
CS722/822: Machine Learning

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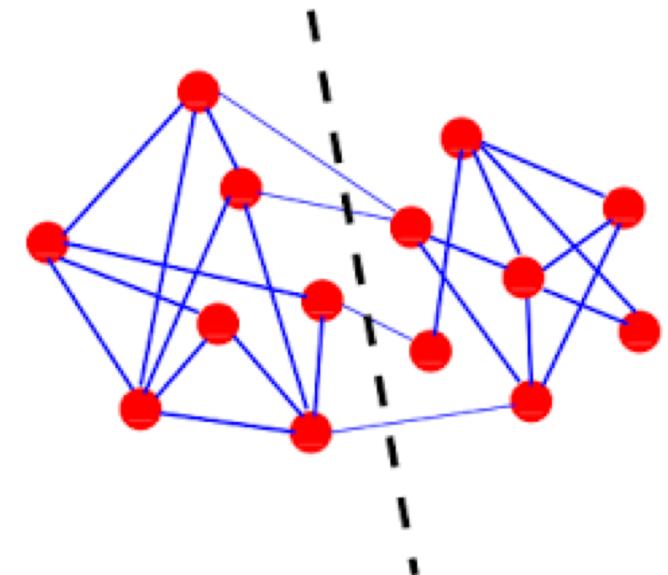
Spectral clustering

Outline

- Basics of graph cut
- Two-way cut
 - RatioCut
- k -way cut
 - RatioCut
 - Ncut
- Accuracy of spectral clustering

Graph Cut

- $G = (V, E)$ is a graph with n vertices and m edges
- Partition vertices into two groups, such that
 - Many within-group edges/vertices
 - Few cross-group edges

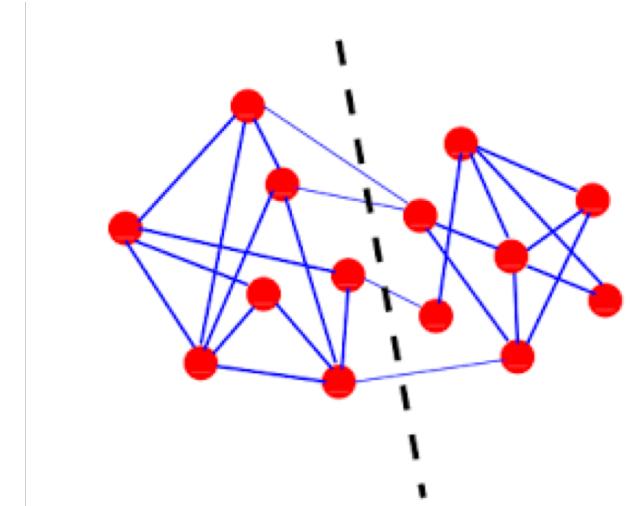


- Search all possible cuts
- Optimization problem

Formulating Graph Cut

- $V = A \cup \bar{A}$
- Weight of crossing-edges between A and \bar{A} :

$$W(A, \bar{A}) = \sum_{\substack{i \in A \\ j \in \bar{A}}} w_{ij}$$



- $\text{cut}(A, \bar{A}) = W(A, \bar{A}) = \frac{1}{2} \sum_{i=1}^2 W(A_i, \bar{A}_i)$
- $\text{cut}(A_1, A_2, \dots, A_k) = \frac{1}{2} \sum_{i=1}^k W(A_i, \bar{A}_i)$

RatioCut

- Minimize crossing edges
 - Maximize within-group vertices
-
- $\text{RatioCut}(A_1, \dots, A_k) = \sum_{i=1}^k \frac{\text{cut}(A_i, \overline{A_i})}{|A_i|}$

