# 2. INTRODUCTIONTO DISTRIBUTED SYSTEMS

Network Distributed Systems

### ROAD MAP

- Fundamentals of DSs
- Basic Communication Protocols
  - ARQs
  - Medium Access Control

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# WHAT IS A DISTRIBUTED SYSTEM?

A collection of independent computers that appear to the users of the system as a single computer.

Tanenbaum

### EXAMPLES OF DSS

**Cluster:** "A type of parallel or distributed processing system, which consists of a collection of interconnected stand-alone computers cooperatively working together as a single, integrated computing resource" [Buyya].

**Grid:** "A type of parallel and distributed system that enables the sharing, selection, and aggregation of geographically distributed autonomous resources dynamically at runtime depending on their availability, capability, performance, cost, and users' quality-of-service requirements" [Buyya].

### NETWORKS VS DS

**Networks:** A media for interconnecting local and wide area computers and exchange messages based on protocols. Network entities are visible and they are explicitly addressed (IP address).

Distributed System: existence of multiple autonomous computers is transparent

However, many problems (e.g., openness, reliability) in common, but at different levels.

Networks focuses on packets, routing, etc., whereas distributed systems focus on applications.

Every distributed system relies on services provided by a computer network.

### WHY DO WE NEED DS?

### **Functional Separation:**

- Existence of computers with different capabilities and purposes:
  - Clients and Servers
  - Data collection and data processing

### Inherent distribution:

- Information:
  - Different information is created and maintained by different people (e.g., Web pages)
- People
  - Computer supported collaborative work (virtual teams, engineering, virtual surgery)
- Retail store and inventory systems for supermarket chains (e.g., tesco, dunnes)

### WHY DO WE NEED DS?

### Power imbalance and load variation:

Distribute computational load among different computers.

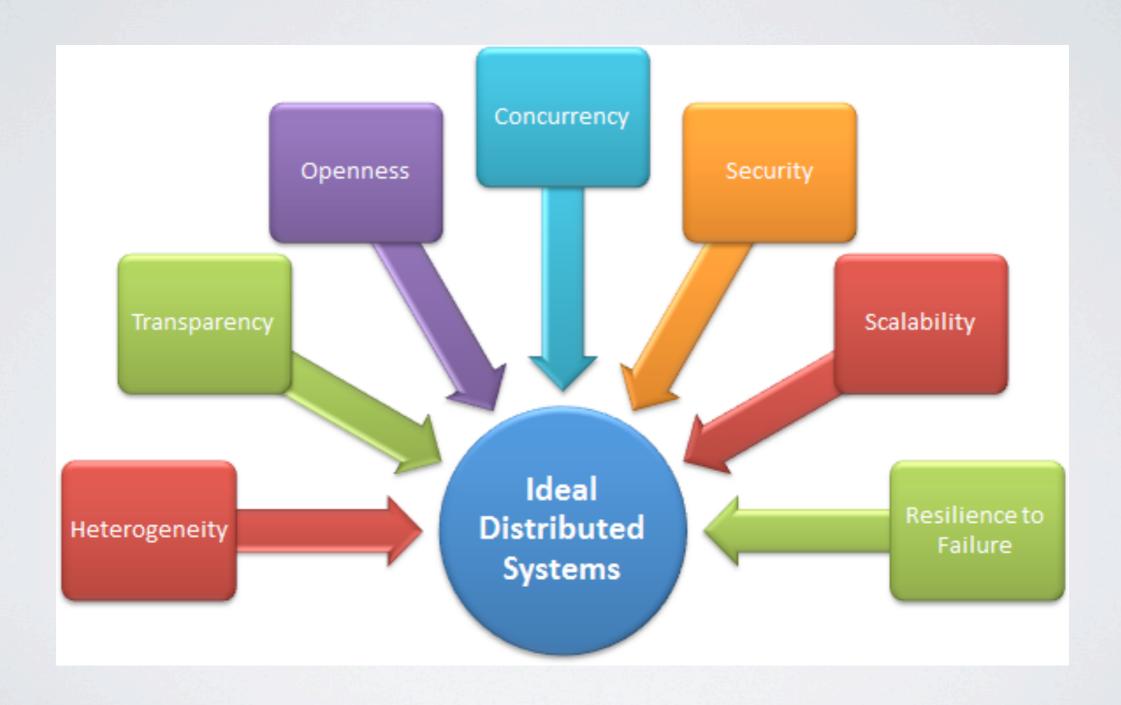
### Reliability:

