Dynamic Scheduling of Network Updates

Xin Jin

Hongqiang Harry Liu, Rohan Gandhi, Srikanth Kandula, Ratul Mahajan, Ming Zhang, Jennifer Rexford, Roger Wattenhofer









SDN: Paradigm Shift in Networking

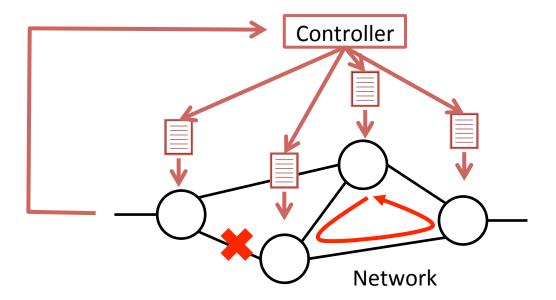
 Direct, centralized updates of forwarding rules in switches



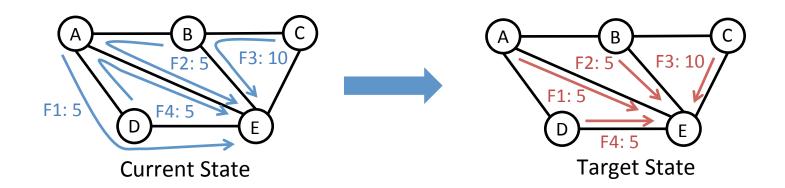
- Many benefits
 - Traffic engineering [B4, SWAN]
 - Flow scheduling [Hedera, DevoFlow]
 - Access control [Ethane, vCRIB]
 - Device power management [ElasticTree]

Network Update is Challenging

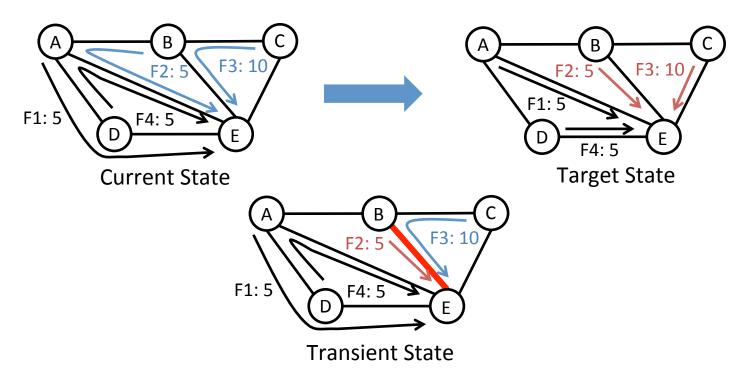
- Requirement 1: fast
 - The agility of control loop
- Requirement 2: consistent
 - No congestion, no blackhole, no loop, etc.



What is Consistent Network Update



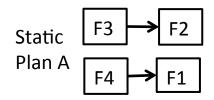
What is Consistent Network Update

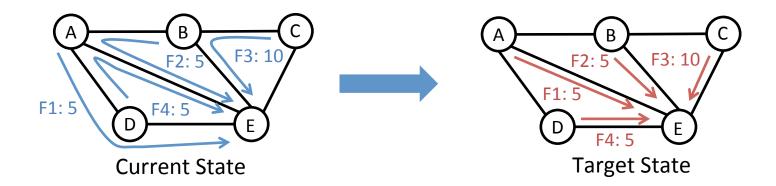


- Asynchronous updates can cause congestion
- Need to carefully order update operations

Existing Solutions are Slow

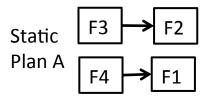
- Existing solutions are static [ConsistentUpdate'12, SWAN'13, zUpdate'13]
 - Pre-compute an order for update operations





Existing Solutions are Slow

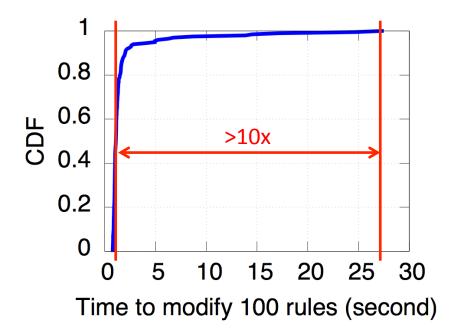
- Existing solutions are static [ConsistentUpdate'12, SWAN'13, zUpdate'13]
 - Pre-compute an order for update operations



