# MSP430x2xx Family

# **User's Guide**



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# **Contents**

Prefa	асе			21
1	Intro	duction	١	23
	1.1		ecture	
	1.2	Flexibl	e Clock System	24
	1.3		dded Emulation	
	1.4		ss Space	
		1.4.1	Flash/ROM	
		1.4.2	RAM	
		1.4.3	Peripheral Modules	
		1.4.4	Special Function Registers (SFRs)	
		1.4.5	Memory Organization	
	1.5	MSP4	30x2xx Family Enhancements	
2	Syste		sets, Interrupts, and Operating Modes	
_	2.1		n Reset and Initialization	
	2.1	2.1.1	Brownout Reset (BOR)	
		2.1.2	Device Initial Conditions After System Reset	
	2.2		pts	
	2.2	2.2.1	(Non)-Maskable Interrupts (NMI)	
		2.2.2	Maskable Interrupts	
		2.2.3	Interrupt Processing	
		2.2.4	Interrupt Vectors	
	2.3		ting Modes	
		2.3.1	Entering and Exiting Low-Power Modes	
	2.4	Princip	ples for Low-Power Applications	
	2.5	•	ection of Unused Pins	
3	CPU			
3	3.1		ntroduction	
	3.2		Registers	
	5.2	3.2.1	Program Counter (PC)	
		3.2.2	Stack Pointer (SP)	
		3.2.3	Status Register (SR)	
		3.2.4	Constant Generator Registers CG1 and CG2	
		3.2.5	General-Purpose Registers R4 to R15	
	3.3	-	ssing Modes	
	0.0	3.3.1	Register Mode	
			Indexed Mode	
		3.3.3	Symbolic Mode	
		3.3.4	Absolute Mode	
		3.3.5	Indirect Register Mode	
		3.3.6	Indirect Autoincrement Mode	
		3.3.7	Immediate Mode	
	3.4		tion Set	
		3.4.1	Double-Operand (Format I) Instructions	
		3.4.2	Single-Operand (Format II) Instructions	
		3.4.3	Jumps	

2



		3.4.4	Instruction Cycles and Lengths	. 60
		3.4.5	Instruction Set Description	. 62
		3.4.6	Instruction Set Details	. 64
4	CPUX	ζ		115
	4.1		ntroduction	
	4.2		pts	
	4.3		Registers	
		4.3.1	Program Counter (PC)	
		4.3.2	Stack Pointer (SP)	
		4.3.3	Status Register (SR)	
		4.3.4	Constant Generator Registers (CG1 and CG2)	
		4.3.5	General-Purpose Registers (R4 to R15)	
	4.4		ssing Modes	
		4.4.1	Register Mode	
		4.4.2	Indexed Mode	
		4.4.3	Symbolic Mode	
		4.4.4	Absolute Mode	
		4.4.5	Indirect Register Mode	
		4.4.6	Indirect Autoincrement Mode	
		4.4.7	Immediate Mode	
	4.5		30 and MSP430X Instructions	
		4.5.1	MSP430 Instructions	
		4.5.2	MSP430X Extended Instructions	
	4.6		tion Set Description	
		4.6.1	Extended Instruction Binary Descriptions	
		4.6.2	MSP430 Instructions	
		4.6.3	MSP430X Extended Instructions	
		4.6.4	MSP430X Address Instructions	
5	Racio	_	Module+	
5	5.1		Clock Module+ Introduction	
	5.1		Clock Module+ Operation	
	5.2	5.2.1	Basic Clock Module+ Features for Low-Power Applications	
		_		
		5.2.2	Internal Very-Low-Power Low-Frequency Oscillator (VLO)	
		5.2.3	XT2 Oscillator	
			Digitally-Controlled Oscillator (DCO)	
		5.2.5		
		5.2.6	DCO Modulator	
		5.2.7	Basic Clock Module+ Fail-Safe Operation	
	F 2	5.2.8	Clock Module+ Registers	
	5.3		· · · · · · · · · · · · · · · · · · ·	
		5.3.1	DCOCTL, DCO Control Register	
		5.3.2	BCSCTL1, Basic Clock System Control Register 1	
		5.3.3	BCSCTL2, Basic Clock System Control Register 2	
		5.3.4	BCSCTL3, Basic Clock System Control Register 3	
		5.3.5	IE1, Interrupt Enable Register 1	
		5.3.6	IFG1, Interrupt Flag Register 1	
6	DMA			
	6.1		ntroduction	
	6.2	DMA (	Operation	
		6.2.1	DMA Addressing Modes	290
		6.2.2	DMA Transfer Modes	291
		6.2.3	Initiating DMA Transfers	297





		6.2.4	Stopping DMA Transfers	298
		6.2.5	DMA Channel Priorities	299
		6.2.6	DMA Transfer Cycle Time	299
		6.2.7	Using DMA With System Interrupts	299
		6.2.8	DMA Controller Interrupts	
		6.2.9	Using the USCI_B I <sup>2</sup> C Module with the DMA Controller	
		6.2.10	3	
		6.2.11	Using DAC12 With the DMA Controller	301
		6.2.12	Writing to Flash With the DMA Controller	301
	6.3	DMA F	Registers	302
		6.3.1	DMACTL0, DMA Control Register 0	303
		6.3.2	DMACTL1, DMA Control Register 1	
		6.3.3	DMAxCTL, DMA Channel x Control Register	304
		6.3.4	DMAxSA, DMA Source Address Register	305
		6.3.5	DMAxDA, DMA Destination Address Register	306
		6.3.6	DMAxSZ, DMA Size Address Register	306
		6.3.7	DMAIV, DMA Interrupt Vector Register	307
7	Flash	Memo	ory Controller	308
	7.1		Memory Introduction	
	7.2		Memory Segmentation	
		7.2.1	SegmentA	
	7.3	Flash I	Memory Operation	
		7.3.1	Flash Memory Timing Generator	
		7.3.2	Erasing Flash Memory	
		7.3.3	Writing Flash Memory	
		7.3.4	Flash Memory Access During Write or Erase	
		7.3.5	Stopping a Write or Erase Cycle	
		7.3.6	Marginal Read Mode	
		7.3.7	Configuring and Accessing the Flash Memory Controller	321
		7.3.8	Flash Memory Controller Interrupts	
		7.3.9	Programming Flash Memory Devices	
	7.4	Flash I	Memory Registers	
		7.4.1	FCTL1, Flash Memory Control Register	324
		7.4.2	FCTL2, Flash Memory Control Register	
		7.4.3	FCTL3, Flash Memory Control Register	
		7.4.4	FCTL4, Flash Memory Control Register	326
		7.4.5	IE1, Interrupt Enable Register 1	326
8	Digita	al I/O		327
•	8.1		I/O Introduction	
	8.2	•	I/O Operation	
	0.2	8.2.1	Input Register PxIN	
		8.2.2	Output Registers PxOUT	
		8.2.3	Direction Registers PxDIR	
		8.2.4	Pullup/Pulldown Resistor Enable Registers PxREN	
		8.2.5	Function Select Registers PxSEL and PxSEL2	
		8.2.6	Pin Oscillator	
		8.2.7	P1 and P2 Interrupts	
		8.2.8	Configuring Unused Port Pins	
	8.3		I/O Registers	
^		•	•	
9		-	age Supervisor (SVS)	
	9.1		Voltage Supervisor (SVS) Introduction	
	9.2		peration	
		9.2.1	Configuring the SVS	337

		9.2.2	SVS Comparator Operation	337
		9.2.3	Changing the VLDx Bits	337
			SVS Operating Range	
	9.3		egisters	
			SVSCTL, SVS Control Register	
10	Watc	hdog T	imer+ (WDT+)	341
	10.1	Watcho	dog Timer+ (WDT+) Introduction	342
	10.2	Watcho	dog Timer+ Operation	344
		10.2.1	Watchdog Timer+ Counter	
		10.2.2	Watchdog Mode	
		10.2.3	Interval Timer Mode	344
		10.2.4	Watchdog Timer+ Interrupts	
		10.2.5	Watchdog Timer+ Clock Fail-Safe Operation	345
		10.2.6	Operation in Low-Power Modes	345
		10.2.7	Software Examples	345
	10.3	Watcho	dog Timer+ Registers	346
		10.3.1	WDTCTL, Watchdog Timer+ Register	347
			,	
		10.3.3	IFG1, Interrupt Flag Register 1	348
11	Hard	ware M	ultiplier	349
	11.1	Hardwa	are Multiplier Introduction	350
	11.2	Hardwa	are Multiplier Operation	350
		11.2.1	Operand Registers	
		11.2.2	Result Registers	
		11.2.3	Software Examples	352
		11.2.4	Indirect Addressing of RESLO	353
		11.2.5	Using Interrupts	353
	11.3	Hardwa	are Multiplier Registers	354
12	Time	r <b>A</b>		355
	12.1	_	A Introduction	
	12.2		A Operation	
			16-Bit Timer Counter	
			Starting the Timer	
		12.2.3	Timer Mode Control	
		12.2.4	Capture/Compare Blocks	
			Output Unit	
		12.2.6	Timer_A Interrupts	
	12.3	Timer	A Registers	
		12.3.1	TACTL, Timer_A Control Register	
		12.3.2	TAR, Timer_A Register	
		12.3.3	TACCRx, Timer_A Capture/Compare Register x	
		12.3.4	TACCTLx, Capture/Compare Control Register	
		12.3.5	TAIV, Timer_A Interrupt Vector Register	373
13	Time			
13	13.1		B Introduction	
	13.1		Similarities and Differences From Timer_A	
	13.2		B Operation	
	10.2	13.2.1	16-Bit Timer Counter	
		13.2.1	Starting the Timer	
		13.2.2	Timer Mode Control	
		13.2.4	Capture/Compare Blocks	
		-	Output Unit	
		10.2.0	Output Offit	JU4





		13.2.6 T	Fimer_B Interrupts	388
	13.3		Registers	
			Fimer_B Control Register TBCTL	
			TBR, Timer_B Register	
			FBCCRx, Timer_B Capture/Compare Register x	
			FBCCTLx, Capture/Compare Control Register	
			FBIV, Timer_B Interrupt Vector Register	
4.4	l lusis ca		•	
14			al Interface (USI)	
	14.1		duction	
	14.2	•	ration	
			JSI Initialization	
			JSI Clock Generation	
			SPI Mode	
			<sup>2</sup> C Mode	
	14.3		sters	
			JSICTL0, USI Control Register 0	
			JSICTL1, USI Control Register 1	
			JSICKCTL, USI Clock Control Register	
			JSICNT, USI Bit Counter Register	
		14.3.5 L	JSISRL, USI Low Byte Shift Register	409
		14.3.6 L	JSISRH, USI High Byte Shift Register	409
15	Unive	rsal Seri	al Communication Interface, UART Mode	410
	15.1		erview	
	15.2		oduction: UART Mode	
	15.3		eration: UART Mode	
		•	JSCI Initialization and Reset	
			Character Format	
			Asynchronous Communication Formats	
			Automatic Baud Rate Detection	
			rDA Encoding and Decoding	
			Automatic Error Detection	
			JSCI Receive Enable	
			JSCI Transmit Enable	
			JART Baud Rate Generation	
			Setting a Baud Rate	
			Transmit Bit Timing	
			Receive Bit Timing	
			Typical Baud Rates and Errors	
			Using the USCI Module in UART Mode with Low Power Modes	
			<del>-</del>	
	45.4		USCI Interrupts	
	15.4	-	gisters: UART Mode	
			JCAxCTL0, USCI_Ax Control Register 0	
			JCAxCTL1, USCI_Ax Control Register 1	
			JCAxBR0, USCI_Ax Baud Rate Control Register 0	
			JCAxBR1, USCI_Ax Baud Rate Control Register 1	
			JCAxMCTL, USCI_Ax Modulation Control Register	
			JCAxSTAT, USCI_Ax Status Register	
			JCAxRXBUF, USCI_Ax Receive Buffer Register	
			JCAxTXBUF, USCI_Ax Transmit Buffer Register	
			JCAxIRTCTL, USCI_Ax IrDA Transmit Control Register	
			UCAxIRRCTL, USCI_Ax IrDA Receive Control Register	
			UCAxABCTL, USCI_Ax Auto Baud Rate Control Register	
		15.4.12	IE2, Interrupt Enable Register 2	433



		15.4.13 IFG2, Interrupt Flag Register 2	133
		15.4.14 UC1IE, USCI_A1 Interrupt Enable Register	
		15.4.15 UC1IFG, USCI_A1 Interrupt Flag Register	
16	Unive	rsal Serial Communication Interface, SPI Mode4	35
	16.1	USCI Overview	
	16.2	USCI Introduction: SPI Mode	136
	16.3	USCI Operation: SPI Mode	138
		16.3.1 USCI Initialization and Reset	138
		16.3.2 Character Format	139
		16.3.3 Master Mode	139
		16.3.4 Slave Mode	140
		16.3.5 SPI Enable	141
		16.3.6 Serial Clock Control	141
		16.3.7 Using the SPI Mode With Low-Power Modes	
		16.3.8 SPI Interrupts	
	16.4	USCI Registers: SPI Mode	
		16.4.1 UCAxCTL0, USCI_Ax Control Register 0, UCBxCTL0, USCI_Bx Control Register 0 4	
		16.4.2 UCAxCTL1, USCI_Ax Control Register 1, UCBxCTL1, USCI_Bx Control Register 1 4	45
		16.4.3 UCAxBR0, USCI_Ax Bit Rate Control Register 0, UCBxBR0, USCI_Bx Bit Rate Control Register	
			146
		16.4.4 UCAxBR1, USCI_Ax Bit Rate Control Register 1, UCBxBR1, USCI_Bx Bit Rate Control Register 1	146
			146
		16.4.6 UCAxRXBUF, USCI_Ax Receive Buffer Register, UCBxRXBUF, USCI_Bx Receive Buffer	
		Register4	146
		16.4.7 UCAxTXBUF, USCI_Ax Transmit Buffer Register, UCBxTXBUF, USCI_Bx Transmit Buffer Register	147
		16.4.8 IE2, Interrupt Enable Register 2	
		16.4.9 IFG2, Interrupt Flag Register 2	
		16.4.10 UC1IE, USCI_A1/USCI_B1 Interrupt Enable Register	
		16.4.11 UC1IFG, USCI_A1/USCI_B1 Interrupt Flag Register	48
17	Unive	rsal Serial Communication Interface, I <sup>2</sup> C Mode4	
	17.1	USCI Overview4	
	17.2	USCI Introduction: I <sup>2</sup> C Mode	
	17.3	USCI Operation: I <sup>2</sup> C Mode	
		17.3.1 USCI Initialization and Reset	
		17.3.2   <sup>2</sup> C Serial Data	
		17.3.3 I <sup>2</sup> C Addressing Modes	
		17.3.4 I <sup>2</sup> C Module Operating Modes	
		17.3.5 I <sup>2</sup> C Clock Generation and Synchronization	
		17.3.6 Using the USCI Module in I <sup>2</sup> C Mode with Low-Power Modes	
		17.3.7 USCI Interrupts in I <sup>2</sup> C Mode	
	17.4	USCI Registers: I <sup>2</sup> C Mode	
		17.4.1 UCBxCTL0, USCI_Bx Control Register 0	
		17.4.2 UCBxCTL1, USCI_Bx Control Register 1	
		17.4.3 UCBxBR0, USCI_Bx Baud Rate Control Register 0	
		17.4.4 UCBxBR1, USCI_Bx Baud Rate Control Register 1	
		17.4.5 UCBxSTAT, USCI_Bx Status Register	
		17.4.6 UCBxRXBUF, USCI_Bx Receive Buffer Register	
		17.4.7 UCBxTXBUF, USCI_Bx Transmit Buffer Register	
		17.4.8 UCBxI2COA, USCIBx I <sup>2</sup> C Own Address Register	
		17.4.9 UCBxl2CSA, USCI_Bx I <sup>2</sup> C Slave Address Register	
		17.4.10 OCBXI2CIE, OSCI_BX I C Interrupt Enable Register	
		17 11 122, Interrupt Enable Register 2	r 1 Z



		17.4.12	IFG2, Interrupt Flag Register 2	472
		17.4.13	UC1IE, USCI_B1 Interrupt Enable Register	472
		17.4.14	UC1IFG, USCI_B1 Interrupt Flag Register	473
18	USAF	RT Perin	heral Interface, UART Mode	474
	18.1	_	Introduction: UART Mode	
	18.2		Operation: UART Mode	
	10.2		USART Initialization and Reset	
			Character Format	
			Asynchronous Communication Formats	
			USART Receive Enable	
			USART Transmit Enable	
			USART Baud Rate Generation	
			USART Interrupts	
	18.3		Registers: UART Mode	
	10.5		UxCTL, USART Control Register	
			UxTCTL, USART Transmit Control Register	
			UxRCTL, USART Receive Control Register	
			UxBR0, USART Baud Rate Control Register 0	
			UxBR1, USART Baud Rate Control Register 1	
			UxMCTL, USART Modulation Control Register	
			UxRXBUF, USART Receive Buffer Register	
			UxTXBUF, USART Transmit Buffer Register	
			IE2, Interrupt Enable Register 2	
		18.3.11	IFG1, Interrupt Flag Register 1	
			IFG2, Interrupt Flag Register 2	
19		RT Perip	heral Interface, SPI Mode	497
19	19.1	<b>RT Perip</b> USART	heral Interface, SPI Mode Introduction: SPI Mode	<b>497</b> 498
19		<b>RT Perip</b> USART USART	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode	<b>497</b> 498 499
19	19.1	RT Perip USART USART 19.2.1	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset	<b>497</b> 498 499 499
19	19.1	RT Perip USART USART 19.2.1 19.2.2	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode	497 498 499 499 500
19	19.1	RT Perip USART USART 19.2.1 19.2.2 19.2.3	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode	497 498 499 499 500 500
19	19.1	USART USART 19.2.1 19.2.2 19.2.3 19.2.4	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable	497 498 499 499 500 500
19	19.1	USART USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control	497 498 499 500 500 501 502
19	19.1	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts	497 498 499 500 500 501 502 504
19	19.1	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode	497 498 499 500 500 501 502 504 506
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register	497 498 499 500 500 501 502 504 506 507
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register UxTCTL, USART Transmit Control Register	497 498 499 500 500 501 502 504 506 507
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2 19.3.3	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register UxTCTL, USART Transmit Control Register UxRCTL, USART Receive Control Register	497 498 499 500 501 502 504 506 507 507 508
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2 19.3.3 19.3.4	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register UxTCTL, USART Transmit Control Register UxRCTL, USART Receive Control Register UxBR0, USART Baud Rate Control Register 0	497 498 499 500 501 502 504 506 507 507 508 508
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register UxTCTL, USART Transmit Control Register UxRCTL, USART Receive Control Register UxBR0, USART Baud Rate Control Register 0 UxBR1, USART Baud Rate Control Register 1	497 498 499 500 501 502 504 506 507 508 508 508
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6	Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register UxTCTL, USART Transmit Control Register UxRCTL, USART Receive Control Register UxBR0, USART Baud Rate Control Register 1 UxMCTL, USART Baud Rate Control Register 1 UxMCTL, USART Modulation Control Register	497 498 499 500 501 502 504 506 507 508 508 508 508
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UXCTL, USART Control Register UXTCTL, USART Transmit Control Register UXRCTL, USART Receive Control Register UXBR0, USART Baud Rate Control Register 0 UXBR1, USART Baud Rate Control Register UXMCTL, USART Modulation Control Register UXRCTL, USART Modulation Control Register	497 498 499 500 501 502 504 506 507 507 508 508 508 508
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8	Introduction: SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register UxTCTL, USART Transmit Control Register UxRCTL, USART Receive Control Register UxBR0, USART Baud Rate Control Register 1 UxBR1, USART Modulation Control Register UxRXBUF, USART Receive Buffer Register UxRXBUF, USART Receive Buffer Register	497 498 499 500 501 502 504 506 507 507 508 508 508 508 508
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8 19.3.9	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UxCTL, USART Control Register UxTCTL, USART Transmit Control Register UxRCTL, USART Baud Rate Control Register UxBR0, USART Baud Rate Control Register 1 UxMCTL, USART Modulation Control Register UxRXBUF, USART Receive Buffer Register UxRXBUF, USART Receive Buffer Register UxTXBUF, USART Transmit Buffer Register ME1, Module Enable Register 1	497 498 499 500 501 502 504 506 507 508 508 508 508 508 508 509 509
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8 19.3.9	Introduction: SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UXCTL, USART Control Register UXTCTL, USART Transmit Control Register UXRCTL, USART Baud Rate Control Register UXBRO, USART Baud Rate Control Register 1 UXMCTL, USART Modulation Control Register UXRCTL, USART Receive Buffer Register UXRCTL, USART Receive Buffer Register	497 498 499 500 501 502 504 506 507 508 508 508 508 508 509 509
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8 19.3.9 19.3.10 19.3.11	Introduction: SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UXCTL, USART Control Register UXTCTL, USART Transmit Control Register UXRCTL, USART Receive Control Register UXBRO, USART Baud Rate Control Register 0 UXBR1, USART Baud Rate Control Register 1 UXMCTL, USART Receive Buffer Register UXRXBUF, USART Receive Buffer Register UXRXBUF, USART Receive Buffer Register UXTXBUF, USART Transmit Buffer Register UXTXBUF, USART Transmit Buffer Register ME1, Module Enable Register 1 ME2, Module Enable Register 2 IE1, Interrupt Enable Register 1	497 498 499 500 501 502 504 506 507 507 508 508 508 508 509 509 509
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8 19.3.9 19.3.10 19.3.11 19.3.12	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UXCTL, USART Control Register UXTCTL, USART Transmit Control Register UXRCTL, USART Receive Control Register UXBRO, USART Baud Rate Control Register 0 UXBR1, USART Baud Rate Control Register 1 UXMCTL, USART Receive Buffer Register UXRXBUF, USART Receive Buffer Register UXTXBUF, USART Transmit Buffer Register ME1, Module Enable Register 1 ME2, Module Enable Register 2 IE1, Interrupt Enable Register 2	497 498 499 500 501 502 504 506 507 507 508 508 508 508 509 509 509 510
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8 19.3.9 19.3.10 19.3.11 19.3.12 19.3.13	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UXCTL, USART Control Register UXTCTL, USART Transmit Control Register UXRCTL, USART Baud Rate Control Register UXBRO, USART Baud Rate Control Register 1 UXMCTL, USART Modulation Control Register UXRXBUF, USART Meceive Buffer Register UXTXBUF, USART Transmit Buffer Register ME1, Module Enable Register 1 ME2, Module Enable Register 2 IE1, Interrupt Enable Register 2 IFG1, Interrupt Flag Register 1	497 498 499 500 501 502 504 506 507 507 508 508 508 508 509 509 509 510 510
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8 19.3.9 19.3.10 19.3.11 19.3.12 19.3.13	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UXCTL, USART Control Register UXTCTL, USART Transmit Control Register UXRCTL, USART Receive Control Register UXBRO, USART Baud Rate Control Register 0 UXBR1, USART Baud Rate Control Register 1 UXMCTL, USART Receive Buffer Register UXRXBUF, USART Receive Buffer Register UXTXBUF, USART Transmit Buffer Register ME1, Module Enable Register 1 ME2, Module Enable Register 2 IE1, Interrupt Enable Register 2	497 498 499 500 501 502 504 506 507 507 508 508 508 508 509 509 509 510 510
19	19.1 19.2	USART USART 19.2.1 19.2.2 19.2.3 19.2.4 19.2.5 19.2.6 USART 19.3.1 19.3.2 19.3.3 19.3.4 19.3.5 19.3.6 19.3.7 19.3.8 19.3.1 19.3.13 19.3.11 19.3.13 19.3.14	heral Interface, SPI Mode Introduction: SPI Mode Operation: SPI Mode USART Initialization and Reset Master Mode Slave Mode SPI Enable Serial Clock Control SPI Interrupts Registers: SPI Mode UXCTL, USART Control Register UXTCTL, USART Transmit Control Register UXRCTL, USART Baud Rate Control Register UXBRO, USART Baud Rate Control Register 1 UXMCTL, USART Modulation Control Register UXRXBUF, USART Meceive Buffer Register UXTXBUF, USART Transmit Buffer Register ME1, Module Enable Register 1 ME2, Module Enable Register 2 IE1, Interrupt Enable Register 2 IFG1, Interrupt Flag Register 1	497 498 499 500 501 502 504 506 507 508 508 508 508 509 509 509 510 510



	20.2	OA Ope	eration	513
		20.2.1	OA Amplifier	514
		20.2.2	OA Input	514
		20.2.3	OA Output and Feedback Routing	514
			OA Configurations	
	20.3	_	gisters	
			OAxCTL0, Opamp Control Register 0	
		20.3.2	OAxCTL1, Opamp Control Register 1	522
21	Com	oarator_	_A+	<b>523</b>
	21.1	Compa	rator_A+ Introduction	524
	21.2	Compa	rator_A+ Operation	525
		21.2.1	Comparator	525
		21.2.2	Input Analog Switches	525
		21.2.3	Input Short Switch	526
		21.2.4	Output Filter	526
		21.2.5	Voltage Reference Generator	527
		21.2.6	Comparator_A+, Port Disable Register CAPD	527
		21.2.7	Comparator_A+ Interrupts	528
		21.2.8	Comparator_A+ Used to Measure Resistive Elements	528
	21.3	Compa	rator_A+ Registers	530
		21.3.1	CACTL1, Comparator_A+ Control Register 1	531
		21.3.2	CACTL2, Comparator_A+, Control Register	532
		21.3.3	CAPD, Comparator_A+, Port Disable Register	532
22	ADC1	0		533
	22.1	ADC10	Introduction	534
	22.2	ADC10	Operation	536
		22.2.1	10-Bit ADC Core	536
		22.2.2	ADC10 Inputs and Multiplexer	536
		22.2.3	Voltage Reference Generator	537
		22.2.4	Auto Power-Down	537
		22.2.5	Sample and Conversion Timing	538
		22.2.6	Conversion Modes	539
		22.2.7	ADC10 Data Transfer Controller	544
		22.2.8	Using the Integrated Temperature Sensor	549
		22.2.9	ADC10 Grounding and Noise Considerations	550
		22.2.10	ADC10 Interrupts	551
	22.3	ADC10	Registers	552
		22.3.1	ADC10CTL0, ADC10 Control Register 0	553
		22.3.2	ADC10CTL1, ADC10 Control Register 1	555
		22.3.3	ADC10AE0, Analog (Input) Enable Control Register 0	
		22.3.4	ADC10AE1, Analog (Input) Enable Control Register 1 (MSP430F22xx only)	556
		22.3.5	ADC10MEM, Conversion-Memory Register, Binary Format	556
		22.3.6	ADC10MEM, Conversion-Memory Register, 2s Complement Format	557
		22.3.7	ADC10DTC0, Data Transfer Control Register 0	557
		22.3.8	ADC10DTC1, Data Transfer Control Register 1	557
		22.3.9	ADC10SA, Start Address Register for Data Transfer	558
23	ADC1	2		559
	23.1		Introduction	
	23.2		Operation	
		23.2.1	12-Bit ADC Core	
		23.2.2	ADC12 Inputs and Multiplexer	
		23.2.3	Voltage Reference Generator	
		23.2.4	Sample and Conversion Timing	
			. ·	

9





		23.2.5	Conversion Memory	565
		23.2.6	ADC12 Conversion Modes	
		23.2.7	Using the Integrated Temperature Sensor	570
			ADC12 Grounding and Noise Considerations	
			ADC12 Interrupts	
	23.3		Registers	
		23.3.1	ADC12CTL0, ADC12 Control Register 0	575
			ADC12CTL1, ADC12 Control Register 1	
		23.3.3	ADC12MEMx, ADC12 Conversion Memory Registers	578
		23.3.4	ADC12MCTLx, ADC12 Conversion Memory Control Registers	578
		23.3.5	ADC12IE, ADC12 Interrupt Enable Register	579
			ADC12IFG, ADC12 Interrupt Flag Register	
		23.3.7	ADC12IV, ADC12 Interrupt Vector Register	580
24	TLV S	Structur	e	<b>581</b>
	24.1	TLV Inti	roduction	582
	24.2	Support	ted Tags	583
		24.2.1	DCO Calibration TLV Structure	583
		24.2.2	TAG_ADC12_1 Calibration TLV Structure	584
	24.3		ng Integrity of SegmentA	
	24.4		TLV Structure of Segment A	
25	DAC1	2		588
	25.1		Introduction	
	25.2		Operation	
	20.2		DAC12 Core	
			DAC12 Reference	
			Updating the DAC12 Voltage Output	
			DAC12_xDAT Data Format	
		25.2.5	DAC12 Output Amplifier Offset Calibration	
		25.2.6	Grouping Multiple DAC12 Modules	
		25.2.7	DAC12 Interrupts	
	25.3	-	Registers	
		25.3.1	DAC12_xCTL, DAC12 Control Register	
		25.3.2	DAC12_xDAT, DAC12 Data Register	
26	SD16		- · · ·	
20	26.1	_	A Introduction	
	26.2	_	A Operation	601
	20.2	26.2.1	ADC Core	
		_	Analog Input Range and PGA	
			Voltage Reference Generator	
		26.2.4	Auto Power-Down	
		-	Analog Input Pair Selection	
			Analog Input Characteristics	
		26.2.7	Digital Filter	
		26.2.8	Conversion Memory Register: SD16MEM0	
			Conversion Modes	
		26.2.10		
		26.2.10		
	26.3		A Registers	611
	20.0	26.3.1	SD16CTL, SD16_A Control Register	
			SD16CCTL0, SD16_A Control Register 0	
		26.3.3	SD16INCTL0, SD16_A Input Control Register	
			SD16MEM0, SD16_A Conversion Memory Register	
				٥.٠



		26.3.5 SI	D16AE, SD16_A Analog Input Enable Register	615
		26.3.6 SI	D16IV, SD16_A Interrupt Vector Register	615
27	SD24	Α		616
	27.1	SD24_A Ir	ntroduction	617
	27.2	SD24_A C	Operation	619
		27.2.1 Al	DC Core	619
		27.2.2 Aı	nalog Input Range and PGA	619
		27.2.3 V	oltage Reference Generator	619
		27.2.4 A	uto Power-Down	619
		27.2.5 Aı	nalog Input Pair Selection	619
		27.2.6 Aı	nalog Input Characteristics	620
		27.2.7 Di	igital Filter	621
		27.2.8 C	onversion Memory Register: SD24MEMx	625
		27.2.9 C	onversion Modes	626
		27.2.10	Conversion Operation Using Preload	628
		27.2.11 l	Using the Integrated Temperature Sensor	629
		27.2.12 I	Interrupt Handling	630
	27.3	SD24_A R	Registers	632
		27.3.1 SI	D24CTL, SD24_A Control Register	633
		27.3.2 SI	D24CCTLx, SD24_A Channel x Control Register	634
		27.3.3 SI	D24INCTLx, SD24_A Channel x Input Control Register	635
		27.3.4 SI	D24MEMx, SD24_A Channel x Conversion Memory Register	636
		27.3.5 SI	D24PREx, SD24_A Channel x Preload Register	636
		27.3.6 SI	D24AE, SD24_A Analog Input Enable Register	636
		27.3.7 SI	D24IV, SD24_A Interrupt Vector Register	637
28	Embe	dded Em	ulation Module (EEM)	638
	28.1		duction	
	28.2	EEM Build	ding Blocks	641
		28.2.1 Tr	riggers	641
		28.2.2 Tr	rigger Sequencer	641
		28.2.3 St	tate Storage (Internal Trace Buffer)	641
			lock Control	
	28.3		figurations	
Revis	ion His	story		643



# **List of Figures**

1-1.	MSP430 Architecture	24
1-2.	Memory Map	25
1-3.	Bits, Bytes, and Words in a Byte-Organized Memory	26
2-1.	Power-On Reset and Power-Up Clear Schematic	29
2-2.	Brownout Timing	30
2-3.	Interrupt Priority	31
2-4.	Block Diagram of (Non)-Maskable Interrupt Sources	32
2-5.	NMI Interrupt Handler	34
2-6.	Interrupt Processing	
2-7.	Return From Interrupt	36
2-8.	Typical Current Consumption of 'F21x1 Devices vs Operating Modes	38
2-9.	Operating Modes For Basic Clock System	39
3-1.	CPU Block Diagram	44
3-2.	Program Counter	44
3-3.	Stack Counter	45
3-4.	Stack Usage	45
3-5.	PUSH SP - POP SP Sequence	45
3-6.	Status Register Bits	46
3-7.	Register-Byte/Byte-Register Operations	47
3-8.	Operand Fetch Operation	
3-9.	Double Operand Instruction Format	57
3-10.	Single Operand Instruction Format	58
3-11.	Jump Instruction Format	59
3-12.	Core Instruction Map	62
3-13.	Decrement Overlap	80
3-14.	Main Program Interrupt	100
3-15.	Destination Operand – Arithmetic Shift Left	101
3-16.	Destination Operand - Carry Left Shift	102
3-17.	Destination Operand – Arithmetic Right Shift	103
3-18.	Destination Operand - Carry Right Shift	104
3-19.	Destination Operand - Byte Swap	111
3-20.	Destination Operand - Sign Extension	112
4-1.	MSP430X CPU Block Diagram	117
4-2.	PC Storage on the Stack for Interrupts	118
4-3.	·	119
4-4.	PC Storage on the Stack for CALLA	119
4-5.	Stack Pointer	120
4-6.	Stack Usage	120
4-7.	PUSHX.A Format on the Stack	120
4-8.	PUSH SP, POP SP Sequence	120
4-9.	SR Bits	121
4-10.	Register-Byte/Byte-Register Operation	123
4-11.	Register-Word Operation	123
4-12.	Word-Register Operation	124
4-13.	Register – Address-Word Operation	124
4-14.	Address-Word – Register Operation	125
4-15.	Indexed Mode in Lower 64KB	127



4-16.	Indexed Mode in Upper Memory	128
4-17.	Overflow and Underflow for Indexed Mode	
4-18.	Symbolic Mode Running in Lower 64KB	
4-19.	Symbolic Mode Running in Upper Memory	133
4-20.	Overflow and Underflow for Symbolic Mode	134
4-21.	MSP430 Double-Operand Instruction Format	142
4-22.	MSP430 Single-Operand Instructions	143
4-23.	Format of Conditional Jump Instructions	144
4-24.	Extension Word for Register Modes	147
4-25.	Extension Word for Non-Register Modes	149
4-26.	Example for Extended Register/Register Instruction	150
4-27.	Example for Extended Immediate/Indexed Instruction	150
4-28.	Extended Format I Instruction Formats	152
4-29.	20-Bit Addresses in Memory	152
4-30.	Extended Format II Instruction Format	153
4-31.	PUSHM/POPM Instruction Format	154
4-32.	RRCM, RRAM, RRUM, and RLAM Instruction Format	154
4-33.	BRA Instruction Format	154
4-34.	CALLA Instruction Format	154
4-35.	Decrement Overlap	180
4-36.	Stack After a RET Instruction	199
4-37.	Destination Operand—Arithmetic Shift Left	201
4-38.	Destination Operand—Carry Left Shift	202
4-39.	Rotate Right Arithmetically RRA.B and RRA.W	203
4-40.	Rotate Right Through Carry RRC.B and RRC.W	204
4-41.	Swap Bytes in Memory	211
4-42.	Swap Bytes in a Register	211
4-43.	Rotate Left Arithmetically—RLAM[.W] and RLAM.A	238
4-44.	Destination Operand-Arithmetic Shift Left	239
4-45.	Destination Operand-Carry Left Shift	240
4-46.	Rotate Right Arithmetically RRAM[.W] and RRAM.A	241
4-47.	Rotate Right Arithmetically RRAX(.B,.A) – Register Mode	243
4-48.	Rotate Right Arithmetically RRAX(.B,.A) – Non-Register Mode	243
4-49.	Rotate Right Through Carry RRCM[.W] and RRCM.A	244
4-50.	Rotate Right Through Carry RRCX(.B,.A) – Register Mode	246
4-51.	Rotate Right Through Carry RRCX(.B,.A) – Non-Register Mode	246
4-52.	Rotate Right Unsigned RRUM[.W] and RRUM.A	247
4-53.	Rotate Right Unsigned RRUX(.B,.A) – Register Mode	248
4-54.	Swap Bytes SWPBX.A Register Mode	252
4-55.	Swap Bytes SWPBX.A In Memory	252
4-56.	Swap Bytes SWPBX[.W] Register Mode	253
4-57.	Swap Bytes SWPBX[.W] In Memory	253
4-58.	Sign Extend SXTX.A	254
4-59.	Sign Extend SXTX[.W]	254
5-1.	Basic Clock Module+ Block Diagram - MSP430F2xx	274
5-2.	Basic Clock Module+ Block Diagram - MSP430AFE2xx	275
5-3.	Off Signals for the LFXT1 Oscillator	277
5-4.	Off Signals for Oscillator XT2	277
5-5.	On/Off Control of DCO	278





5-6.	Typical DCOx Range and RSELx Steps	278
5-7.	Modulator Patterns	279
5-8.	Oscillator-Fault Logic	280
5-9.	Switch MCLK from DCOCLK to LFXT1CLK	281
6-1.	DMA Controller Block Diagram	289
6-2.	DMA Addressing Modes	290
6-3.	DMA Single Transfer State Diagram	292
6-4.	DMA Block Transfer State Diagram	294
6-5.	DMA Burst-Block Transfer State Diagram	296
7-1.	Flash Memory Module Block Diagram	309
7-2.	Flash Memory Segments, 32-KB Example	310
7-3.	Flash Memory Timing Generator Block Diagram	311
7-4.	Erase Cycle Timing	312
7-5.	Erase Cycle from Within Flash Memory	313
7-6.	Erase Cycle from Within RAM	314
7-7.	Byte or Word Write Timing	315
7-8.	Initiating a Byte or Word Write From Flash	316
7-9.	Initiating a Byte or Word Write from RAM	317
7-10.	Block-Write Cycle Timing	318
7-11.	Block Write Flow	319
7-12.	User-Developed Programming Solution	322
8-1.	Example Circuitry and Configuration using the Pin Oscillator	330
8-2.	Typical Pin-Oscillation Frequency	331
9-1.	SVS Block Diagram	336
9-2.	Operating Levels for SVS and Brownout/Reset Circuit	338
10-1.	Watchdog Timer+ Block Diagram	343
11-1.	Hardware Multiplier Block Diagram	350
12-1.	Timer_A Block Diagram	357
12-2.	Up Mode	358
12-3.	Up Mode Flag Setting	359
12-4.	Continuous Mode	359
12-5.	Continuous Mode Flag Setting	359
12-6.	Continuous Mode Time Intervals	360
12-7.	Up/Down Mode	360
12-8.	Up/Down Mode Flag Setting	361
12-9.	Output Unit in Up/Down Mode	362
12-10.	Capture Signal (SCS = 1)	362
12-11.	Capture Cycle	363
12-12.	Output Example—Timer in Up Mode	364
12-13.	Output Example—Timer in Continuous Mode	365
12-14.	Output Example—Timer in Up/Down Mode	366
12-15.	Capture/Compare TACCR0 Interrupt Flag	367
13-1.	Timer_B Block Diagram	376
13-2.	Up Mode	378
13-3.	Up Mode Flag Setting	378
13-4.	Continuous Mode	378
13-5.	Continuous Mode Flag Setting	379
13-6.	Continuous Mode Time Intervals	379
13-7.	Up/Down Mode	380



13-8.	Up/Down Mode Flag Setting	380
13-9.	Output Unit in Up/Down Mode	381
13-10.	Capture Signal (SCS = 1)	381
13-11.	Capture Cycle	382
13-12.	Output Example, Timer in Up Mode	385
13-13.	Output Example, Timer in Continuous Mode	386
13-14.	Output Example, Timer in Up/Down Mode	387
13-15.	Capture/Compare TBCCR0 Interrupt Flag	388
14-1.	USI Block Diagram: SPI Mode	397
14-2.	USI Block Diagram: I <sup>2</sup> C Mode	398
14-3.	SPI Timing	400
14-4.	Data Adjustments for 7-Bit SPI Data	401
15-1.	USCI_Ax Block Diagram: UART Mode (UCSYNC = 0)	412
15-2.	Character Format	413
15-3.	Idle-Line Format	414
15-4.	Address-Bit Multiprocessor Format	415
15-5.	Auto Baud Rate Detection - Break/Synch Sequence	416
15-6.	Auto Baud Rate Detection - Synch Field	416
15-7.	UART vs IrDA Data Format	417
15-8.	Glitch Suppression, USCI Receive Not Started	
15-9.	Glitch Suppression, USCI Activated	
	BITCLK Baud Rate Timing With UCOS16 = 0	
15-11.	Receive Error	423
16-1.	USCI Block Diagram: SPI Mode	437
16-2.	USCI Master and External Slave	439
16-3.	USCI Slave and External Master	440
16-4.	USCI SPI Timing with UCMSB = 1	442
17-1.	USCI Block Diagram: I <sup>2</sup> C Mode	
17-2.	I <sup>2</sup> C Bus Connection Diagram	452
17-3.	I <sup>2</sup> C Module Data Transfer	
17-4.	Bit Transfer on the I <sup>2</sup> C Bus	
17-5.	I <sup>2</sup> C Module 7-Bit Addressing Format	
	I <sup>2</sup> C Module 10-Bit Addressing Format	453
17-7.	I <sup>2</sup> C Module Addressing Format with Repeated START Condition	454
17-8.	I <sup>2</sup> C Time Line Legend	454
17-9.	I <sup>2</sup> C Slave Transmitter Mode	455
	I <sup>2</sup> C Slave Receiver Mode	457
	I <sup>2</sup> C Slave 10-bit Addressing Mode	458
	I <sup>2</sup> C Master Transmitter Mode	460
	I <sup>2</sup> C Master Receiver Mode	462
	I <sup>2</sup> C Master 10-bit Addressing Mode	463
	Arbitration Procedure Between Two Master Transmitters	463
17-16.	Synchronization of Two I <sup>2</sup> C Clock Generators During Arbitration	464
18-1.	USART Block Diagram: UART Mode	476
18-2.	Character Format	477
18-3.	Idle-Line Format	478
18-4.	Address-Bit Multiprocessor Format	479
18-5.	State Diagram of Receiver Enable	480
18-6.	State Diagram of Transmitter Enable	481





18-7.	MSP430 Baud Rate Generator	481
18-8.	BITCLK Baud Rate Timing	482
18-9.	Receive Error	485
18-10.	Transmit Interrupt Operation	487
18-11.	Receive Interrupt Operation	487
18-12.	Glitch Suppression, USART Receive Not Started	489
18-13.	Glitch Suppression, USART Activated	489
19-1.	USART Block Diagram: SPI Mode	498
19-2.	USART Master and External Slave	500
19-3.	USART Slave and External Master	501
19-4.	Master Transmit Enable State Diagram	501
19-5.	Slave Transmit Enable State Diagram	502
19-6.	SPI Master Receive-Enable State Diagram	502
19-7.	SPI Slave Receive-Enable State Diagram	502
19-8.	SPI Baud Rate Generator	503
19-9.	USART SPI Timing	503
19-10.	Transmit Interrupt Operation	504
19-11.	Receive Interrupt Operation	505
19-12.	Receive Interrupt State Diagram	505
20-1.	OA Block Diagram	513
20-2.	Two-Opamp Differential Amplifier	516
20-3.	Two-Opamp Differential Amplifier OAx Interconnections	517
20-4.	Three-Opamp Differential Amplifier	518
20-5.	Three-Opamp Differential Amplifier OAx Interconnections	519
21-1.	Comparator_A+ Block Diagram	524
21-2.	Comparator_A+ Sample-And-Hold	526
21-3.	RC-Filter Response at the Output of the Comparator	527
21-4.	Transfer Characteristic and Power Dissipation in a CMOS Inverter/Buffer	527
21-5.	Comparator_A+ Interrupt System	528
21-6.	Temperature Measurement System	
21-7.	Timing for Temperature Measurement Systems	
22-1.	ADC10 Block Diagram	
22-2.	Analog Multiplexer	536
22-3.	Sample Timing	538
22-4.	Analog Input Equivalent Circuit	538
22-5.	Single-Channel Single-Conversion Mode	540
22-6.	Sequence-of-Channels Mode	541
22-7.	Repeat-Single-Channel Mode	542
22-8.	Repeat-Sequence-of-Channels Mode	543
22-9.	One-Block Transfer	545
	State Diagram for Data Transfer Control in One-Block Transfer Mode	546
	Two-Block Transfer	547
	State Diagram for Data Transfer Control in Two-Block Transfer Mode	548
	Typical Temperature Sensor Transfer Function	550
	ADC10 Grounding and Noise Considerations (Internal V <sub>REF</sub> )	550
	ADC10 Grounding and Noise Considerations (External V <sub>REF</sub> )	551
	ADC10 Interrupt System	551
	ADC12 Block Diagram	561
23-2.	Analog Multiplexer	562



23-3.	Extended Sample Mode	564
23-4.	Pulse Sample Mode	564
23-5.	Analog Input Equivalent Circuit	565
23-6.	Single-Channel, Single-Conversion Mode	566
23-7.	Sequence-of-Channels Mode	567
23-8.	Repeat-Single-Channel Mode	568
23-9.	Repeat-Sequence-of-Channels Mode	569
23-10.	Typical Temperature Sensor Transfer Function	571
23-11.	ADC12 Grounding and Noise Considerations	572
25-1.	DAC12 Block Diagram	590
25-2.	Output Voltage vs DAC12 Data, 12-Bit, Straight Binary Mode	592
25-3.	Output Voltage vs DAC12 Data, 12-Bit, 2s-Compliment Mode	592
25-4.	Negative Offset	593
25-5.	Positive Offset	593
25-6.	DAC12 Group Update Example, Timer_A3 Trigger	594
26-1.	SD16_A Block Diagram	600
26-2.	Analog Input Equivalent Circuit	602
26-3.	Comb Filter Frequency Response With OSR = 32	603
26-4.	Digital Filter Step Response and Conversion Points	604
26-5.	Used Bits of Digital Filter Output	606
26-6.	Input Voltage vs Digital Output	607
26-7.	Single Channel Operation	608
26-8.	Typical Temperature Sensor Transfer Function	609
27-1.	Block Diagram of the SD24_A	618
27-2.	Analog Input Equivalent Circuit	620
27-3.	Comb Filter Frequency Response With OSR = 32	622
27-4.	Digital Filter Step Response and Conversion Points	622
27-5.	Used Bits of Digital Filter Output	624
27-6.	Input Voltage vs Digital Output	625
27-7.	Single Channel Operation - Example	626
27-8.	Grouped Channel Operation - Example	627
27-9.	Conversion Delay Using Preload - Example	628
27-10.	Start of Conversion Using Preload - Example	628
27-11.	Preload and Channel Synchronization	629
27-12.	Typical Temperature Sensor Transfer Function	629
28-1.	Large Implementation of the Embedded Emulation Module (EEM)	640

17



# **List of Tables**

1-1.	MSP430x2xx Family Enhancements	27
2-1.	Interrupt Sources, Flags, and Vectors	37
2-2.	Operating Modes For Basic Clock System	39
2-3.	Connection of Unused Pins	41
3-1.	Description of Status Register Bits	46
3-2.	Values of Constant Generators CG1, CG2	46
3-3.	Source/Destination Operand Addressing Modes	48
3-4.	Register Mode Description	49
3-5.	Indexed Mode Description	50
3-6.	Symbolic Mode Description	51
3-7.	Absolute Mode Description	52
3-8.	Indirect Mode Description	53
3-9.	Indirect Autoincrement Mode Description	54
3-10.	Immediate Mode Description	55
3-11.	Double Operand Instructions	
3-12.	Single Operand Instructions	58
3-13.	Jump Instructions	59
3-14.	Interrupt and Reset Cycles	
3-15.	Format-II Instruction Cycles and Lengths	
3-16.	Format 1 Instruction Cycles and Lengths	
3-17.	MSP430 Instruction Set	62
4-1.	SR Bit Description	121
4-2.	Values of Constant Generators CG1, CG2	122
4-3.	Source/Destination Addressing	125
4-4.	MSP430 Double-Operand Instructions	143
4-5.	MSP430 Single-Operand Instructions	143
4-6.	Conditional Jump Instructions	144
4-7.	Emulated Instructions	144
4-8.	Interrupt, Return, and Reset Cycles and Length	145
4-9.	MSP430 Format II Instruction Cycles and Length	145
4-10.	MSP430 Format I Instructions Cycles and Length	146
4-11.	Description of the Extension Word Bits for Register Mode	147
4-12.	Description of Extension Word Bits for Non-Register Modes	149
4-13.	Extended Double-Operand Instructions	151
4-14.	Extended Single-Operand Instructions	153
4-15.	Extended Emulated Instructions	155
4-16.	Address Instructions, Operate on 20-Bit Register Data	156
4-17.	MSP430X Format II Instruction Cycles and Length	157
4-18.	MSP430X Format I Instruction Cycles and Length	158
4-19.	Address Instruction Cycles and Length	159
4-20.	Instruction Map of MSP430X	160
5-1.	Basic Clock Module+ Registers	282
6-1.	DMA Transfer Modes	291
6-2.	DMA Trigger Operation	297
6-3.	Channel Priorities	299
6-4.	Maximum Single-Transfer DMA Cycle Time	299
6-5.	DMA Registers	302



7.4	Free Modes	242
7-1.	Erase Modes.	312
7-2.	Write Modes	315
7-3.	Flash Access While BUSY = 1	320
7-4.	Flash Memory Registers	
8-1.	PxSEL and PxSEL2	329
8-2.	Digital I/O Registers	333
9-1.	SVS Registers	339
10-1.	Watchdog Timer+ Registers	346
11-1.	OP1 Addresses	351
11-2.	RESHI Contents	
11-3.	SUMEXT Contents	
11-4.	Hardware Multiplier Registers	
12-1.	Timer Modes	358
12-2.	Output Modes	364
12-3.	Timer_A3 Registers	369
13-1.	Timer Modes	377
13-2.	TBCLx Load Events	383
13-3.	Compare Latch Operating Modes	383
13-4.	Output Modes	384
13-5.	Timer_B Registers	390
14-1.	USI Registers	405
14-2.	Word Access to USI Registers	405
15-1.	Receive Error Conditions	418
15-2.	BITCLK Modulation Pattern	420
15-3.	BITCLK16 Modulation Pattern	421
15-4.	Commonly Used Baud Rates, Settings, and Errors, UCOS16 = 0	424
15-5.	Commonly Used Baud Rates, Settings, and Errors, UCOS16 = 1	425
15-6.	USCI_A0 Control and Status Registers	428
15-7.	USCI_A1 Control and Status Registers	428
16-1.	UCxSTE Operation	438
16-2.	USCI_A0 and USCI_B0 Control and Status Registers	444
16-3.	USCI_A1 and USCI_B1 Control and Status Registers	444
17-1.	State Change Interrupt Flags	465
17-2.	USCI_B0 Control and Status Registers	467
17-3.	USCI_B1 Control and Status Registers	467
18-1.	Receive Error Conditions	480
18-2.	Commonly Used Baud Rates, Baud Rate Data, and Errors	486
18-3.	USART0 Control and Status Registers	490
18-4.	USART1 Control and Status Registers	490
19-1.	USART0 Control and Status Registers	506
19-2.	USART1 Control and Status Registers	506
20-1.	OA Output Configurations	514
20-2.	OA Mode Select	514
20-3.	Two-Opamp Differential Amplifier Control Register Settings	516
20-4.	Two-Opamp Differential Amplifier Gain Settings	516
20-5.	Three-Opamp Differential Amplifier Control Register Settings	518
20-6.	Three-Opamp Differential Amplifier Gain Settings	518
20-7.	OA Registers	520
21-1.	Comparator_A+ Registers	530
	· = •	



22-1.	Conversion Mode Summary	539
22-2.	Maximum DTC Cycle Time	549
22-3.	ADC10 Registers	552
23-1.	Conversion Mode Summary	565
23-2.	ADC12 Registers	574
24-1.	Example SegmentA Structure	582
24-2.	Supported Tags (Device Specific)	583
24-3.	DCO Calibration Data (Device Specific)	583
24-4.	TAG_ADC12_1 Calibration Data (Device Specific)	584
25-1.	DAC12 Full-Scale Range (V <sub>REF</sub> = V <sub>eREF+</sub> or V <sub>REF+</sub> )	591
25-2.	DAC12 Registers	595
26-1.	High Input Impedance Buffer	602
26-2.	Sampling Capacitance	603
26-3.	Data Format	607
26-4.	Conversion Mode Summary	608
26-5.	SD16_A Registers	611
27-1.	High Input Impedance Buffer	620
27-2.	Sampling Capacitance	621
27-3.	Data Format	625
27-4.	Conversion Mode Summary	626
27-5.	SD24_A Registers	632
28-1	2vv EEM Configurations	6/2



# Read This First

#### **About This Manual**

This manual discusses modules and peripherals of the MSP430x2xx family of devices. Each discussion presents the module or peripheral in a general sense. Not all features and functions of all modules or peripherals are present on all devices. In addition, modules or peripherals may differ in their exact implementation between device families, or may not be fully implemented on an individual device or device family.

Pin functions, internal signal connections, and operational parameters differ from device to device. The user should consult the device-specific datasheet for these details.

#### **Related Documentation From Texas Instruments**

For related documentation see the web site http://www.ti.com/msp430.

# **FCC Warning**

This equipment is intended for use in a laboratory test environment only. It generates, uses, and can radiate radio frequency energy and has not been tested for compliance with the limits of computing devices pursuant to subpart J of part 15 of FCC rules, which are designed to provide reasonable protection against radio frequency interference. Operation of this equipment in other environments may cause interference with radio communications, in which case the user at his own expense will be required to take whatever measures may be required to correct this interference.

#### **Notational Conventions**

Program examples, are shown in a special typeface.

#### **Glossary**

ACLK	Auxiliary Clock	See Basic Clock Module
ADC	Analog-to-Digital Converter	
BOR	Brown-Out Reset	See System Resets, Interrupts, and Operating Modes
BSL	Bootstrap Loader	See www.ti.com/msp430for application reports
CPU	Central Processing Unit	See RISC 16-Bit CPU
DAC	Digital-to-Analog Converter	
DCO	Digitally Controlled Oscillator	See Basic Clock Module
dst	Destination	See RISC 16-Bit CPU
FLL	Frequency Locked Loop	See FLL+in MSP430x4xx Family User's Guide
GIE	General Interrupt Enable	See System Resets, Interrupts, and Operating Modes
INT(N/2)	Integer portion of N/2	
I/O	Input/Output	See Digital I/O
ISR	Interrupt Service Routine	
LSB	Least-Significant Bit	
LSD	Least-Significant Digit	
LPM	Low-Power Mode	See System Resets, Interrupts, and Operating Modes
MAB	Memory Address Bus	
MCLK	Master Clock	See Basic Clock Module



MDB MSB MSD	Memory Data Bus Most-Significant Bit Most-Significant Digit	
NMI	(Non)-Maskable Interrupt	See System Resets, Interrupts, and Operating Modes
PC	Program Counter	See RISC 16-Bit CPU
POR	Power-On Reset	See System Resets, Interrupts, and Operating Modes
PUC	Power-Up Clear	See System Resets, Interrupts, and Operating Modes
RAM	Random Access Memory	
SCG	System Clock Generator	See System Resets, Interrupts, and Operating Modes
SFR	Special Function Register	
SMCLK	Sub-System Master Clock	See Basic Clock Module
SP	Stack Pointer	See RISC 16-Bit CPU
SR	Status Register	See RISC 16-Bit CPU
src	Source	See RISC 16-Bit CPU
TOS	Top-of-Stack	See RISC 16-Bit CPU
WDT	Watchdog Timer	See Watchdog Timer

# **Register Bit Conventions**

Each register is shown with a key indicating the accessibility of the each individual bit, and the initial condition:

# **Register Bit Accessibility and Initial Condition**

Key	Bit Accessibility
rw	Read/write
r	Read only
r0	Read as 0
r1	Read as 1
w	Write only
w0	Write as 0
w1	Write as 1
(w)	No register bit implemented; writing a 1 results in a pulse. The register bit is always read as 0.
h0	Cleared by hardware
h1	Set by hardware
-0,-1	Condition after PUC
-(0),-(1)	Condition after POR



# Introduction

This chapter describes the architecture of the MSP430.

Topic		Page
1.1	Architecture	24
1.2	Flexible Clock System	24
1.3	Embedded Emulation	25
1.4	Address Space	25
1.5	MSP430x2xx Family Enhancements	27
	•	



Architecture www.ti.com

#### 1.1 Architecture

The MSP430 incorporates a 16-bit RISC CPU, peripherals, and a flexible clock system that interconnect using a von-Neumann common memory address bus (MAB) and memory data bus (MDB) (see Figure 1-1). Partnering a modern CPU with modular memory-mapped analog and digital peripherals, the MSP430 offers solutions for demanding mixed-signal applications.

Key features of the MSP430x2xx family include:

- Ultralow-power architecture extends battery life
  - 0.1 µA RAM retention
  - 0.8 µA real-time clock mode
  - 250 µA/MIPS active
- High-performance analog ideal for precision measurement
  - Comparator-gated timers for measuring resistive elements
- 16-bit RISC CPU enables new applications at a fraction of the code size.
  - Large register file eliminates working file bottleneck
  - Compact core design reduces power consumption and cost
  - Optimized for modern high-level programming
  - Only 27 core instructions and seven addressing modes
  - Extensive vectored-interrupt capability
- In-system programmable Flash permits flexible code changes, field upgrades and data logging

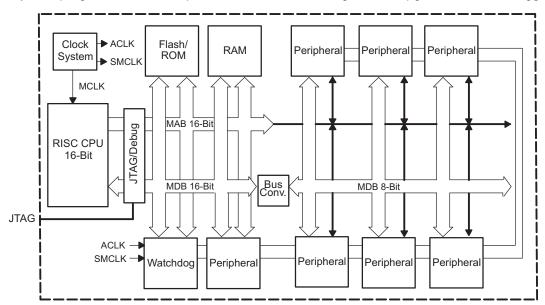


Figure 1-1. MSP430 Architecture

#### 1.2 Flexible Clock System

The clock system is designed specifically for battery-powered applications. A low-frequency auxiliary clock (ACLK) is driven directly from a common 32-kHz watch crystal. The ACLK can be used for a background real-time clock self wake-up function. An integrated high-speed digitally controlled oscillator (DCO) can source the master clock (MCLK) used by the CPU and high-speed peripherals. By design, the DCO is active and stable in less than 2 µs at 1 MHz. MSP430-based solutions effectively use the high-performance 16-bit RISC CPU in very short bursts.

- Low-frequency auxiliary clock = Ultralow-power stand-by mode
- High-speed master clock = High performance signal processing



Embedded Emulation www.ti.com

#### 1.3 **Embedded Emulation**

Dedicated embedded emulation logic resides on the device itself and is accessed via JTAG using no additional system resources.

The benefits of embedded emulation include:

- Unobtrusive development and debug with full-speed execution, breakpoints, and single-steps in an application are supported.
- Development is in-system subject to the same characteristics as the final application.
- Mixed-signal integrity is preserved and not subject to cabling interference.

#### 1.4 **Address Space**

The MSP430 von-Neumann architecture has one address space shared with special function registers (SFRs), peripherals, RAM, and Flash/ROM memory as shown in Figure 1-2. See the device-specific data sheets for specific memory maps. Code access are always performed on even addresses. Data can be accessed as bytes or words.

The addressable memory space is currently 128 KB.

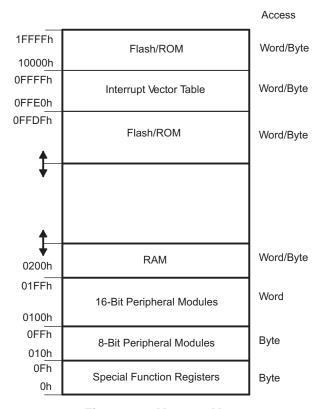


Figure 1-2. Memory Map

# 1.4.1 Flash/ROM

The start address of Flash/ROM depends on the amount of Flash/ROM present and varies by device. The end address for Flash/ROM is 0x0FFFF for devices with less that 60KB of Flash/ROM. Flash can be used for both code and data. Word or byte tables can be stored and used in Flash/ROM without the need to copy the tables to RAM before using them.

The interrupt vector table is mapped into the upper 16 words of Flash/ROM address space, with the highest priority interrupt vector at the highest Flash/ROM word address (0x0FFFE).



Address Space www.ti.com

# 1.4.2 RAM

RAM starts at 0200h. The end address of RAM depends on the amount of RAM present and varies by device. RAM can be used for both code and data.

### 1.4.3 Peripheral Modules

Peripheral modules are mapped into the address space. The address space from 0100 to 01FFh is reserved for 16-bit peripheral modules. These modules should be accessed with word instructions. If byte instructions are used, only even addresses are permissible, and the high byte of the result is always 0.

The address space from 010h to 0FFh is reserved for 8-bit peripheral modules. These modules should be accessed with byte instructions. Read access of byte modules using word instructions results in unpredictable data in the high byte. If word data is written to a byte module only the low byte is written into the peripheral register, ignoring the high byte.

# 1.4.4 Special Function Registers (SFRs)

Some peripheral functions are configured in the SFRs. The SFRs are located in the lower 16 bytes of the address space, and are organized by byte. SFRs must be accessed using byte instructions only. See the device-specific data sheets for applicable SFR bits.

## 1.4.5 Memory Organization

Bytes are located at even or odd addresses. Words are only located at even addresses as shown in Figure 1-3. When using word instructions, only even addresses may be used. The low byte of a word is always an even address. The high byte is at the next odd address. For example, if a data word is located at address xxx4h, then the low byte of that data word is located at address xxx4h, and the high byte of that word is located at address xxx5h.

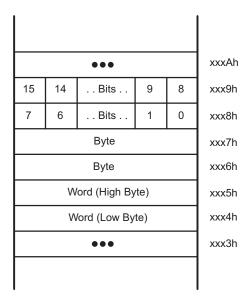


Figure 1-3. Bits, Bytes, and Words in a Byte-Organized Memory



# 1.5 MSP430x2xx Family Enhancements

Table 1-1 highlights enhancements made to the MSP430x2xx family. The enhancements are discussed fully in the following chapters, or in the case of improved device parameters, shown in the device-specific data sheet.

Table 1-1. MSP430x2xx Family Enhancements

Subject	Enhancement
Reset	<ul> <li>Brownout reset is included on all MSP430x2xx devices.</li> <li>PORIFG and RSTIFG flags have been added to IFG1 to indicate the cause of a reset.</li> <li>An instruction fetch from the address range 0x0000 - 0x01FF will reset the device.</li> </ul>
Watchdog Timer	<ul> <li>All MSP430x2xx devices integrate the Watchdog Timer+ module (WDT+). The WDT+ ensures the clock source for the timer is never disabled.</li> </ul>
Basic Clock System	<ul> <li>The LFXT1 oscillator has selectable load capacitors in LF mode.</li> <li>The LFXT1 supports up to 16-MHz crystals in HF mode.</li> <li>The LFXT1 includes oscillator fault detection in LF mode.</li> <li>The XIN and XOUT pins are shared function pins on 20- and 28-pin devices.</li> <li>The external R oscfeature of the DCO not supported on some devices. Software should not set the LSB of the BCSCTL2 register in this case. See the device-specific data sheet for details.</li> <li>The DCO operating frequency has been significantly increased.</li> <li>The DCO temperature stability has been significantly improved.</li> </ul>
Flash Memory	<ul> <li>The information memory has 4 segments of 64 bytes each.</li> <li>SegmentA is individually locked with the LOCKA bit.</li> <li>All information if protected from mass erase with the LOCKA bit.</li> <li>Segment erases can be interrupted by an interrupt.</li> <li>Flash updates can be aborted by an interrupt.</li> <li>Flash programming voltage has been lowered to 2.2 V</li> <li>Program/erase time has been reduced.</li> <li>Clock failure aborts a flash update.</li> </ul>
Digital I/O	<ul> <li>All ports have integrated pullup/pulldown resistors.</li> <li>P2.6 and P2.7 functions have been added to 20- and 28- pin devices. These are shared functions with XIN and XOUT. Software must not clear the P2SELx bits for these pins if crystal operation is required.</li> </ul>
Comparator_A	Comparator_A has expanded input capability with a new input multiplexer.
Low Power	<ul> <li>Typical LPM3 current consumption has been reduced almost 50% at 3 V.</li> <li>DCO startup time has been significantly reduced.</li> </ul>
Operating frequency	The maximum operating frequency is 16 MHz at 3.3 V.
BSL	<ul> <li>An incorrect password causes a mass erase.</li> <li>BSL entry sequence is more robust to prevent accidental entry and erasure.</li> </ul>



# System Resets, Interrupts, and Operating Modes

This chapter describes the MSP430x2xx system resets, interrupts, and operating modes.

 Topic
 Page

 2.1
 System Reset and Initialization
 29

 2.2
 Interrupts
 31

 2.3
 Operating Modes
 38

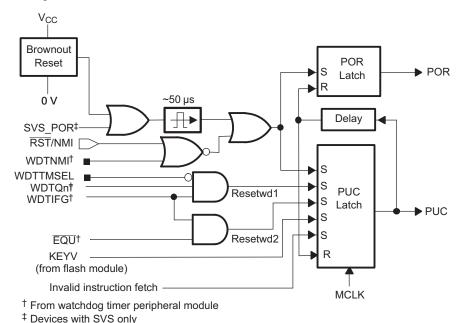
 2.4
 Principles for Low-Power Applications
 40

 2.5
 Connection of Unused Pins
 41



# 2.1 System Reset and Initialization

The system reset circuitry shown in Figure 2-1 sources both a power-on reset (POR) and a power-up clear (PUC) signal. Different events trigger these reset signals and different initial conditions exist depending on which signal was generated.



+ Devices with 3v3 only

A POR is a device reset. A POR is only generated by the following three events:

- Powering up the device
- A low signal on the RST/NMI pin when configured in the reset mode
- An SVS low condition when PORON = 1.

A PUC is always generated when a POR is generated, but a POR is not generated by a PUC. The following events trigger a PUC:

Figure 2-1. Power-On Reset and Power-Up Clear Schematic

- A POR signal
- · Watchdog timer expiration when in watchdog mode only
- Watchdog timer security key violation
- A Flash memory security key violation
- A CPU instruction fetch from the peripheral address range 0h to 01FFh

# 2.1.1 Brownout Reset (BOR)

The brownout reset circuit detects low supply voltages such as when a supply voltage is applied to or removed from the  $V_{\text{CC}}$  terminal. The brownout reset circuit resets the device by triggering a POR signal when power is applied or removed. The operating levels are shown in Figure 2-2.

The POR signal becomes active when  $V_{CC}$  crosses the  $V_{CC(start)}$  level. It remains active until  $V_{CC}$  crosses the  $V_{(B\_IT+)}$  threshold and the delay  $t_{(BOR)}$  elapses. The delay  $t_{(BOR)}$  is adaptive being longer for a slow ramping  $V_{CC}$ . The hysteresis  $V_{hys(B\_IT-)}$  ensures that the supply voltage must drop below  $V_{(B\_IT-)}$  to generate another POR signal from the brownout reset circuitry.



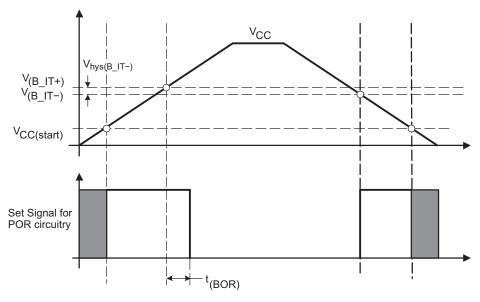


Figure 2-2. Brownout Timing

As the  $V_{(B\_IT-)}$  level is significantly above the  $V_{min}$  level of the POR circuit, the BOR provides a reset for power failures where  $V_{CC}$  does not fall below  $V_{min}$ . See device-specific data sheet for parameters.

# 2.1.2 Device Initial Conditions After System Reset

After a POR, the initial MSP430 conditions are:

- The RST/NMI pin is configured in the reset mode.
- I/O pins are switched to input mode as described in the *Digital I/O* chapter.
- Other peripheral modules and registers are initialized as described in their respective chapters in this manual.
- · Status register (SR) is reset.
- The watchdog timer powers up active in watchdog mode.
- Program counter (PC) is loaded with address contained at reset vector location (0FFFEh). If the reset vectors content is 0FFFFh the device will be disabled for minimum power consumption.

#### 2.1.2.1 Software Initialization

After a system reset, user software must initialize the MSP430 for the application requirements. The following must occur:

- Initialize the SP, typically to the top of RAM.
- Initialize the watchdog to the requirements of the application.
- Configure peripheral modules to the requirements of the application.

Additionally, the watchdog timer, oscillator fault, and flash memory flags can be evaluated to determine the source of the reset.



www.ti.com Interrupts

### 2.2 Interrupts

The interrupt priorities are fixed and defined by the arrangement of the modules in the connection chain as shown in Figure 2-3. The nearer a module is to the CPU/NMIRS, the higher the priority. Interrupt priorities determine what interrupt is taken when more than one interrupt is pending simultaneously.

There are three types of interrupts:

- System reset
- (Non)-maskable NMI
- Maskable

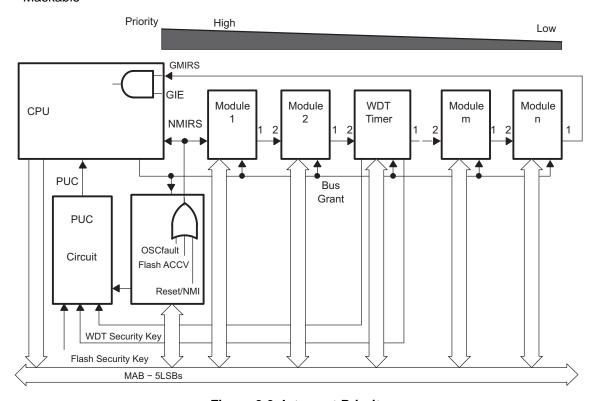


Figure 2-3. Interrupt Priority

# 2.2.1 (Non)-Maskable Interrupts (NMI)

(Non)-maskable NMI interrupts are not masked by the general interrupt enable bit (GIE), but are enabled by individual interrupt enable bits (NMIIE, ACCVIE, OFIE). When a NMI interrupt is accepted, all NMI interrupt enable bits are automatically reset. Program execution begins at the address stored in the (non)-maskable interrupt vector, 0FFFCh. User software must set the required NMI interrupt enable bits for the interrupt to be re-enabled. The block diagram for NMI sources is shown in Figure 2-4.

A (non)-maskable NMI interrupt can be generated by three sources:

- An edge on the RST/NMI pin when configured in NMI mode
- · An oscillator fault occurs
- An access violation to the flash memory



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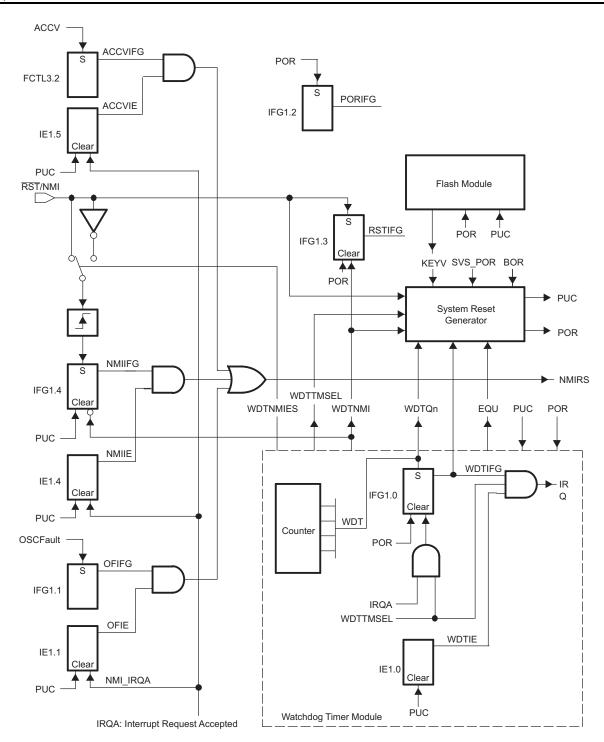


Figure 2-4. Block Diagram of (Non)-Maskable Interrupt Sources



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### 2.2.1.1 Reset/NMI Pin

At power-up, the RST/NMI pin is configured in the reset mode. The function of the RST/NMI pins is selected in the watchdog control register WDTCTL. If the RST/NMI pin is set to the reset function, the CPU is held in the reset state as long as the RST/NMI pin is held low. After the input changes to a high state, the CPU starts program execution at the word address stored in the reset vector, 0FFFEh, and the RSTIFG flag is set.

If the RST/NMI pin is configured by user software to the NMI function, a signal edge selected by the WDTNMIES bit generates an NMI interrupt if the NMIIE bit is set. The RST/NMI flag NMIIFG is also set.

#### NOTE: Holding RST/NMI Low

When configured in the NMI mode, a signal generating an NMI event should not hold the RST/NMI pin low. If a PUC occurs from a different source while the NMI signal is low, the device will be held in the reset state because a PUC changes the RST/NMI pin to the reset function.

#### **NOTE: Modifying WDTNMIES**

When NMI mode is selected and the WDTNMIES bit is changed, an NMI can be generated, depending on the actual level at the RST/NMI pin. When the NMI edge select bit is changed before selecting the NMI mode, no NMI is generated.

#### 2.2.1.2 Flash Access Violation

The flash ACCVIFG flag is set when a flash access violation occurs. The flash access violation can be enabled to generate an NMI interrupt by setting the ACCVIE bit. The ACCVIFG flag can then be tested by the NMI interrupt service routine to determine if the NMI was caused by a flash access violation.

#### 2.2.1.3 Oscillator Fault

The oscillator fault signal warns of a possible error condition with the crystal oscillator. The oscillator fault can be enabled to generate an NMI interrupt by setting the OFIE bit. The OFIFG flag can then be tested by NMI the interrupt service routine to determine if the NMI was caused by an oscillator fault.

A PUC signal can trigger an oscillator fault, because the PUC switches the LFXT1 to LF mode, therefore switching off the HF mode. The PUC signal also switches off the XT2 oscillator.



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## 2.2.1.4 Example of an NMI Interrupt Handler

The NMI interrupt is a multiple-source interrupt. An NMI interrupt automatically resets the NMIIE, OFIE and ACCVIE interrupt-enable bits. The user NMI service routine resets the interrupt flags and re-enables the interrupt-enable bits according to the application needs as shown in Figure 2-5.

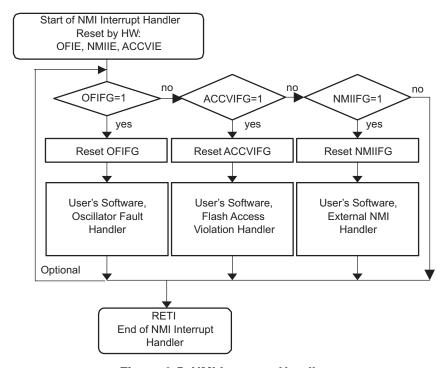


Figure 2-5. NMI Interrupt Handler

## NOTE: Enabling NMI Interrupts with ACCVIE, NMIIE, and OFIE

To prevent nested NMI interrupts, the ACCVIE, NMIIE, and OFIE enable bits should not be set inside of an NMI interrupt service routine.

### 2.2.2 Maskable Interrupts

Maskable interrupts are caused by peripherals with interrupt capability including the watchdog timer overflow in interval-timer mode. Each maskable interrupt source can be disabled individually by an interrupt enable bit, or all maskable interrupts can be disabled by the general interrupt enable (GIE) bit in the status register (SR).

Each individual peripheral interrupt is discussed in the associated peripheral module chapter in this manual.



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## 2.2.3 Interrupt Processing

When an interrupt is requested from a peripheral and the peripheral interrupt enable bit and GIE bit are set, the interrupt service routine is requested. Only the individual enable bit must be set for (non)-maskable interrupts to be requested.

#### 2.2.3.1 Interrupt Acceptance

The interrupt latency is 5 cycles (CPUx) or 6 cycles (CPU), starting with the acceptance of an interrupt request and lasting until the start of execution of the first instruction of the interrupt-service routine, as shown in Figure 2-6. The interrupt logic executes the following:

- 1. Any currently executing instruction is completed.
- 2. The PC, which points to the next instruction, is pushed onto the stack.
- 3. The SR is pushed onto the stack.
- 4. The interrupt with the highest priority is selected if multiple interrupts occurred during the last instruction and are pending for service.
- 5. The interrupt request flag resets automatically on single-source flags. Multiple source flags remain set for servicing by software.
- 6. The SR is cleared. This terminates any low-power mode. Because the GIE bit is cleared, further interrupts are disabled.
- 7. The content of the interrupt vector is loaded into the PC: the program continues with the interrupt service routine at that address.

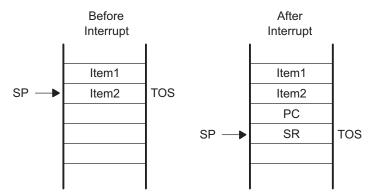


Figure 2-6. Interrupt Processing



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### 2.2.3.2 Return From Interrupt

The interrupt handling routine terminates with the instruction:

RETI (return from an interrupt service routine)

The return from the interrupt takes 5 cycles (CPU) or 3 cycles (CPUx) to execute the following actions and is illustrated in Figure 2-7.

- 1. The SR with all previous settings pops from the stack. All previous settings of GIE, CPUOFF, etc. are now in effect, regardless of the settings used during the interrupt service routine.
- 2. The PC pops from the stack and begins execution at the point where it was interrupted.

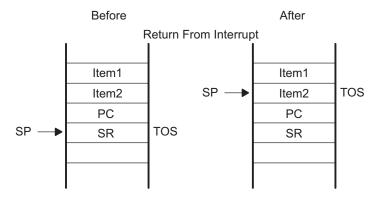


Figure 2-7. Return From Interrupt

### 2.2.3.3 Interrupt Nesting

Interrupt nesting is enabled if the GIE bit is set inside an interrupt service routine. When interrupt nesting is enabled, any interrupt occurring during an interrupt service routine will interrupt the routine, regardless of the interrupt priorities.



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## 2.2.4 Interrupt Vectors

The interrupt vectors and the power-up starting address are located in the address range 0FFFFh to 0FFC0h, as described in Table 2-1. A vector is programmed by the user with the 16-bit address of the corresponding interrupt service routine. See the device-specific data sheet for the complete interrupt vector list.

It is recommended to provide an interrupt service routine for each interrupt vector that is assigned to a module. A dummy interrupt service routine can consist of just the RETI instruction and several interrupt vectors can point to it.

Unassigned interrupt vectors can be used for regular program code if necessary.

Some module enable bits, interrupt enable bits, and interrupt flags are located in the SFRs. The SFRs are located in the lower address range and are implemented in byte format. SFRs must be accessed using byte instructions. See the device-specific data sheet for the SFR configuration.

Table 2-1. Interrupt Sources, Flags, and Vectors

Interrupt Source	Interrupt Flag	System Interrupt	Word Address	Priority
Power-up, external reset, watchdog, flash password, illegal instruction fetch	PORIFG RSTIFG WDTIFG KEYV	Reset	0FFFEh	31, highest
NMI, oscillator fault, flash memory access violation	NMIIFG OFIFG ACCVIFG	(non)-maskable (non)-maskable (non)-maskable	0FFFCh	30
device-specific			0FFFAh	29
device-specific			0FFF8h	28
device-specific			0FFF6h	27
Watchdog timer	WDTIFG	maskable	0FFF4h	26
device-specific			0FFF2h	25
device-specific			0FFF0h	24
device-specific			0FFEEh	23
device-specific			0FFECh	22
device-specific			0FFEAh	21
device-specific			0FFE8h	20
device-specific			0FFE6h	19
device-specific			0FFE4h	18
device-specific			0FFE2h	17
device-specific			0FFE0h	16
device-specific			0FFDEh	15
device-specific			0FFDCh	14
device-specific			0FFDAh	13
device-specific			0FFD8h	12
device-specific			0FFD6h	11
device-specific			0FFD4h	10
device-specific			0FFD2h	9
device-specific			0FFD0h	8
device-specific			0FFCEh	7
device-specific			0FFCCh	6
device-specific			0FFCAh	5
device-specific			0FFC8h	4
device-specific			0FFC6h	3
device-specific			0FFC4h	2
device-specific			0FFC2h	1
device-specific			0FFC0h	0, lowest



Operating Modes www.ti.com

# 2.3 Operating Modes

The MSP430 family is designed for ultralow-power applications and uses different operating modes shown in Figure 2-9.

The operating modes take into account three different needs:

- Ultralow-power
- Speed and data throughput
- Minimization of individual peripheral current consumption

The MSP430 typical current consumption is shown in Figure 2-8.

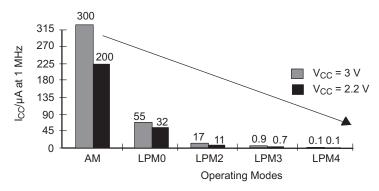


Figure 2-8. Typical Current Consumption of 'F21x1 Devices vs Operating Modes

The low-power modes 0 to 4 are configured with the CPUOFF, OSCOFF, SCG0, and SCG1 bits in the status register The advantage of including the CPUOFF, OSCOFF, SCG0, and SCG1 mode-control bits in the status register is that the present operating mode is saved onto the stack during an interrupt service routine. Program flow returns to the previous operating mode if the saved SR value is not altered during the interrupt service routine. Program flow can be returned to a different operating mode by manipulating the saved SR value on the stack inside of the interrupt service routine. The mode-control bits and the stack can be accessed with any instruction.

When setting any of the mode-control bits, the selected operating mode takes effect immediately (see Figure 2-9). Peripherals operating with any disabled clock are disabled until the clock becomes active. The peripherals may also be disabled with their individual control register settings. All I/O port pins and RAM/registers are unchanged. Wake up is possible through all enabled interrupts.



www.ti.com Operating Modes

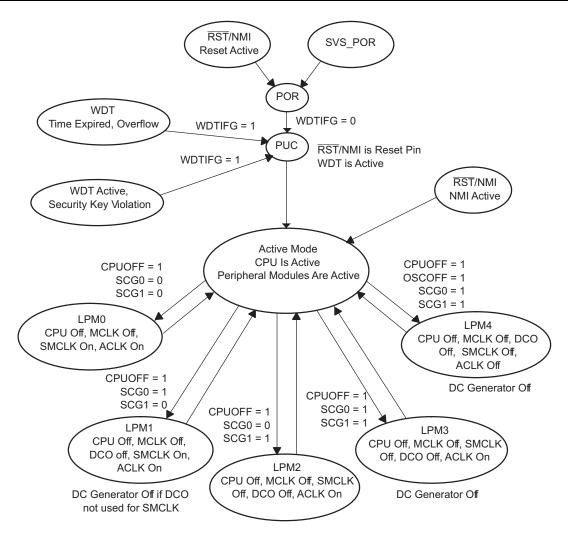


Figure 2-9. Operating Modes For Basic Clock System

Table 2-2. Operating Modes For Basic Clock System

SCG1	SCG0	OSCOFF	CPUOFF	Mode	CPU and Clocks Status
0	0	0	0	Active	CPU is active, all enabled clocks are active
0	0	0	1	LPM0	CPU, MCLK are disabled, SMCLK, ACLK are active
0	1	0	1	LPM1	CPU, MCLK are disabled. DCO and DC generator are disabled if the DCO is not used for SMCLK. ACLK is active.
1	0	0	1	LPM2	CPU, MCLK, SMCLK, DCO are disabled. DC generator remains enabled. ACLK is active.
1	1	0	1	LPM3	CPU, MCLK, SMCLK, DCO are disabled. DC generator disabled. ACLK is active.
1	1	1	1	LPM4	CPU and all clocks disabled



# 2.3.1 Entering and Exiting Low-Power Modes

An enabled interrupt event wakes the MSP430 from any of the low-power operating modes. The program flow is:

- Enter interrupt service routine:
  - The PC and SR are stored on the stack
  - The CPUOFF, SCG1, and OSCOFF bits are automatically reset
- Options for returning from the interrupt service routine:
  - The original SR is popped from the stack, restoring the previous operating mode.
  - The SR bits stored on the stack can be modified within the interrupt service routine returning to a different operating mode when the RETI instruction is executed.

```
Enter LPM0 Example
 BIS
      #GIE+CPUOFF,SR
                                  ; Enter LPM0
                                  ; Program stops here
; Exit LPMO Interrupt Service Routine
 BIC
        #CPUOFF,0(SP)
                          ; Exit LPM0 on RETI
 RETI
; Enter LPM3 Example
        #GIE+CPUOFF+SCG1+SCG0,SR ; Enter LPM3
                                  ; Program stops here
; Exit LPM3 Interrupt Service Routine
        #CPUOFF+SCG1+SCG0,0(SP) ; Exit LPM3 on RETI
 RETI
```

# 2.4 Principles for Low-Power Applications

Often, the most important factor for reducing power consumption is using the MSP430 clock system to maximize the time in LPM3. LPM3 power consumption is less than 2 µA typical with both a real-time clock function and all interrupts active. A 32-kHz watch crystal is used for the ACLK and the CPU is clocked from the DCO (normally off) which has a 1-µs wake-up.

- Use interrupts to wake the processor and control program flow.
- · Peripherals should be switched on only when needed.
- Use low-power integrated peripheral modules in place of software driven functions. For example Timer\_A and Timer\_B can automatically generate PWM and capture external timing, with no CPU resources.
- Calculated branching and fast table look-ups should be used in place of flag polling and long software calculations.
- Avoid frequent subroutine and function calls due to overhead.
- For longer software routines, single-cycle CPU registers should be used.



# 2.5 Connection of Unused Pins

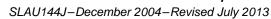
The correct termination of all unused pins is listed in Table 2-3.

**Table 2-3. Connection of Unused Pins** 

Pin	Potential	Comment
AV <sub>CC</sub>	DVCC	
AV <sub>SS</sub>	DVSS	
V <sub>REF+</sub>	Open	
V <sub>eREF+</sub>	DVSS	
V <sub>REF-</sub> /V <sub>eREF-</sub>	DVSS	
XIN	DVCC	For dedicated XIN pins only. XIN pins with shared GPIO functions should be programmed to GPIO and follow Px.0 to Px.7 recomendations.
XOUT	Open	For dedicated XOUT pins only. XOUT pins with shared GPIO functions should be programmed to GPIO and follow Px.0 to Px.7 recomendations.
XT2IN	DVSS	For dedicated X2IN pins only. X2IN pins with shared GPIO functions should be programmed to GPIO and follow Px.0 to Px.7 recomendations.
XT2OUT	Open	For dedicated X2OUT pins only. X2OUT pins with shared GPIO functions should be programmed to GPIO and follow Px.0 to Px.7 recomendations.
Px.0 to Px.7	Open	Switched to port function, output direction or input with pullup/pulldown enabled
RST/NMI	DVCC or VCC	47 kΩ pullup with 10 nF (2.2 nF <sup>(1)</sup> ) pulldown
Test	Open	20xx, 21xx, 22xx devices
TDO	Open	
TDI	Open	
TMS	Open	
TCK	Open	

<sup>(1)</sup> The pulldown capacitor should not exceed 2.2 nF when using devices with Spy-Bi-Wire interface in Spy-Bi-Wire mode or in 4-wire JTAG mode with TI tools like FET interfaces or GANG programmers.







# **CPU**

This chapter describes the MSP430 CPU, addressing modes, and instruction set.

Topic		Page
2.1	CPU Introduction	12
3.2	CPU Registers	44
3.3	Addressing Modes	47
3.4	Instruction Set	<b>56</b>

42



www.ti.com CPU Introduction

#### 3.1 CPU Introduction

The CPU incorporates features specifically designed for modern programming techniques such as calculated branching, table processing, and the use of high-level languages such as C. The CPU can address the complete address range without paging.

The CPU features include:

- RISC architecture with 27 instructions and 7 addressing modes.
- Orthogonal architecture with every instruction usable with every addressing mode.
- Full register access including program counter, status registers, and stack pointer.
- Single-cycle register operations.
- Large 16-bit register file reduces fetches to memory.
- 16-bit address bus allows direct access and branching throughout entire memory range.
- 16-bit data bus allows direct manipulation of word-wide arguments.
- Constant generator provides six most used immediate values and reduces code size.
- · Direct memory-to-memory transfers without intermediate register holding.
- Word and byte addressing and instruction formats.

The block diagram of the CPU is shown in Figure 3-1.



CPU Registers www.ti.com

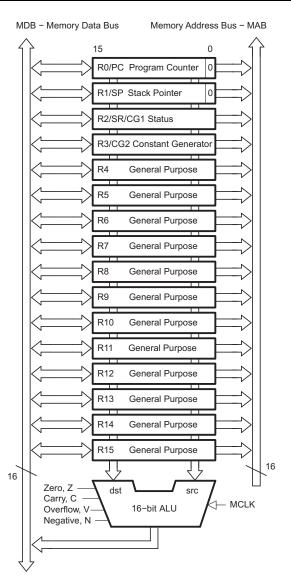


Figure 3-1. CPU Block Diagram

## 3.2 CPU Registers

The CPU incorporates sixteen 16-bit registers. R0, R1, R2, and R3 have dedicated functions. R4 to R15 are working registers for general use.

## 3.2.1 Program Counter (PC)

The 16-bit program counter (PC/R0) points to the next instruction to be executed. Each instruction uses an even number of bytes (two, four, or six), and the PC is incremented accordingly. Instruction accesses in the 64-KB address space are performed on word boundaries, and the PC is aligned to even addresses. Figure 3-2 shows the program counter.

#### Figure 3-2. Program Counter





**CPU Registers** www.ti.com

The PC can be addressed with all instructions and addressing modes. A few examples:

#LABEL,PC ; Branch to address LABEL LABEL, PC ; Branch to address contained in LABEL VOM @R14,PC ; Branch indirect to address in R14 VOM

# 3.2.2 Stack Pointer (SP)

The stack pointer (SP/R1) is used by the CPU to store the return addresses of subroutine calls and interrupts. It uses a predecrement, postincrement scheme. In addition, the SP can be used by software with all instructions and addressing modes. Figure 3-3 shows the SP. The SP is initialized into RAM by the user, and is aligned to even addresses.

Figure 3-4 shows stack usage.

## Figure 3-3. Stack Counter



MOV 2(SP),R6 ; Item I2 -> R6 VOM R7,0(SP) ; Overwrite TOS with R7 PUSH #0123h ; Put 0123h onto TOS ; R8 = 0123hPOP R8

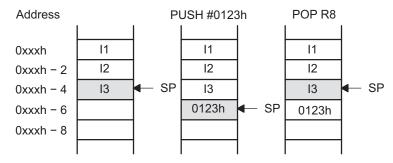


Figure 3-4. Stack Usage

The special cases of using the SP as an argument to the PUSH and POP instructions are described and shown in Figure 3-5.



a PUSH SP instruction.

The stack pointer is changed after The stack pointer is not changed after a POP SP instruction. The POP SP instruction places SP1 into the stack pointer SP (SP2=SP1)

Figure 3-5. PUSH SP - POP SP Sequence

## 3.2.3 Status Register (SR)

The status register (SR/R2), used as a source or destination register, can be used in the register mode only addressed with word instructions. The remaining combinations of addressing modes are used to support the constant generator. Figure 3-6 shows the SR bits.



CPU Registers www.ti.com

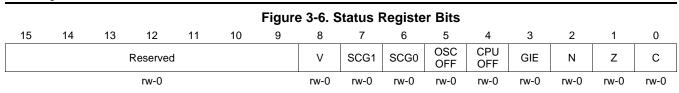


Table 3-1 describes the status register bits.

Table 3-1. Description of Status Register Bits

Bit		Description		
V	Overflow bit. This bit is set when the result of an arithmetic operation overflows the signed-variable range.			
	ADD(.B),ADDC(.B)	Set when:		
		Positive + Positive = Negative		
		Negative + Negative = Positive		
		Otherwise reset		
	SUB(.B),SUBC(.B),CMP(.B)	Set when:		
		Positive – Negative = Negative		
		Negative – Positive = Positive		
		Otherwise reset		
SCG1	System clock generator 1. When set, to	urns off the SMCLK.		
SCG0	System clock generator 0. When set, to	urns off the DCO dc generator, if DCOCLK is not used for MCLK or SMCLK.		
OSCOFF	Oscillator Off. When set, turns off the L	LFXT1 crystal oscillator, when LFXT1CLK is not use for MCLK or SMCLK.		
CPUOFF	CPU off. When set, turns off the CPU.			
GIE	General interrupt enable. When set, er	nables maskable interrupts. When reset, all maskable interrupts are disabled.		
N	Negative bit. Set when the result of a b	byte or word operation is negative and cleared when the result is not negative.		
	Word operation: N is set to the value	ue of bit 15 of the result.		
	Byte operation: N is set to the valu	e of bit 7 of the result.		
Z	Zero bit. Set when the result of a byte or word operation is 0 and cleared when the result is not 0.			
С	Carry bit. Set when the result of a byte	or word operation produced a carry and cleared when no carry occurred.		

# 3.2.4 Constant Generator Registers CG1 and CG2

Six commonly-used constants are generated with the constant generator registers R2 and R3, without requiring an additional 16-bit word of program code. The constants are selected with the source-register addressing modes (As), as described in Table 3-2.

Table 3-2. Values of Constant Generators CG1, CG2

Register	As	Constant	Remarks	
R2	00		Register mode	_
R2	01	(0)	Absolute address mode	
R2	10	00004h	+4, bit processing	
R2	11	00008h	+8, bit processing	
R3	00	00000h	0, word processing	
R3	01	00001h	+1	
R3	10	00002h	+2, bit processing	
R3	11	0FFFFh	-1, word processing	

The constant generator advantages are:

- No special instructions required
- No additional code word for the six constants
- No code memory access required to retrieve the constant



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The assembler uses the constant generator automatically if one of the six constants is used as an immediate source operand. Registers R2 and R3, used in the constant mode, cannot be addressed explicitly; they act as source-only registers.

## 3.2.4.1 Constant Generator - Expanded Instruction Set

The RISC instruction set of the MSP430 has only 27 instructions. However, the constant generator allows the MSP430 assembler to support 24 additional, emulated instructions. For example, the single-operand instruction

```
CLR dst
```

is emulated by the double-operand instruction with the same length:

```
MOV R3,dst
```

where the #0 is replaced by the assembler, and R3 is used with As=00.

INC dst

is replaced by:

ADD 0(R3),dst

## 3.2.5 General-Purpose Registers R4 to R15

The twelve registers, R4-R15, are general-purpose registers. All of these registers can be used as data registers, address pointers, or index values and can be accessed with byte or word instructions as shown in Figure 3-7.

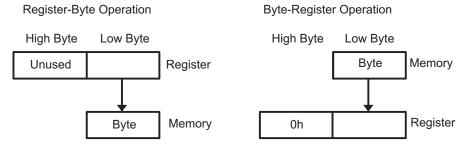
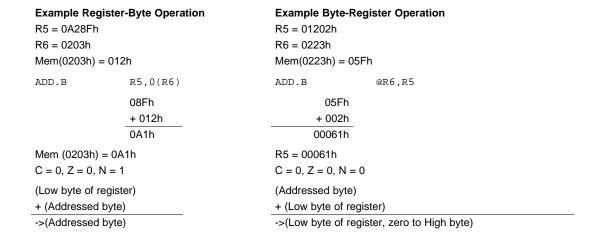


Figure 3-7. Register-Byte/Byte-Register Operations



## 3.3 Addressing Modes

Seven addressing modes for the source operand and four addressing modes for the destination operand can address the complete address space with no exceptions. The bit numbers in Table 3-3 describe the contents of the As (source) and Ad (destination) mode bits.



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**Table 3-3. Source/Destination Operand Addressing Modes** 

As/Ad	Addressing Mode	Syntax	Description
00/0	Register mode	Rn	Register contents are operand
01/1	Indexed mode	X(Rn)	(Rn + X) points to the operand. X is stored in the next word.
01/1	Symbolic mode	ADDR	(PC + X) points to the operand. X is stored in the next word. Indexed mode $X(PC)$ is used.
01/1	Absolute mode	&ADDR	The word following the instruction contains the absolute address. X is stored in the next word. Indexed mode X(SR) is used.
10/-	Indirect register mode	@Rn	Rn is used as a pointer to the operand.
11/-	Indirect autoincrement	@Rn+	Rn is used as a pointer to the operand. Rn is incremented afterwards by 1 for .B instructions and by 2 for .W instructions
11/-	Immediate mode	#N	The word following the instruction contains the immediate constant N. Indirect autoincrement mode @PC+ is used.

The seven addressing modes are explained in detail in the following sections. Most of the examples show the same addressing mode for the source and destination, but any valid combination of source and destination addressing modes is possible in an instruction.

# NOTE: Use of Labels EDE, TONI, TOM, and LEO

Throughout MSP430 documentation EDE, TONI, TOM, and LEO are used as generic labels. They are only labels. They have no special meaning.



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# 3.3.1 Register Mode

The register mode is described in Table 3-4.

**Table 3-4. Register Mode Description** 

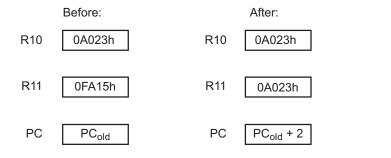
	Assembler Code	Content of ROM			
MOV	R10,R11	MOV R10,R11			

Length: One or two words

Operation: Move the content of R10 to R11. R10 is not affected.

Comment: Valid for source and destination

Example: MOV R10,R11



## NOTE: Data in Registers

The data in the register can be accessed using word or byte instructions. If byte instructions are used, the high byte is always 0 in the result. The status bits are handled according to the result of the byte instructions.



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## 3.3.2 Indexed Mode

The indexed mode is described in Table 3-5.

**Table 3-5. Indexed Mode Description** 

Assembler Code	Content of ROM
MOV 2(R5),6(R6)	MOV X(R5),Y(R6)
	X = 2
	Y = 6

Length: Two or three words

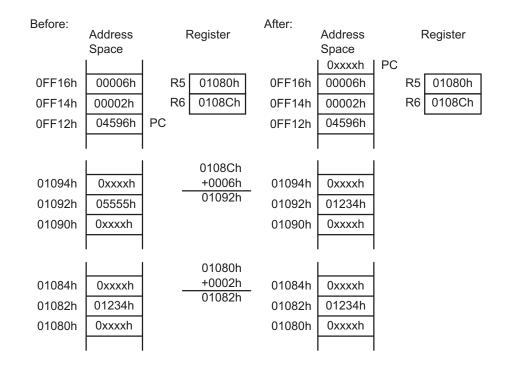
Operation: Move the contents of the source address (contents of R5 + 2) to the destination address (contents of R6 + 6). The

source and destination registers (R5 and R6) are not affected. In indexed mode, the program counter is

incremented automatically so that program execution continues with the next instruction.

Comment: Valid for source and destination

Example: MOV 2(R5), 6(R6);





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# 3.3.3 Symbolic Mode

The symbolic mode is described in Table 3-6.

**Table 3-6. Symbolic Mode Description** 

Assembler Code	Content of ROM
MOV EDE, TONI	MOV X(PC),Y(PC)
	X = EDE - PC
	Y = TONI - PC

Length: Two or three words

Operation: Move the contents of the source address EDE (contents of PC + X) to the destination address TONI (contents of

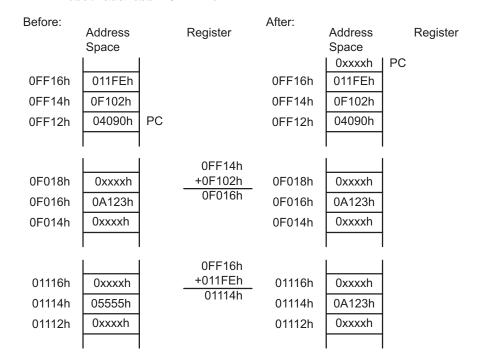
PC + Y). The words after the instruction contain the differences between the PC and the source or destination addresses. The assembler computes and inserts offsets X and Y automatically. With symbolic mode, the program

counter (PC) is incremented automatically so that program execution continues with the next instruction.

Comment: Valid for source and destination

Example:

MOV EDE, TONI ; Source address EDE = 0F016h ; Dest. address TONI = 01114h





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## 3.3.4 Absolute Mode

The absolute mode is described in Table 3-7.

**Table 3-7. Absolute Mode Description** 

Assembler Code	Content of ROM
MOV &EDE,&TONI	MOV X(0),Y(0)
	X = EDE
	Y = TONI

Length: Two or three words

Operation: Move the contents of the source address EDE to the destination address TONI. The words after the instruction

contain the absolute address of the source and destination addresses. With absolute mode, the PC is

incremented automatically so that program execution continues with the next instruction.

Comment: Valid for source and destination

Example:

MOV &EDE,&TONI ;Source address EDE = 0F016h ;Dest. address TONI = 01114h

Before:	Address Space		Register	After:	Address Space		Register
	I	ı			Opace	РС	
0FF16h	01114h			0FF16h	01114h	10	
0FF14h	0F016h			0FF14h	0F016h		
0FF12h	04292h	PC		0FF12h	04292h		
	1				1 1		
0F018h	0xxxxh			0F018h	0xxxxh		
0F016h	0A123h			0F016h	0A123h		
0F014h	0xxxxh			0F014h	0xxxxh		
	1				1 1		
01116h	0xxxxh			01116h	0xxxxh		
01114h	01234h			01114h	0A123h		
01112h	0xxxxh			01112h	0xxxxh		

This address mode is mainly for hardware peripheral modules that are located at an absolute, fixed address. These are addressed with absolute mode to ensure software transportability (for example, position-independent code).



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# 3.3.5 Indirect Register Mode

The indirect register mode is described in Table 3-8.

**Table 3-8. Indirect Mode Description** 

Assembler Code	Content of ROM
MOV @R10,0(R11)	MOV @R10,0(R11)

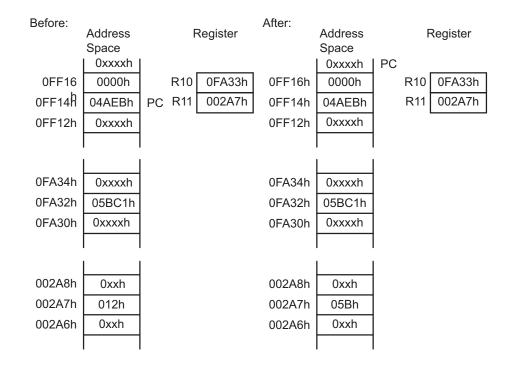
Length: One or two words

Operation: Move the contents of the source address (contents of R10) to the destination address (contents of R11). The

registers are not modified.

Comment: Valid only for source operand. The substitute for destination operand is 0(Rd).

Example: MOV.B @R10,0(R11)





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#### 3.3.6 Indirect Autoincrement Mode

The indirect autoincrement mode is described in Table 3-9.

**Table 3-9. Indirect Autoincrement Mode Description** 

Assembler Code	Content of ROM
MOV @R10+,0(R11)	MOV @R10+,0(R11)

Length: One or two words

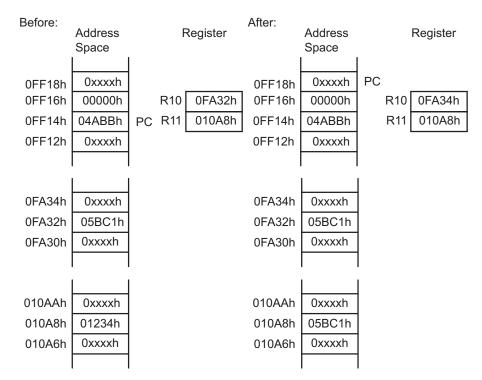
Operation: Move the contents of the source address (contents of R10) to the destination address (contents of R11). Register

R10 is incremented by 1 for a byte operation, or 2 for a word operation after the fetch; it points to the next address

without any overhead. This is useful for table processing.

Comment: Valid only for source operand. The substitute for destination operand is 0(Rd) plus second instruction INCD Rd.

Example: MOV @R10+,0(R11)



The auto-incrementing of the register contents occurs after the operand is fetched. This is shown in Figure 3-8.

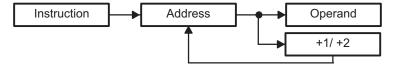


Figure 3-8. Operand Fetch Operation



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## 3.3.7 Immediate Mode

The immediate mode is described in Table 3-10.

**Table 3-10. Immediate Mode Description** 

Assembler Code	Content of ROM
MOV #45h,TONI	MOV @PC+,X(PC)
	45
	X = TONI - PC

Length: Two or three words

It is one word less if a constant of CG1 or CG2 can be used.

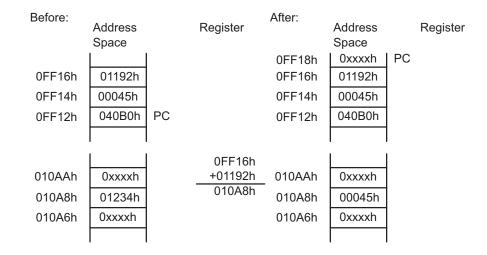
Operation: Move the immediate constant 45h, which is contained in the word following the instruction, to destination address

TONI. When fetching the source, the program counter points to the word following the instruction and moves the

contents to the destination.

Comment: Valid only for a source operand.

Example: MOV #45h, TONI





# 3.4 Instruction Set

The complete MSP430 instruction set consists of 27 core instructions and 24 emulated instructions. The core instructions are instructions that have unique op-codes decoded by the CPU. The emulated instructions are instructions that make code easier to write and read, but do not have op-codes themselves, instead they are replaced automatically by the assembler with an equivalent core instruction. There is no code or performance penalty for using emulated instruction.

There are three core-instruction formats:

- Dual-operand
- Single-operand
- Jump

All single-operand and dual-operand instructions can be byte or word instructions by using .B or .W extensions. Byte instructions are used to access byte data or byte peripherals. Word instructions are used to access word data or word peripherals. If no extension is used, the instruction is a word instruction.

The source and destination of an instruction are defined by the following fields:

src The source operand defined by As and S-reg
dst The destination operand defined by Ad and D-reg

As The addressing bits responsible for the addressing mode used for the source (src)

S-reg The working register used for the source (src)

Ad The addressing bits responsible for the addressing mode used for the destination (dst)

D-reg The working register used for the destination (dst)

B/W Byte or word operation:

0: word operation1: byte operation

#### NOTE: Destination Address

Destination addresses are valid anywhere in the memory map. However, when using an instruction that modifies the contents of the destination, the user must ensure the destination address is writable. For example, a masked-ROM location would be a valid destination address, but the contents are not modifiable, so the results of the instruction would be lost.



# 3.4.1 Double-Operand (Format I) Instructions

Figure 3-9 illustrates the double-operand instruction format.



Figure 3-9. Double Operand Instruction Format

Table 3-11 lists and describes the double operand instructions.

**Table 3-11. Double Operand Instructions** 

Maranania	S-Reg,	Onematica		Statu	s Bits	
Mnemonic	D-Reg	Operation	V	N	Z	С
MOV(.B)	src,dst	$src \rightarrow dst$	-	-	-	-
ADD(.B)	src,dst	$src + dst \rightarrow dst$	*	*	*	*
ADDC(.B)	src,dst	$src + dst + C \rightarrow dst$	*	*	*	*
SUB(.B)	src,dst	$dst + .not.src + 1 \rightarrow dst$	*	*	*	*
SUBC(.B)	src,dst	$dst + .not.src + C \to dst$	*	*	*	*
CMP(.B)	src,dst	dst - src	*	*	*	*
DADD(.B)	src,dst	$src + dst + C \rightarrow dst (decimally)$	*	*	*	*
BIT(.B)	src,dst	src .and. dst	0	*	*	*
BIC(.B)	src,dst	not.src .and. $dst \rightarrow dst$	-	-	-	-
BIS(.B)	src,dst	$\text{src .or. dst} \rightarrow \text{dst}$	-	-	-	-
XOR(.B)	src,dst	$src.xor.\ dst \rightarrow dst$	*	*	*	*
AND(.B)	src,dst	$src.and. dst \rightarrow dst$	0	*	*	*

<sup>\*</sup> The status bit is affected

#### NOTE: Instructions CMP and SUB

The instructions CMP and SUB are identical except for the storage of the result. The same is true for the BIT and AND instructions.

The status bit is not affected

<sup>0</sup> The status bit is cleared

<sup>1</sup> The status bit is set



# 3.4.2 Single-Operand (Format II) Instructions

Figure 3-10 illustrates the single-operand instruction format.

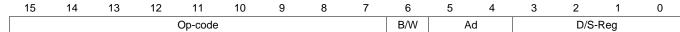


Figure 3-10. Single Operand Instruction Format

Table 3-12 lists and describes the single operand instructions.

**Table 3-12. Single Operand Instructions** 

Mnemonic	S-Reg,	Onematica		Status Bits			
Minemonic	D-Reg	D-Reg Operation	V	N	Z	С	
RRC(.B)	dst	$C \to MSB \to \!\! \dots \!\! \dots \!\! LSB \to C$	*	*	*	*	
RRA(.B)	dst	$MSB \to MSB \to LSB \to C$	0	*	*	*	
PUSH(.B)	src	$SP - 2 \rightarrow SP$ , $src \rightarrow @SP$	-	-	-	-	
SWPB	dst	Swap bytes	-	-	-	-	
CALL	dst	$SP - 2 \rightarrow SP, PC+2 \rightarrow @SP$	-	-	-	-	
		$dst \rightarrow PC$					
RETI		$TOS \rightarrow SR, SP + 2 \rightarrow SP$	*	*	*	*	
		$TOS \to PC, SP + 2 \to SP$					
SXT	dst	Bit $7 \rightarrow$ Bit 8Bit 15	0	*	*	*	

<sup>\*</sup> The status bit is affected

All addressing modes are possible for the CALL instruction. If the symbolic mode (ADDRESS), the immediate mode (#N), the absolute mode (&EDE) or the indexed mode x(RN) is used, the word that follows contains the address information.

The status bit is not affected

<sup>0</sup> The status bit is cleared

<sup>1</sup> The status bit is set



# 3.4.3 **Jumps**

Figure 3-11 shows the conditional-jump instruction format.



Figure 3-11. Jump Instruction Format

Table 3-13 lists and describes the jump instructions

**Table 3-13. Jump Instructions** 

Mnemonic	S-Reg, D-Reg	Operation
JEQ/JZ	Label	Jump to label if zero bit is set
JNE/JNZ	Label	Jump to label if zero bit is reset
JC	Label	Jump to label if carry bit is set
JNC	Label	Jump to label if carry bit is reset
JN	Label	Jump to label if negative bit is set
JGE	Label	Jump to label if $(N.XOR. V) = 0$
JL	Label	Jump to label if (N .XOR. V) = 1
JMP	Label	Jump to label unconditionally

Conditional jumps support program branching relative to the PC and do not affect the status bits. The possible jump range is from –511 to +512 words relative to the PC value at the jump instruction. The 10-bit program-counter offset is treated as a signed 10-bit value that is doubled and added to the program counter:

$$PC_{new} = PC_{old} + 2 + PC_{offset} \times 2$$



# 3.4.4 Instruction Cycles and Lengths

The number of CPU clock cycles required for an instruction depends on the instruction format and the addressing modes used - not the instruction itself. The number of clock cycles refers to the MCLK.

### 3.4.4.1 Interrupt and Reset Cycles

Table 3-14 lists the CPU cycles for interrupt overhead and reset.

Table 3-14. Interrupt and Reset Cycles

Action	No. of Cycles	Length of Instruction
Return from interrupt (RETI)	5	1
Interrupt accepted	6	-
WDT reset	4	-
Reset (RST/NMI)	4	-

# 3.4.4.2 Format-II (Single Operand) Instruction Cycles and Lengths

Table 3-15 lists the length and CPU cycles for all addressing modes of format-II instructions.

Table 3-15. Format-II Instruction Cycles and Lengths

	N	No. of Cycles			
Addressing Mode	RRA, RRC SWPB, SXT	PUSH	PUSH CALL		Example
Rn	1	3	4	1	SWPB R5
@Rn	3	4	4	1	RRC @R9
@Rn+	3	5	5	1	SWPB @R10+
#N	(See note)	4	5	2	CALL #0F000h
X(Rn)	4	5	5	2	CALL 2(R7)
EDE	4	5	5	2	PUSH EDE
&EDE	4	5	5	2	SXT &EDE

## NOTE: Instruction Format II Immediate Mode

Do not use instruction RRA, RRC, SWPB, and SXT with the immediate mode in the destination field. Use of these in the immediate mode results in an unpredictable program operation.

# 3.4.4.3 Format-III (Jump) Instruction Cycles and Lengths

All jump instructions require one code word, and take two CPU cycles to execute, regardless of whether the jump is taken or not.



# 3.4.4.4 Format-I (Double Operand) Instruction Cycles and Lengths

Table 3-16 lists the length and CPU cycles for all addressing modes of format-I instructions.

Table 3-16. Format 1 Instruction Cycles and Lengths

Addressing Mode			Length of		
Src	Dst	No. of Cycles	Instruction		Example
Rn	Rm	1	1	VOM	R5,R8
	PC	2	1	BR	R9
	x(Rm)	4	2	ADD	R5,4(R6)
	EDE	4	2	XOR	R8,EDE
	&EDE	4	2	MOV	R5,&EDE
@Rn	Rm	2	1	AND	@R4,R5
	PC	2	1	BR	@R8
	x(Rm)	5	2	XOR	@R5,8(R6)
	EDE	5	2	MOV	@R5,EDE
	&EDE	5	2	XOR	@R5,&EDE
@Rn+	Rm	2	1	ADD	@R5+,R6
	PC	3	1	BR	@R9+
	x(Rm)	5	2	XOR	@R5,8(R6)
	EDE	5	2	MOV	@R9+,EDE
	&EDE	5	2	MOV	@R9+,&EDE
#N	Rm	2	2	MOV	#20,R9
	PC	3	2	BR	#2AEh
	x(Rm)	5	3	MOV	#0300h,0(SP)
	EDE	5	3	ADD	#33,EDE
	&EDE	5	3	ADD	#33,&EDE
x(Rn)	Rm	3	2	MOV	2(R5),R7
	PC	3	2	BR	2(R6)
	TONI	6	3	MOV	4(R7),TONI
	x(Rm)	6	3	ADD	4(R4),6(R9)
	&TONI	6	3	MOV	2(R4),&TONI
EDE	Rm	3	2	AND	EDE,R6
	PC	3	2	BR	EDE
	TONI	6	3	CMP	EDE, TONI
	x(Rm)	6	3	MOV	EDE,0(SP)
	&TONI	6	3	MOV	EDE,&TONI
&EDE	Rm	3	2	MOV	&EDE,R8
	PC	3	2	BRA	&EDE
	TONI	6	3	MOV	&EDE,TONI
	x(Rm)	6	3	MOV	&EDE,0(SP)
	&TONI	6	3	MOV	&EDE,&TONI

61



# 3.4.5 Instruction Set Description

The instruction map is shown in Figure 3-12 and the complete instruction set is summarized in Table 3-17.

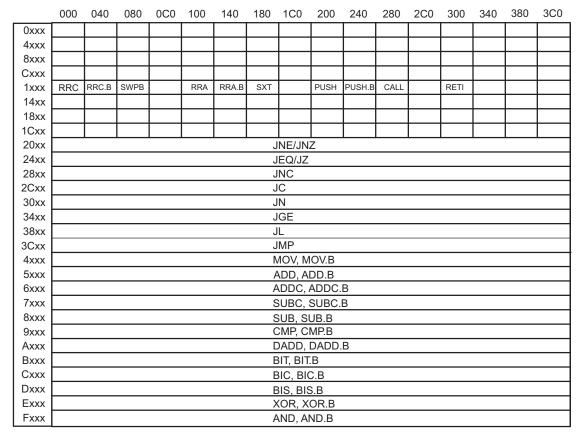


Figure 3-12. Core Instruction Map

Table 3-17. MSP430 Instruction Set

Mner	nonic	Descri	Description V N Z		Z	С	
ADC(.B) (1)	dst	Add C to destination	$dst + C \rightarrow dst$	*	*	*	*
ADD(.B)	src,dst	Add source to destination	$src + dst \rightarrow dst$	*	*	*	*
ADDC(.B)	src,dst	Add source and C to destination	$src + dst + C \rightarrow dst$	*	*	*	*
AND(.B)	src,dst	AND source and destination	$src.and.dst \rightarrow dst$	0	*	*	*
BIC(.B)	src,dst	Clear bits in destination	not.src .and. $dst \rightarrow dst$	-	-	-	-
BIS(.B)	src,dst	Set bits in destination	$\text{src .or. dst} \rightarrow \text{dst}$	-	-	-	-
BIT(.B)	src,dst	Test bits in destination	src .and. dst	0	*	*	*
BR <sup>(1)</sup>	dst	Branch to destination	$dst \to PC$	-	-	-	-
CALL	dst	Call destination	$\text{PC+2} \rightarrow \text{stack, dst} \rightarrow \text{PC}$	-	-	-	-
CLR(.B) (1)	dst	Clear destination	$0 \rightarrow dst$	-	-	-	-
CLRC (1)		Clear C	$0 \rightarrow C$	-	-	-	0
CLRN <sup>(1)</sup>		Clear N	$0 \rightarrow N$	-	0	-	-
CLRZ (1)		Clear Z	$0 \rightarrow Z$	-	-	0	-
CMP(.B)	src,dst	Compare source and destination	dst - src	*	*	*	*
DADC(.B) (1)	dst	Add C decimally to destination	$dst + C \to dst \ (decimally)$	*	*	*	*
DADD(.B)	src,dst	Add source and C decimally to dst	$\text{src} + \text{dst} + \text{C} \rightarrow \text{dst (decimally)}$	*	*	*	*
DEC(.B) (1)	dst	Decrement destination	$dst - 1 \rightarrow dst$	*	*	*	*

<sup>(1)</sup> Emulated Instruction



# Table 3-17. MSP430 Instruction Set (continued)

Mnem	nonic	Descript	ion	٧	N	Z	С
DECD(.B) (1)	dst	Double-decrement destination	$dst - 2 \rightarrow dst$	*	*	*	*
DINT (1)		Disable interrupts	$0 \rightarrow GIE$	-	-	-	-
EINT (1)		Enable interrupts	$1 \rightarrow GIE$	-	-	-	-
INC(.B) (1)	dst	Increment destination	$dst +1 \rightarrow dst$	*	*	*	*
INCD(.B) (1)	dst	Double-increment destination	$dst+2 \rightarrow dst$	*	*	*	*
INV(.B) (1)	dst	Invert destination	$. not. dst \rightarrow dst$	*	*	*	*
JC/JHS	label	Jump if C set/Jump if higher or same		-	-	-	-
JEQ/JZ	label	Jump if equal/Jump if Z set		-	-	-	-
JGE	label	Jump if greater or equal		-	-	-	-
JL	label	Jump if less		-	-	-	-
JMP	label	Jump	$PC + 2 \times offset \rightarrow PC$	-	-	-	-
JN	label	Jump if N set		-	-	-	-
JNC/JLO	label	Jump if C not set/Jump if lower		-	-	-	-
JNE/JNZ	label	Jump if not equal/Jump if Z not set		-	-	-	-
MOV(.B)	src,dst	Move source to destination	$src \rightarrow dst$	-	-	-	-
NOP (2)		No operation		-	-	-	-
POP(.B) (2)	dst	Pop item from stack to destination	$@SP \to dst, SP+2 \to SP$	-	-	-	-
PUSH(.B)	src	Push source onto stack	$SP - 2 \rightarrow SP, src \rightarrow @SP$	-	-	-	-
RET <sup>(2)</sup>		Return from subroutine	$@SP \rightarrow PC, SP + 2 \rightarrow SP$	-	-	-	-
RETI		Return from interrupt		*	*	*	*
RLA(.B) (2)	dst	Rotate left arithmetically		*	*	*	*
RLC(.B) (2)	dst	Rotate left through C		*	*	*	*
RRA(.B)	dst	Rotate right arithmetically		0	*	*	*
RRC(.B)	dst	Rotate right through C		*	*	*	*
SBC(.B) (2)	dst	Subtract not(C) from destination	$dst + 0FFFFh + C \rightarrow dst$	*	*	*	*
SETC (2)		Set C	$1 \rightarrow C$	-	-	-	1
SETN <sup>(2)</sup>		Set N	$1 \rightarrow N$	-	1	-	-
SETZ (2)		Set Z	$1 \rightarrow Z$	-	-	1	-
SUB(.B)	src,dst	Subtract source from destination	$dst + .not.src + 1 \rightarrow dst$	*	*	*	*
SUBC(.B)	src,dst	Subtract source and not(C) from dst	$dst + .not.src + C \rightarrow dst$	*	*	*	*
SWPB	dst	Swap bytes		-	-	-	-
SXT	dst	Extend sign		0	*	*	*
TST(.B) (2)	dst	Test destination	dst + 0FFFFh + 1	0	*	*	1
XOR(.B)	src,dst	Exclusive OR source and destination	$src.xor.\ dst \rightarrow dst$	*	*	*	*

<sup>(2)</sup> Emulated Instruction



#### 3.4.6 Instruction Set Details

#### 3.4.6.1 ADC

\*ADC[.W] Add carry to destination

\*ADC.B Add carry to destination

Syntax ADC dst or ADC.W dst

ADC.B dst

**Operation**  $dst + C \rightarrow dst$ 

Emulation ADDC #0,dst

ADDC.B #0,dst

**Description** The carry bit (C) is added to the destination operand. The previous contents of the

destination are lost.

Status Bit N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if dst was incremented from 0FFFFh to 0000, reset otherwise

Set if dst was incremented from 0FFh to 00, reset otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 16-bit counter pointed to by R13 is added to a 32-bit counter pointed to by R12.

ADD @R13,0(R12) ; Add LSDs

ADC 2(R12) ; Add carry to MSD

**Example** The 8-bit counter pointed to by R13 is added to a 16-bit counter pointed to by R12.

ADD.B @R13,0(R12) ; Add LSDs

ADC.B 1(R12) ; Add carry to MSD

64



#### 3.4.6.2 ADD

ADD[.W] Add source to destination

ADD.B Add source to destination

Syntax ADD src,dst or ADD.W src,dst

ADD.B src,dst

**Operation**  $\operatorname{src} + \operatorname{dst} \to \operatorname{dst}$ 

**Description** The source operand is added to the destination operand. The source operand is not

affected. The previous contents of the destination are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the result, cleared if not

V:Set if an arithmetic overflow occurs, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** R5 is increased by 10. The jump to TONI is performed on a carry.

ADD #10,R5

JC TONI ; Carry occurred

.....; No carry

**Example** R5 is increased by 10. The jump to TONI is performed on a carry.

ADD.B #10,R5 ; Add 10 to Lowbyte of R5

JC TONI ; Carry occurred, if (R5) ≥ 246 [0Ah+0F6h]

.....; No carry

65



#### 3.4.6.3 ADDC

ADDC[.W] Add source and carry to destination

ADDC.B Add source and carry to destination

Syntax 1 4 1 ADDC src,dst or ADDC.W src,dst

ADDC.B src,dst

Operation  $src + dst + C \rightarrow dst$ 

The source operand and the carry bit (C) are added to the destination operand. The Description

source operand is not affected. The previous contents of the destination are lost.

**Status Bits** N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

Example The 32-bit counter pointed to by R13 is added to a 32-bit counter, eleven words

(20/2 + 2/2) above the pointer in R13.

ADD @R13+,20(R13) ; ADD LSDs with no carry in ADDC @R13+,20(R13) ; ADD MSDs with carry ; resulting from the LSDs

**Example** The 24-bit counter pointed to by R13 is added to a 24-bit counter, eleven words above

the pointer in R13.

. . .

ADD.B @R13+,10(R13) ; ADD LSDs with no carry in ADDC.B @R13+,10(R13) ; ADD medium Bits with carry @R13+,10(R13) ADDC.B ; ADD MSDs with carry ; resulting from the LSDs



#### 3.4.6.4 AND

AND[.W] Source AND destination

AND.B Source AND destination

Syntax AND src,dst or AND.W src,dst

AND.B src,dst

Operation src .AND. dst → dst

**Description** The source operand and the destination operand are logically ANDed. The result is

placed into the destination.

Status Bits N: Set if result MSB is set, reset if not set

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise ( = .NOT. Zero)

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The bits set in R5 are used as a mask (#0AA55h) for the word addressed by TOM. If the

result is zero, a branch is taken to label TONI.

MOV #0AA55h,R5 ; Load mask into register R5

AND R5,TOM ; mask word addressed by TOM with R5

JZ TONI ;

.....; Result is not zero

; ; ; or ;

AND #0AA55h,TOM

JZ TONI

**Example** The bits of mask #0A5h are logically ANDed with the low byte TOM. If the result is zero, a branch is taken to label TONI.

AND.B #0A5h,TOM; mask Lowbyte TOM with 0A5h

JZ TONI

.....; Result is not zero



#### 3.4.6.5 BIC

BIC[.W] Clear bits in destination

BIC.B Clear bits in destination

Syntax BIC src,dst or BIC.W src,dst

BIC.B src,dst

**Operation** .NOT.src .AND.  $dst \rightarrow dst$ 

**Description** The inverted source operand and the destination operand are logically ANDed. The

result is placed into the destination. The source operand is not affected.

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The six MSBs of the RAM word LEO are cleared.

BIC #0FC00h, LEO; Clear 6 MSBs in MEM(LEO)

**Example** The five MSBs of the RAM byte LEO are cleared.

BIC.B #0F8h,LEO ; Clear 5 MSBs in Ram location LEO

68



#### 3.4.6.6 BIS

BIS[.W] Set bits in destination

BIS.B Set bits in destination

Syntax BIS src,dst or BIS.W src,dst

BIS.B src,dst

**Operation**  $\operatorname{src.OR.} \operatorname{dst} \to \operatorname{dst}$ 

**Description** The source operand and the destination operand are logically ORed. The result is placed

into the destination. The source operand is not affected.

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The six LSBs of the RAM word TOM are set.

BIS #003Fh, TOM ; set the six LSBs in RAM location TOM

**Example** The three MSBs of RAM byte TOM are set.

BIS.B #0E0h,TOM ; set the 3 MSBs in RAM location TOM

69



#### 3.4.6.7 BIT

BIT[.W] Test bits in destination

BIT.B Test bits in destination

Syntax BIT src,dst or BIT.W src,dst

Operation src .AND. dst

**Description** The source and destination operands are logically ANDed. The result affects only the

status bits. The source and destination operands are not affected.

Status Bits N: Set if MSB of result is set, reset otherwise

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise (.NOT. Zero)

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** If bit 9 of R8 is set, a branch is taken to label TOM.

BIT #0200h,R8 ; bit 9 of R8 set?
JNZ TOM ; Yes, branch to TOM
... ; No, proceed

**Example** If bit 3 of R8 is set, a branch is taken to label TOM.

BIT.B #8,R8 JC TOM

**Example** 

A serial communication receive bit (RCV) is tested. Because the carry bit is equal to the state of the tested bit while using the BIT instruction to test a single bit, the carry bit is used by the subsequent instruction; the read information is shifted into register RECBUF.

```
; Serial communication with LSB is shifted first:
                     ; xxxx xxxx xxxx
BIT.B
        #RCV,RCCTL
                     ; Bit info into carry
RRC
        RECBUF
                     ; Carry -> MSB of RECBUF
                     ; cxxx
                               XXXX
                     ; repeat previous two instructions
. . . . . .
                     ; 8 times
                     ; cccc cccc
                     ; MSB
                                 LSB
; Serial communication with MSB shifted first:
BIT.B #RCV,RCCTL ; Bit info into carry
RLC.B RECBUF
                     ; Carry -> LSB of RECBUF
                     ; xxxx
                              XXXC
                     ; repeat previous two instructions
. . . . . .
                     ; 8 times
. . . . . .
                     ; cccc cccc
                     ;
                     ; MSB
                                 LSB
```



## 3.4.6.8 BR, BRANCH

**Emulation** 

\*BR, BRANCH Branch to ...... destination

MOV dst,PC

**Description** An unconditional branch is taken to an address anywhere in the 64K address space. All

source addressing modes can be used. The branch instruction is a word instruction.

Status Bits Status bits are not affected.

**Example** Examples for all addressing modes are given.

BR #EXEC ; Branch to label EXEC or direct branch (e.g. #0A4h)

; Core instruction MOV @PC+,PC

BR EXEC ; Branch to the address contained in EXEC

; Core instruction MOV X(PC),PC

; Indirect address

BR &EXEC ; Branch to the address contained in absolute

; address EXEC

; Core instruction MOV X(0),PC

; Indirect address

BR R5 ; Branch to the address contained in R5

; Core instruction MOV R5,PC

; Indirect R5

BR @R5 ; Branch to the address contained in the word

; pointed to by R5.

; Core instruction MOV @R5+,PC

; Indirect, indirect R5

BR @R5+ ; Branch to the address contained in the word pointed

; to by R5 and increment pointer in R5 afterwards. ; The next time--S/W flow uses R5 pointer--it can

; alter program execution due to access to ; next address in a table pointed to by R5

; Core instruction MOV @R5,PC

; Indirect, indirect R5 with autoincrement

BR X(R5) ; Branch to the address contained in the address

; pointed to by R5 + X (e.g. table with address; starting at X). X can be an address or a label

; Core instruction MOV X(R5),PC

; Indirect, indirect R5 + X

71



#### 3.4.6.9 CALL

CALL Subroutine

Syntax

CALL dst

Operation

dst → tmp dst is evaluated and stored

 $SP - 2 \rightarrow SP$ 

PC → @SP PC updated to TOS

 $tmp \rightarrow PC$  dst saved to PC

Description

A subroutine call is made to an address anywhere in the 64K address space. All addressing modes can be used. The return address (the address of the following instruction) is stored on the stack. The call instruction is a word instruction.

**Status Bits** 

Status bits are not affected.

Example

Examples for all addressing modes are given.

```
CALL
      #EXEC
             ; Call on label EXEC or immediate address (e.g. #0A4h)
              ; SP-2 -> SP, PC+2 -> @SP, @PC+ -> PC
CALL
      EXEC
              ; Call on the address contained in EXEC
             ; SP-2 \rightarrow SP, PC+2 \rightarrow SP, X(PC) \rightarrow PC
              ; Indirect address
CALL
      &EXEC
             ; Call on the address contained in absolute address
             ; EXEC
             ; SP-2 \rightarrow SP, PC+2 \rightarrow @SP, X(0) \rightarrow PC
             ; Indirect address
CALL R5
             ; Call on the address contained in R5
              ; SP-2 -> SP, PC+2 -> @SP, R5 -> PC
              ; Indirect R5
             ; Call on the address contained in the word
CALL @R5
              ; pointed to by R5
              ; SP-2 -> SP, PC+2 -> @SP, @R5 -> PC
              ; Indirect, indirect R5
CALL @R5+
             ; Call on the address contained in the word
              ; pointed to by R5 and increment pointer in R5.
             ; The next time S/W flow uses R5 pointer
             ; it can alter the program execution due to
             ; access to next address in a table pointed to by R5
             ; SP-2 -> SP, PC+2 -> @SP, @R5 -> PC
             ; Indirect, indirect R5 with autoincrement
CALL X(R5) ; Call on the address contained in the address pointed
             ; to by R5 + X (e.g. table with address starting at X)
              ; X can be an address or a label
              ; SP-2 -> SP, PC+2 -> @SP, X(R5) -> PC
```

; Indirect, indirect R5 + X



# 3.4.6.10 CLR

\*CLR[.W] Clear destination

\*CLR.B Clear destination

Syntax CLR dst or CLR.W dst

CLR.B dst

**Operation**  $0 \rightarrow dst$ 

**Emulation** MOV #0,dst

MOV.B #0,dst

**Description** The destination operand is cleared.

Status Bits Status bits are not affected.

**Example** RAM word TONI is cleared.

CLR TONI ; 0 -> TONI

**Example** Register R5 is cleared.

CLR R5

**Example** RAM byte TONI is cleared.

CLR.B TONI ; 0 -> TONI



# 3.4.6.11 CLRC

\*CLRC Clear carry bit

Syntax CLRC

Emulation BIC #1,SR

**Description** The carry bit (C) is cleared. The clear carry instruction is a word instruction.

Status Bits N: Not affected

Z: Not affected

C: Cleared

V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 16-bit decimal counter pointed to by R13 is added to a 32-bit counter pointed to by

R12.

CLRC ; C=0: defines start

DADD @R13,0(R12) ; add 16=bit counter to low word of 32=bit counter

DADC 2(R12) ; add carry to high word of 32=bit counter



## 3.4.6.12 CLRN

\*CLRN Clear negative bit

Syntax CLRN

or

(.NOT.src .AND.  $dst \rightarrow dst$ )

Emulation BIC #4,SR

**Description** The constant 04h is inverted (0FFFBh) and is logically ANDed with the destination

operand. The result is placed into the destination. The clear negative bit instruction is a

word instruction.

Status Bits N: Reset to 0

Z: Not affectedC: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The Negative bit in the status register is cleared. This avoids special treatment with

negative numbers of the subroutine called.

```
CLRN
CALL SUBR
.....
SUBR JN SUBRET ; If input is negative: do nothing and return
.....
SUBR RET RET
```



# 3.4.6.13 CLRZ

\*CLRZ Clear zero bit

Syntax CLRZ

or

(.NOT.src .AND.  $dst \rightarrow dst$ )

Emulation BIC #2,SR

**Description** The constant 02h is inverted (0FFFDh) and logically ANDed with the destination

operand. The result is placed into the destination. The clear zero bit instruction is a word

instruction.

