

Tetris

Multi-Resource Packing for Cluster Schedulers

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Performance of cluster schedulers

We find that:

- Resources are **fragmented** i.e. machines run below capacity
- Even at 100% usage, goodput is smaller due to **over-allocation**
- Pareto-efficient multi-resource fair schemes do not lead to good avg. performance

Tetris

Up to 40% improvement in makespan¹ and job completion time with near-perfect fairness

¹Time to finish a set of jobs

Findings from Bing and Facebook traces analysis

Applications have (very) diverse resource needs

- Tasks need *varying amounts of each resource*
- *Demands* for resources *are weakly correlated*

Multiple resources become tight

This matters, because no single bottleneck resource in the cluster:

- E.g., enough cross-rack network bandwidth to use all cores

Upper bound on potential gains

- *Makespan reduces by $\approx 49\%$*
- *Avg. job completion time reduces by $\approx 46\%$*

Why so bad #1

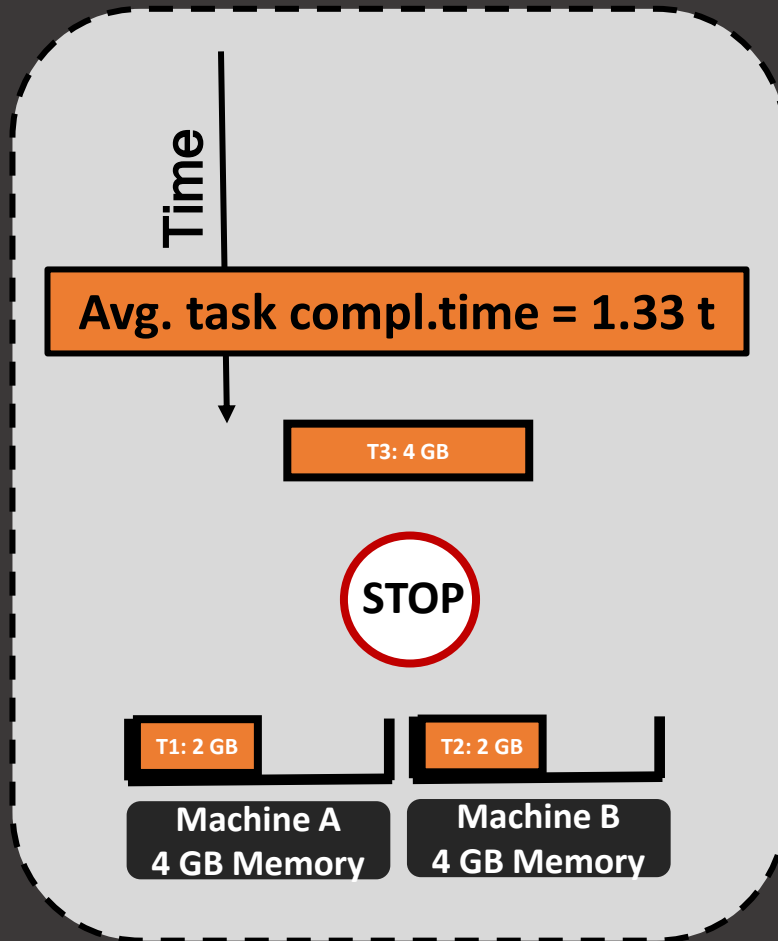
*Production schedulers **neither pack tasks nor consider all their relevant resource demands***