- Downside: Do not adapt to runtime conditions
 - Slow in face of highly variable operation completion time

Operation Completion Times are Highly Variable

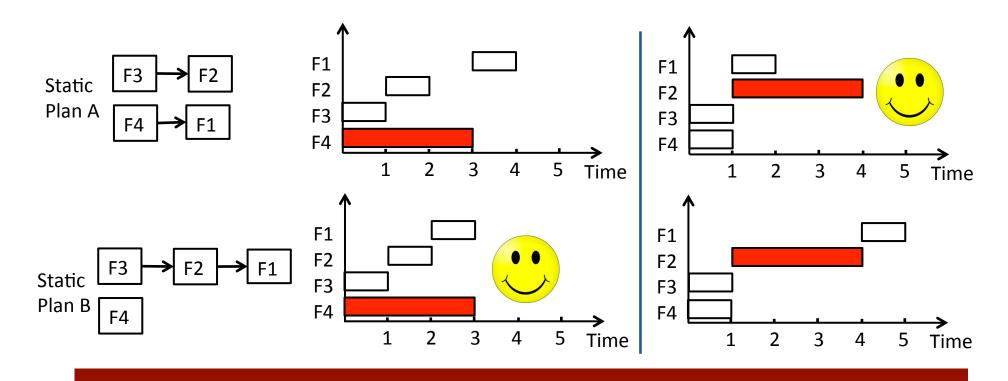
Measurement on commodity switches



Operation Completion Times are Highly Variable

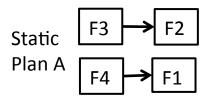
- Measurement on commodity switches
- Contributing factors
 - Control-plane load
 - Number of rules
 - Priority of rules
 - Type of operations (insert vs. modify)

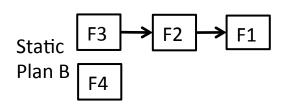
Static Schedules can be Slow

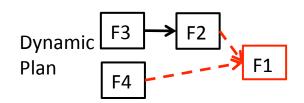


No static schedule is a clear winner under all conditions!

Dynamic Schedules are Adaptive and Fast







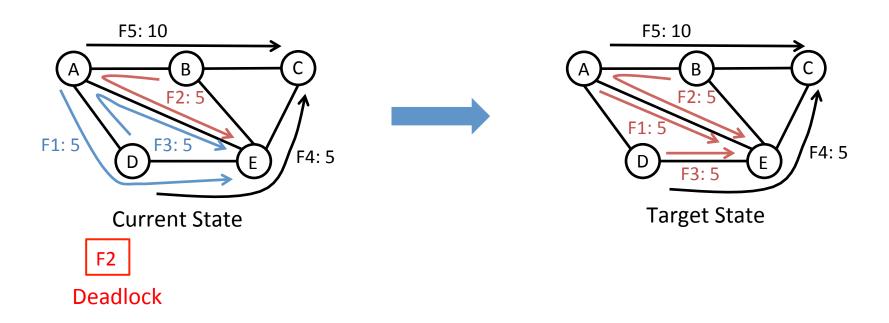
Adapts to actual conditions!

No static schedule is a clear winner under all conditions!

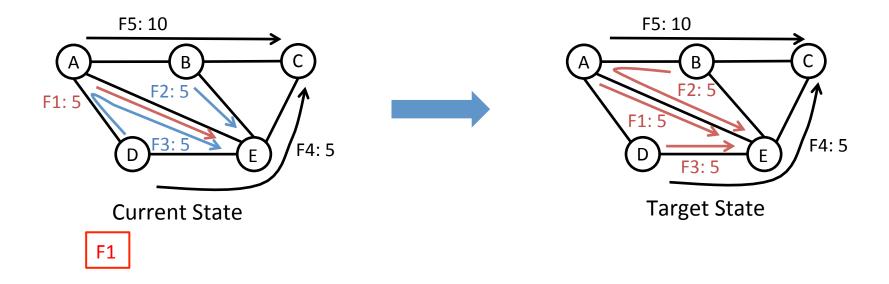
- Exponential number of orderings
- Cannot completely avoid planning



