

Systems Security

COMSM1500

Buffer overflow



Assembly language

Small refresher

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`_start`

- `.text` *# code segment*
- `.global _start` *# export*
- `_start:`
- `movl $0x01, %eax` *# exit*
- `movl $0x00, %ebx` *# return code*
- `int $0x80` *# syscall*

hello world

- *# int write(int fd, char* buf, int len)*
- movl *\$0x04*, %eax *# write*
- movl *\$0x01*, %ebx *# stdout*
- movl *\$str*, %ecx *# buffer*
- movl *\$14*, %edx *# length*
- int *\$0x80*
- .data
- *str*:
- .ascii "Hello, World!\n"

Stack

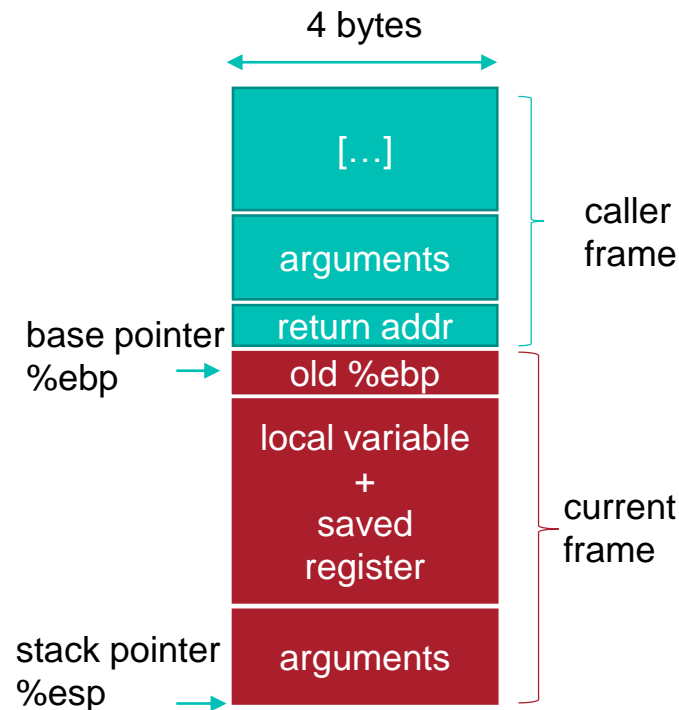
Small refresher

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The stack

- Current stack frame (“top” to bottom)
 - arguments for function about to be called
 - saved register context (if reused)
 - local variable
 - old base pointer
- Caller’s stack frame
 - return address
 - pushed by call instruction
 - Arguments for this call



Example

```
▪ int zip1 = 15213;
▪ int zip2 = 98915;

▪ void call_swap() {
▪     swap(&zip1, &zip2);
▪ }

▪ void swap(int *xp, int *yp) {
▪     int t0 = *xp;
▪     int t1 = *yp;
▪     *xp = t1;
▪     *yp = t0;
▪ }
```


Example

- `int zip1 = 15213;`
- `int zip2 = 98915;`
- `void call_swap() {`
- `swap(&zip1, &zip2);`
- `}`

Example

- **int** zip1 = 15213;
- **int** zip2 = 98915;
- **void** call_swap() {
- swap(&zip1, &zip2);
- }

- *# void call_swap()*
- ...
- pushl \$zip1
- pushl \$zip2
- call swap
- ...

Example

- `int zip1 = 15213;`
- `int zip2 = 98915;`
- `void call_swap() {`
- `swap(&zip1, &zip2);`
- `}`

- `# void call_swap()`

- ...

- `pushl $zip1`

- `pushl $zip2`

- `call swap`

- ...

