

i) let  $t \in \{-1, 0, 1\}^n$  be a vector with at least  $m$  non-zero entries and let  $u \in \{0, 1\}^n$  be a uniformly random vector.

We want to show that:

$$\Pr(|u \cdot t| \leq \sqrt{m}/4) \leq C, \quad (C)$$

because  $u$  is uniformly random, and  $m$  entries of  $t$  are non zero,  $|u \cdot t|$  would be the sum of approximately  $m/2$  independent variables with value  $\{-1, 0, 1\}$ , essentially a random walk. And by the CLT, with  $n$  and  $m$  sufficiently large,  $|u \cdot t|$  is approximately gaussian.

We can thus use a normal approximation for the simple random walk

$$\hookrightarrow \text{Var}( |u \cdot t| ) = m/2$$

thus, the standard deviation of  $|u \cdot t|$  is approximately  $\sqrt{m}/2$ .

thus, the probability of it falling between  $[-\sqrt{m}/4, \sqrt{m}/4]$  (z-score =  $\frac{\sqrt{m}}{2}$ ) is approximately 0.52, which is a constant less than 1.

another explanation is that the  
 the r.v.  $U+1$  behaves like a  $N(0, m/2)$ ,  
 and the STD is  $\sqrt{m/2}$ , we can say the STD is  
 at least  $\sqrt{m/2}$  with high probability. This would  
 also make  $\Pr(U+1 \leq \sqrt{m/2}) \leq \frac{3}{10}$ . In other words, it  
 is bounded away from 1, with a constant strictly  
 less than 1.

However, we do have error since this problem is discrete  
 and we also cannot guarantee exactly  $m/2$ , but due to  
 large  $n, m$ , and  $WT_1$ , we can be pretty confident it  
 is close, so it is a gaussian as well. So the error of  
 $c = 0.9$ , when we had stricter constants like 0.52 and 0.7,  
 comes from the fact that the variance can change based  
 on the # of "random walks" we take, but is still at  
 least  $\sqrt{m/2}$  with high probability (larger bound for error)  
 Thus, we can still say that  $\Pr(U+1 \leq \sqrt{m/2}) < c$ , and  
 and in this case we set  $c + \epsilon$  be the generous term  $\frac{9}{10}$

Back to the problem, if  $D_0 \in B$ , it differs from  $D$  by at least  $256\alpha^2 n^2$  indices. If  $t = D_0 - D$ ,  $t \in \{-1, 0, 1\}^n$  and  $t$  has at least  $256\alpha^2 n^2$  non zero entries. And since queries are chosen at random,  $s_j$  is chosen at random uniformly  $s_j \in \{0, 1\}^n$ ,  $s_j \cdot t \in F_j$

$$\text{Then, } \Pr[|F_j \cdot (D - D_0)| \leq \frac{\sqrt{256\alpha^2 n^2}}{4}] \leq \frac{9}{10}$$

From the previous part, which simplifies  $P$

$$\Pr[|q_{s_j}(D_0) - q_{s_j}(D)| \leq 4\alpha n] \leq \frac{9}{10}$$

2) let  $y = \{M(s(x_i, I))\}$

we can use.

$$\Pr(B|A) = \frac{\Pr(A, B)}{\Pr(A)}$$

$$\text{then, } \Pr(y \in T | I \notin I) = \frac{\Pr(y \in T \cap I \notin I)}{\Pr(I \notin I)}$$

$$= \frac{\sum_{I \notin I} \Pr(y \in T \cap I \notin I)}{m/n}$$

$$\text{Similarly, } \Pr(y | I \notin I) = \frac{\Pr(y \cap I \notin I)}{\Pr(I \notin I)} = \frac{\sum_{I \notin I} \Pr(y \cap T \notin I)}{\frac{m}{n}}$$

Then, we have

$$(n-m) \sum_{I \notin I} \Pr(y \in T, I) \leq e^{\epsilon} m \sum_{I \notin I} \Pr(y \notin T, I)$$

$$= (n-m) \left(\frac{m}{n}\right) \leq (m) \left(\frac{n-m}{n}\right) e^{\epsilon} \quad , \text{since } I \in U \\ \text{mapped in time out}$$

$$= 1 \leq e^{\epsilon} , \text{ which is always true} \quad , \quad I \notin U \text{ mapped} \\ \frac{n-m}{n} \text{ time}$$

