



# Dynamic Price Floors Final Deliverable

05 • 04 • 18

# Team



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# Agenda

- 01 Introduction
- 02 Strategies
- 03 Results
- 04 Hand-off
- 05 Q & A

# 01

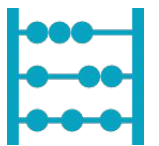
# Introduction

- a. Recap
- b. Methodology

## Data



We had 5 days of data, containing roughly one billion RTB auctions



We investigated several strategies for setting price floors based off of incoming bids

```
vi_cnt : 0, r_timestamp : 2018-02-13T15:15:12.084622Z", "r_cnt": 1, "r_num_ads_thir
ice": 0.00191}], "bid_requests": [21943,
6:15:12.084622Z", "r_cnt": 1, "r_num_ads_thir
46.589004Z", "r_cnt": 1, "r_num_ads_thir
55.853514Z", "r_cnt": 1, "r_num_ads_thir
4586, 20578, 24275, 20579, 20868, 24276,
ce": 0.000219999, "buyer_seat_id": "52840
ice": 0.0026141, "buyer_seat_id": ""}, {"
"vi_cnt": 0, "r_timestamp": "2018-02-13T
ice": 0.00191}], "bid_requests": [22041,
ice": 0.00191}], "bid_requests": [22041,
:08:12.757231Z", "r_cnt": 1, "r_num_ads_t
ice": 9.59e-05, "buyer_seat_id": ""}, {"i
ce": 0.00034054, "buyer_seat_id": "acc-57
:11.204947Z", "r_cnt": 1, "r_num_ads_thir
:34.630270Z", "r_cnt": 1, "r_num_ads_thir
tamp": "2018-02-13T15:47:19.063809Z", "r
16:40:56.391823Z", "r_cnt": 1, "r_num_ads
27:48.790074Z", "r_cnt": 1, "r_num_ads_th
ice": 0.00043}], "bid_requests": [19587,
ice": 0.0069166, "buyer_seat_id": ""}, {"
```



## Assumptions & Caveats



We assumed we had access to every possible bid (if there was no price floor)



We were unable to calculate true revenue or even a baseline revenue, but we were able to compare strategies in a meaningful way



The vast majority of bid requests were not followed by bids (above the unknown price floor)



```
vi_cnt": 0, "r_timestamp": "2018-02-13T15:15:12.084622Z", "r_cnt": 1, "r_num_ads_thir": 46.589004Z", "r_cnt": 1, "r_num_ads_thir": 55.853514Z", "r_cnt": 1, "r_num_ads_thir": 4586, 20578, 24275, 20579, 20868, 24276, "ce": 0.000219999, "buyer_seat_id": "52840", "ce": 0.0026141, "buyer_seat_id": ""}, {"vi_cnt": 0, "r_timestamp": "2018-02-13T15:08:12.757231Z", "r_cnt": 1, "r_num_ads_thir": 9.59e-05, "buyer_seat_id": ""}, {"ce": 0.00034054, "buyer_seat_id": "acc-57", "ce": 11.204947Z", "r_cnt": 1, "r_num_ads_thir": 34.630270Z", "r_cnt": 1, "r_num_ads_thir": 16:40:56.391823Z", "r_cnt": 1, "r_num_ads_thir": 27:48.790074Z", "r_cnt": 1, "r_num_ads_thir": 0.00043}], "bid_requests": [19587, "ce": 0.0069166, "buyer_seat_id": ""}, {"
```

## Simulator



We tuned each strategy on 2/11-13, and tested on 2/14-15



Each strategy must return a price floor based on a set of input features



The bids *above this price floor* are then revealed to the strategy

```
@abstractmethod
def calculate_price_floor(self, input_features):
    pass
```

```
@abstractmethod
def process_line(self, line, input_features, bids):
    pass
```

```
INPUT_FEATURES = {
    'geo_dma_code', 'geo_timezone', 'ua_name', 'ua_device',
    'r_num_ads_third_party', 'created_at', 'r_num_ads_requested',
    'rate_metric', 'geo_code', 'r_timestamp', 'ua_os',
    'ua_device_type', 'ua_ip'
}
```

# 02

# Pricing Strategies

- a. OneShot
- b. Weighted Running Average
- c. Optimization
- d. Vowpal Wabbit
- e. Random Forest
- f. Running Average