`%ebp` →

`%esp` →



Example

```
▪ int zip1 = 15213;
▪ int zip2 = 98915;

▪ void call_swap() {
  ▪ swap(&zip1, &zip2);
  ▪ }
```

```
▪ # void call_swap()
```

```
▪ ...
```

```
▪ pushl $zip1
```

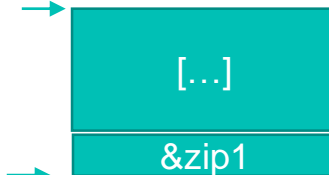
```
▪ pushl $zip2
```

```
▪ call swap
```

```
▪ ...
```

%ebp →

%esp →



Example

```
▪ int zip1 = 15213;
▪ int zip2 = 98915;

▪ void call_swap() {
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▪ }
```

```
▪ # void call_swap()
```

```
▪ ...
```

```
▪ pushl $zip1
```

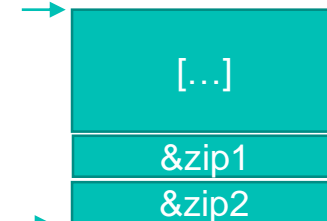
```
▪ pushl $zip2
```

```
▪ call swap
```

```
▪ ...
```

%ebp →

%esp →



Example

- `int zip1 = 15213;`
- `int zip2 = 98915;`
- `void call_swap() {`
 - `swap(&zip1, &zip2);`
 - `}`

- `# void call_swap()`

- ...

- `pushl $zip1`

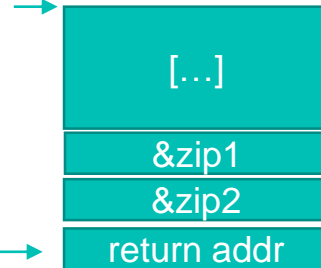
- `pushl $zip2`

- `call swap`

- ...

`%ebp` →

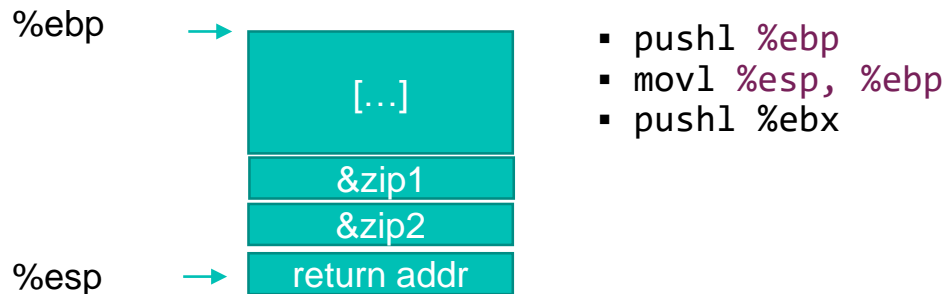
`%esp` →



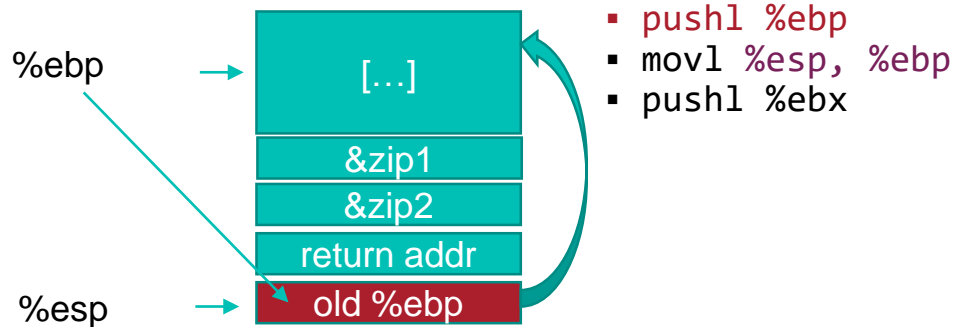
Example

```
▪ void swap(int *xp, int *yp) {  
  ▪   int t0 = *xp;  
  ▪   int t1 = *yp;  
  ▪   *xp = t1;  
  ▪   *yp = t0;  
  ▪ }  
  
▪ # void swap(int *xp, int *yp)  
  
  ▪ pushl %ebp  
  ▪ movl %esp, %ebp  
  ▪ pushl %ebx  
  } Set up  
  
  ▪ movl 12(%ebp), %ecx  
  ▪ movl 8(%ebp), %edx  
  ▪ movl (%ecx), %eax  
  ▪ movl (%edx), %ebx  
  ▪ movl %eax, (%edx)  
  ▪ movl %ebx, (%ecx)  
  } body  
  
  ▪ movl -4(%ebp), %ebx  
  ▪ movl %ebp, %esp  
  ▪ popl %ebp  
  ▪ ret  
  } finish
```