Long term preservation and data backup (replication) at different locations.

### • Economies:

Sharing a printer by many users and reduce the cost of ownership. Building a supercomputer out of a network of computers.

# CHARACTERISTICS OF DS



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### RELIABLETRANSMISSION

### Why?

Frame can be corrupted and lost

CRCs are not enough - CRCs don't correct errors

### Two fundamental mechanisms

Acknowledgment

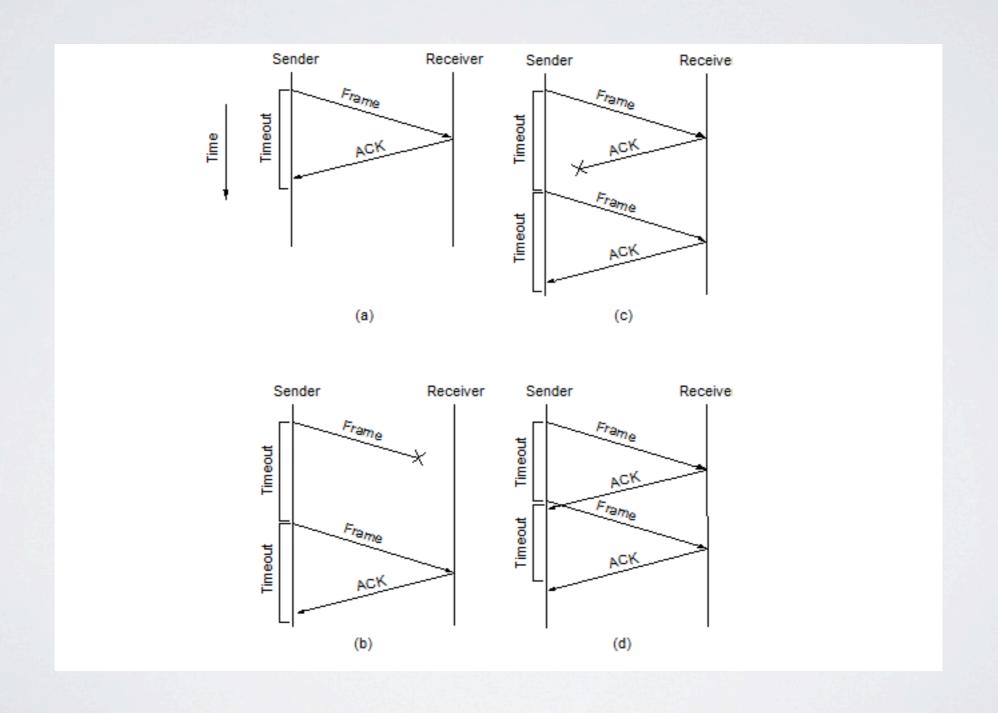
Timeout

General idea is called ARQ (Automatic Repeat Request)

