Quick Tour of ClickHouse Internals

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ClickHouse use cases

A stream of events

- > Actions of website visitors **Yandex** Metrica
- > Ad impressions
- > Financial transactions
- > DNS queries
- **>** ...
- We want to save info about these events and then glean some insights from it



ClickHouse philosophy

- > Interactive queries on data updated in real time
- > Cleaned structured data is needed
- > Try hard not to pre-aggregate anything
- Query language: a dialect of SQL + extensions



Sample query in a web analytics system

Top-10 referers for a website for the last week.

```
SELECT Referer, count(*) AS count
FROM hits
WHERE CounterID = 111 AND Date >= today() - 7
GROUP BY Referer
ORDER BY count DESC
LIMIT 10
```



How to execute a query fast?

Read data fast

- > Only needed columns: CounterID, Date, Referer
- > Locality of reads (an index is needed!)
- Data compression



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- > Data compression

Process data fast

- > Vectorized execution (block-based processing)
- > Parallelize to all available cores and machines
- > Specialization and low-level optimizations



Index needed!

The principle is the same as with classic DBMSes

A majority of queries will contain conditions on CounterID and (possibly) Date

(CounterID, Date) fits the bill

Check this by mentally sorting the table by primary key

Differences

- > The table will be physically sorted on disk
- > Is not a unique constraint



Index internals

N

N+8192

N+16384

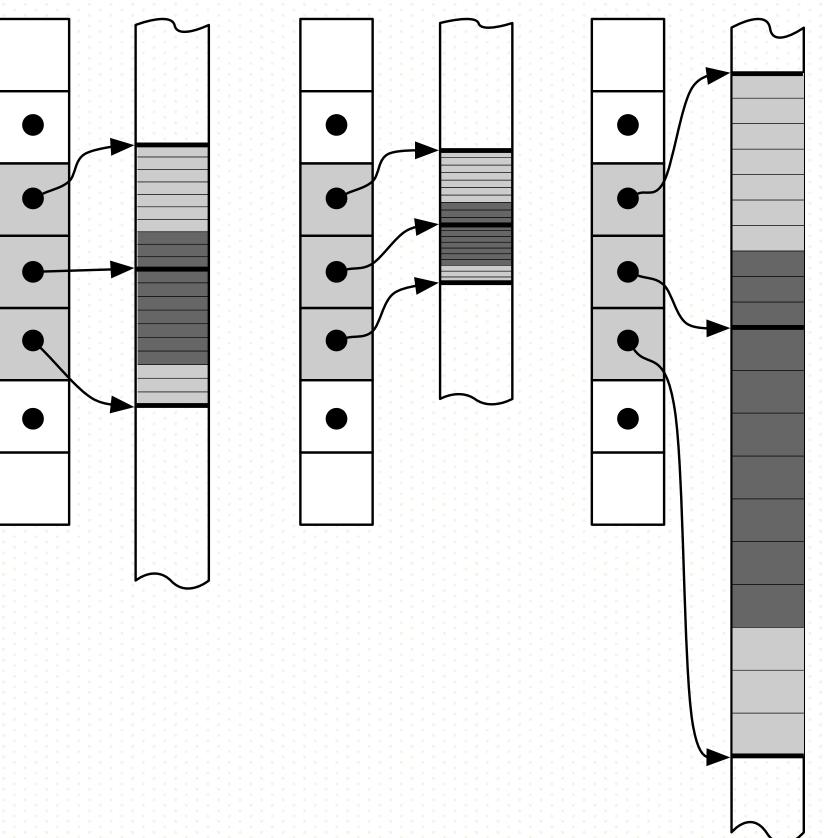
(CounterID, Date) CounterID Referer Date

primary.idx

•••	•••	
111	2017-09-11	
111	2017-09-19	
111	2017-09-27	
222	2013-02-16	
222	2013-03-12	
•••	•••	

(One entry each 8192 rows)







Things to remember about indexes

Index is sparse

- > Must fit into memory
- Default value of granularity (8192) is good enough
- > Does not create a unique constraint
- > Performance of point queries is not stellar