- Exponential number of orderings
- Cannot completely avoid planning



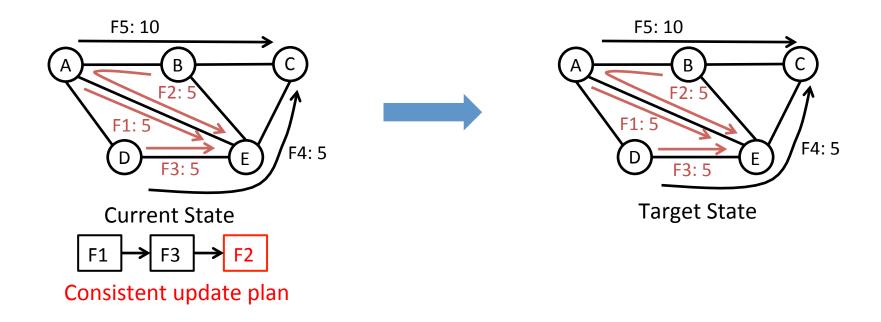
- Exponential number of orderings
- Cannot completely avoid planning



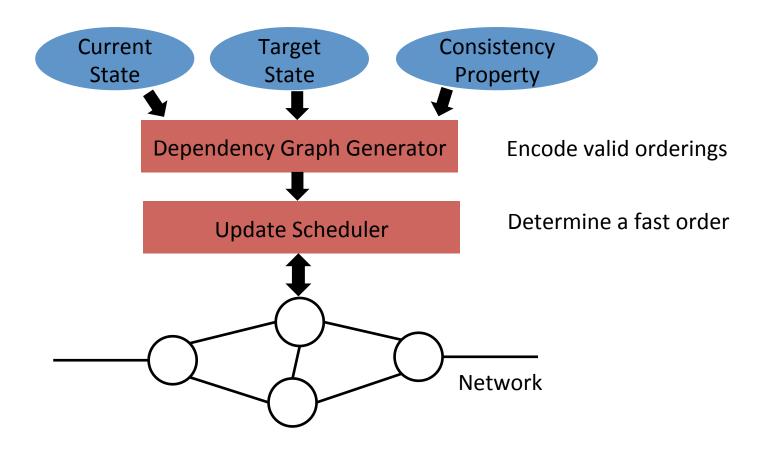
- Exponential number of orderings
- Cannot completely avoid planning

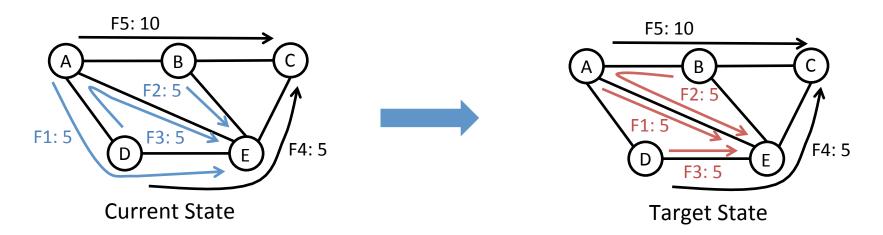


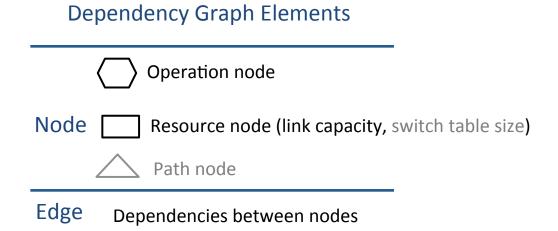
- Exponential number of orderings
- Cannot completely avoid planning

