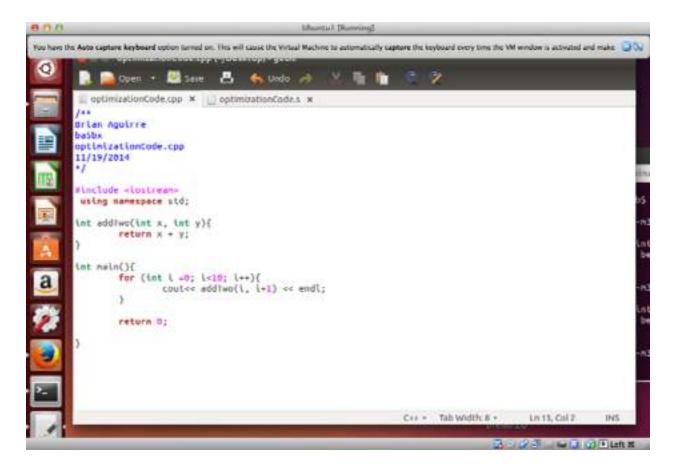
Brian Aguirre ba5bx postlab.pdf 11/21/2014

Optimization:

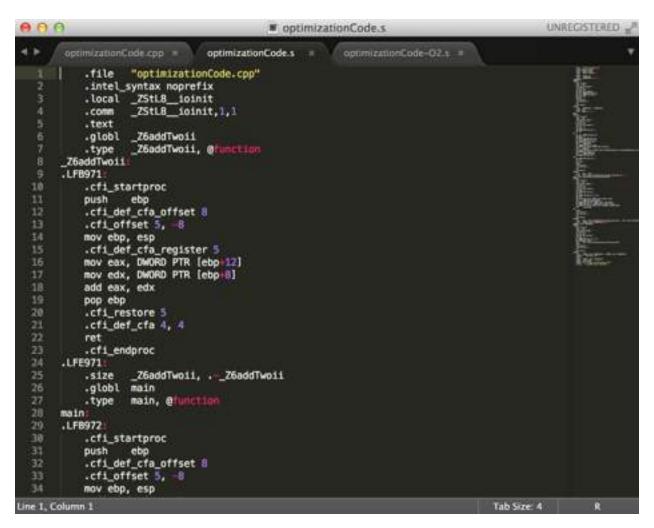
For the in-lab portion of lab 9, I chose to do question 3, pertaining to optimization within loops and function calls. The following screen capture is the code that was then converted into x86 using the -S tag:

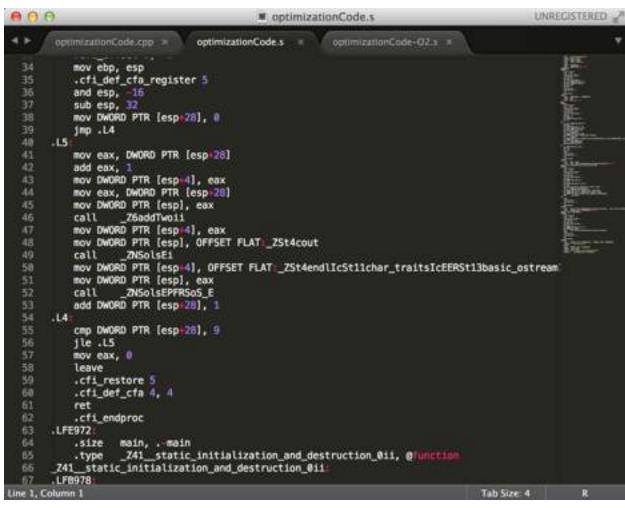


Notice how within the code, a method "addTwo" is created, which takes in two int types, x and y, and their sum is returned. This method is then called in the main method, where there is a for loop feeding it numbers from 0-10 (technically 0-11 as y is x + 1).

