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Security through Distrusting

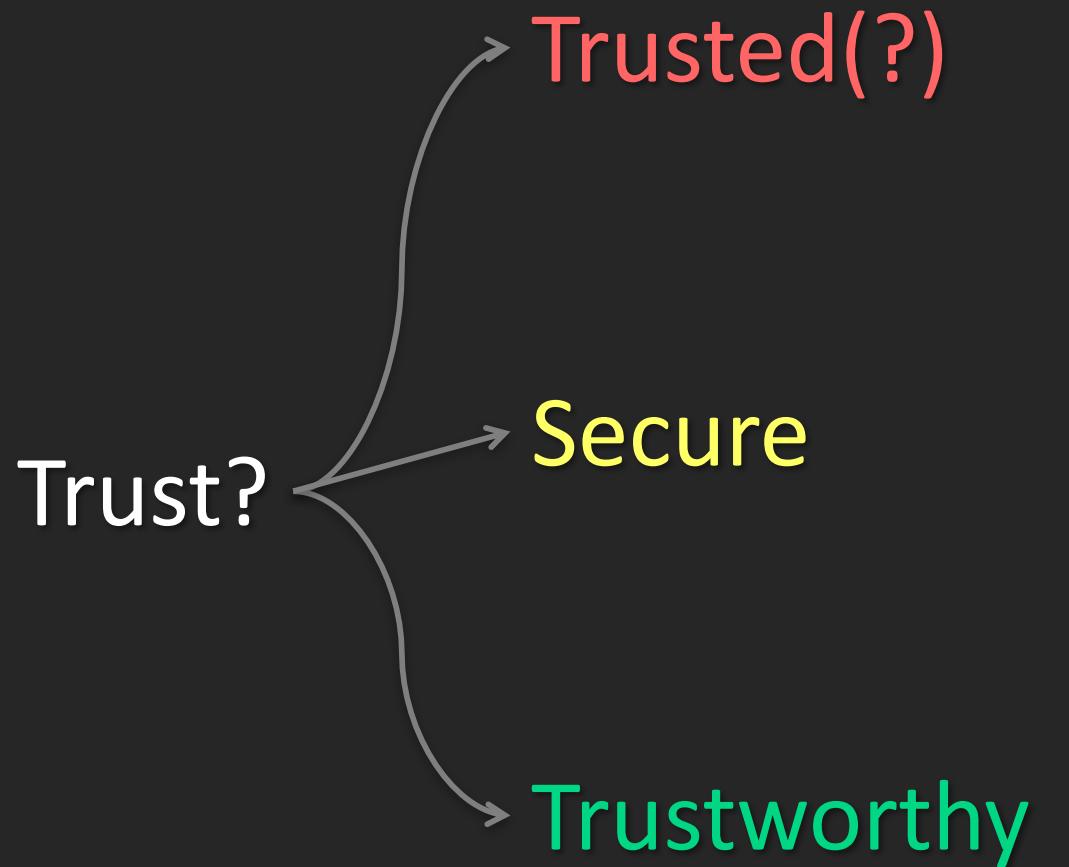
Joanna Rutkowska

Invisible Things Lab & [Qubes OS Project](#)

Black Hat EU, London, UK, December 7, 2017



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more desired



Trust consider harmful!

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Security through Distrusting examples

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Example #1: Pesky microphones

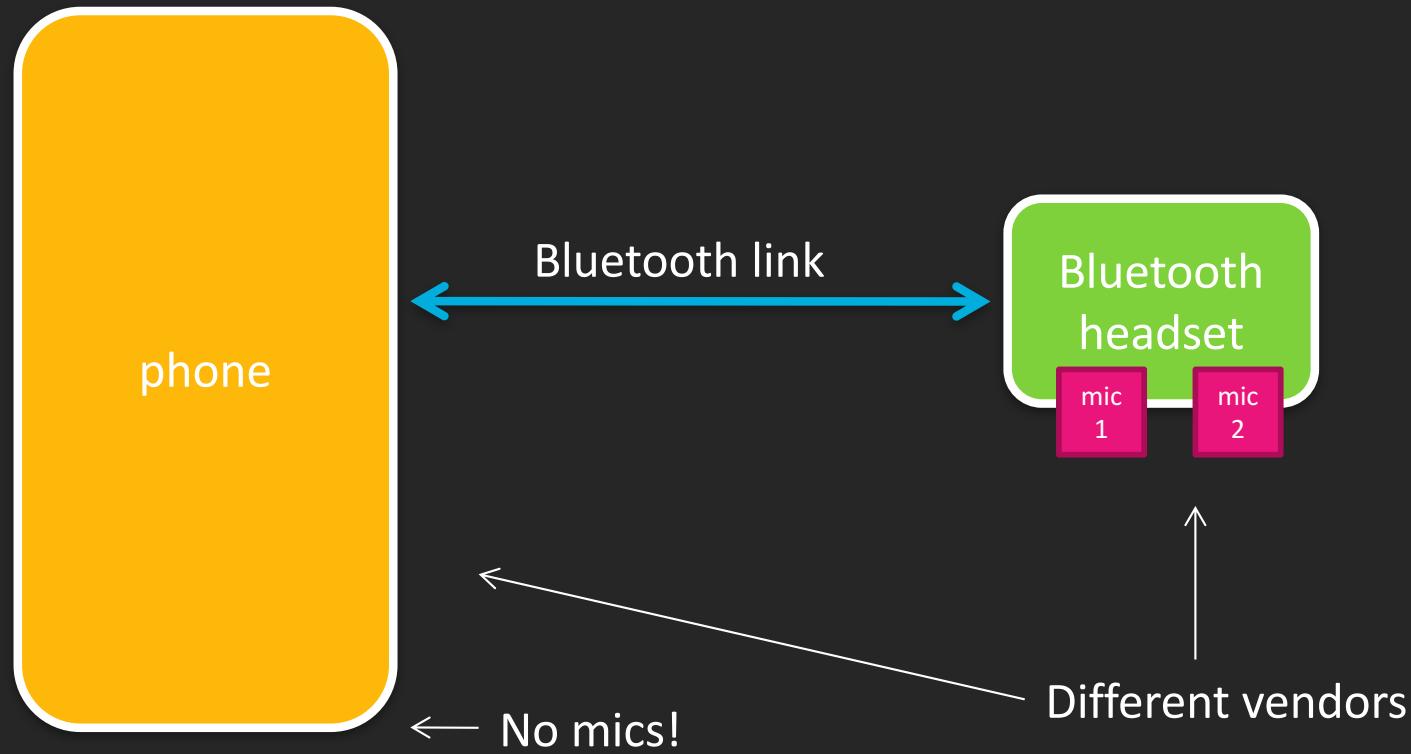




- Mics sniff our activities, including keystrokes, etc.
- Mics are difficult to neutralize
- Mics naturally “cross” security boundaries



- Mics sniff our activities, including keystrokes, etc.
- Mics are difficult to neutralize
- Mics naturally “cross” security boundaries



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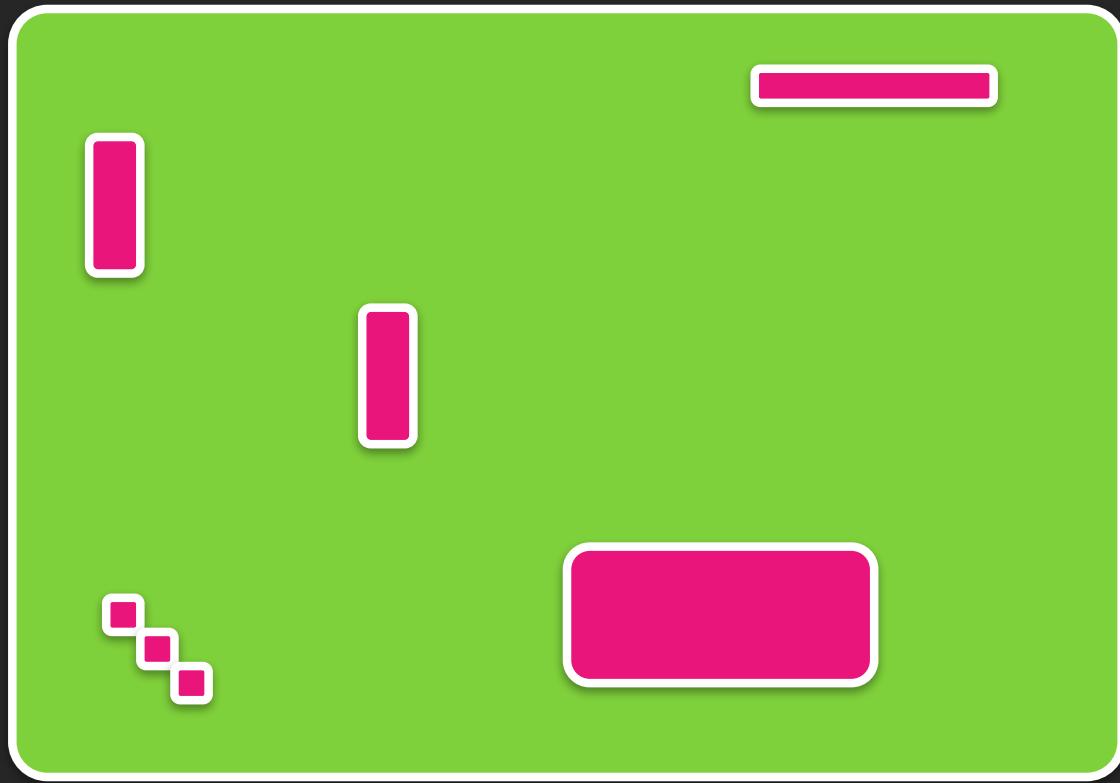
Example #2: Stateless laptop