Table is sorted according to the index

- > There can be only one
- > Using the index is always beneficial



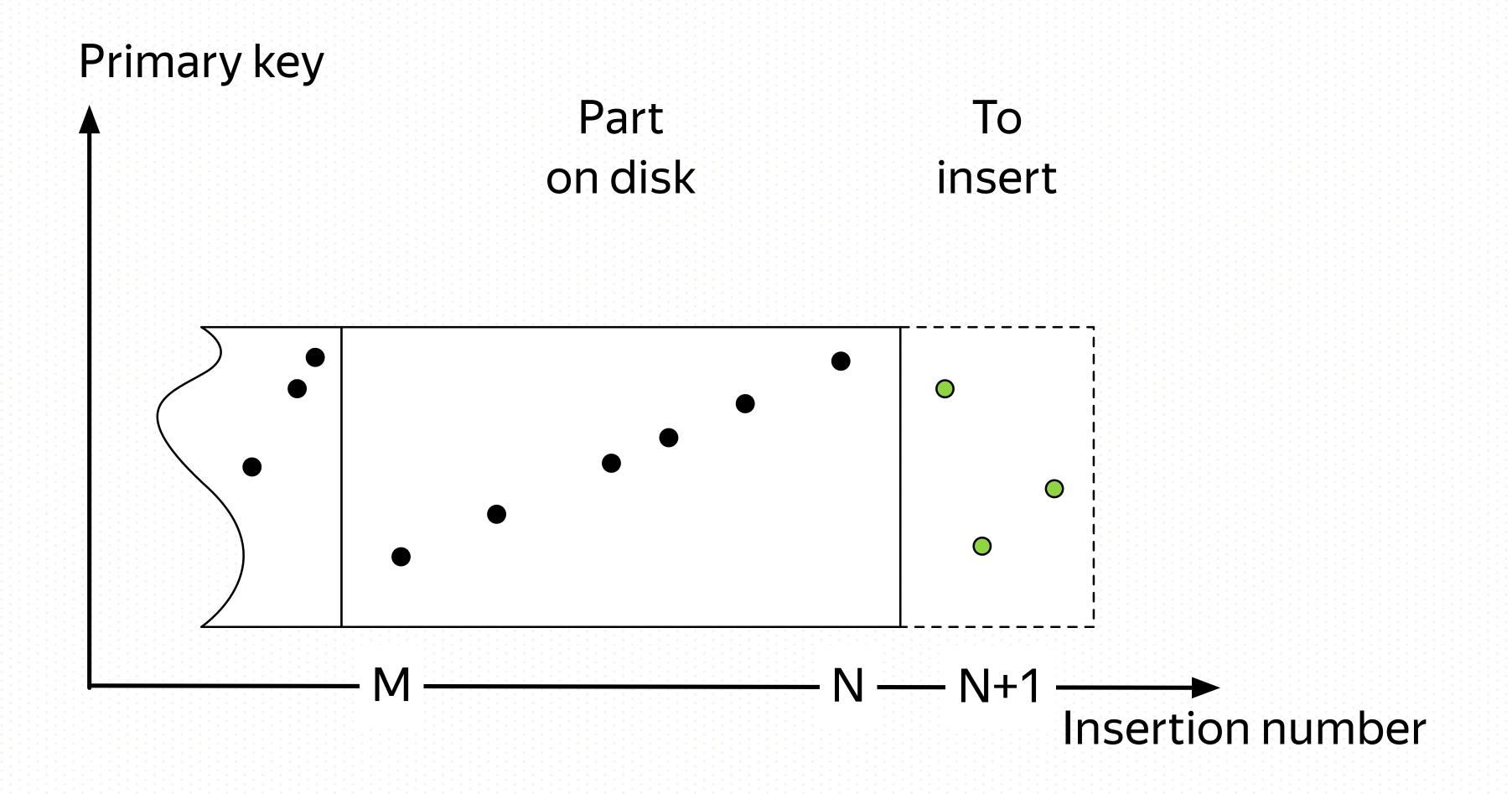
Inserted events are (almost) sorted by time

But we need to sort by primary key!

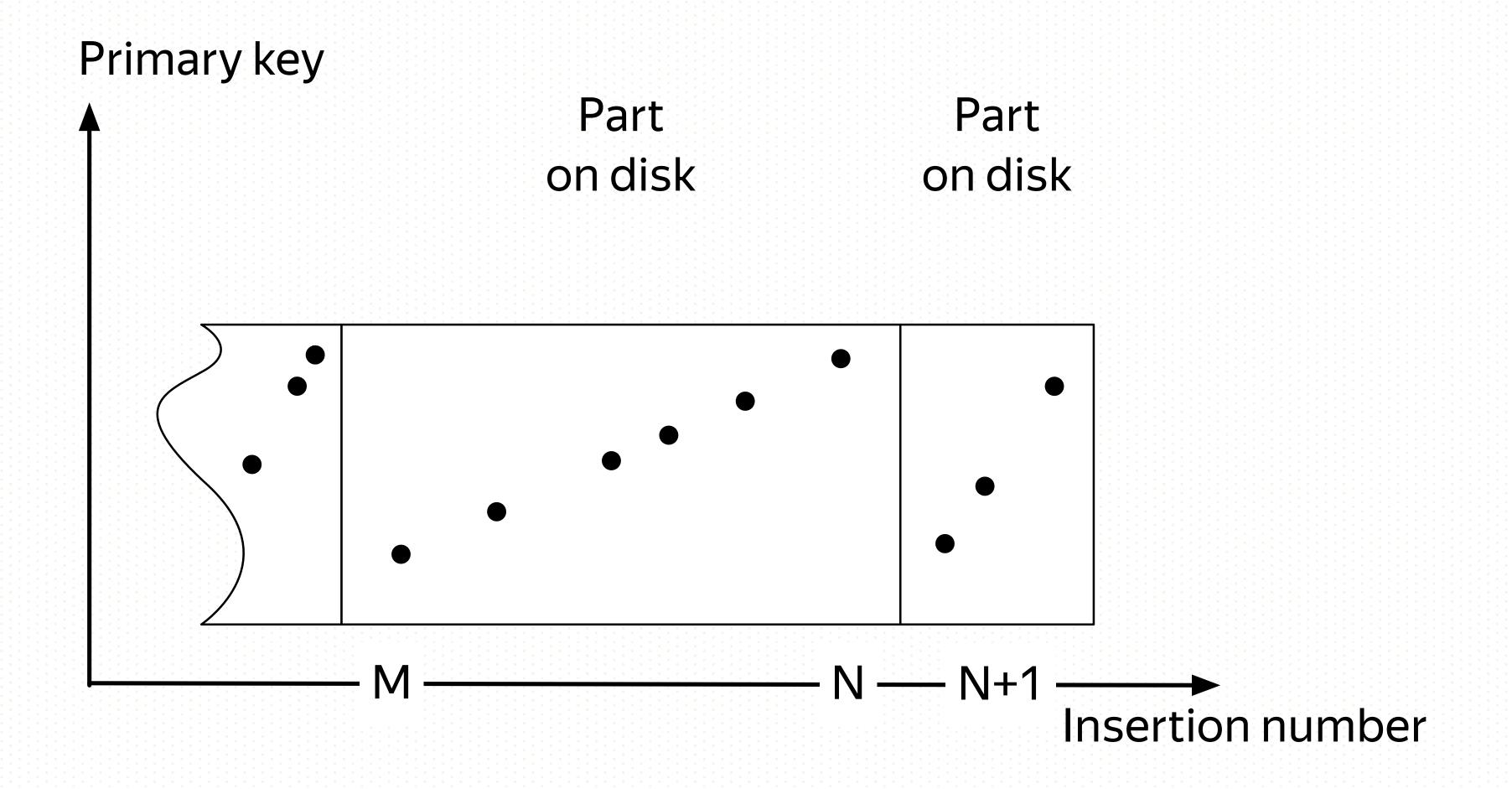
MergeTree: maintain a small set of sorted parts