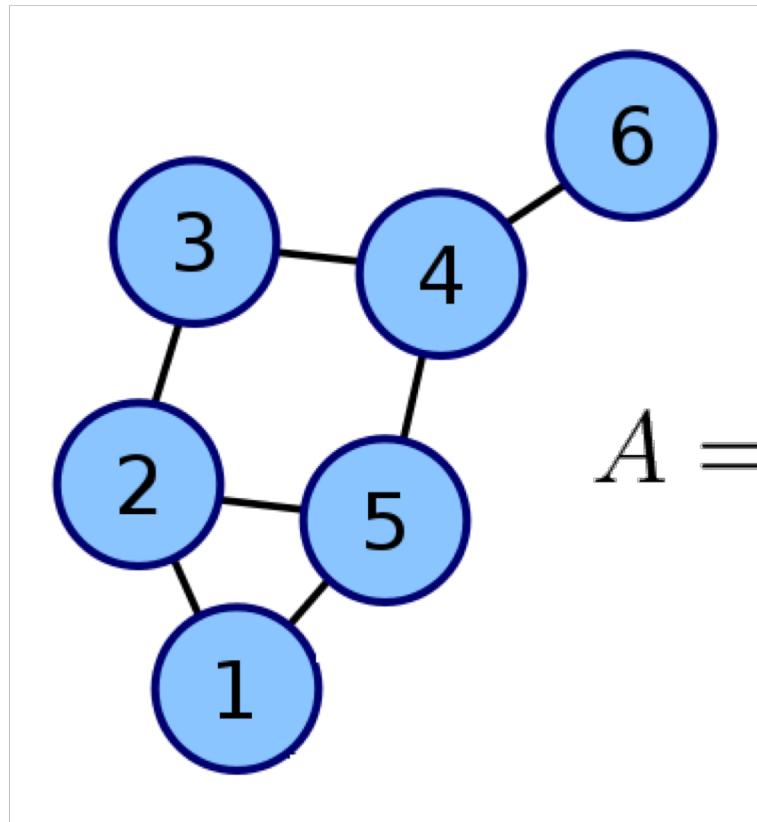
Insights into RatioCut

- Want Balanced Partition
- NP-Hard problem
 - For integer cluster labels $1, 2, \dots, k$
- Spectral Clustering solves the *relaxed* version
 - Reformulate with real-valued vectors
 - Find optimal solution
 - Round to integer labels



GRAPH LAPLACIAN

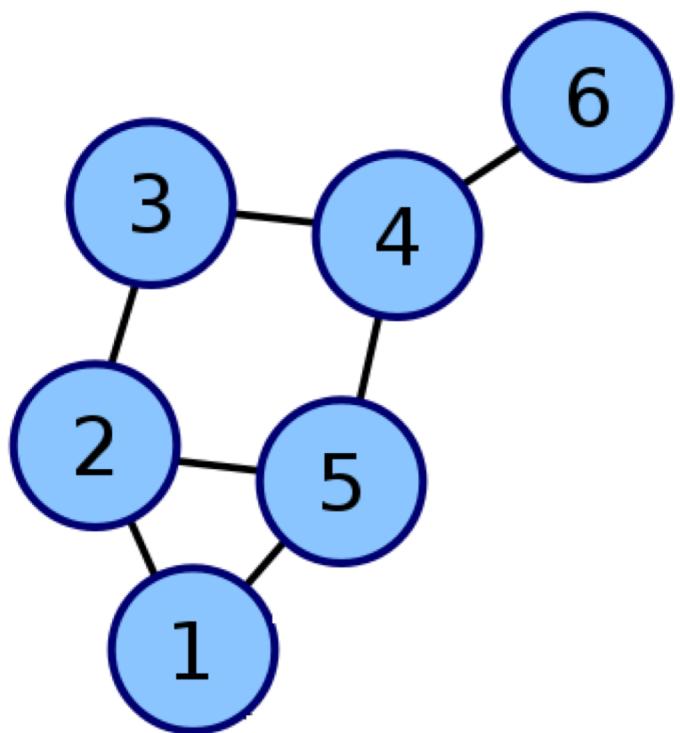
Adjacency Matrix



$$A =$$

$$\begin{bmatrix} 0 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix}$$

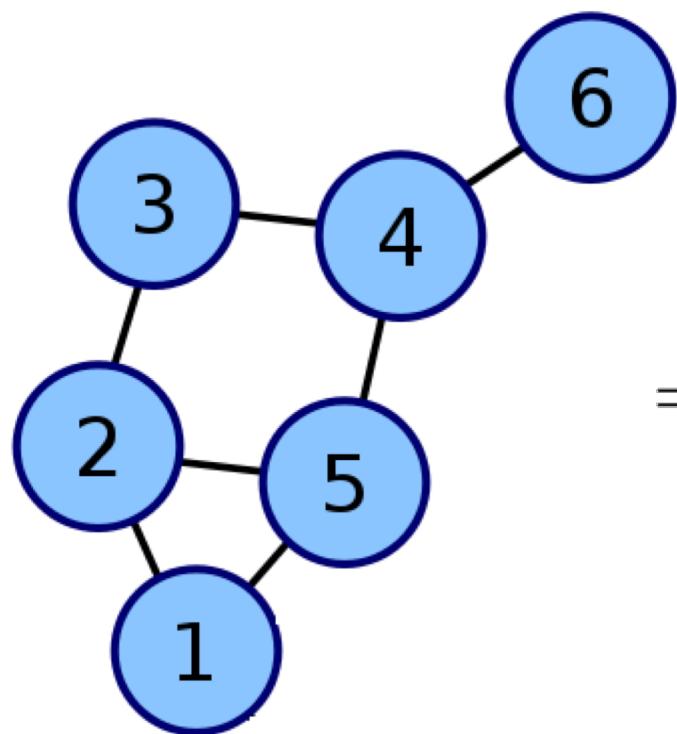
Degree Matrix



$$D = \begin{bmatrix} 2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 3 & 0 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 0 & 0 \\ 0 & 0 & 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$d_{ij} = \begin{cases} \sum_k a_{ik}, & i = j \\ 0, & i \neq j \end{cases}$$

