

# Strong Monads



*[github.com/fdilke/bewl](https://github.com/fdilke/bewl)*



from the set  $\mathcal{A}$  of right inverses  
to  $f$ , Section 17.  $f: P \xleftarrow{S} X \xrightarrow{I} X$ .  
There is an  $A \hookrightarrow P^X$ , transpose to  $A \times X \rightarrow P$   
Given we have to construct a map  
 $A^X \rightarrow A$  ; given  $h: X \rightarrow A$ ,  
define  $h^* \in A^X$  by:  $h^*(x) = h(x)(x) = \alpha(h(x), x)$   
\* Check that this is a right inverse of  $f$ :  
have to show  $x = f(h^*(x)) = f(h(x)(x))$   
which it is because  $h(x) \in A$ . So  $h^*(x) \in A$ .  
So this gives a structure map  $v: A^X \rightarrow A$   
\* Check monad laws: (1) it must respect  
constants: If  $h: X \rightarrow A$  is constantly  $a_0 \in A$ ,  
 $\neq h^*: x \mapsto h(x)(x) = \alpha(a_0, x)$ , so  $h^* = \alpha$ .  
(2) Must satisfy the "diagonal law"  
 $\iint f(x, y) dx dy = \int f(z, z) dz$ ,  
where  $f: X^2 \rightarrow A$  is a bifunction on  $X$  to  $A$ .  
But with  $\int_{x \in X} h(x) dx = [x \mapsto h(x)(x)] \in A$ ,  
RHS =  $\int p(y) dy$  where  $p(y) = \int f(x, y) dx$   
 $= [z \mapsto f(z, y)(z)]$  if  $y \in A \subseteq P^X$   
LHS =  $[w \mapsto \int f(w, y) p(y) dy]$   
 $= [w \mapsto \int \int f(w, y)(w) \alpha(y, w) dy]$   
 $= [w \mapsto \int f(w, w)(w) dw]$  But can't apply on



# The usual disclaimer

*It's hard core math, dumbed down into 10 slides*

# Monads and strong monads

- ❖ Why it took me a while to figure out that strength is even necessary
- ❖ Recap on monads
- ❖ Why this isn't enough - they need an extra structure
- ❖ What is the extra structure?
- ❖ How do you express this in code?



# It took a while to realise I needed strong monads, because there are many misconceptions

- ❖ When computer scientists talk about monads, they usually mean “strong monads” in the sense of category theory
- ❖ In a category with exponentials, a strong monad is a monad that has extra structure respecting the exponential
- ❖ So it’s an abuse of language (which I’ll use anyway) to ask if a given monad is strong. The question should really be “does it come with a naturally occurring strength”?
- ❖ The monads occurring in CS are all strong
- ❖ All monads on the category of sets are automatically strong
- ❖ In fact you have to go to some trouble to find a monad that is **NOT** strong (i.e. can’t be the monad part of a strong monad)
- ❖ Actually, the subset of functionality of a strong monad that makes them work with for comprehensions is the applicative functor, of which more anon

# Monads explained in one slide

- ❖ In Scala, classes like **List**, **Option** and **Future** can be used in *for comprehensions*
- ❖ This is because they have certain features in common
- ❖ For example, given an instance of *X*, you can create an instance of a **List[X]**
- ❖ and you can collapse a **List[List[X]]** into a **List[X]**
- ❖ Same with the others. They are all *monads*



# But... We want more

- ❖ Consider the set  $\mathbb{R}$  of real numbers (usually viewed as a line)
- ❖ It has a *ring* structure - the operators  $0, 1, +, -, *$  are defined on  $\mathbb{R}$
- ❖ Given a bunch of functions with real values, adding and multiplying them is the most natural thing in the world. So...
- ❖ Given some other object  $X$ , the mappings from  $X$  to  $\mathbb{R}$  inherit the ring structure. You can do arithmetic with them!
- ❖ Now generalise out of sight

# Why monads need to be strong

- ❖ If  $R$  is a ring, we can make  $R^X$  into a ring
- ❖ More generally, if  $R$  is an object in a category with exponentials, and  $R$  has been given a structure,  $R^X$  should have that structure too
- ❖ A lot of algebraic structures can be encoded as monads. For example, there is a monad whose algebras are precisely rings
- ❖ For algebras over a monad to have this nice  $R \rightarrow R^X$  property, we need the monad to be equipped with a *strength*



# What is the exact definition?

## Strong monad

From Wikipedia, the free encyclopedia

In [category theory](#), a **strong monad** over a [monoidal category](#)  $(C, \otimes, I)$  is a [monad](#)  $(T, \eta, \mu)$  together with a [natural transformation](#)  $t_{A,B} : A \otimes TB \rightarrow T(A \otimes B)$ , called (*tensorial strength*), such that the [diagrams](#)

$$\begin{array}{ccc}
 I \otimes TA & \xrightarrow{t_{I,A}} & T(I \otimes A) \\
 & \searrow \lambda_{TA} & \downarrow T(\lambda_A) \\
 & & TA
 \end{array}
 , \quad
 \begin{array}{ccc}
 (A \otimes B) \otimes TC & \xrightarrow{t_{A \otimes B, C}} & T((A \otimes B) \otimes C) \\
 \downarrow \alpha_{A,B,TC} & & \downarrow T(\alpha_{A,B,C}) \\
 A \otimes (B \otimes TC) & \xrightarrow{1_A \otimes t_{B,C}} A \otimes T(B \otimes C) \xrightarrow{t_{A, B \otimes C}} & T(A \otimes (B \otimes C))
 \end{array}
 ,$$
  

$$\begin{array}{ccc}
 A \otimes B & \xrightarrow{1_A \otimes \eta_B} & A \otimes TB \\
 & \searrow \eta_{A \otimes B} & \downarrow t_{A,B} \\
 & & T(A \otimes B)
 \end{array}
 , \text{ and }
 \begin{array}{ccc}
 A \otimes T^2 B & \xrightarrow{t_{A, TB}} & T(A \otimes TB) \xrightarrow{T(t_{A,B})} & T^2(A \otimes B) \\
 \downarrow 1_A \otimes \mu_B & & \downarrow \mu_{A \otimes B} \\
 A \otimes TB & \xrightarrow{t_{A,B}} & T(A \otimes B)
 \end{array}$$

These 4 rather horrific diagrams have to commute



# Luckily, all this is easy to code up in Bewl

```
object StrongMonad {  
  trait At[  
    M[X <: ~] <: ~,  
    X <: ~  
  ] extends Monad.At[M, X] {  
    def tensorialStrength[  
      Y <: ~  
    ](  
      dash: DOT[Y]  
    ):  
      X × M[Y] > M[X × Y]  
  }  
}
```

```
def sanityTest4[  
  A <: ~,  
  B <: ~  
](  
  a: DOT[A],  
  b: DOT[B]  
) =  
  tensorialStrength(a, b) o (  
    (a *- b) × (  
      η(b) o (a -* b)  
    )  
  ) shouldBe  
  η(a × b)
```

...and we can test-drive the  
construction of strong monads.

Next time, I will explain the point of all this

*“If people do not believe that mathematics is simple, it is only because they do not realize how complicated life is.”*

*—John von Neumann*

**THANK YOU**

<http://github.com/fdilke/bewl>