Similar idea to an LSM tree

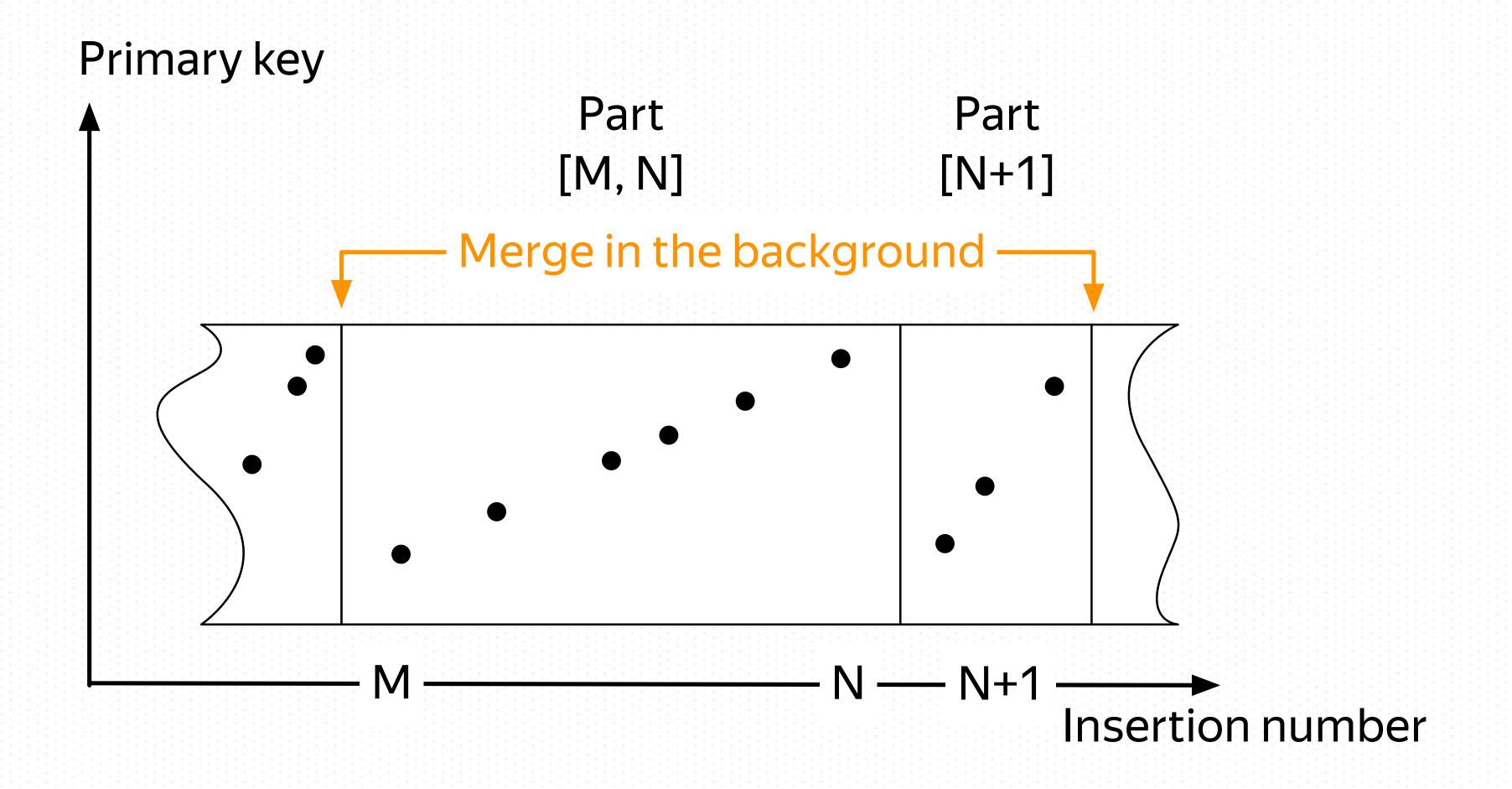




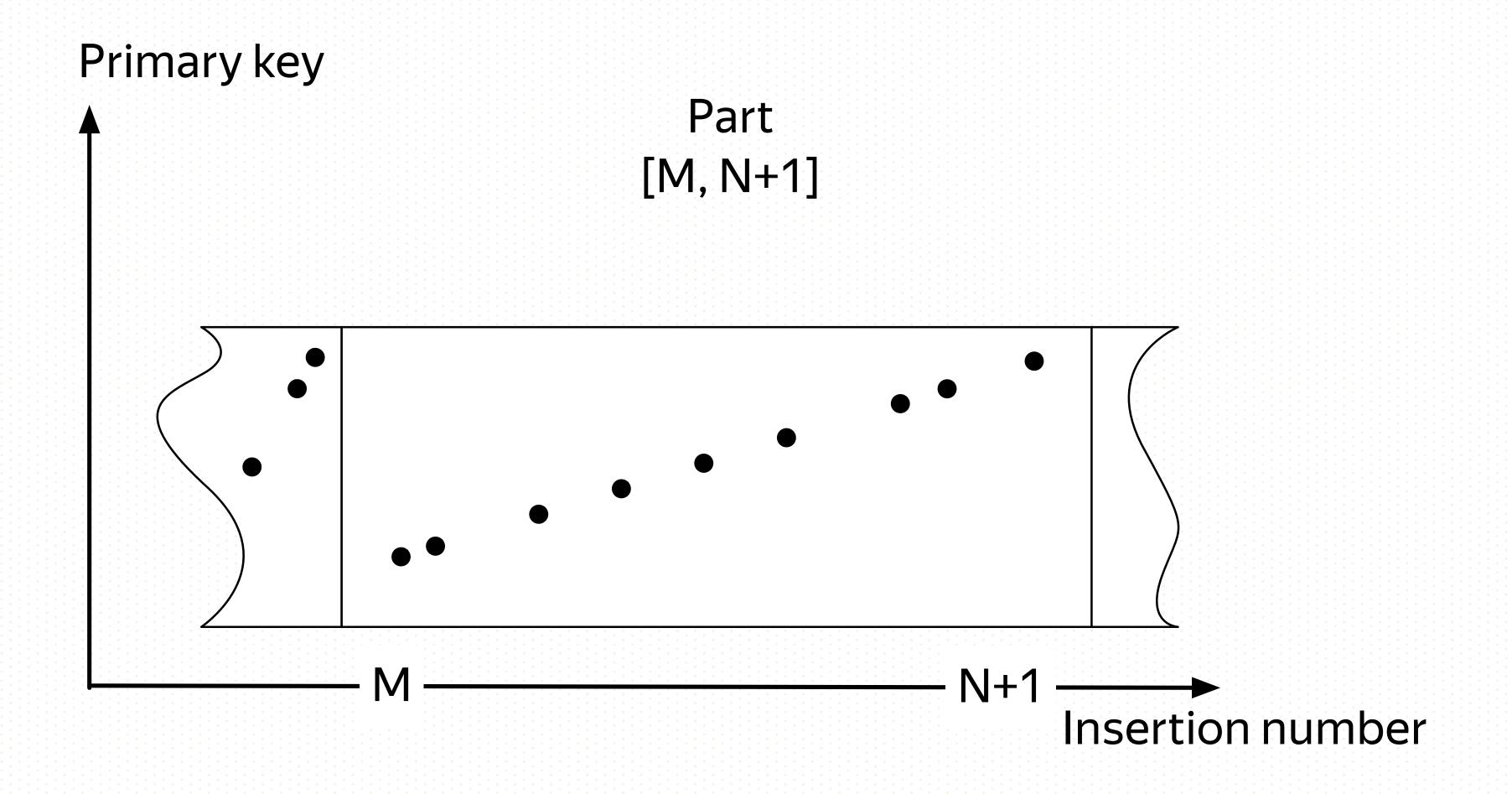














Things to do while merging

- Replace/update records
 - > ReplacingMergeTree
 - CollapsingMergeTree
- Pre-aggregate data
 - > AggregatingMergeTree
- Metrics rollup
 - > GraphiteMergeTree



MergeTree partitioning

ENGINE = MergeTree(Date,...)

- Table is partitioned by month or (soon) by any expression
- > Parts from different partitions are not merged
- > Easy manipulation of partitions

ALTER TABLE DROP PARTITION ALTER TABLE DETACH/ATTACH PARTITION