## Summary



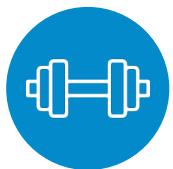
Dynamically adjust price floor based on previous bid



$\lambda_h$ : Price floor too high

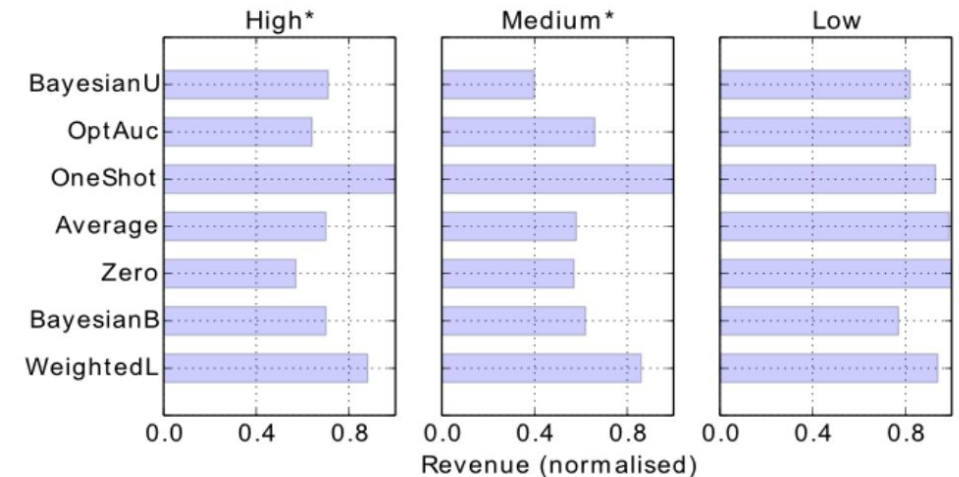
$\lambda_e$ : Price floor between max and second bid (ideal)

$\lambda_l$ : Price floor too low



Responsively evolves over time to produce results

$$\begin{cases} \alpha(t+1) = (1 - \epsilon^t \lambda_h) \alpha(t) & \text{if } \alpha(t) > b_1(t) \\ \alpha(t+1) = (1 + \epsilon^t \lambda_e) \alpha(t) & \text{if } b_1(t) \geq \alpha(t) \geq b_2(t) \\ \alpha(t+1) = (1 + \epsilon^t \lambda_l) \alpha(t) & \text{if } b_2(t) > \alpha(t) \end{cases}$$



## Improvements and Experiments



Tuned hyperparameters for low overshoot rate



Adaptive price floor ceiling



Overshooting adjustments ( $\lambda_h$ )



```
def update(self, bids, price_floor):
    first, second = self.max2(bids)
    revenue = self.calculate_revenue_helper(first, second, price_floor)
    diff = self.calculate_differential(first, second, price_floor)
    if len(bids) >= self.oneshot_min_n:
        self.oneshot(first, second, diff)
    else:
        if len(self.revenues) == 5:
            self.revenues.pop(0)
        self.revenues.append(revenue)
    return revenue
```

## Results / Numbers

| Hyperparameter     | Max revenue | Min overshoot | Mixed |
|--------------------|-------------|---------------|-------|
| $\lambda_h$        | 0.02        | 0.60          | 0.05  |
| $\lambda_e$        | 0.90        | 0.10          | 0.60  |
| $\lambda_l$        | 0.88        | 0.30          | 0.70  |
| Number of clusters | 84          | 4             | 80    |
| Revenue            | 250.6       | 14.6          | 153.5 |
| Overshoot          | 90.6%       | 0.15%         | 80.7% |

Trade off between high revenue and low overshoot rate

## Takeaways

Biased towards revenue in lieu of overshoot

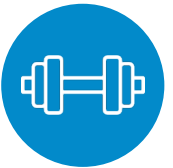
Multishot method is extremely promising for quick results



## Summary and Recap



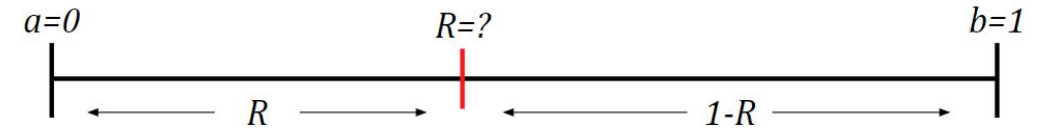
Derived optimal price floor for two uniformly distributed bidders



Found optimal  $R=0.5$ , but tried weighing lower bid more to decrease overshoot rate



Split averages by features, keeping separate averages for site\_id, pub\_network\_id, etc.



