (5,"a")] == fromList [(3,
fromListWithKey f [] == empty
```

Ordered lists

```
toAscList :: Map k a -> [(k, a)]
```

Source

$O(n)$. Convert the map to a list of key/value pairs where the keys are in ascending order. Subject to list fusion.

```
toAscList (fromList [(5,"a"), (3,"b")]) == [(3,"b"), (5,"a")]
```

```
toDescList :: Map k a -> [(k, a)]
```

Source

$O(n)$. Convert the map to a list of key/value pairs where the keys are in descending order. Subject to list fusion.

```
toDescList (fromList [(5,"a"), (3,"b")]) == [(5,"a"), (3,"b")]
```

```
fromAscList :: Eq k => [(k, a)] -> Map k a
```

Source

$O(n)$. Build a map from an ascending list in linear time. *The precondition (input list is ascending) is not checked.*

```
fromAscList [(3,"b"), (5,"a")] == fromList [(3, "b"), (5, "a")]
fromAscList [(3,"b"), (5,"a"), (5,"b")] == fromList [(3, "b"), (5, "b")]
valid (fromAscList [(3,"b"), (5,"a"), (5,"b")]) == True
valid (fromAscList [(5,"a"), (3,"b"), (5,"b")]) == False
```

```
fromAscListWith :: Eq k => (a -> a -> a) -> [(k, a)] -> Map k a
```

Source

$O(n)$. Build a map from an ascending list in linear time with a combining function for equal keys. *The precondition (input list is ascending) is not checked.*

```
fromAscListWith (++) [(3,"b"), (5,"a"), (5,"b")] == fromList [(3, "b"), (5, "ba")]
valid (fromAscListWith (++) [(3,"b"), (5,"a"), (5,"b")]) == True
valid (fromAscListWith (++) [(5,"a"), (3,"b"), (5,"b")]) == False
```

fromAscListWithKey :: Eq k => (k -> a -> a -> a) -> [(k, a)] -> Map k a | # Source

$O(n)$. Build a map from an ascending list in linear time with a combining function for equal keys. *The precondition (input list is ascending) is not checked.*

```
let f k a1 a2 = (show k) ++ ":" ++ a1 ++ a2
fromAscListWithKey f [(3,"b"), (5,"a"), (5,"b"), (5,"b")] == fromList [(3, "b"),
valid (fromAscListWithKey f [(3,"b"), (5,"a"), (5,"b"), (5,"b")]) == True
valid (fromAscListWithKey f [(5,"a"), (3,"b"), (5,"b"), (5,"b")]) == False
```

fromDistinctAscList :: [(k, a)] -> Map k a | # Source

$O(n)$. Build a map from an ascending list of distinct elements in linear time. *The precondition is not checked.*

```
fromDistinctAscList [(3,"b"), (5,"a")] == fromList [(3, "b"), (5, "a")]
valid (fromDistinctAscList [(3,"b"), (5,"a")]) == True
valid (fromDistinctAscList [(3,"b"), (5,"a"), (5,"b")]) == False
```

fromDescList :: Eq k => [(k, a)] -> Map k a | # Source

$O(n)$. Build a map from a descending list in linear time. *The precondition (input list is descending) is not checked.*

```
fromDescList [(5,"a"), (3,"b")] == fromList [(3, "b"), (5, "a")]
fromDescList [(5,"a"), (5,"b"), (3,"a")] == fromList [(3, "b"), (5, "b")]
valid (fromDescList [(5,"a"), (5,"b"), (3,"b")]) == True
valid (fromDescList [(5,"a"), (3,"b"), (5,"b")]) == False
```

fromDescListWith :: Eq k => (a -> a -> a) -> [(k, a)] -> Map k a | # Source

$O(n)$. Build a map from a descending list in linear time with a combining function for equal keys. *The precondition (input list is descending) is not checked.*