When using the -S tag for the .s file and then the -O2 tag for the optimized version of the .s file, the result of each is the following:

Normal Output:

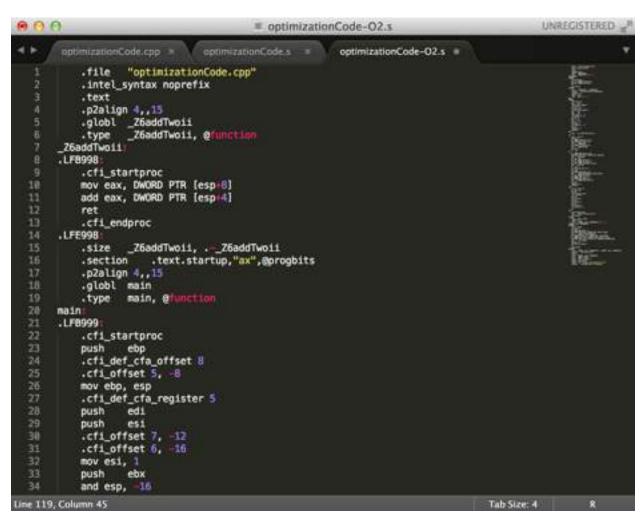


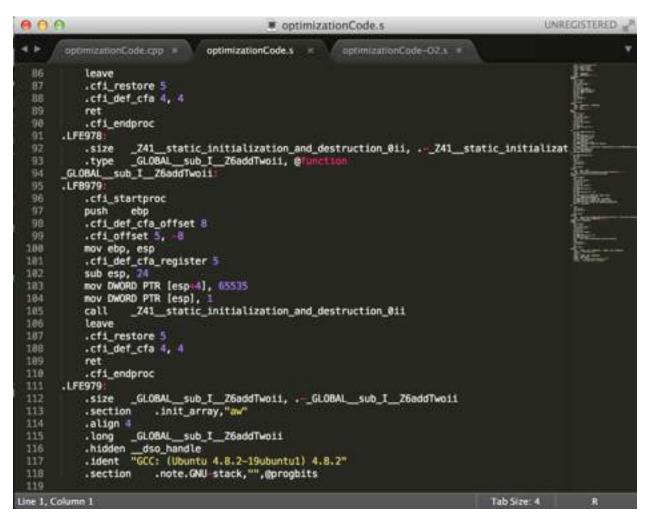


```
0.0

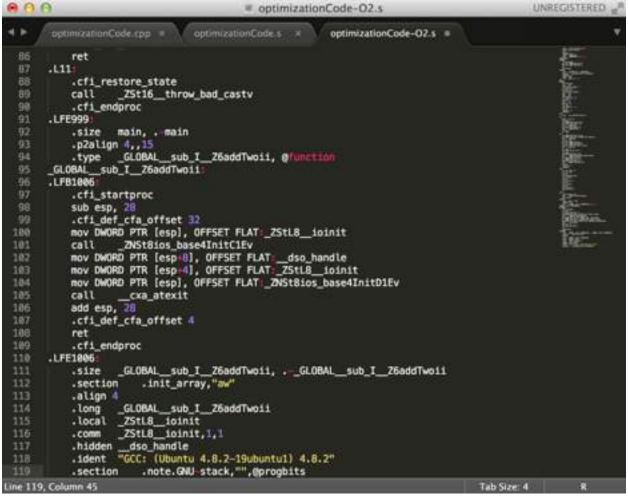
    optimizationCode.s

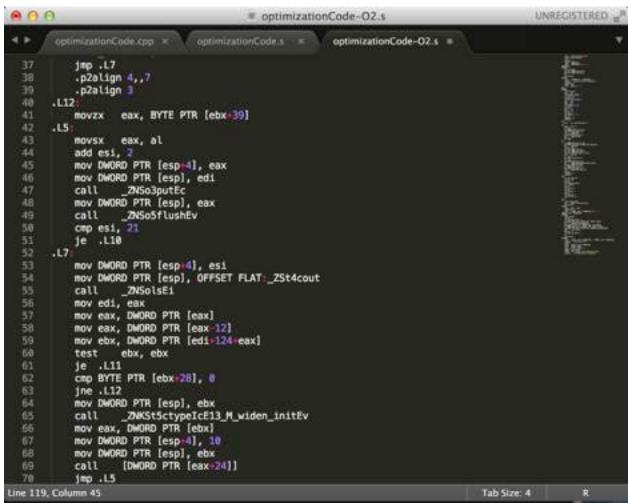
                                                                                                 UNREGISTERED J
        optimizationCode.cpp × optimizationCode.s × optimizationCode-O2.s ×
        Z41_static_initialization_and_destruction_0ii
       .LFB978
           .cfi_startproc
                  ebp
           push
           .cfi_def_cfa_offset 8
  7a
           .cfi_offset 5, 8
           mov ebp, esp
           .cfi_def_cfa_register 5
           sub esp, 24
cmp DWORD PTR [ebp=8], 1
                                                                                                       #1050 Ath.
           jne .L7
           cmp DWORD PTR [ebp-12], 65535
           jne .L7
           mov DWORD PTR [esp], OFFSET FLAT:_ZStL8__ioinit
  79
           call
                     ZNSt8ios_base4InitC1Ev
           mov DMORD PTR [esp 8], OFFSET FLAT:__dso_handle
           mov DWORD PTR [espid], OFFSET FLAT: ZStL8_ioinit
           mov DWORD PTR [esp], OFFSET FLAT _ZNSt8ios_base4InitD1Ev
  84
           call
                   __cxa_atexit
       .L7
           leave.
           .cfi_restore 5
           .cfi_def_cfa 4, 4
           ret
           .cfi_endproc
  98
       .LFE978
       .size _Z41_static_initialization_and_destruction_011, .-Z41_static_initializat
.type _GLOBAL_sub_I_Z6addTwo11, @function
GLOBAL_sub_I_Z6addTwo11;
       .LFB979
            .cfi_startproc
           push ebp
            .cfi_def_cfa_offset B
Line 1, Column 1
                                                                                       Tab Size: 4
```