Status Bits N: Not affected

Z: Reset to 0C: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The zero bit in the status register is cleared.

CLRZ



### 3.4.6.14 CMP

**CMP[.W]** Compare source and destination

CMP.B Compare source and destination

Syntax CMP src,dst or CMP.W src,dst

CMP.B src,dst

Operation dst + .NOT.src + 1

or

(dst - src)

### Description

The source operand is subtracted from the destination operand. This is accomplished by adding the 1s complement of the source operand plus 1. The two operands are not affected and the result is not stored; only the status bits are affected.

#### **Status Bits**

N: Set if result is negative, reset if positive (src ≥ dst)

Z: Set if result is zero, reset otherwise (src = dst)

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

#### Mode Bits

OSCOFF, CPUOFF, and GIE are not affected.

### **Example**

R5 and R6 are compared. If they are equal, the program continues at the label EQUAL.

```
CMP R5,R6 ; R5 = R6?
JEQ EQUAL ; YES, JUMP
```

# Example

Two RAM blocks are compared. If they are not equal, the program branches to the label ERROR.

```
MOV
           #NUM,R5
                        ; number of words to be compared
    MOV
           #BLOCK1,R6
                        ; BLOCK1 start address in R6
    MOV
           #BLOCK2,R7
                        ; BLOCK2 start address in R7
L$1 CMP
           @R6+,0(R7)
                        ; Are Words equal? R6 increments
     JNZ
           ERROR
                        ; No, branch to ERROR
     INCD
          R7
                        ; Increment R7 pointer
    DEC
                        ; Are all words compared?
    JNZ
                        ; No, another compare
```

### **Example**

The RAM bytes addressed by EDE and TONI are compared. If they are equal, the program continues at the label EQUAL.

```
CMP.B EDE,TONI ; MEM(EDE) = MEM(TONI)?
JEQ EQUAL ; YES, JUMP
```



### 3.4.6.15 DADC

\*DADC[.W] Add carry decimally to destination

\*DADC.B Add carry decimally to destination

Syntax DADC dst or DADC.W src,dst

DADC.B dst

**Operation**  $dst + C \rightarrow dst (decimally)$ 

Emulation DADD #0,dst

DADD.B #0,dst

**Description** The carry bit (C) is added decimally to the destination.

Status Bits N: Set if MSB is 1

Z: Set if dst is 0, reset otherwise

C: Set if destination increments from 9999 to 0000, reset otherwise Set if destination increments from 99 to 00, reset otherwise

V: Undefined

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The four-digit decimal number contained in R5 is added to an eight-digit decimal number

pointed to by R8.

CLRC ; Reset carry

; next instruction's start condition is defined

DADD R5,0(R8); Add LSDs + C
DADC 2(R8); Add carry to MSD

**Example** The two-digit decimal number contained in R5 is added to a four-digit decimal number

pointed to by R8.

CLRC ; Reset carry

; next instruction's start condition is defined

DADD.B R5,0(R8) ; Add LSDs + C DADC.B 1(R8) ; Add carry to MSDs



### 3.4.6.16 DADD

**DADD[.W]** Source and carry added decimally to destination

**DADD.B** Source and carry added decimally to destination

Syntax DADD src,dst or DADD.W src,dst

DADD.B src,dst

**Operation**  $\operatorname{src} + \operatorname{dst} + \operatorname{C} \to \operatorname{dst} (\operatorname{decimally})$ 

**Description** The source operand and the destination operand are treated as four binary coded

decimals (BCD) with positive signs. The source operand and the carry bit (C) are added decimally to the destination operand. The source operand is not affected. The previous contents of the destination are lost. The result is not defined for non-BCD numbers.

Status Bits N: Set if the MSB is 1, reset otherwise

Z: Set if result is zero, reset otherwiseC: Set if the result is greater than 9999Set if the result is greater than 99

V: Undefined

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The eight-digit BCD number contained in R5 and R6 is added decimally to an eight-digit

BCD number contained in R3 and R4 (R6 and R4 contain the MSDs).

CLRC ; clear carry DADD R5,R3 ; add LSDs

DADD R6,R4 ; add MSDs with carry

JC OVERFLOW ; If carry occurs go to error handling routine

**Example** The two-digit decimal counter in the RAM byte CNT is incremented by one.

CLRC ; clear carry

DADD.B #1,CNT

or SETC

DADD.B #0,CNT ; equivalent to DADC.B CNT



### 3.4.6.17 DEC

\*DEC[.W] Decrement destination

\*DEC.B Decrement destination

Syntax DEC dst or DEC.W dst

DEC.B dst

**Operation**  $dst - 1 \rightarrow dst$ 

**Emulation** SUB #1,dst

SUB.B #1,dst

**Description** The destination operand is decremented by one. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 1, reset otherwise

C: Reset if dst contained 0, set otherwise

V: Set if an arithmetic overflow occurs, otherwise reset.

Set if initial value of destination was 08000h, otherwise reset.

Set if initial value of destination was 080h, otherwise reset.

### Mode Bits OSCOFF,

OSCOFF, CPUOFF, and GIE are not affected.

# Example

R10 is decremented by 1.

DEC R10 ; Decrement R10

; Move a block of 255 bytes from memory location starting with EDE to memory location starting with

; TONI. Tables should not overlap: start of destination address TONI must not be within the range  ${\tt EDE}$ 

; to EDE+0FEh

MOV #EDE,R6 MOV #255,R10 L\$1 MOV.B @R6+,TONI-EDE-1(R6)

DEC R10 JNZ L\$1

Do not transfer tables using the routine above with the overlap shown in Figure 3-13.

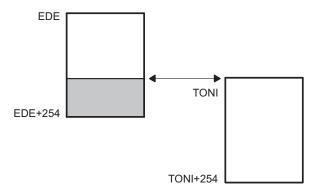


Figure 3-13. Decrement Overlap



### 3.4.6.18 DECD

\*DECD[.W] Double-decrement destination

\*DECD.B Double-decrement destination

Syntax DECD dst or DECD.W dst

DECD.B dst

**Operation**  $dst - 2 \rightarrow dst$ 

**Emulation** SUB #2,dst

Emulation SUB.B #2,dst

**Description** The destination operand is decremented by two. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 2, reset otherwise

C: Reset if dst contained 0 or 1, set otherwise

V: Set if an arithmetic overflow occurs, otherwise reset.

Set if initial value of destination was 08001 or 08000h, otherwise reset.

Set if initial value of destination was 081 or 080h, otherwise reset.

### Mode Bits

OSCOFF, CPUOFF, and GIE are not affected.

# Example

# R10 is decremented by 2.

```
DECD R10 ; Decrement R10 by two
```

; Move a block of 255 words from memory location starting with EDE to  $\,$ 

; memory location starting with TONI

; Tables should not overlap: start of destination address TONI must not be

; within the range EDE to EDE+0FEh  $\,$ 

MOV #EDE,R6 MOV #510,R10

L\$1 MOV @R6+,TONI-EDE-2(R6)

DECD R10 JNZ L\$1

# Example

## Memory at location LEO is decremented by two.

DECD.B LEO ; Decrement MEM(LEO)

Decrement status byte STATUS by two.

DECD.B STATUS



## 3.4.6.19 DINT

\*DINT Disable (general) interrupts

Syntax DINT

**Operation**  $0 \rightarrow GIE$ 

or

(0FFF7h .AND. SR  $\rightarrow$  SR / .NOT.src .AND. dst  $\rightarrow$  dst)

Emulation BIC #8,SR

**Description** All interrupts are disabled.

The constant 08h is inverted and logically ANDed with the status register (SR). The

result is placed into the SR.

Status Bits Status bits are not affected.

Mode Bits GIE is reset. OSCOFF and CPUOFF are not affected.

**Example** The general interrupt enable (GIE) bit in the status register is cleared to allow a

nondisrupted move of a 32-bit counter. This ensures that the counter is not modified

during the move by any interrupt.

DINT ; All interrupt events using the GIE bit are disabled

NOP

MOV COUNTHI, R5 ; Copy counter

MOV COUNTLO, R6

EINT ; All interrupt events using the GIE bit are enabled

### NOTE: Disable Interrupt

If any code sequence needs to be protected from interruption, the DINT should be executed at least one instruction before the beginning of the uninterruptible sequence, or should be followed by a NOP instruction.



## 3.4.6.20 EINT

\*EINT Enable (general) interrupts

Syntax EINT

**Operation**  $1 \rightarrow GIE$ 

or

(0008h .OR. SR  $\rightarrow$  SR / .src .OR. dst  $\rightarrow$  dst)

Emulation BIS #8,SR

**Description** All interrupts are enabled.

The constant #08h and the status register SR are logically ORed. The result is placed

into the SR.

Status Bits Status bits are not affected.

Mode Bits GIE is set. OSCOFF and CPUOFF are not affected.

**Example** The general interrupt enable (GIE) bit in the status register is set.

```
; Interrupt routine of ports P1.2 to P1.7 ; P1IN is the address of the register where all port bits are read. P1IFG is ; the address of the register where all interrupt events are latched.
```

```
PUSH.B
                  &P1TN
         BIC.B
                  @SP,&P1IFG
                               ; Reset only accepted flags
         EINT
                                ; Preset port 1 interrupt flags stored on stack
                                ; other interrupts are allowed
         BIT
                  #Mask,@SP
                  MaskOK
         JEO
                                ; Flags are present identically to mask: jump
         . . . . . .
MaskOK
         BIC
                  #Mask,@SP
         TNCD
                                ; Housekeeping: inverse to PUSH instruction
                  SP
                                ; at the start of interrupt subroutine. Corrects
                                ; the stack pointer.
         RETI
```

## NOTE: Enable Interrupt

The instruction following the enable interrupt instruction (EINT) is always executed, even if an interrupt service request is pending when the interrupts are enable.



## 3.4.6.21 INC

\*INC[.W] Increment destination

\*INC.B Increment destination

Syntax INC dst or INC.W dst

INC.B dst

**Operation**  $dst + 1 \rightarrow dst$ 

Emulation ADD #1,dst

**Description** The destination operand is incremented by one. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFh, reset otherwiseSet if dst contained 0FFh, reset otherwiseC: Set if dst contained 0FFFFh, reset otherwise

Set if dst contained 0FFh, reset otherwise V: Set if dst contained 07FFFh, reset otherwise

Set if dst contained 07Fh, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The status byte, STATUS, of a process is incremented. When it is equal to 11, a branch

to OVFL is taken.

INC.B STATUS
CMP.B #11,STATUS

JEQ OVFL



### 3.4.6.22 INCD

\*INCD[.W] Double-increment destination

\*INCD.B Double-increment destination

Syntax INCD dst or INCD.W dst

INCD.B dst

**Operation**  $dst + 2 \rightarrow dst$ 

Emulation ADD #2,dst

ADD.B #2,dst

**Example** The destination operand is incremented by two. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFEh, reset otherwise Set if dst contained 0FEh, reset otherwise

C: Set if dst contained 0FFFEh or 0FFFFh, reset otherwise

Set if dst contained 0FEh or 0FFh, reset otherwise

V: Set if dst contained 07FFEh or 07FFFh, reset otherwise

Set if dst contained 07Eh or 07Fh, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The item on the top of the stack (TOS) is removed without using a register.

PUSH R5 ; R5 is the result of a calculation, which is stored

; in the system stack

INCD SP ; Remove TOS by double-increment from stack

; Do not use INCD.B, SP is a word-aligned register

RET

**Example** The byte on the top of the stack is incremented by two.

INCD.B 0(SP); Byte on TOS is increment by two



### 3.4.6.23 INV

\*INV[.W] Invert destination

\*INV.B Invert destination

Syntax INV dst

INV.B dst

**Emulation** XOR #0FFFFh,dst

XOR.B #0FFh,dst

**Description** The destination operand is inverted. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFh, reset otherwise Set if dst contained 0FFh, reset otherwise

C: Set if result is not zero, reset otherwise ( = .NOT. Zero)

Set if result is not zero, reset otherwise ( = .NOT. Zero)

V: Set if initial destination operand was negative, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** Content of R5 is negated (twos complement).

MOV #00AEh,R5 ; R5 = 000AEh INV R5 ; Invert R5, R5 = 0FF51h INC R5 ; R5 is now negated, R5 = 0FF52h

**Example** Content of memory byte LEO is negated.

MOV.B #0AEh,LEO ; MEM(LEO) = 0AEh
INV.B LEO ; Invert LEO, MEM(LEO) = 051h
INC.B LEO ; MEM(LEO) is negated, MEM(LEO) = 052h



# 3.4.6.24 JC, JHS

JC Jump if carry set

**JHS** Jump if higher or same

Syntax JC label JHS label

**Operation** If C = 1: PC + 2 offset  $\rightarrow PC$ 

If C = 0: execute following instruction

**Description** The status register carry bit (C) is tested. If it is set, the 10-bit signed offset contained in

the instruction LSBs is added to the program counter. If C is reset, the next instruction following the jump is executed. JC (jump if carry/higher or same) is used for the

comparison of unsigned numbers (0 to 65536).

Status Bits Status bits are not affected.

**Example** The P1IN.1 signal is used to define or control the program flow.

BIT.B #02h,&P1IN ; State of signal -> Carry

JC PROGA ; If carry=1 then execute program routine A

.....; Carry=0, execute program here

**Example** R5 is compared to 15. If the content is higher or the same, branch to LABEL.

CMP #15,R5

JHS LABEL ; Jump is taken if R5 >= 15 ..... ; Continue here if R5 < 15



# 3.4.6.25 JEQ, JZ

JEQ, JZ Jump if equal, jump if zero

Syntax JEQ label

JZ label

**Operation** If Z = 1: PC + 2 offset  $\rightarrow$  PC

If Z = 0: execute following instruction

**Description** The status register zero bit (Z) is tested. If it is set, the 10-bit signed offset contained in

the instruction LSBs is added to the program counter. If Z is not set, the instruction

following the jump is executed.

Status Bits Status bits are not affected.

**Example** Jump to address TONI if R7 contains zero.

TST R7 ; Test R7

JZ TONI ; if zero: JUMP

**Example** Jump to address LEO if R6 is equal to the table contents.

CMP R6, Table(R5) ; Compare content of R6 with content of

; MEM (table address + content of R5)

JEQ LEO ; Jump if both data are equal

.....; No, data are not equal, continue here

**Example** Branch to LABEL if R5 is 0.

TST R5 JZ LABEL

. . . . . .

CPU



## 3.4.6.26 JGE

JGE Jump if greater or equal

Syntax JGE label

**Operation** If (N . XOR. V) = 0 then jump to label: PC + 2 P offset  $\rightarrow PC$ 

If (N.XOR. V) = 1 then execute the following instruction

**Description** The status register negative bit (N) and overflow bit (V) are tested. If both N and V are

set or reset, the 10-bit signed offset contained in the instruction LSBs is added to the program counter. If only one is set, the instruction following the jump is executed.

This allows comparison of signed integers.

Status Bits Status bits are not affected.

**Example** When the content of R6 is greater or equal to the memory pointed to by R7, the program

continues at label EDE.

CMP @R7,R6 ; R6 >= (R7)?, compare on signed numbers

. . . . . .



## 3.4.6.27 JL

JL Jump if less

Syntax JL label

**Operation** If (N . XOR. V) = 1 then jump to label: PC + 2 offset  $\rightarrow PC$ 

If (N.XOR. V) = 0 then execute following instruction

**Description** The status register negative bit (N) and overflow bit (V) are tested. If only one is set, the

10-bit signed offset contained in the instruction LSBs is added to the program counter. If

both N and V are set or reset, the instruction following the jump is executed.

This allows comparison of signed integers.

Status Bits Status bits are not affected.

**Example** When the content of R6 is less than the memory pointed to by R7, the program

continues at label EDE.

CMP @R7,R6 ; R6 < (R7)?, compare on signed numbers

JL EDE ; Yes, R6 < (R7) ..... ; No, proceed

. . . . . .



# 3.4.6.28 JMP

JMP Jump unconditionally

Syntax JMP label

**Operation**  $PC + 2 \times offset \rightarrow PC$ 

**Description** The 10-bit signed offset contained in the instruction LSBs is added to the program

counter.

Status Bits Status bits are not affected.

Hint This one-word instruction replaces the BRANCH instruction in the range of –511 to +512

words relative to the current program counter.



## 3.4.6.29 JN

JN Jump if negative

Syntax JN label

**Operation** if N = 1:  $PC + 2 \times offset \rightarrow PC$ 

if N = 0: execute following instruction

**Description** The negative bit (N) of the status register is tested. If it is set, the 10-bit signed offset

contained in the instruction LSBs is added to the program counter. If N is reset, the next

instruction following the jump is executed.

Status Bits Status bits are not affected.

**Example** The result of a computation in R5 is to be subtracted from COUNT. If the result is negative, COUNT is to be cleared and the program continues execution in another path.

SUB R5,COUNT ; COUNT - R5 -> COUNT
JN L\$1 ; If negative continue with COUNT=0 at PC=L\$1

.....; Continue with COUNT>=0

L\$1 CLR COUNT

. . . . . .



# 3.4.6.30 JNC, JLO

JNC Jump if carry not set

JLO Jump if lower

Syntax JNC label

JLO label

**Operation** if C = 0: PC + 2 offset  $\rightarrow PC$ 

if C = 1: execute following instruction

**Description** The status register carry bit (C) is tested. If it is reset, the 10-bit signed offset contained

in the instruction LSBs is added to the program counter. If C is set, the next instruction following the jump is executed. JNC (jump if no carry/lower) is used for the comparison

of unsigned numbers (0 to 65536).

Status Bits Status bits are not affected.

**Example** The result in R6 is added in BUFFER. If an overflow occurs, an error handling routine at

address ERROR is used.

ADD R6,BUFFER ; BUFFER + R6 -> BUFFER
JNC CONT ; No carry, jump to CONT
ERROR ..... ; Error handler start

.....

. . . . . .

CONT ..... ; Continue with normal program flow

. . . . . .

**Example** Branch to STL2 if byte STATUS contains 1 or 0.

CMP.B #2,STATUS

JLO STL 2 ; STATUS < 2

.....; STATUS >= 2, continue here



# 3.4.6.31 JNE, JNZ

JNE Jump if not equal

JNZ Jump if not zero

Syntax JNE label

JNZ label

**Operation** If Z = 0: PC + 2 a offset  $\rightarrow$  PC

If Z = 1: execute following instruction

**Description** The status register zero bit (Z) is tested. If it is reset, the 10-bit signed offset contained in

the instruction LSBs is added to the program counter. If Z is set, the next instruction

following the jump is executed.

Status Bits Status bits are not affected.

**Example** Jump to address TONI if R7 and R8 have different contents.

CMP R7,R8 ; COMPARE R7 WITH R8 JNE TONI ; if different: jump ..... ; if equal, continue



### 3.4.6.32 MOV

MOV[.W] Move source to destination

MOV.B Move source to destination

Syntax MOV src,dst or MOV.W src,dst

MOV.B src,dst

**Operation**  $\operatorname{src} \to \operatorname{dst}$ 

**Description** The source operand is moved to the destination.

The source operand is not affected. The previous contents of the destination are lost.

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The contents of table EDE (word data) are copied to table TOM. The length of the tables

must be 020h locations.

```
MOV #EDE,R10 ; Prepare pointer
MOV #020h,R9 ; Prepare counter
```

Loop MOV @R10+,TOM-EDE-2(R10) ; Use pointer in R10 for both tables

DEC R9 ; Decrement counter

JNZ Loop ; Counter not 0, continue copying

.....; Copying completed

. . . . . .

**Example** The contents of table EDE (byte data) are copied to table TOM. The length of the tables should be 020h locations

```
MOV #EDE,R10 ; Prepare pointer
MOV #020h,R9 ; Prepare counter
```

Loop MOV.B @R10+,TOM-EDE-1(R10) ; Use pointer in R10 for

; both tables

DEC R9 ; Decrement counter

JNZ Loop ; Counter not 0, continue

; copying

.....; Copying completed

. . . . . .



## 3.4.6.33 NOP

\*NOP No operation

Syntax NOP

**Operation** None

**Emulation** MOV #0, R3

**Description** No operation is performed. The instruction may be used for the elimination of instructions

during the software check or for defined waiting times.

Status Bits Status bits are not affected.

The NOP instruction is mainly used for two purposes:

- · To fill one, two, or three memory words
- To adjust software timing

# NOTE: Emulating No-Operation Instruction

Other instructions can emulate the NOP function while providing different numbers of instruction cycles and code words. Some examples are:

```
MOV
    #0,R3
                  ; 1 cycle, 1 word
VOM
    O(R4), O(R4); 6 cycles, 3 words
    @R4,0(R4)
                  ; 5 cycles, 2 words
    #0,EDE(R4)
BIC
                  ; 4 cycles, 2 words
JMP
    $+2
                  ; 2 cycles, 1 word
BIC
    #0,R5
                  ; 1 cycle, 1 word
```

However, care should be taken when using these examples to prevent unintended results. For example, if MOV 0(R4), 0(R4) is used and the value in R4 is 120h, then a security violation occurs with the watchdog timer (address 120h), because the security key was not used.



### 3.4.6.34 POP

Pop word from stack to destination \*POP[.W]

\*POP.B Pop byte from stack to destination

**Syntax** POP dst

POP.B dst

Operation @SP → temp

> $SP + 2 \rightarrow SP$  $temp \rightarrow dst$

**Emulation** MOV @SP+,dst or MOV.W @SP+,dst

MOV.B @SP+,dst

The stack location pointed to by the stack pointer (TOS) is moved to the destination. The Description

stack pointer is incremented by two afterwards.

**Status Bits** Status bits are not affected.

The contents of R7 and the status register are restored from the stack. Example

> POP ; Restore R7

POP SR ; Restore status register

Example The contents of RAM byte LEO is restored from the stack.

> POP.B ; The low byte of the stack is moved to LEO. LEO

The contents of R7 is restored from the stack. Example

> POP.B R7 ; The low byte of the stack is moved to R7,

> > ; the high byte of R7 is 00h

Example The contents of the memory pointed to by R7 and the status register are restored from

the stack.

POP.B ; The low byte of the stack is moved to the 0(R7)

; the byte which is pointed to by R7

R7 = 203h; Example:

Mem(R7) = low byte of system stack

; Example: R7 = 20Ah

Mem(R7) = low byte of system stack

; Last word on stack moved to the SR POP SR

NOTE: The System Stack Pointer

The system stack pinter (SP) is always incremented by two, independent

of the byte suffix.



## 3.4.6.35 PUSH

PUSH[.W] Push word onto stack

PUSH.B Push byte onto stack

Syntax PUSH src or PUSH.W src

PUSH.B src

**Operation** SP - 2  $\rightarrow$  SP

 $src \rightarrow @SP$ 

**Description** The stack pointer is decremented by two, then the source operand is moved to the RAM

word addressed by the stack pointer (TOS).

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The contents of the status register and R8 are saved on the stack.

PUSH SR ; save status register

PUSH R8 ; save R8

**Example** The contents of the peripheral TCDAT is saved on the stack.

PUSH.B &TCDAT ; save data from 8-bit peripheral module,

; address TCDAT, onto stack

## NOTE: System Stack Pointer

The System stack pointer (SP) is always decremented by two, independent of the byte suffix.



# 3.4.6.36 RET

\*RET Return from subroutine

Syntax RET

 $\textbf{Operation} \qquad \qquad @SP \rightarrow PC$ 

 $SP + 2 \rightarrow SP$ 

Emulation MOV @SP+,PC

**Description** The return address pushed onto the stack by a CALL instruction is moved to the

program counter. The program continues at the code address following the subroutine

call.

Status Bits Status bits are not affected.



## 3.4.6.37 RETI

**RETI** Return from interrupt

Syntax RETI

Operation  $TOS \rightarrow SR$ 

 $SP + 2 \rightarrow SP$ 

 $\mathsf{TOS} \to \mathsf{PC}$ 

 $SP + 2 \rightarrow SP$ 

**Description** The status register is restored to the value at the beginning of the interrupt service

routine by replacing the present SR contents with the TOS contents. The stack pointer

(SP) is incremented by two.

The program counter is restored to the value at the beginning of interrupt service. This is the consecutive step after the interrupted program flow. Restoration is performed by replacing the present PC contents with the TOS memory contents. The stack pointer

(SP) is incremented.

Status Bits N: Restored from system stack

Z: Restored from system stack

C: Restored from system stack

V: Restored from system stack

Mode Bits OSCOFF, CPUOFF, and GIE are restored from system stack.

**Example** Figure 3-14 illustrates the main program interrupt.

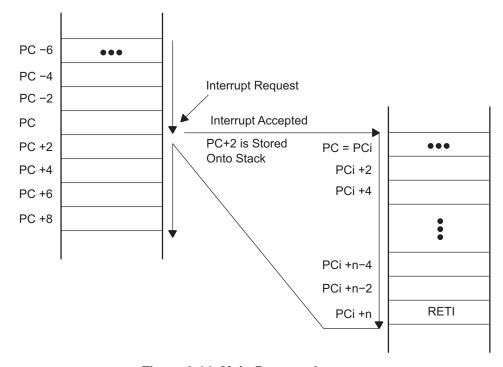


Figure 3-14. Main Program Interrupt



## 3.4.6.38 RLA

\*RLA[.W] Rotate left arithmetically \*RLA.B Rotate left arithmetically Syntax 1 4 1

RLA dst or RLA.W dst

RLA.B dst

Operation C <- MSB <- MSB-1 .... LSB+1 <- LSB <- 0

**Emulation** ADD dst,dst ADD.B dst,dst

### Description

The destination operand is shifted left one position as shown in Figure 3-15. The MSB is shifted into the carry bit (C) and the LSB is filled with 0. The RLA instruction acts as a signed multiplication by 2.

An overflow occurs if dst ≥ 04000h and dst < 0C000h before operation is performed: the result has changed sign.

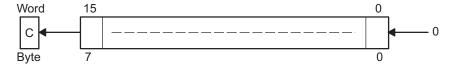


Figure 3-15. Destination Operand – Arithmetic Shift Left

An overflow occurs if dst ≥ 040h and dst < 0C0h before the operation is performed: the result has changed sign.

### **Status Bits**

N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB

V: Set if an arithmetic overflow occurs:

the initial value is 04000h ≤ dst < 0C000h; reset otherwise

Set if an arithmetic overflow occurs:

the initial value is 040h ≤ dst < 0C0h; reset otherwise

### **Mode Bits**

OSCOFF, CPUOFF, and GIE are not affected.

R7 is multiplied by 2. Example

> RLA R7 ; Shift left R7 (x 2)

## Example

The low byte of R7 is multiplied by 4.

RLA.B ; Shift left low byte of R7 (x 2)RLA.B ; Shift left low byte of R7 (x 4)

### NOTE: RLA Substitution

The assembler does not recognize the instruction:

RLA @R5+, RLA.B @R5+, or RLA(.B) @R5

It must be substituted by:

ADD @R5+,-2(R5), ADD.B @R5+,-1(R5), or ADD(.B) @R5



### 3.4.6.39 RLC

\*RLC[.W] Rotate left through carry \*RLC.B Rotate left through carry Syntax 1 4 1 RLC dst or RLC.W dst

RLC.B dst

Operation C <- MSB <- MSB-1 .... LSB+1 <- LSB <- C

**Emulation** ADDC dst, dst

The destination operand is shifted left one position as shown in Figure 3-16. The carry Description

bit (C) is shifted into the LSB and the MSB is shifted into the carry bit (C).

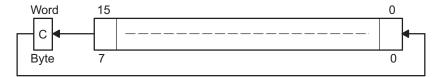


Figure 3-16. Destination Operand - Carry Left Shift

### **Status Bits**

N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB

V: Set if an arithmetic overflow occurs

the initial value is 04000h ≤ dst < 0C000h; reset otherwise

Set if an arithmetic overflow occurs:

the initial value is 040h ≤ dst < 0C0h; reset otherwise

**Mode Bits** 

OSCOFF, CPUOFF, and GIE are not affected.

Example

R5 is shifted left one position.

RLC

 $; (R5 \times 2) + C -> R5$ 

Example

The input P1IN.1 information is shifted into the LSB of R5.

```
BIT.B
        #2,&P1IN
                   ; Information -> Carry
                   ; Carry=P0in.1 -> LSB of R5
```

Example

The MEM(LEO) content is shifted left one position.

```
RLC.B
           LEO
                           ; Mem(LEO) \times 2 + C \rightarrow Mem(LEO)
```

## NOTE: RLC and RLC.B Substitution

The assembler does not recognize the instruction:

```
RLC @R5+, RLC @R5, or RLC(.B) @R5
```

It must be substitued by:

```
ADDC @R5+,-2(R5), ADDC.B @R5+,-1(R5), or ADDC(.B) @R5
```



### 3.4.6.40 RRA

RRA[.W] Rotate right arithmetically

RRA.B Rotate right arithmetically

Syntax RRA dst or RRA.W dst

RRA.B dst

**Operation** MSB  $\rightarrow$  MSB, MSB  $\rightarrow$  MSB-1, ... LSB+1  $\rightarrow$  LSB, LSB  $\rightarrow$  C

## **Description**

The destination operand is shifted right one position as shown in Figure 3-17. The MSB is shifted into the MSB, the MSB is shifted into the MSB-1, and the LSB+1 is shifted into the LSB.

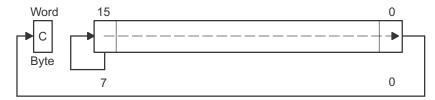


Figure 3-17. Destination Operand – Arithmetic Right Shift

### **Status Bits**

N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

# **Mode Bits**

OSCOFF, CPUOFF, and GIE are not affected.

# Example

R5 is shifted right one position. The MSB retains the old value. It operates equal to an arithmetic division by 2.

```
RRA
             ; R5/2 -> R5
       R5
    The value in R5 is multiplied by 0.75 (0.5 + 0.25).
;
;
PUSH
         R5
                     ; Hold R5 temporarily using stack
                     ; R5 x 0.5 -> R5
RRA
ADD
         @SP+,R5
                     ; R5 \times 0.5 + R5 = 1.5 \times R5 -> R5
                     ; (1.5 \times R5) \times 0.5 = 0.75 \times R5 \rightarrow R5
RRA
         R5
```

### Example

The low byte of R5 is shifted right one position. The MSB retains the old value. It operates equal to an arithmetic division by 2.

```
RRA.B R5 ; R5/2 -> R5: operation is on low byte only ; High byte of R5 is reset PUSH.B R5 ; R5 x 0.5 -> TOS RRA.B @SP ; TOS x 0.5 = 0.5 x R5 x 0.5 = 0.25 x R5 -> TOS ADD.B @SP+,R5 ; R5 x 0.5 + R5 x 0.25 = 0.75 x R5 -> R5 .....
```



## 3.4.6.41 RRC

RRC[.W] Rotate right through carry
RRC.B Rotate right through carry

Syntax RRC dst or RRC.W dst

RRC dst

**Description** The destination operand is shifted right one position as shown in Figure 3-18. The carry

bit (C) is shifted into the MSB, the LSB is shifted into the carry bit (C).

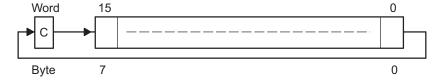


Figure 3-18. Destination Operand - Carry Right Shift

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIEare not affected.

**Example** R5 is shifted right one position. The MSB is loaded with 1.

SETC ; Prepare carry for MSB RRC R5 ; R5/2 + 8000h  $\rightarrow$  R5

**Example** R5 is shifted right one position. The MSB is loaded with 1.

SETC ; Prepare carry for MSB

RRC.B R5 ; R5/2 + 80h -> R5; low byte of R5 is used



## 3.4.6.42 SBC

Subtract source and borrow/.NOT. carry from destination \*SBC[.W]

Subtract source and borrow/.NOT. carry from destination \*SBC.B

SBC dst or SBC.W dst **Syntax** 

SBC.B dst

Operation  $dst + 0FFFFh + C \rightarrow dst$ 

 $dst + 0FFh + C \rightarrow dst$ 

**Emulation** SUBC #0,dst

SUBC.B #0,dst

Description The carry bit (C) is added to the destination operand minus one. The previous contents

of the destination are lost.

**Status Bits** N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise.

Set to 1 if no borrow, reset if borrow.

V: Set if an arithmetic overflow occurs, reset otherwise.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

Example The 16-bit counter pointed to by R13 is subtracted from a 32-bit counter pointed to by

R12.

SUB @R13,0(R12) ; Subtract LSDs

SBC 2(R12) ; Subtract carry from MSD

The 8-bit counter pointed to by R13 is subtracted from a 16-bit counter pointed to by Example

R12.

SUB.B @R13,0(R12) ; Subtract LSDs

SBC.B 1(R12) ; Subtract carry from MSD

**NOTE:** Borrow Implementation

Carry bit The borrow is treated as a .NOT. carry: **Borrow** 

> 0 Yes 1

No



# 3.4.6.43 SETC

\*SETC Set carry bit

Emulation BIS #1,SR

**Description** The carry bit (C) is set.

Status Bits N: Not affected

Z: Not affected

C: Set

V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** Emulation of the decimal subtraction:

Subtract R5 from R6 decimally

Assume that R5 = 03987h and R6 = 04137h

DSUB ADD #06666h,R5 ; Move content R5 from 0-9 to 6-0Fh

; R5 = 03987h + 06666h = 09FEDhINV R5 ; Invert this (result back to 0-9)

; R5 = .NOT. R5 = 06012h

SETC ; Prepare carry = 1

DADD R5,R6 ; Emulate subtraction by addition of:

; (010000h - R5 - 1)

; R6 = R6 + R5 + 1

; R6 = 0150h



# 3.4.6.44 SETN

\*SETN Set negative bit

Syntax SETN

Emulation BIS #4,SR

**Description** The negative bit (N) is set.

Status Bits N: Set

Z: Not affectedC: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.



# 3.4.6.45 SETZ

\*SETZ Set zero bit

Syntax SETZ

Emulation BIS #2,SR

**Description** The zero bit (Z) is set.

Status Bits N: Not affected

Z: Set

C: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.



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#### 3.4.6.46 SUB

SUB[.W] Subtract source from destination

SUB.B Subtract source from destination

Syntax SUB src,dst or SUB.W src,dst

SUB.B src,dst

**Operation**  $dst + .NOT.src + 1 \rightarrow dst$ 

or

 $[(dst - src \rightarrow dst)]$ 

**Description** The source operand is subtracted from the destination operand by adding the source

operand's 1s complement and the constant 1. The source operand is not affected. The

previous contents of the destination are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise.

Set to 1 if no borrow, reset if borrow.

V: Set if an arithmetic overflow occurs, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** See example at the SBC instruction.

**Example** See example at the SBC.B instruction.

NOTE: Borrow Is Treated as a .NOT.

The borrow is treated as a .NOT. carry:

Borrow Carry bit

Yes 0

No 1

CPU



Instruction Set www.ti.com

#### 3.4.6.47 SUBC, SBB

**SUBC[.W], SBB[.W]** Subtract source and borrow/.NOT. carry from destination

SUBC.B, SBB.B Subtract source and borrow/.NOT. carry from destination

Syntax SUBC src,dst or SUBC.W src,dst or

SBB src,dst or SBB.W src,dst SUBC.B src,dst or SBB.B src,dst

**Operation**  $dst + .NOT.src + C \rightarrow dst$ 

or

 $(dst - src - 1 + C \rightarrow dst)$ 

**Description** The source operand is subtracted from the destination operand by adding the source

operand's 1s complement and the carry bit (C). The source operand is not affected. The

previous contents of the destination are lost.

Status Bits N: Set if result is negative, reset if positive.

Z: Set if result is zero, reset otherwise.

C: Set if there is a carry from the MSB of the result, reset otherwise.

Set to 1 if no borrow, reset if borrow.

V: Set if an arithmetic overflow occurs, reset otherwise.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** Two floating point mantissas (24 bits) are subtracted.

LSBs are in R13 and R10, MSBs are in R12 and R9.

SUB.W R13,R10 ; 16-bit part, LSBs SUBC.B R12,R9 ; 8-bit part, MSBs

Example The 16-bit counter pointed to by R13 is subtracted from a 16-bit counter in R10 and

R11(MSD).

SUB.B @R13+,R10 ; Subtract LSDs without carry SUBC.B @R13,R11 ; Subtract MSDs with carry ... ; resulting from the LSDs

#### **NOTE:** Borrow Implementation

The borrow is treated as a .NOT. carry:

Borrow

Yes

No

1



Instruction Set www.ti.com

# 3.4.6.48 SWPB

**SWPB** Swap bytes **Syntax** SWPB dst

Operation Bits 15 to  $8 \leftrightarrow \text{bits 7 to 0}$ 

Description The destination operand high and low bytes are exchanged as shown in Figure 3-19.

OSCOFF, CPUOFF, and GIE are not affected. **Mode Bits** 

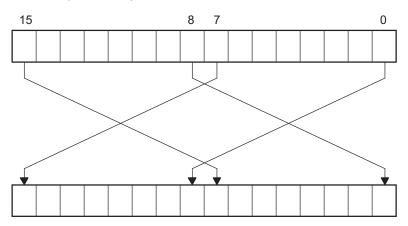


Figure 3-19. Destination Operand - Byte Swap

Example MOV #040BFh,R7 ; 0100000010111111 -> R7 R7 ; 10111111101000000 in R7 SWPB

The value in R5 is multiplied by 256. The result is stored in R5,R4. Example

> SWPB R5

R5,R4 MOV ; Copy the swapped value to R4

BIC #0FF00h,R5 ; Correct the result BIC #00FFh,R4 ; Correct the result

CPU



Instruction Set www.ti.com

#### 3.4.6.49 SXT

SXT Extend Sign

Syntax SXT dst

**Operation** Bit  $7 \rightarrow$  Bit 8 ...... Bit 15

**Description** The sign of the low byte is extended into the high byte as shown in Figure 3-20.

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise (.NOT. Zero)

V: Reset

**Mode Bits** 

OSCOFF, CPUOFF, and GIE are not affected.

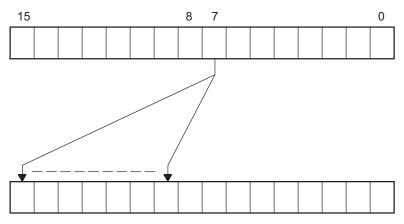


Figure 3-20. Destination Operand - Sign Extension

# **Example**

R7 is loaded with the P1IN value. The operation of the sign-extend instruction expands bit 8 to bit 15 with the value of bit 7.

R7 is then added to R6.

```
MOV.B &P1IN,R7 ; P1IN = 080h: .... 1000 0000 SXT R7 ; R7 = 0FF80h: 1111 1111 1000 0000
```



www.ti.com Instruction Set

#### 3.4.6.50 TST

\*TST[.W] Test destination

\*TST.B Test destination

Syntax TST dst or TST.W dst

TST.B dst

**Operation** dst + 0FFFFh + 1

dst + 0FFh + 1

**Emulation** CMP #0,dst

CMP.B #0,dst

**Description** The destination operand is compared with zero. The status bits are set according to the

result. The destination is not affected.

Status Bits N: Set if destination is negative, reset if positive

Z: Set if destination contains zero, reset otherwise

C: Set

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** R7 is tested. If it is negative, continue at R7NEG; if it is positive but not zero, continue at

R7POS.

TST R7 ; Test R7

JN R7NEG ; R7 is negative

JZ R7ZERO ; R7 is zero

R7POS ..... ; R7 is positive but not zero

R7NEG ..... ; R7 is negative R7ZERO ..... ; R7 is zero

**Example** The low byte of R7 is tested. If it is negative, continue at R7NEG; if it is positive but not zero, continue at R7POS.

TST.B R7 ; Test low byte of R7

JN R7NEG ; Low byte of R7 is negative JZ R7ZERO ; Low byte of R7 is zero

R7POS ..... ; Low byte of R7 is positive but not zero

R7NEG .... ; Low byte of R7 is negative R7ZERO ..... ; Low byte of R7 is zero



Instruction Set www.ti.com

#### 3.4.6.51 XOR

**XOR[.W]** Exclusive OR of source with destination

**XOR.B** Exclusive OR of source with destination

Syntax XOR src,dst or XOR.W src,dst

XOR.B src,dst

**Operation**  $\operatorname{src} .XOR. \operatorname{dst} \to \operatorname{dst}$ 

**Description** The source and destination operands are exclusive ORed. The result is placed into the

destination. The source operand is not affected.

Status Bits N: Set if result MSB is set, reset if not set

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise ( = .NOT. Zero)

V: Set if both operands are negative

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The bits set in R6 toggle the bits in the RAM word TONI.

XOR R6,TONI ; Toggle bits of word TONI on the bits set in R6

**Example** The bits set in R6 toggle the bits in the RAM byte TONI.

 ${\tt XOR.B}$  R6,  ${\tt TONI}$  ; Toggle bits of byte TONI on the bits set in

; low byte of R6

**Example** Reset to 0 those bits in low byte of R7 that are different from bits in RAM byte EDE.

XOR.B EDE,R7 ; Set different bit to "1s"
INV.B R7 ; Invert Lowbyte, Highbyte is 0h

# Chapter 4





# **CPUX**

This chapter describes the extended MSP430X 16-bit RISC CPU with 1-MB memory access, its addressing modes, and instruction set. The MSP430X CPU is implemented in all MSP430 devices that exceed 64-KB of address space.

Topic Page

4.1	CPU Introduction	
4.2	Interrupts	118
4.3	CPU Registers	119
4.4	Addressing Modes	125
4.5	MSP430 and MSP430X Instructions	142
4.6	Instruction Set Description	160
	-	



CPU Introduction www.ti.com

#### 4.1 CPU Introduction

The MSP430X CPU incorporates features specifically designed for modern programming techniques such as calculated branching, table processing and the use of high-level languages such as C. The MSP430X CPU can address a 1-MB address range without paging. In addition, the MSP430X CPU has fewer interrupt overhead cycles and fewer instruction cycles in some cases than the MSP430 CPU, while maintaining the same or better code density than the MSP430 CPU. The MSP430X CPU is backward compatible with the MSP430 CPU.

The MSP430X CPU features include:

- RISC architecture
- Orthogonal architecture
- · Full register access including program counter, status register and stack pointer
- Single-cycle register operations
- Large register file reduces fetches to memory
- 20-bit address bus allows direct access and branching throughout the entire memory range without paging
- 16-bit data bus allows direct manipulation of word-wide arguments
- Constant generator provides the six most often used immediate values and reduces code size
- · Direct memory-to-memory transfers without intermediate register holding
- Byte, word, and 20-bit address-word addressing

The block diagram of the MSP430X CPU is shown in Figure 4-1.



www.ti.com CPU Introduction

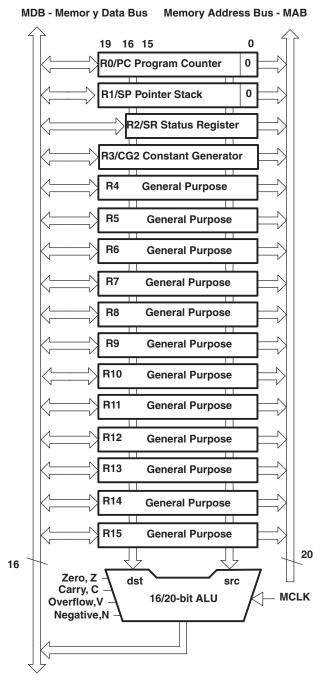


Figure 4-1. MSP430X CPU Block Diagram



Interrupts www.ti.com

# 4.2 Interrupts

The MSP430X uses the same interrupt structure as the MSP430:

- Vectored interrupts with no polling necessary
- Interrupt vectors are located downward from address 0FFFEh

Interrupt operation for both MSP430 and MSP430X CPUs is described in *Chapter 2 System Resets*, *Interrupts*, *and Operating modes*, *Section 2 Interrupts*. The interrupt vectors contain 16-bit addresses that point into the lower 64-KB memory. This means all interrupt handlers must start in the lower 64-KB memory, even in MSP430X devices.

During an interrupt, the program counter and the status register are pushed onto the stack as shown in Figure 4-2. The MSP430X architecture efficiently stores the complete 20-bit PC value by automatically appending the PC bits 19:16 to the stored SR value on the stack. When the RETI instruction is executed, the full 20-bit PC is restored making return from interrupt to any address in the memory range possible.

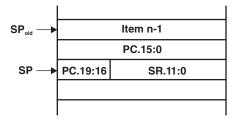


Figure 4-2. PC Storage on the Stack for Interrupts



www.ti.com CPU Registers

# 4.3 CPU Registers

The CPU incorporates 16 registers (R0 through R15). Registers R0, R1, R2, and R3 have dedicated functions. Registers R4 through R15 are working registers for general use.

# 4.3.1 Program Counter (PC)

The 20-bit PC (PC/R0) points to the next instruction to be executed. Each instruction uses an even number of bytes (2, 4, 6, or 8 bytes), and the PC is incremented accordingly. Instruction accesses are performed on word boundaries, and the PC is aligned to even addresses. Figure 4-3 shows the PC.



Figure 4-3. Program Counter

The PC can be addressed with all instructions and addressing modes. A few examples:

```
MOV.W #LABEL,PC; Branch to address LABEL (lower 64KB)

MOVA #LABEL,PC; Branch to address LABEL (1MB memory)

MOV.W LABEL,PC; Branch to address in word LABEL; (lower 64KB)

MOV.W @R14,PC; Branch indirect to address in; R14 (lower 64KB)

ADDA #4,PC; Skip two words (1MB memory)
```

The BR and CALL instructions reset the upper four PC bits to 0. Only addresses in the lower 64-KB address range can be reached with the BR or CALL instruction. When branching or calling, addresses beyond the lower 64-KB range can only be reached using the BRA or CALLA instructions. Also, any instruction to directly modify the PC does so according to the used addressing mode. For example, MOV.W #value, PC clears the upper four bits of the PC, because it is a .W instruction.

The PC is automatically stored on the stack with CALL (or CALLA) instructions and during an interrupt service routine. Figure 4-4 shows the storage of the PC with the return address after a CALLA instruction. A CALL instruction stores only bits 15:0 of the PC.

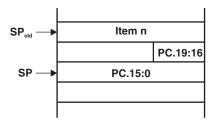


Figure 4-4. PC Storage on the Stack for CALLA

The RETA instruction restores bits 19:0 of the PC and adds 4 to the stack pointer (SP). The RET instruction restores bits 15:0 to the PC and adds 2 to the SP.

# 4.3.2 Stack Pointer (SP)

The 20-bit SP (SP/R1) is used by the CPU to store the return addresses of subroutine calls and interrupts. It uses a predecrement, postincrement scheme. In addition, the SP can be used by software with all instructions and addressing modes. Figure 4-5 shows the SP. The SP is initialized into RAM by the user, and is always aligned to even addresses.



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Figure 4-6 shows the stack usage. Figure 4-7 shows the stack usage when 20-bit address words are pushed.

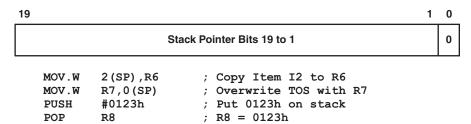


Figure 4-5. Stack Pointer

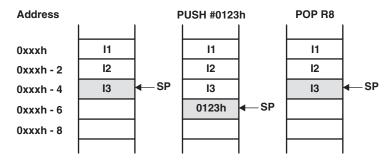


Figure 4-6. Stack Usage

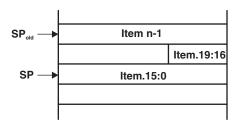
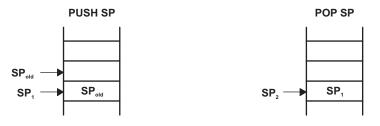


Figure 4-7. PUSHX.A Format on the Stack

The special cases of using the SP as an argument to the PUSH and POP instructions are described and shown in Figure 4-8.



The stack pointer is changed after a PUSH SP instruction.

The stack pointer is not changed after a POP SP instruction. The POP SP instruction places SP1 into the stack pointer SP (SP2 = SP1)

Figure 4-8. PUSH SP, POP SP Sequence



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# 4.3.3 Status Register (SR)

The 16-bit SR (SR/R2), used as a source or destination register, can only be used in register mode addressed with word instructions. The remaining combinations of addressing modes are used to support the constant generator. Figure 4-9 shows the SR bits. Do not write 20-bit values to the SR. Unpredictable operation can result.

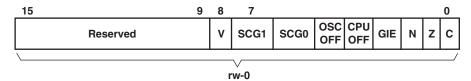


Figure 4-9. SR Bits

Table 4-1 describes the SR bits.

**Table 4-1. SR Bit Description** 

Bit	Description					
Reserved	Reserved					
V	Overflow. This bit is set when the result of an arithmetic operation	on overflows the signed-variable range.				
	ADD(.B), ADDX(.B,.A), ADDC(.B), ADDCX(.B.A), ADDA	Set when: positive + positive = negative negative + negative = positive otherwise reset				
	SUB(.B), SUBX(.B,.A), SUBC(.B), SUBCX(.B,.A), SUBA, CMP(.B), CMPX(.B,.A), CMPA	Set when: positive – negative = negative negative – positive = positive otherwise reset				
SCG1	System clock generator 1. This bit may be to enable/disable functions in the clock system depending on the device family; for example, DCO bias enable/disable					
SCG0	System clock generator 0. This bit may be used to enable/disable device family; for example, FLL disable/enable	le functions in the clock system depending on the				
OSCOFF	Oscillator off. This bit, when set, turns off the LFXT1 crystal osc SMCLK.	illator when LFXT1CLK is not used for MCLK or				
CPUOFF	CPU off. This bit, when set, turns off the CPU.					
GIE	General interrupt enable. This bit, when set, enables maskable disabled.	interrupts. When reset, all maskable interrupts are				
N	Negative. This bit is set when the result of an operation is negative.	tive and cleared when the result is positive.				
Z	Zero. This bit is set when the result of an operation is 0 and clear	ared when the result is not 0.				
С	Carry. This bit is set when the result of an operation produced a	carry and cleared when no carry occurred.				



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# 4.3.4 Constant Generator Registers (CG1 and CG2)

Six commonly-used constants are generated with the constant generator registers R2 (CG1) and R3 (CG2), without requiring an additional 16-bit word of program code. The constants are selected with the source register addressing modes (As), as described in Table 4-2.

Table 4-2. Values of Constant Generators CG1, CG2

Register	As	Constant	Remarks
R2	00	_	Register mode
R2	01	(0)	Absolute address mode
R2	10	00004h	+4, bit processing
R2	11	00008h	+8, bit processing
R3	00	00000h	0, word processing
R3	01	00001h	+1
R3	10	00002h	+2, bit processing
R3	11	FFh, FFFFh, FFFFFh	-1, word processing

The constant generator advantages are:

- No special instructions required
- No additional code word for the six constants
- No code memory access required to retrieve the constant

The assembler uses the constant generator automatically if one of the six constants is used as an immediate source operand. Registers R2 and R3, used in the constant mode, cannot be addressed explicitly; they act as source-only registers.

#### 4.3.4.1 Constant Generator – Expanded Instruction Set

The RISC instruction set of the MSP430 has only 27 instructions. However, the constant generator allows the MSP430 assembler to support 24 additional emulated instructions. For example, the single-operand instruction:

CLR dst

is emulated by the double-operand instruction with the same length:

MOV R3,dst

where the #0 is replaced by the assembler, and R3 is used with As = 00.

INC dst

is replaced by:

ADD #1,dst



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# 4.3.5 General-Purpose Registers (R4 to R15)

The 12 CPU registers (R4 to R15) contain 8-bit, 16-bit, or 20-bit values. Any byte-write to a CPU register clears bits 19:8. Any word-write to a register clears bits 19:16. The only exception is the SXT instruction. The SXT instruction extends the sign through the complete 20-bit register.

The following figures show the handling of byte, word, and address-word data. Note the reset of the leading most significant bits (MSBs) if a register is the destination of a byte or word instruction.

Figure 4-10 shows byte handling (8-bit data, .B suffix). The handling is shown for a source register and a destination memory byte and for a source memory byte and a destination register.

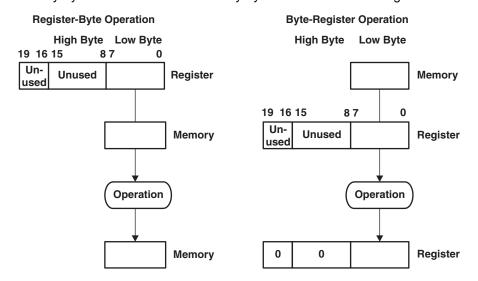


Figure 4-10. Register-Byte/Byte-Register Operation

Figure 4-11 and Figure 4-12 show 16-bit word handling (.W suffix). The handling is shown for a source register and a destination memory word and for a source memory word and a destination register.

#### **Register-Word Operation**

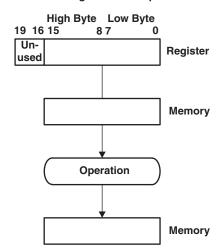


Figure 4-11. Register-Word Operation



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# Word-Register Operation High Byte Low Byte Memory 19 16 15 8 7 0 Unused Operation Register

Figure 4-12. Word-Register Operation

Figure 4-13 and Figure 4-14 show 20-bit address-word handling (.A suffix). The handling is shown for a source register and a destination memory address-word and for a source memory address-word and a destination register.

#### Register - Ad dress-Word Operation

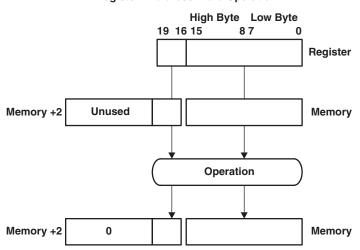


Figure 4-13. Register - Address-Word Operation



# Address-Word - Register Operation High Byte Low Byte 19 16 15 8 7 0 Memory +2 Unused Memory Register Operation

Figure 4-14. Address-Word - Register Operation

Register

# 4.4 Addressing Modes

Seven addressing modes for the source operand and four addressing modes for the destination operand use 16-bit or 20-bit addresses (see Table 4-3). The MSP430 and MSP430X instructions are usable throughout the entire 1MB memory range.

Table 4-3. Source/Destination Addressing

As/Ad	Addressing Mode	Syntax	Description
00/0	Register	Rn	Register contents are operand.
01/1	Indexed	X(Rn)	(Rn + X) points to the operand. X is stored in the next word, or stored in combination of the preceding extension word and the next word.
01/1	Symbolic	ADDR	(PC + X) points to the operand. X is stored in the next word, or stored in combination of the preceding extension word and the next word. Indexed mode X(PC) is used.
01/1	Absolute	&ADDR	The word following the instruction contains the absolute address. X is stored in the next word, or stored in combination of the preceding extension word and the next word. Indexed mode X(SR) is used.
10/-	Indirect Register	@Rn	Rn is used as a pointer to the operand.
11/–	Indirect Autoincrement	@Rn+	Rn is used as a pointer to the operand. Rn is incremented afterwards by 1 for .B instructions. by 2 for .W instructions, and by 4 for .A instructions.
11/–	Immediate	#N	N is stored in the next word, or stored in combination of the preceding extension word and the next word. Indirect autoincrement mode @PC+ is used.

The seven addressing modes are explained in detail in the following sections. Most of the examples show the same addressing mode for the source and destination, but any valid combination of source and destination addressing modes is possible in an instruction.

#### NOTE: Use of Labels EDE, TONI, TOM, and LEO

Throughout MSP430 documentation, EDE, TONI, TOM, and LEO are used as generic labels. They are only labels and have no special meaning.



# 4.4.1 Register Mode

Operation: The operand is the 8-, 16-, or 20-bit content of the used CPU register.

Length: One, two, or three words

Comment: Valid for source and destination

Byte operation: Byte operation reads only the eight least significant bits (LSBs) of the source

register Rsrc and writes the result to the eight LSBs of the destination register Rdst.

The bits Rdst.19:8 are cleared. The register Rsrc is not modified.

Word operation: Word operation reads the 16 LSBs of the source register Rsrc and writes the result

to the 16 LSBs of the destination register Rdst. The bits Rdst.19:16 are cleared.

The register Rsrc is not modified.

Address-word operation:

Address-word operation reads the 20 bits of the source register Rsrc and writes the

result to the 20 bits of the destination register Rdst. The register Rsrc is not

modified

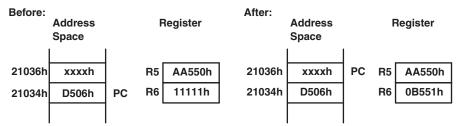
SXT exception: The SXT instruction is the only exception for register operation. The sign of the low

byte in bit 7 is extended to the bits Rdst.19:8.

Example: BIS.W R5,R6;

This instruction logically ORs the 16-bit data contained in R5 with the 16-bit

contents of R6. R6.19:16 is cleared.

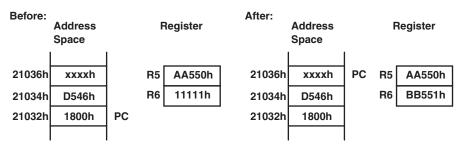


A550h.or.1111h = B551h

Example: BISX.A R5,R6;

This instruction logically ORs the 20-bit data contained in R5 with the 20-bit contents of R6.

The extension word contains the A/L bit for 20-bit data. The instruction word uses byte mode with bits A/L:B/W = 01. The result of the instruction is:



AA550h.or.11111h = BB551h



#### 4.4.2 Indexed Mode

The Indexed mode calculates the address of the operand by adding the signed index to a CPU register. The Indexed mode has three addressing possibilities:

- Indexed mode in lower 64-KB memory
- MSP430 instruction with Indexed mode addressing memory above the lower 64-KB memory
- MSP430X instruction with Indexed mode

# 4.4.2.1 Indexed Mode in Lower 64-KB Memory

If the CPU register Rn points to an address in the lower 64KB of the memory range, the calculated memory address bits 19:16 are cleared after the addition of the CPU register Rn and the signed 16-bit index. This means the calculated memory address is always located in the lower 64KB and does not overflow or underflow out of the lower 64-KB memory space. The RAM and the peripheral registers can be accessed this way and existing MSP430 software is usable without modifications as shown in Figure 4-15.

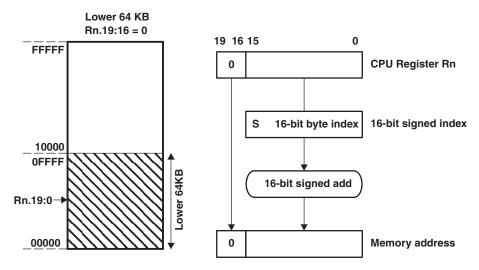


Figure 4-15. Indexed Mode in Lower 64KB

Length: Two or three words

Operation: The signed 16-bit index is located in the next word after the instruction and is added to

the CPU register Rn. The resulting bits 19:16 are cleared giving a truncated 16-bit memory address, which points to an operand address in the range 00000h to 0FFFFh.

The operand is the content of the addressed memory location.

Comment: Valid for source and destination. The assembler calculates the register index and inserts

it.

Example: ADD.B 1000h(R5), 0F000h(R6);

This instruction adds the 8-bit data contained in source byte 1000h(R5) and the

destination byte 0F000h(R6) and places the result into the destination byte. Source and destination bytes are both located in the lower 64KB due to the cleared bits 19:16 of

registers R5 and R6.

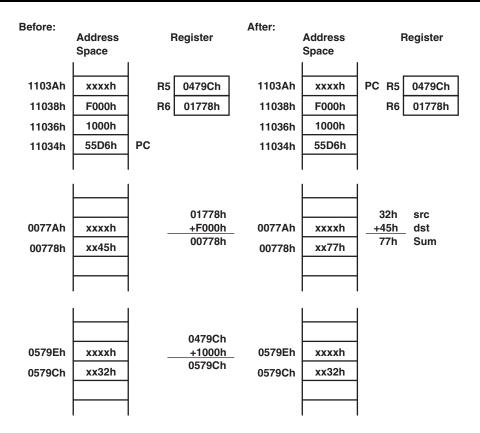
Source: The byte pointed to by R5 + 1000h results in address 0479Ch + 1000h = 0579Ch after

truncation to a 16-bit address.

Destination: The byte pointed to by R6 + F000h results in address 01778h + F000h = 00778h after

truncation to a 16-bit address.





#### 4.4.2.2 MSP430 Instruction With Indexed Mode in Upper Memory

If the CPU register Rn points to an address above the lower 64-KB memory, the Rn bits 19:16 are used for the address calculation of the operand. The operand may be located in memory in the range Rn ±32KB, because the index, X, is a signed 16-bit value. In this case, the address of the operand can overflow or underflow into the lower 64-KB memory space (see Figure 4-16 and Figure 4-17).

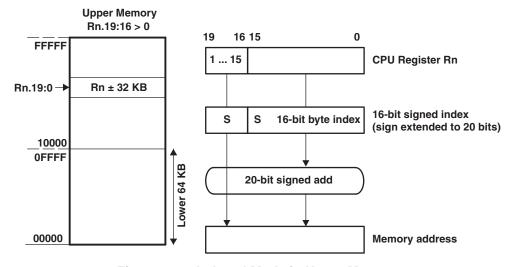


Figure 4-16. Indexed Mode in Upper Memory



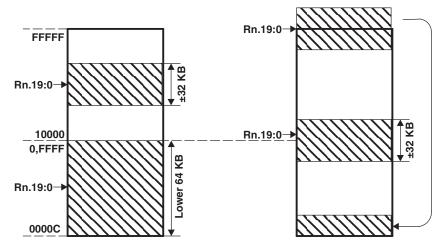


Figure 4-17. Overflow and Underflow for Indexed Mode

Length: Two or three words

Operation: The sign-extended 16-bit index in the next word after the instruction is added to the

20 bits of the CPU register Rn. This delivers a 20-bit address, which points to an address in the range 0 to FFFFFh. The operand is the content of the addressed

memory location.

Comment: Valid for source and destination. The assembler calculates the register index and

inserts it.

Example: ADD.W 8346h(R5),2100h(R6);

This instruction adds the 16-bit data contained in the source and the destination addresses and places the 16-bit result into the destination. Source and destination

operand can be located in the entire address range.

Source: The word pointed to by R5 + 8346h. The negative index 8346h is sign extended,

which results in address 23456h + F8346h = 1B79Ch.

Destination: The word pointed to by R6 + 2100h results in address 15678h + 2100h = 17778h.



Before:	Address Space	Register	After:	Address Space	Register
1103Ah 11038h 11036h 11034h	xxxxh 2100h 8346h 5596h	R5 23456h R6 15678h	1103Ah 11038h 11036h 11034h	xxxxh 2100h 8346h 5596h	PC R5 23456h R6 15678h
1777Ah 17778h	xxxxh 2345h	15678h <u>+02100h</u> 17778h	1777Ah 17778h	xxxxh 7777h	05432h src +02345h dst 07777h Sum
1B79Eh 1B79Ch	xxxxh 5432h	23456h <u>+F8346h</u> 1B79Ch	1B79Eh 1B79Ch	xxxxh 5432h	

#### 4.4.2.3 MSP430X Instruction With Indexed Mode

When using an MSP430X instruction with Indexed mode, the operand can be located anywhere in the range of Rn + 19 bits.

Length: Three or four words

Operation: The operand address is the sum of the 20-bit CPU register content and the 20-bit

index. The 4 MSBs of the index are contained in the extension word; the 16 LSBs are contained in the word following the instruction. The CPU register is not modified

Comment: Valid for source and destination. The assembler calculates the register index and

inserts it.

Example: ADDX.A 12346h(R5),32100h(R6);

This instruction adds the 20-bit data contained in the source and the destination

addresses and places the result into the destination.

Source: Two words pointed to by R5 + 12346h which results in address 23456h + 12346h =

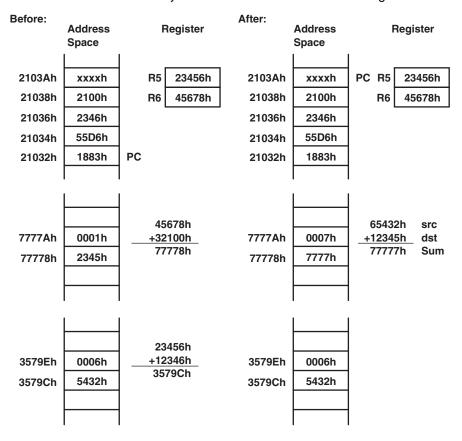
3579Ch.

Destination: Two words pointed to by R6 + 32100h which results in address 45678h + 32100h =

77778h.



The extension word contains the MSBs of the source index and of the destination index and the A/L bit for 20-bit data. The instruction word uses byte mode due to the 20-bit data length with bits A/L:B/W = 01.



# 4.4.3 Symbolic Mode

The Symbolic mode calculates the address of the operand by adding the signed index to the PC. The Symbolic mode has three addressing possibilities:

- Symbolic mode in lower 64-KB memory
- MSP430 instruction with Symbolic mode addressing memory above the lower 64-KB memory.
- MSP430X instruction with Symbolic mode

#### 4.4.3.1 Symbolic Mode in Lower 64KB

If the PC points to an address in the lower 64KB of the memory range, the calculated memory address bits 19:16 are cleared after the addition of the PC and the signed 16-bit index. This means the calculated memory address is always located in the lower 64KB and does not overflow or underflow out of the lower 64-KB memory space. The RAM and the peripheral registers can be accessed this way and existing MSP430 software is usable without modifications as shown in Figure 4-18.



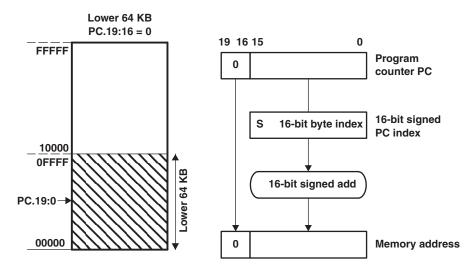


Figure 4-18. Symbolic Mode Running in Lower 64KB

Operation: The signed 16-bit index in the next word after the instruction is added temporarily to

the PC. The resulting bits 19:16 are cleared giving a truncated 16-bit memory address, which points to an operand address in the range 00000h to 0FFFFh. The

operand is the content of the addressed memory location.

Length: Two or three words

Comment: Valid for source and destination. The assembler calculates the PC index and

inserts it.

Example: ADD.B EDE, TONI ;

This instruction adds the 8-bit data contained in source byte EDE and destination byte TONI and places the result into the destination byte TONI. Bytes EDE and

TONI and the program are located in the lower 64KB.

Source: Byte EDE located at address 0579Ch, pointed to by PC + 4766h, where the PC

index 4766h is the result of 0579Ch - 01036h = 04766h. Address 01036h is the

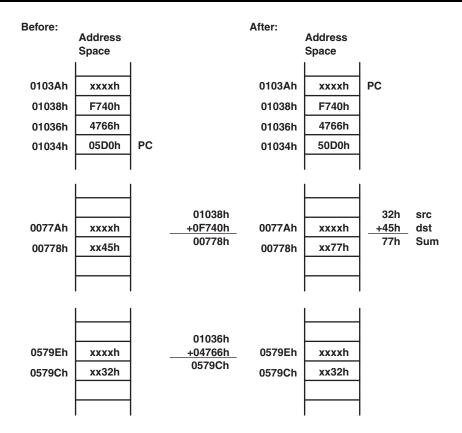
location of the index for this example.

Destination: Byte TONI located at address 00778h, pointed to by PC + F740h, is the truncated

16-bit result of 00778h - 1038h = FF740h. Address 01038h is the location of the

index for this example.





# 4.4.3.2 MSP430 Instruction With Symbolic Mode in Upper Memory

If the PC points to an address above the lower 64-KB memory, the PC bits 19:16 are used for the address calculation of the operand. The operand may be located in memory in the range PC ± 32KB, because the index, X, is a signed 16-bit value. In this case, the address of the operand can overflow or underflow into the lower 64-KB memory space as shown in Figure 4-19 and Figure 4-20.

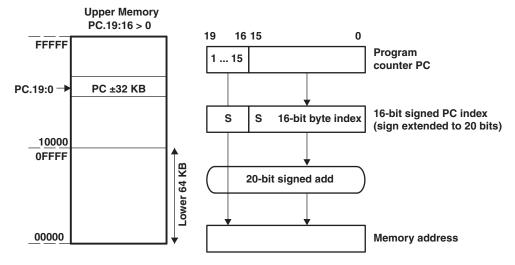


Figure 4-19. Symbolic Mode Running in Upper Memory

133



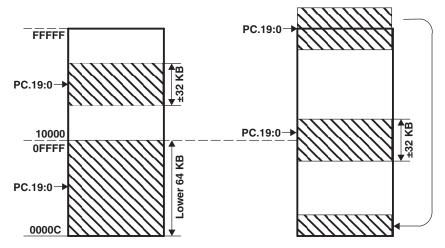


Figure 4-20. Overflow and Underflow for Symbolic Mode

Length: Two or three words

Operation: The sign-extended 16-bit index in the next word after the instruction is added to the

20 bits of the PC. This delivers a 20-bit address, which points to an address in the range 0 to FFFFFh. The operand is the content of the addressed memory location.

Comment: Valid for source and destination. The assembler calculates the PC index and

inserts it

Example: ADD.W EDE,&TONI;

This instruction adds the 16-bit data contained in source word EDE and destination

word TONI and places the 16-bit result into the destination word TONI. For this

example, the instruction is located at address 2F034h.

Source: Word EDE at address 3379Ch, pointed to by PC + 4766h, which is the 16-bit result

of 3379Ch – 2F036h = 04766h. Address 2F036h is the location of the index for this

example.

Destination: Word TONI located at address 00778h pointed to by the absolute address 00778h



Before:	Address Space	l			Address Space	ı	
2F03Ah 2F038h 2F036h 2F034h	xxxxh 0778h 4766h 5092h	PC		2F03Ah 2F038h 2F036h 2F034h	xxxxh 0778h 4766h 5092h	PC	
3379Eh 3379Ch	xxxxh 5432h		2F036h +04766h 3379Ch	3379Eh 3379Ch	xxxxh 5432h		
0077Ah 00778h	xxxxh 2345h			0077Ah 00778h	xxxxh 7777h	5432h <u>+2345h</u> 7777h	src dst Sum

# 4.4.3.3 MSP430X Instruction With Symbolic Mode

When using an MSP430X instruction with Symbolic mode, the operand can be located anywhere in the range of PC + 19 bits.

Length: Three or four words

Operation: The operand address is the sum of the 20-bit PC and the 20-bit index. The 4 MSBs

of the index are contained in the extension word; the 16 LSBs are contained in the

word following the instruction.

Comment: Valid for source and destination. The assembler calculates the register index and

inserts it.

Example: ADDX.B EDE, TONI ;

This instruction adds the 8-bit data contained in source byte EDE and destination

byte TONI and places the result into the destination byte TONI.

Source: Byte EDE located at address 3579Ch, pointed to by PC + 14766h, is the 20-bit

result of 3579Ch - 21036h = 14766h. Address 21036h is the address of the index

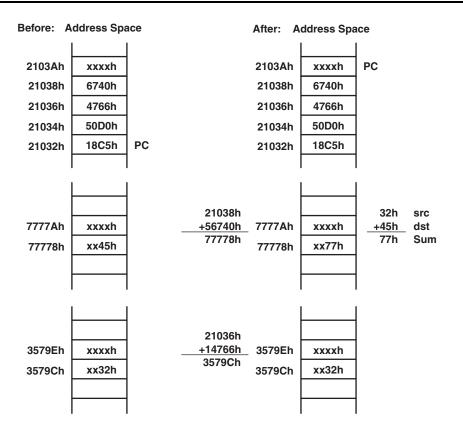
in this example.

Destination: Byte TONI located at address 77778h, pointed to by PC + 56740h, is the 20-bit

result of 77778h – 21038h = 56740h. Address 21038h is the address of the index in

this example.





#### 4.4.4 Absolute Mode

The Absolute mode uses the contents of the word following the instruction as the address of the operand. The Absolute mode has two addressing possibilities:

- · Absolute mode in lower 64-KB memory
- MSP430X instruction with Absolute mode

# 4.4.4.1 Absolute Mode in Lower 64KB

If an MSP430 instruction is used with Absolute addressing mode, the absolute address is a 16-bit value and, therefore, points to an address in the lower 64KB of the memory range. The address is calculated as an index from 0 and is stored in the word following the instruction The RAM and the peripheral registers can be accessed this way and existing MSP430 software is usable without modifications.

Length: Two or three words

Operation: The operand is the content of the addressed memory location.

Comment: Valid for source and destination. The assembler calculates the index from 0 and

inserts it.

Example: ADD.W &EDE,&TONI;

This instruction adds the 16-bit data contained in the absolute source and

destination addresses and places the result into the destination.

Source: Word at address EDE

Destination: Word at address TONI



Before: Address Space		After: Ac	ldress Spa	ce		
2103Ah 21038h 21036h 21034h	xxxxh 7778h 579Ch 5292h	PC	2103Ah 21038h 21036h 21034h	xxxxh 7778h 579Ch 5292h	PC	
0777Ah 07778h	xxxxh 2345h		0777Ah 07778h	xxxxh 7777h	5432h +2345h 7777h	src dst Sum
0579Eh 0579Ch	xxxxh 5432h		0579Eh 0579Ch	xxxxh 5432h		

#### 4.4.4.2 MSP430X Instruction With Absolute Mode

If an MSP430X instruction is used with Absolute addressing mode, the absolute address is a 20-bit value and, therefore, points to any address in the memory range. The address value is calculated as an index from 0. The 4 MSBs of the index are contained in the extension word, and the 16 LSBs are contained in the word following the instruction.

Length: Three or four words

Operation: The operand is the content of the addressed memory location.

Comment: Valid for source and destination. The assembler calculates the index from 0 and

inserts it.

Example: ADDX.A &EDE,&TONI ;

This instruction adds the 20-bit data contained in the absolute source and

destination addresses and places the result into the destination.

Source: Two words beginning with address EDE

Destination: Two words beginning with address TONI



Before:	Address Space	After:		Address Space	I	
2103Ah	xxxxh	210	3 <b>A</b> h	xxxxh	PC	
21038h	7778h	210	38h	7778h		
21036h	579Ch	210	36h	579Ch		
21034h	52D2h	210	34h	52D2h		
21032h	1987h	PC 210	32h	1987h		
7777Ah 77778h	0001h 2345h	777 777		0007h 7777h	65432h <u>+12345h</u> 77777h	src dst Sum
3579Eh 3579Ch	0006h 5432h	357 <sup>9</sup>		0006h 5432h		

# 4.4.5 Indirect Register Mode

The Indirect Register mode uses the contents of the CPU register Rsrc as the source operand. The Indirect Register mode always uses a 20-bit address.

Length: One, two, or three words

Operation: The operand is the content the addressed memory location. The source register

Rsrc is not modified.

Comment: Valid only for the source operand. The substitute for the destination operand is

0(Rdst).

Example: ADDX.W @R5,2100h(R6)

This instruction adds the two 16-bit operands contained in the source and the

destination addresses and places the result into the destination.

Source: Word pointed to by R5. R5 contains address 3579Ch for this example.

Destination: Word pointed to by R6 + 2100h, which results in address 45678h + 2100h = 7778h



Before:	Address Space	Register	After:	Address Space	Register
21038h 21036h 21034h	2100h 55A6h	R5 3579Ch R6 45678h PC	21038h 21036h 21034h	2100h 55A6h	PC R5 3579Ch R6 45678h
4777Ah 47778h	xxxxh 2345h	45678h <u>+02100h</u> 47778h	4777Ah 47778h	xxxxh 7777h	5432h src +2345h dst 7777h Sum
3579Eh 3579Ch	xxxxh 5432h	R5	3579Eh 3579Ch	xxxxh 5432h	R5

#### 4.4.6 Indirect Autoincrement Mode

The Indirect Autoincrement mode uses the contents of the CPU register Rsrc as the source operand. Rsrc is then automatically incremented by 1 for byte instructions, by 2 for word instructions, and by 4 for address-word instructions immediately after accessing the source operand. If the same register is used for source and destination, it contains the incremented address for the destination access. Indirect Autoincrement mode always uses 20-bit addresses.

Length: One, two, or three words

Operation: The operand is the content of the addressed memory location.

Comment: Valid only for the source operand

Example: ADD.B @R5+,0(R6)

This instruction adds the 8-bit data contained in the source and the destination

addresses and places the result into the destination.

Source: Byte pointed to by R5. R5 contains address 3579Ch for this example.

Destination: Byte pointed to by R6 + 0h, which results in address 0778h for this example



Before:	Address Space	Register	After:	Address Space	Register
21038h 21036h 21034h	xxxxh 0000h 55F6h	R5 3579Ch R6 00778h PC	21038h 21036h 21034h	xxxxh 0000h 55F6h	PC R5 3579Dh R6 00778h
0077Ah 00778h	xxxxh xx45h	00778h +0000h 00778h	0077Ah 00778h	xxxxh xx77h	32h src <u>+45h</u> dst 77h Sum
3579Dh 3579Ch	xxh 32h	R5	3579Dh 3579Ch	xxh xx32h	R5

#### 4.4.7 Immediate Mode

The Immediate mode allows accessing constants as operands by including the constant in the memory location following the instruction. The PC is used with the Indirect Autoincrement mode. The PC points to the immediate value contained in the next word. After the fetching of the immediate operand, the PC is incremented by 2 for byte, word, or address-word instructions. The Immediate mode has two addressing possibilities:

- 8-bit or 16-bit constants with MSP430 instructions
- 20-bit constants with MSP430X instruction

#### 4.4.7.1 MSP430 Instructions With Immediate Mode

If an MSP430 instruction is used with Immediate addressing mode, the constant is an 8- or 16-bit value and is stored in the word following the instruction.

Length: Two or three words. One word less if a constant of the constant generator can be

used for the immediate operand.

Operation: The 16-bit immediate source operand is used together with the 16-bit destination

operand.

Comment: Valid only for the source operand

Example: ADD #3456h,&TONI

This instruction adds the 16-bit immediate operand 3456h to the data in the

destination address TONI.

Source: 16-bit immediate value 3456h

Destination: Word at address TONI



Before:	Address Space	Af	ter:	Address Space	ı	
2103Ah 21038h 21036h 21034h	xxxxh 0778h 3456h 50B2h	:	2103Ah 21038h 21036h 21034h	xxxxh 0778h 3456h 50B2h	PC	
0077Ah 00778h	xxxxh 2345h		0077Ah 00778h	xxxxh 579Bh	3456h +2345h 579Bh	src dst Sum

# 4.4.7.2 MSP430X Instructions With Immediate Mode

If an MSP430X instruction is used with Immediate addressing mode, the constant is a 20-bit value. The 4 MSBs of the constant are stored in the extension word, and the 16 LSBs of the constant are stored in the word following the instruction.

Length: Three or four words. One word less if a constant of the constant generator can be

used for the immediate operand.

Operation: The 20-bit immediate source operand is used together with the 20-bit destination

operand.

Comment: Valid only for the source operand

Example: ADDX.A #23456h,&TONI ;

This instruction adds the 20-bit immediate operand 23456h to the data in the

destination address TONI.

Source: 20-bit immediate value 23456h

Destination: Two words beginning with address TONI

Before:	Address	A	fter:	Address		
	Space			Space		
2103Ah	xxxxh		2103Ah	xxxxh	PC	
21038h	7778h		21038h	7778h		
21036h	3456h		21036h	3456h		
21034h	50F2h		21034h	50F2h		
21032h	1907h	PC	21032h	1907h		
	1					
					23456h	src
7777Ah	0001h		7777Ah	0003h	+12345h	dst
77778h	2345h		77778h	579Bh	3579Bh	Sum



# 4.5 MSP430 and MSP430X Instructions

MSP430 instructions are the 27 implemented instructions of the MSP430 CPU. These instructions are used throughout the 1MB memory range unless their 16-bit capability is exceeded. The MSP430X instructions are used when the addressing of the operands, or the data length exceeds the 16-bit capability of the MSP430 instructions.

There are three possibilities when choosing between an MSP430 and MSP430X instruction:

- To use only the MSP430 instructions The only exceptions are the CALLA and the RETA instruction. This can be done if a few, simple rules are met:
  - Placement of all constants, variables, arrays, tables, and data in the lower 64KB. This allows the
    use of MSP430 instructions with 16-bit addressing for all data accesses. No pointers with 20-bit
    addresses are needed.
  - Placement of subroutine constants immediately after the subroutine code. This allows the use of the symbolic addressing mode with its 16-bit index to reach addresses within the range of PC + 32KB.
- To use only MSP430X instructions The disadvantages of this method are the reduced speed due to the additional CPU cycles and the increased program space due to the necessary extension word for any double operand instruction.
- Use the best fitting instruction where needed.

The following sections list and describe the MSP430 and MSP430X instructions.

#### 4.5.1 MSP430 Instructions

The MSP430 instructions can be used, regardless if the program resides in the lower 64KB or beyond it. The only exceptions are the instructions CALL and RET, which are limited to the lower 64-KB address range. CALLA and RETA instructions have been added to the MSP430X CPU to handle subroutines in the entire address range with no code size overhead.

#### 4.5.1.1 MSP430 Double-Operand (Format I) Instructions

Figure 4-21 shows the format of the MSP430 double-operand instructions. Source and destination words are appended for the Indexed, Symbolic, Absolute, and Immediate modes. Table 4-4 lists the 12 MSP430 double-operand instructions.

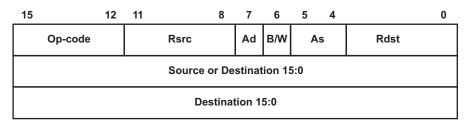


Figure 4-21. MSP430 Double-Operand Instruction Format



Table 4-4. MSP430 Double-Operand Instructions

Mnemonic	S-Reg, D- Reg	Operation	Status Bits <sup>(1)</sup>					
			٧	N	Z	С		
MOV(.B)	src,dst	src → dst	-	-	-	-		
ADD(.B)	src,dst	src + dst → dst	*	*	*	*		
ADDC(.B)	src,dst	$src + dst + C \rightarrow dst$	*	*	*	*		
SUB(.B)	src,dst	dst + .not.src + 1 → dst	*	*	*	*		
SUBC(.B)	src,dst	$dst + .not.src + C \rightarrow dst$	*	*	*	*		
CMP(.B)	src,dst	dst → src	*	*	*	*		
DADD(.B)	src,dst	$src + dst + C \rightarrow dst (decimally)$	*	*	*	*		
BIT(.B)	src,dst	src .and. dst	0	*	*	Z		
BIC(.B)	src,dst	.not.src .and. dst → dst	-	-	-	-		
BIS(.B)	src,dst	$src.or. dst \rightarrow dst$	-	-	-	-		
XOR(.B)	src,dst	$src.xor. dst \rightarrow dst$	*	*	*	Z		
AND(.B)	src,dst	src .and. dst → dst	0	*	*	Z		

<sup>(1) \* =</sup> Status bit is affected.

# 4.5.1.2 MSP430 Single-Operand (Format II) Instructions

Figure 4-22 shows the format for MSP430 single-operand instructions, except RETI. The destination word is appended for the Indexed, Symbolic, Absolute, and Immediate modes. Table 4-5 lists the seven single-operand instructions.

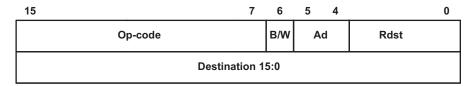


Figure 4-22. MSP430 Single-Operand Instructions

Table 4-5. MSP430 Single-Operand Instructions

Mnemonic	S-Reg, D- Reg	Operation	Status Bits (1)			
Willemonic		Operation	V	N	Z	С
RRC(.B)	dst	$C \to MSB \to \ldots LSB \to C$	*	*	*	*
RRA(.B)	dst	$MSB \to MSB \to LSB \to C$	0	*	*	*
PUSH(.B)	src	$SP - 2 \rightarrow SP$ , $src \rightarrow SP$	-	_	-	_
SWPB	dst	bit 15bit 8 ↔ bit 7bit 0	_	_	_	_
CALL	dst	Call subroutine in lower 64KB	_	_	-	_
RETI		$TOS \rightarrow SR, SP + 2 \rightarrow SP$	*	*	*	*
		$TOS \to PC, SP + 2 \to SP$				
SXT	dst	Register mode: bit 7 → bit 8bit 19 Other modes: bit 7 → bit 8bit 15	0	*	*	Z

<sup>\* =</sup> Status bit is affected.

<sup>- =</sup> Status bit is not affected.

<sup>0 =</sup> Status bit is cleared.

<sup>1 =</sup> Status bit is set.

<sup>-</sup> = Status bit is not affected.

<sup>0 =</sup> Status bit is cleared.

<sup>1 =</sup> Status bit is set.



# 4.5.1.3 Jump Instructions

Figure 4-23 shows the format for MSP430 and MSP430X jump instructions. The signed 10-bit word offset of the jump instruction is multiplied by two, sign-extended to a 20-bit address, and added to the 20-bit PC. This allows jumps in a range of -511 to +512 words relative to the PC in the full 20-bit address space. Jumps do not affect the status bits. Table 4-6 lists and describes the eight jump instructions.

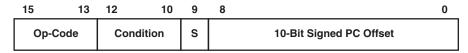


Figure 4-23. Format of Conditional Jump Instructions

**Table 4-6. Conditional Jump Instructions** 

Mnemonic	S-Reg, D-Reg	Operation
JEQ/JZ	Label	Jump to label if zero bit is set
JNE/JNZ	Label	Jump to label if zero bit is reset
JC	Label	Jump to label if carry bit is set
JNC	Label	Jump to label if carry bit is reset
JN	Label	Jump to label if negative bit is set
JGE	Label	Jump to label if (N .XOR. V) = 0
JL	Label	Jump to label if (N .XOR. V) = 1
JMP	Label	Jump to label unconditionally

#### 4.5.1.4 Emulated Instructions

In addition to the MSP430 and MSP430X instructions, emulated instructions are instructions that make code easier to write and read, but do not have op-codes themselves. Instead, they are replaced automatically by the assembler with a core instruction. There is no code or performance penalty for using emulated instructions. The emulated instructions are listed in Table 4-7.

Table 4-7. Emulated Instructions

Instruction	Explanation	Emulation	Status Bits (1)				
instruction		Elliulation	٧	N	Z	С	
ADC(.B) dst	Add Carry to dst	ADDC(.B) #0,dst	*	*	*	*	
BR dst	Branch indirectly dst	MOV dst,PC	_	_	_	_	
CLR(.B) dst	Clear dst	MOV(.B) #0,dst	1	_	_	I	
CLRC	Clear Carry bit	BIC #1,SR	_	_	-	0	
CLRN	Clear Negative bit	BIC #4,SR	-	0	_	ı	
CLRZ	Clear Zero bit	BIC #2,SR	-	_	0	ı	
DADC(.B) dst	Add Carry to dst decimally	DADD(.B) #0,dst	*	*	*	*	
DEC(.B) dst	Decrement dst by 1	SUB(.B) #1,dst	*	*	*	*	
DECD(.B) dst	Decrement dst by 2	SUB(.B) #2,dst	*	*	*	*	
DINT	Disable interrupt	BIC #8,SR	-	-	-	_	
EINT	Enable interrupt	BIS #8,SR	-	-	-	_	
INC(.B) dst	Increment dst by 1	ADD(.B) #1,dst	*	*	*	*	
INCD(.B) dst	Increment dst by 2	ADD(.B) #2,dst	*	*	*	*	

<sup>\* =</sup> Status bit is affected; - = Status bit is not affected; 0 = Status bit is cleared; 1 = Status bit is set.



Instruction	Evalenation	Emulation	Status Bits (1)					
instruction	Explanation	Emulation	٧	N	Z	С		
INV(.B) dst	Invert dst	XOR(.B) #-1,dst	*	*	*	*		
NOP	No operation	MOV R3,R3	-	-	-	_		
POP dst	Pop operand from stack	MOV @SP+,dst	-	-	-	_		
RET	Return from subroutine	MOV @SP+,PC	-	-	-	-		
RLA(.B) dst	Shift left dst arithmetically	ADD(.B) dst,dst	*	*	*	*		
RLC(.B) dst	Shift left dst logically through Carry	ADDC(.B) dst,dst	*	*	*	*		
SBC(.B) dst	Subtract Carry from dst	SUBC(.B) #0,dst	*	*	*	*		
SETC	Set Carry bit	BIS #1,SR	-	-	_	1		
SETN	Set Negative bit	BIS #4,SR	-	1	-	_		
SETZ	Set Zero bit	BIS #2,SR	-	-	1	-		
TST(.B) dst	Test dst (compare with 0)	CMP(.B) #0,dst	0	*	*	1		

Table 4-7. Emulated Instructions (continued)

#### 4.5.1.5 MSP430 Instruction Execution

The number of CPU clock cycles required for an instruction depends on the instruction format and the addressing modes used – not the instruction itself. The number of clock cycles refers to MCLK.

### 4.5.1.5.1 Instruction Cycles and Length for Interrupt, Reset, and Subroutines

Table 4-8 lists the length and the CPU cycles for reset, interrupts, and subroutines.

Table 4-8. Interrupt, Return, and Reset Cycles and Length

Action	Execution Time (MCLK Cycles)	Length of Instruction (Words)
Return from interrupt RETI	3 (1)	1
Return from subroutine RET	3	1
Interrupt request service (cycles needed before first instruction)	5 (2)	_
WDT reset	4	_
Reset ( RST/NMI)	4	_

<sup>(1)</sup> The cycle count in MSP430 CPU is 5.

### 4.5.1.5.2 Format II (Single-Operand) Instruction Cycles and Lengths

Table 4-9 lists the length and the CPU cycles for all addressing modes of the MSP430 single-operand instructions.

Table 4-9. MSP430 Format II Instruction Cycles and Length

	No.	of Cycles		Length of			
Addressing Mode	RRA, RRC SWPB, SXT	PUSH	CALL	Instruction	Example		
Rn	1	3	3 (1)	1	SWPB R5		
@Rn	3	3 (1)	4	1	RRC @R9		
@Rn+	3	3 (1)	4 (2)	1	SWPB @R10+		

<sup>(1)</sup> The cycle count in MSP430 CPU is 4.

The cycle count in MSP430 CPU is 6.

The cycle count in MSP430 CPU is 5. Also, the cycle count is 5 for X(Rn) addressing mode, when Rn = SP.



Table 4-9. MSP430 Format II Instruction Cycles and Length (continued)

	No.	of Cycles		Length of	Example	
Addressing Mode	RRA, RRC SWPB, SXT	PUSH	CALL	Instruction		
#N	N/A	3 (1)	4 (2)	2	CALL #LABEL	
X(Rn)	4	4 (2)	4 (2)	2	CALL 2(R7)	
EDE	4	4 (2)	4 (2)	2	PUSH EDE	
&EDE	4	4 (2)	4 (2)	2	SXT &EDE	

# 4.5.1.5.3 Jump Instructions Cycles and Lengths

All jump instructions require one code word and take two CPU cycles to execute, regardless of whether the jump is taken or not.

### 4.5.1.5.4 Format I (Double-Operand) Instruction Cycles and Lengths

Table 4-10 lists the length and CPU cycles for all addressing modes of the MSP430 Format I instructions.

Table 4-10. MSP430 Format I Instructions Cycles and Length

Address	Addressing Mode		Length of	Example			
Src	Dst	No. of Cycles	Instruction	Example			
	Rm	1	1	MOV	R5,R8		
	PC	2	1	BR	R9		
Rn	x(Rm)	4 <sup>(1)</sup>	2	ADD	R5,4(R6)		
	EDE	4 <sup>(1)</sup>	2	XOR	R8,EDE		
	&EDE	4 <sup>(1)</sup>	2	MOV	R5,&EDE		
	Rm	2	1	AND	@R4,R5		
	PC	3	1	BR	@R8		
@Rn	x(Rm)	5 <sup>(1)</sup>	2	XOR	@R5,8(R6)		
	EDE	5 <sup>(1)</sup>	2	MOV	@R5,EDE		
	&EDE	5 <sup>(1)</sup>	2	XOR	@R5,&EDE		
	Rm	2	1	ADD	@R5+,R6		
	PC	3	1	BR	@R9+		
@Rn+	x(Rm)	5 <sup>(1)</sup>	2	XOR	@R5,8(R6)		
	EDE	5 <sup>(1)</sup>	2	MOV	@R9+,EDE		
	&EDE	5 <sup>(1)</sup>	2	MOV	@R9+,&EDE		
	Rm	2	2	MOV	#20,R9		
	PC	3	2	BR	#2AEh		
#N	x(Rm)	5 <sup>(1)</sup>	3	MOV	#0300h,0(SP)		
	EDE	5 <sup>(1)</sup>	3	ADD	#33,EDE		
	&EDE	5 <sup>(1)</sup>	3	ADD	#33,&EDE		
	Rm	3	2	MOV	2(R5),R7		
	PC	3	2	BR	2(R6)		
x(Rn)	TONI	6 <sup>(1)</sup>	3	MOV	4(R7),TONI		
	x(Rm)	6 <sup>(1)</sup>	3	ADD	4(R4),6(R9)		
	&TONI	6 <sup>(1)</sup>	3	MOV	2(R4),&TONI		

<sup>(1)</sup> MOV, BIT, and CMP instructions execute in one fewer cycle.



Address	ing Mode	No. of Cycles	Length of	Furanta		
Src	Dst	No. of Cycles	Instruction		Example	
	Rm	3	2	AND	EDE,R6	
	PC	3	2	BR	EDE	
EDE	TONI	6 <sup>(1)</sup>	3	CMP	EDE,TONI	
	x(Rm)	6 <sup>(1)</sup>	3	MOV	EDE,0(SP)	
	&TONI	6 <sup>(1)</sup>	3	MOV	EDE,&TONI	
	Rm	3	2	MOV	&EDE,R8	
	PC	3	2	BR	&EDE	
&EDE	TONI	6 <sup>(1)</sup>	3	MOV	&EDE,TONI	
	x(Rm)	6 <sup>(1)</sup>	3	MOV	&EDE,0(SP)	
	&TONI	6 <sup>(1)</sup>	3	MOV	&EDE,&TONI	

Table 4-10. MSP430 Format I Instructions Cycles and Length (continued)

#### 4.5.2 MSP430X Extended Instructions

The extended MSP430X instructions give the MSP430X CPU full access to its 20-bit address space. Most MSP430X instructions require an additional word of op-code called the extension word. Some extended instructions do not require an additional word and are noted in the instruction description. All addresses, indexes, and immediate numbers have 20-bit values when preceded by the extension word.

There are two types of extension words:

- Register/register mode for Format I instructions and register mode for Format II instructions
- · Extension word for all other address mode combinations

### 4.5.2.1 Register Mode Extension Word

The register mode extension word is shown in Figure 4-24 and described in Table 4-11. An example is shown in Figure 4-26.

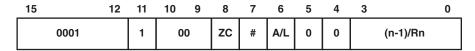


Figure 4-24. Extension Word for Register Modes

Table 4-11. Description of the Extension Word Bits for Register Mode

Bit	Desc	Description				
15:11	Exte	Extension word op-code. Op-codes 1800h to 1FFFh are extension words.				
10:9	Rese	Reserved				
ZC	Zero	Zero carry				
	0	The executed instruction uses the status of the carry bit C.				
	1	The executed instruction uses the carry bit as 0. The carry bit is defined by the result of the final operation after instruction execution.				
#	Repe	etition				
	0	The number of instruction repetitions is set by extension word bits 3:0.				
	1	The number of instruction repetitions is defined by the value of the four LSBs of Rn. See description for bits 3:0.				



# Table 4-11. Description of the Extension Word Bits for Register Mode (continued)

Bit	Descri	Description							
A/L		Data length extension. Together with the B/W bits of the following MSP430 instruction, the AL bit defines the used data length of the instruction.							
	A/L	A/L B/W Comment							
	0	0 0 Reserved							
	0 1 20-bit address word								
	1	0	16-bit word						
	1	1	8-bit byte						
5:4	Reserv	ed							
3:0	Repetit	ion cou	nt						
	# = 0	#=0 These four bits set the repetition count n. These bits contain $n-1$ .							
	# = 1	These	four bits define the CPU register whose bits $3:0$ set the number of repetitions. Rn. $3:0$ contain $n-1$ .						

148



# 4.5.2.2 Non-Register Mode Extension Word

The extension word for non-register modes is shown in Figure 4-25 and described in Table 4-12. An example is shown in Figure 4-27.



Figure 4-25. Extension Word for Non-Register Modes

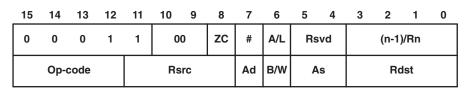
Table 4-12. Description of Extension Word Bits for Non-Register Modes

Bit	Desc	Description						
15:11	Exten	Extension word op-code. Op-codes 1800h to 1FFFh are extension words.						
Source Bits 19:16		The four MSBs of the 20-bit source. Depending on the source addressing mode, these four MSBs may belong to an immediate operand, an index or to an absolute address.						
A/L		Data length extension. Together with the B/W bits of the following MSP430 instruction, the AL bit defines the used data length of the instruction.						
	A/L	B/W	Comment					
	0	0	Reserved					
	0	1	20-bit address word					
	1	0	16-bit word					
	1	1	8-bit byte					
5:4	Rese	Reserved						
Destination Bits 19:16		The four MSBs of the 20-bit destination. Depending on the destination addressing mode, these four MSBs may elong to an index or to an absolute address.						

### NOTE: B/W and A/L bit settings for SWPBX and SXTX

A/L	B/W	
0	0	SWPBX.A, SXTX.A
0	1	N/A
1	0	SWPB.W, SXTX.W
1	1	N/A





XORX.A R9,R8

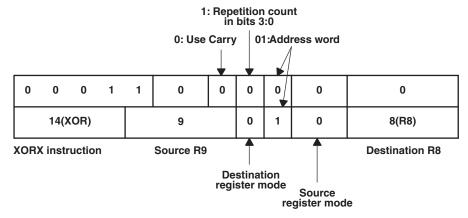
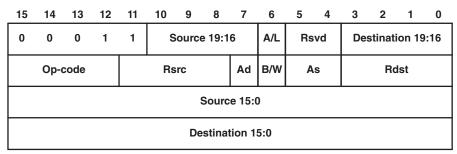


Figure 4-26. Example for Extended Register/Register Instruction



XORX.A #12345h, 45678h(R15)

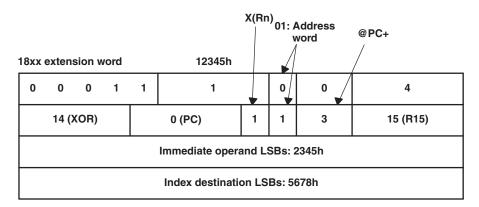


Figure 4-27. Example for Extended Immediate/Indexed Instruction



#### **Extended Double-Operand (Format I) Instructions** 4.5.2.3

All 12 double-operand instructions have extended versions as listed in Table 4-13.

**Table 4-13. Extended Double-Operand Instructions** 

Managaria	On a nom do	Onevetion		Status	Bits (1)	
Mnemonic	Operands	Operation	٧	N	Z	С
MOVX(.B,.A)	src,dst	src → dst	-	_	-	-
ADDX(.B,.A)	src,dst	src + dst → dst	*	*	*	*
ADDCX(.B,.A)	src,dst	$src + dst + C \rightarrow dst$	*	*	*	*
SUBX(.B,.A)	src,dst	dst + .not.src + 1 → dst	*	*	*	*
SUBCX(.B,.A)	src,dst	$dst + .not.src + C \rightarrow dst$	*	*	*	*
CMPX(.B,.A)	src,dst	dst - src	*	*	*	*
DADDX(.B,.A)	src,dst	$src + dst + C \rightarrow dst (decimal)$	*	*	*	*
BITX(.B,.A)	src,dst	src .and. dst	0	*	*	Z
BICX(.B,.A)	src,dst	.not.src .and. $dst \rightarrow dst$	-	_	-	-
BISX(.B,.A)	src,dst	$src.or. dst \rightarrow dst$	-	_	-	-
XORX(.B,.A)	src,dst	$src.xor. dst \rightarrow dst$	*	*	*	Z
ANDX(.B,.A)	src,dst	$src.and. dst \rightarrow dst$	0	*	*	Z

<sup>\* =</sup> Status bit is affected.

- = Status bit is not affected.

<sup>0 =</sup> Status bit is cleared.

<sup>1 =</sup> Status bit is set.



The four possible addressing combinations for the extension word for Format I instructions are shown in Figure 4-28.

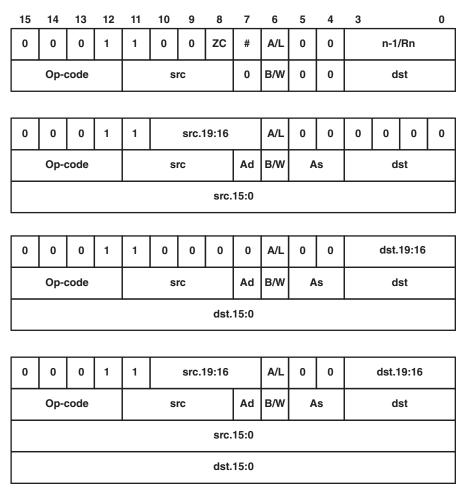


Figure 4-28. Extended Format I Instruction Formats

If the 20-bit address of a source or destination operand is located in memory, not in a CPU register, then two words are used for this operand as shown in Figure 4-29.

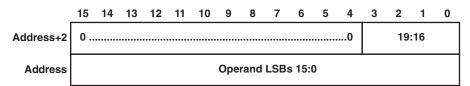


Figure 4-29. 20-Bit Addresses in Memory



# 4.5.2.4 Extended Single-Operand (Format II) Instructions

Extended MSP430X Format II instructions are listed in Table 4-14.

Table 4-14. Extended Single-Operand Instructions

Mnomonio	Onerende	Operation		Status Bits (1)				
Mnemonic	Operands	Operation	n	٧	N	Z	С	
CALLA	dst	Call indirect to subroutine (20-bit address)		-	-	_	-	
POPM.A	#n,Rdst	Pop n 20-bit registers from stack	1 to 16	-	-	_	-	
POPM.W	#n,Rdst	Pop n 16-bit registers from stack	1 to 16	-	-	_	-	
PUSHM.A	#n,Rsrc	Push n 20-bit registers to stack	1 to 16	-	-	_	-	
PUSHM.W	#n,Rsrc	Push n 16-bit registers to stack	1 to 16	-	ı	_	-	
PUSHX(.B,.A)	src	Push 8/16/20-bit source to stack		-	ı	_	-	
RRCM(.A)	#n,Rdst	Rotate right Rdst n bits through carry (16-/20-bit register)	1 to 4	0	*	*	*	
RRUM(.A)	#n,Rdst	Rotate right Rdst n bits unsigned (16-/20-bit register)	1 to 4	0	*	*	*	
RRAM(.A)	#n,Rdst	Rotate right Rdst n bits arithmetically (16-/20-bit register)	1 to 4	*	*	*	*	
RLAM(.A)	#n,Rdst	Rotate left Rdst n bits arithmetically (16-/20-bit register)	1 to 4	*	*	*	*	
RRCX(.B,.A)	dst	Rotate right dst through carry (8-/16-/20-bit data)	1	0	*	*	*	
RRUX(.B,.A)	Rdst	Rotate right dst unsigned (8-/16-/20-bit)	1	0	*	*	*	
RRAX(.B,.A)	dst	Rotate right dst arithmetically	1	*	*	*	*	
SWPBX(.A)	dst	Exchange low byte with high byte	1	-	-	_	-	
SXTX(.A)	Rdst	Bit7 → bit8 bit19	1	0	*	*	*	
SXTX(.A)	dst	Bit7 → bit8 MSB	1	0	*	*	*	

<sup>(1) \* =</sup> Status bit is affected.

The three possible addressing mode combinations for Format II instructions are shown in Figure 4-30.

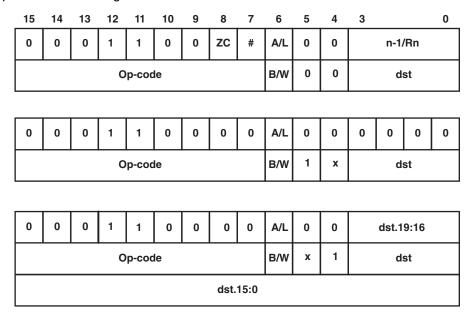


Figure 4-30. Extended Format II Instruction Format

<sup>-=</sup> Status bit is not affected.

<sup>0 =</sup> Status bit is cleared.

<sup>1 =</sup> Status bit is set.



# 4.5.2.4.1 Extended Format II Instruction Format Exceptions

Exceptions for the Format II instruction formats are shown in Figure 4-31 through Figure 4-34.

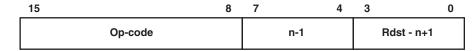


Figure 4-31. PUSHM/POPM Instruction Format



Figure 4-32. RRCM, RRAM, RRUM, and RLAM Instruction Format

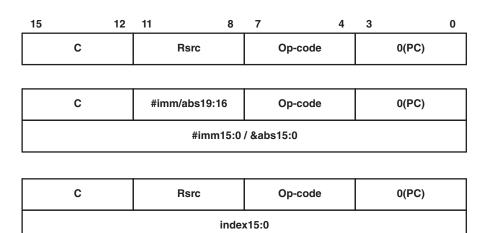


Figure 4-33. BRA Instruction Format

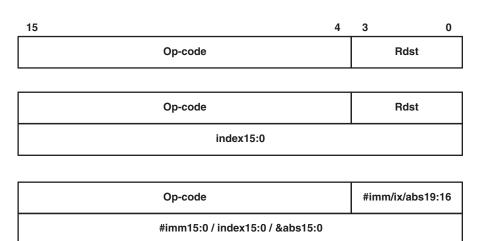


Figure 4-34. CALLA Instruction Format



# 4.5.2.5 Extended Emulated Instructions

The extended instructions together with the constant generator form the extended emulated instructions. Table 4-15 lists the emulated instructions.

**Table 4-15. Extended Emulated Instructions** 

Instruction	Explanation	Emulation
ADCX(.B,.A) dst	Add carry to dst	ADDCX(.B,.A) #0,dst
BRA dst	Branch indirect dst	MOVA dst,PC
RETA	Return from subroutine	MOVA @SP+,PC
CLRA Rdst	Clear Rdst	MOV #0,Rdst
CLRX(.B,.A) dst	Clear dst	MOVX(.B,.A) #0,dst
DADCX(.B,.A) dst	Add carry to dst decimally	DADDX(.B,.A) #0,dst
DECX(.B,.A) dst	Decrement dst by 1	SUBX(.B,.A) #1,dst
DECDA Rdst	Decrement Rdst by 2	SUBA #2,Rdst
DECDX(.B,.A) dst	Decrement dst by 2	SUBX(.B,.A) #2,dst
INCX(.B,.A) dst	Increment dst by 1	ADDX(.B,.A) #1,dst
INCDA Rdst	Increment Rdst by 2	ADDA #2,Rdst
INCDX(.B,.A) dst	Increment dst by 2	ADDX(.B,.A) #2,dst
INVX(.B,.A) dst	Invert dst	XORX(.B,.A) #-1,dst
RLAX(.B,.A) dst	Shift left dst arithmetically	ADDX(.B,.A) dst,dst
RLCX(.B,.A) dst	Shift left dst logically through carry	ADDCX(.B,.A) dst,dst
SBCX(.B,.A) dst	Subtract carry from dst	SUBCX(.B,.A) #0,dst
TSTA Rdst	Test Rdst (compare with 0)	CMPA #0,Rdst
TSTX(.B,.A) dst	Test dst (compare with 0)	CMPX(.B,.A) #0,dst
POPX dst	Pop to dst	MOVX(.B, .A) @SP+,dst



#### 4.5.2.6 MSP430X Address Instructions

MSP430X address instructions are instructions that support 20-bit operands but have restricted addressing modes. The addressing modes are restricted to the Register mode and the Immediate mode, except for the MOVA instruction as listed in Table 4-16. Restricting the addressing modes removes the need for the additional extension-word op-code improving code density and execution time. Address instructions should be used any time an MSP430X instruction is needed with the corresponding restricted addressing mode.

Table 4-16. Address Instructions, Operate on 20-Bit Register Data

Mnemonic	Onevende	Operation	S	tatus	Bits (1)	
winemonic	Operands	Operation	٧	N	Z	С
ADDA	Rsrc,Rdst	Add source to destination register	*	*	*	*
ADDA	#imm20,Rdst	Add source to destination register				
	Rsrc,Rdst					
	#imm20,Rdst					
	z16(Rsrc),Rdst					
	EDE,Rdst					
MOVA	&abs20,Rdst	Move source to destination	_	_	-	- 1
	@Rsrc,Rdst					
	@Rsrc+,Rdst					
	Rsrc,z16(Rdst)					
	Rsrc,&abs20					
CMPA	Rsrc,Rdst	Compare course to destination register	*	*	*	*
CMPA	#imm20,Rdst	Compare source to destination register				
SUBA	Rsrc,Rdst	Culturat course from destination register	*	*	*	*
SUDA	#imm20,Rdst	Subtract source from destination register				

<sup>&</sup>lt;sup>1)</sup> \* = Status bit is affected.

<sup>-=</sup> Status bit is not affected.

<sup>0 =</sup> Status bit is cleared.

<sup>1 =</sup> Status bit is set.



#### 4.5.2.7 MSP430X Instruction Execution

The number of CPU clock cycles required for an MSP430X instruction depends on the instruction format and the addressing modes used, not the instruction itself. The number of clock cycles refers to MCLK.

# 4.5.2.7.1 MSP430X Format II (Single-Operand) Instruction Cycles and Lengths

Table 4-17 lists the length and the CPU cycles for all addressing modes of the MSP430X extended single-operand instructions.

Table 4-17. MSP430X Format II Instruction Cycles and Length

Instruction		Execution	Cycles/Ler	ngth of Inst	ruction (Wo	rds)	
instruction	Rn	@Rn	@Rn+	#N	X(Rn)	EDE	&EDE
RRAM	n/1	_	_	-	_	-	_
RRCM	n/1	-	_	-	_	_	_
RRUM	n/1	-	_	-	_	_	_
RLAM	n/1	-	_	-	-	-	-
PUSHM	2+n/1	_	_	-	_	-	-
PUSHM.A	2+2n/1	-	_	-	_	_	_
POPM	2+n/1	-	_	-	_	_	_
POPM.A	2+2n/1	-	_	-	_	_	_
CALLA	4/1	5/1	5/1	4/2	6 <sup>(1)</sup> /2	6/2	6/2
RRAX(.B)	1+n/2	4/2	4/2	-	5/3	5/3	5/3
RRAX.A	1+n/2	6/2	6/2	-	7/3	7/3	7/3
RRCX(.B)	1+n/2	4/2	4/2	-	5/3	5/3	5/3
RRCX.A	1+n/2	6/2	6/2	-	7/3	7/3	7/3
PUSHX(.B)	4/2	4/2	4/2	4/3	5 <sup>(1)</sup> /3	5/3	5/3
PUSHX.A	5/2	6/2	6/2	6/3	7 (1)/3	7/3	7/3
POPX(.B)	3/2	-	-	-	5/3	5/3	5/3
POPX.A	4/2	_	_	_	7/3	7/3	7/3

<sup>(1)</sup> Add one cycle when Rn = SP



# 4.5.2.7.2 MSP430X Format I (Double-Operand) Instruction Cycles and Lengths

Table 4-18 lists the length and CPU cycles for all addressing modes of the MSP430X extended Format I instructions.

Table 4-18. MSP430X Format I Instruction Cycles and Length

Address	sing Mode	No. o	of Cycles	Length of Instruction	Examples
Source	Destination	.B/.W	.A	.B/.W/.A	
Rn	Rm <sup>(1)</sup>	2	2	2	BITX.B R5,R8
	PC	3	3	2	ADDX R9,PC
	X(Rm)	5 <sup>(2)</sup>	7 <sup>(3)</sup>	3	ANDX.A R5,4(R6)
	EDE	5 <sup>(2)</sup>	7 <sup>(3)</sup>	3	XORX R8,EDE
	&EDE	5 <sup>(2)</sup>	7 <sup>(3)</sup>	3	BITX.W R5,&EDE
@Rn	Rm	3	4	2	BITX @R5,R8
	PC	3	4	2	ADDX @R9,PC
	X(Rm)	6 <sup>(2)</sup>	9(3)	3	ANDX.A @R5,4(R6)
	EDE	6 <sup>(2)</sup>	9(3)	3	XORX @R8,EDE
	&EDE	6 <sup>(2)</sup>	9(3)	3	BITX.B @R5,&EDE
@Rn+	Rm	3	4	2	BITX @R5+,R8
	PC	4	5	2	ADDX.A @R9+,PC
	X(Rm)	6 <sup>(2)</sup>	9(3)	3	ANDX @R5+,4(R6)
	EDE	6 <sup>(2)</sup>	9(3)	3	XORX.B @R8+,EDE
	&EDE	6 <sup>(2)</sup>	9(3)	3	BITX @R5+,&EDE
#N	Rm	3	3	3	BITX #20,R8
	PC <sup>(4)</sup>	4	4	3	ADDX.A #FE000h,PC
	X(Rm)	6 <sup>(2)</sup>	8 <sup>(3)</sup>	4	ANDX #1234,4(R6)
	EDE	6 <sup>(2)</sup>	8(3)	4	XORX #A5A5h,EDE
	&EDE	6 <sup>(2)</sup>	8(3)	4	BITX.B #12,&EDE
X(Rn)	Rm	4	5	3	BITX 2(R5),R8
	PC <sup>(4)</sup>	5	6	3	SUBX.A 2(R6),PC
	X(Rm)	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	ANDX 4(R7),4(R6)
	EDE	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	XORX.B 2(R6),EDE
	&EDE	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	BITX 8(SP),&EDE
EDE	Rm	4	5	3	BITX.B EDE,R8
	PC <sup>(4)</sup>	5	6	3	ADDX.A EDE,PC
	X(Rm)	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	ANDX EDE,4(R6)
	EDE	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	ANDX EDE,TONI
	&TONI	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	BITX EDE,&TONI
&EDE	Rm	4	5	3	BITX &EDE,R8
	PC <sup>(4)</sup>	5	6	3	ADDX.A &EDE,PC
	X(Rm)	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	ANDX.B &EDE,4(R6)
	TONI	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	XORX &EDE,TONI
	&TONI	7 <sup>(2)</sup>	10 <sup>(3)</sup>	4	BITX &EDE,&TONI

<sup>(1)</sup> Repeat instructions require n + 1 cycles, where n is the number of times the instruction is executed.

<sup>(2)</sup> Reduce the cycle count by one for MOV, BIT, and CMP instructions.

<sup>(3)</sup> Reduce the cycle count by two for MOV, BIT, and CMP instructions.

<sup>(4)</sup> Reduce the cycle count by one for MOV, ADD, and SUB instructions.



# 4.5.2.7.3 MSP430X Address Instruction Cycles and Lengths

Table 4-19 lists the length and the CPU cycles for all addressing modes of the MSP430X address instructions.

Table 4-19. Address Instruction Cycles and Length

Addr	essing Mode		on Time Cycles)		Instruction ords)	
Source	Destination	MOVA BRA	CMPA ADDA SUBA	MOVA	CMPA ADDA SUBA	Example
Rn	Rn	1	1	1	1	CMPA R5,R8
	PC	2	2	1	1	SUBA R9,PC
	x(Rm)	4	_	2	-	MOVA R5,4(R6)
	EDE	4	_	2	-	MOVA R8,EDE
	&EDE	4	_	2	-	MOVA R5,&EDE
@Rn	Rm	3	_	1	-	MOVA @R5,R8
	PC	3	_	1	_	MOVA @R9,PC
@Rn+	Rm	3	_	1	_	MOVA @R5+,R8
	PC	3	_	1	-	MOVA @R9+,PC
#N	Rm	2	3	2	2	CMPA #20,R8
	PC	3	3	2	2	SUBA #FE000h,PC
x(Rn)	Rm	4	_	2	-	MOVA 2(R5),R8
	PC	4	_	2	-	MOVA 2(R6),PC
EDE	Rm	4	_	2	-	MOVA EDE, R8
	PC	4	_	2	-	MOVA EDE, PC
&EDE	Rm	4	_	2	_	MOVA &EDE, R8
	PC	4	_	2	-	MOVA &EDE,PC



# 4.6 Instruction Set Description

Table 4-20 shows all available instructions:

# Table 4-20. Instruction Map of MSP430X

	000	040	080	0C0	100	140	180	1C0	200	240	280	2C0	300	340	380	3C0
0xxx					MOVA,	CMPA, A	DDA, SL	JBA, RRO	CM, RRA	M, RLAM	, RRUM					
10xx	RRC	RRC.B	SWPB		RRA	RRA.B	SXT		PUSH	PUSH. B	CALL		RETI	CALL A		
14xx			ļ	Į.	Į.	PUS	SHM.A, F	OPM.A,	PUSHM.	W, POPI	Л.W	Į.	ļ		Į.	
18xx						Evtonoio	a word fo	r Format	Land Fo	rmat II in	atructions					
1Cxx	Extension word for Format I and Format II instructions															
20xx	JNE/JNZ															
24xx	JEQ/JZ															
28xx	JNC															
2Cxx		JC														
30xx		JN														
34xx	JGE															
38xx									L							
3Cxx								JN	ЛP							
4xxx								MOV,	MOV.B							
5xxx									ADD.B							
6xxx									ADDC.B							
7xxx								SUBC,	SUBC.B							
8xxx									SUB.B							
9xxx									CMP.B							
Axxx									DADD.B							
Bxxx									BIT.B							
Cxxx									BIC.B							
Dxxx									BIS.B							
Exxx									XOR.B							
Fxxx								AND,	AND.B							



# 4.6.1 Extended Instruction Binary Descriptions

Detailed MSP430X instruction binary descriptions are shown in the following tables.

Instruction	ı		uction	n	src or data.19:16	ı	nstru Iden	ıctio tifier		dst	
	15			12	11 8	7			4	3 0	=
MOVA	0	0	0	0	src	0	0	0	0	dst	MOVA @Rsrc,Rdst
	0	0	0	0	src	0	0	0	1	dst	MOVA @Rsrc+,Rdst
	0	0	0	0	&abs.19:16	0	0	1	0	dst	MOVA &abs20,Rdst
					&abs	s.15:0	)				
	0	0	0	0	src	0	0	1	1	dst	MOVA x(Rsrc),Rdst
					x.1	5:0					±15-bit index x
	0	0	0	0	src	0	1	1	0	&abs.19:16	MOVA Rsrc,&abs20
					&abs	s.15:0	)				
	0	0	0	0	src	0	1	1	1	dst	MOVA Rsrc,X(Rdst)
					x.1	5:0					±15-bit index x
	0	0	0	0	imm.19:16	1	0	0	0	dst	MOVA #imm20,Rdst
					imm	.15:0					
CMPA	0	0	0	0	imm.19:16	1	0	0	1	dst	CMPA #imm20,Rdst
					imm	.15:0					
ADDA	0	0	0	0	imm.19:16	1	0	1	0	dst	ADDA #imm20,Rdst
					imm	.15:0					
SUBA	0	0	0	0	imm.19:16	1	0	1	1	dst	SUBA #imm20,Rdst
					imm	.15:0					
MOVA	0	0	0	0	src	1	1	0	0	dst	MOVA Rsrc,Rdst
CMPA	0	0	0	0	src	1	1	0	1	dst	CMPA Rsrc,Rdst
ADDA	0	0	0	0	src	1	1	1	0	dst	ADDA Rsrc,Rdst
SUBA	0	0	0	0	src	1	1	1	1	dst	SUBA Rsrc,Rdst

Instruction	Instruction Group			Bit Loc.	Ins	t. ID	Instruction Identifier			dst				
	15			12	11 10	9	8	7			4	3	0	
RRCM.A	0	0	0	0	n – 1	0	0	0	1	0	0	dst		RRCM.A #n,Rdst
RRAM.A	0	0	0	0	n – 1	0	1	0	1	0	0	dst		RRAM.A #n,Rdst
RLAM.A	0	0	0	0	n – 1	1	0	0	1	0	0	dst		RLAM.A #n,Rdst
RRUM.A	0	0	0	0	n – 1	1	1	0	1	0	0	dst		RRUM.A #n,Rdst
RRCM.W	0	0	0	0	n – 1	0	0	0	1	0	1	dst		RRCM.W #n,Rdst
RRAM.W	0	0	0	0	n – 1	0	1	0	1	0	1	dst		RRAM.W #n,Rdst
RLAM.W	0	0	0	0	n – 1	1	0	0	1	0	1	dst		RLAM.W #n,Rdst
RRUM.W	0	0	0	0	n – 1	1	1	0	1	0	1	dst		RRUM.W #n,Rdst



Instruction			ı	nstru	ıctioı	n Ide	ntifie	r						d	st		
instruction	15			12	11			8	7	6	5	4	3			0	
RETI	0	0	0	1	0	0	1	1	0	0	0	0	0	0	0	0	
CALLA	0	0	0	1	0	0	1	1	0	1	0	0		d	st		CALLA Rdst
	0	0	0	1	0	0	1	1	0	1	0	1		d	st		CALLA x(Rdst)
			l	l			l	x.1	5:0	l							
	0	0	0	1	0	0	1	1	0	1	1	0		d	st		CALLA @Rdst
	0	0	0	1	0	0	1	1	0	1	1	1		d	st		CALLA @Rdst+
	0	0	0	1	0	0	1	1	1	0	0	0	8	&abs	19:16	3	CALLA &abs20
								&abs	.15:0								
	0	0	0	1	0	0	1	1	1	0	0	1		x.19	9:16		CALLA EDE
							ļ.	x.1	5:0								CALLA x(PC)
	0	0	0	1	0	0	1	1	1	0	1	1		imm.	19:16	;	CALLA #imm20
			I	I	I	I	I	imm	.15:0	I		I					
Reserved	0	0	0	1	0	0	1	1	1	0	1	0	Х	Х	Х	х	
Reserved	0	0	0	1	0	0	1	1	1	1	Х	Х	х	Х	Х	Х	
PUSHM.A	0	0	0	1	0	1	0	0		n -	- 1			d	st		PUSHM.A #n,Rdst
PUSHM.W	0	0	0	1	0	1	0	1		n -	n – 1		dst				PUSHM.W #n,Rdst
POPM.A	0	0	0	1	0	1	1	0		n -	- 1		dst - n + 1				POPM.A #n,Rdst
POPM.W	0	0	0	1	0	1	1	1		n -	- 1		(	dst –	n + 1		POPM.W #n,Rdst



# 4.6.2 MSP430 Instructions

The MSP430 instructions are described in the following sections.

See Section 4.6.3 for MSP430X extended instructions and Section 4.6.4 for MSP430X address instructions.



#### 4.6.2.1 ADC

\* ADC[.W] Add carry to destination
\* ADC.B Add carry to destination

Syntax ADC dst Or ADC.W dst

ADC.B dst

Operation  $dst + C \rightarrow dst$ Emulation ADDC #0, dst ADDC.B #0, dst

**Description** The carry bit (C) is added to the destination operand. The previous contents of the

destination are lost.

**Status Bits** N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if dst was incremented from 0FFFh to 0000, reset otherwise Set if dst was incremented from 0FFh to 00, reset otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 16-bit counter pointed to by R13 is added to a 32-bit counter pointed to by R12.

ADD @R13,0(R12) ; Add LSDs

ADC 2(R12) ; Add carry to MSD

**Example** The 8-bit counter pointed to by R13 is added to a 16-bit counter pointed to by R12.

ADD.B @R13,0(R12) ; Add LSDs

ADC.B 1(R12) ; Add carry to MSD



#### 4.6.2.2 ADD

ADD[.W] Add source word to destination word ADD.B Add source byte to destination byte **Syntax** ADD src,dst Of ADD.W src,dst

ADD.B src,dst

Operation  $src + dst \rightarrow dst$ 

. . .

Description The source operand is added to the destination operand. The previous content of the

destination is lost.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

> Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

Set if the result of two positive operands is negative, or if the result of two negative

numbers is positive, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Ten is added to the 16-bit counter CNTR located in lower 64KB.

ADD.W #10,&CNTR ; Add 10 to 16-bit counter

**Example** A table word pointed to by R5 (20-bit address in R5) is added to R6. The jump to label

TONI is performed on a carry.

; Add table word to R6. R6.19:16 = 0ADD.W @R5,R6 JC TONI ; Jump if carry ; No carry

A table byte pointed to by R5 (20-bit address) is added to R6. The jump to label TONI is **Example** 

performed if no carry occurs. The table pointer is auto-incremented by 1. R6.19:8 = 0

; Add byte to R6. R5 + 1. R6: 000xxh ADD.B @R5+,R6 JNC TONI ; Jump if no carry ; Carry occurred . . .



#### 4.6.2.3 ADDC

ADDC.W Add source word and carry to destination word ADDC.B Add source byte and carry to destination byte

Syntax ADDC src,dst Of ADDC.W src,dst

ADDC.B src,dst

**Operation**  $\operatorname{src} + \operatorname{dst} + \operatorname{C} \to \operatorname{dst}$ 

**Description** The source operand and the carry bit C are added to the destination operand. The

previous content of the destination is lost.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative

numbers is positive, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Constant value 15 and the carry of the previous instruction are added to the 16-bit

counter CNTR located in lower 64KB.

ADDC.W #15,&CNTR ; Add 15 + C to 16-bit CNTR

**Example** A table word pointed to by R5 (20-bit address) and the carry C are added to R6. The

jump to label TONI is performed on a carry. R6.19:16 = 0

ADDC.W @R5,R6 ; Add table word + C to R6

JC TONI ; Jump if carry ... ; No carry

**Example** A table byte pointed to by R5 (20-bit address) and the carry bit C are added to R6. The

jump to label TONI is performed if no carry occurs. The table pointer is auto-incremented

by 1. R6.19:8 = 0

ADDC.B @R5+,R6 ; Add table byte + C to R6. R5 + 1

JNC TONI ; Jump if no carry ... ; Carry occurred



#### 4.6.2.4 AND

AND[.W] Logical AND of source word with destination word AND.B Logical AND of source byte with destination byte

Syntax AND src,dst Or AND.W src,dst

AND.B src,dst

**Operation** src .and.  $dst \rightarrow dst$ 

**Description** The source operand and the destination operand are logically ANDed. The result is

placed into the destination. The source operand is not affected.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if the result is not zero, reset otherwise. C = (.not. Z)

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The bits set in R5 (16-bit data) are used as a mask (AA55h) for the word TOM located in

the lower 64KB. If the result is zero, a branch is taken to label TONI. R5.19:16 = 0

MOV #AA55h,R5 ; Load 16-bit mask to R5
AND R5,&TOM ; TOM .and. R5 -> TOM
JZ TONI ; Jump if result 0
... ; Result > 0

or shorter:

AND #AA55h,&TOM ; TOM .and. AA55h -> TOM JZ TONI ; Jump if result 0

**Example** A table byte pointed to by R5 (20-bit address) is logically ANDed with R6. R5 is

incremented by 1 after the fetching of the byte. R6.19:8 = 0

AND.B @R5+,R6 ; AND table byte with R6. R5 + 1



### 4.6.2.5 BIC

BIC[.W] Clear bits set in source word in destination word Clear bits set in source byte in destination byte

Syntax BIC src,dst Of BIC.W src,dst

BIC.B src,dst

**Operation** (.not. src) .and.  $dst \rightarrow dst$ 

**Description** The inverted source operand and the destination operand are logically ANDed. The

result is placed into the destination. The source operand is not affected.

Status Bits N: Not affected

Z: Not affectedC: Not affectedV: Not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The bits 15:14 of R5 (16-bit data) are cleared. R5.19:16 = 0

BIC #0C000h,R5 ; Clear R5.19:14 bits

**Example** A table word pointed to by R5 (20-bit address) is used to clear bits in R7. R7.19:16 = 0

BIC.W @R5,R7 ; Clear bits in R7 set in @R5

**Example** A table byte pointed to by R5 (20-bit address) is used to clear bits in Port1.

BIC.B @R5,&P1OUT ; Clear I/O port P1 bits set in @R5



#### 4.6.2.6 BIS

BIS[.W] Set bits set in source word in destination word BIS.B Set bits set in source byte in destination byte

Syntax BIS src,dst Or BIS.W src,dst

BIS.B src,dst

**Description** The source operand and the destination operand are logically ORed. The result is placed

into the destination. The source operand is not affected.

Status Bits N: Not affected

Z: Not affectedC: Not affectedV: Not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Bits 15 and 13 of R5 (16-bit data) are set to one. R5.19:16 = 0

BIS #A000h,R5 ; Set R5 bits

**Example** A table word pointed to by R5 (20-bit address) is used to set bits in R7. R7.19:16 = 0

BIS.W @R5,R7 ; Set bits in R7

**Example** A table byte pointed to by R5 (20-bit address) is used to set bits in Port1. R5 is

incremented by 1 afterwards.

BIS.B @R5+,&P1OUT ; Set I/O port P1 bits. R5 + 1



#### 4.6.2.7 BIT

BIT[.W] Test bits set in source word in destination word
Test bits set in source byte in destination byte

Syntax BIT src,dst Or BIT.W src,dst

BIT.B src,dst

**Operation** src .and. dst

**Description** The source operand and the destination operand are logically ANDed. The result affects

only the status bits in SR.

Register mode: the register bits Rdst.19:16 (.W) resp. Rdst. 19:8 (.B) are not cleared!

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if the result is not zero, reset otherwise. C = (.not. Z)

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Test if one (or both) of bits 15 and 14 of R5 (16-bit data) is set. Jump to label TONI if this

is the case. R5.19:16 are not affected.

BIT #C000h,R5 ; Test R5.15:14 bits

JNZ TONI ; At least one bit is set in R5

... ; Both bits are reset

**Example** A table word pointed to by R5 (20-bit address) is used to test bits in R7. Jump to label

TONI if at least one bit is set. R7.19:16 are not affected.

BIT.W @R5,R7 ; Test bits in R7

JC TONI ; At least one bit is set

... ; Both are reset

**Example** A table byte pointed to by R5 (20-bit address) is used to test bits in output Port1. Jump

to label TONI if no bit is set. The next table byte is addressed.

BIT.B @R5+,&P1OUT ; Test I/O port P1 bits. R5 + 1
JNC TONI ; No corresponding bit is set
... ; At least one bit is set



#### 4.6.2.8 BR, BRANCH

\* BR, Branch to destination in lower 64K address space

BRANCH

**Description** An unconditional branch is taken to an address anywhere in the lower 64K address

space. All source addressing modes can be used. The branch instruction is a word

instruction.

Status Bits Status bits are not affected.

**Example** Examples for all addressing modes are given.

BR ; Branch to label EXEC or direct branch (e.g. #0A4h) ; Core instruction MOV @PC+,PC EXEC ; Branch to the address contained in EXEC BR ; Core instruction MOV X(PC),PC ; Indirect address ; Branch to the address contained in absolute BR &EXEC ; address EXEC ; Core instruction MOV X(0),PC ; Indirect address ; Branch to the address contained in R5 BRR5 ; Core instruction MOV R5,PC ; Indirect R5 BR @R 5 ; Branch to the address contained in the word ; pointed to by R5.

; Core instruction MOV @R5,PC

; Indirect, indirect R5

BR @R5+ ; Branch to the address contained in the word pointed

; to by R5 and increment pointer in R5 afterwards.

; The next time-S/W flow uses R5 pointer-it can

; alter program execution due to access to  $% \left( 1\right) =\left( 1\right) \left( 1\right) \left($ 

; next address in a table pointed to by R5

; Core instruction MOV @R5,PC

; Indirect, indirect R5 with autoincrement

BR X(R5); Branch to the address contained in the address

; pointed to by R5 +  $\rm X$  (e.g. table with address

; starting at X). X can be an address or a label

; Core instruction MOV X(R5),PC

; Indirect, indirect R5 + X

171



#### 4.6.2.9 CALL

CALL Call a subroutine in lower 64KB

Syntax CALL dst

**Operation**  $dst \rightarrow PC$  16-bit dst is evaluated and stored

 $SP - 2 \rightarrow SP$ 

 $PC \rightarrow @SP$  updated PC with return address to TOS

tmp → PC saved 16-bit dst to PC

**Description** A subroutine call is made from an address in the lower 64KB to a subroutine address in

the lower 64KB. All seven source addressing modes can be used. The call instruction is

a word instruction. The return is made with the RET instruction.

**Status Bits** Status bits are not affected.

PC.19:16 cleared (address in lower 64KB)

Mode Bits OSCOFF, CPUOFF, and GIE are not affected. Examples Examples for all addressing modes are given.

Immediate Mode: Call a subroutine at label EXEC (lower 64KB) or call directly to

address.

CALL #EXEC ; Start address EXEC
CALL #0AA04h ; Start address 0AA04h

Symbolic Mode: Call a subroutine at the 16-bit address contained in address EXEC.

EXEC is located at the address (PC + X) where X is within PC + 32 K.

CALL EXEC ; Start address at @EXEC. z16(PC)

Absolute Mode: Call a subroutine at the 16-bit address contained in absolute address

EXEC in the lower 64KB.

CALL &EXEC ; Start address at @EXEC

Register mode: Call a subroutine at the 16-bit address contained in register R5.15:0.

CALL R5 ; Start address at R5

Indirect Mode: Call a subroutine at the 16-bit address contained in the word pointed to by

register R5 (20-bit address).

CALL @R5 ; Start address at @R5



# 4.6.2.10 CLR

\* CLR[.W] Clear destination
\* CLR.B Clear destination

Syntax CLR dst Or CLR.W dst

CLR.B dst

**Description** The destination operand is cleared.

Status Bits Status bits are not affected.

Example RAM word TONI is cleared.