```
fromDescListWith (++) [(5,"a"), (5,"b"), (3,"b")] == fromList [(3, "b"), (5, "ba")
valid (fromDescListWith (++) [(5,"a"), (5,"b"), (3,"b")]) == True
valid (fromDescListWith (++) [(5,"a"), (3,"b"), (5,"b")]) == False
```

fromDescListWithKey :: Eq k => (k -> a -> a -> a) -> [(k, a)] -> Map k a | # Source

$O(n)$. Build a map from a descending list in linear time with a combining function for equal keys. *The precondition (input list is descending) is not checked.*

```
let f k a1 a2 = (show k) ++ ":" ++ a1 ++ a2
fromDescListWithKey f [(5,"a"), (5,"b"), (5,"b"), (3,"b")] == fromList [(3, "b"),
valid (fromDescListWithKey f [(5,"a"), (5,"b"), (5,"b"), (3,"b")]) == True
valid (fromDescListWithKey f [(5,"a"), (3,"b"), (5,"b"), (5,"b")]) == False
```

fromDistinctDescList :: [(k, a)] -> Map k a | # Source

$O(n)$. Build a map from a descending list of distinct elements in linear time. *The precondition is not checked.*

```
fromDistinctDescList [(5,"a"), (3,"b")] == fromList [(3, "b"), (5, "a")]
valid (fromDistinctDescList [(5,"a"), (3,"b")]) == True
valid (fromDistinctDescList [(5,"a"), (3,"b"), (3,"a")]) == False
```

Filter

```
filter :: (a -> Bool) -> Map k a -> Map k a
```

Source

$O(n)$. Filter all values that satisfy the predicate.

```
filter (> "a") (fromList [(5,"a"), (3,"b")]) == singleton 3 "b"
filter (> "x") (fromList [(5,"a"), (3,"b")]) == empty
filter (< "a") (fromList [(5,"a"), (3,"b")]) == empty
```

```
filterWithKey :: (k -> a -> Bool) -> Map k a -> Map k a
```

Source

$O(n)$. Filter all keys/values that satisfy the predicate.

```
filterWithKey (\k _ -> k > 4) (fromList [(5,"a"), (3,"b")]) == singleton 5 "a"
```

```
restrictKeys :: Ord k => Map k a -> Set k -> Map k a
```

Source

$O(m \log(nm + 1))$, $m \leq n$. Restrict a `Map` to only those keys found in a `Set`.

```
m restrictKeys s = filterWithKey (k _ -> k `member` s) m
```

Since: 0.5.8

```
withoutKeys :: Ord k => Map k a -> Set k -> Map k a
```

Source

$O(m \log(nm + 1))$, $m \leq n$. Remove all keys in a `Set` from a `Map`.

```
m withoutKeys s = filterWithKey (k _ -> k `notMember` s) m
```

Since: 0.5.8

```
partition :: (a -> Bool) -> Map k a -> (Map k a, Map k a)
```

Source

$O(n)$. Partition the map according to a predicate. The first map contains all elements that satisfy the predicate, the second all elements that fail the predicate. See also `split`.

```
partition (> "a") (fromList [(5,"a"), (3,"b")]) == (singleton 3 "b", singleton 5 "a")
partition (< "x") (fromList [(5,"a"), (3,"b")]) == (fromList [(3, "b"), (5, "a")])
partition (> "x") (fromList [(5,"a"), (3,"b")]) == (empty, fromList [(3, "b"), (5, "a")])
```

```
partitionWithKey :: (k -> a -> Bool) -> Map k a -> (Map k a, Map k a)
```

Source

$O(n)$. Partition the map according to a predicate. The first map contains all elements that satisfy the predicate, the second all elements that fail the predicate. See also `split`.

```
partitionWithKey (\k _ -> k > 3) (fromList [(5,"a"), (3,"b")]) == (singleton 5 "a", singleton 3 "b")
partitionWithKey (\k _ -> k < 7) (fromList [(5,"a"), (3,"b")]) == (fromList [(3, "b"), (5, "a")])
partitionWithKey (\k _ -> k > 7) (fromList [(5,"a"), (3,"b")]) == (empty, fromList [(3, "b"), (5, "a")])
```

```
takeWhileAntitone :: (k -> Bool) -> Map k a -> Map k a
```

Source

$O(\log n)$. Take while a predicate on the keys holds. The user is responsible for ensuring that for all keys j and k in the map, $j < k \implies p\ j \geq p\ k$. See note at `spanAntitone`.


