

## Systems Security COMSM1500

### RE/PLAY issue

- No sound
- I checked with IT
  - Microphone not properly connected to the computer
  - Should be fixed
- Next lectures should be properly recorded
- Please do check for issue
- Thanks to Rajat who spotted the problem



### Safety



### Safety terminology

- Accident: unplanned and undesired event that leads to a loss
- Hazard: condition that can lead to an accident in a particular environment
- Incident: sometimes, a hazard that does not lead to an accident

• In security only "security incident", it does not "happen by accident"

### Safety and Security

Safety: freedom from accident (caused without intent)

Security: freedom from incident (caused maliciously)

# Most effective

### Safety measure

eliminate hazard

2. reduce likelihood of occurrence

3. control

4. mitigate (reduce damage in case of accident)

### Active/Passive

- Passive safety measure: presence alone contributes to safety
  - -e.g. fences
- Fail-safe: fail in such a way that the system end-up in a safer state
  - -e.g. fuse
- Active safety measure: monitor the system and intervene
  - e.g. fire tower (forest fire)

#### Risk

- risk = loss \* probability of occurrence
  - Favored: reduce probability of accident (to zero if possible)
  - Reduce loss
- In most case it should not be acceptable to not reduce occurrence of hazard if possible
- Safety industry ALARP (as low as reasonably practical)
  - Must demonstrate further reduction is either impractical...
  - ... or unreasonably costly



### Security

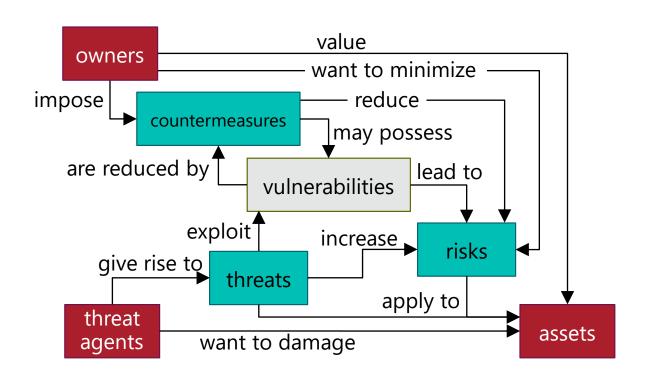


- Protection afforded to an automated information system in order to attain the applicable objectives of preserving the integrity, availability and confidentiality of information system resources (includes hardware, software, firmware, information/data and telecommunications).
  - US National Institute of Standards and Technology 1995

### terminology

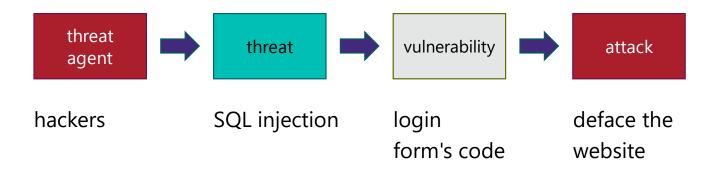
safety	security
hazard	vulnerability
(no equivalent)	threat
safety measure	countermeasure
risk	risk
accident / incident	incident

### **Terminology (IETF RFC 2828)**



#### attack

- realisation of a threat
- sequence of actions that results in a loss



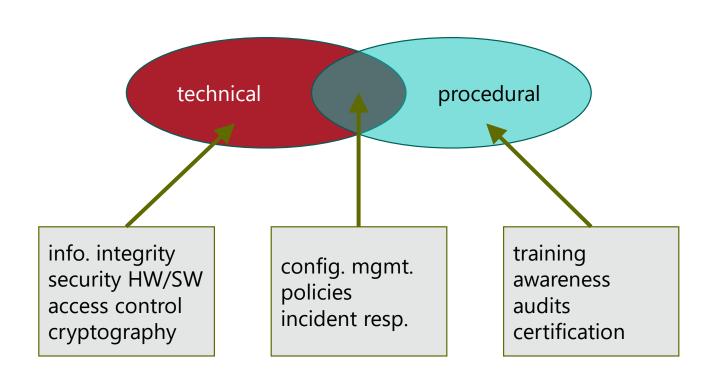
#### countermeasures

prevent

detect

recover

#### countermeasures



### Risk (security economics side notes)

- risk = loss \* probability of occurrence
- Think about CC fraud
  - Cost of preventive measure (i.e. reduce probably)
    - Paying for damage
  - Bank prefer to pay for damage (simple cost calculation)
    - > Simple risk calculation
    - > Risk < cost of preventive measure
- In other industry:
  - Data leaks, maybe some damage to your image
  - GDPR fines for personal data leak (up to 4% of annual global turnover)
  - Loss sky rocket
  - Incentive to bring down probability of occurrence