CLR TONI ; 0 -> TONI

**Example** Register R5 is cleared.

CLR R5

**Example** RAM byte TONI is cleared.

CLR.B TONI ; 0 -> TONI

173



### 4.6.2.11 CLRC

\* CLRC Clear carry bit

**Description** The carry bit (C) is cleared. The clear carry instruction is a word instruction.

Status Bits N: Not affected

Z: Not affectedC: ClearedV: Not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 16-bit decimal counter pointed to by R13 is added to a 32-bit counter pointed to by

R12.

CLRC ; C=0: defines start

DADD @R13,0(R12) ; add 16-bit counter to low word of 32-bit counter

DADC 2(R12) ; add carry to high word of 32-bit counter



#### 4.6.2.12 CLRN

\* CLRN Clear negative bit

or

(.NOT.src .AND.  $dst \rightarrow dst$ )

Emulation BIC #4,SR

**Description** The constant 04h is inverted (0FFFBh) and is logically ANDed with the destination

operand. The result is placed into the destination. The clear negative bit instruction is a

word instruction.

Status Bits N: Reset to 0

Z: Not affectedC: Not affectedV: Not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The negative bit in the SR is cleared. This avoids special treatment with negative

numbers of the subroutine called.

CLRN
CALL SUBR
...
...
JN SUBRET ; If input is negative: do nothing and return
...
...

SUBRET RET

SUBR

175



#### 4.6.2.13 CLRZ

\* CLRZ Clear zero bit

or

(.NOT.src .AND.  $dst \rightarrow dst$ )

Emulation BIC #2,SR

**Description** The constant 02h is inverted (0FFFDh) and logically ANDed with the destination

operand. The result is placed into the destination. The clear zero bit instruction is a word

instruction.

Status Bits N: Not affected

Z: Reset to 0C: Not affectedV: Not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The zero bit in the SR is cleared.

CLRZ

Indirect, Auto-Increment mode: Call a subroutine at the 16-bit address contained in the word pointed to by register R5 (20-bit address) and increment the 16-bit address in R5 afterwards by 2. The next time the software uses R5 as a pointer, it can alter the program execution due to access to the next word address in the table pointed to by R5.

CALL @R5+ ; Start address at @R5. R5 + 2

Indexed mode: Call a subroutine at the 16-bit address contained in the 20-bit address pointed to by register (R5 + X), for example, a table with addresses starting at X. The address is within the lower 64KB. X is within +32KB.

CALL X(R5) ; Start address at @(R5+X). z16(R5)



#### 4.6.2.14 CMP

CMP[.W] Compare source word and destination word CMP.B Compare source byte and destination byte

**Syntax** CMP src,dst Of CMP.W src,dst

CMP.B src,dst

Operation (.not.src) + 1 + dst

or

dst - src

**Emulation** BIC #2,SR

Description

The source operand is subtracted from the destination operand. This is made by adding the 1s complement of the source + 1 to the destination. The result affects only the status

bits in SR.

Register mode: the register bits Rdst.19:16 (.W) resp. Rdst. 19:8 (.B) are not cleared.

**Status Bits** 

Set if result is negative (src > dst), reset if positive (src = dst)

Z: Set if result is zero (src = dst), reset otherwise (src  $\neq$  dst)

C: Set if there is a carry from the MSB, reset otherwise

Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow).

**Mode Bits** 

OSCOFF, CPUOFF, and GIE are not affected.

Example

Compare word EDE with a 16-bit constant 1800h. Jump to label TONI if EDE equals the constant. The address of EDE is within PC + 32 K.

CMP #01800h, EDE ; Compare word EDE with 1800h TONI ; EDE contains 1800h JEQ

; Not equal . . .

Example

A table word pointed to by (R5 + 10) is compared with R7. Jump to label TONI if R7 contains a lower, signed 16-bit number. R7.19:16 is not cleared. The address of the source operand is a 20-bit address in full memory range.

```
CMP.W
         10(R5),R7
                            ; Compare two signed numbers
JL
         TONI
                            ; R7 < 10(R5)
                            ; R7 >= 10(R5)
. . .
```

#### **Example**

A table byte pointed to by R5 (20-bit address) is compared to the value in output Port1. Jump to label TONI if values are equal. The next table byte is addressed.

```
CMP.B
         @R5+,&P10UT
                           ; Compare P1 bits with table. R5 + 1
JEO
         TONI
                            ; Equal contents
                            ; Not equal
. . .
```



### 4.6.2.15 DADC

\* DADC[.W] Add carry decimally to destination

\* DADC.B Add carry decimally to destination

Syntax DADC dst or DADC.W dst

DADC.B dst

**Operation**  $dst + C \rightarrow dst (decimally)$ 

Emulation DADD #0,dst DADD.B #0,dst

**Description** The carry bit (C) is added decimally to the destination.

Status Bits N: Set if MSB is 1

Z: Set if dst is 0, reset otherwise

C: Set if destination increments from 9999 to 0000, reset otherwise Set if destination increments from 99 to 00, reset otherwise

V: Undefined

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The four-digit decimal number contained in R5 is added to an eight-digit decimal number

pointed to by R8.

CLRC ; Reset carry

; next instruction's start condition is defined

**Example** The two-digit decimal number contained in R5 is added to a four-digit decimal number

pointed to by R8.

CLRC ; Reset carry

; next instruction's start condition is defined

DADD.B R5,0(R8) ; Add LSDs + C DADC 1(R8) ; Add carry to MSDs



#### 4.6.2.16 DADD

\* **DADD[.W]** Add source word and carry decimally to destination word \* **DADD.B** Add source byte and carry decimally to destination byte

Syntax DADD src,dst Or DADD.W src,dst

DADD.B src,dst

**Operation**  $\operatorname{src} + \operatorname{dst} + \operatorname{C} \to \operatorname{dst} (\operatorname{decimally})$ 

**Description** The source operand and the destination operand are treated as two (.B) or four (.W)

binary coded decimals (BCD) with positive signs. The source operand and the carry bit C are added decimally to the destination operand. The source operand is not affected. The

previous content of the destination is lost. The result is not defined for non-BCD

numbers.

Status Bits N: Set if MSB of result is 1 (word > 7999h, byte > 79h), reset if MSB is 0

Z: Set if result is zero, reset otherwise

C: Set if the BCD result is too large (word > 9999h, byte > 99h), reset otherwise

V: Undefined

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** Decimal 10 is added to the 16-bit BCD counter DECCNTR.

```
DADD #10h, &DECCNTR ; Add 10 to 4-digit BCD counter
```

#### **Example**

The eight-digit BCD number contained in 16-bit RAM addresses BCD and BCD+2 is added decimally to an eight-digit BCD number contained in R4 and R5 (BCD+2 and R5 contain the MSDs). The carry C is added, and cleared.

```
CLRC ; Clear carry
DADD.W &BCD,R4 ; Add LSDs. R4.19:16 = 0
DADD.W &BCD+2,R5 ; Add MSDs with carry. R5.19:16 = 0
JC OVERFLOW ; Result >9999,9999: go to error routine
... ; Result ok
```

#### Example

The two-digit BCD number contained in word BCD (16-bit address) is added decimally to a two-digit BCD number contained in R4. The carry C is added, also. R4.19:8 = 0CLRC; Clear carryDADD.B &BCD,R4; Add BCD to R4 decimally. R4: 0,00ddh

```
CLRC ; Clear carry
DADD.B &BCD,R4 ; Add BCD to R4 decimally.
R4: 0,00ddh
```



#### 4.6.2.17 DEC

\* **DEC[.W]** Decrement destination \* **DEC.B** Decrement destination

Syntax DEC dst Or DEC.W dst

DEC.B dst

Operation  $dst - 1 \rightarrow dst$ Emulation SUB #1, dst SUB.B #1, dst

**Description** The destination operand is decremented by one. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 1, reset otherwiseC: Reset if dst contained 0, set otherwise

V: Set if an arithmetic overflow occurs, otherwise reset.

Set if initial value of destination was 08000h, otherwise reset.

Set if initial value of destination was 080h, otherwise reset.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** R10 is decremented by 1.

```
DEC R10 ; Decrement R10
```

; Move a block of 255 bytes from memory location starting with EDE to ; memory location starting with TONI. Tables should not overlap: start of ; destination address TONI must not be within the range EDE to EDE+0FEh

```
MOV #EDE,R6

MOV #510,R10

L$1 MOV @R6+,TONI-EDE-1(R6)

DEC R10

JNZ L$1
```

Do not transfer tables using the routine above with the overlap shown in Figure 4-35.

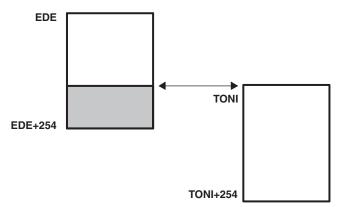


Figure 4-35. Decrement Overlap



#### 4.6.2.18 DECD

\* **DECD[.W]** Double-decrement destination Double-decrement destination

Syntax DECD dst Or DECD.W dst

DECD.B dst

Operation  $dst - 2 \rightarrow dst$ Emulation SUB #2,dst SUB.B #2,dst

**Description** The destination operand is decremented by two. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 2, reset otherwise

C: Reset if dst contained 0 or 1, set otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

Set if initial value of destination was 08001 or 08000h, otherwise reset Set if initial value of destination was 081 or 080h, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** R10 is decremented by 2.

```
DECD R10 ; Decrement R10 by two
```

; Move a block of 255 bytes from memory location starting with EDE to

; memory location starting with TONI.

; Tables should not overlap: start of destination address TONI must not

; be within the range EDE to EDE+0FEh

MOV #EDE,R6 MOV #255,R10

L\$1 MOV.B @R6+,TONI-EDE-2(R6)

DECD R10 JNZ L\$1

## **Example** Memory at location LEO is decremented by two.

DECD.B LEO ; Decrement MEM(LEO)

## Decrement status byte STATUS by two

DECD.B STATUS



#### 4.6.2.19 DINT

\* **DINT** Disable (general) interrupts

or

(0FFF7h .AND. SR  $\rightarrow$  SR / .NOT. src .AND. dst  $\rightarrow$  dst)

Emulation BIC #8,SR

**Description** All interrupts are disabled.

The constant 08h is inverted and logically ANDed with the SR. The result is placed into

the SR.

**Status Bits** Status bits are not affected.

Mode Bits GIE is reset. OSCOFF and CPUOFF are not affected.

**Example** The general interrupt enable (GIE) bit in the SR is cleared to allow a nondisrupted move

of a 32-bit counter. This ensures that the counter is not modified during the move by any

interrupt.

DINT ; All interrupt events using the GIE bit are disabled

NOP

MOV COUNTHI, R5 ; Copy counter

MOV COUNTLO, R6

EINT ; All interrupt events using the GIE bit are enabled

#### **NOTE:** Disable interrupt

If any code sequence needs to be protected from interruption, DINT should be executed at least one instruction before the beginning of the uninterruptible sequence, or it should be followed by a NOP instruction.



#### 4.6.2.20 EINT

\* **EINT** Enable (general) interrupts

or

(0008h .OR. SR  $\rightarrow$  SR / .src .OR. dst  $\rightarrow$  dst)

Emulation BIS #8,SR

**Description** All interrupts are enabled.

The constant #08h and the SR are logically ORed. The result is placed into the SR.

**Status Bits** Status bits are not affected.

Mode Bits GIE is set. OSCOFF and CPUOFF are not affected.Example The general interrupt enable (GIE) bit in the SR is set.

```
; Interrupt routine of ports P1.2 to P1.7
```

; P1IN is the address of the register where all port bits are read.

; PlIFG is the address of the register where all interrupt events are latched.

```
PUSH.B
                  &P1IN
         BIC.B
                  @SP,&P1IFG ; Reset only accepted flags
         EINT
                               ; Preset port 1 interrupt flags stored on stack
                               ; other interrupts are allowed
         BIT
                  #Mask,@SP
         JEQ
                  MaskOK
                               ; Flags are present identically to mask: jump
         . . .
MaskOK
         BIC
                  #Mask,@SP
         . . .
         INCD
                  SP
                               ; Housekeeping: inverse to PUSH instruction
                               ; at the start of interrupt subroutine. Corrects
                               ; the stack pointer.
         RETI
```

## NOTE: Enable interrupt

The instruction following the enable interrupt instruction (EINT) is always executed, even if an interrupt service request is pending when the interrupts are enabled.



#### 4.6.2.21 INC

\* INC[.W] Increment destination \* INC.B Increment destination

**Syntax** INC dst Or INC.W dst

INC.B dst

 $dst + 1 \rightarrow dst$ Operation **Emulation** ADD #1,dst

Description The destination operand is incremented by one. The original contents are lost.

**Status Bits** Set if result is negative, reset if positive N:

> Z: Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFh, reset otherwise C: Set if dst contained 0FFFFh, reset otherwise

Set if dst contained 0FFh, reset otherwise Set if dst contained 07FFFh, reset otherwise Set if dst contained 07Fh, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

The status byte, STATUS, of a process is incremented. When it is equal to 11, a branch **Example** 

to OVFL is taken.

INC.B STATUS CMP.B #11,STATUS OVFL

JEQ



#### 4.6.2.22 INCD

\* INCD[.W] Double-increment destination
\* INCD.B Double-increment destination

Syntax INCD dst Or INCD.W dst

INCD.B dst

Operation  $dst + 2 \rightarrow dst$ Emulation ADD #2, dst ADD.B #2, dst

**Description** The destination operand is incremented by two. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFEh, reset otherwise Set if dst contained 0FEh, reset otherwise

C: Set if dst contained 0FFFEh or 0FFFFh, reset otherwise Set if dst contained 0FEh or 0FFh, reset otherwise

V: Set if dst contained 07FFEh or 07FFFh, reset otherwise Set if dst contained 07Eh or 07Fh, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The item on the top of the stack (TOS) is removed without using a register.

PUSH R5 ; R5 is the result of a calculation, which is stored

; in the system stack

INCD SP ; Remove TOS by double-increment from stack

; Do not use INCD.B, SP is a word-aligned register

RET

**Example** The byte on the top of the stack is incremented by two.

INCD.B 0(SP) ; Byte on TOS is increment by two

185



#### 4.6.2.23 INV

\* INV[.W] Invert destination
\* INV.B Invert destination

Syntax INV dst Or INV.W dst

INV.B dst

**Operation** .not.dst  $\rightarrow$  dst **Emulation** XOR #0FFFFh,dst XOR.B #0FFh,dst

**Description** The destination operand is inverted. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFh, reset otherwise Set if dst contained 0FFh, reset otherwise

C: Set if result is not zero, reset otherwise ( = .NOT. Zero)

V: Set if initial destination operand was negative, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected. Example Content of R5 is negated (2s complement).

MOV #00AEh,R5 ; R5 = 000AEh
INV R5 ; Invert R5, R5 = 0FF51h
INC R5 ; R5 is now negated, R5 = 0FF52h

## **Example** Content of memory byte LEO is negated.

MOV.B #0AEh,LEO ; MEM(LEO) = 0AEh
INV.B LEO ; Invert LEO, MEM(LEO) = 051h
INC.B LEO ; MEM(LEO) is negated, MEM(LEO) = 052h



#### 4.6.2.24 JC, JHS

JC Jump if carry

**JHS** Jump if higher or same (unsigned)

Syntax JC label

JHS label

**Operation** If C = 1:  $PC + (2 \times Offset) \rightarrow PC$ 

If C = 0: execute the following instruction

**Description** The carry bit C in the SR is tested. If it is set, the signed 10-bit word offset contained in

the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range -511 to +512 words relative to the PC in the full memory

range. If C is reset, the instruction after the jump is executed.

JC is used for the test of the carry bit C.

JHS is used for the comparison of unsigned numbers.

Status Bits Status bits are not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The state of the port 1 pin P1IN.1 bit defines the program flow.

```
BIT.B #2,&P1IN ; Port 1, bit 1 set? Bit -> C

JC Label1 ; Yes, proceed at Label1

... ; No, continue
```

## **Example** If $R5 \ge R6$ (unsigned), the program continues at Label2.

```
CMP R6,R5 ; Is R5 \Rightarrow R6? Info to C JHS Label2 ; Yes, C = 1 ; No, R5 < R6. Continue
```

## **Example** If $R5 \ge 12345h$ (unsigned operands), the program continues at Label2.

```
CMPA #12345h,R5 ; Is R5 >= 12345h? Info to C

JHS Label2 ; Yes, 12344h < R5 <= F,FFFFh. C = 1

... ; No, R5 < 12345h. Continue
```



#### 4.6.2.25 JEQ, JZ

JEQ Jump if equal
JZ Jump if zero
Syntax JEQ label
JZ label

**Operation** If Z = 1: PC +  $(2 \times Offset) \rightarrow PC$ 

If Z = 0: execute following instruction

**Description** The zero bit Z in the SR is tested. If it is set, the signed 10-bit word offset contained in

the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range –511 to +512 words relative to the PC in the full memory

range. If Z is reset, the instruction after the jump is executed.

JZ is used for the test of the zero bit Z.

JEQ is used for the comparison of operands.

Status Bits Status bits are not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The state of the P2IN.0 bit defines the program flow.

BIT.B #1,&P2IN ; Port 2, bit 0 reset?

JZ Label1 ; Yes, proceed at Label1

... ; No, set, continue

**Example** If R5 = 15000h (20-bit data), the program continues at Label2.

CMPA #15000h,R5 ; Is R5 = 15000h? Info to SR JEQ Label2 ; Yes, R5 = 15000h. Z = 1

... ; No, R5 not equal 15000h. Continue

**Example** R7 (20-bit counter) is incremented. If its content is zero, the program continues at

Label4.

ADDA #1,R7 ; Increment R7

JZ Label4 ; Zero reached: Go to Label4
... ; R7 not equal 0. Continue here.



#### 4.6.2.26 JGE

**JGE** Jump if greater or equal (signed)

Syntax JGE label

**Operation** If (N .xor. V) = 0: PC +  $(2 \times Offset) \rightarrow PC$ 

If  $(N \cdot xor. V) = 1$ : execute following instruction

**Description** The negative bit N and the overflow bit V in the SR are tested. If both bits are set or both

are reset, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range -511 to +512 words relative to the PC in full Memory range. If only one bit is set, the instruction after

the jump is executed.

JGE is used for the comparison of signed operands: also for incorrect results due to

overflow, the decision made by the JGE instruction is correct.

Note: JGE emulates the nonimplemented JP (jump if positive) instruction if used after the

instructions AND, BIT, RRA, SXTX, and TST. These instructions clear the V bit.

Status Bits Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** If byte EDE (lower 64KB) contains positive data, go to Label1. Software can run in the

full memory range.

```
TST.B &EDE ; Is EDE positive? V <- 0

JGE Label1 ; Yes, JGE emulates JP

... ; No. 80h <= EDE <= FFh
```

#### **Example**

If the content of R6 is greater than or equal to the memory pointed to by R7, the program continues a Label5. Signed data. Data and program in full memory range.

## **Example**

If R5  $\geq$  12345h (signed operands), the program continues at Label2. Program in full memory range.

```
CMPA #12345h,R5 ; Is R5 >= 12345h?
JGE Label2 ; Yes, 12344h < R5
```

JGE Label2 ; Yes, 12344h < R5 <= 7FFFFh
... ; No, 80000h <= R5 < 12345h</pre>



#### 4.6.2.27 JL

**JL** Jump if less (signed)

Syntax JL label

**Operation** If (N .xor. V) = 1: PC +  $(2 \times Offset) \rightarrow PC$ 

If  $(N \cdot xor. V) = 0$ : execute following instruction

**Description** The negative bit N and the overflow bit V in the SR are tested. If only one is set, the

signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range -511 to +512 words relative

to the PC in full memory range. If both bits N and V are set or both are reset, the

instruction after the jump is executed.

JL is used for the comparison of signed operands: also for incorrect results due to

overflow, the decision made by the JL instruction is correct.

Status Bits Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** If byte EDE contains a smaller, signed operand than byte TONI, continue at Label1. The

address EDE is within PC ± 32 K.

CMP.B &TONI,EDE ; IS EDE < TONI

JL Label1 ; Yes

... ; No, TONI <= EDE

**Example** If the signed content of R6 is less than the memory pointed to by R7 (20-bit address), the

program continues at Label5. Data and program in full memory range.

**Example** If R5 < 12345h (signed operands), the program continues at Label2. Data and program

in full memory range.

CMPA #12345h,R5 ; Is R5 < 12345h?

JL Label2 ; Yes, 80000h = R5 < 12345h ... ; No, 12344h < R5 <= 7FFFFh



#### 4.6.2.28 JMP

**JMP** Jump unconditionally

JMP label **Syntax** 

Operation  $PC + (2 \times Offset) \rightarrow PC$ 

**Description** The signed 10-bit word offset contained in the instruction is multiplied by two, sign

> extended, and added to the 20-bit PC. This means an unconditional jump in the range -511 to +512 words relative to the PC in the full memory. The JMP instruction may be

used as a BR or BRA instruction within its limited range relative to the PC.

**Status Bits** Status bits are not affected

OSCOFF, CPUOFF, and GIE are not affected. **Mode Bits** 

**Example** The byte STATUS is set to 10. Then a jump to label MAINLOOP is made. Data in lower

64KB, program in full memory range.

#10,&STATUS ; Set STATUS to 10 MOV.B JMP MAINLOOP ; Go to main loop

**Example** The interrupt vector TAIV of Timer A3 is read and used for the program flow. Program in

full memory range, but interrupt handlers always starts in lower 64KB.

; Add Timer\_A interrupt vector to PC ADD &TAIV, PC RETI ; No Timer\_A interrupt pending

JMP IHCCR1 ; Timer block 1 caused interrupt ; Timer block 2 caused interrupt JMP IHCCR2 ; No legal interrupt, return RETI

191



#### 4.6.2.29 JN

**JN** Jump if negative

Syntax JN label

**Operation** If N = 1:  $PC + (2 \times Offset) \rightarrow PC$ 

If N = 0: execute following instruction

**Description** The negative bit N in the SR is tested. If it is set, the signed 10-bit word offset contained

in the instruction is multiplied by two, sign extended, and added to the 20-bit program PC. This means a jump in the range -511 to +512 words relative to the PC in the full

memory range. If N is reset, the instruction after the jump is executed.

**Status Bits** Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The byte COUNT is tested. If it is negative, program execution continues at Label0. Data

in lower 64KB, program in full memory range.

```
TST.B &COUNT ; Is byte COUNT negative?

JN Label0 ; Yes, proceed at Label0
... ; COUNT >= 0
```

**Example** R6 is subt

R6 is subtracted from R5. If the result is negative, program continues at Label2. Program

in full memory range.

#### **Example**

R7 (20-bit counter) is decremented. If its content is below zero, the program continues at Label4. Program in full memory range.

```
SUBA #1,R7 ; Decrement R7 
JN Label4 ; R7 < 0: Go to Label4 
... ; R7 >= 0. Continue here.
```



#### 4.6.2.30 JNC, JLO

**JNC** Jump if no carry

**JLO** Jump if lower (unsigned)

JNC label **Syntax** 

JLO label

Operation If C = 0: PC +  $(2 \times Offset) \rightarrow PC$ 

If C = 1: execute following instruction

Description The carry bit C in the SR is tested. If it is reset, the signed 10-bit word offset contained in

> the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range -511 to +512 words relative to the PC in the full memory

range. If C is set, the instruction after the jump is executed.

JNC is used for the test of the carry bit C.

JLO is used for the comparison of unsigned numbers.

**Status Bits** Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** If byte EDE < 15, the program continues at Label2. Unsigned data. Data in lower 64KB,

program in full memory range.

#15,&EDE ; Is EDE < 15? Info to C CMP.B Label2 ; Yes, EDE < 15. C = 0 JLO ; No, EDE >= 15. Continue . . .

The word TONI is added to R5. If no carry occurs, continue at Label0. The address of Example

TONI is within PC ± 32 K.

ADD TONI,R5 ; TONI + R5 -> R5. Carry -> C

JNC Label0 ; No carry

; Carry = 1: continue here . . .

193



# 4.6.2.31 JNZ, JNE

JNZ Jump if not zero
JNE Jump if not equal

Syntax JNZ label

JNE label

**Operation** If Z = 0: PC +  $(2 \times Offset) \rightarrow PC$ 

If Z = 1: execute following instruction

**Description** The zero bit Z in the SR is tested. If it is reset, the signed 10-bit word offset contained in

the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range –511 to +512 words relative to the PC in the full memory

range. If  $\boldsymbol{Z}$  is set, the instruction after the jump is executed.

JNZ is used for the test of the zero bit Z. JNE is used for the comparison of operands.

Status Bits Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The byte STATUS is tested. If it is not zero, the program continues at Label3. The

address of STATUS is within PC ± 32 K.

TST.B STATUS ; Is STATUS = 0?

JNZ Label3 ; No, proceed at Label3
... ; Yes, continue here

**Example** If word EDE ≠ 1500, the program continues at Label2. Data in lower 64KB, program in

full memory range.

CMP #1500,&EDE ; Is EDE = 1500? Info to SR
JNE Label2 ; No, EDE not equal 1500.
... ; Yes, R5 = 1500. Continue

**Example** R7 (20-bit counter) is decremented. If its content is not zero, the program continues at

Label4. Program in full memory range.

SUBA #1,R7 ; Decrement R7

JNZ Label4 ; Zero not reached: Go to Label4 ... ; Yes, R7 = 0. Continue here.



#### 4.6.2.32 MOV

MOV[.W]Move source word to destination wordMOV.BMove source byte to destination byteSyntaxMOV src,dst or MOV.W src,dst

MOV.B src,dst

**Operation**  $\operatorname{src} \to \operatorname{dst}$ 

**Description** The source operand is copied to the destination. The source operand is not affected.

Status Bits N: Not affected Z: Not affected

C: Not affected V: Not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

Example Move a 16-bit constant 1800h to absolute address-word EDE (lower 64KB)

MOV #01800h, &EDE ; Move 1800h to EDE

#### **Example**

The contents of table EDE (word data, 16-bit addresses) are copied to table TOM. The length of the tables is 030h words. Both tables reside in the lower 64KB.

```
MOV
                #EDE,R10
                                        ; Prepare pointer (16-bit address)
       MOV
               @R10+,TOM-EDE-2(R10)
                                        ; R10 points to both tables.
Loop
                                        ; R10+2
       CMP
                #EDE+60h,R10
                                        ; End of table reached?
       JLO
               Loop
                                        ; Not yet
                                        ; Copy completed
       . . .
```

## Example

The contents of table EDE (byte data, 16-bit addresses) are copied to table TOM. The length of the tables is 020h bytes. Both tables may reside in full memory range, but must be within R10  $\pm$  32 K.

```
#EDE,R10
       MOVA
                                       ; Prepare pointer (20-bit)
               #20h,R9
                                       ; Prepare counter
       MOV
               @R10+,TOM-EDE-1(R10)
                                       ; R10 points to both tables.
good
       MOV.B
                                       ; R10+1
       DEC
               R9
                                       ; Decrement counter
       JNZ
               Loop
                                       ; Not yet done
                                       ; Copy completed
```



# 4.6.2.33 NOP

\* NOP No operation

**Emulation** MOV #0, R3

**Description** No operation is performed. The instruction may be used for the elimination of instructions

during the software check or for defined waiting times.

Status Bits Status bits are not affected.



#### 4.6.2.34 POP

\* POP[.W] Pop word from stack to destination
\* POP.B Pop byte from stack to destination

Syntax POP dst

POP.B dst

 $\textbf{Operation} \qquad @\, \mathsf{SP} \to \mathsf{temp}$ 

 $SP + 2 \rightarrow SP$ temp  $\rightarrow$  dst

Emulation Mov @SP+,dst or Mov.W @SP+,dst

MOV.B @SP+,dst

**Description** The stack location pointed to by the SP (TOS) is moved to the destination. The SP is

incremented by two afterwards.

**Status Bits** Status bits are not affected.

**Example** The contents of R7 and the SR are restored from the stack.

POP R7 ; Restore R7

POP SR ; Restore status register

**Example** The contents of RAM byte LEO is restored from the stack.

POP.B LEO ; The low byte of the stack is moved to LEO.

**Example** The contents of R7 is restored from the stack.

```
POP.B R7; The low byte of the stack is moved to R7, ; the high byte of R7 is 00h
```

**Example** The contents of the memory pointed to by R7 and the SR are restored from the stack.

```
POP.B 0(R7) ; The low byte of the stack is moved to the ; the byte which is pointed to by R7 : Example: R7 = 203h ; Mem(R7) = low byte of system stack : Example: R7 = 20Ah ; Mem(R7) = low byte of system stack POP SR ; Last word on stack moved to the SR
```

## NOTE: System stack pointer

The system SP is always incremented by two, independent of the byte suffix.



#### 4.6.2.35 PUSH

PUSH[.W] Save a word on the stack
PUSH.B Save a byte on the stack

Syntax PUSH dst Or PUSH.W dst

PUSH.B dst

**Operation**  $SP - 2 \rightarrow SP$ 

 $dst \rightarrow @SP$ 

**Description** The 20-bit SP SP is decremented by two. The operand is then copied to the RAM word

addressed by the SP. A pushed byte is stored in the low byte; the high byte is not

affected.

Status Bits Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Save the two 16-bit registers R9 and R10 on the stack

PUSH R9 ; Save R9 and R10 XXXXh

PUSH R10 ; YYYYh

**Example** Save the two bytes EDE and TONI on the stack. The addresses EDE and TONI are

within PC ± 32 K.

PUSH.B EDE ; Save EDE xxXXh PUSH.B TONI ; Save TONI xxYYh



#### 4.6.2.36 RET

**RET** Return from subroutine

Syntax RET

**Operation** @SP  $\rightarrow$  PC.15:0 Saved PC to PC.15:0. PC.19:16  $\leftarrow$  0

 $SP + 2 \rightarrow SP$ 

**Description** The 16-bit return address (lower 64KB), pushed onto the stack by a CALL instruction is

restored to the PC. The program continues at the address following the subroutine call.

The four MSBs of the PC.19:16 are cleared.

Status Bits Status bits are not affected.

PC.19:16: Cleared

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

Example Call a subroutine SUBR in the lower 64KB and return to the address in the lower 64KB

after the CALL.

CALL #SUBR ; Call subroutine starting at SUBR ; Return by RET to here . . . SUBR PUSH R14 ; Save R14 (16 bit data) ; Subroutine code . . . POP R14 ; Restore R14 RET ; Return to lower 64KB

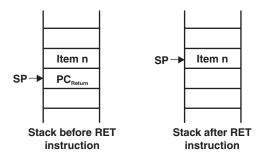


Figure 4-36. Stack After a RET Instruction



#### 4.6.2.37 RETI

**RETI** Return from interrupt

Syntax RETI

**Operation** @SP → SR.15:0 Restore saved SR with PC.19:16

 $SP + 2 \rightarrow SP$ 

 $@SP \rightarrow PC.15:0$  Restore saved PC.15:0

 $SP + 2 \rightarrow SP$  Housekeeping

**Description** The SR is restored to the value at the beginning of the interrupt service routine. This

includes the four MSBs of the PC.19:16. The SP is incremented by two afterward. The 20-bit PC is restored from PC.19:16 (from same stack location as the status bits) and PC.15:0. The 20-bit PC is restored to the value at the beginning of the interrupt service routine. The program continues at the address following the last executed instruction when the interrupt was granted. The SP is incremented by two afterward.

Status Bits N: Restored from stack

C: Restored from stackZ: Restored from stackV: Restored from stack

**Mode Bits** OSCOFF, CPUOFF, and GIE are restored from stack.

**Example** Interrupt handler in the lower 64KB. A 20-bit return address is stored on the stack.

INTRPT PUSHM.A #2,R14 ; Save R14 and R13 (20-bit data)

... ; Interrupt handler code

POPM.A #2,R14 ; Restore R13 and R14 (20-bit data)

RETI ; Return to 20-bit address in full memory range



#### 4.6.2.38 RLA

\* RLA[.W] Rotate left arithmetically
\* RLA.B Rotate left arithmetically

Syntax RLA dst Or RLA.W dst

RLA.B dst

**Operation**  $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow 0$ 

**Emulation** ADD dst,dst ADD.B dst,dst

#### **Description**

The destination operand is shifted left one position as shown in Figure 4-37. The MSB is shifted into the carry bit (C) and the LSB is filled with 0. The RLA instruction acts as a signed multiplication by 2.

An overflow occurs if  $dst \ge 04000h$  and dst < 0C000h before operation is performed; the result has changed sign.



Figure 4-37. Destination Operand—Arithmetic Shift Left

An overflow occurs if  $dst \ge 040h$  and dst < 0C0h before the operation is performed; the result has changed sign.

#### **Status Bits**

- N: Set if result is negative, reset if positive
- Z: Set if result is zero, reset otherwise
- C: Loaded from the MSB
- V: Set if an arithmetic overflow occurs; the initial value is 04000h ≤ dst < 0C000h, reset otherwise

Set if an arithmetic overflow occurs; the initial value is 040h ≤ dst < 0C0h, reset otherwise

# Mode Bits

OSCOFF, CPUOFF, and GIE are not affected.

# Example

R7 is multiplied by 2.

```
RLA R7; Shift left R7 (x 2)
```

#### **Example**

The low byte of R7 is multiplied by 4.

```
RLA.B R7; Shift left low byte of R7 (x 2) RLA.B R7; Shift left low byte of R7 (x 4)
```

#### **NOTE: RLA substitution**

The assembler does not recognize the instructions:

RLA @R5+ RLA.B @R5+ RLA(.B) @R5

They must be substituted by:

ADD @R5+,-2(R5) ADD.B @R5+,-1(R5) ADD(.B) @R5



#### 4.6.2.39 RLC

\* RLC[.W] Rotate left through carry
\* RLC.B Rotate left through carry

Syntax RLC dst Or RLC.W dst

RLC.B dst

**Operation**  $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow C$ 

Emulation ADDC dst,dst

**Description** The destination operand is shifted left one position as shown in Figure 4-38. The carry bit (C) is shifted into the LSB, and the MSB is shifted into the carry bit (C).

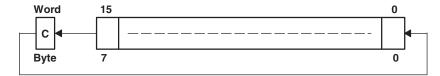


Figure 4-38. Destination Operand—Carry Left Shift

**Status Bits** 

- N: Set if result is negative, reset if positive
- Z: Set if result is zero, reset otherwise
- C: Loaded from the MSB
- V: Set if an arithmetic overflow occurs; the initial value is 04000h ≤ dst < 0C000h, reset otherwise

Set if an arithmetic overflow occurs; the initial value is 040h ≤ dst < 0C0h, reset otherwise

**Mode Bits** 

OSCOFF, CPUOFF, and GIE are not affected.

**Example** 

R5 is shifted left one position.

RLC R5

;  $(R5 \times 2) + C -> R5$ 

#### Example

The input P1IN.1 information is shifted into the LSB of R5.

BIT.B #2,&P1IN RLC R5

; Information -> Carry
; Carry=P0in.1 -> LSB of R5

## **Example**

The MEM(LEO) content is shifted left one position.

RLC.B LEO

;  $Mem(LEO) \times 2 + C \rightarrow Mem(LEO)$ 

#### NOTE: RLA substitution

The assembler does not recognize the instructions:

RLC @R5+

RLC.B @R5+

RLC(.B) @R5

They must be substituted by:

ADDC @R5+,-2(R5)

ADDC.B @R5+,-1(R5)

ADDC(.B) @R5



#### 4.6.2.40 RRA

**RRA[.W]** Rotate right arithmetically destination word RRA.B Rotate right arithmetically destination byte

Syntax RRA.B dst Or RRA.W dst

 $\textbf{Operation} \qquad \text{MSB} \rightarrow \text{MSB} \rightarrow \text{MSB-1} \rightarrow ... \ \text{LSB+1} \rightarrow \text{LSB} \rightarrow \text{C}$ 

**Description**The destination operand is shifted right arithmetically by one bit position as shown in Figure 4-39. The MSB retains its value (sign). RRA operates equal to a signed division by 2. The MSB is retained and shifted into the MSB–1. The LSB+1 is shifted into the

LSB. The previous LSB is shifted into the carry bit C.

**Status Bits** N: Set if result is negative (MSB = 1), reset otherwise (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The signed 16-bit number in R5 is shifted arithmetically right one position.

RRA R5 ; R5/2 -> R5

**Example** The signed RAM byte EDE is shifted arithmetically right one position.

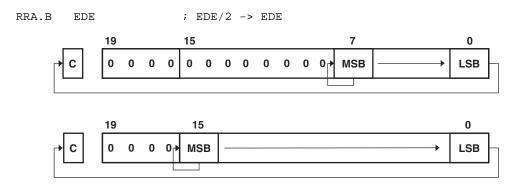


Figure 4-39. Rotate Right Arithmetically RRA.B and RRA.W



#### 4.6.2.41 RRC

RRC[.W] Rotate right through carry destination word RRC.B Rotate right through carry destination byte

Syntax RRC dst Or RRC.W dst

RRC.B dst

**Operation**  $C \rightarrow MSB \rightarrow MSB-1 \rightarrow ... LSB+1 \rightarrow LSB \rightarrow C$ 

**Description** The destination operand is shifted right by one bit position as shown in Figure 4-40. The

carry bit C is shifted into the MSB and the LSB is shifted into the carry bit C.

**Status Bits** N: Set if result is negative (MSB = 1), reset otherwise (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** RAM word EDE is shifted right one bit position. The MSB is loaded with 1.

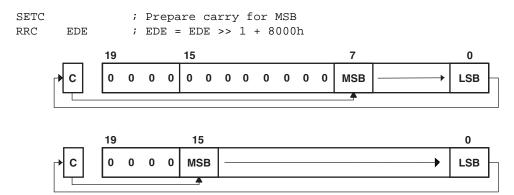


Figure 4-40. Rotate Right Through Carry RRC.B and RRC.W



#### 4.6.2.42 SBC

\* **SBC[.W]** Subtract borrow (.NOT. carry) from destination \* **SBC.B** Subtract borrow (.NOT. carry) from destination

Syntax SBC dst Or SBC.W dst

SBC.B dst

**Operation**  $dst + 0FFFFh + C \rightarrow dst$ 

 $dst + 0FFh + C \rightarrow dst$ 

Emulation SUBC #0,dst

SUBC.B #0,dst

**Description** The carry bit (C) is added to the destination operand minus one. The previous contents

of the destination are lost.

**Status Bits** N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

Set to 1 if no borrow, reset if borrow

V: Set if an arithmetic overflow occurs, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 16-bit counter pointed to by R13 is subtracted from a 32-bit counter pointed to by

R12.

SUB @R13,0(R12) ; Subtract LSDs

SBC 2(R12) ; Subtract carry from MSD

**Example** The 8-bit counter pointed to by R13 is subtracted from a 16-bit counter pointed to by

R12.

SUB.B @R13,0(R12) ; Subtract LSDs

SBC.B 1(R12) ; Subtract carry from MSD

#### **NOTE:** Borrow implementation

The borrow is treated as a .NOT. carry:

Borrow Carry Bit Yes 0 No 1



#### 4.6.2.43 SETC

\* **SETC** Set carry bit

**Description** The carry bit (C) is set. **Status Bits** N: Not affected

Z: Not affected

C: Set

V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** Emulation of the decimal subtraction:

Subtract R5 from R6 decimally.

Assume that R5 = 03987h and R6 = 04137h.

```
DSUB
       ADD
              #06666h,R5
                             ; Move content R5 from 0-9 to 6-0Fh
                              ; R5 = 03987h + 06666h = 09FEDh
       INV
              R5
                              ; Invert this (result back to 0-9)
                             ; R5 = .NOT. R5 = 06012h
       SETC
                              ; Prepare carry = 1
       DADD
                             ; Emulate subtraction by addition of:
              R5,R6
                              ; (010000h - R5 - 1)
                              ; R6 = R6 + R5 + 1
                              ; R6 = 0150h
```



# 4.6.2.44 SETN

\* **SETN** Set negative bit

**Description** The negative bit (N) is set.

Status Bits N: Set

Z: Not affectedC: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

207



# 4.6.2.45 SETZ

\* **SETZ** Set zero bit

Description The zero bit (Z) is set.Status Bits N: Not affected

Z: Set

C: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.



#### 4.6.2.46 SUB

SUB[.W] Subtract source word from destination word SUB.B Subtract source byte from destination byte

Syntax SUB src,dst Of SUB.W src,dst

SUB.B src,dst

**Operation** (.not.src) + 1 + dst  $\rightarrow$  dst or dst – src  $\rightarrow$  dst

**Description** The source operand is subtracted from the destination operand. This is made by adding

the 1s complement of the source + 1 to the destination. The source operand is not

affected, the result is written to the destination operand.

Status Bits N: Set if result is negative (src > dst), reset if positive (src ≤ dst)

Z: Set if result is zero (src = dst), reset otherwise (src  $\neq$  dst)

C: Set if there is a carry from the MSB, reset otherwise

V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no

overflow)

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** A 16-bit constant 7654h is subtracted from RAM word EDE.

SUB #7654h, &EDE ; Subtract 7654h from EDE

**Example** A table word pointed to by R5 (20-bit address) is subtracted from R7. Afterwards, if R7

contains zero, jump to label TONI. R5 is then auto-incremented by 2. R7.19:16 = 0.

SUB @R5+,R7 ; Subtract table number from R7. R5 + 2

JZ TONI ; R7 = @R5 (before subtraction) ... ; R7 <> @R5 (before subtraction)

**Example** Byte CNT is subtracted from byte R12 points to. The address of CNT is within PC ± 32K.

The address R12 points to is in full memory range.

SUB.B CNT,0(R12) ; Subtract CNT from @R12



#### 4.6.2.47 SUBC

SUBC[.W] Subtract source word with carry from destination word SUBC.B Subtract source byte with carry from destination byte

Syntax SUBC src,dst Of SUBC.W src,dst

SUBC.B src,dst

Operation

 $(.not.src) + C + dst \rightarrow dst$  or  $dst - (src - 1) + C \rightarrow dst$ 

Description

The source operand is subtracted from the destination operand. This is done by adding the 1s complement of the source + carry to the destination. The source operand is not affected, the result is written to the destination operand. Used for 32, 48, and 64-bit operands.

**Status Bits** 

N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB, reset otherwise

V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)

**Mode Bits** 

OSCOFF, CPUOFF, and GIE are not affected.

Example

A 16-bit constant 7654h is subtracted from R5 with the carry from the previous instruction. R5.19:16 = 0

SUBC.W #7654h,R5 ; Subtract 7654h + C from R5

#### **Example**

A 48-bit number (3 words) pointed to by R5 (20-bit address) is subtracted from a 48-bit counter in RAM, pointed to by R7. R5 points to the next 48-bit number afterwards. The address R7 points to is in full memory range.

#### Example

Byte CNT is subtracted from the byte, R12 points to. The carry of the previous instruction is used. The address of CNT is in lower 64KB.

SUBC.B &CNT,0(R12) ; Subtract byte CNT from @R12



## 4.6.2.48 SWPB

SWPB Swap bytes
Syntax SWPB dst

**Operation**  $dst.15:8 \leftrightarrow dst.7:0$ 

**Description** The high and the low byte of the operand are exchanged. PC.19:16 bits are cleared in

register mode.

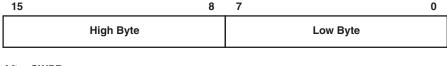
Status Bits Status bits are not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Exchange the bytes of RAM word EDE (lower 64KB)

MOV #1234h,&EDE ; 1234h -> EDE SWPB &EDE ; 3412h -> EDE

# Before SWPB



## After SWPB

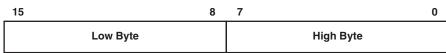


Figure 4-41. Swap Bytes in Memory

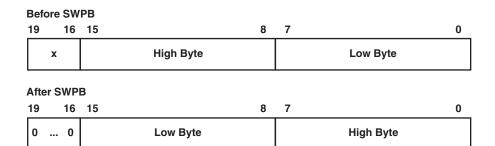


Figure 4-42. Swap Bytes in a Register



#### 4.6.2.49 SXT

SXT Extend sign Syntax SXT dst

**Operation**  $dst.7 \rightarrow dst.15:8$ ,  $dst.7 \rightarrow dst.19:8$  (register mode)

**Description** Register mode: the sign of the low byte of the operand is extended into the bits

Rdst.19:8.

Rdst.7 = 0: Rdst.19:8 = 000h afterwards Rdst.7 = 1: Rdst.19:8 = FFFh afterwards

Other modes: the sign of the low byte of the operand is extended into the high byte.

dst.7 = 0: high byte = 00h afterwards dst.7 = 1: high byte = FFh afterwards

Status Bits N: Set if result is negative, reset otherwise

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise (C = .not.Z)

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The signed 8-bit data in EDE (lower 64KB) is sign extended and added to the 16-bit

signed data in R7.

MOV.B &EDE,R5 ; EDE -> R5. 00XXh

SXT R5 ; Sign extend low byte to R5.19:8

ADD R5,R7 ; Add signed 16-bit values

**Example** The signed 8-bit data in EDE (PC +32 K) is sign extended and added to the 20-bit data

in R7.

MOV.B EDE,R5 ; EDE -> R5. 00XXh

SXT R5; Sign extend low byte to R5.19:8

ADDA R5,R7 ; Add signed 20-bit values



#### 4.6.2.50 TST

\* TST[.W] Test destination \* TST.B Test destination

TST dst or **Syntax** TST.W dst

TST.B dst

dst + 0FFFFh + 1 Operation

dst + 0FFh + 1

**Emulation** CMP #0,dst

CMP.B #0,dst

Description The destination operand is compared with zero. The status bits are set according to the

result. The destination is not affected.

**Status Bits** Set if destination is negative, reset if positive

> Z: Set if destination contains zero, reset otherwise

C: Set V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** R7 is tested. If it is negative, continue at R7NEG; if it is positive but not zero, continue at

R7POS.

TST R7 ; Test R7 R7NEG ; R7 is negative JN

JZR7ZERO ; R7 is zero

R7POS ; R7 is positive but not zero . . . . . .

R7NEG ; R7 is negative . . . . . . R7ZERO ; R7 is zero . . . . . .

**Example** The low byte of R7 is tested. If it is negative, continue at R7NEG; if it is positive but not zero, continue at R7POS.

> TST.B R7 ; Test low byte of R7

JN R7NEG ; Low byte of R7 is negative R7ZERO ; Low byte of R7 is zero JZ

R7POS ; Low byte of R7 is positive but not zero . . . . . .

R7NEG . . . . . ; Low byte of R7 is negative R7ZERO ; Low byte of R7 is zero . . . . . .



#### 4.6.2.51 XOR

**XOR[.W]** Exclusive OR source word with destination word **XOR.B** Exclusive OR source byte with destination byte

Syntax XOR src,dst Or XOR.W src,dst

XOR.B src,dst

**Operation**  $\operatorname{src} .\operatorname{xor} .\operatorname{dst} \to \operatorname{dst}$ 

**Description** The source and destination operands are exclusively ORed. The result is placed into the

destination. The source operand is not affected. The previous content of the destination

is lost.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise (C = .not. Z)

V: Set if both operands are negative before execution, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Toggle bits in word CNTR (16-bit data) with information (bit = 1) in address-word TONI.

Both operands are located in lower 64KB.

XOR &TONI, &CNTR ; Toggle bits in CNTR

**Example** A table word pointed to by R5 (20-bit address) is used to toggle bits in R6. R6.19:16 = 0.

XOR @R5,R6 ; Toggle bits in R6

**Example** Reset to zero those bits in the low byte of R7 that are different from the bits in byte EDE.

R7.19:8 = 0. The address of EDE is within PC  $\pm$  32 K.

XOR.B EDE,R7 ; Set different bits to 1 in R7.

INV.B R7; Invert low byte of R7, high byte is 0h



## 4.6.3 MSP430X Extended Instructions

The MSP430X extended instructions give the MSP430X CPU full access to its 20-bit address space. MSP430X instructions require an additional word of op-code called the extension word. All addresses, indexes, and immediate numbers have 20-bit values when preceded by the extension word. The MSP430X extended instructions are described in the following sections. For MSP430X instructions that do not require the extension word, it is noted in the instruction description.

See Section 4.6.2 for standard MSP430 instructions and Section 4.6.4 for MSP430X address instructions.

Operation

**Emulation** 



## 4.6.3.1 ADCX

\* ADCX.A Add carry to destination address-word

\* ADCX.[W] Add carry to destination word ADCX.B Add carry to destination byte

Syntax ADCX.A dst

ADCX dst Of ADCX.W dst

ADCX.B dst  $dst + C \rightarrow dst$ 

ADDCX.A #0,dst ADDCX #0,dst ADDCX.B #0,dst

**Description** The carry bit (C) is added to the destination operand. The previous contents of the

destination are lost.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative

numbers is positive, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 40-bit counter, pointed to by R12 and R13, is incremented.

INCX.A @R12 ; Increment lower 20 bits ADCX.A @R13 ; Add carry to upper 20 bits



#### 4.6.3.2 ADDX

ADDX.A Add source address-word to destination address-word

ADDX.[W] Add source word to destination word ADDX.B Add source byte to destination byte

Syntax ADDX.A src,dst

ADDX src,dst Of ADDX.W src,dst

ADDX.B src,dst

**Operation**  $\operatorname{src} + \operatorname{dst} \to \operatorname{dst}$ 

**Description** The source operand is added to the destination operand. The previous contents of the

destination are lost. Both operands can be located in the full address space.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative

numbers is positive, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Ten is added to the 20-bit pointer CNTR located in two words CNTR (LSBs) and

CNTR+2 (MSBs).

```
ADDX.A #10,CNTR ; Add 10 to 20-bit pointer
```

## Example

A table word (16-bit) pointed to by R5 (20-bit address) is added to R6. The jump to label TONI is performed on a carry.

```
ADDX.W @R5,R6 ; Add table word to R6
JC TONI ; Jump if carry
... ; No carry
```

## Example

A table byte pointed to by R5 (20-bit address) is added to R6. The jump to label TONI is performed if no carry occurs. The table pointer is auto-incremented by 1.

```
ADDX.B @R5+,R6 ; Add table byte to R6. R5 + 1. R6: 000xxh JNC TONI ; Jump if no carry ... ; Carry occurred
```

Note: Use ADDA for the following two cases for better code density and execution.

```
ADDX.A Rsrc,Rdst
ADDX.A #imm20,Rdst
```



## 4.6.3.3 ADDCX

ADDCX.A Add source address-word and carry to destination address-word

ADDCX.[W] Add source word and carry to destination word ADDCX.B Add source byte and carry to destination byte

Syntax ADDCX.A src,dst

ADDCX src,dst Of ADDCX.W src,dst

ADDCX.B src,dst

**Operation**  $\operatorname{src} + \operatorname{dst} + \operatorname{C} \to \operatorname{dst}$ 

**Description** The source operand and the carry bit C are added to the destination operand. The

previous contents of the destination are lost. Both operands may be located in the full

address space.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative

numbers is positive, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Constant 15 and the carry of the previous instruction are added to the 20-bit counter

CNTR located in two words.

```
ADDCX.A #15,&CNTR ; Add 15 + C to 20-bit CNTR
```

# Example

A table word pointed to by R5 (20-bit address) and the carry C are added to R6. The jump to label TONI is performed on a carry.

```
ADDCX.W @R5,R6 ; Add table word + C to R6
JC TONI ; Jump if carry
... ; No carry
```

#### **Example**

A table byte pointed to by R5 (20-bit address) and the carry bit C are added to R6. The jump to label TONI is performed if no carry occurs. The table pointer is auto-incremented by 1.

```
ADDCX.B @R5+,R6 ; Add table byte + C to R6. R5 + 1 JNC ; Jump if no carry ... ; Carry occurred
```



## 4.6.3.4 ANDX

ANDX.A Logical AND of source address-word with destination address-word

ANDX.[W] Logical AND of source word with destination word ANDX.B Logical AND of source byte with destination byte

Syntax ANDX.A src,dst

ANDX src,dst Of ANDX.W src,dst

ANDX.B src,dst

**Description** The source operand and the destination operand are logically ANDed. The result is

placed into the destination. The source operand is not affected. Both operands may be

located in the full address space.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if the result is not zero, reset otherwise. C = (.not. Z)

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The bits set in R5 (20-bit data) are used as a mask (AAA55h) for the address-word TOM

located in two words. If the result is zero, a branch is taken to label TONI.

```
MOVA #AAA55h,R5 ; Load 20-bit mask to R5 ANDX.A R5,TOM ; TOM .and. R5 -> TOM JZ TONI ; Jump if result 0 .... ; Result > 0
```

# or shorter:

```
ANDX.A #AAA55h,TOM ; TOM .and. AAA55h -> TOM 
JZ TONI ; Jump if result 0
```

#### **Example**

A table byte pointed to by R5 (20-bit address) is logically ANDed with R6. R6.19:8 = 0. The table pointer is auto-incremented by 1.

```
ANDX.B @R5+,R6 ; AND table byte with R6. R5 + 1
```



## 4.6.3.5 BICX

**BICX.A** Clear bits set in source address-word in destination address-word

BICX.[W] Clear bits set in source word in destination word Clear bits set in source byte in destination byte

Syntax BICX.A src,dst

BICX src,dst Of BICX.W src,dst

BICX.B src,dst

**Operation** (.not. src) .and.  $dst \rightarrow dst$ 

**Description** The inverted source operand and the destination operand are logically ANDed. The

result is placed into the destination. The source operand is not affected. Both operands

may be located in the full address space.

Status Bits N: Not affected

Z: Not affectedC: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected. Example The bits 19:15 of R5 (20-bit data) are cleared.

BICX.A #0F8000h,R5 ; Clear R5.19:15 bits

**Example** A table word pointed to by R5 (20-bit address) is used to clear bits in R7. R7.19:16 = 0.

BICX.W @R5,R7 ; Clear bits in R7

**Example** A table byte pointed to by R5 (20-bit address) is used to clear bits in output Port1.

BICX.B @R5,&P1OUT ; Clear I/O port P1 bits



## 4.6.3.6 BISX

**BISX.A** Set bits set in source address-word in destination address-word

BISX.[W] Set bits set in source word in destination word BISX.B Set bits set in source byte in destination byte

Syntax BISX.A src,dst

BISX src,dst Of BISX.W src,dst

BISX.B src,dst

**Description** The source operand and the destination operand are logically ORed. The result is placed

into the destination. The source operand is not affected. Both operands may be located

in the full address space.

Status Bits N: Not affected

Z: Not affectedC: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Bits 16 and 15 of R5 (20-bit data) are set to one.

BISX.A #018000h,R5 ; Set R5.16:15 bits

**Example** A table word pointed to by R5 (20-bit address) is used to set bits in R7.

BISX.W @R5,R7 ; Set bits in R7

**Example** A table byte pointed to by R5 (20-bit address) is used to set bits in output Port1.

BISX.B @R5,&P1OUT ; Set I/O port P1 bits



#### 4.6.3.7 BITX

BITX.A Test bits set in source address-word in destination address-word

BITX.[W] Test bits set in source word in destination word
BITX.B Test bits set in source byte in destination byte

Syntax BITX.A src,dst

BITX src,dst Of BITX.W src,dst

BITX.B src,dst

**Description** The source operand and the destination operand are logically ANDed. The result affects

only the status bits. Both operands may be located in the full address space.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if the result is not zero, reset otherwise. C = (.not. Z)

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Test if bit 16 or 15 of R5 (20-bit data) is set. Jump to label TONI if so.

```
BITX.A #018000h,R5 ; Test R5.16:15 bits
JNZ TONI ; At least one bit is set
... ; Both are reset
```

# **Example**

A table word pointed to by R5 (20-bit address) is used to test bits in R7. Jump to label TONI if at least one bit is set.

```
BITX.W @R5,R7 ; Test bits in R7: C = .not.Z

JC TONI ; At least one is set

... ; Both are reset
```

# Example

A table byte pointed to by R5 (20-bit address) is used to test bits in input Port1. Jump to label TONI if no bit is set. The next table byte is addressed.

```
BITX.B @R5+,&P1IN ; Test input P1 bits. R5 + 1

JNC TONI ; No corresponding input bit is set

... ; At least one bit is set
```



# 4.6.3.8 CLRX

\* CLRX.A Clear destination address-word

\* CLRX.[W] Clear destination word \* CLRX.B Clear destination byte

Syntax CLRX.A dst

CLRX dst Of CLRX.W dst

CLRX.B dst

**Emulation** MOVX.A #0,dst

MOVX #0,dst MOVX.B #0,dst

**Description** The destination operand is cleared.

Status Bits Status bits are not affected.

**Example** RAM address-word TONI is cleared.

CLRX.A TONI ; 0 -> TONI



#### 4.6.3.9 CMPX

**CMPX.A** Compare source address-word and destination address-word

**CMPX.[W]** Compare source word and destination word **CMPX.B** Compare source byte and destination byte

Syntax CMPX.A src,dst

CMPX src,dst Of CMPX.W src,dst

CMPX.B src,dst

**Operation** (.not. src) + 1 + dst or dst - src

**Description** The source operand is subtracted from the destination operand by adding the 1s

complement of the source + 1 to the destination. The result affects only the status bits.

Both operands may be located in the full address space.

Status Bits N: Set if result is negative (src > dst), reset if positive (src ≤ dst)

Z: Set if result is zero (src = dst), reset otherwise (src  $\neq$  dst)

C: Set if there is a carry from the MSB, reset otherwise

V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Compare EDE with a 20-bit constant 18000h. Jump to label TONI if EDE equals the

constant.

```
CMPX.A #018000h,EDE ; Compare EDE with 18000h

JEQ TONI ; EDE contains 18000h

... ; Not equal
```

# **Example**

A table word pointed to by R5 (20-bit address) is compared with R7. Jump to label TONI if R7 contains a lower, signed, 16-bit number.

## Example

A table byte pointed to by R5 (20-bit address) is compared to the input in I/O Port1. Jump to label TONI if the values are equal. The next table byte is addressed.

```
CMPX.B @R5+,&P1IN ; Compare P1 bits with table. R5 + 1
JEQ TONI ; Equal contents
... ; Not equal
```

Note: Use CMPA for the following two cases for better density and execution.

```
CMPA Rsrc,Rdst
CMPA #imm20,Rdst
```



## 4.6.3.10 DADCX

\* DADCX.A Add carry decimally to destination address-word

\* **DADCX.[W]** Add carry decimally to destination word \* **DADCX.B** Add carry decimally to destination byte

Syntax DADCX.A dst

DADCX dst Or DADCX.W dst

DADCX.B dst

**Operation**  $dst + C \rightarrow dst (decimally)$ 

Emulation DADDX.A #0,dst

DADDX #0,dst
DADDX.B #0,dst

**Description** The carry bit (C) is added decimally to the destination.

Status Bits N: Set if MSB of result is 1 (address-word > 79999h, word > 7999h, byte > 79h), reset

if MSB is 0

Z: Set if result is zero, reset otherwise

C: Set if the BCD result is too large (address-word > 99999h, word > 9999h, byte >

99h), reset otherwise

V: Undefined

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 40-bit counter, pointed to by R12 and R13, is incremented decimally.

DADDX.A #1,0(R12) ; Increment lower 20 bits DADCX.A 0(R13) ; Add carry to upper 20 bits



#### 4.6.3.11 DADDX

**DADDX.A** Add source address-word and carry decimally to destination address-word

**DADDX.[W]** Add source word and carry decimally to destination word Add source byte and carry decimally to destination byte

Syntax DADDX.A src,dst

DADDX src,dst Of DADDX.W src,dst

DADDX.B src,dst

**Operation**  $\operatorname{src} + \operatorname{dst} + \operatorname{C} \to \operatorname{dst} (\operatorname{decimally})$ 

**Description** The source operand and the destination operand are treated as two (.B), four (.W), or

five (.A) binary coded decimals (BCD) with positive signs. The source operand and the carry bit C are added decimally to the destination operand. The source operand is not affected. The previous contents of the destination are lost. The result is not defined for

non-BCD numbers. Both operands may be located in the full address space.

Status Bits N: Set if MSB of result is 1 (address-word > 79999h, word > 7999h, byte > 79h), reset if MSB is 0.

Z: Set if result is zero, reset otherwise

C: Set if the BCD result is too large (address-word > 99999h, word > 9999h, byte >

99h), reset otherwise

V: Undefined

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Decimal 10 is added to the 20-bit BCD counter DECCNTR located in two words.

```
DADDX.A #10h,&DECCNTR ; Add 10 to 20-bit BCD counter
```

# **Example**

The eight-digit BCD number contained in 20-bit addresses BCD and BCD+2 is added decimally to an eight-digit BCD number contained in R4 and R5 (BCD+2 and R5 contain the MSDs).

```
CLRC ; Clear carry

DADDX.W BCD,R4 ; Add LSDs

DADDX.W BCD+2,R5 ; Add MSDs with carry

JC OVERFLOW ; Result >999999999 go to error routine

... ; Result ok
```

#### **Example**

The two-digit BCD number contained in 20-bit address BCD is added decimally to a two-digit BCD number contained in R4.

```
CLRC ; Clear carry
DADDX.B BCD,R4 ; Add BCD to R4 decimally.
; R4: 000ddh
```



Operation

**Emulation** 

## 4.6.3.12 DECX

\* **DECX.A** Decrement destination address-word

\* **DECX.[W]** Decrement destination word \* **DECX.B** Decrement destination byte

Syntax DECX.A dst

DECX dst Or DECX.W dst

DECX.B dst  $dst - 1 \rightarrow dst$  SUBX.A #1,dst

SUBX #1,dst SUBX.B #1,dst

**Description** The destination operand is decremented by one. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 1, reset otherwiseC: Reset if dst contained 0, set otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example RAM address-word TONI is decremented by one.

DECX.A TONI ; Decrement TONI

Operation

**Emulation** 



# 4.6.3.13 DECDX

\* **DECDX.A** Double-decrement destination address-word

\* **DECDX.[W]** Double-decrement destination word Double-decrement destination byte

Syntax DECDX.A dst

DECDX dst Or DECDX.W dst

DECDX.B dst  $dst - 2 \rightarrow dst$  SUBX.A #2,dst

SUBX #2,dst SUBX.B #2,dst

**Description** The destination operand is decremented by two. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 2, reset otherwise

C: Reset if dst contained 0 or 1, set otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example RAM address-word TONI is decremented by two.

DECDX.A TONI ; Decrement TONI



## 4.6.3.14 INCX

\* INCX.A Increment destination address-word

\* INCX.[W] Increment destination word
\* INCX.B Increment destination byte

Syntax INCX.A dst

INCX dst Or INCX.W dst

INCX.B dst

Operation  $dst + 1 \rightarrow dst$ Emulation ADDX.A #1, dstADDX #1, dst

ADDX.B #1,dst

**Description** The destination operand is incremented by one. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFFh, reset otherwise Set if dst contained 0FFh, reset otherwise

C: Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFFh, reset otherwise Set if dst contained 0FFh, reset otherwise

V: Set if dst contained 07FFFh, reset otherwise Set if dst contained 07FFFh, reset otherwise Set if dst contained 07Fh, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example RAM address-word TONI is incremented by one.

INCX.A TONI ; Increment TONI (20-bits)



# 4.6.3.15 INCDX

\* **INCDX.A** Double-increment destination address-word

\* INCDX.[W] Double-increment destination word Double-increment destination byte

Syntax INCDX.A dst

INCDX dst Of INCDX.W dst

INCDX.B dst

ADDX.B #2,dst

**Description** The destination operand is incremented by two. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFEh, reset otherwise Set if dst contained 0FFFEh, reset otherwise Set if dst contained 0FEh, reset otherwise

C: Set if dst contained 0FFFEh or 0FFFFh, reset otherwise Set if dst contained 0FFFEh or 0FFFh, reset otherwise Set if dst contained 0FEh or 0FFh, reset otherwise

V: Set if dst contained 07FFFEh or 07FFFFh, reset otherwise Set if dst contained 07FFEh or 07FFFh, reset otherwise Set if dst contained 07Eh or 07Fh, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** RAM byte LEO is incremented by two; PC points to upper memory.

INCDX.B LEO ; Increment LEO by two



## 4.6.3.16 INVX

\* INVX.A Invert destination
\* INVX.[W] Invert destination
\* INVX.B Invert destination

Syntax INVX.A dst

INVX dst Or INVX.W dst

INVX.B dst

**Operation** .NOT.dst  $\rightarrow$  dst

Emulation XORX.A #0FFFFFh,dst

XORX #0FFFFh,dst XORX.B #0FFh,dst

**Description** The destination operand is inverted. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFFh, reset otherwise Set if dst contained 0FFh, reset otherwise

C: Set if result is not zero, reset otherwise ( = .NOT. Zero)

V: Set if initial destination operand was negative, otherwise reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example 20-bit content of R5 is negated (2s complement).

INVX.A R5 ; Invert R5

INCX.A R5 ; R5 is now negated

**Example** Content of memory byte LEO is negated. PC is pointing to upper memory.

INVX.B LEO ; Invert LEO

INCX.B LEO ; MEM(LEO) is negated



#### 4.6.3.17 MOVX

MOVX.A Move source address-word to destination address-word

MOVX.[W] Move source word to destination word MOVX.B Move source byte to destination byte

Syntax MOVX.A src,dst

MOVX src,dst Of MOVX.W src,dst

MOVX.B src,dst

**Description** The source operand is copied to the destination. The source operand is not affected.

Both operands may be located in the full address space.

Status Bits N: Not affected

Z: Not affectedC: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** Move a 20-bit constant 18000h to absolute address-word EDE

MOVX.A #018000h, &EDE ; Move 18000h to EDE

## **Example**

The contents of table EDE (word data, 20-bit addresses) are copied to table TOM. The length of the table is 030h words.

```
MOVA #EDE,R10 ; Prepare pointer (20-bit address)

Loop MOVX.W @R10+,TOM-EDE-2(R10) ; R10 points to both tables.
; R10+2

CMPA #EDE+60h,R10 ; End of table reached?
JLO Loop ; Not yet
... ; Copy completed
```

#### Example

The contents of table EDE (byte data, 20-bit addresses) are copied to table TOM. The length of the table is 020h bytes.

```
MOVA
               #EDE,R10
                                       ; Prepare pointer (20-bit)
      MOV
               #20h,R9
                                       ; Prepare counter
               @R10+,TOM-EDE-2(R10)
Loop
      MOVX.W
                                       ; R10 points to both tables.
                                       ; R10+1
      DEC
               R9
                                       ; Decrement counter
      JNZ
               Loop
                                       ; Not yet done
                                       ; Copy completed
```

Ten of the 28 possible addressing combinations of the MOVX.A instruction can use the MOVA instruction. This saves two bytes and code cycles. Examples for the addressing combinations are:

MOVX.A	Rsrc,Rdst	MOVA	Rsrc,Rdst	;	Reg/Reg
MOVX.A	#imm20,Rdst	MOVA	#imm20,Rdst	;	Immediate/Reg
MOVX.A	&abs20,Rdst	MOVA	&abs20,Rdst	;	Absolute/Reg
MOVX.A	@Rsrc,Rdst	MOVA	@Rsrc,Rdst	;	Indirect/Reg
MOVX.A	@Rsrc+,Rdst	MOVA	@Rsrc+,Rdst	;	Indirect, Auto/Reg
MOVX.A	Rsrc,&abs20	MOVA	Rsrc,&abs20	;	Reg/Absolute



# The next four replacements are possible only if 16-bit indexes are sufficient for the addressing:

MOVX.A	z20(Rsrc),Rdst	MOVA	z16(Rsrc),Rdst	;	Indexed/Reg
MOVX.A	Rsrc,z20(Rdst)	MOVA	Rsrc,z16(Rdst)	;	Reg/Indexed
MOVX.A	symb20,Rdst	MOVA	symb16,Rdst	;	Symbolic/Reg
MOVX.A	Rsrc,svmb20	MOVA	Rsrc,symb16	;	Reg/Symbolic



#### 4.6.3.18 POPM

**POPM.A** Restore n CPU registers (20-bit data) from the stack **POPM.[W]** Restore n CPU registers (16-bit data) from the stack

Syntax POPM.A #n,Rdst  $1 \le n \le 16$ POPM.W #n,Rdst Or POPM #n,Rdst  $1 \le n \le 16$ 

Operation POPM.A: Restore the register values from stack to the specified CPU registers. The SP

is incremented by four for each register restored from stack. The 20-bit values from

stack (two words per register) are restored to the registers.

POPM.W: Restore the 16-bit register values from stack to the specified CPU registers. The SP is incremented by two for each register restored from stack. The 16-bit values

from stack (one word per register) are restored to the CPU registers.

Note: This instruction does not use the extension word.

**Description** POPM.A: The CPU registers pushed on the stack are moved to the extended CPU

registers, starting with the CPU register (Rdst - n + 1). The SP is incremented by (n  $\times$ 

4) after the operation.

POPM.W: The 16-bit registers pushed on the stack are moved back to the CPU registers, starting with CPU register (Rdst - n + 1). The SP is incremented by (n × 2) after the instruction. The MSBs (Rdst.19:16) of the restored CPU registers are cleared.

**Status Bits** Status bits are not affected, except SR is included in the operation.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Restore the 20-bit registers R9, R10, R11, R12, R13 from the stack

POPM.A #5,R13 ; Restore R9, R10, R11, R12, R13

**Example** Restore the 16-bit registers R9, R10, R11, R12, R13 from the stack.

POPM.W #5,R13 ; Restore R9, R10, R11, R12, R13



## 4.6.3.19 PUSHM

**PUSHM.A** Save n CPU registers (20-bit data) on the stack **PUSHM.[W]** Save n CPU registers (16-bit words) on the stack

Syntax PUSHM.A #n,Rdst  $1 \le n \le 16$ PUSHM.W #n,Rdst Or PUSHM #n,Rdst  $1 \le n \le 16$ 

**Operation** PUSHM.A: Save the 20-bit CPU register values on the stack. The SP is decremented

by four for each register stored on the stack. The MSBs are stored first (higher

address).

PUSHM.W: Save the 16-bit CPU register values on the stack. The SP is decremented

by two for each register stored on the stack.

**Description** PUSHM.A: The n CPU registers, starting with Rdst backwards, are stored on the stack.

The SP is decremented by (n × 4) after the operation. The data (Rn.19:0) of the pushed

CPU registers is not affected.

PUSHM.W: The n registers, starting with Rdst backwards, are stored on the stack. The SP is decremented by  $(n \times 2)$  after the operation. The data (Rn.19:0) of the pushed

CPU registers is not affected.

Note: This instruction does not use the extension word.

Status Bits Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Save the five 20-bit registers R9, R10, R11, R12, R13 on the stack

PUSHM.A #5,R13 ; Save R13, R12, R11, R10, R9

**Example** Save the five 16-bit registers R9, R10, R11, R12, R13 on the stack

PUSHM.W #5,R13; Save R13, R12, R11, R10, R9



## 4.6.3.20 POPX

\* **POPX.A** Restore single address-word from the stack

\* POPX.[W] Restore single word from the stack
\* POPX.B Restore single byte from the stack

Syntax POPX.A dst

POPX dst Or POPX.W dst

POPX.B dst

**Operation** Restore the 8-/16-/20-bit value from the stack to the destination. 20-bit addresses are

possible. The SP is incremented by two (byte and word operands) and by four

(address-word operand).

**Emulation** MOVX(.B,.A) @SP+,dst

**Description** The item on TOS is written to the destination operand. Register mode, Indexed mode,

Symbolic mode, and Absolute mode are possible. The SP is incremented by two or

four.

Note: The SP is incremented by two also for byte operations.

**Status Bits** Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Write the 16-bit value on TOS to the 20-bit address &EDE

POPX.W &EDE ; Write word to address EDE

**Example** Write the 20-bit value on TOS to R9

POPX.A R9 ; Write address-word to R9



## 4.6.3.21 PUSHX

**PUSHX.A** Save single address-word to the stack

PUSHX.[W] Save single word to the stack **PUSHX.B** Save single byte to the stack

PUSHX.A src **Syntax** 

> PUSHX src Or PUSHX.W src

PUSHX.B src

Operation Save the 8-/16-/20-bit value of the source operand on the TOS. 20-bit addresses are

possible. The SP is decremented by two (byte and word operands) or by four (address-

word operand) before the write operation.

The SP is decremented by two (byte and word operands) or by four (address-word Description

operand). Then the source operand is written to the TOS. All seven addressing modes

are possible for the source operand.

**Status Bits** Status bits are not affected.

OSCOFF, CPUOFF, and GIE are not affected. **Mode Bits** 

Example Save the byte at the 20-bit address &EDE on the stack

> PUSHX.B &EDE ; Save byte at address EDE

**Example** Save the 20-bit value in R9 on the stack.

> ; Save address-word in R9 PUSHX.A R9



## 4.6.3.22 RLAM

RLAM.A Rotate left arithmetically the 20-bit CPU register content RLAM.[W] Rotate left arithmetically the 16-bit CPU register content

RLAM.A #n,Rdst **Syntax**  $1 \le n \le 4$ RLAM.W #n, Rdst Of RLAM #n, Rdst  $1 \le n \le 4$ 

Operation  $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow 0$ 

Description The destination operand is shifted arithmetically left one, two, three, or four positions as

shown in Figure 4-43. RLAM works as a multiplication (signed and unsigned) with 2, 4,

8, or 16. The word instruction RLAM.W clears the bits Rdst.19:16.

Note: This instruction does not use the extension word.

**Status Bits** Set if result is negative

> .A: Rdst.19 = 1, reset if Rdst.19 = 0.W: Rdst.15 = 1, reset if Rdst.15 = 0

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB (n = 1), MSB-1 (n = 2), MSB-2 (n = 3), MSB-3 (n = 4)

V: Undefined

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** 

The 20-bit operand in R5 is shifted left by three positions. It operates equal to an

arithmetic multiplication by 8.

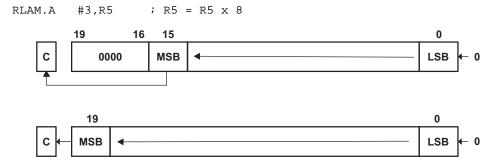


Figure 4-43. Rotate Left Arithmetically—RLAM[.W] and RLAM.A



## 4.6.3.23 RLAX

\* RLAX.A Rotate left arithmetically address-word

\* RLAX.[W] Rotate left arithmetically word
\* RLAX.B Rotate left arithmetically byte

Syntax RLAX.A dst

RLAX dst Of RLAX.W dst

RLAX.B dst

**Operation**  $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow 0$ 

**Emulation** ADDX.A dst,dst

ADDX dst,dst
ADDX.B dst,dst

**Description** The destination operand is shifted left one position as shown in Figure 4-44. The MSB

is shifted into the carry bit (C) and the LSB is filled with 0. The RLAX instruction acts as

a signed multiplication by 2.

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB

V: Set if an arithmetic overflow occurs: the initial value is 040000h ≤ dst < 0C0000h;

reset otherwise

Set if an arithmetic overflow occurs: the initial value is 04000h ≤ dst < 0C000h;

reset otherwise

Set if an arithmetic overflow occurs: the initial value is 040h ≤ dst < 0C0h; reset

otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 20-bit value in R7 is multiplied by 2

RLAX.A R7 ; Shift left R7 (20-bit)

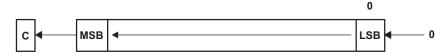


Figure 4-44. Destination Operand-Arithmetic Shift Left



## 4.6.3.24 RLCX

\* RLCX.A Rotate left through carry address-word

\* RLCX.[W] Rotate left through carry word \* RLCX.B Rotate left through carry byte

Syntax RLCX.A dst

RLCX dst Of RLCX.W dst

RLCX.B dst

**Operation**  $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow C$ 

Emulation ADDCX.A dst,dst

ADDCX dst,dst ADDCX.B dst,dst

**Description** The destination operand is shifted left one position as shown in Figure 4-45. The carry

bit (C) is shifted into the LSB and the MSB is shifted into the carry bit (C).

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB

V: Set if an arithmetic overflow occurs: the initial value is 040000h ≤ dst < 0C0000h; reset otherwise

Set if an arithmetic overflow occurs: the initial value is 04000h ≤ dst < 0C000h; reset otherwise

Set if an arithmetic overflow occurs: the initial value is 040h ≤ dst < 0C0h; reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The 20-bit value in R5 is shifted left one position.

RLCX.A R5 ;  $(R5 \times 2) + C -> R5$ 

**Example** The RAM byte LEO is shifted left one position. PC is pointing to upper memory.

RLCX.B LEO ; RAM(LEO) x 2 + C -> RAM(LEO)

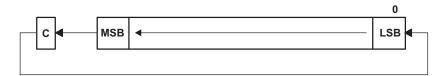


Figure 4-45. Destination Operand-Carry Left Shift



## 4.6.3.25 RRAM

Description

RRAM.A Rotate right arithmetically the 20-bit CPU register content RRAM.[W] Rotate right arithmetically the 16-bit CPU register content

RRAM.A #n,Rdst **Syntax**  $1 \le n \le 4$ RRAM.W #n, Rdst Or RRAM #n, Rdst  $1 \le n \le 4$ 

Operation  $MSB \rightarrow MSB \rightarrow MSB-1 \dots LSB+1 \rightarrow LSB \rightarrow C$ 

> The destination operand is shifted right arithmetically by one, two, three, or four bit positions as shown in Figure 4-46. The MSB retains its value (sign). RRAM operates equal to a signed division by 2/4/8/16. The MSB is retained and shifted into MSB-1. The LSB+1 is shifted into the LSB, and the LSB is shifted into the carry bit C. The word

instruction RRAM.W clears the bits Rdst.19:16.

Note: This instruction does not use the extension word.

**Status Bits** Set if result is negative

> .A: Rdst.19 = 1, reset if Rdst.19 = 0.W: Rdst.15 = 1, reset if Rdst.15 = 0

Z: Set if result is zero, reset otherwise

Loaded from the LSB (n = 1), LSB+1 (n = 2), LSB+2 (n = 3), or LSB+3 (n = 4)

V: Reset

19

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

15

**MSB** 

**Example** The signed 20-bit number in R5 is shifted arithmetically right two positions.

RRAM.A #2,R5 ; R5/4 -> R5

16

0000

#### **Example** The signed 20-bit value in R15 is multiplied by 0.75. $(0.5 + 0.25) \times R15$ .

```
PUSHM.A
             #1,R15
                               ; Save extended R15 on stack
RRAM.A
             #1,R15
                               ; R15 y 0.5 -> R15
ADDX.A
             @SP+,R15
                               ; R15 y 0.5 + R15 = 1.5 y R15 \rightarrow R15
             #1,R15
                               ; (1.5 \text{ y R15}) \text{ y } 0.5 = 0.75 \text{ y R15} \rightarrow \text{R15}
RRAM.A
```

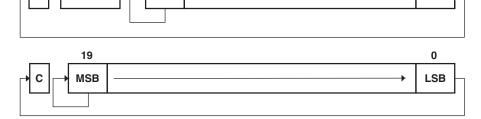


Figure 4-46. Rotate Right Arithmetically RRAM[.W] and RRAM.A

0

LSB



#### 4.6.3.26 RRAX

RRAX.A Rotate right arithmetically the 20-bit operand RRAX.[W] Rotate right arithmetically the 16-bit operand RRAX.B Rotate right arithmetically the 8-bit operand

Syntax RRAX.A Rdst

RRAX.W Rdst RRAX Rdst RRAX.B Rdst RRAX.A dst

RRAX dst Of RRAX.W dst

RRAX.B dst

Operation

 $\mathsf{MSB} \to \mathsf{MSB} \to \mathsf{MSB-1} \; ... \; \mathsf{LSB+1} \to \mathsf{LSB} \to \mathsf{C}$ 

Description

Register mode for the destination: the destination operand is shifted right by one bit position as shown in Figure 4-47. The MSB retains its value (sign). The word instruction RRAX.W clears the bits Rdst.19:16, the byte instruction RRAX.B clears the bits Rdst.19:8. The MSB retains its value (sign), the LSB is shifted into the carry bit. RRAX here operates equal to a signed division by 2.

All other modes for the destination: the destination operand is shifted right arithmetically by one bit position as shown in Figure 4-48. The MSB retains its value (sign), the LSB is shifted into the carry bit. RRAX here operates equal to a signed division by 2. All addressing modes, with the exception of the Immediate mode, are possible in the full memory.

**Status Bits** 

N: Set if result is negative, reset if positive

.A: dst.19 = 1, reset if dst.19 = 0 .W: dst.15 = 1, reset if dst.15 = 0 .B: dst.7 = 1, reset if dst.7 = 0

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

Mode Bits

OSCOFF, CPUOFF, and GIE are not affected.

Example

The signed 20-bit number in R5 is shifted arithmetically right four positions.

RPT #4
RRAX.A R5 ; R5/16 -> R5



# **Example** The signed 8-bit value in EDE is multiplied by 0.5.

RRAX.B &EDE ; EDE/2 -> EDE

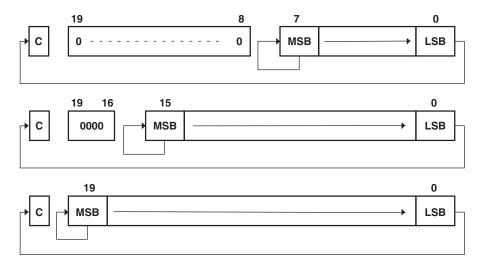


Figure 4-47. Rotate Right Arithmetically RRAX(.B,.A) – Register Mode

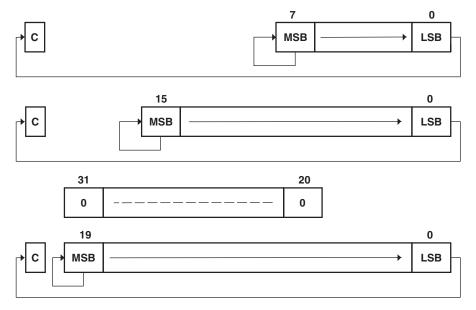


Figure 4-48. Rotate Right Arithmetically RRAX(.B,.A) - Non-Register Mode



## 4.6.3.27 RRCM

RRCM.A Rotate right through carry the 20-bit CPU register content RRCM.[W] Rotate right through carry the 16-bit CPU register content

Syntax RRCM.A #n, Rdst  $1 \le n \le 4$  RRCM.W #n, Rdst 0f RRCM #n, Rdst  $1 \le n \le 4$ 

**Operation**  $C \rightarrow MSB \rightarrow MSB-1 \dots LSB+1 \rightarrow LSB \rightarrow C$ 

**Description** The destination operand is shifted right by one, two, three, or four bit positions as

shown in Figure 4-49. The carry bit C is shifted into the MSB, the LSB is shifted into the

carry bit. The word instruction RRCM.W clears the bits Rdst.19:16.

Note: This instruction does not use the extension word.

Status Bits N: Set if result is negative

.A: Rdst.19 = 1, reset if Rdst.19 = 0 .W: Rdst.15 = 1, reset if Rdst.15 = 0

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB (n = 1), LSB+1 (n = 2), LSB+2 (n = 3), or LSB+3 (n = 4)

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** The address-word in R5 is shifted right by three positions. The MSB–2 is loaded with 1.

```
SETC ; Prepare carry for MSB-2 RRCM.A \#3,R5 ; R5 = R5 \gg 3 + 20000h
```

## Example

The word in R6 is shifted right by two positions. The MSB is loaded with the LSB. The MSB–1 is loaded with the contents of the carry flag.

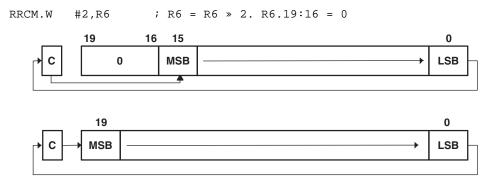


Figure 4-49. Rotate Right Through Carry RRCM[.W] and RRCM.A



## 4.6.3.28 RRCX

RRCX.A Rotate right through carry the 20-bit operand RRCX.[W] Rotate right through carry the 16-bit operand RRCX.B Rotate right through carry the 8-bit operand

Syntax RRCX.A Rdst

RRCX.W Rdst RRCX.B Rdst RRCX.B dst

RRCX dst Of RRCX.W dst

RRCX.B dst

Operation

 $C \rightarrow MSB \rightarrow MSB-1 \dots LSB+1 \rightarrow LSB \rightarrow C$ 

Description

Register mode for the destination: the destination operand is shifted right by one bit position as shown in Figure 4-50. The word instruction RRCX.W clears the bits Rdst.19:16, the byte instruction RRCX.B clears the bits Rdst.19:8. The carry bit C is shifted into the MSB, the LSB is shifted into the carry bit.

All other modes for the destination: the destination operand is shifted right by one bit position as shown in Figure 4-51. The carry bit C is shifted into the MSB, the LSB is shifted into the carry bit. All addressing modes, with the exception of the Immediate mode, are possible in the full memory.

**Status Bits** 

N: Set if result is negative

.A: dst.19 = 1, reset if dst.19 = 0.W: dst.15 = 1, reset if dst.15 = 0.B: dst.7 = 1, reset if dst.7 = 0

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

Mode Bits

OSCOFF, CPUOFF, and GIE are not affected.

Example

The 20-bit operand at address EDE is shifted right by one position. The MSB is loaded with 1.

SETC ; Prepare carry for MSB RRCX.A EDE ; EDE = EDE » 1 + 80000h



# **Example** The word in R6 is shifted right by 12 positions.

RPT #12 RRCX.W R6 ; R6 = R6 » 12. R6.19:16 = 0

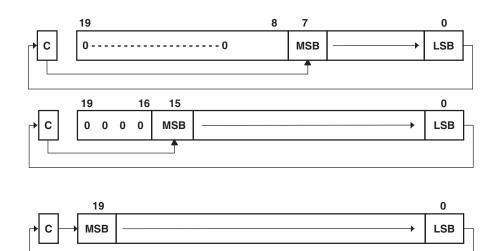


Figure 4-50. Rotate Right Through Carry RRCX(.B,.A) - Register Mode

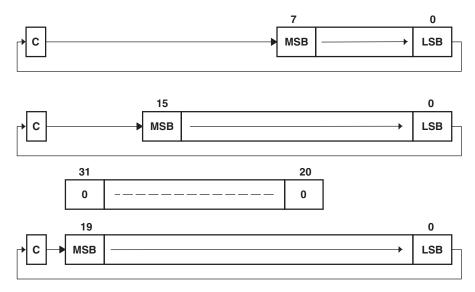


Figure 4-51. Rotate Right Through Carry RRCX(.B,.A) - Non-Register Mode



## 4.6.3.29 RRUM

Rotate right through carry the 20-bit CPU register content RRUM.A RRUM.[W] Rotate right through carry the 16-bit CPU register content

RRUM.A #n,Rdst **Syntax**  $1 \le n \le 4$ RRUM.W #n, Rdst Of RRUM #n, Rdst  $1 \le n \le 4$ 

Operation  $0 \rightarrow MSB \rightarrow MSB-1 \dots LSB+1 \rightarrow LSB \rightarrow C$ 

Description The destination operand is shifted right by one, two, three, or four bit positions as

shown in Figure 4-52. Zero is shifted into the MSB, the LSB is shifted into the carry bit. RRUM works like an unsigned division by 2, 4, 8, or 16. The word instruction RRUM.W

clears the bits Rdst.19:16.

Note: This instruction does not use the extension word.

**Status Bits** Set if result is negative

> .A: Rdst.19 = 1, reset if Rdst.19 = 0.W: Rdst.15 = 1, reset if Rdst.15 = 0Set if result is zero, reset otherwise

Z: C: Loaded from the LSB (n = 1), LSB+1 (n = 2), LSB+2 (n = 3), or LSB+3 (n = 4)

V:

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected. The unsigned address-word in R5 is divided by 16. **Example** 

> RRUM.A #4,R5  $; R5 = R5 \gg 4. R5/16$

**Example** The word in R6 is shifted right by one bit. The MSB R6.15 is loaded with 0.

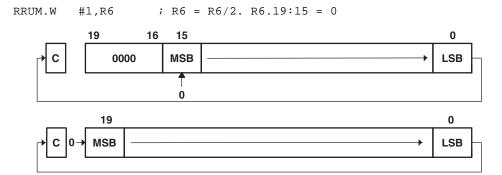


Figure 4-52. Rotate Right Unsigned RRUM[.W] and RRUM.A



## 4.6.3.30 RRUX

RRUX.A Shift right unsigned the 20-bit CPU register content Shift right unsigned the 16-bit CPU register content Shift right unsigned the 8-bit CPU register content Shift right unsigned the 8-bit CPU register content

Syntax RRUX.A Rdst

RRUX.W Rdst RRUX Rdst RRUX.B Rdst

 $\textbf{Operation} \hspace{1cm} \textbf{C=0} \rightarrow \textbf{MSB} \rightarrow \textbf{MSB-1} \; ... \; \textbf{LSB+1} \rightarrow \textbf{LSB} \rightarrow \textbf{C}$ 

**Description** RRUX is valid for register mode only: the destination operand is shifted right by one bit

position as shown in Figure 4-53. The word instruction RRUX.W clears the bits Rdst.19:16. The byte instruction RRUX.B clears the bits Rdst.19:8. Zero is shifted into

the MSB, the LSB is shifted into the carry bit.

Status Bits N: Set if result is negative

.A: dst.19 = 1, reset if dst.19 = 0 .W: dst.15 = 1, reset if dst.15 = 0 .B: dst.7 = 1, reset if dst.7 = 0

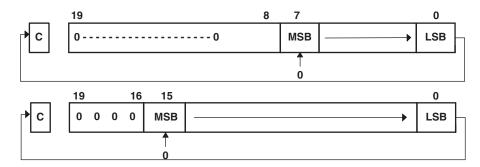
Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected. **Example** The word in R6 is shifted right by 12 positions.

RPT #12 RRUX.W R6 ; R6 = R6 » 12. R6.19:16 = 0



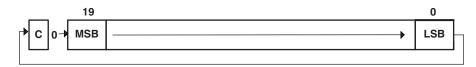


Figure 4-53. Rotate Right Unsigned RRUX(.B,.A) - Register Mode



## 4.6.3.31 SBCX

\* **SBCX.A** Subtract borrow (.NOT. carry) from destination address-word

\* **SBCX.[W]** Subtract borrow (.NOT. carry) from destination word \* **SBCX.B** Subtract borrow (.NOT. carry) from destination byte

Syntax SBCX.A dst

SBCX dst Of SBCX.W dst

SBCX.B dst

**Operation**  $dst + 0FFFFFh + C \rightarrow dst$ 

 $dst + 0FFFFh + C \rightarrow dst$  $dst + 0FFh + C \rightarrow dst$ 

Emulation SBCX.A #0,dst

SBCX #0,dst SBCX.B #0,dst

**Description** The carry bit (C) is added to the destination operand minus one. The previous contents

of the destination are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

Set to 1 if no borrow, reset if borrow

/: Set if an arithmetic overflow occurs, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 8-bit counter pointed to by R13 is subtracted from a 16-bit counter pointed to by

R12.

SUBX.B @R13,0(R12) ; Subtract LSDs

SBCX.B 1(R12) ; Subtract carry from MSD

## NOTE: Borrow implementation

The borrow is treated as a .NOT. carry:

Borrow	Carry Bit			
Yes	0			
No	1			



#### 4.6.3.32 SUBX

**SUBX.A** Subtract source address-word from destination address-word

**SUBX.[W]** Subtract source word from destination word **SUBX.B** Subtract source byte from destination byte

Syntax SUBX.A src,dst

SUBX src,dst Of SUBX.W src,dst

SUBX.B src,dst

Operation (.

 $(.not. src) + 1 + dst \rightarrow dst$  or  $dst - src \rightarrow dst$ 

**Description** 

The source operand is subtracted from the destination operand. This is done by adding the 1s complement of the source + 1 to the destination. The source operand is not affected. The result is written to the destination operand. Both operands may be located in the full address space.

**Status Bits** 

N: Set if result is negative (src > dst), reset if positive (src ≤ dst)

Z: Set if result is zero (src = dst), reset otherwise (src  $\neq$  dst)

C: Set if there is a carry from the MSB, reset otherwise

V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)

**Mode Bits** 

OSCOFF, CPUOFF, and GIE are not affected.

**Example** 

A 20-bit constant 87654h is subtracted from EDE (LSBs) and EDE+2 (MSBs).

```
SUBX.A #87654h,EDE ; Subtract 87654h from EDE+2|EDE
```

# **Example**

A table word pointed to by R5 (20-bit address) is subtracted from R7. Jump to label TONI if R7 contains zero after the instruction. R5 is auto-incremented by two. R7.19:16 = 0.

```
SUBX.W @R5+,R7 ; Subtract table number from R7. R5 + 2 JZ TONI ; R7 = @R5 (before subtraction) ... ; R7 <> @R5 (before subtraction)
```

#### **Example**

Byte CNT is subtracted from the byte R12 points to in the full address space. Address of CNT is within PC  $\pm$  512 K.

```
SUBX.B CNT,0(R12) ; Subtract CNT from @R12
```

Note: Use SUBA for the following two cases for better density and execution.

```
SUBX.A Rsrc,Rdst
SUBX.A #imm20,Rdst
```



# 4.6.3.33 SUBCX

SUBCX.A Subtract source address-word with carry from destination address-word

SUBCX.[W] SUBCX.B Subtract source word with carry from destination word Subtract source byte with carry from destination byte

Syntax SUBCX.A src,dst

SUBCX src,dst Of SUBCX.W src,dst

SUBCX.B src,dst

Operation

 $(.not. src) + C + dst \rightarrow dst$  or  $dst - (src - 1) + C \rightarrow dst$ 

Description

The source operand is subtracted from the destination operand. This is made by adding the 1s complement of the source + carry to the destination. The source operand is not affected, the result is written to the destination operand. Both operands may be located in the full address space.

**Status Bits** 

N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB, reset otherwise

V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow).

**Mode Bits** 

OSCOFF, CPUOFF, and GIE are not affected.

**Example** 

A 20-bit constant 87654h is subtracted from R5 with the carry from the previous

instruction.

```
SUBCX.A #87654h,R5 ; Subtract 87654h + C from R5
```

# **Example**

A 48-bit number (3 words) pointed to by R5 (20-bit address) is subtracted from a 48-bit counter in RAM, pointed to by R7. R5 auto-increments to point to the next 48-bit number.

```
SUBX.W @R5+,0(R7) ; Subtract LSBs. R5 + 2 
SUBCX.W @R5+,2(R7) ; Subtract MIDs with C. R5 + 2 
SUBCX.W @R5+,4(R7) ; Subtract MSBs with C. R5 + 2
```

#### **Example**

Byte CNT is subtracted from the byte R12 points to. The carry of the previous instruction is used. 20-bit addresses.

```
SUBCX.B &CNT,0(R12); Subtract byte CNT from @R12
```



## 4.6.3.34 SWPBX

**SWPBX.A** Swap bytes of lower word SWPBX.[W] Swap bytes of word

SWPBX.A dst **Syntax** 

> SWPBX dst Or SWPBX.W dst

Operation dst.15:8 ↔ dst.7:0

**Description** Register mode: Rn.15:8 are swapped with Rn.7:0. When the .A extension is used,

Rn.19:16 are unchanged. When the .W extension is used, Rn.19:16 are cleared.

Other modes: When the .A extension is used, bits 31:20 of the destination address are cleared, bits 19:16 are left unchanged, and bits 15:8 are swapped with bits 7:0. When the .W extension is used, bits 15:8 are swapped with bits 7:0 of the addressed word.

**Status Bits** Status bits are not affected.

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected. **Example** Exchange the bytes of RAM address-word EDE

> MOVX.A #23456h, &EDE ; 23456h -> EDE SWPBX.A EDE ; 25634h -> EDE

#### **Example** Exchange the bytes of R5

; 23456h -> R5 MOVA #23456h,R5 SWPBX.W ; 05634h -> R5

# **Before SWPBX.A**

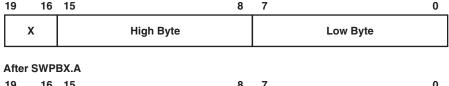
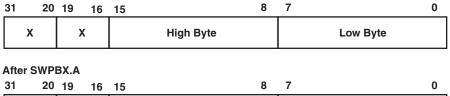




Figure 4-54. Swap Bytes SWPBX.A Register Mode

#### **Before SWPBX.A**

31



8

7



Figure 4-55. Swap Bytes SWPBX.A In Memory



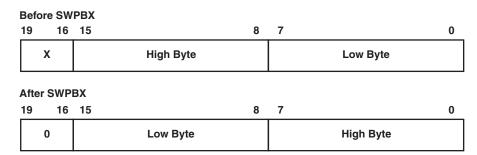


Figure 4-56. Swap Bytes SWPBX[.W] Register Mode

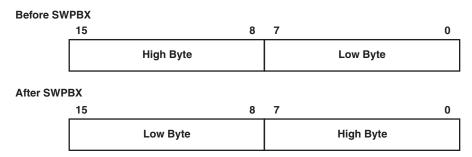


Figure 4-57. Swap Bytes SWPBX[.W] In Memory



# 4.6.3.35 SXTX

**SXTX.A** Extend sign of lower byte to address-word

**SXTX.[W]** Extend sign of lower byte to word

Syntax SXTX.A dst

SXTX dst Or SXTX.W dst

**Operation**  $dst.7 \rightarrow dst.15:8$ , Rdst.7  $\rightarrow$  Rdst.19:8 (Register mode)

**Description** Register mode: The sign of the low byte of the operand (Rdst.7) is extended into the bits

Rdst.19:8.

Other modes: SXTX.A: the sign of the low byte of the operand (dst.7) is extended into

8 7 6

dst.19:8. The bits dst.31:20 are cleared.

SXTX[.W]: the sign of the low byte of the operand (dst.7) is extended into dst.15:8.

Status Bits N: Set if result is negative, reset otherwise

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise (C = .not.Z)

V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

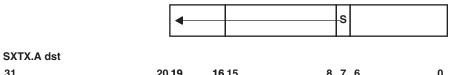
**Example** The signed 8-bit data in EDE.7:0 is sign extended to 20 bits: EDE.19:8. Bits 31:20

located in EDE+2 are cleared.

SXTX.A &EDE ; Sign extended EDE -> EDE+2/EDE

19

SXTX.A Rdst



1615

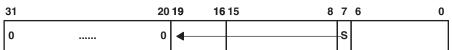
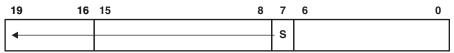


Figure 4-58. Sign Extend SXTX.A

SXTX[.W] Rdst



SXTX[.W] dst

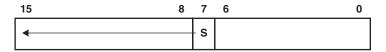


Figure 4-59. Sign Extend SXTX[.W]



# 4.6.3.36 TSTX

\* TSTX.A Test destination address-word

\* TSTX.[W] Test destination word
\* TSTX.B Test destination byte

Syntax TSTX.A dst

TSTX dst Or TSTX.W dst

TSTX.B dst

**Operation** dst + 0FFFFFh + 1

dst + 0FFFFh + 1dst + 0FFh + 1

Emulation CMPX.A #0,dst

CMPX #0,dst
CMPX.B #0,dst

**Description** The destination operand is compared with zero. The status bits are set according to the

result. The destination is not affected.

**Status Bits** N: Set if destination is negative, reset if positive

Z: Set if destination contains zero, reset otherwise

C: Set V: Reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** RAM byte LEO is tested; PC is pointing to upper memory. If it is negative, continue at

LEONEG; if it is positive but not zero, continue at LEOPOS.

TSTX.B LEO ; Test LEO
JN LEONEG ; LEO is negative
JZ LEOZERO ; LEO is zero

LEOPOS ..... ; LEO is positive but not zero

LEONEG : LEO is negative LEOZERO : LEO is zero



# 4.6.3.37 XORX

XORX.A Exclusive OR source address-word with destination address-word

**XORX.[W]** Exclusive OR source word with destination word **XORX.B** Exclusive OR source byte with destination byte

Syntax XORX.A src,dst

XORX src,dst Of XORX.W src,dst

XORX.B src,dst

**Description** The source and destination operands are exclusively ORed. The result is placed into

the destination. The source operand is not affected. The previous contents of the destination are lost. Both operands may be located in the full address space.

**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise (carry = .not. Zero)

V: Set if both operands are negative (before execution), reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** Toggle bits in address-word CNTR (20-bit data) with information in address-word TONI

(20-bit address)

XORX.A TONI, &CNTR ; Toggle bits in CNTR

**Example** A table word pointed to by R5 (20-bit address) is used to toggle bits in R6.

XORX.W @R5,R6 ; Toggle bits in R6. R6.19:16 = 0

**Example** Reset to zero those bits in the low byte of R7 that are different from the bits in byte EDE

(20-bit address)



# 4.6.4 MSP430X Address Instructions

MSP430X address instructions are instructions that support 20-bit operands but have restricted addressing modes. The addressing modes are restricted to the Register mode and the Immediate mode. except for the MOVA instruction. Restricting the addressing modes removes the need for the additional extension-word op-code, which improves code density and execution time. The MSP430X address instructions are described in the following sections.

See Section 4.6.3 for MSP430X extended instructions and Section 4.6.2 for standard MSP430 instructions.



# 4.6.4.1 ADDA

ADDA Add 20-bit source to a 20-bit destination register

Syntax ADDA Rsrc,Rdst

ADDA #imm20,Rdst

**Description** The 20-bit source operand is added to the 20-bit destination CPU register. The previous

contents of the destination are lost. The source operand is not affected.

**Status Bits** N: Set if result is negative (Rdst.19 = 1), reset if positive (Rdst.19 = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the 20-bit result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative

numbers is positive, reset otherwise

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** R5 is increased by 0A4320h. The jump to TONI is performed if a carry occurs.

ADDA #0A4320h,R5 ; Add A4320h to 20-bit R5

JC TONI ; Jump on carry ... ; No carry occurred



# 4.6.4.2 BRA

\* **BRA** Branch to destination

**Description** An unconditional branch is taken to a 20-bit address anywhere in the full address

space. All seven source addressing modes can be used. The branch instruction is an address-word instruction. If the destination address is contained in a memory location

X, it is contained in two ascending words: X (LSBs) and (X + 2) (MSBs).

Status Bits N: Not affected

Z: Not affectedC: Not affectedV: Not affected

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Examples** Examples for all addressing modes are given.

Immediate mode: Branch to label EDE located anywhere in the 20-bit address space or

branch directly to address.

BRA #EDE ; MOVA #imm20,PC

BRA #01AA04h

Symbolic mode: Branch to the 20-bit address contained in addresses EXEC (LSBs) and EXEC+2 (MSBs). EXEC is located at the address (PC + X) where X is within +32 K.

Indirect addressing.

BRA EXEC ; MOVA z16(PC),PC

Note: If the 16-bit index is not sufficient, a 20-bit index may be used with the following

instruction.

 ${\tt MOVX.A}$   ${\tt EXEC,PC}$  ; 1M byte range with 20-bit index

Absolute mode: Branch to the 20-bit address contained in absolute addresses EXEC

(LSBs) and EXEC+2 (MSBs). Indirect addressing.

BRA &EXEC ; MOVA &abs20,PC

Register mode: Branch to the 20-bit address contained in register R5. Indirect R5.

BRA R5 ; MOVA R5, PC

Indirect mode: Branch to the 20-bit address contained in the word pointed to by register

R5 (LSBs). The MSBs have the address (R5 + 2). Indirect, indirect R5.

BRA @R5 ; MOVA @R5,PC



Indirect, Auto-Increment mode: Branch to the 20-bit address contained in the words pointed to by register R5 and increment the address in R5 afterwards by 4. The next time the S/W flow uses R5 as a pointer, it can alter the program execution due to access to the next address in the table pointed to by R5. Indirect, indirect R5.

BRA @R5+ ; MOVA @R5+,PC. R5 + 4

Indexed mode: Branch to the 20-bit address contained in the address pointed to by register (R5 + X) (for example, a table with addresses starting at X). (R5 + X) points to the LSBs, (R5 + X + 2) points to the MSBs of the address. X is within R5 + 32 K. Indirect, indirect (R5 + X).

BRA X(R5) ; MOVA z16(R5), PC

Note: If the 16-bit index is not sufficient, a 20-bit index X may be used with the following instruction:

MOVX.A X(R5),PC ; 1M byte range with 20-bit index



# 4.6.4.3 CALLA

CALLA Call a subroutine

Syntax CALLA dst

**Operation** dst → tmp 20-bit dst is evaluated and stored

 $SP - 2 \rightarrow SP$ 

PC.19:16 → @SP updated PC with return address to TOS (MSBs)

 $SP - 2 \rightarrow SP$ 

PC.15:0 → @SP updated PC to TOS (LSBs)

tmp → PC saved 20-bit dst to PC

**Description** A subroutine call is made to a 20-bit address anywhere in the full address space. All

seven source addressing modes can be used. The call instruction is an address-word instruction. If the destination address is contained in a memory location X, it is contained in two ascending words, X (LSBs) and (X + 2) (MSBs). Two words on the stack are needed for the return address. The return is made with the instruction RETA.

Status Bits N: Not affected

Z: Not affectedC: Not affectedV: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Examples** Examples for all addressing modes are given.

Immediate mode: Call a subroutine at label EXEC or call directly an address.

CALLA #EXEC ; Start address EXEC CALLA #01AA04h ; Start address 01AA04h

Symbolic mode: Call a subroutine at the 20-bit address contained in addresses EXEC (LSBs) and EXEC+2 (MSBs). EXEC is located at the address (PC + X) where X is

within +32 K. Indirect addressing.

CALLA EXEC ; Start address at @EXEC. z16(PC)

Absolute mode: Call a subroutine at the 20-bit address contained in absolute addresses

EXEC (LSBs) and EXEC+2 (MSBs). Indirect addressing.

CALLA &EXEC ; Start address at @EXEC

Register mode: Call a subroutine at the 20-bit address contained in register R5. Indirect

R5.

CALLA R5 ; Start address at @R5

Indirect mode: Call a subroutine at the 20-bit address contained in the word pointed to by register R5 (LSBs). The MSBs have the address (R5 + 2). Indirect, indirect R5.

CALLA @R5 ; Start address at @R5



Indirect, Auto-Increment mode: Call a subroutine at the 20-bit address contained in the words pointed to by register R5 and increment the 20-bit address in R5 afterwards by 4. The next time the S/W flow uses R5 as a pointer, it can alter the program execution due to access to the next word address in the table pointed to by R5. Indirect, indirect R5.

CALLA @R5+ ; Start address at @R5. R5 + 4

Indexed mode: Call a subroutine at the 20-bit address contained in the address pointed to by register (R5 + X); for example, a table with addresses starting at X. (R5 + X) points to the LSBs, (R5 + X + 2) points to the MSBs of the word address. X is within R5 + 32 K. Indirect, indirect (R5 + X).

CALLA X(R5) ; Start address at @(R5+X). z16(R5)



# 4.6.4.4 CLRA

\* CLRA Clear 20-bit destination register

**Description** The destination register is cleared.

**Status Bits** Status bits are not affected.

**Example** The 20-bit value in R10 is cleared.

CLRA R10 ; 0 -> R10

Example



#### 4.6.4.5 CMPA

CMPA Compare the 20-bit source with a 20-bit destination register

Syntax CMPA Rsrc,Rdst

CMPA #imm20,Rdst

**Operation** (.not. src) + 1 + Rdst or Rdst – src

**Description** The 20-bit source operand is subtracted from the 20-bit destination CPU register. This

is made by adding the 1s complement of the source + 1 to the destination register. The

result affects only the status bits.

Status Bits N: Set if result is negative (src > dst), reset if positive (src ≤ dst)

Z: Set if result is zero (src = dst), reset otherwise (src  $\neq$  dst)

C: Set if there is a carry from the MSB, reset otherwise

V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset

The 20-bit values in R5 and R6 are compared. If R5 is greater than (signed) or equal to

otherwise (no overflow)

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

**Example** A 20-bit immediate operand and R6 are compared. If they are equal, the program

continues at label EQUAL.

CMPA #12345h,R6 ; Compare R6 with 12345h

R6, the program continues at label GRE.

CMPA R6,R5 ; Compare R6 with R5 (R5 - R6)

JGE GRE ; R5 >= R6 ... ; R5 < R6



# 4.6.4.6 DECDA

\* DECDA Double-decrement 20-bit destination register

DECDA Rdst **Syntax** Operation  $Rdst - 2 \rightarrow Rdst$ SUBA #2,Rdst **Emulation** 

Description The destination register is decremented by two. The original contents are lost.

**Status Bits** Set if result is negative, reset if positive

> Z: Set if Rdst contained 2, reset otherwise C:

Reset if Rdst contained 0 or 1, set otherwise

V: Set if an arithmetic overflow occurs, otherwise reset

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected. **Example** The 20-bit value in R5 is decremented by 2.

> R5 DECDA ; Decrement R5 by two



# 4.6.4.7 INCDA

\* **INCDA** Double-increment 20-bit destination register

**Description** The destination register is incremented by two. The original contents are lost.

Status Bits N: Set if result is negative, reset if positive

Z: Set if Rdst contained 0FFFEh, reset otherwise Set if Rdst contained 0FFFEh, reset otherwise Set if Rdst contained 0FEh, reset otherwise

C: Set if Rdst contained 0FFFEh or 0FFFFh, reset otherwise Set if Rdst contained 0FFFEh or 0FFFh, reset otherwise Set if Rdst contained 0FEh or 0FFh, reset otherwise

V: Set if Rdst contained 07FFFEh or 07FFFFh, reset otherwise Set if Rdst contained 07FFEh or 07FFFh, reset otherwise Set if Rdst contained 07Eh or 07Fh, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected. **Example** The 20-bit value in R5 is incremented by two.

INCDA R5 ; Increment R5 by two



# 4.6.4.8 MOVA

**MOVA** Move the 20-bit source to the 20-bit destination

**Syntax** MOVA Rsrc, Rdst

> MOVA #imm20, Rdst MOVA z16(Rsrc), Rdst

> MOVA EDE, Rdst MOVA &abs20, Rdst MOVA @Rsrc,Rdst MOVA @Rsrc+, Rdst MOVA Rsrc,z16(Rdst) MOVA Rsrc, &abs20

Operation  $src \rightarrow Rdst$ 

 $Rsrc \rightarrow dst$ 

Description The 20-bit source operand is moved to the 20-bit destination. The source operand is not

affected. The previous content of the destination is lost.

**Status Bits** N: Not affected

> Z: Not affected C: Not affected

Not affected V:

**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.

Copy 20-bit value in R9 to R8 **Examples** 

> MOVA R9,R8 ; R9 -> R8

> > Write 20-bit immediate value 12345h to R12

#12345h,R12 ; 12345h -> R12 MOVA

Copy 20-bit value addressed by (R9 + 100h) to R8. Source operand in addresses (R9 +

100h) LSBs and (R9 + 102h) MSBs.

; Index: + 32 K. 2 words transferred MOVA 100h(R9),R8

Move 20-bit value in 20-bit absolute addresses EDE (LSBs) and EDE+2 (MSBs) to R12

&EDE,R12 ; &EDE -> R12. 2 words transferred MOVA

Move 20-bit value in 20-bit addresses EDE (LSBs) and EDE+2 (MSBs) to R12. PC

index ± 32 K.

MOVA EDE,R12 ; EDE -> R12. 2 words transferred

Copy 20-bit value R9 points to (20 bit address) to R8. Source operand in addresses

@R9 LSBs and @(R9 + 2) MSBs.

MOVA @R9,R8 ; @R9 -> R8. 2 words transferred



Copy 20-bit value R9 points to (20 bit address) to R8. R9 is incremented by four afterwards. Source operand in addresses @R9 LSBs and @(R9 + 2) MSBs.

MOVA @R9+,R8 ; @R9 -> R8. R9 + 4. 2 words transferred.

Copy 20-bit value in R8 to destination addressed by (R9 + 100h). Destination operand in addresses @(R9 + 100h) LSBs and @(R9 + 102h) MSBs.

MOVA R8,100h(R9) ; Index: +- 32 K. 2 words transferred

Move 20-bit value in R13 to 20-bit absolute addresses EDE (LSBs) and EDE+2 (MSBs)

MOVA R13, &EDE ; R13 -> EDE. 2 words transferred

Move 20-bit value in R13 to 20-bit addresses EDE (LSBs) and EDE+2 (MSBs). PC index  $\pm$  32 K.

MOVA R13,EDE ; R13 -> EDE. 2 words transferred



# 4.6.4.9 RETA

\* **RETA** Return from subroutine

Syntax RETA

**Operation** @SP  $\rightarrow$  PC.15:0 LSBs (15:0) of saved PC to PC.15:0

 $SP + 2 \rightarrow SP$ 

@SP  $\rightarrow$  PC.19:16 MSBs (19:16) of saved PC to PC.19:16

 $SP + 2 \rightarrow SP$ 

Emulation MOVA @SP+,PC

**Description** The 20-bit return address information, pushed onto the stack by a CALLA instruction, is

restored to the PC. The program continues at the address following the subroutine call. The SR bits SR.11:0 are not affected. This allows the transfer of information with these

bits.

Status Bits N: Not affected

Z: Not affectedC: Not affected

V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** Call a subroutine SUBR from anywhere in the 20-bit address space and return to the

address after the CALLA

CALLA #SUBR ; Call subroutine starting at SUBR

... ; Return by RETA to here

SUBR PUSHM.A #2,R14 ; Save R14 and R13 (20 bit data)

... ; Subroutine code

POPM.A #2,R14 ; Restore R13 and R14 (20 bit data) RETA ; Return (to full address space)



# 4.6.4.10 TSTA

\* **TSTA** Test 20-bit destination register

Syntax TSTA Rdst

Operation dst + 0FFFFFh + 1

dst + 0FFFFh + 1dst + 0FFh + 1

Emulation CMPA #0,Rdst

**Description** The destination register is compared with zero. The status bits are set according to the

result. The destination register is not affected.

**Status Bits** N: Set if destination register is negative, reset if positive

Z: Set if destination register contains zero, reset otherwise

C: Set V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 20-bit value in R7 is tested. If it is negative, continue at R7NEG; if it is positive but

not zero, continue at R7POS.

TSTA R7 ; Test R7
JN R7NEG ; R7 is negative
JZ R7ZERO ; R7 is zero

R7POS ..... ; R7 is positive but not zero

R7NEG ..... ; R7 is negative R7ZERO ..... ; R7 is zero



# 4.6.4.11 SUBA

SUBA Subtract 20-bit source from 20-bit destination register

Syntax SUBA Rsrc,Rdst

SUBA #imm20,Rdst

**Operation** (.not.src) + 1 + Rdst  $\rightarrow$  Rdst or Rdst – src  $\rightarrow$  Rdst

**Description** The 20-bit source operand is subtracted from the 20-bit destination register. This is

made by adding the 1s complement of the source + 1 to the destination. The result is

written to the destination register, the source is not affected.

Status Bits N: Set if result is negative (src > dst), reset if positive (src ≤ dst)

Z: Set if result is zero (src = dst), reset otherwise (src  $\neq$  dst)

C: Set if there is a carry from the MSB (Rdst.19), reset otherwise

V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset

otherwise (no overflow)

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

**Example** The 20-bit value in R5 is subtracted from R6. If a carry occurs, the program continues at

label TONI.

SUBA R5,R6 ; R6 - R5 -> R6

JC TONI ; Carry occurred

... ; No carry



# **Basic Clock Module+**

The basic clock module+ provides the clocks for MSP430x2xx devices. This chapter describes the operation of the basic clock module+ of the MSP430x2xx device family.

Topic Page

5.1	Basic Clock Module+ Introduction	273
5.2	Basic Clock Module+ Operation	275
5.3	Basic Clock Module+ Registers	282



# 5.1 Basic Clock Module+ Introduction

The basic clock module+ supports low system cost and ultralow power consumption. Using three internal clock signals, the user can select the best balance of performance and low power consumption. The basic clock module+ can be configured to operate without any external components, with one external resistor, with one or two external crystals, or with resonators, under full software control.

The basic clock module+ includes two, three or four clock sources:

- LFXT1CLK: Low-frequency/high-frequency oscillator that can be used with low-frequency watch
  crystals or external clock sources of 32768 Hz or with standard crystals, resonators, or external clock
  sources in the 400-kHz to 16-MHz range.
- XT2CLK: Optional high-frequency oscillator that can be used with standard crystals, resonators, or external clock sources in the 400-kHz to 16-MHz range.
- DCOCLK: Internal digitally controlled oscillator (DCO).
- VLOCLK: Internal very low power, low frequency oscillator with 12-kHz typical frequency.

Three clock signals are available from the basic clock module+:

- ACLK: Auxiliary clock. ACLK is software selectable as LFXT1CLK or VLOCLK. ACLK is divided by 1, 2, 4, or 8. ACLK is software selectable for individual peripheral modules.
- MCLK: Master clock. MCLK is software selectable as LFXT1CLK, VLOCLK, XT2CLK (if available onchip), or DCOCLK. MCLK is divided by 1, 2, 4, or 8. MCLK is used by the CPU and system.
- SMCLK: Sub-main clock. SMCLK is software selectable as LFXT1CLK, VLOCLK, XT2CLK (if available on-chip), or DCOCLK. SMCLK is divided by 1, 2, 4, or 8. SMCLK is software selectable for individual peripheral modules.

The block diagram of the basic clock module+ in the MSP430F2xx devices is shown in Figure 5-1.

The block diagram of the basic clock module+ in the MSP430AFE2xx devices is shown in Figure 5-2.



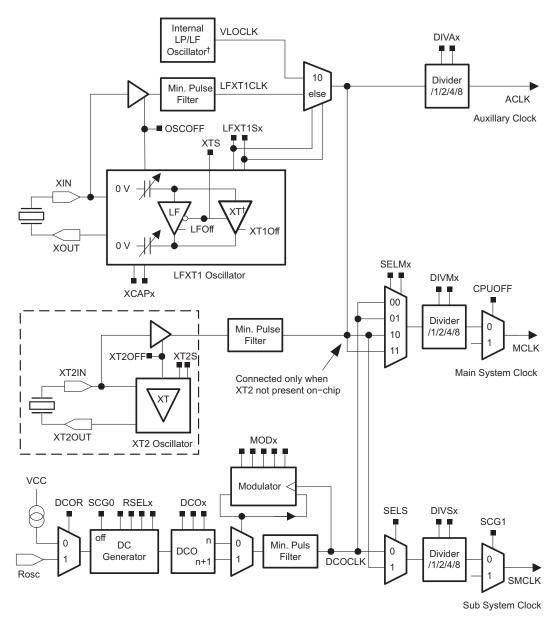


Figure 5-1. Basic Clock Module+ Block Diagram - MSP430F2xx

# NOTE: † Device-Specific Clock Variations

Not all clock features are available on all MSP430x2xx devices: MSP430G22x0: LFXT1 is not present, XT2 is not present, ROSC is not supported.

MSP430F20xx, MSP430G2xx1, MSP430G2xx2, MSP430G2xx3: LFXT1 does not support HF mode, XT2 is not present, ROSC is not supported.

MSP430x21x1: Internal LP/LF oscillator is not present, XT2 is not present, ROSC is not supported.

MSP430x21x2: XT2 is not present.

MSP430F22xx, MSP430x23x0: XT2 is not present.



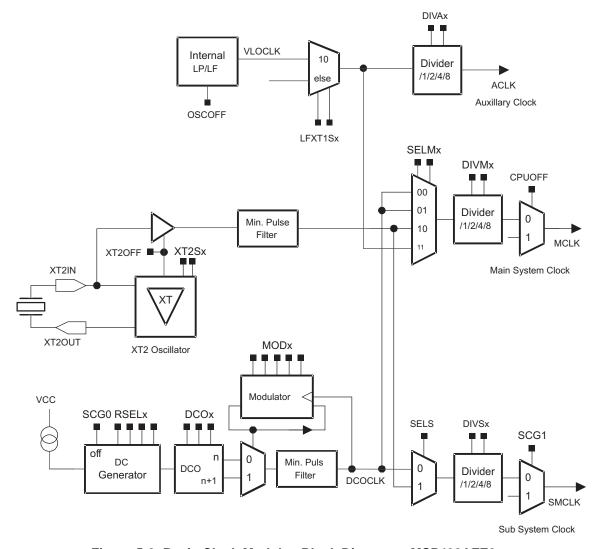


Figure 5-2. Basic Clock Module+ Block Diagram - MSP430AFE2xx

**NOTE:** LFXT1 is not present in MSP430AFE2xx devices.

# 5.2 Basic Clock Module+ Operation

After a PUC, MCLK and SMCLK are sourced from DCOCLK at ~1.1 MHz (see the device-specific data sheet for parameters) and ACLK is sourced from LFXT1CLK in LF mode with an internal load capacitance of 6 pF.

Status register control bits SCG0, SCG1, OSCOFF, and CPUOFF configure the MSP430 operating modes and enable or disable portions of the basic clock module+ (see the *System Resets, Interrupts and Operating Modes* chapter). The DCOCTL, BCSCTL1, BCSCTL2, and BCSCTL3 registers configure the basic clock module+.

The basic clock module+ can be configured or reconfigured by software at any time during program execution, for example:

CLR.B &DCOCTL ; Select lowest DCOx ; and MODx settings
BIS.B #RSEL2+RSEL1+RSEL0,&BCSCTL1 ; Select range 7
BIS.B #DCO2+DCO1+DCO0,&DCOCTL ; Select max DCO tap



# 5.2.1 Basic Clock Module+ Features for Low-Power Applications

Conflicting requirements typically exist in battery-powered applications:

- · Low clock frequency for energy conservation and time keeping
- High clock frequency for fast reaction to events and fast burst processing capability
- Clock stability over operating temperature and supply voltage

The basic clock module+ addresses the above conflicting requirements by allowing the user to select from the three available clock signals: ACLK, MCLK, and SMCLK. For optimal low-power performance, ACLK can be sourced from a low-power 32768-Hz watch crystal (if available), providing a stable time base for the system and low-power standby operation, or from the internal low-frequency oscillator when crystal-accurate time keeping is not required. The MCLK can be configured to operate from the on-chip DCO that can be activated when requested by interrupt-driven events. The SMCLK can be configured to operate from a crystal or the DCO, depending on peripheral requirements. A flexible clock distribution and divider system is provided to fine tune the individual clock requirements.

# 5.2.2 Internal Very-Low-Power Low-Frequency Oscillator (VLO)

The internal very-low-power low-frequency oscillator (VLO) provides a typical frequency of 12 kHz (see device-specific data sheet for parameters) without requiring a crystal. VLOCLK source is selected by setting LFXT1Sx = 10 when XTS = 0. The OSCOFF bit disables the VLO for LPM4. The LFXT1 crystal oscillators are disabled when the VLO is selected reducing current consumption. The VLO consumes no power when not being used.

Devices without LFXT1 (for example, the MSP430G22x0) should be configured to use the VLO as ACLK.

#### 5.2.3 LFXT1 Oscillator

The LFXT1 oscillator is not implemented in the MSP430G22x0 device family.

The LFXT1 oscillator supports ultra-low current consumption using a 32768-Hz watch crystal in LF mode (XTS = 0). A watch crystal connects to XIN and XOUT without any other external components. The software-selectable XCAPx bits configure the internally provided load capacitance for the LFXT1 crystal in LF mode. This capacitance can be selected as 1 pF, 6 pF, 10 pF, or 12.5 pF typical. Additional external capacitors can be added if necessary.

The LFXT1 oscillator also supports high-speed crystals or resonators when in HF mode (XTS = 1, XCAPx = 00). The high-speed crystal or resonator connects to XIN and XOUT and requires external capacitors on both terminals. These capacitors should be sized according to the crystal or resonator specifications. When LFXT1 is in HF mode, the LFXT1Sx bits select the range of operation.

LFXT1 may be used with an external clock signal on the XIN pin in either LF or HF mode when LFXT1Sx = 11, OSCOFF = 0, and XCAPx = 00. When used with an external signal, the external frequency must meet the data sheet parameters for the chosen mode. When the input frequency is below the specified lower limit, the LFXT1OF bit may be set preventing the CPU from being clocked with LFXT1CLK.

Software can disable LFXT1 by setting OSCOFF, if LFXT1CLK does not source SMCLK or MCLK, as shown in Figure 5-3.



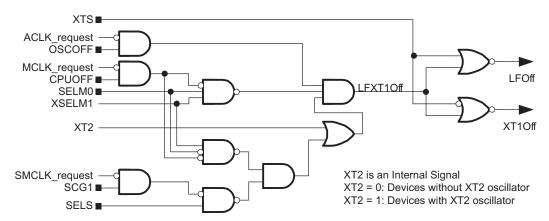


Figure 5-3. Off Signals for the LFXT1 Oscillator

#### NOTE: LFXT1 Oscillator Characteristics

Low-frequency crystals often require hundreds of milliseconds to start up, depending on the crystal.

Ultralow-power oscillators such as the LFXT1 in LF mode should be guarded from noise coupling from other sources. The crystal should be placed as close as possible to the MSP430 with the crystal housing grounded and the crystal traces guarded with ground traces.

# 5.2.4 XT2 Oscillator

Some devices have a second crystal oscillator, XT2. XT2 sources XT2CLK and its characteristics are identical to LFXT1 in HF mode. The XT2Sx bits select the range of operation of XT2. The XT2OFF bit disables the XT2 oscillator if XT2CLK is not used for MCLK or SMCLK as shown in Figure 5-4.

XT2 may be used with external clock signals on the XT2IN pin when XT2Sx = 11 and XT2OFF = 0. When used with an external signal, the external frequency must meet the data sheet parameters for XT2. When the input frequency is below the specified lower limit, the XT2OF bit may be set to prevent the CPU from being clocked with XT2CLK.

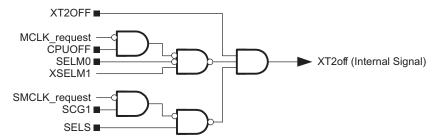


Figure 5-4. Off Signals for Oscillator XT2

# 5.2.5 Digitally-Controlled Oscillator (DCO)

The DCO is an integrated digitally controlled oscillator. The DCO frequency can be adjusted by software using the DCOx, MODx, and RSELx bits.

# 5.2.5.1 Disabling the DCO

Software can disable DCOCLK by setting SCG0 when it is not used to source SMCLK or MCLK in active mode, as shown in Figure 5-5.



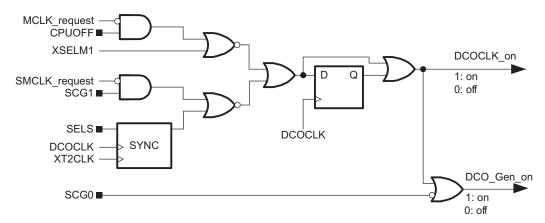


Figure 5-5. On/Off Control of DCO

# 5.2.5.2 Adjusting the DCO Frequency

After a PUC, RSELx = 7 and DCOx = 3, allowing the DCO to start at a mid-range frequency. MCLK and SMCLK are sourced from DCOCLK. Because the CPU executes code from MCLK, which is sourced from the fast-starting DCO, code execution typically begins from PUC in less than 2  $\mu$ s. The typical DCOx and RSELx ranges and steps are shown in Figure 5-6.

The frequency of DCOCLK is set by the following functions:

- The four RSELx bits select one of sixteen nominal frequency ranges for the DCO. These ranges are defined for an individual device in the device-specific data sheet.
- The three DCOx bits divide the DCO range selected by the RSELx bits into 8 frequency steps, separated by approximately 10%.
- The five MODx bits, switch between the frequency selected by the DCOx bits and the next higher frequency set by DCOx+1. When DCOx = 07h, the MODx bits have no effect because the DCO is already at the highest setting for the selected RSELx range.

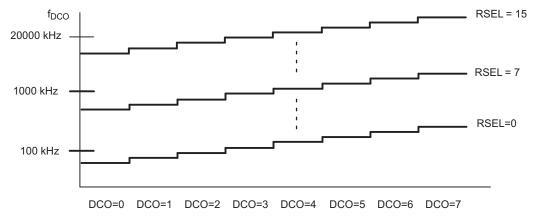


Figure 5-6. Typical DCOx Range and RSELx Steps

Each MSP430F2xx device (and most MSP430G2xx devices; see device-specific data sheets) has calibrated DCOCTL and BCSCTL1 register settings for specific frequencies stored in information memory segment A. To use the calibrated settings, the information is copied into the DCOCTL and BCSCTL1 registers. The calibrated settings affect the DCOx, MODx, and RSELx bits, and clear all other bits, except XT2OFF which remains set. The remaining bits of BCSCTL1 can be set or cleared as needed with BIS.B or BIC.B instructions.

```
; Set DCO to 1 MHz:
CLR.B &DCOCTL ; Select lowest DCOx ; and MODx settings
```



MOV.B &CALBC1\_1MHZ,&BCSCTL1 ; Set range
MOV.B &CALDCO 1MHZ,&DCOCTL ; Set DCO step + modulation

# 5.2.5.3 Using an External Resistor (Rosc) for the DCO

Some MSP430F2xx devices provide the option to source the DCO current through an external resistor,  $R_{\rm OSC}$ , tied to  $DV_{\rm CC}$ , when DCOR = 1. In this case, the DCO has the same characteristics as MSP430x1xx devices, and the RSELx setting is limited to 0 to 7 with the RSEL3 ignored. This option provides an additional method to tune the DCO frequency by varying the resistor value. See the device-specific data sheet for parameters.

# 5.2.6 DCO Modulator

The modulator mixes two DCO frequencies,  $f_{DCO}$  and  $f_{DCO+1}$  to produce an intermediate effective frequency between  $f_{DCO}$  and  $f_{DCO+1}$  and spread the clock energy, reducing electromagnetic interference (EMI). The modulator mixes  $f_{DCO}$  and  $f_{DCO+1}$  for 32 DCOCLK clock cycles and is configured with the MODx bits. When MODx = 0 the modulator is off.

The modulator mixing formula is:

$$t = (32 - MODx) \times t_{DCO} + MODx \times t_{DCO+1}$$

Because  $f_{DCO}$  is lower than the effective frequency and  $f_{DCO+1}$  is higher than the effective frequency, the error of the effective frequency integrates to zero. It does not accumulate. The error of the effective frequency is zero every 32 DCOCLK cycles. Figure 5-7 shows the modulator operation.

The modulator settings and DCO control are configured with software. The DCOCLK can be compared to a stable frequency of known value and adjusted with the DCOx, RSELx, and MODx bits. See <a href="http://www.msp430.com">http://www.msp430.com</a> for application notes and example code on configuring the DCO.

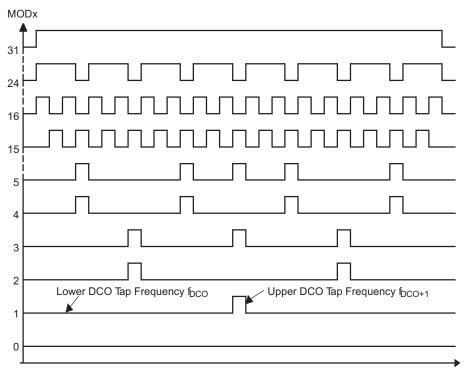


Figure 5-7. Modulator Patterns

# 5.2.7 Basic Clock Module+ Fail-Safe Operation

The basic clock module+ incorporates an oscillator-fault fail-safe feature. This feature detects an oscillator fault for LFXT1 and XT2 as shown in Figure 5-8. The available fault conditions are:

Low-frequency oscillator fault (LFXT1OF) for LFXT1 in LF mode



- High-frequency oscillator fault (LFXT1OF) for LFXT1 in HF mode
- High-frequency oscillator fault (XT2OF) for XT2

The crystal oscillator fault bits LFXT1OF, and XT2OF are set if the corresponding crystal oscillator is turned on and not operating properly. The fault bits remain set as long as the fault condition exists and are automatically cleared if the enabled oscillators function normally.

The OFIFG oscillator-fault flag is set and latched at POR or when an oscillator fault (LFXT1OF, or XT2OF) is detected. When OFIFG is set, MCLK is sourced from the DCO, and if OFIE is set, the OFIFG requests an NMI interrupt. When the interrupt is granted, the OFIE is reset automatically. The OFIFG flag must be cleared by software. The source of the fault can be identified by checking the individual fault bits.

If a fault is detected for the crystal oscillator sourcing the MCLK, the MCLK is automatically switched to the DCO for its clock source. This does not change the SELMx bit settings. This condition must be handled by user software.

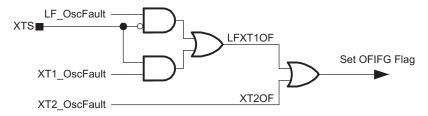


Figure 5-8. Oscillator-Fault Logic

# 5.2.7.1 Sourcing MCLK from a Crystal

After a PUC, the basic clock module+ uses DCOCLK for MCLK. If required, MCLK may be sourced from LFXT1 or XT2 - if available.

The sequence to switch the MCLK source from the DCO clock to the crystal clock (LFXT1CLK or XT2CLK) is:

- 1. Turn on the crystal oscillator and select the appropriate mode
- 2. Clear the OFIFG flag
- 3. Wait at least 50 µs
- 4. Test OFIFG, and repeat steps 2 through 4 until OFIFG remains cleared.

```
; Select LFXT1 (HF mode) for MCLK
    BIC.W #OSCOFF,SR
                                        ; Turn on osc.
    BIS.B
            #XTS,&BCSCTL1
                                        ; HF mode
    MOV.B
            #LFXT1S0,&BCSCTL3
                                        ; 1-3MHz Crystal
L1
    BIC.B
            #OFIFG,&IFG1
                                        ; Clear OFIFG
    MOV.W
            #0FFh,R15
                                        ; Delay
T.2
    DEC. W
            R15
    JNZ
            L_2
           #OFIFG,&IFG1
    BIT.B
                                       ; Re-test OFIFG
    JNZ
                                       ; Repeat test if needed
    BIS.B #SELM1+SELM0,&BCSCTL2
                                        ; Select LFXT1CLK
```

# 5.2.8 Synchronization of Clock Signals

When switching MCLK or SMCLK from one clock source to another, the switch is synchronized to avoid critical race conditions as shown in Figure 5-9:

- The current clock cycle continues until the next rising edge.
- The clock remains high until the next rising edge of the new clock.
- The new clock source is selected and continues with a full high period.



