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CS 576 - Lab 11

To successfully exploit the vulnerable format string for completion of this assignment, we need to find the return address of the test_fmts function and the stack address of the stack canary once the format string is printed.

First, to find the return address of test_fmts, we use gdb to set a breakpoint at the beginning of the function and use the command `info frame`. The output gives us the saved eip which is the return address: 0x8049233.

Next, to find the stack canary we need to set a breakpoint at the second printf in the test_fmts function. Once we reach the second printf, we can print out the stack to save what it will look like during execution of the command. Next, we disassemble the function and look for the instruction in which the stack canary is loaded into eax. Once we break there (at the instruction containing gs:), we can use the command `p $eax` which gives us the current stack canary. If we scroll back to the stack that we printed before, we can search for the stack canary and then we can write it's address which is 0xffffcfcc. A little further down we can find the return address as well.

Earlier when we disassembled the test_fmts function, we found that the first argument on the stack was 0x40 so if we check the stack once more, we can find the location of a single 0x40. Then, we can count the number of doublewords from the 0x40 to the stack canary stored at 0xffffcfcc. After this, we do the same for the return address which yields us 23 and 27 respectively (as these are the argument #'s for the canary and return address).

Now, we can input the string `canary: %23$x ret addr: %27$x` which gives the following result (note: this was run twice to verify stack canary value changed between processes):