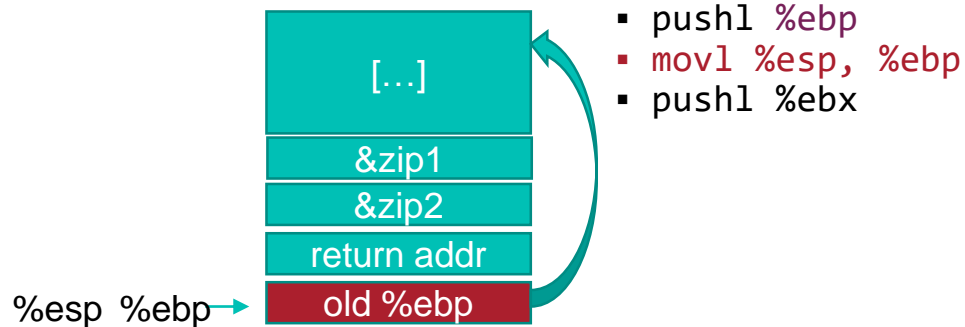
Example



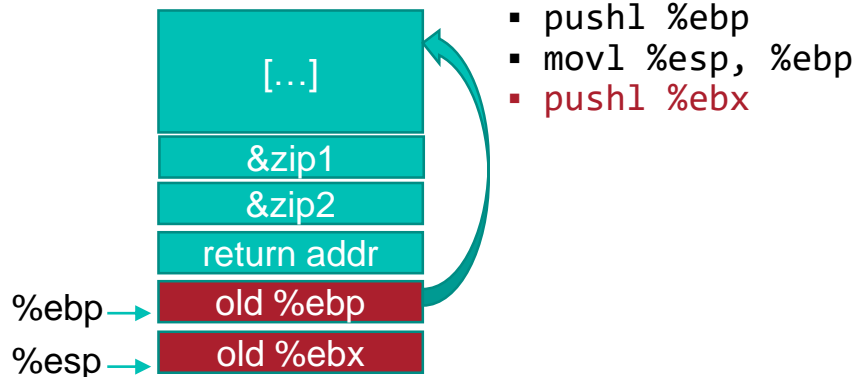
Example



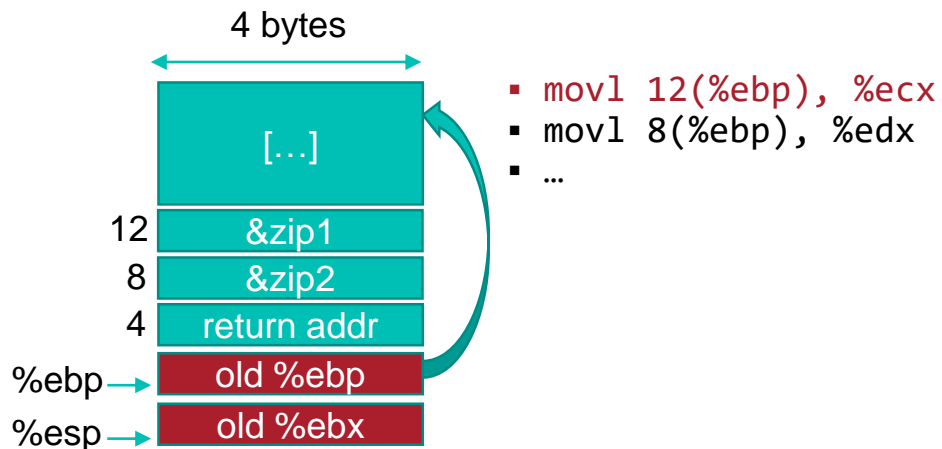
Example



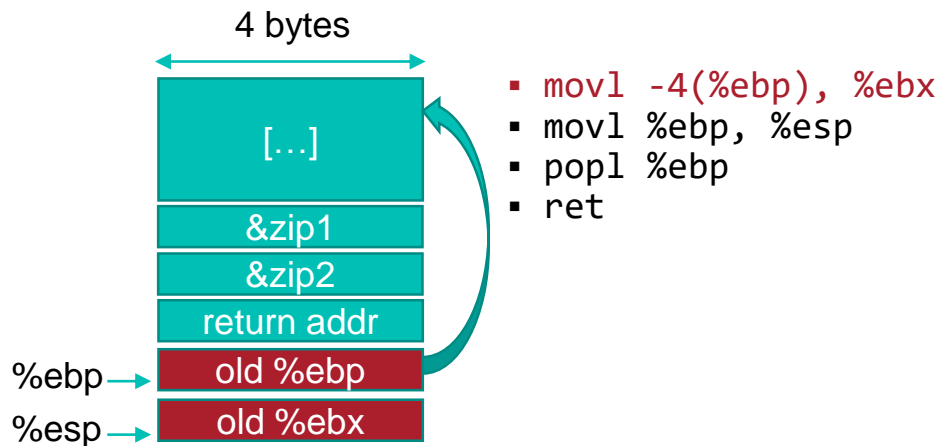
Example



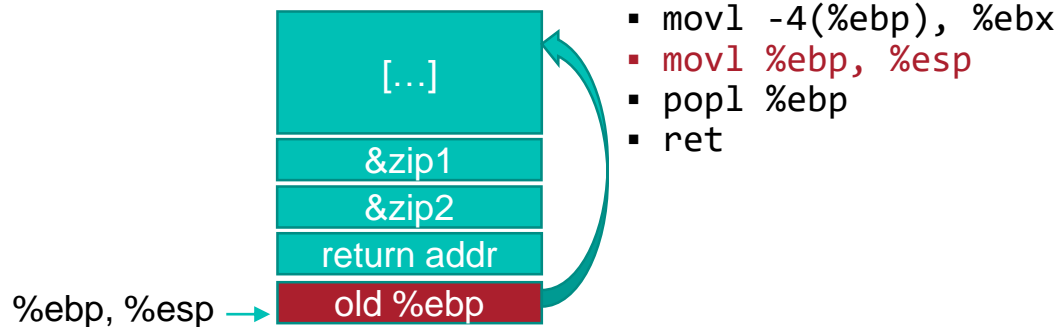
Example



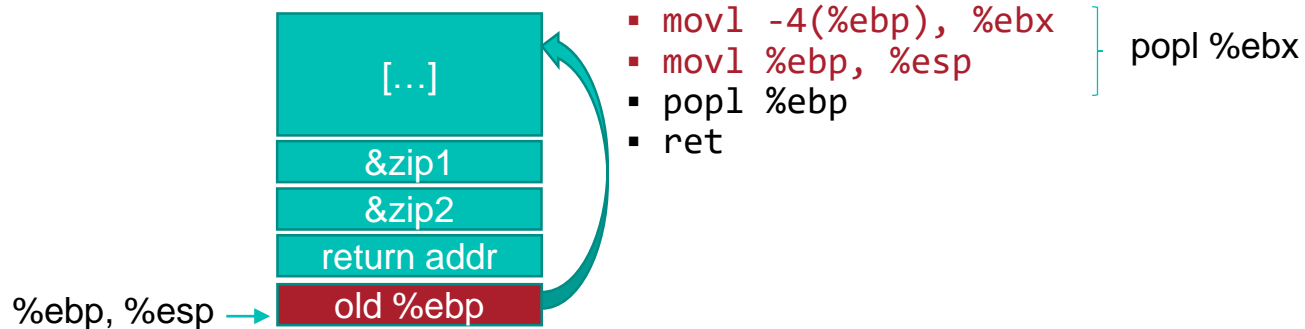
Example



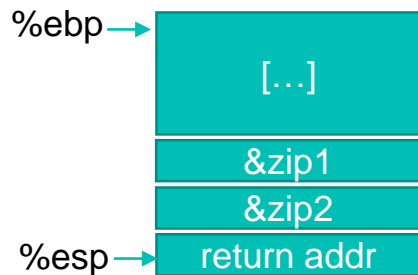
Example