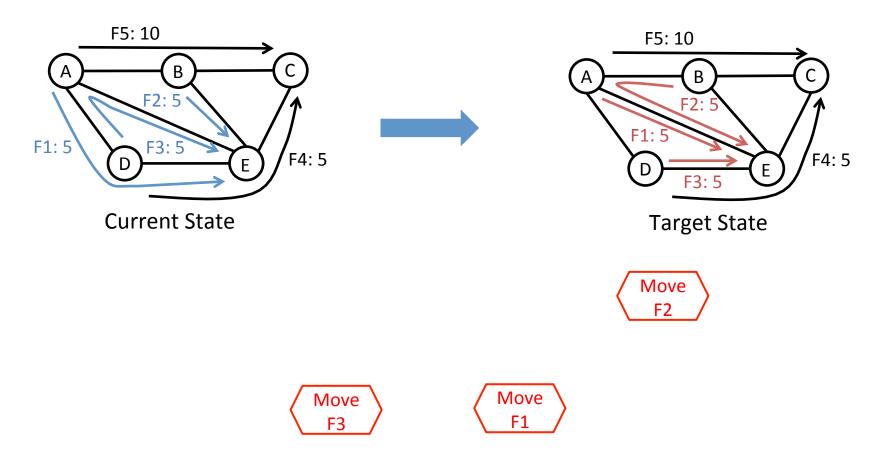


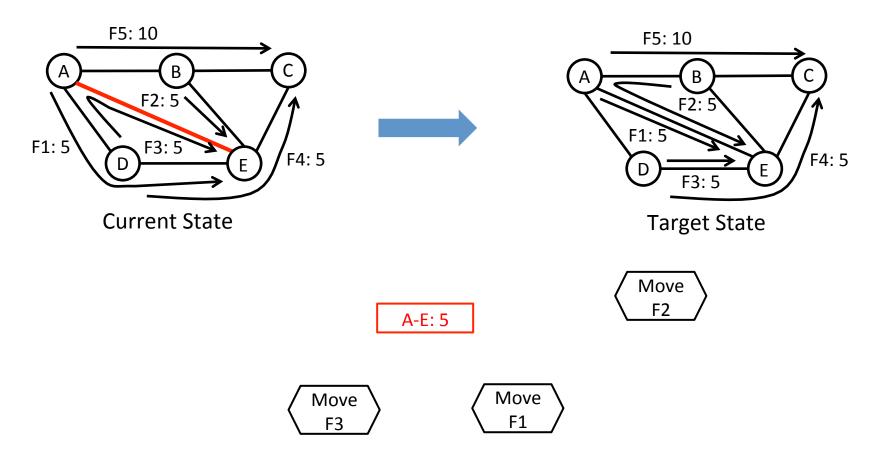
Dionysus Pipeline

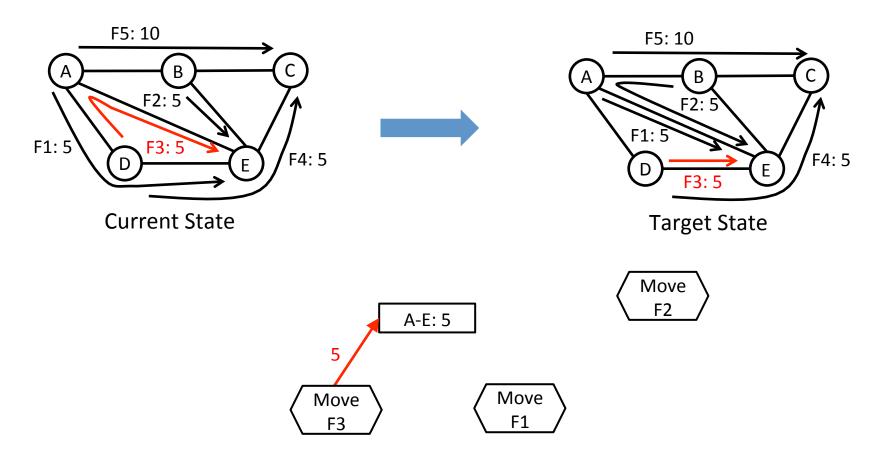


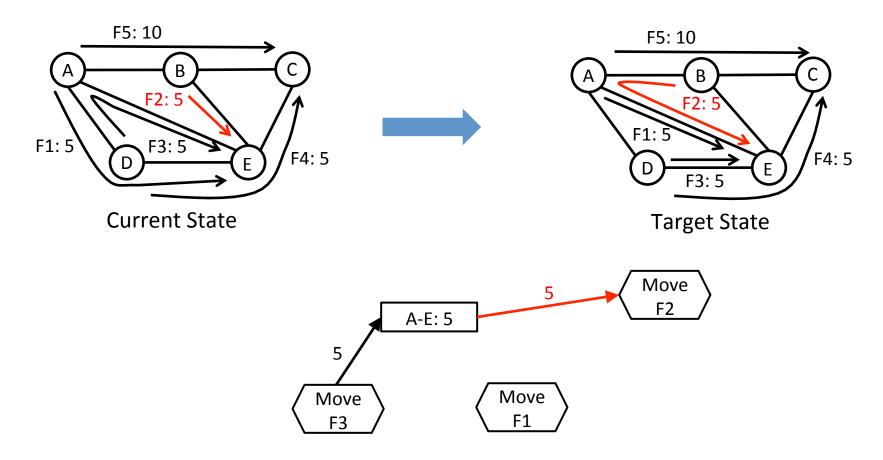


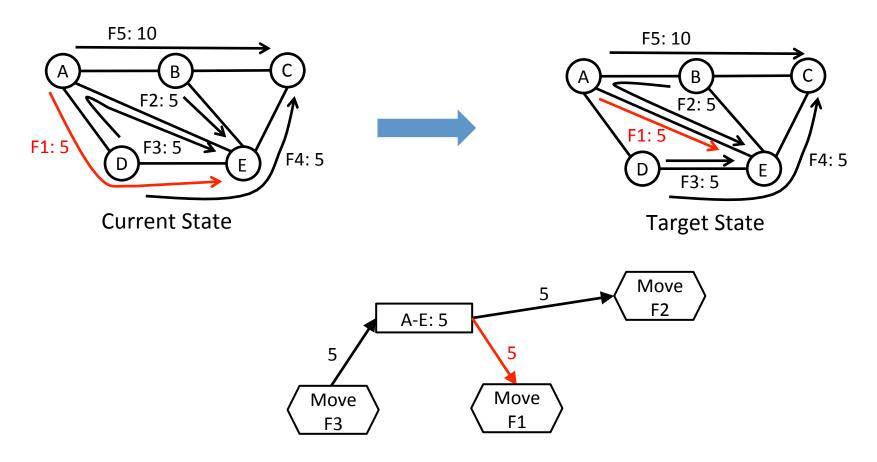


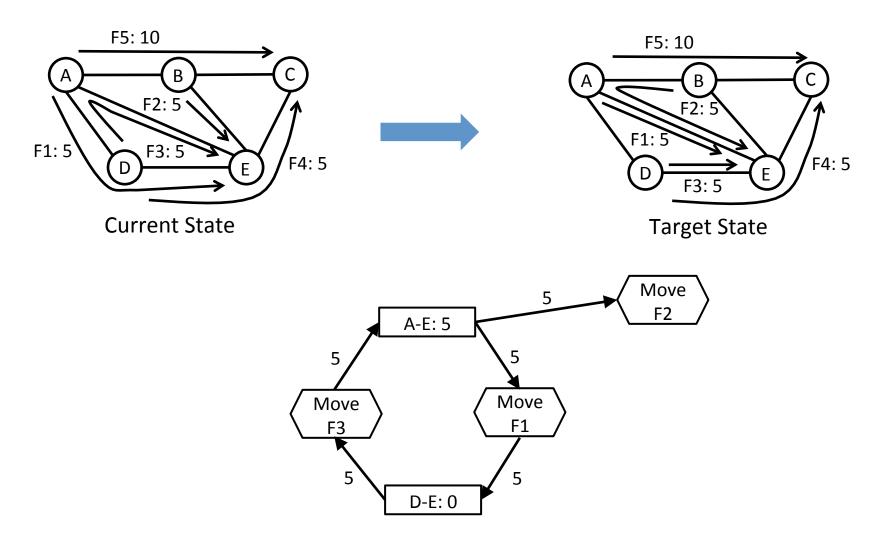






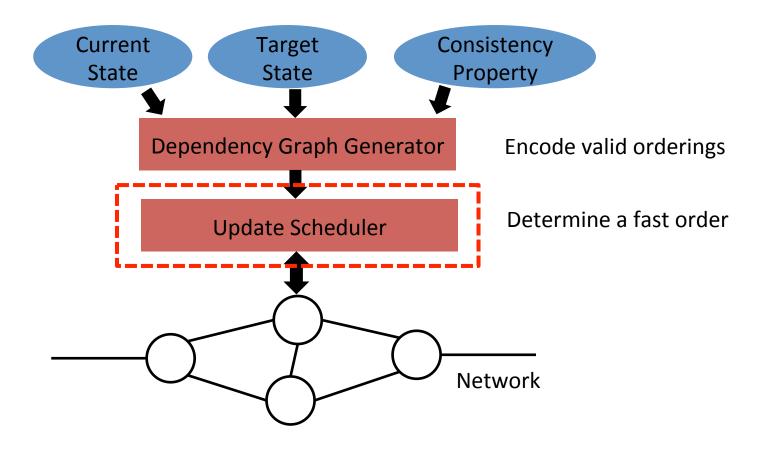


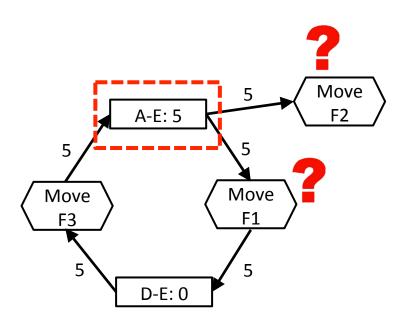


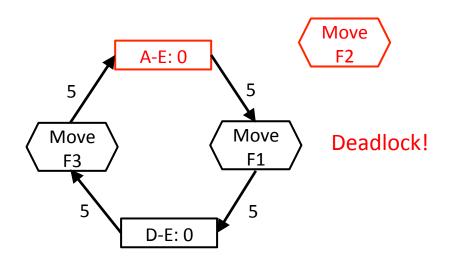


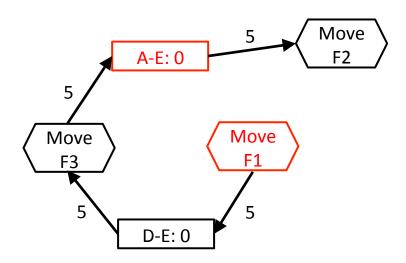
- Supported scenarios
 - Tunnel-based forwarding: WANs
