

Auslander-Buchweitz Approximation Theory for Extriangulated Categories

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Abstract

Extriangulated categories were introduced by Nakaoka and Palu as a simultaneous generalization of exact categories and triangulated categories. In this paper, we introduce and develop an analogous theory of Auslander-Buchweitz approximations for extriangulated categories. We establish the existence of precovers (and preenvelopes) and obtain characterizations of relative homological dimensions, which are based on certain subcategories under finiteness of resolutions. Finally, we give a description of cotorsion pairs on extriangulated categories under some conditions, and provide a characterization of silting subcategories on stable categories.

Keywords: Extriangulated category; Homological dimension; Cogenerator; Cotor-sion pair.

1. Introduction

Originated from the concept of injective envelopes, the approximation theory has attracted increasing interest and, hence, obtained considerable development especially in the context of module categories (see, for example [3, 9]). Auslander and Buchweitz [2] studied the ideals of injective envelopes and projective covers in terms of maximal Cohen-Macaulay approximations for certain modules. Indeed, they established their theory in the context of abelian categories and provided important applications. Inspired by their work, Mendoza Hernández, Sáenz Valadez, Santiago Vargas and Souta Salorio developed in [13, 14] an analogous theory of approximations for triangulated categories.

Triangulated categories and exact categories are two fundamental structures in mathematics. They are also important tools in many mathematical branches. It is well known that these two kinds of categories have some similarities, there are even direct connections between them. By extracting the similarities between triangulated categories and exact categories, Nakaoka and palu [15] recently introduced the notion of extriangulated categories, whose extriangulated structures are given by \mathbb{E} -triangles with some axioms. Except triangulated categories and exact categories, there are many examples for extriangulated categories [15, 16]. Hence, many results known to hold on exact categories and triangulated categories can be unified in the same framework [11, 12, 15, 16]. Motivated by the ideal, we introduce and develop an analogous theory of approximations in the sense of Auslander and Buchweitz [2] for extriangulated categories. Let \mathcal{C} be an extriangulated category with enough projectives and injectives. The main results deal with a pair $(\mathcal{X}, \mathcal{W})$ of subcategories of \mathcal{C} , where \mathcal{X} is closed under extensions and \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . We consider the subcategory $\hat{\mathcal{X}}$ of \mathcal{C} consisting of all

objects with a finite resolution by objects of \mathcal{X} . Moreover, a notion of \mathcal{X} -resolution dimensions is also introduced, which is compared with other relative homological dimensions. We prove that any object of $\widehat{\mathcal{X}}$ admits two \mathbb{E} -triangles: one giving rise to an \mathcal{X} -precover and the other to a $\widehat{\mathcal{W}}$ -preenvelope, which is used to construct a cotorsion pair on the extriangulated category $\widehat{\mathcal{X}}$. Whenever \mathcal{C} is a Frobenius extriangulated category, we give a characterization of hereditary cotorsion pairs $(\mathcal{U}, \mathcal{V})$ on the extriangulated category \mathcal{C} with $\widehat{\mathcal{U}} = \check{\mathcal{V}} = \mathcal{C}$. As an application, we also obtain a characterization of silting subcategories on the stable category $\underline{\mathcal{C}}$.

The paper is organized as follows. In Section 2, we recall the definition of an extriangulated category and outline some basic properties that will be used later. In Section 3, we study the notion of \mathcal{X} -resolution dimensions and give some relationships between the relative projective dimension and the \mathcal{X} -resolution dimension. Moreover, we focus our attention on the notion of \mathcal{X} -injective cogenerators for \mathcal{X} and establish the existence of \mathcal{X} -precovers and $\widehat{\mathcal{W}}$ -preenvelopes. If \mathcal{X} is closed under extensions, we also obtain that $\widehat{\mathcal{X}}$ is closed under extensions, hence it is an extriangulated category, which is essential for our main result. In Section 4, we define hereditary cotorsion pairs on extriangulated categories with enough projectives and injectives. If \mathcal{C} is a Frobenius extriangulated category, we establish a bijective correspondence between hereditary cotorsion pairs $(\mathcal{U}, \mathcal{V})$ on the extriangulated category \mathcal{C} with $\widehat{\mathcal{U}} = \check{\mathcal{V}} = \mathcal{C}$ and that of certain specially precovering classes. Finally, in Corollary 4.10, we give a characterization of silting subcategories on the stable category $\underline{\mathcal{C}}$ if \mathcal{C} is a Frobenius extriangulated category.

2. Preliminaries

Throughout this paper, \mathcal{C} denotes an additive category, by the term “*subcategory*” we always mean a full additive subcategory of an additive category closed under isomorphisms and direct summands. We denote by $\mathcal{C}(A, B)$ the set of morphisms from A to B in \mathcal{C} .

Let \mathcal{X} and \mathcal{Y} be two subcategories of \mathcal{C} , a morphism $f : X \rightarrow C$ in \mathcal{C} is said to be an \mathcal{X} -precover of C if $X \in \mathcal{X}$ and $\mathcal{C}(X', f) : \mathcal{C}(X', X) \rightarrow \mathcal{C}(X', C)$ is surjective, $\forall X' \in \mathcal{X}$. If any $C \in \mathcal{Y}$ admits an \mathcal{X} -precover, then \mathcal{X} is called a precovering class in \mathcal{Y} . By dualizing the definition above, we get the notion of an \mathcal{X} -preenvelope of C and a preenveloping class in \mathcal{Y} , for details, see [3, 9].

Let us briefly recall some definitions and basic properties of extriangulated categories from [15].

Definition 2.1. [15, Definition 2.1] Suppose that \mathcal{C} is equipped with an additive bifunctor

$$\mathbb{E} : \mathcal{C}^{op} \times \mathcal{C} \rightarrow \text{Ab},$$

where Ab is the category of abelian groups. For any objects $A, C \in \mathcal{C}$, an element $\delta \in \mathbb{E}(C, A)$ is called an \mathbb{E} -extension. Thus formally, an \mathbb{E} -extension is a triple (A, δ, C) . For any $A, C \in \mathcal{C}$, the zero element $0 \in \mathbb{E}(C, A)$ is called the split \mathbb{E} -extension.

Let $\delta \in \mathbb{E}(C, A)$ be any \mathbb{E} -extension. By the functoriality, for any $a \in \mathcal{C}(A, A')$ and $c \in \mathcal{C}(C', C)$, we have \mathbb{E} -extensions

$$\mathbb{E}(C, a)(\delta) \in \mathbb{E}(C, A') \text{ and } \mathbb{E}(c, A)(\delta) \in \mathbb{E}(C', A).$$

We abbreviately denote them by $a_*\delta$ and $c^*\delta$. In this terminology, we have

$$\mathbb{E}(c, a)(\delta) = c^*a_*\delta = a_*c^*\delta$$

in $\mathbb{E}(C', A')$.

Definition 2.2. [15, Definition 2.3] Let $\delta \in \mathbb{E}(C, A)$ and $\delta' \in \mathbb{E}(C', A')$ be any pair of \mathbb{E} -extensions. A morphism $(a, c) : \delta \rightarrow \delta'$ of \mathbb{E} -extensions is a pair of morphisms $a \in \mathcal{C}(A, A')$ and $c \in \mathcal{C}(C, C')$ in \mathcal{C} satisfying the equality

$$a_*\delta = c^*\delta'.$$

We simply denote it as $(a, c) : \delta \rightarrow \delta'$.

Definition 2.3. [15, Definition 2.6] Let $\delta = (A, \delta, C)$ and $\delta' = (A', \delta', C')$ be any pair of \mathbb{E} -extensions. Let

$$C \xrightarrow{\iota_C} C \oplus C' \xleftarrow{\iota_{C'}} C' \quad \text{and} \quad A \xleftarrow{p_A} A \oplus A' \xrightarrow{p_{A'}} A'$$

be coproduct and product in \mathcal{C} , respectively. Remark that, by the additivity of \mathbb{E} , we have a natural isomorphism

$$\mathbb{E}(C \oplus C', A \oplus A') \simeq \mathbb{E}(C, A) \oplus \mathbb{E}(C, A') \oplus \mathbb{E}(C', A) \oplus \mathbb{E}(C', A').$$

Let $\delta \oplus \delta' \in \mathbb{E}(C \oplus C', A \oplus A')$ be the element corresponding to $(\delta, 0, 0, \delta')$ through this isomorphism. This is the unique element which satisfies

$$\mathbb{E}(\iota_C, p_A)(\delta \oplus \delta') = \delta, \quad \mathbb{E}(\iota_C, p_{A'})(\delta \oplus \delta') = 0, \quad \mathbb{E}(\iota_{C'}, p_A)(\delta \oplus \delta') = 0, \quad \mathbb{E}(\iota_{C'}, p_{A'})(\delta \oplus \delta') = \delta'.$$

Definition 2.4. [15, Definition 2.7] Let $A, C \in \mathcal{C}$ be any pair of objects. Two sequences of morphisms in \mathcal{C}

$$A \xrightarrow{x} B \xrightarrow{y} C \quad \text{and} \quad A \xrightarrow{x'} B' \xrightarrow{y'} C$$

are said to be equivalent if there exists an isomorphism $b \in \mathcal{C}(B, B')$ which makes the following diagram commutative

$$\begin{array}{ccccc} & & B & & \\ & \nearrow x & \downarrow \simeq b & \searrow y & \\ A & & & & C \\ & \searrow x' & \downarrow b & \nearrow y' & \\ & & B' & & \end{array}$$

We denote the equivalence class of $A \xrightarrow{x} B \xrightarrow{y} C$ by $[A \xrightarrow{x} B \xrightarrow{y} C]$.

Definition 2.5. [15, Definition 2.8] (1) For any $A, C \in \mathcal{C}$, we denote

$$0 = [A \xrightarrow{\begin{bmatrix} 1 \\ 0 \end{bmatrix}} A \oplus C \xrightarrow{\begin{bmatrix} 0 & 1 \end{bmatrix}} C]$$

(2) For any two classes $[A \xrightarrow{x} B \xrightarrow{y} C]$ and $[A' \xrightarrow{x'} B' \xrightarrow{y'} C']$, we denote

$$[A \xrightarrow{x} B \xrightarrow{y} C] \oplus [A' \xrightarrow{x'} B' \xrightarrow{y'} C'] = [A \oplus A' \xrightarrow{x \oplus x'} B \oplus B' \xrightarrow{y \oplus y'} C \oplus C'].$$

Definition 2.6. [15, Definition 2.9] Let \mathfrak{s} be a correspondence which associates an equivalence class

$$\mathfrak{s}(\delta) = [A \xrightarrow{x} B \xrightarrow{y} C]$$

to any \mathbb{E} -extension $\delta \in \mathbb{E}(C, A)$. This \mathfrak{s} is called a *realization* of \mathbb{E} , if it satisfies the following condition (\star) . In this case, we say that the sequence $A \xrightarrow{x} B \xrightarrow{y} C$ realizes δ , whenever it satisfies $\mathfrak{s}(\delta) = [A \xrightarrow{x} B \xrightarrow{y} C]$.

