
Protocol Joint – Feldman DKG

Maliciously secure threshold DKG protocol from [GJKR07] for a group $\mathbb{G}(q, G)$, using t -out-of- n Pedersen VSS and a ZKPoK of the discrete log (e.g., Fischlin)

Players: $\mathcal{P}_1, \dots, \mathcal{P}_i, \dots, \mathcal{P}_n$, a set of n share holders.

Inputs: sid , a unique session identifier (e.g., obtained from Protocol ??)

Outputs: A public key Y and n secret shares x_i of the private key x .

$\mathcal{P}_i.\text{Round1}() \dashrightarrow (x_i, x'_i, B_i)$

- 1: Sample $a_{i,0} \xleftarrow{\$} \mathbb{Z}_q$
- 2: Run $(x_i, x'_i, C_i, B_i) \leftarrow \text{Pedersen.Split}(a_{i,0})$ as shares x_i , blinding shares x'_i , public key shares C_i and blinded public key shares B_i
- 3: Run $\pi_i \leftarrow \{\text{Fischlin.Prove}(s)\} \forall s \in \{a_{i,0}, x_{(i,1)}, \dots, x_{(i,n)}\}$
- 4: $\text{Send}(x_{(i,j)}, x'_{(i,j)}) \rightarrow \mathcal{P}_j \quad \forall j \in [n]$
- 5: $\mathcal{F}^{\text{Broadcast}}(B_i)$

$\mathcal{P}_i.\text{Round2}(\{B_j\}_{j \in [n]}, \{x_{(j,i)}, x'_{(j,i)}\}_{j \in [n]}) \dashrightarrow (C_i, \pi)_i$

- 1: Check $\text{Pedersen.Verify}(j, x_{(j,i)}, x'_{(j,i)}, B_j) \quad \forall j \in [n]$; **ABORT** if it fails
- 2: $x_i \leftarrow \sum_{j \in [n]} x_{(j,i)}$ as the private key share of \mathcal{P}_i
- 3: $\mathcal{F}^{\text{Broadcast}}(C_i, \pi_i)$

$\mathcal{P}_i.\text{Round3}(\{C_j\}_{j \in [n]}, \{\pi_j\}_{j \in [n]}) \dashrightarrow (x_i, Y)$

- 1: Run $\text{Fischlin.Verify}(j, \pi_j) \quad \forall j \in [n]$; **ABORT** if it fails
 - 2: Run $\text{Feldman.Verify}(j, x_{(j,i)}, C_j) \quad \forall j \in [n]$; **ABORT** if it fails
 - 3: $Y \leftarrow \sum_{j=1}^n C_{(j,0)}$ as the public key
 return (x_i, Y)
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References

- [GJKR07] Rosario Gennaro, Stanislaw Jarecki, Hugo Krawczyk, and Tal Rabin. Secure distributed key generation for discrete-log based cryptosystems. *Journal of Cryptology*, 20:51–83, 2007.