# AUTOMATIC REPEAT REQUEST PROTOCOLS

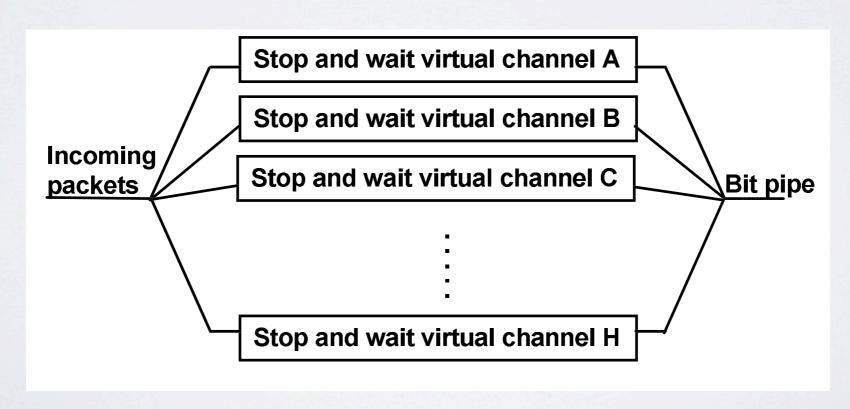
- Causes of message loss
  - Buffer overflow
  - Error detection in a pkt
- ARQ Protocols
  - Stop & Wait
  - Arpanet
  - Go back n
  - Selective Repeat

