

Homework 3 Simple Solution

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1. (Binary hypothesis testing) [16]

Let X_1, X_2, \dots be a sequence of i.i.d. Bernoulli p random variables, that is,

$$\Pr\{X_i = 1\} = 1 - \Pr\{X_i = 0\} = p.$$

Based on the observations so far, the goal is of a decision maker to determine which of the following two hypotheses is true:

$$\mathcal{H}_0 : p = p_0$$

$$\mathcal{H}_1 : p = p_1$$

where $0 < p_0 < p_1 \leq 1/2$.

- a) (Warm-up) Consider the problem of making the decision based on X_1 .

Draw the optimal $(\pi_{1|0}, \pi_{0|1})$ trade-off curve. [4]

- b) Suppose the decision maker waits until an 1 appears and makes the decision based on the whole observed sequence. Sketch the optimal $(\pi_{1|0}, \pi_{0|1})$ trade-off curve. [4]

- c) Now suppose the decision maker waits until in total n 1's appear and makes the decision based on the whole observed sequence. Let $\varpi_{0|1}^*(n, \epsilon)$ denote the minimum type-II error probability subject to the constraint that the type-I error probability is not greater than ϵ , $0 < \epsilon < 1$. Does $\lim_{n \rightarrow \infty} \frac{1}{n} \log \frac{1}{\varpi_{0|1}^*(n, \epsilon)}$ exist? If so, find it. Otherwise, show that the limit does not exist. [8]

Solution:

- a) By Neyman-Pearson theorem, the optimal test is randomized LRT. Note that the likelihood ratio can only take two values: $\frac{p_1}{p_0}, \frac{1-p_1}{1-p_0}$. Therefore, discuss the range of τ we get

$$\begin{cases} \pi_{1|0} = 1, \pi_{0|1} = 0, & 0 \leq \tau < \frac{1-p_1}{1-p_0} \\ \pi_{1|0} = p_0 + \gamma(1-p_0), \pi_{0|1} = (1-\gamma)(1-p_1) = \frac{1-p_1}{1-p_0}(1-\pi_{1|0}), & \tau = \frac{1-p_1}{1-p_0} \\ \pi_{1|0} = p_0, \pi_{0|1} = 1-p_1, & \frac{1-p_1}{1-p_0} < \tau < \frac{p_1}{p_0} \\ \pi_{1|0} = \gamma p_0, \pi_{0|1} = (1-\gamma)p_1 + (1-p_1) = 1 - \frac{p_1}{p_0}\pi_{1|0}, & \tau = \frac{p_1}{p_0} \\ \pi_{1|0} = 0, \pi_{0|1} = 1, & \tau > \frac{p_1}{p_0}. \end{cases}$$

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We can then draw the trade-off curve using the equations derived above.

- b) Note that our observation can only be $1, 01, 001, 0001, \dots$, let L be the length of the observation, we have

$$\begin{aligned}\mathcal{H}_0 : & L \sim \text{Geo}(p_0) \\ \mathcal{H}_1 : & L \sim \text{Geo}(p_1)\end{aligned}$$

Similar to a), we can discuss the range of τ and get:

$$\begin{cases} \pi_{1|0} = 0, \pi_{0|1} = 1, \tau > \frac{p_1}{p_0} \\ \pi_{1|0} = \sum_{i=1}^{n-1} (1-p_0)^{i-1} p_0 + \gamma (1-p_0)^{n-1} p_0, \\ \pi_{0|1} = \sum_{i=n+1}^{\infty} (1-p_1)^{i-1} p_1 + (1-\gamma)(1-p_1)^{n-1} p_1, \tau = \frac{(1-p_1)^{n-1} p_1}{(1-p_0)^{n-1} p_0} \\ \pi_{1|0} = \sum_{i=1}^n (1-p_0)^{i-1} p_0, \pi_{0|1} = \sum_{i=n+1}^{\infty} (1-p_1)^{i-1} p_1, \frac{(1-p_1)^{n-1} p_1}{(1-p_0)^{n-1} p_0} > \tau > \frac{(1-p_1)^n p_1}{(1-p_0)^n p_0}. \end{cases}$$

And we can draw the trade-off curve using the equations derived above.