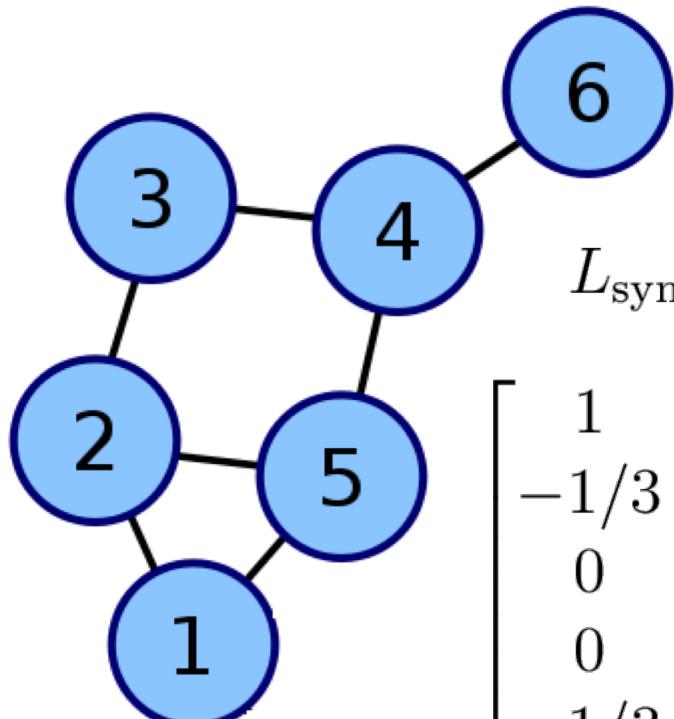
Laplacian Matrix



$$L = D - A$$

$$= \begin{bmatrix} 2 & -1 & 0 & 0 & -1 & 0 \\ -1 & 3 & -1 & 0 & -1 & 0 \\ 0 & -1 & 2 & -1 & 0 & 0 \\ 0 & 0 & -1 & 3 & -1 & -1 \\ -1 & -1 & 0 & -1 & 3 & 0 \\ 0 & 0 & 0 & -1 & 0 & 1 \end{bmatrix}$$

Normalized Laplacian Matrix



$$L_{\text{sym}} = D^{-1/2} L D^{-1/2} =$$

$$\begin{bmatrix} 1 & -1/2 & 0 & 0 & -1/2 & 0 \\ -1/3 & 1 & -1/3 & 0 & -1/3 & 0 \\ 0 & -1/2 & 1 & -1/2 & 0 & 0 \\ 0 & 0 & -1/3 & 1 & -1/3 & -1/3 \\ -1/3 & -1/3 & 0 & -1/3 & 1 & 0 \\ 0 & 0 & 0 & -1 & 0 & 1 \end{bmatrix}$$

Edge Vector

$$b_e = b_{ij} = \begin{bmatrix} \vdots \\ 1 \\ \vdots \\ -1 \\ \vdots \end{bmatrix} \begin{array}{l} \leftarrow i \text{ th entry} \\ \leftarrow j \text{ th entry} \end{array}$$

$$\begin{aligned} b_e^T x &= [\cdots \quad 1 \quad \cdots \quad -1 \quad \cdots] \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \\ &= (x_i - x_j) \end{aligned}$$

Edge Laplacians

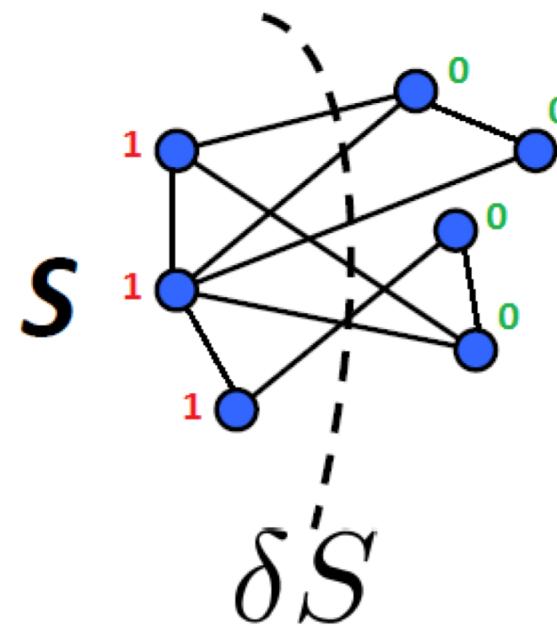
$$L_e = b_e b_e^T = \begin{bmatrix} \ddots & & \cdots & \cdots & \cdots & & \vdots \\ \vdots & 1 & \cdots & -1 & & & \vdots \\ \vdots & \vdots & \ddots & & \vdots & & \vdots \\ \vdots & -1 & \cdots & 1 & & & \vdots \\ \vdots & \cdots & \cdots & \cdots & \ddots & & \end{bmatrix}$$

$$\begin{aligned} L &= \sum_e w_e L_e \\ &= \sum_e w_e b_e b_e^T \end{aligned}$$

Laplacian Quadratic Form, $x^T Lx$

$$\begin{aligned} & x^T Lx \\ &= x^T \left(\sum_{e=(i,j)} w_e b_e b_e^T \right) x \\ &= \sum_{e=(i,j)} (x^T b_e) w_e (b_e^T x) \\ &= \sum_{e=(i,j)} (b_e^T x)^T w_e (b_e^T x) \\ &= \sum_{e=(i,j)} (x_i - x_j) w_e (x_i - x_j) \\ &= \sum_{e=(i,j)} w_e (x_i - x_j)^2 \end{aligned}$$

If $x \in \{0,1\}^n$ denotes a vertex partition,
 $x^T Lx = \text{cut-weight}$



How this relates to the RatioCut?