```
takeWhileAntitone p = fromDistinctAscList . takeWhile (p . fst) . toList
takeWhileAntitone p = filterWithKey (k _ -> p k)
```

```
dropWhileAntitone :: (k -> Bool) -> Map k a -> Map k a
```

Source

$O(\log n)$. Drop while a predicate on the keys holds. The user is responsible for ensuring that for all keys j and k in the map, $j < k \implies p\ j \geq p\ k$. See note at [spanAntitone](#).

```
dropWhileAntitone p = fromDistinctAscList . dropWhile (p . fst) . toList
dropWhileAntitone p = filterWithKey (k -> not (p k))
```

```
spanAntitone :: (k -> Bool) -> Map k a -> (Map k a, Map k a)
```

Source

$O(\log n)$. Divide a map at the point where a predicate on the keys stops holding. The user is responsible for ensuring that for all keys j and k in the map, $j < k \implies p\ j \geq p\ k$.

```
spanAntitone p xs = (takeWhileAntitone p xs, dropWhileAntitone p xs)
spanAntitone p xs = partition p xs
```

Note: if p is not actually antitone, then `spanAntitone` will split the map at some *unspecified* point where the predicate switches from holding to not holding (where the predicate is seen to hold before the first key and to fail after the last key).

```
mapMaybe :: (a -> Maybe b) -> Map k a -> Map k b
```

Source

$O(n)$. Map values and collect the **Just** results.

```
let f x = if x == "a" then Just "new a" else Nothing
mapMaybe f (fromList [(5,"a"), (3,"b")]) == singleton 5 "new a"
```

```
mapMaybeWithKey :: (k -> a -> Maybe b) -> Map k a -> Map k b
```

Source

$O(n)$. Map keys/values and collect the **Just** results.

```
let f k _ = if k < 5 then Just ("key : " ++ (show k)) else Nothing
mapMaybeWithKey f (fromList [(5,"a"), (3,"b")]) == singleton 3 "key : 3"
```

```
mapEither :: (a -> Either b c) -> Map k a -> (Map k b, Map k c)
```

Source

$O(n)$. Map values and separate the **Left** and **Right** results.

```
let f a = if a < "c" then Left a else Right a
mapEither f (fromList [(5,"a"), (3,"b"), (1,"x"), (7,"z")])
  == (fromList [(3,"b"), (5,"a")], fromList [(1,"x"), (7,"z")])

mapEither (\ a -> Right a) (fromList [(5,"a"), (3,"b"), (1,"x"), (7,"z")])
  == (empty, fromList [(5,"a"), (3,"b"), (1,"x"), (7,"z")])
```

```
mapEitherWithKey :: (k -> a -> Either b c) -> Map k a -> (Map k b, Map k c)
```

Source
#

$O(n)$. Map keys/values and separate the **Left** and **Right** results.

```
let f k a = if k < 5 then Left (k * 2) else Right (a ++ a)
mapEitherWithKey f (fromList [(5,"a"), (3,"b"), (1,"x"), (7,"z")])
  == (fromList [(1,2), (3,6)], fromList [(5,"aa"), (7,"zz")])

mapEitherWithKey (\_ a -> Right a) (fromList [(5,"a"), (3,"b"), (1,"x"), (7,"z")])
  == (empty, fromList [(1,"x"), (3,"b"), (5,"a"), (7,"z")])
```

```
split :: Ord k => k -> Map k a -> (Map k a, Map k a) | # Source
```

$O(\log n)$. The expression (`split k map`) is a pair (`map1, map2`) where the keys in `map1` are smaller than `k` and the keys in `map2` larger than `k`. Any key equal to `k` is found in neither `map1` nor `map2`.