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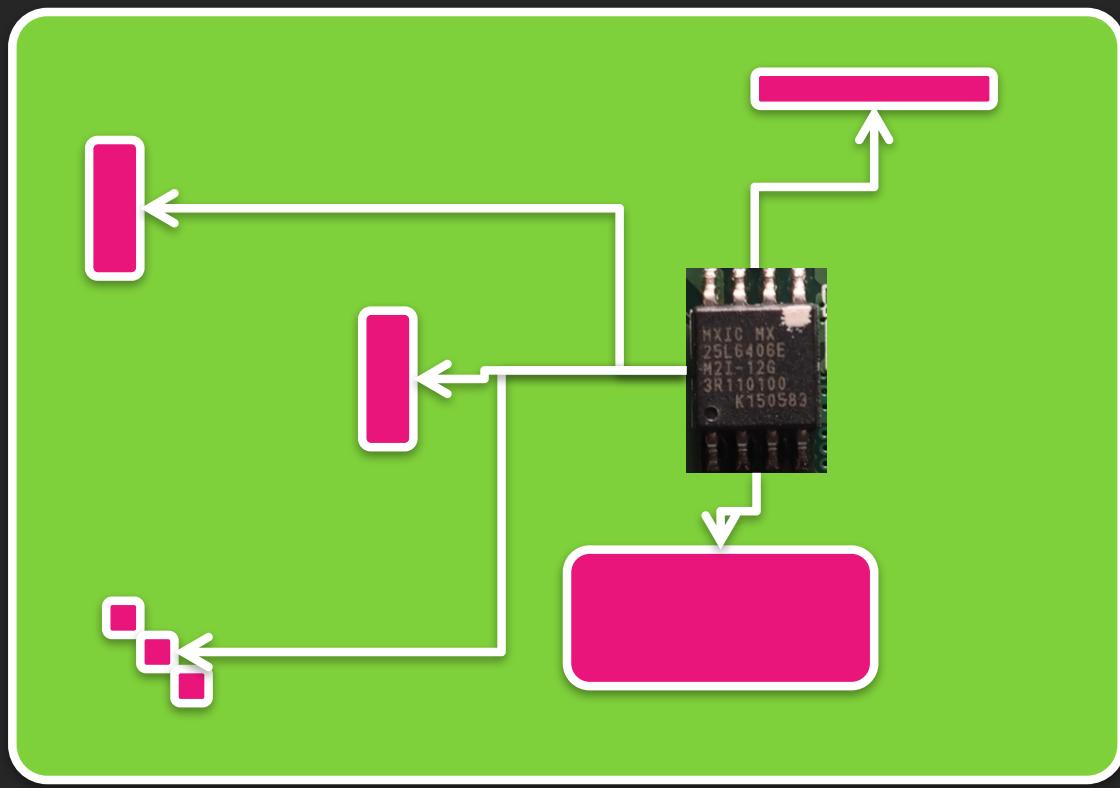
Persistent laptop compromises...

- Persist
- Store secrets
- PII



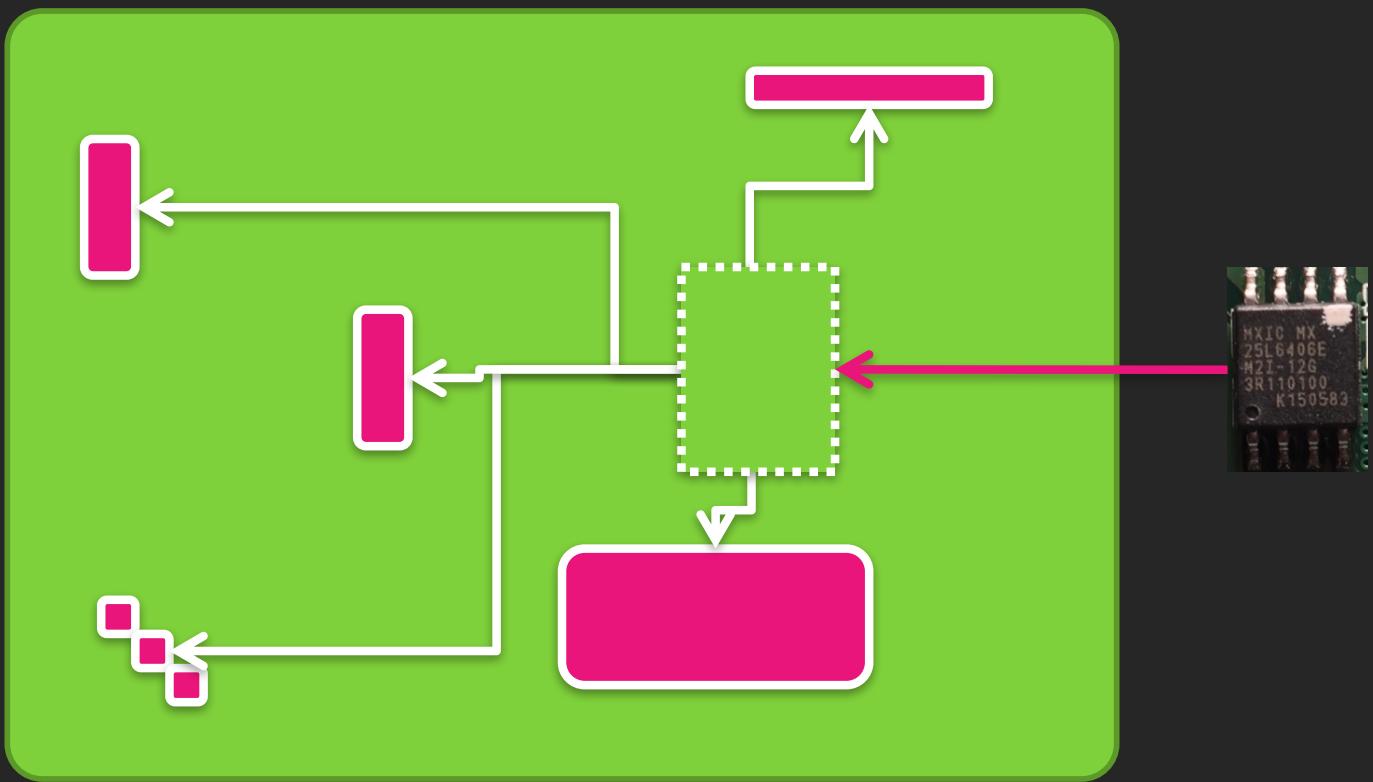
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- Persist
- Store secrets
- PII



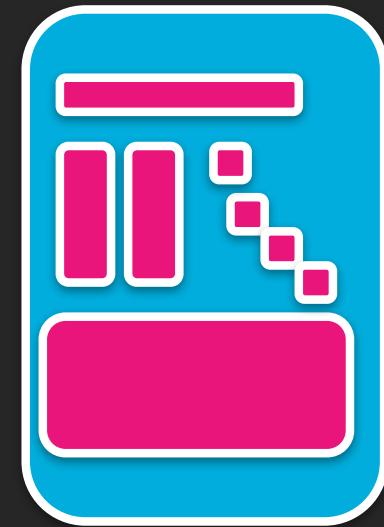
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- Persist
- Store secrets
- PII



Stateless Hardware (persistent state eliminated)

- Firmware infections prevented
- No places to store stolen secrets
- Reliable way to verify firmware
- Reliable way to *choose* firmware
- Boot multiple environments
- Share laptops with others



Trusted Stick



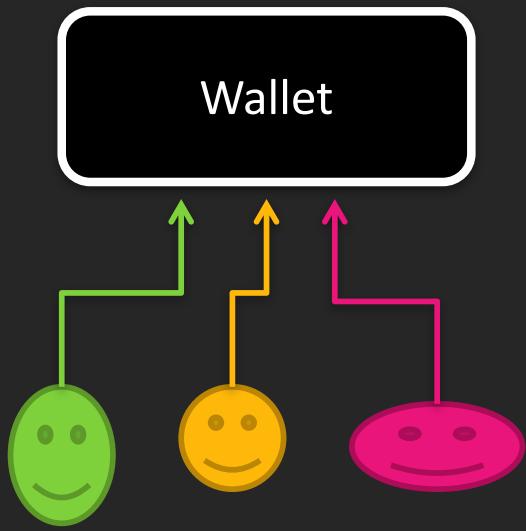
Example #3: Multi-party signatures



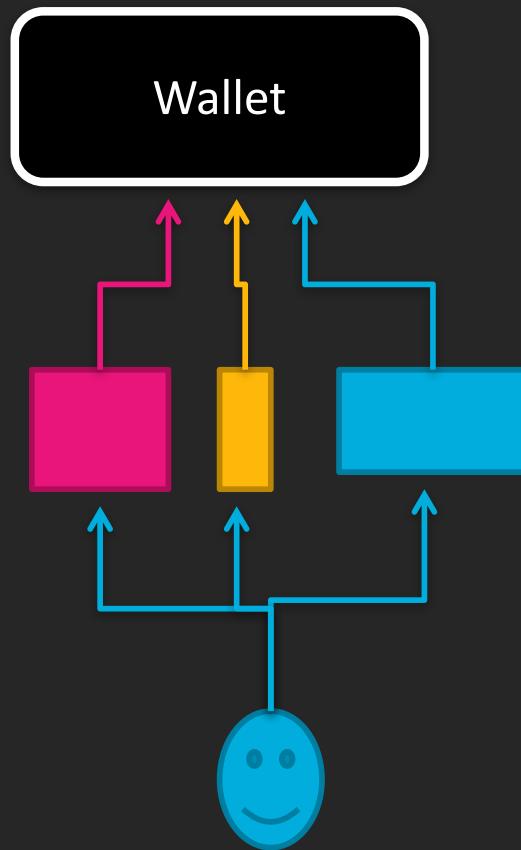
Photo via Peter Todd (@petertoddbtc)



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vs



Multi-sig does not need to involve multiple users!



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Also: not just Bitcoin wallets...



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Example #4: Binary (multi) signing



Why care about binary (multi-) signing?

