## Graph Theory HW4

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November 7, 2023

**Problem 2.** Let  $uv \in E(G)$ , H := G - uv,  $N_H(w) := \{x | wx \in E(H)\}$ .

Let  $S := \{p | p \text{ is a Hamiltonian path in } H \text{ with } u \text{ being one of the endpoint } \}, p_i :=$ the *i*-th vertex on the path p  $(p = p_1(=u)p_2p_3\cdots p_n)$ , and  $T_w := \{p | p \in S, p_n = w\}$ .

Let  $f: S \to 2^S$  (the power set of S), where  $f(p) := \{p_1 p_2 \cdots p_i p_n p_{n-1} \cdots p_{i+1} | p_i p_n \in E(H), i \neq n-1\}.$ 

$$\Rightarrow |f(p)| = |\{i|p_ip_n \in E(H), i \neq n-1\}| = \deg_H(p_n) - 1.$$

Suppose that  $q \in f(p)$ , where  $q = p_1 p_2 \cdots p_i p_n p_{n-1} \cdots p_{i+1}$ .

$$\Rightarrow p = q_1 q_2 \cdots q_i q_n q_{n-1} \cdots q_{i+1}.$$

$$\therefore q_i q_n = p_i p_{i+1} \in E(H).$$

 $\therefore p \in f(q)$  by the definition of f and Hamiltonian path.

$$\therefore \forall p,q \in S, \ p \in f(q) \iff q \in f(p).$$

Consider a graph H' with vertex set S, and  $pq \in E(H') \iff p \in f(q) (\iff q \in f(p))$ .

$$\deg_{H'}(p) = |f(p)| = \deg_H(p_n) - 1.$$

The number of Hamiltonian cycles of G that contain  $uv \equiv$  the number of Hamiltonian paths that starts at u and ends at v in  $H \equiv \sum_{n=v} 1^{\deg_H(v) \equiv \deg_G(v) - 1 \equiv 0 \pmod{2}}$ 

$$\sum_{p_n=v} (\deg_H(v)-1)^{\deg_H(w)\equiv \deg_G(v)\equiv 1 \pmod{2}} \stackrel{(\text{mod } 2) \text{ for } w\neq v}{\equiv} \sum_{w\in V} \sum_{p_n=w}^{p_n=v} (\deg_H(w)-1) \equiv \sum_{p\in S} (\deg_H(p_n)-1) = \sum_{p\in S} (\deg_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes this problem}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject finishes}}{\equiv} (\log_H(v)-1) \stackrel{(\text{mod } 2) \text{ subject$$

1)  $\equiv \sum_{p \in S} \deg_{H'}(p) \equiv 0 \pmod{2}$ , which finishes this problem.

**Problem 3.** Let  $m := \frac{n}{2}$ .

Consider the graph  $K_{m+1,m-1}$  (where  $V(K_{m-1,m+1})$  has bipartition  $A \cup B$ , and

|A| = m + 1), it does not have a perfect matching because  $|N_S| < |S|$  for S = A by Hall theorem.

$$d(n) > \delta(K_{m+1,m-1}) = m-1.$$

Claim: d(n) = m.

Proof: Let G be a graph with  $\delta(G) \geq m$ , and  $S \subseteq V(G)$ .

For 
$$|S| > m$$
,  $q(G \setminus S) \le |V(G \setminus S)| \le m - 1 \le |S|$ .

For  $|S| \leq m$ :

The size of a connected component in  $G \setminus S$  is at least  $\delta(G \setminus S) + 1$ , because every vertex is connected to at least  $\delta(G \setminus S)$  vertices.

 $\therefore \text{the number of connected components of } G \backslash S \leq \big\lfloor \frac{n-|S|}{\delta(G \backslash S)+1} \big\rfloor \stackrel{\delta(G \backslash S) \geq \delta(G)-|S| \geq m-|S|}{\leq}$ 

$$\left\lfloor \frac{2m - |S|}{m - |S| + 1} \right\rfloor.$$

 $\lfloor \frac{2m - |S|}{m - |S| + 1} \rfloor.$ Let  $f(x) := \frac{2m - x}{(m - x + 1)x}$  for  $x \in (0, m + 1)$  (so that (m - x + 1)x > 0).  $f(x) \le 1 \iff 2m - x \le (m - x + 1)x \iff x^2 - (m + 2)x + 2m \le 0 \iff -\infty \quad -(1)$ 

$$f(x) \le 1 \iff 2m - x \le (m - x + 1)x \iff x^2 - (m + 2)x + 2m \le 0 \iff$$

$$(x-m)(x-2) \le 0 \iff 2 \le x \le m. -(1)$$

$$\Rightarrow \frac{q(G \setminus S)}{|S|} \le \frac{\lfloor \frac{2m-|S|}{m-|S|+1} \rfloor}{|S|} \begin{cases} = \lfloor \frac{2m-1}{m} \rfloor = 1, & \text{if } x = 1\\ \frac{2m-|S|}{m} \le \frac{2m-|S|}{m} \le 1, & \text{if } x \ge 2 \end{cases}$$

$$\Rightarrow q(G \setminus S) \le |S|.$$

 $\therefore$  by Tutte theorem, G has a perfect matching, which finishes the proof.

**Problem 6.** Consider G with bipartition  $V(G) = A \cup B$ .

 $A = \{u_i | i \in [n]\}, B = \{v_i | i \in [n]\}, \text{ where } u_i \in A, v_j \in B \text{ are adjacent } \iff \text{ the } i\text{-the } i$ column of L does not contain j.

- $\therefore$  the size of a row is n, each element of [n] appears exactly once in every row.
- $\Rightarrow$  there are r js for all  $j \in [n]$ .
- $\therefore$  no column contains two js.
- $\Rightarrow$  deg $(v_i) = n$  the number of columns that contain j = n r.
- $\therefore$  every column contains r different elements.
- $\therefore \deg(u_i) = n r \forall i \in [n].$
- $\therefore G$  is k-regular.

By corollary 5.10 in class, G has a perfect matching, and suppose that M is a perfect

Author: 許博翔 B10902085 Teammate: 黃芊禕 B10902029 matching.

Let f(j) denote the index of the other endpoint of the edge containing  $v_j$  in M (that is,  $u_{f(j)}v_j \in M$ ).

Put j at the intersection of the f(j)-th column and the (r+1)-th row for all  $j \in [n]$ .

—(1)

By the definition of G, since  $u_{f(j)}v_j \in M \subseteq E(G)$ , j does not appears in the f(i)-th column of L.

Also, since M is a perfect matching, f is a bijection, which means the (r + 1)-th row contains each element in [n] exactly once.

 $\therefore$  (1) is a valid way to extend L.

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