- no sale if both bids below  $R$  - happens with probability  $R^2$  and revenue=0
- sale at price  $R$  if one bid above reserve and other below - happens with probability  $2(1 - R)R$  and revenue  $= R$
- sale at second highest bid if both bids above reserve - happens with probability  $(1 - R)^2$  and revenue  $= E[\min v_i | \min v_i \geq R] = \frac{1+2R}{3}$
- Expected revenue  $= 2(1 - R)R^2 + (1 - R)^2 \frac{1+2R}{3}$
- Expected revenue  $= \frac{1+3R^2-4R^3}{3}$
- Maximizing:  $0 = 2R - 4R^2$ , or  $R = \frac{1}{2}$ .

## Simulating Lost Bids



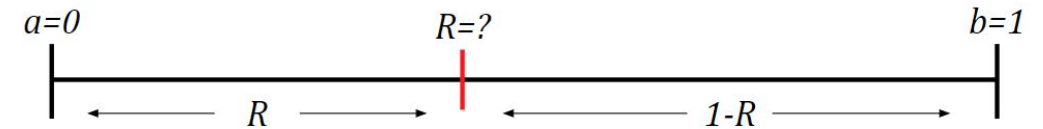
In real-life auction, bids below the price floor are lost



Lost bid information causes average strategy to strictly increase to unreasonable levels



Simulate lost bidders by deriving expected ratio between first and second bid using original assumptions



Expected high bid  $\approx 0.66$

Expected low bid  $\approx 0.33$

Ratio  $= 0.66/0.33 = 0.5$

**Simulated low bid**  $= 0.5 \cdot \text{high bid}$



## Implementation



### Parameters

- window: number of past auctions to incorporate into average
- weight: shift average toward low or high bidder
- sim\_bid: multiplier for generating simulated second bid using first
- param\_id: feature used to track separate running averages

```
class AverageSingleID(sim.Simulator):  
    """  
    Only works for input features with a single value (e.g. works with site_id,  
    pub_network_id, but not bid_requests).  
    """  
  
    def __init__(self, window, weight, sim_bid, param_id, **kwargs):  
        sim.Simulator.__init__(self, **kwargs)  
        self.window = window  
        self.weight = weight  
        self.simBid = sim_bid  
        self.id = param_id  
        self.averages = {}  
        self.counts = {}  
  
    def calculate_price_floor(self, input_features):  
        key = input_features[self.id]  
        if key in self.averages:  
            return self.averages[key]  
        else:  
            return DEFAULT_FLOOR  
  
    def process_line(self, line, input_features, bids):  
        key = input_features[self.id]  
        if key in self.averages:  
            reserve = self.averages[key]  
        else:  
            # If site_id has never been seen before and auction has 0 bids  
            reserve = DEFAULT_FLOOR  
  
        low, high = 0, 0  
        if len(bids) == 0:  
            high = reserve  
            low = high * self.simBid  
        elif len(bids) == 1:  
            high = bids[0]  
            low = high * self.simBid  
        else:  
            high = bids[0]  
            low = bids[1]  
  
        weighted_avg = low * self.weight + high * (1 - self.weight)  
        if key in self.averages:  
            n = min(self.counts[key] + 1, self.window)  
            self.averages[key] = self.averages[key] * (n - 1) / n + weighted_avg * 1.0 / n  
            self.counts[key] += 1  
        else:  
            self.averages[key] = weighted_avg
```

# Weighted Running Average



## Results

window=500, weight=0.6, sim\_bid=0.3

```
-----
GT site_id
-----
Total Revenue: 424276.93237573805
Auction Count: 518420532
Auction Count (non-null): 300064887
Price Floor Engaged (non-null): 30.63%
Price Floor Too High (non-null): 64.70%
Average Revenue: 0.0008184030264753057
Average Revenue (not-null): 0.0014139506178733204
Average Bid Count: 0.954231081650138
Average Bid Count (non-null): 1.648620036638942
Average Bid Amount: 0.0028150392288430067
Average Bid Amount (non-null): 0.004863528516142781
Average Price Floor: 0.002432278846190381
```

