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CECS 329 Concepts of CS Theory

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NP-Complete Problems in Practice

Work Cited

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I’ve chosen to cover a problem in networking. Graph coloring is useful in wireless networking, despite it also being useful in other things like coloring maps, solving sudokus, and optimizing compilers. Graph coloring helps many wireless technologies assign channels to clients. A wireless channel is basically a frequency, so there’s problems because there can be many clients, but the access point can only broadcast to a finite number of frequencies. In this case clients can be thought of as a receiver and is generally a device that needs communication. My cell phone can be a client for a cell tower, a speaker can be a Bluetooth client for my phone, and my computer can be a Wi-Fi client for my wireless access point. Furthermore, if too many clients get on similar frequencies the signals may interfere with each other, especially when they’re nearby each other.

Given a graph G and an integer K, decide if there’s a coloring such that no two adjacent vertices share a color in G using K colors (Chauhan, 2020). The smallest number of colors required to color a graph is also simply referred to as the graph’s chromatic number. The question can be rephrased as ‘is the graphs chromatic number at most k?’

Given G(V, E), K ∈ ℤ

Is there a coloring such that

Colors C = (c1, c2, … ck), |C| = K, Colors vertices such that

The color of and vertex v = vx ∈ C and given any edge E=(va, vb)

The color of va ≠ vb

Graph coloring is used to model wireless networks by making a graph of the network. Each client becomes a vertex, and edges are added between the vertices of clients that may interfere with each other. K is the number of frequencies available (Bhuyan, 2023). The problem can also have special cases that the graph coloring problem doesn’t model well but can be altered to do so. Examples could be devices that work best on certain frequencies and if edges can have weights depicting how much interference is acceptable and trying to minimize the sum of edge weights of same-color adjacent vertices.

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Description automatically generated

Call this graph G.

Color(G, 3) is a True instance. Here’s an example of a valid coloring:

c1 -> v1

c1 -> v4

c2 -> v2

c3 -> v3­

Simply, color vertices 1 and 4 red, vertex 3 green, and vertex 2 red.

Color(G, 2) is a False instance since there does not exist a valid coloring without having at least one edge with nodes sharing a color.