# STOP & WAIT ARQ

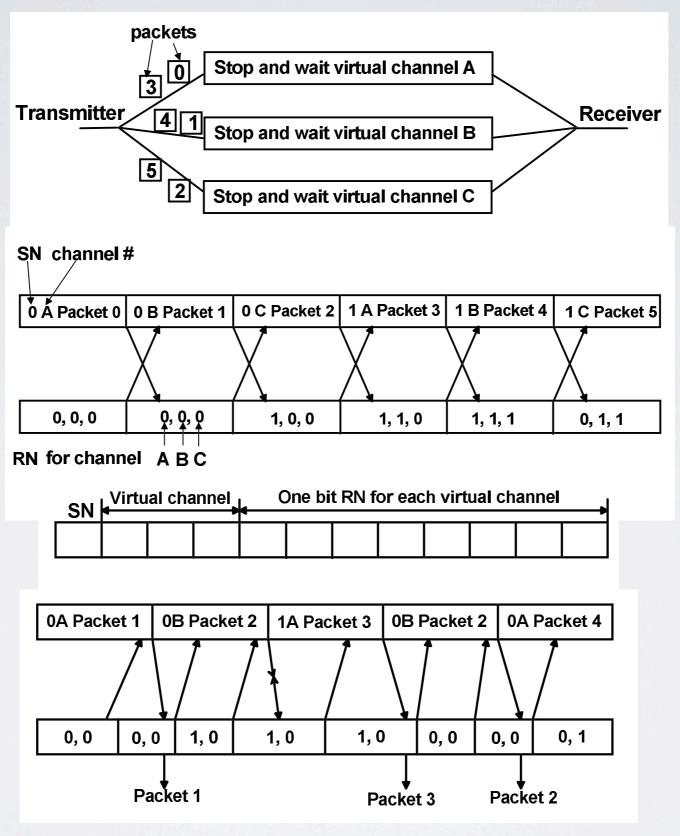


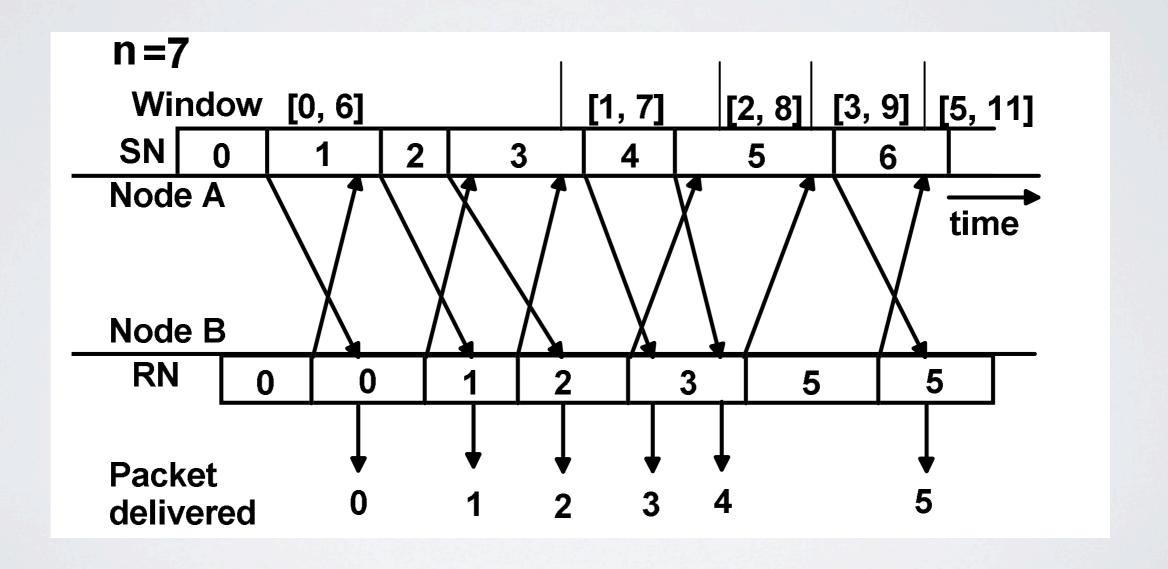
# ARPANET ARQ

- Better line utilisation
- Unlimited memory required in theory
- Packets are assigned to one of 8 virtual channels which are served in a round robin order
- If the virtual channel's turn comes before an ACK for the virtual channel is received, the packet is retransmitted