```
split 2 (fromList [(5,"a"), (3,"b")]) == (empty, fromList [(3,"b"), (5,"a")])
split 3 (fromList [(5,"a"), (3,"b")]) == (empty, singleton 5 "a")
split 4 (fromList [(5,"a"), (3,"b")]) == (singleton 3 "b", singleton 5 "a")
split 5 (fromList [(5,"a"), (3,"b")]) == (singleton 3 "b", empty)
split 6 (fromList [(5,"a"), (3,"b")]) == (fromList [(3,"b"), (5,"a")], empty)
```

```
splitLookup :: Ord k => k -> Map k a -> (Map k a, Maybe a, Map k a) | # Source
```

$O(\log n)$. The expression (`splitLookup k map`) splits a map just like `split` but also returns `lookup k map`.

```
splitLookup 2 (fromList [(5,"a"), (3,"b")]) == (empty, Nothing, fromList [(3,"b")])
splitLookup 3 (fromList [(5,"a"), (3,"b")]) == (empty, Just "b", singleton 5 "a")
splitLookup 4 (fromList [(5,"a"), (3,"b")]) == (singleton 3 "b", Nothing, singleton 5 "a")
splitLookup 5 (fromList [(5,"a"), (3,"b")]) == (singleton 3 "b", Just "a", empty)
splitLookup 6 (fromList [(5,"a"), (3,"b")]) == (fromList [(3,"b"), (5,"a")], Nothing, empty)
```

```
splitRoot :: Map k b -> [Map k b] | # Source
```

$O(1)$. Decompose a map into pieces based on the structure of the underlying tree. This function is useful for consuming a map in parallel.

No guarantee is made as to the sizes of the pieces; an internal, but deterministic process determines this. However, it is guaranteed that the pieces returned will be in ascending order (all elements in the first submap less than all elements in the second, and so on).

Examples:

```
splitRoot (fromList (zip [1..6] ['a'..f])) ==
  [fromList [(1,'a'),(2,'b'),(3,'c')],fromList [(4,'d')],fromList [(5,'e'),(6,'f')]]
```

```
splitRoot empty == []
```

Note that the current implementation does not return more than three submaps, but you should not depend on this behaviour because it can change in the future without notice.

Submap

```
isSubmapOf :: (Ord k, Eq a) => Map k a -> Map k a -> Bool | # Source
```

$O(m \cdot \log(nm + 1))$, $m \leq n$. This function is defined as (`isSubmapOf = isSubmapOfBy (==)`).

```
isSubmapOfBy :: Ord k => (a -> b -> Bool) -> Map k a -> Map k b -> Bool | # Source
```

$O(m \cdot \log(nm + 1))$, $m \leq n$. The expression (`isSubmapOfBy f t1 t2`) returns `True` if all keys in `t1` are in tree `t2`, and when `f` returns `True` when applied to their respective values. For example, the following expressions are all `True`:

```
isSubmapOfBy (==) (fromList [('a',1)]) (fromList [('a',1),('b',2)])
isSubmapOfBy (<=) (fromList [('a',1)]) (fromList [('a',1),('b',2)])
```

```
isSubmapOfBy (==) (fromList [('a',1),('b',2)]) (fromList [('a',1),('b',2)])
```

But the following are all **False**:

```
isSubmapOfBy (==) (fromList [('a',2)]) (fromList [('a',1),('b',2)])
isSubmapOfBy (<) (fromList [('a',1)]) (fromList [('a',1),('b',2)])
isSubmapOfBy (==) (fromList [('a',1),('b',2)]) (fromList [('a',1)])
```

```
isProperSubmapOf :: (Ord k, Eq a) => Map k a -> Map k a -> Bool
```

Source

$O(m \log(nm + 1))$, $m \leq n$. Is this a proper submap? (ie. a submap but not equal). Defined as
isProperSubmapOf = **isProperSubmapOfBy** (==).