3) if we don't know anything about the size of the sets  $G_i$ , the sensitivity is trivially 2. This is because if we have a 2 clusters,  $i_1 \neq i_2$ , and  $G_{i_1}$  contains  $\|x\|_1 = 1$ , then if we change it so that  $\|x\|_1 = 1$  but now belongs to  $G_{i_2}$ , we have a change of 1 in both  $G_{i_1}'$  and  $G_{i_2}'$ , with the  $l_1$ -sensitivity then being 2.

more formally, since changing one value  $x$  can shift one data point from one to another changing the  $|G_i|$  by 1 by a maximum of 1 and  $G_i$  by  $\frac{1}{|G_i|}$

$$\max_i (f_j \left( 0 + \frac{1}{|G_{i_1}|}_{\min} \right) + \left( 0 + \frac{1}{|G_{i_2}|}_{\max} \right), i \notin$$

$$\leq \frac{2}{\min_i |G_i|}$$

4) a) define  $A = M(D) \cap S$

$B = M(D') \cap S$

we can break this problem down for all subset  $S$

1) if  $S$  is empty or doesn't contain  $D, D'$ , or  $\perp$

$$|Pr(A) - Pr(B)| = |0 - 0| \leq \delta$$

2) if  $S$  contains  $D, D' \perp$

$$|Pr(A) - Pr(B)| = |1 - 1| = 0 \leq \delta$$

3)  $S$  contains  $D, D'$  but not  $\perp$

$$|Pr(A) - Pr(B)| = |\delta - \delta| = 0 \leq \delta$$

4)  $S$  contains  $D$  and not  $D, \perp$

$$|Pr(A) - Pr(B)| = |\delta - 0| = \delta \leq \delta$$

5)  $D \in S, D, \perp \notin S$

$$|Pr(A) - Pr(B)| = |0 - \delta| = \delta \leq \delta$$

6)  $\perp \in S, D, D' \notin S$

$$|Pr(A) - Pr(B)| = |1 - \delta - (1 - \delta)| = 0 \leq \delta$$

7)  $\perp, D \in S, D' \notin S$  and  $\perp, D' \in S, D \notin S$

$$|Pr(A) - Pr(B)| = |1 - (1 - \delta)| = \delta = \delta \leq \delta$$

Since  $|Pr(A) - Pr(B)| \leq \delta$ ,  $M$  is  $\delta$ -DP

4) b) without loss of generality, let  $X \sim X'$  differ at  $X_1$ , with  $X_1^{(0)} = x^*$   
 for  $S = \{S_1, \dots, S_n\}$

$$\text{Then, } \Pr(M(X)=S) = \Pr(Y_1=S_1) \prod_{j=2}^n \Pr(Y_j=S_j)$$

$$\Pr(M(X')=S) = \Pr(Y'_1=S_1) \prod_{j=2}^n \Pr(Y'_j=S_j)$$

Since  $Y_i = S_j$  for  $2 \leq j \leq n$

$$\prod_{j=2}^n \Pr(Y_j=S_j) = \prod_{j=2}^n \Pr(Y'_j=S_j) \approx 1$$

subtracting

$$\Pr(M(X)=S) - \Pr(M(X')=S) = \left( \prod_{j=2}^n \Pr(Y_j=S_j) \right) \cdot (\Pr(Y_1=S_1) - \Pr(Y'_1=S_1))$$

$$\leq \Pr(Y_1=S_1) - \Pr(Y'_1=S_1)$$

if  $S_1 = X_1$ , then

$$\Pr(Y_1=S_1) - \Pr(Y'_1=S_1) = \delta - 0 = \delta$$

if  $S_1 = X'_1$ , then  $\Pr(Y_1=S_1) - \Pr(Y'_1=S_1) = 0 - \delta = -\delta$

if  $S_1 \neq X_1, X'_1$ , then  $\Pr(Y_1=S_1) - \Pr(Y'_1=S_1) = 0 - 0 = 0$

if  $S_1 = \perp$ , then  $\Pr(Y_1=S_1) - \Pr(Y'_1=S_1) = 1 - (1 - \delta) = \delta$