Recall Ratio Cut

- Minimize crossing edges
 - Maximize in-part vertices
-
- $\text{RatioCut}(A_1, \dots, A_k): \min \sum_{i=1}^k \frac{\text{cut}(A_i, \overline{A_i})}{|A_i|}$



RATIOCUT FOR $k = 2$

RatioCut to Laplacian Quadratic Form (1)

$$\begin{aligned} & \text{RatioCut}(A, \bar{A}) \cdot |V| \\ &= \left(\frac{\text{cut}(A, \bar{A})}{|A|} + \frac{\text{cut}(\bar{A}, A)}{|\bar{A}|} \right) \cdot |V| \\ &= \text{cut}(A, \bar{A}) \cdot \left(\frac{|V|}{|A|} + \frac{|V|}{|\bar{A}|} \right) \\ &= \text{cut}(A, \bar{A}) \cdot \left(\frac{|A| + |\bar{A}|}{|A|} + \frac{|A| + |\bar{A}|}{|\bar{A}|} \right) \\ &= \text{cut}(A, \bar{A}) \cdot \left(\frac{|\bar{A}|}{|A|} + \frac{|A|}{|\bar{A}|} + 2 \right) \end{aligned}$$

RatioCut to Laplacian Quadratic Form (2)

$$\begin{aligned} & \text{RatioCut}(A, \bar{A}) \cdot |V| \\ = & \left(\sum_{i \in A, j \in \bar{A}} w_{ij} \right) \cdot \left(\frac{|\bar{A}|}{|A|} + \frac{|A|}{|\bar{A}|} + 2 \right) \\ = & \left(\frac{1}{2} \sum_{i \in A, j \in \bar{A}} w_{ij} + \frac{1}{2} \sum_{j \in \bar{A}, i \in A} w_{ij} \right) \cdot \left(\sqrt{\frac{|\bar{A}|}{|A|}} + \sqrt{\frac{|A|}{|\bar{A}|}} \right)^2 \\ = & \frac{1}{2} \sum_{i \in A, j \in \bar{A}} w_{ij} \left(\sqrt{\frac{|\bar{A}|}{|A|}} + \sqrt{\frac{|A|}{|\bar{A}|}} \right)^2 + \frac{1}{2} \sum_{j \in \bar{A}, i \in A} w_{ij} \left(-\sqrt{\frac{|\bar{A}|}{|A|}} - \sqrt{\frac{|A|}{|\bar{A}|}} \right)^2 \end{aligned}$$

RatioCut to Laplacian Quadratic Form (3)

$$\text{Let } f_i = \begin{cases} \sqrt{\frac{|\bar{A}|}{|A|}} & \text{if } v_i \in A \\ -\sqrt{\frac{|A|}{|\bar{A}|}} & \text{if } v_i \in \bar{A} \end{cases}$$

Then

$$\begin{aligned} & \text{RatioCut}(A, \bar{A}) \cdot |V| \\ &= \frac{1}{2} \sum_{i \in A, j \in \bar{A}} w_{ij} (f_i - f_j)^2 + \frac{1}{2} \sum_{i \in \bar{A}, j \in A} w_{ij} (f_j - f_i)^2 \\ &= \frac{1}{2} \sum_{i,j=1}^n w_{ij} (f_i - f_j)^2 \\ &= \frac{1}{2} \cdot \left(2 \cdot \sum_{i \leq j} w_{ij} (f_i - f_j)^2 \right) \\ &= f^T L f \end{aligned}$$

What Did We Do?

- We showed that $\min \text{RatioCut} \equiv \min f^T L f$

- Properties of f :

- f is orthogonal to the constant, $f \perp 1$

$$f^T 1 = \sum f_i = |A| \cdot \sqrt{\frac{|\bar{A}|}{|A|}} + |\bar{A}| \cdot \left(-\sqrt{\frac{|A|}{|\bar{A}|}}\right) = \sqrt{|A||\bar{A}|} - \sqrt{|A||\bar{A}|} = 0$$

- Length of f is \sqrt{n}

$$\|f\|^2 = f^T f = |A| \cdot \frac{|\bar{A}|}{|A|} + |\bar{A}| \cdot \frac{|A|}{|\bar{A}|} = |A| + |\bar{A}| = n$$