(\star) Let $\delta \in \mathbb{E}(C, A)$ and $\delta' \in \mathbb{E}(C', A')$ be any pair of \mathbb{E} -extensions, with

$$\mathfrak{s}(\delta) = [A \xrightarrow{x} B \xrightarrow{y} C] \text{ and } \mathfrak{s}(\delta') = [A' \xrightarrow{x'} B' \xrightarrow{y'} C'] .$$

Then, for any morphism $(a, c) : \delta \rightarrow \delta'$, there exists $b \in \mathcal{C}(B, B')$ which makes the following diagram commutative

$$\begin{array}{ccccc} A & \xrightarrow{x} & B & \xrightarrow{y} & C \\ a \downarrow & & b \downarrow & & c \downarrow \\ A' & \xrightarrow{x'} & B' & \xrightarrow{y'} & C' . \end{array}$$

In the above situation, we say that the triplet (a, b, c) realizes (a, c) .

Definition 2.7. [15, Definition 2.10] Let \mathcal{C}, \mathbb{E} be as above. A realization of \mathbb{E} is said to be *additive*, if it satisfies the following conditions.

- (i) For any $A, C \in \mathcal{C}$, the split \mathbb{E} -extension $0 \in \mathbb{E}(C, A)$ satisfies $\mathfrak{s}(0) = 0$.
- (ii) For any pair of \mathbb{E} -extensions $\delta \in \mathbb{E}(C, A)$ and $\delta' \in \mathbb{E}(C', A')$, we have

$$\mathfrak{s}(\delta \oplus \delta') = \mathfrak{s}(\delta) \oplus \mathfrak{s}(\delta') .$$

Definition 2.8. [15, Definition 2.12] A triplet $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ is called an *extriangulated category* if it satisfies the following conditions.

- (ET1) $\mathbb{E} : \mathcal{C}^{op} \times \mathcal{C} \rightarrow \mathbf{Ab}$ is an additive bifunctor.
- (ET2) \mathfrak{s} is an additive realization of \mathbb{E} .
- (ET3) Let $\delta \in \mathbb{E}(C, A)$ and $\delta' \in \mathbb{E}(C', A')$ be any pair of \mathbb{E} -extensions, realized as

$$\mathfrak{s}(\delta) = [A \xrightarrow{x} B \xrightarrow{y} C] \text{ and } \mathfrak{s}(\delta') = [A' \xrightarrow{x'} B' \xrightarrow{y'} C'] .$$

For any commutative square

$$\begin{array}{ccccc} A & \xrightarrow{x} & B & \xrightarrow{y} & C \\ a \downarrow & & b \downarrow & & \\ A' & \xrightarrow{x'} & B' & \xrightarrow{y'} & C' . \end{array}$$

in \mathcal{C} , there exists a morphism $(a, c) : \delta \rightarrow \delta'$ which is realized by (a, b, c) .

$(ET3)^{op}$ Let $\delta \in \mathbb{E}(C, A)$ and $\delta' \in \mathbb{E}(C', A')$ be any pair of \mathbb{E} -extensions, realized by

$$A \xrightarrow{x} B \xrightarrow{y} C \text{ and } A' \xrightarrow{x'} B' \xrightarrow{y'} C'$$

respectively. For any commutative square

$$\begin{array}{ccccc} A & \xrightarrow{x} & B & \xrightarrow{y} & C \\ & & \downarrow b & & \downarrow c \\ A' & \xrightarrow{x'} & B' & \xrightarrow{y'} & C' \end{array}$$

in \mathcal{C} , there exists a morphism $(a, c) : \delta \rightarrow \delta'$ which is realized by (a, b, c) .

(ET_4) Let $\delta \in \mathbb{E}(D, A)$ and $\delta' \in \mathbb{E}(F, B)$ be \mathbb{E} -extensions realized by

$$A \xrightarrow{f} B \xrightarrow{f'} D \quad \text{and} \quad B \xrightarrow{g} C \xrightarrow{g'} F$$

respectively. Then there exist an object $E \in \mathcal{C}$, a commutative diagram

$$\begin{array}{ccccc} A & \xrightarrow{f} & B & \xrightarrow{f'} & D \\ \parallel & & \downarrow g & & \downarrow d \\ A & \xrightarrow{h} & C & \xrightarrow{h'} & E \\ & & \downarrow g' & & \downarrow e \\ & & F & \xlongequal{\quad} & F \end{array}$$

in \mathcal{C} , and an \mathbb{E} -extension $\delta'' \in \mathbb{E}(E, A)$ realized by $A \xrightarrow{h} C \xrightarrow{h'} E$, which satisfy the following compatibilities.

- (i) $D \xrightarrow{d} E \xrightarrow{e} F$ realizes $f'_*\delta'$,
- (ii) $d^*\delta'' = \delta$,
- (iii) $f_*\delta'' = e^*\delta'$.

(ET_4)^{op} Let $\delta \in \mathbb{E}(B, D)$ and $\delta' \in \mathbb{E}(C, F)$ be \mathbb{E} -extensions realized by

$$D \xrightarrow{f'} A \xrightarrow{f} B \quad \text{and} \quad F \xrightarrow{g'} B \xrightarrow{g} C$$

respectively. Then there exist an object $E \in \mathcal{C}$, a commutative diagram

$$\begin{array}{ccccc} D & \xrightarrow{d} & E & \xrightarrow{e} & F \\ \parallel & & \downarrow h' & & \downarrow g' \\ D & \xrightarrow{f'} & A & \xrightarrow{f} & B \\ & & \downarrow h & & \downarrow g \\ & & C & \xlongequal{\quad} & C \end{array}$$

in \mathcal{C} , and an \mathbb{E} -extension $\delta'' \in \mathbb{E}(C, E)$ realized by $E \xrightarrow{h'} A \xrightarrow{h} C$, which satisfy the following compatibilities.

- (i) $D \xrightarrow{d} E \xrightarrow{e} F$ realizes $g'_*\delta$,
- (ii) $\delta' = e_*\delta''$,
- (iii) $d_*\delta = g^*\delta''$.

For an extriangulated category \mathcal{C} , we use the following notation:

- A sequence $A \xrightarrow{x} B \xrightarrow{y} C$ is called a conflation if it realizes some \mathbb{E} -extension $\delta \in \mathbb{E}(C, A)$.
- A morphism $f \in \mathcal{C}(A, B)$ is called an inflation if it admits some conflation $A \xrightarrow{f} B \longrightarrow C$.
- A morphism $f \in \mathcal{C}(A, B)$ is called a deflation if it admits some conflation $K \longrightarrow A \xrightarrow{f} B$.
- If a conflation $A \xrightarrow{x} B \xrightarrow{y} C$ realizes $\delta \in \mathbb{E}(C, A)$, we call the pair $(A \xrightarrow{x} B \xrightarrow{y} C, \delta)$ an \mathbb{E} -triangle, and write it in the following way.

$$A \xrightarrow{x} B \xrightarrow{y} C - \overset{\delta}{-} \triangleright$$

We usually do not write this “ δ ” if it is not used in the argument.

- Given an \mathbb{E} -triangle $A \xrightarrow{x} B \xrightarrow{y} C - \overset{\delta}{-} \triangleright$, we call A the CoCone of $y : B \longrightarrow C$, and denote it by $CoCone(B \longrightarrow C)$, or $CoCone(y)$; we call C the Cone of $x : A \longrightarrow B$, and denote it by $Cone(A \longrightarrow B)$, or $Cone(x)$.

- Let $A \xrightarrow{x} B \xrightarrow{y} C - \overset{\delta}{-} \triangleright$ and $A' \xrightarrow{x'} B' \xrightarrow{y'} C' - \overset{\delta'}{-} \triangleright$ be any pair of \mathbb{E} -triangles. If a triplet (a, b, c) realizes $(a, c) : \delta \rightarrow \delta'$, then we write it as

$$\begin{array}{ccccc} A & \xrightarrow{x} & B & \xrightarrow{y} & C - \overset{\delta}{-} \triangleright \\ \downarrow a & & \downarrow b & & \downarrow c \\ A' & \xrightarrow{x'} & B' & \xrightarrow{y'} & C' - \overset{\delta'}{-} \triangleright \end{array}$$

and call (a, b, c) a morphism of \mathbb{E} -triangles.

- A subcategory \mathcal{T} of \mathcal{C} is called extension-closed if \mathcal{T} is closed under extensions, i.e. for any \mathbb{E} -triangle $A \xrightarrow{x} B \xrightarrow{y} C - \overset{\delta}{-} \triangleright$ with $A, C \in \mathcal{T}$, we have $B \in \mathcal{T}$.

Example 2.9. (1) Exact category \mathcal{B} can be viewed as an extriangulated category. For the definition and basic properties of an exact category, see [6]. In fact, a biadditive functor $\mathbb{E} := \text{Ext}_{\mathcal{B}}^1 : \mathcal{B}^{\text{op}} \times \mathcal{B} \rightarrow \text{Ab}$. Let $A, C \in \mathcal{B}$ be any pair of objects. Define $\text{Ext}_{\mathcal{B}}^1(C, A)$ to be the collection of all equivalence classes of short exact sequences of the form $A \xrightarrow{x} B \xrightarrow{y} C$. We denote the equivalence class by $[A \xrightarrow{x} B \xrightarrow{y} C]$ as before. For any $\delta = [A \xrightarrow{x} B \xrightarrow{y} C] \in \text{Ext}_{\mathcal{B}}^1(C, A)$, define the realization $\mathfrak{s}(\delta)$ of $[A \xrightarrow{x} B \xrightarrow{y} C]$ to be δ itself. For more details, see [15, Example 2.13].

(2) Let \mathcal{C} be a triangulated category with shift functor $[1]$. Put $\mathbb{E} := \mathcal{C}(-, -[1])$. For any $\delta \in \mathbb{E}(C, A) = \mathcal{C}(C, A[1])$, take a triangle

$$A \xrightarrow{x} B \xrightarrow{y} C \xrightarrow{\delta} A[1]$$

and define as $\mathfrak{s}(\delta) = [A \xrightarrow{x} B \xrightarrow{y} C]$. Then $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ is an extriangulated category. It is easy to see that extension closed subcategories of triangulated categories are also extriangulated categories. For more details, see [15, Proposition 3.22].