Example

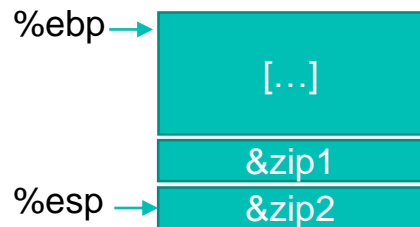


Example



- `movl -4(%ebp), %ebx`
- `movl %ebp, %esp`
- `popl %ebp`
- `ret`

Example



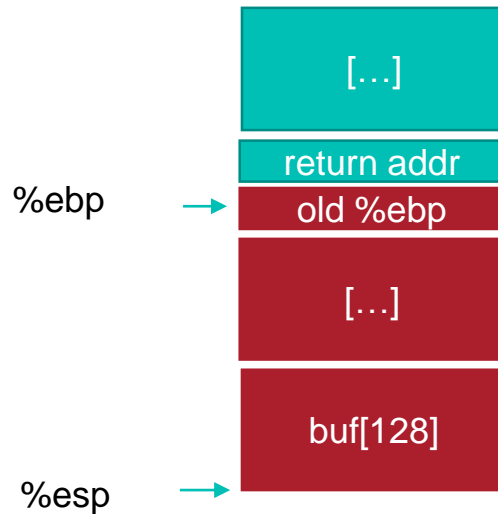
- `movl -4(%ebp), %ebx`
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- `ret`

Buffer overflow



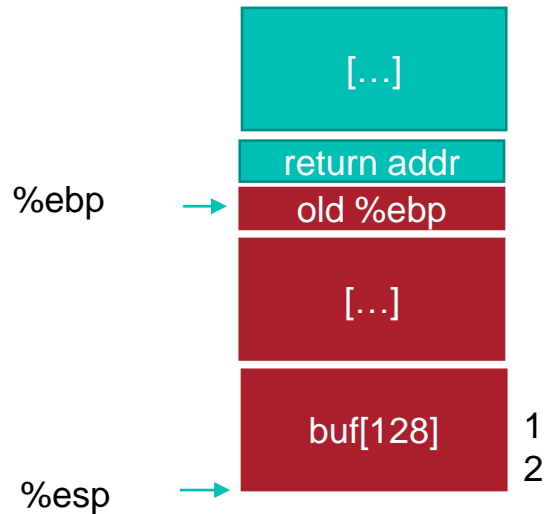
Example

```
▪ int read_get(void) {  