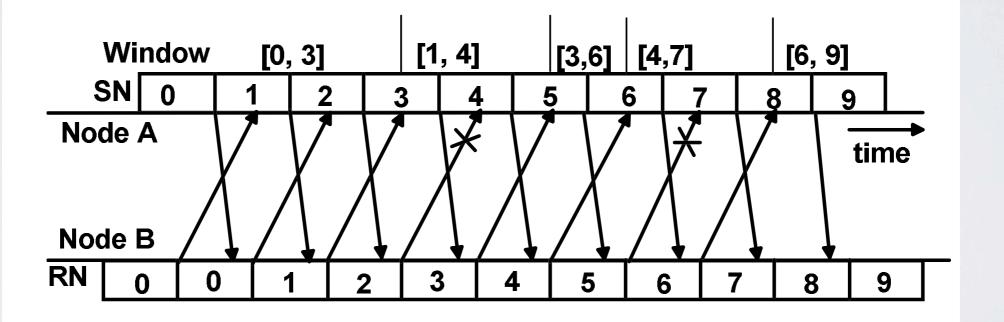


### ARPANET

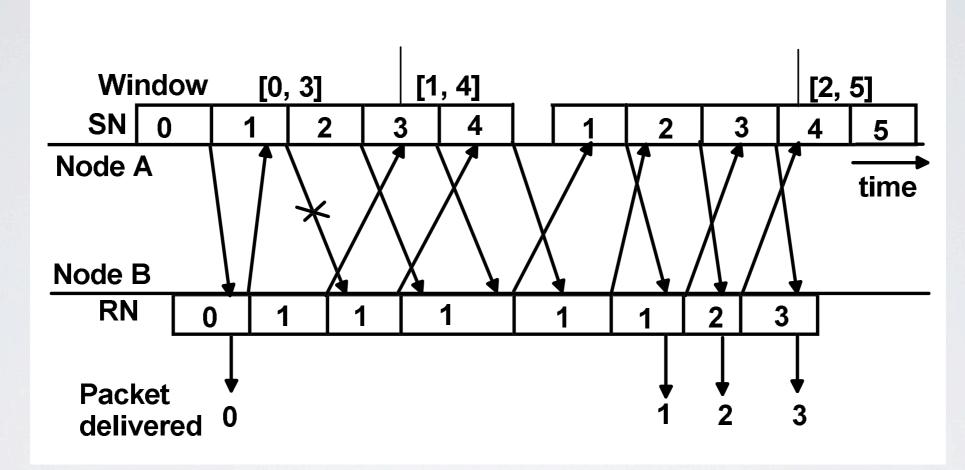






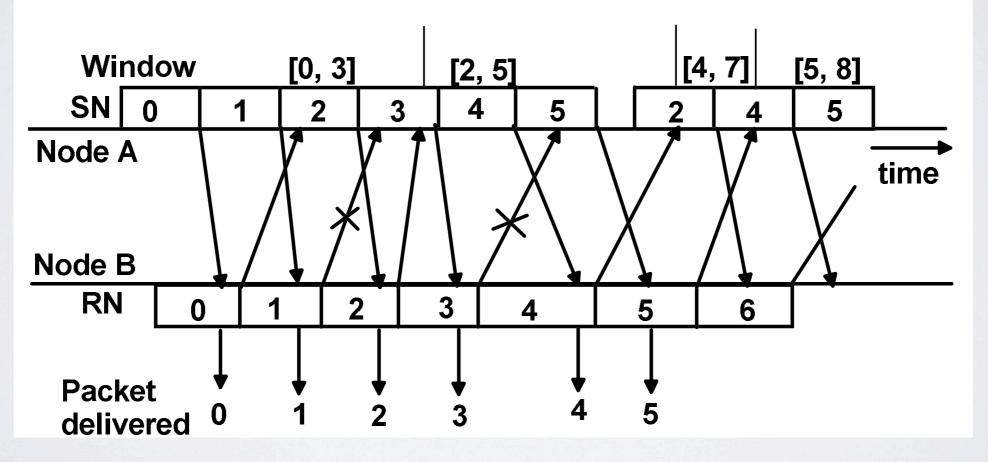


**Example 2**: (retransmissions because of errors for go back 4)

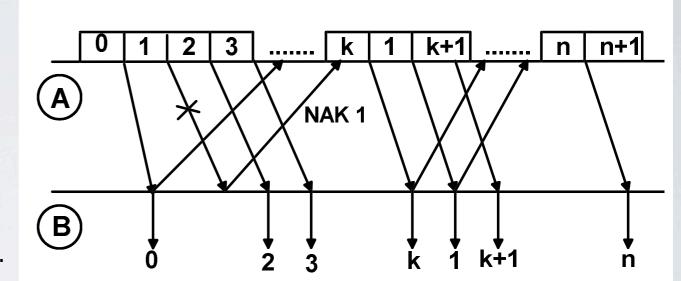


- Good utilisation
- · limited memory required (one pkt)
- Full window is retransmitted in case of one error

Example 3: (retransmissions due to feedback errors for go back n)



# SELECTIVE REPEAT ARQ



Sliding window technique (as Go back n).