$$\text{then } |\Pr(M(X)=S) - \Pr(M(X')=S)| \leq \delta,$$

which is  $\delta$ -DP

5) Take  $n=1$ , and  $X=\{1\}$ . Then,  $X'=\{0\}$

then, for all  $\epsilon$ , there exists an interval  $S$  where  
 $\Pr(\tilde{f}(X') \in S) = 0$ , and  $\Pr(\tilde{f}(X) \in S) > 0$ ,

where

$$\Pr(\tilde{f}(X') \in S) \leq e^{\epsilon} \Pr(\tilde{f}(X) \in S)$$
$$= c \leq e^{\epsilon}(0) , \text{ when } c > 0$$

which fails the definition of  $\epsilon$ -DP

for all  $\epsilon \in \mathbb{R}$ , the interval  $[-3/\epsilon, 3/\epsilon]$ , where the last interval  $[3/\epsilon - 1, 3/\epsilon]$  that  $\tilde{f}(X')$  applies has probability 0, since  $0+3/\epsilon$  (highest thw  $\tilde{f}(X')$  can produce is always  $\leq 1+3/\epsilon$  which is what  $\tilde{f}(X)$  can produce. Thus,  $\Pr(\tilde{f}(X') \in [-3/\epsilon - 1, 3/\epsilon]) = \frac{\epsilon}{\epsilon} = 1$ , but  $\Pr(\tilde{f}(X) \in [3/\epsilon - 1, 3/\epsilon]) = 0$ , if  $\epsilon > 0$  the intervals don't overlap, and if  $\epsilon = 0$  so  $\frac{\epsilon}{\epsilon} \leq e^{\epsilon}(0)$  doesn't satisfy the definition of  $\epsilon$ -DP

if  $\epsilon \geq 6$ , then the intervals are completely disjoint, giving zero privacy.

6) (M1)  $M_1(x) = \lceil \bar{x} + z \rceil^b$ ,  $z \sim \text{Lap}(2/n)$

Clamping  $\bar{x} + z$  is a post processing step, and we prove post processing do not reveal additional information. More formally, if  $g(x) = \bar{x} + z$ , which is  $\epsilon$ -DP, then  $M_1(x) = |g(x)|^b$  is also  $\epsilon$ -DP by post processing

Thus,  $M_1$  is still  $\epsilon$ -DP

To find the smallest value of  $\epsilon_1$  we know  $\text{Lap}(\Delta/\epsilon_1)$  is  $\epsilon$ -DP so we find  $\Delta$ .

$$\Delta \bar{x} = \frac{1}{n} \quad (\text{we can change one value from } 0 \text{ to } 1)$$

$$\Rightarrow \text{Lap}\left(\frac{1}{n}/\epsilon_1\right) = \text{Lap}\left(\frac{1}{2n}\right)$$

$M_1(x)$  is  $\frac{1}{2}-\text{DP}$

6) (M2) from the prev part we have  $\alpha = \frac{1}{n}$

$$\frac{\Pr(M2(x)=1)}{\Pr(M2(x')=1)} \leq e^\epsilon \text{ to satisfy OP.}$$

However if  $\bar{x}$  change by  $\frac{1}{n}$ , the above ratio is

$$\frac{\bar{x} + \frac{1}{n}}{\bar{x}} \leq e^\epsilon. \quad \text{If } \bar{x} = 0 \rightarrow \text{, then it is infinite and } 2e^\epsilon. \quad \left( \begin{array}{l} x \text{ is all zero} \\ \text{and } x \neq 0 \text{ or 1} \end{array} \right)$$

$$\text{Similarly, } \frac{\Pr(M2(x)=0)}{\Pr(M2(x')=0)} \leq e^\epsilon$$

$$= \frac{1-\bar{x}}{1-\bar{x}-\frac{1}{n}} \leq e^\epsilon \text{ when if } \bar{x} = 1 - \frac{1}{n}, \text{ it goes to infinity again, as there is no bound on } n.$$

Thus, there is no  $\epsilon$  when  $\frac{\Pr(M2(x)=0 \text{ or } 1)}{\Pr(M2(x')=0 \text{ or } 1)}$  when  $e^\epsilon$  is bounded by infinity, so M2 is not  $\epsilon$ -DP for any finite  $\epsilon$ .