Reformulated Minimization Problem

- Discrete optimization

$$\min_A f' L f, \quad \text{subject to } f \perp 1, \|f\| = \sqrt{n}$$

- Relax to arbitrary real vectors
 - because Integer programming is NP-hard

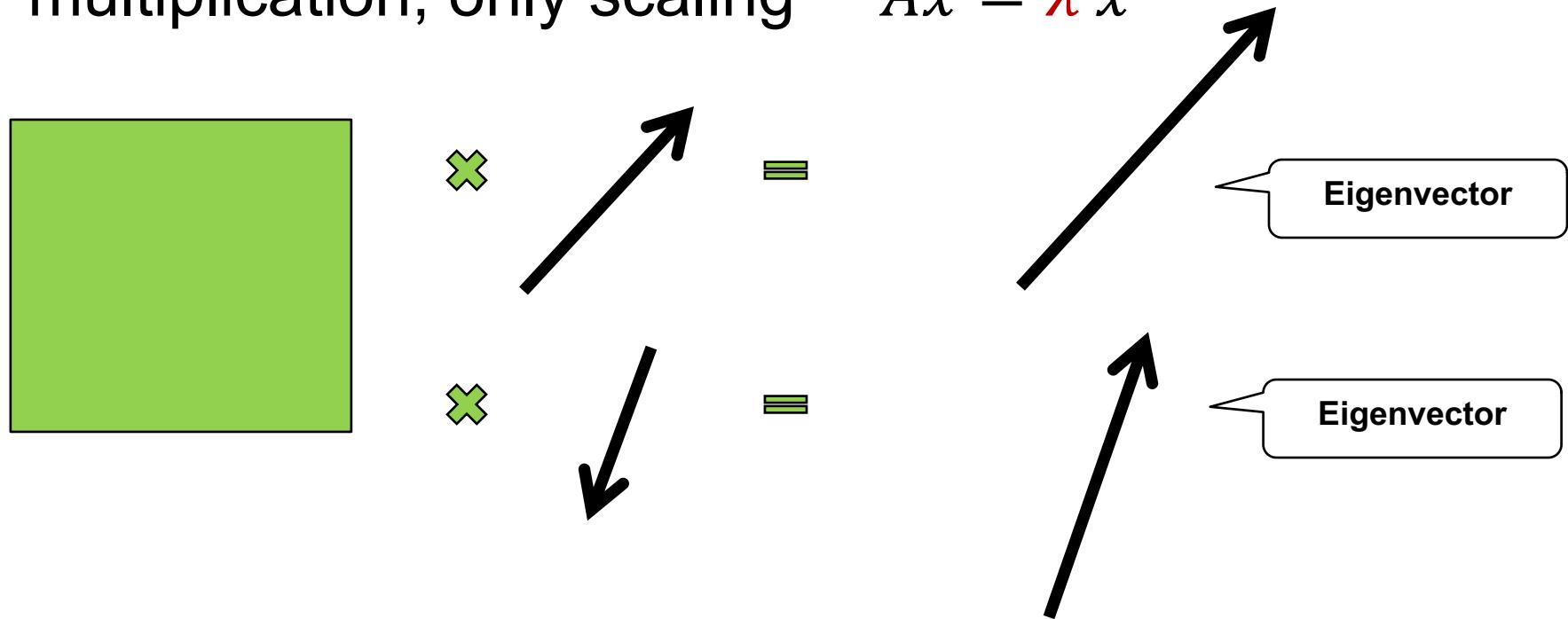
$$\min_{f \in \mathbb{R}^n} f' L f, \quad \text{subject to } f \perp 1, \|f\| = \sqrt{n}$$

Eigenvalues and Eigenvectors

$A \leftarrow n \times n$ matrix,

$x \leftarrow n \times 1$ vector

Eigenvector: direction same/negative after multiplication, only scaling $Ax = \lambda x$



Eigenvalue: amount of distortion $\lambda = x^T A x / x^T x$

Rayleigh Quotient and Eigenvalues

$$q(A, x) = \frac{x^T A x}{x^T x}$$

- S_1 = Set of all vectors
 - Vector v_1 attains minimum $\lambda_1 = q(A, x) \quad \forall x \in S_1$
 - Set of first smallest eigenvectors
- S_2 = Set of all vectors $\perp \{v_1\}$
 - Vector v_2 attains minimum $\lambda_2 = q(A, x) \quad \forall x \in S_2$
 - Set of second smallest eigenvectors
- ...
- S_n = Set of all vectors $\perp \{v_1, v_2, \dots, v_{n-1}\}$
 - v_n attains minimum $\lambda_n = q(A, x) \quad \forall x \in S_n$
 - Set of n-th smallest eigenvectors

$$\lambda_1 < \lambda_2 < \dots < \lambda_n$$

Rayleigh-Ritz Theorem (simplified)

- the k -th smallest eigenvector v_k is the n -dimensional vector x
- such that
 - x is perpendicular to all previous eigenvectors and
 - x minimizes the quantity $\frac{x^T Ax}{x^T x}$

Smallest Eigenvalue of the Laplacian

- Rows/columns of L sum to 0

$$\text{[green square]} \times \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix} = 0 \times \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}$$

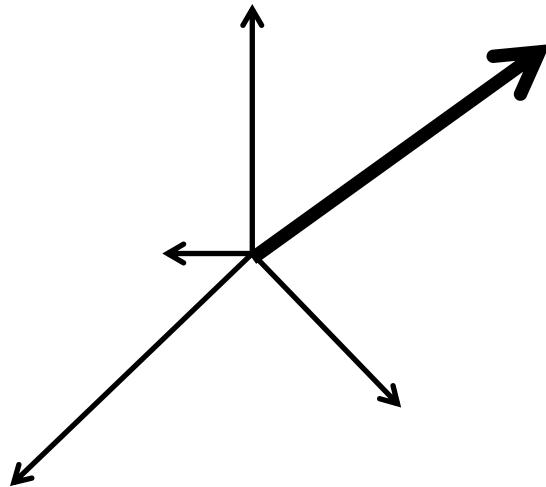
- $\lambda_1 = 0$ and $v_1 = 1$, the constant vector
- Any constant vector is the first eigenvector