(3) Let \mathcal{C} be an extriangulated category, and \mathcal{J} a subcategory of \mathcal{C} . If $\mathcal{J} \subseteq \text{Proj}(\mathcal{C}) \cap \text{Inj}(\mathcal{C})$, where $\text{Proj}(\mathcal{C})$ is the full category of projective objects in \mathcal{C} and $\text{Inj}(\mathcal{C})$ is the full category of

injective objects in \mathcal{C} , then \mathcal{C}/\mathcal{J} is an extriangulated category. This construction gives extriangulated categories which are neither exact nor triangulated in general. For more details, see [15, Proposition 3.30]

Lemma 2.10. [15, Proposition 3.15] *Let $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ be an extriangulated category. Then the following hold.*

(1) *Let C be any object, and let $A_1 \xrightarrow{x_1} B_1 \xrightarrow{y_1} C \xrightarrow{\delta_1} \triangleright$ and $A_2 \xrightarrow{x_2} B_2 \xrightarrow{y_2} C \xrightarrow{\delta_2} \triangleright$ be any pair of \mathbb{E} -triangles. Then there is a commutative diagram in \mathcal{C}*

$$\begin{array}{ccccc} & & A_2 & \xlongequal{\quad} & A_2 \\ & & \downarrow m_2 & & \downarrow x_2 \\ A_1 & \xrightarrow{m_1} & M & \xrightarrow{e_1} & B_2 \\ \parallel & & \downarrow e_2 & & \downarrow y_2 \\ A_1 & \xrightarrow{x_1} & B_1 & \xrightarrow{y_1} & C \end{array}$$

which satisfies $\mathfrak{s}(y_2^ \delta_1) = [A_1 \xrightarrow{m_1} M \xrightarrow{e_1} B_2]$ and $\mathfrak{s}(y_1^* \delta_2) = [A_2 \xrightarrow{m_2} M \xrightarrow{e_2} B_1]$.*

(2) *Let A be any object, and let $A \xrightarrow{x_1} B_1 \xrightarrow{y_1} C_1 \xrightarrow{\delta_1} \triangleright$ and $A \xrightarrow{x_2} B_2 \xrightarrow{y_2} C_2 \xrightarrow{\delta_2} \triangleright$ be any pair of \mathbb{E} -triangles. Then there is a commutative diagram in \mathcal{C}*

$$\begin{array}{ccccc} A & \xrightarrow{x_1} & B_1 & \xrightarrow{y_1} & C_1 \\ \downarrow x_2 & & \downarrow m_2 & & \parallel \\ B_2 & \xrightarrow{m_1} & M & \xrightarrow{e_1} & C_1 \\ \downarrow y_2 & & \downarrow e_2 & & \\ C_2 & \xlongequal{\quad} & C_2 & & \end{array}$$

which satisfies $\mathfrak{s}(x_2^ \delta_1) = [B_2 \xrightarrow{m_1} M \xrightarrow{e_1} C_1]$ and $\mathfrak{s}(x_1^* \delta_2) = [B_1 \xrightarrow{m_2} M \xrightarrow{e_2} C_2]$.*

Assume that $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ is an extriangulated category. By Yoneda's Lemma, any \mathbb{E} -extension $\delta \in \mathbb{E}(C, A)$ induces natural transformations

$$\delta_{\#} : \mathcal{C}(-, C) \Rightarrow \mathbb{E}(-, A) \text{ and } \delta^{\#} : \mathcal{C}(A, -) \Rightarrow \mathbb{E}(C, -).$$

For any $X \in \mathcal{C}$, these $(\delta_{\#})_X$ and $\delta_X^{\#}$ are given as follows:

- (1) $(\delta_{\#})_X : \mathcal{C}(X, C) \Rightarrow \mathbb{E}(X, A); f \mapsto f^* \delta.$
- (2) $\delta_X^{\#} : \mathcal{C}(A, X) \Rightarrow \mathbb{E}(C, X); g \mapsto g_* \delta.$

Lemma 2.11. [15, Corollary 3.12] *Let $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ be an extriangulated category, and*

$$A \xrightarrow{x} B \xrightarrow{y} C \xrightarrow{\delta} \triangleright$$

an \mathbb{E} -triangle. Then we have the following long exact sequences:

$$\begin{aligned} \mathcal{C}(C, -) &\xrightarrow{\mathcal{C}(y, -)} \mathcal{C}(B, -) \xrightarrow{\mathcal{C}(x, -)} \mathcal{C}(A, -) \xrightarrow{\delta_{\#}} \mathbb{E}(C, -) \xrightarrow{\mathbb{E}(y, -)} \mathbb{E}(B, -) \xrightarrow{\mathbb{E}(x, -)} \mathbb{E}(A, -); \\ \mathcal{C}(-, A) &\xrightarrow{\mathcal{C}(-, x)} \mathcal{C}(-, B) \xrightarrow{\mathcal{C}(-, y)} \mathcal{C}(-, C) \xrightarrow{\delta^{\#}} \mathbb{E}(-, A) \xrightarrow{\mathbb{E}(-, x)} \mathbb{E}(-, B) \xrightarrow{\mathbb{E}(-, y)} \mathbb{E}(-, C). \end{aligned}$$

The following lemma was given in [12, Proposition 1.20], which is another version of [15, Corollary 3.16].

Lemma 2.12. [12, Proposition 1.20] *Let $A \xrightarrow{x} B \xrightarrow{y} C \xrightarrow{\delta} \gg$ be an \mathbb{E} -triangle and $f : A \rightarrow B$ any morphism, and let $D \xrightarrow{d} E \xrightarrow{e} C \xrightarrow{f_*\delta} \gg$ be any \mathbb{E} -triangle realizing $f_*\delta$. Then there is a morphism g which gives a morphisms of \mathbb{E} -triangles*

$$\begin{array}{ccccc} A & \xrightarrow{x} & B & \xrightarrow{y} & C \xrightarrow{\delta} \gg \\ \downarrow f & & \downarrow g & & \parallel \\ D & \xrightarrow{d} & E & \xrightarrow{e} & C \xrightarrow{f_*(\delta)} \gg \end{array}$$

Moreover, $A \xrightarrow{\begin{bmatrix} -f \\ x \end{bmatrix}} D \oplus B \xrightarrow{\begin{bmatrix} d & g \end{bmatrix}} E \xrightarrow{e_*\delta} \gg$ is an \mathbb{E} -triangle.

We recall some concepts from [15]. Let \mathcal{C}, \mathbb{E} be as above. An object $P \in \mathcal{C}$ is called projective if it satisfies the following condition.

- For any \mathbb{E} -triangle $A \xrightarrow{x} B \xrightarrow{y} C \xrightarrow{\delta} \gg$ and any morphism $c \in \mathcal{C}(P, C)$, there exists $b \in \mathcal{C}(P, B)$ satisfying $y \circ b = c$.

Injective objects are defined dually.

We denote the subcategory consisting of projective objects in \mathcal{C} by $Proj(\mathcal{C})$. Dually, the subcategory of injective objects in \mathcal{C} is denoted by $Inj(\mathcal{C})$.

- We say \mathcal{C} has enough projectives (enough injectives, resp.) if for any object $C \in \mathcal{C}$ ($A \in \mathcal{C}$, resp.), there exists an \mathbb{E} -triangle

$$A \xrightarrow{x} P \xrightarrow{y} C \xrightarrow{\delta} \gg \quad (A \xrightarrow{x} I \xrightarrow{y} C \xrightarrow{\delta} \gg)$$

satisfying $P \in Proj(\mathcal{C})$ ($I \in Inj(\mathcal{C})$, resp.).

In this case, A is called the syzygy of C (C is called the cosyzygy of A , resp.) and is denoted by $\Omega(C)$ ($\Sigma(A)$, resp.).

- \mathcal{C} is said to be *Frobenius* if \mathcal{C} has enough projectives and injectives and if moreover the projectives coincide with the injectives. In this case one has the quotient category $\underline{\mathcal{C}}$ of \mathcal{C} by projectives, which is a triangulated category by [15]. We refer to this as the stable category of \mathcal{C} .

Remark 2.13. (1) If $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ is an exact category, then the definitions of enough projectives and enough injectives agree with the usual definitions.

(2) If $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ is a triangulated category, then $Proj(\mathcal{C})$ and $Inj(\mathcal{C})$ consist of zero objects. Moreover $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$ is Frobenius as an extriangulated category.

Suppose \mathcal{C} is an extriangulated category with enough projectives and injectives. For a subcategory $\mathcal{B} \subseteq \mathcal{C}$, put $\Omega^0 \mathcal{B} = \mathcal{B}$, and for $i > 0$ we define $\Omega^i \mathcal{B}$ inductively to be the subcategory consisting of syzygies of objects in Ω^{i-1} , i.e.

$$\Omega^i \mathcal{B} = \Omega(\Omega^{i-1} \mathcal{B}).$$

We call $\Omega^i \mathcal{B}$ the i -th syzygy of \mathcal{B} . Dually we define the i -th cosyzygy $\Sigma^i \mathcal{B}$ by $\Sigma^0 \mathcal{B} = \mathcal{B}$ and $\Sigma^i \mathcal{B} = \Sigma(\Sigma^{i-1} \mathcal{B})$ for $i > 0$.

In [12] the authors defined higher extension groups in an extriangulated category having enough projectives and injectives as $\mathbb{E}^{i+1}(X, Y) \cong \mathbb{E}(X, \Sigma^i Y) \cong \mathbb{E}(\Omega^i X, Y)$ for $i \geq 0$, and they showed the following result:

Lemma 2.14. *Let $A \xrightarrow{x} B \xrightarrow{y} C \xrightarrow{\delta} \triangleright$ be an \mathbb{E} -triangle. For any object $X \in \mathcal{B}$, there are long exact sequences*

$$\cdots \longrightarrow \mathbb{E}^i(X, A) \xrightarrow{x^*} \mathbb{E}^i(X, B) \xrightarrow{y^*} \mathbb{E}^i(X, C) \longrightarrow \mathbb{E}^{i+1}(X, A) \xrightarrow{x^*} \mathbb{E}^{i+1}(X, B) \xrightarrow{y^*} \cdots (i \geq 1),$$

$$\cdots \longrightarrow \mathbb{E}^i(C, X) \xrightarrow{y^*} \mathbb{E}^i(B, X) \xrightarrow{x^*} \mathbb{E}^i(A, X) \longrightarrow \mathbb{E}^{i+1}(C, X) \xrightarrow{y^*} \mathbb{E}^{i+1}(B, X) \xrightarrow{x^*} \cdots (i \geq 1).$$

An \mathbb{E} -triangle sequence in \mathcal{C} is displayed as a sequence

$$\cdots \longrightarrow X_{n+1} \xrightarrow{d_{n+1}} X_n \xrightarrow{d_n} X_{n-1} \longrightarrow \cdots$$

over \mathcal{C} such that for any n , there are \mathbb{E} -triangles $K_{n+1} \xrightarrow{g_n} X_n \xrightarrow{f_n} K_n \xrightarrow{\delta^n} \triangleright$ and the differential $d_n = g_{n-1} f_n$.

From now on to the end of the paper, we always suppose that extriangulated category \mathcal{C} has enough projectives and injectives.

3. Relative Homological Dimensions

Let \mathcal{X} be a subcategory of \mathcal{C} . The symbol $\widehat{\mathcal{X}}_n$ ($\widetilde{\mathcal{X}}_n$, resp.) denotes the subcategory of objects $A \in \mathcal{C}$ such that there exists an \mathbb{E} -triangle sequence $X_n \rightarrow X_{n-1} \rightarrow \cdots \rightarrow X_0 \rightarrow C$ ($C \rightarrow X_0 \rightarrow \cdots \rightarrow X_{n-1} \rightarrow X_n$, resp.) with each $X_i \in \mathcal{X}$. We denote by $\widehat{\mathcal{X}}$ ($\widetilde{\mathcal{X}}$, resp.) the union of all $\widehat{\mathcal{X}}_n$ ($\widetilde{\mathcal{X}}_n$, resp.) for some nonnegative n . That is to say $\widehat{\mathcal{X}} = \bigcup_{n=0}^{\infty} \widehat{\mathcal{X}}_n$, $\widetilde{\mathcal{X}} = \bigcup_{n=0}^{\infty} \widetilde{\mathcal{X}}_n$.