Feature: site\_id  
Average Revenue: \$0.0014  
Engage / Too High: 30% /  
64%

```
-----
GT pub_network_id
-----
Total Revenue: 227656.77194918026
Auction Count: 518420532
Auction Count (non-null): 300064887
Price Floor Engaged (non-null): 25.24%
Price Floor Too High (non-null): 69.67%
Average Revenue: 0.00043913533106204264
Average Revenue (not-null): 0.000758691809045889
Average Bid Count: 0.954231081650138
Average Bid Count (non-null): 1.648620036638942
Average Bid Amount: 0.0028150392288430067
Average Bid Amount (non-null): 0.004863528516142781
Average Price Floor: 0.0028786333704231702
```

Feature: pub\_network\_id  
Average Revenue: \$0.0007  
Engage / Too High: 25% /  
69%



## Underlying Model

$X, Y$  first and second highest bid distributions

$$r^* = \arg \max_r E(\text{rev}_r(X, Y))$$



$$\text{rev}_r(X, Y) = \begin{cases} 0 & : X < r \\ r & : Y < r \leq X \\ Y & : r \leq Y \end{cases}$$

set reserve price to maximize expected revenue

## Example

$$\begin{array}{lll} X_1 \sim \text{uniform}(0, 1) & \text{assume uniformly} & X = \max(X_1, X_2) \\ X_2 \sim \text{uniform}(0, 1) & \text{distributed bidders} & Y = \min(X_1, X_2) \end{array}$$

$$\begin{aligned} E(\text{rev}_r(X, Y)) &= r^2(0) + 2r(1-r)(r) + (1-r)^2 \frac{1+2r}{3} \\ &= \frac{1 + 3r^2 - 4r^3}{3} \end{aligned}$$

[function to maximize over  $r$ ]