Something I immediately notices was that the optimized version was much harder to follow. Although there are still things I cannot understand about the un-optimized version, the normal output, I can generally understand and see what it is doing. While in the optimized version, there are small little snippets and overall the structure is different. After reading a bit more about why the optimized version looks considerably different, I found two explanations. The first stated that there are less offsets within the optimized version and more use of registers. This can be seen in the comparison of .L5 in the non-optimized version of the .s file and the .L10 in the optimized version. In .L5 there are several lines in which offsets of ESP are being moved or called in order to do the calculations. In comparison, the optimized code in the .L10 section, it shows that there were at least 3 registers used, and throughout the rest of the code, it has less offsets since it instead uses the registers. Which means that probably using these registers to store values is much after than going to the stack and fetching the data based on their locations off of offsets.

Another reason the structure is different is because as I was looking at the two, the unoptimized version has a call for the _Z6addTwoii which is the method I defined in order to add the two numbers but the optimized version does not have a call line for it. As I read more, I found out that the optimized version does a different version of the loop which is called loop unwinding or loop unrolling. According to a UMD professor (included within the citations), this method allows for the program to be faster as it neglects to be referencing back for some kind of pointer or out bounds for the loop and rather just counts. This method is indeed faster but it also does not do much in terms of the size of the code as it is clear that the optimized version is literally as long (in terms of lines) as the un-optimized version.

