

Systems Security COMSM1500



Anonymous communication



The problem?

- An attacker could observe network traffic
- Even without access to message data can learn a lot
- Service accessed, usage pattern, etc...

- Internet Service Provider
- Know the domain and IP address you want to visit
- Port (i.e. can infer service), timestamps etc...
- Packet size can leak information about what you do

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- Should I care?

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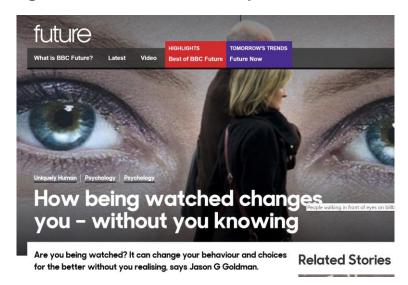


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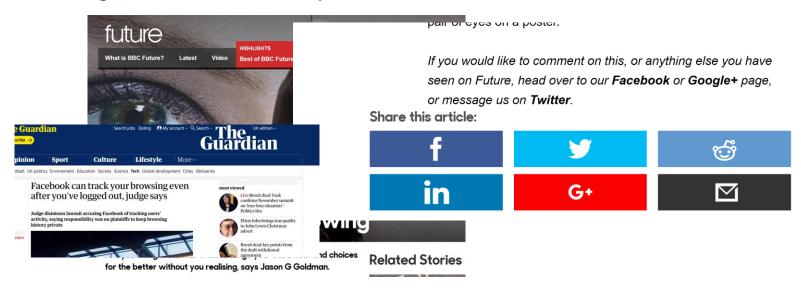
Observation is problematic

Being observed affect your behavior



Observation is problematic Some irony...

Being observed affect your behavior



Many reason for anonymity

- Means to communicate anonymously in some circumstances
 - Law enforcement to not tip their targets
 - Minority groups
 - Journalists
 - Political militants
 - Lawyers
- There is a few technology to achieve anonymity
- Some usages are less acceptable (more on that later...)

Many reason for anonymity

Homework/exam question:
Discuss why anonymity is important
even in a democratic society

- Means to communicate anonymously in some circumstances
 - Law enforcement to not tip their targets
 - Minority groups
 - Journalists
 - Political militants
 - Lawyers
- There is a few technology to achieve anonymity
- Some usages are less acceptable (more on that later...)

Plan

- Anonymity
- Unlinkability
- Unobservability
- VPN
- TOR
- TOR Circuit
- TOR Directory Authority
- TOR vulnerabilities

Anonymity

Preventing an observer on a network to link a participant to an action

Anonymity

Preventing an observer on a network to link a participant to an action

- We saw "private browsing" in previous lecture
 - Goal: do not leave trace on your local machine
 - This is not the same
 - You may want both

Anonymity

Preventing an observer on a network to link a participant to an action



- Observer can now Alice is doing something
- Observer can now someone is buying a t-shirt
- Observer cannot say Alice in particular is buying a t-shirt
 - Absolutely or probabilistically

Other important concepts

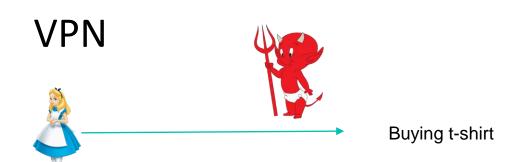
- Unlinkability
 - Cannot link Alice to some online identity/profile
- Unobservability
 - Cannot tell Alice is on Internet
 - More realistic cannot tell Alice is using some anonymity tool
- Confidentiality != Anonymity