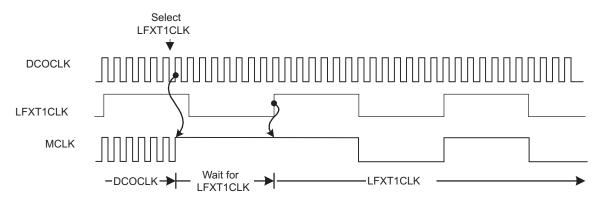


Figure 5-9. Switch MCLK from DCOCLK to LFXT1CLK



# 5.3 Basic Clock Module+ Registers

The basic clock module+ registers are listed in Table 5-1.

Table 5-1. Basic Clock Module+ Registers

Register	Short Form	Register Type	Address	Initial State
DCO control register	DCOCTL	Read/write	056h	060h with PUC
Basic clock system control 1	BCSCTL1	Read/write	057h	087h with POR <sup>(1)</sup>
Basic clock system control 2	BCSCTL2	Read/write	058h	Reset with PUC
Basic clock system control 3	BCSCTL3	Read/write	053h	005h with PUC <sup>(2)</sup>
SFR interrupt enable register 1	IE1	Read/write	000h	Reset with PUC
SFR interrupt flag register 1	IFG1	Read/write	002h	Reset with PUC

<sup>(1)</sup> Some of the register bits are also PUC initialized (see Section 5.3.2).

<sup>&</sup>lt;sup>(2)</sup> The initial state of BCSCTL3 is 000h in the MSP430AFE2xx devices.



# 5.3.1 DCOCTL, DCO Control Register

7	6	5	4	3	2	1	0		
	DCOx				MODx				
rw-0	rw-1	rw-1	rw-0	rw-0	rw-0	rw-0	rw-0		
DCOx	Bits 7-5	. ,	DCO frequency select. These bits select which of the eight discrete DCO frequencies within the range defined by the RSELx setting is selected.						
MODx	Bits 4-0	Modulator selection. These bits define how often the f $_{DCO+1}$ frequency is used within a period of 32 DCOCLK cycles. During the remaining clock cycles (32-MOD) the f $_{DCO}$ frequency is used. Not useable when DCOx = 7.							

# 5.3.2 BCSCTL1, Basic Clock System Control Register 1

7	6		5	4	3	2	1	0	
XT2OFF	XTS <sup>(1)(2</sup>	2)	DIVAx			RSELx			
rw-(1)	rw-(0)		rw-(0)	rw-(0)	rw-0	rw-1	rw-1	rw-1	
XT2OFF	Bit 7	XT2	off. This bit turns	off the XT2 oscilla	ator				
		0	XT2 is on						
		1	XT2 is off if it is	off if it is not used for MCLK or SMCLK.					
XTS	Bit 6	LFX.	T1 mode select.						
		0	Low-frequency r	node					
		1	High-frequency	node					
DIVAx	Bits 5-4	Divid	der for ACLK						
		00	/1						
		01	/2						
		10	/4						
		11	/8						
RSELx	Bits 3-0		ge select. Sixteen ng RSELx = 0. RS			ilable. The lowest f	requency range i	s selected by	

<sup>(1)</sup> XTS = 1 is not supported in MSP430x20xx and MSP430G2xx devices (see Figure 5-1 and Figure 5-2 for details on supported settings for all devices).

<sup>&</sup>lt;sup>(2)</sup> This bit is reserved in the MSP430AFE2xx devices.



# 5.3.3 BCSCTL2, Basic Clock System Control Register 2

7	7 6		5	4	3	2	1	0
\$	SELMx		DIVMx		SELS	DIV	/Sx	DCOR (1)(2)
rw-0	rw-0	)	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
SELMx	Bits 7-6	Sele	ct MCLK. These b	its select the MC	LK source.			
		00	DCOCLK					
		01	DCOCLK					
		10	XT2CLK when a on-chip.	XT2 oscillator pre	esent on-chip. LFX	T1CLK or VLOCLI	K when XT2 osci	llator not present
		11	LFXT1CLK or V	/LOCLK				
DIVMx	Bits 5-4	Divid	ler for MCLK					
		00	/1					
		01	/2					
		10	/4					
		11	/8					
SELS	Bit 3	Sele	ct SMCLK. This bi	t selects the SM	CLK source.			
		0	DCOCLK					
		1	XT2CLK when 2	XT2 oscillator pre	esent. LFXT1CLK	or VLOCLK when 2	XT2 oscillator no	t present
DIVSx	Bits 2-1	Divid	ler for SMCLK					
		00	/1					
		01	/2					
		10	/4					
		11	/8					
DCOR	Bit 0	DCO	resistor select. N	ot available in all	devices. See the	device-specific dat	a sheet.	
		0	Internal resistor					
		1	External resisto	r				

Does not apply to MSP430x20xx or MSP430x21xx devices.

<sup>(2)</sup> This bit is reserved in the MSP430AFE2xx devices.



# 5.3.4 BCSCTL3, Basic Clock System Control Register 3

7	6		5	4	3	2	1	0
XT2Sx			LFXT1Sx <sup>(1)</sup>		XC	CAPx <sup>(2)</sup>	XT2OF <sup>(3)</sup>	LFXT10F <sup>(2)</sup>
rw-0	rw-0 rw-0		rw-0	rw-0	rw-0	rw-1	rO	r-(1)
XT2Sx	Bits 7-6	XT2	range select. The	se bits select the	frequency range	e for XT2.		
		00	0.4- to 1-MHz	crystal or resonato	or			
		01	1- to 3-MHz cry	stal or resonator				
		10	3- to 16-MHz c	rystal or resonator	r			
		11	Digital external	0.4- to 16-MHz cl	ock source			
LFXT1Sx	Bits 5-4	Low-	frequency clock s	select and LFXT1 lency range for LF	range select. Th	nese bits select between	veen LFXT1 and '	VLO when XTS =
			n XTS = 0:	leffcy farige for Li	XII WIIEII XI3	= 1.		
		00	32768-Hz crys	tal on LEXT1				
		01	Reserved	lai Oii Li XI I				
		10		erved in MSP430F	21x1 devices)			
		11	Digital external		2171 (601665)			
			· ·		430x20xx device	es, MSP430G2xx1/	2/3)	
		00	` '	crystal or resonato			_, _,	
		01		stal or resonator				
		10	· · · · · · · · · · · · · · · · · · ·	rystal or resonator	r			
		11	Digital external	0.4- to 16-MHz cl	lock source			
		LFX	Γ1Sx definition for	MSP430AFE2xx	devices:			
		00	Reserved					
		01	Reserved					
		10	VLOCLK					
		11	Reserved					
XCAPx	Bits 3-2			election. These bits r if LFXT1Sx = 11		ctive capacitance se be 00.	een by the LFXT1	crystal when
		00	~1 pF					
		01	~6 pF					
		10	~10 pF					
		11	~12.5 pF					
XT2OF	Bit 1	XT2	oscillator fault					
		0	No fault conditi	on present				
		1	Fault condition	present				
LFXT10F	Bit 0	LFX	Γ1 oscillator fault					
		0	No fault conditi	on present				
		1	Fault condition	present				

<sup>(1)</sup> MSP430G22x0: The LFXT1Sx bits should be programmed to 10b during the initialization and start-up code to select VLOCLK (for more details refer to Digital I/O chapter). The other bits are reserved and should not be altered.

<sup>&</sup>lt;sup>(2)</sup> This bit is reserved in the MSP430AFE2xx devices.

Does not apply to MSP430x2xx, MSP430x21xx, or MSP430x22xx devices.



# 5.3.5 IE1, Interrupt Enable Register 1

7	6	5	4	3	2	1	0
						OFIE <sup>(1)</sup>	
						rw-0	_

Bits 7-2 These bits may be used by other modules. See device-specific data sheet.

OFIE Bit 1

Oscillator fault interrupt enable. This bit enables the OFIFG interrupt. Because other bits in IE1 may be used for other modules, it is recommended to set or clear this bit using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.

0 Interrupt not enabled

1 Interrupt enabled

Bits 0 This bit may be used by other modules. See device-specific data sheet.

# 5.3.6 IFG1, Interrupt Flag Register 1

1	О	5	4	3	2	1	U		
						OFIFG <sup>(1)</sup>			
						rw-1			
	Bits 7-2	These bits ma	y be used by other	modules. See device	-specific data sheet.				
OFIFG	Bit 1		Oscillator fault interrupt flag. Because other bits in IFG1 may be used for other modules, it is recommended to set or clear this bit using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.						
		0 No inter	upt pending						
		1 Interrupt	pending						

Bits 0 This bit may be used by other modules. See device-specific data sheet.

<sup>(1)</sup> MSP430G22x0: This bit should not be set.

<sup>(1)</sup> MSP430G22x0: The LFXT1 oscillator pins are not available in this device. The oscillator fault flag will always be set by hardware. The interrupt enable bit should not be set.



# **DMA Controller**

The DMA controller module transfers data from one address to another without CPU intervention. This chapter describes the operation of the DMA controller of the MSP430x2xx device family.

Topic Page

6.1	DMA Introduction	288
6.2	DMA Operation	290
6.3	DMA Registers	302
	3	



DMA Introduction www.ti.com

# 6.1 DMA Introduction

The direct memory access (DMA) controller transfers data from one address to another, without CPU intervention, across the entire address range. For example, the DMA controller can move data from the ADC12 conversion memory to RAM.

Devices that contain a DMA controller may have one, two, or three DMA channels available. Therefore, depending on the number of DMA channels available, some features described in this chapter are not applicable to all devices.

Using the DMA controller can increase the throughput of peripheral modules. It can also reduce system power consumption by allowing the CPU to remain in a low-power mode without having to awaken to move data to or from a peripheral.

The DMA controller features include:

- Up to three independent transfer channels
- Configurable DMA channel priorities
- · Requires only two MCLK clock cycles per transfer
- Byte or word and mixed byte/word transfer capability
- Block sizes up to 65535 bytes or words
- Configurable transfer trigger selections
- Selectable edge or level-triggered transfer
- Four addressing modes
- Single, block, or burst-block transfer modes

The DMA controller block diagram is shown in Figure 6-1.



www.ti.com DMA Introduction

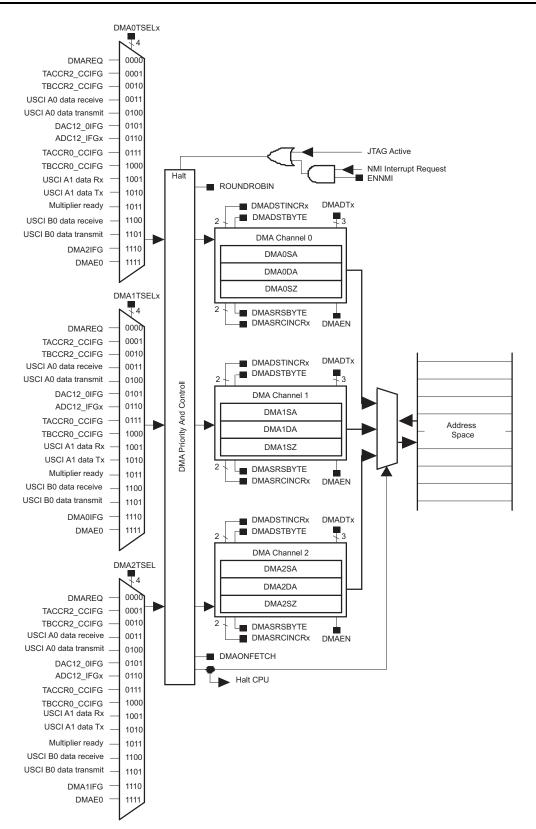


Figure 6-1. DMA Controller Block Diagram



DMA Operation www.ti.com

#### 6.2 DMA Operation

The DMA controller is configured with user software. The setup and operation of the DMA is discussed in the following sections.

#### 6.2.1 DMA Addressing Modes

The DMA controller has four addressing modes. The addressing mode for each DMA channel is independently configurable. For example, channel 0 may transfer between two fixed addresses, while channel 1 transfers between two blocks of addresses. The addressing modes are shown in Figure 6-2. The addressing modes are:

- · Fixed address to fixed address
- Fixed address to block of addresses
- Block of addresses to fixed address
- Block of addresses to block of addresses

The addressing modes are configured with the DMASRCINCRx and DMADSTINCRx control bits. The DMASRCINCRx bits select if the source address is incremented, decremented, or unchanged after each transfer. The DMADSTINCRx bits select if the destination address is incremented, decremented, or unchanged after each transfer.

Transfers may be byte-to-byte, word-to-word, byte-to-word, or word-to-byte. When transferring word-to-byte, only the lower byte of the source-word transfers. When transferring byte-to-word, the upper byte of the destination-word is cleared when the transfer occurs.

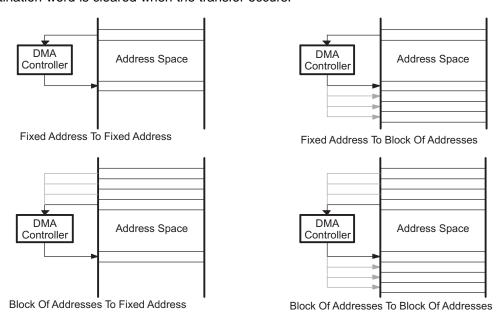


Figure 6-2. DMA Addressing Modes



www.ti.com DMA Operation

#### 6.2.2 DMA Transfer Modes

The DMA controller has six transfer modes selected by the DMADTx bits as listed in Table 6-1. Each channel is individually configurable for its transfer mode. For example, channel 0 may be configured in single transfer mode, while channel 1 is configured for burst-block transfer mode, and channel 2 operates in repeated block mode. The transfer mode is configured independently from the addressing mode. Any addressing mode can be used with any transfer mode.

Two types of data can be transferred selectable by the DMAxCTL DSTBYTE and SRCBYTE fields. The source and/or destination location can be either byte or word data. It is also possible to transfer byte to byte, word to word or any combination.

**Table 6-1. DMA Transfer Modes** 

DMADTx	Transfer Mode	Description
000	Single transfer	Each transfer requires a trigger. DMAEN is automatically cleared when DMAxSZ transfers have been made.
001	Block transfer	A complete block is transferred with one trigger. DMAEN is automatically cleared at the end of the block transfer.
010, 011	Burst-block transfer	CPU activity is interleaved with a block transfer. DMAEN is automatically cleared at the end of the burst-block transfer.
100	Repeated single transfer	Each transfer requires a trigger. DMAEN remains enabled.
101	Repeated block transfer	A complete block is transferred with one trigger. DMAEN remains enabled.
110, 111	Repeated burst-block transfer	CPU activity is interleaved with a block transfer. DMAEN remains enabled.



DMA Operation www.ti.com

#### 6.2.2.1 Single Transfer

In single transfer mode, each byte/word transfer requires a separate trigger. The single transfer state diagram is shown in Figure 6-3.

The DMAxSZ register is used to define the number of transfers to be made. The DMADSTINCRx and DMASRCINCRx bits select if the destination address and the source address are incremented or decremented after each transfer. If DMAxSZ = 0, no transfers occur.

The DMAxSA, DMAxDA, and DMAxSZ registers are copied into temporary registers. The temporary values of DMAxSA and DMAxDA are incremented or decremented after each transfer. The DMAxSZ register is decremented after each transfer. When the DMAxSZ register decrements to zero it is reloaded from its temporary register and the corresponding DMAIFG flag is set. When DMADTx = 0, the DMAEN bit is cleared automatically when DMAxSZ decrements to zero and must be set again for another transfer to occur.

In repeated single transfer mode, the DMA controller remains enabled with DMAEN = 1, and a transfer occurs every time a trigger occurs.

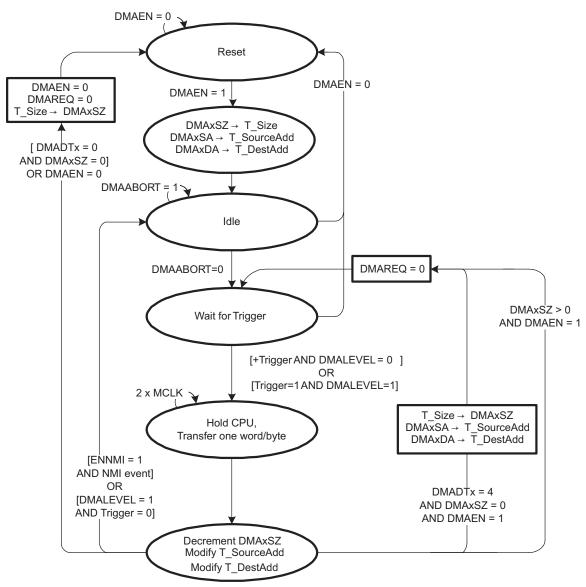


Figure 6-3. DMA Single Transfer State Diagram



www.ti.com DMA Operation

#### 6.2.2.2 Block Transfers

In block transfer mode, a transfer of a complete block of data occurs after one trigger. When DMADTx = 1, the DMAEN bit is cleared after the completion of the block transfer and must be set again before another block transfer can be triggered. After a block transfer has been triggered, further trigger signals occurring during the block transfer are ignored. The block transfer state diagram is shown in Figure 6-4.

The DMAxSZ register is used to define the size of the block and the DMADSTINCRx and DMASRCINCRx bits select if the destination address and the source address are incremented or decremented after each transfer of the block. If DMAxSZ = 0, no transfers occur.

The DMAxSA, DMAxDA, and DMAxSZ registers are copied into temporary registers. The temporary values of DMAxSA and DMAxDA are incremented or decremented after each transfer in the block. The DMAxSZ register is decremented after each transfer of the block and shows the number of transfers remaining in the block. When the DMAxSZ register decrements to zero it is reloaded from its temporary register and the corresponding DMAIFG flag is set.

During a block transfer, the CPU is halted until the complete block has been transferred. The block transfer takes 2 x MCLK x DMAxSZ clock cycles to complete. CPU execution resumes with its previous state after the block transfer is complete.

In repeated block transfer mode, the DMAEN bit remains set after completion of the block transfer. The next trigger after the completion of a repeated block transfer triggers another block transfer.



DMA Operation www.ti.com

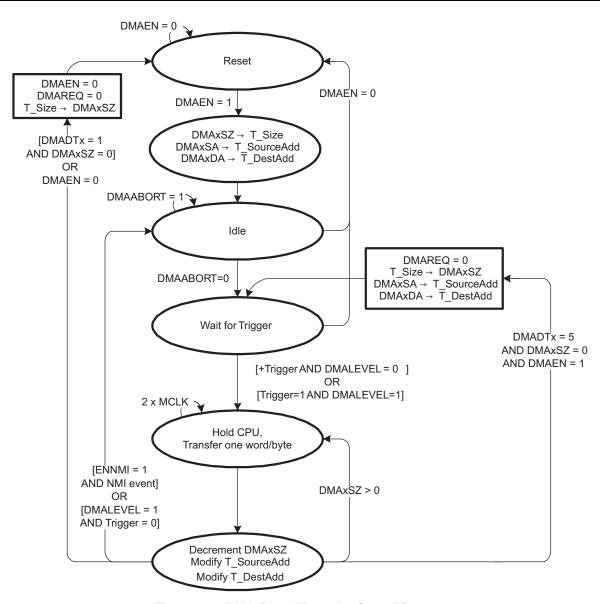


Figure 6-4. DMA Block Transfer State Diagram



www.ti.com DMA Operation

#### 6.2.2.3 Burst-Block Transfers

In burst-block mode, transfers are block transfers with CPU activity interleaved. The CPU executes 2 MCLK cycles after every four byte/word transfers of the block resulting in 20% CPU execution capacity. After the burst-block, CPU execution resumes at 100% capacity and the DMAEN bit is cleared. DMAEN must be set again before another burst-block transfer can be triggered. After a burst-block transfer has been triggered, further trigger signals occurring during the burst-block transfer are ignored. The burst-block transfer state diagram is shown in Figure 6-5.

The DMAxSZ register is used to define the size of the block and the DMADSTINCRx and DMASRCINCRx bits select if the destination address and the source address are incremented or decremented after each transfer of the block. If DMAxSZ = 0, no transfers occur.

The DMAxSA, DMAxDA, and DMAxSZ registers are copied into temporary registers. The temporary values of DMAxSA and DMAxDA are incremented or decremented after each transfer in the block. The DMAxSZ register is decremented after each transfer of the block and shows the number of transfers remaining in the block. When the DMAxSZ register decrements to zero it is reloaded from its temporary register and the corresponding DMAIFG flag is set.

In repeated burst-block mode the DMAEN bit remains set after completion of the burst-block transfer and no further trigger signals are required to initiate another burst-block transfer. Another burst-block transfer begins immediately after completion of a burst-block transfer. In this case, the transfers must be stopped by clearing the DMAEN bit, or by an NMI interrupt when ENNMI is set. In repeated burst-block mode the CPU executes at 20% capacity continuously until the repeated burst-block transfer is stopped.



DMA Operation www.ti.com

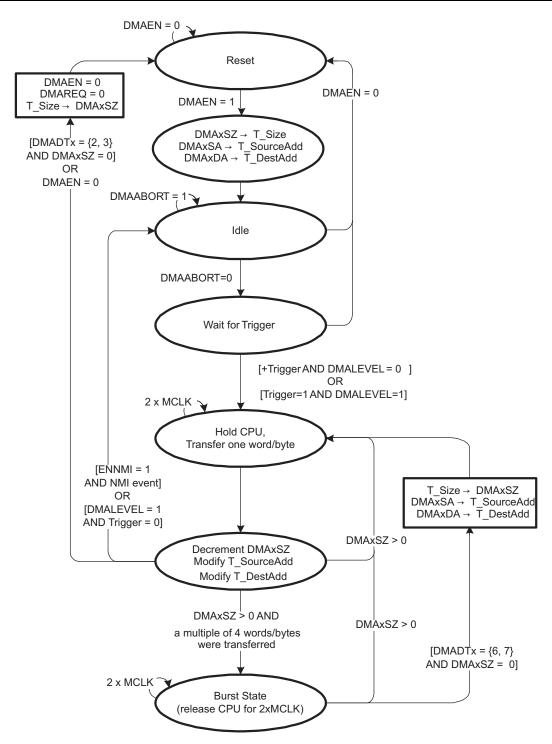


Figure 6-5. DMA Burst-Block Transfer State Diagram



www.ti.com DMA Operation

#### 6.2.3 Initiating DMA Transfers

Each DMA channel is independently configured for its trigger source with the DMAxTSELx bits as described in Table 6-2. The DMAxTSELx bits should be modified only when the DMACTLx DMAEN bit is 0. Otherwise, unpredictable DMA triggers may occur.

When selecting the trigger, the trigger must not have already occurred, or the transfer will not take place. For example, if the TACCR2 CCIFG bit is selected as a trigger, and it is already set, no transfer will occur until the next time the TACCR2 CCIFG bit is set.

#### 6.2.3.1 Edge-Sensitive Triggers

When DMALEVEL = 0, edge-sensitive triggers are used and the rising edge of the trigger signal initiates the transfer. In single-transfer mode, each transfer requires its own trigger. When using block or burst-block modes, only one trigger is required to initiate the block or burst-block transfer.

#### 6.2.3.2 Level-Sensitive Triggers

When DMALEVEL = 1, level-sensitive triggers are used. For proper operation, level-sensitive triggers can only be used when external trigger DMAE0 is selected as the trigger. DMA transfers are triggered as long as the trigger signal is high and the DMAEN bit remains set.

The trigger signal must remain high for a block or burst-block transfer to complete. If the trigger signal goes low during a block or burst-block transfer, the DMA controller is held in its current state until the trigger goes back high or until the DMA registers are modified by software. If the DMA registers are not modified by software, when the trigger signal goes high again, the transfer resumes from where it was when the trigger signal went low.

When DMALEVEL = 1, transfer modes selected when DMADTx =  $\{0, 1, 2, 3\}$  are recommended because the DMAEN bit is automatically reset after the configured transfer.

#### 6.2.3.3 Halting Executing Instructions for DMA Transfers

The DMAONFETCH bit controls when the CPU is halted for a DMA transfer. When DMAONFETCH = 0, the CPU is halted immediately and the transfer begins when a trigger is received. When DMAONFETCH = 1, the CPU finishes the currently executing instruction before the DMA controller halts the CPU and the transfer begins.

#### NOTE: DMAONFETCH Must Be Used When The DMA Writes To Flash

If the DMA controller is used to write to flash memory, the DMAONFETCH bit must be set. Otherwise, unpredictable operation can result.

#### Table 6-2. DMA Trigger Operation

DMAxTSELx	Operation
0000	A transfer is triggered when the DMAREQ bit is set. The DMAREQ bit is automatically reset when the transfer starts
0001	A transfer is triggered when the TACCR2 CCIFG flag is set. The TACCR2 CCIFG flag is automatically reset when the transfer starts. If the TACCR2 CCIE bit is set, the TACCR2 CCIFG flag will not trigger a transfer.
0010	A transfer is triggered when the TBCCR2 CCIFG flag is set. The TBCCR2 CCIFG flag is automatically reset when the transfer starts. If the TBCCR2 CCIE bit is set, the TBCCR2 CCIFG flag will not trigger a transfer.
0011	A transfer is triggered when serial interface receives new data.  Devices with USCI_A0 module: A transfer is triggered when USCI_A0 receives new data. UCA0RXIFG is automatically reset when the transfer starts. If UCA0RXIE is set, the UCA0RXIFG flag will not trigger a transfer.
0100	A transfer is triggered when serial interface is ready to transmit new data.  Devices with USCI_A0 module: A transfer is triggered when USCI_A0 is ready to transmit new data.  UCA0TXIFG is automatically reset when the transfer starts. If UCA0TXIE is set, the UCA0TXIFG flag will not trigger a transfer.



DMA Operation www.ti.com

# **Table 6-2. DMA Trigger Operation (continued)**

DMAxTSELx	Operation
0101	A transfer is triggered when the DAC12_0CTL DAC12IFG flag is automatically cleared when the transfer starts. If the DAC12_0CTL DAC12IE bit is set, the DAC12_0CTL DAC12IFG flag will not trigger a transfer.
0110	A transfer is triggered by an ADC12IFGx flag. When single-channel conversions are performed, the corresponding ADC12IFGx is the trigger. When sequences are used, the ADC12IFGx for the last conversion in the sequence is the trigger. A transfer is triggered when the conversion is completed and the ADC12IFGx is set. Setting the ADC12IFGx with software will not trigger a transfer. All ADC12IFGx flags are automatically reset when the associated ADC12MEMx register is accessed by the DMA controller.
0111	A transfer is triggered when the TACCR0 CCIFG flag is set. The TACCR0 CCIFG flag is automatically reset when the transfer starts. If the TACCR0 CCIE bit is set, the TACCR0 CCIFG flag will not trigger a transfer.
1000	A transfer is triggered when the TBCCR0 CCIFG flag is set. The TBCCR0 CCIFG flag is automatically reset when the transfer starts. If the TBCCR0 CCIE bit is set, the TBCCR0 CCIFG flag will not trigger a transfer.
1001	A transfer is triggered when the UCA1RXIFG flag is set. UCA1RXIFG is automatically reset when the transfer starts. If URXIE1 is set, the UCA1RXIFG flag will not trigger a transfer.
1010	A transfer is triggered when the UCA1TXIFG flag is set. UCA1TXIFG is automatically reset when the transfer starts. If UTXIE1 is set, the UCA1TXIFG flag will not trigger a transfer.
1011	A transfer is triggered when the hardware multiplier is ready for a new operand.
1100	No transfer is triggered.  Devices with USCI_B0 module: A transfer is triggered when USCI_B0 receives new data. UCB0RXIFG is automatically reset when the transfer starts. If UCB0RXIE is set, the UCB0RXIFG flag will not trigger a transfer.
1101	No transfer is triggered.  Devices with USCI_B0 module: A transfer is triggered when USCI_B0 is ready to transmit new data.  UCB0TXIFG is automatically reset when the transfer starts. If UCB0TXIE is set, the UCB0TXIFG flag will not trigger a transfer.
1110	A transfer is triggered when the DMAxIFG flag is set. DMA0IFG triggers channel 1, DMA1IFG triggers channel 2, and DMA2IFG triggers channel 0. None of the DMAxIFG flags are automatically reset when the transfer starts.
1111	A transfer is triggered by the external trigger DMAE0.

# 6.2.4 Stopping DMA Transfers

There are two ways to stop DMA transfers in progress:

- A single, block, or burst-block transfer may be stopped with an NMI interrupt, if the ENNMI bit is set in register DMACTL1.
- A burst-block transfer may be stopped by clearing the DMAEN bit.



www.ti.com DMA Operation

#### 6.2.5 DMA Channel Priorities

The default DMA channel priorities are DMA0-DMA1-DMA2. If two or three triggers happen simultaneously or are pending, the channel with the highest priority completes its transfer (single, block or burst-block transfer) first, then the second priority channel, then the third priority channel. Transfers in progress are not halted if a higher priority channel is triggered. The higher priority channel waits until the transfer in progress completes before starting.

The DMA channel priorities are configurable with the ROUNDROBIN bit. When the ROUNDROBIN bit is set, the channel that completes a transfer becomes the lowest priority. The *order* of the priority of the channels always stays the same, DMA0-DMA1-DMA2 (see Table 6-3).

DMA Priority	Transfer Occurs	New DMA Priority		
DMA0 - DMA1 - DMA2	DMA1	DMA2 - DMA0 - DMA1		
DMA2 - DMA0 - DMA1	DMA2	DMA0 - DMA1 - DMA2		
DMA0 - DMA1 - DMA2	DMA0	DMA1 - DMA2 - DMA0		

**Table 6-3. Channel Priorities** 

When the ROUNDROBIN bit is cleared the channel priority returns to the default priority.

#### 6.2.6 DMA Transfer Cycle Time

The DMA controller requires one or two MCLK clock cycles to synchronize before each single transfer or complete block or burst-block transfer. Each byte/word transfer requires two MCLK cycles after synchronization, and one cycle of wait time after the transfer. Because the DMA controller uses MCLK, the DMA cycle time is dependent on the MSP430 operating mode and clock system setup.

If the MCLK source is active, but the CPU is off, the DMA controller will use the MCLK source for each transfer, without re-enabling the CPU. If the MCLK source is off, the DMA controller will temporarily restart MCLK, sourced with DCOCLK, for the single transfer or complete block or burst-block transfer. The CPU remains off, and after the transfer completes, MCLK is turned off. The maximum DMA cycle time for all operating modes is shown in Table 6-4.

CPU Operating Mode	Clock Source	Maximum DMA Cycle Time
Active mode	MCLK = DCOCLK	4 MCLK cycles
Active mode	MCLK = LFXT1CLK	4 MCLK cycles
Low-power mode LPM0/1	MCLK = DCOCLK	5 MCLK cycles
Low-power mode LPM3/4	MCLK = DCOCLK	5 MCLK cycles + 6 µs <sup>(1)</sup>
Low-power mode LPM0/1	MCLK = LFXT1CLK	5 MCLK cycles
Low-power mode LPM3	MCLK = LFXT1CLK	5 MCLK cycles
Low-power mode LPM4	MCLK = LFXT1CLK	5 MCLK cycles + 6 μs <sup>(1)</sup>

Table 6-4. Maximum Single-Transfer DMA Cycle Time

#### 6.2.7 Using DMA With System Interrupts

DMA transfers are not interruptible by system interrupts. System interrupts remain pending until the completion of the transfer. NMI interrupts can interrupt the DMA controller if the ENNMI bit is set.

System interrupt service routines are interrupted by DMA transfers. If an interrupt service routine or other routine must execute with no interruptions, the DMA controller should be disabled prior to executing the routine.

The additional 6  $\mu s$  are needed to start the DCOCLK. It is the  $t_{(LPMx)}$  parameter in the data sheet.



DMA Operation www.ti.com

#### 6.2.8 DMA Controller Interrupts

Each DMA channel has its own DMAIFG flag. Each DMAIFG flag is set in any mode, when the corresponding DMAxSZ register counts to zero. If the corresponding DMAIE and GIE bits are set, an interrupt request is generated.

All DMAIFG flags source only one DMA controller interrupt vector and, on some devices, the interrupt vector may be shared with other modules. Please refer to the device specific datasheet for further details. For these devices, software must check the DMAIFG and respective module flags to determine the source of the interrupt. The DMAIFG flags are not reset automatically and must be reset by software.

Additionally, some devices utilize the DMAIV register. All DMAIFG flags are prioritized, with DMA0IFG being the highest, and combined to source a single interrupt vector. The highest priority enabled interrupt generates a number in the DMAIV register. This number can be evaluated or added to the program counter to automatically enter the appropriate software routine. Disabled DMA interrupts do not affect the DMAIV value.

Any access, read or write, of the DMAIV register automatically resets the highest pending interrupt flag. If another interrupt flag is set, another interrupt is immediately generated after servicing the initial interrupt. For example, assume that DMA0 has the highest priority. If the DMA0IFG and DMA2IFG flags are set when the interrupt service routine accesses the DMAIV register, DMA0IFG is reset automatically. After the RETI instruction of the interrupt service routine is executed, the DMA2IFG will generate another interrupt.

The following software example shows the recommended use of DMAIV and the handling overhead. The DMAIV value is added to the PC to automatically jump to the appropriate routine.

The numbers at the right margin show the necessary CPU cycles for each instruction. The software overhead for different interrupt sources includes interrupt latency and return-from-interrupt cycles, but not the task handling itself.

#### Example 6-1. DMAIV Software Example

```
;Interrupt handler for DMA0IFG, DMA1IFG, DMA2IFG
                                               Cycles
DMA_HND ...
                ; Interrupt latency
             &DMAIV,PC ; Add offset to Jump table
       ADD
                                                  3
       RETI ; Vector 0: No interrupt
       JMP DMA0_HND ; Vector 2: DMA channel 0 2
        JMP DMA1_HND ; Vector 4: DMA channel 1 2
        JMP DMA2_HND ; Vector 6: DMA channel 2 2
        RETI ; Vector 8: Reserved
        RETI
                     ; Vector 10: Reserved
                    ; Vector 12: Reserved
        RETI
                       ; Vector 14: Reserved
                      ; Vector 6: DMA channel 2
DMA2 HND
                   ; Task starts nel; Back to main program
       RETI
                    ; Vector 4: DMA channel 1
DMA1_HND
       ... ; Task starts here
RETI ; Back to main program
                                                  5
DMA0_HND
                      ; Vector 2: DMA channel 0
                       ; Task starts here
       RETT
                       ; Back to main program
```

### 6.2.9 Using the USCI\_B &C Module with the DMA Controller

The USCI\_B I<sup>2</sup>C module provides two trigger sources for the DMA controller. The USCI\_B I<sup>2</sup>C module can trigger a transfer when new I<sup>2</sup>C data is received and when data is needed for transmit.

A transfer is triggered if UCB0RXIFG is set. The UCB0RXIFG is cleared automatically when the DMA controller acknowledges the transfer. If UCB0RXIE is set, UCB0RXIFG will not trigger a transfer.



www.ti.com DMA Operation

A transfer is triggered if UCB0TXIFG is set. The UCB0TXIFG is cleared automatically when the DMA controller acknowledges the transfer. If UCB0TXIE is set, UCB0TXIFG will not trigger a transfer.

#### 6.2.10 Using ADC12 with the DMA Controller

MSP430 devices with an integrated DMA controller can automatically move data from any ADC12MEMx register to another location. DMA transfers are done without CPU intervention and independently of any low-power modes. The DMA controller increases throughput of the ADC12 module, and enhances low-power applications allowing the CPU to remain off while data transfers occur.

DMA transfers can be triggered from any ADC12IFGx flag. When CONSEQx = {0,2} the ADC12IFGx flag for the ADC12MEMx used for the conversion can trigger a DMA transfer. When CONSEQx = {1,3}, the ADC12IFGx flag for the last ADC12MEMx in the sequence can trigger a DMA transfer. Any ADC12IFGx flag is automatically cleared when the DMA controller accesses the corresponding ADC12MEMx.

#### 6.2.11 Using DAC12 With the DMA Controller

MSP430 devices with an integrated DMA controller can automatically move data to the DAC12\_xDAT register. DMA transfers are done without CPU intervention and independently of any low-power modes. The DMA controller increases throughput to the DAC12 module, and enhances low-power applications allowing the CPU to remain off while data transfers occur.

Applications requiring periodic waveform generation can benefit from using the DMA controller with the DAC12. For example, an application that produces a sinusoidal waveform may store the sinusoid values in a table. The DMA controller can continuously and automatically transfer the values to the DAC12 at specific intervals creating the sinusoid with zero CPU execution. The DAC12\_xCTL DAC12IFG flag is automatically cleared when the DMA controller accesses the DAC12\_xDAT register.

#### 6.2.12 Writing to Flash With the DMA Controller

MSP430 devices with an integrated DMA controller can automatically move data to the Flash memory. DMA transfers are done without CPU intervention and independent of any low-power modes. The DMA controller performs the move of the data word/byte to the Flash. The write timing control is done by the Flash controller. Write transfers to the Flash memory succeed if the Flash controller is set up prior to the DMA transfer and if the Flash is not busy. To set up the Flash controller for write accesses, see the *Flash Memory Controller* chapter.



DMA Registers www.ti.com

# 6.3 DMA Registers

The DMA registers are listed in Table 6-5.

# **Table 6-5. DMA Registers**

Register	Short Form	Register Type	Address	Initial State
DMA control 0	DMACTL0	Read/write	0122h	Reset with POR
DMA control 1	DMACTL1	Read/write	0124h	Reset with POR
DMA interrupt vector	DMAIV	Read only	0126h	Reset with POR
DMA channel 0 control	DMA0CTL	Read/write	01D0h	Reset with POR
DMA channel 0 source address	DMA0SA	Read/write	01D2h	Unchanged
DMA channel 0 destination address	DMA0DA	Read/write	01D6h	Unchanged
DMA channel 0 transfer size	DMA0SZ	Read/write	01DAh	Unchanged
DMA channel 1 control	DMA1CTL	Read/write	01DCh	Reset with POR
DMA channel 1 source address	DMA1SA	Read/write	01DEh	Unchanged
DMA channel 1 destination address	DMA1DA	Read/write	01E2h	Unchanged
DMA channel 1 transfer size	DMA1SZ	Read/write	01E6h	Unchanged
DMA channel 2 control	DMA2CTL	Read/write	01E8h	Reset with POR
DMA channel 2 source address	DMA2SA	Read/write	01EAh	Unchanged
DMA channel 2 destination address	DMA2DA	Read/write	01EEh	Unchanged
DMA-channel 2 transfer size	DMA2SZ	Read/write	01F2h	Unchanged



www.ti.com DMA Registers

### 6.3.1 DMACTL0, DMA Control Register 0

15	14	13	12	11	10	9	8
	Rese	Reserved DMA2TSELx					
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
7	6	5	4	3	2	1	0
	DMA1TSELx				DMA0	TSELx	
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)

Reserved Bits 15-12 Reserved

**DMA2TSELx** Bits 11-8 DMA trigger select. These bits select the DMA transfer trigger.

0000 DMAREQ bit (software trigger)

0001 TACCR2 CCIFG bit

0010 TBCCR2 CCIFG bit

0011 Serial data received UCA0RXIFG

0100 Serial data transmit ready UCA0TXIFG

0101 DAC12\_0CTL DAC12IFG bit

0110 ADC12 ADC12IFGx bit

0111 TACCR0 CCIFG bit

1000 TBCCR0 CCIFG bit

1001 Serial data received UCA1RXIFG

1010 Serial data transmit ready UCA1TXIFG

1011 Multiplier ready

1100 Serial data received UCB0RXIFG

1101 Serial data transmit ready UCB0TXIFG

DMA0IFG bit triggers DMA channel 1 DMA1IFG bit triggers DMA channel 2

DMA2IFG bit triggers DMA channel 0

External trigger DMAE0

DMA1TSELx Bits 7-4 Same as DMA2TSELx DMA0TSELx Bits 3-0 Same as DMA2TSELx

#### 6.3.2 DMACTL1, DMA Control Register 1

1110

1111

15	14	13	12	11	10	9	8
0	0	0	0	0	0	0	0
r0	r0	rO	rO	rO	rO	rO	rO
7	6	5	4	3	2	1	0
0	0	0	0	0	DMAON FETCH	ROUND ROBIN	ENNMI
r0	r0	r0	r0	r0	rw-(0)	rw-(0)	rw-(0)

Reserved Bits 15-3 Reserved. Read only. Always read as 0.

**DMAONFETCH** Bit 2 DMA on fetch

0 The DMA transfer occurs immediately.

1 The DMA transfer occurs on next instruction fetch after the trigger.

ROUNDROBIN Bit 1 Round robin. This bit enables the round-robin DMA channel priorities.

0 DMA channel priority is DMA0 - DMA1 - DMA2

1 DMA channel priority changes with each transfer

ENNMI Bit 0 Enable NMI. This bit enables the interruption of a DMA transfer by an NMI interrupt. When an NMI

interrupts a DMA transfer, the current transfer is completed normally, further transfers are stopped, and

DMAABORT is set.

0 NMI interrupt does not interrupt DMA transfer

1 NMI interrupt interrupts a DMA transfer



DMA Registers www.ti.com

# 6.3.3 DMAxCTL, DMA Channel x Control Register

15	14	13	12	11	10	9	8	
Reserved		DMADTx		DMADS	TINCRx	DMASR	CINCRx	
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	
7	6	5	4	3	2	1	0	
DMADST BYTE	DMASRC BYTE	DMALEVEL	DMAEN	DMAIFG	DMAIE	DMAABORT	DMAREQ	
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	
Reserved	Bit 15	Reserve						
DMADTx	Bits 14-12	DMA Tra	ansfer mode.					
		000	Single transfer					
		001	Block transfer					
		010	Burst-block transfe	er				
		011	Burst-block transfe	er				
		100	Repeated single to	ansfer				
		101	Repeated block tra	ansfer				
		110	Repeated burst-bl	ock transfer				
		111	Repeated burst-bl	ock transfer				
DMADSTINCRx	Bits 11-10	destinati address increme tempora	ion address after e increments/decrer nts/decrements by ry register is increr	ach byte or word to ments by one. Whe two. The DMAxDA mented or decreme	ransfer. When DN n DMADSTBYTE A is copied into a	enting or decrement MADSTBYTE = 1, to E = 0, the destination temporary register is not incremented of	he destination on address and the	
		00	Destination addres	ss is unchanged				
		01	Destination address	ss is unchanged				
		10	Destination address	ss is decremented				
		11	Destination address	ss is incremented				
DMASRCINCRX	Bits 9-8	DMA source increment. This bit selects automatic incrementing or decrementing of the source address for each byte or word transfer. When DMASRCBYTE = 1, the source address increments/decrements by one. When DMASRCBYTE = 0, the source address increments/decrements by two. The DMAxSA is copied into a temporary register and the temporary register is incremented or decremented. DMAxSA is not incremented or decremented.						
		00	Source address is	unchanged				
		01	Source address is	unchanged				
		10	Source address is	decremented				
		11	Source address is	incremented				
DMADSTBYTE	Bit 7	DMA de	stination byte. This	bit selects the des	stination as a byte	e or word.		
		0	Word					
		1	Byte					
DMASRCBYTE	Bit 6	DMA so	urce byte. This bit	selects the source	as a byte or wor	d.		
		0	Word					
		1	Byte					
DMALEVEL	Bit 5	DMA lev		between edge-ser	nsitive and level-s	sensitive triggers.		
		0	Edge sensitive (ris	• • •				
		1	Level sensitive (hi	gh level)				
DMAEN	Bit 4	DMA en						
		0	Disabled					
		1	Enabled					
DMAIFG	Bit 3		errupt flag					
		0	No interrupt pendi	ng				
	_	1	Interrupt pending					
DMAIE	Bit 2		errupt enable					
		0	Disabled					
		1	Enabled					



www.ti.com DMA Registers

DMAABORT

Bit 1

DMA Abort. This bit indicates if a DMA transfer was interrupt by an NMI.

0

DMA transfer not interrupted

1

DMA transfer was interrupted by NMI

DMAREQ

Bit 0

DMA request. Software-controlled DMA start. DMAREQ is reset automatically.

0

No DMA start

1

Start DMA

#### 6.3.4 DMAxSA, DMA Source Address Register

15	14	13	12	11	10	9	8			
	Reserved									
r0	r0	rO	rO	r0	rO	rO	rO			
7	6	5	4	3	2	1	0			
	Rese	erved			DMA	xSAx				
r0	r0	r0	r0	rw	rw	rw	rw			
15	14	13	12	11	10	9	8			
			DMA	xSAx						
rw	rw	rw	rw	rw	rw	rw	rw			
7	6	5	4	3	2	1	0			
	DMAxSAx									
rw	rw	rw	rw	rw	rw	rw	rw			

# DMAxSA Bits 15-0 DMA source address

The source address register points to the DMA source address for single transfers or the first source address for block transfers. The source address register remains unchanged during block and burst-block transfers.

Devices that have addressable memory range 64 KB or below contain a single word for the DMAxSA. The upper word is automatically cleared when writing using word operations. Reads from this location are always read as zero.

Devices that have addressable memory range beyond 64 KB contain an additional word for the source address. Bits 15-4 of this additional word are reserved and always read as zero. When writing to DMAxSA with word formats, this additional word is automatically cleared. Reads of this additional word using word formats, are always read as zero.



DMA Registers www.ti.com

### 6.3.5 DMAxDA, DMA Destination Address Register

15	14	13	12	11	10	9	8			
	Reserved									
rO	rO	r0	r0	r0	rO	rO	rO			
7	6	5	4	3	2	1	0			
	Rese	erved			DMA	xDAx				
r0	r0	r0	r0	rw	rw	rw	rw			
15	14	13	12	11	10	9	8			
			DMA	xDAx						
rw	rw	rw	rw	rw	rw	rw	rw			
7	6	5	4	3	2	1	0			
DMAxDAx										
rw	rw	rw	rw	rw	rw	rw	rw			

#### DMAxDA Bits 15-0

DMA destination address

The destination address register points to the DMA destination address for single transfers or the first destination address for block transfers. The destination address register remains unchanged during block and burst-block transfers.

Devices that have addressable memory range 64 KB or below contain a single word for the DMAxDA. Devices that have addressable memory range beyond 64 KB contain an additional word for the destination address. Bits 15-4 of this additional word are reserved and always read as zero. When writing to DMAxDA with word formats, this additional word is automatically cleared. Reads of this additional word using word formats, are always read as zero.

# 6.3.6 DMAxSZ, DMA Size Address Register

15	14	13	12	11	10	9	8			
DMAxSZx										
rw	rw	rw	rw	rw	rw	rw	rw			
7	6	5	4	3	2	1	0			
	DMAxSZx									
rw	rw	rw	rw	rw	rw	rw	rw			
DMAxSZx	Bits 15-0	DMA size. The DMA	size register defin	es the number of	byte/word data pe	er block transfer. D	DMAxSZ register			

DMA size. The DMA size register defines the number of byte/word data per block transfer. DMAxSZ register decrements with each word or byte transfer. When DMAxSZ decrements to 0, it is immediately and automatically reloaded with its previously initialized value.

00000h Transfer is disabled

00001h One byte or word to be transferred

00002h Two bytes or words have to be transferred

:

0FFFh 65535 bytes or words have to be transferred



www.ti.com DMA Registers

# 6.3.7 DMAIV, DMA Interrupt Vector Register

15	14	13	12	11	10	9	8
0	0	0	0	0	0	0	0
rO	r0	r0	r0	rO	rO	rO	rO
7	6	5	4	3	2	1	0
0	0	0	0		DMAIVx		0
r0	r0	r0	r0	r(0)	r(0)	r(0)	r0

**DMAIVx** Bits 15-0 DMA interrupt vector value

DMAIV Contents	Interrupt Source	Interrupt Flag	Interrupt Priority
00h	No interrupt pending	-	
02h	DMA channel 0	DMA0IFG	Highest
04h	DMA channel 1	DMA1IFG	
06h	DMA channel 2	DMA2IFG	
08h	Reserved	-	
0Ah	Reserved	-	
0Ch	Reserved	-	
0Eh	Reserved	-	Lowest



# Flash Memory Controller

This chapter describes the operation of the MSP430x2xx flash memory controller.

TopicPage7.1Flash Memory Introduction3097.2Flash Memory Segmentation3097.3Flash Memory Operation3117.4Flash Memory Registers323



#### 7.1 Flash Memory Introduction

The MSP430 flash memory is bit-, byte-, and word-addressable and programmable. The flash memory module has an integrated controller that controls programming and erase operations. The controller has four registers, a timing generator, and a voltage generator to supply program and erase voltages.

MSP430 flash memory features include:

- Internal programming voltage generation
- Bit, byte, or word programmable
- Ultralow-power operation
- Segment erase and mass erase
- Marginal 0 and marginal 1 read mode (optional, see the device-specific data sheet)

Figure 7-1 shows the block diagram of the flash memory and controller.

#### NOTE: Minimum V<sub>cc</sub> during flash write or erase

The minimum  $V_{cc}$  voltage during a flash write or erase operation is 2.2 V. If  $V_{cc}$  falls below 2.2 V during write or erase, the result of the write or erase is unpredictable.

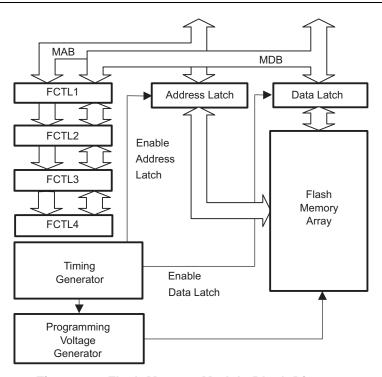


Figure 7-1. Flash Memory Module Block Diagram

#### 7.2 Flash Memory Segmentation

MSP430 flash memory is partitioned into segments. Single bits, bytes, or words can be written to flash memory, but the segment is the smallest size of flash memory that can be erased.

The flash memory is partitioned into main and information memory sections. There is no difference in the operation of the main and information memory sections. Code or data can be located in either section. The differences between the two sections are the segment size and the physical addresses.

The information memory has four 64-byte segments. The main memory has one or more 512-byte segments. See the device-specific data sheet for the complete memory map of a device.

The segments are further divided into blocks.



Figure 7-2 shows the flash segmentation using an example of 32-KB flash that has eight main segments and four information segments.

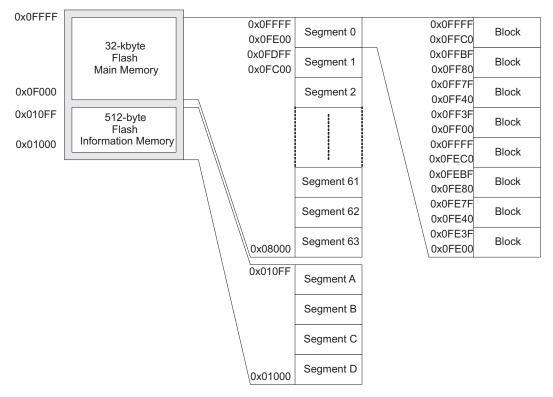


Figure 7-2. Flash Memory Segments, 32-KB Example

#### 7.2.1 SegmentA

SegmentA of the information memory is locked separately from all other segments with the LOCKA bit. When LOCKA = 1, SegmentA cannot be written or erased and all information memory is protected from erasure during a mass erase or production programming. When LOCKA = 0, SegmentA can be erased and written as any other flash memory segment, and all information memory is erased during a mass erase or production programming.

The state of the LOCKA bit is toggled when a 1 is written to it. Writing a 0 to LOCKA has no effect. This allows existing flash programming routines to be used unchanged.

```
; Unlock SegmentA
BIT #LOCKA,&FCTL3 ; Test LOCKA
JZ SEGA_UNLOCKED ; Already unlocked?
MOV #FWKEY+LOCKA,&FCTL3 ; No, unlock SegmentA
SEGA_UNLOCKED ; Yes, continue
; SegmentA is unlocked

; Lock SegmentA
BIT #LOCKA,&FCTL3 ; Test LOCKA
JNZ SEGA_LOCKED ; Already locked?
MOV #FWKEY+LOCKA,&FCTL3 ; No, lock SegmentA
SEGA_LOCKED ; Yes, continue
; SegmentA is locked
```



#### 7.3 Flash Memory Operation

The default mode of the flash memory is read mode. In read mode, the flash memory is not being erased or written, the flash timing generator and voltage generator are off, and the memory operates identically to ROM.

MSP430 flash memory is in-system programmable (ISP) without the need for additional external voltage. The CPU can program its own flash memory. The flash memory write and erase modes are selected with the BLKWRT, WRT, MERAS, and ERASE bits and are:

- Byte or word write
- Block write
- Segment erase
- Mass erase (all main memory segments)
- All erase (all segments)

Reading from or writing to flash memory while it is being programmed or erased is prohibited. If CPU execution is required during the write or erase, the code to be executed must be in RAM. Any flash update can be initiated from within flash memory or RAM.

#### 7.3.1 Flash Memory Timing Generator

Write and erase operations are controlled by the flash timing generator shown in Figure 7-3. The flash timing generator operating frequency,  $f_{FTG}$ , must be in the range from approximately 257 kHz to approximately 476 kHz (see device-specific data sheet).

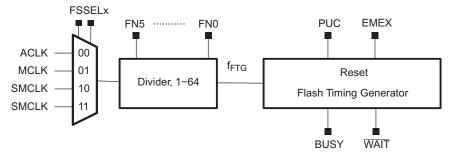


Figure 7-3. Flash Memory Timing Generator Block Diagram

#### 7.3.1.1 Flash Timing Generator Clock Selection

The flash timing generator can be sourced from ACLK, SMCLK, or MCLK. The selected clock source should be divided using the FNx bits to meet the frequency requirements for  $f_{FTG}$ . If the  $f_{FTG}$  frequency deviates from the specification during the write or erase operation, the result of the write or erase may be unpredictable, or the flash memory may be stressed above the limits of reliable operation.

If a clock failure is detected during a write or erase operation, the operation is aborted, the FAIL flag is set, and the result of the operation is unpredictable.

While a write or erase operation is active the selected clock source can not be disabled by putting the MSP430 into a low-power mode. The selected clock source remains active until the operation is completed before being disabled.



#### 7.3.2 Erasing Flash Memory

The erased level of a flash memory bit is 1. Each bit can be programmed from 1 to 0 individually but to reprogram from 0 to 1 requires an erase cycle. The smallest amount of flash that can be erased is a segment. There are three erase modes selected with the ERASE and MERAS bits listed in Table 7-1.

Table 7-1. Erase Modes

MERAS	ERASE	Erase Mode			
0	1	Segment erase			
1	0	Mass erase (all main memory segments)			
1	1	LOCKA = 0: Erase main and information flash memory.  LOCKA = 1: Erase only main flash memory.			

Any erase is initiated by a dummy write into the address range to be erased. The dummy write starts the flash timing generator and the erase operation. Figure 7-4 shows the erase cycle timing. The BUSY bit is set immediately after the dummy write and remains set throughout the erase cycle. BUSY, MERAS, and ERASE are automatically cleared when the cycle completes. The erase cycle timing is not dependent on the amount of flash memory present on a device. Erase cycle times are equivalent for all MSP430F2xx and MSP430G2xx devices.

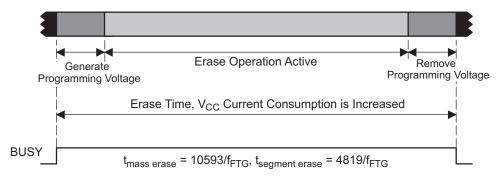


Figure 7-4. Erase Cycle Timing

A dummy write to an address not in the range to be erased does not start the erase cycle, does not affect the flash memory, and is not flagged in any way. This errant dummy write is ignored.



### 7.3.2.1 Initiating an Erase from Within Flash Memory

Any erase cycle can be initiated from within flash memory or from RAM. When a flash segment erase operation is initiated from within flash memory, all timing is controlled by the flash controller, and the CPU is held while the erase cycle completes. After the erase cycle completes, the CPU resumes code execution with the instruction following the dummy write.

When initiating an erase cycle from within flash memory, it is possible to erase the code needed for execution after the erase. If this occurs, CPU execution is unpredictable after the erase cycle.

The flow to initiate an erase from flash is shown in Figure 7-5.

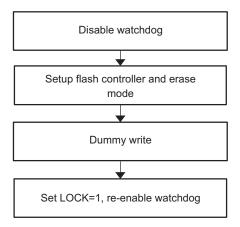


Figure 7-5. Erase Cycle from Within Flash Memory

```
; Segment Erase from flash. 514 kHz < SMCLK < 952 kHz
; Assumes ACCVIE = NMIIE = OFIE = 0.
         #WDTPW+WDTHOLD,&WDTCTL ; Disable WDT
  MOV
            #FWKEY+FSSEL1+FN0,&FCTL2 ; SMCLK/2
  MOV
            #FWKEY, &FCTL3 ; Clear LOCK
#FWKEY+ERASE, &FCTL1 ; Enable segr
   VOM
   VOM
                                        ; Enable segment erase
   CLR
            &0FC10h
                                        ; Dummy write, erase S1
            #FWKEY+LOCK, &FCTL3
   VOM
                                       ; Done, set LOCK
                                        ; Re-enable WDT?
   . . .
```



#### 7.3.2.2 Initiating an Erase from RAM

Any erase cycle may be initiated from RAM. In this case, the CPU is not held and can continue to execute code from RAM. The BUSY bit must be polled to determine the end of the erase cycle before the CPU can access any flash address again. If a flash access occurs while BUSY = 1, it is an access violation, ACCVIFG is set, and the erase results are unpredictable.

The flow to initiate an erase from flash from RAM is shown in Figure 7-6.

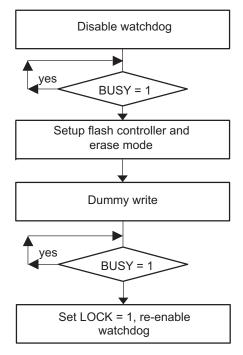


Figure 7-6. Erase Cycle from Within RAM

```
; Segment Erase from RAM. 514 kHz < SMCLK < 952 kHz
; Assumes ACCVIE = NMIIE = OFIE = 0.
    MOV
          #WDTPW+WDTHOLD, &WDTCTL
                                      ; Disable WDT
          #BUSY, &FCTL3
    BIT
                                      ; Test BUSY
    JNZ
                                      ; Loop while busy
          #FWKEY+FSSEL1+FN0, &FCTL2 ; SMCLK/2
    MOV
          #FWKEY&FCTL3
                                      ; Clear LOCK
    VOM
    MOV
          #FWKEY+ERASE, &FCTL1
                                     ; Enable erase
          &0FC10h
    CLR
                                     ; Dummy write, erase S1
L2
    BIT
          #BUSY, &FCTL3
                                     ; Test BUSY
                                     ; Loop while busy
    JNZ
          L2
          #FWKEY+LOCK&FCTL3
    MOV
                                      ; Done, set LOCK
                                      ; Re-enable WDT?
     . . .
```



# 7.3.3 Writing Flash Memory

The write modes, selected by the WRT and BLKWRT bits, are listed in Table 7-2.

Table 7-2. Write Modes

BLKWRT	WRT	Write Mode
0	1	Byte or word write
1	1	Block write

Both write modes use a sequence of individual write instructions, but using the block write mode is approximately twice as fast as byte or word mode, because the voltage generator remains on for the complete block write. Any instruction that modifies a destination can be used to modify a flash location in either byte or word write mode or block write mode. A flash word (low and high bytes) must not be written more than twice between erasures. Otherwise, damage can occur.

The BUSY bit is set while a write operation is active and cleared when the operation completes. If the write operation is initiated from RAM, the CPU must not access flash while BUSY = 1. Otherwise, an access violation occurs, ACCVIFG is set, and the flash write is unpredictable.

#### 7.3.3.1 Byte or Word Write

A byte or word write operation can be initiated from within flash memory or from RAM. When initiating from within flash memory, all timing is controlled by the flash controller, and the CPU is held while the write completes. After the write completes, the CPU resumes code execution with the instruction following the write. The byte or word write timing is shown in Figure 7-7.

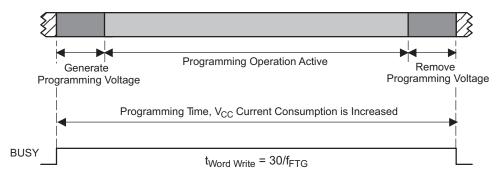


Figure 7-7. Byte or Word Write Timing

When a byte or word write is executed from RAM, the CPU continues to execute code from RAM. The BUSY bit must be zero before the CPU accesses flash again, otherwise an access violation occurs, ACCVIFG is set, and the write result is unpredictable.

In byte or word mode, the internally-generated programming voltage is applied to the complete 64-byte block, each time a byte or word is written, for 27 of the 30  $f_{FTG}$  cycles. With each byte or word write, the amount of time the block is subjected to the programming voltage accumulates. The cumulative programming time,  $t_{CPT}$ , must not be exceeded for any block. If the cumulative programming time is met, the block must be erased before performing any further writes to any address within the block. See the device-specific data sheet for specifications.



#### 7.3.3.2 Initiating a Byte or Word Write From Within Flash Memory

The flow to initiate a byte or word write from flash is shown in Figure 7-8.

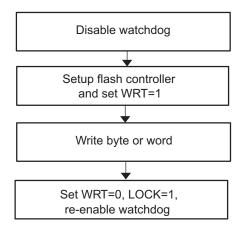


Figure 7-8. Initiating a Byte or Word Write From Flash

```
; Byte/word write from flash. 514 kHz < SMCLK < 952 kHz
; Assumes OFF1Eh is already erased
; Assumes ACCVIE = NMIIE = OFIE = 0.
  MOV #WDTPW+WDTHOLD, &WDTCTL
                                ; Disable WDT
        #FWKEY+FSSEL1+FN0,&FCTL2 ; SMCLK/2
  MOV
  VOM
        #FWKEY,&FCTL3
                                ; Clear LOCK
  MOV
        #FWKEY+WRT,&FCTL1
                                ; Enable write
        #0123h,&0FF1Eh
                                ; 0123h
  MOV
                                          -> OFF1Eh
  VOM
        #FWKEY,&FCTL1
                                 ; Done. Clear WRT
  VOM
        #FWKEY+LOCK,&FCTL3
                                 ; Set LOCK
                                  ; Re-enable WDT?
  . . .
```



#### 7.3.3.3 Initiating a Byte or Word Write From RAM

The flow to initiate a byte or word write from RAM is shown in Figure 7-9.

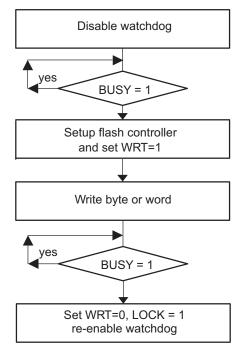


Figure 7-9. Initiating a Byte or Word Write from RAM

```
; Byte/word write from RAM. 514 kHz < SMCLK < 952 kHz
; Assumes OFF1Eh is already erased
; Assumes ACCVIE = NMIIE = OFIE = 0.
          #WDTPW+WDTHOLD, &WDTCTL
                                    ; Disable WDT
          #BUSY,&FCTL3
L1
    BIT
                                    ; Test BUSY
    JNZ
                                    ; Loop while busy
          T.1
    MOV
          #FWKEY+FSSEL1+FN0,&FCTL2 ; SMCLK/2
    MOV
          #FWKEY,&FCTL3
                                    ; Clear LOCK
    MOV
          #FWKEY+WRT,&FCTL1
                                    ; Enable write
    MOV
          #0123h,&0FF1Eh
                                    ; 0123h -> 0FF1Eh
          #BUSY,&FCTL3
L2
                                    ; Test BUSY
    BIT
    JNZ
          L2
                                    ; Loop while busy
    MOV
          #FWKEY,&FCTL1
                                    ; Clear WRT
          #FWKEY+LOCK,&FCTL3
    MOV
                                    ; Set LOCK
                                     ; Re-enable WDT?
```



#### 7.3.3.4 Block Write

The block write can be used to accelerate the flash write process when many sequential bytes or words need to be programmed. The flash programming voltage remains on for the duration of writing the 64-byte block. The cumulative programming time t<sub>CPT</sub> must not be exceeded for any block during a block write.

A block write cannot be initiated from within flash memory. The block write must be initiated from RAM only. The BUSY bit remains set throughout the duration of the block write. The WAIT bit must be checked between writing each byte or word in the block. When WAIT is set the next byte or word of the block can be written. When writing successive blocks, the BLKWRT bit must be cleared after the current block is complete. BLKWRT can be set initiating the next block write after the required flash recovery time given by t<sub>end</sub>. BUSY is cleared following each block write completion indicating the next block can be written. Figure 7-10 shows the block write timing.

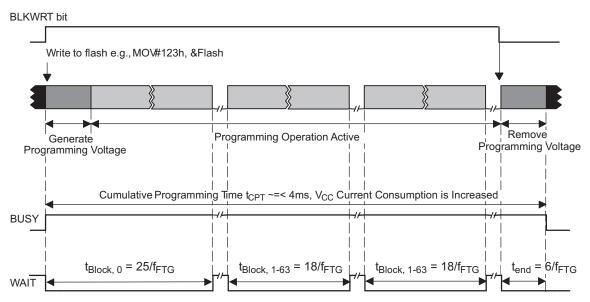


Figure 7-10. Block-Write Cycle Timing



#### **Block Write Flow and Example** 7.3.3.5

A block write flow is shown in Figure 7-11 and the following example.

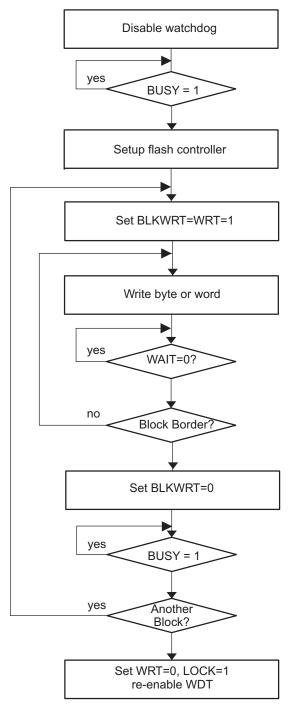


Figure 7-11. Block Write Flow

319



```
; Write one block starting at OF000h.
; Must be executed from RAM, Assumes Flash is already erased.
; 514 kHz < SMCLK < 952 kHz
; Assumes ACCVIE = NMIIE = OFIE = 0.
                                  ; Use as write counter
    MOM
        #32.R5
                                ; Write pointer
        #0F000h,R6
    VOM
    MOV #WDTPW+WDTHOLD,&WDTCTL ; Disable WDT
L1 BIT #BUSY,&FCTL3
                                ; Test BUSY
    JNZ L1
                                 ; Loop while busy
    MOV #FWKEY+FSSEL1+FN0,&FCTL2 ; SMCLK/2
          #FWKEY,&FCTL3 ; Clear LOCK
#FWKEY+BLKWRT+WRT,&FCTL1 ; Enable block write
    MOV
         #FWKEY,&FCTL3
    MOV
          MOV
   BIT
          #WAIT,&FCTL3
L3
                                 ; Loop while WAIT = 0
    JZ.
          T.3
    INCD R6
                                 ; Point to next word
    DEC
                                 ; Decrement write counter
    JNZ
                                 ; End of block?
    MOV #FWKEY,&FCTL1 ; Clear WRT,BLKWRT
BIT #BUSY,&FCTL3 ; Test BUSY
   BIT
L4
                                ; Loop while busy
    JNZ
          T.4
         #FWKEY+LOCK,&FCTL3 ; Set LOCK
    MOV
                                  ; Re-enable WDT if needed
```

#### 7.3.4 Flash Memory Access During Write or Erase

When any write or any erase operation is initiated from RAM and while BUSY = 1, the CPU may not read or write to or from any flash location. Otherwise, an access violation occurs, ACCVIFG is set, and the result is unpredictable. Also if a write to flash is attempted with WRT = 0, the ACCVIFG interrupt flag is set, and the flash memory is unaffected.

When a byte or word write or any erase operation is initiated from within flash memory, the flash controller returns op-code 03FFFh to the CPU at the next instruction fetch. Op-code 03FFFh is the JMP PC instruction. This causes the CPU to loop until the flash operation is finished. When the operation is finished and BUSY = 0, the flash controller allows the CPU to fetch the proper op-code and program execution resumes.

The flash access conditions while BUSY = 1 are listed in Table 7-3.

Flash Operation Flash Access		WAIT	Result
	Read	0	ACCVIFG = 0. 03FFFh is the value read.
Any erase, or byte or word write	Write	Write 0 ACCVIFG = 1. Write is ignored.	
byte of word write	Instruction fetch	0	ACCVIFG = 0. CPU fetches 03FFFh. This is the JMP PC instruction.
	Any	0	ACCVIFG = 1, LOCK = 1
Block write	Read	1	ACCVIFG = 0. 03FFFh is the value read.
DIOCK WITE	Write	1	ACCVIFG = 0. Write is written.
	Instruction fetch	1	ACCVIFG = 1, LOCK = 1

Table 7-3. Flash Access While BUSY = 1

Interrupts are automatically disabled during any flash operation when EEI = 0 and EEIEX = 0 and on MSP430x20xx and MSP430G2xx devices where EEI and EEIEX are not present. After the flash operation has completed, interrupts are automatically re-enabled. Any interrupt that occurred during the operation has its associated flag set and generates an interrupt request when re-enabled.

When EEIEX = 1 and GIE = 1, an interrupt immediately aborts any flash operation and the FAIL flag is set. When EEI = 1, GIE = 1, and EEIEX = 0, a segment erase is interrupted by a pending interrupt every 32  $f_{\text{FTG}}$  cycles. After servicing the interrupt, the segment erase is continued for at least 32  $f_{\text{FTG}}$  cycles or until it is complete. During the servicing of the interrupt, the BUSY bit remains set but the flash memory can be accessed by the CPU without causing an access violation occurs. Nested interrupts and using the RETI instruction inside interrupt service routines are not supported.



The watchdog timer (in watchdog mode) should be disabled before a flash erase cycle. A reset aborts the erase and the results are unpredictable. After the erase cycle has completed, the watchdog may be reenabled.

#### 7.3.5 Stopping a Write or Erase Cycle

Any write or erase operation can be stopped before its normal completion by setting the emergency exit bit EMEX. Setting the EMEX bit stops the active operation immediately and stops the flash controller. All flash operations cease, the flash returns to read mode, and all bits in the FCTL1 register are reset. The result of the intended operation is unpredictable.

#### 7.3.6 Marginal Read Mode

The marginal read mode can be used to verify the integrity of the flash memory contents. This feature is implemented in selected 2xx devices; see the device-specific data sheet for availability. During marginal read mode marginally programmed flash memory bit locations can be detected. Events that could produce this situation include improper  $f_{FTG}$  settings, or violation of minimum  $V_{CC}$  during erase or program operations. One method for identifying such memory locations would be to periodically perform a checksum calculation over a section of flash memory (for example, a flash segment) and repeating this procedure with the marginal read mode enabled. If they do not match, it could indicate an insufficiently programmed flash memory location. It is possible to refresh the affected Flash memory segment by disabling marginal read mode, copying to RAM, erasing the flash segment, and writing back to it from RAM.

The program checking the flash memory contents must be executed from RAM. Executing code from flash automatically disables the marginal read mode. The marginal read modes are controlled by the MRG0 and MRG1 register bits. Setting MRG1 is used to detect insufficiently programmed flash cells containing a 1 (erased bits). Setting MRG0 is used to detect insufficiently programmed flash cells containing a 0 (programmed bits). Only one of these bits should be set at a time. Therefore, a full marginal read check requires two passes of checking the flash memory content's integrity. During marginal read mode, the flash access speed (MCLK) must be limited to 1 MHz (see the device-specific data sheet).

#### 7.3.7 Configuring and Accessing the Flash Memory Controller

The FCTLx registers are 16-bit password-protected read/write registers. Any read or write access must use word instructions and write accesses must include the write password 0A5h in the upper byte. Any write to any FCTLx register with any value other than 0A5h in the upper byte is a security key violation, sets the KEYV flag and triggers a PUC system reset. Any read of any FCTLx registers reads 096h in the upper byte.

Any write to FCTL1 during an erase or byte or word write operation is an access violation and sets ACCVIFG. Writing to FCTL1 is allowed in block write mode when WAIT = 1, but writing to FCTL1 in block write mode when WAIT = 0 is an access violation and sets ACCVIFG.

Any write to FCTL2 when the BUSY = 1 is an access violation.

Any FCTLx register may be read when BUSY = 1. A read does not cause an access violation.

### 7.3.8 Flash Memory Controller Interrupts

The flash controller has two interrupt sources, KEYV, and ACCVIFG. ACCVIFG is set when an access violation occurs. When the ACCVIE bit is re-enabled after a flash write or erase, a set ACCVIFG flag generates an interrupt request. ACCVIFG sources the NMI interrupt vector, so it is not necessary for GIE to be set for ACCVIFG to request an interrupt. ACCVIFG may also be checked by software to determine if an access violation occurred. ACCVIFG must be reset by software.

The key violation flag KEYV is set when any of the flash control registers are written with an incorrect password. When this occurs, a PUC is generated immediately resetting the device.

#### 7.3.9 Programming Flash Memory Devices

There are three options for programming an MSP430 flash device. All options support in-system programming:



- Program via JTAG
- Program via the bootstrap loader
- · Program via a custom solution

#### 7.3.9.1 Programming Flash Memory via JTAG

MSP430 devices can be programmed via the JTAG port. The JTAG interface requires four signals (five signals on 20- and 28-pin devices), ground and, optionally, V<sub>CC</sub> and RST/NMI.

The JTAG port is protected with a fuse. Blowing the fuse completely disables the JTAG port and is not reversible. Further access to the device via JTAG is not possible. For details, see the *MSP430 Programming Via the JTAG Interface User's Guide* (SLAU320).

#### 7.3.9.2 Programming Flash Memory via the Bootstrap Loader (BSL)

Most MSP430 flash devices contain a bootstrap loader. See the device-specific data sheet for implementation details. The BSL enables users to read or program the flash memory or RAM using a UART serial interface. Access to the MSP430 flash memory via the BSL is protected by a 256-bit user-defined password. For more details see the MSP430 Programming Via the Bootstrap Loader User's Guide (SLAU319).

#### 7.3.9.3 Programming Flash Memory via a Custom Solution

The ability of the MSP430 CPU to write to its own flash memory allows for in-system and external custom programming solutions as shown in Figure 7-12. The user can choose to provide data to the MSP430 through any means available (UART, SPI, etc.). User-developed software can receive the data and program the flash memory. Since this type of solution is developed by the user, it can be completely customized to fit the application needs for programming, erasing, or updating the flash memory.

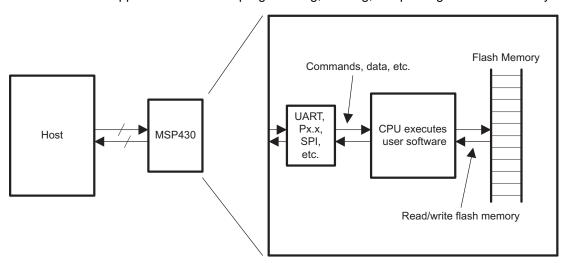
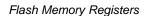


Figure 7-12. User-Developed Programming Solution



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# 7.4 Flash Memory Registers

The flash memory registers are listed in Table 7-4.

# **Table 7-4. Flash Memory Registers**

Register	Short Form	Register Type	Address	Initial State
Flash memory control register 1	FCTL1	Read/write	0x0128	0x9600 with PUC
Flash memory control register 2	FCTL2	Read/write	0x012A	0x9642 with PUC
Flash memory control register 3	FCTL3	Read/write	0x012C	0x9658 with PUC <sup>(1)</sup>
Flash memory control register 4 <sup>(2)</sup>	FCTL4	Read/write	0x01BE	0x0000 with PUC
Interrupt Enable 1	IE1	Read/write	0x0000	Reset with PUC
Interrupt Flag 1	IFG1	Read/write	0x0002	

<sup>(1)</sup> KEYV is reset with POR.

<sup>(2)</sup> Not present in all devices. See device-specific data sheet.



Flash Memory Registers www.ti.com

# 7.4.1 FCTL1, Flash Memory Control Register

15	14	13	12	11	10	9	8		
FRKEY, Read as 096h									
FWKEY, Must be written as 0A5h									

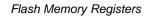
FWKEY, Must be written as 0A5h								
7	6	5	4	3	2	1	0	
BLKWRT	WRT	Reserved	EEIEX <sup>(1)</sup>	EEI <sup>(1)</sup>	MERAS	ERASE	Reserved	
rw-0	rw-0	r0	rw-0	rw-0	rw-0	rw-0	r0	
FRKEY FWKEY	Bits 15-8	FCTLx password. Alw	ays reads as 096	h. Must be written	as 0A5h. Writing	any other value g	enerates a PUC.	
BLKWRT	Bit 7	Block write mode. WF set.	RT must also be se	et for block write r	node. BLKWRT is	automatically res	et when EMEX is	
		0 Block-write mod	e is off					
		1 Block-write mod	e is on					
WRT	Bit 6	Write. This bit is used	to select any writ	e mode. WRT is a	automatically reset	when EMEX is s	et.	
		0 Write mode is o	ff					
		1 Write mode is o	n					
Reserved	Bit 5	Reserved. Always rea	ıd as 0.					
EEIEX	Bit 4	Enable Emergency Interrupt Exit. Setting this bit enables an interrupt to cause an emergency exit from a floperation when GIE = 1. EEIEX is automatically reset when EMEX is set.				exit from a flash		
		0 Exit interrupt dis	abled.					
		1 Exit on interrupt	enabled.					
EEI	Bits 3	Enable Erase Interrupts. Setting this bit allows a segment erase to be interrupted by an interrupt reque After the interrupt is serviced the erase cycle is resumed.						
		0 Interrupts during	segment erase d	lisabled.				
		1 Interrupts during	segment erase e	enabled.				
MERAS ERASE	Bit 2 Bit 1	Mass erase and erase automatically reset wi		used together to s	elect the erase mo	de. MERAS and	ERASE are	
MERAS	ERASE	Erase Cycle						
0	0	No erase						
0	1	Erase individual segn	nent only					
1	0	Erase all main memo	rase all main memory segments					

WERAS	EKASE	Erase Cycle
0	0	No erase
0	1	Erase individual segment only
1	0	Erase all main memory segments
1	1	LOCKA = 0: Erase main and information flash memory. LOCKA = 1: Erase only main flash memory.
Reserved	Bit 0	Reserved. Always read as 0.

Not present on MSP430x20xx and MSP430G2xx devices.

# 7.4.2 FCTL2, Flash Memory Control Register

15	14	13	12	11	10	9	8	
			,	ead as 096h tten as 0A5h				
7	6	5	4	3	2	1	0	
F	FSSELx FNx							
rw-0	rw-1	rw-0	rw-0	rw-0	rw-0	rw-1	rw-0	
<b>FWKEYx</b>	Bits 15-8	FCTLx password. Always reads as 096h. Must be written as 0A5h. Writing any other value generates a PUC.						
<b>FSSEL</b> x	Bits 7-6	Flash controller clock	source select					
		00 ACLK						
		01 MCLK						
		10 SMCLK						
		11 SMCLK						
FNx	Bits 5-0	Flash controller clock is FNx + 1. For examp						





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## 7.4.3 FCTL3, Flash Memory Control Register

15	14		13	12	11	10	9	8
					ead as 096h tten as 0A5h			
7	6		5	4	3	2	1	0
FAIL	LOCK	(A	EMEX	LOCK	WAIT	ACCVIFG	KEYV	BUSY
r(w)-0	r(w)-	1	rw-0	rw-1	r-1	rw-0	rw-(0)	r(w)-0
FWKEYx	Bits 15-8	FCTLx pa	assword. Alv	vays reads as 096	h. Must be writter	n as 0A5h. Writing a	ny other value ge	enerates a PUC.
FAIL	Bit 7			s bit is set if the fF L must be reset wi		fails, or a flash oper	ration is aborted f	rom an interrupt
		0 No	failure					
		1 Fai	lure					
LOCKA	Bit 6	Segment	A and Info Id	ock. Write a 1 to th	is bit to change it	s state. Writing 0 ha	as no effect.	
		0 Seg	gment A unio	ocked and all inforr	mation memory is	erased during a ma	ass erase.	
		1 Se	gment A lock	ed and all informa	tion memory is p	rotected from erasu	re during a mass	erase.
EMEX	Bit 5	Emergen	cy exit					
		0 No	emergency	exit				
		1 Em	ergency exit					
LOCK	Bit 4	byte or w	ord write or	erase operation, a	nd the operation	ising. The LOCK bit completes normally RT and WAIT are re	. In the block writ	e mode if the
		0 Unl	ocked					
		1 Loc	ked					
WAIT	Bit 3	Wait. Ind	icates the fla	ish memory is beir	ng written to.			
		0 The	e flash memo	ory is not ready for	the next byte/wo	rd write		
		1 The	e flash memo	ory is ready for the	next byte/word w	vrite		
ACCVIFG	Bit 2	Access v	iolation inter	rupt flag				
		0 No	interrupt per	nding				
		1 Inte	errupt pendin	g				
KEYV	Bit 1	Flash security key violation. This bit indicates an incorrect FCTLx password was written to any flas register and generates a PUC when set. KEYV must be reset with software.						
		0 FC	TLx passwor	d was written corre	ectly			
		1 FC	TLx passwor	d was written inco	rrectly			
BUSY	Bit 0	Busy. Th	is bit indicate	es the status of the	flash timing gen	erator.		
		0 Not	Busy					
		1 Bus	sy					



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## 7.4.4 FCTL4, Flash Memory Control Register

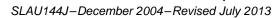
This register is not available in all devices. See the device-specific data sheet for details.

15	14	13	12	11	10	9	8		
			<b>FWKEYx</b> , Re Must be writt						
7	6	5	4	3	2	1	0		
		MRG1	MRG0						
r-0	r-0	rw-0	rw-0	r-0	r-0	r-0	r-0		
<b>FWKEYx</b>	Bits 15-8	FCTLx password. Alwa	ays reads as 096h	. Must be written	as 0A5h. Writing a	any other value ge	enerates a PUC.		
Reserved	Bits 7-6	Reserved. Always read	d as 0.						
MRG1	Bit 5		Marginal read 1 mode. This bit enables the marginal 1 read mode. The marginal read 1 bit is cleared if the CPU starts execution from the flash memory. If both MRG1 and MRG0 are set MRG1 is active and MRG0 is ignored.						
		0 Marginal 1 read i	mode is disabled.						
		1 Marginal 1 read i	mode is enabled.						
MRG0	Bit 4	•	Marginal read 0 mode. This bit enables the marginal 0 read mode. The marginal mode 0 is cleared if the CPU starts execution from the flash memory. If both MRG1 and MRG0 are set MRG1 is active and MRG0 is ignored.						
		0 Marginal 0 read i	mode is disabled.						
		1 Marginal 0 read i	mode is enabled.						
Reserved	Bits 3-0	Reserved. Always read	d as 0.						

## 7.4.5 IE1, Interrupt Enable Register 1

7	6		5	4	3	2	1	0						
			ACCVIE											
			rw-0											
	Bits 7-6 These bits may be used by other modules. See the device-specific data sheet.													
ACCVIE	Bit 5	IE1 r	Flash memory access violation interrupt enable. This bit enables the ACCVIFG interrupt. Because other bits in IE1 may be used for other modules, it is recommended to set or clear this bit using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.											
		0	0 Interrupt not enabled											
		1	Interrupt enabled											
	Bits 4-0	Thes	se bits may be use	d by other module	es. See the device	e-specific data she	These bits may be used by other modules. See the device-specific data sheet.							

## Chapter 8





# Digital I/O

This chapter describes the operation of the digital I/O ports.

Topic Page

8.1	Digital I/O Introduction	328
8.2	Digital I/O Operation	328
8.3	Digital I/O Registers	333



Digital I/O Introduction www.ti.com

#### 8.1 **Digital I/O Introduction**

MSP430 devices have up to eight digital I/O ports implemented, P1 to P8. Each port has up to eight I/O pins. Every I/O pin is individually configurable for input or output direction, and each I/O line can be individually read or written to.

Ports P1 and P2 have interrupt capability. Each interrupt for the P1 and P2 I/O lines can be individually enabled and configured to provide an interrupt on a rising edge or falling edge of an input signal. All P1 I/O lines source a single interrupt vector, and all P2 I/O lines source a different, single interrupt vector.

The digital I/O features include:

- Independently programmable individual I/Os
- Any combination of input or output
- Individually configurable P1 and P2 interrupts
- Independent input and output data registers
- Individually configurable pullup or pulldown resistors
- Individually configurable pin-oscillator function (some MSP430 devices)

NOTE: MSP430G22x0: These devices feature digital I/O pins P1.2, P1.5, P1.6 and P1.7. The GPIOs P1.0, P1.1, P1.3, P1.4, P2.6, and P2.7 are implemented on this device but not available on the device pin-out. To avoid floating inputs, these GPIOs, these digital I/Os should be properly initialized by running a start-up code. See initialization code below: mov.b #0x1B, P1REN; ; Terminate unavailable Port1 pins properly ; Config as Input with pulldown enabled

xor.b #0x20, BCSCTL3; ; Select VLO as low freq clock

The initialization code configures GPIOs P1.0, P1.1, P1.3, and P1.4 as inputs with pull-down resistor enabled (that is, P1REN.x = 1) and GPIOs P2.6 and P2.7 are terminated by selecting VLOCLK as ACLK - see the Basic Clock System chapter for details. The register bits of P1.0, P1.1, P1.3, and P1.4 in registers P1OUT, P1DIR, P1IFG, P1IE, P1IES, P1SEL and P1REN should not be altered after the initialization code is executed. Also, all Port2 registers are should not be altered.

#### 8.2 **Digital I/O Operation**

The digital I/O is configured with user software. The setup and operation of the digital I/O is discussed in the following sections.

#### 8.2.1 Input Register PxIN

Each bit in each PxIN register reflects the value of the input signal at the corresponding I/O pin when the pin is configured as I/O function.

Bit = 0: The input is low

Bit = 1: The input is high

#### NOTE: Writing to Read-Only Registers PxIN

Writing to these read-only registers results in increased current consumption while the write attempt is active.

#### 8.2.2 Output Registers PxOUT

Each bit in each PxOUT register is the value to be output on the corresponding I/O pin when the pin is configured as I/O function, output direction, and the pullup/down resistor is disabled.

Bit = 0: The output is low

Bit = 1: The output is high



www.ti.com Digital I/O Operation

If the pin's pullup/pulldown resistor is enabled, the corresponding bit in the PxOUT register selects pullup or pulldown.

Bit = 0: The pin is pulled down

Bit = 1: The pin is pulled up

#### 8.2.3 Direction Registers PxDIR

Each bit in each PxDIR register selects the direction of the corresponding I/O pin, regardless of the selected function for the pin. PxDIR bits for I/O pins that are selected for other functions must be set as required by the other function.

Bit = 0: The port pin is switched to input direction

Bit = 1: The port pin is switched to output direction

#### 8.2.4 Pullup/Pulldown Resistor Enable Registers PxREN

Each bit in each PxREN register enables or disables the pullup/pulldown resistor of the corresponding I/O pin. The corresponding bit in the PxOUT register selects if the pin is pulled up or pulled down.

Bit = 0: Pullup/pulldown resistor disabled

Bit = 1: Pullup/pulldown resistor enabled

### 8.2.5 Function Select Registers PxSEL and PxSEL2

Port pins are often multiplexed with other peripheral module functions. See the device-specific data sheet to determine pin functions. Each PxSEL and PxSEL2 bit is used to select the pin function - I/O port or peripheral module function.

Table 8-1. PxSEL and PxSEL2

PxSEL2	PxSEL	Pin Function
0	0	I/O function is selected.
0	1	Primary peripheral module function is selected.
1	0	Reserved. See device-specific data sheet.
1	1	Secondary peripheral module function is selected.

Setting PxSELx = 1 does not automatically set the pin direction. Other peripheral module functions may require the PxDIRx bits to be configured according to the direction needed for the module function. See the pin schematics in the device-specific data sheet.

#### NOTE: Setting PxREN = 1 When PxSEL = 1

On some I/O ports on the MSP430F261x and MSP430F2416/7/8/9, enabling the pullup/pulldown resistor (PxREN = 1) while the module function is selected (PxSEL = 1) does not disable the logic output driver. This combination is not recommended and may result in unwanted current flow through the internal resistor. See the device-specific data sheet pin schematics for more information.

```
;Output ACLK on P2.0 on MSP430F21x1
BIS.B #01h,&P2SEL ; Select ACLK function for pin
BIS.B #01h,&P2DIR ; Set direction to output *Required*
```

#### NOTE: P1 and P2 Interrupts Are Disabled When PxSEL = 1

When any P1SELx or P2SELx bit is set, the corresponding pin's interrupt function is disabled. Therefore, signals on these pins will not generate P1 or P2 interrupts, regardless of the state of the corresponding P1IE or P2IE bit.



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When a port pin is selected as an input to a peripheral, the input signal to the peripheral is a latched representation of the signal at the device pin. While PxSELx = 1, the internal input signal follows the signal at the pin. However, if the PxSELx = 0, the input to the peripheral maintains the value of the input signal at the device pin before the PxSELx bit was reset.

#### 8.2.6 Pin Oscillator

Some MSP430 devices have a pin oscillator function built-in to some pins. The pin oscillator function may be used in capacitive touch sensing applications to eliminate external passive components. Additionally, the pin oscillator may be used in sensor applications.

No external components to create the oscillation

Capacitive sensors can be connected directly to MSP430 pin

Robust, typical built-in hysteresis of ~0.7 V

When the pin oscillator function is enabled, other pin configurations are overwritten. The output driver is turned off while the weak pullup/pulldown is enabled and controlled by the voltage level on the pin itself. The voltage on the I/O is fed into the Schmitt trigger of the pin and then routed to a timer. The connection to the timer is device specific and, thus, defined in the device-specific data sheet. The Schmitt-trigger output is inverted and then decides if the pullup or the pulldown is enabled. Due to the inversion, the pin starts to oscillate as soon as the pin oscillator pin configuration is selected. Some of the pin-oscillator outputs are combined by a logical OR before routing to a timer clock input or timer capture channel. Therefore, only one pin oscillator should be enabled at a time. The oscillation frequency of each pin is defined by the load on the pin and by the I/O type. I/Os with analog functions typically show a lower oscillation frequency than pure digital I/Os. See the device-specific data sheet for details. Pins without external load show typical oscillation frequencies of 1 MHz to 3 MHz.

#### Pin oscillator in a cap touch application

A typical touch pad application using the pin oscilator is shown in Figure 8-1.

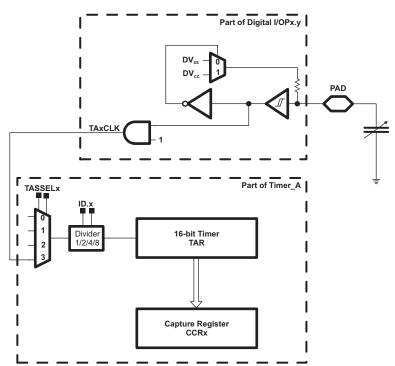


Figure 8-1. Example Circuitry and Configuration using the Pin Oscillator

A change of the capacitance of the touch pad (external capacitive load) has an effect on the pin oscillator frequency. An approaching finger tip increases the capacitance of the touch pad thus leads to a lower self-oscillation frequency due to the longer charging time. The oscillation frequency can directly be captured in a built-in Timer channel. The typical sensitivity of a pin is shown in Figure 8-2.



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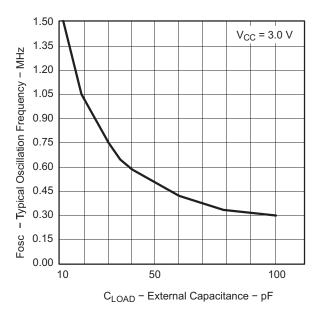


Figure 8-2. Typical Pin-Oscillation Frequency

## 8.2.7 P1 and P2 Interrupts

Each pin in ports P1 and P2 have interrupt capability, configured with the PxIFG, PxIE, and PxIES registers. All P1 pins source a single interrupt vector, and all P2 pins source a different single interrupt vector. The PxIFG register can be tested to determine the source of a P1 or P2 interrupt.

#### 8.2.7.1 Interrupt Flag Registers P1IFG, P2IFG

Each PxIFGx bit is the interrupt flag for its corresponding I/O pin and is set when the selected input signal edge occurs at the pin. All PxIFGx interrupt flags request an interrupt when their corresponding PxIE bit and the GIE bit are set. Each PxIFG flag must be reset with software. Software can also set each PxIFG flag, providing a way to generate a software initiated interrupt.

Bit = 0: No interrupt is pending

Bit = 1: An interrupt is pending

Only transitions, not static levels, cause interrupts. If any PxIFGx flag becomes set during a Px interrupt service routine, or is set after the RETI instruction of a Px interrupt service routine is executed, the set PxIFGx flag generates another interrupt. This ensures that each transition is acknowledged.

#### NOTE: PxIFG Flags When Changing PxOUT or PxDIR

Writing to P10UT, P1DIR, P20UT, or P2DIR can result in setting the corresponding P1IFG or P2IFG flags.

#### 8.2.7.2 Interrupt Edge Select Registers P1IES, P2IES

Each PxIES bit selects the interrupt edge for the corresponding I/O pin.

Bit = 0: The PxIFGx flag is set with a low-to-high transition

Bit = 1: The PxIFGx flag is set with a high-to-low transition



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#### NOTE: Writing to PxIESx

Writing to P1IES, or P2IES can result in setting the corresponding interrupt flags.

PxIESx	PxINx	PxIFGx
$0 \rightarrow 1$	0	May be set
$0 \rightarrow 1$	1	Unchanged
$1 \rightarrow 0$	0	Unchanged
$1 \rightarrow 0$	1	May be set

#### 8.2.7.3 Interrupt Enable P1IE, P2IE

Each PxIE bit enables the associated PxIFG interrupt flag.

Bit = 0: The interrupt is disabled.

Bit = 1: The interrupt is enabled.

### 8.2.8 Configuring Unused Port Pins

Unused I/O pins should be configured as I/O function, output direction, and left unconnected on the PC board, to prevent a floating input and reduce power consumption. The value of the PxOUT bit is irrelevant, since the pin is unconnected. Alternatively, the integrated pullup/pulldown resistor can be enabled by setting the PxREN bit of the unused pin to prevent the floating input. See the *System Resets, Interrupts, and Operating Modes* chapter for termination of unused pins.



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## 8.3 Digital I/O Registers

The digital I/O registers are listed in Table 8-2.

Table 8-2. Digital I/O Registers

Port	Register	Short Form	Address	Register Type	Initial State
	Input	P1IN	020h	Read only	-
	Output	P1OUT	021h	Read/write	Unchanged
	Direction	P1DIR	022h	Read/write	Reset with PUC
	Interrupt Flag	P1IFG	023h	Read/write	Reset with PUC
P1	Interrupt Edge Select	P1IES	024h	Read/write	Unchanged
	Interrupt Enable	P1IE	025h	Read/write	Reset with PUC
	Port Select	P1SEL	026h	Read/write	Reset with PUC
	Port Select 2	P1SEL2	041h	Read/write	Reset with PUC
	Resistor Enable	P1REN	027h	Read/write	Reset with PUC
	Input	P2IN	028h	Read only	-
	Output	P2OUT	029h	Read/write	Unchanged
	Direction	P2DIR	02Ah	Read/write	Reset with PUC
	Interrupt Flag	P2IFG	02Bh	Read/write	Reset with PUC
P2	Interrupt Edge Select	P2IES	02Ch	Read/write	Unchanged
	Interrupt Enable	P2IE	02Dh	Read/write	Reset with PUC
	Port Select	P2SEL	02Eh	Read/write	0C0h with PUC
	Port Select 2	P2SEL2	042h	Read/write	Reset with PUC
	Resistor Enable	P2REN	02Fh	Read/write	Reset with PUC
	Input	P3IN	018h	Read only	-
	Output	P3OUT	019h	Read/write	Unchanged
P3	Direction	P3DIR	01Ah	Read/write	Reset with PUC
PS	Port Select	P3SEL	01Bh	Read/write	Reset with PUC
	Port Select 2	P3SEL2	043h	Read/write	Reset with PUC
	Resistor Enable	P3REN	010h	Read/write	Reset with PUC
	Input	P4IN	01Ch	Read only	-
	Output	P4OUT	01Dh	Read/write	Unchanged
P4	Direction	P4DIR	01Eh	Read/write	Reset with PUC
Γ4	Port Select	P4SEL	01Fh	Read/write	Reset with PUC
	Port Select 2	P4SEL2	044h	Read/write	Reset with PUC
	Resistor Enable	P4REN	011h	Read/write	Reset with PUC
	Input	P5IN	030h	Read only	-
	Output	P5OUT	031h	Read/write	Unchanged
P5	Direction	P5DIR	032h	Read/write	Reset with PUC
F3	Port Select	P5SEL	033h	Read/write	Reset with PUC
	Port Select 2	P5SEL2	045h	Read/write	Reset with PUC
	Resistor Enable	P5REN	012h	Read/write	Reset with PUC
	Input	P6IN	034h	Read only	-
	Output	P6OUT	035h	Read/write	Unchanged
P6	Direction	P6DIR	036h	Read/write	Reset with PUC
rσ	Port Select	P6SEL	037h	Read/write	Reset with PUC
	Port Select 2	P6SEL2	046h	Read/write	Reset with PUC
	Resistor Enable	P6REN	013h	Read/write	Reset with PUC



Digital I/O Registers www.ti.com

## Table 8-2. Digital I/O Registers (continued)

Port	Register	Short Form	Address	Register Type	Initial State
	Input	P7IN	038h	Read only	-
	Output	P7OUT	03Ah	Read/write	Unchanged
P7	Direction	P7DIR	03Ch	Read/write	Reset with PUC
Ρ/	Port Select	P7SEL	03Eh	Read/write	Reset with PUC
	Port Select 2	P7SEL2	047h	Read/write	Reset with PUC
	Resistor Enable	P7REN	014h	Read/write	Reset with PUC
	Input	P8IN	039h	Read only	-
	Output	P8OUT	03Bh	Read/write	Unchanged
P8	Direction	P8DIR	03Dh	Read/write	Reset with PUC
Po	Port Select	P8SEL	03Fh	Read/write	Reset with PUC
	Port Select 2	P8SEL2	048h	Read/write	Reset with PUC
	Resistor Enable	P8REN	015h	Read/write	Reset with PUC



# Supply Voltage Supervisor (SVS)

This chapter describes the operation of the SVS. The SVS is implemented in selected MSP430x2xx devices.

Topic Page

9.1	Supply Voltage Supervisor (SVS) Introduction	336
9.2	SVS Operation	337
9.3	SVS Registers	339



### 9.1 Supply Voltage Supervisor (SVS) Introduction

The SVS is used to monitor the  $AV_{CC}$  supply voltage or an external voltage. The SVS can be configured to set a flag or generate a POR reset when the supply voltage or external voltage drops below a user-selected threshold.

The SVS features include:

- AV<sub>CC</sub> monitoring
- · Selectable generation of POR
- Output of SVS comparator accessible by software
- Low-voltage condition latched and accessible by software
- 14 selectable threshold levels
- External channel to monitor external voltage

The SVS block diagram is shown in Figure 9-1.

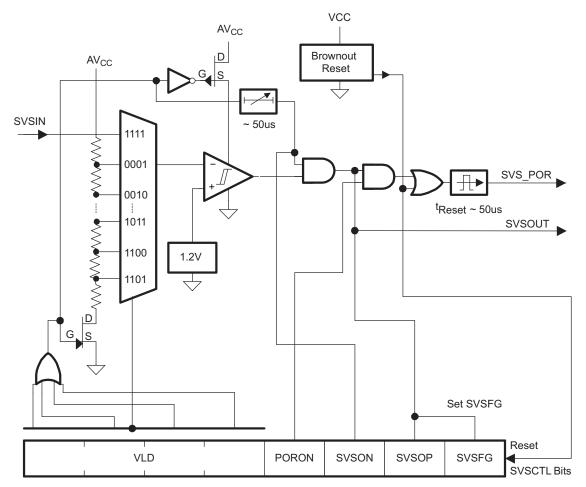


Figure 9-1. SVS Block Diagram



www.ti.com SVS Operation

#### 9.2 SVS Operation

The SVS detects if the  $AV_{CC}$  voltage drops below a selectable level. It can be configured to provide a POR or set a flag, when a low-voltage condition occurs. The SVS is disabled after a brownout reset to conserve current consumption.

#### 9.2.1 Configuring the SVS

The VLDx bits are used to enable/disable the SVS and select one of 14 threshold levels ( $V_{(SVS\_IT-)}$ ) for comparison with AV<sub>CC</sub>. The SVS is off when VLDx = 0 and on when VLDx > 0. The SVSON bit does not turn on the SVS. Instead, it reflects the on/off state of the SVS and can be used to determine when the SVS is on.

When VLDx = 1111, the external SVSIN channel is selected. The voltage on SVSIN is compared to an internal level of approximately 1.25 V.

#### 9.2.2 SVS Comparator Operation

A low-voltage condition exists when AV<sub>CC</sub> drops below the selected threshold or when the external voltage drops below its 1.25-V threshold. Any low-voltage condition sets the SVSFG bit.

The PORON bit enables or disables the device-reset function of the SVS. If PORON = 1, a POR is generated when SVSFG is set. If PORON = 0, a low-voltage condition sets SVSFG, but does not generate a POR.

The SVSFG bit is latched. This allows user software to determine if a low-voltage condition occurred previously. The SVSFG bit must be reset by user software. If the low-voltage condition is still present when SVSFG is reset, it will be immediately set again by the SVS.

#### 9.2.3 Changing the VLDx Bits

When the VLDx bits are changed from zero to any non-zero value there is a automatic settling delay  $t_{d(SVSon)}$  implemented that allows the SVS circuitry to settle. The  $t_{d(SVSon)}$  delay is approximately 50  $\mu$ s. During this delay, the SVS will not flag a low-voltage condition or reset the device, and the SVSON bit is cleared. Software can test the SVSON bit to determine when the delay has elapsed and the SVS is monitoring the voltage properly. Writing to SVSCTL while SVSON = 0 will abort the SVS automatic settling delay,  $t_{d(SVSon)}$ , and switch the SVS to active mode immediately. In doing so, the SVS circuitry might not be settled, resulting in unpredictable behavior.

When the VLDx bits are changed from any non-zero value to any other non-zero value the circuitry requires the time  $t_{\text{settle}}$  to settle. The settling time  $t_{\text{settle}}$  is a maximum of ~12  $\mu$ s. See the device-specific data sheet. There is no automatic delay implemented that prevents SVSFG to be set or to prevent a reset of the device. The recommended flow to switch between levels is shown in the following code.



SVS Operation www.ti.com

#### 9.2.4 SVS Operating Range

Each SVS level has hysteresis to reduce sensitivity to small supply voltage changes when  $AV_{CC}$  is close to the threshold. The SVS operation and SVS/Brownout interoperation are shown in Figure 9-2.

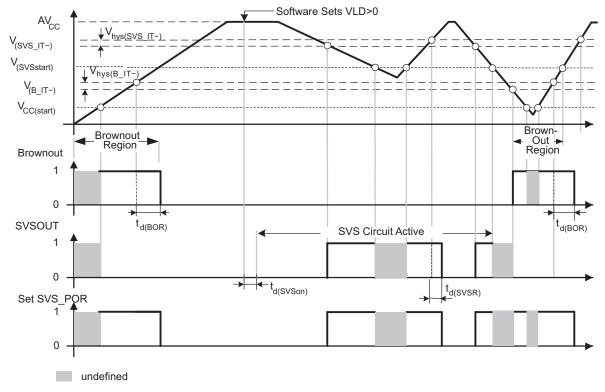


Figure 9-2. Operating Levels for SVS and Brownout/Reset Circuit



www.ti.com SVS Registers

## 9.3 SVS Registers

The SVS registers are listed in Table 9-1.

## Table 9-1. SVS Registers

Register	Short Form	Register Type	Address	Initial State
SVS Control Register	SVSCTL	Read/write	055h	Reset with BOR



SVS Registers www.ti.com

## 9.3.1 SVSCTL, SVS Control Register

7	6		5	4	3	2	1	0	
		VLDx			PORON	SVSON	SVSOP	SVSFG	
rw-0 <sup>(1)</sup>	rw-0 <sup>(1)</sup>		rw-0 <sup>(1)</sup>	rw-0 <sup>(1)</sup>	rw-0 <sup>(1)</sup>	r <sup>(1)</sup>	r <sup>(1)</sup>	rw-0 <sup>(1)</sup>	
VLDx	Bits 7-4	Voltage level detect. These bits turn on the SVS and select the nominal SVS threshold voltage level. See the device-specific data sheet for parameters.							
		0000	SVS is off						
		0001	1.9 V						
		0010	2.1 V						
		0011	2.2 V						
		0100	2.3 V						
		0101	2.4 V						
		0110	2.5 V						
		0111	2.65 V						
		1000	2.8 V						
		1001	2.9 V						
		1010	3.05 V						
		1011	3.2 V						
		1100	3.35 V						
		1101	3.5 V						
		1110	3.7 V						
		1111	•		age SVSIN to 1.25				
PORON	Bit 3	POR o	n. This bit en	ables the SVSFG	flag to cause a PC	OR device reset.			
		0		es not cause a PO	R				
		1	SVSFG cau	ises a POR					
SVSON	Bit 2		n. This bit refl on by setting		SVS operation. Th	nis bit DOES NOT	turn on the SVS.	The SVS is	
		0	SVS is Off						
		1	SVS is On						
SVSOP	Bit 1	SVS o	utput. This bit	reflects the output	it value of the SVS	comparator.			
		0	SVS compa	rator output is lov	<i>I</i>				
		1	SVS compa	rator output is hig	h				
SVSFG	Bit 0		ag. This bit in by software.	dicates a low volta	age condition. SVS	FG remains set at	ter a low voltage	condition until	
		0	No low volta	age condition occu	urred				
		1	A low condi	tion is present or	has occurred				

<sup>(1)</sup> Reset by a brownout reset only, not by a POR or PUC.



# Watchdog Timer+ (WDT+)

The watchdog timer+ (WDT+) is a 16-bit timer that can be used as a watchdog or as an interval timer. This chapter describes the WDT+ The WDT+ is implemented in all MSP430x2xx devices.

Topic Page

10.1	Watchdog Timer+ (WDT+) Introduction	342
10.2	Watchdog Timer+ Operation	344
10.3	Watchdog Timer+ Registers	346



#### 10.1 Watchdog Timer+ (WDT+) Introduction

The primary function of the WDT+ module is to perform a controlled system restart after a software problem occurs. If the selected time interval expires, a system reset is generated. If the watchdog function is not needed in an application, the module can be configured as an interval timer and can generate interrupts at selected time intervals.

Features of the watchdog timer+ module include:

- Four software-selectable time intervals
- Watchdog mode
- Interval mode
- Access to WDT+ control register is password protected
- Control of RST/NMI pin function
- · Selectable clock source
- Can be stopped to conserve power
- Clock fail-safe feature

The WDT+ block diagram is shown in Figure 10-1.

#### NOTE: Watchdog Timer+ Powers Up Active

After a PUC, the WDT+ module is automatically configured in the watchdog mode with an initial 32768 clock cycle reset interval using the DCOCLK. The user must setup or halt the WDT+ prior to the expiration of the initial reset interval.



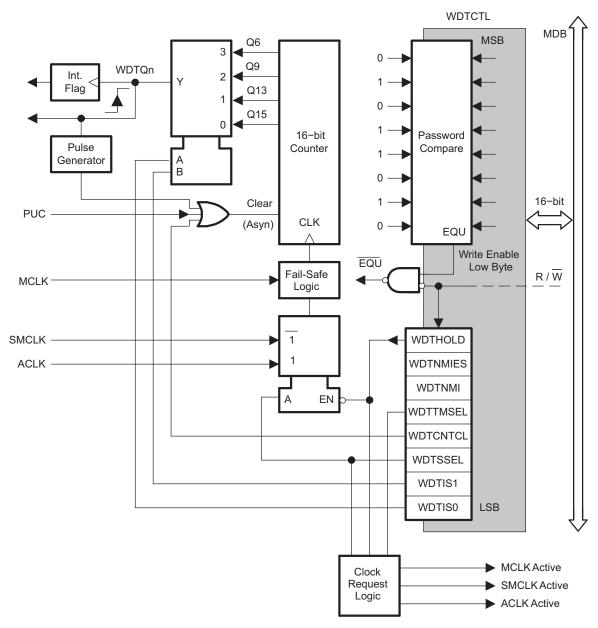


Figure 10-1. Watchdog Timer+ Block Diagram



#### 10.2 Watchdog Timer+ Operation

The WDT+ module can be configured as either a watchdog or interval timer with the WDTCTL register. The WDTCTL register also contains control bits to configure the RST/NMI pin. WDTCTL is a 16-bit, password-protected, read/write register. Any read or write access must use word instructions and write accesses must include the write password 05Ah in the upper byte. Any write to WDTCTL with any value other than 05Ah in the upper byte is a security key violation and triggers a PUC system reset regardless of timer mode. Any read of WDTCTL reads 069h in the upper byte. The WDT+ counter clock should be slower or equal than the system (MCLK) frequency.

## 10.2.1 Watchdog Timer+ Counter

The watchdog timer+ counter (WDTCNT) is a 16-bit up-counter that is not directly accessible by software. The WDTCNT is controlled and time intervals selected through the watchdog timer+ control register WDTCTL.

The WDTCNT can be sourced from ACLK or SMCLK. The clock source is selected with the WDTSSEL bit.

#### 10.2.2 Watchdog Mode

After a PUC condition, the WDT+ module is configured in the watchdog mode with an initial 32768 cycle reset interval using the DCOCLK. The user must setup, halt, or clear the WDT+ prior to the expiration of the initial reset interval or another PUC will be generated. When the WDT+ is configured to operate in watchdog mode, either writing to WDTCTL with an incorrect password, or expiration of the selected time interval triggers a PUC. A PUC resets the WDT+ to its default condition and configures the RST/NMI pin to reset mode.

#### 10.2.3 Interval Timer Mode

Setting the WDTTMSEL bit to 1 selects the interval timer mode. This mode can be used to provide periodic interrupts. In interval timer mode, the WDTIFG flag is set at the expiration of the selected time interval. A PUC is not generated in interval timer mode at expiration of the selected timer interval and the WDTIFG enable bit WDTIE remains unchanged.

When the WDTIE bit and the GIE bit are set, the WDTIFG flag requests an interrupt. The WDTIFG interrupt flag is automatically reset when its interrupt request is serviced, or may be reset by software. The interrupt vector address in interval timer mode is different from that in watchdog mode.

#### NOTE: Modifying the Watchdog Timer+

The WDT+ interval should be changed together with WDTCNTCL = 1 in a single instruction to avoid an unexpected immediate PUC or interrupt.

The WDT+ should be halted before changing the clock source to avoid a possible incorrect interval.

#### 10.2.4 Watchdog Timer+ Interrupts

The WDT+ uses two bits in the SFRs for interrupt control.

- The WDT+ interrupt flag, WDTIFG, located in IFG1.0
- The WDT+ interrupt enable, WDTIE, located in IE1.0

When using the WDT+ in the watchdog mode, the WDTIFG flag sources a reset vector interrupt. The WDTIFG can be used by the reset interrupt service routine to determine if the watchdog caused the device to reset. If the flag is set, then the watchdog timer+ initiated the reset condition either by timing out or by a security key violation. If WDTIFG is cleared, the reset was caused by a different source.

When using the WDT+ in interval timer mode, the WDTIFG flag is set after the selected time interval and requests a WDT+ interval timer interrupt if the WDTIE and the GIE bits are set. The interval timer interrupt vector is different from the reset vector used in watchdog mode. In interval timer mode, the WDTIFG flag is reset automatically when the interrupt is serviced, or can be reset with software.



### 10.2.5 Watchdog Timer+ Clock Fail-Safe Operation

The WDT+ module provides a fail-safe clocking feature assuring the clock to the WDT+ cannot be disabled while in watchdog mode. This means the low-power modes may be affected by the choice for the WDT+ clock. For example, if ACLK is the WDT+ clock source, LPM4 will not be available, because the WDT+ will prevent ACLK from being disabled. Also, if ACLK or SMCLK fail while sourcing the WDT+, the WDT+ clock source is automatically switched to MCLK. In this case, if MCLK is sourced from a crystal, and the crystal has failed, the fail-safe feature will activate the DCO and use it as the source for MCLK.

When the WDT+ module is used in interval timer mode, there is no fail-safe feature for the clock source.

#### 10.2.6 Operation in Low-Power Modes

The MSP430 devices have several low-power modes. Different clock signals are available in different low-power modes. The requirements of the user's application and the type of clocking used determine how the WDT+ should be configured. For example, the WDT+ should not be configured in watchdog mode with SMCLK as its clock source if the user wants to use low-power mode 3 because the WDT+ will keep SMCLK enabled for its clock source, increasing the current consumption of LPM3. When the watchdog timer+ is not required, the WDTHOLD bit can be used to hold the WDTCNT, reducing power consumption.

#### 10.2.7 Software Examples

Any write operation to WDTCTL must be a word operation with 05Ah (WDTPW) in the upper byte:

```
; Periodically clear an active watchdog
MOV #WDTPW+WDTCNTCL,&WDTCTL
;
; Change watchdog timer+ interval
MOV #WDTPW+WDTCNTL+WDTSSEL,&WDTCTL
;
; Stop the watchdog
MOV #WDTPW+WDTHOLD,&WDTCTL
;
; Change WDT+ to interval timer mode, clock/8192 interval
MOV #WDTPW+WDTCNTCL+WDTTMSEL+WDTISO,&WDTCTL
```



## 10.3 Watchdog Timer+ Registers

The WDT+ registers are listed in Table 10-1.

## Table 10-1. Watchdog Timer+ Registers

Register	Short Form	Register Type	Address	Initial State
Watchdog timer+ control register	WDTCTL	Read/write	0120h	06900h with PUC
SFR interrupt enable register 1	IE1	Read/write	0000h	Reset with PUC
SFR interrupt flag register 1	IFG1	Read/write	0002h	Reset with PUC (1)

<sup>(1)</sup> WDTIFG is reset with POR.



## 10.3.1 WDTCTL, Watchdog Timer+ Register

15	14	13	12	11	10	9	8
			,	ead as 069h tten as 05Ah			
7	6	5	4	3	2	1	0
WDTHOLD	WDTNMIE		WDTTMSEL	WDTCNTCL	WDTSSEL		TISx
rw-0	rw-0	rw-0	rw-0	r0(w)	rw-0	rw-0	rw-0
WDTPW	Bits 15-8	Watchdog timer+ pa	ssword. Always re	ad as 069h. Must	be written as 05A	h, or a PUC is ge	nerated.
WDTHOLD	Bit 7	Watchdog timer+ house conserves power		he watchdog timer	+. Setting WDTHO	DLD = 1 when the	WDT+ is not in
		0 Watchdog t	mer+ is not stoppe	ed			
		1 Watchdog t	mer+ is stopped				
WDTNMIES	Bit 6	Watchdog timer+ NN 1. Modifying this bit NMI.					
		0 NMI on risir	ig edge				
		1 NMI on falli	ng edge				
WDTNMI	Bit 5	Watchdog timer+ N	/II select. This bit s	selects the function	n for the RST/NMI	pin.	
		0 Reset funct	on				
		1 NMI function	า				
WDTTMSEL	Bit 4	Watchdog timer+ mo	ode select				
		0 Watchdog r	node				
		1 Interval time	er mode				
WDTCNTCL	Bit 3	Watchdog timer+ co automatically reset.	unter clear. Setting	g WDTCNTCL = 1	clears the count v	/alue to 0000h. W	DTCNTCL is
		0 No action					
		1 WDTCNT =	0000h				
WDTSSEL	Bit 2	Watchdog timer+ clo	ck source select				
		0 SMCLK					
		1 ACLK					
WDTISx	Bits 1-0	Watchdog timer+ int and/or generate a P		e bits select the wa	atchdog timer+ int	erval to set the W	DTIFG flag
		00 Watchdog o	lock source /3276	8			
		01 Watchdog of	lock source /8192				
		· ·	lock source /512				
		11 Watchdog o	lock source /64				



### 10.3.2 IE1, Interrupt Enable Register 1

7	6	5	4	3	2	1	0
			NMIIE				WDTIE
·							

rw-0

Bits 7-5

NMIIE Bit 4

These bits may be used by other modules. See device-specific data sheet.

NMI interrupt enable. This bit enables the NMI interrupt. Because other bits in IE1 may be used for other modules, it is recommended to set or clear this bit using BIS.B or BIC.B instructions, rather than MOV.B or

CLR.B instructions.

0 Interrupt not enabled

1 Interrupt enabled

Bits 3-1

These bits may be used by other modules. See device-specific data sheet.

WDTIE Bit 0

Watchdog timer+ interrupt enable. This bit enables the WDTIFG interrupt for interval timer mode. It is not necessary to set this bit for watchdog mode. Because other bits in IE1 may be used for other modules, it is recommended to set or clear this bit using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.

0 Interrupt not enabled

Interrupt pending

1 Interrupt enabled

#### 10.3.3 IFG1, Interrupt Flag Register 1

7	6	5	4	3	2	1	0
			NMIIFG				WDTIFG
			rw-0				rw-(0)
	Bits 7-5	These bits may be u	sed by other mode	ules. See device-s	specific data sheet		
NMIIFG	Bit 4	NMI interrupt flag. N modules, it is recom CLR.B instructions.					
		0 No interrupt	pending				
		1 Interrupt pe	nding				
	Bits 3-1	These bits may be u	sed by other mode	ules. See device-s	specific data sheet		
WDTIFG	Bit 0	Watchdog timer+ int mode, WDTIFG is re bits in IFG1 may be instructions, rather t	eset automatically used for other mo	by servicing the in dules, it is recomn	terrupt, or can be	reset by software.	Because other
		0 No interrupt	pending				



# Hardware Multiplier

This chapter describes the hardware multiplier. The hardware multiplier is implemented in some MSP430x2xx devices.

Topic Page

11.1 Hardware Multiplier Introduction	350
11.2 Hardware Multiplier Operation 3	350
11.3 Hardware Multiplier Registers	354



#### 11.1 Hardware Multiplier Introduction

The hardware multiplier is a peripheral and is not part of the MSP430 CPU. This means, its activities do not interfere with the CPU activities. The multiplier registers are peripheral registers that are loaded and read with CPU instructions.

The hardware multiplier supports:

- Unsigned multiply
- · Signed multiply
- · Unsigned multiply accumulate
- · Signed multiply accumulate
- 16x16 bits, 16x8 bits, 8x16 bits, 8x8 bits

The hardware multiplier block diagram is shown in Figure 11-1.

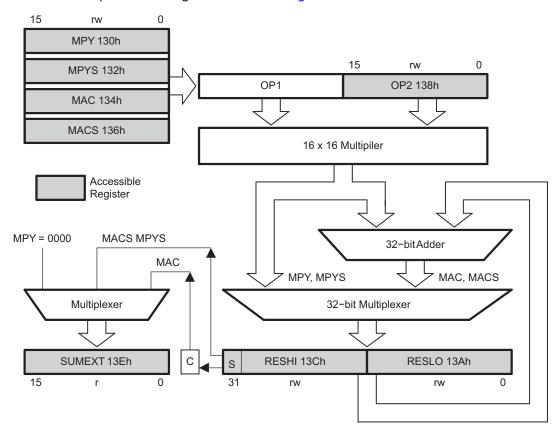


Figure 11-1. Hardware Multiplier Block Diagram

#### 11.2 Hardware Multiplier Operation

The hardware multiplier supports unsigned multiply, signed multiply, unsigned multiply accumulate, and signed multiply accumulate operations. The type of operation is selected by the address the first operand is written to.

The hardware multiplier has two 16-bit operand registers, OP1 and OP2, and three result registers, RESLO, RESHI, and SUMEXT. RESLO stores the low word of the result, RESHI stores the high word of the result, and SUMEXT stores information about the result. The result is ready in three MCLK cycles and can be read with the next instruction after writing to OP2, except when using an indirect addressing mode to access the result. When using indirect addressing for the result, a NOP is required before the result is ready.



### 11.2.1 Operand Registers

The operand one register OP1 has four addresses, shown in Table 11-1, used to select the multiply mode. Writing the first operand to the desired address selects the type of multiply operation but does not start any operation. Writing the second operand to the operand two register OP2 initiates the multiply operation. Writing OP2 starts the selected operation with the values stored in OP1 and OP2. The result is written into the three result registers RESLO, RESHI, and SUMEXT.

Repeated multiply operations may be performed without reloading OP1 if the OP1 value is used for successive operations. It is not necessary to re-write the OP1 value to perform the operations.

 OP1 Address
 Register Name
 Operation

 0130h
 MPY
 Unsigned multiply

 0132h
 MPYS
 Signed multiply

 0134h
 MAC
 Unsigned multiply accumulate

 0136h
 MACS
 Signed multiply accumulate

Table 11-1. OP1 Addresses

## 11.2.2 Result Registers

The result low register RESLO holds the lower 16-bits of the calculation result. The result high register RESHI contents depend on the multiply operation and are listed in Table 11-2.

Mode	RESHI Contents			
MPY	Upper 16-bits of the result			
MPYS	The MSB is the sign of the result. The remaining bits are the upper 15-bits of the result. Two's complement notation is used for the result.			
MAC	Upper 16-bits of the result			
MACS	Upper 16-bits of the result. Two's complement notation is used for the result.			

**Table 11-2. RESHI Contents** 

The sum extension registers SUMEXT contents depend on the multiply operation and are listed in Table 11-3.

Mode	SUMEXT
MPY	SUMEXT is always 0000h
	SUMEXT contains the extended sign of the result
MPYS	00000h = Result was positive or zero
	0FFFFh = Result was negative
	SUMEXT contains the carry of the result
MAC	0000h = No carry for result
	0001h = Result has a carry
	SUMEXT contains the extended sign of the result
MACS	00000h = Result was positive or zero
	0FFFFh = Result was negative

**Table 11-3. SUMEXT Contents** 



#### 11.2.2.1 MACS Underflow and Overflow

The multiplier does not automatically detect underflow or overflow in the MACS mode. The accumulator range for positive numbers is 0 to 7FFF FFFFh and for negative numbers is 0FFFF FFFFh to 8000 0000h. An underflow occurs when the sum of two negative numbers yields a result that is in the range for a positive number. An overflow occurs when the sum of two positive numbers yields a result that is in the range for a negative number. In both of these cases, the SUMEXT register contains the sign of the result, 0FFFFh for overflow and 0000h for underflow. User software must detect and handle these conditions appropriately.

#### 11.2.3 Software Examples

Examples for all multiplier modes follow. All 8x8 modes use the absolute address for the registers because the assembler will not allow .B access to word registers when using the labels from the standard definitions file.

There is no sign extension necessary in software. Accessing the multiplier with a byte instruction during a signed operation will automatically cause a sign extension of the byte within the multiplier module.

```
; 16x16 Unsigned Multiply
        #01234h,&MPY ; Load first operand
  MOV
  VOM
          #05678h,&OP2 ; Load second operand
                         ; Process results
; 8x8 Unsigned Multiply. Absolute addressing.
  MOV.B #012h,&0130h ; Load first operand
  MOV.B #034h,&0138h ; Load 2nd operand
                         ; Process results
; 16x16 Signed Multiply
  MOV #01234h, &MPYS ; Load first operand
  MOV
          #05678h,&OP2 ; Load 2nd operand
                         ; Process results
  . . .
; 8x8 Signed Multiply. Absolute addressing.
  MOV.B #012h,&0132h ; Load first operand
  MOV.B #034h,&0138h ; Load 2nd operand
                       ; Process results
; 16x16 Unsigned Multiply Accumulate
          #01234h,&MAC ; Load first operand
  MOV
  VOM
          #05678h,&OP2
                         ; Load 2nd operand
                         ; Process results
; 8x8 Unsigned Multiply Accumulate. Absolute addressing
  MOV.B #012h,&0134h ; Load first operand
  MOV.B #034h,&0138h ; Load 2nd operand
                         ; Process results
; 16x16 Signed Multiply Accumulate
          #01234h,&MACS ; Load first operand
  VOM
  VOM
          \#05678h,\&OP2; Load 2nd operand
                         ; Process results
; 8x8 Signed Multiply Accumulate. Absolute addressing
  MOV.B #012h,&0136h ; Load first operand
  \mbox{MOV.B} \ \ \ \mbox{\#034h,R5} \ \ \ \ \mbox{; Temp. location for 2nd operand}
                       ; Load 2nd operand
  MOV R5,&OP2
                         ; Process results
 . . .
```



### 11.2.4 Indirect Addressing of RESLO

When using indirect or indirect autoincrement addressing mode to access the result registers, At least one instruction is needed between loading the second operand and accessing one of the result registers:

```
; Access multiplier results with indirect addressing
  MOV
        #RESLO,R5
                   ; RESLO address in R5 for indirect
  MOV
        &OPER1,&MPY
                     ; Load 1st operand
                    ; Load 2nd operand
  MOV
        &OPER2,&OP2
  NOP
                     ; Need one cycle
                    ; Move RESLO
  VOM
        @R5+,&xxx
  MOV
      @R5,&xxx
                    ; Move RESHI
```

### 11.2.5 Using Interrupts

If an interrupt occurs after writing OP1, but before writing OP2, and the multiplier is used in servicing that interrupt, the original multiplier mode selection is lost and the results are unpredictable. To avoid this, disable interrupts before using the hardware multiplier or do not use the multiplier in interrupt service routines.

```
; Disable interrupts before using the hardware multiplier
DINT ; Disable interrupts
NOP ; Required for DINT
MOV #xxh,&MPY ; Load 1st operand
MOV #xxh,&OP2 ; Load 2nd operand
EINT ; Interrupts may be enable before
; Process results
```



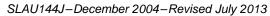
## 11.3 Hardware Multiplier Registers

The hardware multiplier registers are listed in Table 11-4.

**Table 11-4. Hardware Multiplier Registers** 

Register	Short Form	Register Type	Address	Initial State
Operand one - multiply	MPY	Read/write	0130h	Unchanged
Operand one - signed multiply	MPYS	Read/write	0132h	Unchanged
Operand one - multiply accumulate	MAC	Read/write	0134h	Unchanged
Operand one - signed multiply accumulate	MACS	Read/write	0136h	Unchanged
Operand two	OP2	Read/write	0138h	Unchanged
Result low word	RESLO	Read/write	013Ah	Undefined
Result high word	RESHI	Read/write	013Ch	Undefined
Sum extension register	SUMEXT	Read	013Eh	Undefined

## Chapter 12





## Timer\_A

 $\label{timer_A} \mbox{Timer_A is a 16-bit timer/counter with multiple capture/compare registers. This chapter describes the operation of the Timer_A of the MSP430x2xx device family.}$ 

Topic Page

12.1	Timer_A Introduction	356
12.2	Timer_A Operation	357
12.3	Timer_A Registers	369



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#### 12.1 Timer\_A Introduction

Timer\_A is a 16-bit timer/counter with three capture/compare registers. Timer\_A can support multiple capture/compares, PWM outputs, and interval timing. Timer\_A also has extensive interrupt capabilities. Interrupts may be generated from the counter on overflow conditions and from each of the capture/compare registers.

Timer\_A features include:

- · Asynchronous 16-bit timer/counter with four operating modes
- · Selectable and configurable clock source
- Two or three configurable capture/compare registers
- Configurable outputs with PWM capability
- · Asynchronous input and output latching
- Interrupt vector register for fast decoding of all Timer\_A interrupts

The block diagram of Timer\_A is shown in Figure 12-1.

#### NOTE: Use of the Word Count

Count is used throughout this chapter. It means the counter must be in the process of counting for the action to take place. If a particular value is directly written to the counter, then an associated action will not take place.



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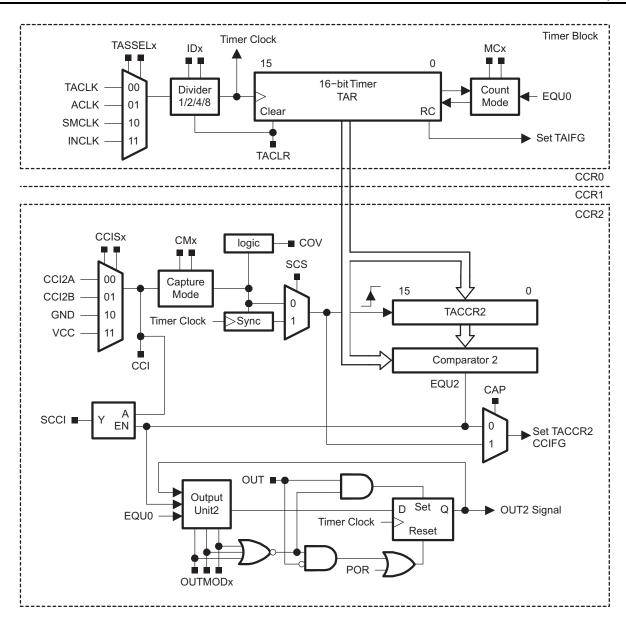


Figure 12-1. Timer\_A Block Diagram

#### 12.2 Timer\_A Operation

The Timer\_A module is configured with user software. The setup and operation of Timer\_A is discussed in the following sections.

#### 12.2.1 16-Bit Timer Counter

The 16-bit timer/counter register, TAR, increments or decrements (depending on mode of operation) with each rising edge of the clock signal. TAR can be read or written with software. Additionally, the timer can generate an interrupt when it overflows.

TAR may be cleared by setting the TACLR bit. Setting TACLR also clears the clock divider and count direction for up/down mode.



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#### NOTE: Modifying Timer\_A Registers

It is recommended to stop the timer before modifying its operation (with exception of the interrupt enable, and interrupt flag) to avoid errant operating conditions.

When the timer clock is asynchronous to the CPU clock, any read from TAR should occur while the timer is not operating or the results may be unpredictable. Alternatively, the timer may be read multiple times while operating, and a majority vote taken in software to determine the correct reading. Any write to TAR will take effect immediately.

#### 12.2.1.1 Clock Source Select and Divider

The timer clock can be sourced from ACLK, SMCLK, or externally via TACLK or INCLK. The clock source is selected with the TASSELx bits. The selected clock source may be passed directly to the timer or divided by 2, 4, or 8, using the IDx bits. The timer clock divider is reset when TACLR is set.

#### 12.2.2 Starting the Timer

The timer may be started, or restarted in the following ways:

- The timer counts when MCx > 0 and the clock source is active.
- When the timer mode is either up or up/down, the timer may be stopped by writing 0 to TACCR0. The
  timer may then be restarted by writing a nonzero value to TACCR0. In this scenario, the timer starts
  incrementing in the up direction from zero.

#### 12.2.3 Timer Mode Control

The timer has four modes of operation as described in Table 12-1: stop, up, continuous, and up/down. The operating mode is selected with the MCx bits.

MCx	Mode	Description
00	Stop	The timer is halted.
01	Up	The timer repeatedly counts from zero to the value of TACCR0.
10	Continuous	The timer repeatedly counts from zero to 0FFFFh.
11	Up/down	The timer repeatedly counts from zero up to the value of TACCR0 and back down to zero.

**Table 12-1. Timer Modes** 

#### 12.2.3.1 Up Mode

The up mode is used if the timer period must be different from 0FFFFh counts. The timer repeatedly counts up to the value of compare register TACCR0, which defines the period, as shown in Figure 12-2. The number of timer counts in the period is TACCR0+1. When the timer value equals TACCR0 the timer restarts counting from zero. If up mode is selected when the timer value is greater than TACCR0, the timer immediately restarts counting from zero.

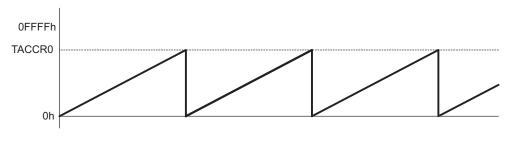


Figure 12-2. Up Mode

The TACCR0 CCIFG interrupt flag is set when the timer *counts* to the TACCR0 value. The TAIFG interrupt flag is set when the timer *counts* from TACCR0 to zero. Figure 12-3 shows the flag set cycle.



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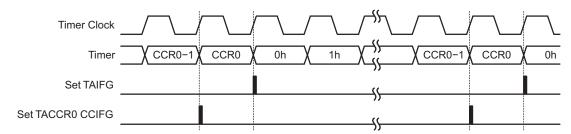


Figure 12-3. Up Mode Flag Setting

#### 12.2.3.2 Changing the Period Register TACCR0

When changing TACCR0 while the timer is running, if the new period is greater than or equal to the old period, or greater than the current count value, the timer counts up to the new period. If the new period is less than the current count value, the timer rolls to zero. However, one additional count may occur before the counter rolls to zero.

#### 12.2.3.3 Continuous Mode

In the continuous mode, the timer repeatedly counts up to 0FFFFh and restarts from zero as shown in Figure 12-4. The capture/compare register TACCR0 works the same way as the other capture/compare registers.

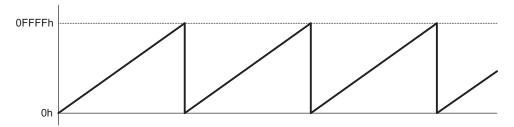


Figure 12-4. Continuous Mode

The TAIFG interrupt flag is set when the timer *counts* from 0FFFFh to zero. Figure 12-5 shows the flag set cycle.