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- Threat Model -> assumptions about the adversary, the system, the environment
- Mechanism -> Software/Hardware

### Why is it hard?

- A faculty can access coursework
  - Easy just need to test!
- Negative goal
  - This is much harder
- Only a faculty can access coursework
  - It is hard need to think about what the adversary may do
  - Guess password?
  - Steal a laptop?
- It is an iterative process (we saw it during the last lecture)

### What can go wrong?

- Policy
  - Last lecture
    - > e-mail access
    - > Amazon example

### What can go wrong?

- Threat models
  - Human factor! (making assumption about what people would do)
  - Wrong assumptions about attacker knowledge (last lecture Apple)
  - Change over time (DES)
  - DARPA secure OS / similar issue more recently in China

HACK BRIEF: MALWARE SNEAKS INTO THE CHINESE IOS APP STORE

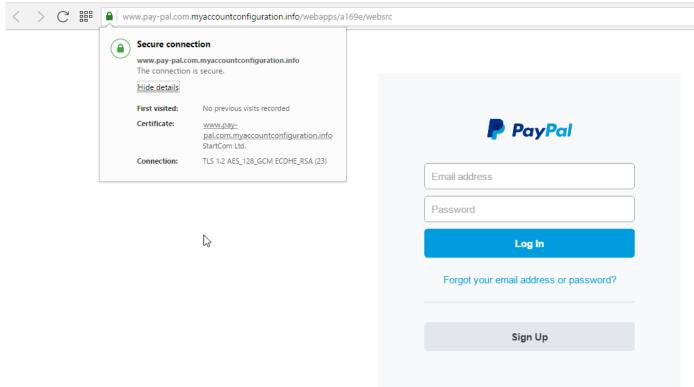


### Jennifer Lawrence Phishing

- Simple and easy
- Send e-mail that appears to come from a legitimate party
- Ask for information
- Or link to a fake website looking like the real thing (i.e. please reset your password)



### A legitimate fake site (untrustworthy CA)



### What can go wrong?

- Mechanisms
  - That is the source of most vulnerabilities

- Late 2017/ Early 2018
- Lots of chatter online
- Fixes to the Linux Kernel being pushed

 Family of exploit to the same underlying vulnerability



- Problem is much larger
- Every OS is affected
- Modern CPU architecture is the problem
- Intel, AMD, ARM are affected

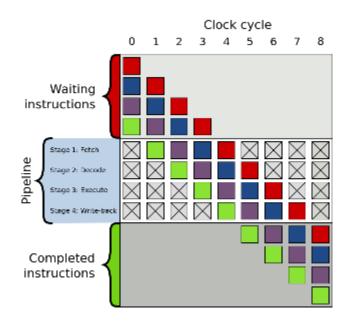
- Problem is much larger
- Every OS is affected
- Modern CPU architecture is the problem
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- Can be exploited by a simple piece of javascript



### High risk!



### Instructions pipeline





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- Evaluate conditions
- Need to wait to get the instructions to put in the pipeline
- This create a "bubble"

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    - > CPU state is reverted (that is good)
    - ... but CPU cache is not (that is bad)

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  - If guessed wrong we may have (half) executed instructions in the pipeline

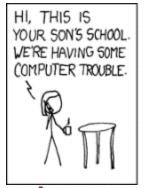
Exploit branch prediction

- Exploit branch prediction
- You can get access to data in the cache
- Access check is done when instructions are "executed"
- You can read kernel memory, leak password, keys etc...
  - -if (x < array1.size)
  - y = array2[array1[x]]



### SQL injection

- Attacker can pass any command to the database
  - SELECT, INSERT, UPDATE, DELETE









### Relatively easy fix

#### **COURSEWORK GROUPS!**

- Groups of 4 students
- (only) 1 e-mail per group
  - thomas.pasquier@bristol.ac.uk
    - > Header: [Systems Security] Group
    - ➤ Body: id1, id2, id3, id4
- If you cannot form a group you will get one, no need to e-mail me (>80)
- I will circulate the list
- FRIDAY 7PM
- IF NOT IN A GROUP RANDOM ALLOCATION



### Thank you

Office MVB 3.26