Second Eigenvector of the Laplacian

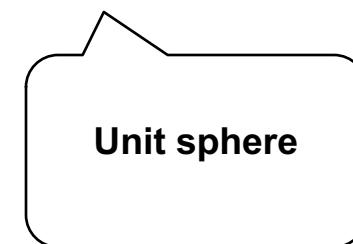
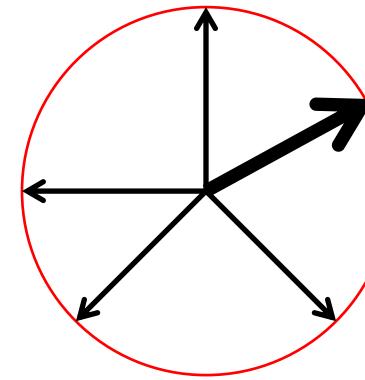
- Since eigenvectors are perpendicular to each other,
 - $v_2 \perp 1$
 - “Fiedler vector”... named after Miroslav Fiedler
- $\lambda_2 = \operatorname{argmin} (x^T L x / x^T x)$ over all $x \in R^n, x \perp 1$
- Condition similar to the vector f chosen when minimizing RatioCut

Eigenvector and Scaling

- Consider an eigenvector



- Scale the vector to length 1
 - Still eigenvector
 - With the same eigenvalue
- Find unit-length eigenvectors on the unit-sphere



ν_2 and Laplacian Quadratic Form

- $\lambda_2 = \operatorname{argmin} (x^T L x / x^T x)$ over all $x \in R^n, \textcolor{blue}{x} \perp 1$



- $\lambda_2 = \operatorname{argmin} (x^T L x)$ over all $x \in R^n, \textcolor{blue}{x} \perp 1, x^T x = 1$

- $\min_{f \in \mathbb{R}^n} f^T L f$, subject to $f \perp 1, \|f\| = 1$



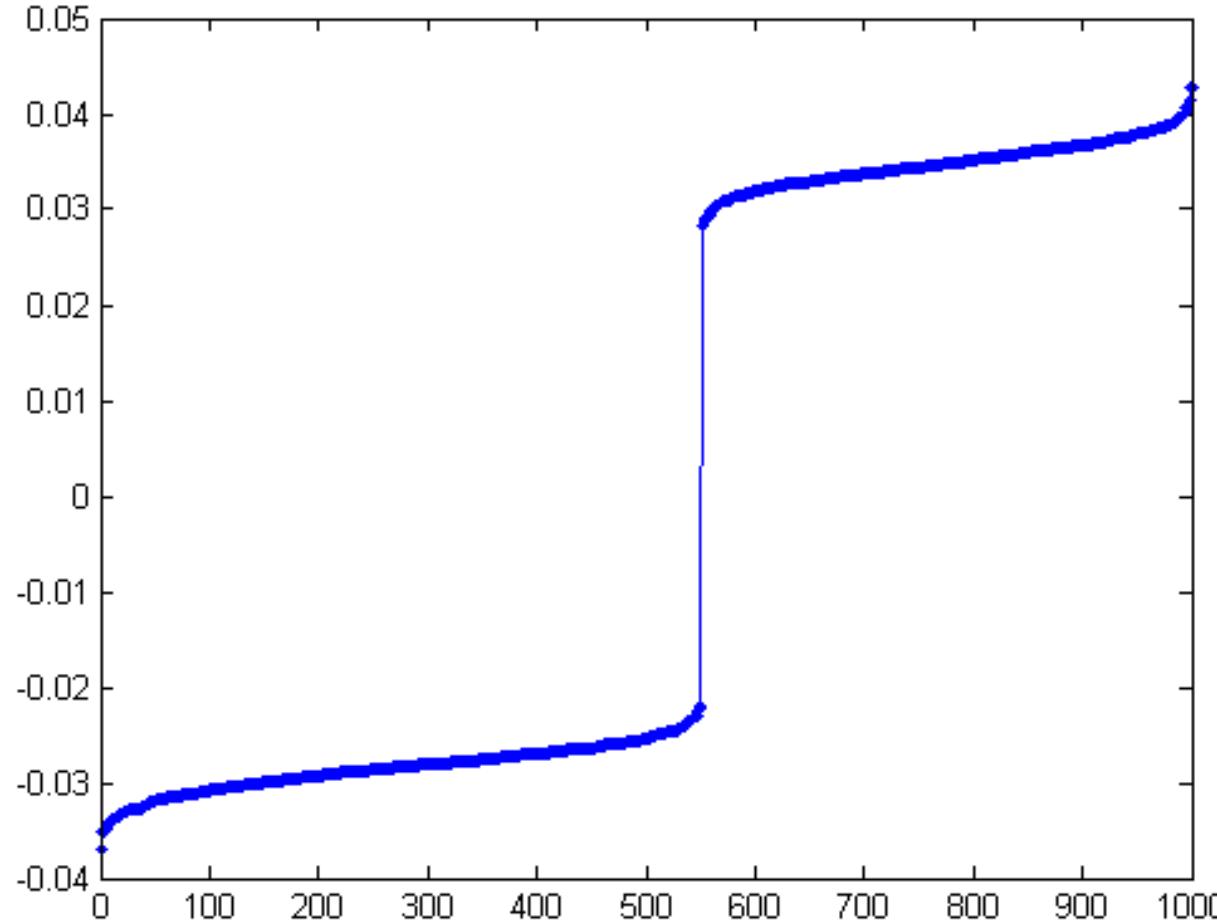
- $\min_{f \in \mathbb{R}^n} f^T L f$, subject to $f \perp 1, \|f\| = \sqrt{n}$