- c) The observation can be viewed as n i.i.d. geometric random variables. To see this, for any realization of observation, insert a “—” symbol in front of the sequence, also insert a “—” right after a “1”. For example, if $n = 4$ and the realization is 010001101, we write it as |01|0001|1|01|. Apparently, the length of the subsequence between two — is a geometric random variable. Hence, in this subproblem, we are testing $\text{Geo}(p_0)^{\otimes n}$ and $\text{Geo}(p_1)^{\otimes n}$. By Chernoff-Stein lemma,

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \frac{1}{\varpi_{0|1}^*(n, \epsilon)} = D(\text{Geo}(p_0) \parallel \text{Geo}(p_1)) = \log \frac{p_0}{p_1} + \left(\frac{1-p_0}{p_0} \right) \log \frac{1-p_0}{1-p_1}.$$

Grading Policy:

- Specify the optimal trade-off curve [2]; invoke Neyman-Pearson Theorem or directly argue optimality [2].
- Formulate the problem as detecting geometric random variables (or directly used the mass function) [2] and invoke Neyman-Pearson's to specify the curve [2].
- Formulate the problem as a hypothesis testing with n instances [3], Chernoff-Stein lemma [2], and calculation [3].

2. (Asymptotic behavior of posterior probability [12])

Consider a binary hypothesis testing problem

$$\begin{cases} \mathcal{H}_0 : X_i \stackrel{\text{i.i.d.}}{\sim} P_0, i = 1, 2, \dots, n \\ \mathcal{H}_1 : X_i \stackrel{\text{i.i.d.}}{\sim} P_1, i = 1, 2, \dots, n \end{cases}.$$

Under a Bayes setup, the unknown binary parameter Θ is assumed to be *random* and follow a prior distribution defined by the *prior probabilities*

$$\pi_0^{(0)} := \Pr \{ \Theta = 0 \}, \pi_1^{(0)} := \Pr \{ \Theta = 1 \}.$$

Let the *posterior probabilities* be the conditional distribution of Θ given $X^n = x^n$:

$$\pi_0^{(n)}(x^n) = \Pr\{\Theta = 0 | X^n = x^n\}, \quad \pi_1^{(n)}(x^n) = \Pr\{\Theta = 1 | X^n = x^n\}.$$

- a) Derive the expressions of $\pi_0^{(n)}(x^n)$ and $\pi_1^{(n)}(x^n)$ in terms of $\pi_0^{(0)}, \pi_1^{(0)}, P_0, P_1$. [4]
- b) Consider $\pi_0^{(n)}(X^n)$ and $\pi_1^{(n)}(X^n)$ as random variables, because they are functions of the random sequence X^n . Use the Strong Law of Large Numbers to show that if \mathcal{H}_0 is true, then with probability 1,

$$\pi_0^{(n)}(X^n) \rightarrow 1, \quad -\frac{1}{n} \log \pi_1^{(n)}(X^n) \rightarrow D(P_0 \| P_1) \quad \text{as } n \rightarrow \infty. \quad [8]$$

Solution:

- a) Denote $x^n = (x_1, x_2, \dots, x_i, \dots, x_n)$

$$\begin{aligned} \pi_0^{(n)}(X^n) &= \Pr\{\Theta = 0 | X^n = x^n\} = \frac{\Pr\{\Theta = 0, X^n = x^n\}}{\Pr\{X^n = x^n\}} \\ &= \frac{\pi_0^{(0)} \Pr\{X^n = x^n | \Theta = 0\}}{\pi_0^{(0)} \Pr\{X^n = x^n | \Theta = 0\} + \pi_1^{(0)} \Pr\{X^n = x^n | \Theta = 1\}} \\ &= \frac{\pi_0^{(0)} \prod_{i=1}^n P_0(x_i)}{\pi_0^{(0)} \prod_{i=1}^n P_0(x_i) + \pi_1^{(0)} \prod_{i=1}^n P_1(x_i)} \\ \pi_1^{(n)}(X^n) &= \frac{\pi_1^{(0)} \prod_{i=1}^n P_1(x_i)}{\pi_0^{(0)} \prod_{i=1}^n P_0(x_i) + \pi_1^{(0)} \prod_{i=1}^n P_1(x_i)} \end{aligned}$$

- b) Recall the basic Bayes' theorem.