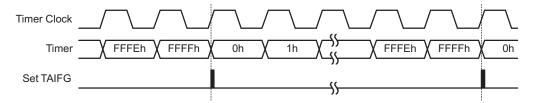


Figure 12-5. Continuous Mode Flag Setting

#### 12.2.3.4 Use of the Continuous Mode

The continuous mode can be used to generate independent time intervals and output frequencies. Each time an interval is completed, an interrupt is generated. The next time interval is added to the TACCRx register in the interrupt service routine. Figure 12-6 shows two separate time intervals  $t_0$  and  $t_1$  being added to the capture/compare registers. In this usage, the time interval is controlled by hardware, not software, without impact from interrupt latency. Up to three independent time intervals or output frequencies can be generated using all three capture/compare registers.



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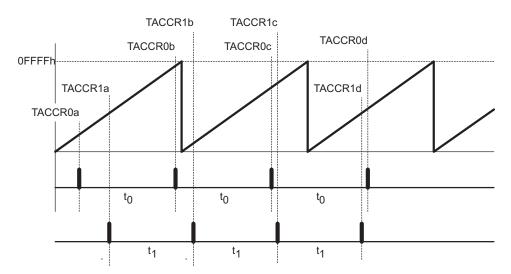


Figure 12-6. Continuous Mode Time Intervals

Time intervals can be produced with other modes as well, where TACCR0 is used as the period register. Their handling is more complex since the sum of the old TACCRx data and the new period can be higher than the TACCR0 value. When the previous TACCRx value plus  $t_x$  is greater than the TACCR0 data, TACCR0 + 1 must be subtracted to obtain the correct time interval.

#### 12.2.3.5 Up/Down Mode

The up/down mode is used if the timer period must be different from 0FFFFh counts, and if a symmetrical pulse generation is needed. The timer repeatedly counts up to the value of compare register TACCR0 and back down to zero, as shown in Figure 12-7. The period is twice the value in TACCR0.

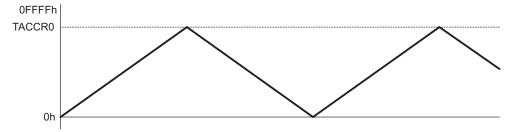


Figure 12-7. Up/Down Mode

The count direction is latched. This allows the timer to be stopped and then restarted in the same direction it was counting before it was stopped. If this is not desired, the TACLR bit must be set to clear the direction. The TACLR bit also clears the TAR value and the timer clock divider.

In up/down mode, the TACCR0 CCIFG interrupt flag and the TAIFG interrupt flag are set only once during a period, separated by 1/2 the timer period. The TACCR0 CCIFG interrupt flag is set when the timer *counts* from TACCR0 – 1 to TACCR0, and TAIFG is set when the timer completes *counting* down from 0001h to 0000h. Figure 12-8 shows the flag set cycle.



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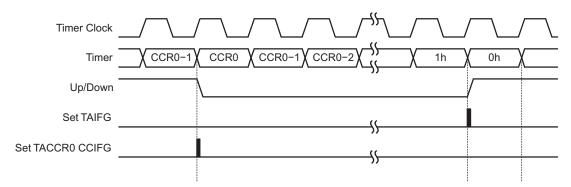


Figure 12-8. Up/Down Mode Flag Setting

## 12.2.3.6 Changing the Period Register TACCR0

When changing TACCR0 while the timer is running, and counting in the down direction, the timer continues its descent until it reaches zero. The value in TACCR0 is latched into TACL0 immediately, however the new period takes effect after the counter counts down to zero.

When the timer is counting in the up direction, and the new period is greater than or equal to the old period, or greater than the current count value, the timer counts up to the new period before counting down. When the timer is counting in the up direction, and the new period is less than the current count value, the timer begins counting down. However, one additional count may occur before the counter begins counting down.

## 12.2.3.7 Use of the Up/Down Mode

The up/down mode supports applications that require dead times between output signals (See section Timer A Output Unit). For example, to avoid overload conditions, two outputs driving an H-bridge must never be in a high state simultaneously. In the example shown in Figure 12-9 the t<sub>dead</sub> is:

 $t_{dead} = t_{timer} (TACCR1 - TACCR2)$ 

Where.

 $t_{dead}$  = Time during which both outputs need to be inactive

 $t_{timer}$  = Cycle time of the timer clock

TACCRx = Content of capture/compare register x

The TACCRx registers are not buffered. They update immediately when written to. Therefore, any required dead time will not be maintained automatically.



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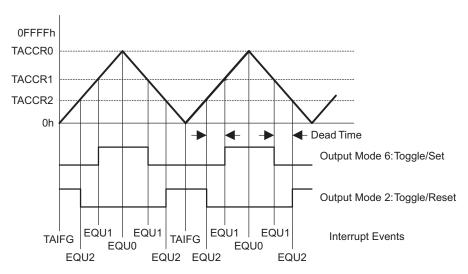


Figure 12-9. Output Unit in Up/Down Mode

## 12.2.4 Capture/Compare Blocks

Two or three identical capture/compare blocks, TACCRx, are present in Timer\_A. Any of the blocks may be used to capture the timer data, or to generate time intervals.

#### Capture Mode

The capture mode is selected when CAP = 1. Capture mode is used to record time events. It can be used for speed computations or time measurements. The capture inputs CCIxA and CCIxB are connected to external pins or internal signals and are selected with the CCISx bits. The CMx bits select the capture edge of the input signal as rising, falling, or both. A capture occurs on the selected edge of the input signal. If a capture occurs:

- The timer value is copied into the TACCRx register
- The interrupt flag CCIFG is set

The input signal level can be read at any time via the CCI bit. MSP430x2xx family devices may have different signals connected to CCIxA and CCIxB. See the device-specific data sheet for the connections of these signals.

The capture signal can be asynchronous to the timer clock and cause a race condition. Setting the SCS bit will synchronize the capture with the next timer clock. Setting the SCS bit to synchronize the capture signal with the timer clock is recommended. This is illustrated in Figure 12-10.

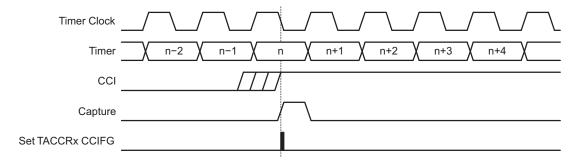


Figure 12-10. Capture Signal (SCS = 1)

Overflow logic is provided in each capture/compare register to indicate if a second capture was performed before the value from the first capture was read. Bit COV is set when this occurs as shown in Figure 12-11. COV must be reset with software.



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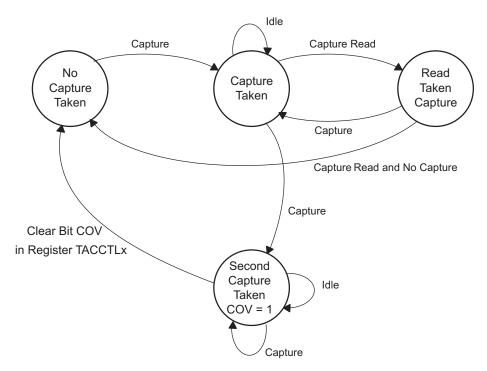


Figure 12-11. Capture Cycle

## 12.2.4.1 Capture Initiated by Software

Captures can be initiated by software. The CMx bits can be set for capture on both edges. Software then sets CCIS1 = 1 and toggles bit CCIS0 to switch the capture signal between  $V_{CC}$  and GND, initiating a capture each time CCIS0 changes state:

```
MOV #CAP+SCS+CCIS1+CM_3,&TACCTLx ; Setup TACCTLx XOR #CCIS0,&TACCTLx ; TACCTLx = TAR
```

#### 12.2.4.2 Compare Mode

The compare mode is selected when CAP = 0. The compare mode is used to generate PWM output signals or interrupts at specific time intervals. When TAR *counts* to the value in a TACCRx:

- Interrupt flag CCIFG is set
- Internal signal EQUx = 1
- EQUx affects the output according to the output mode
- · The input signal CCI is latched into SCCI

#### 12.2.5 Output Unit

Each capture/compare block contains an output unit. The output unit is used to generate output signals such as PWM signals. Each output unit has eight operating modes that generate signals based on the EQU0 and EQUx signals.

#### 12.2.5.1 Output Modes

The output modes are defined by the OUTMODx bits and are described in Table 12-2. The OUTx signal is changed with the rising edge of the timer clock for all modes except mode 0. Output modes 2, 3, 6, and 7 are not useful for output unit 0, because EQUx = EQU0.



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## **Table 12-2. Output Modes**

OUTMODx	Mode	Description
000	Output	The output signal OUTx is defined by the OUTx bit. The OUTx signal updates immediately when OUTx is updated.
001	Set	The output is set when the timer <i>counts</i> to the TACCRx value. It remains set until a reset of the timer, or until another output mode is selected and affects the output.
010	Toggle/Reset	The output is toggled when the timer <i>counts</i> to the TACCRx value. It is reset when the timer <i>counts</i> to the TACCR0 value.
011	Set/Reset	The output is set when the timer <i>counts</i> to the TACCRx value. It is reset when the timer <i>counts</i> to the TACCR0 value.
100	Toggle	The output is toggled when the timer <i>counts</i> to the TACCRx value. The output period is double the timer period.
101	Reset	The output is reset when the timer <i>counts</i> to the TACCRx value. It remains reset until another output mode is selected and affects the output.
110	Toggle/Set	The output is toggled when the timer <i>counts</i> to the TACCRx value. It is set when the timer <i>counts</i> to the TACCR0 value.
111	Reset/Set	The output is reset when the timer <i>counts</i> to the TACCRx value. It is set when the timer <i>counts</i> to the TACCR0 value.

## 12.2.5.2 Output Example — Timer in Up Mode

The OUTx signal is changed when the timer *counts* up to the TACCRx value, and rolls from TACCR0 to zero, depending on the output mode. An example is shown in Figure 12-12 using TACCR0 and TACCR1.

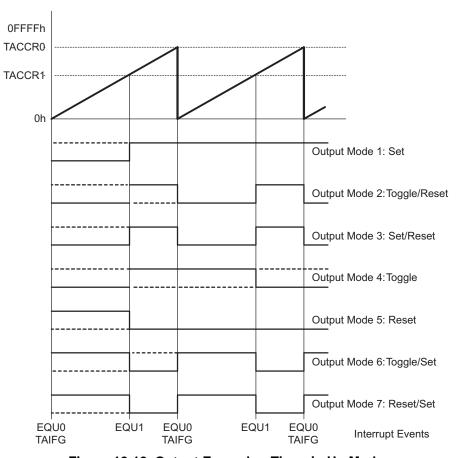


Figure 12-12. Output Example—Timer in Up Mode



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## 12.2.5.3 Output Example — Timer in Continuous Mode

The OUTx signal is changed when the timer reaches the TACCRx and TACCR0 values, depending on the output mode. An example is shown in Figure 12-13 using TACCR0 and TACCR1.

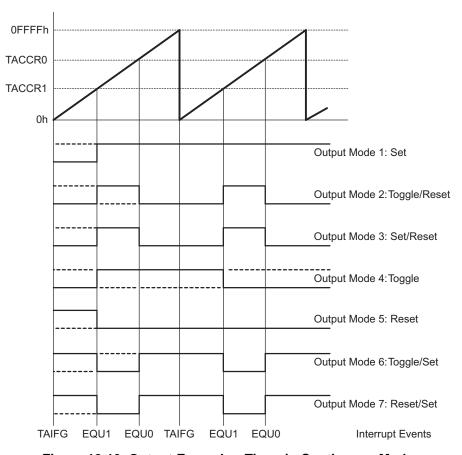


Figure 12-13. Output Example—Timer in Continuous Mode



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## 12.2.5.4 Output Example — Timer in Up/Down Mode

The OUTx signal changes when the timer equals TACCRx in either count direction and when the timer equals TACCR0, depending on the output mode. An example is shown in Figure 12-14 using TACCR0 and TACCR2.

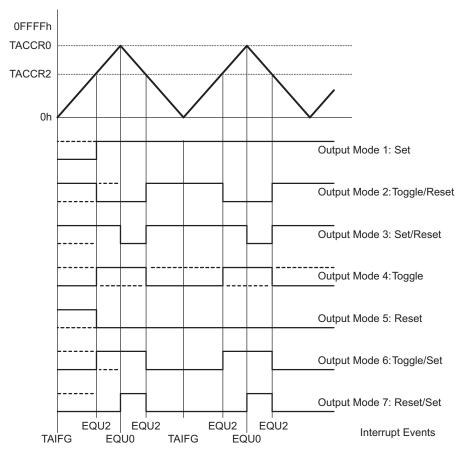


Figure 12-14. Output Example—Timer in Up/Down Mode

#### **NOTE: Switching Between Output Modes**

When switching between output modes, one of the OUTMODx bits should remain set during the transition, unless switching to mode 0. Otherwise, output glitching can occur because a NOR gate decodes output mode 0. A safe method for switching between output modes is to use output mode 7 as a transition state:

```
BIS #OUTMOD_7,&TACCTLx ; Set output mode=7
BIC #OUTMODx, &TACCTLx ; Clear unwanted bits
```



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## 12.2.6 Timer A Interrupts

Two interrupt vectors are associated with the 16-bit Timer\_A module:

- TACCR0 interrupt vector for TACCR0 CCIFG
- TAIV interrupt vector for all other CCIFG flags and TAIFG

In capture mode any CCIFG flag is set when a timer value is captured in the associated TACCRx register. In compare mode, any CCIFG flag is set if TAR *counts* to the associated TACCRx value. Software may also set or clear any CCIFG flag. All CCIFG flags request an interrupt when their corresponding CCIE bit and the GIE bit are set.

## 12.2.6.1 TACCR0 Interrupt

The TACCR0 CCIFG flag has the highest Timer\_A interrupt priority and has a dedicated interrupt vector as shown in Figure 12-15. The TACCR0 CCIFG flag is automatically reset when the TACCR0 interrupt request is serviced.

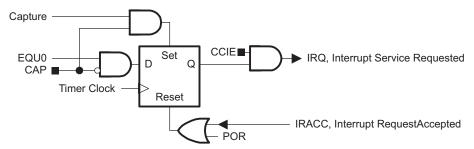


Figure 12-15. Capture/Compare TACCR0 Interrupt Flag

## 12.2.6.2 TAIV, Interrupt Vector Generator

The TACCR1 CCIFG, TACCR2 CCIFG, and TAIFG flags are prioritized and combined to source a single interrupt vector. The interrupt vector register TAIV is used to determine which flag requested an interrupt.

The highest priority enabled interrupt generates a number in the TAIV register (see register description). This number can be evaluated or added to the program counter to automatically enter the appropriate software routine. Disabled Timer\_A interrupts do not affect the TAIV value.

Any access, read or write, of the TAIV register automatically resets the highest pending interrupt flag. If another interrupt flag is set, another interrupt is immediately generated after servicing the initial interrupt. For example, if the TACCR1 and TACCR2 CCIFG flags are set when the interrupt service routine accesses the TAIV register, TACCR1 CCIFG is reset automatically. After the RETI instruction of the interrupt service routine is executed, the TACCR2 CCIFG flag will generate another interrupt.



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## 12.2.6.3 TAIV Software Example

The following software example shows the recommended use of TAIV and the handling overhead. The TAIV value is added to the PC to automatically jump to the appropriate routine.

The numbers at the right margin show the necessary CPU cycles for each instruction. The software overhead for different interrupt sources includes interrupt latency and return-from-interrupt cycles, but not the task handling itself. The latencies are:

- Capture/compare block TACCR0: 11 cycles
- Capture/compare blocks TACCR1, TACCR2: 16 cycles
- Timer overflow TAIFG: 14 cycles

; Interrupt CCIFG 0 HND	handler for TAC	CR0	CCIFG	Cycles
; RETI	; Start of har	ndle	r Interrupt latency	6 5
; Interrupt TA_HND	handler for TAI	FG,	TACCR1 and TACCR2 CCIFG	
		;	Interrupt latency	6
ADD	&TAIV,PC	;	Add offset to Jump table	3
RETI		;	Vector 0: No interrupt	5
JMP	CCIFG_1_HND	;	Vector 2: TACCR1	2
JMP	CCIFG_2_HND	;	Vector 4: TACCR2	2
RETI		;	Vector 6: Reserved	5
RETI		;	Vector 8: Reserved	5
TAIFG_HND		;	Vector 10: TAIFG Flag	
		;	Task starts here	
RETI				5
CCIFG_2_HND		;	Vector 4: TACCR2	
		;	Task starts here	
RETI		;	Back to main program	5
CCIFG_1_HND		;	Vector 2: TACCR1	
		;	Task starts here	
RETI		;	Back to main program	5



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## 12.3 Timer\_A Registers

The Timer\_A registers are listed in Table 12-3.

## Table 12-3. Timer\_A3 Registers

Register	Short Form	Register Type	Address	Initial State
Timer_A control	TACTL	Read/write	0160h	Reset with POR
Timer_A counter	TAR	Read/write	0170h	Reset with POR
Timer_A capture/compare control 0	TACCTL0	Read/write	0162h	Reset with POR
Timer_A capture/compare 0	TACCR0	Read/write	0172h	Reset with POR
Timer_A capture/compare control 1	TACCTL1	Read/write	0164h	Reset with POR
Timer_A capture/compare 1	TACCR1	Read/write	0174h	Reset with POR
Timer_A capture/compare control 2	TACCTL2 <sup>(1)</sup>	Read/write	0166h	Reset with POR
Timer_A capture/compare 2	TACCR2 <sup>(1)</sup>	Read/write	0176h	Reset with POR
Timer_A interrupt vector	TAIV	Read only	012Eh	Reset with POR

Not present on MSP430 devices with Timer\_A2 like MSP430F20xx and other devices.



Timer\_A Registers www.ti.com

## 12.3.1 TACTL, Timer\_A Control Register

15	14		13	12	11	10	9	8
			Un	used			TAS	SELx
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
7	6		5	4	3	2	1	0
	IDx		N	1Cx	Unused	TACLR	TAIE	TAIFG
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
Unused	Bits 15-10	Unus	ed					
<b>TASSEL</b> x	Bits 9-8	Timer	r_A clock source	ce select				
		00	TACLK					
		01	ACLK					
		10	SMCLK					
		11	INCLK (INC specific dat		ific and is often as	signed to the inver	ted TBCLK) (see	the device-
IDx	Bits 7-6	Input	divider. These	bits select the div	ider for the input cl	ock.		
		00	/1					
		01	/2					
		10	/4					
		11	/8					
MCx	Bits 5-4	Mode	control. Setting	g MCx = 00h whe	n Timer_A is not in	use conserves po	wer.	
		00	Stop mode:	the timer is halted	d.			
		01	Up mode: t	he timer counts up	to TACCR0.			
		10	Continuous	mode: the timer of	ounts up to 0FFFF	h.		
		11	Up/down m	ode: the timer cou	nts up to TACCR0	then down to 000	0h.	
Unused	Bit 3	Unus	ed					
TACLR	Bit 2			ng this bit resets T and is always read	AR, the clock divided as zero.	der, and the count	direction. The TA	CLR bit is
TAIE	Bit 1	Timer	r_A interrupt er	nable. This bit ena	bles the TAIFG inte	errupt request.		
		0	Interrupt dis	sabled				
		1	Interrupt en	abled				
TAIFG	Bit 0	Timer	r_A interrupt fla	ag				
		0	No interrup	t pending				
		1	Interrupt pe	nding				



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## 12.3.2 TAR, Timer\_A Register

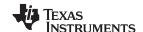
15	14	13	12	11	10	9	8		
	TARx								
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)		
7	6	5	4	3	2	1	0		
			TA	.Rx					
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)		
TARx Bits 15-0 Timer_A register. The TAR register is the count of Timer_A.									

## 12.3.3 TACCRx, Timer\_A Capture/Compare Register x

15	14	13	12	11	10	9	8			
	TACCRX									
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)			
7	6	5	4	3	2	1	0			
			TAC	CRx						
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)			
TACCRx	Bits 15-0 Tir	mer_A capture/con	npare register.							

Compare mode: TACCRx holds the data for the comparison to the timer value in the Timer\_A Register, TAR.

Capture mode: The Timer\_A Register, TAR, is copied into the TACCRx register when a capture is performed.



Timer\_A Registers www.ti.com

## 12.3.4 TACCTLx, Capture/Compare Control Register

15	14		13	12	11	10	9	8
	CMx		CC	SISx	scs	SCCI	Unused	CAP
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	r	r0	rw-(0)
7	6		5	4	3	2	1	0
	OUTMOD	x		CCIE	CCI	OUT	cov	CCIFG
rw-(0)	rw-(0)		rw-(0)	rw-(0)	r	rw-(0)	rw-(0)	rw-(0)
CMx	Bit 15-14	Capture		(0)		(2)	(0)	(0)
CIVIX	DIC 13-14	00	No capture					
		01	Capture on	risina edae				
		10	Capture on					
		11	•	both rising and fa	llina edaes			
CCISx	Bit 13-12			_		CCRx input signal.	See the device-sp	pecific data
				nal connections.			•	
		00	CCIxA					
		01	CCIxB					
		10	GND					
		11	$V_{CC}$					
SCS	Bit 11	Synchro			s used to synchror	nize the capture in	put signal with the	timer clock.
		0	Asynchrono	•				
		1	Synchronou	•				
SCCI	Bit 10		onized captur I via this bit	e/compare input.	The selected CCI	input signal is latc	hed with the EQU	x signal and can
Unused	Bit 9	Unused	d. Read only.	Always read as 0				
CAP	Bit 8	Capture						
		0	Compare m					
		1	Capture mo					
OUTMODx	Bits 7-5	•			e not useful for TA	CCR0, because E	QUx = EQU0.	
		000	OUT bit valu	ie				
		001	Set					
		010	Toggle/rese	i				
		011	Set/reset					
		100	Toggle					
		101	Reset					
		110 111	Toggle/set					
CCIE	Bit 4		Reset/set	orrunt anable. Thi	is hit anables the i	nterrupt request of	the corresponding	a CCIEC floa
CCIE	Dit 4	0	Interrupt dis		s bit enables the i	interrupt request of	the corresponding	y CCIFG flag.
		1	Interrupt dis					
CCI	Bit 3		•		input signal can be	e read by this hit		
OUT	Bit 2				. •	state of the output.		
	DI( 2	0	Output low	iodo o, uno bit dire	only controls the c	nate of the output.		
		1	Output high					
COV	Bit 1		-	is bit indicates a	capture overflow o	ccurred. COV mus	st be reset with so	ftware.
		0		overflow occurred				
		1		rflow occurred				
CCIFG	Bit 0		e/compare int					
-	-	0	No interrupt					
		1	Interrupt per	-				
			-1-1	J				



www.ti.com Timer\_A Registers

## 12.3.5 TAIV, Timer\_A Interrupt Vector Register

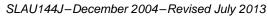
15	14	13	12	11	10	9	8
0	0	0	0	0	0	0	0
r0	r0	r0	r0	rO	rO	rO	r0
7	6	5	4	3	2	1	0
0	0	0	0		TAIVx		0
r0	r0	r0	r0	r-(0)	r-(0)	r-(0)	r0

TAIVx Bits 15-0 Timer\_A interrupt vector value

TAIV Contents	Interrupt Source	Interrupt Flag	Interrupt Priority
00h	No interrupt pending	-	
02h	Capture/compare 1	TACCR1 CCIFG	Highest
04h	Capture/compare 2 <sup>(1)</sup>	TACCR2 CCIFG	
06h	Reserved	-	
08h	Reserved	-	
0Ah	Timer overflow	TAIFG	
0Ch	Reserved	-	
0Eh	Reserved	-	Lowest

<sup>(1)</sup> Not implemented in MSP430x20xx devices

## Chapter 13





# Timer\_B

Timer\_B is a 16-bit timer/counter with multiple capture/compare registers. This chapter describes the operation of the Timer\_B of the MSP430x2xx device family.

**Topic** Page

13.1	Timer_B Introduction	375
13.2	Timer_B Operation	377
13.3	Timer_B Registers	390
	<b>~</b>	



www.ti.com Timer\_B Introduction

## 13.1 Timer\_B Introduction

Timer\_B is a 16-bit timer/counter with three or seven capture/compare registers. Timer\_B can support multiple capture/compares, PWM outputs, and interval timing. Timer\_B also has extensive interrupt capabilities. Interrupts may be generated from the counter on overflow conditions and from each of the capture/compare registers.

Timer\_B features include :

- Asynchronous 16-bit timer/counter with four operating modes and four selectable lengths
- Selectable and configurable clock source
- Three or seven configurable capture/compare registers
- Configurable outputs with PWM capability
- Double-buffered compare latches with synchronized loading
- Interrupt vector register for fast decoding of all Timer\_B interrupts

The block diagram of Timer\_B is shown in Figure 13-1.

#### NOTE: Use of the Word Count

*Count* is used throughout this chapter. It means the counter must be in the process of counting for the action to take place. If a particular value is directly written to the counter, then an associated action does not take place.

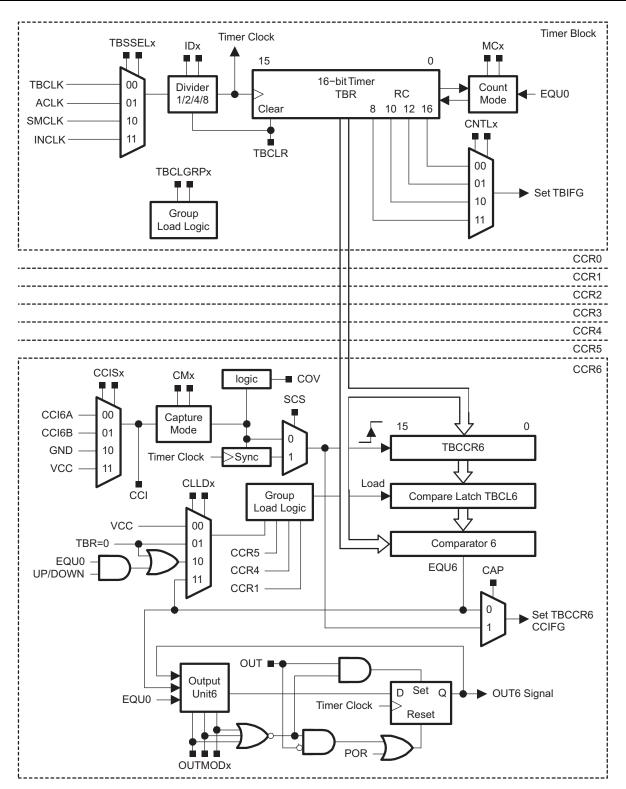
## 13.1.1 Similarities and Differences From Timer A

Timer\_B is identical to Timer\_A with the following exceptions:

- The length of Timer\_B is programmable to be 8, 10, 12, or 16 bits.
- Timer B TBCCRx registers are double-buffered and can be grouped.
- All Timer\_B outputs can be put into a high-impedance state.
- The SCCI bit function is not implemented in Timer\_B.



Timer\_B Introduction www.ti.com



NOTE: INCLK is device-specific, often assigned to the inverted TBCLK, refer to device-specific data sheet.

Figure 13-1. Timer\_B Block Diagram



www.ti.com Timer\_B Operation

#### 13.2 Timer B Operation

The Timer\_B module is configured with user software. The setup and operation of Timer\_B is discussed in the following sections.

#### 13.2.1 16-Bit Timer Counter

The 16-bit timer/counter register, TBR, increments or decrements (depending on mode of operation) with each rising edge of the clock signal. TBR can be read or written with software. Additionally, the timer can generate an interrupt when it overflows.

TBR may be cleared by setting the TBCLR bit. Setting TBCLR also clears the clock divider and count direction for up/down mode.

## NOTE: Modifying Timer\_B Registers

It is recommended to stop the timer before modifying its operation (with exception of the interrupt enable, interrupt flag, and TBCLR) to avoid errant operating conditions.

When the timer clock is asynchronous to the CPU clock, any read from TBR should occur while the timer is not operating or the results may be unpredictable. Alternatively, the timer may be read multiple times while operating, and a majority vote taken in software to determine the correct reading. Any write to TBR will take effect immediately.

## 13.2.1.1 TBR Length

Timer\_B is configurable to operate as an 8-, 10-, 12-, or 16-bit timer with the CNTLx bits. The maximum count value, TBR<sub>(max)</sub>, for the selectable lengths is 0FFh, 03FFh, 0FFFh, and 0FFFFh, respectively. Data written to the TBR register in 8-, 10-, and 12-bit mode is right-justified with leading zeros.

## 13.2.1.2 Clock Source Select and Divider

The timer clock can be sourced from ACLK, SMCLK, or externally via TBCLK or INCLK (INCLK is device-specific, often assigned to the inverted TBCLK, refer to device-specific data sheet). The clock source is selected with the TBSSELx bits. The selected clock source may be passed directly to the timer or divided by 2,4, or 8, using the IDx bits. The clock divider is reset when TBCLR is set.

#### 13.2.2 Starting the Timer

The timer may be started or restarted in the following ways:

- The timer counts when MCx > 0 and the clock source is active.
- When the timer mode is either up or up/down, the timer may be stopped by loading 0 to TBCL0. The
  timer may then be restarted by loading a nonzero value to TBCL0. In this scenario, the timer starts
  incrementing in the up direction from zero.

#### 13.2.3 Timer Mode Control

The timer has four modes of operation as described in Table 13-1: stop, up, continuous, and up/down. The operating mode is selected with the MCx bits.

 MCx
 Mode
 Description

 00
 Stop
 The timer is halted.

 01
 Up
 The timer repeatedly counts from zero to the value of compare register TBCL0.

 10
 Continuous
 The timer repeatedly counts from zero to the value selected by the CNTLx bits.

 11
 Up/down
 The timer repeatedly counts from zero up to the value of TBCL0 and then back down to zero.

Table 13-1. Timer Modes



Timer\_B Operation www.ti.com

#### 13.2.3.1 Up Mode

The up mode is used if the timer period must be different from TBR<sub>(max)</sub> counts. The timer repeatedly counts up to the value of compare latch TBCL0, which defines the period, as shown in Figure 13-2. The number of timer counts in the period is TBCL0+1. When the timer value equals TBCL0 the timer restarts counting from zero. If up mode is selected when the timer value is greater than TBCL0, the timer immediately restarts counting from zero.

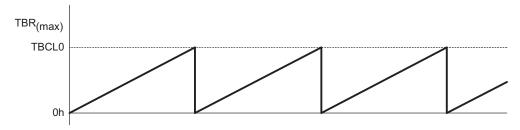


Figure 13-2. Up Mode

The TBCCR0 CCIFG interrupt flag is set when the timer counts to the TBCL0 value. The TBIFG interrupt flag is set when the timer counts from TBCL0 to zero. Figure 13-3 shows the flag set cycle.

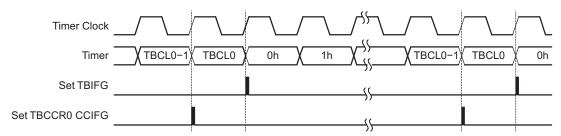


Figure 13-3. Up Mode Flag Setting

## 13.2.3.2 Changing the Period Register TBCL0

When changing TBCL0 while the timer is running and when the TBCL0 load event is immediate, CLLD0 = 00, if the new period is greater than or equal to the old period, or greater than the current count value, the timer counts up to the new period. If the new period is less than the current count value, the timer rolls to zero. However, one additional count may occur before the counter rolls to zero.

#### 13.2.3.3 Continuous Mode

In continuous mode the timer repeatedly counts up to  $\mathsf{TBR}_{(\mathsf{max})}$  and restarts from zero as shown in Figure 13-4. The compare latch TBCL0 works the same way as the other capture/compare registers.

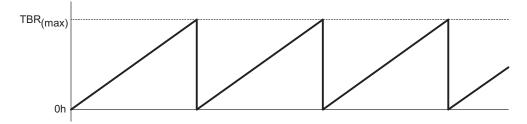


Figure 13-4. Continuous Mode

The TBIFG interrupt flag is set when the timer counts from TBR<sub>(max)</sub> to zero. Figure 13-5 shows the flag set cycle.



www.ti.com Timer\_B Operation

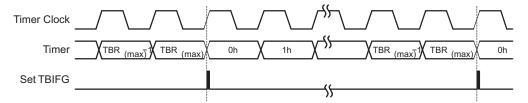


Figure 13-5. Continuous Mode Flag Setting

#### 13.2.3.4 Use of the Continuous Mode

The continuous mode can be used to generate independent time intervals and output frequencies. Each time an interval is completed, an interrupt is generated. The next time interval is added to the TBCLx latch in the interrupt service routine. Figure 13-6 shows two separate time intervals  $t_0$  and  $t_1$  being added to the capture/compare registers. The time interval is controlled by hardware, not software, without impact from interrupt latency. Up to three (Timer\_B3) or 7 (Timer\_B7) independent time intervals or output frequencies can be generated using capture/compare registers.

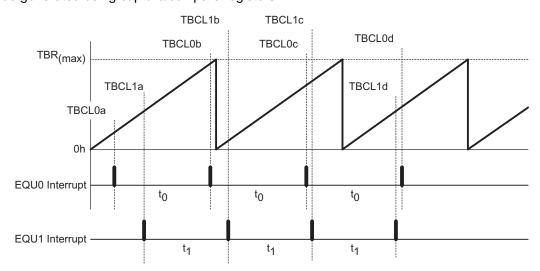


Figure 13-6. Continuous Mode Time Intervals

Time intervals can be produced with other modes as well, where TBCL0 is used as the period register. Their handling is more complex since the sum of the old TBCLx data and the new period can be higher than the TBCL0 value. When the sum of the previous TBCLx value plus  $t_x$  is greater than the TBCL0 data, TBCL0 + 1 must be subtracted to obtain the correct time interval.

#### 13.2.3.5 Up/Down Mode

The up/down mode is used if the timer period must be different from  $TBR_{(max)}$  counts, and if a symmetrical pulse generation is needed. The timer repeatedly counts up to the value of compare latch TBCL0, and back down to zero, as shown in Figure 13-7. The period is twice the value in TBCL0.

## NOTE: TBCL0 > TBR(max)

If TBCL0 > TBR $_{\text{(max)}}$ , the counter operates as if it were configured for continuous mode. It does not count down from TBR $_{\text{(max)}}$  to zero.



Timer B Operation www.ti.com

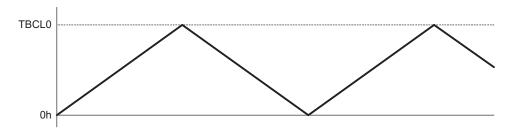


Figure 13-7. Up/Down Mode

The count direction is latched. This allows the timer to be stopped and then restarted in the same direction it was counting before it was stopped. If this is not desired, the TBCLR bit must be used to clear the direction. The TBCLR bit also clears the TBR value and the clock divider.

In up/down mode, the TBCCR0 CCIFG interrupt flag and the TBIFG interrupt flag are set only once during the period, separated by 1/2 the timer period. The TBCCR0 CCIFG interrupt flag is set when the timer *counts* from TBCL0-1 to TBCL0, and TBIFG is set when the timer completes *counting* down from 0001h to 0000h. Figure 13-8 shows the flag set cycle.

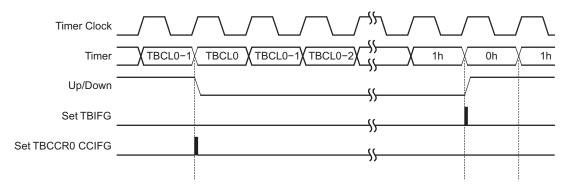


Figure 13-8. Up/Down Mode Flag Setting

## 13.2.3.6 Changing the Value of Period Register TBCL0

When changing TBCL0 while the timer is running, and counting in the down direction, and when the TBCL0 load event is *immediate*, the timer continues its descent until it reaches zero. The value in TBCCR0 is latched into TBCL0 immediately; however, the new period takes effect after the counter counts down to zero.

If the timer is counting in the up direction when the new period is latched into TBCL0, and the new period is greater than or equal to the old period, or greater than the current count value, the timer counts up to the new period before counting down. When the timer is counting in the up direction, and the new period is less than the current count value when TBCL0 is loaded, the timer begins counting down. However, one additional count may occur before the counter begins counting down.

## 13.2.3.7 Use of the Up/Down Mode

The up/down mode supports applications that require dead times between output signals (see section *Timer\_B Output Unit*). For example, to avoid overload conditions, two outputs driving an H-bridge must never be in a high state simultaneously. In the example shown in Figure 13-9 the t<sub>dead</sub> is:

$$t_{dead} = t_{timer} x (TBCL1 - TBCL3)$$

Where.

 $t_{dead}$  = Time during which both outputs need to be inactive

 $t_{timer}$  = Cycle time of the timer clock

TBCLx = Content of compare latch x



www.ti.com Timer\_B Operation

The ability to simultaneously load grouped compare latches assures the dead times.

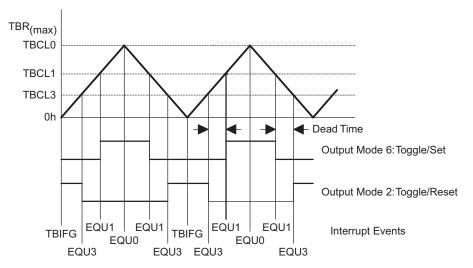


Figure 13-9. Output Unit in Up/Down Mode

## 13.2.4 Capture/Compare Blocks

Three or seven identical capture/compare blocks, TBCCRx, are present in Timer\_B. Any of the blocks may be used to capture the timer data or to generate time intervals.

## 13.2.4.1 Capture Mode

The capture mode is selected when CAP = 1. Capture mode is used to record time events. It can be used for speed computations or time measurements. The capture inputs CCIxA and CCIxB are connected to external pins or internal signals and are selected with the CCISx bits. The CMx bits select the capture edge of the input signal as rising, falling, or both. A capture occurs on the selected edge of the input signal. If a capture is performed:

- The timer value is copied into the TBCCRx register
- · The interrupt flag CCIFG is set

The input signal level can be read at any time via the CCI bit. MSP430x2xx family devices may have different signals connected to CCIxA and CCIxB. Refer to the device-specific data sheet for the connections of these signals.

The capture signal can be asynchronous to the timer clock and cause a race condition. Setting the SCS bit will synchronize the capture with the next timer clock. Setting the SCS bit to synchronize the capture signal with the timer clock is recommended. This is illustrated in Figure 13-10.

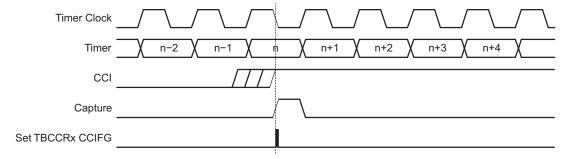


Figure 13-10. Capture Signal (SCS = 1)



Timer B Operation www.ti.com

Overflow logic is provided in each capture/compare register to indicate if a second capture was performed before the value from the first capture was read. Bit COV is set when this occurs as shown in Figure 13-11. COV must be reset with software.

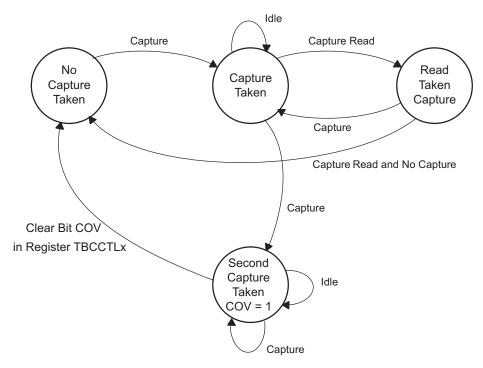


Figure 13-11. Capture Cycle

## 13.2.4.1.1 Capture Initiated by Software

Captures can be initiated by software. The CMx bits can be set for capture on both edges. Software then sets bit CCIS1=1 and toggles bit CCIS0 to switch the capture signal between  $V_{CC}$  and GND, initiating a capture each time CCIS0 changes state:

```
MOV #CAP+SCS+CCIS1+CM_3,&TBCCTLx ; Setup TBCCTLx XOR #CCIS0, &TBCCTLx ; TBCCTLx = TBR
```



www.ti.com Timer\_B Operation

## 13.2.4.2 Compare Mode

The compare mode is selected when CAP = 0. Compare mode is used to generate PWM output signals or interrupts at specific time intervals. When TBR *counts* to the value in a TBCLx:

- Interrupt flag CCIFG is set
- Internal signal EQUx = 1
- · EQUx affects the output according to the output mode

## 13.2.4.2.1 Compare Latch TBCLx

The TBCCRx compare latch, TBCLx, holds the data for the comparison to the timer value in compare mode. TBCLx is buffered by TBCCRx. The buffered compare latch gives the user control over when a compare period updates. The user cannot directly access TBCLx. Compare data is written to each TBCCRx and automatically transferred to TBCLx. The timing of the transfer from TBCCRx to TBCLx is user-selectable with the CLLDx bits as described in Table 13-2.

Table 13-2. TBCLx Load Events

CLLDx	Description
00	New data is transferred from TBCCRx to TBCLx immediately when TBCCRx is written to.
01	New data is transferred from TBCCRx to TBCLx when TBR counts to 0
10	New data is transferred from TBCCRx to TBCLx when TBR <i>counts</i> to 0 for up and continuous modes. New data is transferred to from TBCCRx to TBCLx when TBR <i>counts</i> to the old TBCL0 value or to 0 for up/down mode
11	New data is transferred from TBCCRx to TBCLx when TBR counts to the old TBCLx value.

#### 13.2.4.2.2 Grouping Compare Latches

Multiple compare latches may be grouped together for simultaneous updates with the TBCLGRPx bits. When using groups, the CLLDx bits of the lowest numbered TBCCRx in the group determine the load event for each compare latch of the group, except when TBCLGRP = 3, as shown in Table 13-3. The CLLDx bits of the controlling TBCCRx must not be set to zero. When the CLLDx bits of the controlling TBCCRx are set to zero, all compare latches update immediately when their corresponding TBCCRx is written; no compare latches are grouped.

Two conditions must exist for the compare latches to be loaded when grouped. First, all TBCCRx registers of the group must be updated, even when new TBCCRx data = old TBCCRx data. Second, the load event must occur.

**Table 13-3. Compare Latch Operating Modes** 

TBCLGRPx	Grouping	Update Control
00	None	Individual
01	TBCL1+TBCL2 TBCL3+TBCL4 TBCL5+TBCL6	TBCCR1 TBCCR3 TBCCR5
10	TBCL1+TBCL2+TBCL3 TBCL4+TBCL5+TBCL6	TBCCR1 TBCCR4
11	TBCL0+TBCL1+TBCL2+TBCL3+TBCL4+TBCL5+TBCL6	TBCCR1



Timer\_B Operation www.ti.com

## 13.2.5 Output Unit

Each capture/compare block contains an output unit. The output unit is used to generate output signals such as PWM signals. Each output unit has eight operating modes that generate signals based on the EQU0 and EQUx signals. The TBOUTH pin function can be used to put all Timer\_B outputs into a high-impedance state. When the TBOUTH pin function is selected for the pin, and when the pin is pulled high, all Timer\_B outputs are in a high-impedance state.

## 13.2.5.1 Output Modes

The output modes are defined by the OUTMODx bits and are described in Table 13-4. The OUTx signal is changed with the rising edge of the timer clock for all modes except mode 0. Output modes 2, 3, 6, and 7 are not useful for output unit 0 because EQUx = EQU0.

Table 13-4. Output Modes

OUTMODx	Mode	Description
000	Output	The output signal OUTx is defined by the OUTx bit. The OUTx signal updates immediately when OUTx is updated.
001	Set	The output is set when the timer <i>counts</i> to the TBCLx value. It remains set until a reset of the timer, or until another output mode is selected and affects the output.
010	Toggle/Reset	The output is toggled when the timer <i>counts</i> to the TBCLx value. It is reset when the timer <i>counts</i> to the TBCL0 value.
011	Set/Reset	The output is set when the timer <i>counts</i> to the TBCLx value. It is reset when the timer <i>counts</i> to the TBCL0 value.
100	Toggle	The output is toggled when the timer <i>counts</i> to the TBCLx value. The output period is double the timer period.
101	Reset	The output is reset when the timer <i>counts</i> to the TBCLx value. It remains reset until another output mode is selected and affects the output.
110	Toggle/Set	The output is toggled when the timer <i>counts</i> to the TBCLx value. It is set when the timer <i>counts</i> to the TBCL0 value.
111	Reset/Set	The output is reset when the timer <i>counts</i> to the TBCLx value. It is set when the timer <i>counts</i> to the TBCL0 value.



www.ti.com Timer\_B Operation

## 13.2.5.1.1 Output Example, Timer in Up Mode

The OUTx signal is changed when the timer *counts* up to the TBCLx value, and rolls from TBCL0 to zero, depending on the output mode. An example is shown in Figure 13-12 using TBCL0 and TBCL1.

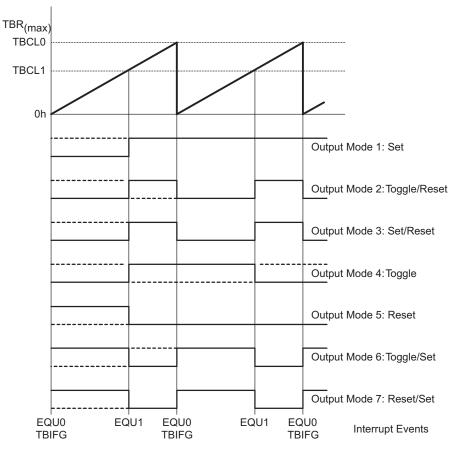


Figure 13-12. Output Example, Timer in Up Mode



Timer\_B Operation www.ti.com

## 13.2.5.1.2 Output Example, Timer in Continuous Mode

The OUTx signal is changed when the timer reaches the TBCLx and TBCL0 values, depending on the output mode, An example is shown in Figure 13-13 using TBCL0 and TBCL1.

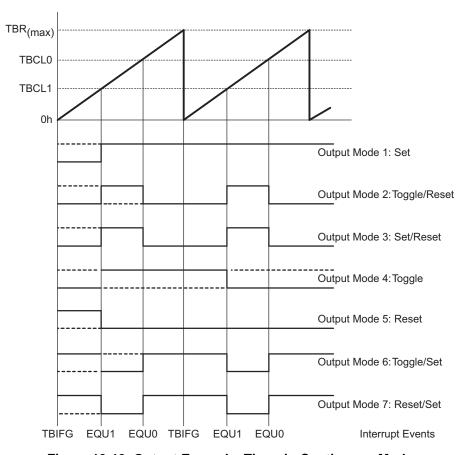


Figure 13-13. Output Example, Timer in Continuous Mode



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## 13.2.5.1.3 Output Example, Timer in Up/Down Mode

The OUTx signal changes when the timer equals TBCLx in either count direction and when the timer equals TBCL0, depending on the output mode. An example is shown in Figure 13-14 using TBCL0 and TBCL3.

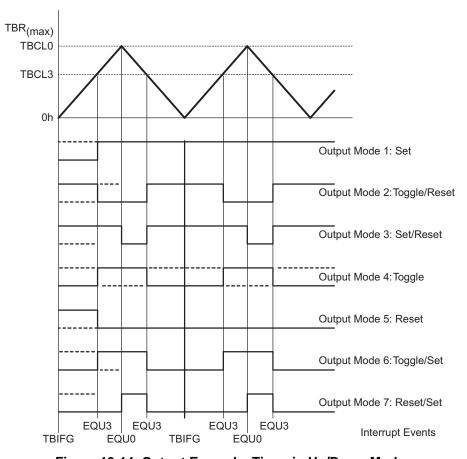


Figure 13-14. Output Example, Timer in Up/Down Mode

## NOTE: Switching Between Output Modes

When switching between output modes, one of the OUTMODx bits should remain set during the transition, unless switching to mode 0. Otherwise, output glitching can occur because a NOR gate decodes output mode 0. A safe method for switching between output modes is to use output mode 7 as a transition state:

```
BIS #OUTMOD_7,&TBCCTLx ; Set output mode=7
BIC #OUTMODx, &TBCCTLx ; Clear unwanted bits
```



Timer B Operation www.ti.com

#### 13.2.6 Timer B Interrupts

Two interrupt vectors are associated with the 16-bit Timer B module:

- TBCCR0 interrupt vector for TBCCR0 CCIFG
- TBIV interrupt vector for all other CCIFG flags and TBIFG

In capture mode, any CCIFG flag is set when a timer value is captured in the associated TBCCRx register. In compare mode, any CCIFG flag is set when TBR *counts* to the associated TBCLx value. Software may also set or clear any CCIFG flag. All CCIFG flags request an interrupt when their corresponding CCIE bit and the GIE bit are set.

## 13.2.6.1 TBCCR0 Interrupt Vector

The TBCCR0 CCIFG flag has the highest Timer\_B interrupt priority and has a dedicated interrupt vector as shown in Figure 13-15. The TBCCR0 CCIFG flag is automatically reset when the TBCCR0 interrupt request is serviced.

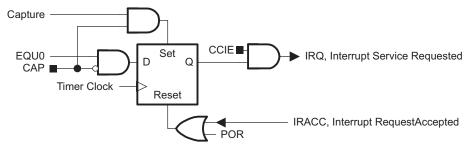


Figure 13-15. Capture/Compare TBCCR0 Interrupt Flag

## 13.2.6.2 TBIV, Interrupt Vector Generator

The TBIFG flag and TBCCRx CCIFG flags (excluding TBCCR0 CCIFG) are prioritized and combined to source a single interrupt vector. The interrupt vector register TBIV is used to determine which flag requested an interrupt.

The highest priority enabled interrupt (excluding TBCCR0 CCIFG) generates a number in the TBIV register (see register description). This number can be evaluated or added to the program counter to automatically enter the appropriate software routine. Disabled Timer\_B interrupts do not affect the TBIV value.

Any access, read or write, of the TBIV register automatically resets the highest pending interrupt flag. If another interrupt flag is set, another interrupt is immediately generated after servicing the initial interrupt. For example, if the TBCCR1 and TBCCR2 CCIFG flags are set when the interrupt service routine accesses the TBIV register, TBCCR1 CCIFG is reset automatically. After the RETI instruction of the interrupt service routine is executed, the TBCCR2 CCIFG flag will generate another interrupt.

## 13.2.6.3 TBIV, Interrupt Handler Examples

The following software example shows the recommended use of TBIV and the handling overhead. The TBIV value is added to the PC to automatically jump to the appropriate routine.

The numbers at the right margin show the necessary CPU clock cycles for each instruction. The software overhead for different interrupt sources includes interrupt latency and return-from-interrupt cycles, but not the task handling itself. The latencies are:

- Capture/compare block CCR0: 11 cycles
- Capture/compare blocks CCR1 to CCR6: 16 cycles
- Timer overflow TBIFG: 14 cycles

Example 13-1 shows the recommended use of TBIV for Timer\_B3.



www.ti.com Timer\_B Operation

## Example 13-1. Recommended Use of TBIV

```
; Interrupt handler for TBCCR0 CCIFG.
                                                     Cycles
CCIFG_0_HND
              ; Start of handler Interrupt latency
                                                        6
                                                        5
  RETI
; Interrupt handler for TBIFG, TBCCR1 and TBCCR2 CCIFG.
TB_HND
                       ; Interrupt latency
           . . .
                        ; Add offset to Jump table
           &TBIV,PC
  ADD
                                                        3
                        ; Vector 0: No interrupt
  RETI
          CCIFG_1_HND ; Vector 2: Module 1
  JMP
                                                        2
           CCIFG_2_HND ; Vector 4: Module 2
  JMP
  RETI
                         ; Vector 6
  RETI
                         ; Vector 8
  RETI
                         ; Vector 10
  RETI
                         ; Vector 12
TBIFG_HND
                         ; Vector 14: TIMOV Flag
                         ; Task starts here
  RETI
                                                        5
CCIFG_2_HND
                         ; Vector 4: Module 2
                         ; Task starts here
  RETI
                         ; Back to main program
; The Module 1 handler shows a way to look if any other
; interrupt is pending: 5 cycles have to be spent, but
; 9 cycles may be saved if another interrupt is pending
CCIFG_1_HND
                         ; Vector 6: Module 3
                         ; Task starts here
  . . .
  JMP
           TB_HND
                         ; Look for pending ints
```



Timer\_B Registers www.ti.com

## 13.3 Timer\_B Registers

The Timer\_B registers are listed in Table 13-5:

## Table 13-5. Timer\_B Registers

Register	Short Form	Register Type	Address	Initial State
Timer_B control	TBCTL	Read/write	0180h	Reset with POR
Timer_B counter	TBR	Read/write	0190h	Reset with POR
Timer_B capture/compare control 0	TBCCTL0	Read/write	0182h	Reset with POR
Timer_B capture/compare 0	TBCCR0	Read/write	0192h	Reset with POR
Timer_B capture/compare control 1	TBCCTL1	Read/write	0184h	Reset with POR
Timer_B capture/compare 1	TBCCR1	Read/write	0194h	Reset with POR
Timer_B capture/compare control 2	TBCCTL2	Read/write	0186h	Reset with POR
Timer_B capture/compare 2	TBCCR2	Read/write	0196h	Reset with POR
Timer_B capture/compare control 3	TBCCTL3	Read/write	0188h	Reset with POR
Timer_B capture/compare 3	TBCCR3	Read/write	0198h	Reset with POR
Timer_B capture/compare control 4	TBCCTL4	Read/write	018Ah	Reset with POR
Timer_B capture/compare 4	TBCCR4	Read/write	019Ah	Reset with POR
Timer_B capture/compare control 5	TBCCTL5	Read/write	018Ch	Reset with POR
Timer_B capture/compare 5	TBCCR5	Read/write	019Ch	Reset with POR
Timer_B capture/compare control 6	TBCCTL6	Read/write	018Eh	Reset with POR
Timer_B capture/compare 6	TBCCR6	Read/write	019Eh	Reset with POR
Timer_B interrupt vector	TBIV	Read only	011Eh	Reset with POR



www.ti.com Timer\_B Registers

## 13.3.1 Timer\_B Control Register TBCTL

15	14		13	12	11	10	9	8
Unused	1	BCLG	RPx	CI	NTLx	Unused	TBS	SELx
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
7	6		5	4	3	2	1	0
	IDx MCx		Unused	TBCLR	TBIE	TBIFG		
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	w-(0)	rw-(0)	rw-(0)
Unused	Bit 15	Unuse	ed					
TBCLGRP	Bit 14-13	TBCL	x group					
		00	Each TBCLx	latch loads inde	pendently			
		01	TBCL3+TBC	L4 (TBCCR3 CL L6 (TBCCR5 CL	LDx bits control th LDx bits control th LDx bits control th	e update)		
		10	TBCL4+TBC TBCL0 indep	L5+TBCL6 (TBC pendent	CCR1 CLLDx bits c CCR4 CLLDx bits c	ontrol the update)		
		11		L1+TBCL2+TBC	CL3+TBCL4+TBCL	5+TBCL6 (TBCCR	1 CLLDx bits cor	trol the update)
CNTLx	Bits 12-11	š						
		00		max) = 0FFFFh				
		01		max) = 0FFFh				
		10 11	8-bit, TBR(m	max) = 03FFh				
Unused	Bit 10	Unuse	•	ax) = 0FF11				
TBSSELx	Bits 9-8		_B clock source	select				
IBOOLLX	Dito o o	00	_B Glock Source TBCLK	, 30,000.				
		01	ACLK					
		10	SMCLK					
		11	INCLK (INCL specific data	•	cific and is often as	ssigned to the inve	rted TBCLK) (see	the device-
IDx	Bits 7-6	Input	divider. These b	oits select the div	vider for the input c	lock.00 /101 /210	/411 /8	
MCx	Bits 5-4	Mode	control. Setting	MCx = 00h  whe	en Timer_B is not in	n use conserves p	ower.	
		00	Stop mode: t	he timer is halte	d			
		01	Up mode: the	e timer counts up	to TBCL0			
		10	Continuous r	mode: the timer o	counts up to the va	lue set by CNTLx		
		11	Up/down mo	de: the timer cou	unts up to TBCL0 a	and down to 0000h	1	
Unused	Bit 3	Unuse	ed					
TBCLR	Bit 2			g this bit resets and is always read	TBR, the clock divided as zero.	der, and the count	direction. The TE	BCLR bit is
TBIE	Bit 1	Timer	_B interrupt ena	able. This bit ena	ables the TBIFG int	errupt request.		
		0	Interrupt disa					
		1	Interrupt ena					
TBIFG	Bit 0		_B interrupt flag					
		0	No interrupt	-				
		1	Interrupt pen	ding				



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## 13.3.2 TBR, Timer\_B Register

15	14	13	12	11	10	9	8		
			TB	Rx					
rw-(0)									
7	6	5	4	3	2	1	0		
	TBRx								
rw-(0)									

TBRx Bits 15-0 Timer\_B register. The TBR register is the count of Timer\_B.

## 13.3.3 TBCCRx, Timer\_B Capture/Compare Register x

15	14	13	12	11	10	9	8			
	TBCCRx									
rw-(0)										
7	6	5	4	3	2	1	0			
	TBCCRx									
rw-(0)										

**TBCCRx** Bits 15-0 Timer\_B capture/compare register.

Compare mode: Compare data is written to each TBCCRx and automatically transferred to TBCLx. TBCLx holds the data for the comparison to the timer value in the Timer\_B Register, TBR.

Capture mode: The Timer\_B Register, TBR, is copied into the TBCCRx register when a capture is performed.



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## 13.3.4 TBCCTLx, Capture/Compare Control Register

15	14		13	12	11	10	9	8			
	CMx		C	CISx	SCS	CL	LDx	CAP			
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	r-(0)	rw-(0)			
7	6		5	4	3	2	1	0			
	OUTMOI	Эx		CCIE	CCI	OUT	cov	CCIFG			
rw-(0)	rw-(0)		rw-(0)	rw-(0)	r	rw-(0)	rw-(0)	rw-(0)			
CMx	Bit 15-14	Conti	ire mode	(-)	•	(5)	(5)	(5)			
CIVIX	DIL 13-14	00	No capture								
		01	Capture on	risina edae							
		10	•	falling edge							
		11	•	both rising and fa	allina edaes						
CCISx	Bit 13-12	Captu	re/compare in	=		CCRx input signal.	See the device-s	pecific data			
		00	CCIxA	,							
		01	CCIxB								
		10	GND								
		11	V <sub>CC</sub>								
SCS	Bit 11		00	source. This bit	is used to synchro	nize the capture in	put signal with the	e timer clock.			
		0	Synchronize capture source. This bit is used to synchronize the capture input signal with the timer clock.  O Asynchronous capture								
		1	Synchronou	is capture							
CLLDx	Bit 10-9	Comp	are latch load.	These bits selec	t the compare latc	h load event.					
		00									
		01 TBCLx loads when TBR counts to 0									
		10			ntsto 0 (up or cont ntsto TBCL0 or to	inuous mode) 0 (up/down mode)					
		11	TBCLx load	ls when TBR coul	ntsto TBCLx						
CAP	Bit 8	Captu	ıre mode								
		0	Compare m								
		1	Capture mo								
OUTMODx	Bits 7-5				re not useful for TE	BCL0 because EQ	Ux = EQU0.				
		000	OUT bit val	ue							
		001	Set								
		010	Toggle/rese	et							
		011	Set/reset								
		100	Toggle								
		101	Reset								
		110 111	Toggle/set Reset/set								
CCIE	Bit 4			terrunt enable. Th	is hit anables the	interrupt request of	f the correspondin	og CCIEG flag			
COIL	Dit 4	0 0	Interrupt dis		iis bit eriables trie	interrupt request o	i trie correspondir	ig Coil G liag.			
		1	Interrupt en								
CCI	Bit 3				input signal can b	e read by this bit.					
OUT	Bit 2				-	state of the output.					
	2.1. 2	0	Output low			olalo ol lilo oulpuli					
		1	Output high	ı							
COV	Bit 1				capture overflow	occurred. COV mu	st be reset with so	oftware.			
		0		overflow occurred							
		1		erflow occurred							
CCIFG	Bit 0		re/compare in								
		0	No interrupt								
		1	Interrupt pe								



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## 13.3.5 TBIV, Timer\_B Interrupt Vector Register

15	14	13	12	11	10	9	8
0	0	0	0	0	0	0	0
r0	r0	r0	r0	r0	r0	rO	r0
7	6	5	4	3	2	1	0
0	0	0	0		TBIVx		0
r0	r0	r0	r0	r-(0)	r-(0)	r-(0)	r0

TBIVx Bits 15-0 Timer\_B interrupt vector value

TBIV Contents	Interrupt Source	Interrupt Flag	Interrupt Priority
00h	No interrupt pending	=	
02h	Capture/compare 1	TBCCR1 CCIFG	Highest
04h	Capture/compare 2	TBCCR2 CCIFG	
06h	Capture/compare 3 <sup>(1)</sup>	TBCCR3 CCIFG	
08h	Capture/compare 4 <sup>(1)</sup>	TBCCR4 CCIFG	
0Ah	Capture/compare 5 <sup>(1)</sup>	TBCCR5 CCIFG	
0Ch	Capture/compare 6 <sup>(1)</sup>	TBCCR6 CCIFG	
0Eh	Timer overflow	TBIFG	Lowest

<sup>(1)</sup> Not available on all devices



# Universal Serial Interface (USI)

The Universal Serial Interface (USI) module provides SPI and I<sup>2</sup>C serial communication with one hardware module. This chapter discusses both modes.

Topic Page

14.1	USI Introduction	396
14.2	USI Operation	399
14.3	USI Registers	405
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USI Introduction www.ti.com

#### 14.1 USI Introduction

The USI module provides the basic functionality to support synchronous serial communication. In its simplest form, it is an 8- or 16-bit shift register that can be used to output data streams, or when combined with minimal software, can implement serial communication. In addition, the USI includes built-in hardware functionality to ease the implementation of SPI and I<sup>2</sup>C communication. The USI module also includes interrupts to further reduce the necessary software overhead for serial communication and to maintain the ultra-low-power capabilities of the MSP430.

The USI module features include:

- Three-wire SPI mode support
- I<sup>2</sup>C mode support
- · Variable data length
- Slave operation in LPM4; no internal clock required
- Selectable MSB or LSB data order
- START and STOP detection for I<sup>2</sup>C mode with automatic SCL control
- Arbitration lost detection in master mode
- Programmable clock generation
- · Selectable clock polarity and phase control

Figure 14-1 shows the USI module in SPI mode. Figure 14-2 shows the USI module in I<sup>2</sup>C mode.



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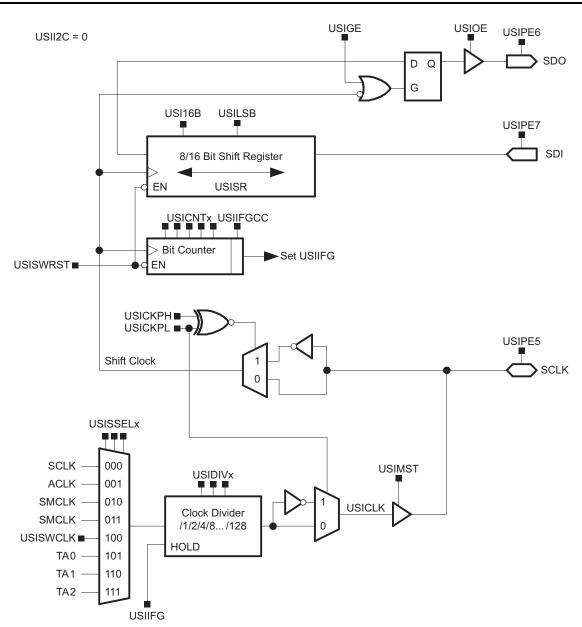


Figure 14-1. USI Block Diagram: SPI Mode



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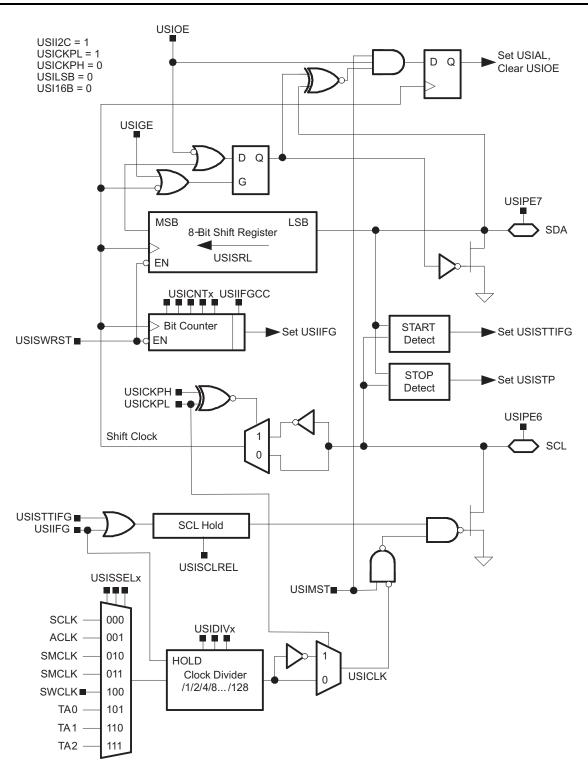


Figure 14-2. USI Block Diagram: I<sup>2</sup>C Mode



www.ti.com USI Operation

### 14.2 USI Operation

The USI module is a shift register and bit counter that includes logic to support SPI and I<sup>2</sup>C communication. The USI shift register (USISR) is directly accessible by software and contains the data to be transmitted or the data that has been received.

The bit counter counts the number of sampled bits and sets the USI interrupt flag USIFG when the USICNTx value becomes zero, either by decrementing or by directly writing zero to the USICNTx bits. Writing USICNTx with a value > 0 automatically clears USIFG when USIFGCC = 0, otherwise USIFG is not affected. The USICNTx bits stop decrementing when they become 0. They will not underflow to 0FFh.

Both the counter and the shift register are driven by the same shift clock. On a rising shift clock edge, USICNTx decrements and USISR samples the next bit input. The latch connected to the shift register's output delays the change of the output to the falling edge of shift clock. It can be made transparent by setting the USIGE bit. This setting will immediately output the MSB or LSB of USISR to the SDO pin, depending on the USILSB bit.

#### 14.2.1 USI Initialization

While the USI software reset bit, USISWRST, is set, the flags USIIFG, USISTTIFG, USISTP, and USIAL will be held in their reset state. USISR and USICNTx are not clocked and their contents are not affected. In I<sup>2</sup>C mode, the SCL line is also released to the idle state by the USI hardware.

To activate USI port functionality the corresponding USIPEx bits in the USI control register must be set. This will select the USI function for the pin and maintains the PxIN and PxIFG functions for the pin as well. With this feature, the port input levels can be read via the PxIN register by software and the incoming data stream can generate port interrupts on data transitions. This is useful, for example, to generate a port interrupt on a START edge.

#### 14.2.2 USI Clock Generation

The USI clock generator contains a clock selection multiplexer, a divider, and the ability to select the clock polarity as shown in the block diagrams Figure 14-1 and Figure 14-2.

The clock source can be selected from the internal clocks ACLK or SMCLK, from an external clock SCLK, as well as from the capture/compare outputs of Timer\_A. In addition, it is possible to clock the module by software using the USISWCLK bit when USISSELx = 100.

The USIDIVx bits can be used to divide the selected clock by a power of 2 up to 128. The generated clock, USICLK, is stopped when USIIFG = 1 or when the module operates in slave mode.

The USICKPL bit is used to select the polarity of USICLK. When USICKPL = 0, the inactive level of USICLK is low. When USICKPL = 1 the inactive level of USICLK is high.



USI Operation www.ti.com

#### 14.2.3 SPI Mode

The USI module is configured in SPI mode when USII2C = 0. Control bit USICKPL selects the inactive level of the SPI clock while USICKPH selects the clock edge on which SDO is updated and SDI is sampled. Figure 14-3 shows the clock/data relationship for an 8-bit, MSB-first transfer. USIPE5, USIPE6, and USIPE7 must be set to enable the SCLK, SDO, and SDI port functions.

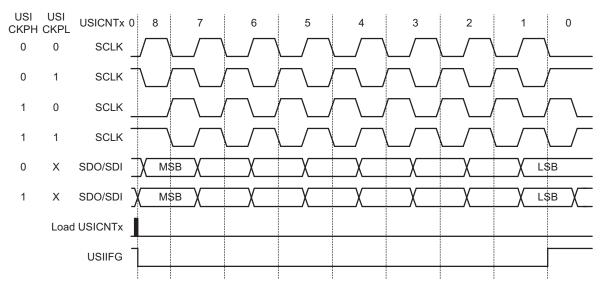


Figure 14-3. SPI Timing

#### 14.2.3.1 SPI Master Mode

The USI module is configured as SPI master by setting the master bit USIMST and clearing the  $I^2C$  bit USII2C. Since the master provides the clock to the slave(s) an appropriate clock source needs to be selected and SCLK configured as output. When USIPE5 = 1, SCLK is automatically configured as an output.

When USIIFG = 0 and USICNTx > 0, clock generation is enabled and the master will begin clocking in/out data using USISR.

Received data must be read from the shift register before new data is written into it for transmission. In a typical application, the USI software will read received data from USISR, write new data to be transmitted to USISR, and enable the module for the next transfer by writing the number of bits to be transferred to USICNTx.

#### 14.2.3.2 SPI Slave Mode

The USI module is configured as SPI slave by clearing the USIMST and the USII2C bits. In this mode, when USIPE5 = 1 SCLK is automatically configured as an input and the USI receives the clock externally from the master.

If the USI is to transmit data, the shift register must be loaded with the data before the master provides the first clock edge. The output must be enabled by setting USIOE. When USICKPH = 1, the MSB will be visible on SDO immediately after loading the shift register.

The SDO pin can be disabled by clearing the USIOE bit. This is useful if the slave is not addressed in an environment with multiple slaves on the bus.

Once all bits are received, the data must be read from USISR and new data loaded into USISR before the next clock edge from the master. In a typical application, after receiving data, the USI software will read the USISR register, write new data to USISR to be transmitted, and enable the USI module for the next transfer by writing the number of bits to be transferred to USICNTx.



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#### 14.2.3.3 USISR Operation

The 16-bit USISR is made up of two 8-bit registers, USISRL and USISRH. Control bit USI16B selects the number of bits of USISR that are used for data transmit and receive. When USI16B = 0, only the lower 8 bits, USISRL, are used.

To transfer < 8 bits, the data must be loaded into USISRL such that unused bits are not shifted out. The data must be MSB- or LSB-aligned depending on USILSB. Figure 14-4 shows an example of 7-bit data handling.

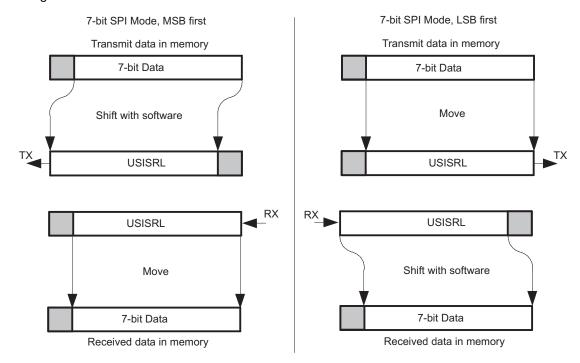


Figure 14-4. Data Adjustments for 7-Bit SPI Data

When USI16B = 1, all 16 bits are used for data handling. When using USISR to access both USISRL and USISRH, the data needs to be properly adjusted when < 16 bits are used in the same manner as shown in Figure 14-4.

### 14.2.3.4 SPI Interrupts

There is one interrupt vector associated with the USI module, and one interrupt flag, USIIFG, relevant for SPI operation. When USIE and the GIE bit are set, the interrupt flag will generate an interrupt request.

USIIFG is set when USICNTx becomes zero, either by counting or by directly writing 0 to the USICNTx bits. USIIFG is cleared by writing a value > 0 to the USICNTx bits when USIIFGCC = 0, or directly by software.



USI Operation www.ti.com

#### 14.2.4 fC Mode

The USI module is configured in I<sup>2</sup>C mode when USII2C =1, USICKPL = 1, and USICKPH = 0. For I<sup>2</sup>C data compatibility, USILSB and USI16B must be cleared. USIPE6 and USIPE7 must be set to enable the SCL and SDA port functions.

#### 14.2.4.1 I<sup>2</sup>C Master Mode

To configure the USI module as an  $I^2C$  master the USIMST bit must be set. In master mode, clocks are generated by the USI module and output to the SCL line while USIIFG = 0. When USIIFG = 1, the SCL will stop at the idle, or high, level. Multi-master operation is supported as described in the Arbitration section.

The master supports slaves that are holding the SCL line low only when USIDIVx > 0. When USIDIVx is set to /1 clock division (USIDIVx = 0), connected slaves must not hold the SCL line low during data transmission. Otherwise the communication may fail.

#### 14.2.4.2 I<sup>2</sup>C Slave Mode

To configure the USI module as an  $I^2C$  slave the USIMST bit must be cleared. In slave mode, SCL is held low if USIFFG = 1, USISTTIFG = 1 or if USICNTx = 0. USISTTIFG must be cleared by software after the slave is setup and ready to receive the slave address from a master.

#### 14.2.4.3 I<sup>2</sup>C Transmitter

In transmitter mode, data is first loaded into USISRL. The output is enabled by setting USIOE and the transmission is started by writing 8 into USICNTx. This clears USIIFG and SCL is generated in master mode or released from being held low in slave mode. After the transmission of all 8 bits, USIIFG is set, and the clock signal on SCL is stopped in master mode or held low at the next low phase in slave mode.

To receive the I<sup>2</sup>C acknowledgment bit, the USIOE bit is cleared with software and USICNTx is loaded with 1. This clears USIIFG and one bit is received into USISRL. When USIIFG becomes set again, the LSB of USISRL is the received acknowledge bit and can be tested in software.

```
; Receive ACK/NACK
BIC.B #USIOE,&USICTL0 ; SDA input
MOV.B #01h,&USICNT ; USICNTx = 1
TEST_USIIFG
BIT.B #USIIFG,&USICTL1 ; Test USIIFG
JZ TEST_USIIFG
BIT.B #01h,&USISRL ; Test received ACK bit
JNZ HANDLE_NACK ; Handle if NACK
...Else, handle ACK
```



www.ti.com USI Operation

#### 14.2.4.4 I<sup>2</sup>C Receiver

In I<sup>2</sup>C receiver mode the output must be disabled by clearing USIOE and the USI module is prepared for reception by writing 8 into USICNTx. This clears USIIFG and SCL is generated in master mode or released from being held low in slave mode. The USIIFG bit will be set after 8 clocks. This stops the clock signal on SCL in master mode or holds SCL low at the next low phase in slave mode.

To transmit an acknowledge or no-acknowledge bit, the MSB of the shift register is loaded with 0 or 1, the USIOE bit is set with software to enable the output, and 1 is written to the USICNTx bits. As soon as the MSB bit is shifted out, USIIFG will be become set and the module can be prepared for the reception of the next I<sup>2</sup>C data byte.

```
; Generate ACK
BIS.B #USIOE,&USICTL0 ; SDA output
MOV.B #00h,&USISRL ; MSB = 0
MOV.B #01h,&USICNT ; USICNTx = 1
TEST_USIIFG
BIT.B #USIIFG,&USICTL1 ; Test USIIFG
JZ TEST_USIIFG
...continue...
; Generate NACK
BIS.B #USIOE,&USICTL0 ; SDA output
MOV.B #0FFh,&USISRL ; MSB = 1
MOV.B #01h,&USICNT ; USICNTx = 1
TEST_USIIFG
BIT.B #USIIFG,&USICTL1 ; Test USIIFG
JZ TEST_USIIFG
...continue...
```

#### 14.2.4.5 START Condition

A START condition is a high-to-low transition on SDA while SCL is high. The START condition can be generated by setting the MSB of the shift register to 0. Setting the USIGE and USIOE bits makes the output latch transparent and the MSB of the shift register is immediately presented to SDA and pulls the line low. Clearing USIGE resumes the clocked-latch function and holds the 0 on SDA until data is shifted out with SCL.

```
; Generate START
  MOV.B #000h,&USISRL ; MSB = 0
BIS.B #USIGE+USIOE,&USICTL0 ; Latch/SDA output enabled
BIC.B #USIGE,&USICTL0 ; Latch disabled
...continue...
```

#### 14.2.4.6 STOP Condition

A STOP condition is a low-to-high transition on SDA while SCL is high. To finish the acknowledgment bit and pull SDA low to prepare the STOP condition generation requires clearing the MSB in the shift register and loading 1 into USICNTx. This will generate a low pulse on SCL and during the low phase SDA is pulled low. SCL stops in the idle, or high, state since the module is in master mode. To generate the low-to-high transition, the MSB is set in the shift register and USICNTx is loaded with 1. Setting the USIGE and USIOE bits makes the output latch transparent and the MSB of USISRL releases SDA to the idle state. Clearing USIGE stores the MSB in the output latch and the output is disabled by clearing USIOE. SDA remains high until a START condition is generated because of the external pullup.

```
; Generate STOP
BIS.B #USIOE,&USICTL0 ; SDA=output
MOV.B #000h,&USISRL ; MSB = 0
MOV.B #001h,&USICNT ; USICNT = 1 for one clock
TEST_USIIFG
BIT.B #USIIFG,&USICTL1 ; Test USIIFG
JZ test_USIIFG ;
MOV.B #0FFh,&USISRL ; USISRL = 1 to drive SDA high
BIS.B #USIGE,&USICTL0 ; Transparent latch enabled
BIC.B #USIGE+USIOE,&USICTL; Latch/SDA output disabled
...continue...
```



USI Operation www.ti.com

### 14.2.4.7 Releasing SCL

Setting the USISCLREL bit will release SCL if it is being held low by the USI module without requiring USIIFG to be cleared. The USISCLREL bit will be cleared automatically if a START condition is received and the SCL line will be held low on the next clock.

In slave operation this bit should be used to prevent SCL from being held low when the slave has detected that it was not addressed by the master. On the next START condition USISCLREL will be cleared and the USISTTIFG will be set.

#### 14.2.4.8 Arbitration

The USI module can detect a lost arbitration condition in multi-master I<sup>2</sup>C systems. The I<sup>2</sup>C arbitration procedure uses the data presented on SDA by the competing transmitters. The first master transmitter that generates a logic high loses arbitration to the opposing master generating a logic low. The loss of arbitration is detected in the USI module by comparing the value presented to the bus and the value read from the bus. If the values are not equal arbitration is lost and the arbitration lost flag, USIAL, is set. This also clears the output enable bit USIOE and the USI module no longer drives the bus. In this case, user software must check the USIAL flag together with USIIFG and configure the USI to slave receiver when arbitration is lost. The USIAL flag must be cleared by software.

To prevent other faster masters from generating clocks during the arbitration procedure SCL is held low if another master on the bus drives SCL low and USIIFG or USISTTIFG is set, or if USICNTx = 0.

### 14.2.4.9 I<sup>2</sup>C Interrupts

There is one interrupt vector associated with the USI module with two interrupt flags relevant for I<sup>2</sup>C operation, USIIFG and USISTTIFG. Each interrupt flag has its own interrupt enable bit, USIIE and USISTTIE. When an interrupt is enabled, and the GIE bit is set, a set interrupt flag will generate an interrupt request.

USIIFG is set when USICNTx becomes zero, either by counting or by directly writing 0 to the USICNTx bits. USIIFG is cleared by writing a value > 0 to the USICNTx bits when USIIFGCC = 0, or directly by software.

USISTTIFG is set when a START condition is detected. The USISTTIFG flag must be cleared by software.

The reception of a STOP condition is indicated with the USISTP flag but there is no interrupt function associated with the USISTP flag. USISTP is cleared by writing a value > 0 to the USICNTx bits when USIIFGCC = 0 or directly by software.



www.ti.com USI Registers

# 14.3 USI Registers

The USI registers are listed in Table 14-1.

Table 14-1. USI Registers

Register	Short Form	Register Type	Address	Initial State
USI control register 0	USICTL0	Read/write	078h	01h with PUC
USI control register 1	USICTL1	Read/write	079h	01h with PUC
USI clock control	USICKCTL	Read/write	07Ah	Reset with PUC
USI bit counter	USICNT	Read/write	07Bh	Reset with PUC
USI low byte shift register	USISRL	Read/write	07Ch	Unchanged
USI high byte shift register	USISRH	Read/write	07Dh	Unchanged

The USI registers can be accessed with word instructions as shown in Table 14-2.

### Table 14-2. Word Access to USI Registers

Register	Short Form	High-Byte Register	Low-Byte Register	Address
USI control register	USICTL	USICTL1	USICTL0	078h
USI clock and counter control register	USICCTL	USICNT	USICKCTL	07Ah
USI shift register	USISR	USISRH	USISRL	07Ch



USI Registers www.ti.com

# 14.3.1 USICTL0, USI Control Register 0

7	6	5	4	3	2	1	0
USIPE7	USIPE6	USIPE5	USILSB	USIMST	USIGE	USIOE	USISWRST
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-1
USIPE7	Bit 7	USI SDI/SDA port	enable. Input in SP	I mode, input or op	oen drain output ir	ı l <sup>2</sup> C mode.	
		0 USI funct	ion disabled				
		1 USI funct	ion enabled				
USIPE6	Bit 6	USI SDO/SCL po	t enable. Output in S	SPI mode, input or	open drain outpu	t in I <sup>2</sup> C mode.	
		0 USI funct	ion disabled				
		1 USI funct	ion enabled				
USIPE5	Bit 5	USI SCLK port en	able. Input in SPI sla	ave mode, or I <sup>2</sup> C n	node, output in SF	PI master mode.	
		0 USI funct	ion disabled				
		1 USI funct	ion enabled				
USILSB	Bit 4	LSB first select. T	his bit controls the d	irection of the rece	eive and transmit	shift register.	
		0 MSB first					
		1 LSB first					
USIMST	Bit 3	Master select					
		0 Slave mo	de				
		1 Master m	ode				
USIGE	Bit 2	Output latch contr	ol				
		0 Output la	tch enable depends	on shift clock			
		1 Output la	tch always enabled	and transparent			
USIOE	Bit 1	Data output enabl	е				
		0 Output di	sabled				
		1 Output er	nabled				
USISWRST	Bit 0	USI software rese	t				
		0 USI relea	sed for operation.				
		1 USI logic	held in reset state.				



www.ti.com USI Registers

# 14.3.2 USICTL1, USI Control Register 1

7	6	5	4	3	2	1	0
USICKPH	USII2C	USISTTIE	USIIE	USIAL	USISTP	USISTTIFG	USIIFG
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-1
USICKPH	Bit 7	Clock phase select					
		0 Data is chan	ged on the first S	CLK edge and cap	otured on the follo	wing edge.	
		1 Data is captu	ired on the first S	CLK edge and ch	anged on the follo	wing edge.	
USII2C	Bit 6	I <sup>2</sup> C mode enable					
		0 I <sup>2</sup> C mode dis	abled				
		1 I <sup>2</sup> C mode en	abled				
USISTTIE	Bit 5	START condition inte	rrupt-enable				
		0 Interrupt on S	START condition	disabled			
		1 Interrupt on S	START condition	enabled			
USIIE	Bit 4	USI counter interrupt	enable				
		0 Interrupt disa	abled				
		1 Interrupt ena	bled				
USIAL	Bit 3	Arbitration lost					
		0 No arbitration	n lost condition				
		<ol> <li>Arbitration lo</li> </ol>	st				
USISTP	Bit 2	STOP condition recei USIIFGCC = 0.	ved. USISTP is a	automatically clear	ed if USICNTx is	loaded with a value	: > 0 when
		0 No STOP co	ndition received				
		1 STOP condit	ion received				
USISTTIFG	Bit 1	START condition inte	rrupt flag				
		0 No START c	ondition received	. No interrupt pend	ding.		
		1 START cond	ition received. Int	terrupt pending.			
USIIFG	Bit 0	USI counter interrupt value > 0 when USIIF		ie USICNTx = 0. A	utomatically clear	red if USICNTx is lo	paded with a
		0 No interrupt	pending				
		1 Interrupt pen	ding				



USI Registers www.ti.com

# 14.3.3 USICKCTL, USI Clock Control Register

7	6		5	4	3	2	1	0
	USIDIV	x			USISSELx		USICKPL	USISWCLK
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
USIDIVx	Bits 7-5	Clock	divider select					
		000	Divide by 1					
		001	Divide by 2					
		010	Divide by 4					
		011	Divide by 8					
		100	Divide by 16					
		101	Divide by 32					
		110	Divide by 64					
		111	Divide by 128	1				
USISSELx	Bits 4-2	Clock	source select. N	lot used in slave	mode.			
		000	SCLK (Not us	ed in SPI mode)				
		001	ACLK					
		010	SMCLK					
		011	SMCLK					
		100	USISWCLK b	it				
		101	TACCR0					
		110	TACCR1					
		111	TACCR2 (Re	served on MSP4	30F20xx devices)			
USICKPL	Bit 1	Clock	polarity select					
		0	Inactive state	is low				
		1	Inactive state	is high				
USISWCLK	Bit 0	Softwa	are clock					
		0	Input clock is	low				
		1	Input clock is	high				

# 14.3.4 USICNT, USI Bit Counter Register

7	6		5	4	3	2	1	0			
USISCLREL	USI16B	3	USIIFGCC			USICNTx					
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0			
USISCLREL	Bit 7	SCL r		line is released	from low to idle.	JSISCLREL is clea	ared if a START c	ondition is			
		0	SCL line is he	eld low if USIIFG	is set						
		1	1 SCL line is released								
USI16B	Bit 6	16-bit	shift register en	able							
		0	8-bit shift register mode. Low byte register USISRL is used.								
		1		gister mode. Botl 16 bits simultan		e registers USISR	L and USISRH ar	e used. USISR			
USIIFGCC	Bit 5		nterrupt flag clea NTx is written wi		JSIIFGCC = 1 the	USIIFG will not be	e cleared automat	ically when			
		0	USIIFG autor	natically cleared	on USICNTx upda	ate					
		1	USIIFG is not	cleared automa	tically						
USICNTx	Bits 4-0	USI b	it count. The US	ICNTx bits set th	e number of bits t	o be received or to	ansmitted.				



www.ti.com USI Registers

# 14.3.5 USISRL, USI Low Byte Shift Register



# 14.3.6 USISRH, USI High Byte Shift Register

7	6	5	4	3	2	1	0
			USIS	SRHx			
rw	rw	rw	rw	rw	rw	rw	rw
USISRHx	Bits 7-0	Contents of the USI h	nigh byte shift regi	ster. Ignored whe	n USI16B = 0.		



# Universal Serial Communication Interface, UART Mode

The universal serial communication interface (USCI) supports multiple serial communication modes with one hardware module. This chapter discusses the operation of the asynchronous UART mode.

Topic Page

15.1 USCI Overview	411
15.2 USCI Introduction: UART Mode	411
15.3 USCI Operation: UART Mode	413
15.4 USCI Registers: UART Mode	



www.ti.com USCI Overview

#### 15.1 USCI Overview

The universal serial communication interface (USCI) modules support multiple serial communication modes. Different USCI modules support different modes. Each different USCI module is named with a different letter. For example, USCI\_A is different from USCI\_B, etc. If more than one identical USCI module is implemented on one device, those modules are named with incrementing numbers. For example, if one device has two USCI\_A modules, they are named USCI\_A0 and USCI\_A1. See the device-specific data sheet to determine which USCI modules, if any, are implemented on which devices.

The USCI\_Ax modules support:

- UART mode
- · Pulse shaping for IrDA communications
- · Automatic baud rate detection for LIN communications
- SPI mode

The USCI\_Bx modules support:

- I<sup>2</sup>C mode
- SPI mode

#### 15.2 USCI Introduction: UART Mode

In asynchronous mode, the USCI\_Ax modules connect the MSP430 to an external system via two external pins, UCAxRXD and UCAxTXD. UART mode is selected when the UCSYNC bit is cleared.

UART mode features include:

- 7- or 8-bit data with odd, even, or non-parity
- Independent transmit and receive shift registers
- · Separate transmit and receive buffer registers
- LSB-first or MSB-first data transmit and receive
- Built-in idle-line and address-bit communication protocols for multiprocessor systems
- · Receiver start-edge detection for auto-wake up from LPMx modes
- Programmable baud rate with modulation for fractional baud rate support
- Status flags for error detection and suppression
- · Status flags for address detection
- Independent interrupt capability for receive and transmit

Figure 15-1 shows the USCI\_Ax when configured for UART mode.



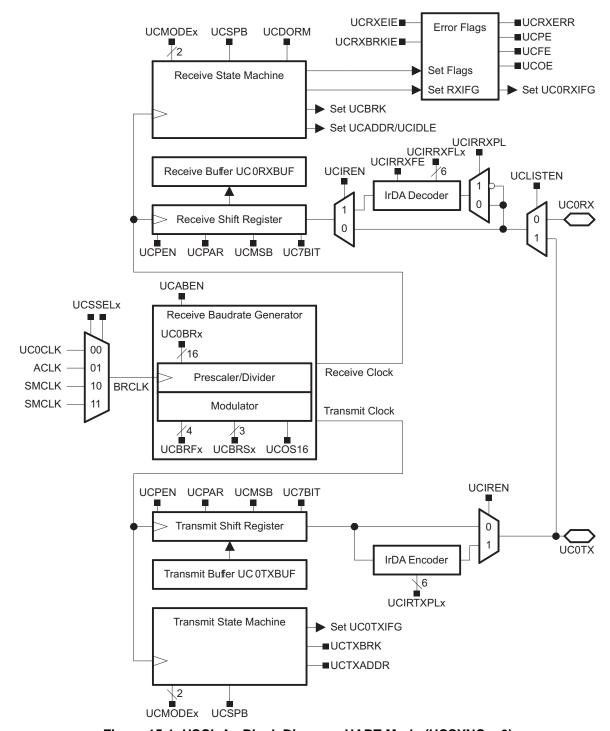


Figure 15-1. USCI\_Ax Block Diagram: UART Mode (UCSYNC = 0)



### 15.3 USCI Operation: UART Mode

In UART mode, the USCI transmits and receives characters at a bit rate asynchronous to another device. Timing for each character is based on the selected baud rate of the USCI. The transmit and receive functions use the same baud rate frequency.

#### 15.3.1 USCI Initialization and Reset

The USCI is reset by a PUC or by setting the UCSWRST bit. After a PUC, the UCSWRST bit is automatically set, keeping the USCI in a reset condition. When set, the UCSWRST bit resets the UCAXRXIE, UCAXRXIFG, UCRXERR, UCBRK, UCPE, UCOE, UCFE, UCSTOE and UCBTOE bits and sets the UCAXTXIFG bit. Clearing UCSWRST releases the USCI for operation.

### NOTE: Initializing or Re-Configuring the USCI Module

The recommended USCI initialization/re-configuration process is:

- 1. Set UCSWRST (BIS.B #UCSWRST,&UCAxCTL1)
- 2. Initialize all USCI registers with UCSWRST = 1 (including UCAxCTL1)
- 3. Configure ports.
- 4. Clear UCSWRST via software (BIC.B #UCSWRST,&UCAxCTL1)
- Enable interrupts (optional) via UCAxRXIE and/or UCAxTXIE

#### 15.3.2 Character Format

The UART character format, shown in Figure 15-2, consists of a start bit, seven or eight data bits, an even/odd/no parity bit, an address bit (address-bit mode), and one or two stop bits. The UCMSB bit controls the direction of the transfer and selects LSB or MSB first. LSB-first is typically required for UART communication.

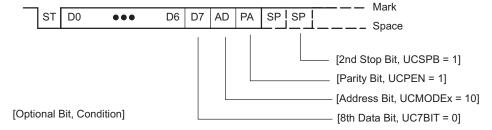


Figure 15-2. Character Format

#### 15.3.3 Asynchronous Communication Formats

When two devices communicate asynchronously, no multiprocessor format is required for the protocol. When three or more devices communicate, the USCI supports the idle-line and address-bit multiprocessor communication formats.

#### 15.3.3.1 Idle-Line Multiprocessor Format

When UCMODEx = 01, the idle-line multiprocessor format is selected. Blocks of data are separated by an idle time on the transmit or receive lines as shown in Figure 15-3. An idle receive line is detectedwhen 10 or more continuous ones (marks) are received after the one or two stop bits of a character. The baud rate generator is switched off after reception of an idle line until the next start edge is detected. When an idle line is detected the UCIDLE bit is set.

The first character received after an idle period is an address character. The UCIDLE bit is used as an address tag for each block of characters. In idle-line multiprocessor format, this bit is set when a received character is an address.

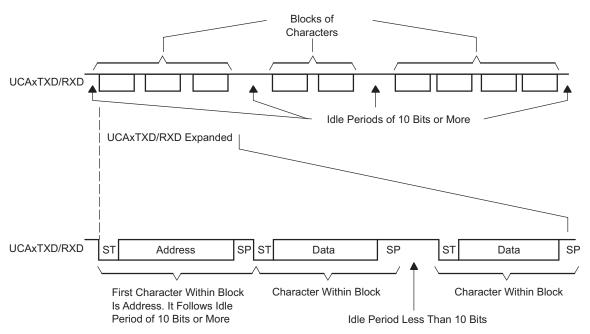


Figure 15-3. Idle-Line Format

The UCDORM bit is used to control data reception in the idle-line multiprocessor format. When UCDORM = 1, all non-address characters are assembled but not transferred into the UCAxRXBUF, and interrupts are not generated. When an address character is received, the character is transferred into UCAxRXBUF, UCAxRXIFG is set, and any applicable error flag is set when UCRXEIE = 1. When UCRXEIE = 0 and an address character is received but has a framing error or parity error, the character is not transferred into UCAxRXBUF and UCAxRXIFG is not set.

If an address is received, user software can validate the address and must reset UCDORM to continue receiving data. If UCDORM remains set, only address characters will be received. When UCDORM is cleared during the reception of a character the receive interrupt flag will be set after the reception completed. The UCDORM bit is not modified by the USCI hardware automatically.

For address transmission in idle-line multiprocessor format, a precise idle period can be generated by the USCI to generate address character identifiers on UCAxTXD. The double-buffered UCTXADDR flag indicates if the next character loaded into UCAxTXBUF is preceded by an idle line of 11 bits. UCTXADDR is automatically cleared when the start bit is generated.

#### 15.3.3.2 Transmitting an Idle Frame

The following procedure sends out an idle frame to indicate an address character followed by associated data:

- Set UCTXADDR, then write the address character to UCAxTXBUF. UCAxTXBUF must be ready for new data (UCAxTXIFG = 1).
  - This generates an idle period of exactly 11 bits followed by the address character. UCTXADDR is reset automatically when the address character is transferred from UCAxTXBUF into the shift register.
- Write desired data characters to UCAxTXBUF. UCAxTXBUF must be ready for new data (UCAxTXIFG = 1).

The data written to UCAxTXBUF is transferred to the shift register and transmitted as soon as the shift register is ready for new data.

The idle-line time must not be exceeded between address and data transmission or between data transmissions. Otherwise, the transmitted data will be misinterpreted as an address.



### 15.3.3.3 Address-Bit Multiprocessor Format

When UCMODEx = 10, the address-bit multiprocessor format is selected. Each processed character contains an extra bit used as an address indicator shown in Figure 15-4. The first character in a block of characters carries a set address bit which indicates that the character is an address. The USCI UCADDR bit is set when a received character has its address bit set and is transferred to UCAxRXBUF.

The UCDORM bit is used to control data reception in the address-bit multiprocessor format. When UCDORM is set, data characters with address bit = 0 are assembled by the receiver but are not transferred to UCAxRXBUF and no interrupts are generated. When a character containing a set address bit is received, the character is transferred into UCAxRXBUF, UCAxRXIFG is set, and any applicable error flag is set when UCRXEIE = 1. When UCRXEIE = 0 and a character containing a set address bit is received, but has a framing error or parity error, the character is not transferred into UCAxRXBUF and UCAxRXIFG is not set.

If an address is received, user software can validate the address and must reset UCDORM to continue receiving data. If UCDORM remains set, only address characters with address bit = 1 will be received. The UCDORM bit is not modified by the USCI hardware automatically.

When UCDORM = 0 all received characters will set the receive interrupt flag UCAxRXIFG. If UCDORM is cleared during the reception of a character the receive interrupt flag will be set after the reception is completed.

For address transmission in address-bit multiprocessor mode, the address bit of a character is controlled by the UCTXADDR bit. The value of the UCTXADDR bit is loaded into the address bit of the character transferred from UCAxTXBUF to the transmit shift register. UCTXADDR is automatically cleared when the start bit is generated.

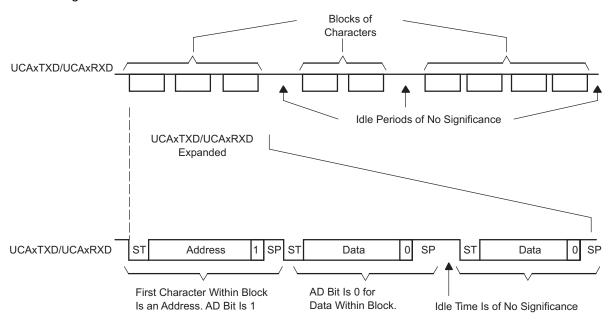


Figure 15-4. Address-Bit Multiprocessor Format

#### 15.3.3.4 Break Reception and Generation

When UCMODEx = 00, 01, or 10 the receiver detects a break when all data, parity, and stop bits are low, regardless of the parity, address mode, or other character settings. When a break is detected, the UCBRK bit is set. If the break interrupt enable bit, UCBRKIE, is set, the receive interrupt flag UCAxRXIFG will also be set. In this case, the value in UCAxRXBUF is 0h since all data bits were zero.

To transmit a break set the UCTXBRK bit, then write 0h to UCAxTXBUF. UCAxTXBUF must be ready for new data (UCAxTXIFG = 1). This generates a break with all bits low. UCTXBRK is automatically cleared when the start bit is generated.



#### 15.3.4 Automatic Baud Rate Detection

When UCMODEx = 11 UART mode with automatic baud rate detection is selected. For automatic baud rate detection, a data frame is preceded by a synchronization sequence that consists of a break and a synch field. A break is detected when 11 or more continuous zeros (spaces) are received. If the length of the break exceeds 22 bit times the break timeout error flag UCBTOE is set. The synch field follows the break as shown in Figure 15-5.



Figure 15-5. Auto Baud Rate Detection - Break/Synch Sequence

For LIN conformance the character format should be set to 8 data bits, LSB first, no parity and one stop bit. No address bit is available.

The synch field consists of the data 055h inside a byte field as shown in Figure 15-6. The synchronization is based on the time measurement between the first falling edge and the last falling edge of the pattern. The transmit baud rate generator is used for the measurement if automatic baud rate detection is enabled by setting UCABDEN. Otherwise, the pattern is received but not measured. The result of the measurement is transferred into the baud rate control registers UCAxBR0, UCAxBR1, and UCAxMCTL. If the length of the synch field exceeds the measurable time the synch timeout error flag UCSTOE is set.

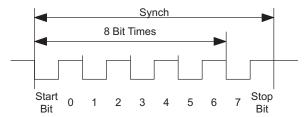


Figure 15-6. Auto Baud Rate Detection - Synch Field

The UCDORM bit is used to control data reception in this mode. When UCDORM is set, all characters are received but not transferred into the UCAxRXBUF, and interrupts are not generated. When a break/synch field is detected the UCBRK flag is set. The character following the break/synch field is transferred into UCAxRXBUF and the UCAxRXIFG interrupt flag is set. Any applicable error flag is also set. If the UCBRKIE bit is set, reception of the break/synch sets the UCAxRXIFG. The UCBRK bit is reset by user software or by reading the receive buffer UCAxRXBUF.

When a break/synch field is received, user software must reset UCDORM to continue receiving data. If UCDORM remains set, only the character after the next reception of a break/synch field will be received. The UCDORM bit is not modified by the USCI hardware automatically.

When UCDORM = 0 all received characters will set the receive interrupt flag UCAxRXIFG. If UCDORM is cleared during the reception of a character the receive interrupt flag will be set after the reception is complete.

The automatic baud rate detection mode can be used in a full-duplex communication system with some restrictions. The USCI can not transmit data while receiving the break/sync field and if a 0h byte with framing error is received any data transmitted during this time gets corrupted. The latter case can be discovered by checking the received data and the UCFE bit.



### 15.3.4.1 Transmitting a Break/Synch Field

The following procedure transmits a break/synch field:

- Set UCTXBRK with UMODEx = 11.
- Write 055h to UCAxTXBUF. UCAxTXBUF must be ready for new data (UCAxTXIFG = 1).

This generates a break field of 13 bits followed by a break delimiter and the synch character. The length of the break delimiter is controlled with the UCDELIMx bits. UCTXBRK is reset automatically when the synch character is transferred from UCAxTXBUF into the shift register.

 Write desired data characters to UCAxTXBUF. UCAxTXBUF must be ready for new data (UCAxTXIFG = 1).

The data written to UCAxTXBUF is transferred to the shift register and transmitted as soon as the shift register is ready for new data.

### 15.3.5 IrDA Encoding and Decoding

When UCIREN is set the IrDA encoder and decoder are enabled and provide hardware bit shaping for IrDA communication.

### 15.3.5.1 IrDA Encoding

The encoder sends a pulse for every zero bit in the transmit bit stream coming from the UART as shown in Figure 15-7. The pulse duration is defined by UCIRTXPLx bits specifying the number of half clock periods of the clock selected by UCIRTXCLK.

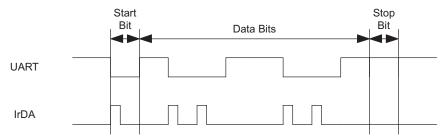


Figure 15-7. UART vs IrDA Data Format

To set the pulse time of 3/16 bit period required by the IrDA standard the BITCLK16 clock is selected with UCIRTXCLK = 1 and the pulse length is set to 6 half clock cycles with UCIRTXPLx = 6 - 1 = 5.

When UCIRTXCLK = 0, the pulse length  $t_{PULSE}$  is based on BRCLK and is calculated as follows:

UCIRTXPLx = 
$$t_{PULSE} \times 2 \times f_{BRCLK} - 1$$

When the pulse length is based on BRCLK the prescaler UCBRx must to be set to a value greater or equal to 5.

#### 15.3.5.2 IrDA Decoding

The decoder detects high pulses when UCIRRXPL = 0. Otherwise it detects low pulses. In addition to the analog deglitch filter an additional programmable digital filter stage can be enabled by setting UCIRRXFE. When UCIRRXFE is set, only pulses longer than the programmed filter length are passed. Shorter pulses are discarded. The equation to program the filter length UCIRRXFLx is:

UCIRRXFLx = 
$$(t_{PULSE} - t_{WAKE}) \times 2 \times f_{BRCLK} - 4$$

Where,

t<sub>PULSE</sub> = Minimum receive pulse width

 $t_{WAKE}$  = Wake time from any low power mode. Zero when MSP430 is in active mode.



#### 15.3.6 Automatic Error Detection

Glitch suppression prevents the USCI from being accidentally started. Any pulse on UCAxRXD shorter than the deglitch time  $t_{\tau}$  (approximately 150 ns) will be ignored. See the device-specific data sheet for parameters.

When a low period on UCAxRXD exceeds  $t_{\tau}$  a majority vote is taken for the start bit. If the majority vote fails to detect a valid start bit the USCI halts character reception and waits for the next low period on UCAxRXD. The majority vote is also used for each bit in a character to prevent bit errors.

The USCI module automatically detects framing errors, parity errors, overrun errors, and break conditions when receiving characters. The bits UCFE, UCPE, UCOE, and UCBRK are set when their respective condition is detected. When the error flags UCFE, UCPE or UCOE are set, UCRXERR is also set. The error conditions are described in Table 15-1.

<b>Error Condition</b>	Error Flag	Description
Framing error	UCFE	A framing error occurs when a low stop bit is detected. When two stop bits are used, both stop bits are checked for framing error. When a framing error is detected, the UCFE bit is set.
Parity error	UCPE	A parity error is a mismatch between the number of 1s in a character and the value of the parity bit. When an address bit is included in the character, it is included in the parity calculation. When a parity error is detected, the UCPE bit is set.
Receive overrun	UCOE	An overrun error occurs when a character is loaded into UCAxRXBUF before the prior character has been read. When an overrun occurs, the UCOE bit is set.
Break condition	UCBRK	When not using automatic baud rate detection, a break is detected when all data, parity, and stop bits are low. When a break condition is detected, the UCBRK bit is set. A break condition can also set the interrupt flag UCAxRXIFG if the break interrupt enable UCBRKIE bit is set.

Table 15-1. Receive Error Conditions

When UCRXEIE = 0 and a framing error, or parity error is detected, no character is received into UCAxRXBUF. When UCRXEIE = 1, characters are received into UCAxRXBUF and any applicable error bit is set.

When UCFE, UCPE, UCOE, UCBRK, or UCRXERR is set, the bit remains set until user software resets it or UCAxRXBUF is read. UCOE must be reset by reading UCAxRXBUF. Otherwise it will not function properly. To detect overflows reliably, the following flow is recommended. After a character is received and UCAxRXIFG is set, first read UCAxSTAT to check the error flags including the overflow flag UCOE. Read UCAxRXBUF next. This will clear all error flags except UCOE, if UCAxRXBUF was overwritten between the read access to UCAxSTAT and to UCAxRXBUF. The UCOE flag should be checked after reading UCAxRXBUF to detect this condition. Note that, in this case, the UCRXERR flag is not set.

#### 15.3.7 USCI Receive Enable

The USCI module is enabled by clearing the UCSWRST bit and the receiver is ready and in an idle state. The receive baud rate generator is in a ready state but is not clocked nor producing any clocks.

The falling edge of the start bit enables the baud rate generator and the UART state machine checks for a valid start bit. If no valid start bit is detected the UART state machine returns to its idle state and the baud rate generator is turned off again. If a valid start bit is detected a character will be received.

When the idle-line multiprocessor mode is selected with UCMODEx = 01 the UART state machine checks for an idle line after receiving a character. If a start bit is detected another character is received. Otherwise the UCIDLE flag is set after 10 ones are received and the UART state machine returns to its idle state and the baud rate generator is turned off.



### 15.3.7.1 Receive Data Glitch Suppression

Glitch suppression prevents the USCI from being accidentally started. Any glitch on UCAxRXD shorter than the deglitch time t<sub>r</sub> (approximately 150 ns) will be ignored by the USCI and further action will be initiated as shown in Figure 15-8. See the device-specific data sheet for parameters.



Figure 15-8. Glitch Suppression, USCI Receive Not Started

When a glitch is longer than  $t_{\tau}$  or a valid start bit occurs on UCAxRXD, the USCI receive operation is started and a majority vote is taken as shown in Figure 15-9. If the majority vote fails to detect a start bit the USCI halts character reception.

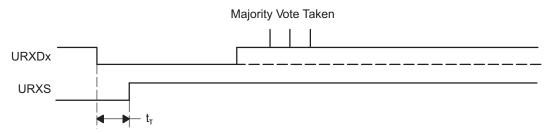


Figure 15-9. Glitch Suppression, USCI Activated

### 15.3.8 USCI Transmit Enable

The USCI module is enabled by clearing the UCSWRST bit and the transmitter is ready and in an idle state. The transmit baud rate generator is ready but is not clocked nor producing any clocks.

A transmission is initiated by writing data to UCAxTXBUF. When this occurs, the baud rate generator is enabled and the data in UCAxTXBUF is moved to the transmit shift register on the next BITCLK after the transmit shift register is empty. UCAxTXIFG is set when new data can be written into UCAxTXBUF.

Transmission continues as long as new data is available in UCAxTXBUF at the end of the previous byte transmission. If new data is not in UCAxTXBUF when the previous byte has transmitted, the transmitter returns to its idle state and the baud rate generator is turned off.

### 15.3.9 UART Baud Rate Generation

The USCI baud rate generator is capable of producing standard baud rates from non-standard source frequencies. It provides two modes of operation selected by the UCOS16 bit.

### 15.3.9.1 Low-Frequency Baud Rate Generation

The low-frequency mode is selected when UCOS16 = 0. This mode allows generation of baud rates from low frequency clock sources (for example, 9600 baud from a 32768-Hz crystal). By using a lower input frequency the power consumption of the module is reduced. Using this mode with higher frequencies and higher prescaler settings will cause the majority votes to be taken in an increasingly smaller window and thus decrease the benefit of the majority vote.

In low-frequency mode the baud rate generator uses one prescaler and one modulator to generate bit clock timing. This combination supports fractional divisors for baud rate generation. In this mode, the maximum USCI baud rate is one-third the UART source clock frequency BRCLK.



Timing for each bit is shown in Figure 15-10. For each bit received, a majority vote is taken to determine the bit value. These samples occur at the N/2 - 1/2, N/2, and N/2 + 1/2 BRCLK periods, where N is the number of BRCLKs per BITCLK.

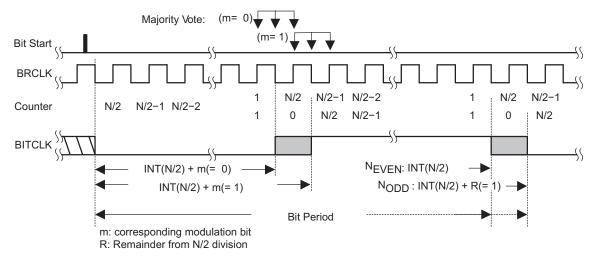


Figure 15-10. BITCLK Baud Rate Timing With UCOS16 = 0

Modulation is based on the UCBRSx setting as shown in Table 15-2. A 1 in the table indicates that m = 1 and the corresponding BITCLK period is one BRCLK period longer than a BITCLK period with m = 0. The modulation wraps around after 8 bits but restarts with each new start bit.

UCBRSx	Bit 0 (Start Bit)	Bit 1	Bit 2	Bit 3	Bit 4	Bit 5	Bit 6	Bit 7
0	0	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0
2	0	1	0	0	0	1	0	0
3	0	1	0	1	0	1	0	0
4	0	1	0	1	0	1	0	1
5	0	1	1	1	0	1	0	1
6	0	1	1	1	0	1	1	1
7	0	1	1	1	1	1	1	1

**Table 15-2. BITCLK Modulation Pattern** 

#### 15.3.9.2 Oversampling Baud Rate Generation

The oversampling mode is selected when UCOS16 = 1. This mode supports sampling a UART bit stream with higher input clock frequencies. This results in majority votes that are always 1/16 of a bit clock period apart. This mode also easily supports IrDA pulses with a 3/16 bit-time when the IrDA encoder and decoder are enabled.

This mode uses one prescaler and one modulator to generate the BITCLK16 clock that is 16 times faster than the BITCLK. An additional divider and modulator stage generates BITCLK from BITCLK16. This combination supports fractional divisions of both BITCLK16 and BITCLK for baud rate generation. In this mode, the maximum USCI baud rate is 1/16 the UART source clock frequency BRCLK. When UCBRx is set to 0 or 1 the first prescaler and modulator stage is bypassed and BRCLK is equal to BITCLK16.

Modulation for BITCLK16 is based on the UCBRFx setting as shown in Table 15-3. A 1 in the table indicates that the corresponding BITCLK16 period is one BRCLK period longer than the periods m=0. The modulation restarts with each new bit timing.

Modulation for BITCLK is based on the UCBRSx setting as shown in Table 15-2 as previously described.



#### Table 15-3. BITCLK16 Modulation Pattern

UCBRFx					No. of	BITCL	(16 Clo	cks Aft	er Last	Falling	BITCL	K Edge	l			
UCBRFX	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
00h	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
01h	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0
02h	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	1
03h	0	1	1	0	0	0	0	0	0	0	0	0	0	0	0	1
04h	0	1	1	0	0	0	0	0	0	0	0	0	0	0	1	1
05h	0	1	1	1	0	0	0	0	0	0	0	0	0	0	1	1
06h	0	1	1	1	0	0	0	0	0	0	0	0	0	1	1	1
07h	0	1	1	1	1	0	0	0	0	0	0	0	0	1	1	1
08h	0	1	1	1	1	0	0	0	0	0	0	0	1	1	1	1
09h	0	1	1	1	1	1	0	0	0	0	0	0	1	1	1	1
0Ah	0	1	1	1	1	1	0	0	0	0	0	1	1	1	1	1
0Bh	0	1	1	1	1	1	1	0	0	0	0	1	1	1	1	1
0Ch	0	1	1	1	1	1	1	0	0	0	1	1	1	1	1	1
0Dh	0	1	1	1	1	1	1	1	0	0	1	1	1	1	1	1
0Eh	0	1	1	1	1	1	1	1	0	1	1	1	1	1	1	1
0Fh	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1

### 15.3.10 Setting a Baud Rate

For a given BRCLK clock source, the baud rate used determines the required division factor N:

$$N = \frac{f_{BRCLK}}{Baud rate}$$

The division factor N is often a non-integer value thus at least one divider and one modulator stage is used to meet the factor as closely as possible.

If N is equal or greater than 16 the oversampling baud rate generation mode can be chosen by setting UCOS16.

#### 15.3.10.1 Low-Frequency Baud Rate Mode Setting

In the low-frequency mode, the integer portion of the divisor is realized by the prescaler:

$$UCBRx = INT(N)$$

and the fractional portion is realized by the modulator with the following nominal formula:

$$UCBRSx = round((N - INT(N)) \times 8)$$

Incrementing or decrementing the UCBRSx setting by one count may give a lower maximum bit error for any given bit. To determine if this is the case, a detailed error calculation must be performed for each bit for each UCBRSx setting.

#### 15.3.10.2 Oversampling Baud Rate Mode Setting

In the oversampling mode the prescaler is set to:

$$UCBRx = INT(\frac{N}{16})$$

and the first stage modulator is set to:

UCBRFx = round ( 
$$(\frac{N}{16} - INT(\frac{N}{16})) \times 16$$
)



When greater accuracy is required, the UCBRSx modulator can also be implemented with values from 0 to 7. To find the setting that gives the lowest maximum bit error rate for any given bit, a detailed error calculation must be performed for all settings of UCBRSx from 0 to 7 with the initial UCBRFx setting and with the UCBRFx setting incremented and decremented by one.

### 15.3.11 Transmit Bit Timing

The timing for each character is the sum of the individual bit timings. Using the modulation features of the baud rate generator reduces the cumulative bit error. The individual bit error can be calculated using the following steps.

#### 15.3.11.1 Low-Frequency Baud Rate Mode Bit Timing

In low-frequency mode, calculate the length of bit i T<sub>bit.TX</sub>[i] based on the UCBRx and UCBRSx settings:

$$T_{bit,TX}[i] = \frac{1}{f_{BRCLK}} (UCBRx + m_{UCBRSx}[i])$$

Where,

m<sub>UCBRSx</sub>[i] = Modulation of bit i from Table 15-2

### 15.3.11.2 Oversampling Baud Rate Mode Bit Timing

In oversampling baud rate mode calculate the length of bit i  $T_{bit,TX}[i]$  based on the baud rate generator UCBRx, UCBRFx and UCBRSx settings:

$$T_{bit,TX}[i] = \frac{1}{f_{BRCLK}} \left( (16 + m_{UCBRSx}[i]) \times UCBRx + \sum_{j=0}^{15} m_{UCBRFx}[j] \right)$$

Where,

$$\sum_{i=0}^{15} m_{\text{UCBRFx}}[j]$$

= Sum of ones from the corresponding row in Table 15-3

 $m_{UCBRSx}[i] = Modulation of bit i from Table 15-2$ 

This results in an end-of-bit time  $t_{bit,TX}[i]$  equal to the sum of all previous and the current bit times:

$$t_{bit,TX}[i] = \sum_{i=0}^{i} T_{bit,TX}[j]$$

To calculate bit error, this time is compared to the ideal bit time t<sub>bit.ideal.TX</sub>[i]:

$$t_{bit,ideal,TX}[i] = \frac{1}{Baud rate} (i + 1)$$

This results in an error normalized to one ideal bit time (1/baudrate):

$$Error_{TX}[i] = (t_{bit,TX}[i] - t_{bit,ideal,TX}[i]) \times Baudrate \times 100\%$$

### 15.3.12 Receive Bit Timing

Receive timing error consists of two error sources. The first is the bit-to-bit timing error similar to the transmit bit timing error. The second is the error between a start edge occurring and the start edge being accepted by the USCI module. Figure 15-11 shows the asynchronous timing errors between data on the UCAxRXD pin and the internal baud-rate clock. This results in an additional synchronization error. The synchronization error  $t_{\text{SYNC}}$  is between -0.5 BRCLKs and +0.5 BRCLKs independent of the selected baud rate generation mode.



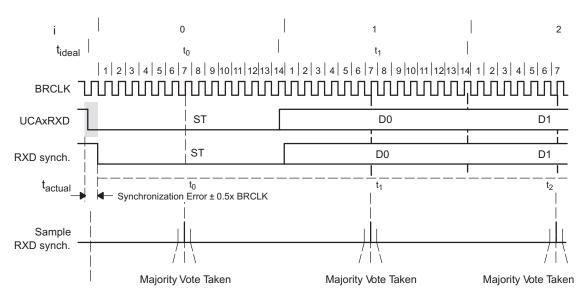


Figure 15-11. Receive Error

The ideal sampling time is in the middle of a bit period:

$$t_{bit,ideal,RX}[i] = \frac{1}{Baud\ rate}$$
 (i + 0.5)

The real sampling time is equal to the sum of all previous bits according to the formulas shown in the transmit timing section, plus one half BITCLK for the current bit i, plus the synchronization error t<sub>SYNC</sub>.

This results in the following for the low-frequency baud rate mode:

$$t_{bit,RX}[i] = t_{SYNC} + \sum_{j=0}^{i-1} T_{bit,RX}[j] + \frac{1}{f_{BRCLK}} \left( INT \left( \frac{1}{2} UCBRx \right) + m_{UCBRSx}[i] \right)$$

Where,

$$T_{bit,RX}[i] = \frac{1}{f_{BRCLK}} \left( UCBRx + m_{UCBRSx}[i] \right)$$

 $m_{UCBRSx}[i] = Modulation of bit i from Table 15-2$ 

For the oversampling baud rate mode the sampling time of bit i is calculated by:

$$t_{bit,RX}[i] = t_{SYNC} + \sum_{j=0}^{i-1} T_{bit,RX}[j] + \frac{1}{f_{BRCLK}} \left( \left(8 + m_{UCBRSx}[i]\right) \times UCBRx + \sum_{j=0}^{7 + m_{UCBRSx}[i]} m_{UCBRFx}[j] \right)$$

Where,

$$T_{bit,RX}[i] = \frac{1}{f_{BRCLK}} \left( (16 + m_{UCBRSx}[i]) \times UCBRx + \sum_{j=0}^{15} m_{UCBRFx}[j] \right)$$

$$\sum_{i=0}^{7+m_{\text{UCBRSx}}[i]} m_{\text{UCBRFx}}[j]$$

= Sum of ones from columns 0 - from the corresponding row in Table 15-3

 $m_{UCBRSx}[i] = Modulation of bit i from Table 15-2$ 

This results in an error normalized to one ideal bit time (1/baudrate) according to the following formula:

$$Error_{RX}[i] = (t_{bit,RX}[i] - t_{bit,ideal,RX}[i]) \times Baudrate \times 100\%$$



### 15.3.13 Typical Baud Rates and Errors

Standard baud rate data for UCBRx, UCBRSx and UCBRFx are listed in Table 15-4 and Table 15-5 for a 32768-Hz crystal sourcing ACLK and typical SMCLK frequencies. Ensure that the selected BRCLK frequency does not exceed the device-specific maximum USCI input frequency (see the device-specific data sheet).

The receive error is the accumulated time versus the ideal scanning time in the middle of each bit. The worst case error is given for the reception of an 8-bit character with parity and one stop bit including synchronization error.

The transmit error is the accumulated timing error versus the ideal time of the bit period. The worst case error is given for the transmission of an 8-bit character with parity and stop bit.

Table 15-4. Commonly Used Baud Rates, Settings, and Errors, UCOS16 = 0

BRCLK Frequency [Hz]	Baud Rate [Baud]	UCBRx	UCBRSx	UCBRFx	Maximum <sup>1</sup>	TX Error [%]	Maximum F	RX Error [%]
32,768	1200	27	2	0	-2.8	1.4	-5.9	2.0
32,768	2400	13	6	0	-4.8	6.0	-9.7	8.3
32,768	4800	6	7	0	-12.1	5.7	-13.4	19.0
32,768	9600	3	3	0	-21.1	15.2	-44.3	21.3
1,048,576	9600	109	2	0	-0.2	0.7	-1.0	0.8
1,048,576	19200	54	5	0	-1.1	1.0	-1.5	2.5
1,048,576	38400	27	2	0	-2.8	1.4	-5.9	2.0
1,048,576	56000	18	6	0	-3.9	1.1	-4.6	5.7
1,048,576	115200	9	1	0	-1.1	10.7	-11.5	11.3
1,048,576	128000	8	1	0	-8.9	7.5	-13.8	14.8
1,048,576	256000	4	1	0	-2.3	25.4	-13.4	38.8
1,000,000	9600	104	1	0	-0.5	0.6	-0.9	1.2
1,000,000	19200	52	0	0	-1.8	0	-2.6	0.9
1,000,000	38400	26	0	0	-1.8	0	-3.6	1.8
1,000,000	56000	17	7	0	-4.8	0.8	-8.0	3.2
1,000,000	115200	8	6	0	-7.8	6.4	-9.7	16.1
1,000,000	128000	7	7	0	-10.4	6.4	-18.0	11.6
1,000,000	256000	3	7	0	-29.6	0	-43.6	5.2
4,000,000	9600	416	6	0	-0.2	0.2	-0.2	0.4
4,000,000	19200	208	3	0	-0.2	0.5	-0.3	0.8
4,000,000	38400	104	1	0	-0.5	0.6	-0.9	1.2
4,000,000	56000	71	4	0	-0.6	1.0	-1.7	1.3
4,000,000	115200	34	6	0	-2.1	0.6	-2.5	3.1
4,000,000	128000	31	2	0	-0.8	1.6	-3.6	2.0
4,000,000	256000	15	5	0	-4.0	3.2	-8.4	5.2
8,000,000	9600	833	2	0	-0.1	0	-0.2	0.1
8,000,000	19200	416	6	0	-0.2	0.2	-0.2	0.4
8,000,000	38400	208	3	0	-0.2	0.5	-0.3	0.8
8,000,000	56000	142	7	0	-0.6	0.1	-0.7	0.8
8,000,000	115200	69	4	0	-0.6	0.8	-1.8	1.1
8,000,000	128000	62	4	0	-0.8	0	-1.2	1.2
8,000,000	256000	31	2	0	-0.8	1.6	-3.6	2.0
12,000,000	9600	1250	0	0	0	0	-0.05	0.05
12,000,000	19200	625	0	0	0	0	-0.2	0
12,000,000	38400	312	4	0	-0.2	0	-0.2	0.2
12,000,000	56000	214	2	0	-0.3	0.2	-0.4	0.5



Table 15-4. Commonly Used Baud Rates, Settings, and Errors, UCOS16 = 0 (continued)

BRCLK Frequency [Hz]	Baud Rate [Baud]	UCBRx	UCBRSx	UCBRFx	Maximum TX Error [%]		Maximum RX Error [%]	
12,000,000	115200	104	1	0	-0.5	0.6	-0.9	1.2
12,000,000	128000	93	6	0	-0.8	0	-1.5	0.4
12,000,000	256000	46	7	0	-1.9	0	-2.0	2.0
16,000,000	9600	1666	6	0	-0.05	0.05	-0.05	0.1
16,000,000	19200	833	2	0	-0.1	0.05	-0.2	0.1
16,000,000	38400	416	6	0	-0.2	0.2	-0.2	0.4
16,000,000	56000	285	6	0	-0.3	0.1	-0.5	0.2
16,000,000	115200	138	7	0	-0.7	0	-0.8	0.6
16,000,000	128000	125	0	0	0	0	-0.8	0
16,000,000	256000	62	4	0	-0.8	0	-1.2	1.2

Table 15-5. Commonly Used Baud Rates, Settings, and Errors, UCOS16 = 1

BRCLK Frequency [Hz]	Baud Rate [Baud]	UCBRx	UCBRSx	UCBRFx	Maximum TX Error [%]		Maximum F	RX Error [%]
1,048,576	9600	6	0	13	-2.3	0	-2.2	0.8
1,048,576	19200	3	1	6	-4.6	3.2	-5.0	4.7
1,000,000	9600	6	0	8	-1.8	0	-2.2	0.4
1,000,000	19200	3	0	4	-1.8	0	-2.6	0.9
1,000,000	57600	1	7	0	-34.4	0	-33.4	0
4,000,000	9600	26	0	1	0	0.9	0	1.1
4,000,000	19200	13	0	0	-1.8	0	-1.9	0.2
4,000,000	38400	6	0	8	-1.8	0	-2.2	0.4
4,000,000	57600	4	5	3	-3.5	3.2	-1.8	6.4
4,000,000	115200	2	3	2	-2.1	4.8	-2.5	7.3
4,000,000	230400	1	7	0	-34.4	0	-33.4	0
8,000,000	9600	52	0	1	-0.4	0	-0.4	0.1
8,000,000	19200	26	0	1	0	0.9	0	1.1
8,000,000	38400	13	0	0	-1.8	0	-1.9	0.2
8,000,000	57600	8	0	11	0	0.88	0	1.6
8,000,000	115200	4	5	3	-3.5	3.2	-1.8	6.4
8,000,000	230400	2	3	2	-2.1	4.8	-2.5	7.3
8,000,000	460800	1	7	0	-34.4	0	-33.4	0
12,000,000	9600	78	0	2	0	0	-0.05	0.05
12,000,000	19200	39	0	1	0	0	0	0.2
12,000,000	38400	19	0	8	-1.8	0	-1.8	0.1
12,000,000	57600	13	0	0	-1.8	0	-1.9	0.2
12,000,000	115200	6	0	8	-1.8	0	-2.2	0.4
12,000,000	230400	3	0	4	-1.8	0	-2.6	0.9
16,000,000	9600	104	0	3	0	0.2	0	0.3
16,000,000	19200	52	0	1	-0.4	0	-0.4	0.1
16,000,000	38400	26	0	1	0	0.9	0	1.1
16,000,000	57600	17	0	6	0	0.9	-0.1	1.0
16,000,000	115200	8	0	11	0	0.9	0	1.6
16,000,000	230400	4	5	3	-3.5	3.2	-1.8	6.4
16,000,000	460800	2	3	2	-2.1	4.8	-2.5	7.3



### 15.3.14 Using the USCI Module in UART Mode with Low Power Modes

The USCI module provides automatic clock activation for SMCLK for use with low-power modes. When SMCLK is the USCI clock source, and is inactive because the device is in a low-power mode, the USCI module automatically activates it when needed, regardless of the control-bit settings for the clock source. The clock remains active until the USCI module returns to its idle condition. After the USCI module returns to the idle condition, control of the clock source reverts to the settings of its control bits. Automatic clock activation is not provided for ACLK.

When the USCI module activates an inactive clock source, the clock source becomes active for the whole device and any peripheral configured to use the clock source may be affected. For example, a timer using SMCLK will increment while the USCI module forces SMCLK active.

### 15.3.15 USCI Interrupts

The USCI has one interrupt vector for transmission and one interrupt vector for reception.

#### 15.3.15.1 USCI Transmit Interrupt Operation

The UCAxTXIFG interrupt flag is set by the transmitter to indicate that UCAxTXBUF is ready to accept another character. An interrupt request is generated if UCAxTXIE and GIE are also set. UCAxTXIFG is automatically reset if a character is written to UCAxTXBUF.

UCAxTXIFG is set after a PUC or when UCSWRST = 1. UCAxTXIE is reset after a PUC or when UCSWRST = 1.

#### 15.3.15.2 USCI Receive Interrupt Operation

The UCAxRXIFG interrupt flag is set each time a character is received and loaded into UCAxRXBUF. An interrupt request is generated if UCAxRXIE and GIE are also set. UCAxRXIFG and UCAxRXIE are reset by a system reset PUC signal or when UCSWRST = 1. UCAxRXIFG is automatically reset when UCAxRXBUF is read.

Additional interrupt control features include:

- When UCAxRXEIE = 0 erroneous characters will not set UCAxRXIFG.
- When UCDORM = 1, non-address characters will not set UCAxRXIFG in multiprocessor modes. In plain UART mode, no characters will set UCAxRXIFG.
- When UCBRKIE = 1 a break condition will set the UCBRK bit and the UCAxRXIFG flag.

#### 15.3.15.3 USCI Interrupt Usage

USCI\_Ax and USCI\_Bx share the same interrupt vectors. The receive interrupt flags UCAxRXIFG and UCBxRXIFG are routed to one interrupt vector, the transmit interrupt flags UCAxTXIFG and UCBxTXIFG share another interrupt vector.

Example 15-1 shows an extract of an interrupt service routine to handle data receive interrupts from USCI\_A0 in either UART or SPI mode and USCI\_B0 in SPI mode.

### Example 15-1. Shared Interrupt Vectors Software Example, Data Receive

```
USCIAO_RX_USCIBO_RX_ISR

BIT.B #UCAORXIFG, &IFG2; USCI_AO Receive Interrupt?

JNZ USCIAO_RX_ISR

USCIBO_RX_ISR?

; Read UCBORXBUF (clears UCBORXIFG)

...

RETI

USCIAO_RX_ISR

; Read UCAORXBUF (clears UCAORXIFG)

...

RETI

RETI
```



Example 15-2 shows an extract of an interrupt service routine to handle data transmit interrupts from USCI\_A0 in either UART or SPI mode and USCI\_B0 in SPI mode.

### Example 15-2. Shared Interrupt Vectors Software Example, Data Transmit

```
USCIAO_TX_USCIBO_TX_ISR
BIT.B #UCAOTXIFG, &IFG2 ; USCI_AO Transmit Interrupt?

JNZ USCIAO_TX_ISR
USCIBO_TX_ISR
; Write UCBOTXBUF (clears UCBOTXIFG)
...
RETI
USCIAO_TX_ISR
; Write UCAOTXBUF (clears UCAOTXIFG)
...
RETI
```

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### 15.4 USCI Registers: UART Mode

The USCI registers applicable in UART mode are listed in Table 15-6 and Table 15-7.

Table 15-6. USCI\_A0 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USCI_A0 control register 0	UCA0CTL0	Read/write	060h	Reset with PUC
USCI_A0 control register 1	UCA0CTL1	Read/write	061h	001h with PUC
USCI_A0 Baud rate control register 0	UCA0BR0	Read/write	062h	Reset with PUC
USCI_A0 baud rate control register 1	UCA0BR1	Read/write	063h	Reset with PUC
USCI_A0 modulation control register	UCA0MCTL	Read/write	064h	Reset with PUC
USCI_A0 status register	UCA0STAT	Read/write	065h	Reset with PUC
USCI_A0 receive buffer register	UCA0RXBUF	Read	066h	Reset with PUC
USCI_A0 transmit buffer register	UCA0TXBUF	Read/write	067h	Reset with PUC
USCI_A0 Auto baud control register	UCA0ABCTL	Read/write	05Dh	Reset with PUC
USCI_A0 IrDA transmit control register	UCA0IRTCTL	Read/write	05Eh	Reset with PUC
USCI_A0 IrDA receive control register	UCA0IRRCTL	Read/write	05Fh	Reset with PUC
SFR interrupt enable register 2	IE2	Read/write	001h	Reset with PUC
SFR interrupt flag register 2	IFG2	Read/write	003h	00Ah with PUC

### NOTE: Modifying SFR bits

To avoid modifying control bits of other modules, it is recommended to set or clear the IEx and IFGx bits using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.