7) a)  $X \sim X'$ , without loss of generality, let's say they differ at first element,  $x_1$  and  $x'_1$

if  $x'_1 = x_1 + c$ , where  $c > 2\epsilon n$

$$f(x') = f(x) + \frac{c}{n} = f(x) + 2\epsilon$$

for 1-OP, we have

$\overline{\epsilon}$

$$\Pr(M(x) \in S) \leq e \Pr(M(x') \in S)$$

Suppose that  $\Pr(|M(x) - f(x)| \leq \epsilon) \geq 0.9$  for  $t \in \mathbb{R}, x$ .

Let  $S = [f(x) - t, f(x) + t]$ ,  $T = [f(x') - t, f(x') + t]$

$\Pr(M(x) \in S) \geq 0.9$  since  $S$  and  $T$  disjoint

then  $\Pr(M(x') \in S) \leq \underbrace{1 - \Pr(|M(x') - f(x')| \leq t)}_{\text{if } t \geq 0.9} \leq 0.1$

$$\Pr(M(x') \in S) \leq 0.1$$

from OP,  $0.9 \leq \Pr(M(x) \in S) \leq e \Pr(M(x') \in S)$

plugging our numbers in

$0.9 \leq e \cdot 0.1$ , which is a contradiction. Thus, for any finite  $t > 0$ , there is no 1-OP algorithm  $M$  s.t.  $\Pr(|M(x) - f(x)| \leq t) \geq 0.9$

7) b) here, the max A happens if we change one value in  $X$  from  $-R$  to  $R$  or vice versa, resulting in an  $A$  of  $\frac{2R}{n}$ .

then, this is just a scaled case of  $M(X)$  in the  $X \in \{0, 1\}^n$  case,  
 where  $A_2(x) = f(x) + \text{Lap}(\frac{2R}{n\epsilon})$ , when  $2R/n = 1$ ,  
 which we already proved is  $\epsilon$ -DP

Then,  $\Pr(|A_2(x) - f(x)| \leq \frac{CR}{n\epsilon}) \geq 0.9$  is just

$\Pr(|Z| \leq \frac{CR}{n\epsilon}) \geq 0.9$ , where  $Z \sim \text{Lap}(\frac{2R}{n\epsilon})$  is a r.v.

using the cdf for  $\text{Lap}(c, b)$   $\rightarrow \frac{1}{2} \cdot \exp(-\frac{|x|}{b})$ . Since we are taking the absolute value, our interval is  $2x$  thus  $\rightarrow$

$$\Pr(|\text{Lap}(c, b)| \leq t) = 1 - e^{-t/b} \quad \rightarrow \begin{cases} \text{CDF of Lap}(c, b) \\ \frac{1}{2} \exp(-\frac{|x|}{b}) \end{cases} \quad x \in \mathbb{R}$$

then,  $\Pr(|\text{Lap}(2R/n\epsilon)| \leq CR/n\epsilon) \geq 0.9$

$$\rightarrow 1 - e^{-CR/n\epsilon} / (2R/n\epsilon) =$$

$$= 1 - e^{-C/2} \geq 0.9. \quad = e^{-C/2} \leq 0.1$$

$(\geq 2\ln(10))$ . This value of  $C$

satisfies  $\Pr(|A_2(x) - f(x)| \leq \frac{CR}{2\epsilon}) \geq 0.9$ .

8) a) for each  $X_i$ ,  $Y_i = X_i$  w.p.  $\frac{e^\epsilon}{e^\epsilon + k - 1}$ . This is proportional to  $e^\epsilon$

for  $s \in \{1, \dots, k\} \setminus \{X_i\}$ ,  $Y_i = s$  w.p.  $\frac{1}{e^\epsilon + k - 1}$ .

$$\text{Then, for all } s \in \{1, \dots, k\}, \Pr(Y_i = s) = \frac{e^\epsilon}{e^\epsilon + k - 1} + \frac{k-1}{e^\epsilon + k - 1} = 1$$

$$X \sim X' \quad \Pr(Y_i = y | X_i = x_i) = \begin{cases} e^\epsilon / e^\epsilon + k - 1 & y_i = x_i \\ 1 / e^\epsilon + k - 1 & y_i \neq x_i \end{cases}$$

Without loss of generality, say  $X \sim X'$  differs at  $X_i$ .