$$\begin{aligned} \pi_0^{(n)}(X^n) &= \frac{\pi_0^{(0)} \prod_{i=1}^n P_0(x_i)}{\pi_0^{(0)} \prod_{i=1}^n P_0(x_i) + \pi_1^{(0)} \prod_{i=1}^n P_1(x_i)} \\ &= \frac{\pi_0^{(0)} 2^{-n(-\frac{1}{n} \sum_{i=1}^n \log P_0(x_i))}}{\pi_0^{(0)} 2^{-n(-\frac{1}{n} \sum_{i=1}^n \log P_0(x_i))} + \pi_1^{(0)} 2^{-n(-\frac{1}{n} \sum_{i=1}^n \log P_1(x_i))}} \\ &= \frac{1}{1 + \left(\frac{\pi_1^{(0)}}{\pi_0^{(0)}}\right) 2^{-n\left(\frac{1}{n} \sum_{i=1}^n \log \frac{P_0(x_i)}{P_1(x_i)}\right)}} = \frac{1}{1 + c 2^{-n\left(\frac{1}{n} \sum_{i=1}^n \log \frac{P_0(x_i)}{P_1(x_i)}\right)}}, \end{aligned}$$

where we let $c = \pi_1^{(0)}/\pi_0^{(0)}$. For any $\epsilon \in (0, D(P_0 \| P_1))$, there is an $N_\epsilon \in \mathbb{N}$ such that, for every $n \geq N_\epsilon$

$$\Pr\left\{\left|\frac{1}{n} \sum_{i=1}^n \log \frac{P_0(X_i)}{P_1(X_i)} - D(P_0 \| P_1)\right| < \epsilon\right\} = 1$$

$$\Pr \left\{ \frac{1}{1 + c2^{-n(D(P_0\|P_1) - \epsilon)}} < \pi_0^{(n)}(X^n) < \frac{1}{1 + c2^{-n(D(P_0\|P_1) + \epsilon)}} \right\} = 1$$

Hence, with probability 1, $\pi_0^{(n)}(X^n) \rightarrow 1$ as $n \rightarrow \infty$.

Following this derivation, for every $n \geq N_\epsilon$, with probability 1,

$$-\frac{1}{n} \log \left(\frac{c2^{-n(D(P_0\|P_1) - \epsilon)}}{1 + c2^{-n(D(P_0\|P_1) - \epsilon)}} \right) < -\frac{1}{n} \log \pi_1^{(n)}(X^n) < -\frac{1}{n} \log \left(\frac{c2^{-n(D(P_0\|P_1) + \epsilon)}}{1 + c2^{-n(D(P_0\|P_1) + \epsilon)}} \right)$$

Furthermore,

$$\begin{aligned} -\frac{1}{n} \log \left(\frac{c2^{-n(D(P_0\|P_1) + \epsilon)}}{1 + c2^{-n(D(P_0\|P_1) + \epsilon)}} \right) &= D(P_0\|P_1) + \epsilon - \frac{\log c}{n} + \frac{1}{n} \log (1 + c2^{-n(D(P_0\|P_1) + \epsilon)}) \\ &\leq D(P_0\|P_1) + \epsilon - \frac{\log c}{n} + \frac{1}{n} \log (1 + c2^{-N_\epsilon(D(P_0\|P_1) + \epsilon)}) \\ &= D(P_0\|P_1) + \epsilon + O\left(\frac{1}{n}\right) \\ -\frac{1}{n} \log \left(\frac{c2^{-n(D(P_0\|P_1) - \epsilon)}}{1 + c2^{-n(D(P_0\|P_1) - \epsilon)}} \right) &= D(P_0\|P_1) - \epsilon - \frac{\log c}{n} + \frac{1}{n} \log (1 + c2^{-n(D(P_0\|P_1) - \epsilon)}) \\ &\geq D(P_0\|P_1) - \epsilon - \frac{\log c}{n} + \frac{1}{n} \log (1 + c2^{-N_\epsilon(D(P_0\|P_1) - \epsilon)}) \\ &= D(P_0\|P_1) - \epsilon + O\left(\frac{1}{n}\right) \end{aligned}$$

Hence, with probability 1,

$$D(P_0\|P_1) - \epsilon \leq \lim_{n \rightarrow \infty} -\frac{1}{n} \log \pi_1^{(n)}(X^n) \leq D(P_0\|P_1) + \epsilon,$$

where $\epsilon \in (0, D(P_0\|P_1))$ arbitrarily. Hence,

$$\Pr \left\{ \lim_{n \rightarrow \infty} -\frac{1}{n} \log \pi_1^{(n)}(X^n) = D(P_0\|P_1) \right\} = 1$$

Grading Policy:

- $\pi_0^{(n)}(x^n)$ [2] and $\pi_1^{(n)}(x^n)$ [2].
- Application of SLLN and argumnet for almost sure convergence [4], convergence for $\pi_0^{(n)}(X^n)$ [2], convergence for $-(1/n) \log \pi_1^{(n)}(X^n)$ [2].