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graph TD; A["Production schedulers neither pack tasks nor consider all their relevant resource demands"] --> B["#1 Resource Fragmentation"]; A --> C["#2 Over-allocation"]
```

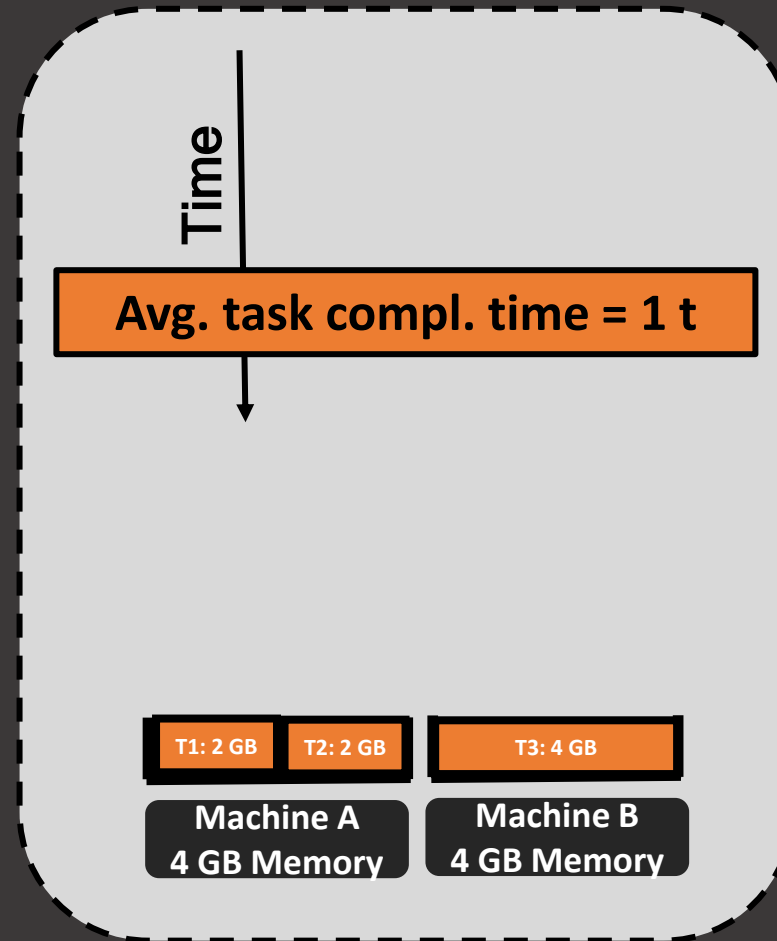
#1 Resource Fragmentation

#2 Over-allocation

Resource Fragmentation (RF)



Current Schedulers



"Packer" Scheduler

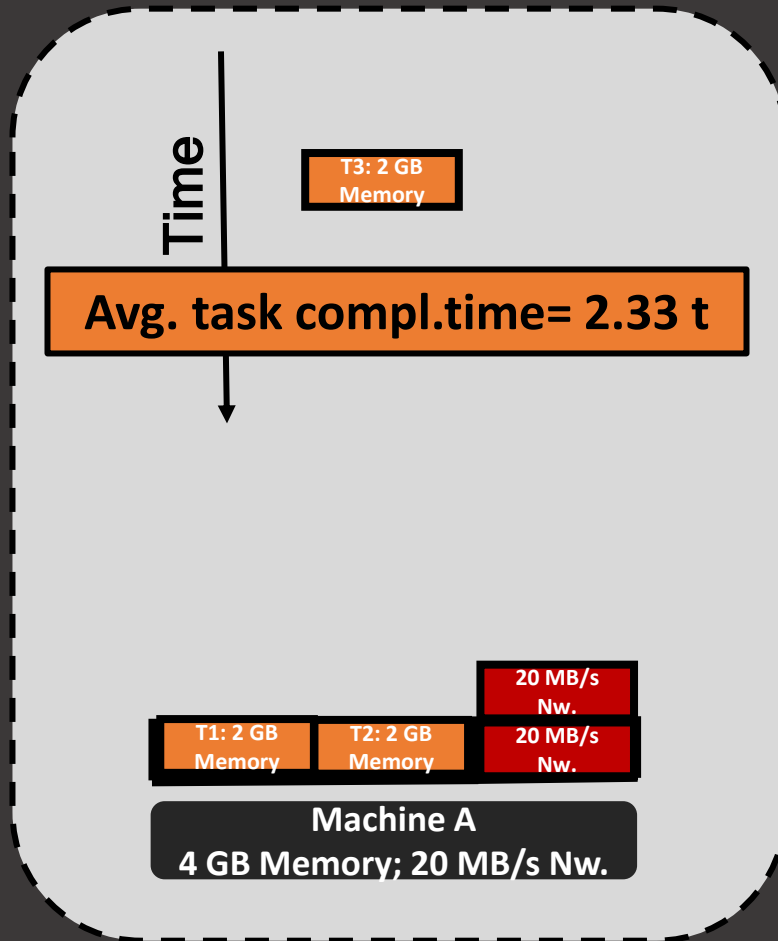
Current Schedulers

Are not explicit about packing.

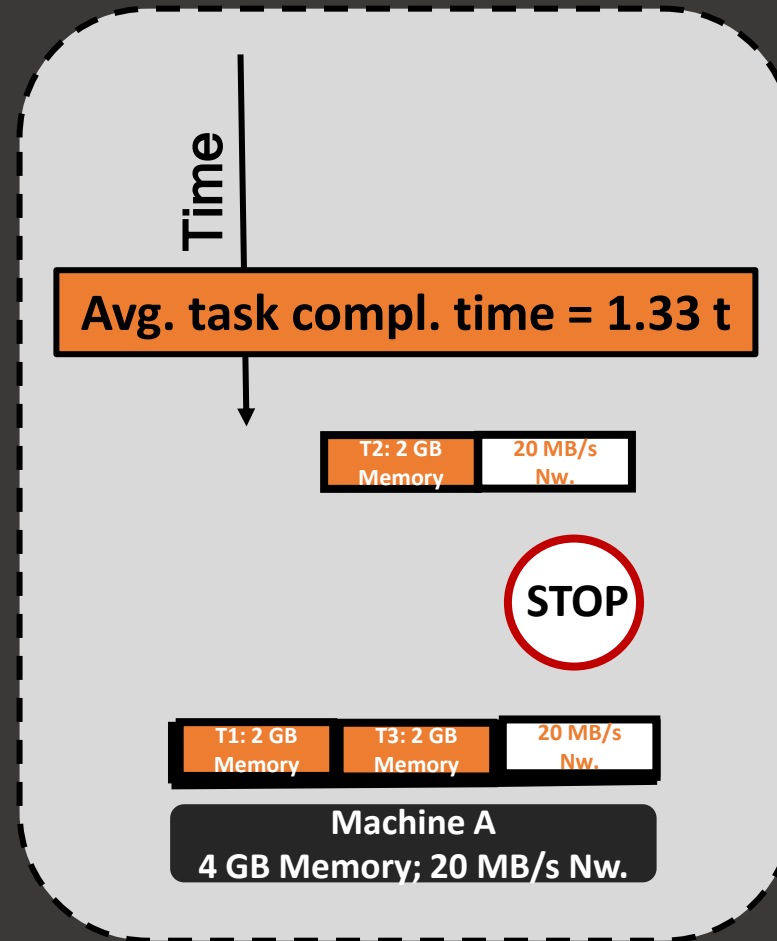
Allocate resources per slots, fairness.

RF increase with the number of resources being allocated !

Over-Allocation



Current Schedulers



"Packer" Scheduler

Current Schedulers

Not all of the resources are explicitly allocated

E.g., disk and network can be over-allocated

Why so bad #2

Multi-resource Fairness Schemes do not solve the problem

Example in paper

Packer vs. DRF: makespan and avg. completion time improve by over 30%

Work Conserving != no fragmentation, over-allocation

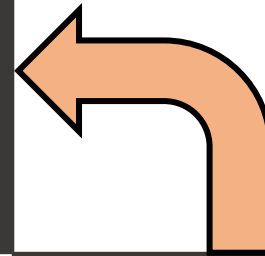
Pareto¹ efficient != performant