Definition 3.1. (1) For any $C \in \mathcal{C}$, the \mathcal{X} -resolution dimension of C is

$$\text{resdim}_{\mathcal{X}}(C) := \min\{n \in \mathbb{N} : C \in \widehat{\mathcal{X}}_n\}.$$

If $C \notin \widehat{\mathcal{X}}_n$ for any $n \in \mathbb{N}$, then $\text{resdim}_{\mathcal{X}}(C) = \infty$. Dually, we also have the \mathcal{X} -coresolution dimension of C denoted by $\text{coresdim}_{\mathcal{X}}(C)$.

(2) For any subcategory \mathcal{Y} of \mathcal{C} , we set

$$\text{resdim}_{\mathcal{X}}(\mathcal{Y}) := \sup\{\text{resdim}_{\mathcal{X}}(M) : M \in \mathcal{Y}\}.$$

Dually, we also have $\text{coresdim}_{\mathcal{X}}(\mathcal{Y})$.

(3) The \mathcal{X} -projective dimension of C is

$$\text{pd}_{\mathcal{X}}(C) := \min\{n \in \mathbb{N} : \mathbb{E}^i(C, -)|_{\mathcal{X}} = 0, \forall i > n\}.$$

(4) The \mathcal{X} -injective dimension of \mathcal{C} is

$$\text{id}_{\mathcal{X}}(C) := \min\{n \in \mathbb{N} : \mathbb{E}^i(-, C)|_{\mathcal{X}} = 0, \forall i > n\}.$$

(5) For any subcategory \mathcal{Y} of \mathcal{C} , we set

$$\text{pd}_{\mathcal{X}}(\mathcal{Y}) := \sup\{\text{pd}_{\mathcal{X}}(M) : M \in \mathcal{Y}\} \quad \text{and} \quad \text{id}_{\mathcal{X}}(\mathcal{Y}) := \sup\{\text{id}_{\mathcal{X}}(M) : M \in \mathcal{Y}\}.$$

Lemma 3.2. *Let \mathcal{X} and \mathcal{Y} be subcategories of \mathcal{C} . Then $\text{pd}_{\mathcal{X}}(\mathcal{Y}) = \text{id}_{\mathcal{Y}}(\mathcal{X})$. Furthermore, for any \mathbb{E} -triangle $A \xrightarrow{x} B \xrightarrow{y} C \xrightarrow{\delta} \succ$ in \mathcal{C} , we have*

- (1) $\text{id}_{\mathcal{X}}(B) \leq \max\{\text{id}_{\mathcal{X}}(A), \text{id}_{\mathcal{X}}(C)\}$;
- (2) $\text{id}_{\mathcal{X}}(A) \leq \max\{\text{id}_{\mathcal{X}}(B), \text{id}_{\mathcal{X}}(C) + 1\}$;
- (3) $\text{id}_{\mathcal{X}}(C) \leq \max\{\text{id}_{\mathcal{X}}(B), \text{id}_{\mathcal{X}}(A) - 1\}$.

Proof. It is straightforward. \square

For a subcategory \mathcal{X} of \mathcal{C} , define $\mathcal{X}^{\perp} = \{Y \in \mathcal{C} | \mathbb{E}^i(X, Y) = 0, \forall i \geq 1, X \in \mathcal{X}\}$. Similarly, we can define ${}^{\perp}\mathcal{X}$. Now we give a relationship between the relative projective dimension and the resolution dimension.

Theorem 3.3. *Let \mathcal{X} and \mathcal{Y} be subcategories of \mathcal{C} . Then, the following statements hold.*

- (1) $\text{pd}_{\mathcal{X}}(L) \leq \text{pd}_{\mathcal{X}}(\mathcal{Y}) + \text{resdim}_{\mathcal{Y}}(L)$, $\forall L \in \mathcal{C}$.
- (2) If $\mathcal{Y} \subseteq \mathcal{X} \cap {}^{\perp}\mathcal{X}$, then $\text{pd}_{\mathcal{X}}(L) = \text{resdim}_{\mathcal{Y}}(L)$, $\forall L \in \hat{\mathcal{Y}}$.

Proof. (1) Let $d := \text{resdim}_{\mathcal{Y}}(L)$ and $\alpha := \text{pd}_{\mathcal{X}}(\mathcal{Y})$. We may assume that d and α are finite. We prove (1) by induction on d . If $d = 0$, it follows that $L \in \mathcal{Y}$, hence (1) holds in this case.

Assume $d \geq 1$. So we have an \mathbb{E} -triangle

$$K \xrightarrow{x} Y \xrightarrow{y} L \xrightarrow{\delta} \succ$$

in \mathcal{C} with $Y \in \mathcal{Y}$ and $\text{resdim}_{\mathcal{Y}}(K) = d - 1$. Applying $\text{Hom}_{\mathcal{C}}(-, M)$, with $M \in \mathcal{X}$, to the \mathbb{E} -triangle δ , we get an exact sequence $\mathbb{E}^{i-1}(K, M) \rightarrow \mathbb{E}^i(L, M) \rightarrow \mathbb{E}^i(Y, M)$. By induction, we know that $\text{pd}_{\mathcal{X}}(K) \leq \alpha + d - 1$. Therefore $\mathbb{E}^i(L, M) = 0$ for $i > \alpha + d$, and so $\text{pd}_{\mathcal{X}}(L) \leq \alpha + d$.

(2) Let $\mathcal{Y} \subseteq \mathcal{X} \cap {}^{\perp}\mathcal{X}$. Consider $L \in \hat{\mathcal{Y}}$ and let $d = \text{resdim}_{\mathcal{Y}}(L)$. Since $\text{pd}_{\mathcal{X}}(\mathcal{Y}) = 0$, it follows that $\text{pd}_{\mathcal{X}}(L) \leq \text{resdim}_{\mathcal{Y}}(L) = d$ by (1). We prove, by induction on d , that the equality given in (2) holds. For $d = 0$, it is obvious.

Suppose that $d = 1$. Then we have an \mathbb{E} -triangle

$$Y_1 \xrightarrow{x} Y_0 \xrightarrow{y} L \xrightarrow{\delta} \succ$$

in \mathcal{C} with $Y_i \in \mathcal{Y}, i = 1, 2$. If $\text{pd}_{\mathcal{X}}(L) = 0$, then $L \in {}^{\perp}\mathcal{X}$. Since $\mathcal{Y} \subseteq \mathcal{X}$, $\mathbb{E}(L, Y_1) = 0$, therefore the \mathbb{E} -triangle δ splits giving us that $L \in \mathcal{Y}$, which is a contradiction as $d = 1$. So $\text{pd}_{\mathcal{X}}(L) > 0$ proving (2) for $d = 1$.

Assume now that $d \geq 2$. Thus we have an \mathbb{E} -triangle

$$K \xrightarrow{d} Y \xrightarrow{e} L \xrightarrow{\theta} \succ$$

in \mathcal{C} with $Y \in \mathcal{Y}$, $\text{resdim}_{\mathcal{Y}}(K) = d - 1$. Hence $\text{pd}_{\mathcal{X}}(K) = d - 1$ by inductive hypothesis. For any $X \in \mathcal{X}$, there is an exact sequence

$$\mathbb{E}^{d-1}(Y, X) \rightarrow \mathbb{E}^{d-1}(K, X) \rightarrow \mathbb{E}^d(L, X).$$

If $\text{pd}_{\mathcal{X}}(L) \leq d - 1$, then $\mathbb{E}^{d-1}(K, X) = 0$ contradicting that $\text{pd}_{\mathcal{X}}(K) = d - 1$. This means that $\text{pd}_{\mathcal{X}}(L) > d - 1$; proving (2). \square

Now, we begin to focus our attention on pairs $(\mathcal{X}, \mathcal{W})$ of subcategories of \mathcal{C} and study the relationship between \mathcal{X} -injective cogenerators for \mathcal{X} and $\hat{\mathcal{X}}$.

Definition 3.4. Let \mathcal{X} and \mathcal{W} be two subcategories of \mathcal{C} . We say that

- (1) \mathcal{W} is a cogenerator for \mathcal{X} , if $\mathcal{W} \subseteq \mathcal{X}$ and for each object $X \in \mathcal{X}$, there exists an \mathbb{E} -triangle $X \xrightarrow{x} W \xrightarrow{y} X' \xrightarrow{\delta} \triangleright$. The term generator is defined dually.
- (2) \mathcal{W} is \mathcal{X} -injective if $\text{id}_{\mathcal{X}}(\mathcal{W})=0$. The term \mathcal{X} -projective is defined dually.
- (3) \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} if \mathcal{W} is a cogenerator for \mathcal{X} and $\text{id}_{\mathcal{X}}(\mathcal{W})=0$. The term \mathcal{X} -projective generator for \mathcal{X} is defined dually.

In the following, let \mathcal{X} and \mathcal{W} be two subcategories of \mathcal{C} such that $\mathcal{W} \subseteq \mathcal{X}$.

Lemma 3.5. Suppose that \mathcal{X} is an extension closed subcategory of \mathcal{C} . Consider two \mathbb{E} -triangles

$$N \xrightarrow{a} X_1 \xrightarrow{b} D \xrightarrow{\delta} \triangleright \quad \text{and} \quad D \xrightarrow{c} X_0 \xrightarrow{d} M \xrightarrow{\theta} \triangleright$$

in \mathcal{C} with $X_0, X_1 \in \mathcal{X}$. If \mathcal{W} is a cogenerator for \mathcal{X} , then there exist two \mathbb{E} -triangles

$$N \longrightarrow W_1 \longrightarrow D' \dashrightarrow \triangleright \quad \text{and} \quad D' \longrightarrow X'_0 \longrightarrow M \dashrightarrow \triangleright$$

with $W_1 \in \mathcal{W}$ and $X'_0 \in \mathcal{X}$.

Proof. We have an \mathbb{E} -triangle $X_1 \longrightarrow W_1 \longrightarrow X'_1 \dashrightarrow \triangleright$ with $W_1 \in \mathcal{W}$ and $X'_1 \in \mathcal{X}$ as \mathcal{W} is a cogenerator for \mathcal{X} . By (ET4), we obtain a commutative diagram in \mathcal{C}

$$\begin{array}{ccccc} N & \longrightarrow & X_1 & \longrightarrow & D \\ \parallel & & \downarrow & & \downarrow \\ N & \longrightarrow & W_1 & \longrightarrow & D' \\ & & \downarrow & & \downarrow \\ & & X'_1 & = & X'_1. \end{array}$$

By Lemma 2.10 (2), we have the following commutative diagram in \mathcal{C}

$$\begin{array}{ccccc} D & \longrightarrow & X_0 & \longrightarrow & M \\ \downarrow & & \downarrow & & \parallel \\ D' & \longrightarrow & X'_0 & \longrightarrow & M \\ \downarrow & & \downarrow & & \\ X'_1 & = & X'_1. & & \end{array}$$

Since \mathcal{X} is closed under extensions, it follows that $X'_0 \in \mathcal{X}$. The second rows in the above two diagrams are desired \mathbb{E} -triangles. \square

Lemma 3.6. Suppose \mathcal{X} is closed under extensions and \mathcal{W} is a cogenerator for \mathcal{X} . Then for any $X \in \mathcal{X}$ and nonnegative integer n , $C \in \widehat{\mathcal{X}}_n$ if and only if there exists an \mathbb{E} -triangle

$$W_n \rightarrow \cdots \rightarrow W_2 \rightarrow W_1 \rightarrow X_0 \rightarrow C$$

with $X_0 \in \mathcal{X}$ and $W_i \in \mathcal{W}$ for $1 \leq i \leq n$.