Table 15-7. USCI\_A1 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USCI_A1 control register 0	UCA1CTL0	Read/write	0D0h	Reset with PUC
USCI_A1 control register 1	UCA1CTL1	Read/write	0D1h	001h with PUC
USCI_A1 baud rate control register 0	UCA1BR0	Read/write	0D2h	Reset with PUC
USCI_A1 baud rate control register 1	UCA1BR1	Read/write	0D3h	Reset with PUC
USCI_A1 modulation control register	UCA1MCTL	Read/write	0D4h	Reset with PUC
USCI_A1 status register	UCA1STAT	Read/write	0D5h	Reset with PUC
USCI_A1 receive buffer register	UCA1RXBUF	Read	0D6h	Reset with PUC
USCI_A1 transmit buffer register	UCA1TXBUF	Read/write	0D7h	Reset with PUC
USCI_A1 auto baud control register	UCA1ABCTL	Read/write	0CDh	Reset with PUC
USCI_A1 IrDA transmit control register	UCA1IRTCTL	Read/write	0CEh	Reset with PUC
USCI_A1 IrDA receive control register	UCA1IRRCTL	Read/write	0CFh	Reset with PUC
USCI_A1/B1 interrupt enable register	UC1IE	Read/write	006h	Reset with PUC
USCI_A1/B1 interrupt flag register	UC1IFG	Read/write	007h	00Ah with PUC



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# 15.4.1 UCAxCTL0, USCI\_Ax Control Register 0

7	6		5	4	3	2	1	0
UCPEN	UCPA	₹	UCMSB	UC7BIT	UCSPB	UCM	ODEx	UCSYNC
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
UCPEN	Bit 7	Parity	y enable					
		0	Parity disabl	ed.				
		1			enerated (UCAxTXI ress bit is included			ddress-bit
UCPAR	Bit 6	Parity	y select. UCPAF	R is not used whe	n parity is disabled			
		0	Odd parity					
		1	Even parity					
UCMSB	Bit 5	MSB	first select. Con	trols the direction	of the receive and	l transmit shift re	gister.	
		0	LSB first					
		1	MSB first					
UC7BIT	Bit 4	Char	acter length. Se	lects 7-bit or 8-bit	character length.			
		0	8-bit data					
		1	7-bit data					
UCSPB	Bit 3	Stop	bit select. Numb	per of stop bits.				
		0	One stop bit					
		1	Two stop bit	s				
UCMODEx	Bits 2-1	USCI	I mode. The UC	MODEx bits selec	ct the asynchronou	s mode when UC	SYNC = 0.	
		00	UART mode	!				
		01	Idle-line mul	tiprocessor mode				
		10	Address-bit	multiprocessor mo	ode			
		11	UART mode	with automatic b	aud rate detection			
UCSYNC	Bit 0	Sync	hronous mode e	enable				
		0	Asynchrono	us mode				
		1	Synchronous	s mode				



# 15.4.2 UCAxCTL1, USCI\_Ax Control Register 1

7	6		5	4	3	2	1	0
UC	SSELx		UCRXEIE	UCBRKIE	UCDORM	UCTXADDR	UCTXBRK	UCSWRST
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-1
UCSSELx	Bits 7-6	USCI	clock source se	elect. These bits s	elect the BRCLK	source clock.		
		00	UCLK					
		01	ACLK					
		10	SMCLK					
		11	SMCLK					
UCRXEIE	Bit 5	Rece	ive erroneous-cl	haracter interrupt-	enable			
		0	Erroneous cl	naracters rejected	and UCAxRXIFG	is not set		
		1	Erroneous cl	naracters received	will set UCAxRX	IFG		
UCBRKIE	Bit 4	Rece	ive break chara	cter interrupt-enab	ole			
		0	Received bre	eak characters do	not set UCAxRX	FG.		
		1	Received bre	eak characters set	UCAxRXIFG.			
UCDORM	Bit 3	Dorm	ant. Puts USCI	into sleep mode.				
		0	Not dormant	. All received char	acters will set UC	AxRXIFG.		
		1	UCAxRXIFG		vith automatic ba	an idle-line or with ud rate detection o		
UCTXADDR	Bit 2		smit address. Ne processor mode		nsmitted will be n	narked as address	depending on the	e selected
		0	Next frame t	ransmitted is data				
		1	Next frame t	ransmitted is an a	ddress			
UCTXBRK	Bit 1	baud	rate detection 0		en into UCAxTXB	the transmit buffer UF to generate the		
		0	Next frame t	ransmitted is not a	a break			
		1	Next frame t	ransmitted is a bre	eak or a break/sy	nch		
UCSWRST	Bit 0	Softw	are reset enable	е				
		0	Disabled. US	CI reset released	for operation.			
		1	Enabled. US	CI logic held in re	set state.			

# 15.4.3 UCAxBR0, USCI\_Ax Baud Rate Control Register 0

7	6	5	4	3	2	1	0				
	UCBRx										
rw	rw	rw	rw	rw	rw	rw	rw				

# 15.4.4 UCAxBR1, USCI\_Ax Baud Rate Control Register 1

7	6	5	4	3	2	1	0
			UC	BRx			
rw	rw	rw	rw	rw	rw	rw	rw
UCBRx		ck prescaler settir	ng of the Baud rat	e generator. The 1	6-bit value of (UC	AxBR0 + UCAxB	R1 × 256) forms

www.ti.com USCI Registers: UART Mode

# 15.4.5 UCAxMCTL, USCI\_Ax Modulation Control Register

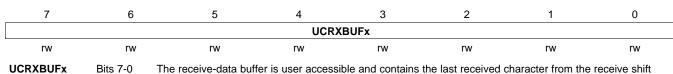
7	6	5	4	3	2	1	0
		UCBRFx			UCOS16		
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
UCBRFx	Bits 7-4	First modulation stag 1. Ignored with UCO	,			for BITCLK16 w	hen UCOS16 =
UCBRSx	Bits 3-1	Second modulation s the modulation patte	0	se bits determine t	the modulation patt	ern for BITCLK.	Table 15-2 shows
UCOS16	Bit 0	Oversampling mode	enabled				
		0 Disabled					
		1 Enabled					

# 15.4.6 UCAxSTAT, USCI\_Ax Status Register

7	6		5	4	3	2	1	0
UCLISTEN	UCFE		UCOE	UCPE	UCBRK	UCRXERR	UCADDR UCIDLE	UCBUSY
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	r-0
UCLISTEN	Bit 7	Listen e	enable. The U	CLISTEN bit sele	cts loopback mod	le.		
		0	Disabled					
		1	Enabled. UC	AxTXD is interna	lly fed back to the	receiver.		
UCFE	Bit 6	Framing	g error flag					
		0	No error					
		1	Character re	ceived with low s	top bit			
UCOE	Bit 5	charact	er was read. l		automatically whe	nsferred into UCA n UCxRXBUF is re		
		0	No error					
		1	Overrun erro	r occurred				
UCPE	Bit 4	Parity e	error flag. Whe	n UCPEN = $0$ , U	CPE is read as 0.			
		0	No error					
		1	Character re	ceived with parity	error			
UCBRK	Bit 3		letect flag					
		0	No break cor					
		1	Break conditi					
UCRXERR	Bit 2					ceived with error(s RXERR is cleared		
		0	No receive e	rrors detected				
		1	Receive erro	r detected				
UCADDR	Bit 1	Address	s received in a	address-bit multip	rocessor mode.			
		0	Received cha	aracter is data				
		1	Received cha	aracter is an addı	ress			
UCIDLE		Idle line	detected in id	dle-line multiproce	essor mode.			
		0	No idle line o	letected				
		1	Idle line dete	cted				
UCBUSY	Bit 0	USCI b	usy. This bit ir	ndicates if a trans	mit or receive ope	eration is in progre	SS.	
		0	USCI inactive	е				
		1	USCI transm	itting or receiving	1			



### 15.4.7 UCAxRXBUF, USCI\_Ax Receive Buffer Register



register. Reading UCAxRXBUF resets the receive-error bits, the UCADDR or UCIDLE bit, and UCAxRXIFG. In 7-bit data mode, UCAxRXBUF is LSB justified and the MSB is always reset.

### 15.4.8 UCAxTXBUF, USCI\_Ax Transmit Buffer Register

7	6	5	4	3	2	1	0			
			UCTX	(BUFx						
rw	rw	rw	rw	rw	rw	rw	rw			
UCTXBUFx	Bits 7-0	ts 7-0 The transmit data buffer is user accessible and holds the data waiting to be moved into the transmit shift register and transmitted on UCAxTXD. Writing to the transmit data buffer clears UCAxTXIFG. The MSB of UCAxTXBUF is not used for 7-bit data and is reset.								

### 15.4.9 UCAxIRTCTL, USCI\_Ax IrDA Transmit Control Register

7	6	5	4	3	2	1	0
		UCIR	TXPLx			UCIRTXCLK	UCIREN
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
UCIRTXPLx	Bits 7-2	Transmit pulse lengt	h. Pulse length t <sub>Pl</sub>	JLSE = (UCIRTXPL)	x + 1) / (2 × f <sub>IRTX</sub>	clk)	
UCIRTXCLK	Bit 1	IrDA transmit pulse of	clock select				
		0 BRCLK					
		1 BITCLK16 v	when UCOS16 = 1	. Otherwise, BRC	LK		
UCIREN	Bit 0	IrDA encoder/decode	er enable.				
		0 IrDA encode	er/decoder disable	ed			
		1 IrDA encode	er/decoder enable	d			

### 15.4.10 UCAxIRRCTL, USCI\_Ax IrDA Receive Control Register

7	6	5	4	3	2	1	0
UCIRRXFLx						UCIRRXPL	UCIRRXFE
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
UCIRRXFLx	Bits 7-2	Receive filter length. The minimum pulse length for receive is given by: $t_{MIN} = (UCIRRXFLx + 4) / (2 \times f_{IRTXCLK})$					
UCIRRXPL	Bit 1	IrDA receive input UCAxRXD polarity  IrDA transceiver delivers a high pulse when a light pulse is seen  IrDA transceiver delivers a low pulse when a light pulse is seen					
UCIRRXFE	Bit 0	IrDA receive filter enabled  Receive filter disabled  Receive filter enabled					



www.ti.com USCI Registers: UART Mode

## 15.4.11 UCAxABCTL, USCI\_Ax Auto Baud Rate Control Register

7	6		5	4	3	2	1	0
Res	served		UCDE	LIMx	UCSTOE	UCBTOE	Reserved	UCABDEN
r-0	r-0		rw-0	rw-0	rw-0	rw-0	r-0	rw-0
Reserved	Bits 7-6	Rese	rved					
UCDELIMx	Bits 5-4	Breal	k/synch delimite	length				
		00	1 bit time					
		01	2 bit times					
		10	3 bit times					
		11	4 bit times					
UCSTOE	Bit 3	Sync	h field time out e	error				
		0	No error					
		1	Length of syr	nch field exceede	ed measurable time			
UCBTOE	Bit 2	Breal	k time out error					
		0	No error					
		1	Length of bre	eak field exceede	ed 22 bit times.			
Reserved	Bit 1	Rese	rved					
UCABDEN	Bit 0	Autor	matic baud rate	detect enable				
		0	Baud rate de	tection disabled.	Length of break ar	nd synch field is no	ot measured.	
		1	Baud rate de are changed		Length of break an	d synch field is me	easured and baud	d rate settings

## 15.4.12 IE2, Interrupt Enable Register 2

7	6	5	4	3	2	1	0
						UCA0TXIE	UCA0RXIE
						rw-0	rw-0
	Bits 7-2	These bits may be u	sed by other modu	ules (see the devic	ce-specific data sh	eet).	
<b>UCA0TXIE</b>	Bit 1	USCI_A0 transmit in	terrupt enable				
		0 Interrupt dis	abled				
		1 Interrupt en	abled				
<b>UCA0RXIE</b>	Bit 0	USCI_A0 receive in	errupt enable				
		0 Interrupt dis	abled				
		1 Interrupt en	abled				

## 15.4.13 IFG2, Interrupt Flag Register 2

7	6	5	4	3	2	1	0	
						UCA0TXIFG	UCA0RXIFG	
						rw-1	rw-0	
	Bits 7-2	These bits may	be used by other mod	ules (see the devic	ce-specific data sh	eet).		
<b>UCA0TXIFG</b>	Bit 1	USCI_A0 trans	USCI_A0 transmit interrupt flag. UCA0TXIFG is set when UCA0TXBUF is empty.					
		0 No into	errupt pending					
		1 Interru	pt pending					
UCA0RXIFG	Bit 0	USCI_A0 recei	ve interrupt flag. UCA0	RXIFG is set when	UCA0RXBUF ha	s received a comp	olete character.	
		0 No inte	errupt pending					
		1 Interru	pt pending					



## 15.4.14 UC1IE, USCI\_A1 Interrupt Enable Register

7	6	5	4	3	2	1	0		
		Unused				UCA1TXIE	UCA1RXIE		
rw-0	rw-0	rw-0	rw-0			rw-0	rw-0		
Unused	Bits 7-4	Unused							
	Bits 3-2	These bits may be us	ese bits may be used by other USCI modules (see the device-specific data sheet).						
<b>UCA1TXIE</b>	Bit 1	USCI_A1 transmit int	errupt enable						
		0 Interrupt disa	abled						
		1 Interrupt ena	bled						
<b>UCA1RXIE</b>	Bit 0	USCI_A1 receive inte	errupt enable						
		0 Interrupt disa	abled						
		1 Interrupt ena	bled						

## 15.4.15 UC1IFG, USCI\_A1 Interrupt Flag Register

7	6	5	4	3	2	1	0	
		Unused				UCA1TXIFG	UCA1RXIFG	
rw-0	rw-0	rw-0	rw-0			rw-1	rw-0	
Unused	Bits 7-4	Unused						
	Bits 3-2	These bits may be us	ed by other USC	I modules (see the	e device-specific d	lata sheet).		
<b>UCA1TXIFG</b>	Bit 1	USCI_A1 transmit int	SCI_A1 transmit interrupt flag. UCA1TXIFG is set when UCA1TXBUF is empty.					
		0 No interrupt	pending					
		1 Interrupt pen	ding					
<b>UCA1RXIFG</b>	Bit 0	USCI_A1 receive inte	errupt flag. UCA1I	RXIFG is set wher	UCA1RXBUF ha	s received a comp	olete character.	
		0 No interrupt	pending					
		1 Interrupt pen	ding					



## Universal Serial Communication Interface, SPI Mode

The universal serial communication interface (USCI) supports multiple serial communication modes with one hardware module. This chapter discusses the operation of the synchronous peripheral interface or SPI mode.

Topic Page

16.1	USCI Overview	436
16.2	USCI Introduction: SPI Mode	436
16.3	USCI Operation: SPI Mode	438
16.4	USCI Registers: SPI Mode	444



USCI Overview www.ti.com

### 16.1 USCI Overview

The universal serial communication interface (USCI) modules support multiple serial communication modes. Different USCI modules support different modes. Each different USCI module is named with a different letter (for example, USCI\_A is different from USCI\_B). If more than one identical USCI module is implemented on one device, those modules are named with incrementing numbers. For example, if one device has two USCI\_A modules, they are named USCI\_A0 and USCI\_A1. See the device-specific data sheet to determine which USCI modules, if any, are implemented on each device.

The USCI\_Ax modules support:

- UART mode
- Pulse shaping for IrDA communications
- Automatic baud rate detection for LIN communications
- SPI mode

The USCI\_Bx modules support:

- I<sup>2</sup>C mode
- SPI mode

### 16.2 USCI Introduction: SPI Mode

In synchronous mode, the USCI connects the MSP430 to an external system via three or four pins: UCxSIMO, UCxSOMI, UCxCLK, and UCxSTE. SPI mode is selected when the UCSYNC bit is set and SPI mode (3-pin or 4-pin) is selected with the UCMODEx bits.

SPI mode features include:

- 7- or 8-bit data length
- LSB-first or MSB-first data transmit and receive
- 3-pin and 4-pin SPI operation
- · Master or slave modes
- Independent transmit and receive shift registers
- Separate transmit and receive buffer registers
- Continuous transmit and receive operation
- Selectable clock polarity and phase control
- Programmable clock frequency in master mode
- Independent interrupt capability for receive and transmit
- Slave operation in LPM4

Figure 16-1 shows the USCI when configured for SPI mode.



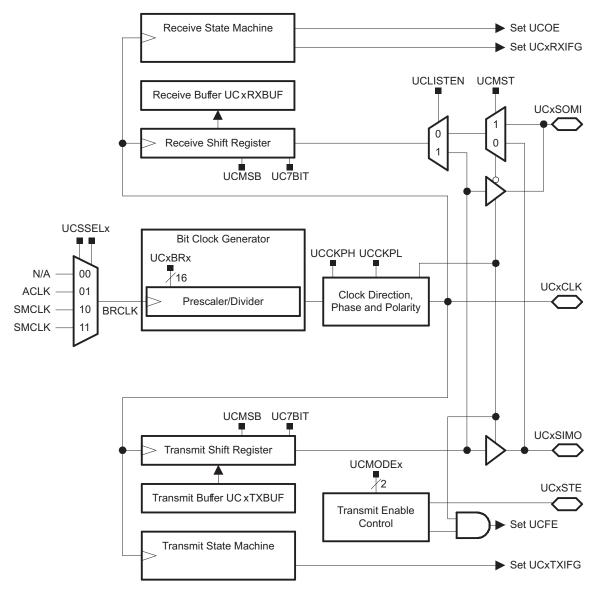


Figure 16-1. USCI Block Diagram: SPI Mode



### 16.3 USCI Operation: SPI Mode

In SPI mode, serial data is transmitted and received by multiple devices using a shared clock provided by the master. An additional pin, UCxSTE, is provided to enable a device to receive and transmit data and is controlled by the master.

Three or four signals are used for SPI data exchange:

- UCxSIMO: Slave in, master out
  - Master mode: UCxSIMO is the data output line.
  - Slave mode: UCxSIMO is the data input line.
- UCxSOMI: Slave out, master in
  - Master mode: UCxSOMI is the data input line.
  - Slave mode: UCxSOMI is the data output line.
- UCxCLK: USCI SPI clock
  - Master mode: UCxCLK is an output.
  - Slave mode: UCxCLK is an input.
- UCxSTE: Slave transmit enable

Used in 4-pin mode to allow multiple masters on a single bus. Not used in 3-pin mode. Table 16-1 describes the UCxSTE operation.

UCMODEx	UCxSTE Active State	UCxSTE	Slave	Master
01	Lliab	0	Inactive	Active
O1	High	1	Active	Inactive
40	Low	0	Active	Inactive
10	Low	1	Inactive	Active

Table 16-1. UCxSTE Operation

### 16.3.1 USCI Initialization and Reset

The USCI is reset by a PUC or by the UCSWRST bit. After a PUC, the UCSWRST bit is automatically set, keeping the USCI in a reset condition. When set, the UCSWRST bit resets the UCxRXIE, UCxTXIE, UCxRXIFG, UCOE, and UCFE bits and sets the UCxTXIFG flag. Clearing UCSWRST releases the USCI for operation.

### NOTE: Initializing or Re-Configuring the USCI Module

The recommended USCI initialization/re-configuration process is:

- 1. Set UCSWRST (BIS.B #UCSWRST,&UCxCTL1)
- 2. Initialize all USCI registers with UCSWRST=1 (including UCxCTL1)
- 3. Configure ports
- 4. Clear UCSWRST via software (BIC.B #UCSWRST,&UCxCTL1)
- 5. Enable interrupts (optional) via UCxRXIE and/or UCxTXIE



### 16.3.2 Character Format

The USCI module in SPI mode supports 7-bit and 8-bit character lengths selected by the UC7BIT bit. In 7-bit data mode, UCxRXBUF is LSB justified and the MSB is always reset. The UCMSB bit controls the direction of the transfer and selects LSB or MSB first.

### **NOTE: Default Character Format**

The default SPI character transmission is LSB first. For communication with other SPI interfaces it MSB-first mode may be required.

### **NOTE: Character Format for Figures**

Figures throughout this chapter use MSB first format.

### 16.3.3 Master Mode

Figure 16-2 shows the USCI as a master in both 3-pin and 4-pin configurations. The USCI initiates data transfer when data is moved to the transmit data buffer UCxTXBUF. The UCxTXBUF data is moved to the TX shift register when the TX shift register is empty, initiating data transfer on UCxSIMO starting with either the most-significant or least-significant bit depending on the UCMSB setting. Data on UCxSOMI is shifted into the receive shift register on the opposite clock edge. When the character is received, the receive data is moved from the RX shift register to the received data buffer UCxRXBUF and the receive interrupt flag, UCxRXIFG, is set, indicating the RX/TX operation is complete.

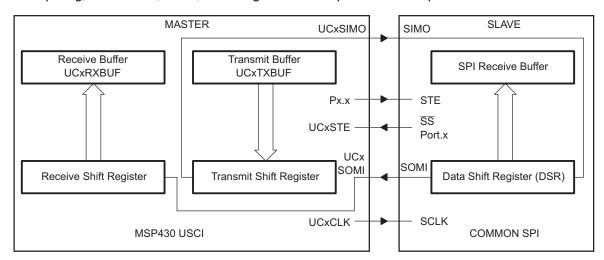


Figure 16-2. USCI Master and External Slave

A set transmit interrupt flag, UCxTXIFG, indicates that data has moved from UCxTXBUF to the TX shift register and UCxTXBUF is ready for new data. It does not indicate RX/TX completion.

To receive data into the USCI in master mode, data must be written to UCxTXBUF because receive and transmit operations operate concurrently.



### 16.3.3.1 Four-Pin SPI Master Mode

In 4-pin master mode, UCxSTE is used to prevent conflicts with another master and controls the master as described in Table 16-1. When UCxSTE is in the master-inactive state:

- UCxSIMO and UCxCLK are set to inputs and no longer drive the bus
- The error bit UCFE is set indicating a communication integrity violation to be handled by the user.
- The internal state machines are reset and the shift operation is aborted.

If data is written into UCxTXBUF while the master is held inactive by UCxSTE, it will be transmitted as soon as UCxSTE transitions to the master-active state. If an active transfer is aborted by UCxSTE transitioning to the master-inactive state, the data must be re-written into UCxTXBUF to be transferred when UCxSTE transitions back to the master-active state. The UCxSTE input signal is not used in 3-pin master mode.

### 16.3.4 Slave Mode

Figure 16-3 shows the USCI as a slave in both 3-pin and 4-pin configurations. UCxCLK is used as the input for the SPI clock and must be supplied by the external master. The data-transfer rate is determined by this clock and not by the internal bit clock generator. Data written to UCxTXBUF and moved to the TX shift register before the start of UCxCLK is transmitted on UCxSOMI. Data on UCxSIMO is shifted into the receive shift register on the opposite edge of UCxCLK and moved to UCxRXBUF when the set number of bits are received. When data is moved from the RX shift register to UCxRXBUF, the UCxRXIFG interrupt flag is set, indicating that data has been received. The overrun error bit, UCOE, is set when the previously received data is not read from UCxRXBUF before new data is moved to UCxRXBUF.

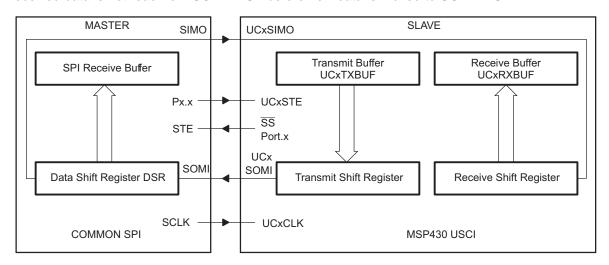


Figure 16-3. USCI Slave and External Master

### 16.3.4.1 Four-Pin SPI Slave Mode

In 4-pin slave mode, UCxSTE is used by the slave to enable the transmit and receive operations and is provided by the SPI master. When UCxSTE is in the slave-active state, the slave operates normally. When UCxSTE is in the slave- inactive state:

- Any receive operation in progress on UCxSIMO is halted
- UCxSOMI is set to the input direction
- The shift operation is halted until the UCxSTE line transitions into the slave transmit active state.

The UCxSTE input signal is not used in 3-pin slave mode.

www.ti.com USCI Operation: SPI Mode

### 16.3.5 SPI Enable

When the USCI module is enabled by clearing the UCSWRST bit it is ready to receive and transmit. In master mode the bit clock generator is ready, but is not clocked nor producing any clocks. In slave mode the bit clock generator is disabled and the clock is provided by the master.

A transmit or receive operation is indicated by UCBUSY = 1.

A PUC or set UCSWRST bit disables the USCI immediately and any active transfer is terminated.

#### 16.3.5.1 Transmit Enable

In master mode, writing to UCxTXBUF activates the bit clock generator and the data will begin to transmit.

In slave mode, transmission begins when a master provides a clock and, in 4-pin mode, when the UCxSTE is in the slave-active state.

#### 16.3.5.2 Receive Enable

The SPI receives data when a transmission is active. Receive and transmit operations operate concurrently.

#### 16.3.6 Serial Clock Control

UCxCLK is provided by the master on the SPI bus. When UCMST = 1, the bit clock is provided by the USCI bit clock generator on the UCxCLK pin. The clock used to generate the bit clock is selected with the UCSSELx bits. When UCMST = 0, the USCI clock is provided on the UCxCLK pin by the master, the bit clock generator is not used, and the UCSSELx bits are don't care. The SPI receiver and transmitter operate in parallel and use the same clock source for data transfer.

The 16-bit value of UCBRx in the bit rate control registers UCxxBR1 and UCxxBR0 is the division factor of the USCI clock source, BRCLK. The maximum bit clock that can be generated in master mode is BRCLK. Modulation is not used in SPI mode and UCAxMCTL should be cleared when using SPI mode for USCI\_A. The UCAxCLK/UCBxCLK frequency is given by:

$$f_{BitClock} = \frac{f_{BRCLK}}{UCBRx}$$



### 16.3.6.1 Serial Clock Polarity and Phase

The polarity and phase of UCxCLK are independently configured via the UCCKPL and UCCKPH control bits of the USCI. Timing for each case is shown in Figure 16-4.

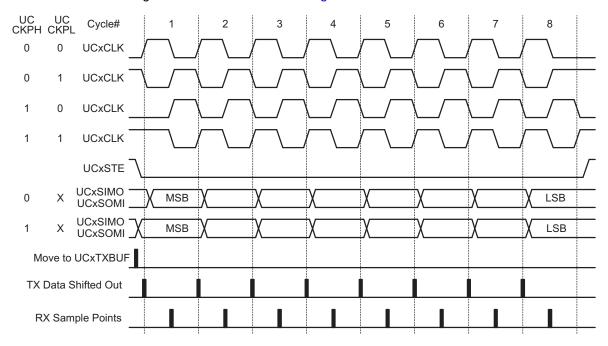


Figure 16-4. USCI SPI Timing with UCMSB = 1

### 16.3.7 Using the SPI Mode With Low-Power Modes

The USCI module provides automatic clock activation for SMCLK for use with low-power modes. When SMCLK is the USCI clock source, and is inactive because the device is in a low-power mode, the USCI module automatically activates it when needed, regardless of the control-bit settings for the clock source. The clock remains active until the USCI module returns to its idle condition. After the USCI module returns to the idle condition, control of the clock source reverts to the settings of its control bits. Automatic clock activation is not provided for ACLK.

When the USCI module activates an inactive clock source, the clock source becomes active for the whole device and any peripheral configured to use the clock source may be affected. For example, a timer using SMCLK increments while the USCI module forces SMCLK active.

In SPI slave mode, no internal clock source is required because the clock is provided by the external master. It is possible to operate the USCI in SPI slave mode while the device is in LPM4 and all clock sources are disabled. The receive or transmit interrupt can wake up the CPU from any low power mode.

### 16.3.8 SPI Interrupts

The USCI has one interrupt vector for transmission and one interrupt vector for reception.

### 16.3.8.1 SPI Transmit Interrupt Operation

The UCxTXIFG interrupt flag is set by the transmitter to indicate that UCxTXBUF is ready to accept another character. An interrupt request is generated if UCxTXIE and GIE are also set. UCxTXIFG is automatically reset if a character is written to UCxTXBUF. UCxTXIFG is set after a PUC or when UCSWRST = 1. UCxTXIE is reset after a PUC or when UCSWRST = 1.

NOTE: Writing to UCxTXBUF in SPI Mode

Data written to UCxTXBUF when UCxTXIFG = 0 may result in erroneous data transmission.



### 16.3.8.2 SPI Receive Interrupt Operation

The UCxRXIFG interrupt flag is set each time a character is received and loaded into UCxRXBUF. An interrupt request is generated if UCxRXIE and GIE are also set. UCxRXIFG and UCxRXIE are reset by a system reset PUC signal or when UCSWRST = 1. UCxRXIFG is automatically reset when UCxRXBUF is read.

### 16.3.8.3 USCI Interrupt Usage

USCI\_Ax and USCI\_Bx share the same interrupt vectors. The receive interrupt flags UCAxRXIFG and UCBxRXIFG are routed to one interrupt vector, the transmit interrupt flags UCAxTXIFG and UCBxTXIFG share another interrupt vector.

Example 16-1 shows an extract of an interrupt service routine to handle data receive interrupts from USCI\_A0 in either UART or SPI mode and USCI\_B0 in SPI mode.

### Example 16-1. Shared Receive Interrupt Vectors Software Example

```
USCIAO_RX_USCIBO_RX_ISR

BIT.B #UCAORXIFG, &IFG2 ; USCI_AO Receive Interrupt?

JNZ USCIAO_RX_ISR

USCIBO_RX_ISR?

; Read UCBORXBUF (clears UCBORXIFG)

...

RETI

USCIAO_RX_ISR

; Read UCAORXBUF (clears UCAORXIFG)

...

RETI

RETI
```

Example 16-2 shows an extract of an interrupt service routine to handle data transmit interrupts from USCI A0 in either UART or SPI mode and USCI B0 in SPI mode.

### Example 16-2. Shared Transmit Interrupt Vectors Software Example

```
USCIAO_TX_USCIBO_TX_ISR

BIT.B #UCAOTXIFG, &IFG2 ; USCI_AO Transmit Interrupt?

JNZ USCIAO_TX_ISR

USCIBO_TX_ISR

; Write UCBOTXBUF (clears UCBOTXIFG)

...

RETI

USCIAO_TX_ISR

; Write UCAOTXBUF (clears UCAOTXIFG)

...

RETI
```

USCI Registers: SPI Mode www.ti.com

### 16.4 USCI Registers: SPI Mode

The USCI registers applicable in SPI mode for USCI\_A0 and USCI\_B0 are listed in Table 16-2. Registers applicable in SPI mode for USCI\_A1 and USCI\_B1 are listed in Table 16-3.

Table 16-2. USCI\_A0 and USCI\_B0 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USCI_A0 control register 0	UCA0CTL0	Read/write	060h	Reset with PUC
USCI_A0 control register 1	UCA0CTL1	Read/write	061h	001h with PUC
USCI_A0 baud rate control register 0	UCA0BR0	Read/write	062h	Reset with PUC
USCI_A0 baud rate control register 1	UCA0BR1	Read/write	063h	Reset with PUC
USCI_A0 modulation control register	UCA0MCTL	Read/write	064h	Reset with PUC
USCI_A0 status register	UCA0STAT	Read/write	065h	Reset with PUC
USCI_A0 receive buffer register	UCA0RXBUF	Read	066h	Reset with PUC
USCI_A0 transmit buffer register	UCA0TXBUF	Read/write	067h	Reset with PUC
USCI_B0 control register 0	UCB0CTL0	Read/write	068h	001h with PUC
USCI_B0 control register 1	UCB0CTL1	Read/write	069h	001h with PUC
USCI_B0 bit rate control register 0	UCB0BR0	Read/write	06Ah	Reset with PUC
USCI_B0 bit rate control register 1	UCB0BR1	Read/write	06Bh	Reset with PUC
USCI_B0 status register	UCB0STAT	Read/write	06Dh	Reset with PUC
USCI_B0 receive buffer register	UCB0RXBUF	Read	06Eh	Reset with PUC
USCI_B0 transmit buffer register	UCB0TXBUF	Read/write	06Fh	Reset with PUC
SFR interrupt enable register 2	IE2	Read/write	001h	Reset with PUC
SFR interrupt flag register 2	IFG2	Read/write	003h	00Ah with PUC

### NOTE: Modifying SFR bits

To avoid modifying control bits of other modules, it is recommended to set or clear the IEx and IFGx bits using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.

Table 16-3. USCI\_A1 and USCI\_B1 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USCI_A1 control register 0	UCA1CTL0	Read/write	0D0h	Reset with PUC
USCI_A1 control register 1	UCA1CTL1	Read/write	0D1h	001h with PUC
USCI_A1 baud rate control register 0	UCA1BR0	Read/write	0D2h	Reset with PUC
USCI_A1 baud rate control register 1	UCA1BR1	Read/write	0D3h	Reset with PUC
USCI_A1 modulation control register	UCA10MCTL	Read/write	0D4h	Reset with PUC
USCI_A1 status register	UCA1STAT	Read/write	0D5h	Reset with PUC
USCI_A1 receive buffer register	UCA1RXBUF	Read	0D6h	Reset with PUC
USCI_A1 transmit buffer register	UCA1TXBUF	Read/write	0D7h	Reset with PUC
USCI_B1 control register 0	UCB1CTL0	Read/write	0D8h	001h with PUC
USCI_B1 control register 1	UCB1CTL1	Read/write	0D9h	001h with PUC
USCI_B1 bit rate control register 0	UCB1BR0	Read/write	0DAh	Reset with PUC
USCI_B1 bit rate control register 1	UCB1BR1	Read/write	0DBh	Reset with PUC
USCI_B1 status register	UCB1STAT	Read/write	0DDh	Reset with PUC
USCI_B1 receive buffer register	UCB1RXBUF	Read	0DEh	Reset with PUC
USCI_B1 transmit buffer register	UCB1TXBUF	Read/write	0DFh	Reset with PUC
USCI_A1/B1 interrupt enable register	UC1IE	Read/write	006h	Reset with PUC
USCI_A1/B1 interrupt flag register	UC1IFG	Read/write	007h	00Ah with PUC



www.ti.com USCI Registers: SPI Mode

# 16.4.1 UCAxCTL0, USCI\_Ax Control Register 0, UCBxCTL0, USCI\_Bx Control Register 0

7	6		5	4	3	2	1	0
UCCKPH	UCCKF	L	UCMSB	UC7BIT	UCMST	UCM	DDEx	UCSYNC=1
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	
UCCKPH	Bit 7	Clock	c phase select.					
		0	Data is char	ged on the first U	CLK edge and cap	tured on the follo	wing edge.	
		1	Data is capt	ured on the first U	CLK edge and cha	nged on the follo	wing edge.	
UCCKPL	Bit 6	Clock	polarity select.					
		0	The inactive	state is low.				
		1	The inactive	state is high.				
UCMSB	Bit 5	MSB	first select. Cor	trols the direction	of the receive and	transmit shift reg	ister.	
		0	LSB first					
		1	MSB first					
UC7BIT	Bit 4	Char	acter length. Se	lects 7-bit or 8-bit	character length.			
		0	8-bit data					
		1	7-bit data					
UCMST	Bit 3	Mast	er mode select					
		0	Slave mode					
		1	Master mode	е				
UCMODEx	Bits 2-1	USCI	mode. The UC	MODEx bits selec	t the synchronous	mode when UCS	YNC = 1.	
		00	3-pin SPI					
		01	4-pin SPI wi	th UCxSTE active	high: slave enable	d when UCxSTE	= 1	
		10	4-pin SPI wi	th UCxSTE active	low: slave enabled	when UCxSTE :	= 0	
		11	I <sup>2</sup> C mode					
UCSYNC	Bit 0	Sync	hronous mode e	enable				
		0	Asynchrono	us mode				
		1	Synchronous	s mode				

# 16.4.2 UCAxCTL1, USCI\_Ax Control Register 1, UCBxCTL1, USCI\_Bx Control Register 1

7	6	5	4	3	2	1	0	
UC	SSELx		Unused					
rw-0	rw-0	rw-0 <sup>(1)</sup> r0 <sup>(2)</sup>	rw-0	rw-0	rw-0	rw-0	rw-1	
UCSSELx	Bits 7-6	USCI clock source se used in slave mode.	elect. These bits s	select the BRCLK s	source clock in ma	aster mode. UCx	CLK is always	
		00 NA						
		01 ACLK						
		10 SMCLK						
		11 SMCLK						
Unused	Bits 5-1	Unused						
UCSWRST	Bit 0	Software reset enable	e					
		0 Disabled. US	CI reset released	d for operation.				
		1 Enabled. US	CI logic held in re	eset state.				

<sup>(1)</sup> UCAxCTL1 (USCI\_Ax)

<sup>(2)</sup> UCBxCTL1 (USCI\_Bx)



16.4.3	UCAxBR0, USCI_Ax Bit Rate Control Registe	er 0,
<b>UCBxB</b>	R0, USCI_Bx Bit Rate Control Register 0	•

7	6	5	4	3	2	1	0
			UCBRx -	low byte			
rw	rw	rw	rw	rw	rw	rw	rw

## 16.4.4 UCAxBR1, USCI\_Ax Bit Rate Control Register 1, UCBxBR1, USCI\_Bx Bit Rate Control Register 1

7	6	5	4	3	2	1	0	
			UCBRx -	high byte				
rw	rw	rw	rw	rw	rw	rw	rw	
UCBRx	UCBRx Bit clock prescaler setting. The 16-bit value of (UCxxBR0 + UCxxBR1 × 256) forms the prescaler value.							

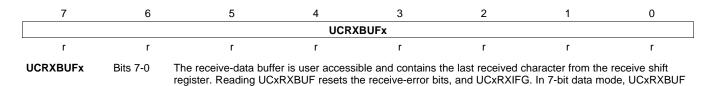
## 16.4.5 UCAxSTAT, USCI\_Ax Status Register, UCBxSTAT, USCI\_Bx Status Register

7	6		5	4	3	2	1	0		
UCLISTEN	UCFE		UCOE		Unu	ısed		UCBUSY		
rw-0	rw-0		rw-0	rw-0 <sup>(1)</sup> r0 <sup>(2)</sup>	rw-0	rw-0	rw-0	r-0		
UCLISTEN	Bit 7	Listen	Listen enable. The UCLISTEN bit selects loopback mode.							
		0	Disabled							
		1	Enabled. The	e transmitter outpu	ut is internally fed	back to the receive	er.			
UCFE	Bit 6		ng error flag. Tl er or any slave r	nis bit indicates a mode.	bus conflict in 4-w	vire master mode.	UCFE is not used	d in 3-wire		
		0	No error							
		1	Bus conflict	occurred						
UCOE	Bit 5	chara	cter was read. I	nis bit is set when JCOE is cleared a it will not function	utomatically whe					
		0	No error							
		1	Overrun erro	r occurred						
Unused	Bits 4-1	Unuse	ed							
UCBUSY	Bit 0	USCI	busy. This bit in	ndicates if a transr	mit or receive ope	ration is in progre	SS.			
		0	USCI inactiv	е						
		1	USCI transm	itting or receiving						

<sup>(1)</sup> UCAxSTAT (USCI\_Ax)

## 16.4.6 UCAxRXBUF, USCI\_Ax Receive Buffer Register, UCBxRXBUF, USCI\_Bx Receive Buffer Register

is LSB justified and the MSB is always reset.



<sup>(2)</sup> UCBxSTAT (USCI\_Bx)



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# 16.4.7 UCAxTXBUF, USCI\_Ax Transmit Buffer Register, UCBxTXBUF, USCI\_Bx Transmit Buffer Register



The transmit data buffer is user accessible and holds the data waiting to be moved into the transmit shift register and transmitted. Writing to the transmit data buffer clears UCxTXIFG. The MSB of UCxTXBUF is not used for 7-bit data and is reset

### 16.4.8 IE2, Interrupt Enable Register 2

7	6	5	4	3	2	1	0
				UCB0TXIE	UCB0RXIE	UCA0TXIE	UCA0RXIE
				rw-0	rw-0	rw-0	rw-0
	Bits 7-4	These bits may be	used by other mod	dules (see the devi	ce-specific data sh	eet).	
<b>UCB0TXIE</b>	Bit 3	USCI_B0 transmit	interrupt enable				
		0 Interrupt d	isabled				
		1 Interrupt e	nabled				
<b>UCB0RXIE</b>	Bit 2	USCI_B0 receive i	nterrupt enable				
		0 Interrupt d	isabled				
		1 Interrupt e	nabled				
<b>UCA0TXIE</b>	Bit 1	USCI_A0 transmit	interrupt enable				
		0 Interrupt d	isabled				
		1 Interrupt e	nabled				
<b>UCA0RXIE</b>	Bit 0	USCI_A0 receive i	nterrupt enable				
		0 Interrupt d	isabled				
		1 Interrupt e	nabled				

### 16.4.9 IFG2, Interrupt Flag Register 2

7	6	5	4	3	2	1	0
				UCB0TXIFG	UCB0RXIFG	UCA0TXIFG	UCA0RXIFG
				rw-1	rw-0	rw-1	rw-0
	Bits 7-4	These bits may be us	sed by other modu	ules (see the devic	e-specific data sh	eet).	
<b>UCB0TXIFG</b>	Bit 3	USCI_B0 transmit inf	errupt flag. UCB0	TXIFG is set wher	n UCB0TXBUF is	empty.	
		0 No interrupt	pending				
		1 Interrupt per	iding				
UCB0RXIFG	Bit 2	USCI_B0 receive inte	errupt flag. UCB0F	RXIFG is set when	UCB0RXBUF ha	s received a comp	olete character.
		0 No interrupt	pending				
		1 Interrupt per	iding				
<b>UCA0TXIFG</b>	Bit 1	USCI_A0 transmit in	errupt flag. UCA0	TXIFG is set wher	n UCA0TXBUF en	npty.	
		0 No interrupt	pending				
		1 Interrupt per	iding				
<b>UCA0RXIFG</b>	Bit 0	USCI_A0 receive into	errupt flag. UCA0F	RXIFG is set when	UCA0RXBUF ha	s received a comp	olete character.
		0 No interrupt	pending				
		1 Interrupt per	ding				

USCI Registers: SPI Mode www.ti.com

## 16.4.10 UC1IE, USCI\_A1/USCI\_B1 Interrupt Enable Register

7	6	5	4	3	2	1	0
		Unused		UCB1TXIE	UCB1RXIE	UCA1TXIE	UCA1RXIE
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
Unused	Bits 7-4	Unused					
UCB1TXIE	Bit 3	USCI_B1 transmit inte	errupt enable				
		0 Interrupt disa	bled				
		1 Interrupt enal	oled				
UCB1RXIE	Bit 2	USCI_B1 receive inte	rrupt enable				
		0 Interrupt disa	bled				
		1 Interrupt enal	oled				
UCA1TXIE	Bit 1	USCI_A1 transmit inte	errupt enable				
		0 Interrupt disa	bled				
		1 Interrupt enal	oled				
UCA1RXIE	Bit 0	USCI_A1 receive inte	rrupt enable				
		0 Interrupt disa	bled				
		1 Interrupt enal	oled				

## 16.4.11 UC1IFG, USCI\_A1/USCI\_B1 Interrupt Flag Register

7	6	5	4	3	2	1	0
		Unused		UCB1TXIFG	UCB1RXIFG	UCA1TXIFG	UCA1RXIFG
rw-0	rw-0	rw-0	rw-0	rw-1	rw-0	rw-1	rw-0
Unused	Bits 7-4	Unused					
UCB1TXIFG	Bit 3	USCI_B1 transmit inte  No interrupt p  Interrupt pend	ending	TXIFG is set when	n UCB1TXBUF is	empty.	
UCB1RXIFG	Bit 2	USCI_B1 receive inter  No interrupt p  Interrupt pend	ending	RXIFG is set wher	uCB1RXBUF ha	s received a comp	olete character.
UCA1TXIFG	Bit 1	USCI_A1 transmit inte 0 No interrupt p	ending	TXIFG is set when	n UCA1TXBUF en	npty.	
UCA1RXIFG	Bit 0	USCI_A1 receive inter  No interrupt pend  Interrupt pend	ending	RXIFG is set wher	n UCA1RXBUF ha	s received a comp	olete character.



## Universal Serial Communication Interface, I<sup>2</sup>C Mode

The universal serial communication interface (USCI) supports multiple serial communication modes with one hardware module. This chapter discusses the operation of the I<sup>2</sup>C mode.

Topic Page

17.1 USCI Overview	450
17.2 USCI Introduction: I <sup>2</sup> C Mode	450
17.3 USCI Operation: I <sup>2</sup> C Mode	451
17.4 USCI Registers: I <sup>2</sup> C Mode	



USCI Overview www.ti.com

### 17.1 USCI Overview

The universal serial communication interface (USCI) modules support multiple serial communication modes. Different USCI modules support different modes. Each different USCI module is named with a different letter. For example, USCI\_A is different from USCI\_B, etc. If more than one identical USCI module is implemented on one device, those modules are named with incrementing numbers. For example, if one device has two USCI\_A modules, they are named USCI\_A0 and USCI\_A1. See the device-specific data sheet to determine which USCI modules, if any, are implemented on which devices.

The USCI\_Ax modules support:

- UART mode
- · Pulse shaping for IrDA communications
- Automatic baud rate detection for LIN communications
- SPI mode

The USCI\_Bx modules support:

- I2C mode
- SPI mode

### 17.2 USCI Introduction: I<sup>2</sup>C Mode

In I<sup>2</sup>C mode, the USCI module provides an interface between the MSP430 and I<sup>2</sup>C-compatible devices connected by way of the two-wire I<sup>2</sup>C serial bus. External components attached to the I<sup>2</sup>C bus serially transmit and/or receive serial data to/from the USCI module through the 2-wire I<sup>2</sup>C interface.

The I2C mode features include:

- Compliance to the Philips Semiconductor I<sup>2</sup>C specification v2.1
  - 7-bit and 10-bit device addressing modes
  - General call
  - START/RESTART/STOP
  - Multi-master transmitter/receiver mode
  - Slave receiver/transmitter mode
  - Standard mode up to 100 kbps and fast mode up to 400 kbps support
- Programmable UCxCLK frequency in master mode
- Designed for low power
- Slave receiver START detection for auto-wake up from LPMx modes
- Slave operation in LPM4

Figure 17-1 shows the USCI when configured in I<sup>2</sup>C mode.



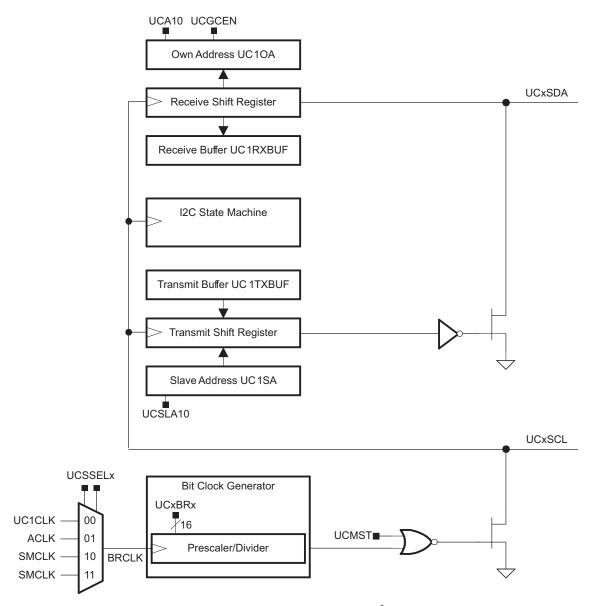


Figure 17-1. USCI Block Diagram: I<sup>2</sup>C Mode

### 17.3 USCI Operation: I<sup>2</sup>C Mode

The I<sup>2</sup>C mode supports any slave or master I<sup>2</sup>C-compatible device. Figure 17-2 shows an example of an I<sup>2</sup>C bus. Each I<sup>2</sup>C device is recognized by a unique address and can operate as either a transmitter or a receiver. A device connected to the I<sup>2</sup>C bus can be considered as the master or the slave when performing data transfers. A master initiates a data transfer and generates the clock signal SCL. Any device addressed by a master is considered a slave.

I<sup>2</sup>C data is communicated using the serial data pin (SDA) and the serial clock pin (SCL). Both SDA and SCL are bidirectional, and must be connected to a positive supply voltage using a pullup resistor.

### NOTE: SDA and SCL Levels

The MSP430 SDA and SCL pins must not be pulled up above the MSP430 V<sub>cc</sub> level.

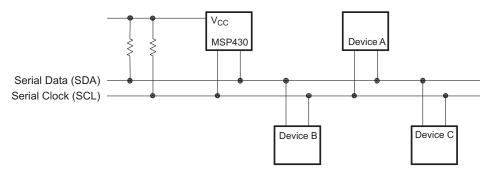


Figure 17-2. I<sup>2</sup>C Bus Connection Diagram

### 17.3.1 USCI Initialization and Reset

The USCI is reset by a PUC or by setting the UCSWRST bit. After a PUC, the UCSWRST bit is automatically set, keeping the USCI in a reset condition. To select I<sup>2</sup>C operation the UCMODEx bits must be set to 11. After module initialization, it is ready for transmit or receive operation. Clearing UCSWRST releases the USCI for operation.

Configuring and reconfiguring the USCI module should be done when UCSWRST is set to avoid unpredictable behavior. Setting UCSWRST in I<sup>2</sup>C mode has the following effects:

- I<sup>2</sup>C communication stops
- SDA and SCL are high impedance
- UCBxI2CSTAT, bits 6-0 are cleared
- UCBxTXIE and UCBxRXIE are cleared
- UCBxTXIFG and UCBxRXIFG are cleared
- All other bits and registers remain unchanged.

### NOTE: Initializing or Reconfiguring the USCI Module

The recommended USCI initialization or reconfiguration process is:

- Set UCSWRST (BIS.B #UCSWRST,&UCxCTL1)
- 2. Initialize all USCI registers with UCSWRST=1 (including UCxCTL1)
- Configure ports.
- 4. Clear UCSWRST via software (BIC.B #UCSWRST,&UCxCTL1)
- 5. Enable interrupts (optional) via UCxRXIE and/or UCxTXIE

### 17.3.2 fC Serial Data

One clock pulse is generated by the master device for each data bit transferred. The I<sup>2</sup>C mode operates with byte data. Data is transferred most significant bit first as shown in Figure 17-3.

The first byte after a START condition consists of a 7-bit slave address and the R/W bit. When R/W = 0, the master transmits data to a slave. When R/W = 1, the master receives data from a slave. The ACK bit is sent from the receiver after each byte on the 9th SCL clock.

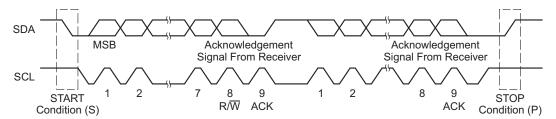


Figure 17-3. I<sup>2</sup>C Module Data Transfer



START and STOP conditions are generated by the master and are shown in Figure 17-3. A START condition is a high-to-low transition on the SDA line while SCL is high. A STOP condition is a low-to-high transition on the SDA line while SCL is high. The bus busy bit, UCBBUSY, is set after a START and cleared after a STOP.

Data on SDA must be stable during the high period of SCL as shown in Figure 17-4. The high and low state of SDA can only change when SCL is low, otherwise START or STOP conditions will be generated.

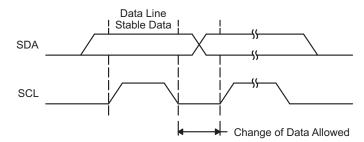


Figure 17-4. Bit Transfer on the I<sup>2</sup>C Bus

### 17.3.3 FC Addressing Modes

The I<sup>2</sup>C mode supports 7-bit and 10-bit addressing modes.

### 17.3.3.1 7-Bit Addressing

In the 7-bit addressing format, shown in Figure 17-5, the first byte is the 7-bit slave address and the R/W bit. The ACK bit is sent from the receiver after each byte.

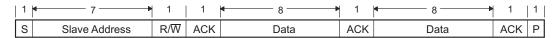


Figure 17-5. I<sup>2</sup>C Module 7-Bit Addressing Format

### 17.3.3.2 10-Bit Addressing

In the 10-bit addressing format, shown in Figure 17-6, the first byte is made up of 11110b plus the two MSBs of the 10-bit slave address and the R/W bit. The ACK bit is sent from the receiver after each byte. The next byte is the remaining 8 bits of the 10-bit slave address, followed by the ACK bit and the 8-bit data.

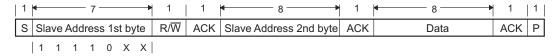


Figure 17-6. I<sup>2</sup>C Module 10-Bit Addressing Format

### 17.3.3.3 Repeated Start Conditions

The direction of data flow on SDA can be changed by the master, without first stopping a transfer, by issuing a repeated START condition. This is called a RESTART. After a RESTART is issued, the slave address is again sent out with the new data direction specified by the R/W bit. The RESTART condition is shown in Figure 17-7.

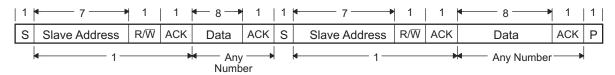


Figure 17-7. I<sup>2</sup>C Module Addressing Format with Repeated START Condition

### 17.3.4 fC Module Operating Modes

In I<sup>2</sup>C mode the USCI module can operate in master transmitter, master receiver, slave transmitter, or slave receiver mode. The modes are discussed in the following sections. Time lines are used to illustrate the modes.

Figure 17-8 shows how to interpret the time line figures. Data transmitted by the master is represented by grey rectangles, data transmitted by the slave by white rectangles. Data transmitted by the USCI module, either as master or slave, is shown by rectangles that are taller than the others.

Actions taken by the USCI module are shown in grey rectangles with an arrow indicating where in the data stream the action occurs. Actions that must be handled with software are indicated with white rectangles with an arrow pointing to where in the data stream the action must take place.

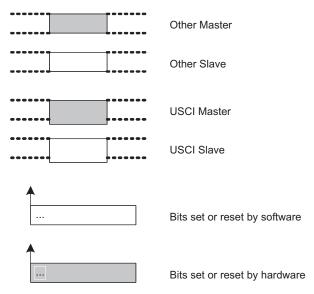


Figure 17-8. I<sup>2</sup>C Time Line Legend

### 17.3.4.1 Slave Mode

The USCI module is configured as an  $I^2C$  slave by selecting the  $I^2C$  mode with UCMODEx = 11 and UCSYNC = 1 and clearing the UCMST bit.

Initially the USCI module must to be configured in receiver mode by clearing the UCTR bit to receive the  $I^2C$  address. Afterwards, transmit and receive operations are controlled automatically depending on the R/W bit received together with the slave address.

The USCI slave address is programmed with the UCBxI2COA register. When UCA10 = 0, 7-bit addressing is selected. When UCA10 = 1, 10-bit addressing is selected. The UCGCEN bit selects if the slave responds to a general call.

When a START condition is detected on the bus, the USCI module will receive the transmitted address and compare it against its own address stored in UCBxI2COA. The UCSTTIFG flag is set when address received matches the USCI slave address.



### 17.3.4.1.1 fC Slave Transmitter Mode

Slave transmitter mode is entered when the slave address transmitted by the master is identical to its own address with a set R/W bit. The slave transmitter shifts the serial data out on SDA with the clock pulses that are generated by the master device. The slave device does not generate the clock, but it will hold SCL low while intervention of the CPU is required after a byte has been transmitted.

If the master requests data from the slave the USCI module is automatically configured as a transmitter and UCTR and UCBxTXIFG become set. The SCL line is held low until the first data to be sent is written into the transmit buffer UCBxTXBUF. Then the address is acknowledged, the UCSTTIFG flag is cleared, and the data is transmitted. As soon as the data is transferred into the shift register the UCBxTXIFG is set again. After the data is acknowledged by the master the next data byte written into UCBxTXBUF is transmitted or if the buffer is empty the bus is stalled during the acknowledge cycle by holding SCL low until new data is written into UCBxTXBUF. If the master sends a NACK succeeded by a STOP condition the UCSTPIFG flag is set. If the NACK is succeeded by a repeated START condition the USCI I<sup>2</sup>C state machine returns to its address-reception state.

Figure 17-9 shows the slave transmitter operation.

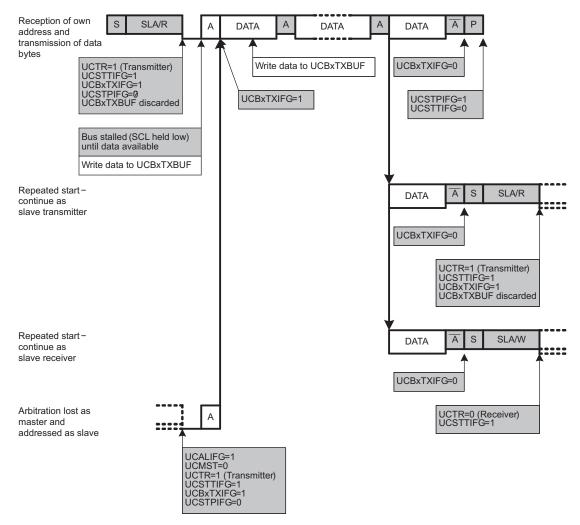


Figure 17-9. I<sup>2</sup>C Slave Transmitter Mode



### 17.3.4.1.2 fC Slave Receiver Mode

Slave receiver mode is entered when the slave address transmitted by the master is identical to its own address and a cleared R/W bit is received. In slave receiver mode, serial data bits received on SDA are shifted in with the clock pulses that are generated by the master device. The slave device does not generate the clock, but it can hold SCL low if intervention of the CPU is required after a byte has been received.

If the slave should receive data from the master the USCI module is automatically configured as a receiver and UCTR is cleared. After the first data byte is received the receive interrupt flag UCBxRXIFG is set. The USCI module automatically acknowledges the received data and can receive the next data byte.

If the previous data was not read from the receive buffer UCBxRXBUF at the end of a reception, the bus is stalled by holding SCL low. As soon as UCBxRXBUF is read the new data is transferred into UCBxRXBUF, an acknowledge is sent to the master, and the next data can be received.

Setting the UCTXNACK bit causes a NACK to be transmitted to the master during the next acknowledgment cycle. A NACK is sent even if UCBxRXBUF is not ready to receive the latest data. If the UCTXNACK bit is set while SCL is held low the bus will be released, a NACK is transmitted immediately, and UCBxRXBUF is loaded with the last received data. Since the previous data was not read that data will be lost. To avoid loss of data the UCBxRXBUF needs to be read before UCTXNACK is set.

When the master generates a STOP condition the UCSTPIFG flag is set.

If the master generates a repeated START condition the USCI I<sup>2</sup>C state machine returns to its address reception state.

Figure 17-10 shows the I<sup>2</sup>C slave receiver operation.



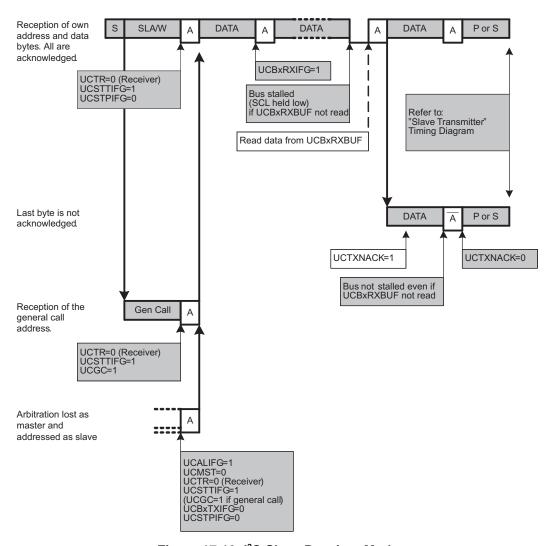


Figure 17-10. I<sup>2</sup>C Slave Receiver Mode



### 17.3.4.1.3 fC Slave 10-bit Addressing Mode

The 10-bit addressing mode is selected when UCA10 = 1 and is as shown in Figure 17-11. In 10-bit addressing mode, the slave is in receive mode after the full address is received. The USCI module indicates this by setting the UCSTTIFG flag while the UCTR bit is cleared. To switch the slave into transmitter mode the master sends a repeated START condition together with the first byte of the address but with the R/W bit set. This will set the UCSTTIFG flag if it was previously cleared by software and the USCI modules switches to transmitter mode with UCTR = 1.

Slave Receiver Reception of own S 11110 xx/W DATA P or S SLA (2.) DATA address and data bytes. All are acknowledged. UCBxRXIFG=1 UCTR=0 (Receiver) UCSTTIFG=1 UCSTPIFG=0 Reception of the Gen Call DATA DATA  $\mathsf{P}\,\mathsf{or}\,\mathsf{S}$ Α Α general call address. UCTR=0 (Receiver) UCSTTIFG=1 UCBxRXIFG=1 UCGC=1 Slave Transmitter Reception of own S 11110 xx/W SLA (2.) S 11110 xx/R Α P or S DATA address and transmission of data bytes UCTR=0 (Receiver) UCSTTIFG=1 UCSTPIFG=0 UCSTTIFG=0 UCTR=1 (Transmitter) UCSTTIFG=1 UCBxTXIFG=1 UCSTPIFG=0

Figure 17-11. I<sup>2</sup>C Slave 10-bit Addressing Mode

Universal Serial Communication Interface, f<sup>2</sup>C Mode



#### 17.3.4.2 Master Mode

The USCI module is configured as an  $I^2C$  master by selecting the  $I^2C$  mode with UCMODEx = 11 and UCSYNC = 1 and setting the UCMST bit. When the master is part of a multi-master system, UCMM must be set and its own address must be programmed into the UCBxI2COA register. When UCA10 = 0, 7-bit addressing is selected. When UCA10 = 1, 10-bit addressing is selected. The UCGCEN bit selects if the USCI module responds to a general call.

### 17.3.4.2.1 PC Master Transmitter Mode

After initialization, master transmitter mode is initiated by writing the desired slave address to the UCBxI2CSA register, selecting the size of the slave address with the UCSLA10 bit, setting UCTR for transmitter mode, and setting UCTXSTT to generate a START condition.

The USCI module checks if the bus is available, generates the START condition, and transmits the slave address. The UCBxTXIFG bit is set when the START condition is generated and the first data to be transmitted can be written into UCBxTXBUF. As soon as the slave acknowledges the address the UCTXSTT bit is cleared.

The data written into UCBxTXBUF is transmitted if arbitration is not lost during transmission of the slave address. UCBxTXIFG is set again as soon as the data is transferred from the buffer into the shift register. If there is no data loaded to UCBxTXBUF before the acknowledge cycle, the bus is held during the acknowledge cycle with SCL low until data is written into UCBxTXBUF. Data is transmitted or the bus is held as long as the UCTXSTP bit or UCTXSTT bit is not set.

Setting UCTXSTP will generate a STOP condition after the next acknowledge from the slave. If UCTXSTP is set during the transmission of the slave's address or while the USCI module waits for data to be written into UCBxTXBUF, a STOP condition is generated even if no data was transmitted to the slave. When transmitting a single byte of data, the UCTXSTP bit must be set while the byte is being transmitted, or anytime after transmission begins, without writing new data into UCBxTXBUF. Otherwise, only the address will be transmitted. When the data is transferred from the buffer to the shift register, UCBxTXIFG will become set indicating data transmission has begun and the UCTXSTP bit may be set.

Setting UCTXSTT will generate a repeated START condition. In this case, UCTR may be set or cleared to configure transmitter or receiver, and a different slave address may be written into UCBxI2CSA if desired.

If the slave does not acknowledge the transmitted data the not-acknowledge interrupt flag UCNACKIFG is set. The master must react with either a STOP condition or a repeated START condition. If data was already written into UCBxTXBUF it will be discarded. If this data should be transmitted after a repeated START it must be written into UCBxTXBUF again. Any set UCTXSTT is discarded, too. To trigger a repeated start UCTXSTT needs to be set again.

Figure 17-12 shows the I<sup>2</sup>C master transmitter operation.



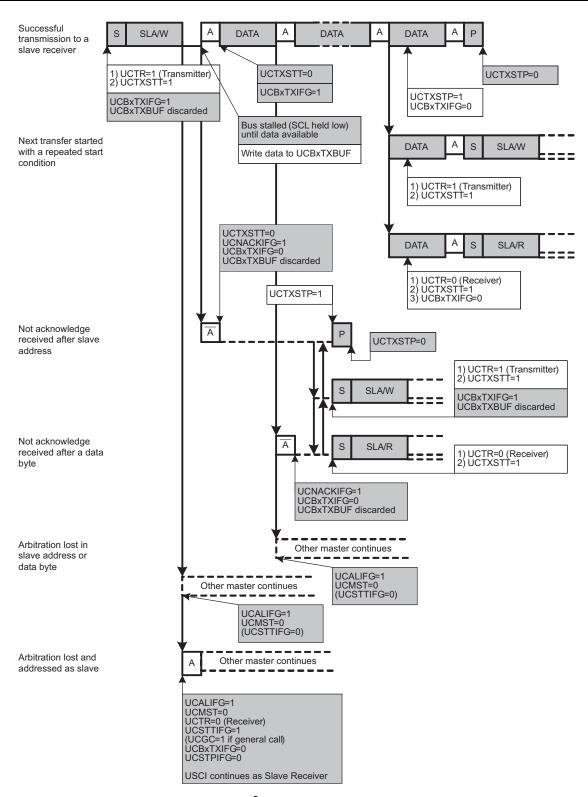


Figure 17-12. I<sup>2</sup>C Master Transmitter Mode



### 17.3.4.2.2 PC Master Receiver Mode

After initialization, master receiver mode is initiated by writing the desired slave address to the UCBxI2CSA register, selecting the size of the slave address with the UCSLA10 bit, clearing UCTR for receiver mode, and setting UCTXSTT to generate a START condition.

The USCI module checks if the bus is available, generates the START condition, and transmits the slave address. As soon as the slave acknowledges the address the UCTXSTT bit is cleared.

After the acknowledge of the address from the slave the first data byte from the slave is received and acknowledged and the UCBxRXIFG flag is set. Data is received from the slave ss long as UCTXSTP or UCTXSTT is not set. If UCBxRXBUF is not read the master holds the bus during reception of the last data bit and until the UCBxRXBUF is read.

If the slave does not acknowledge the transmitted address the not-acknowledge interrupt flag UCNACKIFG is set. The master must react with either a STOP condition or a repeated START condition.

Setting the UCTXSTP bit will generate a STOP condition. After setting UCTXSTP, a NACK followed by a STOP condition is generated after reception of the data from the slave, or immediately if the USCI module is currently waiting for UCBxRXBUF to be read.

If a master wants to receive a single byte only, the UCTXSTP bit must be set while the byte is being received. For this case, the UCTXSTT may be polled to determine when it is cleared:

```
BIS.B #UCTXSTT,&UCBOCTL1 ;Transmit START cond.

POLL_STT BIT.B #UCTXSTT,&UCBOCTL1 ;Poll UCTXSTT bit

JC POLL_STT ;When cleared,

BIS.B #UCTXSTP,&UCBOCTL1 ;transmit STOP cond.
```

Setting UCTXSTT will generate a repeated START condition. In this case, UCTR may be set or cleared to configure transmitter or receiver, and a different slave address may be written into UCBxI2CSA if desired.

Figure 17-13 shows the I<sup>2</sup>C master receiver operation.

#### NOTE: Consecutive Master Transactions Without Repeated Start

When performing multiple consecutive  $I^2C$  master transactions without the repeated start feature, the current transaction must be completed before the next one is initiated. This can be done by ensuring that the transmit stop condition flag UCTXSTP is cleared before the next  $I^2C$  transaction is initiated with setting UCTXSTT = 1. Otherwise, the current transaction might be affected.



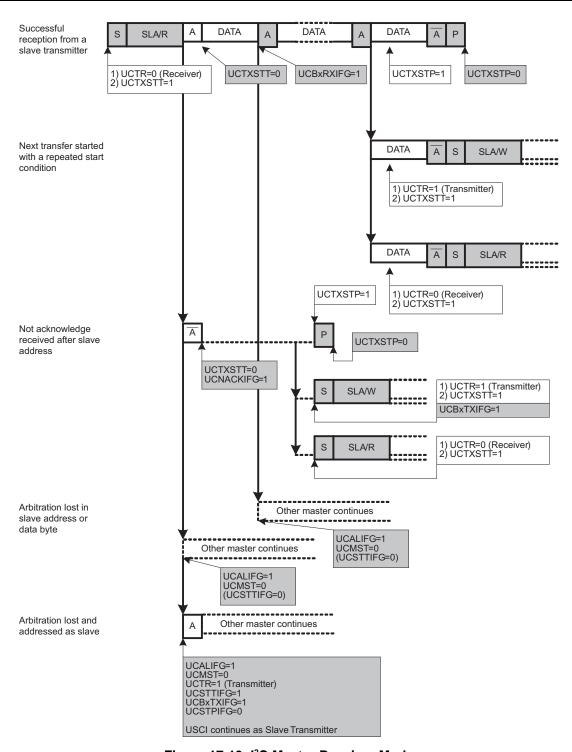


Figure 17-13. I<sup>2</sup>C Master Receiver Mode



### 17.3.4.2.3 PC Master 10-Bit Addressing Mode

The 10-bit addressing mode is selected when UCSLA10 = 1 and is shown in Figure 17-14.

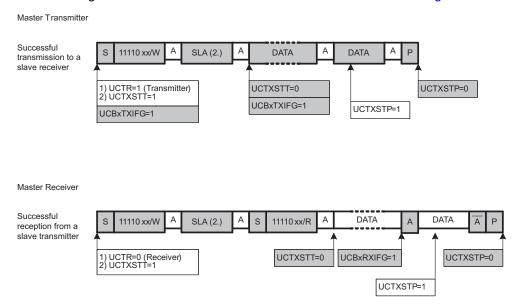


Figure 17-14. I<sup>2</sup>C Master 10-bit Addressing Mode

### 17.3.4.2.4 Arbitration

If two or more master transmitters simultaneously start a transmission on the bus, an arbitration procedure is invoked. Figure 17-15 shows the arbitration procedure between two devices. The arbitration procedure uses the data presented on SDA by the competing transmitters. The first master transmitter that generates a logic high is overruled by the opposing master generating a logic low. The arbitration procedure gives priority to the device that transmits the serial data stream with the lowest binary value. The master transmitter that lost arbitration switches to the slave receiver mode, and sets the arbitration lost flag UCALIFG. If two or more devices send identical first bytes, arbitration continues on the subsequent bytes.

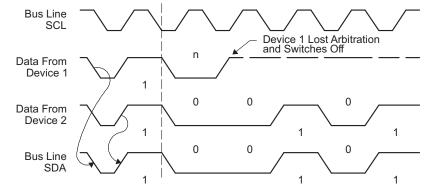


Figure 17-15. Arbitration Procedure Between Two Master Transmitters

If the arbitration procedure is in progress when a repeated START condition or STOP condition is transmitted on SDA, the master transmitters involved in arbitration must send the repeated START condition or STOP condition at the same position in the format frame. Arbitration is not allowed between:

- A repeated START condition and a data bit
- A STOP condition and a data bit
- A repeated START condition and a STOP condition



### 17.3.5 fC Clock Generation and Synchronization

The I<sup>2</sup>C clock SCL is provided by the master on the I<sup>2</sup>C bus. When the USCI is in master mode, BITCLK is provided by the USCI bit clock generator and the clock source is selected with the UCSSELx bits. In slave mode the bit clock generator is not used and the UCSSELx bits are don't care.

The 16-bit value of UCBRx in registers UCBxBR1 and UCBxBR0 is the division factor of the USCI clock source, BRCLK. The maximum bit clock that can be used in single master mode is  $f_{BRCLK}/4$ . In multi-master mode the maximum bit clock is  $f_{BRCLK}/8$ . The BITCLK frequency is given by:

$$f_{BitClock} = \frac{f_{BRCLK}}{UCBRx}$$

The minimum high and low periods of the generated SCL are

$$\begin{split} t_{LOW,MIN} &= t_{HIGH,MIN} = \frac{\text{UCBRx} \, / \, 2}{f_{BRCLK}} & \text{when UCBRx is even and} \\ t_{LOW,MIN} &= t_{HIGH,MIN} = \frac{(\text{UCBRx} - 1) \, / \, 2}{f_{BRCLK}} & \text{when UCBRx is odd.} \end{split}$$

The USCI clock source frequency and the prescaler setting UCBRx must to be chosen such that the minimum low and high period times of the I<sup>2</sup>C specification are met.

During the arbitration procedure the clocks from the different masters must be synchronized. A device that first generates a low period on SCL overrules the other devices forcing them to start their own low periods. SCL is then held low by the device with the longest low period. The other devices must wait for SCL to be released before starting their high periods. Figure 17-16 shows the clock synchronization. This allows a slow slave to slow down a fast master.

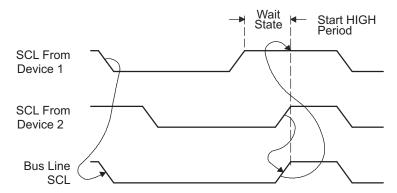


Figure 17-16. Synchronization of Two I<sup>2</sup>C Clock Generators During Arbitration

### 17.3.5.1 Clock Stretching

The USCI module supports clock stretching and also makes use of this feature as described in the operation mode sections.

The UCSCLLOW bit can be used to observe if another device pulls SCL low while the USCI module already released SCL due to the following conditions:

- USCI is acting as master and a connected slave drives SCL low.
- USCI is acting as master and another master drives SCL low during arbitration.

The UCSCLLOW bit is also active if the USCI holds SCL low because it is waiting as transmitter for data being written into UCBxTXBUF or as receiver for the data being read from UCBxRXBUF.

The UCSCLLOW bit might get set for a short time with each rising SCL edge because the logic observes the external SCL and compares it to the internally generated SCL.



### 17.3.6 Using the USCI Module in fC Mode with Low-Power Modes

The USCI module provides automatic clock activation for SMCLK for use with low-power modes. When SMCLK is the USCI clock source, and is inactive because the device is in a low-power mode, the USCI module automatically activates it when needed, regardless of the control-bit settings for the clock source. The clock remains active until the USCI module returns to its idle condition. After the USCI module returns to the idle condition, control of the clock source reverts to the settings of its control bits. Automatic clock activation is not provided for ACLK.

When the USCI module activates an inactive clock source, the clock source becomes active for the whole device and any peripheral configured to use the clock source may be affected. For example, a timer using SMCLK will increment while the USCI module forces SMCLK active.

In I<sup>2</sup>C slave mode no internal clock source is required because the clock is provided by the external master. It is possible to operate the USCI in I<sup>2</sup>C slave mode while the device is in LPM4 and all internal clock sources are disabled. The receive or transmit interrupts can wake up the CPU from any low power mode.

### 17.3.7 USCI Interrupts in fC Mode

There are two interrupt vectors for the USCI module in I<sup>2</sup>C mode. One interrupt vector is associated with the transmit and receive interrupt flags. The other interrupt vector is associated with the four state change interrupt flags. Each interrupt flag has its own interrupt enable bit. When an interrupt is enabled, and the GIE bit is set, the interrupt flag will generate an interrupt request. DMA transfers are controlled by the UCBxTXIFG and UCBxRXIFG flags on devices with a DMA controller.

### 17.3.7.1 I<sup>2</sup>C Transmit Interrupt Operation

The UCBxTXIFG interrupt flag is set by the transmitter to indicate that UCBxTXBUF is ready to accept another character. An interrupt request is generated if UCBxTXIE and GIE are also set. UCBxTXIFG is automatically reset if a character is written to UCBxTXBUF or if a NACK is received. UCBxTXIFG is set when UCSWRST = 1 and the I<sup>2</sup>C mode is selected. UCBxTXIE is reset after a PUC or when UCSWRST = 1.

### 17.3.7.2 I2C Receive Interrupt Operation

The UCBxRXIFG interrupt flag is set when a character is received and loaded into UCBxRXBUF. An interrupt request is generated if UCBxRXIE and GIE are also set. UCBxRXIFG and UCBxRXIE are reset after a PUC signal or when UCSWRST = 1. UCxRXIFG is automatically reset when UCxRXBUF is read.

### 17.3.7.3 I<sup>2</sup>C State Change Interrupt Operation

Table 17-1 describes the I<sup>2</sup>C state change interrupt flags.

Table 17-1. State Change Interrupt Flags

Interrupt Flag	Interrupt Condition
UCALIFG	Arbitration-lost. Arbitration can be lost when two or more transmitters start a transmission simultaneously, or when the USCI operates as master but is addressed as a slave by another master in the system. The UCALIFG flag is set when arbitration is lost. When UCALIFG is set the UCMST bit is cleared and the I <sup>2</sup> C controller becomes a slave.
UCNACKIFG	Not-acknowledge interrupt. This flag is set when an acknowledge is expected but is not received. UCNACKIFG is automatically cleared when a START condition is received.
UCSTTIFG	Start condition detected interrupt. This flag is set when the I <sup>2</sup> C module detects a START condition together with its own address while in slave mode. UCSTTIFG is used in slave mode only and is automatically cleared when a STOP condition is received.
UCSTPIFG	Stop condition detected interrupt. This flag is set when the I <sup>2</sup> C module detects a STOP condition while in slave mode. UCSTPIFG is used in slave mode only and is automatically cleared when a START condition is received.



### 17.3.7.4 Interrupt Vector Assignment

USCI\_Ax and USCI\_Bx share the same interrupt vectors. In I<sup>2</sup>C mode the state change interrupt flags UCSTTIFG, UCSTPIFG, UCNACKIFG, UCALIFG from USCI\_Bx and UCAxRXIFG from USCI\_Ax are routed to one interrupt vector. The I<sup>2</sup>C transmit and receive interrupt flags UCBxTXIFG and UCBxRXIFG from USCI\_Bx and UCAxTXIFG from USCI\_Ax share another interrupt vector.

Example 17-1 shows an extract of the interrupt service routine to handle data receive interrupts from USCI\_A0 in either UART or SPI mode and state change interrupts from USCI\_B0 in I<sup>2</sup>C mode.

### Example 17-1. Shared Receive Interrupt Vectors Software Example

```
USCIAO_RX_USCIBO_I2C_STATE_ISR

BIT.B #UCAORXIFG, &IFG2 ; USCI_AO Receive Interrupt?

JNZ USCIAO_RX_ISR

USCIBO_I2C_STATE_ISR
; Decode I2C state changes ...
; Decode I2C state changes ...

RETI

USCIAO_RX_ISR
; Read UCAORXBUF ... - clears UCAORXIFG

...

RETI
```

Example 17-2 shows an extract of the interrupt service routine that handles data transmit interrupts from USCI A0 in either UART or SPI mode and the data transfer interrupts from USCI B0 in I<sup>2</sup>C mode.

### Example 17-2. Shared Transmit Interrupt Vectors Software Example

```
USCIA0_TX_USCIB0_I2C_DATA_ISR
  BIT.B #UCAOTXIFG, &IFG2 ; USCI_AO Transmit Interrupt?
  JNZ
          USCIAO_TX_ISR
USCIBO_I2C_DATA_ISR
  BIT.B #UCBORXIFG, &IFG2
  JNZ USCIB0_I2C_RX
USCIBO_I2C_TX
  ; Write UCBOTXBUF... - clears UCBOTXIFG
  RETI
USCIBO_I2C_RX
  ; Read UCBORXBUF... - clears UCBORXIFG
  . . .
  RETI
USCIAO TX ISR
  ; Write UCAOTXBUF ... - clears UCAOTXIFG
  RETI
```



www.ti.com USCI Registers: FC Mode

### 17.4 USCI Registers: I<sup>2</sup>C Mode

The USCI registers applicable in I<sup>2</sup>C mode for USCI\_B0 are listed in Table 17-2, and for USCI\_B1 in Table 17-3.

Table 17-2. USCI\_B0 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USCI_B0 control register 0	UCB0CTL0	Read/write	068h	001h with PUC
USCI_B0 control register 1	UCB0CTL1	Read/write	069h	001h with PUC
USCI_B0 bit rate control register 0	UCB0BR0	Read/write	06Ah	Reset with PUC
USCI_B0 bit rate control register 1	UCB0BR1	Read/write	06Bh	Reset with PUC
USCI_B0 I <sup>2</sup> C interrupt enable register	UCB0I2CIE	Read/write	06Ch	Reset with PUC
USCI_B0 status register	UCB0STAT	Read/write	06Dh	Reset with PUC
USCI_B0 receive buffer register	UCB0RXBUF	Read	06Eh	Reset with PUC
USCI_B0 transmit buffer register	UCB0TXBUF	Read/write	06Fh	Reset with PUC
USCI_B0 I <sup>2</sup> C own address register	UCB0I2COA	Read/write	0118h	Reset with PUC
USCI_B0 I <sup>2</sup> C slave address register	UCB0I2CSA	Read/write	011Ah	Reset with PUC
SFR interrupt enable register 2	IE2	Read/write	001h	Reset with PUC
SFR interrupt flag register 2	IFG2	Read/write	003h	00Ah with PUC

### NOTE: Modifying SFR bits

To avoid modifying control bits of other modules, it is recommended to set or clear the IEx and IFGx bits using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.

Table 17-3. USCI\_B1 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USCI_B1 control register 0	UCB1CTL0	Read/write	0D8h	Reset with PUC
USCI_B1 control register 1	UCB1CTL1	Read/write	0D9h	001h with PUC
USCI_B1 baud rate control register 0	UCB1BR0	Read/write	0DAh	Reset with PUC
USCI_B1 baud rate control register 1	UCB1BR1	Read/write	0DBh	Reset with PUC
USCI_B1 I <sup>2</sup> C interrupt enable register	UCB1I2CIE	Read/write	0DCh	Reset with PUC
USCI_B1 status register	UCB1STAT	Read/write	0DDh	Reset with PUC
USCI_B1 receive buffer register	UCB1RXBUF	Read	0DEh	Reset with PUC
USCI_B1 transmit buffer register	UCB1TXBUF	Read/write	0DFh	Reset with PUC
USCI_B1 I <sup>2</sup> C own address register	UCB1I2COA	Read/write	017Ch	Reset with PUC
USCI_B1 I <sup>2</sup> C slave address register	UCB1I2CSA	Read/write	017Eh	Reset with PUC
USCI_A1/B1 interrupt enable register	UC1IE	Read/write	006h	Reset with PUC
USCI_A1/B1 interrupt flag register	UC1IFG	Read/write	007h	00Ah with PUC



## 17.4.1 UCBxCTL0, USCI\_Bx Control Register 0

7	6		5	4	3	2	1	0
UCA10	UCSLA	10	UCMM	Unused	UCMST	UCMOI	DEx=11	UCSYNC=1
rw-0	rw-0	·	rw-0	rw-0	rw-0	rw-0	rw-0	r-1
UCA10	Bit 7	Own	addressing mod	le select				
		0	Own addres	s is a 7-bit addres	SS			
		1	Own addres	s is a 10-bit addre	ess			
UCSLA10	Bit 6	Slave	addressing mo	de select				
		0	Address slav	e with 7-bit addre	ess			
		1	Address slav	e with 10-bit add	ress			
UCMM	Bit 5	Multi-	master environ	ment select				
		0	Single maste disabled.	er environment. T	here is no other ma	aster in the syster	n. The address o	compare unit is
		1	Multi-master	environment				
Unused	Bit 4	Unus	ed					
UCMST	Bit 3				oses arbitration in a d the module acts		vironment (UCMI	M = 1) the
		0	Slave mode					
		1	Master mode	e				
<b>UCMODE</b> x	Bits 2-1	USCI	Mode. The UC	MODEx bits selec	t the synchronous	mode when UCS	YNC = 1.	
		00	3-pin SPI					
		01	4-pin SPI (m	aster/slave enabl	ed if STE = 1)			
		10	4-pin SPI (m	aster/slave enabl	ed if STE = 0)			
		11	I <sup>2</sup> C mode					
UCSYNC	Bit 0	Sync	hronous mode e	enable				
		0	Asynchronou	us mode				
		1	Synchronous	s mode				



www.ti.com USCI Registers: fC Mode

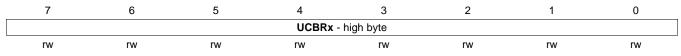
# 17.4.2 UCBxCTL1, USCI\_Bx Control Register 1

7	6		5	4	3	2	1	0
UCS	SSELx		Unused	UCTR	UCTXNACK	UCTXSTP	UCTXSTT	UCSWRST
rw-0	rw-0		rO	rw-0	rw-0	rw-0	rw-0	rw-1
UCSSELx	Bits 7-6	USCI	USCI clock source select. These bits select the BRCLK source clock.					
		00	UCLKI					
		01	ACLK					
		10	SMCLK					
		11	SMCLK					
Unused	Bit 5	Unus	ed					
UCTR	Bit 4	Trans	mitter/receiver					
		0	Receiver					
		1	Transmitter					
UCTXNACK	Bit 3	Trans	smit a NACK. U	CTXNACK is auto	matically cleared a	after a NACK is tra	ansmitted.	
		0	Acknowledge	e normally				
		1	Generate NA	ACK				
UCTXSTP	Bit 2				ode. Ignored in slav TXSTP is automati			
		0	No STOP ge	enerated				
		1	Generate ST	OP				
UCTXSTT	Bit 1	STAR	RT condition is p	receded by a NA	node. Ignored in sla CK. UCTXSTT is a ored in slave mode	utomatically clear		
		0	Do not gene	rate START cond	ition			
		1	Generate ST	ART condition				
UCSWRST	Bit 0	Softw	are reset enable	е				
		0	Disabled. US	SCI reset released	d for operation.			
		1	Enabled. US	CI logic held in re	eset state.			

# 17.4.3 UCBxBR0, USCI\_Bx Baud Rate Control Register 0

7	6	5	4	3	2	1	0	
UCBRx - low byte								
rw rw rw rw rw rw rw								

# 17.4.4 UCBxBR1, USCI\_Bx Baud Rate Control Register 1



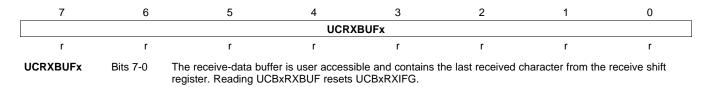
UCBRx Bit clock prescaler setting. The 16-bit value of (UCBxBR0 + UCBxBR1 x 256) forms the prescaler value.



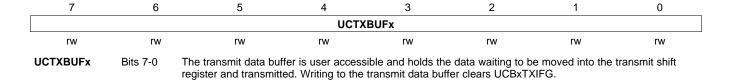
# 17.4.5 UCBxSTAT, USCI\_Bx Status Register

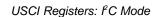
7	6	5	4	3	2	1	0
Unused	UCSCLLO	W UCGC	UCBBUSY	UCNACKIFG	UCSTPIFG	UCSTTIFG	UCALIFG
rw-0	r-0	rw-0	r-0	rw-0	rw-0	rw-0	rw-0
Unused	Bit 7	Unused.					
UCSCLLOW	Bit 6	SCL low					
		0 SCL is not	held low				
		1 SCL is held	l low				
UCGC	Bit 5	General call addres	s received. UCGC	is automatically cle	eared when a STA	ART condition is re	ceived.
		0 No general	call address receive	ved			
		1 General ca	ll address received	l			
UCBBUSY	Bit 4	Bus busy					
		0 Bus inactive	Э				
		1 Bus busy					
UCNACKIFG	Bit 3	Not-acknowledge re received.	ceived interrupt fla	g. UCNACKIFG is	automatically clea	ared when a STAF	RT condition is
		0 No interrup	t pending				
		1 Interrupt pe	ending				
UCSTPIFG	Bit 2	Stop condition inter	upt flag. UCSTPIF	G is automatically	cleared when a S	TART condition is	received.
		0 No interrup	t pending				
		1 Interrupt per	ending				
UCSTTIFG	Bit 1	Start condition inter	rupt flag. UCSTTIF	G is automatically	cleared if a STOF	condition is recei	ved.
		0 No interrup	t pending				
		1 Interrupt pe	ending				
UCALIFG	Bit 0	Arbitration lost inter	rupt flag				
		0 No interrup	t pending				
		1 Interrupt pe	ending				

# 17.4.6 UCBxRXBUF, USCI\_Bx Receive Buffer Register



# 17.4.7 UCBxTXBUF, USCI\_Bx Transmit Buffer Register









# 17.4.8 UCBxI2COA, USCIBx fC Own Address Register

15	14	13	12	11	10	9	8		
UCGCEN	0	0	0	0	0	I2C	OAx		
rw-0	r0	r0	r0	r0	r0	rw-0	rw-0		
7	6	5	4	3	2	1	0		
I2COAx									
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0		
UCGCEN	Bit 15	General call respon	se enable						
		0 Do not resp	ond to a general	call					
		1 Respond to	a general call						
I2COAx	I2COAx  Bits 9-0  I <sup>2</sup> C own address. The I2COAx bits contain the local address of the USCI_Bx I <sup>2</sup> C controller. The address right-justified. In 7-bit addressing mode, bit 6 is the MSB, and bits 9-7 are ignored. In 10-bit addressing mode, bit 9 is the MSB.								

# 17.4.9 UCBxl2CSA, USCI\_Bx FC Slave Address Register

15	14	13	12	11	10	9	8
0	0	0	0	0	0	I2C	SAx
r0	r0	r0	r0	r0	r0	rw-0	rw-0
7	6	5	4	3	2	1	0
			I2C	SAx			
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
I2CSAx Bits 9-0 I <sup>2</sup> C slave address. The I2CSAx bits contain the slave address of the external device to be addressed by the							

I<sup>2</sup>C slave address. The I2CSAx bits contain the slave address of the external device to be addressed by the USCI\_Bx module. It is only used in master mode. The address is right-justified. In 7-bit slave addressing mode, bit 6 is the MSB, and bits 9-7 are ignored. In 10-bit slave addressing mode, bit 9 is the MSB.

# 17.4.10 UCBxl2CIE, USCI\_Bx fC Interrupt Enable Register

7	6	5	4	3	2	1	0
		Reserved		UCNACKIE	UCSTPIE	UCSTTIE	UCALIE
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
Reserved	Bits 7-4	Reserved					
UCNACKIE	Bit 3	Not-acknowledge inte	rrupt enable				
		0 Interrupt disa	bled				
		<ol> <li>Interrupt enal</li> </ol>	oled				
UCSTPIE	Bit 2	Stop condition interrup	ot enable				
		0 Interrupt disa	bled				
		<ol> <li>Interrupt enal</li> </ol>	oled				
UCSTTIE	Bit 1	Start condition interru	ot enable				
		0 Interrupt disa	bled				
		<ol> <li>Interrupt enal</li> </ol>	oled				
UCALIE	Bit 0	Arbitration lost interru	ot enable				
		0 Interrupt disa	bled				
		1 Interrupt enal	oled				



# 17.4.11 IE2, Interrupt Enable Register 2

7	6	5	4	3	2	1	0
				UCB0TXIE	UCB0RXIE		
				rw-0	rw-0		
	Bits 7-4	These bits may be u	ised by other modi	ules (see the devi	ce-specific data sh	eet).	
<b>UCB0TXIE</b>	Bit 3	USCI_B0 transmit in	nterrupt enable				
		0 Interrupt dis	sabled				
		1 Interrupt en	abled				
<b>UCB0RXIE</b>	Bit 2	USCI_B0 receive in	terrupt enable				
		0 Interrupt dis	sabled				
		1 Interrupt er	abled				
	Bits 1-0	These bits may be u	ised by other modi	ules (see the devi	ce-specific data sh	eet).	

# 17.4.12 IFG2, Interrupt Flag Register 2

7	6	5	4	3	2	1	0			
				UCB0TXIFG	UCB0RXIFG					
				rw-1	rw-0					
	Bits 7-4 These bits may be used by other modules (see the device-specific data sheet).									
<b>UCB0TXIFG</b>	Bit 3	USCI_B0 transmit interrupt flag. UCB0TXIFG is set when UCB0TXBUF is empty.								
		0 No interru	0 No interrupt pending							
		1 Interrupt p	ending							
<b>UCB0RXIFG</b>	Bit 2	USCI_B0 receive	nterrupt flag. UCB0	RXIFG is set when	UCB0RXBUF ha	s received a com	plete character.			
		0 No interru	pt pending							
		1 Interrupt p	ending							
	Bits 1-0	These bits may be	used by other mod	ules (see the devic	ce-specific data sh	neet).				

# 17.4.13 UC1IE, USCI\_B1 Interrupt Enable Register

7	6	5	4	3	2	1	0
		Unused		UCB1TXIE	UCB1RXIE		
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0		
Unused	Bits 7-4	Unused					
UCB1TXIE	Bit 3	USCI_B1 transmit int	errupt enable				
		0 Interrupt disa	abled				
		1 Interrupt ena	bled				
UCB1RXIE	Bit 2	USCI_B1 receive inte	errupt enable				
		0 Interrupt disa	abled				
		1 Interrupt ena	bled				
	Bits 1-0	These bits may be us	sed by other USO	CI modules (see the	e device-specific da	ta sheet).	



www.ti.com USCI Registers: FC Mode

# 17.4.14 UC1IFG, USCI\_B1 Interrupt Flag Register

7	6	5	4	3	2	1	0	
		Unused		UCB1TXIFG	UCB1RXIFG			
rw-0	rw-0	rw-0	rw-0	rw-1	rw-0			
Unused	Bits 7-4	Unused.						
UCB1TXIFG	Bit 3	USCI_B1 transmit interrupt flag. UCB1TXIFG is set when UCB1TXBUF is empty.  No interrupt pending  Interrupt pending						
UCB1RXIFG	Bit 2	USCI_B1 receive interrupt flag. UCB1RXIFG is set when UCB1RXBUF has received a complete character.  No interrupt pending  Interrupt pending  These bits may be used by other modules (see the device-specific data sheet).						



# USART Peripheral Interface, UART Mode

The universal synchronous/asynchronous receive/transmit (USART) peripheral interface supports two serial modes with one hardware module. This chapter discusses the operation of the asynchronous UART mode. USART0 is implemented on the MSP430AFE2xx devices.

Topic Page

18.1	USART Introduction: UART Mode	475
18.2	USART Operation: UART Mode	476
18.3	USART Registers: UART Mode	490



#### 18.1 USART Introduction: UART Mode

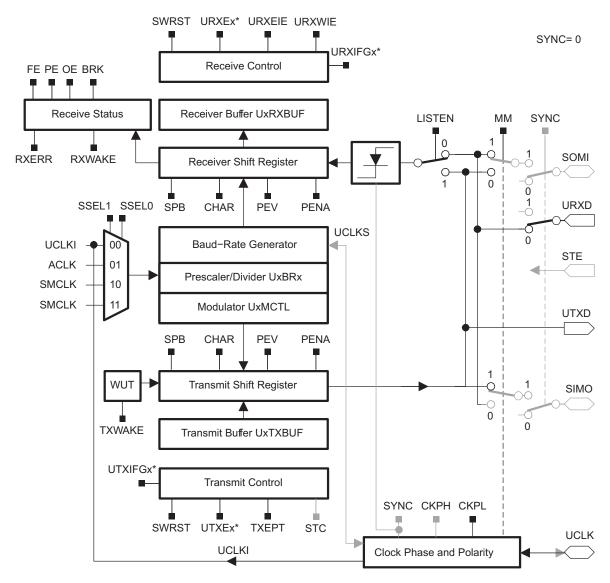
In asynchronous mode, the USART connects the MSP430 to an external system via two external pins, URXD and UTXD. UART mode is selected when the SYNC bit is cleared.

**UART** mode features include:

- 7- or 8-bit data with odd parity, even parity, or non-parity
- Independent transmit and receive shift registers
- · Separate transmit and receive buffer registers
- · LSB-first data transmit and receive
- · Built-in idle-line and address-bit communication protocols for multiprocessor systems
- Receiver start-edge detection for auto-wake up from LPMx modes
- Programmable baud rate with modulation for fractional baud rate support
- Status flags for error detection and suppression and address detection
- · Independent interrupt capability for receive and transmit

Figure 18-1 shows the USART when configured for UART mode.





<sup>\*</sup> See the device-specific data sheet for SFR locations.

Figure 18-1. USART Block Diagram: UART Mode

### 18.2 USART Operation: UART Mode

In UART mode, the USART transmits and receives characters at a bit rate asynchronous to another device. Timing for each character is based on the selected baud rate of the USART. The transmit and receive functions use the same baud rate frequency.

### 18.2.1 USART Initialization and Reset

The USART is reset by a PUC or by setting the SWRST bit. After a PUC, the SWRST bit is automatically set, keeping the USART in a reset condition. When set, the SWRST bit resets the URXIEx, UTXIEx, URXIFGx, RXWAKE, TXWAKE, RXERR, BRK, PE, OE, and FE bits and sets the UTXIFGx and TXEPT bits. The receive and transmit enable flags, URXEx and UTXEx, are not altered by SWRST. Clearing SWRST releases the USART for operation. See also chapter *USART Module*, *PC mode* for USART0 when reconfiguring from I<sup>2</sup>C mode to UART mode.



### NOTE: Initializing or Reconfiguring the USART Module

The required USART initialization/reconfiguration process is:

- 1. Set SWRST (BIS.B #SWRST,&UxCTL)
- 2. Initialize all USART registers with SWRST = 1 (including UxCTL)
- 3. Enable USART module via the MEx SFRs (URXEx and/or UTXEx)
- 4. Clear SWRST via software (BIC.B #SWRST,&UxCTL)
- 5. Enable interrupts (optional) via the IEx SFRs (URXIEx and/or UTXIEx)

Failure to follow this process may result in unpredictable USART behavior.

#### 18.2.2 Character Format

The UART character format, shown in Figure 18-2, consists of a start bit, seven or eight data bits, an even/odd/no parity bit, an address bit (address-bit mode), and one or two stop bits. The bit period is defined by the selected clock source and setup of the baud rate registers.

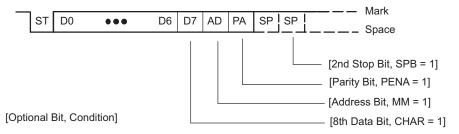


Figure 18-2. Character Format

#### 18.2.3 Asynchronous Communication Formats

When two devices communicate asynchronously, the idle-line format is used for the protocol. When three or more devices communicate, the USART supports the idle-line and address-bit multiprocessor communication formats.

#### 18.2.3.1 Idle-Line Multiprocessor Format

When MM = 0, the idle-line multiprocessor format is selected. Blocks of data are separated by an idle time on the transmit or receive lines as shown in Figure 18-3. An idle receive line is detectedwhen 10 or more continuous ones (marks) are received after the first stop bit of a character. When two stop bits are used for the idle line the second stop bit is counted as the first mark bit of the idle period.

The first character received after an idle period is an address character. The RXWAKE bit is used as an address tag for each block of characters. In the idle-line multiprocessor format, this bit is set when a received character is an address and is transferred to UxRXBUF.

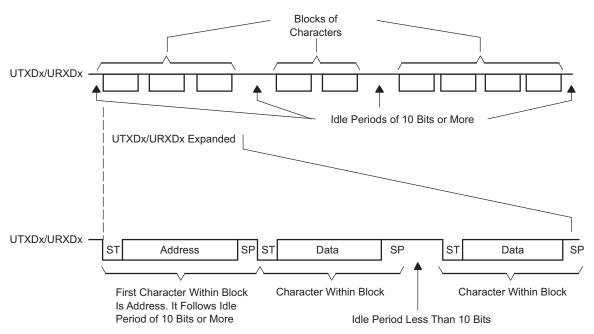


Figure 18-3. Idle-Line Format

The URXWIE bit is used to control data reception in the idle-line multiprocessor format. When the URXWIE bit is set, all non-address characters are assembled but not transferred into the UxRXBUF, and interrupts are not generated. When an address character is received, the receiver is temporarily activated to transfer the character to UxRXBUF and sets the URXIFGx interrupt flag. Any applicable error flag is also set. The user can then validate the received address.

If an address is received, user software can validate the address and must reset URXWIE to continue receiving data. If URXWIE remains set, only address characters are received. The URXWIE bit is not modified by the USART hardware automatically.

For address transmission in idle-line multiprocessor format, a precise idle period can be generated by the USART to generate address character identifiers on UTXDx. The wake-up temporary (WUT) flag is an internal flag double-buffered with the user-accessible TXWAKE bit. When the transmitter is loaded from UxTXBUF, WUT is also loaded from TXWAKE resetting the TXWAKE bit.

The following procedure sends out an idle frame to indicate an address character follows:

- Set TXWAKE, then write any character to UxTXBUF. UxTXBUF must be ready for new data (UTXIFGx = 1).
  - The TXWAKE value is shifted to WUT and the contents of UxTXBUF are shifted to the transmit shift register when the shift register is ready for new data. This sets WUT, which suppresses the start, data, and parity bits of a normal transmission, then transmits an idle period of exactly 11 bits. When two stop bits are used for the idle line, the second stop bit is counted as the first mark bit of the idle period. TXWAKE is reset automatically.
- 2. Write desired address character to UxTXBUF. UxTXBUF must be ready for new data (UTXIFGx = 1). The new character representing the specified address is shifted out following the address-identifying idle period on UTXDx. Writing the first "don't care" character to UxTXBUF is necessary in order to shift the TXWAKE bit to WUT and generate an idle-line condition. This data is discarded and does not appear on UTXDx.

#### 18.2.3.2 Address-Bit Multiprocessor Format

When MM = 1, the address-bit multiprocessor format is selected. Each processed character contains an extra bit used as an address indicator shown in Figure 18-4. The first character in a block of characters carries a set address bit which indicates that the character is an address. The USART RXWAKE bit is set when a received character is a valid address character and is transferred to UxRXBUF.



The URXWIE bit is used to control data reception in the address-bit multiprocessor format. If URXWIE is set, data characters (address bit = 0) are assembled by the receiver but are not transferred to UxRXBUF and no interrupts are generated. When a character containing a set address bit is received, the receiver is temporarily activated to transfer the character to UxRXBUF and set URXIFGx. All applicable error status flags are also set.

If an address is received, user software must reset URXWIE to continue receiving data. If URXWIE remains set, only address characters (address bit = 1) are received. The URXWIE bit is not modified by the USART hardware automatically.

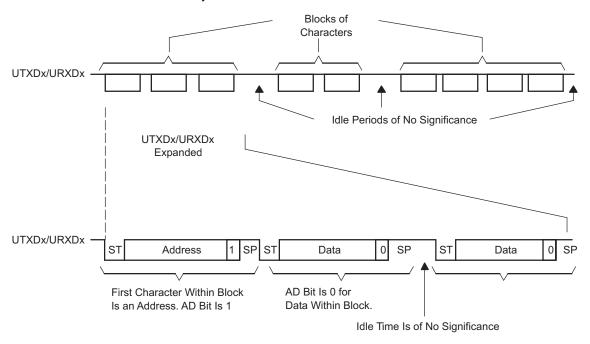


Figure 18-4. Address-Bit Multiprocessor Format

For address transmission in address-bit multiprocessor mode, the address bit of a character can be controlled by writing to the TXWAKE bit. The value of the TXWAKE bit is loaded into the address bit of the character transferred from UxTXBUF to the transmit shift register, automatically clearing the TXWAKE bit. TXWAKE must not be cleared by software. It is cleared by USART hardware after it is transferred to WUT or by setting SWRST.

#### 18.2.3.3 Automatic Error Detection

Glitch suppression prevents the USART from being accidentally started. Any low-level on URXDx shorter than the deglitch time  $t_{\tau}$  (approximately 300 ns) is ignored. See the device-specific data sheet for parameters.

When a low period on URXDx exceeds  $t_{\tau}$  a majority vote is taken for the start bit. If the majority vote fails to detect a valid start bit the USART halts character reception and waits for the next low period on URXDx. The majority vote is also used for each bit in a character to prevent bit errors.

The USART module automatically detects framing errors, parity errors, overrun errors, and break conditions when receiving characters. The bits FE, PE, OE, and BRK are set when their respective condition is detected. When any of these error flags are set, RXERR is also set. The error conditions are described in Table 18-1.



Table	18-1	Receive	Frror	Conditions

Error Condition	Description					
Framing error	A framing error occurs when a low stop bit is detected. When two stop bits are used, only the first stop bit is checked for framing error. When a framing error is detected, the FE bit is set.					
Parity error	A parity error is a mismatch between the number of 1s in a character and the value of the parity bit. When an address bit is included in the character, it is included in the parity calculation. When a parity error is detected, the PE bit is set.					
Receive overrun error	An overrun error occurs when a character is loaded into UxRXBUF before the prior character has been read. When an overrun occurs, the OE bit is set.					
Break condition	A break condition is a period of 10 or more low bits received on URXDx after a missing stop bit. When a break condition is detected, the BRK bit is set. A break condition can also set the interrupt flag URXIFGx when URXEIE = 0.					

When URXEIE = 0 and a framing error, parity error, or break condition is detected, no character is received into UxRXBUF. When URXEIE = 1, characters are received into UxRXBUF and any applicable error bit is set.

When any of the FE, PE, OE, BRK, or RXERR bits are set, the bit remains set until user software resets it or UxRXBUF is read.

#### 18.2.4 USART Receive Enable

The receive enable bit, URXEx, enables or disables data reception on URXDx as shown in Figure 18-5. Disabling the USART receiver stops the receive operation following completion of any character currently being received or immediately if no receive operation is active. The receive-data buffer, UxRXBUF, contains the character moved from the RX shift register after the character is received.

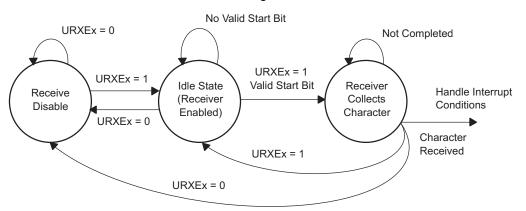


Figure 18-5. State Diagram of Receiver Enable

### NOTE: Re-Enabling the Receiver (Setting URXEx): UART Mode

When the receiver is disabled (URXEx = 0), re-enabling the receiver (URXEx = 1) is asynchronous to any data stream that may be present on URXDx at the time. Synchronization can be performed by testing for an idle line condition before receiving a valid character (see URXWIE).

#### 18.2.5 USART Transmit Enable

When UTXEx is set, the UART transmitter is enabled. Transmission is initiated by writing data to UxTXBUF. The data is then moved to the transmit shift register on the next BITCLK after the TX shift register is empty, and transmission begins. This process is shown in Figure 18-6.

When the UTXEx bit is reset the transmitter is stopped. Any data moved to UxTXBUF and any active transmission of data currently in the transmit shift register prior to clearing UTXEx continue until all data transmission is completed.



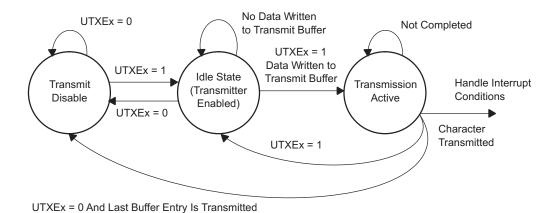


Figure 18-6. State Diagram of Transmitter Enable

When the transmitter is enabled (UTXEx = 1), data should not be written to UxTXBUF unless it is ready for new data indicated by UTXIFGx = 1. Violation can result in an erroneous transmission if data in UxTXBUF is modified as it is being moved into the TX shift register.

It is recommended that the transmitter be disabled (UTXEx = 0) only after any active transmission is complete. This is indicated by a set transmitter empty bit (TXEPT = 1). Any data written to UxTXBUF while the transmitter is disabled are held in the buffer but are not moved to the transmit shift register or transmitted. Once UTXEx is set, the data in the transmit buffer is immediately loaded into the transmit shift register and character transmission resumes.

### 18.2.6 USART Baud Rate Generation

The USART baud rate generator is capable of producing standard baud rates from non-standard source frequencies. The baud rate generator uses one prescaler/divider and a modulator as shown in Figure 18-7. This combination supports fractional divisors for baud rate generation. The maximum USART baud rate is one-third the UART source clock frequency BRCLK.

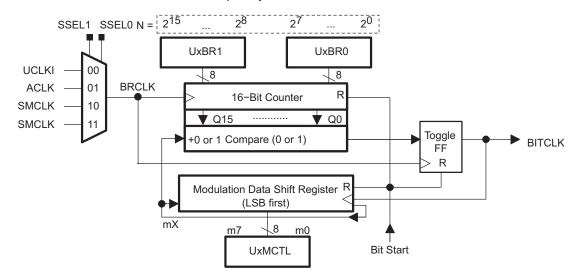


Figure 18-7. MSP430 Baud Rate Generator

Timing for each bit is shown in Figure 18-8. For each bit received, a majority vote is taken to determine the bit value. These samples occur at the N/2-1, N/2, and N/2+1 BRCLK periods, where N is the number of BRCLKs per BITCLK.

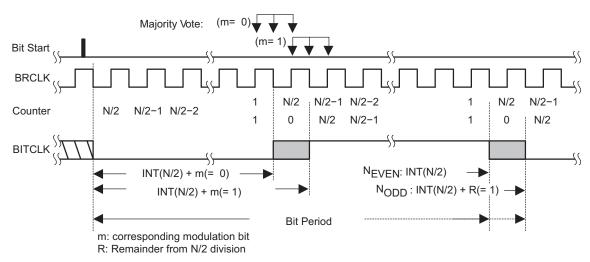


Figure 18-8. BITCLK Baud Rate Timing

### 18.2.6.1 Baud Rate Bit Timing

The first stage of the baud rate generator is the 16-bit counter and comparator. At the beginning of each bit transmitted or received, the counter is loaded with INT(N/2) where N is the value stored in the combination of UxBR0 and UxBR1. The counter reloads INT(N/2) for each bit period half-cycle, giving a total bit period of N BRCLKs. For a given BRCLK clock source, the baud rate used determines the required division factor N:

$$N = \frac{BRCLK}{Baud Rate}$$

The division factor N is often a non-integer value of which the integer portion can be realized by the prescaler/divider. The second stage of the baud rate generator, the modulator, is used to meet the fractional part as closely as possible. The factor N is then defined as:

N = UxBR + 
$$\frac{1}{n}\sum_{i=0}^{n-1} m_i$$

Where,

N = Target division factor

UxBR = 16-bit representation of registers UxBR0 and UxBR1

i = Bit position in the character

n = Total number of bits in the character

 $m_i$  = Data of each corresponding modulation bit (1 or 0)

Baud rate = 
$$\frac{BRCLK}{N} + \frac{BRCLK}{UxBR + \frac{1}{n} \sum_{i=0}^{n-1} m_i}$$

The BITCLK can be adjusted from bit to bit with the modulator to meet timing requirements when a non-integer divisor is needed. Timing of each bit is expanded by one BRCLK clock cycle if the modulator bit m<sub>i</sub> is set. Each time a bit is received or transmitted, the next bit in the modulation control register determines the timing for that bit. A set modulation bit increases the division factor by one while a cleared modulation bit maintains the division factor given by UxBR.

The timing for the start bit is determined by UxBR plus m0, the next bit is determined by UxBR plus m1, and so on. The modulation sequence begins with the LSB. When the character is greater than 8 bits, the modulation sequence restarts with m0 and continues until all bits are processed.



#### 18.2.6.2 Determining the Modulation Value

Determining the modulation value is an interactive process. Using the timing error formula provided, beginning with the start bit, the individual bit errors are calculated with the corresponding modulator bit set and cleared. The modulation bit setting with the lower error is selected and the next bit error is calculated. This process is continued until all bit errors are minimized. When a character contains more than 8 bits, the modulation bits repeat. For example, the ninth bit of a character uses modulation bit 0.

#### 18.2.6.3 Transmit Bit Timing

The timing for each character is the sum of the individual bit timings. By modulating each bit, the cumulative bit error is reduced. The individual bit error can be calculated by:

Error [%] = 
$$\left\{ \frac{\text{baud rate}}{\text{BRCLK}} \times \left[ (j+1) \times \text{UxBR} + \sum_{i=0}^{j} m_i \right] - (j+1) \right\} \times 100\%$$

Where,

baud rate = Desired baud rate

BRCLK = Input frequency - UCLKI, ACLK, or SMCLK

j = Bit position - 0 for the start bit, 1 for data bit D0, and so on

UxBR = Division factor in registers UxBR1 and UxBR0

For example, the transmit errors for the following conditions are calculated:

Baud rate = 2400

BRCLK = 32 768 Hz (ACLK)

UxBR = 13, since the ideal division factor is 13.65

UxMCTL = 6Bh: m7 = 0, m6 = 1, m5 = 1, m4 = 0, m3 = 1, m2 = 0, m1 = 1, and m0 = 1. The LSB of

UxMCTL is used first.



Start bit Error [%]= 
$$\binom{\text{baud rate}}{\text{BRCLK}}$$
.  $((0+1)\cdot \text{UxBR}+1)-1$ .  $100\%=2.54\%$ 

Data bit D0 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((1+1)\cdot \text{UxBR}+2)-2$ .  $100\%=5.08\%$ 

Data bit D1 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((2+1)\cdot \text{UxBR}+2)-3$ .  $100\%=0.29\%$ 

Data bit D2 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((3+1)\cdot \text{UxBR}+3)-4$ .  $100\%=2.83\%$ 

Data bit D3 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((4+1)\cdot \text{UxBR}+3)-5$ .  $100\%=-1.95\%$ 

Data bit D4 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((5+1)\cdot \text{UxBR}+4)-6$ .  $100\%=0.59\%$ 

Data bit D5 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((6+1)\cdot \text{UxBR}+5)-7$ .  $100\%=3.13\%$ 

Data bit D6 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((6+1)\cdot \text{UxBR}+5)-8$ .  $100\%=-1.66\%$ 

Data bit D7 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((6+1)\cdot \text{UxBR}+6)-9$ .  $100\%=0.88\%$ 

Parity bit Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((9+1)\cdot \text{UxBR}+7)-10$ .  $100\%=3.42\%$ 

Stop bit 1 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$ .  $((10+1)\cdot \text{UxBR}+7)-11$ .  $100\%=-1.37\%$ 

The results show the maximum per-bit error to be 5.08% of a BITCLK period.



#### 18.2.6.4 Receive Bit Timing

Receive timing is subject to two error sources. The first is the bit-to-bit timing error. The second is the error between a start edge occurring and the start edge being accepted by the USART. Figure 18-9 shows the asynchronous timing errors between data on the URXDx pin and the internal baud-rate clock.

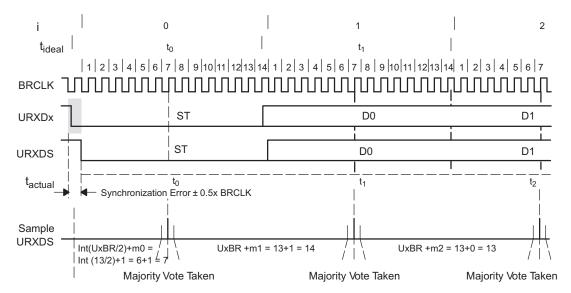


Figure 18-9. Receive Error

The ideal start bit timing  $t_{ideal(0)}$  is half the baud-rate timing  $t_{baudrate}$ , because the bit is tested in the middle of its period. The ideal baud-rate timing  $t_{ideal(i)}$  for the remaining character bits is the baud rate timing  $t_{baudrate}$ . The individual bit errors can be calculated by:

Error [%] = 
$$\left\{ \frac{\text{baud rate}}{\text{BRCLK}} \times \left( 2 \times \left[ \text{m0 + int} \left( \frac{\text{UxBR}}{2} \right) \right] + \left[ i \times \text{UxBR} + \sum_{i=1}^{j} m_i \right] \right) - 1 - j \right\} \times 100\%$$

Where.

baud rate = the required baud rate

BRCLK = the input frequency; selected for UCLK, ACLK, or SMCLK

j = 0 for the start bit, 1 for data bit D0, and so on

UxBR = the division factor in registers UxBR1 and UxBR0

For example, the receive errors for the following conditions are calculated:

Baud rate = 2400

BRCLK = 32 768 Hz (ACLK)

UxBR = 13, since the ideal division factor is 13.65

UxMCTL = 6B: m7 = 0, m6 = 1, m5 = 1, m4 = 0, m3 = 1, m2 = 0, m1 = 1 and m0 = 1. The LSB of UxMCTL is used first.



Data bit D1 Error [%]= 
$$\binom{\text{baud rate}}{\text{BRCLK}}$$
 [2x(1+6)+2·UxBR+1]-1-2)·100%=0.29%

Data bit D2 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+3·UxBR+2]-1-3)·100%=2.83%

Data bit D3 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+4·UxBR+2]-1-4)·100%=-1.95%

Data bit D4 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+5·UxBR+3]-1-5)·100%=0.59%

Data bit D5 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+6·UxBR+4]-1-6)·100%=3.13%

Data bit D6 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+7·UxBR+4]-1-7)·100%=-1.66%

Data bit D7 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+8·UxBR+5]-1-8)·100%=0.88%

Parity bit Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+9·UxBR+6]-1-9)·100%=3.42%

Stop bit 1 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+10·UxBR+6]-1-10)·100%=2.54%

Data bit D0 Error [%]=  $\binom{\text{baud rate}}{\text{BRCLK}}$  [2x(1+6)+0·UxBR+0]-1-0)·100%=2.54%

The results show the maximum per-bit error to be 5.08% of a BITCLK period.

### 18.2.6.5 Typical Baud Rates and Errors

Standard baud rate frequency data for UxBRx and UxMCTL are listed in Table 18-2 for a 32 768-Hz watch crystal (ACLK) and a typical 1 048 576-Hz SMCLK.

The receive error is the accumulated time versus the ideal scanning time in the middle of each bit. The transmit error is the accumulated timing error versus the ideal time of the bit period.

Divide by A: BRCLK = 32 768 Hz B: BRCLK = 1 048 576 Hz Synch **Baud Rate** Max TX Max RX Max TX Max RX UxBR1 UxBR0 UxMCTL UxBR1 UxBR0 UxMCTL A: B: **RX Error** Error % Error % Error % Error % -4/3 ±2 69 FF 0/0.3 1200 27.31 873.81 0 1B 03 -4/3 03 ±2 0/0.3 01 FF 2400 13.65 436.91 0 0D 6B -6/3-6/3±4 **B4** ±2 4800 6.83 218.45 0 06 6F -9/11 -9/11 0 DA 55 0/0.4 ±7 ±2 9600 3.41 109.23 0 03 4A -21/12 -21/12 ±15 0 6D 03 -0.4/1 ±2 19 200 54.61 0 36 6B -0.2/2+2 38 400 27.31 0 1B 03 -4/3 ±2 76 800 13.65 0 0D 6B -6/3 ±4

Table 18-2. Commonly Used Baud Rates, Baud Rate Data, and Errors

9.1

115 200

-5/7

±7



#### 18.2.7 USART Interrupts

The USART has one interrupt vector for transmission and one interrupt vector for reception.

#### 18.2.7.1 USART Transmit Interrupt Operation

The UTXIFGx interrupt flag is set by the transmitter to indicate that UxTXBUF is ready to accept another character. An interrupt request is generated if UTXIEx and GIE are also set. UTXIFGx is automatically reset if the interrupt request is serviced or if a character is written to UxTXBUF.

UTXIFGx is set after a PUC or when SWRST = 1. UTXIEx is reset after a PUC or when SWRST = 1. The operation is shown is Figure 18-10.

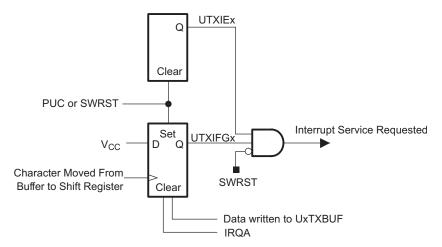


Figure 18-10. Transmit Interrupt Operation

### 18.2.7.2 USART Receive Interrupt Operation

The URXIFGx interrupt flag is set each time a character is received and loaded into UxRXBUF. An interrupt request is generated if URXIEx and GIE are also set. URXIFGx and URXIEx are reset by a system reset PUC signal or when SWRST = 1. URXIFGx is automatically reset if the pending interrupt is served (when URXSE = 0) or when UxRXBUF is read. The operation is shown in Figure 18-11.

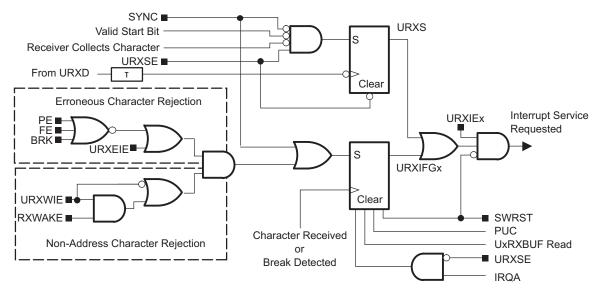


Figure 18-11. Receive Interrupt Operation



URXEIE is used to enable or disable erroneous characters from setting URXIFGx. When using multiprocessor addressing modes, URXWIE is used to auto-detect valid address characters and reject unwanted data characters.

Two types of characters do not set URXIFGx:

- Erroneous characters when URXEIE = 0
- Non-address characters when URXWIE = 1

When URXEIE = 1 a break condition sets the BRK bit and the URXIFGx flag.

#### 18.2.7.3 Receive-Start Edge Detect Operation

The URXSE bit enables the receive start-edge detection feature. The recommended usage of the receivestart edge feature is when BRCLK is sourced by the DCO and when the DCO is off because of low-power mode operation. The ultra-fast turn-on of the DCO allows character reception after the start edge detection.

When URXSE, URXIEx and GIE are set and a start edge occurs on URXDx, the internal signal URXS is set. When URXS is set, a receive interrupt request is generated but URXIFGx is not set. User software in the receive interrupt service routine can test URXIFGx to determine the source of the interrupt. When URXIFGx = 0 a start edge was detected, and when URXIFGx = 1 a valid character (or break) was received.

When the ISR determines the interrupt request was from a start edge, user software toggles URXSE, and must enable the BRCLK source by returning from the ISR to active mode or to a low-power mode where the source is active. If the ISR returns to a low-power mode where the BRCLK source is inactive, the character is not received. Toggling URXSE clears the URXS signal and re-enables the start edge detect feature for future characters. See chapter *System Resets, Interrupts, and Operating Modes* for information on entering and exiting low-power modes.

The now active BRCLK allows the USART to receive the balance of the character. After the full character is received and moved to UxRXBUF, URXIFGx is set and an interrupt service is again requested. Upon ISR entry, URXIFGx = 1 indicating a character was received. The URXIFGx flag is cleared when user software reads UxRXBUF.

#### NOTE: Break Detect With Halted UART Clock

When using the receive start-edge detect feature, a break condition cannot be detected when the BRCLK source is off.



#### 18.2.7.4 Receive-Start Edge Detect Conditions

When URXSE = 1, glitch suppression prevents the USART from being accidentally started. Any low-level on URXDx shorter than the deglitch time  $t_{\tau}$  (approximately 300 ns) is ignored by the USART and no interrupt request is generated (see Figure 18-12). See the device-specific data sheet for parameters.



Figure 18-12. Glitch Suppression, USART Receive Not Started

When a glitch is longer than  $t_{\tau}$  or a valid start bit occurs on URXDx, the USART receive operation is started and a majority vote is taken as shown in Figure 18-13. If the majority vote fails to detect a start bit, the USART halts character reception.

If character reception is halted, an active BRCLK is not necessary. A time-out period longer than the character receive duration can be used by software to indicate that a character was not received in the expected time, and the software can disable BRCLK.

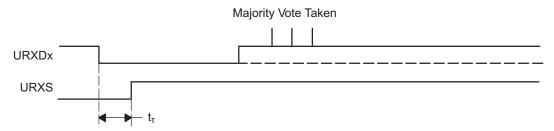


Figure 18-13. Glitch Suppression, USART Activated



### 18.3 USART Registers: UART Mode

Table 18-3 lists the registers for all devices implementing a USART module. Table 18-4 applies only to devices with a second USART module, USART1.

Table 18-3. USART0 Control and Status Registers

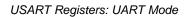
Register	Short Form	Register Type	Address	Initial State
USART control register	U0CTL	Read/write	070h	001h with PUC
Transmit control register	U0TCTL	Read/write	071h	001h with PUC
Receive control register	U0RCTL	Read/write	072h	000h with PUC
Modulation control register	U0MCTL	Read/write	073h	Unchanged
Baud rate control register 0	U0BR0	Read/write	074h	Unchanged
Baud rate control register 1	U0BR1	Read/write	075h	Unchanged
Receive buffer register	U0RXBUF	Read	076h	Unchanged
Transmit buffer register	U0TXBUF	Read/write	077h	Unchanged
SFR interrupt enable register 1	IE1	Read/write	000h	000h with PUC
SFR interrupt flag register 1	IFG1	Read/write	002h	082h with PUC

### Table 18-4. USART1 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USART control register	U1CTL	Read/write	078h	001h with PUC
Transmit control register	U1TCTL	Read/write	079h	001h with PUC
Receive control register	U1RCTL	Read/write	07Ah	000h with PUC
Modulation control register	U1MCTL	Read/write	07Bh	Unchanged
Baud rate control register 0	U1BR0	Read/write	07Ch	Unchanged
Baud rate control register 1	U1BR1	Read/write	07Dh	Unchanged
Receive buffer register	U1RXBUF	Read	07Eh	Unchanged
Transmit buffer register	U1TXBUF	Read/write	07Fh	Unchanged
SFR interrupt enable register 2	IE2	Read/write	001h	000h with PUC
SFR interrupt flag register 2	IFG2	Read/write	003h	020h with PUC

### NOTE: Modifying SFR bits

To avoid modifying control bits of other modules, it is recommended to set or clear the IEx and IFGx bits using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.







# 18.3.1 UxCTL, USART Control Register

7	6		5	4	3	2	1	0				
PENA	PEV	1	SPB	CHAR	LISTEN	SYNC	MM	SWRST				
rw-0	rw-C	)	rw-0	rw-0	rw-0	rw-0	rw-0	rw-1				
PENA	Bit 7	Parity	Parity enable									
		0	Parity disable	ed								
		1			enerated (UTXDx) a dress bit is included			-bit				
PEV	Bit 6	Parity	y select. PEV is	not used when p	arity is disabled.							
		0	Odd parity									
		1	1 Even parity									
SPB	Bit 5	Stop bit select. Number of stop bits transmitted. The receiver always checks for one stop bit.										
		0	One stop bit									
		1	Two stop bits	S								
CHAR	Bit 4	Char	acter length. Sel	lects 7-bit or 8-bit	t character length.							
		0	7-bit data									
		1	8-bit data									
LISTEN	Bit 3	Liste	n enable. The Ll	STEN bit selects	loopback mode.							
		0	Disabled									
		1	Enabled. UT	XDx is internally	fed back to the rec	eiver.						
SYNC	Bit 2	Sync	chronous mode e	enable								
		0	UART mode									
		1	SPI mode									
MM	Bit 1	Multi	processor mode	select								
		0	Idle-line mult	tiprocessor proto	col							
		1	Address-bit ı	multiprocessor pr	otocol							
SWRST	Bit 0	Softv	vare reset enable	е								
		0 Disabled. USART reset released for operation										
		1	Enabled. US	ART logic held in	reset state							



# 18.3.2 UxTCTL, USART Transmit Control Register

7	6		5	4	3	2	1	0
Unused	CKPL	SSEL		ELx	URXSE	TXWAKE	Unused	TXEPT
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-1
Unused	Bit 7	Unuse	ed					
CKPL	Bit 6	Clock	polarity select					
		0	UCLKI = UC	LK				
		1	UCLKI = inve	erted UCLK				
SSELx	Bits 5-4	Sourc	e select. These	bits select the B	RCLK source cloc	ck.		
		00	UCLKI					
		01	ACLK					
		10	SMCLK					
		11	SMCLK					
URXSE	Bit 3	UART	receive start-e	dge. The bit ena	bles the UART red	eive start-edge fea	ture.	
		0	Disabled					
		1	Enabled					
TXWAKE	Bit 2	Trans	mitter wake					
		0	Next frame t	ransmitted is dat	a			
		1	Next frame t	ransmitted is an	address			
Unused	Bit 1	Unuse	ed					
TXEPT	Bit 0	Trans	mitter empty fla	g				
		0	UART is tran	smitting data an	d/or data is waiting	g in UxTXBUF		
		1	Transmitter	shift register and	UxTXBUF are em	pty or SWRST = 1		





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# 18.3.3 UxRCTL, USART Receive Control Register

7	6		5	4	3	2	1	0		
FE	PE		OE	BRK	URXEIE	URXWIE	RXWAKE	RXERR		
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0		
FE	Bit 7	Frami	ng error flag							
		0	No error							
		1	Character re	ceived with low s	top bit					
PE	Bit 6	Parity	error flag. Whe	n PENA = 0, PE	is read as 0.					
		0	No error							
		1	Character re	ceived with parity	error					
OE	Bit 5	Overrun error flag. This bit is set when a character is transferred into UxRXBUF before the previous character was read.								
		0	No error							
		1	Overrun erro	r occurred						
BRK	Bit 4	Break	reak detect flag							
		0	No break cor	ndition						
		1	Break condit	ion occurred						
URXEIE	Bit 3	Recei	ve erroneous-cl	naracter interrupt	-enable					
		0	Erroneous ch	naracters rejected	d and URXIFGx is	not set				
		1	Erroneous ch	naracters receive	d set URXIFGx					
URXWIE	Bit 2				s bit enables URX ess character does					
		0	All received	characters set UF	RXIFGx					
		1	Only receive	d address charac	ters set URXIFGx					
RXWAKE	Bit 1	Recei	ve wake-up flag	1						
		0	Received cha	aracter is data						
		1	Received cha	aracter is an add	ress					
RXERR	Bit 0	Receive error flag. This bit indicates a character was received with error(s). When RXERR error flags (FE, PE, OE, BRK) is also set. RXERR is cleared when UxRXBUF is read.								
		0	No receive e	rrors detected						
		1	Receive erro	r detected						

# 18.3.4 UxBR0, USART Baud Rate Control Register 0

7	6	5	4	3	2	1	0
<b>2</b> <sup>7</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>2</sup>	<b>2</b> <sup>1</sup>	<b>2</b> <sup>0</sup>
rw							

# 18.3.5 UxBR1, USART Baud Rate Control Register 1

7	6	5	4	3	2	1	0
2 <sup>15</sup>	2 <sup>14</sup>	2 <sup>13</sup>	<b>2</b> <sup>12</sup>	2 <sup>11</sup>	2 <sup>10</sup>	<b>2</b> <sup>9</sup>	2 <sup>8</sup>
rw	rw	rw	rw	rw	rw	rw	rw

UxBRx

The valid baud-rate control range is 3 ≤ UxBR ≤ 0FFFFh, where UxBR = (UxBR1 + UxBR0). Unpredictable receive and transmit timing occurs if UxBR < 3.



18.3.6 Ux	MCTL, USAR	T Modulation	Control	Reaister
-----------	------------	--------------	---------	----------

7	6	5	4	3	2	1	0
m7	m6	m5	m4	m3	m2	m1	m0
rw							

UxMCTLx

Modulation bits. These bits select the modulation for BRCLK.

### 18.3.7 UxRXBUF, USART Receive Buffer Register

7	6	5	4	3	2	1	0
<b>2</b> <sup>7</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>2</sup>	<b>2</b> <sup>1</sup>	<b>2</b> <sup>0</sup>
_		_			_		

UxRXBUFx Bits 7-0

The receive-data buffer is user accessible and contains the last received character from the receive shift register. Reading UxRXBUF resets the receive-error bits, the RXWAKE bit, and URXIFGx. In 7-bit data mode, UxRXBUF is LSB justified and the MSB is always reset.

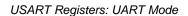
# 18.3.8 UxTXBUF, USART Transmit Buffer Register

7	6	5	4	3	2	1	0
27	<b>2</b> <sup>6</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>2</sup>	<b>2</b> <sup>1</sup>	<b>2</b> <sup>0</sup>
rw	rw	rw	rw	rw	rw	rw	rw

UxTXBUFx

Bits 7-0

The transmit data buffer is user accessible and holds the data waiting to be moved into the transmit shift register and transmitted on UTXDx. Writing to the transmit data buffer clears UTXIFGx. The MSB of UxTXBUF is not used for 7-bit data and is reset.





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# 18.3.9 IE1, Interrupt Enable Register 1

7	6	5	4	3	2	1	0
UTXIE0	URXIE	)					
rw-0	rw-0						
UTXIE0	Bit 7	USART0 transmit in 0 Interrupt no	t enabled	s bit enables the l	JTXIFG0 interrupt		
URXIE0	Bit 6	<ul> <li>Interrupt enabled</li> <li>USART0 receive interrupt enable. This bit enables the URXIFG0 interrupt.</li> <li>Interrupt not enabled</li> </ul>					
	Bits 5-0	1 Interrupt en These bits may be u		ıles. See device-s	pecific data sheet		

# 18.3.10 IE2, Interrupt Enable Register 2

7	6	5	4	3	2	1	0				
		UTXIE1	URXIE1								
		rw-0	rw-0								
	Bits 7-6	These bits may be used by other modules. See device-specific data sheet.									
UTXIE1	<b>TXIE1</b> Bit 5 USART1 transmit interrupt enable. This bit enables the UTXIFG1 interrupt.										
	0 Interrupt not enabled										
		1 Interrupt ena	abled								
URXIE1	Bit 4	USART1 receive interrupt enable. This bit enables the URXIFG1 interrupt.									
		0 Interrupt not	enabled								
		1 Interrupt ena	abled								
	Bits 3-0	These bits may be u	sed by other modu	ules. See device-s	pecific data sheet.						

# 18.3.11 IFG1, Interrupt Flag Register 1

7	6	5	4	3	2	1	0
UTXIFG0	URXIFG	0					
rw-1	rw-0						
UTXIFG0	Bit 7	USART0 transmit in 0 No interrup 1 Interrupt pe	t pending	FG0 is set when U	OTXBUF is empty.		
URXIFG0	Bit 6	USART0 receive int  No interrupt  Interrupt pe  These bits may be used.	t pending ending			·	naracter.



# 18.3.12 IFG2, Interrupt Flag Register 2

7	6	5	4	3	2	1	0		
		UTXIFG1	URXIFG1						
		rw-1	rw-0						
	Bits 7-6	These bits may be us	sed by other modu	les. See device-s	pecific data sheet.				
UTXIFG1	Bit 5	USART1 transmit interrupt flag. UTXIFG1 is set when U1TXBUF empty.							
		0 No interrupt	pending						
		1 Interrupt per	nding						
URXIFG1	Bit 4	USART1 receive interrupt flag. URXIFG1 is set when U1RXBUF has received a complete character.							
		0 No interrupt pending							
		1 Interrupt per	nding						
	Bits 3-0	These bits may be u	sed by other modu	les. See device-s	pecific data sheet.				



# USART Peripheral Interface, SPI Mode

The universal synchronous/asynchronous receive/transmit (USART) peripheral interface supports two serial modes with one hardware module. This chapter discusses the operation of the synchronous peripheral interface or SPI mode. USART0 is implemented on the MSP430AFE2xx devices.

Topic Page

19.1	USART Introduction: SPI Mode	498
19.2	USART Operation: SPI Mode	499
19.3	USART Registers: SPI Mode	<b>506</b>



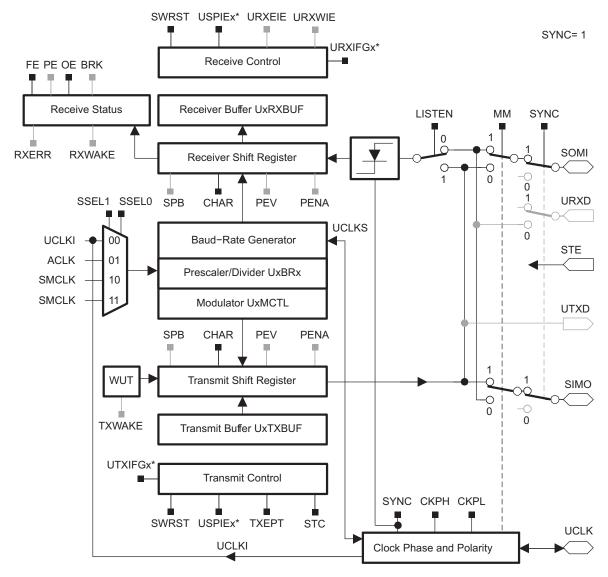
### 19.1 USART Introduction: SPI Mode

In synchronous mode, the USART connects the MSP430 to an external system via three or four pins: SIMO, SOMI, UCLK, and STE. SPI mode is selected when the SYNC bit is set and the I2C bit is cleared.