- OS installation images
- Applications
- Updates
- Firmware

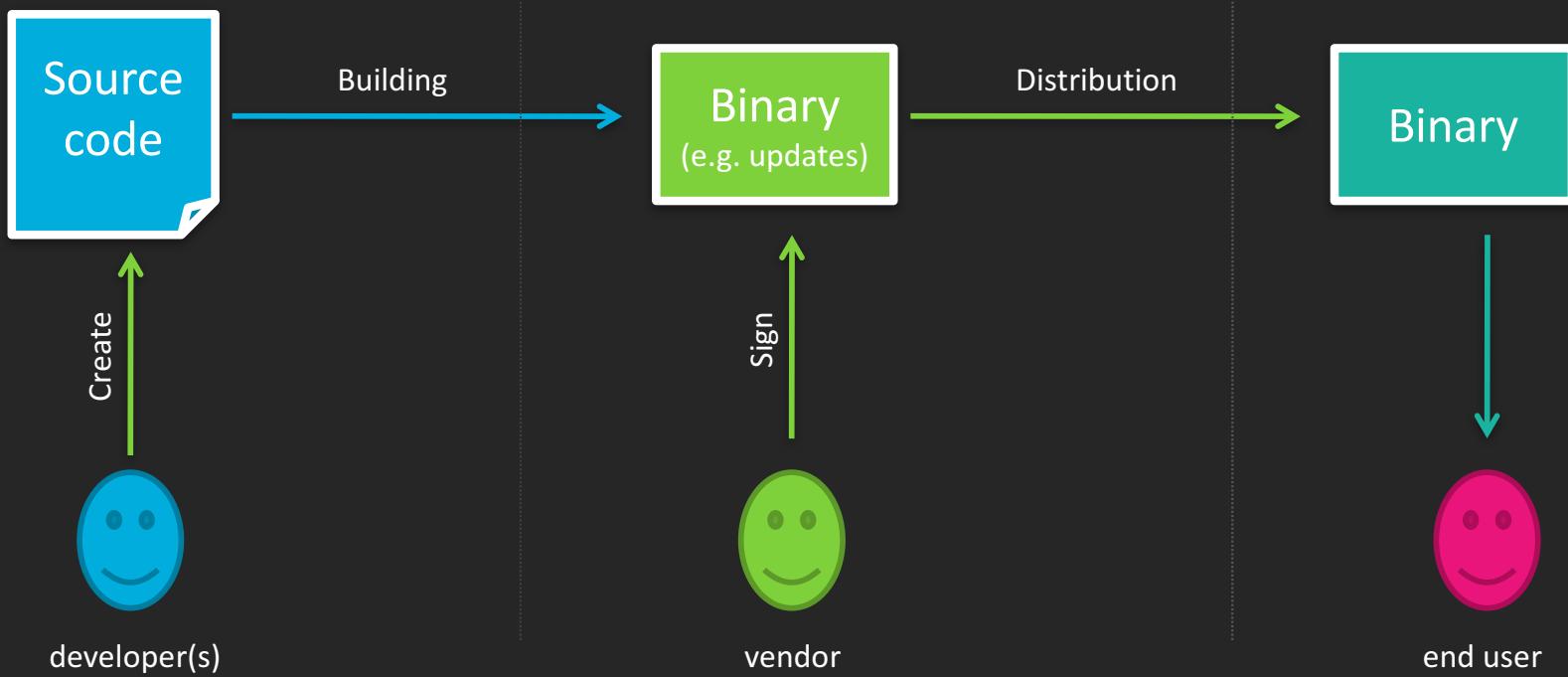


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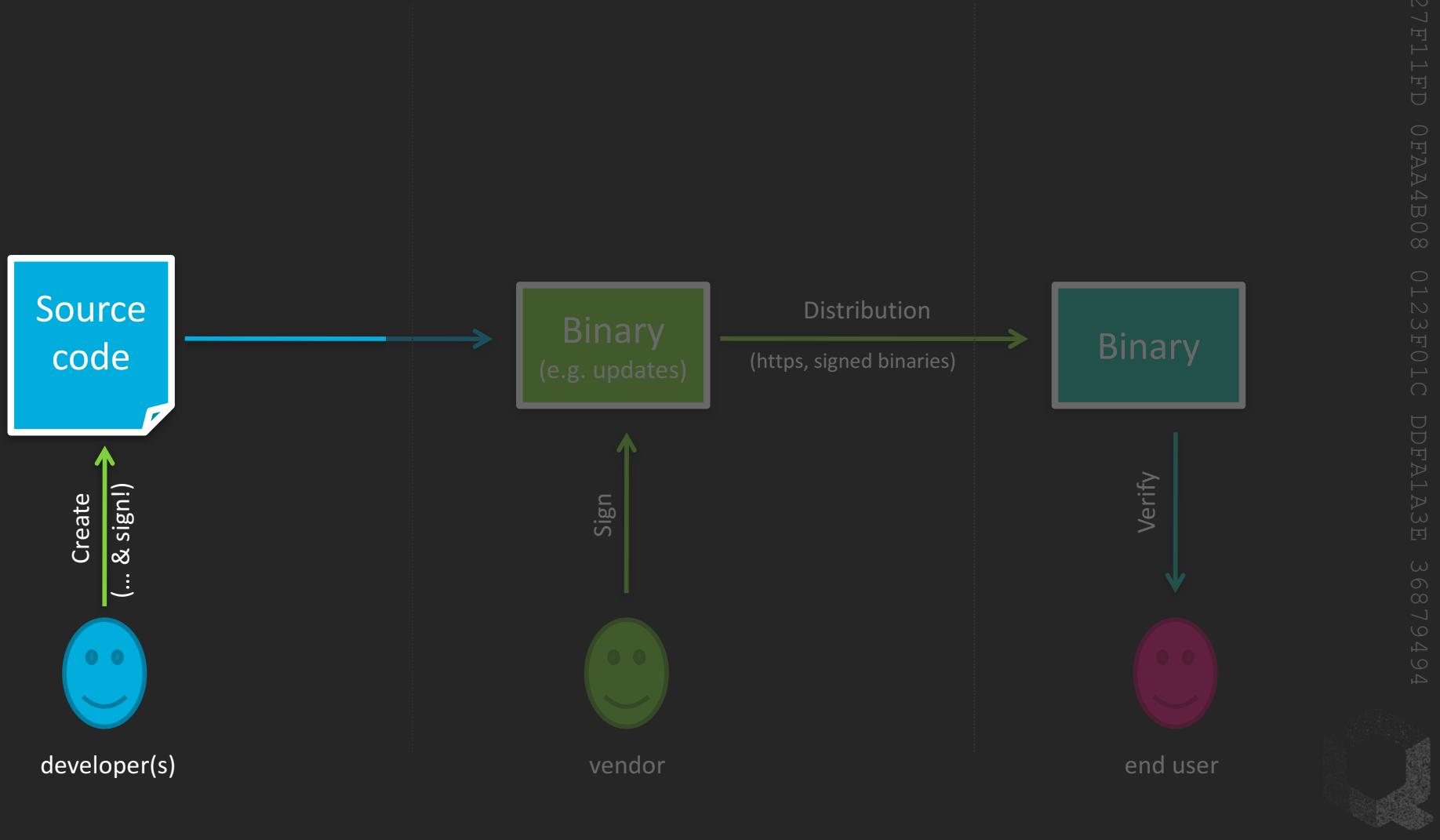
Prime target for backdooring!