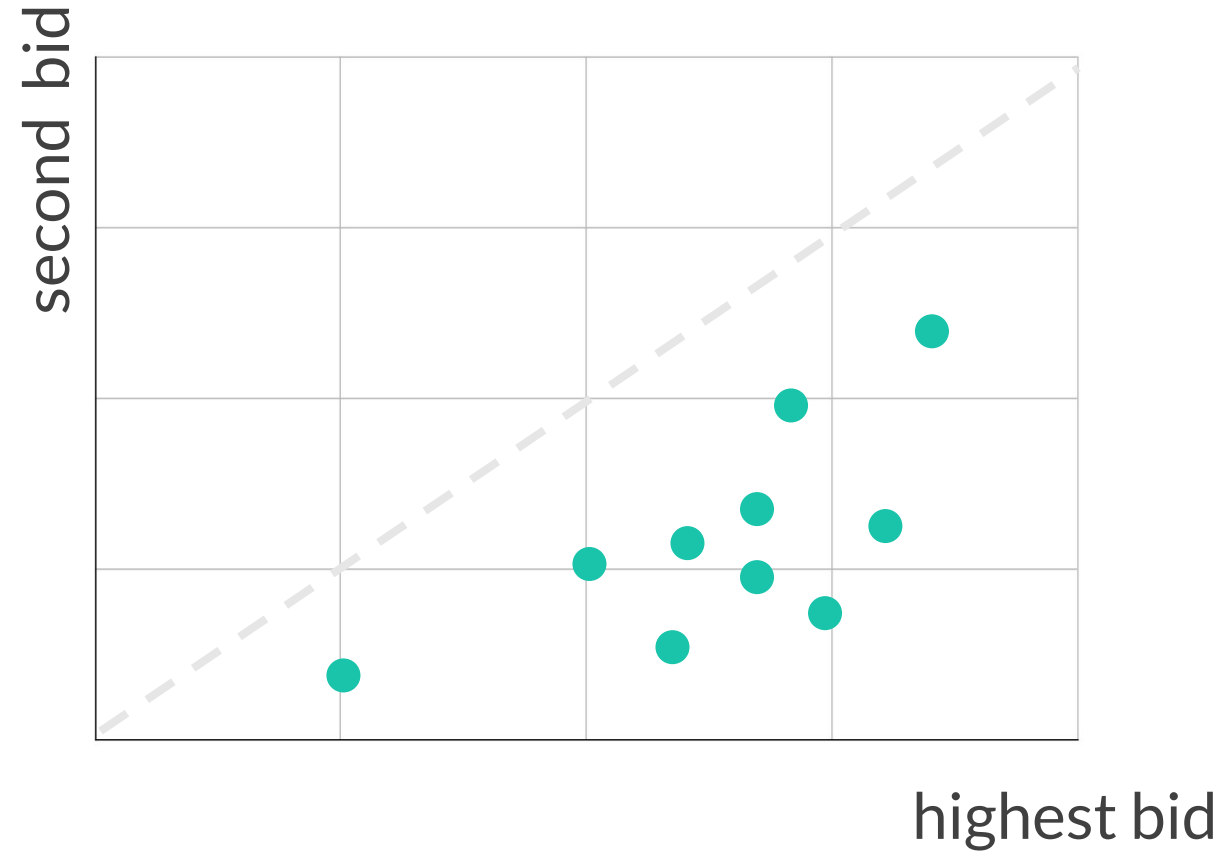
## Problem



what are bid distributions like?



more data, less structure principle





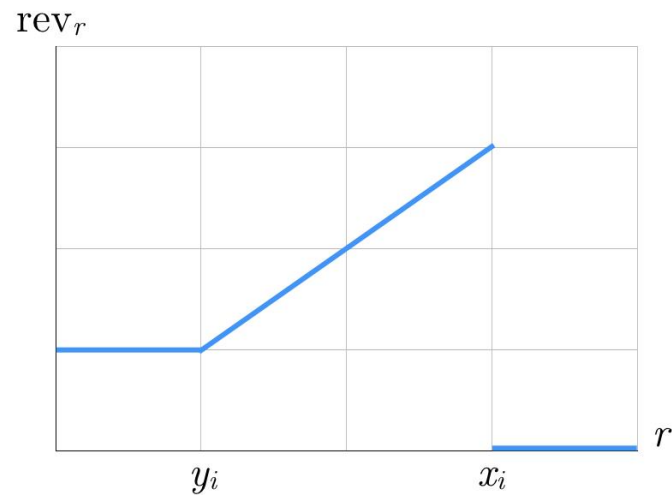
## Solution

● ● ● ● ● ●  $A_n = \{(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)\}$   
empirical joint distribution

$$r^* \approx \arg \max_r \sum_{t=1}^n \text{rev}_r(x_t, y_t)$$

finding the maximum of this function  $f(r; A_n)$

## Linear Heuristic

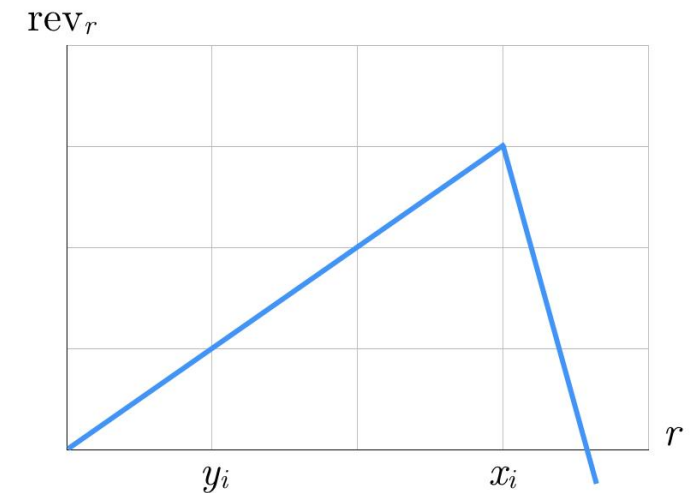


revenue function

$$\begin{aligned} \max \quad & \sum z_i \\ \text{s.t.} \quad & z_i \leq r \\ & z_i \leq -\delta(r - x_i) + x_i \\ & r \geq 0 \end{aligned}$$