 - WCMP forwarding: data center networks
- Supported consistency properties
 - Loop freedom
 - Blackhole freedom
 - Packet coherence
 - Congestion freedom
- Check paper for details

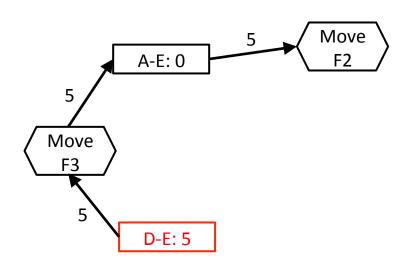
Dionysus Pipeline

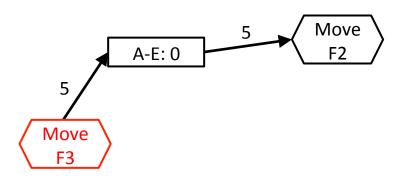


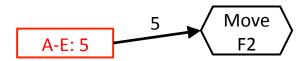














- Scheduling as a resource allocation problem
- NP-complete problems under link capacity and switch table size constraints
- Approach
 - DAG: always feasible, critical-path scheduling
 - General case: covert to a virtual DAG
 - Rate limit flows to resolve deadlocks

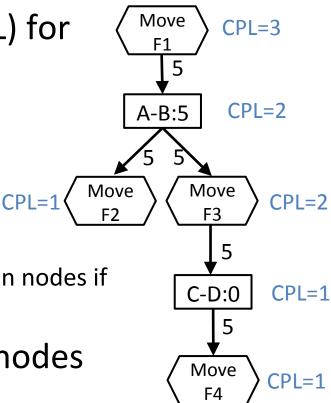
Critical-Path Scheduling

 Calculate critical-path length (CPL) for each node

$$CPL_i = w_i + \max_{j \in children(i)} CPL_j$$

 $w_i = \begin{cases} 1, & if i \text{ is operation node} \\ 0, & otherwise \end{cases}$

- Extension: assign larger weight to operation nodes if we know in advance the switch is slow
- Resource allocated to operation nodes with larger CPLs