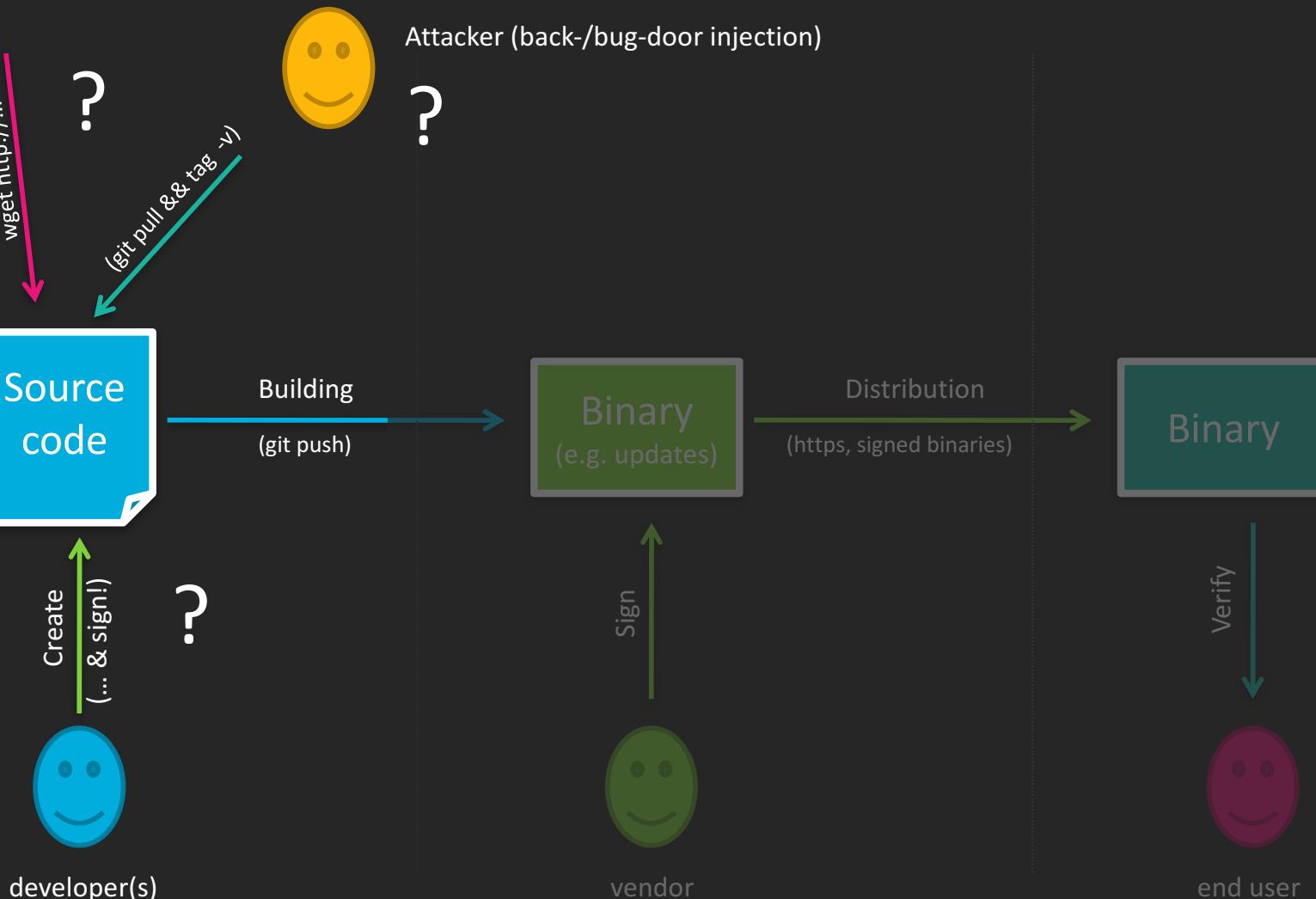
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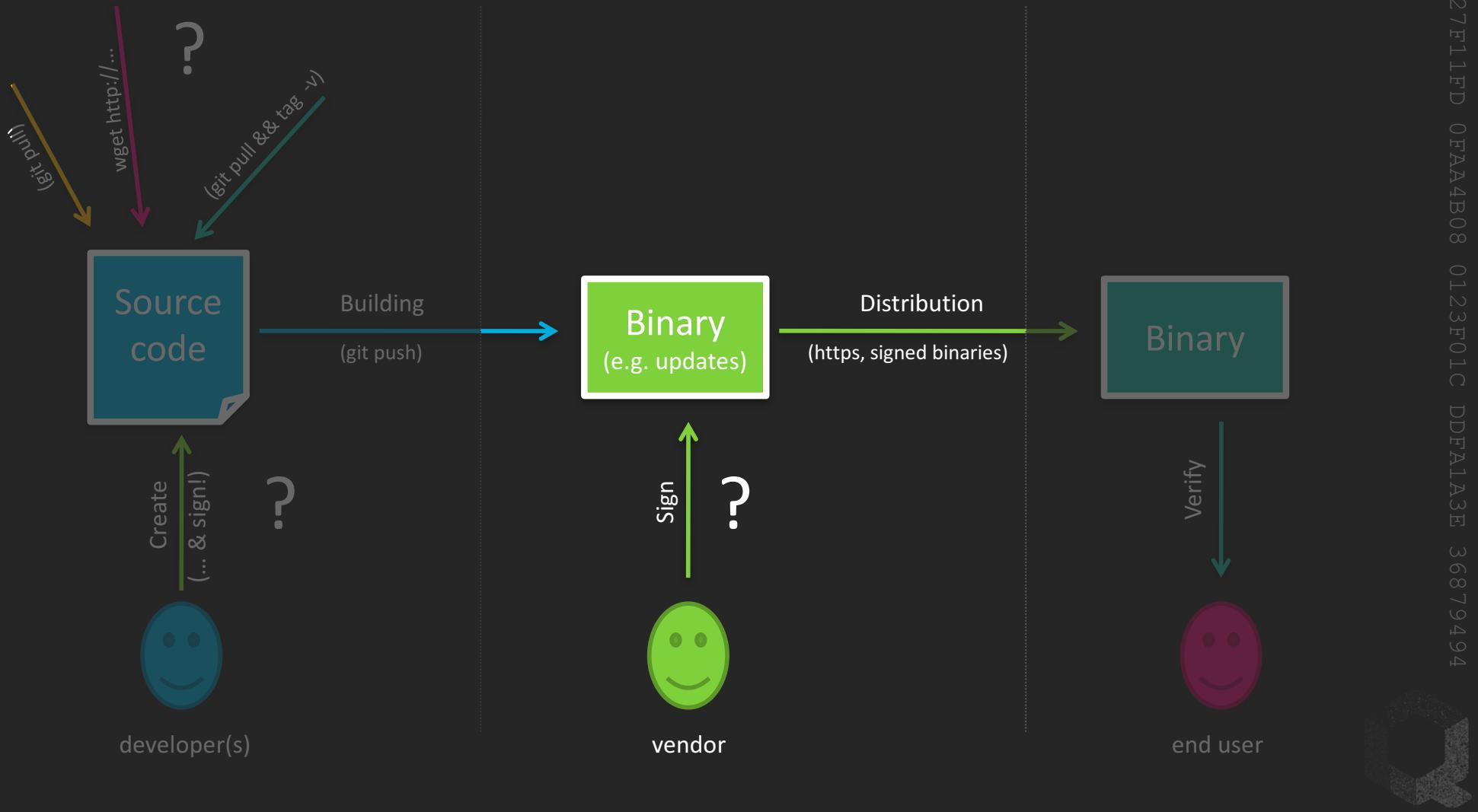
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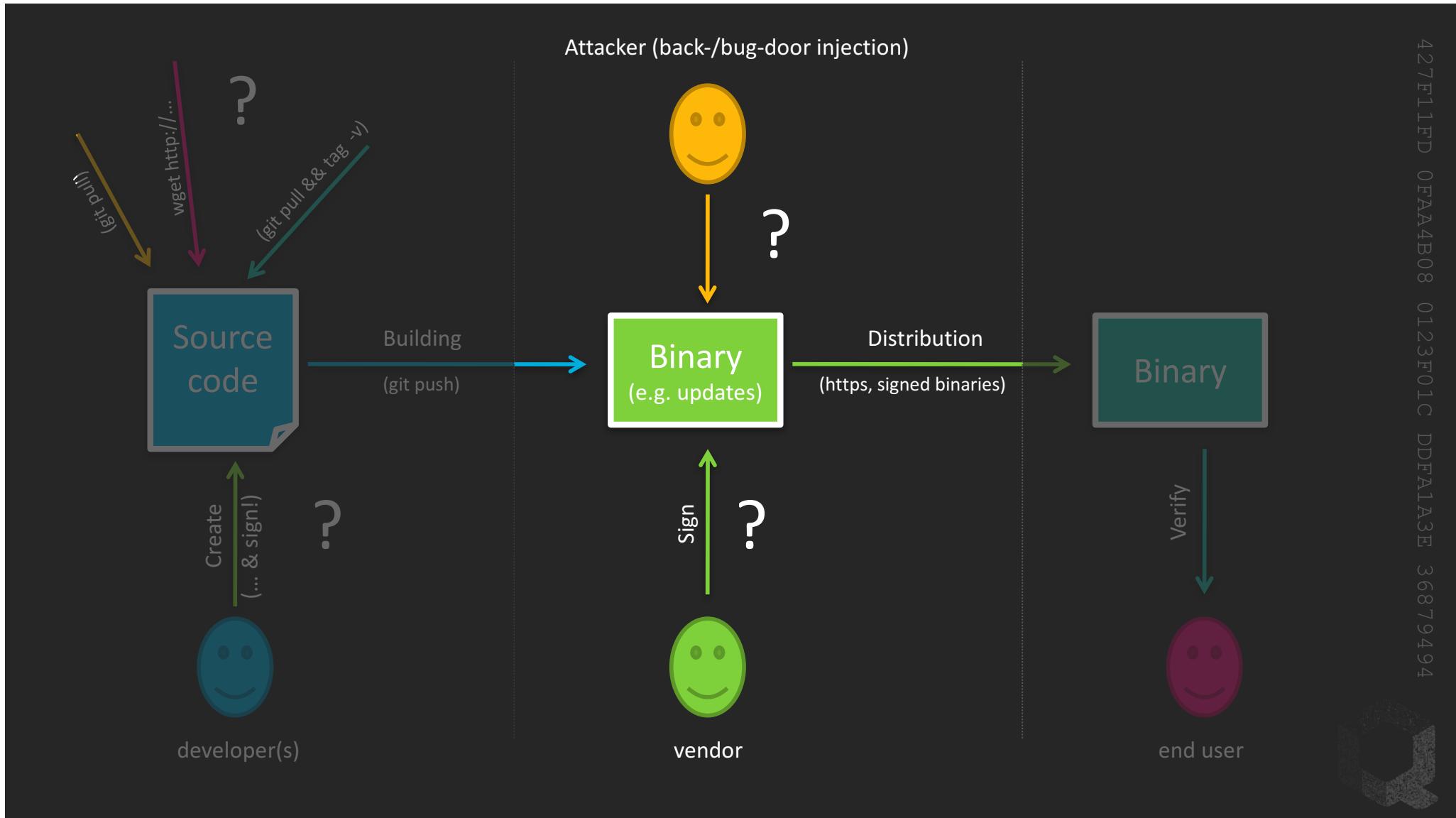
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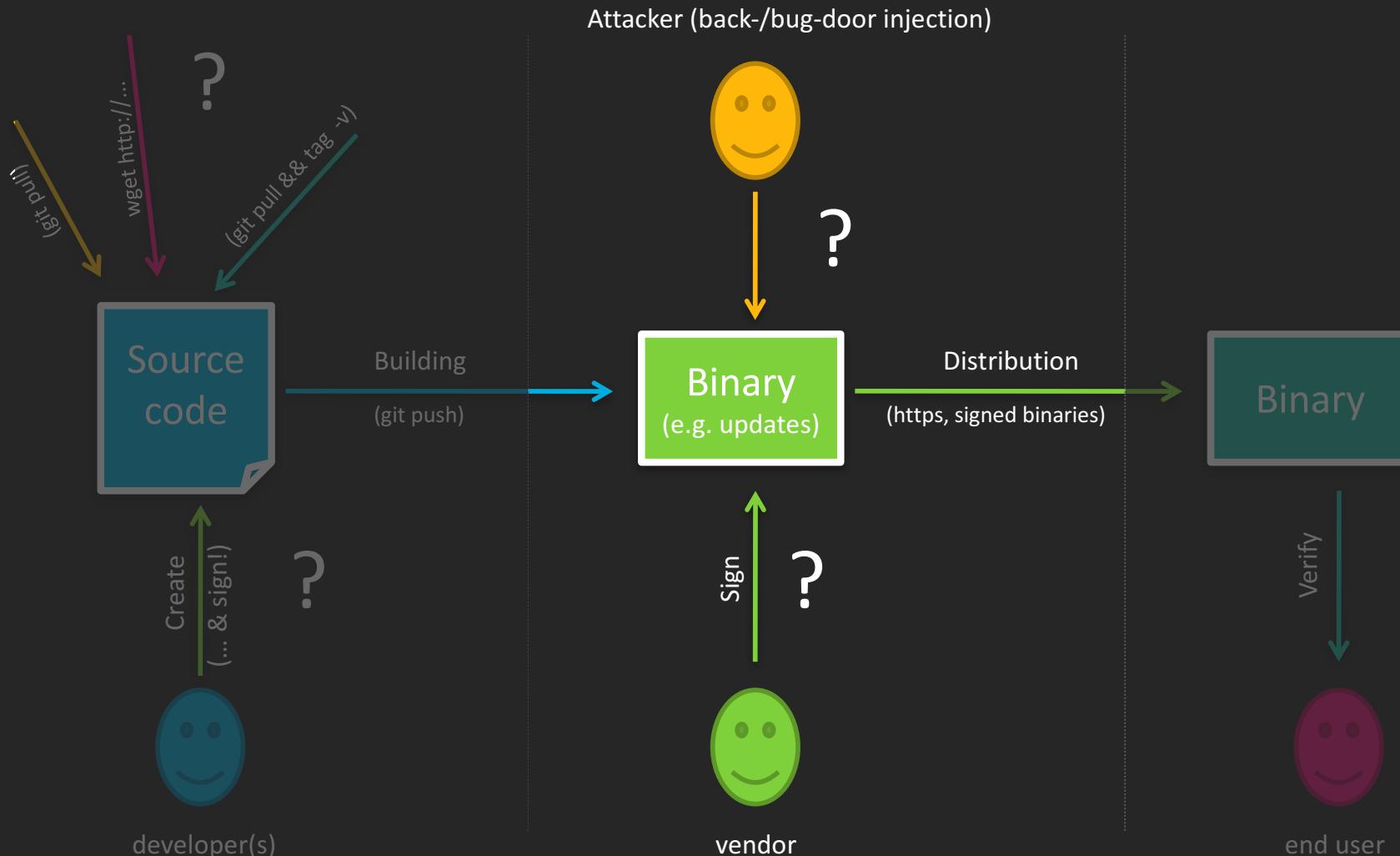
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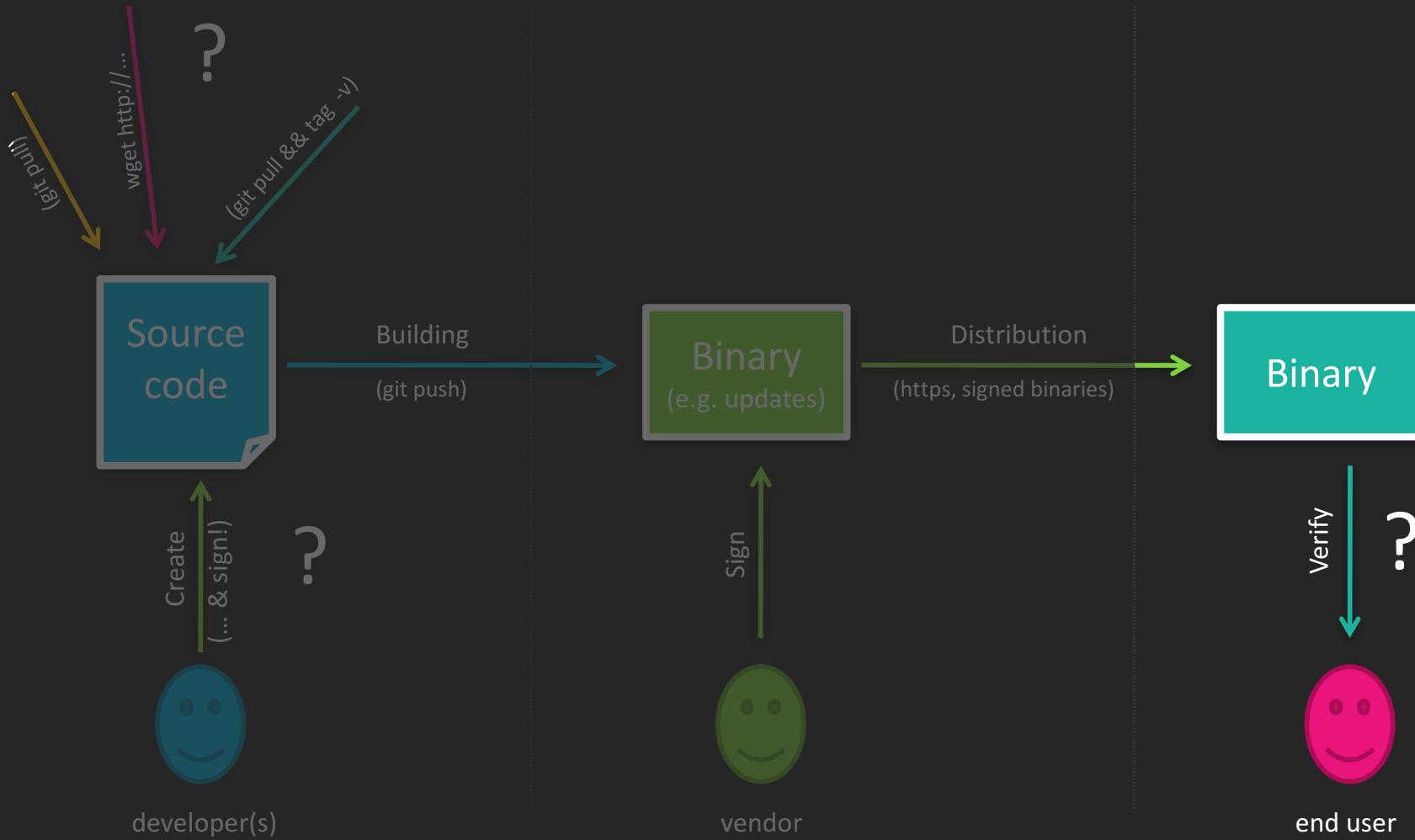