- Treat cluster as a big bag of resources
 - ❑ *Hides the impact of resource fragmentation*
- Assume job has a fixed resource profile
 - ❑ *Different tasks in the same job have different demands*
 - ❑ *How the job is scheduled impacts jobs' current resource profiles*
 - ❑ *Can schedule to create complementarity*

¹no job can increase its share without decreasing the share of another

Current Schedulers

1. Resource Fragmentation
2. Over-Allocation
3. Fair allocations sacrifice performance



Competing objectives

Cluster efficiency

vs.

Job completion time

vs.

Fairness

Tetris



1

Pack tasks along multiple resources to improve
cluster efficiency and reduce makespan

Theory

Multi-Resource Packing of Tasks
similar to
Multi-Dimensional Bin Packing

APX-Hard¹

Avoiding fragmentation looks like:

- ☐ Tight bin packing
- ☐ Reduce # of bins → **reduce makespan**

¹APX-Hard is a strict subset of NP-hard
Balls could be tasks
Bin could be machine, time

Practice

Existing heuristics do not directly apply:

- ☐ Assume balls of a fixed size
- ☐ Assume balls are known apriori



- vary with time / machine placed
- elastic
- cope with online arrival of jobs, dependencies, cluster activity

Tetris

A packing heuristic

- Tasks resources demand vector



- Machine resource vector

Alignment score (A)

“A” works because:

1. Check for fit to ensure **no over-allocation**

✓ Over-Allocation

2. Bigger balls *get bigger scores*

✓ Resource Fragmentation

3. Abundant resources *used first*

1

Packing heuristic

Tetris



2

Faster average job completion time

CHALLENGE

Q: What is the shortest “~~remaining time~~” ?

“remaining work” = $\frac{\text{remaining \# tasks}}{\text{tasks' resource demands} + \text{tasks' durations}}$

A job completion time heuristic

- Gives a score **P** to every job
- Extended SRTF to incorporate multiple resources

2

Job Completion Time Heuristic

Shortest Remaining Time First¹ (SRTF)

schedules jobs in ascending order of their remaining time

¹SRTF – M. Harchol-Balter et al. Connection Scheduling in Web Servers [USITS'99]

CHALLENGE

A: delays job completion time

Packing
Efficiency

?

Completion
Time

P: loss in packing efficiency

Combine A and P scores !