> MinMax index by partition columns



Things to remember about MergeTree

- Merging runs in the background
 - > Even when there are no queries!
- Control total number of parts
 - > Rate of INSERTs
 - MaxPartsCountForPartition and DelayedInserts metrics are your friends



When one server is not enough

- > The data won't fit on a single server...
- You want to increase performance by adding more servers...
- > Multiple simultaneous queries are competing for resources...



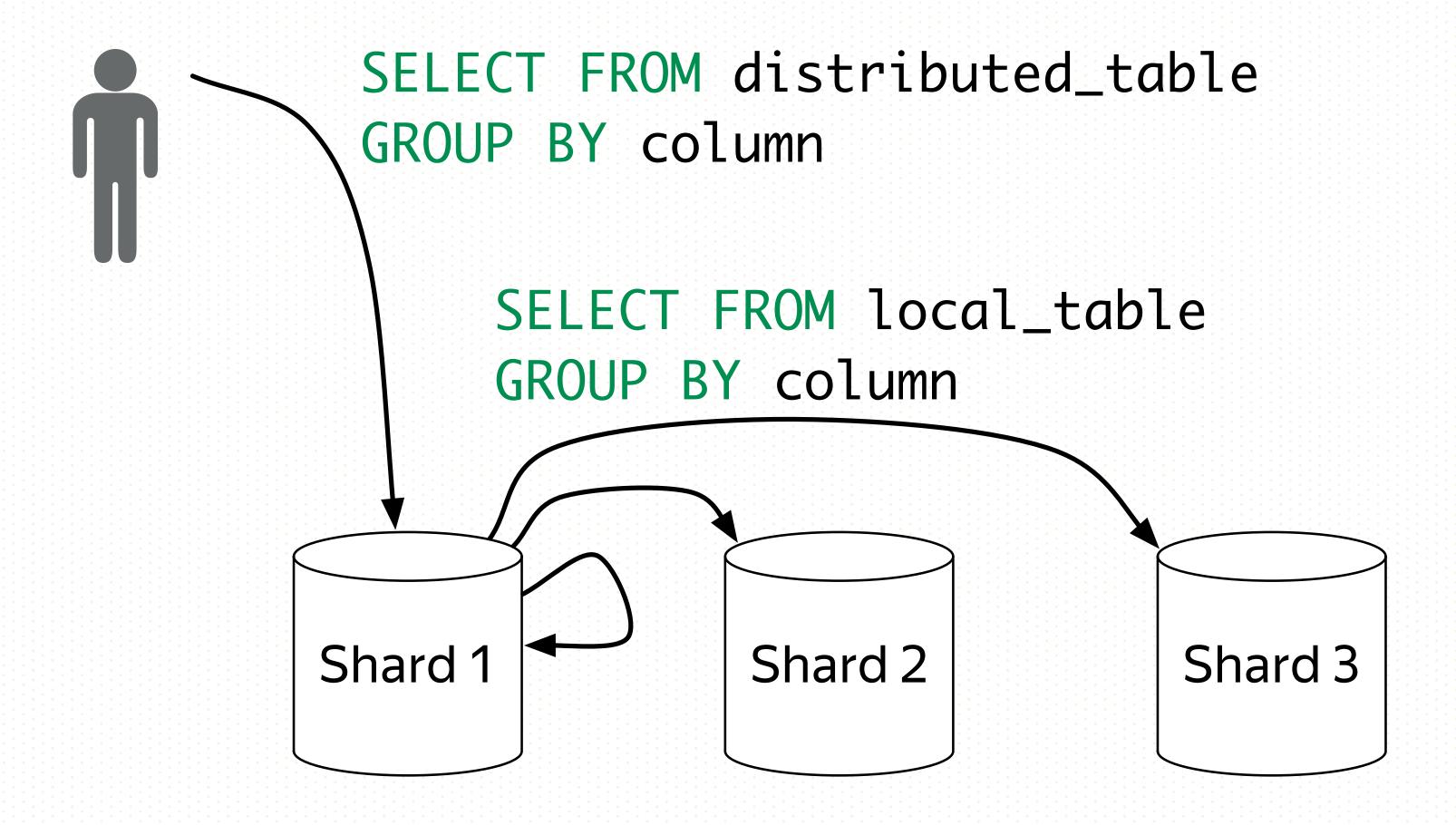
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ClickHouse: Sharding + Distributed tables!