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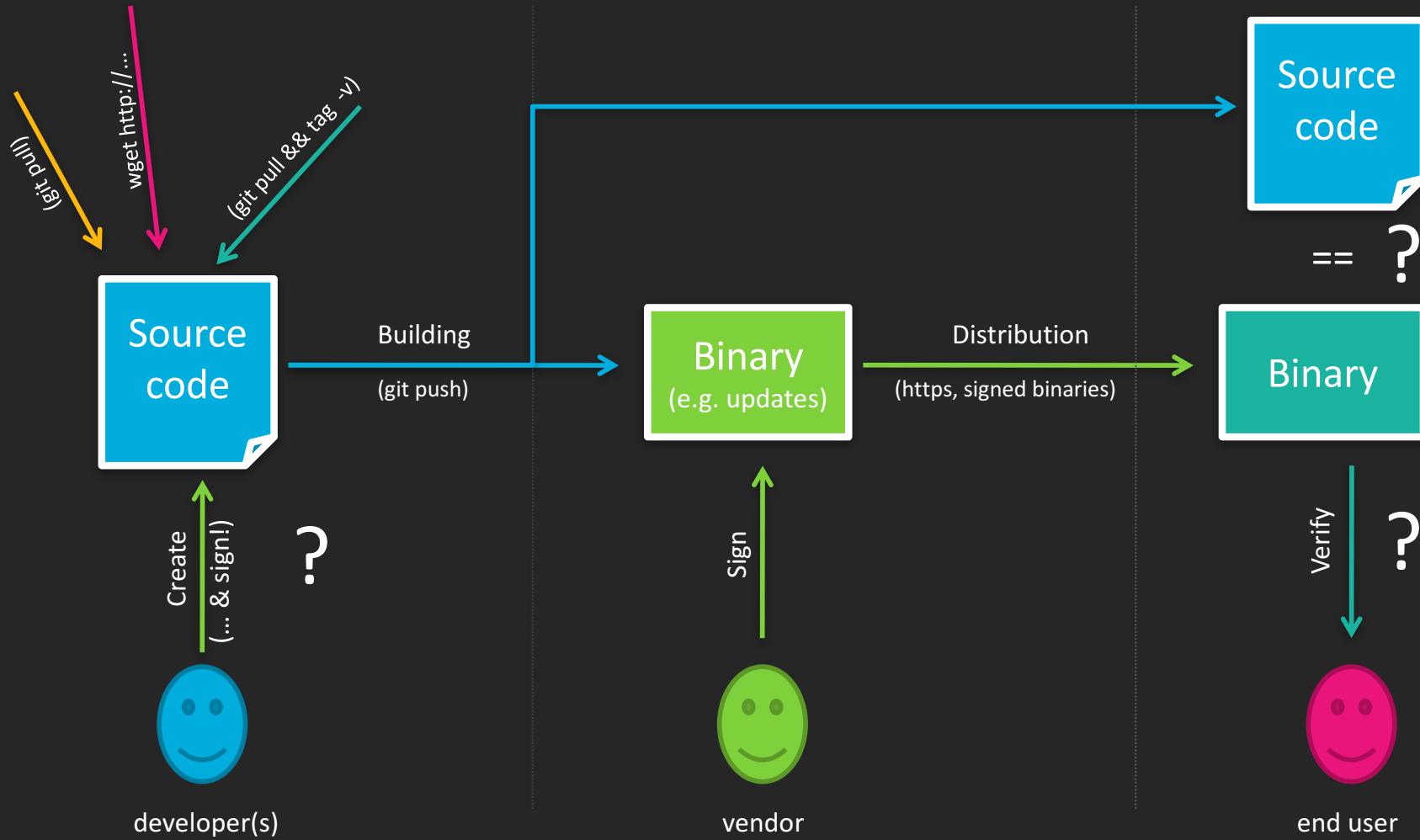
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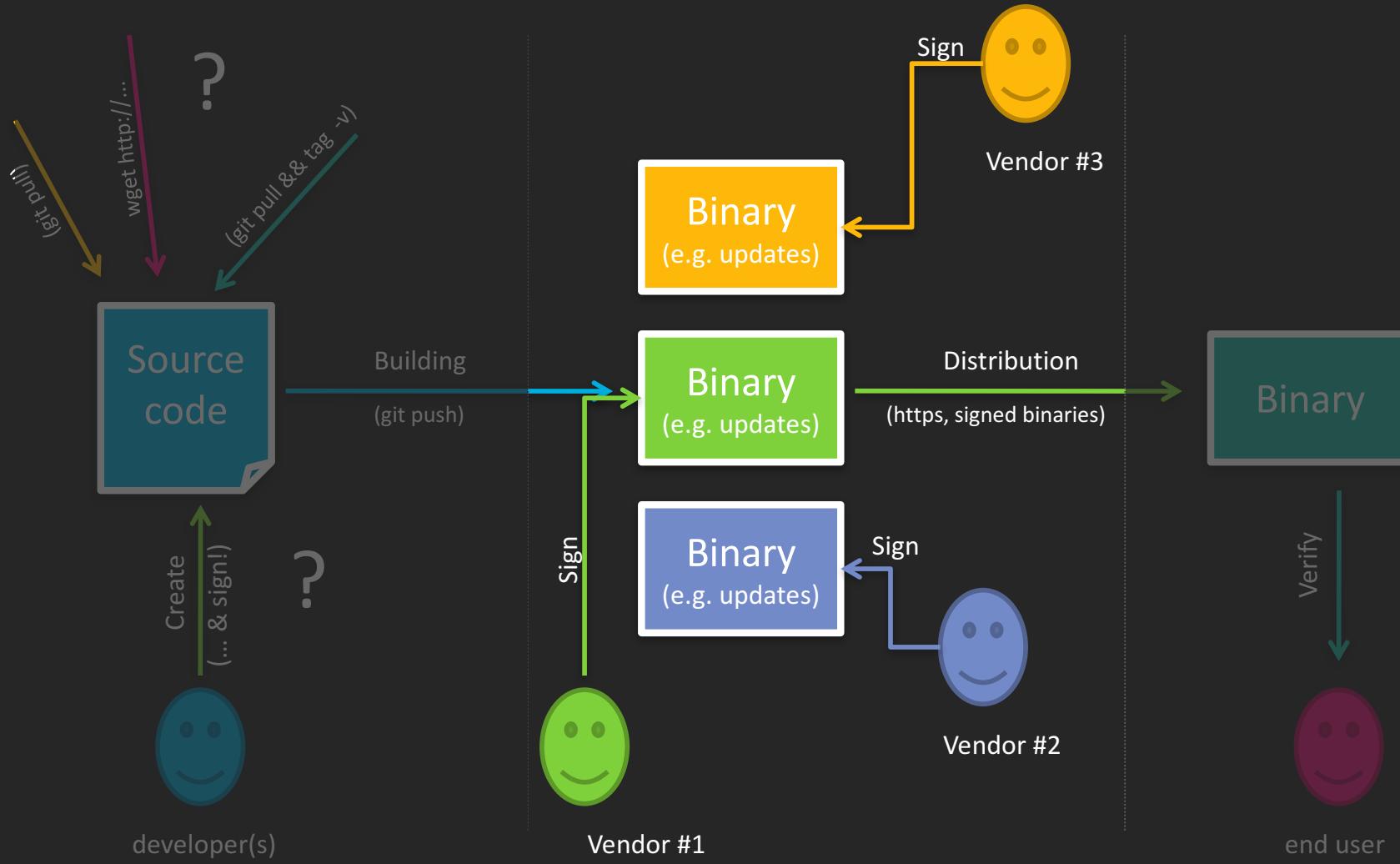
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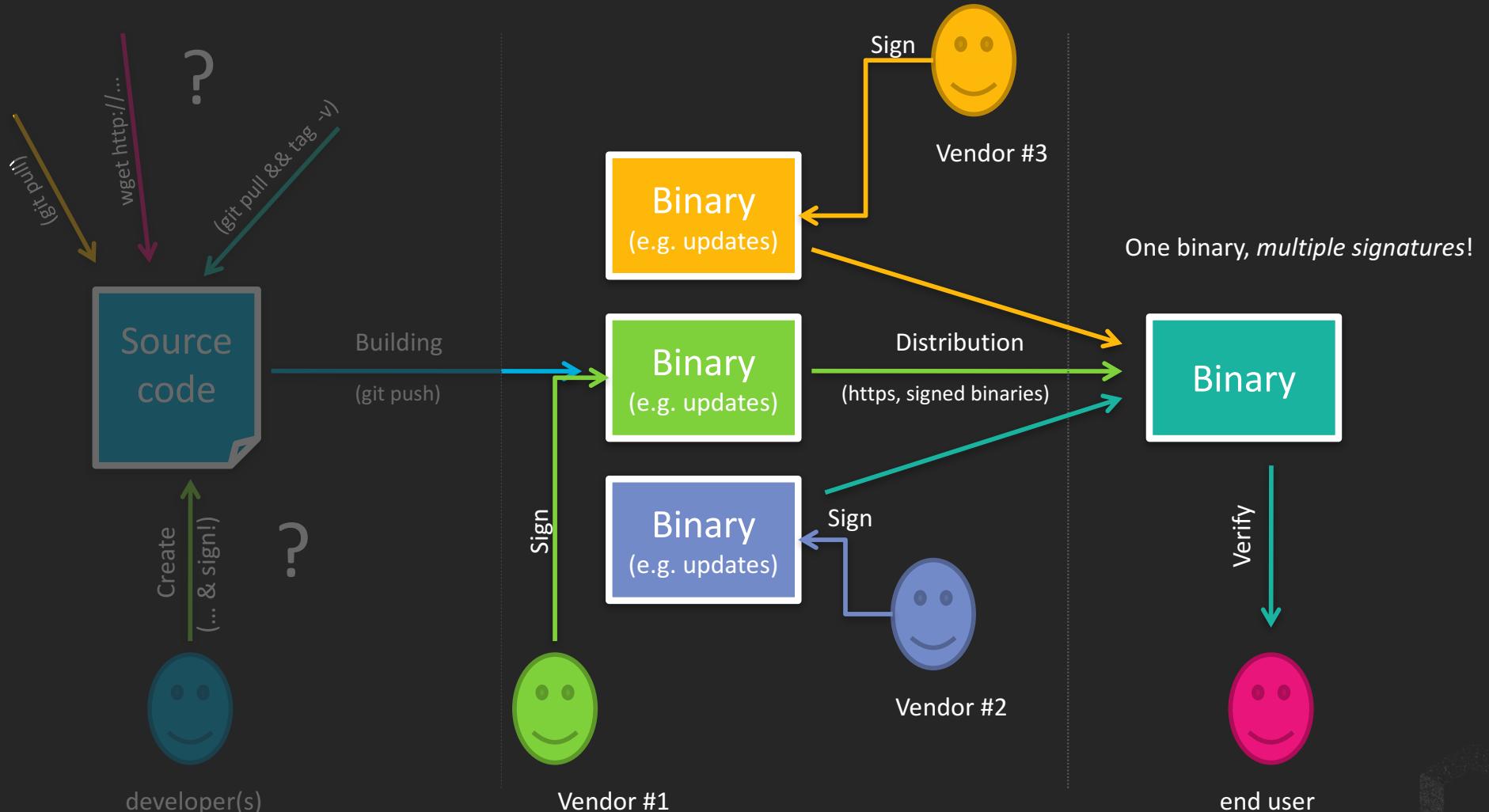
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Multi-signed binaries

