
Protocol $\text{COT}_{\mathbb{G}, \xi, \ell}$

Correlated OT protocol from any ROT protocol and a uniform mapper $H: \mathbb{Z}_2^\kappa \mapsto \mathbb{G}$ (e.g., Hash2Field [FHSS⁺23] for EC), parametrized by a group \mathbb{G}

Players: a sender \mathcal{S} , and receiver \mathcal{R}

Inputs: $\mathcal{S}: \mathbf{a} \in \mathbb{G}^{\xi \times \ell}$, group elements

$\mathcal{R}: \mathbf{x} \in \mathbb{Z}_2^\xi$, the choice bits

Outputs: $\mathcal{S} \leftarrow \mathbf{z}_A \in \mathbb{G}^{\xi \times \ell}$

$\mathcal{R} \leftarrow \mathbf{z}_B \in \mathbb{G}^{\xi \times \ell}$ s.t. $z_{A(i,j)} + z_{B(i,j)} = a_{(i,j)} \cdot x_{(i)} \quad \forall i \in [\xi] \quad \forall j \in [\ell]$

$\mathcal{S} \& \mathcal{R}. \text{RunROT}(b) \dashrightarrow \mathcal{S}: (\mathbf{r}_0, \mathbf{r}_1); \mathcal{R}: \mathbf{r}_b$

1: \mathcal{S} runs $\text{ROT}_{\xi, \ell}$ as sender, obtaining $\mathbf{r}_0, \mathbf{r}_1 \in \mathbb{Z}_2^{\xi \times \ell \times \kappa}$

2: \mathcal{R} runs $\text{ROT}_{\xi, \ell}$ as receiver, obtaining $\mathbf{r}_b \in \mathbb{Z}_2^{\xi \times \ell \times \kappa}$

$\mathcal{S}. \text{CreateCorrelation}(\mathbf{r}_0, \mathbf{r}_1, \mathbf{a}) \dashrightarrow (\boldsymbol{\tau})$

1: $\mathbf{z}_A \leftarrow \{ \{ H(r_{0(i,j)}) \}_{j \in [\ell]} \}_{i \in [\xi]}$

2: $\boldsymbol{\tau} \leftarrow \{ \{ H(r_{1(i,j)}) - z_{A(i,j)} + a_{(i,j)} \}_{j \in [\ell]} \}_{i \in [\xi]}$

3: $\text{Send}(\boldsymbol{\tau}) \rightarrow \mathcal{R}$

return \mathbf{z}_A

$\mathcal{R}. \text{ApplyCorrelation}(\boldsymbol{\tau}) \dashrightarrow \mathbf{z}_B$

return $\mathbf{z}_B \leftarrow \{ \{ \tau_{(i,j)} \cdot b_{(i)} - H(r_{b(i,j)}) \}_{j \in [\ell]} \}_{i \in [\xi]}$

References

- [FHSS⁺23] Armando Faz-Hernandez, Sam Scott, Nick Sullivan, Riad S. Wahby, and Christopher A. Wood. Hashing to Elliptic Curves. RFC 9380, August 2023.