Proof. The “if” part is trivial. We prove the “only if” part by induction on n . If $n = 1$, then there exists an \mathbb{E} -triangles

$$X_1 \longrightarrow X_0 \longrightarrow C \dashrightarrow \quad \text{and} \quad 0 \longrightarrow X_1 \longrightarrow X_1 \dashrightarrow \quad \text{with } X_1, X_0 \in \mathcal{X}.$$

By Lemma 3.5, we obtain an \mathbb{E} -triangle $W_1 \longrightarrow X'_0 \longrightarrow C \dashrightarrow$ that is desired.

Suppose now $n \geq 2$. Then we have an \mathbb{E} -triangle

$$K \longrightarrow X'_0 \longrightarrow C \dashrightarrow$$

with $X_0 \in \mathcal{X}$ and $K \in \widehat{\mathcal{X}_{n-1}}$. By the induction, there exists an \mathbb{E} -triangle sequence

$$W_n \rightarrow \cdots \rightarrow W_2 \xrightarrow{f} X_1 \rightarrow K$$

with $X_1 \in \mathcal{X}$ and $W_i \in \mathcal{W}$ for $2 \leq i \leq n$. Hence we have an \mathbb{E} -triangle sequence

$$W_{n-1} \rightarrow \cdots \rightarrow W_2 \xrightarrow{f_1} K'$$

and a \mathbb{E} -triangle

$$K' \xrightarrow{f_2} X_1 \longrightarrow K \dashrightarrow$$

with $f_2 f_1 = f$. Applying Lemma 3.5(1) to \mathbb{E} -triangles

$$K' \xrightarrow{f_2} X_1 \longrightarrow K \dashrightarrow \quad \text{and} \quad K \longrightarrow X'_0 \longrightarrow C \dashrightarrow,$$

we obtain two \mathbb{E} -triangles

$$K' \longrightarrow W_1 \longrightarrow K'' \dashrightarrow \quad \text{and} \quad K'' \longrightarrow X_0 \longrightarrow C \dashrightarrow.$$

Therefore, we have a \mathbb{E} -triangle sequence

$$W_n \rightarrow \cdots \rightarrow W_2 \rightarrow W_1 \rightarrow X_0 \rightarrow C$$

with $X_0 \in \mathcal{X}$ and $W_i \in \mathcal{W}$ for $1 \leq i \leq n$. □

The following theorem shows that any object in $\widehat{\mathcal{X}}$ admits two \mathbb{E} -triangles: one giving rise to an \mathcal{X} -precover and the other to a $\widehat{\mathcal{W}}$ -preenvelope, which generalizes [13, Theorem 5.4].

Theorem 3.7. *Suppose \mathcal{X} is closed under extensions and \mathcal{W} is a cogenerator for \mathcal{X} . Consider the following conditions:*

- (1) C is in $\widehat{\mathcal{X}_n}$.
- (2) There exists an \mathbb{E} -triangle

$$Y_C \longrightarrow X_C \xrightarrow{\varphi_C} C \dashrightarrow$$

with $X_C \in \mathcal{X}$ and $Y_C \in \widehat{\mathcal{W}_{n-1}}$.

- (3) There exists an \mathbb{E} -triangle

$$C \xrightarrow{\psi^C} Y^C \longrightarrow X^C \dashrightarrow$$

with $X^C \in \mathcal{X}$ and $Y^C \in \widehat{\mathcal{W}_n}$.

Then, (1) \Leftrightarrow (2) \Rightarrow (3). If \mathcal{X} is also closed under CoCones , then (3) \Rightarrow (2), and hence all three conditions are equivalent. If \mathcal{W} is \mathcal{X} -injective, then φ_C is an \mathcal{X} -precover of C and ψ^C is a $\widehat{\mathcal{W}}$ -preenvelope of C .

Proof. (1) \Leftrightarrow (2) follows from Lemma 3.6.

(2) \Rightarrow (3) Since $X_C \in \mathcal{X}$ and \mathcal{W} is a cogenerator for \mathcal{X} , we have an \mathbb{E} -triangle

$$X_C \xrightarrow{x} W \xrightarrow{y} X' \dashrightarrow$$

with $W \in \mathcal{W}$ and $X' \in \mathcal{X}$. By (ET4), we obtain a commutative diagram

$$\begin{array}{ccccc} Y_C & \longrightarrow & X_C & \longrightarrow & C \\ \parallel & & \downarrow & & \downarrow \\ Y_C & \longrightarrow & W & \longrightarrow & Y^C \\ & & \downarrow & & \downarrow \\ & & X' & \xlongequal{\quad} & X'. \end{array}$$

From the second row, it follows that $Y^C \in \widehat{\mathcal{W}}_n$. Hence the third column is the desired one.

Suppose \mathcal{X} is also closed under CoCones. Since $Y^C \in \widehat{\mathcal{W}}_n$, we have an \mathbb{E} -triangle

$$Y_C \xrightarrow{x} W \xrightarrow{y} Y^C \dashrightarrow$$

in \mathcal{C} with $W \in \mathcal{W}$ and $Y_C \in \widehat{\mathcal{W}}_{n-1}$. By (ET4)^{op}, we obtain a commutative diagram

$$\begin{array}{ccccc} Y_C & \longrightarrow & X_C & \longrightarrow & C \\ \parallel & & \downarrow & & \downarrow \\ Y_C & \longrightarrow & W & \longrightarrow & Y^C \\ & & \downarrow & & \downarrow \\ & & X^C & \xlongequal{\quad} & X^C. \end{array}$$

Since \mathcal{X} is closed under CoCones, it follows that $X_C \in \mathcal{X}$. Hence the first row is the desired \mathbb{E} -triangle.

Assume \mathcal{W} is \mathcal{X} -injective. Applying $\mathcal{C}(X, -)$ to the \mathbb{E} -triangle

$$Y_C \longrightarrow X_C \xrightarrow{\varphi_C} C \dashrightarrow^{\delta},$$

we have an exact sequence

$$\mathcal{C}(X, X_C) \xrightarrow{\mathcal{C}(X, \varphi_C)} \mathcal{C}(X, C) \longrightarrow \mathbb{E}(X, Y_C).$$

Since $\mathbb{E}(X, Y_C) = 0$, it follows that $\text{Hom}_{\mathcal{C}}(X, \varphi_C)$ is an epimorphism, hence φ_C is an \mathcal{X} -precover of C . Similarly, we can prove that ψ^C is a $\widehat{\mathcal{W}}$ -preenvelope of C . \square

Lemma 3.8. *Let \mathcal{X} and \mathcal{Y} be subcategories of \mathcal{C} . Then $\text{id}_{\mathcal{Y}}(\widehat{\mathcal{X}}) = \text{id}_{\mathcal{Y}}(\mathcal{X})$.*

Proof. Since $\mathcal{X} \subseteq \widehat{\mathcal{X}}$, it follows that $\text{id}_{\mathcal{Y}}(\mathcal{X}) \leq \text{id}_{\mathcal{Y}}(\widehat{\mathcal{X}})$. It is enough to show that $\text{id}_{\mathcal{Y}}(\widehat{\mathcal{X}}) \leq \text{id}_{\mathcal{Y}}(\mathcal{X})$. We may assume that $\alpha := \text{id}_{\mathcal{Y}}(\mathcal{X})$ is finite. Let $C \in \widehat{\mathcal{X}}$. We prove that, by induction on $n = \text{resdim}_{\mathcal{X}}(C)$. If $n = 0$, then $C \in \mathcal{X}$, there is nothing to prove.

Let $n \geq 1$. Then we have an \mathbb{E} -triangle

$$K \longrightarrow X \longrightarrow C \dashrightarrow^{\delta}$$

with $X \in \mathcal{X}, K \in \widehat{\mathcal{X}_{n-1}}$, $\text{id}_{\mathcal{Y}}(K) \leq \alpha$ by inductive hypothesis. For any $Y \in \mathcal{Y}$, applying $\text{Hom}_{\mathcal{C}}(X, -)$ to the \mathbb{E} -triangle δ , we obtain an exact sequence

$$\mathbb{E}^i(Y, K) \rightarrow \mathbb{E}^i(Y, X) \rightarrow \mathbb{E}^i(Y, C) \rightarrow \mathbb{E}^{i+1}(Y, K).$$

Therefore $\mathbb{E}^i(Y, C) = 0$ for $i > \alpha$ as $\text{id}_{\mathcal{Y}}(K) \leq \alpha$. So we get $\text{id}_{\mathcal{Y}}(\widehat{\mathcal{X}}) \leq \text{id}_{\mathcal{Y}}(\mathcal{X})$. \square

The following result shows that there is a unique \mathcal{X} -injective cogenerator for \mathcal{X} (in case it exists).

Proposition 3.9. *Let \mathcal{X} and \mathcal{W} be two subcategories of \mathcal{C} such that \mathcal{W} is \mathcal{X} -injective. Then the following statements hold.*

- (1) $\widehat{\mathcal{W}}$ is \mathcal{X} -injective.
- (2) If \mathcal{W} is a cogenerator for \mathcal{X} , then $\mathcal{W} = \mathcal{X} \cap \mathcal{X}^\perp = \mathcal{X} \cap \widehat{\mathcal{W}}$.
- (3) If \mathcal{W} is a cogenerator for \mathcal{X} , then $\widehat{\mathcal{W}} = \widehat{\mathcal{X}} \cap \mathcal{X}^\perp$.

Proof. (1) It follows from the Lemma 3.8.

- (2) Let $X \in \mathcal{X} \cap \mathcal{X}^\perp$. We have an \mathbb{E} -triangle

$$X \longrightarrow W \xrightarrow{y} X' \xrightarrow{\theta} \gg$$

with $X' \in \mathcal{X}$ and $W \in \mathcal{W}$. Moreover $X \in \mathcal{X}^\perp$ implies that the \mathbb{E} -triangle θ splits and so $X \in \mathcal{W}$. On the other hand, it is easy to see $\mathcal{W} \subseteq \mathcal{X} \cap \widehat{\mathcal{W}}$. Since $\text{id}_{\mathcal{X}}(\widehat{\mathcal{W}}) = 0$, it follows that $\mathcal{X} \cap \widehat{\mathcal{W}} \subseteq \mathcal{X} \cap \mathcal{X}^\perp$. Hence $\mathcal{W} = \mathcal{X} \cap \widehat{\mathcal{W}}$.