return  $\gamma \cdot r^*$

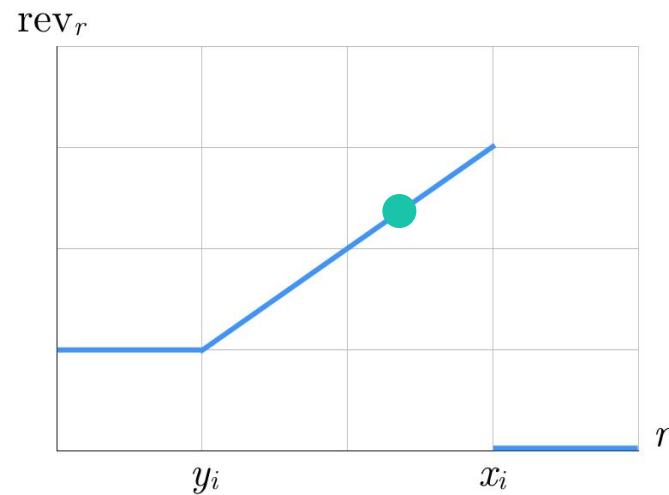
algorithm  $[\gamma, \square, n]$



-approximation

## Brute Force

$$\frac{d}{dr} \sum_{i=1}^n$$



try all drop points

algorithm[ $\gamma, n$ ]: return  $\gamma \cdot \arg \max_{x_i} \sum_{t=1}^n rev_r(x_t, y_t)$

## Tuning

linear heuristic[ $\gamma$ ,  $\square$ ,  $n$ ]: Fix  $n$  as 30,  
find  $\gamma$  and  $\square$  separately

brute force[ $\gamma$ ,  $n$ ]: Fix  $n$  as 100, find  
 $\gamma$  by varying



test on three files, best  
algorithms:

1. linear heuristic[1, -2, 30]
2. brute force[1, 100]

## Results

|                  | Period    | Price Floor Engaged<br>(non-null auctions) | Price Floor Too High<br>(non-null auctions) | Average Revenue<br>(non-null auctions) | Average Price Floor | Latency  |
|------------------|-----------|--|---|--|---------------------|----------|
| linear heuristic | 2 minutes | 6.25%                                      | 93.36%                                      | \$3.00 CPM                             | \$48.34 CPM         | 25.6 ms  |
| brute force      | 6 minutes | 7.79%                                      | 92.20%                                      | \$3.69 CPM                             | \$47.47 CPM         | 8.35 ms  |
| weighted average | 24 hours  | 5.69%                                      | 94.31%                                      | \$2.11 CPM                             | \$39.36 CPM         | 0.004 ms |



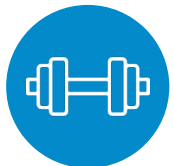
## Recap



Predict optimal price floor using bidder features



Strategy: machine learning framework to develop model



Use multipliers to get as close to highest bid as possible

```
0.001910 | ua_device_iPhone site_id_11257 geo_country_  
0.000020 | ua_device_iPhone site_id_11257 geo_country_  
0.048544 | ua_device_Microsoft_RM-1092 site_id_17615 g  
0.002515 | ua_device_Other site_id_17495 geo_country_c  
0.000314 | ua_device_Other site_id_10361 geo_country_c  
0.000101 | ua_device_Other site_id_31978 geo_country_c  
0.000404 | ua_device_Other site_id_2512 geo_country_co  
0.000278 | ua_device_Other site_id_22719 geo_country_c  
0.000248 | ua_device_iPhone site_id_1978 geo_country_c  
0.001456 | ua_device_Other site_id_3963 geo_country_co
```



## Implementation



Used Vowpal Wabbit: fast ML framework



### Parameters

- 1-pass model: campaign\_id, site\_id, zone\_id
- 1-5 pass models: hashed features
- Trained on Feb 11-13



$\text{price\_floor} = \text{multiplier} * b_1 \text{ prediction}$



**VOWPAL WABBIT**

```
prepared_line = sim.utils.prepare_line(line)
bids, input_features = prepared_line['bids'], prepared_line['input_features']
price_floor = self.calculate_price_floor(input_features, prepared_line, line_num,
if price_floor == None:
    break
price_floor *= multiplier
self.stats.process_line(bids, input_features, price_floor)
```

## Results



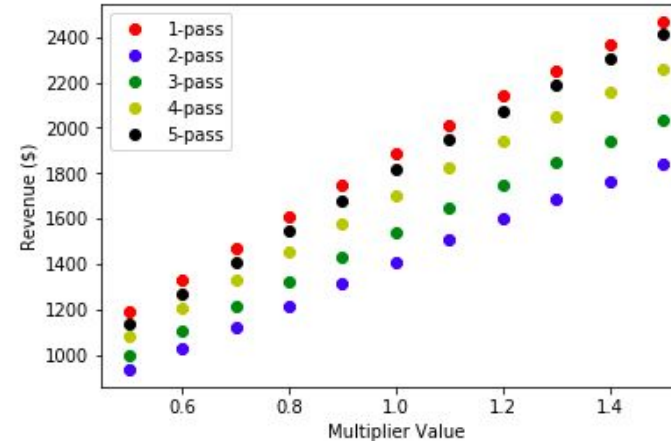
Tested on Feb 14, 12 AM  
Multipliers  $\in [0.5, 1.5]$   
1-5 pass models