▪   char buf[128];  
▪   int i;  
▪   gets(buf);  
▪   i = atoi(buf);  
▪   return i;  
▪ }  
  
▪ int main() {  
▪   x = read_get();  
▪   printf("%s", x);  
▪ }
```



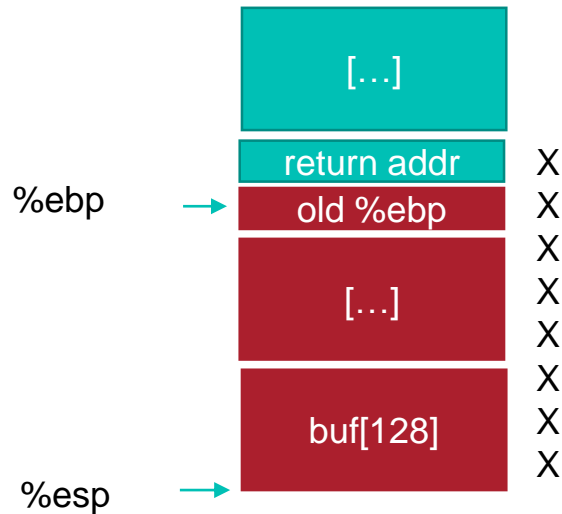
Example

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▪   int i;  
▪   gets(buf);  
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```