```
isProperSubmapOfBy :: Ord k => (a -> b -> Bool) -> Map k a -> Map k b -> Bool
```

Source
#

$O(m \log(nm + 1))$, $m \leq n$. Is this a proper submap? (ie. a submap but not equal). The expression
isProperSubmapOfBy f m1 m2 returns **True** when m1 and m2 are not equal, all keys in m1 are in m2, and
 when f returns **True** when applied to their respective values. For example, the following expressions are all
True:

```
isProperSubmapOfBy (==) (fromList [(1,1)]) (fromList [(1,1),(2,2)])
isProperSubmapOfBy (<=) (fromList [(1,1)]) (fromList [(1,1),(2,2)])
```

But the following are all **False**:

```
isProperSubmapOfBy (==) (fromList [(1,1),(2,2)]) (fromList [(1,1),(2,2)])
isProperSubmapOfBy (==) (fromList [(1,1),(2,2)]) (fromList [(1,1)])
isProperSubmapOfBy (<) (fromList [(1,1)]) (fromList [(1,1),(2,2)])
```

Indexed

```
lookupIndex :: Ord k => k -> Map k a -> Maybe Int
```

Source

$O(\log n)$. Lookup the *index* of a key, which is its zero-based index in the sequence sorted by keys. The index is
 a number from 0 up to, but not including, the **size** of the map.

```
isJust (lookupIndex 2 (fromList [(5,"a"), (3,"b")])) == False
fromJust (lookupIndex 3 (fromList [(5,"a"), (3,"b")])) == 0
fromJust (lookupIndex 5 (fromList [(5,"a"), (3,"b")])) == 1
isJust (lookupIndex 6 (fromList [(5,"a"), (3,"b")])) == False
```

```
findIndex :: Ord k => k -> Map k a -> Int
```

Source

$O(\log n)$. Return the *index* of a key, which is its zero-based index in the sequence sorted by keys. The index is
 a number from 0 up to, but not including, the **size** of the map. Calls **error** when the key is not a **member** of
 the map.

```
findIndex 2 (fromList [(5,"a"), (3,"b")]) Error: element is not in the map
findIndex 3 (fromList [(5,"a"), (3,"b")]) == 0
findIndex 5 (fromList [(5,"a"), (3,"b")]) == 1
findIndex 6 (fromList [(5,"a"), (3,"b")]) Error: element is not in the map
```

```
elemAt :: Int -> Map k a -> (k, a)
```

Source

$O(\log n)$. Retrieve an element by its *index*, i.e. by its zero-based index in the sequence sorted by keys. If the
index is out of range (less than zero, greater or equal to **size** of the map), **error** is called.

```
elemAt 0 (fromList [(5,"a"), (3,"b")]) == (3,"b")
elemAt 1 (fromList [(5,"a"), (3,"b")]) == (5, "a")
elemAt 2 (fromList [(5,"a"), (3,"b")])   Error: index out of range
```

```
updateAt :: (k -> a -> Maybe a) -> Int -> Map k a -> Map k a
```

Source

$O(\log n)$. Update the element at *index*. Calls **error** when an invalid index is used.

```
updateAt (\_ _ -> Just "x") 0 (fromList [(5,"a"), (3,"b")]) == fromList [(3,
updateAt (\_ _ -> Just "x") 1 (fromList [(5,"a"), (3,"b")]) == fromList [(3,
updateAt (\_ _ -> Just "x") 2 (fromList [(5,"a"), (3,"b")])   Error: index o
updateAt (\_ _ -> Just "x") (-1) (fromList [(5,"a"), (3,"b")]) Error: index o
updateAt (\_ _ -> Nothing) 0 (fromList [(5,"a"), (3,"b")]) == singleton 5 "a"
updateAt (\_ _ -> Nothing) 1 (fromList [(5,"a"), (3,"b")]) == singleton 3 "b"
updateAt (\_ _ -> Nothing) 2 (fromList [(5,"a"), (3,"b")])   Error: index o
updateAt (\_ _ -> Nothing) (-1) (fromList [(5,"a"), (3,"b")]) Error: index o
```

```
deleteAt :: Int -> Map k a -> Map k a
```

Source

$O(\log n)$. Delete the element at *index*, i.e. by its zero-based index in the sequence sorted by keys. If the *index* is out of range (less than zero, greater or equal to **size** of the map), **error** is called.