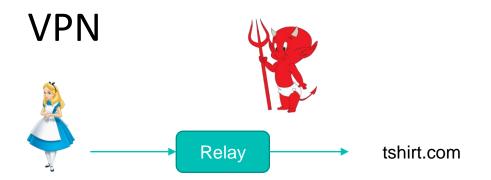


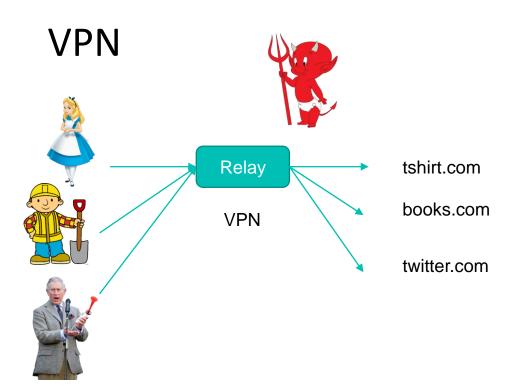
TOR

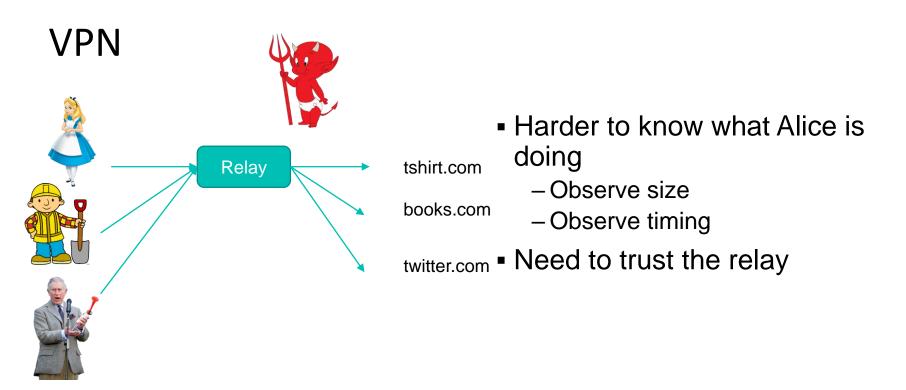
The Onion Router

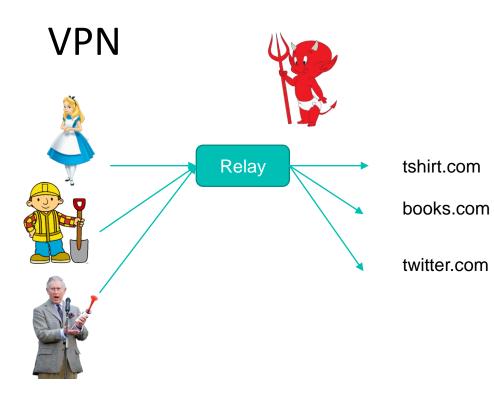












Harder to know what Alice is doing

- Observe size

- Observe timing

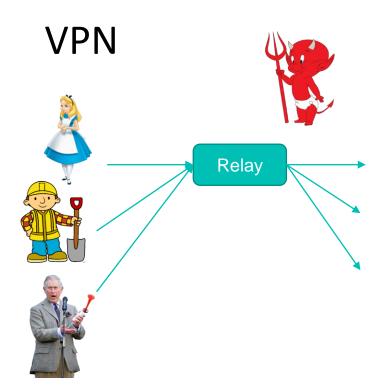
"Mixminion" fix-size request + answer

Batch a number of request together

- Send all at once

- Problem?