- Signed by people from different countries
- Different organizations (vendor & auditing)
- Signed by different *machines*
 - In the same organization
 - In different organization



<https://reproducible-builds.org>

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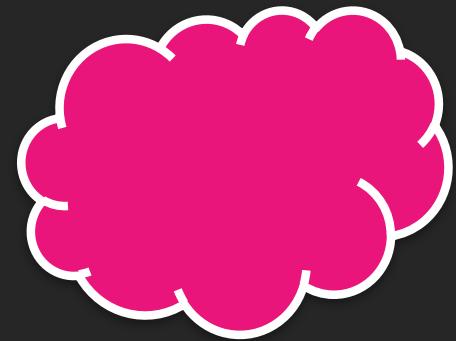
Example 5: Preventing data leaks



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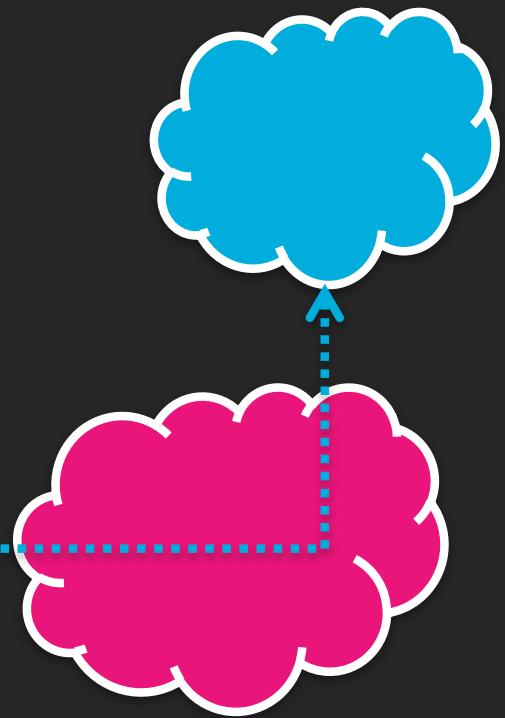
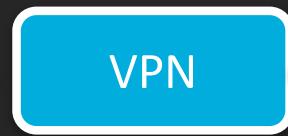
Your data...!

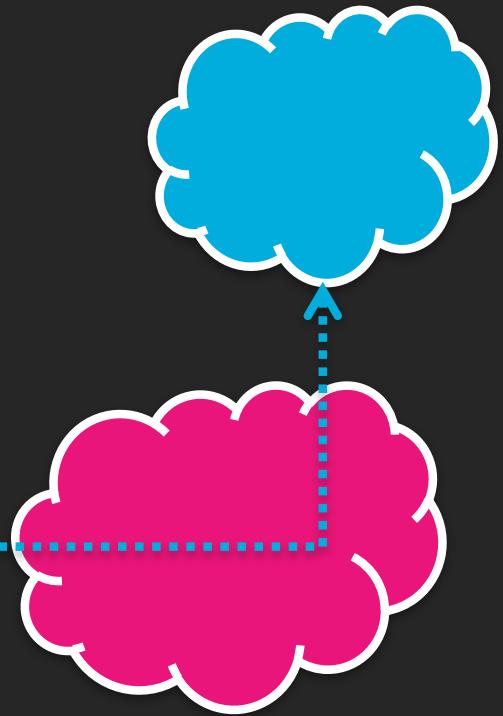
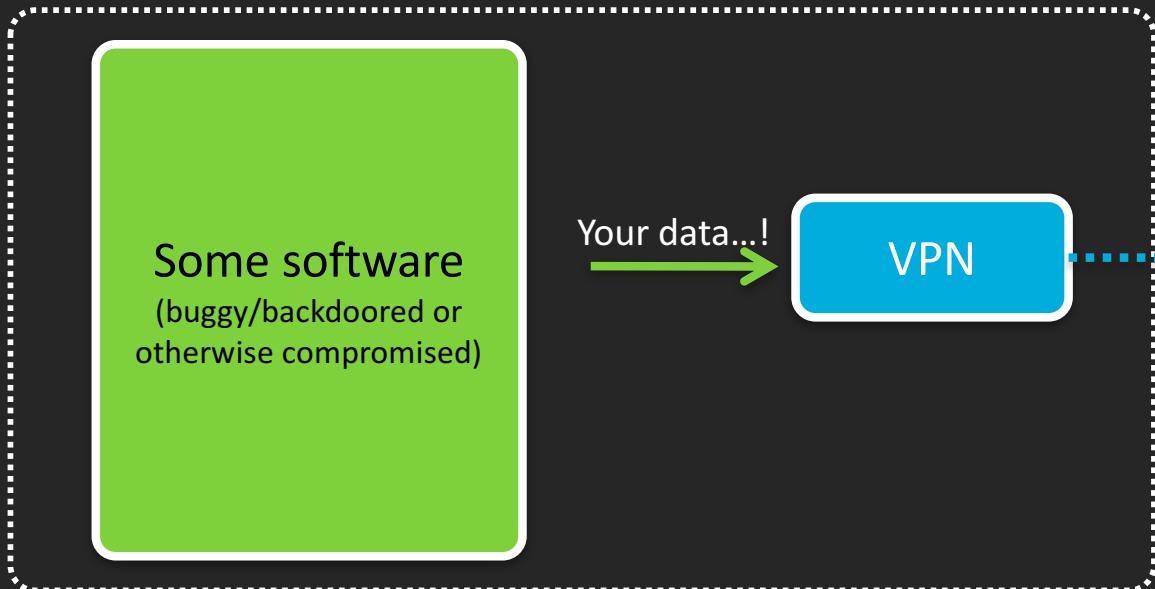
A thick yellow arrow points from the text "Your data...!" towards the pink cloud icon.