```
deleteAt 0 (fromList [(5,"a"), (3,"b")]) == singleton 5 "a"
deleteAt 1 (fromList [(5,"a"), (3,"b")]) == singleton 3 "b"
deleteAt 2 (fromList [(5,"a"), (3,"b")])   Error: index out of range
deleteAt (-1) (fromList [(5,"a"), (3,"b")]) Error: index out of range
```

```
take :: Int -> Map k a -> Map k a
```

Source

Take a given number of entries in key order, beginning with the smallest keys.

```
take n = fromDistinctAscList . take n . toAscList
```

```
drop :: Int -> Map k a -> Map k a
```

Source

Drop a given number of entries in key order, beginning with the smallest keys.

```
drop n = fromDistinctAscList . drop n . toAscList
```

```
splitAt :: Int -> Map k a -> (Map k a, Map k a)
```

Source

$O(\log n)$. Split a map at a particular index.

```
splitAt !n !xs = (take n xs, drop n xs)
```

Min/Max

```
findMin :: Map k a -> (k, a)
```

Source

$O(\log n)$. The minimal key of the map. Calls **error** if the map is empty.

```
findMin (fromList [(5,"a"), (3,"b")]) == (3,"b")
findMin empty                          Error: empty map has no minimal element
```

```
findMax :: Map k a -> (k, a)
```

Source

$O(\log n)$. The maximal key of the map. Calls **error** if the map is empty.

```
findMax (fromList [(5,"a"), (3,"b")]) == (5,"a")
findMax empty                          Error: empty map has no maximal element
```

deleteMin :: **Map** k a -> **Map** k a # Source

$O(\log n)$. Delete the minimal key. Returns an empty map if the map is empty.

```
deleteMin (fromList [(5,"a"), (3,"b"), (7,"c")]) == fromList [(5,"a"), (7,"c")]
deleteMin empty == empty
```

deleteMax :: **Map** k a -> **Map** k a # Source

$O(\log n)$. Delete the maximal key. Returns an empty map if the map is empty.

```
deleteMax (fromList [(5,"a"), (3,"b"), (7,"c")]) == fromList [(3,"b"), (5,"a")]
deleteMax empty == empty
```

deleteFindMin :: **Map** k a -> ((k, a), **Map** k a) # Source

$O(\log n)$. Delete and find the minimal element.

```
deleteFindMin (fromList [(5,"a"), (3,"b"), (10,"c")]) == ((3,"b"), fromList [(5,"a"), (10,"c")])
deleteFindMin empty                                Error: can not return the minimal element
```

deleteFindMax :: **Map** k a -> ((k, a), **Map** k a) # Source

$O(\log n)$. Delete and find the maximal element.

```
deleteFindMax (fromList [(5,"a"), (3,"b"), (10,"c")]) == ((10,"c"), fromList [(5,"a"), (3,"b")])
deleteFindMax empty                                Error: can not return the maximal element
```

updateMin :: (a -> **Maybe** a) -> **Map** k a -> **Map** k a # Source

$O(\log n)$. Update the value at the minimal key.

```
updateMin (\ a -> Just ("X" ++ a)) (fromList [(5,"a"), (3,"b")]) == fromList [(3,"Xb"), (5,"a")]
updateMin (\ _ -> Nothing) (fromList [(5,"a"), (3,"b")]) == singleton 5 "a"
```

updateMax :: (a -> **Maybe** a) -> **Map** k a -> **Map** k a # Source

$O(\log n)$. Update the value at the maximal key.

```
updateMax (\ a -> Just ("X" ++ a)) (fromList [(5,"a"), (3,"b")]) == fromList [(3,"b"), (5,"Xa")]
updateMax (\ _ -> Nothing) (fromList [(5,"a"), (3,"b")]) == singleton 3 "b"
```

updateMinWithKey :: (k -> a -> **Maybe** a) -> **Map** k a -> **Map** k a # Source

$O(\log n)$. Update the value at the minimal key.