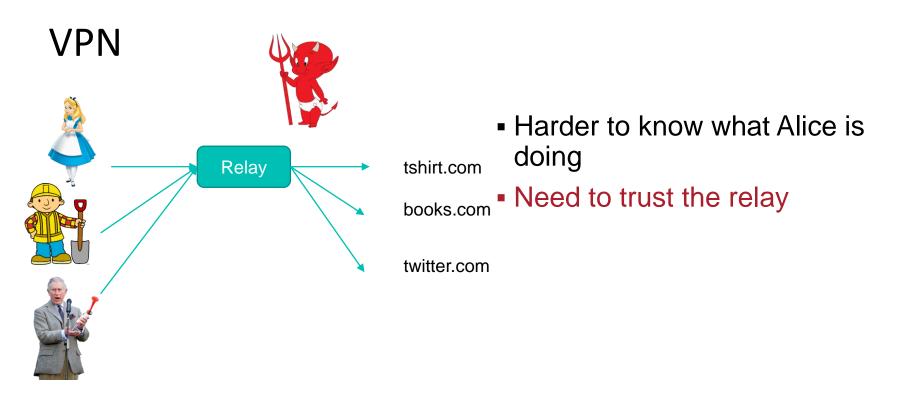
Need to trust the relay

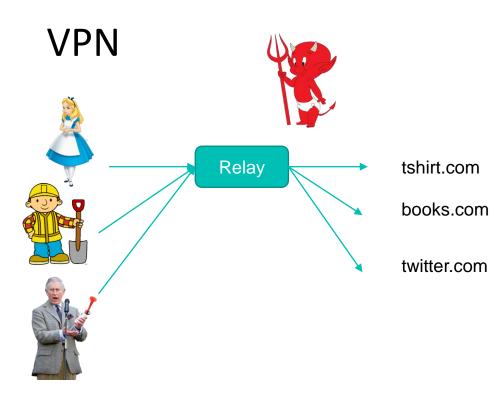


- Harder to know what Alice is doing
 - Observe size
 - Observe timing
- books.com "Mixminion" fix-size request + answer
- twitter.com

tshirt.com

- Batch a number of request together
- Send all at once
- Problem?
 - > Not going to be great to surf online
- Need to trust the relay



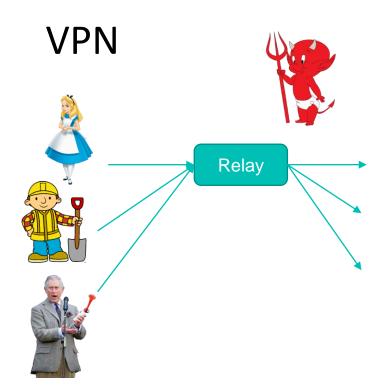


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Single relay is obviously a problem

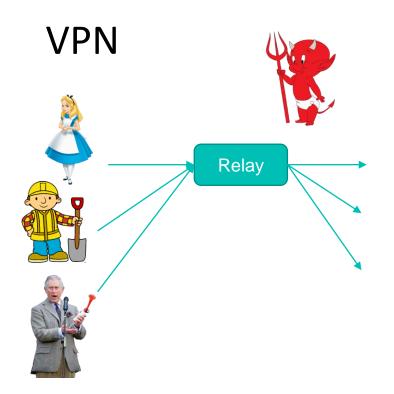
If it is compromised no guarantees



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- twitter.com Trusted VPN are fine

books.com

 e.g. universities run one, if you need to access some info from country that bans some content

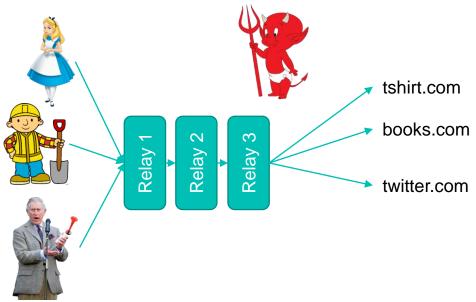


Homework/potential exam question: Discuss: Why VPN do not provide good anonymity.

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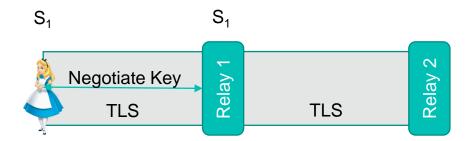
- Harder to know what Alice is doing
- Need to trust the relay
 - Relay 1 now Alice is doing something
 - Relay 3 now some is talking to t-shirt.com
 - Attacker need to control 1 and 3 to be really harmful
 - Hard/Costly to achieve
 - Discussed further later...