Example

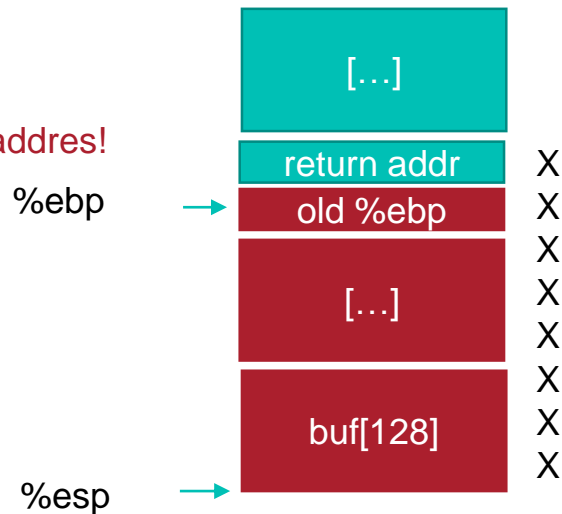
```
▪ int read_get(void) {  
▪   char buf[128];  
▪   int i;  
▪   gets(buf);  
▪   i = atoi(buf);  
▪   return I;  
▪ }  
  
▪ int main() {  
▪   x = read_get();  
▪   printf("%s", x);  
▪ }
```



Example

```
▪ int read_get(void) {  
▪   char buf[128];  
▪   int i;  
▪   gets(buf);  
▪   i = atoi(buf);  
▪   return I;  
▪ }  
  
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```

Changed returned address!
and old ebp.



Example

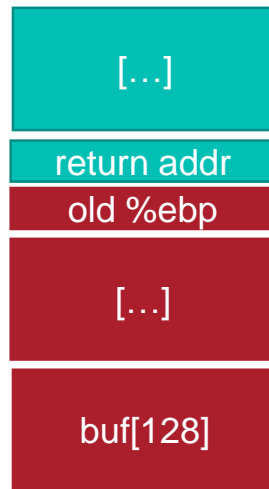
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▪   return i;  
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▪ }
```

Changed return address!
and old ebp.