- (3) Let $C \in \widehat{\mathcal{X}} \cap \mathcal{X}^\perp$. We have an \mathbb{E} -triangle

$$Y_C \longrightarrow X_C \longrightarrow C \xrightarrow{\theta} \gg$$

with $Y_C \in \widehat{\mathcal{W}}$ and $X_C \in \mathcal{X}$ by Theorem 3.7. Since Y_C and C are in \mathcal{X}^\perp , it follows that $X_C \in \mathcal{X} \cap \mathcal{X}^\perp$. Hence $X_C \in \mathcal{W}$ by (2), implying $\widehat{\mathcal{X}} \cap \mathcal{X}^\perp \subseteq \widehat{\mathcal{W}}$. On the other hand, it is obvious that $\widehat{\mathcal{W}} \subseteq \widehat{\mathcal{X}} \cap \mathcal{X}^\perp$. Therefore $\widehat{\mathcal{W}} = \widehat{\mathcal{X}} \cap \mathcal{X}^\perp$. \square

Proposition 3.10. *Let \mathcal{X} be closed under extensions such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . Then $\widehat{\mathcal{X}}$ is closed under extensions, and hence an extriangulated category.*

Proof. Suppose $A \xrightarrow{d} B \xrightarrow{e} C \xrightarrow{\delta} \gg$ is an \mathbb{E} -triangle with A and C in $\widehat{\mathcal{X}}$. Proceed by induction on $n := \text{resdim}_{\mathcal{X}}(C)$. Assume $n = 0$, which means that C is in \mathcal{X} . As A is in $\widehat{\mathcal{X}}$, there exists an \mathbb{E} -triangle

$$Y_A \longrightarrow X_A \xrightarrow{p} A \xrightarrow{\xi} \gg$$

with $X_A \in \mathcal{X}$ and $Y_A \in \widehat{\mathcal{W}_{n-1}}$ by Theorem 3.7. Since $\widehat{\mathcal{W}} \subseteq \mathcal{X}^\perp$, it follows that $\mathbb{E}(C, p) : \mathbb{E}(C, X_A) \rightarrow \mathbb{E}(C, A)$ is an isomorphism. Hence we obtain a commutative diagram

$$\begin{array}{ccccc} X_A & \xrightarrow{x} & Z & \xrightarrow{y} & C \xrightarrow{\theta} \gg \\ \downarrow p & & \downarrow g & & \parallel \\ A & \xrightarrow{d} & B & \xrightarrow{e} & C \xrightarrow{\delta} \gg \end{array}$$

with $p_*\theta = \delta$. By (ET4), we have the following commutative diagram

$$\begin{array}{ccccc}
 Y_A & \longrightarrow & X_A & \xrightarrow{p} & A \\
 \parallel & & \downarrow x & & \downarrow d \\
 Y_A & \longrightarrow & Z & \longrightarrow & B \\
 & & \downarrow y & & \downarrow e \\
 & & C & \xlongequal{\quad} & C.
 \end{array}$$

Since X_A and C are in \mathcal{X} , the object $Z \in \mathcal{X}$ as \mathcal{X} is closed under extensions. Note that $Y_A \in \widehat{\mathcal{W}}$, hence in $\widehat{\mathcal{X}}$, it follows that $B \in \widehat{\mathcal{X}}$.

Assume $n > 0$ and let $L \longrightarrow X_0 \xrightarrow{y} C \dashrightarrow$ be an \mathbb{E} -triangle with $\text{resdim}_{\mathcal{X}}(L) = n - 1$. By Lemma 2.10, we have a commutative diagram

$$\begin{array}{ccccc}
 & & L & \xlongequal{\quad} & L \\
 & & \downarrow & & \downarrow \\
 A & \longrightarrow & V' & \longrightarrow & X_0 \\
 \parallel & & \downarrow & & \downarrow \\
 A & \longrightarrow & B & \longrightarrow & C.
 \end{array}$$

Since $X_0 \in \mathcal{X}$, it follows that $\mathbb{E}(X_0, p) : \mathbb{E}(X_0, X_A) \rightarrow \mathbb{E}(X_0, A)$ is an isomorphism, using (ET4), we obtain a commutative diagram

$$\begin{array}{ccccc}
 Y_A & \longrightarrow & X_A & \xrightarrow{p} & A \\
 \parallel & & \downarrow & & \downarrow \\
 Y_A & \longrightarrow & V & \longrightarrow & V' \\
 & & \downarrow & & \downarrow \\
 & & X_0 & \xlongequal{\quad} & X_0.
 \end{array}$$

Using (ET4)^{op} again, we also have the following commutative diagram

$$\begin{array}{ccccc}
 Y_A & \longrightarrow & U & \xrightarrow{p} & L \\
 \parallel & & \downarrow x & & \downarrow d \\
 Y_A & \longrightarrow & V & \longrightarrow & V' \\
 & & \downarrow & & \downarrow \\
 & & B & \xlongequal{\quad} & B.
 \end{array}$$

By inductive hypothesis, U is in $\widehat{\mathcal{X}}$. Since X_A and X_0 are in \mathcal{X} , $V \in \mathcal{X}$, as \mathcal{X} is closed under extensions. It follows that $B \in \widehat{\mathcal{X}}$. Hence $\widehat{\mathcal{X}}$ is also an extriangulated category by [15, Remark 2.18]. \square

Lemma 3.11. *Let \mathcal{X} be closed under extensions and CoCones such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . Given an \mathbb{E} -triangle $K \xrightarrow{x} X \xrightarrow{y} C \xrightarrow{\delta} \gg$ with $X \in \mathcal{X}$. Then $C \in \widehat{\mathcal{X}}$ if and only if $K \in \widehat{\mathcal{X}}$.*

Proof. By definition, if $K \in \widehat{\mathcal{X}}$, so is $C \in \widehat{\mathcal{X}}$. Now assume that $C \in \widehat{\mathcal{X}}$. By Theorem 3.7, we obtain a commutative diagram

$$\begin{array}{ccccc} K & \xrightarrow{x} & X & \xrightarrow{y} & C \xrightarrow{\delta} \gg \\ \downarrow f & & \downarrow g & & \parallel \\ Y_C & \xrightarrow{l} & X_C & \xrightarrow{p} & C \xrightarrow{f_*\delta} \gg \end{array}$$

with $X_C \in \mathcal{X}$ and $Y_C \in \widehat{\mathcal{W}}$. Since $Y_C \in \widehat{\mathcal{W}}$, there exists an \mathbb{E} -triangle

$$L \xrightarrow{m} W \xrightarrow{e} Y_C \xrightarrow{\delta'} \gg$$

with $W \in \mathcal{W}$. We also have a commutative diagram

$$\begin{array}{ccccc} L & \xrightarrow{u} & V & \xrightarrow{v} & K \xrightarrow{f_*\delta'} \gg \\ \parallel & & \downarrow h & & \downarrow f \\ L & \xrightarrow{m} & W & \xrightarrow{e} & Y_C \xrightarrow{\delta'} \gg. \end{array}$$

By the dual of Lemma 2.12, there exists an \mathbb{E} -triangle $V \xrightarrow{\begin{bmatrix} -v \\ h \end{bmatrix}} K \oplus W \xrightarrow{\begin{bmatrix} f & e \end{bmatrix}} Y_C \xrightarrow{u_*\delta'} \gg$. By (ET4), we obtain a commutative diagram

$$\begin{array}{ccccc} V & \xrightarrow{\begin{bmatrix} -v \\ h \end{bmatrix}} & K \oplus W & \xrightarrow{\begin{bmatrix} f & e \end{bmatrix}} & Y_C \xrightarrow{u_*\delta'} \gg \\ \parallel & & \downarrow \begin{bmatrix} x & 0 \\ 0 & 1 \end{bmatrix} & & \downarrow \\ V & \longrightarrow & X \oplus W & \longrightarrow & H \xrightarrow{\quad} \gg \\ & & \downarrow \begin{bmatrix} y & 0 \end{bmatrix} & & \downarrow \\ & & C & \xlongequal{\quad} & C \\ & & \downarrow \delta \oplus 0 & & \downarrow \begin{bmatrix} f & e \end{bmatrix}_*(\delta \oplus 0) \\ & & \Downarrow & & \Downarrow \end{array} \quad (*)$$

Since $\begin{bmatrix} f & e \end{bmatrix}_*(\delta \oplus 0) = f_*\delta$, we obtain a commutative diagram

$$\begin{array}{ccccc} Y_C & \longrightarrow & H & \longrightarrow & C \xrightarrow{f_*\delta} \gg \\ \parallel & & \downarrow & & \parallel \\ Y_C & \xrightarrow{l} & X_0 & \xrightarrow{p} & C \xrightarrow{f_*\delta} \gg. \end{array}$$

It follows that H is isomorphic to X_C . Hence we can replace H by X_C in $(*)$ and have the commutative diagram

$$\begin{array}{ccccc}
 V & \xrightarrow{\begin{bmatrix} -v \\ h \end{bmatrix}} & K \oplus W & \xrightarrow{\begin{bmatrix} f & e \end{bmatrix}} & Y_C & \xrightarrow{u_*\delta'} & \rightarrow \\
 \parallel & & \downarrow \begin{bmatrix} x & 0 \\ 0 & 1 \end{bmatrix} & & \downarrow l & & \\
 V & \longrightarrow & X \oplus W & \longrightarrow & X_C & \dashrightarrow & \rightarrow \\
 & & \downarrow \begin{bmatrix} y & 0 \end{bmatrix} & & \downarrow p & & \\
 & & C & \xlongequal{\quad} & C & & \\
 & & \downarrow \delta \oplus 0 & & \downarrow f_*\delta & & \\
 & & Y & & Y & &
 \end{array} .$$

Since W and X are both in \mathcal{X} , it follows that V is in \mathcal{X} as \mathcal{X} is closed under Cocones. By Proposition 3.10, $K \oplus W \in \hat{\mathcal{X}}$. So there exists an \mathbb{E} -triangle

$$Y_{K \oplus W} \longrightarrow X_{K \oplus W} \longrightarrow K \oplus W \dashrightarrow$$

with $Y_{K \oplus W} \in \widehat{\mathcal{W}}$, $X_{K \oplus W} \in \mathcal{X}$ by Theorem 3.7. By $(ET4)^{op}$, there is a commutative diagram

$$\begin{array}{ccccc}
 Y_{K \oplus W} & \longrightarrow & Z & \longrightarrow & W \\
 \parallel & & \downarrow & & \downarrow \\
 Y_{K \oplus W} & \longrightarrow & X_{K \oplus W} & \longrightarrow & K \oplus W \\
 & & \downarrow & & \downarrow \\
 & & K & \xlongequal{\quad} & K
 \end{array}$$

implying $Z \in \hat{\mathcal{X}}$ as $Y_{K \oplus W}$ and W are in $\hat{\mathcal{X}}$. Hence $K \in \hat{\mathcal{X}}$. This completes the proof of the lemma. \square

Proposition 3.12. *Let \mathcal{X} be closed under extensions such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . Then $\text{pd}_{\widehat{\mathcal{W}}}(C) = \text{pd}_{\mathcal{W}}(C) = \text{resdim}_{\mathcal{X}}(C)$ for any $C \in \hat{\mathcal{X}}$.*

Proof. Let $C \in \hat{\mathcal{X}}$, $\text{pd}_{\mathcal{W}}(C) = \text{id}_{\{C\}}(\mathcal{W}) = \text{id}_{\{C\}}(\widehat{\mathcal{W}})$ by Lemma 3.2 and 3.8. To prove the last equality, we proceed by induction on $n = \text{resdim}_{\mathcal{X}}(C)$. If $n = 0$, then $C \in \mathcal{X}$ and $\text{pd}_{\mathcal{W}}(C) = \text{resdim}_{\mathcal{X}}(C) = 0$

Let $n = 1$. Then we have an \mathbb{E} -triangle

$$Y_C \longrightarrow X_C \longrightarrow C \xrightarrow{\delta} \rightarrow$$

with $X_C \in \mathcal{X}$ and $Y_C \in \mathcal{W}$ by Theorem 3.7. We claim that $\text{pd}_{\mathcal{W}}(C) > 0$. Indeed, suppose $\text{pd}_{\mathcal{W}}(C) = 0$, it follows that the \mathbb{E} -triangle δ splits. Hence $C \in \mathcal{X}$ contradicting that $\text{resdim}_{\mathcal{X}}(C) = 1$. Hence $\text{pd}_{\mathcal{W}}(C) = 1$ by Theorem 3.3.