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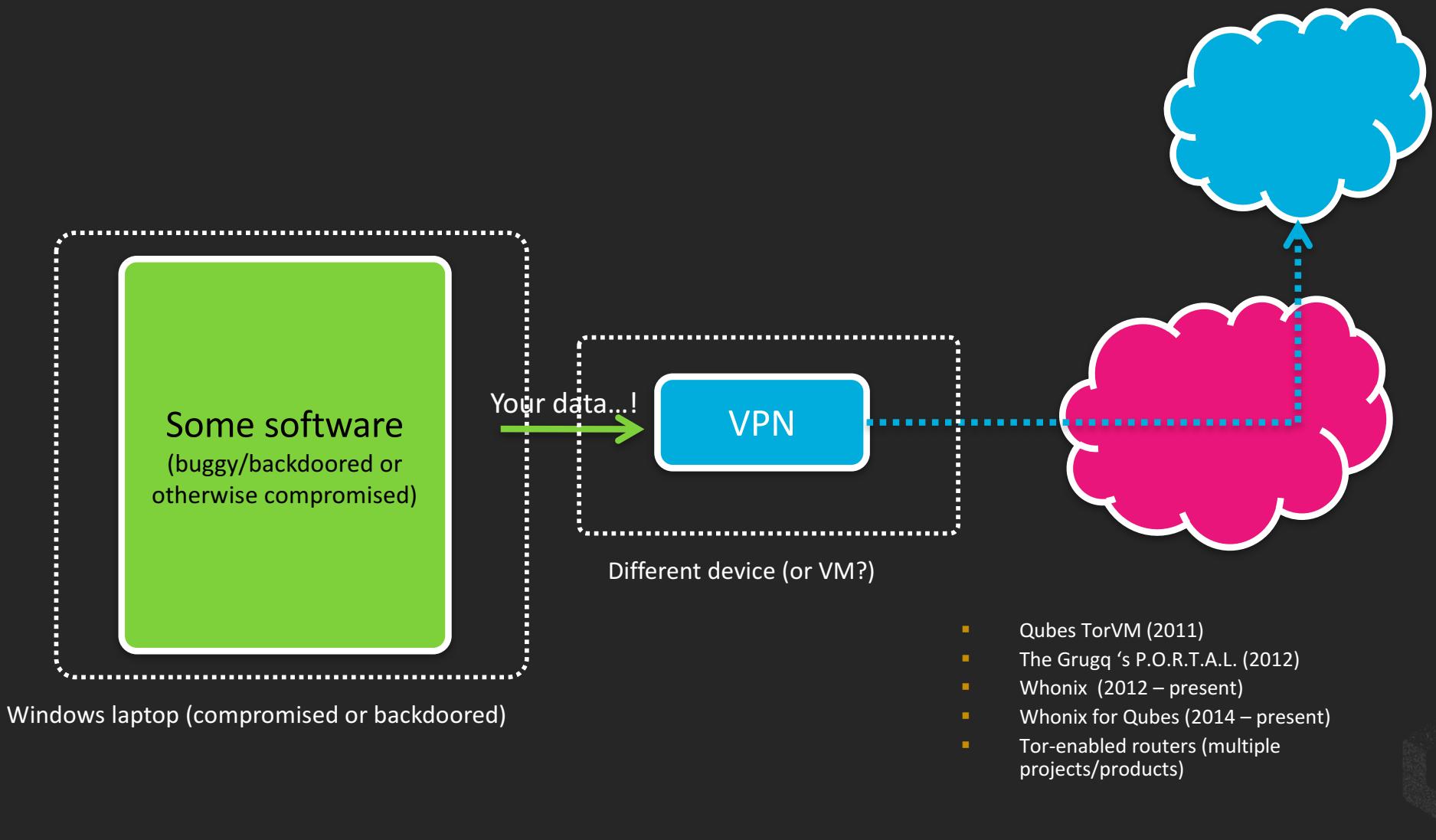
Your data...!



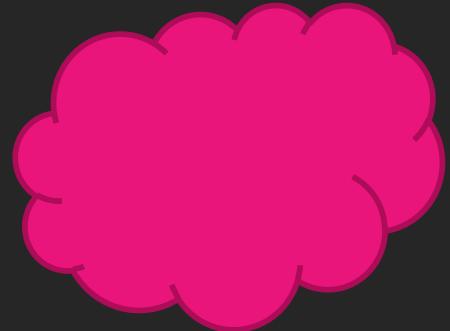
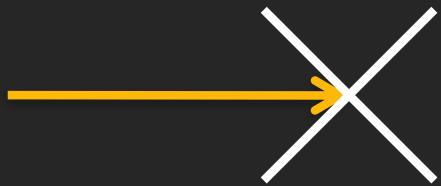
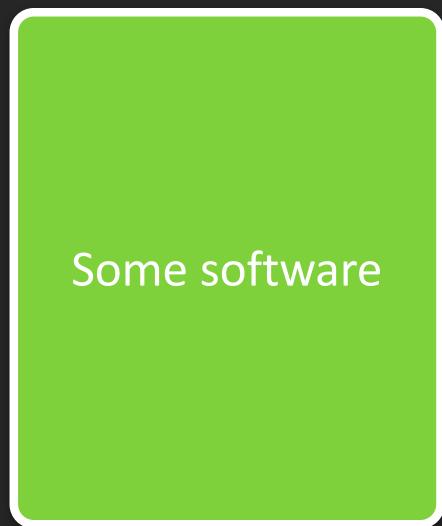


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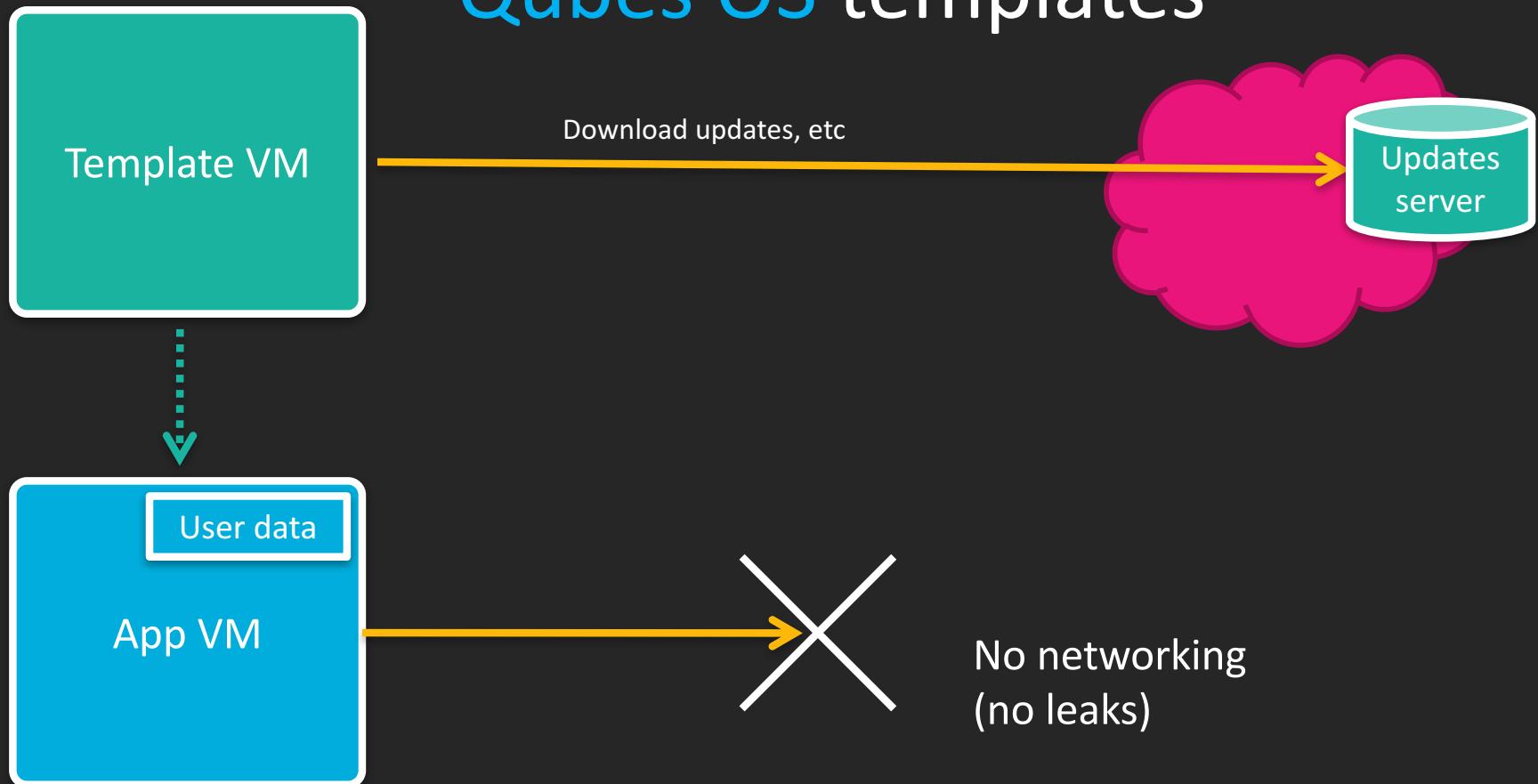
Cut off networking?



Not very useful...



Qubes OS templates



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Example #6: Compartmentalization



“Classic” compartmentalization...

Work VM

Personal VM



“Classic” compartmentalization...

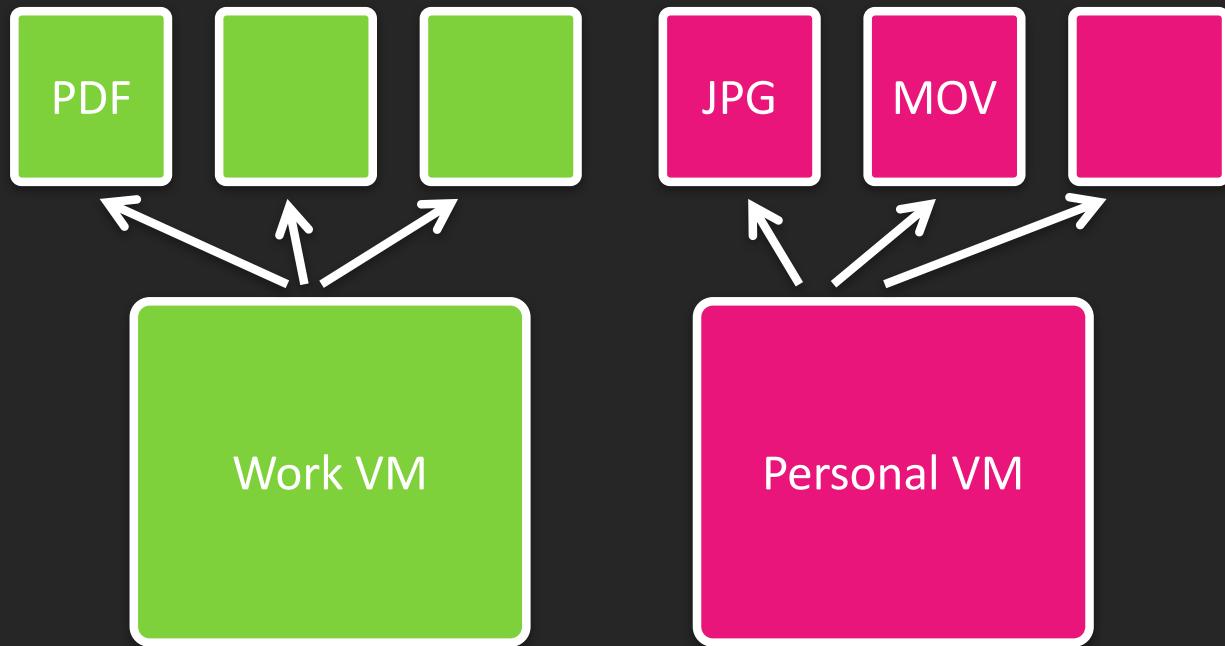
Work VM

Personal VM

...not very useful!