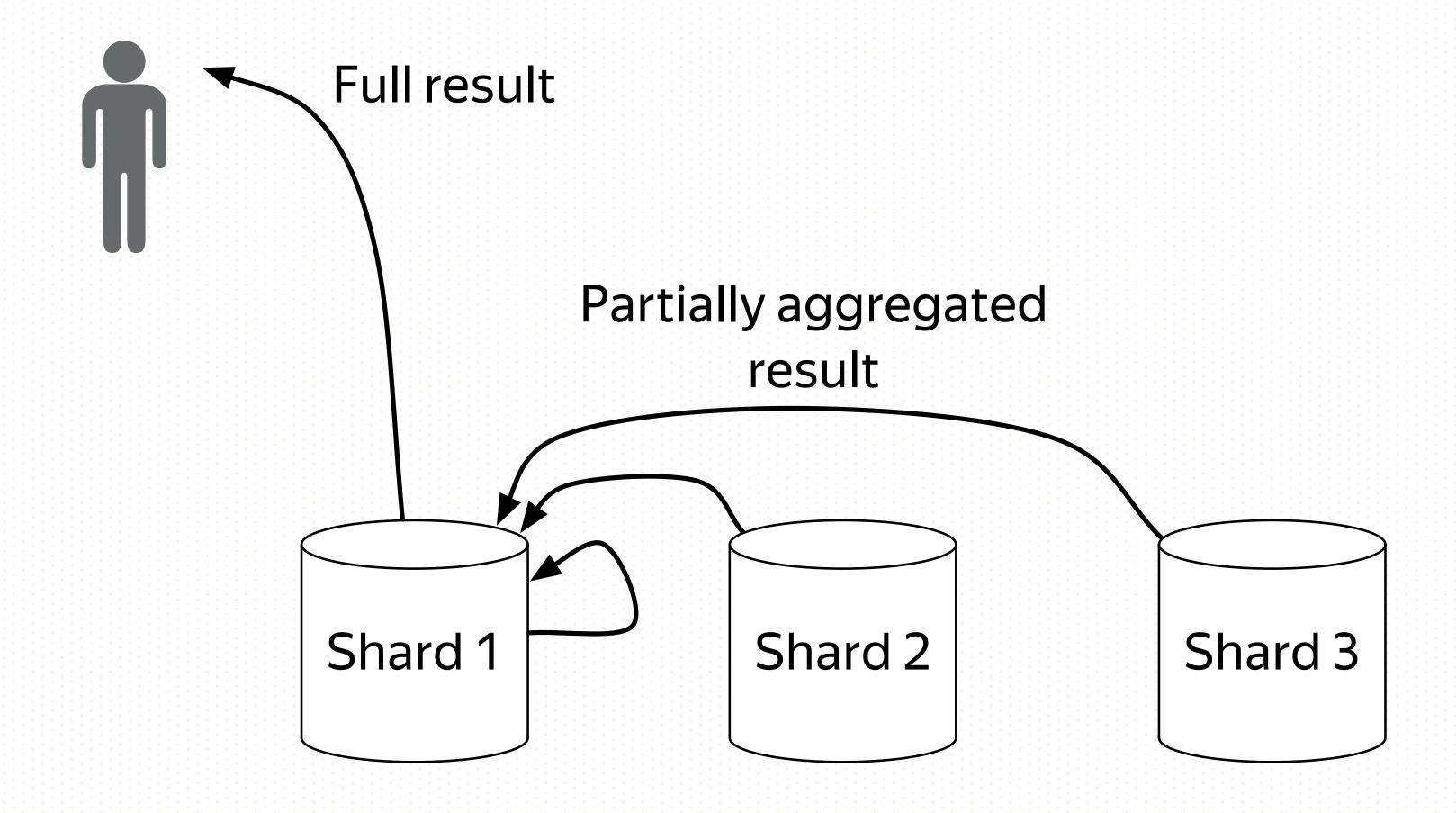


Reading from a Distributed table





Reading from a Distributed table





NYC taxi benchmark

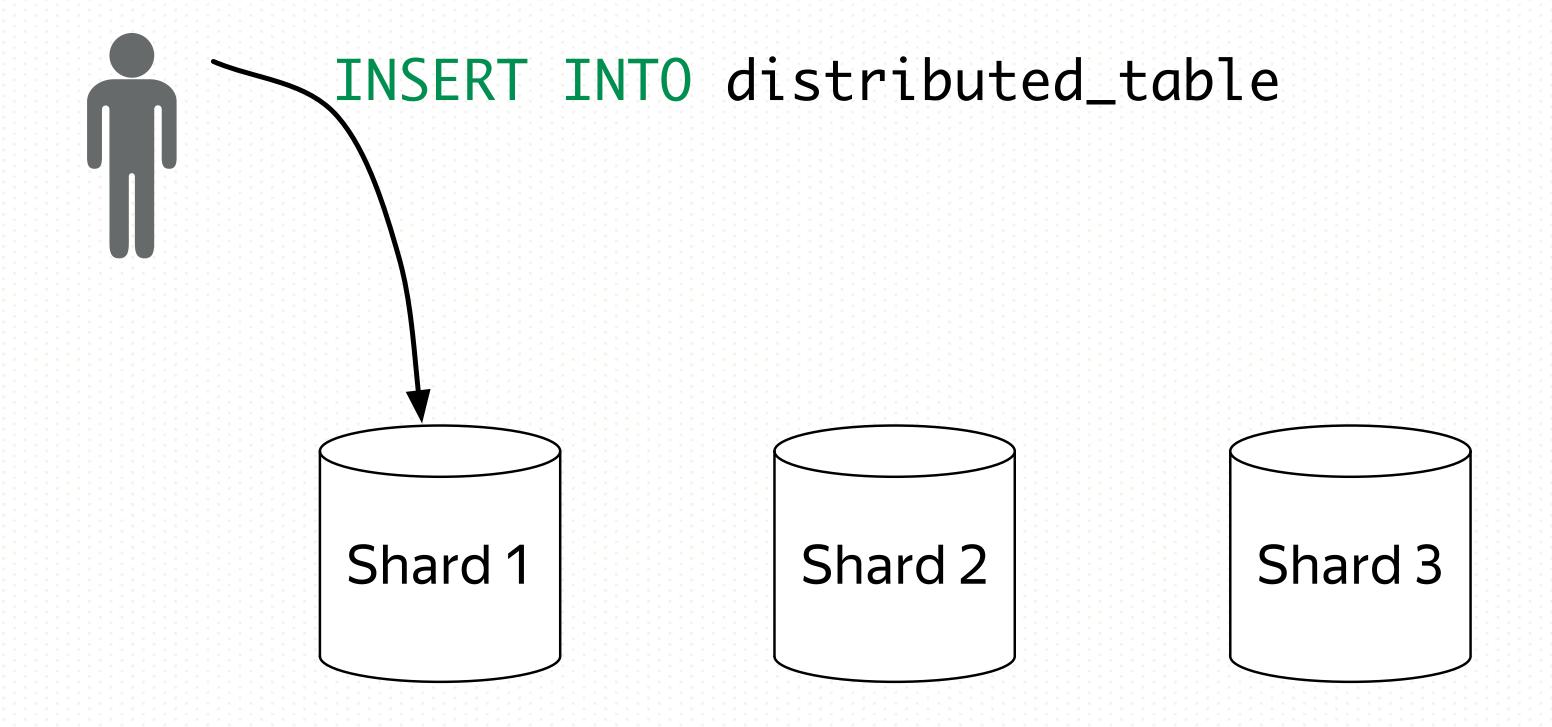
CSV 227 Gb, ~1.3 bln rows

SELECT passenger_count, avg(total_amount)
FROM trips GROUP BY passenger_count

Shards	1	3	140
Time, s.	1,224	0,438	0,043
Speedup		x2.8	x28.5



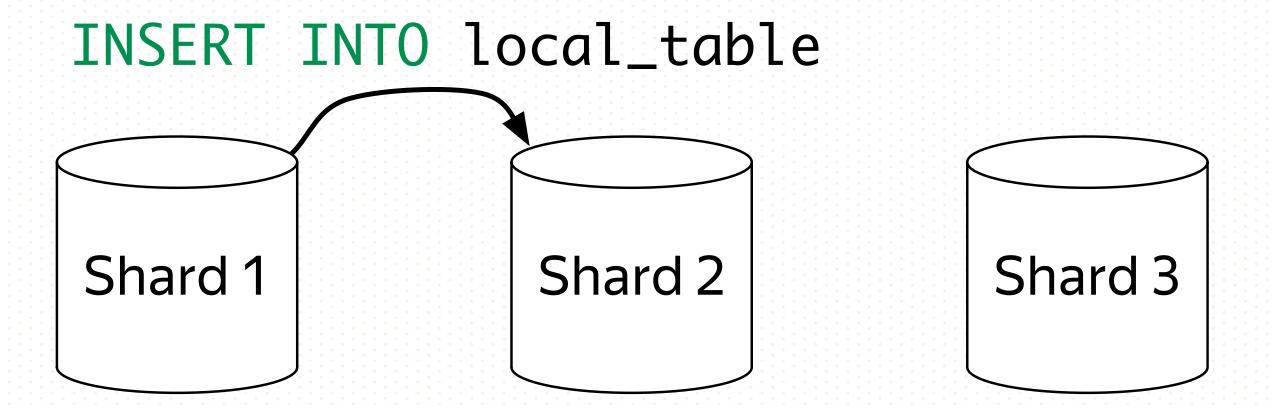
Inserting into a Distributed table