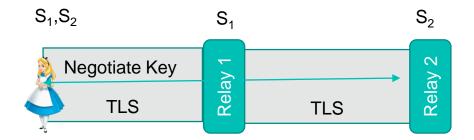




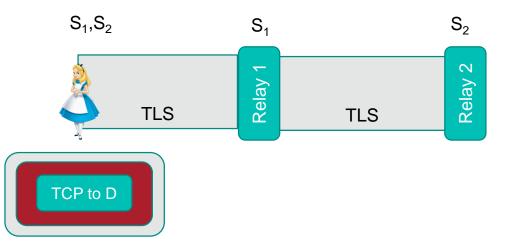




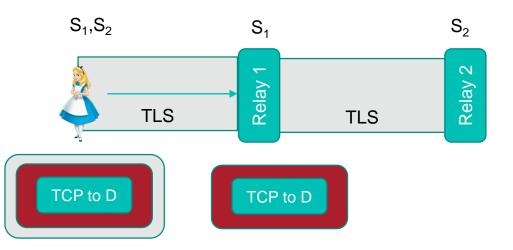




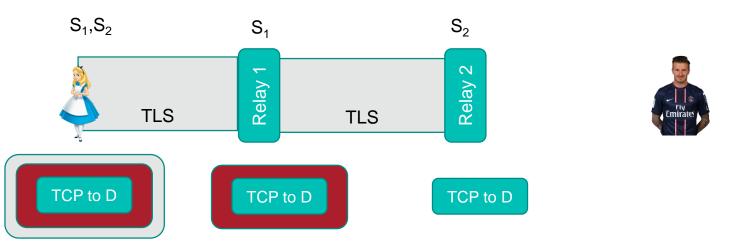


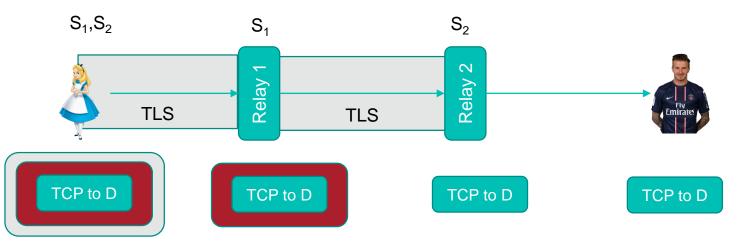


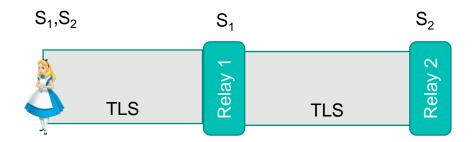






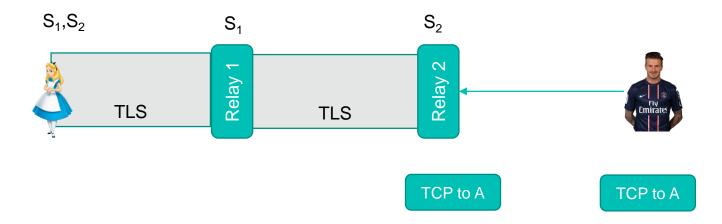


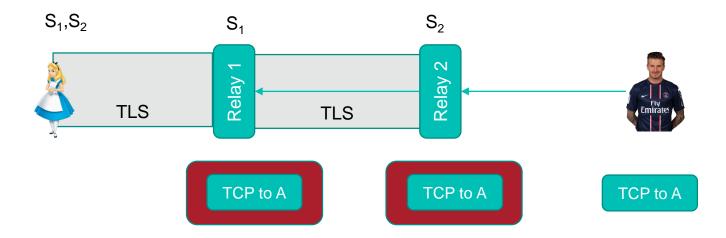


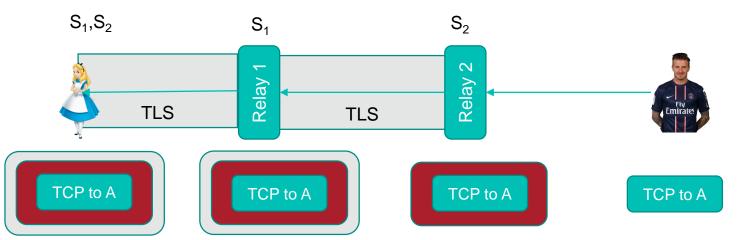




TCP to A





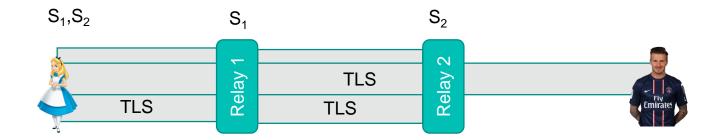




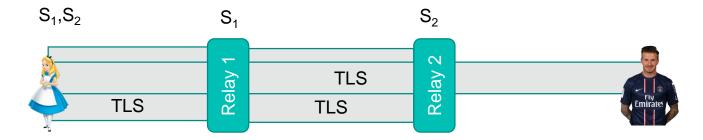
Careful

Messages between end relays and destination is unencrypted!





- Carry TCP packets
- Alice can establish an encrypted connection with David over TOR relays
 - e.g. HTTPS/TLS
- None of the relay can see content exchanged between Alice and David
- Relay 1 knows Alice send Data
- Relay 2 knows someone talk to David
- Things we need to be careful about?



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- Things we need to be careful about?
 - DNS
 - Certificate verification
 - Need to make sure it goes through TOR

Remember end-service can track you!

- End-servers can track you!
 - Cookies
 - Browser/Machine ID etc...
 - Browser used is important!
- That also include advertisements etc...
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Directory Authorities

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- A few of them
- Used to download a list of known relays
- Consensus protocol to decide trusted relays
- A majority of authorities needs to be trustworthy
 - Classic consensus problem

- It is very hard to deanonymize everyone all the time
- however, definitely possible to deanonymize some person sometimes

Passive attacks

- Size, timing (the more you can observe the easier)
 - > Possible if observe in relay and out relay
 - > Either own a lot of relay so you have high change to be picked
 - > ... or be able to observe the network
- Service fingerprint
 - Build pattern of size/timing of a service response (e.g. Facebook)
 - > Observe entry node and try to match
 - > You can learn which users is accessing service you care about

Homework/cool project: Look at fingerprinting as a min to deanonymization.

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Homework/potential exam question: Discuss: why it is a bad idea to have entry and exit nodes in the same country or owned by the same entity?