Big Picture

- Minimizing RatioCut \equiv Finding v_2 of L
 - v_2 is the eigenvector corresponding to the second smallest eigenvalue, λ_2
- Clustering can be made based on v_2
 - Simply by rounding to integer labels, or
 - Running k-means

This is a mapping from n-dimension to 1-dimension

Plot of the Second Eigenvector





RATIOCUT FOR $k \geq 3$

Partition Indicator Vectors (1)

- Partitions A_1, A_2, \dots, A_k
- Partition A_j has its indicator vector h_j

$$h_j = \begin{bmatrix} 0 \\ \vdots \\ 1/\sqrt{|A_j|} \\ 0 \\ \vdots \\ 0 \\ 1/\sqrt{|A_j|} \\ \vdots \\ 0 \end{bmatrix} \leftarrow \begin{array}{l} \text{This vertex in } A_j \\ \text{This vertex in } A_j \end{array}$$

Partition Indicator Vectors (2)

- Note that $h_j^T h_j = |A_j| \cdot (1/\sqrt{|A_j|})^2 = 1$
- Note that $h_i^T h_j = 0$ for all $i \neq j$
- Thus vectors h_1, h_2, \dots, h_k are **orthonormal**

$$h_j = \begin{bmatrix} 0 \\ \vdots \\ 1/\sqrt{|A_j|} \\ 0 \\ \vdots \\ 0 \\ 1/\sqrt{|A_j|} \\ \vdots \\ 0 \end{bmatrix} \leftarrow \begin{array}{l} \text{This vertex in } A_j \\ \text{This vertex in } A_j \end{array}$$

RatioCut for $k > 2$

- It can be shown that

$$h_i^T L h_i = \frac{\text{cut}(A_i, \bar{A}_i)}{|A_i|}$$

- Then

$$\text{RatioCut}(A_1, \dots, A_k) = \sum_{i=1}^k \frac{\text{cut}(A_i, \bar{A}_i)}{|A_i|} = \sum_i h_i^T L h_i = \text{Trace}(H^T L H)$$

RatioCut for $k > 2$

- Now we have a trace minimization problem

$$\min_{A_1, \dots, A_k} \text{Tr}(H'LH) \text{ subject to } H'H = I$$

- Relax to take real values

$$\min_{H \in \mathbb{R}^{n \times k}} \text{Tr}(H'LH) \text{ subject to } H'H = I.$$

- The minimum objective is the sum of the k smallest eigenvectors of L
- So each optimal h_i is the eigenvector associated with the i -th smallest eigenvalue?

RatioCut for $k > 2$

- Now we know the $n \times k$ matrix H
 - by computing k smallest eigenvectors of L
- Treat each row of H as a k -dimensional mapping of the corresponding n -dimensional point.
- Use k-means to partition rows of H
- Assign cluster label to each vertex



NCUT

Normalized Cut

- Minimize # of crossing edges
 - Maximize # of within-group edges
-
- $\text{NCut}(A_1, \dots, A_k) = \sum_{i=1}^k \frac{\text{cut}(A_i, \overline{A_i})}{\text{vol}(A_i)}$

where

- $\text{vol}(v) = \deg(v)$
- $\text{vol}(A) = \sum_{v \in A} \deg(v)$

Cluster Indicator Vectors ($k = 2$ case)

- RatioCut

$$f_i = \begin{cases} \sqrt{|\overline{A}|/|A|} & \text{if } v_i \in A \\ -\sqrt{|A|/|\overline{A}|} & \text{if } v_i \in \overline{A}. \end{cases}$$

- $f \perp 1$
- $f^T f = n$

- Ncut

$$f_i = \begin{cases} \sqrt{\frac{\text{vol}(\overline{A})}{\text{vol } A}} & \text{if } v_i \in A \\ -\sqrt{\frac{\text{vol}(A)}{\text{vol}(\overline{A})}} & \text{if } v_i \in \overline{A}. \end{cases}$$

- $Df \perp 1$
- $f^T Df = \text{vol}(V)$

Minimization Problem

- minimize $f^T L f$
over all $f \in R^n$

subject to

- $Df \perp 1$
- $f^T Df = \text{vol}(V)$

- *Constraints different from RatioCut*

Work with $g = D^{1/2} f$

- Let $d_i = \deg(v_i)$
- Let $\textcolor{red}{g = D^{1/2}f}$, that is $g_i = \sqrt{d_i} \cdot f_i$