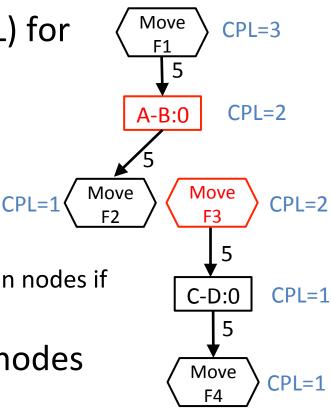
Critical-Path Scheduling

 Calculate critical-path length (CPL) for each node

$$CPL_i = w_i + \max_{j \in children(i)} CPL_j$$

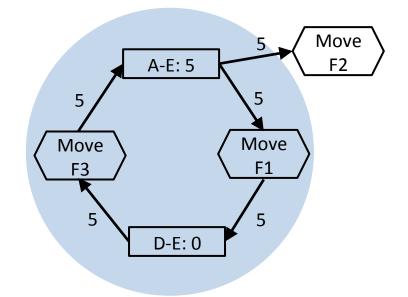
 $w_i = \begin{cases} 1, & \text{if } i \text{ is operation node} \\ 0, & \text{otherwise} \end{cases}$

- Extension: assign larger weight to operation nodes if we know in advance the switch is slow
- Resource allocated to operation nodes with larger CPLs

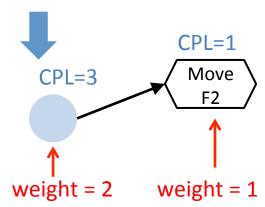


Handling Cycles

- Convert to virtual DAG
 - Consider each strongly connected component (SCC) as a virtual node

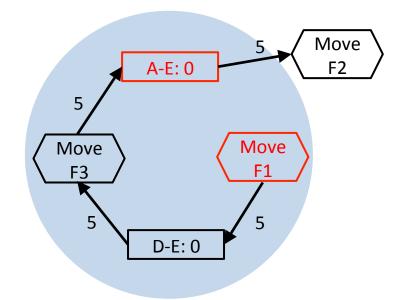


- Critical-path scheduling on virtual DAG
 - Weight w_i of SCC: number of operation nodes

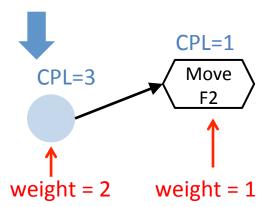


Handling Cycles

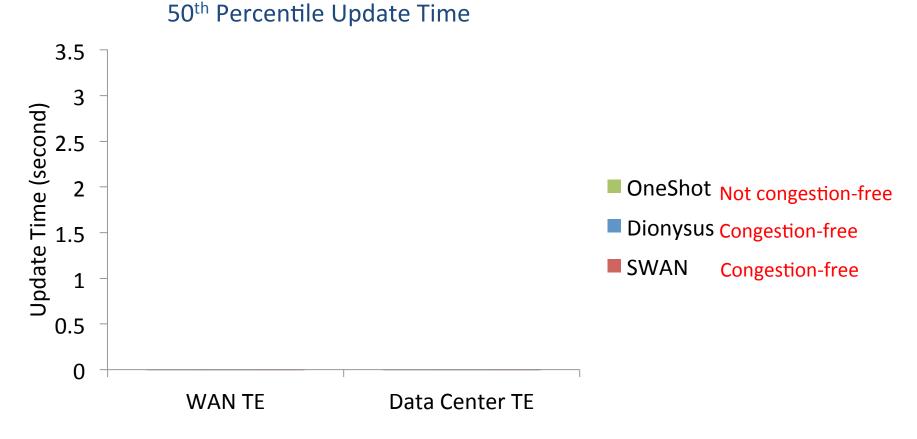
- Convert to virtual DAG
 - Consider each strongly connected component (SCC) as a virtual node



- Critical-path scheduling on virtual DAG
 - Weight w_i of SCC: number of operation nodes

