


Lemma 0.1 (ultrafilter lemma) Let S be a set. Every filter on S is contained in an ultrafilter.

Corollary 0.2 (nonprincipal ultrafilter) For any infinite set S , there is an ultrafilter \mathcal{U} on S that is not generated by a singleton. Such a \mathcal{U} is nonprincipal.

Proof. Consider the cofinite filter on S :

$$F = \{X \subseteq S : |X| \geq \aleph_0\}$$

Then F is contained in a nonprincipal ultrafilter. 

Definition 0.3 (ultrapower) Suppose M_i is a model over language \mathcal{L} for each $i \in I$. Let \mathcal{U} be a nonprincipal ultrafilter on I . Consider

$$\prod M_i / \mathcal{U}$$

where quotienting by \mathcal{U} is understood to mean quotienting by the equivalence relation \sim where

$$(x_i)_{i \in I} \sim (y_i)_{i \in I}$$

iff the set of indices i such that $x_i = y_i$ is in \mathcal{U} . (Equality by vote).

We can make $\prod M_i / \mathcal{U}$ into a model over \mathcal{L} . Interpret the constant symbol c as the equivalence class of

$$(c|_{M_i})_{i \in I}$$

with $c|_{M_i}$ indicating interpretation (of c) in M_i .

Similarly, interpret the function $f(x)$ by

$$f((x_i)_{i \in I}) = [(f(x_i)|_{M_i})]$$

where $[-]$ denotes ‘equivalence class of’.

Finally, interpret the predicate $\phi(x)$ by

$$\phi((x_i)_{i \in I}) \iff \{i \in I : \phi(x_i)\} \in \mathcal{U}$$

which is truth by vote. Because \mathcal{U} is an ultrafilter, the law of the excluded middle will in fact hold as either

$$\{i \in I : p|_i \text{ holds}\} \in \mathcal{U}$$

or

$$\{i \in I : (\neg p)|_i \text{ holds}\} = \mathcal{U} - \{i \in I : p|_i \text{ holds}\} \in \mathcal{U} \in \mathcal{U}$$

Theorem 0.4 (compactness theorem) If every finite subtheory of Φ has a model, then Φ has a model.

Proof. Let $\text{fin}(\Phi)$ be the set of finite subtheories of Φ . By assumption, for each $F \in \text{fin}(\Phi)$, there is a model M_F of F . For appropriate choice of ultrafilter \mathcal{U} on $\text{fin}(\Phi)$, I claim

$$\prod_{F \in \text{fin}(\Phi)} M_F / \mathcal{U}$$

models Φ . To see this, pick an arbitrary $\phi \in \Phi$. Define

$$H_\phi = \{F \subseteq \text{fin}(\Phi) : M_F \models \phi\}$$

to be the set of models where ϕ holds. Consider

$$F = \{H_\phi : \phi \in \Phi\}$$

Then F is a filter on $\text{fin}(\Phi)$. By the [ultrafilter lemma](#), F is contained in some ultrafilter \mathcal{U} . Hence we have chosen a \mathcal{U} such that, for each $\phi \in \Phi$, the set $\{f \in \text{fin}(\Phi) : M_f \models \phi\}$ is in \mathcal{U} . Then

$$\prod_{F \in \text{fin}(\Phi)} M_F / \mathcal{U} \models \Phi$$



Corollary 0.5 The [ultrafilter lemma](#) implies the [compactness theorem](#).

The converse is also true:

Theorem 0.6 The [compactness theorem](#) implies [ultrafilter lemma](#).

Proof. Suppose F is a filter on S . We can encode this assertion in a theory. Let \mathcal{L} be the language of set theory with constant symbols S and F . Let set be the axioms of set theory. To encode the statement F is a filter on S , define the theory

$$\begin{aligned} \Phi = \text{set} \cup \{ \\ & S \in F, \\ & \forall A, B (A \in F \wedge B \in S \implies A \cap B \in F), \\ & \emptyset \notin F, \\ & \forall A, B (A \in F \wedge A \subseteq B \implies B \in F) \\ & \} \end{aligned}$$

Then Φ has a model iff we can find a filter F on S . Define the language \mathcal{L}' by adding a constant symbol x_s for each $s \subseteq S$. Now define the theory Φ' over \mathcal{L}' by

$$\begin{aligned} \Phi' = \Phi \cup \{ & x_s \subseteq S : s \subseteq S \} \\ & \cup \{ x_s \in F \vee S - x_s \in F : s \subseteq S \} \end{aligned}$$

By [compactness theorem](#), Φ' has a model M . This M has a copy of S and \mathcal{U} . They may contain extra elements, but

$$\{X \cap S : X \subseteq \mathcal{U}\} - \emptyset$$

gives an ultrafilter on S containing F .