- Then $f = D^{-\frac{1}{2}}g$

- And

$$f^T L f = \left(D^{-\frac{1}{2}} g \right)^T L \left(D^{-\frac{1}{2}} g \right) = g^T \left(D^{-\frac{1}{2}} L D^{-\frac{1}{2}} \right) g$$

- And $f^T D f = g^T g = \text{vol}(V)$

- $Df \perp 1 \equiv \textcolor{red}{g \perp D^{1/2}1}$

Relaxed Minimization Problem with g

- minimize $g^T (D^{-\frac{1}{2}} L D^{-\frac{1}{2}}) g$
over all $g \in R^n$

subject to

- $g \perp D^{1/2} 1$
- $g^T g = \text{vol}(V)$

- $D^{-\frac{1}{2}} L D^{-\frac{1}{2}}$ is the *normalized* Laplacian, L_{sym}
- $D^{1/2} 1$ is the first eigenvector of L_{sym} with eigenvalue 0

Minimization problem with f

- minimize $f^T L f$
over all $f \in R^n$

subject to

- $Df \perp 1$
- $f^T Df = \text{vol}(V)$

Ncut

- Like RatioCut, the solution to NCut is the second eigenvector of L_{sym}
- Analysis for $k \geq 3$ is similar to that of RatioCut

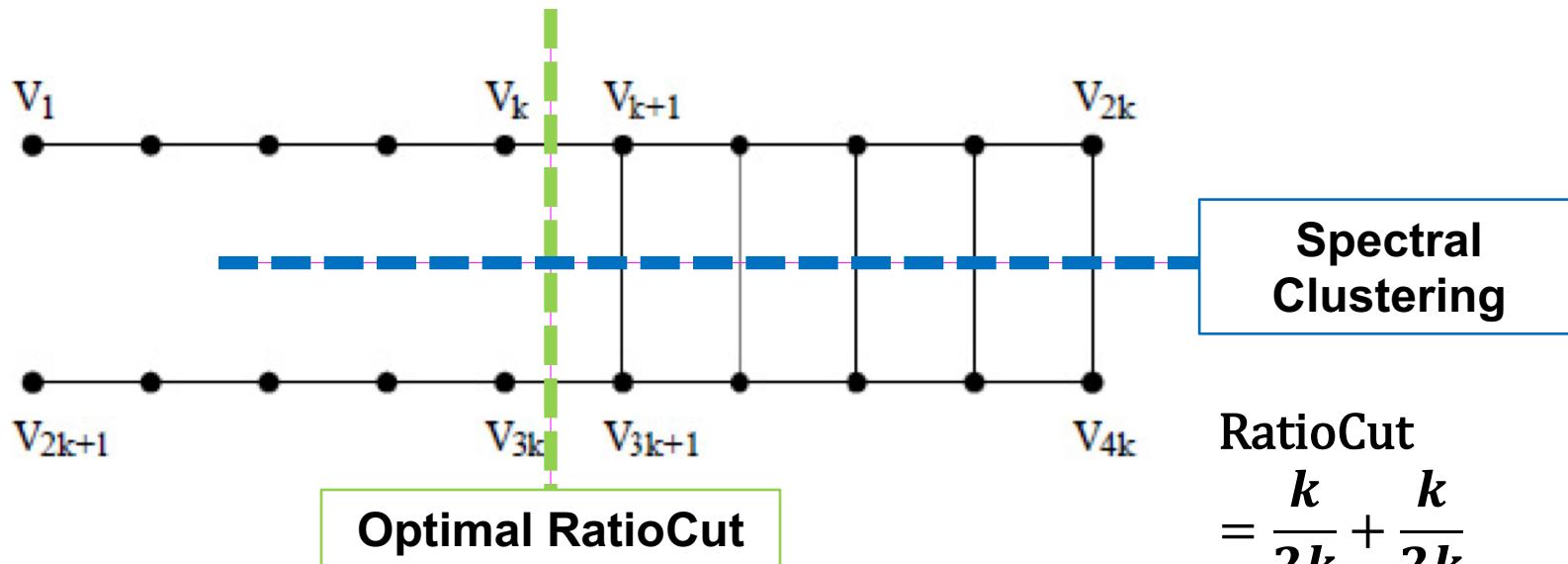


ACCURACY OF SPECTRAL CLUSTERING

Big Picture

- Finding balanced cut is NP-hard
- Relax the integer cluster label vectors to real-valued ones
- Solve the relaxed balanced cut problem, which is equivalent to an eigen decomposition problem
- The relaxation **not** guaranteed to give optimal solution

Badness of Spectral Clustering



$$\begin{aligned}\text{RatioCut} \\ = \frac{k}{2k} + \frac{k}{2k} \\ = 1\end{aligned}$$

$$\begin{aligned}\text{RatioCut} \\ = \frac{2}{2k} + \frac{2}{2k}\end{aligned}$$

$$= \frac{2}{k}$$

Spectral Clustering result is
 $k/2$ times worse than
optimal result