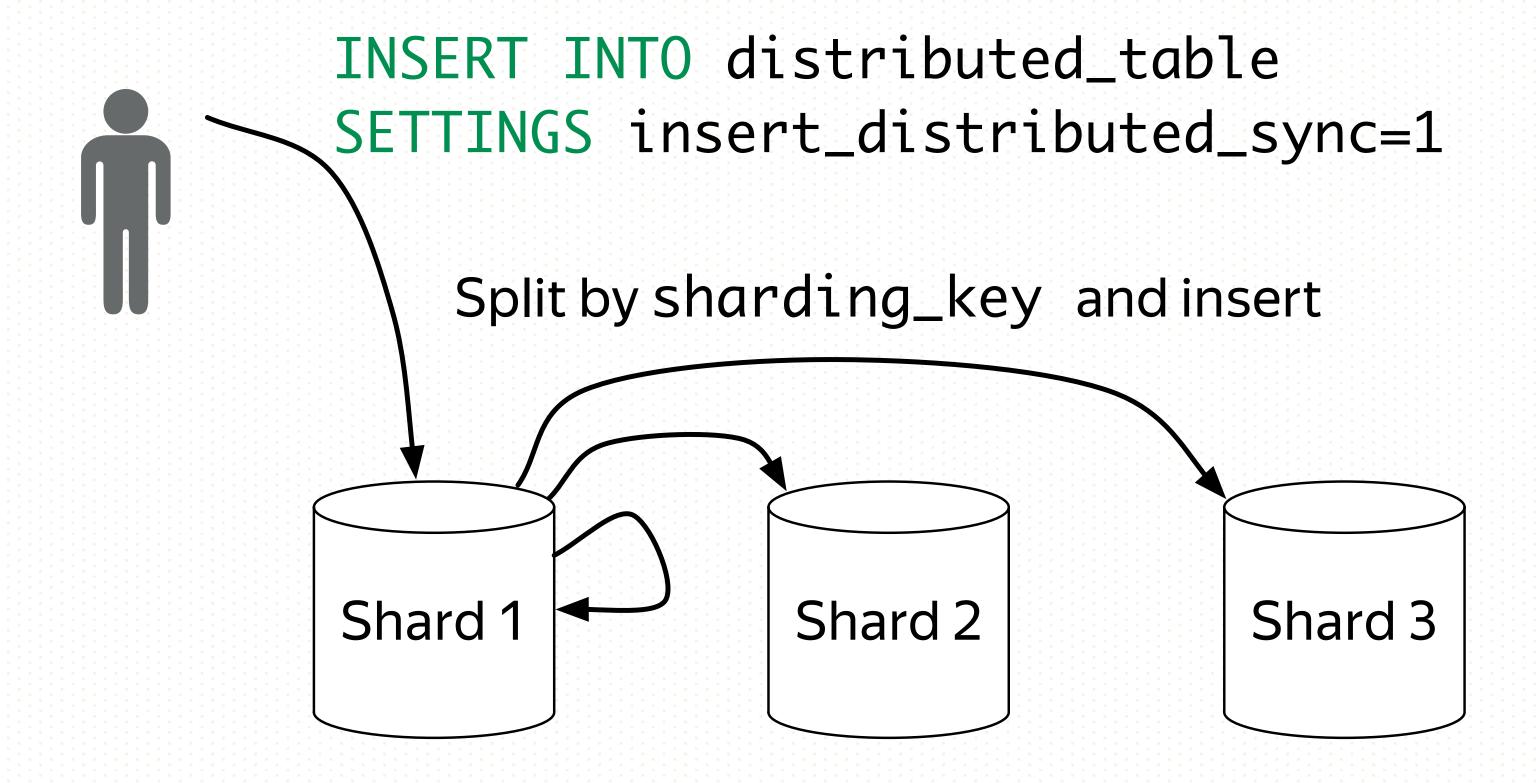
Inserting into a Distributed table

Async insert into shard # sharding_key % 3





Inserting into a Distributed table





Things to remember about Distributed tables

- It is just a view
 - > Doesn't store any data by itself
- Will always query all shards
- Ensure that the data is divided into shards uniformly
 - > either by inserting directly into local tables
 - or let the Distributed table do it