```
1: among J runnable jobs
2: score (j) = A(t, R) + ε P(j)
3:    max task t in j, demand(t) ≤ R (resources free)
4: pick j*, t* = argmax score(j)
```

2

Job Completion Time Heuristic

Tetris

A small cluster of Tetris blocks is positioned above the letter 'i' in the word 'Tetris'. The blocks include an orange horizontal tetromino (I-tetromino) at the top, a red vertical tetromino (I-tetromino) to the left, a black L-shaped tetromino (L-tetromino) in the center, and a grey Z-shaped tetromino (Z-tetromino) to the right.

3

Achieve performance and fairness

Tetris

Performance and fairness do not mix well in general

But

We can get “perfect fairness” and much better performance

- **Packer** says: “*task T should go next to improve **packing efficiency***”
- **SRTF** says: “schedule *job J to improve **avg. completion time***”
- **Fairness** says: “*this set of **jobs** should be scheduled next*”

Possible to satisfy all three
In fact, happens often in practice

3

Fairness Heuristic

Tetris

Fairness is not a tight constraint

- Lose a bit of **fairness** for a lot of **gains in performance**
- **Long term fairness** not short term fairness

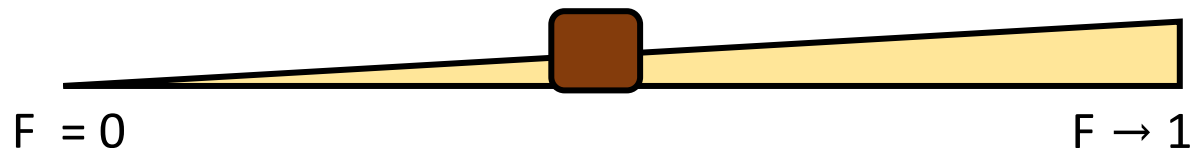
Heuristic

- **Fairness Knob**, $F \in [0, 1)$

Pick the **best-for-perf.** task from among
 $\lceil 1 - F \rceil$ fraction of jobs furthest from fair share

- ☐ Most unfair
- ☐ Most efficient scheduling,

Close to perfect fairness