SPI mode features include:

- 7-bit or 8-bit data length
- 3-pin and 4-pin SPI operation
- · Master or slave modes
- Independent transmit and receive shift registers
- · Separate transmit and receive buffer registers
- Selectable UCLK polarity and phase control
- Programmable UCLK frequency in master mode
- Independent interrupt capability for receive and transmit

Figure 19-1 shows the USART when configured for SPI mode.



<sup>\*</sup> See the device-specific data sheet for SFR locations.

Figure 19-1. USART Block Diagram: SPI Mode







### 19.2 USART Operation: SPI Mode

In SPI mode, serial data is transmitted and received by multiple devices using a shared clock provided by the master. An additional pin, STE, is provided as to enable a device to receive and transmit data and is controlled by the master.

Three or four signals are used for SPI data exchange:

- SIMO: Slave in, master out
  - Master mode: SIMO is the data output line.
  - Slave mode: SIMO is the data input line.
- SOMI: Slave out, master in
  - Master mode: SOMI is the data input line.
  - Slave mode: SOMI is the data output line.
- UCLK: USART SPI clock
  - Master mode: UCLK is an output.
  - Slave mode: UCLK is an input.
- STE: Slave transmit enable. Used in 4-pin mode to allow multiple masters on a single bus. Not used in 3-pin mode.
  - 4-pin master mode:
    - When STE is high, SIMO and UCLK operate normally.
    - When STE is low, SIMO and UCLK are set to the input direction.
  - 4-pin slave mode:
    - When STE is high, RX/TX operation of the slave is disabled and SOMI is forced to the input direction.
    - When STE is low, RX/TX operation of the slave is enabled and SOMI operates normally.

#### 19.2.1 USART Initialization and Reset

The USART is reset by a PUC or by the SWRST bit. After a PUC, the SWRST bit is automatically set, keeping the USART in a reset condition. When set, the SWRST bit resets the URXIEx, UTXIEx, URXIFGx, OE, and FE bits and sets the UTXIFGx flag. The USPIEx bit is not altered by SWRST. Clearing SWRST releases the USART for operation.

#### NOTE: Initializing or Reconfiguring the USART Module

The required USART initialization/reconfiguration process is:

- 1. Set SWRST (BIS.B #SWRST,&UxCTL)
- 2. Initialize all USART registers with SWRST=1 (including UxCTL)
- 3. Enable USART module via the MEx SFRs (USPIEx)
- 4. Clear SWRST via software (BIC.B #SWRST,&UxCTL)
- 5. Enable interrupts (optional) via the IEx SFRs (URXIEx and/or UTXIEx)

Failure to follow this process may result in unpredictable USART behavior.



#### 19.2.2 Master Mode

Figure 19-2 shows the USART as a master in both 3-pin and 4-pin configurations. The USART initiates a data transfer when data is moved to the transmit data buffer UxTXBUF. The UxTXBUF data is moved to the TX shift register when the TX shift register is empty, initiating data transfer on SIMO starting with the most significant bit. Data on SOMI is shifted into the receive shift register on the opposite clock edge, starting with the most significant bit. When the character is received, the receive data is moved from the RX shift register to the received data buffer UxRXBUF and the receive interrupt flag, URXIFGx, is set, indicating the RX/TX operation is complete.

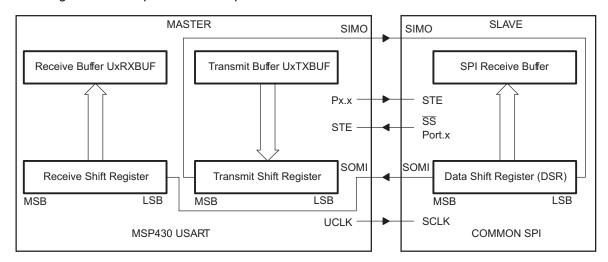


Figure 19-2. USART Master and External Slave

A set transmit interrupt flag, UTXIFGx, indicates that data has moved from UxTXBUF to the TX shift register and UxTXBUF is ready for new data. It does not indicate RX/TX completion. In master mode, the completion of an active transmission is indicated by a set transmitter empty bit TXEPT = 1.

To receive data into the USART in master mode, data must be written to UxTXBUF because receive and transmit operations operate concurrently.

#### 19.2.2.1 Four-Pin SPI Master Mode

In 4-pin master mode, STE is used to prevent conflicts with another master. The master operates normally when STE is high. When STE is low:

- SIMO and UCLK are set to inputs and no longer drive the bus
- The error bit FE is set indicating a communication integrity violation to be handled by the user

A low STE signal does not reset the USART module. The STE input signal is not used in 3-pin master mode.

#### 19.2.3 Slave Mode

Figure 19-3 shows the USART as a slave in both 3-pin and 4-pin configurations. UCLK is used as the input for the SPI clock and must be supplied by the external master. The data transfer rate is determined by this clock and not by the internal baud rate generator. Data written to UxTXBUF and moved to the TX shift register before the start of UCLK is transmitted on SOMI. Data on SIMO is shifted into the receive shift register on the opposite edge of UCLK and moved to UxRXBUF when the set number of bits are received. When data is moved from the RX shift register to UxRXBUF, the URXIFGx interrupt flag is set, indicating that data has been received. The overrun error bit, OE, is set when the previously received data is not read from UxRXBUF before new data is moved to UxRXBUF.



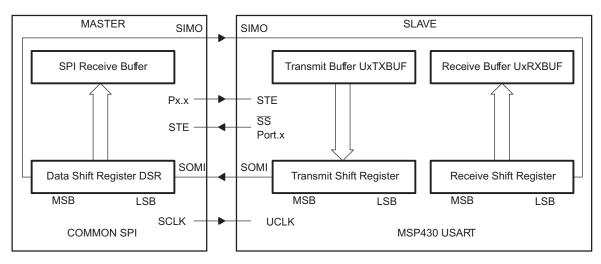


Figure 19-3. USART Slave and External Master

#### 19.2.3.1 Four-Pin SPI Slave Mode

In 4-pin slave mode, STE is used by the slave to enable the transmit and receive operations and is provided by the SPI master. When STE is low, the slave operates normally. When STE is high:

- Any receive operation in progress on SIMO is halted
- SOMI is set to the input direction

A high STE signal does not reset the USART module. The STE input signal is not used in 3-pin slave mode.

#### 19.2.4 SPI Enable

The SPI transmit/receive enable bit USPIEx enables or disables the USART in SPI mode. When USPIEx = 0, the USART stops operation after the current transfer completes, or immediately if no operation is active. A PUC or set SWRST bit disables the USART immediately and any active transfer is terminated.

### 19.2.4.1 Transmit Enable

When USPIEx = 0, any further write to UxTXBUF does not transmit. Data written to UxTXBUF begin to transmit when USPIEx = 1 and the BRCLK source is active. Figure 19-4 and Figure 19-5 show the transmit enable state diagrams.

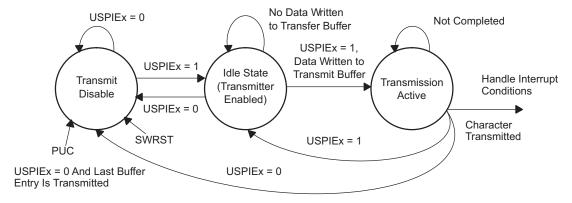


Figure 19-4. Master Transmit Enable State Diagram

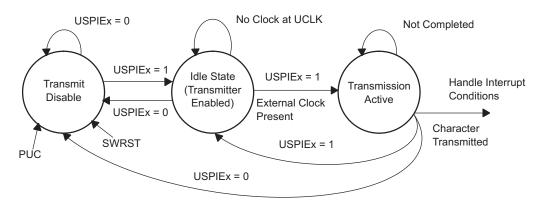


Figure 19-5. Slave Transmit Enable State Diagram

#### 19.2.4.2 Receive Enable

The SPI receive enable state diagrams are shown in Figure 19-6 and Figure 19-7. When USPIEx = 0, UCLK is disabled from shifting data into the RX shift register.

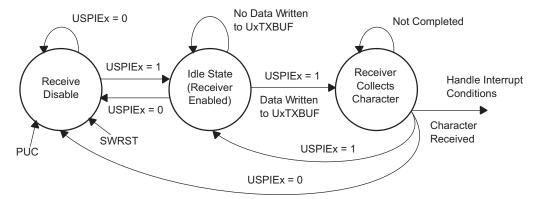


Figure 19-6. SPI Master Receive-Enable State Diagram

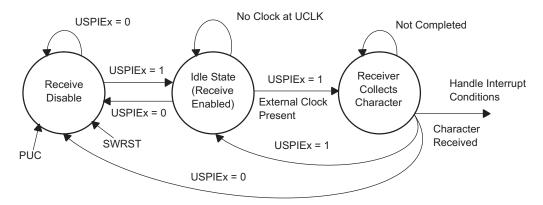


Figure 19-7. SPI Slave Receive-Enable State Diagram

### 19.2.5 Serial Clock Control

UCLK is provided by the master on the SPI bus. When MM = 1, BITCLK is provided by the USART baud rate generator on the UCLK pin as shown in Figure 19-8. When MM = 0, the USART clock is provided on the UCLK pin by the master and, the baud rate generator is not used and the SSELx bits are "don't care". The SPI receiver and transmitter operate in parallel and use the same clock source for data transfer.



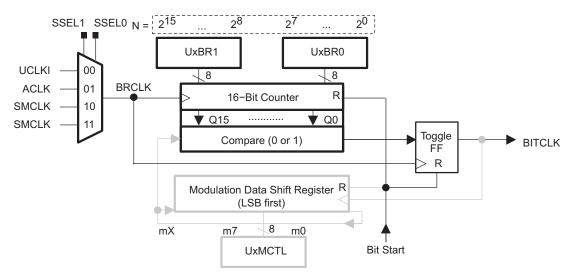


Figure 19-8. SPI Baud Rate Generator

The 16-bit value of UxBR0+UxBR1 is the division factor of the USART clock source, BRCLK. The maximum baud rate that can be generated in master mode is BRCLK/2. The maximum baud rate that can be generated in slave mode is BRCLK The modulator in the USART baud rate generator is not used for SPI mode and is recommended to be set to 000h. The UCLK frequency is given by:

Baud rate = 
$$\frac{BRCLK}{UxBR}$$
 with UxBR= [UxBR1, UxBR0]

#### 19.2.5.1 Serial Clock Polarity and Phase

The polarity and phase of UCLK are independently configured via the CKPL and CKPH control bits of the USART. Timing for each case is shown in Figure 19-9.

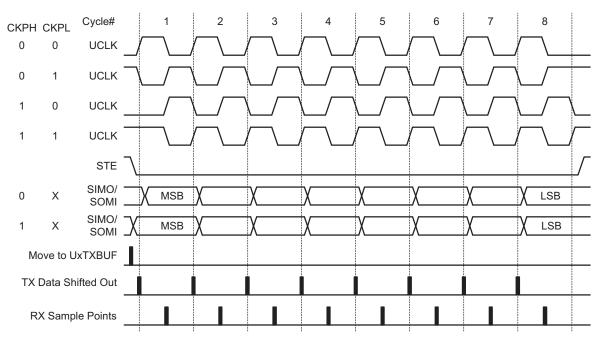


Figure 19-9. USART SPI Timing



### 19.2.6 SPI Interrupts

The USART has one interrupt vector for transmission and one interrupt vector for reception.

#### 19.2.6.1 SPI Transmit Interrupt Operation

The UTXIFGx interrupt flag is set by the transmitter to indicate that UxTXBUF is ready to accept another character. An interrupt request is generated if UTXIEx and GIE are also set. UTXIFGx is automatically reset if the interrupt request is serviced or if a character is written to UxTXBUF.

UTXIFGx is set after a PUC or when SWRST = 1. UTXIEx is reset after a PUC or when SWRST = 1. The operation is shown is Figure 19-10.

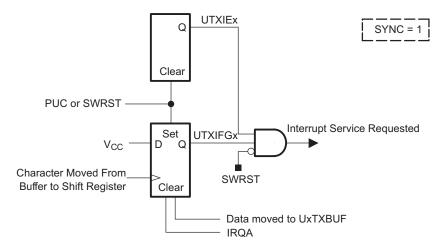


Figure 19-10. Transmit Interrupt Operation

#### NOTE: Writing to UxTXBUF in SPI Mode

Data written to UxTXBUF when UTXIFGx = 0 and USPIEx = 1 may result in erroneous data transmission.



### 19.2.6.2 SPI Receive Interrupt Operation

The URXIFGx interrupt flag is set each time a character is received and loaded into UxRXBUF as shown in Figure 19-11 and Figure 19-12. An interrupt request is generated if URXIEx and GIE are also set. URXIFGx and URXIEx are reset by a system reset PUC signal or when SWRST = 1. URXIFGx is automatically reset if the pending interrupt is served or when UxRXBUF is read.

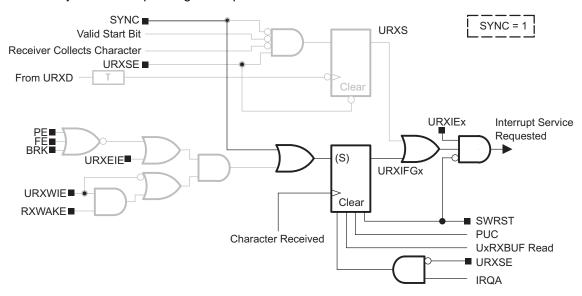


Figure 19-11. Receive Interrupt Operation

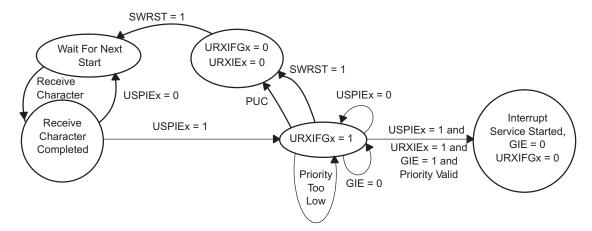


Figure 19-12. Receive Interrupt State Diagram



### 19.3 USART Registers: SPI Mode

Table 19-1 lists the registers for all devices implementing a USART module. Table 19-2 applies only to devices with a second USART module, USART1.

Table 19-1. USART0 Control and Status Registers

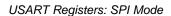
Register	Short Form	Register Type	Address	Initial State
USART control register	U0CTL	Read/write	070h	001h with PUC
Transmit control register	U0TCTL	Read/write	071h	001h with PUC
Receive control register	U0RCTL	Read/write	072h	000h with PUC
Modulation control register	U0MCTL	Read/write	073h	Unchanged
Baud rate control register 0	U0BR0	Read/write	074h	Unchanged
Baud rate control register 1	U0BR1	Read/write	075h	Unchanged
Receive buffer register	U0RXBUF	Read	076h	Unchanged
Transmit buffer register	U0TXBUF	Read/write	077h	Unchanged
SFR module enable register 1	ME1	Read/write	004h	000h with PUC
SFR interrupt enable register 1	IE1	Read/write	000h	000h with PUC
SFR interrupt flag register 1	IFG1	Read/write	002h	082h with PUC

### Table 19-2. USART1 Control and Status Registers

Register	Short Form	Register Type	Address	Initial State
USART control register	U1CTL	Read/write	078h	001h with PUC
Transmit control register	U1TCTL	Read/write	079h	001h with PUC
Receive control register	U1RCTL	Read/write	07Ah	000h with PUC
Modulation control register	U1MCTL	Read/write	07Bh	Unchanged
Baud rate control register 0	U1BR0	Read/write	07Ch	Unchanged
Baud rate control register 1	U1BR1	Read/write	07Dh	Unchanged
Receive buffer register	U1RXBUF	Read	07Eh	Unchanged
Transmit buffer register	U1TXBUF	Read/write	07Fh	Unchanged
SFR module enable register 2	ME2	Read/write	005h	000h with PUC
SFR interrupt enable register 2	IE2	Read/write	001h	000h with PUC
SFR interrupt flag register 2	IFG2	Read/write	003h	020h with PUC

### NOTE: Modifying the SFR bits

To avoid modifying control bits for other modules, it is recommended to set or clear the IEx and IFGx bits using BIS.B or BIC.B instructions, rather than MOV.B or CLR.B instructions.





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### 19.3.1 UxCTL, USART Control Register

7	6		5	4	3	2	1	0			
U	Unused		I2C	CHAR	LISTEN	SYNC	ММ	SWRST			
rw-0 rw-			rw-0	rw-0	rw-0	rw-0	rw-0	rw-1			
Unused	Bits 7-6	Unus	sed								
I2C	Bit 5	I <sup>2</sup> C m	$I^2C$ mode enable. This bit selects $I^2C$ or SPI operation when SYNC = 1.								
		0	SPI mode								
		1	I <sup>2</sup> C mode								
CHAR	Bit 4	Char	Character length								
		0	7-bit data								
		1	8-bit data								
LISTEN	Bit 3	Liste	n enable. The L	ISTEN bit selects	the loopback mod	е					
		0	Disabled								
		1	Enabled. Th	e transmit signal	is internally fed bad	ck to the receiver.					
SYNC	Bit 2	Synchronous mode enable									
		0	UART mode	<b>)</b>							
		1	SPI mode								
мм	Bit 1	Mast	er mode								
		0	USART is sl	ave							
		1	USART is m	naster							
SWRST	Bit 0	Softw	vare reset enabl	е							
		0	Disabled. USART reset released for operation.								
		1	Enabled, US	SART logic held in	reset state.						

## 19.3.2 UxTCTL, USART Transmit Control Register

7	6		5	4	3	2	1	0		
СКРН	PH CKPL		SSELx		Unused		STC	TXEPT		
rw-0	rw-0	•	rw-0	rw-0	w-0 rw-0 rw-0 rw-0					
СКРН	Bit 7	Clock	c phase select.							
		Data is changed on the first UCLK edge and captured on the following edge.								
		1	Data is capt	Pata is captured on the first UCLK edge and changed on the following edge.						
CKPL	Bit 6	Clock	Clock polarity select							
		0								
		1	The inactive	state is high.						
SSELx	Bits 5-4	Sour	ce select. These	e bits select the BF	RCLK source clock	ζ.				
		00 External UCLK (valid for slave mode only)								
		01	ACLK (valid	for master mode	only)					
		10	SMCLK (valid for master mode only)							
		11	SMCLK (val	id for master mode	e only)					
Unused	Bits 3-2	Unus	ed							
STC	Bit 1	Slave	transmit contro	ol.						
		0	4-pin SPI m	ode: STE enabled						
		1	3-pin SPI m	ode: STE disabled	l.					
TXEPT	Bit 0	Trans	smitter empty fla	ag. The TXEPT fla	g is not used in sl	ave mode.				
		0	Transmissio	n active and/or da	ta waiting in UxTX	waiting in UxTXBUF				
		1	UxTXBUF a	nd TX shift registe	r are empty					



### 19.3.3 UxRCTL, USART Receive Control Register

7	6		5	4	3	2	1	0			
FE	Unuse	d	OE			Unused					
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0			
FE	Bit 7	Frami 0 1	Framing error flag. This bit indicates a bus conflict when MM = 1 and STC = 0. FE is unused in slave 0 No conflict detected 1 A negative edge occurred on STE, indicating bus conflict								
Unused	Bit 6	Unus	ed								
OE	Bit 5	chara			a character is trai lly reset when UxR						
		0	No error								
		1	Overrun erro	or occurred							
Unused	Bits 4-0	Unus	ed								

### 19.3.4 UxBR0, USART Baud Rate Control Register 0

7	6	5	4	3	2	1	0
27	<b>2</b> <sup>6</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>2</sup>	<b>2</b> <sup>1</sup>	<b>2</b> <sup>0</sup>
rw	rw	rw	rw	rw	rw	rw	rw

### 19.3.5 UxBR1, USART Baud Rate Control Register 1

7	6	5	4	3	2	1	0
2 <sup>15</sup>	2 <sup>14</sup>	<b>2</b> <sup>13</sup>	<b>2</b> <sup>12</sup>	2 <sup>11</sup>	2 <sup>10</sup>	2 <sup>9</sup>	<b>2</b> <sup>8</sup>
rw	rw	rw	rw	rw	rw	rw	rw

**UxBRx** 

The baud-rate generator uses the content of {UxBR1+UxBR0} to set the baud rate. Unpredictable SPI operation occurs if UxBR < 2.

### 19.3.6 UxMCTL, USART Modulation Control Register

7	6	5	4	3	2	1	0
m7	m6	m5	m4	m3	m2	m1	m0
rw							

**UxMCTLx** Bits 7-0 The modulation control register is not used for SPI mode and should be set to 000h.

### 19.3.7 UxRXBUF, USART Receive Buffer Register

7	6	5	4	3	2	1	0
<b>2</b> <sup>7</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>3</sup>	2 <sup>2</sup>	<b>2</b> <sup>1</sup>	<b>2</b> <sup>0</sup>
r	r	r	r	r	r	r	r

UxRXBUFx Bits 7-0

The receive-data buffer is user accessible and contains the last received character from the receive shift register. Reading UxRXBUF resets the OE bit and URXIFGx flag. In 7-bit data mode, UxRXBUF is LSB justified and the MSB is always reset.





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USPIE1

### 10.2.9 LIVTYPHE HSAPT Transmit Puffor Posistor

19.3.0	UXIXBUF, USAKI	Transiliit buller Register	

7	6	5	4	3	2	1	0
27	<b>2</b> <sup>6</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>2</sup>	<b>2</b> <sup>1</sup>	<b>2</b> <sup>0</sup>
rw	rw	rw	rw	rw	rw	rw	rw

UxTXBUFx Bits 7-0

The transmit data buffer is user accessible and contains current data to be transmitted. When seven-bit character-length is used, the data should be MSB justified before being moved into UxTXBUF. Data is transmitted MSB first. Writing to UxTXBUF clears UTXIFGx.

### 19.3.9 ME1, Module Enable Register 1

7	6		5	4	3	2	1	0
	USPIE	:0						
	rw-0							
	Bit 7	This b	oit may be used	by other modules.	. See device-spe	ecific data sheet.		
USPIE0	Bit 6	USAF	RT0 SPI enable.	This bit enables the	he SPI mode for	USART0.		
		0	Module not e	enabled				
		1	Module enab	oled				
	Bits 5-0	These	e bits may be us	sed by other modu	les. See device-	specific data sheet.		

7	6	5	4	3	2	1	0
			USPIE1				
			rw-0				

Bits 7-5 These bits may be used by other modules. See device-specific data sheet.

Bit 4 USART1 SPI enable. This bit enables the SPI mode for USART1.

0 Module not enabled1 Module enabled

Bits 3-0 These bits may be used by other modules. See device-specific data sheet.

### 19.3.11 IE1, Interrupt Enable Register 1

19.3.10 ME2, Module Enable Register 2

7	6	5	4	3	2	1	0
UTXIE0	URXIE	)					
rw-0	rw-0						
UTXIE0	Bit 7	USART0 transmit int  Interrupt no  Interrupt en	enabled	s bit enables the U	JTXIFG0 interrupt		
URXIE0	Bit 6 Bits 5-0	0 Interrupt no 1 Interrupt en	USART0 receive interrupt enable. This bit enables the URXIFG0 interrupt.				



### 19.3.12 IE2, Interrupt Enable Register 2

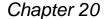
7	6	5	4	3	2	1	0
		UTXIE1	URXIE1				
		rw-0	rw-0				
	Bits 7-6	These bits may be u	sed by other mod	ules. See device-s	pecific data sheet.		
UTXIE1	Bit 5	USART1 transmit in	errupt enable. Thi	is bit enables the l	JTXIFG1 interrupt.		
		0 Interrupt no	enabled				
		1 Interrupt en	abled				
URXIE1	Bit 4	USART1 receive into	errupt enable. This	s bit enables the U	RXIFG1 interrupt.		
		0 Interrupt no	enabled				
		1 Interrupt en	abled				
	Bits 3-0	These bits may be u	sed by other mod	ules. See device-s	pecific data sheet.		

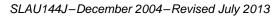
### 19.3.13 IFG1, Interrupt Flag Register 1

7	6	5	4	3	2	1	0
UTXIFG0	URXIFG	0					
rw-1	rw-0						
UTXIFG0	Bit 7	USART0 transmit in 0 No interrup 1 Interrupt pe	t pending	G0 is set when U0	OTXBUF is empty.		
URXIFG0	Bit 6 Bits 5-0	<ul><li>0 No interrup</li><li>1 Interrupt per</li></ul>	SARTO receive interrupt flag. URXIFG0 is set when U0RXBUF has received a complete character.  No interrupt pending Interrupt pending lese bits may be used by other modules. See device-specific data sheet.				

### 19.3.14 IFG2, Interrupt Flag Register 2

7	6	5	4	3	2	1	0	
		UTXIFG1	URXIFG1					
		rw-1	rw-0					
	Bits 7-6	These bits may be u	sed by other mode	ules. See device-s	specific data sheet.			
UTXIFG1	Bit 5	USART1 transmit int	errupt flag. UTXIF	G1 is set when U	1TXBUF empty.			
		0 No interrupt	pending					
		1 Interrupt per	ding					
URXIFG1	Bit 4	USART1 receive inte	rrupt flag. URXIF	G1 is set when U	RXBUF has receiv	red a complete ch	naracter.	
		0 No interrupt	pending					
		1 Interrupt per	ding					
	Bits 3-0	These bits may be used by other modules. See device-specific data sheet.						







OA

The OA is a general purpose operational amplifier. This chapter describes the OA. Two OA modules are implemented in the MSP430x22x4 devices.

Topic Page

20.1	OA Introduction	512
20.2	OA Operation	513
20.3	OA Registers	520



OA Introduction www.ti.com

#### 20.1 OA Introduction

The OA operational amplifiers support front-end analog signal conditioning prior to analog-to-digital conversion.

Features of the OA include:

- Single supply, low-current operation
- Rail-to-rail output
- Programmable settling time vs. power consumption
- Software selectable configurations
- Software selectable feedback resistor ladder for PGA implementations

#### **NOTE: Multiple OA Modules**

Some devices may integrate more than one OA module. If more than one OA is present on a device, the multiple OA modules operate identically.

Throughout this chapter, nomenclature appears such as OAxCTL0 to describe register names. When this occurs, the x is used to indicate which OA module is being discussed. In cases where operation is identical, the register is simply referred to as OAxCTL0.

The block diagram of the OA module is shown in Figure 20-1.



www.ti.com OA Operation

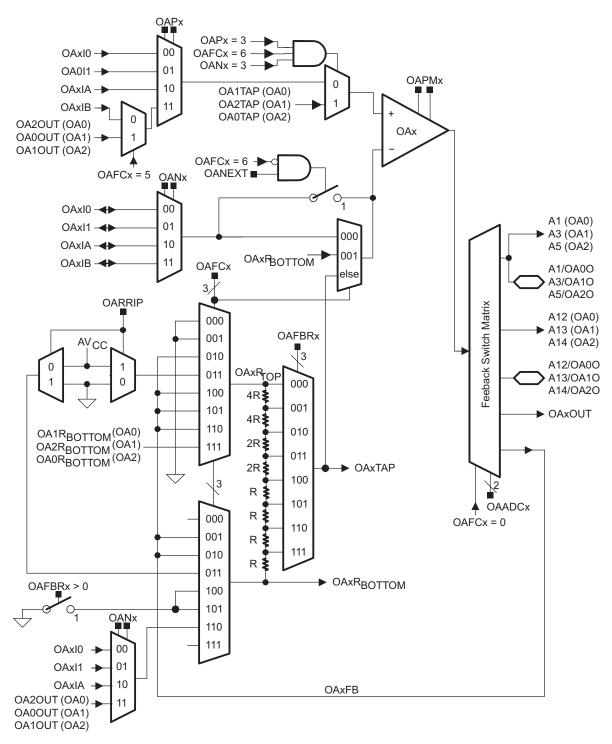


Figure 20-1. OA Block Diagram

### 20.2 OA Operation

The OA module is configured with user software. The setup and operation of the OA is discussed in the following sections.



OA Operation www.ti.com

#### 20.2.1 OA Amplifier

The OA is a configurable, low-current, rail-to-rail output operational amplifier. It can be configured as an inverting amplifier, or a non-inverting amplifier, or can be combined with other OA modules to form differential amplifiers. The output slew rate of the OA can be configured for optimized settling time vs power consumption with the OAPMx bits. When OAPMx = 00 the OA is off and the output is high-impedance. When OAPMx > 0, the OA is on. See the device-specific data sheet for parameters.

#### 20.2.2 OA Input

The OA has configurable input selection. The signals for the + and - inputs are individually selected with the OANx and OAPx bits and can be selected as external signals or internal signals. OAxI0 and OAxI1 are external signals provided for each OA module. OA0I1 provides a non-inverting input that is tied together internally for all OA modules. OAxIA and OAxIB provide device-dependent inputs. See the device data sheet for signal connections.

When the external inverting input is not needed for a mode, setting the OANEXT bit makes the internal inverting input externally available.

### 20.2.3 OA Output and Feedback Routing

The OA has configurable output selection controlled by the OAADCx bits and the OAFCx bits. The OA output signals can be routed to ADC inputs A12 (OA0), A13 (OA1), or A14 (OA2) internally, or can be routed to these ADC inputs and their external pins. The OA output signals can also be routed to ADC inputs A1 (OA0), A3 (OA1), or A5 (OA2) and the corresponding external pin. The OA output is also connected to an internal R-ladder with the OAFCx bits. The R-ladder tap is selected with the OAFBRx bits to provide programmable gain amplifier functionality.

Table 20-1 shows the OA output and feedback routing configurations. When OAFCx = 0 the OA is in general-purpose mode and feedback is achieved externally to the device. When OAFCx > 0 and when OAADCx = 00 or 11, the output of the OA is kept internal to the device. When OAFCx > 0 and OAADCx = 01 or 10, the OA output is routed both internally and externally.

OAFCx	OAADCx	OA Output and Feedback Routing
= 0	х0	OAxOUT connected to external pins and ADC input A1, A3, or A5.
= 0	x1	OAxOUT connected to external pins and ADC input A12, A13, or A14.
> 0	00	OAxOUT used for internal routing only.
> 0	01	OAxOUT connected to external pins and ADC input A12, A13, or A14.
> 0	10	OAxOUT connected to external pins and ADC input A1, A3, or A5.
> 0	11	OAxOUT connected internally to ADC input A12, A13, or A14. External A12, A13, or A14 pin connections are disconnected from the ADC.

**Table 20-1. OA Output Configurations** 

### 20.2.4 OA Configurations

The OA can be configured for different amplifier functions with the OAFCx bits as listed in Table 20-2.

**OAFC**x **OA Mode** 000 General-purpose opamp Unity gain buffer for three-opamp differential amplifier 001 010 Unity gain buffer 011 Comparator 100 Non-inverting PGA amplifier Cascaded non-inverting PGA amplifier 101 110 Inverting PGA amplifier 111 Differential amplifier

Table 20-2. OA Mode Select



www.ti.com OA Operation

### 20.2.4.1 General Purpose Opamp Mode

In this mode the feedback resistor ladder is isolated from the OAx and the OAxCTL0 bits define the signal routing. The OAx inputs are selected with the OAPx and OANx bits. The OAx output is connected to the ADC input channel as selected by the OAxCTL0 bits.

#### 20.2.4.2 Unity Gain Mode for Differential Amplifier

In this mode the output of the OAx is connected to the inverting input of the OAx providing a unity gain buffer. The non-inverting input is selected by the OAPx bits. The external connection for the inverting input is disabled and the OANx bits are don't care. The output of the OAx is also routed through the resistor ladder as part of the three-opamp differential amplifier. This mode is only for construction of the three-opamp differential amplifier.

### 20.2.4.3 Unity Gain Mode

In this mode the output of the OAx is connected to the inverting input of the OAx providing a unity gain buffer. The non-inverting input is selected by the OAPx bits. The external connection for the inverting input is disabled and the OANx bits are don't care. The OAx output is connected to the ADC input channel as selected by the OAxCTL0 bits.

#### 20.2.4.4 Comparator Mode

In this mode the output of the OAx is isolated from the resistor ladder.  $R_{\text{TOP}}$  is connected to  $AV_{\text{SS}}$  and  $R_{\text{BOTTOM}}$  is connected to  $AV_{\text{CC}}$  when OARRIP = 0. When OARRIP = 1, the connection of the resistor ladder is reversed.  $R_{\text{TOP}}$  is connected to  $AV_{\text{CC}}$  and  $R_{\text{BOTTOM}}$  is connected to  $AV_{\text{SS}}$ . The OAxTAP signal is connected to the inverting input of the OAx providing a comparator with a programmable threshold voltage selected by the OAFBRx bits. The non-inverting input is selected by the OAPx bits. Hysteresis can be added by an external positive feedback resistor. The external connection for the inverting input is disabled and the OANx bits are don't care. The OAx output is connected to the ADC input channel as selected by the OAxCTL0 bits.

### 20.2.4.5 Non-Inverting PGA Mode

In this mode the output of the OAx is connected to  $R_{\text{TOP}}$  and  $R_{\text{BOTTOM}}$  is connected to  $AV_{\text{SS}}$ . The OAxTAP signal is connected to the inverting input of the OAx providing a non-inverting amplifier configuration with a programmable gain of [1+OAxTAP ratio]. The OAxTAP ratio is selected by the OAFBRx bits. If the OAFBRx bits = 0, the gain is unity. The non-inverting input is selected by the OAPx bits. The external connection for the inverting input is disabled and the OANx bits are don't care. The OAx output is connected to the ADC input channel as selected by the OAxCTL0 bits.

#### 20.2.4.6 Cascaded Non-Inverting PGA Mode

This mode allows internal routing of the OA signals to cascade two or three OA in non-inverting mode. In this mode the non-inverting input of the OAx is connected to OA2OUT (OA0), OA0OUT (OA1), or OA1OUT (OA2) when OAPx = 11. The OAx outputs are connected to the ADC input channel as selected by the OAxCTL0 bits.

### 20.2.4.7 Inverting PGA Mode

In this mode the output of the OAx is connected to  $R_{\text{TOP}}$  and  $R_{\text{BOTTOM}}$  is connected to an analog multiplexer that multiplexes the OAxI0, OAxI1, OAxIA, or the output of one of the remaining OAs, selected with the OANx bits. The OAxTAP signal is connected to the inverting input of the OAx providing an inverting amplifier with a gain of -OAxTAP ratio. The OAxTAP ratio is selected by the OAFBRx bits. The non-inverting input is selected by the OAPx bits. The OAx output is connected to the ADC input channel as selected by the OAxCTL0 bits.



OA Operation www.ti.com

### NOTE: Using OAx Negative Input Simultaneously as ADC Input

When the pin connected to the negative input multiplexer is also used as an input to the ADC, conversion errors up to 5 mV may be observed due to internal wiring voltage drops.

#### 20.2.4.8 Differential Amplifier Mode

This mode allows internal routing of the OA signals for a two-opamp or three-opamp instrumentation amplifier. Figure 20-2 shows a two-opamp configuration with OA0 and OA1. In this mode the output of the OAx is connected to  $R_{\text{TOP}}$  by routing through another OAx in the Inverting PGA mode.  $R_{\text{BOTTOM}}$  is unconnected providing a unity gain buffer. This buffer is combined with one or two remaining OAx to form the differential amplifier. The OAx output is connected to the ADC input channel as selected by the OAxCTL0 bits.

Figure 20-2 shows an example of a two-opamp differential amplifier using OA0 and OA1. The control register settings and are shown in Table 20-3. The gain for the amplifier is selected by the OAFBRx bits for OA1 and is shown in Table 20-4. The OAx interconnections are shown in Figure 20-3.

Register	Settings (binary)
OA0CTL0	xx xx xx 0 0
OA0CTL1	000 111 0 x
OA1CTL0	11 xx xx x x
OA1CTL1	xxx 110 0 x

Table 20-3. Two-Opamp Differential Amplifier Control Register Settings

Table 20-4. Two-Opamp Differential Amplifier Gain Settings

OA1 OAFBRx	Gain
000	0
001	1/3
010	1
011	1 2/3
100	3
101	4 1/3
110	7
111	15

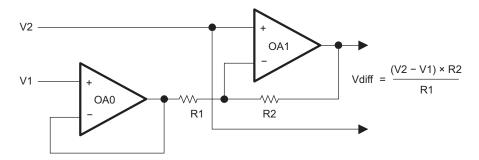


Figure 20-2. Two-Opamp Differential Amplifier



www.ti.com OA Operation

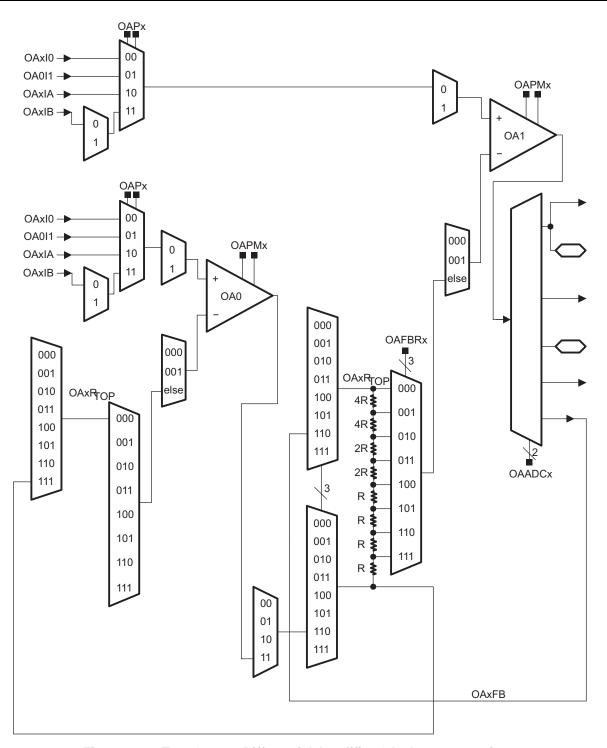


Figure 20-3. Two-Opamp Differential Amplifier OAx Interconnections

OA



OA Operation www.ti.com

Figure 20-4 shows an example of a three-opamp differential amplifier using OA0, OA1 and OA2 (Three opamps are not available on all devices. See device-specific data sheet for implementation.). The control register settings are shown in Table 20-5. The gain for the amplifier is selected by the OAFBRx bits of OA0 and OA2. The OAFBRx settings for both OA0 and OA2 must be equal. The gain settings are shown in Table 20-6. The OAx interconnections are shown in Figure 20-5.

Table 20-5. Three-Opamp Differential Amplifier Control Register Settings

Register	Settings (binary)
OA0CTL0	xx xx xx 0 0
OA0CTL1	xxx 001 0 x
OA1CTL0	xx xx xx 0 0
OA1CTL1	000 111 0 x
OA2CTL0	11 11 xx x x
OA2CTL1	xxx 110 0 x

Table 20-6. Three-Opamp Differential Amplifier Gain Settings

OA0/OA2 OAFBRx	Gain
000	0
001	1/3
010	1
011	1 2/3
100	3
101	4 1/3
110	7
111	15

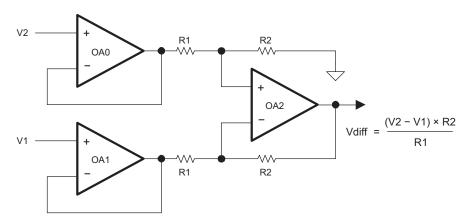


Figure 20-4. Three-Opamp Differential Amplifier



www.ti.com OA Operation

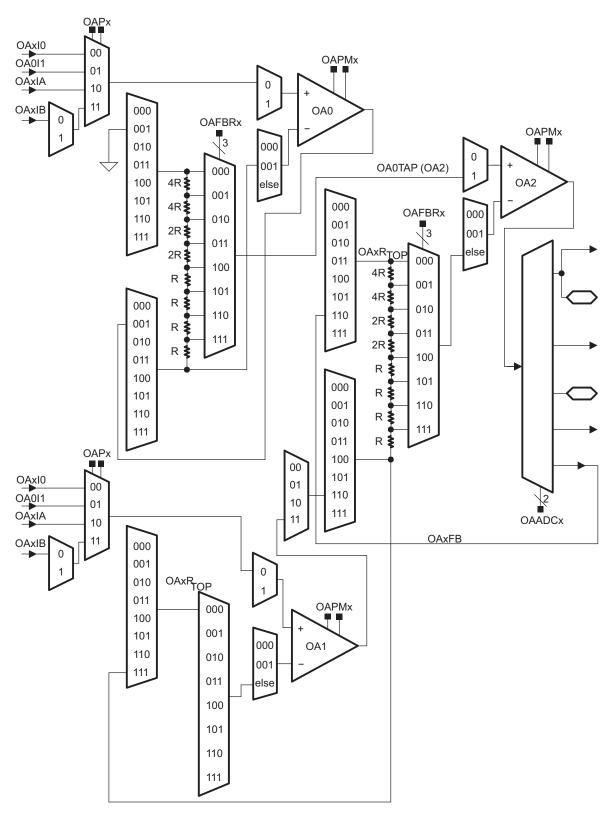


Figure 20-5. Three-Opamp Differential Amplifier OAx Interconnections

OA



OA Registers www.ti.com

### 20.3 OA Registers

The OA registers are listed in Table 20-7.

### Table 20-7. OA Registers

Register	Short Form	Register Type	Address	Initial State
OA0 control register 0	OA0CTL0	Read/write	0C0h	Reset with POR
OA0 control register 1	OA0CTL1	Read/write	0C1h	Reset with POR
OA1 control register 0	OA1CTL0	Read/write	0C2h	Reset with POR
OA1 control register 1	OA1CTL1	Read/write	0C3h	Reset with POR
OA2 control register 0	OA2CTL0	Read/write	0C4h	Reset with POR
OA2 control register 1	OA2CTL1	Read/write	0C5h	Reset with POR



www.ti.com OA Registers

### 20.3.1 OAxCTL0, Opamp Control Register 0

7	6		5	4	3	2	1	0
	OANx		OAPx		OA	OAPMx		ADCx
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
OANx	Bits 7-6	Invert	ing input select	. These bits select	the input signal f	or the OA inverting	j input.	
		00	OAxI0					
		01	OAxI1					
		10	OAxIA (see	the device-specific	data sheet for co	onnected signal)		
		11	OAxIB (see	the device-specific	data sheet for co	onnected signal)		
OAPx	Bits 5-4	Non-i	nverting input s	elect. These bits s	elect the input sig	gnal for the OA nor	n-inverting input.	
		00	OAxI0					
		01	OA0I1					
		10	OAxIA (see	the device-specific	data sheet for co	onnected signal)		
		11	OAxIB (see	the device-specific	data sheet for co	onnected signal)		
OAPMx	Bits 3-2	Slew	rate select. The	ese bits select the	slew rate vs. curre	ent consumption fo	or the OA.	
		00	Off, output h	igh Z				
		01	Slow					
		10	Medium					
		11	Fast					
OAADCx	Bits 1-0		utput select. Th Mx > 0.	ese bits, together	with the OAFCx b	oits, control the rou	ting of the OAx o	utput when
		When	OAFCx = 0:					
		00	OAxOUT co	nnected to externa	al pins and ADC in	nput A1, A3, or A5		
		01	OAxOUT co	nnected to externa	al pins and ADC in	nput A12, A13, or <i>i</i>	<b>A14</b>	
		10	OAxOUT co	nnected to externa	al pins and ADC in	nput A1, A3, or A5		
		11	OAxOUT co	nnected to externa	al pins and ADC in	nput A12, A13, or <i>i</i>	<b>A14</b>	
		When	OAFCx > 0:					
		00	OAxOUT us	ed for internal rout	ing only			
		01	OAxOUT co	nnected to externa	al pins and ADC in	nput A12, A13, or A	<b>A14</b>	
		10	OAxOUT co	nnected to externa	al pins and ADC in	nput A1, A3, or A5		
		11		nnected internally are disconnected		2, A13 , or A14. Ex	ternal A12, A13,	or A14 pin



OA Registers www.ti.com

## 20.3.2 OAxCTL1, Opamp Control Register 1

7	6		5	4	3	2	1	0
	OAFBRx				OAFCx		OANEXT	OARRIP
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
OAFBRx	Bits 7-5	OAx fe	eedback resisto	or select				
		000	Tap 0 - 0R/1	6R				
		001	Tap 1 - 4R/1	2R				
		010	Tap 2 - 8R/8	R				
		011	Tap 3 - 10R/	6R				
		100	Tap 4 - 12R/	4R				
		101	Tap 5 - 13R/	3R				
		110	Tap 6 - 14R/	2R				
		111	Tap 7 - 15R/	1R				
OAFCx	Bits 4-2	OAx fu	unction control.	This bit selects the	ne function of OAx	(		
		000	General purp	oose opamp				
		001	Unity gain bu	uffer for three-opa	mp differential am	plifier		
		010	Unity gain bu	uffer				
		011	Comparator					
		100	Non-inverting	g PGA amplifier				
		101	Cascaded no	on-inverting PGA	amplifier			
		110	Inverting PG	A amplifier				
		111	Differential a	mplifier				
OANEXT	Bit 1			xternally available ed resistor netwo		et, connects the in	verting OAx input	to the external
		0	OAx inverting	g input not extern	ally available			
		1	OAx inverting	g input externally	available			
OARRIP	Bit 0	OAx r	everse resistor	connection in cor	nparator mode			
		0	R <sub>TOP</sub> is conn	ected to AV <sub>SS</sub> and	R <sub>BOTTOM</sub> is conne	cted to AV <sub>CC</sub> wher	OAFCx = 3	
		1	R <sub>TOP</sub> is conn	ected to AV <sub>CC</sub> and	d R <sub>BOTTOM</sub> is conne	ected to AV <sub>SS</sub> wher	OAFCx = 3.	



# Comparator\_A+

Comparator\_A+ is an analog voltage comparator. This chapter describes the operation of the Comparator\_A+ of the 2xx family.

Topic Page

21.1	Comparator_A+ Introduction	<b>524</b>
21.2	Comparator_A+ Operation	525
21.3	Comparator_A+ Registers	530
	• •	



### 21.1 Comparator\_A+ Introduction

The Comparator\_A+ module supports precision slope analog-to-digital conversions, supply voltage supervision, and monitoring of external analog signals.

Features of Comparator\_A+ include:

- Inverting and non-inverting terminal input multiplexer
- Software selectable RC-filter for the comparator output
- Output provided to Timer\_A capture input
- · Software control of the port input buffer
- Interrupt capability
- Selectable reference voltage generator
- Comparator and reference generator can be powered down
- Input Multiplexer

The Comparator\_A+ block diagram is shown in Figure 21-1.

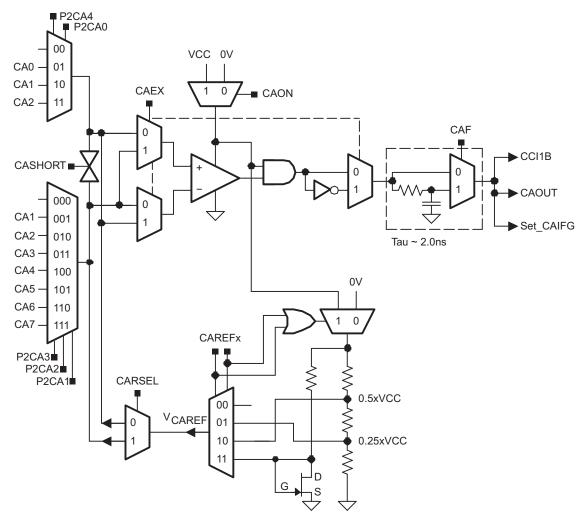


Figure 21-1. Comparator\_A+ Block Diagram

**NOTE: MSP430G2210:** Channels 2, 5, 6, and 7 are available. Other channels should not be enabled.



### 21.2 Comparator\_A+ Operation

The Comparator\_A+ module is configured with user software. The setup and operation of Comparator\_A+ is discussed in the following sections.

### 21.2.1 Comparator

The comparator compares the analog voltages at the + and - input terminals. If the + terminal is more positive than the - terminal, the comparator output CAOUT is high. The comparator can be switched on or off using control bit CAON. The comparator should be switched off when not in use to reduce current consumption. When the comparator is switched off, the CAOUT is always low.

### 21.2.2 Input Analog Switches

The analog input switches connect or disconnect the two comparator input terminals to associated port pins using the P2CAx bits. Both comparator terminal inputs can be controlled individually. The P2CAx bits allow:

- Application of an external signal to the + and terminals of the comparator
- Routing of an internal reference voltage to an associated output port pin

Internally, the input switch is constructed as a T-switch to suppress distortion in the signal path.

### **NOTE:** Comparator Input Connection

When the comparator is on, the input terminals should be connected to a signal, power, or ground. Otherwise, floating levels may cause unexpected interrupts and increased current consumption.

**NOTE: MSP430G2210:** Comparator channels 0, 1, 3, and 4 are implemented but not available at the device pins. To avoid floating inputs, these comparator inputs should not be enabled.

The CAEX bit controls the input multiplexer, exchanging which input signals are connected to the comparator's + and - terminals. Additionally, when the comparator terminals are exchanged, the output signal from the comparator is inverted. This allows the user to determine or compensate for the comparator input offset voltage.



### 21.2.3 Input Short Switch

The CASHORT bit shorts the comparator\_A+ inputs. This can be used to build a simple sample-and-hold for the comparator as shown in Figure 21-2.

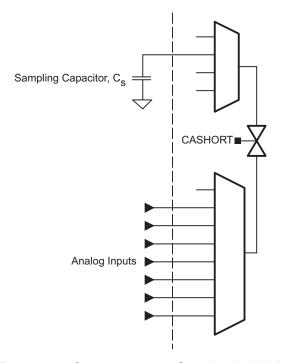


Figure 21-2. Comparator\_A+ Sample-And-Hold

The required sampling time is proportional to the size of the sampling capacitor ( $C_s$ ), the resistance of the input switches in series with the short switch ( $R_i$ ), and the resistance of the external source ( $R_s$ ). The total internal resistance ( $R_i$ ) is typically in the range of 2 to 10 k $\Omega$ . The sampling capacitor  $C_s$  should be greater than 100 pF. The time constant, Tau, to charge the sampling capacitor  $C_s$  can be calculated with the following equation:

$$Tau = (R_I + R_S) \times C_S$$

Depending on the required accuracy 3 to 10 Tau should be used as a sampling time. With 3 Tau the sampling capacitor is charged to approximately 95% of the input signals voltage level, with 5 Tau it is charge to more than 99% and with 10 Tau the sampled voltage is sufficient for 12-bit accuracy.

### 21.2.4 Output Filter

The output of the comparator can be used with or without internal filtering. When control bit CAF is set, the output is filtered with an on-chip RC-filter.

Any comparator output oscillates if the voltage difference across the input terminals is small. Internal and external parasitic effects and cross coupling on and between signal lines, power supply lines, and other parts of the system are responsible for this behavior as shown in Figure 21-3. The comparator output oscillation reduces accuracy and resolution of the comparison result. Selecting the output filter can reduce errors associated with comparator oscillation.



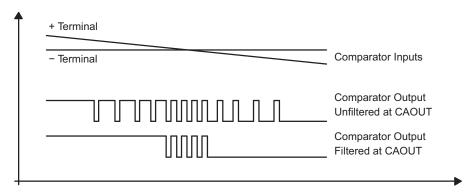


Figure 21-3. RC-Filter Response at the Output of the Comparator

### 21.2.5 Voltage Reference Generator

The voltage reference generator is used to generate V<sub>CAREF</sub>, which can be applied to either comparator input terminal. The CAREFx bits control the output of the voltage generator. The CARSEL bit selects the comparator terminal to which V<sub>CAREE</sub> is applied. If external signals are applied to both comparator input terminals, the internal reference generator should be turned off to reduce current consumption. The voltage reference generator can generate a fraction of the device's V<sub>CC</sub> or a fixed transistor threshold voltage of ~0.55 V.

### 21.2.6 Comparator A+, Port Disable Register CAPD

The comparator input and output functions are multiplexed with the associated I/O port pins, which are digital CMOS gates. When analog signals are applied to digital CMOS gates, parasitic current can flow from VCC to GND. This parasitic current occurs if the input voltage is near the transition level of the gate. Disabling the port pin buffer eliminates the parasitic current flow and therefore reduces overall current consumption.

The CAPDx bits, when set, disable the corresponding pin input and output buffers as shown in Figure 21-4. When current consumption is critical, any port pin connected to analog signals should be disabled with its CAPDx bit.

Selecting an input pin to the comparator multiplexer with the P2CAx bits automatically disables the input and output buffers for that pin, regardless of the state of the associated CAPDx bit.

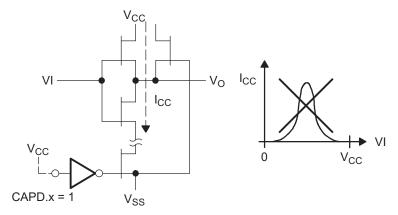


Figure 21-4. Transfer Characteristic and Power Dissipation in a CMOS Inverter/Buffer

NOTE: MSP430G2210:The channels 0, 1, 3, an 4 are implemented by not available at pins. To avoid floating inputs these inputs should not be used.



### 21.2.7 Comparator A+ Interrupts

One interrupt flag and one interrupt vector are associated with the Comparator\_A+ as shown in Figure 21-5. The interrupt flag CAIFG is set on either the rising or falling edge of the comparator output, selected by the CAIES bit. If both the CAIE and the GIE bits are set, then the CAIFG flag generates an interrupt request. The CAIFG flag is automatically reset when the interrupt request is serviced or may be reset with software.

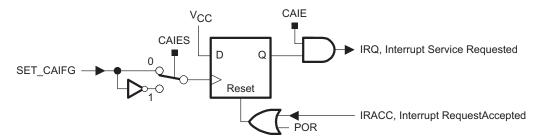


Figure 21-5. Comparator\_A+ Interrupt System

### 21.2.8 Comparator A+ Used to Measure Resistive Elements

The Comparator\_A+ can be optimized to precisely measure resistive elements using single slope analog-to-digital conversion. For example, temperature can be converted into digital data using a thermistor, by comparing the thermistor's capacitor discharge time to that of a reference resistor as shown in Figure 21-6. A reference resister Rref is compared to Rmeas.

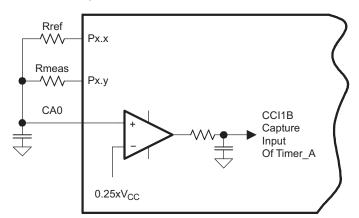


Figure 21-6. Temperature Measurement System

The MSP430 resources used to calculate the temperature sensed by Rmeas are:

- Two digital I/O pins to charge and discharge the capacitor.
- I/O set to output high (V<sub>CC</sub>) to charge capacitor, reset to discharge.
- I/O switched to high-impedance input with CAPDx set when not in use.
- One output charges and discharges the capacitor via Rref.
- · One output discharges capacitor via Rmeas.
- The + terminal is connected to the positive terminal of the capacitor.
- The terminal is connected to a reference level, for example 0.25 x V<sub>cc</sub>.
- The output filter should be used to minimize switching noise.
- CAOUT used to gate Timer\_A CCI1B, capturing capacitor discharge time.

More than one resistive element can be measured. Additional elements are connected to CA0 with available I/O pins and switched to high impedance when not being measured.



The thermistor measurement is based on a ratiometric conversion principle. The ratio of two capacitor discharge times is calculated as shown in Figure 21-7.

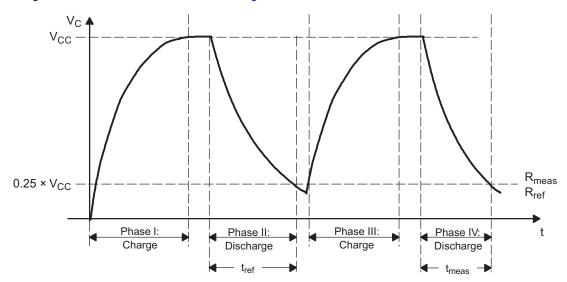


Figure 21-7. Timing for Temperature Measurement Systems

The  $V_{\text{CC}}$  voltage and the capacitor value should remain constant during the conversion, but are not critical since they cancel in the ratio:

$$\frac{N_{meas}}{N_{ref}} = \frac{-R_{meas} \times C \times ln \; \frac{V_{ref}}{V_{CC}}}{-R_{ref} \times C \times ln \; \frac{V_{ref}}{V_{CC}}}$$

$$\frac{N_{meas}}{N_{ref}} = \frac{R_{meas}}{R_{ref}}$$

$$R_{meas} = R_{ref} \times \frac{N_{meas}}{N_{ref}}$$



### 21.3 Comparator\_A+ Registers

The Comparator\_A+ registers are listed in Table 21-1.

Table 21-1. Comparator\_A+ Registers

Register	Short Form	Register Type	Address	Initial State
Comparator_A+ control register 1	CACTL1	Read/write	059h	Reset with POR
Comparator_A+ control register 2	CACTL2	Read/write	05Ah	Reset with POR
Comparator_A+ port disable	CAPD	Read/write	05Bh	Reset with POR



## 21.3.1 CACTL1, Comparator\_A+ Control Register 1

7	6		5	4	3	2	1	0
CAEX	CARSEI	L	CAF	REFx	CAON	CAIES	CAIE	CAIFG
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
CAEX	Bit 7	Compa	arator_A+ exch	nange. This bit ex	changes the comp	arator inputs and i	nverts the compa	rator output.
CARSEL	Bit 6	Compa	arator_A+ refe	rence select. This	bit selects which	terminal the $V_{CAREF}$	is applied to.	
		When	CAEX = 0:					
		0	V <sub>CAREF</sub> is app	olied to the + term	inal			
		1	V <sub>CAREF</sub> is app	olied to the - termi	nal			
		When	CAEX = 1:					
		0	V <sub>CAREF</sub> is app	olied to the - termi	nal			
		1	V <sub>CAREF</sub> is app	olied to the + term	inal			
CAREF	Bits 5-4	Compa	arator_A+ refe	rence. These bits	select the referen	ce voltage V <sub>CAREF.</sub>		
		00	Internal reference off. An external reference can be applied.					
		01	$0.25 \times V_{CC}$					
		10	$0.50 \times V_{CC}$					
		11	Diode refere	nce is selected				
CAON	Bit 3		Comparator_A+ on. This bit turns on the comparator. When the comparator is off it consumes no current. The reference circuitry is enabled or disabled independently.					
		0	Off					
		1	On					
CAIES	Bit 2	Compa	arator_A+ inter	rupt edge select				
		0	Rising edge					
		1	Falling edge					
CAIE	Bit 1	Comparator_A+ interrupt enable						
		0	Disabled					
		1	Enabled					
CAIFG	Bit 0	The Co	omparator_A+	interrupt flag				
		0	No interrupt	pending				
		1	Interrupt per	nding				



### 21.3.2 CACTL2, Comparator\_A+, Control Register

7	6		5	4	3	2	1	0
CASHORT	P2CA4		P2CA3	P2CA2	P2CA1	P2CA0	CAF	CAOUT
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	r-(0)
CASHORT	Bit 7	Input s	short. This bit s	horts the + and -	input terminals.			
		0	Inputs not sh	orted				
		1	Inputs shorte	ed				
P2CA4	Bit 6		select. This bit t when CAEX = 1		A0 selects the + t	erminal input wher	n CAEX = 0 and t	the - terminal
P2CA3 <sup>(1)</sup>	Bits 5-3	Input s	select. These b	its select the - ter	minal input when	CAEX = 0 and the	+ terminal input	when CAEX = 1.
P2CA2 P2CA1		000	No connection	on				
PZCAT		001	CA1					
		010	CA2					
		011	CA3					
		100	CA4					
		101	CA5					
		110	CA6					
		111	CA7					
P2CA0	Bit 2		select. This bit, when CAEX = 1		CA4, selects the +	terminal input whe	en CAEX = 0 and	the - terminal
		00	No connection	on				
		01	CA0					
		10	CA1					
		11	CA2					
CAF	Bit 1	Comp	arator_A+ outp	ut filter				
		0	Comparator_	_A+ output is not	filtered			
		1	Comparator_	_A+ output is filter	red			
CAOUT	Bit 0	Comp	arator_A+ outp	ut. This bit reflect	s the value of the	comparator output	. Writing this bit h	nas no effect.

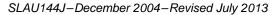
<sup>(1)</sup> **MSP430G2210**: Only channels 2, 5, 6, and 7 are available. Other channels should not be selected.

### 21.3.3 CAPD, Comparator\_A+, Port Disable Register

7	6	5	4	3	2	1	0	
CAPD7	CAPD6	CAPD5	CAPD4	CAPD3	CAPD2	CAPD1	CAPD0	
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	
CAPDx <sup>(1)</sup>	Bits 7-0	Comparator_A+ port disable. These bits individually disable the input buffer for the pins of the port associated with Comparator_A+. For example, if CA0 is on pin P2.3, the CAPDx bits can be used to individually enable or disable each P2.x pin buffer. CAPD0 disables P2.0, CAPD1 disables P2.1, etc.						
		0 The input b	uffer is enabled.					
		1 The input b	uffer is disabled.					

 $<sup>^{(1)}</sup>$  MSP430G2210: Channels 2, 5, 6, and 7 are available. Other channels should not be disabled.

## Chapter 22





# ADC10

The ADC10 module is a high-performance 10-bit analog-to-digital converter. This chapter describes the operation of the ADC10 module of the 2xx family in general. There are device with less than eight external input channels.

Topic Page

22.1	ADC10 Introduction	534
22.2	ADC10 Operation	536
22.3	ADC10 Registers	552



ADC10 Introduction www.ti.com

#### 22.1 ADC10 Introduction

The ADC10 module supports fast, 10-bit analog-to-digital conversions. The module implements a 10-bit SAR core, sample select control, reference generator, and data transfer controller (DTC).

The DTC allows ADC10 samples to be converted and stored anywhere in memory without CPU intervention. The module can be configured with user software to support a variety of applications.

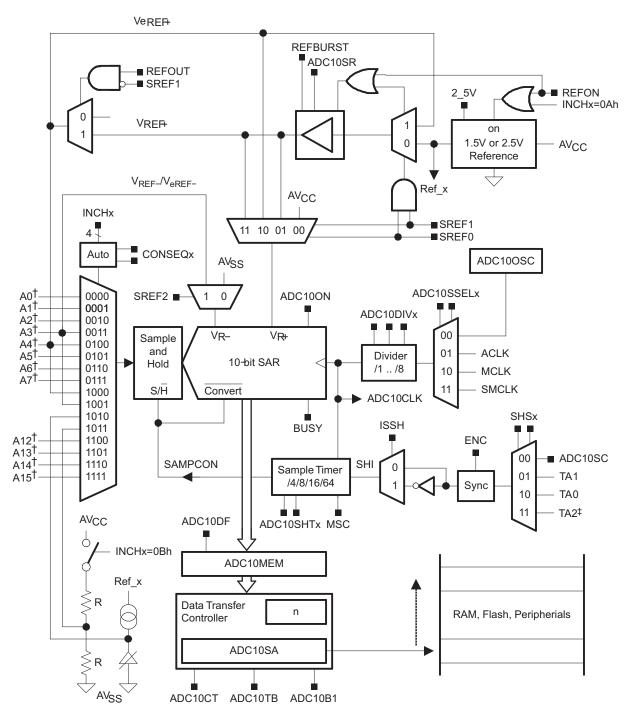
#### ADC10 features include:

- Greater than 200-ksps maximum conversion rate
- Monotonic 10-bit converter with no missing codes
- · Sample-and-hold with programmable sample periods
- Conversion initiation by software or Timer\_A
- Software selectable on-chip reference voltage generation (1.5 V or 2.5 V)
- Software selectable internal or external reference
- Up to eight external input channels (twelve on MSP430F22xx devices)
- Conversion channels for internal temperature sensor, V<sub>CC</sub>, and external references
- · Selectable conversion clock source
- Single-channel, repeated single-channel, sequence, and repeated sequence conversion modes
- ADC core and reference voltage can be powered down separately
- Data transfer controller for automatic storage of conversion results

The block diagram of ADC10 is shown in Figure 22-1.



www.ti.com ADC10 Introduction



<sup>†</sup>Channels A12-A15 are available in MSP430F22xx devices only. Channels A12-A15 tied to channel A11 in other devices. Not all channels are available in all devices.

Figure 22-1. ADC10 Block Diagram

<sup>‡</sup>TA1 on MSP430F20x2, MSP430G2x31, and MSP430G2x30 devices



ADC10 Operation www.ti.com

### 22.2 ADC10 Operation

The ADC10 module is configured with user software. The setup and operation of the ADC10 is discussed in the following sections.

#### 22.2.1 10-Bit ADC Core

The ADC core converts an analog input to its 10-bit digital representation and stores the result in the ADC10MEM register. The core uses two programmable/selectable voltage levels ( $V_{R+}$  and  $V_{R-}$ ) to define the upper and lower limits of the conversion. The digital output ( $N_{ADC}$ ) is full scale (03FFh) when the input signal is equal to or higher than  $V_{R+}$ , and zero when the input signal is equal to or lower than  $V_{R-}$ . The input channel and the reference voltage levels ( $V_{R+}$  and  $V_{R-}$ ) are defined in the conversion-control memory. Conversion results may be in straight binary format or 2s-complement format. The conversion formula for the ADC result when using straight binary format is:

$$N_{ADC} = 1023 \times \frac{V_{IN} - V_{R-}}{V_{R+} - V_{R-}}$$

The ADC10 core is configured by two control registers, ADC10CTL0 and ADC10CTL1. The core is enabled with the ADC10ON bit. With few exceptions the ADC10 control bits can only be modified when ENC = 0. ENC must be set to 1 before any conversion can take place.

#### 22.2.1.1 Conversion Clock Selection

The ADC10CLK is used both as the conversion clock and to generate the sampling period. The ADC10 source clock is selected using the ADC10SSELx bits and can be divided from 1 to 8 using the ADC10DIVx bits. Possible ADC10CLK sources are SMCLK, MCLK, ACLK, and internal oscillator ADC10OSC.

The ADC10OSC, generated internally, is in the 5-MHz range, but varies with individual devices, supply voltage, and temperature. See the device-specific data sheet for the ADC10OSC specification.

The user must ensure that the clock chosen for ADC10CLK remains active until the end of a conversion. If the clock is removed during a conversion, the operation does not complete, and any result is invalid.

### 22.2.2 ADC10 Inputs and Multiplexer

The eight external and four internal analog signals are selected as the channel for conversion by the analog input multiplexer. The input multiplexer is a break-before-make type to reduce input-to-input noise injection that can result from channel switching (see Figure 22-2). The input multiplexer is also a T-switch to minimize the coupling between channels. Channels that are not selected are isolated from the A/D, and the intermediate node is connected to analog ground ( $V_{SS}$ ) so that the stray capacitance is grounded to help eliminate crosstalk.

The ADC10 uses the charge redistribution method. When the inputs are internally switched, the switching action may cause transients on the input signal. These transients decay and settle before causing errant conversion.

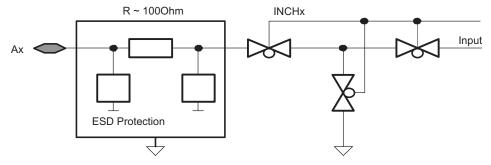


Figure 22-2. Analog Multiplexer



www.ti.com ADC10 Operation

#### 22.2.2.1 Analog Port Selection

The ADC10 external inputs Ax,  $V_{eREF+}$ , and  $V_{REF-}$  share terminals with general purpose I/O ports, which are digital CMOS gates (see the device-specific data sheet). When analog signals are applied to digital CMOS gates, parasitic current can flow from VCC to GND. This parasitic current occurs if the input voltage is near the transition level of the gate. Disabling the port pin buffer eliminates the parasitic current flow and therefore reduces overall current consumption. The ADC10AEx bits provide the ability to disable the port pin input and output buffers.

```
; P2.3 on MSP430F22xx device configured for analog input BIS.B #08h,&ADC10AE0 ; P2.3 ADC10 function and enable
```

Devices which don't have all the ADC10 external inputs channels Ax or  $V_{eREF+}/V_{REF+}$  and  $V_{eREF-}/V_{REF-}$  available at device pins must not alter the default register bit configuration of the not available pins. See device specific data sheet.

### 22.2.3 Voltage Reference Generator

The ADC10 module contains a built-in voltage reference with two selectable voltage levels. Setting REFON = 1 enables the internal reference. When REF2\_5V = 1, the internal reference is 2.5 V. When REF2\_5V = 0, the reference is 1.5 V. The internal reference voltage may be used internally (REFOUT = 0) and, when REFOUT = 1, externally on pin  $V_{REF+}$ . REFOUT = 1 should only be used if the pins  $V_{REF+}$  and  $V_{REF-}$  are available as device pins.

External references may be supplied for  $V_{R+}$  and  $V_{R-}$  through pins A4 and A3 respectively. When external references are used, or when  $V_{CC}$  is used as the reference, the internal reference may be turned off to save power.

An external positive reference  $V_{eREF+}$  can be buffered by setting SREF0 = 1 and SREF1 = 1 (only devices with  $V_{eREF+}$  pin). This allows using an external reference with a large internal resistance at the cost of the buffer current. When REFBURST = 1 the increased current consumption is limited to the sample and conversion period.

External storage capacitance is not required for the ADC10 reference source as on the ADC12.

#### 22.2.3.1 Internal Reference Low-Power Features

The ADC10 internal reference generator is designed for low power applications. The reference generator includes a band-gap voltage source and a separate buffer. The current consumption of each is specified separately in the device-specific data sheet. When REFON = 1, both are enabled and when REFON = 0 both are disabled. The total settling time when REFON becomes set is approximately 30  $\mu$ s.

When REFON = 1, but no conversion is active, the buffer is automatically disabled and automatically reenabled when needed. When the buffer is disabled, it consumes no current. In this case, the bandgap voltage source remains enabled.

When REFOUT = 1, the REFBURST bit controls the operation of the internal reference buffer. When REFBURST = 0, the buffer is on continuously, allowing the reference voltage to be present outside the device continuously. When REFBURST = 1, the buffer is automatically disabled when the ADC10 is not actively converting and is automatically re-enabled when needed.

The internal reference buffer also has selectable speed versus power settings. When the maximum conversion rate is below 50 ksps, setting ADC10SR = 1 reduces the current consumption of the buffer approximately 50%.

#### 22.2.4 Auto Power-Down

The ADC10 is designed for low power applications. When the ADC10 is not actively converting, the core is automatically disabled and is automatically re-enabled when needed. The ADC10OSC is also automatically enabled when needed and disabled when not needed. When the core or oscillator is disabled, it consumes no current.



ADC10 Operation www.ti.com

### 22.2.5 Sample and Conversion Timing

An analog-to-digital conversion is initiated with a rising edge of sample input signal SHI. The source for SHI is selected with the SHSx bits and includes the following:

- The ADC10SC bit
- The Timer\_A Output Unit 1
- The Timer\_A Output Unit 0
- The Timer\_A Output Unit 2

The polarity of the SHI signal source can be inverted with the ISSH bit. The SHTx bits select the sample period t<sub>sample</sub> to be 4, 8, 16, or 64 ADC10CLK cycles. The sampling timer sets SAMPCON high for the selected sample period after synchronization with ADC10CLK. Total sampling time is t<sub>sample</sub> plus t<sub>sync</sub>. The high-to-low SAMPCON transition starts the analog-to-digital conversion, which requires 13 ADC10CLK cycles as shown in Figure 22-3.

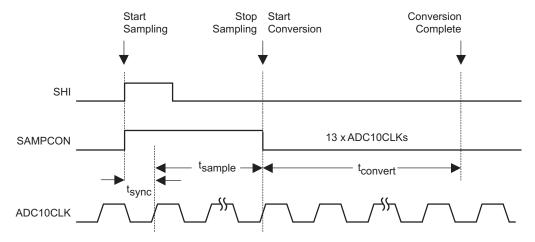


Figure 22-3. Sample Timing

### 22.2.5.1 Sample Timing Considerations

When SAMPCON = 0 all Ax inputs are high impedance. When SAMPCON = 1, the selected Ax input can be modeled as an RC low-pass filter during the sampling time  $t_{\text{sample}}$ , as shown in Figure 22-4. An internal MUX-on input resistance  $R_{\text{I}}$  (2 k $\Omega$  maximum) in series with capacitor  $C_{\text{I}}$  (27 pF maximum) is seen by the source. The capacitor  $C_{\text{I}}$  voltage  $V_{\text{C}}$  must be charged to within ½ LSB of the source voltage  $V_{\text{S}}$  for an accurate 10-bit conversion.

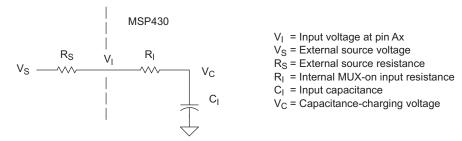


Figure 22-4. Analog Input Equivalent Circuit

The resistance of the source  $R_s$  and  $R_i$  affect  $t_{sample}$ . The following equations can be used to calculate the minimum sampling time for a 10-bit conversion.

$$t_{\text{sample}} > (R_S + R_I) \times \ln(2^{11}) \times C_I$$

Substituting the values for R<sub>1</sub> and C<sub>1</sub> given above, the equation becomes:

$$t_{sample} > (R_S + 2 k\Omega) \times 7.625 \times 27 pF$$



www.ti.com ADC10 Operation

For example, if  $R_S$  is 10 k $\Omega$ ,  $t_{sample}$  must be greater than 2.47  $\mu s$ .

When the reference buffer is used in burst mode, the sampling time must be greater than the sampling time calculated and the settling time of the buffer,  $t_{REFBURST}$ :

$$t_{\text{sample}} > \begin{cases} (R_{\text{S}} + R_{\text{I}}) \times \ln(2^{11}) \times C_{\text{I}} \\ t_{\text{REFBURST}} \end{cases}$$

For example, if  $V_{Ref}$  is 1.5 V and  $R_S$  is 10 k $\Omega$ ,  $t_{sample}$  must be greater than 2.47  $\mu s$  when ADC10SR = 0, or 2.5  $\mu s$  when ADC10SR = 1. See the device-specific data sheet for parameters.

To calculate the buffer settling time when using an external reference, the formula is:

$$t_{REFBURST} = S_R \times V_{Ref} - 0.5 \mu s$$

Where:

SR = Buffer slew rate (~1 μs/V when ADC10SR = 0 and ~2 μs/V when ADC10SR = 1)

 $V_{Ref}$  = External reference voltage

### 22.2.6 Conversion Modes

The ADC10 has four operating modes selected by the CONSEQx bits as discussed in Table 22-1.

**Table 22-1. Conversion Mode Summary** 

CONSEQx	Mode	Operation
00	Single channel single-conversion	A single channel is converted once.
01	Sequence-of-channels	A sequence of channels is converted once.
10	Repeat single channel	A single channel is converted repeatedly.
11	Repeat sequence-of-channels	A sequence of channels is converted repeatedly.



ADC10 Operation www.ti.com

### 22.2.6.1 Single-Channel Single-Conversion Mode

A single channel selected by INCHx is sampled and converted once. The ADC result is written to ADC10MEM. Figure 22-5 shows the flow of the single-channel, single-conversion mode. When ADC10SC triggers a conversion, successive conversions can be triggered by the ADC10SC bit. When any other trigger source is used, ENC must be toggled between each conversion.

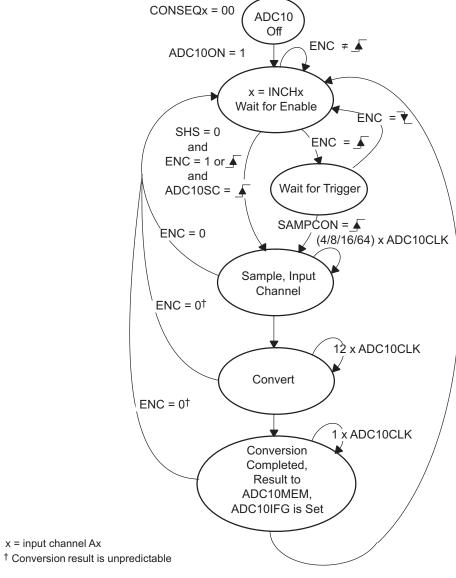


Figure 22-5. Single-Channel Single-Conversion Mode



www.ti.com ADC10 Operation

## 22.2.6.2 Sequence-of-Channels Mode

A sequence of channels is sampled and converted once. The sequence begins with the channel selected by INCHx and decrements to channel A0. Each ADC result is written to ADC10MEM. The sequence stops after conversion of channel A0. Figure 22-6 shows the sequence-of-channels mode. When ADC10SC triggers a sequence, successive sequences can be triggered by the ADC10SC bit. When any other trigger source is used, ENC must be toggled between each sequence.

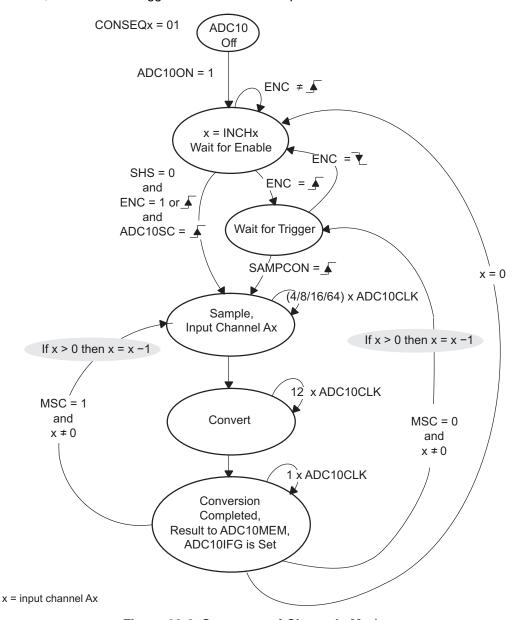


Figure 22-6. Sequence-of-Channels Mode

541



ADC10 Operation www.ti.com

## 22.2.6.3 Repeat-Single-Channel Mode

A single channel selected by INCHx is sampled and converted continuously. Each ADC result is written to ADC10MEM. Figure 22-7 shows the repeat-single-channel mode.

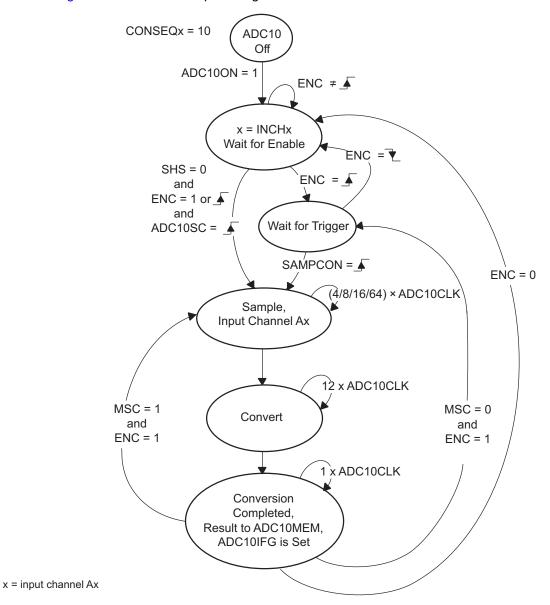


Figure 22-7. Repeat-Single-Channel Mode



www.ti.com ADC10 Operation

## 22.2.6.4 Repeat-Sequence-of-Channels Mode

A sequence of channels is sampled and converted repeatedly. The sequence begins with the channel selected by INCHx and decrements to channel A0. Each ADC result is written to ADC10MEM. The sequence ends after conversion of channel A0, and the next trigger signal re-starts the sequence. Figure 22-8 shows the repeat-sequence-of-channels mode.

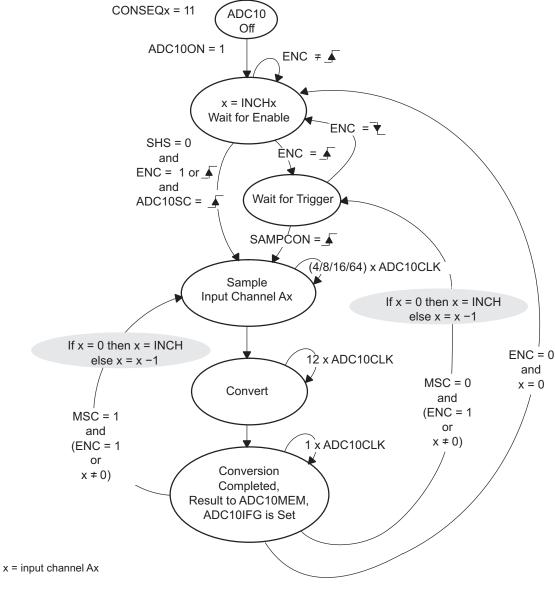


Figure 22-8. Repeat-Sequence-of-Channels Mode



ADC10 Operation www.ti.com

#### 22.2.6.5 Using the MSC Bit

To configure the converter to perform successive conversions automatically and as quickly as possible, a multiple sample and convert function is available. When MSC = 1 and CONSEQx > 0, the first rising edge of the SHI signal triggers the first conversion. Successive conversions are triggered automatically as soon as the prior conversion is completed. Additional rising edges on SHI are ignored until the sequence is completed in the single-sequence mode or until the ENC bit is toggled in repeat-single-channel, or repeated-sequence modes. The function of the ENC bit is unchanged when using the MSC bit.

## 22.2.6.6 Stopping Conversions

Stopping ADC10 activity depends on the mode of operation. The recommended ways to stop an active conversion or conversion sequence are:

- Resetting ENC in single-channel single-conversion mode stops a conversion immediately and the results are unpredictable. For correct results, poll the ADC10BUSY bit until reset before clearing ENC.
- Resetting ENC during repeat-single-channel operation stops the converter at the end of the current conversion.
- Resetting ENC during a sequence or repeat sequence mode stops the converter at the end of the sequence.
- Any conversion mode may be stopped immediately by setting the CONSEQx = 0 and resetting the ENC bit. Conversion data is unreliable.

#### 22.2.7 ADC10 Data Transfer Controller

The ADC10 includes a data transfer controller (DTC) to automatically transfer conversion results from ADC10MEM to other on-chip memory locations. The DTC is enabled by setting the ADC10DTC1 register to a nonzero value.

When the DTC is enabled, each time the ADC10 completes a conversion and loads the result to ADC10MEM, a data transfer is triggered. No software intervention is required to manage the ADC10 until the predefined amount of conversion data has been transferred. Each DTC transfer requires one CPU MCLK. To avoid any bus contention during the DTC transfer, the CPU is halted, if active, for the one MCLK required for the transfer.

A DTC transfer must not be initiated while the ADC10 is busy. Software must ensure that no active conversion or sequence is in progress when the DTC is configured:

```
; ADC10 activity test
BIC.W #ENC,&ADC10CTL0;
busy_test BIT.W #BUSY,&ADC10CTL1;
JNZ busy_test;
MOV.W #xxx,&ADC10SA; Safe
MOV.B #xx,&ADC10DTC1;
; continue setup
```



ADC10 Operation www.ti.com

#### 22.2.7.1 One-Block Transfer Mode

The one-block mode is selected if the ADC10TB is reset. The value n in ADC10DTC1 defines the total number of transfers for a block. The block start address is defined anywhere in the MSP430 address range using the 16-bit register ADC10SA. The block ends at ADC10SA + 2n - 2. The one-block transfer mode is shown in Figure 22-9.

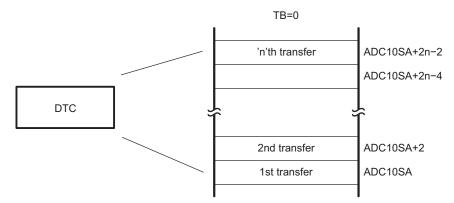


Figure 22-9. One-Block Transfer

The internal address pointer is initially equal to ADC10SA and the internal transfer counter is initially equal to 'n'. The internal pointer and counter are not visible to software. The DTC transfers the word-value of ADC10MEM to the address pointer ADC10SA. After each DTC transfer, the internal address pointer is incremented by two and the internal transfer counter is decremented by one.

The DTC transfers continue with each loading of ADC10MEM, until the internal transfer counter becomes equal to zero. No additional DTC transfers occur until a write to ADC10SA. When using the DTC in the one-block mode, the ADC10IFG flag is set only after a complete block has been transferred. Figure 22-10 shows a state diagram of the one-block mode.

ADC10



ADC10 Operation www.ti.com

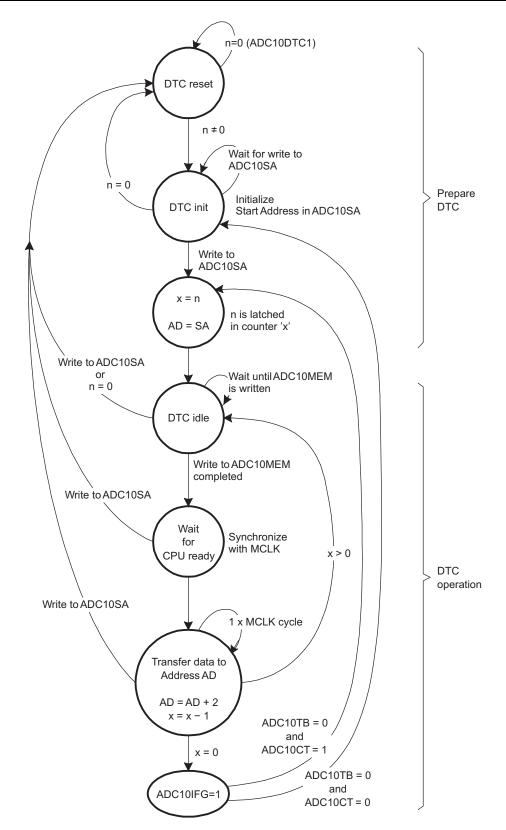


Figure 22-10. State Diagram for Data Transfer Control in One-Block Transfer Mode



www.ti.com ADC10 Operation

#### 22.2.7.2 Two-Block Transfer Mode

The two-block mode is selected if the ADC10TB bit is set. The value n in ADC10DTC1 defines the number of transfers for one block. The address range of the first block is defined anywhere in the MSP430 address range with the 16-bit register ADC10SA. The first block ends at ADC10SA+2n-2. The address range for the second block is defined as SA+2n to SA+4n-2. The two-block transfer mode is shown in Figure 22-11.

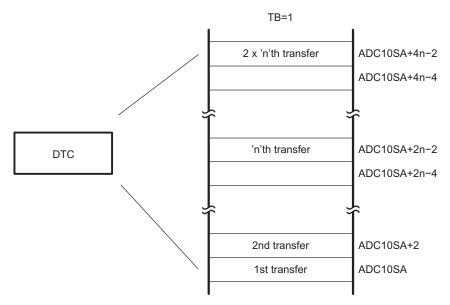


Figure 22-11. Two-Block Transfer

The internal address pointer is initially equal to ADC10SA and the internal transfer counter is initially equal to 'n'. The internal pointer and counter are not visible to software. The DTC transfers the word-value of ADC10MEM to the address pointer ADC10SA. After each DTC transfer the internal address pointer is incremented by two and the internal transfer counter is decremented by one.

The DTC transfers continue, with each loading of ADC10MEM, until the internal transfer counter becomes equal to zero. At this point, block one is full and both the ADC10IFG flag the ADC10B1 bit are set. The user can test the ADC10B1 bit to determine that block one is full.

The DTC continues with block two. The internal transfer counter is automatically reloaded with 'n'. At the next load of the ADC10MEM, the DTC begins transferring conversion results to block two. After n transfers have completed, block two is full. The ADC10IFG flag is set and the ADC10B1 bit is cleared. User software can test the cleared ADC10B1 bit to determine that block two is full. Figure 22-12 shows a state diagram of the two-block mode.

ADC10



ADC10 Operation www.ti.com

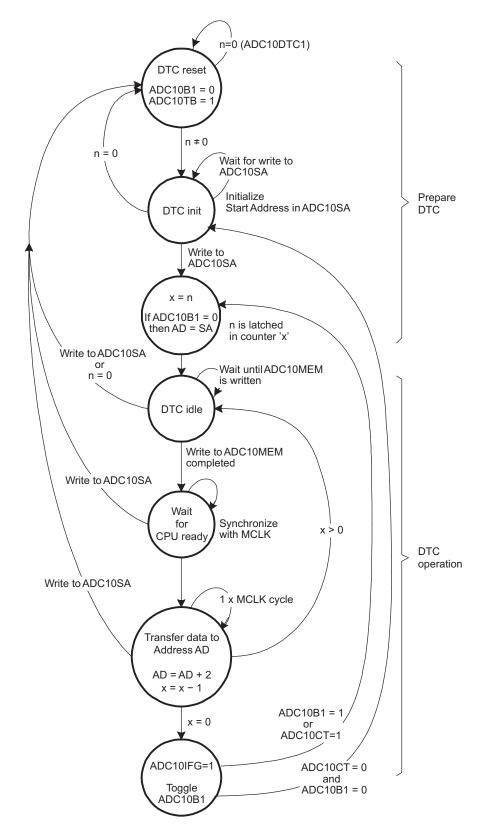


Figure 22-12. State Diagram for Data Transfer Control in Two-Block Transfer Mode



www.ti.com ADC10 Operation

#### 22.2.7.3 Continuous Transfer

A continuous transfer is selected if ADC10CT bit is set. The DTC does not stop after block one (in one-block mode) or block two (in two-block mode) has been transferred. The internal address pointer and transfer counter are set equal to ADC10SA and n respectively. Transfers continue starting in block one. If the ADC10CT bit is reset, DTC transfers cease after the current completion of transfers into block one (in one-block mode) or block two (in two-block mode) have been transferred.

#### 22.2.7.4 DTC Transfer Cycle Time

For each ADC10MEM transfer, the DTC requires one or two MCLK clock cycles to synchronize, one for the actual transfer (while the CPU is halted), and one cycle of wait time. Because the DTC uses MCLK, the DTC cycle time is dependent on the MSP430 operating mode and clock system setup.

If the MCLK source is active but the CPU is off, the DTC uses the MCLK source for each transfer, without re-enabling the CPU. If the MCLK source is off, the DTC temporarily restarts MCLK, sourced with DCOCLK, only during a transfer. The CPU remains off, and MCLK is again turned off after the DTC transfer. The maximum DTC cycle time for all operating modes is show in Table 22-2.

CPU Operating Mode	Clock Source	Maximum DTC Cycle Time
Active mode	MCLK = DCOCLK	3 MCLK cycles
Active mode	MCLK = LFXT1CLK	3 MCLK cycles
Low-power mode LPM0/1	MCLK = DCOCLK	4 MCLK cycles
Low-power mode LPM3/4	MCLK = DCOCLK	4 MCLK cycles + 2 μs <sup>(1)</sup>
Low-power mode LPM0/1	MCLK = LFXT1CLK	4 MCLK cycles
Low-power mode LPM3	MCLK = LFXT1CLK	4 MCLK cycles
Low-power mode LPM4	MCLK = LFXT1CLK	4 MCLK cycles + 2 μs <sup>(1)</sup>

**Table 22-2. Maximum DTC Cycle Time** 

#### 22.2.8 Using the Integrated Temperature Sensor

To use the on-chip temperature sensor, select the analog input channel INCHx = 1010. Any other configuration is done as if an external channel was selected, including reference selection, conversion-memory selection, etc.

The typical temperature sensor transfer function is shown in Figure 22-13. When using the temperature sensor, the sample period must be greater than 30  $\mu$ s. The temperature sensor offset error is large. Deriving absolute temperature values in the application requires calibration. See the device-specific data sheet for the parameters. See Section 24.2.2.1 for the calibration equations.

Selecting the temperature sensor automatically turns on the on-chip reference generator as a voltage source for the temperature sensor. However, it does not enable the  $V_{REF+}$  output or affect the reference selections for the conversion. The reference choices for converting the temperature sensor are the same as with any other channel.

The additional 2 µs are needed to start the DCOCLK. See the device-specific data sheet for parameters.



ADC10 Operation www.ti.com

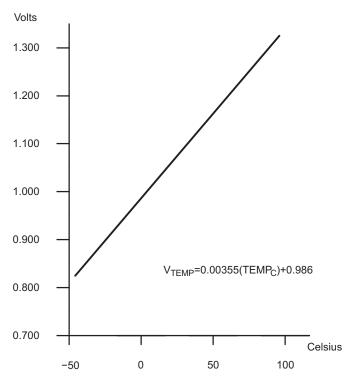


Figure 22-13. Typical Temperature Sensor Transfer Function

## 22.2.9 ADC10 Grounding and Noise Considerations

As with any high-resolution ADC, appropriate printed-circuit-board layout and grounding techniques should be followed to eliminate ground loops, unwanted parasitic effects, and noise.

Ground loops are formed when return current from the A/D flows through paths that are common with other analog or digital circuitry. If care is not taken, this current can generate small, unwanted offset voltages that can add to or subtract from the reference or input voltages of the A/D converter. The connections shown in Figure 22-14 and Figure 22-15 help avoid this.

In addition to grounding, ripple and noise spikes on the power supply lines due to digital switching or switching power supplies can corrupt the conversion result. A noise-free design is important to achieve high accuracy.

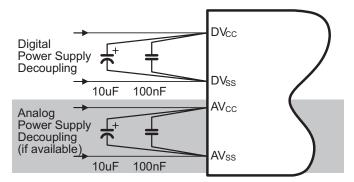


Figure 22-14. ADC10 Grounding and Noise Considerations (Internal V<sub>REF</sub>)



www.ti.com ADC10 Operation

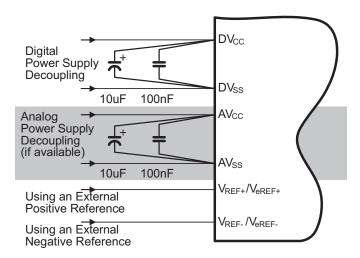


Figure 22-15. ADC10 Grounding and Noise Considerations (External V<sub>REF</sub>)

## 22.2.10 ADC10 Interrupts

One interrupt and one interrupt vector are associated with the ADC10 as shown in Figure 22-16. When the DTC is not used (ADC10DTC1 = 0), ADC10IFG is set when conversion results are loaded into ADC10MEM. When DTC is used (ADC10DTC1 > 0), ADC10IFG is set when a block transfer completes and the internal transfer counter n = 0. If both the ADC10IE and the GIE bits are set, then the ADC10IFG flag generates an interrupt request. The ADC10IFG flag is automatically reset when the interrupt request is serviced, or it may be reset by software.

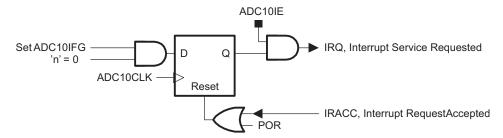


Figure 22-16. ADC10 Interrupt System



ADC10 Registers www.ti.com

# 22.3 ADC10 Registers

The ADC10 registers are listed in Table 22-3.

## Table 22-3. ADC10 Registers

Register	Short Form	Register Type	Address	Initial State
ADC10 input enable register 0	ADC10AE0	Read/write	04Ah	Reset with POR
ADC10 input enable register 1	ADC10AE1	Read/write	04Bh	Reset with POR
ADC10 control register 0	ADC10CTL0	Read/write	01B0h	Reset with POR
ADC10 control register 1	ADC10CTL1	Read/write	01B2h	Reset with POR
ADC10 memory	ADC10MEM	Read	01B4h	Unchanged
ADC10 data transfer control register 0	ADC10DTC0	Read/write	048h	Reset with POR
ADC10 data transfer control register 1	ADC10DTC1	Read/write	049h	Reset with POR
ADC10 data transfer start address	ADC10SA	Read/write	01BCh	0200h with POR



www.ti.com ADC10 Registers

## 22.3.1 ADC10CTL0, ADC10 Control Register 0

15	14	13	12	11	10	9	8	
	SREFx			ADC10SHTx		REFOUT	REFBURST	
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	
7	6	5	4	3	2	1	0	
MSC	REF2_5V	REFON	ADC10ON	ADC10IE	ADC10IFG	ENC	ADC10SC	
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	
	Can be modified only when ENC = 0							

**SREFx** Bits 15-13 Select reference.

 $V_{R+} = V_{CC} \text{ and } V_{R-} = V_{SS}$ 

001  $V_{R+} = V_{REF+}$  and  $V_{R-} = V_{SS}$ 

 $V_{R+} = V_{eREF+}$  and  $V_{R-} = V_{SS}$ . Devices with  $V_{eREF+}$  only.

011  $V_{R+}$  = Buffered  $V_{eREF+}$  and  $V_{R-}$  =  $V_{SS}$ . Devices with  $V_{eREF+}$  pin only.

 $V_{R+} = V_{CC}$  and  $V_{R-} = V_{REF}/V_{eREF-}$ . Devices with  $V_{eREF-}$  pin only.

101  $V_{R+} = V_{REF+}$  and  $V_{R-} = V_{REF-}/V_{eREF-}$ . Devices with  $V_{eREF+/-}$  pins only.

110  $V_{R+} = V_{eREF+}$  and  $V_{R-} = V_{REF}/V_{eREF-}$ . Devices with  $V_{eREF+/-}$  pins only.

111  $V_{R+}$  = Buffered  $V_{eREF+}$  and  $V_{R-}$  =  $V_{REF-}$ . Devices with  $V_{eREF+}$  pins only.

ADC10SHTx Bits 12-11 ADC10 sample-and-hold time

00 4 × ADC10CLKs 01 8 × ADC10CLKs

10 16 × ADC10CLKs
 11 64 × ADC10CLKs

ADC10SR Bit 10 ADC10 sampling rate. This bit selects the reference buffer drive capability for the maximum sampling rate.

Setting ADC10SR reduces the current consumption of the reference buffer.

Reference buffer supports up to ~200 ksps

1 Reference buffer supports up to ~50 ksps

**REFOUT** Bit 9 Reference output

0 Reference output off

1 Reference output on. Devices with  $V_{eREF+} / V_{REF+}$  pin only.

REFBURST Bit 8 Reference burst.

0 Reference buffer on continuously

Reference buffer on only during sample-and-conversion

MSC Bit 7 Multiple sample and conversion. Valid only for sequence or repeated modes.

The sampling requires a rising edge of the SHI signal to trigger each sample-and-conversion.

The first rising edge of the SHI signal triggers the sampling timer, but further sample-and-conversions are performed automatically as soon as the prior conversion is completed

**REF2\_5V** Bit 6 Reference-generator voltage. REFON must also be set.

0 1.5 V 1 2.5 V

**REFON** Bit 5 Reference generator on

0 Reference off

1 Reference on

ADC100N Bit 4 ADC10 on

0 ADC10 off

1 ADC10 on ADC10IE Bit 3 ADC10 interrupt enable

0 Interrupt disabled

1 Interrupt enabled



ADC10 Registers www.ti.com ADC10 interrupt flag. This bit is set if ADC10MEM is loaded with a conversion result. It is automatically reset when the interrupt request is accepted, or it may be reset by software. When using the DTC this flag is set ADC10IFG Bit 2 when a block of transfers is completed. 0 No interrupt pending 1 Interrupt pending **ENC** Bit 1 Enable conversion 0 ADC10 disabled ADC10 enabled ADC10SC Bit 0 Start conversion. Software-controlled sample-and-conversion start. ADC10SC and ENC may be set together with one instruction. ADC10SC is reset automatically.

No sample-and-conversion start

Start sample-and-conversion

0

1



ADC10 Registers www.ti.com

# 22.3.2 ADC10CTL1, ADC10 Control Register 1

15	14	1	3 12	11	10	9	8		
10		INCHx	0 12		SHSx	ADC10DF	ISSH		
rw-(0)	rw-(0)		-(0) rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)		
						` '	, ,		
7	6		5 4	3	2	1	0		
(0)	ADC10DI			DC10SSELx		SEQx	ADC10BUSY		
rw-(0)	rw-(0)		-(0) rw-(0)	rw-(0)	rw-(0)	rw-(0)	r-0		
	Can be mod	dified only whe	en ENC = 0						
INCHx	sequence of conversions. Only available ADC channels should be selected. See device spec								
		0000 A0							
		0001 A1							
		0010 A2							
		0011 A3							
		0100 A4							
		0101 A5							
		0110 A6							
		0111 A7							
		1000 V <sub>eR</sub>							
			F/V <sub>eREF</sub> -						
			nperature sensor						
			c - V <sub>SS</sub> ) / 2	00400500					
			c - V <sub>SS</sub> ) / 2, A12 on M						
			c - V <sub>SS</sub> ) / 2, A13 on M						
			c - V <sub>SS</sub> ) / 2, A14 on M						
0110	D'(- 44 40		c - V <sub>SS</sub> ) / 2, A15 on M	SP430F22xx device	es				
SHSx	Bits 11-10	•							
		00 ADC10SC bit							
		01 Timer_A.OUT1 <sup>(1)</sup> 10 Timer_A.OUT0 <sup>(1)</sup>							
				OUT4 on MCD420	E20v0 MCD420C2v2	1 and MCD 120C	2v20 daviaga)(1)		
A DC40DE	D:+ 0			0011 00 1057430	F20x0, MSP430G2x3	1, and MSP430G	2x30 devices)		
ADC10DF	Bit 9	ADC10 data 0 Stra							
			aight binary						
ISSH	Bit 8		complement sample-and-hold						
Ю	DIL O	-	sample-and-noid sample-input signal	is not inverted					
			sample-input signal						
ADC10DIVx	Bits 7-5	ADC10 cloc		is inverted.					
ADCIODIVA	טונס וי-ט	000 /1	k dividei						
		000 /1							
		010 /3							
		010 /3							
		100 /5							
		100 /5							
		110 /7							
		110 //							
ADC10SSELx	Bits 4-3		k source select						
ADOTOSSELX	DII3 4-3		C10OSC						
		00 AD							
		10 MC							
		TO IVIC	LIX						

**SMCLK**  $^{(1)}$   $\;$  Timer triggers are from Timer0\_Ax if more than one timer module exists on the device.

11

555

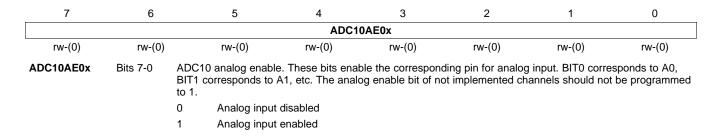
SLAU144J-December 2004-Revised July 2013



ADC10 Registers www.ti.com

**CONSEQ**x Bits 2-1 Conversion sequence mode select 00 Single-channel-single-conversion 01 Sequence-of-channels 10 Repeat-single-channel Repeat-sequence-of-channels 11 **ADC10BUSY** Bit 0 ADC10 busy. This bit indicates an active sample or conversion operation No operation is active. A sequence, sample, or conversion is active.

## 22.3.3 ADC10AE0, Analog (Input) Enable Control Register 0



## 22.3.4 ADC10AE1, Analog (Input) Enable Control Register 1 (MSP430F22xx only)

7	6	5	4	3	2	1	0		
	Α	DC10AE1x			Reserved				
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)		
ADC10AE1x	DC10AE1x Bits 7-4 ADC10 analog enable. These bits enable the corresponding pin for analog input. BIT4 corresponds to A12 BIT5 corresponds to A13, BIT6 corresponds to A14, and BIT7 corresponds to A15. The analog enable bit not implemented channels should not be programmed to 1.								
		0 Analog input	t disabled						
		1 Analog input	t enabled						
Reserved	Bits 3-0	Reserved							

## 22.3.5 ADC10MEM, Conversion-Memory Register, Binary Format

15	14	13	12	11	10	9	8
0	0	0	0	0	0	Conversion Results	
rO	r0	r0	rO	rO	rO	r	r
7	6	5	4	3	2	1	0
			Conversion	n Results			
r	r	r	r	r	r	r	r
<b>Conversion</b> Bits 15-0 The 10-bit conversion results are right justified, straight-binary format. Bit 9 is the MSB. Bits 15-10 are always 0.						15-10 are	



www.ti.com ADC10 Registers

## 22.3.6 ADC10MEM, Conversion-Memory Register, 2s Complement Format

15	14	13	12	11	10	9	8			
	Conversion Results									
r	r	r	r	r	r	r	r			
7	6	5	4	3	2	1	0			
Conversio	n Results	0	0	0	0	0	0			
r	r	r0	r0	r0	r0	r0	r0			

Conversion Results

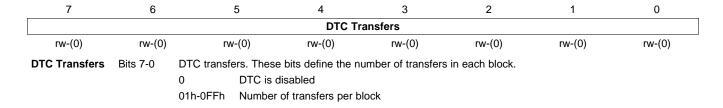
Bits 15-0

The 10-bit conversion results are left-justified, 2s complement format. Bit 15 is the MSB. Bits 5-0 are always

## 22.3.7 ADC10DTC0, Data Transfer Control Register 0

7	6		5	4	3	2	1	0	
		Reserved			ADC10TB	ADC10CT	ADC10B1	ADC10FETCH	
r0	r0		r0	r0	rw-(0)	rw-(0)	r-(0)	rw-(0)	
Reserved	Bits 7-4	Reserved.	Always rea	d as 0.					
ADC10TB	Bit 3	3 ADC10 two-block mode							
		0 O	ne-block tra	nsfer mode					
		1 T	wo-block tra	nsfer mode					
ADC10CT	Bit 2	ADC10 co	ntinuous tra	nsfer					
		Data transfer stops when one block (one-block mode) or two blocks (two-block mode) have completed.							
			ata is transf written to.	erred continuous	sly. DTC operation	is stopped only if	ADC10CT clear	ed, or ADC10SA	
ADC10B1	Bit 1	ADC10 block one. This bit indicates for two-block mode which block is filled with ADC10 conversion results. ADC10B1 is valid only after ADC10IFG has been set the first time during DTC operation. ADC10TB must also be set.							
		0 B	lock 2 is fille	ed					
		1 B	lock 1 is fille	ed					
ADC10FETCH	Bit 0	This bit sh	ould normal	ly be reset.					

## 22.3.8 ADC10DTC1, Data Transfer Control Register 1



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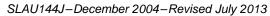
ADC10 Registers www.ti.com

# 22.3.9 ADC10SA, Start Address Register for Data Transfer

15	14	13	12	11	10	9	8				
	ADC10SAx										
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)				
7	6	5	4	3	2	1	0				
			ADC10SAx				0				
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	r0				
ADC10SAx		ADC10 start address to initiate DTC transf		ne start address fo	or the DTC. A write	e to register ADC1	IOSA is required				

Unused Bit 0 Unused, Read only. Always read as 0.







# ADC12

The ADC12 module is a high-performance 12-bit analog-to-digital converter. This chapter describes the ADC12 of the MSP430x2xx device family.

Topic Page

23.1	ADC12 Introduction	<b>560</b>
23.2	ADC12 Operation	<b>562</b>
23.3	ADC12 Registers	574
	<u> </u>	



ADC12 Introduction www.ti.com

#### 23.1 ADC12 Introduction

The ADC12 module supports fast 12-bit analog-to-digital conversions. The module implements a 12-bit SAR core, sample select control, reference generator, and a 16-word conversion-and-control buffer. The conversion-and-control buffer allows up to 16 independent ADC samples to be converted and stored without any CPU intervention.

#### ADC12 features include:

- Greater than 200-ksps maximum conversion rate
- Monotonic 12-bit converter with no missing codes
- Sample-and-hold with programmable sampling periods controlled by software or timers
- Conversion initiation by software, Timer\_A, or Timer\_B
- Software selectable on-chip reference voltage generation (1.5 V or 2.5 V)
- · Software selectable internal or external reference
- · Eight individually configurable external input channels
- Conversion channels for internal temperature sensor, AV<sub>CC</sub>, and external references
- Independent channel-selectable reference sources for both positive and negative references
- · Selectable conversion clock source
- Single-channel, repeat-single-channel, sequence, and repeat-sequence conversion modes
- ADC core and reference voltage can be powered down separately
- Interrupt vector register for fast decoding of 18 ADC interrupts
- 16 conversion-result storage registers

The block diagram of ADC12 is shown in Figure 23-1.



www.ti.com ADC12 Introduction

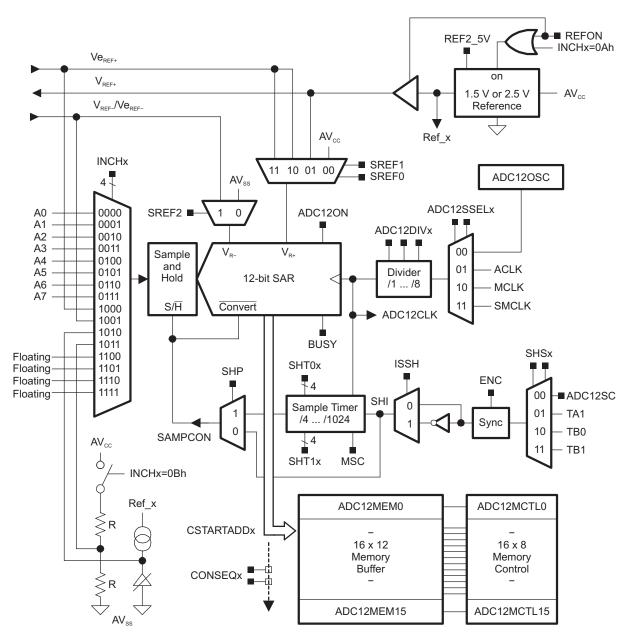


Figure 23-1. ADC12 Block Diagram



ADC12 Operation www.ti.com

#### 23.2 ADC12 Operation

The ADC12 module is configured with user software. The setup and operation of the ADC12 is discussed in the following sections.

#### 23.2.1 12-Bit ADC Core

The ADC core converts an analog input to its 12-bit digital representation and stores the result in conversion memory. The core uses two programmable/selectable voltage levels ( $V_{R+}$  and  $V_{R-}$ ) to define the upper and lower limits of the conversion. The digital output ( $N_{ADC}$ ) is full scale (0FFFh) when the input signal is equal to or higher than  $V_{R+}$ , and the digital output is zero when the input signal is equal to or lower than  $V_{R-}$ . The input channel and the reference voltage levels ( $V_{R+}$  and  $V_{R-}$ ) are defined in the conversion-control memory. The conversion formula for the ADC result  $V_{ADC}$  is:

$$N_{ADC} = 4095 \times \frac{V_{IN} - V_{R-}}{V_{R+} - V_{R-}}$$

The ADC12 core is configured by two control registers, ADC12CTL0 and ADC12CTL1. The core is enabled with the ADC12ON bit. The ADC12 can be turned off when not in use to save power. With few exceptions, the ADC12 control bits can only be modified when ENC = 0. ENC must be set to 1 before any conversion can take place.

#### 23.2.1.1 Conversion Clock Selection

The ADC12CLK is used both as the conversion clock and to generate the sampling period when the pulse sampling mode is selected. The ADC12 source clock is selected using the ADC12SSELx bits and can be divided from 1 through 8 using the ADC12DIVx bits. Possible ADC12CLK sources are SMCLK, MCLK, ACLK, and an internal oscillator ADC12OSC.

The ADC12OSC is generated internally and is in the 5-MHz range, but the frequency varies with individual devices, supply voltage, and temperature. See the device-specific data sheet for the ADC12OSC specification.

The application must ensure that the clock chosen for ADC12CLK remains active until the end of a conversion. If the clock is removed during a conversion, the operation does not complete and any result is invalid.

#### 23.2.2 ADC12 Inputs and Multiplexer

The eight external and four internal analog signals are selected as the channel for conversion by the analog input multiplexer. The input multiplexer is a break-before-make type to reduce input-to-input noise injection that can result from channel switching (see Figure 23-2). The input multiplexer is also a T-switch to minimize the coupling between channels. Channels that are not selected are isolated from the A/D, and the intermediate node is connected to analog ground (AV $_{\rm SS}$ ) so that the stray capacitance is grounded to help eliminate crosstalk.

The ADC12 uses the charge redistribution method. When the inputs are internally switched, the switching action may cause transients on the input signal. These transients decay and settle before causing errant conversion.

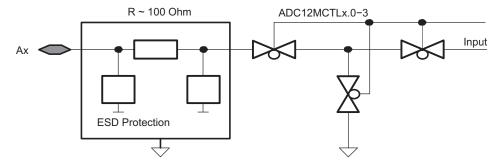


Figure 23-2. Analog Multiplexer



www.ti.com ADC12 Operation

#### 23.2.2.1 Analog Port Selection

The ADC12 inputs are multiplexed with the port P6 pins, which are digital CMOS gates. When analog signals are applied to digital CMOS gates, parasitic current can flow from VCC to GND. This parasitic current occurs if the input voltage is near the transition level of the gate. Disabling the port pin buffer eliminates the parasitic current flow and, therefore, reduces overall current consumption. The P6SELx bits provide the ability to disable the port pin input and output buffers.

```
; P6.0 and P6.1 configured for analog input
BIS.B #3h,&P6SEL ; P6.1 and P6.0 ADC12 function
```

## 23.2.3 Voltage Reference Generator

The ADC12 module contains a built-in voltage reference with two selectable voltage levels, 1.5 V and 2.5 V. Either of these reference voltages may be used internally and externally on pin V<sub>REF+</sub>.

Setting REFON = 1 enables the internal reference. When REF2\_5V = 1, the internal reference is 2.5 V. When REF2\_5V = 0, the reference is 1.5 V. The reference can be turned off to save power when not in use.

For proper operation, the internal voltage reference generator must be supplied with storage capacitance across  $V_{REF+}$  and  $AV_{SS}$ . The recommended storage capacitance is a parallel combination of 10- $\mu$ F and 0.1- $\mu$ F capacitors. From turn-on, a maximum of 17 ms must be allowed for the voltage reference generator to bias the recommended storage capacitors. If the internal reference generator is not used for the conversion, the storage capacitors are not required.

#### NOTE: Reference Decoupling

Approximately 200 μA is required from *any* reference used by the ADC12 while the two LSBs are being resolved during a conversion. A parallel combination of 10-μF and 0.1-μF capacitors is recommended for *any* reference as shown in Figure 23-11.

External references may be supplied for  $V_{R+}$  and  $V_{R-}$  through pins  $V_{eREF+}$  and  $V_{REF-}$  respectively.

## 23.2.4 Sample and Conversion Timing

An analog-to-digital conversion is initiated with a rising edge of the sample input signal SHI. The source for SHI is selected with the SHSx bits and includes the following:

- The ADC12SC bit
- The Timer A Output Unit 1
- The Timer\_B Output Unit 0
- The Timer\_B Output Unit 1

The polarity of the SHI signal source can be inverted with the ISSH bit. The SAMPCON signal controls the sample period and start of conversion. When SAMPCON is high, sampling is active. The high-to-low SAMPCON transition starts the analog-to-digital conversion, which requires 13 ADC12CLK cycles. Two different sample-timing methods are defined by control bit SHP, extended sample mode and pulse mode.



ADC12 Operation www.ti.com

#### 23.2.4.1 Extended Sample Mode

The extended sample mode is selected when SHP = 0. The SHI signal directly controls SAMPCON and defines the length of the sample period  $t_{\text{sample}}$ . When SAMPCON is high, sampling is active. The high-to-low SAMPCON transition starts the conversion after synchronization with ADC12CLK (see Figure 23-3).

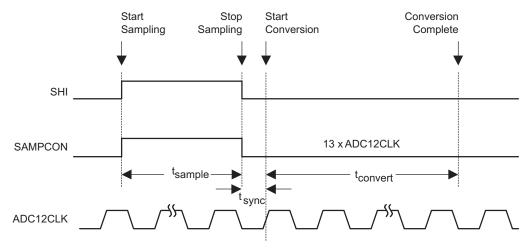


Figure 23-3. Extended Sample Mode

#### 23.2.4.2 Pulse Sample Mode

The pulse sample mode is selected when SHP = 1. The SHI signal is used to trigger the sampling timer. The SHT0x and SHT1x bits in ADC12CTL0 control the interval of the sampling timer that defines the SAMPCON sample period  $t_{\text{sample}}$ . The sampling timer keeps SAMPCON high after synchronization with AD12CLK for a programmed interval  $t_{\text{sample}}$ . The total sampling time is  $t_{\text{sample}}$  plus  $t_{\text{sync}}$  (see Figure 23-4).

The SHTx bits select the sampling time in 4x multiples of ADC12CLK. SHT0x selects the sampling time for ADC12MCTL0 to 7 and SHT1x selects the sampling time for ADC12MCTL8 to 15.

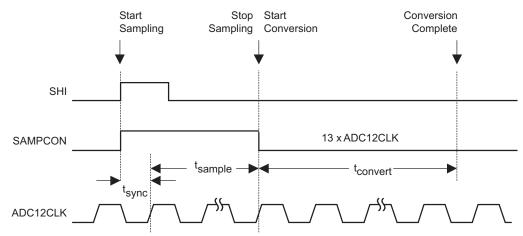


Figure 23-4. Pulse Sample Mode



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#### 23.2.4.3 Sample Timing Considerations

When SAMPCON = 0, all Ax inputs are high impedance. When SAMPCON = 1, the selected Ax input can be modeled as an RC low-pass filter during the sampling time  $t_{\text{sample}}$ , as shown in Figure 23-5. An internal MUX-on input resistance  $R_{\text{I}}$  (maximum of 2 k $\Omega$ ) in series with capacitor  $C_{\text{I}}$  (maximum of 40 pF) is seen by the source. The capacitor  $C_{\text{I}}$  voltage ( $V_{\text{C}}$ ) must be charged to within 1/2 LSB of the source voltage ( $V_{\text{S}}$ ) for an accurate 12-bit conversion.

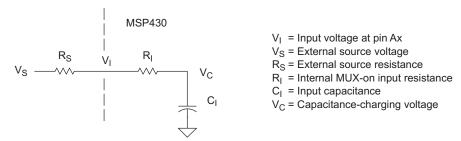


Figure 23-5. Analog Input Equivalent Circuit

The resistance of the source  $R_s$  and  $R_l$  affect  $t_{sample}$ . The following equation can be used to calculate the minimum sampling time  $t_{sample}$  for a 12-bit conversion:

$$t_{sample} > (R_S + R_I) \times ln(2^{13}) \times C_I + 800 \text{ ns}$$

Substituting the values for R<sub>I</sub> and C<sub>I</sub> given above, the equation becomes:

$$t_{sample} > (R_S + 2 k\Omega) \times 9.011 \times 40 pF + 800 ns$$

For example, if  $R_s$  is 10 k $\Omega$ ,  $t_{sample}$  must be greater than 5.13  $\mu s$ .

## 23.2.5 Conversion Memory

There are 16 ADC12MEMx conversion memory registers to store conversion results. Each ADC12MEMx is configured with an associated ADC12MCTLx control register. The SREFx bits define the voltage reference and the INCHx bits select the input channel. The EOS bit defines the end of sequence when a sequential conversion mode is used. A sequence rolls over from ADC12MEM15 to ADC12MEM0 when the EOS bit in ADC12MCTL15 is not set.

The CSTARTADDx bits define the first ADC12MCTLx used for any conversion. If the conversion mode is single-channel or repeat-single-channel the CSTARTADDx points to the single ADC12MCTLx to be used.

If the conversion mode selected is either sequence-of-channels or repeat-sequence-of-channels, CSTARTADDx points to the first ADC12MCTLx location to be used in a sequence. A pointer, not visible to software, is incremented automatically to the next ADC12MCTLx in a sequence when each conversion completes. The sequence continues until an EOS bit in ADC12MCTLx is processed; this is the last control byte processed.

When conversion results are written to a selected ADC12MEMx, the corresponding flag in the ADC12IFGx register is set.

#### 23.2.6 ADC12 Conversion Modes

The ADC12 has four operating modes selected by the CONSEQx bits as shown in Table 23-1.

 CONSEQx
 Mode
 Operation

 00
 Single channel single-conversion
 A single channel is converted once.

 01
 Sequence-of-channels
 A sequence of channels is converted once.

 10
 Repeat-single-channel
 A single channel is converted repeatedly.

 11
 Repeat-sequence-of-channels
 A sequence of channels is converted repeatedly.

**Table 23-1. Conversion Mode Summary** 



ADC12 Operation www.ti.com

## 23.2.6.1 Single-Channel Single-Conversion Mode

A single channel is sampled and converted once. The ADC result is written to the ADC12MEMx defined by the CSTARTADDx bits. Figure 23-6 shows the flow of the single-channel, single-conversion mode. When ADC12SC triggers a conversion, successive conversions can be triggered by the ADC12SC bit. When any other trigger source is used, ENC must be toggled between each conversion.

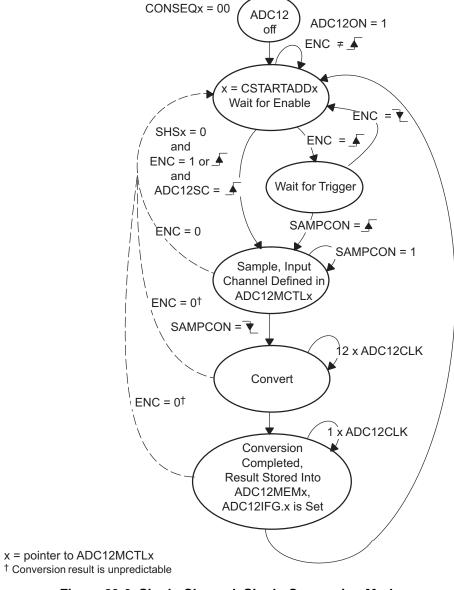


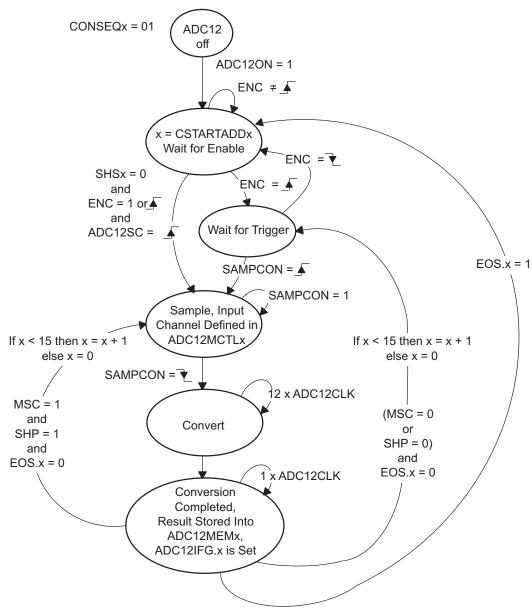
Figure 23-6. Single-Channel, Single-Conversion Mode



www.ti.com ADC12 Operation

## 23.2.6.2 Sequence-of-Channels Mode

A sequence of channels is sampled and converted once. The ADC results are written to the conversion memories starting with the ADCMEMx defined by the CSTARTADDx bits. The sequence stops after the measurement of the channel with a set EOS bit. Figure 23-7 shows the sequence-of-channels mode. When ADC12SC triggers a sequence, successive sequences can be triggered by the ADC12SC bit. When any other trigger source is used, ENC must be toggled between each sequence.



x = pointer to ADC12MCTLx

Figure 23-7. Sequence-of-Channels Mode



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## 23.2.6.3 Repeat-Single-Channel Mode

A single channel is sampled and converted continuously. The ADC results are written to the ADC12MEMx defined by the CSTARTADDx bits. It is necessary to read the result after the completed conversion, because only one ADC12MEMx memory is used and is overwritten by the next conversion. Figure 23-8 shows repeat-single-channel mode.

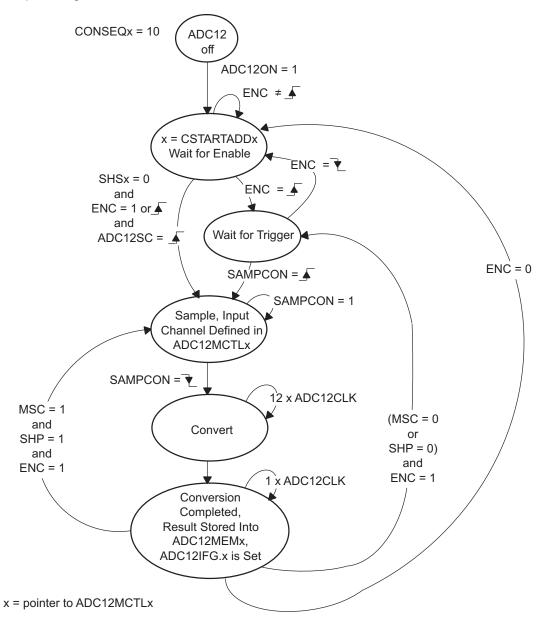


Figure 23-8. Repeat-Single-Channel Mode



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## 23.2.6.4 Repeat-Sequence-of-Channels Mode

A sequence of channels is sampled and converted repeatedly. The ADC results are written to the conversion memories starting with the ADC12MEMx defined by the CSTARTADDx bits. The sequence ends after the measurement of the channel with a set EOS bit, and the next trigger signal re-starts the sequence. Figure 23-9 shows the repeat-sequence-of-channels mode.

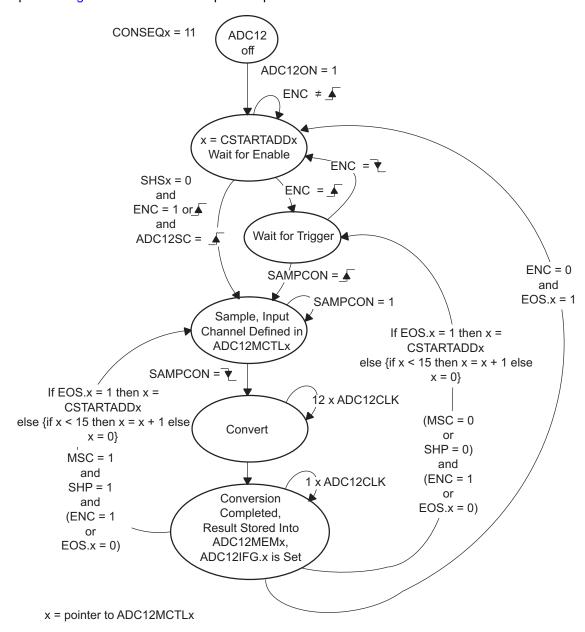


Figure 23-9. Repeat-Sequence-of-Channels Mode

ADC12



ADC12 Operation www.ti.com

#### 23.2.6.5 Using the Multiple Sample and Convert (MSC) Bit

To configure the converter to perform successive conversions automatically and as quickly as possible, a multiple sample and convert function is available. When MSC = 1, CONSEQx > 0, and the sample timer is used, the first rising edge of the SHI signal triggers the first conversion. Successive conversions are triggered automatically as soon as the prior conversion is completed. Additional rising edges on SHI are ignored until the sequence is completed in the single-sequence mode or until the ENC bit is toggled in repeat-single-channel or repeated-sequence modes. The function of the ENC bit is unchanged when using the MSC bit.

## 23.2.6.6 Stopping Conversions

Stopping ADC12 activity depends on the mode of operation. The recommended ways to stop an active conversion or conversion sequence are:

- Resetting ENC in single-channel single-conversion mode stops a conversion immediately and the results are unpredictable. For correct results, poll the busy bit until it is reset before clearing ENC.
- Resetting ENC during repeat-single-channel operation stops the converter at the end of the current conversion.
- Resetting ENC during a sequence or repeat-sequence mode stops the converter at the end of the sequence.
- Any conversion mode may be stopped immediately by setting the CONSEQx = 0 and resetting ENC bit. In this case, conversion data are unreliable.

#### NOTE: No EOS Bit Set For Sequence

If no EOS bit is set and a sequence mode is selected, resetting the ENC bit does not stop the sequence. To stop the sequence, first select a single-channel mode and then reset ENC.

#### 23.2.7 Using the Integrated Temperature Sensor

To use the on-chip temperature sensor, select the analog input channel INCHx = 1010. Any other configuration is done as if an external channel was selected, including reference selection, conversion-memory selection, etc.

The typical temperature sensor transfer function is shown in Figure 23-10. When using the temperature sensor, the sample period must be greater than 30  $\mu$ s. The temperature sensor offset error can be large and needs to be calibrated for most applications. See the device-specific data sheet for parameters. See Section 24.2.2.1 for the calibration equations.

Selecting the temperature sensor automatically turns on the on-chip reference generator as a voltage source for the temperature sensor. However, it does not enable the  $V_{REF+}$  output or affect the reference selections for the conversion. The reference choices for converting the temperature sensor are the same as with any other channel.



www.ti.com ADC12 Operation

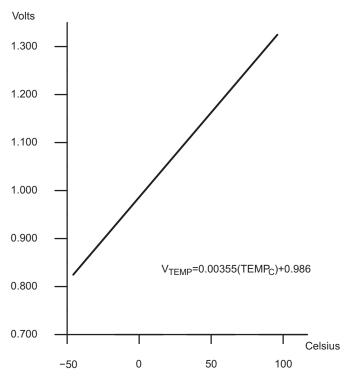


Figure 23-10. Typical Temperature Sensor Transfer Function

## 23.2.8 ADC12 Grounding and Noise Considerations

As with any high-resolution ADC, appropriate printed-circuit-board layout and grounding techniques should be followed to eliminate ground loops, unwanted parasitic effects, and noise.

Ground loops are formed when return current from the A/D flows through paths that are common with other analog or digital circuitry. If care is not taken, this current can generate small unwanted offset voltages that can add to or subtract from the reference or input voltages of the A/D converter. The connections shown in Figure 23-11 help avoid this.

In addition to grounding, ripple and noise spikes on the power supply lines due to digital switching or switching power supplies can corrupt the conversion result. A noise-free design using separate analog and digital ground planes with a single-point connection is recommend to achieve high accuracy.



ADC12 Operation www.ti.com

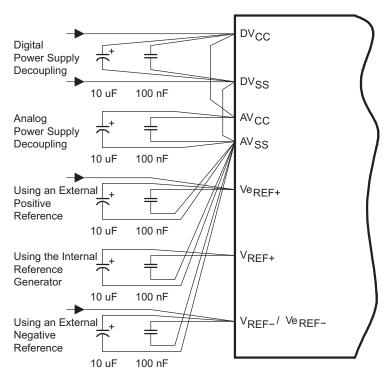


Figure 23-11. ADC12 Grounding and Noise Considerations

## 23.2.9 ADC12 Interrupts

The ADC12 has 18 interrupt sources:

- ADC12IFG0 to ADC12IFG15
- ADC12OV, ADC12MEMx overflow
- ADC12TOV, ADC12 conversion time overflow

The ADC12IFGx bits are set when their corresponding ADC12MEMx memory register is loaded with a conversion result. An interrupt request is generated if the corresponding ADC12IEx bit and the GIE bit are set. The ADC12OV condition occurs when a conversion result is written to any ADC12MEMx before its previous conversion result was read. The ADC12TOV condition is generated when another sample-and-conversion is requested before the current conversion is completed. The DMA is triggered after the conversion in single channel modes or after the completion of a sequence-of-channel modes.

#### 23.2.9.1 ADC12IV, Interrupt Vector Generator

All ADC12 interrupt sources are prioritized and combined to source a single interrupt vector. The interrupt vector register ADC12IV is used to determine which enabled ADC12 interrupt source requested an interrupt.

The highest priority enabled ADC12 interrupt generates a number in the ADC12IV register (see Section 23.3.7). This number can be evaluated or added to the program counter to automatically enter the appropriate software routine. Disabled ADC12 interrupts do not affect the ADC12IV value.

Any access (read or write) of the ADC12IV register automatically resets the ADC12OV condition or the ADC12TOV condition if either was the highest pending interrupt. Neither interrupt condition has an accessible interrupt flag. The ADC12IFGx flags are not reset by an ADC12IV access. ADC12IFGx bits are reset automatically by accessing their associated ADC12MEMx register or may be reset with software.

If another interrupt is pending after servicing of an interrupt, another interrupt is generated. For example, if the ADC12OV and ADC12IFG3 interrupts are pending when the interrupt service routine accesses the ADC12IV register, the ADC12OV interrupt condition is reset automatically. After the RETI instruction of the interrupt service routine is executed, the ADC12IFG3 generates another interrupt.



www.ti.com ADC12 Operation

#### 23.2.9.2 ADC12 Interrupt Handling Software Example

Example 23-1 shows the recommended use of ADC12IV and the handling overhead. The ADC12IV value is added to the PC to automatically jump to the appropriate routine.

The numbers at the right margin show the necessary CPU cycles for each instruction. The software overhead for different interrupt sources includes interrupt latency and return-from-interrupt cycles, but not the task handling itself. The latencies are:

- ADC12IFG0 to ADC12IFG14, ADC12TOV, and ADC12OV: 16 cycles
- ADC12IFG15: 14 cycles

The interrupt handler for ADC12IFG15 shows a way to check immediately if a higher prioritized interrupt occurred during the processing of ADC12IFG15. This saves nine cycles if another ADC12 interrupt is pending.