```
updateMinWithKey (\ k a -> Just ((show k) ++ ":" ++ a)) (fromList [(5,"a"), (3,"b")]) == fromList [(3,"3:b"), (5,"5:a")]
updateMinWithKey (\ _ _ -> Nothing) (fromList [(5,"a"), (3,"b")]) == fromList [(5,"a"), (3,"b")]
```

```
updateMaxWithKey :: (k -> a -> Maybe a) -> Map k a -> Map k a
```

Source

$O(\log n)$. Update the value at the maximal key.

```
updateMaxWithKey (\ k a -> Just ((show k) ++ ":" ++ a)) (fromList [(5,"a"), (3,"b")])
updateMaxWithKey (\ _ _ -> Nothing) (fromList [(5,"a"), (3,"b")])
```

```
minView :: Map k a -> Maybe (a, Map k a)
```

Source

$O(\log n)$. Retrieves the value associated with minimal key of the map, and the map stripped of that element, or **Nothing** if passed an empty map.

```
minView (fromList [(5,"a"), (3,"b")]) == Just ("b", singleton 5 "a")
minView empty == Nothing
```

```
maxView :: Map k a -> Maybe (a, Map k a)
```

Source

$O(\log n)$. Retrieves the value associated with maximal key of the map, and the map stripped of that element, or **Nothing** if passed an empty map.

```
maxView (fromList [(5,"a"), (3,"b")]) == Just ("a", singleton 3 "b")
maxView empty == Nothing
```

```
minViewWithKey :: Map k a -> Maybe ((k, a), Map k a)
```

Source

$O(\log n)$. Retrieves the minimal (key,value) pair of the map, and the map stripped of that element, or **Nothing** if passed an empty map.

```
minViewWithKey (fromList [(5,"a"), (3,"b")]) == Just ((3,"b"), singleton 5 "a")
minViewWithKey empty == Nothing
```

```
maxViewWithKey :: Map k a -> Maybe ((k, a), Map k a)
```

Source

$O(\log n)$. Retrieves the maximal (key,value) pair of the map, and the map stripped of that element, or **Nothing** if passed an empty map.

```
maxViewWithKey (fromList [(5,"a"), (3,"b")]) == Just ((5,"a"), singleton 3 "b")
maxViewWithKey empty == Nothing
```

Debugging

```
showTree :: (Show k, Show a) => Map k a -> String
```

Source

Deprecated: This function is being removed from the public API.

$O(n)$. Show the tree that implements the map. The tree is shown in a compressed, hanging format. See `showTreeWith`.

```
showTreeWith :: (k -> a -> String) -> Bool -> Bool -> Map k a -> String
```

Source

Deprecated: This function is being removed from the public API.

$O(n)$. The expression `(showTreeWith showElem hang wide map)` shows the tree that implements the map. Elements are shown using the `showElem` function. If `hang` is **True**, a *hanging* tree is shown otherwise a rotated tree is shown. If `wide` is **True**, an extra wide version is shown.

```

Map> let t = fromDistinctAscList [(x,()) | x <- [1..5]]
Map> putStrLn $ showTreeWith (\k x -> show (k,x)) True False t
(4,())
+--(2,())
|  +--(1,())
|  +--(3,())
+--(5,())

Map> putStrLn $ showTreeWith (\k x -> show (k,x)) True True t
(4,())
|
+--(2,())
| |
|  +--(1,())
|  |
|  +--(3,())
|  |
+--(5,())

Map> putStrLn $ showTreeWith (\k x -> show (k,x)) False True t
+--(5,())
|
(4,())
|
|  +--(3,())
|  |
+--(2,())
|
|  +--(1,())

```

valid :: Ord k => Map k a -> Bool

Source

$O(n)$. Test if the internal map structure is valid.

```

valid (fromAscList [(3,"b"), (5,"a")]) == True
valid (fromAscList [(5,"a"), (3,"b")]) == False

```