Something else to note about the un-optimized version vs. the optimized version is the manner in which things are done. In the un-optimized version the line mov eax, 0 is done with xor eax, eax. These two do the same but perhaps the mov command is more taxing (albeit by a very small fraction of a second) than xor, but lines like these are found throughout the two, which definitely add up; lines like lea and other quicker ways to perform additions or moving of data. Something else is the fact that there are less call lines in the optimized version than the un-optimized version. This is probably because despite both having defined all of the code I wrote in c++, the optimized version probably considered certain things to be unnecessary and found it's own way to work around it (I don't completely understand it now, but I can definitely tell it performs things differently). But I guess that's something to also note, that even though the optimized version is faster and although in this case it takes the same amount of lines, it is also much more confusing, since the computer has essentially done things it's own way and users are left to try to figure out which way that is, and the unoptimized version, although slower, it's much more what I would write/have written in x86.

Citations:

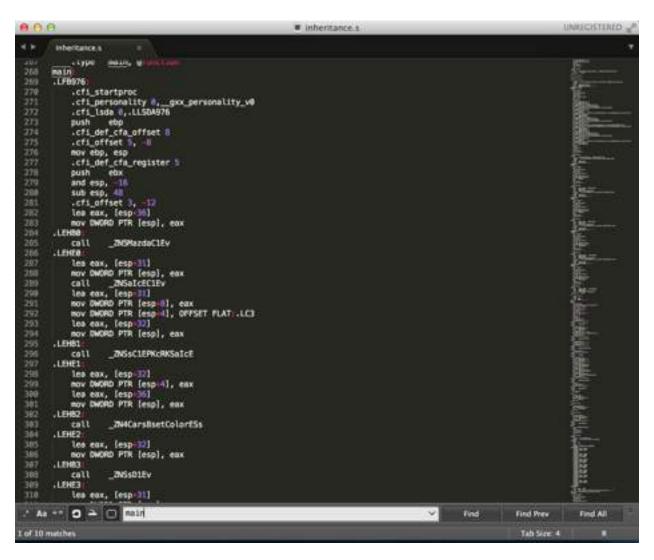
http://www.cs.virginia.edu/~evans/cs216/guides/x86.html http://www.eecg.toronto.edu/~amza/www.mindsec.com/files/x86regs.html http://www.cs.umd.edu/~meesh/cmsc411/website/proj01/proja/loop.html

Inheritance:

For the second portion of the report, I have chosen to do inheritance within C++. In order to test and view what the assembly equivalent for an inherited class and object declaration looked like, I wrote the following code:

```
■ inheritance.cpp
           inheritance.cpp
        #include <iostream>
#include <string>
 10112131557890051555555555600
         Uning namespace std;
         class Cars(
           string color;
               int year;
               int milage;
               word setColor(string c){
               color = c;
               void setYear(int y){
               year = y;
               void setMilage(int m){
               milage = m;
        4:
        class Mazda: millio Cars{
             void driveOneThousand(){
milage==1888;
             void carinfo(){
cout << "Year:" << year << endl;
cout << "Color:" << color << endl;
cout << "Milage:" << milage << endl;</pre>
        }
}:
        Int main(){
Hazda BriansCar;
             BriansCar.setColor("black");
             BriansCar.setYear(2014);
BriansCar.setMilage(0);
BriansCar.driveOneThousand();
             BriansCar.corInfo();
Line 1, Column 1
                                                                                                                                                   Tab Size: 4
```

The code above contains two classes, Cars and Mazda. Mazda is a child child of Cars, as the name would imply, since all Mazda vehicles are cars of some sort. This lab has introduced me on how to declare inherited classes. The key word to do so is class inherited_class: public parentClass{}. Although the keyword public does not have to be necessarily public, but that's a little beyond of what we're trying to achieve here. The code above allows the Class Mazda to utilize the public methods within Cars and the variables in order to set the attributes and data for each specific Mazda car. What I have done is I have written up a Mazda object called BriansCar. Then I set the attributes and and then print them out. The .s file, which contains the x86 assembly for the code above ends up being several lines long; 480 to be exact. For the sake of keeping it brief, only key screen shots will be portrayed.