Let $n \geq 2$. From Theorem 3.3, we have $\text{pd}_{\mathcal{W}}(C) \leq \text{resdim}_{\mathcal{X}}(C) = n$ as $\text{pd}_{\mathcal{W}}(\mathcal{X}) = 0$. Then it is enough to prove that $\mathbb{E}^n(C, W) \neq 0$ for some $W \in \mathcal{W}$. Consider an \mathbb{E} -triangle

$$K \longrightarrow X_0 \longrightarrow C \xrightarrow{\xi} \rightarrow$$

with $X_0 \in \mathcal{X}$ and $\text{resdim}_{\mathcal{X}}(K) = n - 1$. By inductive hypothesis $\text{pd}_{\mathcal{W}}(K) = n - 1$. Applying $\text{Hom}_{\mathcal{C}}(-, W)$ to the \mathbb{E} -triangle ξ with $W \in \mathcal{W}$, we have an exact sequence

$$\mathbb{E}^{n-1}(X_0, W) \rightarrow \mathbb{E}^{n-1}(K, W) \rightarrow \mathbb{E}^n(C, W).$$

Suppose that $\mathbb{E}^n(C, -)|_{\mathcal{W}} = 0$, then $\mathbb{E}^{n-1}(K, -)|_{\mathcal{W}} = 0$ as $\text{id}_{\mathcal{X}}(\mathcal{W}) = 0$ and $n \geq 2$; contradicting that $\text{pd}_{\mathcal{W}}(K) = n - 1$. \square

Now we give a characterizations of \mathcal{X} -resolution dimensions of objects in $\widehat{\mathcal{X}}$.

Theorem 3.13. *Let \mathcal{X} be closed under extensions and CoCones such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . The following are equivalent for any $C \in \widehat{\mathcal{X}}$ and nonnegative integer n .*

- (1) $\text{resdim}_{\mathcal{X}}(C) \leq n$.
- (2) If $U \rightarrow X_{n-1} \rightarrow \cdots \rightarrow X_0 \rightarrow C$ is an \mathbb{E} -triangle sequence with $X_i \in \mathcal{X}$ for $0 \leq i \leq n - 1$, then $U \in \mathcal{X}$.
- (3) $\mathbb{E}^{n+i}(C, Y) = 0$ for any object $Y \in \widehat{\mathcal{W}}$ and $i \geq 1$.
- (4) $\mathbb{E}^{n+i}(C, W) = 0$ for any object $W \in \mathcal{W}$ and $i \geq 1$.
- (5) $\mathbb{E}^{n+1}(C, W) = 0$ for any object $W \in \mathcal{W}$.

Proof. By Proposition 3.12, we have (1) \Rightarrow (3) \Leftrightarrow (4) \Rightarrow (5). (2) \Rightarrow (1) is trivial.

(5) \Rightarrow (1) Since $C \in \widehat{\mathcal{X}}$, there is an \mathbb{E} -triangle sequence

$$U \rightarrow X_{n-1} \rightarrow \cdots \rightarrow X_0 \rightarrow C$$

with $X_i \in \mathcal{X}$ for $0 \leq i \leq n - 1$. We also have $\mathbb{E}(U, \mathcal{W}) \cong \mathbb{E}^{n+1}(C, \mathcal{W}) = 0$ as \mathcal{W} is \mathcal{X} -injective. Note that $U \in \widehat{\mathcal{X}}$ by Lemma 3.11. Hence there exists an \mathbb{E} -triangle

$$Y_U \longrightarrow X_U \longrightarrow U \xrightarrow{\theta} \succ$$

with $X_U \in \mathcal{X}$ and $Y_U \in \widehat{\mathcal{W}}$ by Theorem 3.7. It follows that the \mathbb{E} -triangle θ splits as $\mathbb{E}(U, \mathcal{W}) = 0$. Hence $U \in \mathcal{X}$. \square

Proposition 3.14. *Let \mathcal{X} be closed under extensions and CoCones such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . Then $\widehat{\mathcal{X}}$ is closed under direct summands.*

Proof. Suppose $C_1 \oplus C_2 \in \widehat{\mathcal{X}}$. Proceed by induction on $n = \text{resdim}_{\mathcal{X}}(C_1 \oplus C_2)$. If $n = 0$, then C_1 and C_2 are in \mathcal{X} .

Suppose $n > 0$. There is an \mathbb{E} -triangle

$$K \longrightarrow X \xrightarrow{y} C_1 \oplus C_2 \dashrightarrow$$

with $X \in \mathcal{X}$ and $\text{resdim}_{\mathcal{X}}(K) = n - 1$. By $(ET4)^{op}$, we obtain a commutative diagram

$$\begin{array}{ccccccc}
 K & \longrightarrow & L_2 & \longrightarrow & C_1 & \dashrightarrow & \triangleright \\
 \parallel & & \downarrow x_2 & & \downarrow \begin{bmatrix} 1 \\ 0 \end{bmatrix} & & \\
 K & \xrightarrow{x} & X & \xrightarrow{\begin{bmatrix} y_1 \\ y_2 \end{bmatrix}} & C_1 \oplus C_2 & \dashrightarrow \delta & \triangleright \\
 & & \downarrow y_2 & & \downarrow \begin{bmatrix} 0 & 1 \end{bmatrix} & & \\
 & & C_2 & \xlongequal{\quad} & C_2 & & \\
 & & \downarrow \delta_2 & & \downarrow 0 & & \\
 & & \Psi & & \Psi & &
 \end{array}$$

Similarly, we can obtain an \mathbb{E} -triangle

$$L_1 \xrightarrow{x_1} X \xrightarrow{y_1} C_1 \dashrightarrow^{\delta_1} \triangleright.$$

Hence there is an \mathbb{E} -triangle

$$L_1 \oplus L_2 \xrightarrow{\begin{bmatrix} x_1 & 0 \\ 0 & x_2 \end{bmatrix}} X \oplus X \xrightarrow{\begin{bmatrix} y_1 & 0 \\ 0 & y_2 \end{bmatrix}} C_1 \oplus C_2 \dashrightarrow^{\delta_1 \oplus \delta_2} \triangleright.$$

By Lemma 3.11, $L_1 \oplus L_2 \in \hat{\mathcal{X}}$, and Theorem 3.13 shows that $\text{resdim}_{\mathcal{X}}(L_1 \oplus L_2) \leq n - 1$. By inductive hypothesis, L_1 and L_2 are in $\hat{\mathcal{X}}$. It follows that C_1 and C_2 are in $\hat{\mathcal{X}}$ by Lemma 3.11 \square

The following result will play a key role in Section 4.

Proposition 3.15. *Let \mathcal{X} be closed under extensions and CoCones such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . For any \mathbb{E} -triangle $A \xrightarrow{x} B \xrightarrow{y} C \dashrightarrow^{\delta} \triangleright$, if any two of A, B and C are in $\hat{\mathcal{C}}$, then the third one is in $\hat{\mathcal{X}}$.*

Proof. Since we already knew that $\hat{\mathcal{X}}$ is closed under extensions by Proposition 3.10, it suffices to show that if $B \in \hat{\mathcal{X}}$, then $A \in \hat{\mathcal{X}}$ if and only if $C \in \hat{\mathcal{X}}$. We first show that if A and B are in $\hat{\mathcal{X}}$, then $C \in \hat{\mathcal{X}}$. Since $B \in \hat{\mathcal{X}}$, we have an \mathbb{E} -triangle

$$Y_B \longrightarrow X_B \longrightarrow B \dashrightarrow$$

with $X_B \in \mathcal{X}, Y_B \in \widehat{\mathcal{W}}$. By $(ET4)^{op}$, we obtain a commutative diagram

$$\begin{array}{ccccc}
 Y_B & \longrightarrow & L & \longrightarrow & A \\
 \parallel & & \downarrow & & \downarrow \\
 Y_B & \longrightarrow & X_B & \longrightarrow & B \\
 & & \downarrow & & \downarrow \\
 & & C & \xlongequal{\quad} & C
 \end{array}$$

It follows that $L \in \hat{\mathcal{X}}$ as A and Y_B are $\in \hat{\mathcal{X}}$. Therefore $C \in \hat{\mathcal{X}}$.

Suppose now B and C are $\in \hat{\mathcal{X}}$. By Lemma 3.11, $L \in \hat{\mathcal{X}}$. Applying the just established result to \mathbb{E} -triangle

$$Y_B \longrightarrow L \longrightarrow A \dashrightarrow$$

gives $A \in \hat{\mathcal{X}}$. \square

We say \mathcal{X} is resolving (coresolving, resp.) if it contains $\text{Proj}(\mathcal{C})$ ($\text{Inj}(\mathcal{C})$, resp.), closed under extensions and CoCones (Cones, resp.).

Corollary 3.16. *Let \mathcal{X} be a resolving subcategory of \mathcal{C} such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . If \mathcal{C} is a Frobenius extriangulated category, then $\hat{\mathcal{X}}$ is a thick subcategory of $\underline{\mathcal{C}}$.*

Proof. If \mathcal{C} is a Frobenius extriangulated category, it is easy to see that $\hat{\mathcal{X}}$ is also a Frobenius extriangulated category by Proposition 3.15. By [15, Remark 7.5] and Proposition 3.14, $\hat{\mathcal{X}}$ is a thick subcategory of $\underline{\mathcal{C}}$. \square

4. Cotorsion Pairs in Extriangulated categories

We recall the definition of a cotorsion pair in an extriangulated category from [15].

Definition 4.1. Let \mathcal{U} and \mathcal{V} be subcategories of an extriangulated category \mathcal{C} . The pair $(\mathcal{U}, \mathcal{V})$ is called a cotorsion pair if it satisfies the following conditions.

- (1) $\mathbb{E}(\mathcal{U}, \mathcal{V}) = 0$.
- (2) For any $C \in \mathcal{C}$, there exists an \mathbb{E} -triangle

$$V_C \xrightarrow{x} U_C \xrightarrow{y} C \xrightarrow{\delta} \triangleright$$

satisfying $U_C \in \mathcal{U}, V_C \in \mathcal{V}$.

- (3) For any $C \in \mathcal{C}$, there exists an \mathbb{E} -triangle

$$C \xrightarrow{f} V^C \xrightarrow{g} U^C \xrightarrow{\theta} \triangleright$$

satisfying $U^C \in \mathcal{U}, V^C \in \mathcal{V}$.