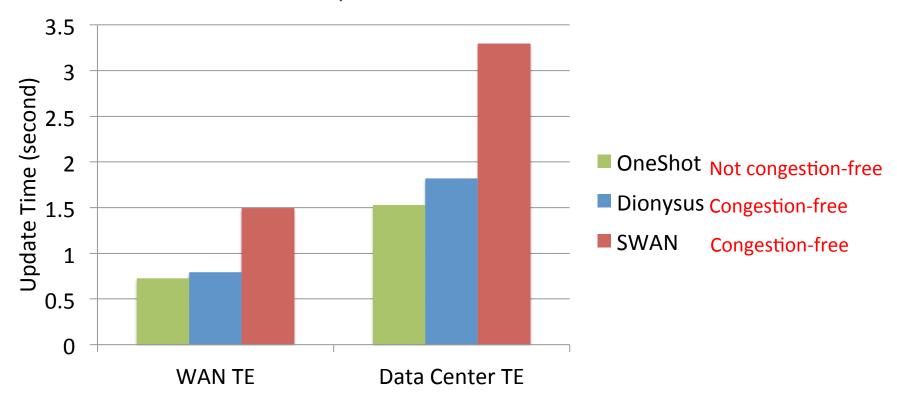


Evaluation: Traffic Engineering



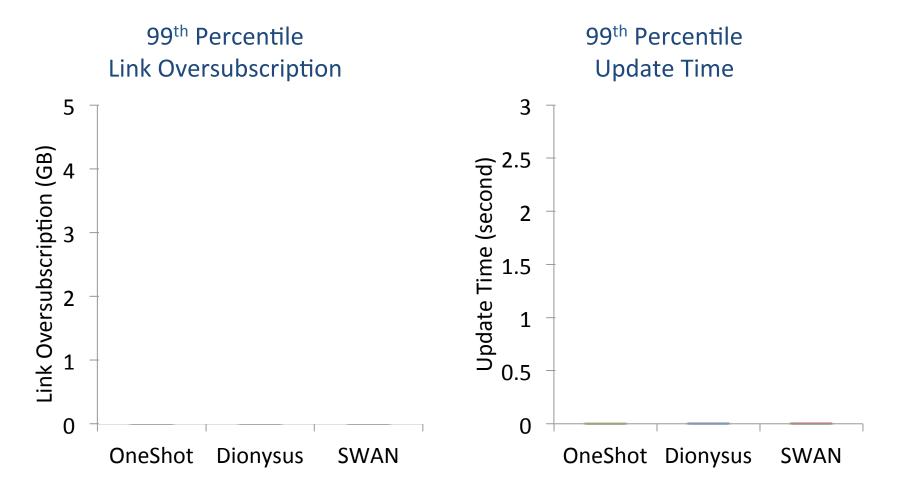
Evaluation: Traffic Engineering

50th Percentile Update Time

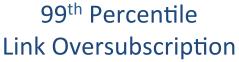


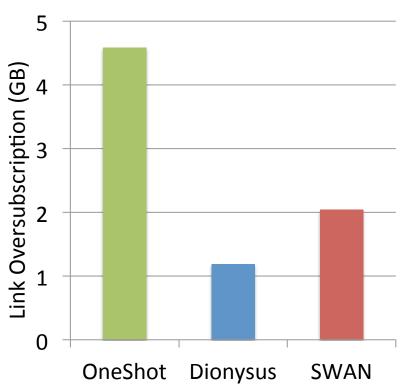
Improve 50th percentile update speed by 80% compared to static scheduling (SWAN), close to OneShot

Evaluation: Failure Recovery

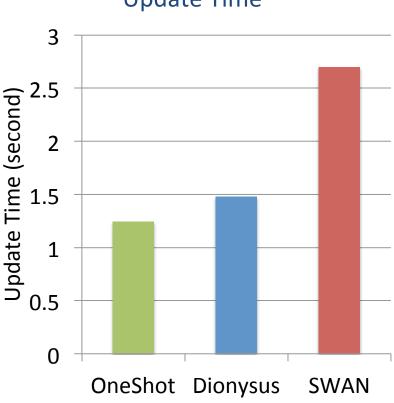


Evaluation: Failure Recovery





99th Percentile Update Time



Reduce 99th percentile link oversubscription by 40% compared to static scheduling (SWAN)

Improve 99th percentile update speed by 80% compared to static scheduling (SWAN)

Conclusion

- Dionysus provides fast, consistent network updates through dynamic scheduling
 - Dependency graph: compactly encode orderings
 - Scheduling: dynamically schedule operations

Dionysus enables more agile SDN control loops

Thanks!