Within the main: portion of the assembly code, we see that data is being offset to make sure it accessed the right information. Once that has been done, it calls the Mazda portion of the x86 since it is setting all the values for the Mazda object. Then when looking at the Mazda bit (shown in the following screen shot) we can see that the registers are being prepared to store data, but the most important part is that the Cars portion of the assembly is then called in line

227. That then allows for the Mazda to access the Cars methods, which then set the values for the space allotted for the variables for Car objects and anything that inherits the class.

```
# inheritance.s
         inheritance s
        ZNS Mazda CZEV
217
218
        .LF8984
             .cfi_startproc
219
220
221
222
223
224
225
226
             .cfs_def_cfa_offset 8
             .cfi_offset 5, 8
            mov ebp, esp
.cfi_def_cfa_register 5
            sub esp, 24
mov eax, DWORD PTR [ebp+8]
            mov DWORD PTR [esp], eax
227
228
             call
                       _ZN4CarsC2Ev
             leave
             .cfi_restore 5
238
231
             .cfi_def_cfa 4, 4
             .cfi_endproc
```

Following the call then takes you to the following snippet:

```
_ZN4CarsC2Ev
      .LFB979
167 ¥
          .cfi_startproc
          push
                 ebp
           .cfi_def_cfa_offset |
           .cfi_offset 5, 8
          mov ebp, esp
.cfi_def_cfa_register §
173
174
          mov eax, DWORD PTR [ebp:8]
175
176
           mov DWORD PTR [esp], eax
                   _ZNSsC1Ev
           call
178
179
           Leave
           .cfi_restore 5
           .cfi_def_cfa 4, 4
           ret
```

It is interesting to point out that the methods are defined before the class comes up. But the methods make references to the Class, and since Mazda makes reference to the Cars bit above, it is able to then find the memory space where the methods and variables are defined. Something interesting to note os that the methods I defined are shown first and then the class name is then shown; meaning setYear, setColor, setMilage, driveOneThousand, etc, are defined first within assembly and the the snippet above is shown. Which is perhaps similar to how ebp works; it stacks things up backwards, showing the methods first and then the class references because that way it can keep track of how things are aligned and it knows where memory is since it is placed linearly on the stack.

Looking back, there are also 2 portions of the Cars class, the one above and the one shown on the next page. I cannot follow it perfectly well, but could one serve as a constructor of the class and the other the actual memory reference for the whole class.

```
_ZN4CarsD2Ev, @
ZN4CarsD2Ev
LFB982
    .cfi_startproc
    push ebp
.cfi_def_cfa_offset 8
    .cfi_offset 5, -8
    mov ebp, esp
.cfi_def_cfa_register 5
   sub esp, 24
mov eax, DWORD PTR [ebp:8]
mov DWORD PTR [esp], eax
    call
              _ZNSsD1Ev
    leave
    .cfi_restore 5
    .cfi_def_cfa 4, 4
    ret
    .cfi_endproc
LFE982
    .size _ZN4CarsD2Ev, _ZN4CarsD2Ev
.weak _ZN4CarsD1Ev
.set _ZN4CarsD1Ev, ZN4CarsD2Ev
                 .text._ZN5MazdaC2Ev,"axG",@progbits,_ZN5MazdaC5Ev,comdat
    .section
    .align 2
              _ZN5MazdaC2Ev
             ZN5MazdaC2Ev, @fur
    .type
```

So that's how x86 is placing the methods and makes calls to each within the code, but something important to point out as well is that the variables, since they are just static ones and they are being set within the subroutine for setter methods, they are treated like global variables as shown below:

```
.section .rodata
.LC3:
.string "black"
.text
.globl main
.type main, @function
```

"Black" here is a .string and it is also being set within the main portion of the code.

Information From the following sources:

http://en.wikipedia.org/wiki/Call_stack#FRAME-POINTER http://en.wikipedia.org/wiki/X86 http://nongnu.askapache.com//pgubook/ProgrammingGroundUp-1-0-booksize.pdf