

Sean Shu (Sys3bb)
Brian Barbu (brb9da)

82-2=80

problem set 2

1) a. There exists an integer that is both prime and not odd.
↳ True, the number 2

b. If x is prime then x is not a perfect square.
↳ True, one is not a prime number

c. There exists an integer that is odd and a perfect square.
↳ True, the number 25, the perfect square of 5 that is odd.

2) a. negation: $\exists x \in \mathbb{Z}, \neg(x < 0)$
English: There exists an integer that is not less than zero

b. negation: $\forall x \in \mathbb{Z}, \neg(x = x + 1)$
English: For all integers x does not equal $x + 1$.

c. negation: $\forall x \in \mathbb{N}, \neg(x > 10)$
English: For all natural numbers, x is not greater than 10

d. negation: $\exists x \in \mathbb{N}, \neg(x > 10)$
English: There is a natural number that is not greater than 10.

3) a) \forall fish x , x has gills

negation: $\exists x \in \text{fish}, \neg(\text{has gills})$ ✓

b) \forall computers c , c has a CPU

negation: $\exists c \in \text{computer}, \neg(\text{has CPU})$ ✓

c) \exists a movie m , such that m is over 6 hours long

negation: $\forall m \in \text{movie}, \neg(m > 6 \text{ hours long})$ ✓

d) \exists a band b , such that b has won at least 10 grammy awards

negation: $\forall b \in \text{band}, \neg(b \geq 10 \text{ grammy awards})$ ✓

4) a) There is a student S that has seen the movie Casablanca ✓

b) If there is a student, then the student S has seen starwars ✓

c) For all students S , there exists a movie m such that all student have seen movie m ✓

d) There is a movie m , such that for all students S , student S has seen movie m . ✓

e) There is a student S , there is a student t , and there is a movie m such that S is not the same as t and both student S has seen movie m and student t has seen movie m . ✓

f) There is a student S , there is a student t , and for all movies m , such that S is not the same as t and if student S has seen movie m , then student t has seen movie m . ✓

5) The absence of error messages during translation of a computer program is a required element but not the only requirement and therefore is not enough for reasonable program correctness.

6) a) For all positive integers x , there exists a positive integer y such that $x = y + 1$.

→ False, goes wrong when $x=1$, y would have to ~~be~~ equal 0, which is not a positive number

b) For all integers x , there exists an integer y such that $x = y + 1$

→ True, if you consider 0 an integer, so when $x=1$, y can equal 0

c) There exists a real number x , such that for all real numbers y , $x = y + 1$

→ False, for a single x value not all y values will be able to make the statement true (only one y -value)

d) For all positive real numbers x , there exists a positive real number y , such that $xy = 1$

→ True, whatever x is, y will equal $1/x$.

e) For all integers x , there exists a real number y such that $xy = 1$.

→ True, because there has to be some sort of real number that will multiply x down to 1 (y will equal $1/x$ just like in D)

6) f. for all positive integers x and for all positive positive integers y , there exists a positive integer z such that $z = x - y$. ✓

↳ False, if x is greater than or equal to y , then z will be 0 or negative.

g. For all integers x and for all integers y , there exists an integer z such that $z = x - y$. ✓

↳ True, since z can be all integers; as the subtraction of two integers will result as an integer

h. There exists a positive real number u such that all positive real numbers v , $uv < v$

True. $u = .1$ $v = 10$ $uv < v$ because $1 < 10$ ✓

(as long as u is between 0 and 1)

7) a) $\forall x (P(x) \rightarrow Q(x))$ and $\forall x P(x) \rightarrow \forall x Q(x)$

Explanation: through an equivalence method, we know that $\forall x (P(x) \rightarrow Q(x)) \equiv \forall x (\sim P \vee Q)$, based off the rules of quantifiers, we know that $\forall x$ cannot be distributed into an "or" statement which makes it impossible for the two to be logically equivalent. In $\forall x (P(x) \rightarrow Q(x))$ $P(x)$ is sufficient for $Q(x)$ to be true, while in $\forall x P(x) \rightarrow \forall x Q(x)$, $P(x)$ is a necessary condition

b) $\exists x (P(x) \rightarrow Q(x))$ and $\forall x P(x) \rightarrow \exists x Q(x)$

1. $\exists x (\sim P(x) \vee Q(x))$ equality 1. $\neg (\forall x P(x)) \rightarrow \exists x Q(x)$ Distribute
2. $\exists x \sim P(x) \vee \exists x Q(x) \equiv$ 2. $\exists x \sim P(x) \vee \exists x Q(x)$ equality

So...

yes they are logically equivalent.

c. No they are not logically equivalent.

$\forall x (P(x) \leftrightarrow Q(x))$ is asking for all "sets" of where $P(x) \leftrightarrow Q(x)$ is valid, where as $\forall x P(x) \leftrightarrow \forall x Q(x)$ is saying all values of $P(x) \leftrightarrow$ all values of $Q(x)$; so one is asking for all within a "set" in x while the other says for all $P(x)$ to be true if and only if all $Q(x)$. So basically $\forall x (P(x) \leftrightarrow Q(x))$ is asking for all values where $(P(x) \leftrightarrow Q(x))$, where as the other says All $P(x)$ if and only if all $Q(x)$.

d) $\forall x(A \rightarrow P(x))$ and $A \rightarrow \forall x P(x)$

since A does not rely on x , the $\forall x$ in $\forall x(A \rightarrow P(x))$ skips over the " $A \rightarrow$ " and goes to the $P(x)$

$A \rightarrow \forall x(P(x))$ which is $A \rightarrow \forall x P(x)$ which is indeed equivalent to $A \rightarrow \forall x P(x)$ by null quantification

e) $\forall x(P(x) \rightarrow A)$ and $\exists x P(x) \rightarrow A$

in $\forall x(P(x) \rightarrow A)$, they are asking for all the values of x that meet the condition $(P(x) \rightarrow A)$; whereas $\exists x P(x) \rightarrow A$ says there exists an x where $(P(x) \rightarrow A)$, which includes all values of x that meet that condition, which makes the two equations logically equivalent.

f) Not equivalent because you cannot distribute universal quantifiers to "or" statements

8) a) True; any number would work because squaring it will make it a real number

b) False; if x is a negative number then you cannot square root y and get a negative

c) True; there will always be a reciprocal y that will make $xy = 1$ because $x \neq 0$.

d) False; there's no one x that can make all y in $xy = 1$.

e) True; for all real numbers $\frac{x+y}{2}$ will result in a real number z .

9) a) $A = \{x \in \mathbb{R} \mid x \leq b\}$ ✓

b) 4, 5, 69 ✓

c) no, because the values of x reach ∞ on the upper range ✓

d) 1.5, 2, 2.5 ✓

e) $A = \{x \in \mathbb{R} \mid x > b\}$ ✓