$$\begin{aligned} \Pr(Y = y | X) &= \frac{\Pr(Y_i = y | X_i)}{\Pr(Y_i = y | X')} = \frac{\prod_{j=1}^n \Pr(Y_j = y_j | X_j)}{\prod_{j=1}^n \Pr(Y_j = y_j | X'_j)} \\ &\quad \underbrace{\qquad}_{=1, \text{ since } X_j = X'_j \text{ for } j \neq i} \\ &= 1, \text{ since } X_j = X'_j \text{ for } j \neq i \end{aligned}$$

$$= \frac{\Pr(Y_i = y_i | X_i = a)}{\Pr(Y_i = y_i | X_i = b)} \quad \begin{matrix} \text{if } X_i = a \\ X_i = b \end{matrix}$$

Then if  $Y_i = a$

$$= \frac{\Pr(Y_i = a | X_i = a)}{\Pr(Y_i = a | X_i = b)} = \frac{e^\epsilon / e^\epsilon + k - 1}{1 / e^\epsilon + k - 1} = e^\epsilon$$

cont

if  $Y_1 = b$

$$\frac{\Pr(Y_1 = b | X_i = a)}{\Pr(Y_1 = b | X_i = b)} = \frac{e^{E[Y_1|X_i=a]}}{e^{E[Y_1|X_i=b]}} = e^{-\epsilon}$$

if  $Y_1 = c$

$$\frac{\Pr(Y_1 = c | X_i = a)}{\Pr(Y_1 = c | X_i = b)} = \frac{e^{E[Y_1|X_i=a]}}{e^{E[Y_1|X_i=b]}} = 1$$

Thus, the probability is bounded by  $e^{\epsilon_i}$  proving E-DP

8) b) convert  $X_i$  to a one-hot vector  $V_i \in \{0, 1\}^k$

$$V_{i,j} = \begin{cases} 1 & \text{if } j = X_i \\ 0 & \text{if } j \neq X_i \end{cases}$$

for all  $V_i$ , if  $V_{i,j} = 1$ . Let  $\varepsilon' = \varepsilon/2$ .

$$Y_{ij} = \begin{cases} v_{ij} & \text{w.p. } e^{\varepsilon'}/(e^{\varepsilon'} + 1) \\ 1 - v_{ij} & \text{w.p. } 1/(e^{\varepsilon'} + 1) \end{cases}$$

similar to the last part suppose differ at  $V_{i,a} = 1, V_{i,b} = 0$

$$V_{i,a}' = 0, V_{i,b}' = 1$$

$$\frac{\Pr(Y=y|X)}{\Pr(Y=y|X')} = \frac{\Pr(Y_i=y|V_i)}{\Pr(Y_i=y|V_i')} \cdot \underbrace{\frac{\prod_{j=2}^n \Pr(Y_j=y|V_j)}{\prod_{j=2}^n \Pr(Y_j=y|V'_j)}}_{=} = 1$$

Then,  $\frac{\Pr(Y_{i,a}|V_{i,a}=1)}{\Pr(Y_{i,a}|V_{i,a}=0)} = \begin{cases} \frac{e^{\varepsilon'}/(e^{\varepsilon'+1})}{1/(e^{\varepsilon'+1})} = e^{\varepsilon'/2} & Y_{i,a}=1 \\ \frac{1/(e^{\varepsilon'+1})}{e^{\varepsilon'}/(e^{\varepsilon'+1})} = e^{-\varepsilon'/2} & Y_{i,a}=0 \end{cases}$

$$\frac{\Pr(Y_{i,b}|V_{i,b}=0)}{\Pr(Y_{i,b}|V_{i,b}=1)} = \begin{cases} \frac{1/(e^{\varepsilon'+1})}{e^{\varepsilon'}/(e^{\varepsilon'+1})} = e^{-\varepsilon'/2} & Y_{i,b}=1 \\ \frac{e^{\varepsilon'}/(e^{\varepsilon'+1})}{1/(e^{\varepsilon'+1})} = e^{\varepsilon'/2} & Y_{i,b}=0 \end{cases}$$

if  $c \neq a, b$ , then  $\frac{\Pr(Y_{i,c} | V_{i,c})}{\Pr(Y_{i,c} | V_{i,a})} = \frac{\Pr(Y_{i,c} | V_{i,c})}{\Pr(Y_{i,c} | V_{i,b})}$