3. Minimizing information divergence) [22]

- Let $\mathcal{P}(\mathbb{N})$ denote the collection of all probability distributions over \mathbb{N} and $G(p) \in \mathcal{P}(\mathbb{N})$ be a geometric distribution with parameter $p \in (0, 1)$:

$$X \sim G(p) \iff \Pr\{X = n\} = (1 - p)p^{n-1}, \quad n \in \mathbb{N} = \{1, 2, \dots\}.$$

Under the constraint that $\mathbf{P} \in \mathcal{P}(\mathbb{N})$ and $\mathbb{E}_{X \sim \mathbf{P}}[X] = \sum_{x=1}^{\infty} x\mathbf{P}(x) = \mu > 1$, find the minimum value of $D(\mathbf{P} \parallel \mathbf{G}(p))$ and a minimizing distribution. [12]

- b) For m discrete probability distributions $\mathbf{P}_1, \mathbf{P}_2, \dots, \mathbf{P}_m$ with the same support \mathcal{X} , consider the following minimization problem:

$$\min_{\mathbf{Q} \in \mathcal{P}(\mathcal{X})} \sum_{i=1}^m D(\mathbf{P}_i \parallel \mathbf{Q}),$$

where $\mathcal{P}(\mathcal{X})$ denotes the collection of probability distributions over \mathcal{X} . Find a minimizer to the above problem. [10]

Solution:

- a) Expanding the Kullback-Leibler divergence over countable alphabet gives:

$$\begin{aligned} D(\mathbf{P} \parallel \mathbf{G}(p)) &= \sum_{x=1}^{\infty} \mathbf{P}(x) \log \frac{\mathbf{P}(x)}{(1-p)p^{x-1}} \\ &= \sum_{x=1}^{\infty} \mathbf{P}(x) \log \mathbf{P}(x) - \log(1-p) - \log p \sum_{x=1}^{\infty} \mathbf{P}(x)(x-1) \\ &= -H(\mathbf{P}) - \log(1-p) - (\mu-1) \log p \\ &\geq -\mu h_b\left(\frac{1}{\mu}\right) - \log(1-p) - (\mu-1) \log p, \end{aligned}$$

where the inequality comes from Problem 1 of Homework 2, and $\mathbf{P}^* = \mathbf{G}(1-\mu^{-1})$ is the minimizing distribution.

- b) Let $\bar{\mathbf{P}} = \frac{1}{m} \sum_{i=1}^m \mathbf{P}_i$, $\bar{\mathbf{P}} \in \mathcal{P}(\mathcal{X})$.

$$\forall \mathbf{Q} \in \mathcal{P}(\mathcal{X}),$$

$$\begin{aligned} &\frac{1}{m} \sum_{i=1}^m D(\mathbf{P}_i \parallel \mathbf{Q}) - \frac{1}{m} \sum_{i=1}^m D(\mathbf{P}_i \parallel \bar{\mathbf{P}}) \\ &= \frac{1}{m} \sum_{i=1}^m \sum_{x \in \mathcal{X}} \mathbf{P}_i(x) \log \frac{\bar{\mathbf{P}}(x)}{\mathbf{Q}(x)} \\ &= \sum_{x \in \mathcal{X}} \left(\frac{1}{m} \sum_{i=1}^m \mathbf{P}_i(x) \right) \log \frac{\bar{\mathbf{P}}(x)}{\mathbf{Q}(x)} \\ &= \sum_{x \in \mathcal{X}} \bar{\mathbf{P}}(x) \log \frac{\bar{\mathbf{P}}(x)}{\mathbf{Q}(x)} \\ &= D(\bar{\mathbf{P}} \parallel \mathbf{Q}) \geq 0. \end{aligned}$$

Hence, $\bar{\mathbf{P}}$ is a minimizer.

Grading Policy

- a) Find the tight lower bound [8] and find minimizer to achieve minimum [4].
- b) Specify the minimizer [3] and justify it [7].