 (but beware of async inserts by default)



When failure is not an option

- > Protection against hardware failure
- Data must be always available for reading and writing

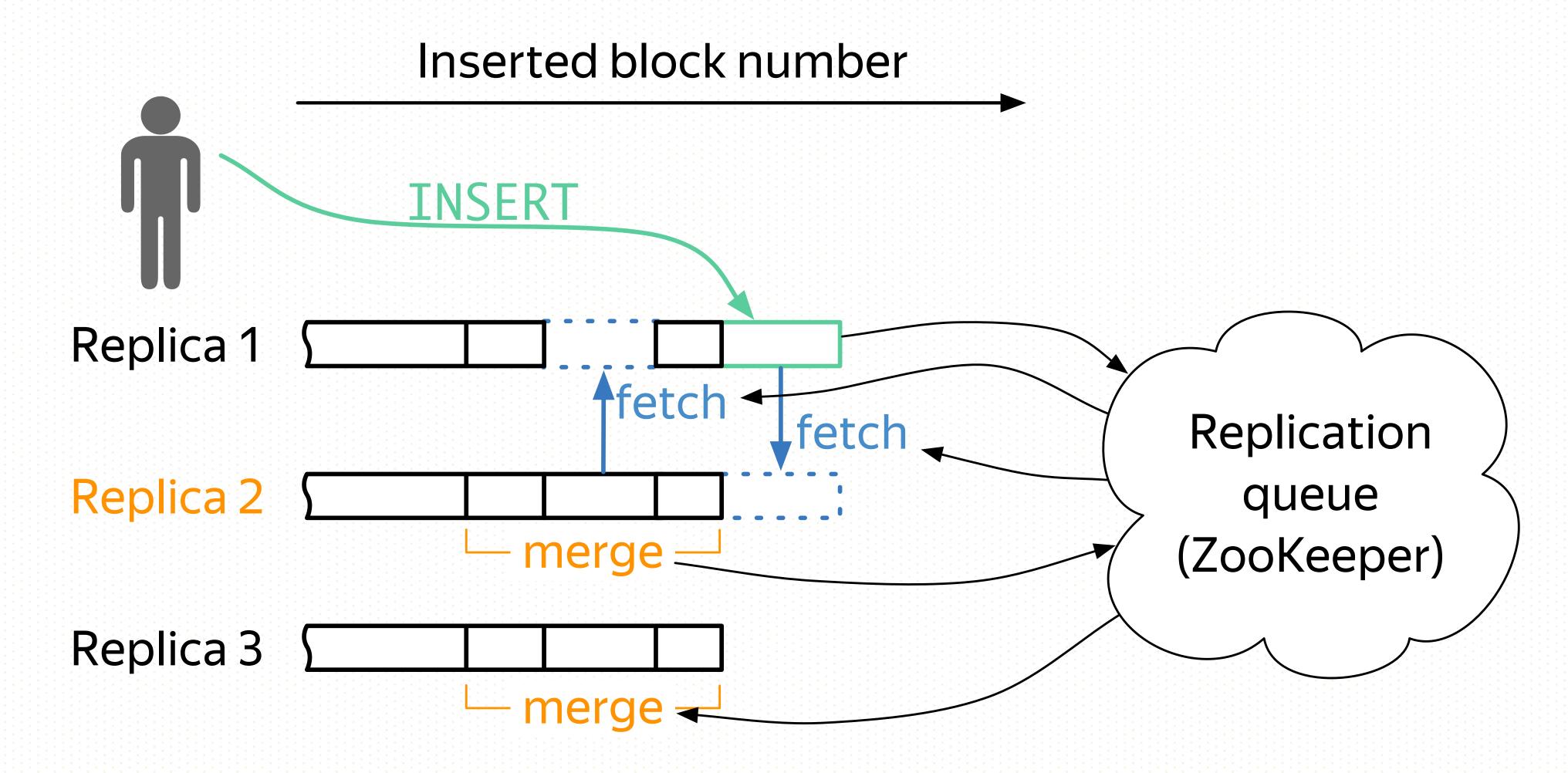


When failure is not an option

- > Protection against hardware failure
- Data must be always available for reading and writing
- ClickHouse: ReplicatedMergeTree engine!
- > Async master-master replication
- > Works on per-table basis