Then for most case  $Y_{i,a} = 1$  and  $Y_{i,b} = 0$ , we

have  $\frac{\Pr(Y_{i,c} | V_j)}{\Pr(Y_{i,c} | V_j')} = e^{\epsilon_j} \cdot e^{-\epsilon_b} = e^\epsilon$ ,

which satisfies  $\epsilon$ -DP

9) (M1). Intuitively, we need to find the probability that  $Z \sim \text{Lap}(2/n) > 1$  or  $<-1$ , and have  $f$  to be that probability, since  $M_1$  is  $\epsilon$ -DP within the interval but not outside (similar to HMTS), so we need  $f = \Pr(|\text{Lap}(b)| > 1)$

The CDF of  $\text{Lap}(b)$  is  $\frac{1}{2} \exp\left(\frac{x-\mu}{b}\right)$  for  $x \leq \mu$ .

$$\Pr(|\text{Lap}(b)| > 1) = e^{-n/2}$$

Then,  $f = e^{-n/2}$ , and  $M_1(x)$  is  $(e, e^{-n/2})$ -DP

9) ( $M_2$ ) This is like a continuation of problem 6.

The issue there was that there exists a case when in the worst case we have

$$\Pr(M_2(x)=1) \leq e^{\epsilon} \cdot \Pr(M_2(x')=1),$$

where worst case  $x$  has an 1, and  $x'$  is all 0s,

we have  $\frac{1}{n} \leq e^{\epsilon} \cdot 0$ , so we need to set the  $f = Y_n$  so it satisfies  $(\epsilon, S)$ -DP, with

$$f = Y_n.$$

Then,  $Y_n \leq e^{\epsilon} \cdot 0 + Y_n$  works, and is

$(\epsilon, Y_n)$ -DP

10) a) we can show this via an example. In the simple case, let

$$n=1, \text{ and } X = (b, b). \text{ Here } \beta(b, b) = \frac{b^2}{b} = 1.$$

Then, if we choose  $X_i \approx X^* \Rightarrow X^* = (0, 1)$  (or  $\lim_{i \rightarrow 0} x_i \downarrow$ ),

then  $\beta(X^*) = \frac{b}{\lim_{i \rightarrow 0} x_i}$ , which is infinite. So  $\Delta$  is

also infinity, and we cannot use regular DP  $\nrightarrow$  a useful estimate

b)  $S_{xy} = \sum_i x_i y_i$ . Suppose we change  $(x_i, y_i) \mapsto (x'_i, y'_i)$

$$\Delta = \max_{x_i, x'_i, y_i, y'_i \in S_{xy}} |x_i y_i - x'_i y'_i|. \text{ Thus } \max \Delta \text{ results}$$

when  $x_i y_i = b^2$  and  $x'_i y'_i = -b^2 \Rightarrow \Delta \approx 2b^2$

Thus,  $S_{xy} + \log(2b^2/\epsilon)$  is  $\epsilon$ -DP

for  $S_{xy}$ ,  $4 = \max_{x_i \in [-b, b]} |x_i^2 - x_i'^2| \leq b^2$ , since

$x^2$  and  $x'^2$  are  $\epsilon$ -DP on  $[0, b]$

$S_{xx} + Lup(\frac{b^2}{\epsilon})$  is  $\epsilon$ -DP

c) we first clip the dataset so for each  $(x_i, y_i)$ ,

$$x_i = \begin{cases} b & x_i > b \\ x_i & -b \leq x_i \leq b \\ -b & x_i < -b \end{cases} \quad y_i = \begin{cases} b & y_i > b \\ y_i & -b \leq y_i \leq b \\ -b & y_i < -b \end{cases}$$

We can just use simple post processing, with  $\epsilon = \epsilon_1 + \epsilon_2$  overall  
 so it is  $\epsilon$  overall

$$\tilde{S}_{xy} = S_{xy} + Lup\left(\frac{2b^2}{\epsilon_2}\right) = S_{xy} + Lup\left(\frac{4b^2}{\epsilon}\right)$$

$$\tilde{S}_{xx} = S_{xx} + Lup\left(\frac{4}{\epsilon_2}\right) = S_{xx} + Lup\left(\frac{2b^2}{\epsilon}\right)$$

$$\text{then } \tilde{\beta} = \frac{\tilde{S}_{xy}}{\tilde{S}_{xx}}$$

and  $\beta$  is  $\epsilon$ -DP overall