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**Protocol**  $\text{COT}_{\mathbb{G}, \xi, \ell}$ 

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Correlated OT protocol from any ROT protocol and a uniform mapper  $H: \mathbb{Z}_2^\kappa \mapsto \mathbb{G}$  (e.g., Hash2Field [FHSS<sup>+</sup>23] for EC), parametrized by a group  $\mathbb{G}$

**Players:** a sender  $\mathcal{S}$ , and receiver  $\mathcal{R}$

**Inputs:**  $\mathcal{S}: \mathbf{a} \in \mathbb{G}^{\xi \times \ell}$ , group elements

$\mathcal{R}: \mathbf{x} \in \mathbb{Z}_2^\xi$ , the choice bits

**Outputs:**  $\mathcal{S} \leftarrow \mathbf{z}_A \in \mathbb{G}^{\xi \times \ell}$

$\mathcal{R} \leftarrow \mathbf{z}_B \in \mathbb{G}^{\xi \times \ell}$  s.t.  $z_{A(i,j)} + z_{B(i,j)} = a_{(i,j)} \cdot x_{(i)}$   $\forall i \in [\xi] \forall j \in [\ell]$

$\mathcal{S} \& \mathcal{R}. \text{RunROT}(b) \dashrightarrow \mathcal{S}: (\mathbf{r}_0, \mathbf{r}_1); \mathcal{R}: \mathbf{r}_b$

1:  $\mathcal{S}$  runs  $\text{ROT}_{\xi, \ell}$  as sender, obtaining  $\mathbf{r}_0, \mathbf{r}_1 \in \mathbb{Z}_2^{\xi \times \ell \times \kappa}$

2:  $\mathcal{R}$  runs  $\text{ROT}_{\xi, \ell}$  as receiver, obtaining  $\mathbf{r}_b \in \mathbb{Z}_2^{\xi \times \ell \times \kappa}$

$\mathcal{S}. \text{CreateCorrelation}(\mathbf{r}_0, \mathbf{r}_1, \mathbf{a}) \dashrightarrow (\boldsymbol{\tau})$

1:  $\mathbf{z}_A \leftarrow \{\{H(r_0(i,j))\}_{j \in [\ell]}\}_{i \in [\xi]}$

2:  $\boldsymbol{\tau} \leftarrow \{\{H(r_1(i,j)) - z_{A(i,j)} + a_{(i,j)}\}_{j \in [\ell]}\}_{i \in [\xi]}$

3:  $\text{Send}(\boldsymbol{\tau}) \rightarrow \mathcal{R}$

return  $\mathbf{z}_A$

$\mathcal{R}. \text{ApplyCorrelation}(\boldsymbol{\tau}) \dashrightarrow \mathbf{z}_B$

return  $\mathbf{z}_B \leftarrow \{\{\tau_{(i,j)} \cdot b_{(i)} - H(r_b(i,j))\}_{j \in [\ell]}\}_{i \in [\xi]}$

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## References

- [FHSS<sup>+</sup>23] Armando Faz-Hernandez, Sam Scott, Nick Sullivan, Riad S. Wahby, and Christopher A. Wood. Hashing to Elliptic Curves. RFC 9380, August 2023.