3

Fairness Heuristic

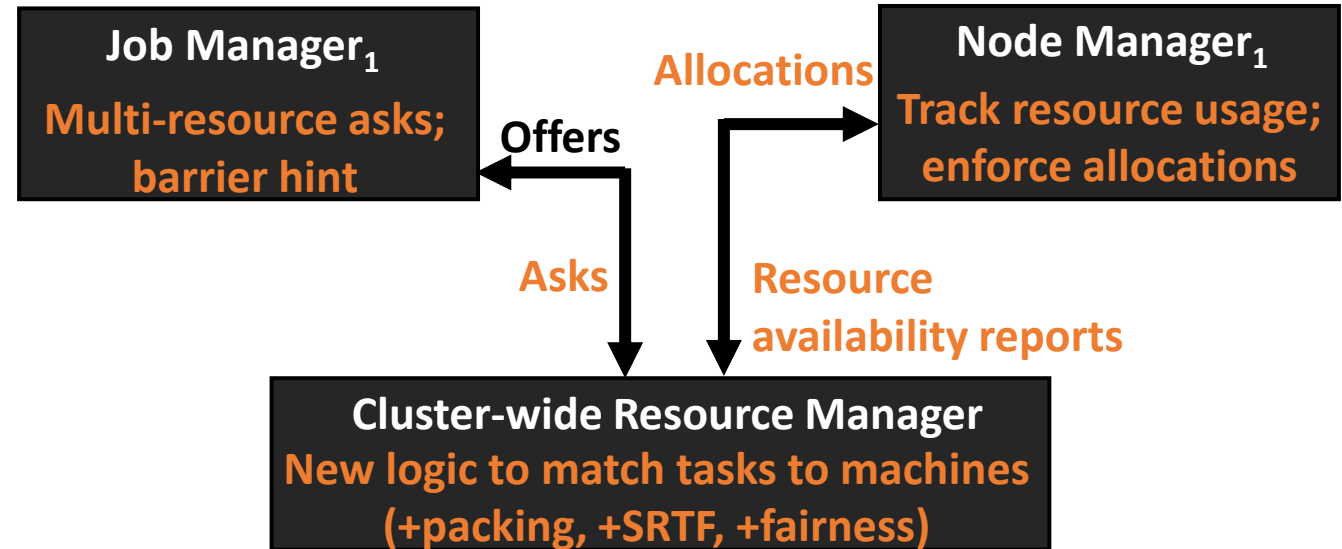
Putting it all together

We saw:

- **Packing** efficiency
- **Prefer small remaining work**
- **Fairness knob**

Other things in the paper:

- Estimate **task demands**
- Deal with **inaccuracies, barriers**
- Other **cluster activities**



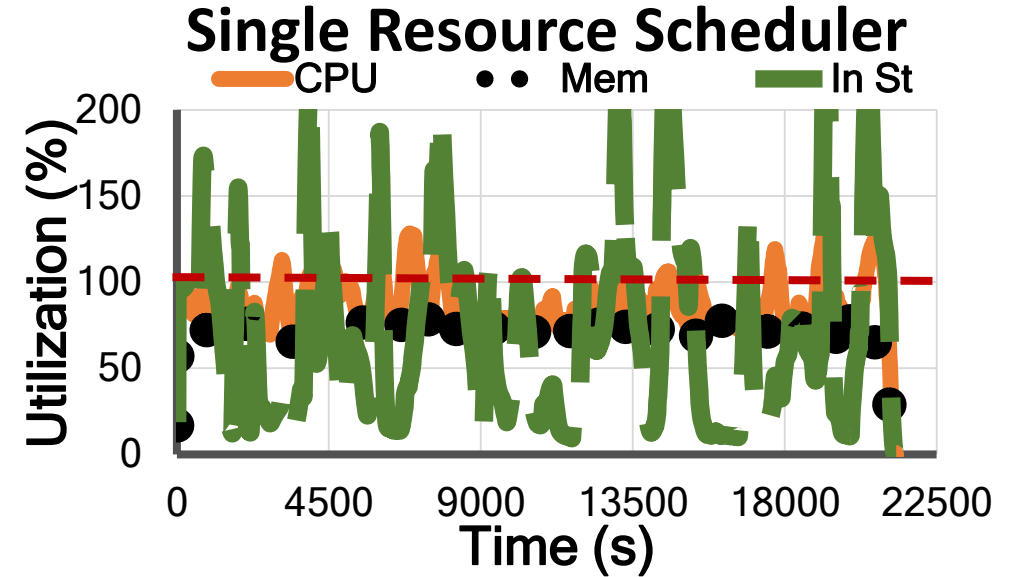
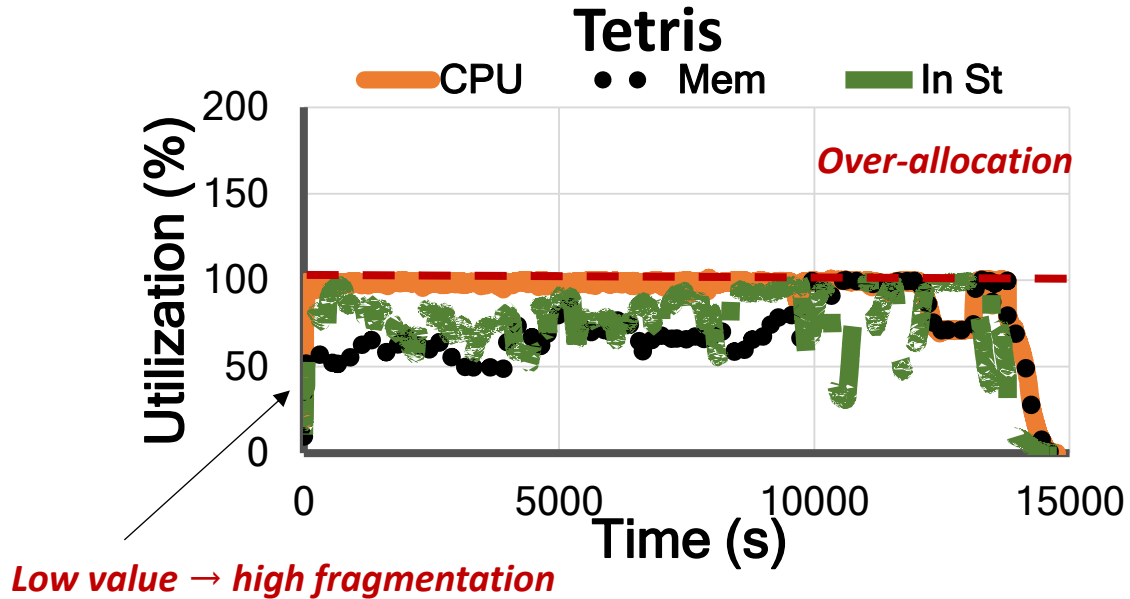
Yarn architecture

Changes to add Tetris(shown in orange)

Evaluation

- Implemented in Yarn 2.4
- 250 machine cluster deployment
- Bing and Facebook workload

Efficiency



Tetris vs.	Makespan	Avg. Job Compl. Time
Single Resource Scheduler	29 %	30 %
Multi-resource Scheduler	28 %	35 %

Gains from

- avoiding *fragmentation*
- avoiding *over-allocation*

Fairness

Fairness Knob

- *quantifies the extent to which Tetris adheres to fair allocation*

	Makespan	Job Compl. Time	Avg. Slowdown [over impacted jobs]
No Fairness $F = 0$	50 %	40 %	25 %
Full Fairness $F \rightarrow 1$	10 %	23 %	2 %
$F = 0.25$	25 %	35 %	5 %

Tetris



Pack efficiently
along multiple
resources

Prefer jobs
with less
“remaining
work”

Incorporate
Fairness

- Combine heuristics that *improve packing efficiency* with those that *lower* average *job completion time*
- Achieving desired amounts of *fairness can coexist* with improving cluster performance

- Imp
sho

We are working towards a Yarn check-in
<http://research.microsoft.com/en-us/UM/redmond/projects/tetris/>