#### Example 23-1. Interrupt Handling

```
; Interrupt handler for ADC12.
INT_ADC12
                      ; Enter Interrupt Service Routine
         &ADC12IV,PC ; Add offset to PC
  ADD
  RETI
                      ; Vector 0: No interrupt
                     ; Vector 2: ADC overflow
  JMP
  JMP
                     ; Vector 4: ADC timing overflow
                                                            2
         ADM0
                     ; Vector 6: ADC12IFG0
                      ; Vectors 8-32
                                                            2
  JMP
       ADM14
                     ; Vector 34: ADC12IFG14
                                                            2
; Handler for ADC12IFG15 starts here. No JMP required.
            &ADC12MEM15,xxx ; Move result, flag is reset
ADM15
       MOV
                              ; Other instruction needed?
        . . .
        JMP
            INT ADC12
                              ; Check other int pending
   ADC12IFG14-ADC12IFG1 handlers go here
       VOM
             &ADC12MEM0,xxx ; Move result, flag is reset
DMOA
                              ; Other instruction needed?
        . . .
        RETI
                               ; Return
ADTOV
                               ; Handle Conv. time overflow
        RETI
                               ; Return
ADOV
                               ; Handle ADCMEMx overflow
        RETI
                               ; Return
                                                            5
```



ADC12 Registers www.ti.com

# 23.3 ADC12 Registers

The ADC12 registers are listed in Table 23-2.

## Table 23-2. ADC12 Registers

Register	Short Form	Register Type	Address	Initial State
ADC12 control register 0	ADC12CTL0	Read/write	01A0h	Reset with POR
ADC12 control register 1	ADC12CTL1	Read/write	01A2h	Reset with POR
ADC12 interrupt flag register	ADC12IFG	Read/write	01A4h	Reset with POR
ADC12 interrupt enable register	ADC12IE	Read/write	01A6h	Reset with POR
ADC12 interrupt vector word	ADC12IV	Read	01A8h	Reset with POR
ADC12 memory 0	ADC12MEM0	Read/write	0140h	Unchanged
ADC12 memory 1	ADC12MEM1	Read/write	0142h	Unchanged
ADC12 memory 2	ADC12MEM2	Read/write	0144h	Unchanged
ADC12 memory 3	ADC12MEM3	Read/write	0146h	Unchanged
ADC12 memory 4	ADC12MEM4	Read/write	0148h	Unchanged
ADC12 memory 5	ADC12MEM5	Read/write	014Ah	Unchanged
ADC12 memory 6	ADC12MEM6	Read/write	014Ch	Unchanged
ADC12 memory 7	ADC12MEM7	Read/write	014Eh	Unchanged
ADC12 memory 8	ADC12MEM8	Read/write	0150h	Unchanged
ADC12 memory 9	ADC12MEM9	Read/write	0152h	Unchanged
ADC12 memory 10	ADC12MEM10	Read/write	0154h	Unchanged
ADC12 memory 11	ADC12MEM11	Read/write	0156h	Unchanged
ADC12 memory 12	ADC12MEM12	Read/write	0158h	Unchanged
ADC12 memory 13	ADC12MEM13	Read/write	015Ah	Unchanged
ADC12 memory 14	ADC12MEM14	Read/write	015Ch	Unchanged
ADC12 memory 15	ADC12MEM15	Read/write	015Eh	Unchanged
ADC12 memory control 0	ADC12MCTL0	Read/write	080h	Reset with POR
ADC12 memory control 1	ADC12MCTL1	Read/write	081h	Reset with POR
ADC12 memory control 2	ADC12MCTL2	Read/write	082h	Reset with POR
ADC12 memory control 3	ADC12MCTL3	Read/write	083h	Reset with POR
ADC12 memory control 4	ADC12MCTL4	Read/write	084h	Reset with POR
ADC12 memory control 5	ADC12MCTL5	Read/write	085h	Reset with POR
ADC12 memory control 6	ADC12MCTL6	Read/write	086h	Reset with POR
ADC12 memory control 7	ADC12MCTL7	Read/write	087h	Reset with POR
ADC12 memory control 8	ADC12MCTL8	Read/write	088h	Reset with POR
ADC12 memory control 9	ADC12MCTL9	Read/write	089h	Reset with POR
ADC12 memory control 10	ADC12MCTL10	Read/write	08Ah	Reset with POR
ADC12 memory control 11	ADC12MCTL11	Read/write	08Bh	Reset with POR
ADC12 memory control 12	ADC12MCTL12	Read/write	08Ch	Reset with POR
ADC12 memory control 13	ADC12MCTL13	Read/write	08Dh	Reset with POR
ADC12 memory control 14	ADC12MCTL14	Read/write	08Eh	Reset with POR
ADC12 memory control 15	ADC12MCTL15	Read/write	08Fh	Reset with POR



SHT0x

www.ti.com ADC12 Registers

## 23.3.1 ADC12CTL0, ADC12 Control Register 0

15	14	13	12	11	10	9	8		
	SH	T1x			SHT	0x			
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)		
7	6	5	4	3	2	1	0		
MSC	REF2_5V	REFON	ADC120N	ADC120VIE	ADC12TOVIE	ENC	ADC12SC		
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)		
	Can be modified only when ENC = 0								

Can be modified only when ENC = 0

0000

SHT1x Bits 15-12 Sample-and-hold time. These bits define the number of ADC12CLK cycles in the sampling period for registers ADC12MEM8 to ADC12MEM15.

> 0000 4 ADC12CLK cycles 0001 8 ADC12CLK cycles 0010 16 ADC12CLK cycles 0011 32 ADC12CLK cycles 0100 64 ADC12CLK cycles 0101 96 ADC12CLK cycles 0110 128 ADC12CLK cycles 0111 192 ADC12CLK cycles 1000 256 ADC12CLK cycles 1001 384 ADC12CLK cycles 1010 512 ADC12CLK cycles 1011 768 ADC12CLK cycles 1100 1024 ADC12CLK cycles 1101 1024 ADC12CLK cycles 1110 1024 ADC12CLK cycles 1111 1024 ADC12CLK cycles

Bits 11-8 Sample-and-hold time. These bits define the number of ADC12CLK cycles in the sampling period for registers ADC12MEM0 to ADC12MEM7.

> 4 ADC12CLK cycles 0001 8 ADC12CLK cycles 0010 16 ADC12CLK cycles 0011 32 ADC12CLK cycles 0100 64 ADC12CLK cycles 0101 96 ADC12CLK cycles 0110 128 ADC12CLK cycles 0111 192 ADC12CLK cycles 1000 256 ADC12CLK cycles 1001 384 ADC12CLK cycles 1010 512 ADC12CLK cycles 1011 768 ADC12CLK cycles 1100 1024 ADC12CLK cycles 1101 1024 ADC12CLK cycles 1110 1024 ADC12CLK cycles 1111 1024 ADC12CLK cycles

**MSC** Bit 7 Multiple sample and conversion. Valid only for sequence or repeated modes.

> 0 The sampling timer requires a rising edge of the SHI signal to trigger each sample-and-conversion.

The first rising edge of the SHI signal triggers the sampling timer, but further sample-and-1 conversions are performed automatically as soon as the prior conversion is completed.

REF2 5V Bit 6 Reference generator voltage. REFON must also be set.

> 1.5 V 0 1 2.5 V



ADC12 Registers www.ti.com

REFON	Bit 5	Reference generator on	
		0 Reference off	
		1 Reference on	
ADC12ON	Bit 4	ADC12 on	
		0 ADC12 off	
		1 ADC12 on	
ADC120VIE	Bit 3	ADC12MEMx overflow-interrupt enable. The GIE bit must also be set to enable the interrupt.	
		0 Overflow interrupt disabled	
		1 Overflow interrupt enabled	
ADC12TOVIE	Bit 2	ADC12 conversion-time-overflow interrupt enable. The GIE bit must also be set to enable the interrupt.	
		O Conversion time overflow interrupt disabled	
		1 Conversion time overflow interrupt enabled	
ENC	Bit 1	Enable conversion	
		0 ADC12 disabled	
		1 ADC12 enabled	
ADC12SC	Bit 0	Start conversion. Software-controlled sample-and-conversion start. ADC12SC and ENC may be set together with one instruction. ADC12SC is reset automatically.	
		0 No sample-and-conversion-start	
		1 Start sample-and-conversion	



www.ti.com ADC12 Registers

# 23.3.2 ADC12CTL1, ADC12 Control Register 1

15	14		13	12	11	10	9	8
	CS	STARTA	ADDx		SI	HSx	SHP	ISSH
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
7	6		5	4	3	2	1	0
	ADC12DI	Vx		ADC12	2SSELx	CON	SEQx	ADC12BUSY
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
	Can be mod	dified or	nly when ENC =	<b>=</b> 0				
CSTARTADDx	Bits 15-12	conve		first conversion in		12 conversion-men e value of CSTART		
SHSx	Bits 11-10	Samp	le-and-hold sou	ırce select				
		00	ADC12SC b	it				
		01	Timer_A.OU	T1				
		10	Timer_B.OU	Т0				
		11	Timer_B.OU	T1				
SHP	Bit 9					e source of the san out signal directly.	npling signal (SA	MPCON) to be
		0	SAMPCON	signal is sourced	from the sample-i	nput signal.		
		1	SAMPCON :	signal is sourced	from the sampling	g timer.		
ISSH	Bit 8	Invert	signal sample-	and-hold				
		0	The sample-	input signal is not	t inverted.			
		1	The sample-	input signal is inv	erted.			
ADC12DIVx	Bits 7-5	ADC1	2 clock divider					
		000	/1					
		001	/2					
		010	/3					
		011	/4					
		100	/5					
		101	/6					
		110	/7					
		111	/8					
ADC12SSELx	Bits 4-3	ADC1	2 clock source	select				
		00	ADC12OSC					
		01	ACLK					
		10	MCLK					
		11	SMCLK					
CONSEQx	Bits 2-1	Conve	ersion sequence	e mode select				
		00		nel, single-conver	sion			
		01	Sequence-of					
		10	Repeat-sing					
		11		uence-of-channels				
ADC12BUSY	Bit 0		•		ve sample or con	version operation.		
		0	No operation	n is active.				

A sequence, sample, or conversion is active.



ADC12 Registers www.ti.com

# 23.3.3 ADC12MEMx, ADC12 Conversion Memory Registers

15	14	13	12	11	10	9	8	
0	0	0	0	Conversion Results				
r0	r0	r0	r0	rw	rw	rw	rw	
7	6	5	4	3	2	1	0	
			Conversi	on Results				
rw	rw	rw	rw	rw	rw	rw	rw	
Conversion	Rits 15-0	The 12-hit conversion	n results are right.	instified Bit 11 is	the MSR Rits 15-	12 are always 0 \	Writing to the	

Conversion Results Bits 15-0

The 12-bit conversion results are right-justified. Bit 11 is the MSB. Bits 15-12 are always 0. Writing to the conversion memory registers corrupts the results.

# 23.3.4 ADC12MCTLx, ADC12 Conversion Memory Control Registers

7	6		5	4	3	2	1	0
EOS			SREFx			INC	Hx	
rw-(0)	rw-(0)		rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
	Can be mod	dified only	y when ENC	= 0				
EOS	Bit 7	End of	sequence. I	ndicates the last con	version in a sec	quence.		
		0	Not end of	sequence				
		1	End of seq	uence				
SREFx	Bits 6-4	Select	reference					
		000		and $V_{R-} = AV_{SS}$				
		001		and $V_{R} = AV_{SS}$				
		010		$_{+}$ and $V_{R-} = AV_{SS}$				
		011		$_{+}$ and $V_{R-} = AV_{SS}$				
		100		and $V_{R-} = V_{REF} / V_{eRE}$				
		101		and $V_{R-} = V_{REF} / V_{eRE}$				
		110	$V_{R+} = V_{eREF}$	$_{+}$ and $V_{R-} = V_{REF} / V_{eR}$	EF-			
		111		$_{+}$ and $V_{R-} = V_{REF} / V_{eR}$	EF-			
INCHx	Bits 3-0		hannel selec	t				
		0000	A0					
		0001	A1					
		0010	A2					
		0011	A3					
		0100	A4					
		0101	A5					
		0110	A6					
		0111	A7					
		1000	$V_{eREF+}$					
		1001	$V_{REF-}/V_{eREF}$	-				
		1010	Temperatu	re diode				
		1011	(AV <sub>CC</sub> - AV	( <sub>SS</sub> ) / 2				
		1100	GND					
		1101	GND					
		1110	GND					
		1111	GND					



www.ti.com ADC12 Registers

# 23.3.5 ADC12IE, ADC12 Interrupt Enable Register

15	14	13	12	11	10	9	8
ADC12IE15	ADC12IE14	ADC12IE13	ADC12IE12	ADC12IE11	ADC12IE10	ADC12IFG9	ADC12IE8
rw-(0)	rw-(0)						
7	6	5	4	3	2	1	0
ADC12IE7	ADC12IE6	ADC12IE5	ADC12IE4	ADC12IE3	ADC12IE2	ADC12IE1	ADC12IE0
rw-(0)	rw-(0)						

ADC12IEx Bits 15-0 Interrupt enable. These bits enable or disable the interrupt request for the ADC12IFGx bits.

0 Interrupt disabled1 Interrupt enabled

# 23.3.6 ADC12IFG, ADC12 Interrupt Flag Register

15	14	13	12	11	10	9	8
ADC12IFG15	ADC12IFG14	ADC12IFG13	ADC12IFG12	ADC12IFG11	ADC12IFG10	ADC12IFG9	ADC12IFG8
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
7	6	5	4	3	2	1	0
ADC12IFG7	ADC12IFG6	ADC12IFG5	ADC12IFG4	ADC12IFG3	ADC12IFG2	ADC12IFG1	ADC12IFG0
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)

ADC12IFGx

Bits 15-0

ADC12MEMx Interrupt flag. These bits are set when corresponding ADC12MEMx is loaded with a conversion result. The ADC12IFGx bits are reset if the corresponding ADC12MEMx is accessed, or may be reset with software.

No interrupt pendingInterrupt pending



ADC12 Registers www.ti.com

# 23.3.7 ADC12IV, ADC12 Interrupt Vector Register

15	14	13	12	11	10	9	8
0	0	0	0	0	0	0	0
rO	r0	r0	r0	r0	r0	r0	rO
7	6	5	4	3	2	1	0
0	0		ADC12IVx				
r0	r0	r-(0)	r-(0)	r-(0)	r-(0)	r-(0)	r0

ADC12IVx Bits 15-0 ADC12 interrupt vector value

ADC12IV Contents	Interrupt Source	Interrupt Flag	Interrupt Priority
000h	No interrupt pending	=	
002h	ADC12MEMx overflow	=	Highest
004h	Conversion time overflow	=	
006h	ADC12MEM0 interrupt flag	ADC12IFG0	
008h	ADC12MEM1 interrupt flag	ADC12IFG1	
00Ah	ADC12MEM2 interrupt flag	ADC12IFG2	
00Ch	ADC12MEM3 interrupt flag	ADC12IFG3	
00Eh	ADC12MEM4 interrupt flag	ADC12IFG4	
010h	ADC12MEM5 interrupt flag	ADC12IFG5	
012h	ADC12MEM6 interrupt flag	ADC12IFG6	
014h	ADC12MEM7 interrupt flag	ADC12IFG7	
016h	ADC12MEM8 interrupt flag	ADC12IFG8	
018h	ADC12MEM9 interrupt flag	ADC12IFG9	
01Ah	ADC12MEM10 interrupt flag	ADC12IFG10	
01Ch	ADC12MEM11 interrupt flag	ADC12IFG11	
01Eh	ADC12MEM12 interrupt flag	ADC12IFG12	
020h	ADC12MEM13 interrupt flag	ADC12IFG13	
022h	ADC12MEM14 interrupt flag	ADC12IFG14	
024h	ADC12MEM15 interrupt flag	ADC12IFG15	Lowest



# **TLV Structure**

The Tag-Length-Value (TLV) structure is used in selected MSP430x2xx devices to provide device-specific information in the device's flash memory SegmentA, such as calibration data. For the device-dependent implementation, see the device-specific data sheet.

Topic Page

24.1	TLV Introduction	<b>582</b>
24.2	Supported Tags	583
24.3	Checking Integrity of SegmentA	586
24.4	Parsing TLV Structure of Segment A	586



TLV Introduction www.ti.com

### 24.1 TLV Introduction

The TLV structure stores device-specific data in SegmentA. The SegmentA content of an example device is shown in Table 24-1.

Table 24-1. Example SegmentA Structure

Word Address	Upper Byte	Lower Byte	Tag Address and Offset
0x10FE	CALBC1_1MHZ	CALDCO_1MHZ	0x10F6 + 0x0008
0x10FC	CALBC1_8MHZ	CALDCO_8MHZ	0x10F6 + 0x0006
0x10FA	CALBC1_12MHZ	CALDCO_12MHZ	0x10F6 + 0x0004
0x10F8	CALBC1_16MHZ	CALDCO_16MHZ	0x10F6 + 0x0002
0x10F6	0x08 (LENGTH)	TAG_DCO_30	0x10F6
0x10F4	0xFF	0xFF	
0x10F2	0xFF	0xFF	
0x10F0	0xFF	0xFF	
0x10EE	0xFF	0xFF	
0x10EC	0x08 (LENGTH)	TAG_EMPTY	0x10EC
0x10EA	CAL_ADO	C_25T85	0x10DA + 0x0010
0x10E8	CAL_ADO	C_25T30	0x10DA + 0x000E
0x10E6	CAL_ADC_25VI	REF_FACTOR	0x10DA + 0x000C
0x10E4	CAL_ADO	C_15T85	0x10DA + 0x000A
0x10E2	CAL_ADO	C_15T30	0x10DA + 0x0008
0x10E0	CAL_ADC_15VI	REF_FACTOR	0x10DA + 0x0006
0x10DE	CAL_ADC_	_OFFSET	0x10DA + 0x0004
0x10DC	CAL_ADC_GA	AIN_FACTOR	0x10DA + 0x0002
0x10DA	0x10 (LENGTH)	TAG_ADC12_1	0x10DA
0x10D8	0xFF	0xFF	
0x10D6	0xFF	0xFF	
0x10D4	0xFF	0xFF	
0x10D2	0xFF	0xFF	
0x10D0	0xFF	0xFF	
0x10CE	0xFF	0xFF	
0x10CC	0xFF	0xFF	
0x10CA	0xFF	0xFF	
0x10C8	0xFF	0xFF	
0x10C6	0xFF	0xFF	
0x10C4	0xFF	0xFF	
0x10C2	0x16 (LENGTH)	TAG_EMPTY	0x10C2
0x10C0	2s complement of	of bit-wise XOR	0x10C0

The first two bytes of SegmentA (0x10C0 and 0x10C1) hold the checksum of the remainder of the segment (addresses 0x10C2 to 0x10FF).

The first tag is located at address 0x10C2 and, in this example, is the TAG\_EMPTY tag. The following byte (0x10C3) holds the length of the following structure. The length of this TAG\_EMPTY structure is 0x16 and, therefore, the next tag, TAG\_ADC12\_1, is found at address 0x10DA. Again, the following byte holds the length of the TAG\_ADC12\_1 structure.

The TLV structure maps the entire address range 0x10C2 to 0x10FF of the SegmentA. A program routine looking for tags starting at the SegmentA address 0x10C2 can extract all information even if it is stored at a different (device-specific) absolute address.



www.ti.com Supported Tags

# 24.2 Supported Tags

Each device contains a subset of the tags shown in Table 24-2. See the device-specific data sheet for details.

Table 24-2. Supported Tags (Device Specific)

Tag	Description	Value
TAG_EMPTY	Identifies an unused memory area	0xFE
TAG_DCO_30	Calibration values for the DCO at room temperature and DV <sub>CC</sub> = 3 V	0x01
TAG_ADC12_1	Calibration values for the ADC12 module	0x08
TAG_ADC10_1	Calibration values for the ADC10 module	0x08

# 24.2.1 DCO Calibration TLV Structure

For DCO calibration, the BCS+ registers (BCSCTL1 and DCOCTL) are used. The values stored in the flash information memory SegmentA are written to the BCS+ registers (see Table 24-3).

Table 24-3. DCO Calibration Data (Device Specific)

Label	Description	Offset
CALBC1_1MHZ	Value for the BCSCTL1 register for 1 MHz, T <sub>A</sub> = 25°C	0x07
CALDCO_1MHZ	Value for the DCOCTL register for 1 MHz, T <sub>A</sub> = 25°C	0x06
CALBC1_8MHZ	Value for the BCSCTL1 register for 8 MHz, T <sub>A</sub> = 25°C	0x05
CALDCO_8MHZ	Value for the DCOCTL register for 8 MHz, T <sub>A</sub> = 25°C	0x04
CALBC1_12MHZ	Value for the BCSCTL1 register for 12 MHz, T <sub>A</sub> = 25°C	0x03
CALDCO_12MHZ	Value for the DCOCTL register for 12 MHz, T <sub>A</sub> = 25°C	0x02
CALBC1_16MHZ	Value for the BCSCTL1 register for 16 MHz, T <sub>A</sub> = 25°C	0x01
CALDCO_16MHZ	Value for the DCOCTL register for 16 MHz, T <sub>A</sub> = 25°C	0x00

The calibration data for the DCO is available in all 2xx devices and is stored at the same absolute addresses. The device-specific SegmentA content is applied using the absolute addressing mode if the sample code shown in Example 24-1 is used.

### Example 24-1. Code Example Using Absolute Addressing Mode

```
; Calibrate the DCO to 1 MHz

CLR.B &DCOCTL ; Select lowest DCOx
; and MODx settings

MOV.B &CALBC1_1MHZ,&BCSCTL1 ; Set RSELx

MOV.B &CALDCO_1MHZ,&DCOCTL ; Set DCOx and MODx
```

The TLV structure allows use of the address of the TAG\_DCO\_30 tag to address the DCO registers. Example 24-2 shows how to address the DCO calibration data using the TAG\_DCO\_30 tag.

# Example 24-2. Code Example Using the TLV Structure



Supported Tags www.ti.com

### 24.2.2 TAG ADC12 1 Calibration TLV Structure

The calibration data for the ADC12 module consists of eight words (see Table 24-4).

Table 24-4. TAG\_ADC12\_1 Calibration Data (Device Specific)

Label	Description	Offset
CAL_ADC_25T85	VREF2_5 = 1, T <sub>A</sub> = 85°C ± 2°C, 12-bit conversion result	0x0E
CAL_ADC_25T30	VREF2_5 = 1, T <sub>A</sub> = 30°C ± 2°C, 12-bit conversion result	0x0C
CAL_ADC_25VREF_FACTOR	VREF2_5 = 1, T <sub>A</sub> = 30°C ± 2°C	0x0A
CAL_ADC_15T85	VREF2_5 = 0, T <sub>A</sub> = 85°C ± 2°C, 12-bit conversion result	0x08
CAL_ADC_15T30	VREF2_5 = 0, T <sub>A</sub> = 30°C ± 2°C, 12-bit conversion result	0x06
CAL_ADC_15VREF_FACTOR	VREF2_5 = 0, T <sub>A</sub> = 30°C ± 2°C	0x04
CAL_ADC_OFFSET	VeREF = 2.5 V, T <sub>A</sub> = 85°C ± 2°C, f <sub>ADC12CLK</sub> = 5 MHz	0x02
CAL_ADC_GAIN_FACTOR	VeREF = 2.5 V, T <sub>A</sub> = 85°C ± 2°C, f <sub>ADC12CLK</sub> = 5 MHz	0x00

### 24.2.2.1 Temperature Sensor Calibration Data

The temperature sensor is calibrated using the internal voltage references. Each reference voltage (1.5 V and 2.5 V) contains a measured value for two temperatures, 30°C±2°C and 85°C±2°C and are stored in the TLV structure at the respective SegmentA location (see Table 24-4). The characteristic equation of the temperature sensor voltage, in mV, is:

$$V_{SENSE} = TC_{SENSOR} \times Temp + V_{SENSOR}$$
 (1)

The temperature coefficient,  $TC_{SENSOR}$  in mV/°C, represents the slope of the equation.  $V_{SENSOR}$ , in mV, represents the y-intercept of the equation. Temp, in °C, is the temperature of interest.

The temperature (Temp, °C) can be computed as follows for each of the reference voltages used in the ADC measurement:

$$Temp = (ADC(raw) - CAL\_ADC\_15T30) \times \left(\frac{85 - 30}{CAL\_ADC\_15T85 - CAL\_ADC\_15T30}\right) + 30$$

$$Temp = (ADC(raw) - CAL\_ADC\_25T30) \times \left(\frac{85 - 30}{CAL\_ADC\_25T85 - CAL\_ADC\_25T30}\right) + 30$$
(2)

# 24.2.2.2 Integrated Voltage Reference Calibration Data

The reference voltages (VREF2\_5 = 0 and 1) are measured at room temperature. The measured value is normalized by 1.5 V or 2.5 V before stored into the flash information memory SegmentA.

CAL ADC 15VREF FACTOR = 
$$(Ve_{REF} / 1.5 V) \times 2^{15}$$

The conversion result is corrected by multiplying it with the CAL\_ADC\_15VREF\_FACTOR (or CAL\_ADC\_25VREF\_FACTOR) and dividing the result by 2<sup>15</sup>.

$$ADC(corrected) = ADC(raw) \times CAL ADC 15VREF FACTOR \times (1 / 2^{15})$$

### 24.2.2.3 Example Using the Reference Calibration

In the following example, the integrated 1.5-V reference voltage is used during a conversion.

- Conversion result: 0x0100
- Reference voltage calibration factor (CAL\_ADC\_15VREF\_FACTOR): 0x7BBB

The following steps show an example of how the ADC12 conversion result can be corrected by using the hardware multiplier:

- 1. Multiply the conversion result by 2 (this step simplifies the final division).
- 2. Multiply the result by CAL\_ADC\_15VREF\_FACTOR.
- 3. Divide the result by 2<sup>16</sup> (use the upper word of the 32-bit multiplication result RESHI).



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In the example:

- 1.  $0x0100 \times 0x0002 = 0x0200$
- 2.  $0x0200 \times 0x7BBB = 0x00F7 7600$
- 3.  $0x00F7_7600 \div 0x0001_0000 = 0x0000_00F7 (= 247)$

The code example using the hardware multiplier follows.

```
; The ADC conversion result is stored in ADC12MEM0
; It is assumed that R9 contains the address of the
; TAG_ADC12_1.
; The corrected value is available in ADC_COR
     MOV.W &ADC12MEM0,R10; move result to R10
                            ; R10 x 2
     RLA.W
            R10
     MOV.W R10,&MPY
                              ; unsigned multiply OP1
     MOV.W CAL_ADC_15VREF_FACTOR(R9),&OP2
                             ; calibration value OP2
     MOV.W
            &RESHI, &ADC_COR ; result: upper 16-bit MPY
```

### 24.2.2.4 Offset and Gain Calibration Data

The offset of the ADC12 is determined and stored as a twos-complement number in SegmentA. The offset error correction is done by adding the CAL\_ADC\_OFFSET to the conversion result.

```
ADC(offset_corrected) = ADC(raw) + CAL_ADC_OFFSET
```

The gain of the ADC12, stored at offset 0x00, is calculated by the following equation.

```
CAL ADC GAIN FACTOR = (1 / GAIN) \times 2^{15}
```

The conversion result is gain corrected by multiplying it with the CAL ADC GAIN FACTOR and dividing the result by  $2^{15}$ .

```
ADC(gain corrected) = ADC(raw) \times CAL ADC GAIN FACTOR \times (1 / 2<sup>15</sup>)
```

If both gain and offset are corrected, the gain correction is done first.

```
ADC(gain\_corrected) = ADC(raw) \times CAL\_ADC\_GAIN\_FACTOR \times (1 / 2^{15})
ADC(final) = ADC(gain corrected) + CAL ADC OFFSET
```

#### 24.2.2.5 Example Using Gain and Offset Calibration

In the following example, an external reference voltage is used during a conversion.

- Conversion result: 0x0800 (= 2048)
- Gain calibration factor: 0x7FE0 (gain error: +2 LSB)
- Offset calibration: 0xFFFE (2s complement of -2)

The following steps show an example of how the ADC12 conversion result is corrected by using the hardware multiplier:

- 1. Multiply the conversion result by 2 (this step simplifies the final division).
- 2. Multiply the result by CAL\_ADC\_GAIN\_FACTOR.
- 3. Divide the result by 2<sup>16</sup> (use the upper word of the 32-bit multiplication result RESHI)
- 4. Add CAL\_ADC\_OFFSET to the result.

In the example:

- 1.  $0x0800 \times 0x0002 = 0x1000$
- 2.  $0x1000 \times 0x8010 = 0x0801 0000$
- 3.  $0x0801_0000 \div 0x0001_0000 = 0x0000_0801 (= 2049)$
- 4. 0x801 + 0xFFFE = 0x07FF (= 2047)

The code example using the hardware multiplier follows.

```
; The ADC conversion result is stored in ADC12MEM0
```

; It is assumed that R9 contains the address of the TAG\_ADC12\_1.



```
; The corrected value is available in ADC_COR
     MOV.W
             &ADC12MEM0,R10 ; move result to R10
     RLA.W
             R10
                                ; R10 * 2
     MOV.W
             R10,&MPY
                                ; unsigned multiply OP1
             CAL_ADC_GAIN_FACTOR(R9),&OP2
     MOV.W
                                ; calibration value OP2
     MOV.W
              &RESHI, &ADC_COR
                                ; use upper 16-bit MPY
     ADD.W
             CAL_ADC_OFFSET(R9), &ADC_COR
                                ; add offset correction
```

# 24.3 Checking Integrity of SegmentA

The 64-byte SegmentA contains a 2-byte checksum of the data stored at 0x10C2 up to 0x10FF at addresses 0x10C0 and 0x10C1. The checksum is a bit-wise XOR of 31 words stored in the twoscomplement data format.

A code example to calculate the checksum follows.

```
; Checking the SegmentA integrity by calculating the 2's
; complement of the 31 words at 0x10C2 - 0x10FE.
; It is assumed that the SegmentA Start Address is stored
; in R10. R11 is initialized to 0x00.
; The label TLV_CHKSUM is set to 0x10C0.
       ADD.W
               #2,R10
                                 ; Skip the checksum
LP0
       XOR.W
               @R10+,R11
                                 ; Add a word to checksum
       CMP.W
               #0x10FF,R10
                                 ; Last word included?
       JN
               LP0
                                 ; No, add more data
              &TLV CHKSUM,R11
                                 ; Add checksum
      ADD.W
       JNZ
              CSNOK
                                 ; Checksum not ok
                                 ; Use SegmentA data
CSNOK
                                 ; Do not use SegmentA Data
      . . .
```

### 24.4 Parsing TLV Structure of Segment A

Example code to analyze SegmentA follows.

```
; It is assumed that the SegmentA start address
; is stored in R10.
LP1
      ADD.W
              #2,R10
                                     ; Skip two bytes
      CMP.W
              #0x10FF,R10
                                     ; SegmentA end reached?
      JGE
              DONE
                                     ; Yes, done
      CMP.B
              #TAG_EMPTY,0(R10)
                                     ; TAG_EMPTY?
      JNZ
              т1
                                     ; No, continue
      JMP
              LP2
                                     ; Yes, done with TAG_EMPTY
т1
      CMP.B
              #TAG_ADC12_1,0(R10) ; TAG_ADC12_1?
      JNZ
                                     ; No, continue
                                     ; Yes, found TAG_ADC12_1
      . . .
      JMP
              LP2
                                     ; Done with TAG_ADC12_1
Т2
      CMP.B
              #DCO_30,0(R10)
                                     ; TAG_DCO_30?
      JNZ
                                     ; No, continue
              &DCOCTL
                                     ; Select lowest DCOx
      CLR.B
              7(R10),&BCSCTL1
      MOV.B
                                     ; Yes, use e.g. 8MHz data and
      MOV.B
              6(R10),&DCOCTL
                                     ; set DCOx and MODx
      JMP
                                     ; Done with TAG_DCO_30
              LP2
Т3
                                     ; Test for "next tag"
      . . .
      . . .
      JMP
              LP2
                                     ; Done with "next tag"
```





LP2 MOV.B 1(R10),R11 ; Store LENGTH in R11 ADD.W R11,R10 ; Add LENGTH to R10

JMP LP1 ; Jump to continue analysis

DONE

# Chapter 25





# DAC12

The DAC12 module is a 12-bit voltage-output digital-to-analog converter (DAC). This chapter describes the operation of the DAC12 module of the MSP430x2xx device family.

**Topic** Page

25.1	DAC12 Introduction	589
25.2	DAC12 Operation	591
25.3	DAC12 Registers	595



www.ti.com DAC12 Introduction

### 25.1 DAC12 Introduction

The DAC12 module is a 12-bit voltage-output DAC. The DAC12 can be configured in 8-bit or 12-bit mode and may be used in conjunction with the DMA controller. When multiple DAC12 modules are present, they may be grouped together for synchronous update operation.

Features of the DAC12 include:

- 12-bit monotonic output
- 8-bit or 12-bit voltage output resolution
- Programmable settling time vs power consumption
- Internal or external reference selection
- Straight binary or 2s compliment data format
- Self-calibration option for offset correction
- Synchronized update capability for multiple DAC12 modules

### **NOTE: Multiple DAC12 Modules**

Some devices may integrate more than one DAC12 module. If more than one DAC12 is present on a device, the multiple DAC12 modules operate identically.

Throughout this chapter, nomenclature appears such as DAC12\_xDAT or DAC12\_xCTL to describe register names. When this occurs, the x is used to indicate which DAC12 module is being discussed. In cases where operation is identical, the register is simply referred to as DAC12\_xCTL.

The block diagram of the DAC12 module is shown in Figure 25-1.



DAC12 Introduction www.ti.com

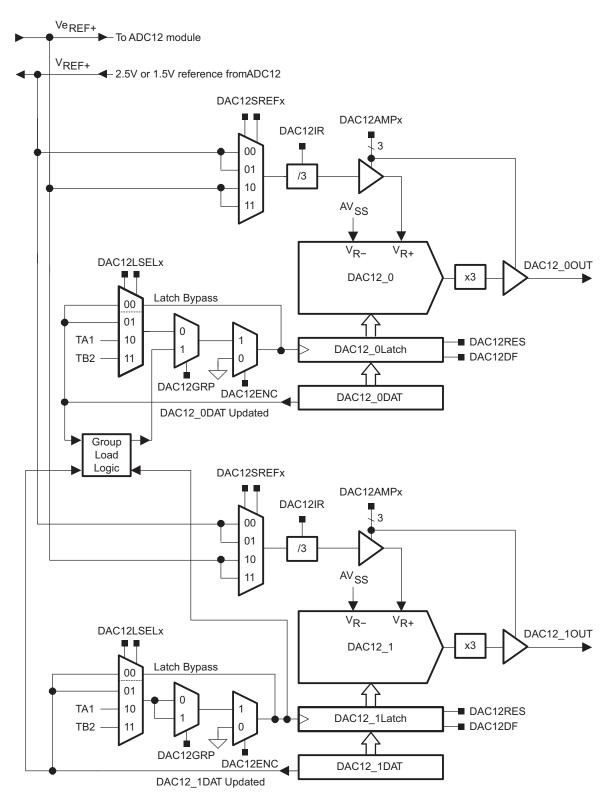


Figure 25-1. DAC12 Block Diagram



www.ti.com DAC12 Operation

### 25.2 DAC12 Operation

The DAC12 module is configured with user software. The setup and operation of the DAC12 is discussed in the following sections.

#### 25.2.1 DAC12 Core

The DAC12 can be configured to operate in 8-bit or 12-bit mode using the DAC12RES bit. The full-scale output is programmable to be 1x or 3x the selected reference voltage via the DAC12IR bit. This feature allows the user to control the dynamic range of the DAC12. The DAC12DF bit allows the user to select between straight binary data and 2s-compliment data for the DAC. When using straight binary data format, the formula for the output voltage is given in Table 25-1.

Resolution	DAC12RES	DAC12IR	Output Voltage Formula
12 bit	0	0	$V_{OUT} = V_{REF} \times 3 \times \frac{DAC12\_xDAT}{4096}$
12 bit	0	1	$V_{OUT} = V_{REF} \times \frac{DAC12\_xDAT}{4096}$
8 bit	1	0	$V_{OUT} = V_{REF} \times 3 \times \frac{DAC12\_xDAT}{256}$
8 bit	1	1	$V_{OUT} = V_{REF} \times \frac{DAC12\_xDAT}{256}$

Table 25-1. DAC12 Full-Scale Range ( $V_{REF} = V_{eREF+}$  or  $V_{REF+}$ )

In 8-bit mode, the maximum useable value for DAC12\_xDAT is 0FFh. In 12-bit mode, the maximum useable value for DAC12\_xDAT is 0FFFh. Values greater than these may be written to the register, but all leading bits are ignored.

#### 25.2.1.1 DAC12 Port Selection

The DAC12 outputs are multiplexed with the port P6 pins and ADC12 analog inputs, and also the  $V_{\text{eREF+}}$  pins. When DAC12AMPx > 0, the DAC12 function is automatically selected for the pin, regardless of the state of the associated PxSELx and PxDIRx bits. The DAC12OPS bit selects between the P6 pins and the  $V_{\text{eREF+}}$  pins for the DAC outputs. For example, when DAC12OPS = 0, DAC12\_0 outputs on P6.6 and DAC12\_1 outputs on P6.7. When DAC12OPS = 1, DAC12\_0 outputs on  $V_{\text{eREF+}}$  and DAC12\_1 outputs on P6.5. See the port pin schematic in the device-specific data sheet for more details.

### 25.2.2 DAC12 Reference

The reference for the DAC12 is configured to use either an external reference voltage or the internal 1.5-V/2.5-V reference from the ADC12 module with the DAC12SREFx bits. When DAC12SREFx =  $\{0,1\}$  the  $V_{REF+}$  signal is used as the reference and when DAC12SREFx =  $\{2,3\}$  the  $V_{eREF+}$  signal is used as the reference.

To use the ADC12 internal reference, it must be enabled and configured via the applicable ADC12 control bits.

# 25.2.2.1 DAC12 Reference Input and Voltage Output Buffers

The reference input and voltage output buffers of the DAC12 can be configured for optimized settling time vs power consumption. Eight combinations are selected using the DAC12AMPx bits. In the low/low setting, the settling time is the slowest, and the current consumption of both buffers is the lowest. The medium and high settings have faster settling times, but the current consumption increases. See the device-specific data sheet for parameters.

# 25.2.3 Updating the DAC12 Voltage Output

The DAC12\_xDAT register can be connected directly to the DAC12 core or double buffered. The trigger for updating the DAC12 voltage output is selected with the DAC12LSELx bits.



DAC12 Operation www.ti.com

When DAC12LSELx = 0 the data latch is transparent and the DAC12\_xDAT register is applied directly to the DAC12 core. the DAC12 output updates immediately when new DAC12 data is written to the DAC12\_xDAT register, regardless of the state of the DAC12ENC bit.

When DAC12LSELx = 1, DAC12 data is latched and applied to the DAC12 core after new data is written to DAC12\_xDAT. When DAC12LSELx = 2 or 3, data is latched on the rising edge from the Timer\_A CCR1 output or Timer\_B CCR2 output respectively. DAC12ENC must be set to latch the new data when DAC12LSELx > 0.

### 25.2.4 DAC12 xDAT Data Format

The DAC12 supports both straight binary and 2s compliment data formats. When using straight binary data format, the full-scale output value is 0FFFh in 12-bit mode (0FFh in 8-bit mode) as shown in Figure 25-2.

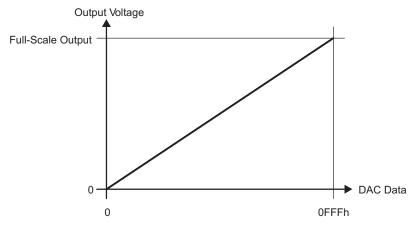


Figure 25-2. Output Voltage vs DAC12 Data, 12-Bit, Straight Binary Mode

When using 2s-compliment data format, the range is shifted such that a DAC12\_xDAT value of 0800h (0080h in 8-bit mode) results in a zero output voltage, 0000h is the mid-scale output voltage, and 07FFh (007Fh for 8-bit mode) is the full-scale voltage output (see Figure 25-3).

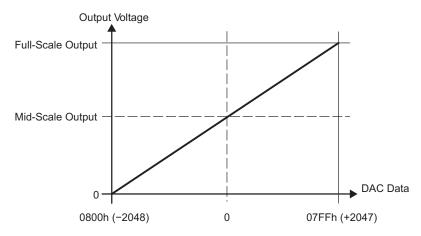


Figure 25-3. Output Voltage vs DAC12 Data, 12-Bit, 2s-Compliment Mode

### 25.2.5 DAC12 Output Amplifier Offset Calibration

The offset voltage of the DAC12 output amplifier can be positive or negative. When the offset is negative, the output amplifier attempts to drive the voltage negative but cannot do so. The output voltage remains at zero until the DAC12 digital input produces a sufficient positive output voltage to overcome the negative offset voltage, resulting in the transfer function shown in Figure 25-4.



www.ti.com DAC12 Operation

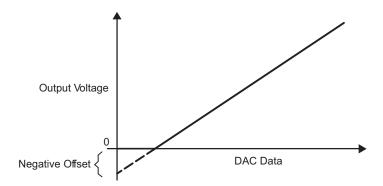


Figure 25-4. Negative Offset

When the output amplifier has a positive offset, a digital input of zero does not result in a zero output voltage. The DAC12 output voltage reaches the maximum output level before the DAC12 data reaches the maximum code. This is shown in Figure 25-5.

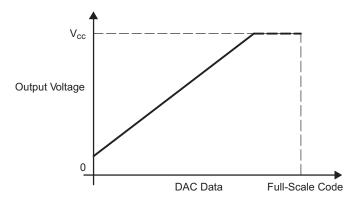


Figure 25-5. Positive Offset

The DAC12 has the capability to calibrate the offset voltage of the output amplifier. Setting the DAC12CALON bit initiates the offset calibration. The calibration should complete before using the DAC12. When the calibration is complete, the DAC12CALON bit is automatically reset. The DAC12AMPx bits should be configured before calibration. For best calibration results, port and CPU activity should be minimized during calibration.

# 25.2.6 Grouping Multiple DAC12 Modules

Multiple DAC12s can be grouped together with the DAC12GRP bit to synchronize the update of each DAC12 output. Hardware ensures that all DAC12 modules in a group update simultaneously independent of any interrupt or NMI event.

DAC12\_0 and DAC12\_1 are grouped by setting the DAC12GRP bit of DAC12\_0. The DAC12GRP bit of DAC12 1 is don't care. When DAC12 0 and DAC12 1 are grouped:

- The DAC12\_1 DAC12LSELx bits select the update trigger for both DACs
- The DAC12LSELx bits for both DACs must be > 0
- The DAC12ENC bits of both DACs must be set to 1

When DAC12 0 and DAC12 1 are grouped, both DAC12 xDAT registers must be written to before the outputs update, even if data for one or both of the DACs is not changed. Figure 25-6 shows a latch-update timing example for grouped DAC12 0 and DAC12 1.

When DAC12 0 DAC12GRP = 1 and both DAC12 x DAC12LSELx > 0 and either DAC12ENC = 0, neither DAC12 updates.



DAC12 Operation www.ti.com

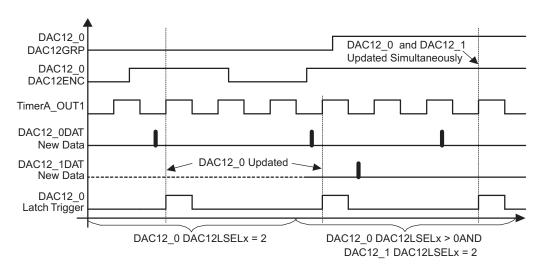


Figure 25-6. DAC12 Group Update Example, Timer\_A3 Trigger

### NOTE: DAC12 Settling Time

The DMA controller is capable of transferring data to the DAC12 faster than the DAC12 output can settle. The user must assure the DAC12 settling time is not violated when using the DMA controller. See the device-specific data sheet for parameters.

# 25.2.7 DAC12 Interrupts

The DAC12 interrupt vector is shared with the DMA controller on some devices (see device-specific data sheet for interrupt assignment). In this case, software must check the DAC12IFG and DMAIFG flags to determine the source of the interrupt.

The DAC12IFG bit is set when DAC12LSELx > 0 and DAC12 data is latched from the DAC12\_xDAT register into the data latch. When DAC12LSELx = 0, the DAC12IFG flag is not set.

A set DAC12IFG bit indicates that the DAC12 is ready for new data. If both the DAC12IE and GIE bits are set, the DAC12IFG generates an interrupt request. The DAC12IFG flag is not reset automatically. It must be reset by software.



www.ti.com DAC12 Registers

# 25.3 DAC12 Registers

The DAC12 registers are listed in Table 25-2.

# Table 25-2. DAC12 Registers

Register	Short Form	Register Type	Address	Initial State
DAC12_0 control	DAC12_0CTL	Read/write	01C0h	Reset with POR
DAC12_0 data	DAC12_0DAT	Read/write	01C8h	Reset with POR
DAC12_1 control	DAC12_1CTL	Read/write	01C2h	Reset with POR
DAC12_1 data	DAC12_1DAT	Read/write	01CAh	Reset with POR



DAC12 Registers www.ti.com

# 25.3.1 DAC12\_xCTL, DAC12 Control Register

15	14	13	12	11	10	9	8
DAC12OPS	DA	AC12SREFx	DAC12RES	DAC1	2LSELx	DAC12CALON	DAC12IR
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
7	6	5	4	3	2	1	0
	DAC12AM	lPx	DAC12DF	DAC12IE	DAC12IFG	DAC12ENC	DAC12GRP
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
	Can be mod	dified only when D	ied only when DAC12ENC = 0				
DAC12OPS	Bit 15	DAC12 output o	oloot				
DACIZOFS	DIL 13	DAC12 output s  0 DAC12	elect _0 output on P6.6, DA0	^12 1 output on	D6 7		
			_0 output on V <sub>eREF+</sub> , DA				
DAC12SREFx	Bits 14-13	DAC12 select re		NOTZ_T output of	11 0.5		
27101201121 X	Dito 11 10	00 V <sub>REF+</sub>	noronoo vollago				
		$V_{REF+}$					
		10 V <sub>eREF+</sub>					
		11 V <sub>eREF+</sub>					
DAC12RES	Bit 12	DAC12 resolution	n select				
		0 12-bit re	esolution				
		1 8-bit res	solution				
DAC12LSELx	Bits 11-10		ect. Selects the load tri when DAC12LSELx = 0		C12 latch. DAC12E	NC must be set for	the DAC to
		00 DAC12	latch loads when DAC	12_xDAT written	(DAC12ENC is ig	nored)	
			latch loads when DAC roup have been writter	_	, or, when grouped	d, when all DAC12_	xDAT registers
		10 Rising	edge of Timer_A.OUT1	(TA1)			
		11 Rising	edge of Timer_B.OUT2	? (TB2)			
DAC12CALON	Bit 9	DAC12 calibration when the calibration	on on. This bit initiates ation completes.	the DAC12 offse	et calibration seque	ence and is automa	tically reset
		0 Calibra	ion is not active				
		1 Initiate	calibration/calibration in	n progress			
DAC12IR	Bit 8	•	nge. This bit sets the re	•	• ,	ange.	
			full-scale output = 3x r	Ū			
			full-scale output = 1x r	Ū			
DAC12AMPx	Bits 7-5	output amplifier	setting. These bits se	lect settling time	vs current consum	iption for the DAC1	2 input and
		DAC12AMPx	Input Buffer	Ou	tput Buffer		
		000	Off	DA	C12 off, output high	gh Z	
		001	Off	DA	C12 off, output 0	V	
		010	Low speed/current	Lov	w speed/current		
		011	Low speed/current	Me	dium speed/curre	nt	
		100	Low speed/current	Hig	gh speed/current		
		101	Medium speed/curren	nt Me	dium speed/curre	nt	
		110	Medium speed/curren	nt Hig	gh speed/current		
		111	High speed/current	Hig	gh speed/current		
DAC12DF	Bit 4	DAC12 data for					
		0 Straight	•				
	<b>B</b> 1		plement				
DAC12IE	Bit 3	DAC12 interrupt					
		0 Disable					
		1 Enable	ג				



www.ti.com DAC12 Registers

DAC12IFG	Bit 2	DAC12 Interrupt flag
		0 No interrupt pending
		1 Interrupt pending
DAC12ENC	Bit 1	DAC12 enable conversion. This bit enables the DAC12 module when DAC12LSELx > 0. when DAC12LSELx = 0, DAC12ENC is ignored.
		0 DAC12 disabled
		1 DAC12 enabled
DAC12GRP	Bit 0	DAC12 group. Groups DAC12_x with the next higher DAC12_x. Not used for DAC12_1.
		0 Not grouped
		1 Grouped

# 25.3.2 DAC12\_xDAT, DAC12 Data Register

15	14	13	12	11	10	9	8
0	0	0	0		DAC1	2 Data	
r(0)	r(0)	r(0)	r(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)
7	6	5	4	3	2	1	0
DAC12 Data							
rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)	rw-(0)

Unused Bits 15-12 Unused. These bits are always 0 and do not affect the DAC12 core.

DAC12 Data Bits 11-0 DAC12 data

DAC12 Data Format	DAC12 Data
12-bit binary	The DAC12 data are right-justified. Bit 11 is the MSB.
12-bit 2s complement	The DAC12 data are right-justified. Bit 11 is the MSB (sign).
8-bit binary	The DAC12 data are right-justified. Bit 7 is the MSB. Bits 11-8 are don't care and do not affect the DAC12 core.
8-bit 2s complement	The DAC12 data are right-justified. Bit 7 is the MSB (sign). Bits 11-8 are don't care and do not affect the DAC12 core.



# SD16\_A

The SD16\_A module is a single-converter 16-bit sigma-delta analog-to-digital conversion module with high impedance input buffer. This chapter describes the SD16\_A. The SD16\_A module is implemented in the MSP430x20x3 devices.

Topic Page

26.1	SD16_A Introduction	<b>599</b>
26.2	SD16_A Operation	601
26.3	SD16_A Registers	611



www.ti.com SD16 A Introduction

### 26.1 SD16\_A Introduction

The SD16\_A module consists of one sigma-delta analog-to-digital converter with a high-impedance input buffer and an internal voltage reference. It has up to eight fully differential multiplexed analog input pairs including a built-in temperature sensor and a divided supply voltage. The converter is based on a second-order oversampling sigma-delta modulator and digital decimation filter. The decimation filter is a comb type filter with selectable oversampling ratios of up to 1024. Additional filtering can be done in software.

The high impedance input buffer is not implemented in MSP430x20x3 devices.

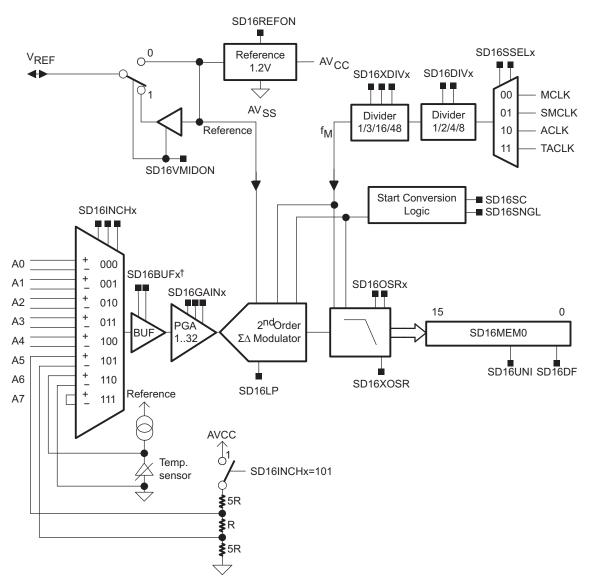
Features of the SD16\_A include:

- 16-bit sigma-delta architecture
- Up to eight multiplexed differential analog inputs per channel(The number of inputs is device dependent, see the device-specific data sheet.)
- Software selectable on-chip reference voltage generation (1.2 V)
- · Software selectable internal or external reference
- Built-in temperature sensor
- · Up to 1.1-MHz modulator input frequency
- · High impedance input buffer(not implemented on all devices, see the device-specific data sheet)
- · Selectable low-power conversion mode

The block diagram of the SD16\_A module is shown in Figure 26-1.



SD16\_A Introduction www.ti.com



<sup>†</sup> Not Implemented in MSP430x20x3 devices

Figure 26-1. SD16\_A Block Diagram



www.ti.com SD16\_A Operation

### 26.2 SD16\_A Operation

The SD16\_A module is configured with user software. The setup and operation of the SD16\_A is discussed in the following sections.

#### 26.2.1 ADC Core

The analog-to-digital conversion is performed by a 1-bit second-order sigma-delta modulator. A single-bit comparator within the modulator quantizes the input signal with the modulator frequency  $f_M$ . The resulting 1-bit data stream is averaged by the digital filter for the conversion result.

# 26.2.2 Analog Input Range and PGA

The full-scale input voltage range for each analog input pair is dependent on the gain setting of the programmable gain amplifier of each channel. The maximum full-scale range is  $\pm V_{FSR}$  where  $V_{FSR}$  is defined by:

$$V_{FSR} = \frac{V_{REF}/2}{GAIN_{PGA}}$$

For a 1.2-V reference, the maximum full-scale input range for a gain of 1 is:

$$\pm V_{FSR} = \frac{1.2 \text{ V/2}}{1} = \pm 0.6 \text{ V}$$

See the device-specific data sheet for full-scale input specifications.

# 26.2.3 Voltage Reference Generator

The SD16\_A module has a built-in 1.2-V reference. It is enabled by the SD16REFON bit. When using the internal reference an external 100-nF capacitor connected from  $V_{REF}$  to  $AV_{SS}$  is recommended to reduce noise. The internal reference voltage can be used off-chip when SD16VMIDON = 1. The buffered output can provide up to 1 mA of drive. When using the internal reference off-chip, a 470-nF capacitor connected from  $V_{REF}$  to  $AV_{SS}$  is required. See the device-specific data sheet for parameters.

An external voltage reference can be applied to the  $V_{\text{REF}}$  input when SD16REFON and SD16VMIDON are both reset.

## 26.2.4 Auto Power-Down

The SD16\_A is designed for low power applications. When the SD16\_A is not actively converting, it is automatically disabled and automatically re-enabled when a conversion is started. The reference is not automatically disabled, but can be disabled by setting SD16REFON = 0. When the SD16\_A or reference are disabled, they consume no current.

### 26.2.5 Analog Input Pair Selection

The SD16\_A can convert up to 8 differential input pairs multiplexed into the PGA. Up to five analog input pairs (A0-A4) are available externally on the device. A resistive divider to measure the supply voltage is available using the A5 multiplexer input. An internal temperature sensor is available using the A6 multiplexer input.

Input A7 is a shorted connection between the + and - input pair and can be used to calibrate the offset of the SD16\_A input stage. Note that the measured offset depends on the impedance of the external circuitry; thus, the actual offset seen at any of the analog inputs may be different.

# 26.2.5.1 Analog Input Setup

The analog input is configured using the SD16INCTL0 and the SD16AE registers. The SD16INCHx bits select one of eight differential input pairs of the analog multiplexer. The gain for the PGA is selected by the SD16GAINx bits. A total of six gain settings are available. The SD16AEx bits enable or disable the analog input pin. Setting any SD16AEx bit disables the multiplexed digital circuitry for the associated pin. See the device-specific data sheet for pin diagrams.



During conversion any modification to the SD16INCHx and SD16GAINx bits will become effective with the next decimation step of the digital filter. After these bits are modified, the next three conversions may be invalid due to the settling time of the digital filter. This can be handled automatically with the SD16INTDLYx bits. When SD16INTDLY = 00h, conversion interrupt requests will not begin until the fourth conversion after a start condition.

On devices implementing the high impedance input buffer it can be enabled using the SD16BUFx bits. The speed settings are selected based on the SD16 A modulator frequency as shown in Table 26-1.

SD16BUFx	Buffer	SD16 Modulator Frequency f <sub>M</sub>
00	Buffer disabled	
01	Low speed/current	f <sub>M</sub> < 200 kHz
10	Medium speed/current	200 kHz < f <sub>M</sub> < 700 kHz
11	High speed/current	700 kHz < f <sub>M</sub> < 1.1 MHz

Table 26-1. High Input Impedance Buffer

An external RC anti-aliasing filter is recommended for the SD16\_A to prevent aliasing of the input signal. The cutoff frequency should be < 10 kHz for a 1-MHz modulator clock and OSR = 256. The cutoff frequency may set to a lower frequency for applications that have lower bandwidth requirements.

# 26.2.6 Analog Input Characteristics

The SD16\_A uses a switched-capacitor input stage that appears as an impedance to external circuitry as shown in Figure 26-2.

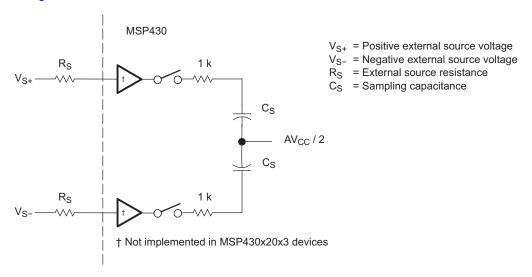


Figure 26-2. Analog Input Equivalent Circuit

When the buffers are used,  $R_S$  does not affect the sampling frequency  $f_S$ . However, when the buffers are not used or are not present on the device, the maximum sampling frequency  $f_S$  may be calculated from the minimum settling time  $t_{Settling}$  of the sampling circuit given by:

$$t_{Settling} \ \geq \ (R_S \ + \ 1 \ k\Omega) \times C_S \ \times \ In \Bigg( \frac{GAIN \times 2^{17} \ \times V_{Ax}}{V_{REF}} \Bigg)$$

where

$$f_{S} = \frac{1}{2 \times t_{Settling}} \quad \text{and} \quad V_{Ax} = max \Biggl( \left| \frac{AV_{CC}}{2} - V_{S+} \right| \; , \; \left| \frac{AV_{CC}}{2} - V_{S-} \right| \Biggr)$$

with  $V_{S+}$  and  $V_{S-}$  referenced to  $AV_{SS}$ .

C<sub>s</sub> varies with the gain setting as shown in Table 26-2.



www.ti.com SD16\_A Operation

Table 26-2.	Sampling	Capacitance
-------------	----------	-------------

PGA Gain	Sampling Capacitance, C <sub>s</sub>
1	1.25 pF
2, 4	2.5 pF
8	5 pF
16, 32	10 pF

# 26.2.7 Digital Filter

The digital filter processes the 1-bit data stream from the modulator using a SINC<sup>3</sup> comb filter. The transfer function is described in the z-Domain by:

$$H(z) = \left(\frac{1}{OSR} \times \frac{1 - z^{-OSR}}{1 - z^{-1}}\right)^3$$

and in the frequency domain by:

$$H(f) = \left[ \frac{\text{sinc} \bigg( \text{OSR} \times \pi \times \frac{f}{f_M} \bigg)}{\text{sinc} \bigg( \pi \times \frac{f}{f_M} \bigg)} \right]^3 = \left[ \frac{1}{\text{OSR}} \times \frac{\text{sin} \bigg( \text{OSR} \times \pi \times \frac{f}{f_M} \bigg)}{\text{sin} \bigg( \pi \times \frac{f}{f_M} \bigg)} \right]^3$$

where the oversampling rate, OSR, is the ratio of the modulator frequency  $f_M$  to the sample frequency  $f_S$ . Figure 26-3 shows the filter's frequency response for an OSR of 32. The first filter notch is at  $f_S = f_M/OSR$ . The notch's frequency can be adjusted by changing the modulator's frequency,  $f_M$ , using SD16SSELx and SD16DIVx and the oversampling rate using the SD16OSRx and SD16XOSR bits.

The digital filter for each enabled ADC channel completes the decimation of the digital bit-stream and outputs new conversion results to the SD16MEM0 register at the sample frequency  $f_s$ .

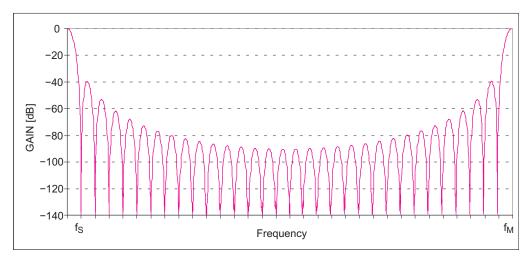


Figure 26-3. Comb Filter Frequency Response With OSR = 32

Figure 26-4 shows the digital filter step response and conversion points. For step changes at the input after start of conversion a settling time must be allowed before a valid conversion result is available. The SD16INTDLYx bits can provide sufficient filter settling time for a full-scale change at the ADC input. If the step occurs synchronously to the decimation of the digital filter the valid data will be available on the third conversion. An asynchronous step will require one additional conversion before valid data is available.



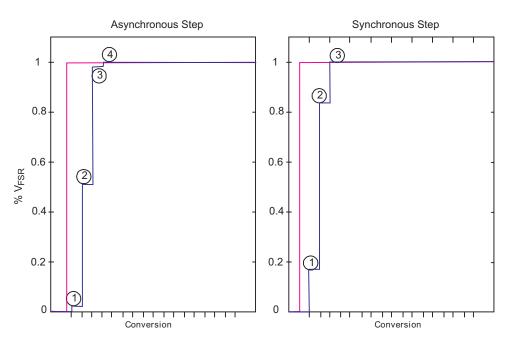


Figure 26-4. Digital Filter Step Response and Conversion Points

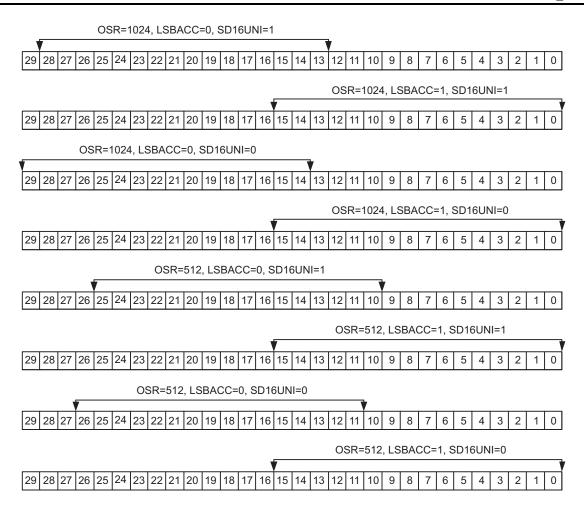
### 26.2.7.1 Digital Filter Output

The number of bits output by the digital filter is dependent on the oversampling ratio and ranges from 15 to 30 bits. Figure 26-5 shows the digital filter output and their relation to SD16MEM0 for each OSR, LSBACC, and SD16UNI setting. For example, for OSR = 1024, LSBACC = 0, and SD16UNI = 1, the SD16MEM0 register contains bits 28 - 13 of the digital filter output. When OSR = 32, the one (SD16UNI = 0) or two (SD16UNI=1) LSBs are always zero.

The SD16LSBACC and SD16LSBTOG bits give access to the least significant bits of the digital filter output. When SD16LSBACC = 1 the 16 least significant bits of the digital filter's output are read from SD16MEM0 using word instructions. The SD16MEM0 register can also be accessed with byte instructions returning only the 8 least significant bits of the digital filter output.

When SD16LSBTOG = 1 the SD16LSBACC bit is automatically toggled each time SD16MEM0 is read. This allows the complete digital filter output result to be read with two reads of SD16MEM0. Setting or clearing SD16LSBTOG does not change SD16LSBACC until the next SD16MEM0 access.







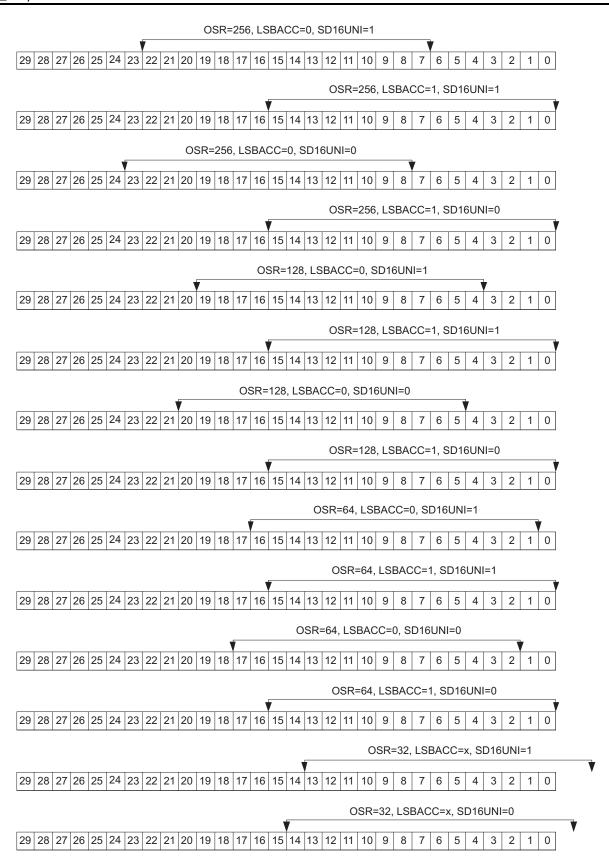


Figure 26-5. Used Bits of Digital Filter Output



# 26.2.8 Conversion Memory Register: SD16MEM0

The SD16MEM0 register is associated with the SD16\_A channel. Conversion results are moved to the SD16MEM0 register with each decimation step of the digital filter. The SD16IFG bit is set when new data is written to SD16MEM0. SD16IFG is automatically cleared when SD16MEM0 is read by the CPU or may be cleared with software.

### 26.2.8.1 Output Data Format

The output data format is configurable in two's complement, offset binary or unipolar mode as shown in Table 26-3. The data format is selected by the SD16DF and SD16UNI bits.

Tabl	ا2 ما	6-3	Data	Forn	nat

SD16UNI	SD16DF	Format	Analog Input	SD16MEM0 <sup>(1)</sup>	Digital Filter Output (OSR = 256)	
0	0	Bipolar Offset Binary	+FSR	FFFF	FFFFF	
			ZERO	8000	800000	
			-FSR	0000	000000	
0	1	Bipolar Twos Compliment	+FSR	7FFF	7FFFF	
			ZERO	0000	000000	
			-FSR	8000	800000	
1	0	Unipolar	+FSR	FFFF	FFFFF	
			ZERO	0000	800000	
			-FSR	0000	000000	

Independent of SD16OSRx and SD16XOSR settings; SD16LSBACC = 0.

#### NOTE: **Offset Measurements and Data Format**

Any offset measurement done either externally or using the internal differential pair A7 would be appropriate only when the channel is operating under bipolar mode with SD16UNI = 0.

Figure 26-6 shows the relationship between the full-scale input voltage range from -V<sub>FSR</sub> to +V<sub>FSR</sub> and the conversion result. The data formats are illustrated.

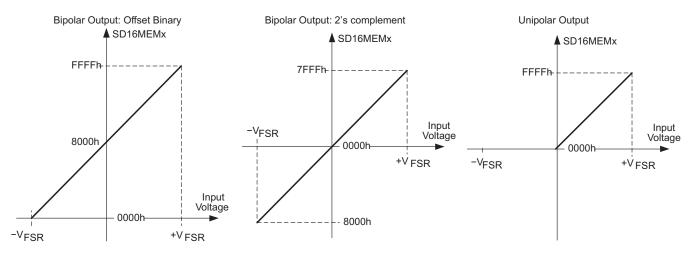


Figure 26-6. Input Voltage vs Digital Output



#### 26.2.9 Conversion Modes

The SD16\_A module can be configured for two modes of operation, listed in Table 26-4. The SD16SNGL bit selects the conversion mode.

Table 26-4. Conversion Mode Summary

SD16SNGL	Mode	Operation		
1	Single conversion	The channel is converted once.		
0	Continuous conversion	The channel is converted continuously.		

### 26.2.9.1 Single Conversion

Setting the SD16SC bit of the channel initiates one conversion on that channel when SD16SNGL = 1. The SD16SC bit will automatically be cleared after conversion completion.

Clearing SD16SC before the conversion is completed immediately stops conversion of the channel, the channel is powered down and the corresponding digital filter is turned off. The value in SD16MEM0 can change when SD16SC is cleared. It is recommended that the conversion data in SD16MEM0 be read prior to clearing SD16SC to avoid reading an invalid result.

#### 26.2.9.2 Continuous Conversion

When SD16SNGL = 0 continuous conversion mode is selected. Conversion of the channel will begin when SD16SC is set and continue until the SD16SC bit is cleared by software.

Clearing SD16SC immediately stops conversion of the selected channel, the channel is powered down and the corresponding digital filter is turned off. The value in SD16MEM0 can change when SD16SC is cleared. It is recommended that the conversion data in SD16MEM0 be read prior to clearing SD16SC to avoid reading an invalid result.

Figure 26-7 shows conversion operation.

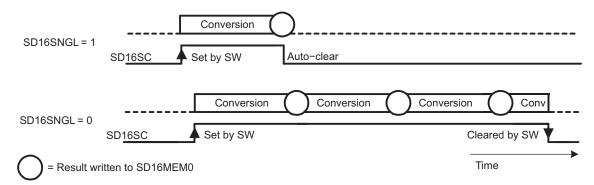


Figure 26-7. Single Channel Operation

### 26.2.10 Using the Integrated Temperature Sensor

To use the on-chip temperature sensor, the user selects the analog input pair SD16INCHx = 110 and sets SD16REFON = 1. Any other configuration is done as if an external analog input pair was selected, including SD16INTDLYx and SD16GAINx settings. Because the internal reference must be on to use the temperature sensor, it is not possible to use an external reference for the conversion of the temperature sensor voltage. Also, the internal reference will be in contention with any used external reference. In this case, the SD16VMIDON bit may be set to minimize the affects of the contention on the conversion.

The typical temperature sensor transfer function is shown in Figure 26-8. When switching inputs of an SD16\_A channel to the temperature sensor, adequate delay must be provided using SD16INTDLYx to allow the digital filter to settle and assure that conversion results are valid. The temperature sensor offset error can be large, and may need to be calibrated for most applications. See device-specific data sheet for temperature sensor parameters.



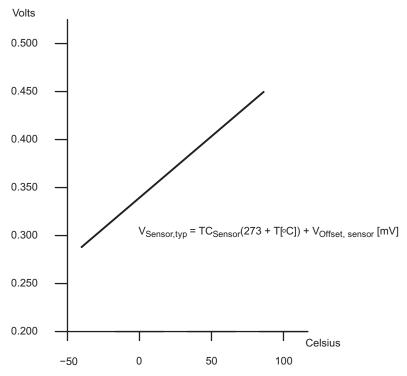


Figure 26-8. Typical Temperature Sensor Transfer Function

### 26.2.11 Interrupt Handling

The SD16\_A has 2 interrupt sources for its ADC channel:

- SD16IFG
- SD16OVIFG

The SD16IFG bit is set when the SD16MEM0 memory register is written with a conversion result. An interrupt request is generated if the corresponding SD16IE bit and the GIE bit are set. The SD16 A overflow condition occurs when a conversion result is written to SD16MEM0 location before the previous conversion result was read.

### 26.2.11.1 SD16IV, Interrupt Vector Generator

All SD16\_A interrupt sources are prioritized and combined to source a single interrupt vector. SD16IV is used to determine which enabled SD16 A interrupt source requested an interrupt. The highest priority SD16\_A interrupt request that is enabled generates a number in the SD16IV register (see register description). This number can be evaluated or added to the program counter to automatically enter the appropriate software routine. Disabled SD16 A interrupts do not affect the SD16IV value.

Any access, read or write, of the SD16IV register has no effect on the SD16OVIFG or SD16IFG flags. The SD16IFG flags are reset by reading the SD16MEM0 register or by clearing the flags in software. SD16OVIFG bits can only be reset with software.

If another interrupt is pending after servicing of an interrupt, another interrupt is generated. For example, if the SD16OVIFG and one or more SD16IFG interrupts are pending when the interrupt service routine accesses the SD16IV register, the SD16OVIFG interrupt condition is serviced first and the corresponding flag(s) must be cleared in software. After the RETI instruction of the interrupt service routine is executed, the highest priority SD16IFG pending generates another interrupt request.

SD16\_A



# 26.2.11.2 Interrupt Delay Operation

The SD16INTDLYx bits control the timing for the first interrupt service request for the corresponding channel. This feature delays the interrupt request for a completed conversion by up to four conversion cycles allowing the digital filter to settle prior to generating an interrupt request. The delay is applied each time the SD16SC bit is set or when the SD16GAINx or SD16INCHx bits for the channel are modified. SD16INTDLYx disables overflow interrupt generation for the channel for the selected number of delay cycles. Interrupt requests for the delayed conversions are not generated during the delay.



www.ti.com SD16\_A Registers

# 26.3 SD16\_A Registers

The SD16\_A registers are listed in Table 26-5.

# Table 26-5. SD16\_A Registers

Register	Short Form	Register Type	Address	Initial State
SD16_A control	SD16CTL	Read/write	0100h	Reset with PUC
SD16_A interrupt vector	SD16IV	Read/write	0110h	Reset with PUC
SD16_A channel 0 control	SD16CCTL0	Read/write	0102h	Reset with PUC
SD16_A conversion memory	SD16MEM0	Read/write	0112h	Reset with PUC
SD16_A input control	SD16INCTL0	Read/write	0B0h	Reset with PUC
SD16_A analog enable	SD16AE	Read/write	0B7h	Reset with PUC



SD16\_A Registers www.ti.com

# 26.3.1 SD16CTL, SD16\_A Control Register

15	14		13	12	11	10	9	8
	Reserved			SD16XDIVx			SD16LP	
r0	r0		r0	r0	rw-0	rw-0	rw-0	rw-0
7	6		5	4	3	2	1	0
SD10	SD16DIVx		SD16	SSELx	SD16VMIDON	SD16REFON	SD160VIE	Reserved
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	r0
Reserved	Bits 15-12	Reser	ved					
SD16XDIVx	Bits 11-9	SD16	_A clock divide	er				
		000	/1					
		001	/3					
		010	/16					
		011	/48					
		1xx	Reserved					
SD16LP	Bit 8	Low power mode. This bit selects a reduced speed, reduced power mode						
		0	0 Low-power mode is disabled					
		1	Low-power	mode is enabled.	The maximum cloc	ck frequency for th	e SD16_A is redu	ced.
SD16DIVx	Bits 7-6	SD16	_A clock divide	er				
		00	/1					
		01	/2					
		10	/4					
		11	/8					
SD16SSELx	Bits 5-4	SD16	_A clock source	e select				
		00	MCLK					
		01	SMCLK					
		10	ACLK					
		11	External TA	CLK				
SD16VMIDON	Bit 3	VMID	buffer on					
		0	Off					
		1	On					
SD16REFON	Bit 2	Reference generator on						
		0	Reference	off				
		1	Reference					
SD16OVIE	Bit 1	SD16_A overflow interrupt enable. The GIE bit must also be set to enable the interrupt.  Overflow interrupt disabled						
		1		terrupt enabled				
Reserved	Bit 0	Reser	ved					



www.ti.com SD16\_A Registers

# 26.3.2 SD16CCTL0, SD16\_A Control Register 0

15	14	13	12	11	10	9	8
Reserved	S	D16BUFx <sup>(1)</sup>	SD16UNI	SD16XOSR	SD16SNGL	SD16	OSRx
r0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
7	6	5	4	3	2	1	0
SD16LSBTOG	SD16LSBA	CC SD160VIFG	SD16DF	SD16IE	SD16IFG	SD16SC	Reserved
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	r-0
Reserved	Bit 15	Reserved					
SD16BUFx	Bits 14-13	High-impedance inpu	ıt buffer mode				
		00 Buffer disab	led				
		01 Slow speed/	current				
		10 Medium spe	ed/current				
		11 High speed/	current				
SD16UNI	Bit 12	Unipolar mode selec	t				
		0 Bipolar mod	е				
		1 Unipolar mo	de				
SD16XOSR	Bit 11	Extended oversampl	ing ratio. This bit,		16OSRx bits, sele	ct the oversampli	ng ratio. See
CD4CCNO!	D:+ 40	SD16OSRx bit descr	-	•			
SD16SNGL	Bit 10	Single conversion me					
			conversion mode				
CD4COCD.	D:+- 0 0	1 Single conve	ersion mode				
SD16OSRx	Bits 9-8	Oversampling ratio	0				
		When SD16XOSR =	U				
		00 256					
		01 128					
		10 64 11 32					
		*-	4				
		When SD16XOSR =	1				
		00 512 01 1024					
		10 Reserved					
		11 Reserved					
SD16LSBTOG	Bit 7	LSB toggle. This bit,	when set source	CD16  CD160 to	togalo ocob timo t	ha SD16MEMO r	ogiator in road
3010130100	DIL 1			e with each SD16l		THE SET TOWNER TO	egister is reau.
				ach SD16MEM0 re			
SD16LSBACC	Bit 6	LSB access. This bit				16 A conversion	recult
ODTOLOBAGO	Dit 0				s of the conversion		rosuit.
				=	s of the conversion		
SD160VIFG	Bit 5	SD16_A overflow int		t digrimount 10 bit	or the conversion	•	
021001110	Dit 0		interrupt pending				
			errupt pending				
SD16DF	Bit 4	SD16 A data format					
		0 Offset binary					
		1 2's complem					
SD16IE	Bit 3	SD16_A interrupt en					
-		0 Disabled	-				
		1 Enabled					
SD16IFG	Bit 2	SD16_A interrupt flag	g. SD16IFG is set	when new convei	rsion results are av	ailable. SD16IFG	is automatically
-		reset when the corre					
		0 No interrupt	pending				
		1 Interrupt per	nding				

<sup>(1)</sup> Reserved in MSP430x20x3 devices



SD16\_A Registers www.ti.com

**SD16SC** Bit 1 SD16\_A start conversion

0 No conversion start

Start conversion

Reserved Bit 0 Reserved

# 26.3.3 SD16INCTL0, SD16\_A Input Control Register

7	6		5	4	3	2	1	0	
SD16IN	NTDLYx		SD16GAINx			SD16INCHx			
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	
SD16INTDLYx	Bits 7-6		ıpt delay gener rsion start.	ation after conversi	on start. These	bits select the dela	y for the first inter	rupt after	
		00	Fourth samp	le causes interrupt					
		01	Third sample	e causes interrupt					
		10	Second sam	ple causes interrup	t				
		11	First sample	causes interrupt					
SD16GAINx	Bits 5-3	SD16	_A preamplifier	gain					
		000	x1						
		001	x2						
		010	x4						
		011	x8						
		100	x16						
		101	x32						
		110	Reserved						
		111	Reserved						
SD16INCHx	Bits 2-0	SD16	_A channel diff	erential pair input					
		000	A0						
		001	A1						
		010	A2						
		011	A3						
		100	A4						
		101	A5 - (AV <sub>CC</sub> -	AV <sub>SS</sub> ) / 11					
		110	A6 - Temper	ature Sensor					
		111	A7 - Short fo	or PGA offset meas	urement				
			AT - SHOILIC	n i On onset meas	aromont.				



www.ti.com SD16\_A Registers

# 26.3.4 SD16MEM0, SD16\_A Conversion Memory Register

15	14	13	12	11	10	9	8		
Conversion Results									
r	r	r	r	r	r	r	r		
7	6	5	4	3	2	1	0		
Conversion Results									
r	r	r	r	r	r	r	r		

Conversion Results

Bits 15-0

Conversion Results. The SD16MEMx register holds the upper or lower 16-bits of the digital filter output, depending on the SD16LSBACC bit.

# 26.3.5 SD16AE, SD16\_A Analog Input Enable Register

7	6	5	4	3	2	1	0
SD16AE7	SD16AE6	SD16AE5	SD16AE4	SD16AE3	SD16AE2	SD16AE1	SD16AE0
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
SD16AEx	Bits 7-0	SD16 A analog enab	ole				

0 External input disabled. Negative inputs are internally connected to VSS.

1 External input enabled.

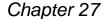
# 26.3.6 SD16IV, SD16\_A Interrupt Vector Register

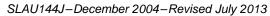
15	14	13	12	11	10	9	8
0	0	0	0	0	0	0	0
r0	r0	r0	r0	r0	rO	rO	rO
7	6	5	4	3	2	1	0
0	0	0	SD16IVx			0	
rO	r0	rO					r0

**SD16IVx** Bits 15-0 SD16\_A interrupt vector value

SD16IV Contents	Interrupt Source	Interrupt Flag	Interrupt Priority
000h	No interrupt pending	=	
002h	SD16MEMx overflow	SD16CCTLx SD16OVIFG	Highest
004h	SD16_A interrupt	SD16CCTL0 SD16IFG	
006h	Reserved	=	
008h	Reserved	-	
00Ah	Reserved	-	
00Ch	Reserved	-	
00Eh	Reserved	-	
010h	Reserved	-	Lowest

SD16\_A







# SD24\_A

The SD24\_A module is a multichannel 24-bit sigma-delta analog-to-digital converter (ADC). This chapter describes the SD24\_A of the MSP430x2xx family.

Topic Page

27.1	SD24_A Introduction	617
27.2	SD24_A Operation	619
27.3	SD24_A Registers	632



www.ti.com SD24\_A Introduction

#### 27.1 **SD24 A Introduction**

The SD24 A module consists of up to seven independent sigma-delta analog-to-digital converters, referred to as channels, and an internal voltage reference. Each channel has up to eight fully differential multiplexed analog input pairs including a built-in temperature sensor and a divided supply voltage. The converters are based on second-order oversampling sigma-delta modulators and digital decimation filters. The decimation filters are comb-type filters with selectable oversampling ratios of up to 1024. Additional filtering can be done in software.

The digital filter output of SD24\_A can range from 15 bits to 30 bits, based on the oversampling ratio. The default oversampling ratio is 256, which results in 24-bit output from the digital filter. The 16 most significant bits of the filter are captured in the SD24 A conversion memory register and, by setting SD24LSBACC = 1, the 16 least significant bits of the filter output can be read (see Section 27.2.7 for details).

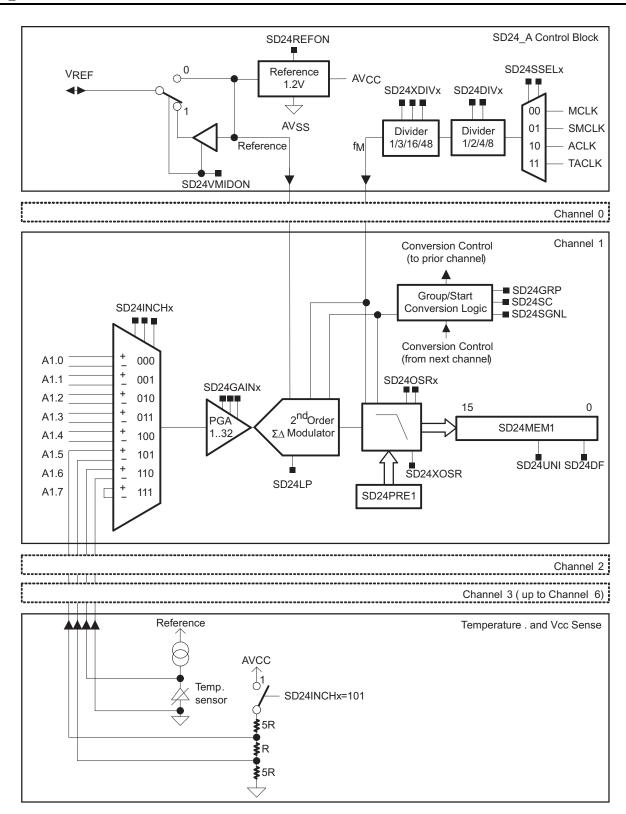
Features of the SD24 A include:

- Up to seven independent, simultaneously-sampling ADC channels (the number of channels is device dependent, see the device-specific data sheet)
- Up to eight multiplexed differential analog inputs per channel (the number of inputs is device dependent, see the device-specific data sheet)
- Software selectable on-chip reference voltage generation (1.2 V)
- Software selectable internal or external reference
- Built-in temperature sensor accessible by all channels
- Up to 1.1-MHz modulator input frequency
- High impedance input buffer (not implemented on all devices, see the device-specific data sheet)
- Selectable low-power conversion mode

The block diagram of the SD24\_A module is shown in Figure 27-1.



SD24\_A Introduction www.ti.com



NOTE: Ax.1 to Ax.4 not available on all devices. See device-specific data sheet.

Figure 27-1. Block Diagram of the SD24\_A



www.ti.com SD24\_A Operation

## 27.2 SD24\_A Operation

The SD24\_A module is configured with user software. The setup and operation of the SD24\_A is discussed in the following sections.

#### 27.2.1 ADC Core

The analog-to-digital conversion is performed by a 1-bit second-order sigma-delta modulator. A single-bit comparator within the modulator quantizes the input signal with the modulator frequency  $f_M$ . The resulting 1-bit data stream is averaged by the digital filter for the conversion result.

### 27.2.2 Analog Input Range and PGA

The full-scale input voltage range for each analog input pair is dependent on the gain setting of the programmable gain amplifier of each channel. The maximum full-scale range is  $\pm V_{FSR}$  where  $V_{FSR}$  is defined by:

$$V_{FSR} = \frac{V_{REF}/2}{GAIN_{PGA}}$$

For a 1.2-V reference, the maximum full-scale input range for a gain of 1 is:

$$\pm V_{FSR} = \frac{1.2 \text{ V/2}}{1} = \pm 0.6 \text{ V}$$

See the device-specific data sheet for full-scale input specifications.

### 27.2.3 Voltage Reference Generator

The SD24\_A module has a built-in 1.2-V reference. It can be used for each SD24\_A channel and is enabled by the SD24REFON bit. When using the internal reference an external 100-nF capacitor connected from  $V_{REF}$  to  $AV_{SS}$  is recommended to reduce noise. The internal reference voltage can be used off-chip when SD24VMIDON = 1. The buffered output can provide up to 1 mA of drive. When using the internal reference off-chip, a 470-nF capacitor connected from  $V_{REF}$  to  $AV_{SS}$  is required. See device-specific data sheet for parameters.

An external voltage reference can be applied to the  $V_{\text{REF}}$  input when SD24REFON and SD24VMIDON are both reset.

#### 27.2.4 Auto Power-Down

The SD24\_A is designed for low-power applications. When the SD24\_A is not actively converting, it is automatically disabled and automatically re-enabled when a conversion is started. The reference is not automatically disabled, but it can be disabled by setting SD24REFON = 0. When the SD24\_A or reference are disabled, they consume no current.

### 27.2.5 Analog Input Pair Selection

The SD24\_A can convert up to eight differential input pairs multiplexed into the PGA. Up to five analog input pairs (A0 to A4) are available externally on the device. A resistive divider to measure the supply voltage is available using the A5 multiplexer input. An internal temperature sensor is available using the A6 multiplexer input.

Input A7 is a shorted connection between the + and – input pair and can be used to calibrate the offset of the SD24\_A input stage. Note that the measured offset depends on the impedance of the external circuitry; thus, the actual offset seen at any of the analog inputs may be different.

#### 27.2.5.1 Analog Input Setup

The analog input of each channel is configured using the SD24INCTLx register. These settings can be independently configured for each SD24 A channel.



SD24 A Operation www.ti.com

The SD24INCHx bits select one of eight differential input pairs of the analog multiplexer. The gain for each PGA is selected by the SD24GAINx bits. A total of six gain settings are available.

On some devices SD24AEx bits are available to enable or disable the analog input pin. Setting any SD24AEx bit disables the multiplexed digital circuitry for the associated pin. See the device-specific data sheet for pin diagrams.

During conversion any modification to the SD24INCHx and SD24GAINx bits will become effective with the next decimation step of the digital filter. After these bits are modified, the next three conversions may be invalid due to the settling time of the digital filter. This can be handled automatically with the SD24INTDLYx bits. When SD24INTDLY = 00h, conversion interrupt requests will not begin until the fourth conversion after a start condition.

On devices implementing the high impedance input buffer it can be enabled using the SD24BUFx bits. The speed settings are selected based on the SD24\_A modulator frequency as shown in Table 27-1.

SD24BUFx	Buffer	SD24 Modulator Frequency, f <sub>M</sub>
00	Buffer disabled	
01	Low speed/current	f <sub>M</sub> < 200 kHz
10	Medium speed/current	200 kHz < f <sub>M</sub> < 700 kHz
11	High speed/current	700 kHz < f <sub>M</sub> < 1.1 MHz

Table 27-1. High Input Impedance Buffer

An external RC anti-aliasing filter is recommended for the SD24\_A to prevent aliasing of the input signal. The cutoff frequency should be less than 10 kHz for a 1-MHz modulator clock and OSR = 256. The cutoff frequency may set to a lower frequency for applications that have lower bandwidth requirements.

### 27.2.6 Analog Input Characteristics

The SD24\_A uses a switched-capacitor input stage that appears as an impedance to external circuitry as shown in Figure 27-2.

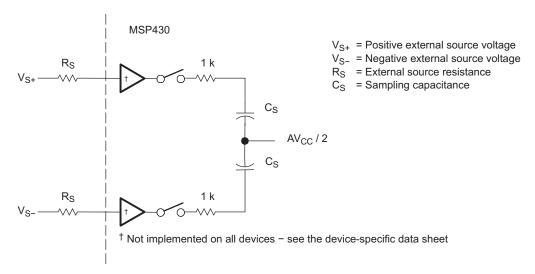


Figure 27-2. Analog Input Equivalent Circuit



www.ti.com SD24\_A Operation

When the buffers are used, R<sub>s</sub> does not affect the sampling frequency f<sub>s</sub>. However, when the buffers are not used or are not present on the device, the maximum modulator frequency f<sub>M</sub> may be calculated from the minimum settling time t<sub>Settling</sub> of the sampling circuit given by:

$$t_{Settling} \ \geq \ (R_S \ + \ 1 \ k\Omega) \times C_S \times In \Bigg( \frac{GAIN \times 2^{17} \ \times V_{AX}}{V_{REF}} \Bigg)$$

Where.

$$f_{M} = \frac{1}{2 \times t_{Settling}} \quad \text{and} \quad V_{Ax} = max \Biggl( \left| \frac{AV_{CC}}{2} - V_{S+} \right| \; , \; \left| \frac{AV_{CC}}{2} - V_{S-} \right| \Biggr)$$

with  $V_{\text{S+}}$  and  $V_{\text{S-}}$  referenced to  $AV_{\text{SS}}.$ 

C<sub>s</sub> varies with the gain setting as shown in Table 27-2.

Table 27-2. Sampling Capacitance

PGA Gain	Sampling Capacitance (C <sub>s</sub> )
1	1.25 pF
2, 4	2.5 pF
8	5 pF
16, 32	10 pF

### 27.2.7 Digital Filter

The digital filter processes the 1-bit data stream from the modulator using a SINC<sup>3</sup> comb filter. The transfer function is described in the z-Domain by:

$$H(z) = \left(\frac{1}{OSR} \times \frac{1 - z^{-OSR}}{1 - z^{-1}}\right)^3$$

and in the frequency domain by:

$$H(f) = \left\lceil \frac{sinc \left( OSR \times \pi \times \frac{f}{f_M} \right)}{sinc \left( \pi \times \frac{f}{f_M} \right)} \right\rceil^3 = \left\lceil \frac{1}{OSR} \times \frac{sin \left( OSR \times \pi \times \frac{f}{f_M} \right)}{sin \left( \pi \times \frac{f}{f_M} \right)} \right\rceil^3$$

where the oversampling rate, OSR, is the ratio of the modulator frequency f<sub>M</sub> to the sample frequency f<sub>S</sub>. Figure 27-3 shows the filter's frequency response for an OSR of 32. The first filter notch is at  $f_S = f_M/OSR$ . The notch frequency can be adjusted by changing the modulator frequency, f<sub>M</sub>, using SD24SSELx and SD24DIVx and the oversampling rate using the SD24OSRx and SD24XOSR bits.

The digital filter for each enabled ADC channel completes the decimation of the digital bit-stream and outputs new conversion results to the corresponding SD24MEMx register at the sample frequency fs.



SD24\_A Operation www.ti.com

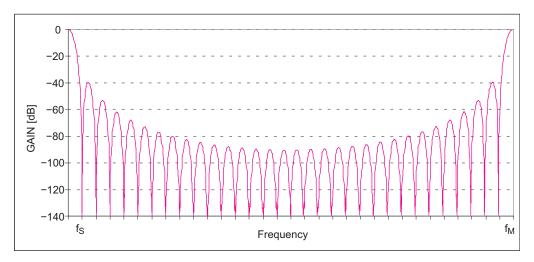


Figure 27-3. Comb Filter Frequency Response With OSR = 32

Figure 27-4 shows the digital filter step response and conversion points. For step changes at the input after start of conversion a settling time must be allowed before a valid conversion result is available. The SD24INTDLYx bits can provide sufficient filter settling time for a full-scale change at the ADC input. If the step occurs synchronously to the decimation of the digital filter the valid data will be available on the third conversion. An asynchronous step will require one additional conversion before valid data is available.

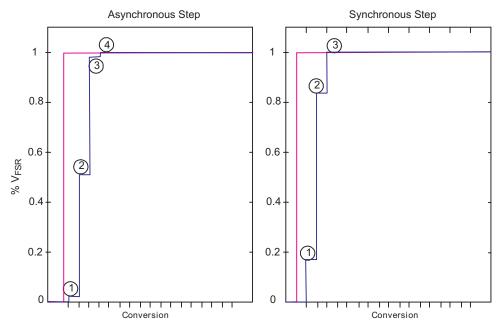


Figure 27-4. Digital Filter Step Response and Conversion Points



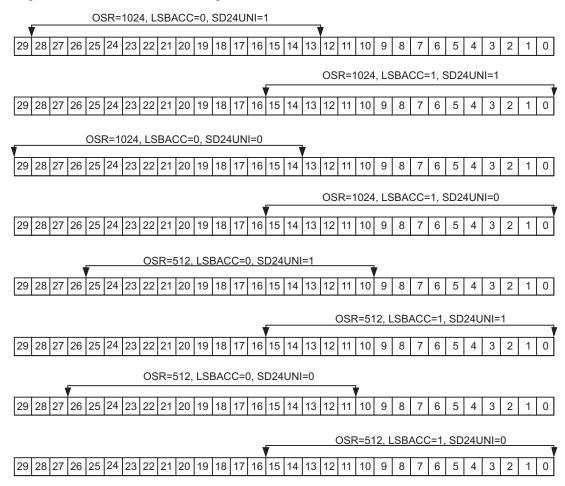
www.ti.com SD24\_A Operation

#### 27.2.7.1 Digital Filter Output

The number of bits output by the digital filter is dependent on the oversampling ratio and ranges from 15 to 30 bits. Figure 27-5 shows the digital filter output and their relation to SD24MEMx for each OSR, LSBACC, and SD24UNI setting. For example, for OSR = 1024, LSBACC = 0, and SD24UNI = 1, the SD24MEMx register contains bits 28 to 13 of the digital filter output. When OSR = 32, the one (SD24UNI = 0) or two (SD24UNI = 1) LSBs are always zero.

The SD24LSBACC and SD24LSBTOG bits give access to the least significant bits of the digital filter output. When SD24LSBACC = 1 the 16 least significant bits of the digital filter's output are read from SD24MEMx using word instructions. The SD24MEMx register can also be accessed with byte instructions returning only the 8 least significant bits of the digital filter output.

When SD24LSBTOG = 1 the SD24LSBACC bit is automatically toggled each time SD24MEMx is read. This allows the complete digital filter output result to be read with two reads of SD24MEMx. Setting or clearing SD24LSBTOG does not change SD24LSBACC until the next SD24MEMx access.





SD24\_A Operation www.ti.com

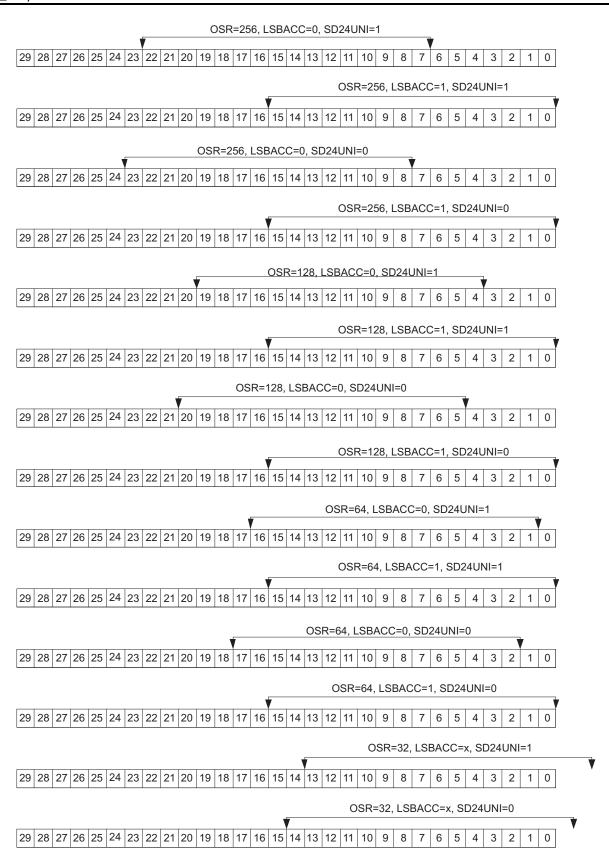


Figure 27-5. Used Bits of Digital Filter Output



www.ti.com SD24 A Operation

# 27.2.8 Conversion Memory Register: SD24MEMx

One SD24MEMx register is associated with each SD24\_A channel. Conversion results are moved to the corresponding SD24MEMx register with each decimation step of the digital filter. The SD24IFG bit is set when new data is written to SD24MEMx. SD24IFG is automatically cleared when SD24MEMx is read by the CPU or may be cleared with software.

### 27.2.8.1 Output Data Format

The output data format is configurable in twos complement, offset binary or unipolar mode as shown in Table 27-3. The data format is selected by the SD24DF and SD24UNI bits.

Tabl	<u>ام</u> 2	7-3	Data	Forr	nat

SD24UNI	SD24DF	Format	Analog Input	SD24MEMx (1)	Digital Filter Output (OSR = 256)
		Bipolar offset binary	+FSR	FFFF	FFFFF
0	0 0		ZERO	8000	800000
			-FSR	0000	000000
			+FSR	7FFF	7FFFF
0	1	Bipolar twos compliment	ZERO	0000	000000
			-FSR	8000	800000
		0 Unipolar	+FSR	FFFF	FFFFF
1	0		ZERO	0000	800000
			-FSR	0000	000000

<sup>(1)</sup> Independent of SD24OSRx and SD24XOSR settings; SD24LSBACC = 0.

#### NOTE: Offset Measurements and Data Format

Any offset measurement done either externally or using the internal differential pair A7 would be appropriate only when the channel is operating under bipolar mode with SD24UNI = 0.

If the measured value is to be used in the unipolar mode for offset correction it needs to be multiplied by two.

Figure 27-6 shows the relationship between the full-scale input voltage range from  $-V_{FSR}$  to  $+V_{FSR}$  and the conversion result. The data formats are illustrated.

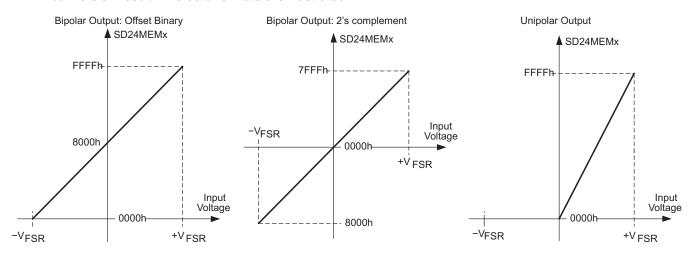


Figure 27-6. Input Voltage vs Digital Output



SD24 A Operation www.ti.com

#### 27.2.9 Conversion Modes

The SD24\_A module can be configured for four modes of operation, listed in Table 27-4. The SD24SNGL and SD24GRP bits for each channel selects the conversion mode.

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SD24SNGL	SD24GRP (1)	Mode	Operation
1	0	Single channel, Single conversion	A single channel is converted once.
0	0	Single channel, Continuous conversion	A single channel is converted continuously.
1	1	Group of channels, Single conversion	A group of channels is converted once.
0	1	Group of channels, Continuous conversion	A group of channels is converted continuously.

<sup>(1)</sup> A channel is grouped and is the master channel of the group when SD24GRP = 0 if SD24GRP for the prior channel(s) is set.

#### 27.2.9.1 Single Channel, Single Conversion

Setting the SD24SC bit of a channel initiates one conversion on that channel when SD24SNGL = 1 and it is not grouped with any other channels. The SD24SC bit will automatically be cleared after conversion completion.

Clearing SD24SC before the conversion is completed immediately stops conversion of the selected channel, the channel is powered down and the corresponding digital filter is turned off. The value in SD24MEMx can change when SD24SC is cleared. It is recommended that the conversion data in SD24MEMx be read prior to clearing SD24SC to avoid reading an invalid result.

## 27.2.9.2 Single Channel, Continuous Conversion

When SD24SNGL = 0 continuous conversion mode is selected. Conversion of the selected channel will begin when SD24SC is set and continue until the SD24SC bit is cleared by software when the channel is not grouped with any other channel.

Clearing SD24SC immediately stops conversion of the selected channel, the channel is powered down and the corresponding digital filter is turned off. The value in SD24MEMx can change when SD24SC is cleared. It is recommended that the conversion data in SD24MEMx be read prior to clearing SD24SC to avoid reading an invalid result.

Figure 27-7 shows single channel operation for single conversion mode and continuous conversion mode.

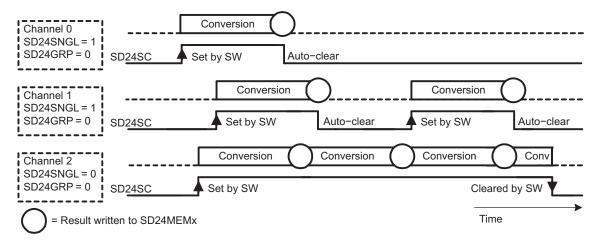


Figure 27-7. Single Channel Operation - Example



www.ti.com SD24\_A Operation

## 27.2.9.3 Group of Channels, Single Conversion

Consecutive SD24\_A channels can be grouped together with the SD24GRP bit to synchronize conversions. Setting SD24GRP for a channel groups that channel with the next channel in the module. For example, setting SD24GRP for channel 0 groups that channel with channel 1. In this case, channel 1 is the master channel, enabling and disabling conversion of all channels in the group with its SD24SC bit. The SD24GRP bit of the master channel is always 0. The SD24GRP bit of last channel in SD24\_A has no function and is always 0.

When SD24SNGL = 1 for a channel in a group, single conversion mode is selected. A single conversion of that channel will occur synchronously when the master channel SD24SC bit is set. The SD24SC bit of all channels in the group will automatically be set and cleared by SD24SC of the master channel. SD24SC for each channel can also be cleared in software independently.

Clearing SD24SC of the master channel before the conversions are completed immediately stops conversions of all channels in the group, the channels are powered down and the corresponding digital filters are turned off. Values in SD24MEMx can change when SD24SC is cleared. It is recommended that the conversion data in SD24MEMx be read prior to clearing SD24SC to avoid reading an invalid result.

#### 27.2.9.4 Group of Channels, Continuous Conversion

When SD24SNGL = 0 for a channel in a group, continuous conversion mode is selected. Continuous conversion of that channel will occur synchronously when the master channel SD24SC bit is set. SD24SC bits for all grouped channels will be automatically set and cleared with the master channel's SD24SC bit. SD24SC for each channel in the group can also be cleared in software independently.

When SD24SC of a grouped channel is set by software independently of the master, conversion of that channel will automatically synchronize to conversions of the master channel. This ensures that conversions for grouped channels are always synchronous to the master.

Clearing SD24SC of the master channel immediately stops conversions of all channels in the group the channels are powered down and the corresponding digital filters are turned off. Values in SD24MEMx can change when SD24SC is cleared. It is recommended that the conversion data in SD24MEMx be read prior to clearing SD24SC to avoid reading an invalid result.

Figure 27-8 shows grouped channel operation for three SD24\_A channels. Channel 0 is configured for single conversion mode, SD24SNGL = 1, and channels 1 and 2 are in continuous conversion mode, SD24SNGL = 0. Channel two, the last channel in the group, is the master channel. Conversions of all channels in the group occur synchronously to the master channel regardless of when each SD24SC bit is set using software.

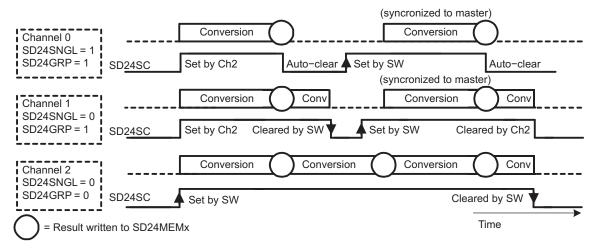


Figure 27-8. Grouped Channel Operation - Example



SD24 A Operation www.ti.com

# 27.2.10 Conversion Operation Using Preload

When multiple channels are grouped the SD24PREx registers can be used to delay the conversion time frame for each channel. Using SD24PREx, the decimation time of the digital filter is increased by the specified number of  $f_{\rm M}$  clock cycles and can range from 0 to 255. Figure 27-9 shows an example using SD24PREx.

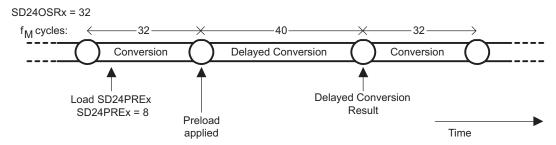


Figure 27-9. Conversion Delay Using Preload - Example

The SD24PREx delay is applied to the beginning of the next conversion cycle after being written. The delay is used on the first conversion after SD24SC is set and on the conversion cycle following each write to SD24PREx. Following conversions are not delayed. After modifying SD24PREx, the next write to SD24PREx should not occur until the next conversion cycle is completed, otherwise the conversion results may be incorrect.

The accuracy of the result for the delayed conversion cycle using SD24PREx is dependent on the length of the delay and the frequency of the analog signal being sampled. For example, when measuring a DC signal, SD24PREx delay has no effect on the conversion result regardless of the duration. The user must determine when the delayed conversion result is useful in their application.

Figure 27-10 shows the operation of grouped channels 0 and 1. The preload register of channel 1 is loaded with zero resulting in immediate conversion whereas the conversion cycle of channel 0 is delayed by setting SD24PRE0 = 8. The first channel 0 conversion uses SD24PREx = 8, shifting all subsequent conversions by eight  $f_M$  clock cycles.

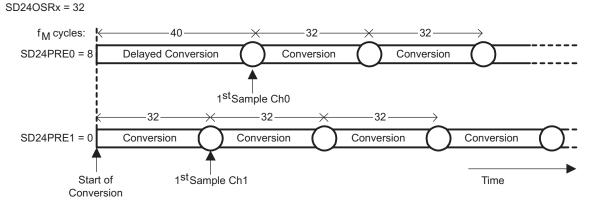


Figure 27-10. Start of Conversion Using Preload - Example

When channels are grouped, care must be taken when a channel or channels operate in single conversion mode or are disabled in software while the master channel remains active. Each time channels in the group are re-enabled and re-synchronize with the master channel, the preload delay for that channel will be reintroduced. Figure 27-11 shows the re-synchronization and preload delays for channels in a group. It is recommended that SD24PREx = 0 for the master channel to maintain a consistent delay between the master and remaining channels in the group when they are re-enabled.



www.ti.com SD24 A Operation

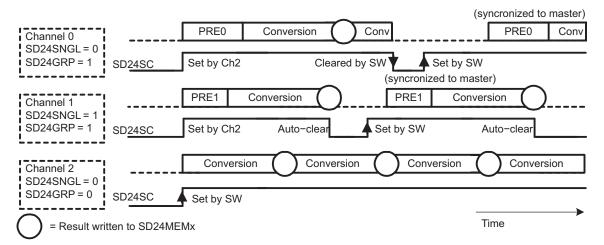


Figure 27-11. Preload and Channel Synchronization

#### 27.2.11 Using the Integrated Temperature Sensor

To use the on-chip temperature sensor, the user selects the analog input pair SD24INCHx = 110 and sets SD24REFON = 1. Any other configuration is done as if an external analog input pair was selected, including SD24INTDLYx and SD24GAINx settings. Because the internal reference must be on to use the temperature sensor, it is not possible to use an external reference for the conversion of the temperature sensor voltage. Also, the internal reference will be in contention with any used external reference. In this case, the SD24VMIDON bit may be set to minimize the affects of the contention on the conversion.

The typical temperature sensor transfer function is shown in Figure 27-12. When switching inputs of an SD24\_A channel to the temperature sensor, adequate delay must be provided using SD24INTDLYx to allow the digital filter to settle and assure that conversion results are valid. The temperature sensor offset error can be large, and may need to be calibrated for most applications. See device-specific data sheet for temperature sensor parameters.

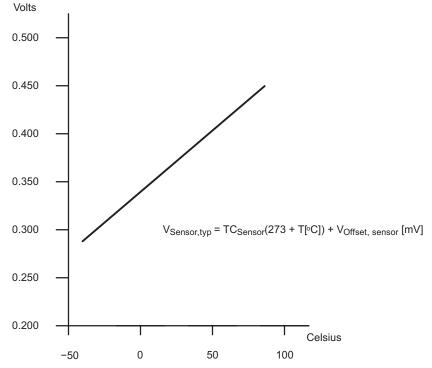


Figure 27-12. Typical Temperature Sensor Transfer Function



SD24 A Operation www.ti.com

#### 27.2.12 Interrupt Handling

The SD24 A has 2 interrupt sources for each ADC channel:

- SD24IFG
- SD24OVIFG

The SD24IFG bits are set when their corresponding SD24MEMx memory register is written with a conversion result. An interrupt request is generated if the corresponding SD24IE bit and the GIE bit are set. The SD24\_A overflow condition occurs when a conversion result is written to any SD24MEMx location before the previous conversion result was read.

#### 27.2.12.1 SD24IV, Interrupt Vector Generator

All SD24\_A interrupt sources are prioritized and combined to source a single interrupt vector. SD24IV is used to determine which enabled SD24\_A interrupt source requested an interrupt. The highest priority SD24\_A interrupt request that is enabled generates a number in the SD24IV register (see register description). This number can be evaluated or added to the program counter to automatically enter the appropriate software routine. Disabled SD24\_A interrupts do not affect the SD24IV value.

Any access, read or write, of the SD24IV register has no effect on the SD24OVIFG or SD24IFG flags. The SD24IFG flags are reset by reading the associated SD24MEMx register or by clearing the flags in software. SD24OVIFG bits can only be reset with software.

If another interrupt is pending after servicing of an interrupt, another interrupt is generated. For example, if the SD24OVIFG and one or more SD24IFG interrupts are pending when the interrupt service routine accesses the SD24IV register, the SD24OVIFG interrupt condition is serviced first and the corresponding flag(s) must be cleared in software. After the RETI instruction of the interrupt service routine is executed, the highest priority SD24IFG pending generates another interrupt request.

#### 27.2.12.2 Interrupt Delay Operation

The SD24INTDLYx bits control the timing for the first interrupt service request for the corresponding channel. This feature delays the interrupt request for a completed conversion by up to four conversion cycles allowing the digital filter to settle prior to generating an interrupt request. The delay is applied each time the SD24SC bit is set or when the SD24GAINx or SD24INCHx bits for the channel are modified. SD24INTDLYx disables overflow interrupt generation for the channel for the selected number of delay cycles. Interrupt requests for the delayed conversions are not generated during the delay.



www.ti.com SD24\_A Operation

### 27.2.12.3 SD24\_A Interrupt Handling Software Example

The following software example shows the recommended use of SD24IV and the handling overhead. The SD24IV value is added to the PC to automatically jump to the appropriate routine.

The numbers at the right margin show the necessary CPU cycles for each instruction. The software overhead for different interrupt sources includes interrupt latency and return-from-interrupt cycles, but not the task handling itself. The latencies are:

- SD24OVIFG, CH0 SD24IFG, CH1 SD24IFG: 16 cycles
- CH2 SD24IFG: 14 cycles

The interrupt handler for channel 2 SD24IFG shows a way to check immediately if a higher prioritized interrupt occurred during the processing of the ISR. This saves nine cycles if another SD24\_A interrupt is pending.

```
; Interrupt handler for SD24_A.
INT_SD24
                       ; Enter Interrupt Service Routine
                                                             6
  ADD
           &SD24IV,PC ; Add offset to PC
                                                             3
  RETI
                      ; Vector 0: No interrupt
                                                             5
  JMP
           ADOV
                       ; Vector 2: ADC overflow
                                                             2
           ADM0
                       ; Vector 4: CH_0 SD24IFG
  JMP
  JMP
           ADM1
                       ; Vector 6: CH_1 SD24IFG
                                                             2
; Handler for CH_2 SD24IFG starts here. No JMP required.
      MOV
            &SD24MEM2,xxx ; Move result, flag is reset
ADM2
                          ; Other instruction needed?
            INT SD24
      JMP
                          ; Check other int pending
                                                             2
;
; Remaining Handlers
            &SD24MEM1,xxx ; Move result, flag is reset
ADM1
      VOM
                           ; Other instruction needed?
      RETI
                           ; Return
                                                             5
ADM0
      VOM
            &SD24MEM0,xxx ; Move result, flag is reset
                                                             5
      RETI
                           ; Return
;
ADOV
                           ; Handle SD24MEMx overflow
      RETI
                                                             5
                            ; Return
```



SD24\_A Registers www.ti.com

# 27.3 SD24\_A Registers

The SD24\_A registers are listed in Table 27-5 (registers for channels not implemented are unavailable; see the device-specific data sheet).

Table 27-5. SD24\_A Registers

Register	Short Form	Register Type	Address	Initial State
SD24_A Control	SD24CTL	Read/write	0100h	Reset with PUC
SD24_A Interrupt Vector	SD24IV	Read/write	0110h	Reset with PUC
SD24_A Analog Enable <sup>(1)</sup>	SD24AE	Read/write	0B7h	Reset with PUC
SD24_A Channel 0 Control	SD24CCTL0	Read/write	0102h	Reset with PUC
SD24_A Channel 0 Conversion Memory	SD24MEM0	Read/write	0112h	Reset with PUC
SD24_A Channel 0 Input Control	SD24INCTL0	Read/write	0B0h	Reset with PUC
SD24_A Channel 0 Preload	SD24PRE0	Read/write	0B8h	Reset with PUC
SD24_A Channel 1 Control	SD24CCTL1	Read/write	0104h	Reset with PUC
SD24_A Channel 1 Conversion Memory	SD24MEM1	Read/write	0114h	Reset with PUC
SD24_A Channel 1 Input Control	SD24INCTL1	Read/write	0B1h	Reset with PUC
SD24_A Channel 1 Preload	SD24PRE1	Read/write	0B9h	Reset with PUC
SD24_A Channel 2 Control	SD24CCTL2	Read/write	0106h	Reset with PUC
SD24_A Channel 2 Conversion Memory	SD24MEM2	Read/write	0116h	Reset with PUC
SD24_A Channel 2 Input Control	SD24INCTL2	Read/write	0B2h	Reset with PUC
SD24_A Channel 2 Preload	SD24PRE2	Read/write	0BAh	Reset with PUC
SD24_A Channel 3 Control	SD24CCTL3	Read/write	0108h	Reset with PUC
SD24_A Channel 3 Conversion Memory	SD24MEM3	Read/write	0118h	Reset with PUC
SD24_A Channel 3 Input Control	SD24INCTL3	Read/write	0B3h	Reset with PUC
SD24_A Channel 3 Preload	SD24PRE3	Read/write	0BBh	Reset with PUC
SD24_A Channel 4 Control	SD24CCTL4	Read/write	010Ah	Reset with PUC
SD24_A Channel 4 Conversion Memory	SD24MEM4	Read/write	011Ah	Reset with PUC
SD24_A Channel 4 Input Control	SD24INCTL4	Read/write	0B4h	Reset with PUC
SD24_A Channel 4 Preload	SD24PRE4	Read/write	0BCh	Reset with PUC
SD24_A Channel 5 Control	SD24CCTL5	Read/write	010Ch	Reset with PUC
SD24_A Channel 5 Conversion Memory	SD24MEM5	Read/write	011Ch	Reset with PUC
SD24_A Channel 5 Input Control	SD24INCTL5	Read/write	0B5h	Reset with PUC
SD24_A Channel 5 Preload	SD24PRE5	Read/write	0BDh	Reset with PUC
SD24_A Channel 6 Control	SD24CCTL6	Read/write	010Eh	Reset with PUC
SD24_A Channel 6 Conversion Memory	SD24MEM6	Read/write	011Eh	Reset with PUC
SD24_A Channel 6 Input Control	SD24INCTL6	Read/write	0B6h	Reset with PUC
SD24_A Channel 6 Preload	SD24PRE6	Read/write	0BEh	Reset with PUC

<sup>(1)</sup> Not implemented on all devices; see the device-specific data sheet.



www.ti.com SD24\_A Registers

# 27.3.1 SD24CTL, SD24\_A Control Register

15	14		13	12	11	10	9	8
		Reserv	/ed			SD24XDIVx		SD24LP
r0	r0		r0	r0	rw-0	rw-0	rw-0	rw-0
7	6		5	4	3	2	1	0
SD24	SD24DIVx SD24SSELx		SD24VMIDON	SD24REFON	SD24OVIE	Reserved		
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	r0
Reserved	Bits 15-12	Rese	rved					
SD24XDIVx	Bits 11-9	SD24	L_A clock divide	er				
		00	/1					
		01	/3					
		10	/16					
		11	/48					
		1xx	Reserved					
SD24LP	Bit 8	Low-	power mode. T	nis bit selects a re	duced-speed redu	ced-power mode		
		0	Low-power	mode is disabled				
		1	Low-power	mode is enabled.	The maximum clod	ck frequency for th	e SD24_A is redu	ced.
SD24DIVx	Bits 7-6	SD24	LA clock divide	er				
		00	/1					
		01	/2					
		10	/4					
		11	/8					
SD24SSELx	Bits 5-4	SD24	LA clock source	e select				
		00	MCLK					
		01	SMCLK					
		10	ACLK					
		11	External TA	CLK				
SD24VMIDON	Bit 3	VMID	buffer on					
		0	Off					
		1	On					
SD24REFON	Bit 2	Refer	rence generator	on				
		0	Reference of	off				
		1	Reference of	on				
SD24OVIE	Bit 1	SD24	LA overflow int	errupt enable. The	e GIE bit must also	be set to enable	the interrupt.	
		0		errupt disabled			-	
		1	Overflow int	errupt enabled				
Reserved	Bit 0	Rese	rved					



SD24\_A Registers www.ti.com

# 27.3.2 SD24CCTLx, SD24\_A Channel x Control Register

15	14		13	12	11	10	9	8
Reserved	S	D24BU	Fx <sup>(1)</sup>	SD24UNI	SD24XOSR	SD24SNGL	SD24	OSRx
r0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
7	6		5	4	3	2	1	0
SD24LSBTOG	SD24LSBA	CC	SD24OVIFG	SD24DF	SD24IE	SD24IFG	SD24SC	SD24GRP
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	r(w)-0
Reserved	Bit 15	Reser	rved					
SD24BUFx	Bits 14-13	High-i	impedance inpu	t buffer mode				
		00	Buffer disabl	ed				
		01	Slow speed/o	current				
		10	Medium spe	ed/current				
		11	High speed/o	current				
SD24UNI	Bit 12	Unipo	lar mode select					
		0	Bipolar mode	•				
		1	Unipolar mod	de				
SD24XOSR	Bit 11			ng ratio. This bit, ption for settings.		24OSRx bits, selec	ct the oversampli	ng ratio. See
SD24SNGL	Bit 10	Single	e conversion mo	de select				
		0	Continuous of	conversion mode				
		1	Single conve	rsion mode				
SD24OSRx	Bits 9-8	Overs	sampling ratio					
		When	SD24XOSR =	0				
		00	256					
		01	128					
		10	64					
		11	32					
			SD24XOSR =	1				
		00	512					
		01	1024					
		10	Reserved					
CD241 CDTOC	D:+ 7	11	Reserved	when oot oou	CD241 CDACC to	togalo coch timo t	ha CD24MEMy r	agistar is road
SD24LSBTOG	Bit 7	0			e with each SD24N	toggle each time t	THE SDZ4IVIEIVIX I	egister is read.
		1			ach SD24MEMx re			
SD24LSBACC	Bit 6					· 16-bits of the SD2	24 A conversion	result
ODZ-IZOBAGO	Dit 0	0				s of the conversion		roodit.
		1			-	s of the conversion		
SD24OVIFG	Bit 5		A overflow inte					
		0	_	interrupt pending				
		1		errupt pending				
SD24DF	Bit 4	SD24	_A data format					
		0	Offset binary					
		1	2s compleme					
SD24IE	Bit 3	SD24	_A interrupt ena					
		0	Disabled					
		1	Enabled					

Not implemented on all devices (see the device-specific data sheet). Reserved with r0 access if high-impedance buffer not implemented.



SD24\_A Registers www.ti.com

SD24IFG	Bit 2	SD24_A interrupt flag. SD24IFG is set when new conversion results are available. SD24IFG is automatically reset when the corresponding SD24MEMx register is read, or may be cleared with software.
		0 No interrupt pending
		1 Interrupt pending
SD24SC	Bit 1	SD24_A start conversion
		0 No conversion start
		1 Start conversion
SD24GRP	Bit 0	SD24_A group. Groups SD24_A channel with next higher channel. Not used for the last channel.
		0 Not grouped
		1 Grouped

# 27.3.3 SD24INCTLx, SD24\_A Channel x Input Control Register

7	6		5	4	3	2	1	0
SD24IN	NTDLYx			SD24GAINx			SD24INCHx	
rw-0	rw-0		rw-0	rw-0	rw-0	rw-0	rw-0	rw-0
SD24INTDLYx	Bits 7-6		ıpt delay gener rsion start.	ation after conversi	on start. These I	bits select the dela	y for the first inter	rupt after
		00	Fourth samp	le causes interrupt				
		01	Third sample	e causes interrupt				
		10	Second sam	ple causes interrup	ot			
		11	First sample	causes interrupt				
SD24GAINx	Bits 5-3	SD24	_A preamplifier	gain				
		000	x1					
		001	x2					
		010	x4					
		011	x8					
		100	x16					
		101	x32					
		110	Reserved					
		111	Reserved					
SD24INCHx	Bits 2-0		_A channel diffiction diffiction data sheet.	erential pair input.	The available sel	lections are device	dependent. See	the device-
		000	Ax.0					
		001	Ax.1 <sup>(1)</sup>					
		010	Ax.2 <sup>(1)</sup>					
		011	Ax.3 <sup>(1)</sup>					
		100	Ax.4 <sup>(1)</sup>					
		101	(AV <sub>CC</sub> - AV <sub>SS</sub>	s) / 11				
		110	Temperature	e sensor				
		111	Short for PG	A offset measurem	ent			

 $<sup>\,^{(1)}\,</sup>$  Ax.1 to Ax.4 not available on all devices (see device-specific data sheet).



SD24\_A Registers www.ti.com

# 27.3.4 SD24MEMx, SD24\_A Channel x Conversion Memory Register

15	14	13	12	11	10	9	8		
	Conversion Results								
r	r	r	r	r	r	r	r		
7	6	5	4	3	2	1	0		
	Conversion Results								
r	r	r	r	r	r	r	r		

Conversion Results

Bits 15-0

Conversion results. The SD24MEMx register holds the upper or lower 16-bits of the digital filter output, depending on the SD24LSBACC bit.

# 27.3.5 SD24PREx, SD24\_A Channel x Preload Register

7	6	5	4	3	2	1	0
			Preloa	d Value			
rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0	rw-0

Preload Value Bits 7-0 SD24\_A digital filter preload value

# 27.3.6 SD24AE, SD24\_A Analog Input Enable Register

7	6	5	4	3	2	1	0
SD24AE7	SD24AE6	SD24AE5	SD24AE4	SD24AE3	SD24AE2	SD24AE1	SD24AE0
rw-0							

**SD24AEx** Bits 7-0 SD24\_A analog enable

0 External input disabled. Negative inputs are internally connected to VSS.

External input enabled



SD24\_A Registers www.ti.com

# 27.3.7 SD24IV, SD24\_A Interrupt Vector Register

15	14	13	12	11	10	9	8
0	0	0	0	0	0	0	0
rO	r0	r0	rO	r0	rO	rO	r0
7	6	5	4	3	2	1	0
0	0	0	SD24IVx				0
rO	r∩	rO	r-O	r-O	r-O	r-O	rΩ

SD24IVx Bits 15-0 SD24\_A interrupt vector value

SD24IV Contents	Interrupt Source	Interrupt Flag	Interrupt Priority
000h	No interrupt pending	-	
002h	SD24MEMx overflow	SD24CCTLx SD24OVIFG <sup>(1)</sup>	Highest
004h	SD24_A Channel 0 Interrupt	SD24CCTL0 SD24IFG	
006h	SD24_A Channel 1 Interrupt	SD24CCTL1 SD24IFG	
008h	SD24_A Channel 2 Interrupt	SD24CCTL2 SD24IFG	
00Ah	SD24_A Channel 3 Interrupt	SD24CCTL3 SD24IFG	
00Ch	SD24_A Channel 4 Interrupt	SD24CCTL4 SD24IFG	
00Eh	SD24_A Channel 5 Interrupt	SD24CCTL5 SD24IFG	
010h	SD24_A Channel 6 Interrupt	SD24CCTL6 SD24IFG	Lowest

When an SD24\_A overflow occurs, the user must check all SD24CCTLx SD24OVIFG flags to determine which channel overflowed.



# Embedded Emulation Module (EEM)

This chapter describes the Embedded Emulation Module (EEM) that is implemented in all MSP430 flash devices.

Topic Page

28.1	EEM Introduction	639
28.2	EEM Building Blocks	641
28.3	EEM Configurations	642
	· · · · · · · · · · · · · · · · · · ·	



www.ti.com EEM Introduction

#### 28.1 EEM Introduction

Every MSP430 flash-based microcontroller implements an embedded emulation module (EEM). It is accessed and controlled through JTAG. Each implementation is device dependent and is described in section 1.3 *EEM Configurations* and the device-specific data sheet.

In general, the following features are available:

- Non-intrusive code execution with real-time breakpoint control
- Single step, step into and step over functionality
- Full support of all low-power modes
- Support for all system frequencies, for all clock sources
- Up to eight (device dependent) hardware triggers/breakpoints on memory address bus (MAB) or memory data bus (MDB)
- Up to two (device dependent) hardware triggers/breakpoints on CPU register write accesses
- MAB, MDB ,and CPU register access triggers can be combined to form up to eight (device dependent) complex triggers/breakpoints
- Trigger sequencing (device dependent)
- · Storage of internal bus and control signals using an integrated trace buffer (device dependent)
- Clock control for timers, communication peripherals, and other modules on a global device level or on a per-module basis during an emulation stop

Figure 28-1 shows a simplified block diagram of the largest currently available 2xx EEM implementation.

For more details on how the features of the EEM can be used together with the IAR Embedded Workbench™ debugger see the application report *Advanced Debugging Using the Enhanced Emulation Module* (SLAA263) at www.msp430.com. Code Composer Essentials (CCE) and most other debuggers supporting MSP430 have the same or a similar feature set. For details see the user's guide of the applicable debugger.



EEM Introduction www.ti.com

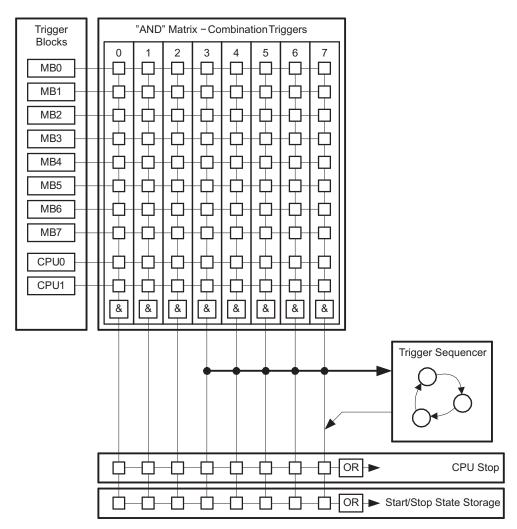


Figure 28-1. Large Implementation of the Embedded Emulation Module (EEM)



www.ti.com EEM Building Blocks

## 28.2 EEM Building Blocks

## 28.2.1 Triggers

The event control in the EEM of the MSP430 system consists of triggers, which are internal signals indicating that a certain event has happened. These triggers may be used as simple breakpoints, but it is also possible to combine two or more triggers to allow detection of complex events and trigger various reactions besides stopping the CPU.

In general, the triggers can be used to control the following functional blocks of the EEM:

- Breakpoints (CPU stop)
- State storage
- Sequencer

There are two different types of triggers, the memory trigger and the CPU register write trigger.

Each memory trigger block can be independently selected to compare either the MAB or the MDB with a given value. Depending on the implemented EEM the comparison can be =,  $\neq$ ,  $\geq$ , or  $\leq$ . The comparison can also be limited to certain bits with the use of a mask. The mask is either bit-wise or byte-wise, depending upon the device. In addition to selecting the bus and the comparison, the condition under which the trigger is active can be selected. The conditions include read access, write access, DMA access, and instruction fetch.

Each CPU register write trigger block can be independently selected to compare what is written into a selected register with a given value. The observed register can be selected for each trigger independently. The comparison can be =,  $\neq$ , or  $\leq$ . The comparison can also be limited to certain bits with the use of a bit mask.

Both types of triggers can be combined to form more complex triggers. For example, a complex trigger can signal when a particular value is written into a user-specified address.

# 28.2.2 Trigger Sequencer

The trigger sequencer allows the definition of a certain sequence of trigger signals before an event is accepted for a break or state storage event. Within the trigger sequencer, it is possible to use the following features:

- Four states (State 0 to State 3)
- Two transitions per state to any other state
- Reset trigger that resets the sequencer to State 0.

The Trigger sequencer always starts at State 0 and must execute to State 3 to generate an action. If State 1 or State 2 are not required, they can be bypassed.

### 28.2.3 State Storage (Internal Trace Buffer)

The state storage function uses a built-in buffer to store MAB, MDB, and CPU control signal information (that is, read, write, or instruction fetch) in a non-intrusive manner. The built-in buffer can hold up to eight entries. The flexible configuration allows the user to record the information of interest very efficiently.

#### 28.2.4 Clock Control

The EEM provides device dependent flexible clock control. This is useful in applications where a running clock is needed for peripherals after the CPU is stopped (for example, to allow a UART module to complete its transfer of a character or to allow a timer to continue generating a PWM signal).

The clock control is flexible and supports both modules that need a running clock and modules that must be stopped when the CPU is stopped due to a breakpoint.



EEM Configurations www.ti.com

### 28.3 EEM Configurations

Table 28-1 gives an overview of the EEM configurations in the MSP430 2xx family. The implemented configuration is device dependent - see the device data sheet.

Table 28-1. 2xx EEM Configurations

Feature	xs	S	М	L
Memory Bus Triggers	2(=, ≠ only)	3	5	8
Memory Bus Trigger Mask for	1) Low byte	1) Low byte	1) Low byte	All 16 or 20 bits
	2) High byte	2) High byte	2) High byte	
CPU Register-Write Triggers	0	1	1	2
Combination Triggers	2	4	6	8
Sequencer	No	No	Yes	Yes
State Storage	No	No	No	Yes

In general the following features can be found on any 2xx device:

- · At least two MAB/MDB triggers supporting:
  - Distinction between CPU, DMA, read, and write accesses
  - =,  $\neq$ ,  $\geq$ , or  $\leq$  comparison (in XS only =,  $\neq$ )
- At least two trigger Combination registers
- Hardware breakpoints using the CPU Stop reaction
- Clock control with individual control of module clocks (in some XS configurations the control of module clocks is hardwired)



www.ti.com Revision History

# **Revision History**

Revision	Comments
SLAU144G	Chapter 5 Basic Clock Module+, Added information specific to the MSP430AFE2xx devices:
	Figure 5-2. Basic Clock Module+ Block Diagram - MSP430AFE2xx
	Section 5.3, Register BCSCTL3 default
	Section 5.3.2, 5.3.3, 5.3.4, Available register bits, defaults, and definitions
	Added chapters:
	Chapter 18 USART Peripheral Interface, UART Mode
	Chapter 19 USART Peripheral Interface, SPI Mode
	Chapter 27 SD24_A
	Made editorial and format changes throughout.
SLAU144H	Section 2.4, Corrected DCO startup time.
	Section 8.2.6, Updated pin oscillator information; added Figure 8-1.
	Section 3.4.6.5, Corrected typo in BIC description.
	Section 7.2.1, Corrected typo in code example.
SLAU144I	Table 2-3, Changed comments on crystal pins.
	Section 1.4.1, Corrected addresses for end of Flash/ROM.
	Section 3.3.5, Changed example figure.
	Updated descriptions in the following sections: Section 5.1, Section 5.2.1, Section 5.2.2, Section 5.2.3, Section 5.2.5.2, Section 5.2.7.1, Section 5.3.3 (DCOR bit), Section 5.3.4 (FLST1Sx bit).
	Section 7.3.2 and Section 7.3.4, Added information regarding MSP430G2xx.
	Section 8.1, Added note regarding MSP430G22x0.
	Chapter 21, Added notes throughout regarding MSP430G2210.
	Figure 22-1, Updated block diagram.
	Section 22.2.2.1, Changed Analog Port Selection description.
	Section 22.2.3, Changed Voltage Reference Generator description.
	Section 22.3.1, Updated SREF bit description.
	Section 22.3.2, Updated INCHx bit description.
	Figure 23-1, Changed four inputs on center left mux from GND to Floating.
	Table 24-1, Corrected CALDCO names.  Made editorial changes throughout.
SLAU144J	
SLAU 1443	Figure 3-17, Corrected bottom left bit number.
	Section 7.2, Corrected minimum number of main memory segments.
	Section 24.2.2.1, Added temperature sensor calibration equations.
	Section 26.2.5, Changed description.
	Section 27.2.5, Changed description.

NOTE: Page numbers for previous revisions may differ from page numbers in the current version.

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