Replication internals



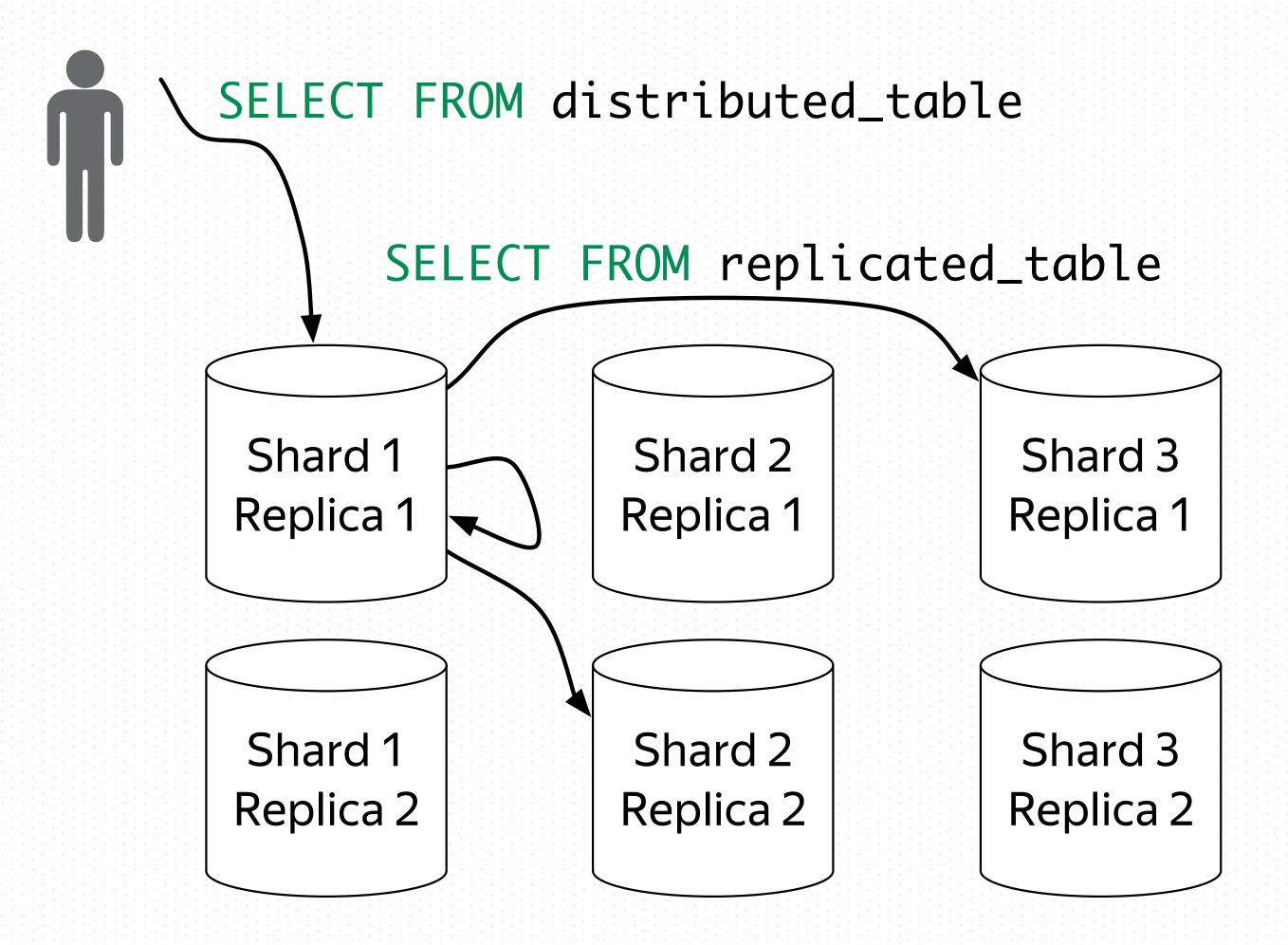


Replication and the CAP-theorem

- What happens in case of network failure (partition)?
- Not consistent
 As is any system with async replication
 But you can turn linearizability on
- Highly available (almost) Tolerates the failure of one datacenter, if ClickHouse replicas are in min 2 DCs and ZK replicas are in 3 DCs.
 - *A server partitioned from ZK quorum is unavailable for writes



Putting it all together





Things to remember about replication

Use it!

- > Replicas check each other
- > Unsure if INSERT went through?
 Simply retry the blocks will be deduplicated
- ZooKeeper needed, but only for INSERTs (No added latency for SELECTs)

Monitor replica lag

> system.replicas and system.replication_queue tables are your friends



Brief recap

- > Column-oriented
- > Fast interactive queries on real time data
- > SQL dialect + extensions
- Bad fit for OLTP, Key-Value, blob storage
- Scales linearly
- > Fault tolerant
- > Open source!



Thank you

Start using ClickHouse today!

- Questions? Or reach us at:
 - > clickhouse-feedback@yandex-team.com
 - > Telegram: https://t.me/clickhouse_en
 - > GitHub: https://github.com/yandex/ClickHouse/
 - > Google group: https://groups.google.com/group/clickhouse

