	15/02/22						
	Tutorial - 06 Page No.:						
	5.5 Date: / /						
	U19CS009						
	Barrens Branch Prost Hold says emoners						
8-1	Explain two pass assembles functions with example.						
And -5	Pass 2 assemblez sequires 2 scans of program to						
	generate machine code						
-3	SIT was data structure defined by passi like symbol tabl						
25 OV. 1	a Bull will paid with the disease a patholic						
	Functioning of two pass assembles						
	P						
3107	Pars!						
	-> separate symbol, mnemonic operade, and operand fields						
	-s build symbol table						
	-> perform Il processing -> Construct intermediate representation (IC)						
-)							
	-> synthesis targer program.						
	Data structure						
ar emba	1) and structure						
1.1	tay nonetups, the second secon						
Noise	100 100 4 100 100 100 100 100 100 100 10						
	The state of the s						
Source	de la set de la						
	Pars II - Target						
35 3 A 1 3 A	beoded w						
	Potermedrate - s Data access						
DAVA.	code sata access -s corrtrol transfer						
•	Eg: - Source code LC IC By Pass 1						
	MOVER AREGIX 100 (15,04) (200)						
LI	ADD BREW, ONE 101 (15:01) (PU102) (S(1)						

	Const No.
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	APD CREG, TEN 102 (15,01) (RG,03) (5,2)
	STOP 103 (15,00)
X	DC '5' (04 (PLO) (C15)
ONE	= DC 11 (05 (D):01) (C1)
TEN	DL 110' 106 (DL,01) (C,10)
	END (AD, 02)
	SYMTAB => INDEX SYMBOLNAME Address
	0 x. 104
	ONE 105
	2 TEN 106
->	Machine code by Pass 2
LC	IC by Pass! Machine code
	(ADIOI) (C(100)
100	(15/04) (RG101) (3/0) 04 01 104
	(15,01) (R07,02) (S,1) 01 02 105
102	(15,01) (PC2103) (S,1) 01 03 106
103	(15,00) 00 00 000
104	(DLig) (Cis)
105	(DLOI) (C(1) 7 They don't change
106	(DLOI) (CII) They don't change (DLOI) (CID) In forming machine code
	(AD, 02)
	3 -5 stands for symbol.
	C -> stemas for constant.

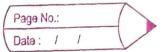
1 0-2 What are some advantages of marking assembly languages over high level language? Ang o Debugging and verifying -> looking at compiler generated assembly code as their die assembly window for a debugger is useful for.
finding exposs and for checking how well a compiler
optimizes a particular prece of code. O Making compilers -> understanding assembly coding is must for making compilers, debuggers and other development tools © Embedded System > small embedded system has fewer Schources. Assembly programming i's necessary for optimizing code for speed or size. O Mardware d'svers 4: System code : accersing hardware and system control registers etc. may sometimes be difficult or impossible with high revel language. O Self-modifying code -> It generally not profitable because it interferes with efficreat code caching. It may be advourtageous for example to include a small compiler in moth programs where a user defined function has to be calculated multiple times

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03	what tools are used for compiler construction?
M ()	Parser Generator -> it provides syntax analyzers form the input that is based on a
	Form the input that is based on a
THE AL	grammatical description of programming language or on a context-free grammar. It is weful as the syntax
1	a context-free grammar. It is weful as the syntax
	analyze phase is highly complex and consumes more
	nanual and compilation time, token
3.3	photograph is common a bound of the property
	(FCr >> Parse >> syntax
i kirone	grammaz analy Zez
	The state will be the specific
	J
6)4	Parse tree
	Parse tree
, A1,	Contraction of the contraction o
, A1,	Scanner Generator > It generates lexical analyzes from Proput that consists of regular express-
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, A1,	Scanner Generator > It generates lexical analyzer from Proput that consists of regular express- pon description based on tokens of a language. It generate
, A1,	Scanner Generator > it generates lexical analyzer from Proput that consists of regular express- por description based on tokens of a language. It generate a finite automation to recognize regular expression.
, A1,	Scanner Generator > it generates lesscal analyzer from Paput that consists of regular express- por description based on tokens of a language. It generate a finite automation to recognize regular expression. Source program
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	Scanner Generator > it generates lexical analyzer from Paput that consists of regular expression Por description based on tokens of a larguage. It generate a finite automation to recognize regular expression Source program Sepcifications of -> scannar -> lexical regular-expression generator analyzer
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	Scanner Generator > it generates lexical analyzer from Paput that consists of regular expression Por description based on tokens of a larguage. It generate a finite automation to recognize regular expression Source program Sepcifications of -> scannar -> lexical regular-expression generator analyzer

code with three address formats from the Propert

that consists of a parse tree. These engines have

routines to focuerse the parse tree and then



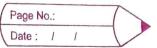
produces the intermediate code. In this, each node of parse tree is associated with one or more (4) Automatic code generator -> Pt generates machine language for a target machine. Each operation of intermediate language is translated using a collection of rules and then is taken as an input by code generator. A template mat-Ching machine process is used. An intermediate language statement is replaced by it's equivalent machine language statement using template. 5 Data flow analysis engine 7 it is used in code optimization. Data flow analysis is key part of the code optimization that gather the information that is the value that flow from one part of a program to another.

6 Compiler Construction toolkit -> it provides an Prtegrated set of moutines that aids in building compiler components cer in construction of various phases of compilers.

on helps in full implementation of high level programming languages 0 +> supports optimization for computer architecture

parallellen o -> Design of new memory-hierarchy of machines · + > Toronslating programs

Is used with other software productivity tods.



	Date: / /
03-5	D/B a marso and submoutine. And explain mar so definition
	and expansion using an example
Ang o	Both permit a group of instructions to be defined as a
	single entity with a unique given label or name called up
	when needed.
0	A subsputine is called by BSR or JSR instruction, while
	a macro is called by simply using it's name.
0	macros are simple to write and use
	Macro are faster
	Marco are not substitute of subsolutine -> as marco is
	substituted with the code and additional code is generated
\	every time a macro is called . Very long macro that
	are used many times in a program will result in an
	enormous expansion of code size.
D	Macroso can be called only in program it is defined, subsou-
	Ene can be called from other priegram also.
	112,16,41
	Macro -s A macro definition is enclosed between a
	marso header statement and a marso end
	statement.
-	-> macro definitions are typically located at the start
	of the program.
	Marno definition consists of parameter
	·
7	o macro prototype - declares name of macro 4 14 ind od,
	o one or more model -s from which cusemply language is
	o preprocessor statements generated
	sued to perform auxillary function
	cluring marro expansion.

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Ø	macro expansion			
	macro expansion <pre>/macro name> [<pormal parameter="" space="">[]]</pormal></pre>			
	The production of the control of the			
, 0	< formal parameter spaces 's of form			
	52			
	<pre><pre><pre></pre></pre></pre>			
	The state of the s			
-	Macro Expansion			
-P				
	a majoro call leads to majoro expansion, in which			
	macro call statement is seplaced by a sequence of			
	assembly statement.			
->	9			
	INCR AIB, AREG			
	MACRO CALL			
	00.000			
	MACRO I MACRO			
A 15	INCR AMEM-VAL LINC-VAL, 4 REG			
	MOVER 4REG 4MEM-VAL			
20	APP 4 REG 4 INC-VAL			
7	MOVEM 4 REG 4 MEM-VAL			
3	MEND			
	+ MOVER AREGO A			
	+ ADD AREGO B			
1.1	+ MOVEM ARELY A			