Higher multipliers  $\in [2.0, 4.0]$   
1-pass model



Optimum multiplier  $\approx 3$   
Revenue (Feb 14-15): \$223,703.77  
Engagement: 43.89%



| Multiplier | Revenue (\$) |
|------------|--------------|
| 2.0        | 11,183.74    |
| 2.5        | 11,694.98    |
| 3.0        | 11,789.79    |
| 3.5        | 11,246.78    |
| 4.0        | 10,333.79    |

## Takeaways



### Disadvantages:

- Slower than other strategies
- Assumptions made
- Must train offline

### Advantage:

- Great potential with finely tuned model
- High price floor engagement
- Can be trivially parallelized

## Classification



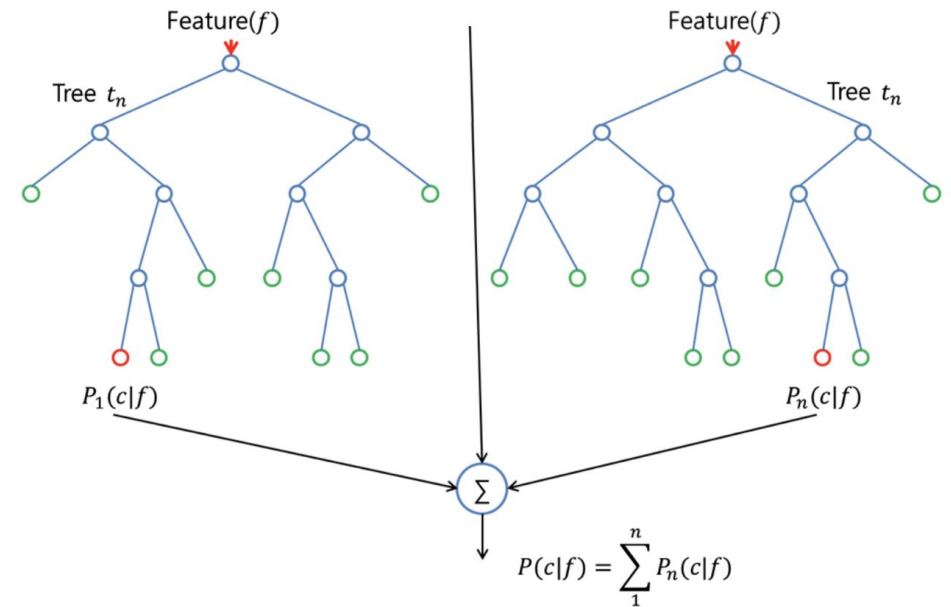
Bids are distributed tri-modally



k-means bins from 1 day of data



RF selected one bin for all auctions





## Regression



Labels derived from game theory



State space is large



83% of predictions above optimal



## Experiments



Global running average



Running average for each site\_id or publisher\_id



Include auctions with no bids?



## Results

Auction Count: 518420532  
Auction Count (non-null): 300064887  
Average Bid Count: 0.954231081650138  
Average Bid Count (non-null): 1.648620036638942  
Average Bid Amount: 0.0028150392288430067  
Average Bid Amount (non-null): 0.004863528516142781

### Running Average Without Zeros

Total Revenue: 485298.0741042547  
Price Floor Engaged (non-null): 7.73%  
Price Floor Too High (non-null): 92.27%  
Average Revenue: 0.0009361089003015311  
Average Revenue (not-null): 0.0016173104389393441  
Average Price Floor: 0.02089341090786549  
Time taken: 1598.51 seconds (0.003 milliseconds per auction)

### Running Average Including Zero-Bids

Total Revenue: 633231.2406333139  
Price Floor Engaged (non-null): 5.69%  
Price Floor Too High (non-null): 94.31%  
Average Revenue: 0.0012214625030192937  
Average Revenue (not-null): 0.0021103143622176425  
Average Price Floor: 0.039357196287724196  
Time taken: 2126.88 seconds (0.004 milliseconds per auction)