%ebp



%esp

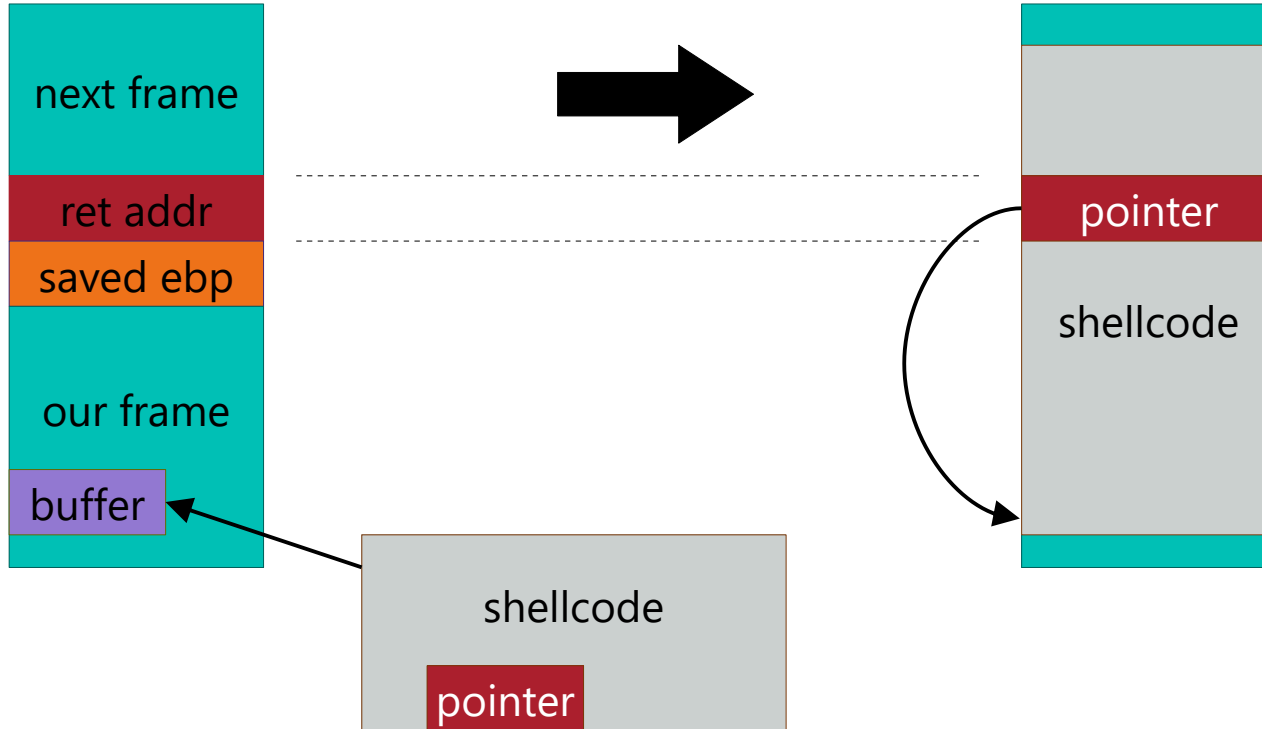


&evil
something
X
E
V
I
L
X

buffer overflows

- Vulnerability in C / assembly programs where the compiler does not enforce array bounds.
 - `a[1]`
 - `*(a+1)`
- Take over a setuid program, get root.

stack overflow



execve

- **NAME**

- `execve` - execute program

- **SYNOPSIS**

- `#include <unistd.h>`

- `int execve(const char *filename,`
- `char *const argv[],`
- `char *const envp[]);`

- **DESCRIPTION**

- `execve()` executes the program pointed to by filename.

execve in assembly

```
▪ .section .data
▪     cmd: .asciz "/bin/sh"
▪     ptr: .int cmd
▪         .int 0

▪ .section .text
▪     .globl _start
▪ _start:
▪     mov $0x0b, %eax    # execve
▪     mov $cmd, %ebx     # command
▪     mov $ptr, %ecx     # args
▪     mov $0, %edx      # env
▪     int $0x80
```

What's the problem?



The problem

- `mov $0x0b, %eax = B8 0B 00 00 00`
 - B8: `mov IMM32, %eax`
 - those null bytes will terminate a `strcpy/scanf/gets` etc.
- challenge is to create shellcode with only “legal” bytes
- also, how to address your payload?
- For **you** to figure out in the coursework ;-)

How to protect from this?



countermeasures



prevent

detect

recover

Prevent!

- Solution A: Avoid bugs in your C code!

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 - Maybe can check usage of problematic C functions?

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 - What about raw pointer manipulations?

Prevent!

- Solution A: Avoid bugs in your C code!
 - Maybe can check usage of problematic C functions?
 - What about raw pointer manipulations?
 - Look at a real large C projects... does not look easy

Prevent!

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- Solution B: build tools

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 - Static analysis

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 - Static analysis
- ```
▪ void foo(int *p) {
 ▪ int off;
 ▪ *z = p + off;
 ▪ if (off > 8)
 ▪ bar(8);
 ▪ }
```

# Prevent!

- Solution A: Avoid bugs in your C code!
  - Solution B: build tools
    - To help find bugs
    - Static analysis
- ```
▪ void foo(int *p) {  
  ▪ int off; NOT INITIALIZED  
  ▪ *z = p + off;  
  ▪ if (off > 8)  
    ▪ bar(8);  
  ▪ }
```


Prevent!

- Solution A: Avoid bugs in your C code!

- Solution B: build tools

- To help find bugs
- Static analysis

```
▪ void foo(int *p) {  
  ▪ int off; NOT INITIALIZED  
  ▪ *z = p + off;  
  ▪ if (off > 8) PROPAGATE ASSUMPTION ABOUT  
    bar(off); off VALUE  
  ▪ }
```

Prevent!

- Solution A: Avoid bugs in your C code!
- Solution B: build tools
 - To help find bugs
 - Static analysis
 - Fuzzing
 - Pushing massive amount of random value to a program
 - See if it crashes

Prevent!

- Solution A: Avoid bugs in your C code!
- Solution B: build tools
 - To help find bugs
 - Static analysis
 - Fuzzing
 - Pushing massive amount of random value to a program
 - See if it crashes
 - Can be a bit smarter and make sure we reach every branch in the program

Prevent!

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- Solution C: use a memory safe language

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 - It used to be a problem, not necessarily anymore

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- Solution A: avoid bugs in your C code!
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- Solution C: use a memory safe language
 - JAVA, C#, Rust etc...
 - Legacy code! (that's how the real world exists)
 - Need low level hardware access?
 - Performance?
 - It used to be a problem, not necessarily anymore
 - Is your program CPU bound anyway?

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countermeasures



prevent

detect

recover

Next lecture

Detect!

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Thank you

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