Remark 4.2. Let $(\mathcal{U}, \mathcal{V})$ be a cotorsion pair on an extriangulated category \mathcal{C} . Then

- (1) \mathcal{U} is a precovering class in \mathcal{C} and \mathcal{V} is a preenveloping class in \mathcal{C} .
- (2) $C \in \mathcal{U}$ if and only if $\mathbb{E}(C, \mathcal{V}) = 0$.
- (3) $C \in \mathcal{V}$ if and only if $\mathbb{E}(\mathcal{U}, C) = 0$.
- (4) \mathcal{U} and \mathcal{V} are closed under extensions.
- (5) $\text{Proj}(\mathcal{C}) \subseteq \mathcal{U}$ and $\text{Inj}(\mathcal{C}) \subseteq \mathcal{V}$.

Lemma 4.3. *For a cotorsion pair $(\mathcal{U}, \mathcal{V})$ on \mathcal{C} , the following conditions are equivalent.*

- (1) \mathcal{U} is resolving.
- (2) \mathcal{V} is coresolving.
- (3) $\text{id}_{\mathcal{U}}(\mathcal{V}) = 0$.

Proof. The proof is similar to that of [10, Lemma 5.24]. \square

Remark 4.4. Note that hereditary cotorsion pairs can only be defined on an extriangulated category with enough projectives and injectives.

Proposition 4.5. *Let \mathcal{X} be closed under extensions and CoCones such that \mathcal{W} is an \mathcal{X} -injective cogenerator for \mathcal{X} . Then $(\mathcal{X}, \widehat{\mathcal{W}})$ is a cotorsion pair on the extriangulated category $\hat{\mathcal{X}}$.*

Proof. This comes from Theorem 3.7, Proposition 3.9, 3.10 and 3.14. \square

Definition 4.6. Let \mathcal{H} be a subcategory of \mathcal{C} . \mathcal{H} is called *specialy precovering* in \mathcal{C} provided that for any $C \in \mathcal{C}$, there is an \mathbb{E} -triangle

$$K \longrightarrow H \longrightarrow C \dashrightarrow$$

with $H \in \mathcal{H}$, $\mathbb{E}(\mathcal{H}, K) = 0$.

Lemma 4.7. Let \mathcal{H} be a subcategory of \mathcal{C} . Suppose that \mathcal{H} is resolving and specialy precovering. Then $\mathcal{H} \cap \mathcal{H}^\perp$ is an \mathcal{H} -injective cogenerator for \mathcal{H} .

Proof. Let $H \in \mathcal{H}$. There is an \mathbb{E} -triangle

$$H \longrightarrow I \longrightarrow X \dashrightarrow$$

with $I \in \text{Inj}(\mathcal{C})$. Since \mathcal{H} is specialy precovering, we have an \mathbb{E} -triangle

$$K \longrightarrow H' \longrightarrow X \dashrightarrow$$

with $H' \in \mathcal{H}$ and $\mathbb{E}(\mathcal{H}, K) = 0$. By Lemma 2.10, we have the following commutative diagram

$$\begin{array}{ccccc} & & H & \xlongequal{\quad} & H \\ & & \downarrow & & \downarrow \\ K & \longrightarrow & M & \longrightarrow & I \\ \parallel & & \downarrow & & \downarrow \\ K & \longrightarrow & H' & \longrightarrow & X. \end{array}$$

Since \mathcal{H} is closed under extensions, it follows that $M \in \mathcal{H}$. Note $\mathbb{E}(\mathcal{H}, I) = \mathbb{E}(\mathcal{H}, K) = 0$. So $\mathbb{E}(\mathcal{H}, M) = 0$. We claim $M \in \mathcal{H} \cap \mathcal{H}^\perp$. Indeed, for any positive integer n and $H' \in \mathcal{H}$, we have the following \mathbb{E} -triangle sequence

$$L \rightarrow P_{n-1} \rightarrow \cdots \rightarrow P_1 \rightarrow P_0 \rightarrow H''$$

with $P_i \in \text{Proj}(\mathcal{C})$ for $0 \leq i \leq n-1$. We have $\mathbb{E}^n(H'', M) \cong \mathbb{E}(L, M) = 0$ as \mathcal{H} is resolving and $\mathbb{E}(\mathcal{H}, M) = 0$. Hence $M \in \mathcal{H} \cap \mathcal{H}^\perp$. The second column in the above diagram implies that $\mathcal{H} \cap \mathcal{H}^\perp$ is an \mathcal{H} -injective cogenerator for \mathcal{H} . \square

Theorem 4.8. Let \mathcal{C} be a Frobenius extriangulated category. The assignments

$$(\mathcal{U}, \mathcal{V}) \mapsto \mathcal{U} \quad \text{and} \quad \mathcal{H} \mapsto (\mathcal{H}, \widehat{\mathcal{M}}),$$

where $\mathcal{M} = \mathcal{H} \cap \mathcal{H}^\perp$, give mutually inverse bijections between the following classes:

- (1) Hereditary cotorsion pairs $(\mathcal{U}, \mathcal{V})$ on \mathcal{C} with $\widehat{\mathcal{U}} = \mathcal{C}$ and $\check{\mathcal{V}} = \mathcal{C}$.
- (2) Subcategories \mathcal{H} of \mathcal{C} , which is specialy precovering and resolving in \mathcal{C} such that $\widehat{\mathcal{H}} = \mathcal{C}$ and for any $H \in \mathcal{H}$, there exists a positive integer $i \geq 1$ making $\mathbb{E}^i(H', H) = 0$ for any $H' \in \mathcal{H}$.

Proof. Let $(\mathcal{U}, \mathcal{V})$ be a hereditary cotorsion pair on \mathcal{C} with $\widehat{\mathcal{U}} = \mathcal{C}$ and $\check{\mathcal{V}} = \mathcal{C}$. Then \mathcal{U} is precovering and resolving in \mathcal{C} by Remark 4.2 and Lemma 4.3.

For any $U \in \mathcal{U}$, $U \in \check{\mathcal{V}}$ by assumption. Hence there exists an \mathbb{E} -triangle sequence

$$U \rightarrow V_0 \rightarrow V_1 \rightarrow \cdots \rightarrow V_n$$

with $V_i \in \mathcal{V}$ for $0 \leq i \leq n$. We have $\mathbb{E}^{n+1}(U', U) \cong \mathbb{E}(U', V_n) = 0$ for any $U' \in \mathcal{U}$. Clearly, $n+1$ is the desired i .

Assume now that \mathcal{H} is a subcategory of \mathcal{C} satisfying the conditions in (2). By Lemma 4.7, $\mathcal{M} = \mathcal{H} \cap \mathcal{H}^\perp$ is an \mathcal{H} -injective cogenerator for \mathcal{H} . By Proposition 3.9, $\widehat{\mathcal{M}} = \widehat{\mathcal{H}} \cap \mathcal{H}^\perp = \mathcal{H}^\perp$, where the second equality is due to $\widehat{\mathcal{H}} = \mathcal{C}$. Therefore $(\mathcal{H}, \widehat{\mathcal{M}})$ is a hereditary cotorsion pair on \mathcal{C} by Proposition 4.5. For any $X \in \mathcal{C}$, we have an \mathbb{E} -triangle

$$X \longrightarrow Y \longrightarrow L \dashrightarrow$$

with $Y \in \widehat{\mathcal{M}}$ and $L \in \mathcal{H}$ by Theorem 3.7. We claim that $L \in \widetilde{\mathcal{M}} := \widetilde{\widehat{\mathcal{M}}}$. Indeed, note that there exists a positive integer i such that $\mathbb{E}^i(H, L) = 0$ for any $H \in \mathcal{H}$. If $i = 1$, then $L \in \widehat{\mathcal{M}}$, as desired. Assume $i > 1$, then $\mathbb{E}^i(H, L) = 0$ for any $H \in \mathcal{H}$ implies that $\Sigma^{i-1}(L) \in \mathcal{H}^\perp$. Meanwhile $\Sigma^{i-1}(L) \in \widehat{\mathcal{H}}$ as \mathcal{H} contains $\text{Inj}(\mathcal{C})$. Hence $\Sigma^{i-1}(L) \in \widehat{\mathcal{H}} \cap \mathcal{H}^\perp = \widehat{\mathcal{M}}$ by Proposition 3.9. Since $\text{Inj}(\mathcal{C}) \subseteq \widehat{\mathcal{M}}$, it follows that $L \in \widetilde{\mathcal{M}}$. Hence $\mathcal{C} = \widetilde{\mathcal{M}}$. This completes the proof. \square

We conclude the paper by the following result, which gives a characterization of silting subcategories on stable categories and also generalizes [8, Corollary 3.7]. For the convenience of the reader, we recall the definition of silting subcategories.

Definition 4.9. [1, Definition 2.1] Let \mathcal{M} be a subcategory of a triangulated category \mathcal{T} . \mathcal{M} is called silting if $\text{Hom}_{\mathcal{T}}(\mathcal{M}, \mathcal{M}[\geq 1]) = 0$ and $\mathcal{T} = \text{thick}(\mathcal{M})$, where $\mathcal{T} = \text{thick}(\mathcal{M})$ is the smallest triangulated subcategory of \mathcal{T} which contains \mathcal{M} and is closed under direct summands and isomorphisms.

Corollary 4.10. *Let \mathcal{C} be a Frobenius extriangulated category. The assignments*

$$\underline{\mathcal{M}} \mapsto {}^\perp \mathcal{M} \quad \text{and} \quad \mathcal{H} \mapsto \underline{\mathcal{H} \cap \mathcal{H}^\perp}$$

give mutually inverse bijections between the following classes:

- (1) *Silting subcategories $\underline{\mathcal{M}}$ of the stable category $\underline{\mathcal{C}}$.*
- (2) *Subcategories \mathcal{H} of \mathcal{C} , which is specially precovering and resolving in \mathcal{C} such that $\widehat{\mathcal{H}} = \mathcal{C}$ and for any $H \in \mathcal{H}$, there exists a positive integer $i \geq 1$ making $\mathbb{E}^i(H', H) = 0$ for any $H' \in \mathcal{H}$.*

Proof. Since \mathcal{C} is a Frobenius extriangulated category, it follows that $\underline{\mathcal{C}}$ is a triangulated category by [15, Corollary 7.4]. By [7, Corollary 3.10], $(\mathcal{U}, \mathcal{V})$ is a cotorsion pair on \mathcal{C} if and only if $(\underline{\mathcal{U}}, \underline{\mathcal{V}})$ is a cotorsion pair on $\underline{\mathcal{C}}$. Note that any distinguished triangle

$$A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{h} A\langle 1 \rangle$$

in $\underline{\mathcal{C}}$ is induced by an \mathbb{E} -triangle

$$A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{\delta} \dashrightarrow .$$

By [14, Remark 4.2], it is easy to see that $(\mathcal{U}, \mathcal{V})$ is a hereditary cotorsion pair on \mathcal{C} with $\widehat{\mathcal{U}} = \widetilde{\mathcal{V}} = \mathcal{C}$ if and only if $(\underline{\mathcal{U}}, \underline{\mathcal{V}})$ is a bounded co-t-structure on $\underline{\mathcal{C}}$. Hence the Corollary follows from [14, Corollary 5.9] and Theorem 4.8. \square

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