### Per-Site Average

**Total Revenue: 585769.6258170931**  
**Price Floor Engaged (non-null): 35.06%**  
**Price Floor Too High (non-null): 55.92%**  
Average Revenue: 0.0011299120880827558  
Average Revenue (not-null): 0.0019521431903396717  
Average Price Floor: 0.003797489056934269  
Time taken: 1838.94 seconds (0.004 milliseconds per auction)

### Per-Site Average Including Zero-Bids

**Total Revenue: 823332.5974897781**  
**Price Floor Engaged (non-null): 6.11%**  
**Price Floor Too High (non-null): 93.88%**  
Average Revenue: 0.0015881558438927302  
Average Revenue (not-null): 0.0027438485246352004  
Average Price Floor: 0.17501760124503202  
Time taken: 2194.64 seconds (0.004 milliseconds per auction)

# 03

## Results

- a. Comparison
- b. Recommendations

# Comparison of Best-Performing Strategies

| Strategy                               | Revenue   | Price Floor Engaged<br>(non-null auctions) | Price Floor Too High<br>(non-null auctions) | Average Revenue<br>(non-null auctions) | Average Price Floor | Latency  |
|--|-----------|--|---|--|---------------------|----------|
| Running Average (per site, with nulls) | \$823,332 | 6.11%                                      | 93.88%                                      | \$2.74 CPM                             | \$175.02 CPM        | 0.004 ms |
| Running Average (grouped by site)      | \$585,769 | 35.06%                                     | 55.92%                                      | \$1.95 CPM                             | \$3.79 CPM          | 0.004 ms |
| OneShot                                | \$339,306 | 5.27%                                      | 94.71%                                      | \$1.13 CPM                             | \$29.85 CPM         | 0.017 ms |
| Weighted Average (grouped by site)     | \$424,276 | 30.63%                                     | 64.70%                                      | \$0.82 CPM                             | \$2.43 CPM          | 0.002 ms |
| Vowpal Wabbit                          | \$223,703 | 43.89%                                     | 25.92%                                      | \$0.75 CPM                             | \$8.93 CPM          | 0.156 ms |
| No Price Floor                         | \$49,862  | 61.56%                                     | 0.00%                                       | \$0.17 CPM                             | \$0.00 CPM          | 0.001 ms |

Results are from 518 million auctions from 02/13/18 12:00 AM - 2/15/18 12:00AM. 300 million auctions received at least one visible bid. 120 million auctions received at least two visible bids. Latency varies by machine, but results show that these methods have good computational performance.



## Implementation Suggestions



Collect a testing dataset that disables price floors and can produce accurate revenue estimates



Implement A/B testing or similar to compare approaches



All strategies are promising, and will behave differently towards a real bid distribution



# 04

## Hand-off

- a. Materials
- b. Next Steps



## Report

### Dynamic Price Floors

CodeBase

University of California, Berkeley

May 4, 2018

### Introduction

Polymorph provides publishers the opportunity to set a price floor to protect



Data Processing

Machine Learning

Takeaways

Hand-off

Q & A

## Code

|                     |  |                |
|---------------------|--|----------------|
| exploratory         | Rename directory                               | 2 months ago   |
| gametheory          | site_id and pub_network_id results             | 12 hours ago   |
| linear_optimization | Make tuning even simpler                       | 11 hours ago   |
| oneshot             | cleaned up oneshot directory, edited tuning.py | 12 minutes ago |
| randomForest        | Add tensorflow version of random forest        | 2 months ago   |
| randomForestModel   | Add RFSimulator                                | a month ago    |
| results             | cleanup  | 2 hours ago    |
| running_average     | move running average to new package            | 3 hours ago    |
| scripts             | update file names                              | 9 days ago     |
| simulator           | add per-auction time                           | 2 days ago     |
| vwprediction        | Delete polytest                                | 2 days ago     |



## Maintenance



We will be available for two additional weeks to answer questions, fix issues, and help you reproduce our results or implement strategies!



The background of the left half of the slide is a solid blue color. Overlaid on this blue background is a faint, light-blue pattern of various scientific and technical sketches. These sketches include a periodic table of elements, a molecular structure, a globe, a DNA double helix, a bar chart, a line graph, a clock face, and various mathematical symbols and formulas. The sketches are drawn in a simple, hand-drawn style.

# 05

# Q & A