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- Active attacks
 - Steal key for TLS encryption between relay
 - ➤ High cost attack
 - > Rotate keys regularly
 - Iterated compromise
 - > i.e. identifying relays one after the other and compromising/coercing them
 - > Change circuit regularly
 - Cross border (make coercion harder)
 - Run Relay
 - > If attackers control a large number of relays it is likely he could have both ends
 - Need to own a significant portions of relays
 - Cost barrier?

- Active attacks
 - Smear attacks
 - > Purpose is to force end-nodes to shutdown (e.g. to increase portion of end-nodes controlled by an attacker)
 - Make request to legally questionable service
 - > End-nodes need to either have policy to filter this...
 - > ... or be able to take the heat
 - > Running other type of relay is ok
 - DOS on directory authority
 - Could stop the network
 - Run/Compromise directory authority
 - List attacker-controlled relays
 - Consensus is used to decide which relays are used
 - Would need large number of directory servers controlled by the attacker
 - > ... but see above?

Homework/potential exam question: Discuss: what is the danger of running a TOR exit relay.

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- Active attacks
 - Block Relay
 - > Everyone can access directory authorities
 - > Filter relays IP in traffic
 - > China does this
 - Countermeasure: TOR bridge (not advertised)
 - Block bridge
 - > Look at SSL traffic
 - Connection to TOR bridge had some recognizable artefact
 - > Try to connect to it and see if it is a TOR bridge
 - > China did it again
 - > Countermeasure: some shared secret between TOR client and Bridge

Homework/potential exam question: Discuss: arm race to prevent access to TOR network.

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- TOR Directory Authority
- TOR vulnerabilities

Conclusion

- Internet anonymity is hard
- Possible to hide from network observation
- Can identify some people sometimes
 - Everyone, all the time is much harder
- Active area of research
 - Check the papers on the github repo
- There is obviously a dark side to TOR-like software
 - Check work by Brian Neil Levine at UMass



Thank you, questions?

Office MVB 3.26