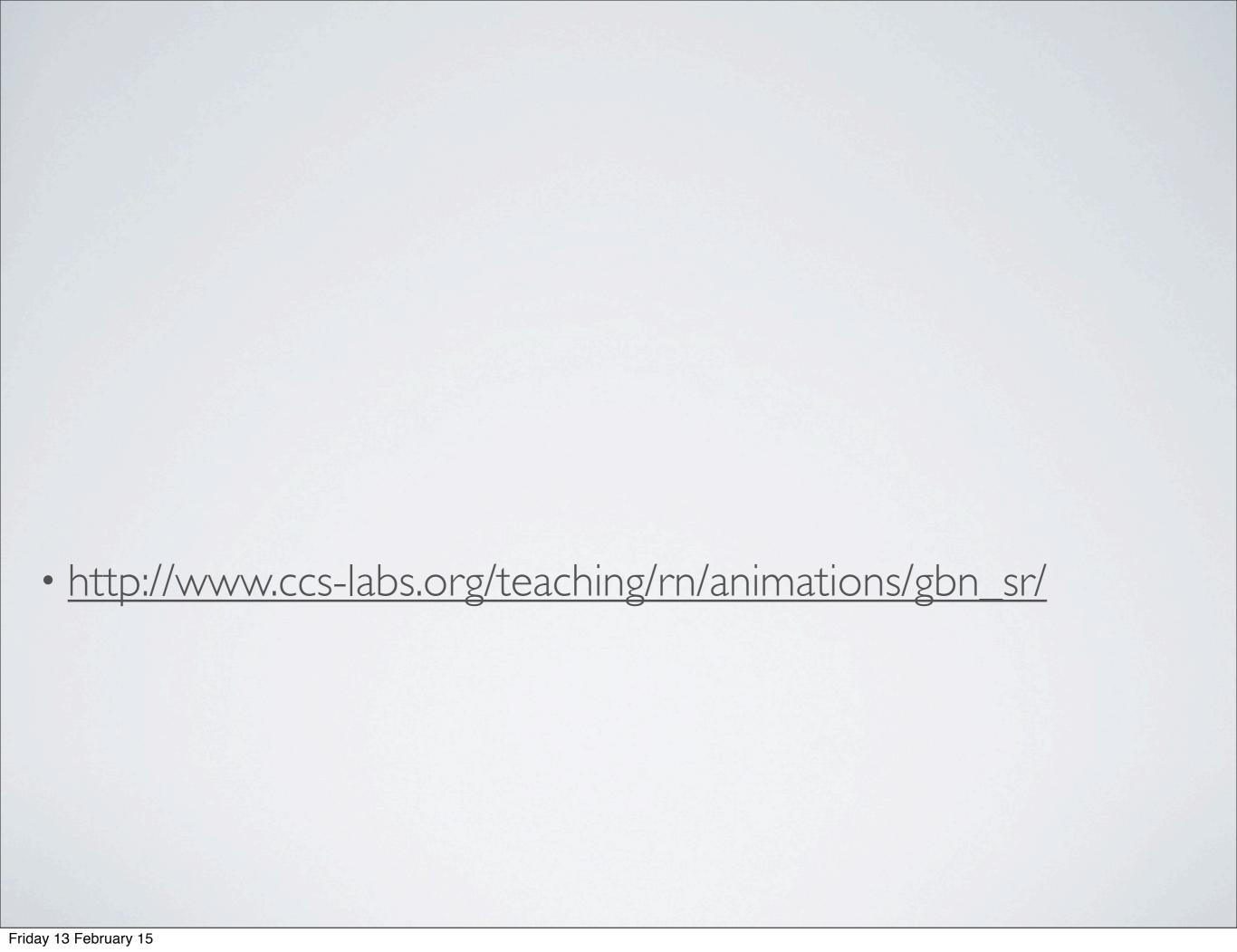
Specifically indicating which packet is missing.

Combines nacks and cumulative acks.

Acks acknowledge all messages with index of up to and including the ack value.

Nacks (negative acknowledgements) specifically request the messages with the indices in the nacks' values.

Limited memory required (a full window).



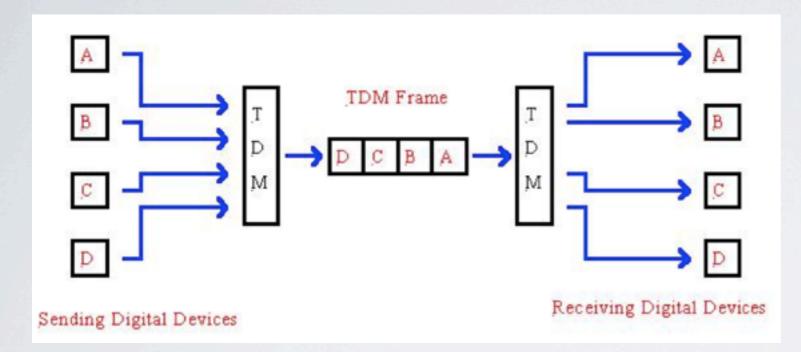
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# MAC FOR MULTI ACCESS

