## REFERENCE: THEOREMS AND LEMMATA

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## Figure 1: Key Theorems

## Type Safety:

- 1. (Preservation) If  $\vdash e : \tau$  and  $e \longmapsto e'$  then  $\vdash e' : \tau$ .
- 2. (Progress) If  $\vdash e : \tau$  then either e val or  $e \longmapsto e'$  for some e'.

<u>Termination</u>: For every  $\vdash e : \tau$  there exists a v val such that  $e \longmapsto^* v$ .

## Figure 2: Key Lemmata

Inversion: Suppose  $\Gamma \vdash e : \tau$ .

- 1. If  $e = plus(e_1; e_2)$  then it must be that
  - $\tau = \text{Num}$
  - $\Gamma \vdash e_1 : \mathsf{Num}$
  - $\Gamma \vdash e_2 : \mathsf{Num}$
- 2. ..

Weakening: If  $\Gamma \vdash e : \tau$  and x is fresh then  $\Gamma, x : \sigma \vdash e : \tau$ .

 $\underline{\overline{\text{Substitution}}}\text{: }\text{If }\Gamma\vdash e:\tau\text{ and }\Gamma,x:\tau\vdash u:\sigma\text{, then }\Gamma\vdash u[e/x]:\sigma\text{.}$ 

Canonical forms: Suppose e val.

- 1. If  $\vdash e$ : Num then  $e = \mathsf{num}[n]$  for some  $n \in \mathbb{N}$ .
- 2. If  $\vdash e$ : Str then e = str[s] for some  $s \in \Sigma^*$ .