

# **Transfer Principle Proof Tactic for Nonstandard Analysis**

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# Introduction

- What is nonstandard analysis (NSA)?
  - NSA extends the reals  $\mathbb{R}$  with new elements to form the “hyperreals”  $\mathbb{R}^*$
  - Hyperreals include standard reals, plus infinite and infinitesimal values
  - Hyperreals inherit many algebraic properties from the reals
- What is it good for?
  - Provides formalization of intuitive notions in calculus
  - Can express concepts like limits more simply, with fewer quantifiers

## Previous Work

- Fleuriot and Paulson, Mechanizing Nonstandard Real Analysis (2000)
  - Includes formalization of reals, hyperreals, and much real analysis in Isabelle/HOL
- My work presented here is based on theirs
  - Their formalization of reals can be reused
  - Their formalization of analysis may be adapted
- My contribution is a new way to define nonstandard types and operations on them, which provides more generality, more abstraction, and more automation

## **Part 1: Nonstandard Extensions of Types**

## Free Ultrafilters

- An ultrafilter  $\mathcal{U}$  is a set of sets of naturals
- Basically,  $\mathcal{U}$  specifies a relation that says when two countable sequences “mostly agree”:

$$(X \sim Y) \equiv \{n. X_n = Y_n\} \in \mathcal{U}$$

- For any ultrafilter  $\mathcal{U}$ , this relation is reflexive, symmetric, and transitive
- For a *free* ultrafilter  $\mathcal{U}$ , if two sequences mostly agree, they must match in infinitely many places (i.e.  $\mathcal{U}$  contains no finite sets)

# Hyperreal Numbers

- The hyperreals are defined by sequences of reals, modulo ( $\sim$ ) relation
- Standard reals are represented by constant sequences
- Algebraic operations ( $+$ ,  $-$ ,  $\times$ ,  $\div$ ,  $\sin$ ,  $\exp$ , etc.) are defined pointwise
- Predicates hold for sequences if they hold for “most” terms:

$$P^*(X) \equiv \{n. P(X_n)\} \in \mathcal{U}$$

- The sequence  $(1, 2, 3, 4, \dots)$  represents an infinite hyperreal

## Type Constructor `star`

- The previous quotient construction is not limited to reals
  - NSA traditionally makes use of hypercomplexes  $\mathbb{C}^*$ , hypernaturals  $\mathbb{N}^*$
  - Why stop there? The construction works the same way over any type!
- It is natural to define a type constructor for nonstandard extensions of types
  - For every Isabelle type `'a`, we have a nonstandard extension `'a star`
  - Constructor functions:
    - (surjective) `star_n :: (nat => 'a) => 'a star`
    - (injective) `star_of :: 'a => 'a star`

# Defining Functions on Nonstandard Types

- Old method: Define functions pointwise over sequences
  - Drawback: Have to prove that this respects equivalence relation ( $\sim$ )
  - Different proofs needed for each function arity
- New method: Use nonstandard extension of function space!
  - Combinator Ifun returns so-called “internal” functions  
`Ifun :: ('a => 'b) star => ('a star => 'b star) (infixl ★)`
  - Definition designed to satisfy the following property:  
`star_n F ★ star_n X = star_n ( $\lambda n. (F\ n)\ (X\ n)$ )`
  - Reasoning about equivalence relation is only needed once!



## Using Ifun

- We can use Ifun to lift functions of any arity

```
Ifun_of::('a => 'b) => 'a star => 'b star
```

```
Ifun_of f x == star_of f ★ x
```

```
Ifun2_of::('a => 'b => 'c) => 'a star => 'b star => 'c star
```

```
Ifun2_of f x y == star_of f ★ x ★ y
```

```
Ifun3_of::('a => 'b => 'c => 'd)
```

```
    => 'a star => 'b star => 'c star => 'd star
```

```
Ifun3_of f x y z == star_of f ★ x ★ y ★ z
```

## Defining Predicates with unstar

- Function unstar is the inverse of star\_of for booleans

```
unstar::bool star => bool
unstar b == (b = star_of True)
```

- Together with Ifun, we can lift predicates of any arity

```
Ipred_of::('a => bool) => 'a star => bool
Ipred_of p x == unstar (star_of p ★ x)
```

```
Ipred2_of::('a => 'b => bool) => 'a star => 'b star => bool
Ipred2_of p x y == unstar (star_of p ★ x ★ y)
```

## **Part 2: The Transfer Principle**

# What is the Transfer Principle?

- Transfer principle states that any first-order statement about reals is logically equivalent to a syntactically similar statement about hyperreals
- Transfer principle is a meta-mathematical theorem, since it quantifies over statements
- Compare with duality principle for statements about sets
- As a meta-mathematical theorem, we can't express the transfer principle as a theorem in Isabelle's logic
- We *can* write an algorithm for generating a proof for any instance of the transfer principle

## Transfer Principle Proof Tactic

- In Isabelle, the transfer tactic replaces a subgoal about nonstandard types with an equivalent subgoal about ordinary types
- The tactic produces a proof of the equivalence, using a syntax-directed procedure

# Conclusions

- Generalization
  - Can make nonstandard extensions of any type
  - Polymorphic functions over nonstandard types are now possible
- Abstraction
  - Combinators `star_of`, `Ifun`, and `unstar` are enough to define nearly all operations on nonstandard types
- Automation
  - The transfer tactic can quickly and effortlessly prove many theorems about nonstandard types