### TIME DIVISION MULTIPLEXING

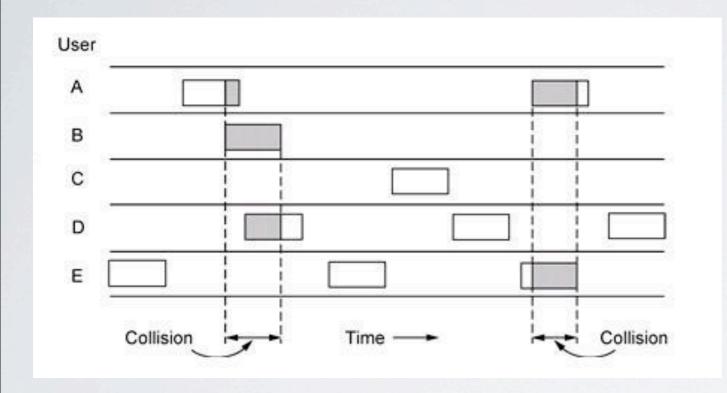
Receiving Digital Devices



- The best utilisation if everyone always has something to send
- $\begin{array}{c|c} Idle & A \\ \hline B \\ \hline Idle & C \\ \hline D \\ \hline D \\ \hline \end{array} \begin{array}{c} TDM \ Frame \\ \hline D \\ \hline M \\ \hline \end{array} \begin{array}{c} T \\ \hline D \\ \hline M \\ \hline \end{array} \begin{array}{c} T \\ \hline D \\ \hline M \\ \hline \end{array} \begin{array}{c} T \\ \hline D \\ \hline \end{array} \begin{array}{c} T \\ \hline \end{array} \begin{array}{c} T \\ \hline D \\ \hline \end{array} \begin{array}{c} T \\ \end{array} \begin{array}{c} T \\$
- Wastes time otherwise
- Slots can be unevenly assigned to improve efficiency

Sending Digital Devices

### ALOHA

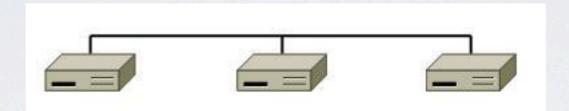


- Send packet immediately
- If collision occurs, back off for a random time, try again
- Simple, works well in load load.
- Improvement: slotted aloha

### CSMA

- · Improve ALOHA by listening to the medium before we send
  - Easier with wired medium
  - Does this completely eliminate collisions?

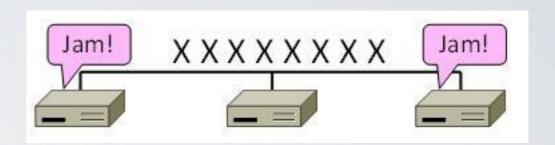
### CSMA

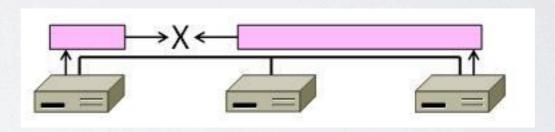


- Propagation delay and simultaneous transmissions can still cause collisions
- · Is quite effective when delay and contention is small

### CSMA CD

- CSMA with Collision
   Detection
- Reduces the cost of a collision by detecting it an aborting the frame transmission
- Sends jamming signal to inform everyone that a collision occurred
- Uses minimum frame length (64 bytes) and persistence







# CSMA CD BINARY EXPONENTIAL BACKOFF

- Estimates the probability
  - Ist collision wait 0 or I frame time
  - 2nd collision wait from 0 to 3 times
  - 3rd collision wait from 0 to 7 times
  - and so on
- BEB doubles interval for each successive collision
  - quickly gets large enough to work
  - very efficient in practice

HTTPS://
WWW.YOUTUBE.COM/
WATCH?V=\_AJV5AQR68