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...more useful...

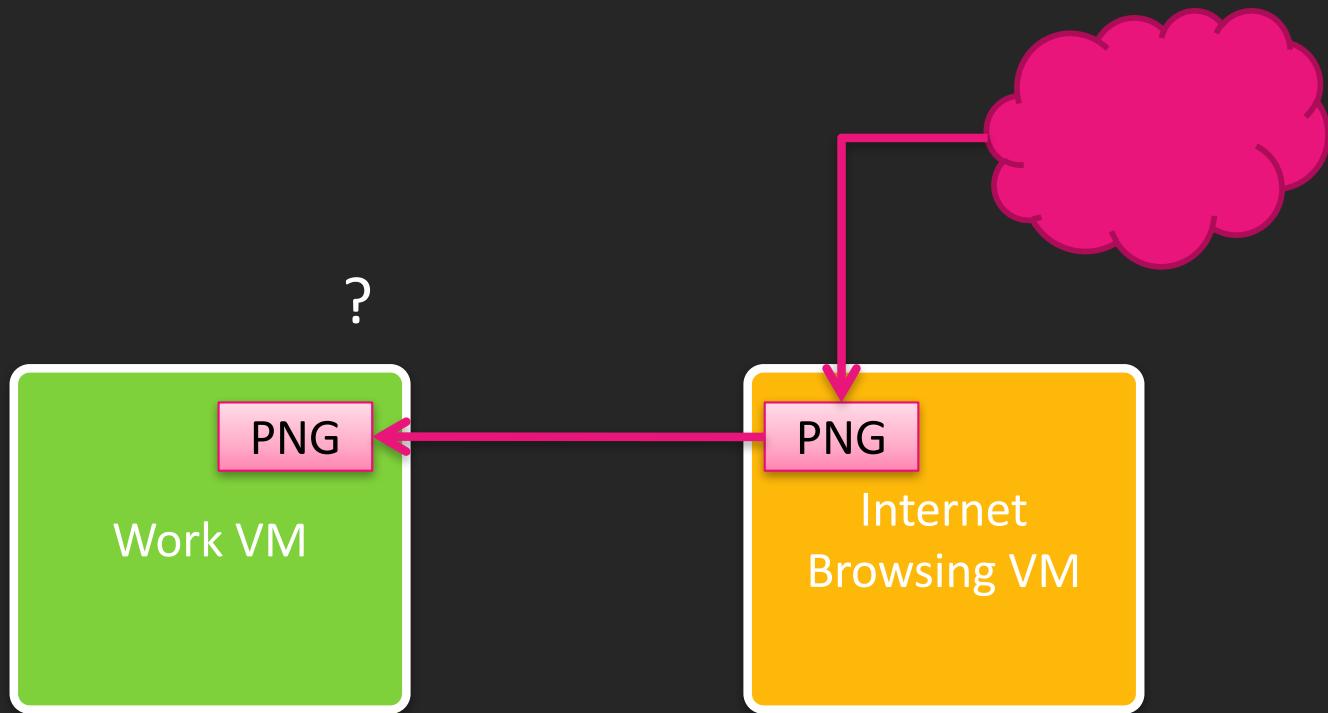


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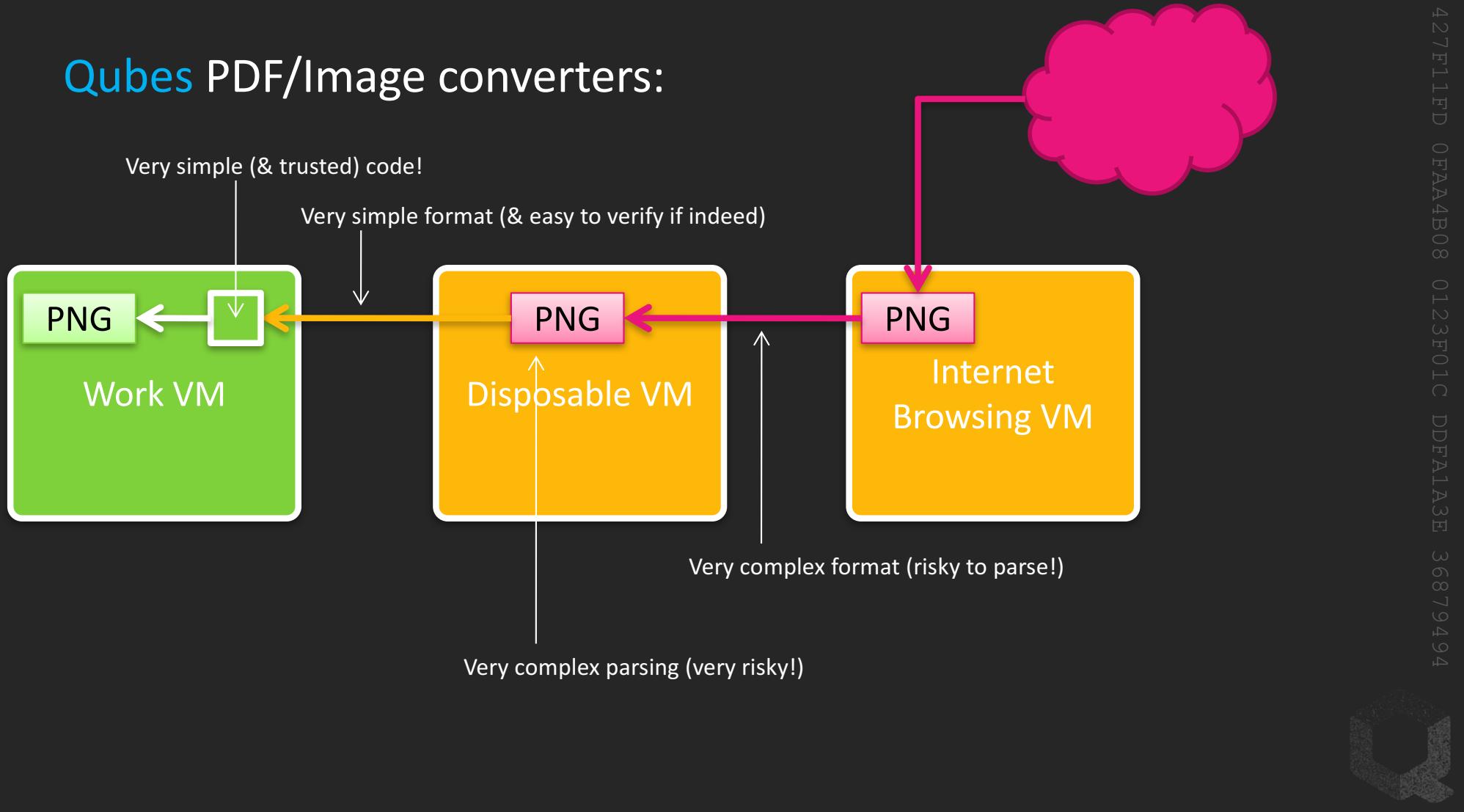
Inter-compartments data transfers



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Qubes PDF/Image converters:



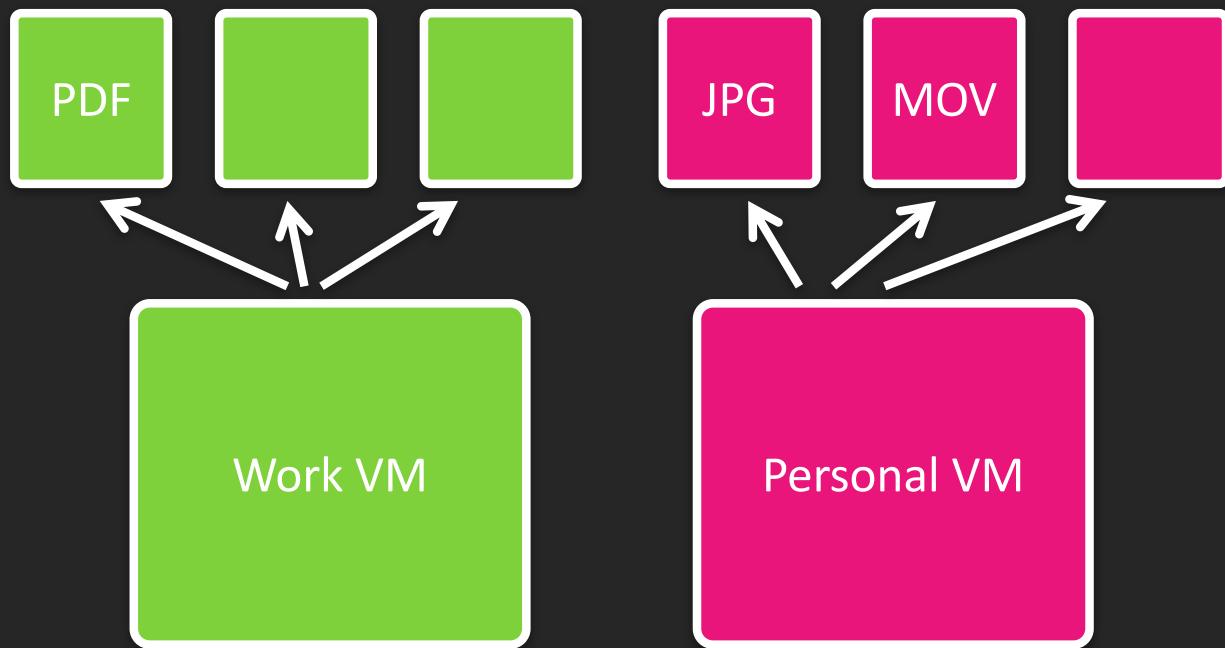
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App sandboxing is just part of the story...

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Networking
stacks

USB &
Bluetooth
stacks

Graphics &
UI

VPNs &
firewalling

Corporate
management

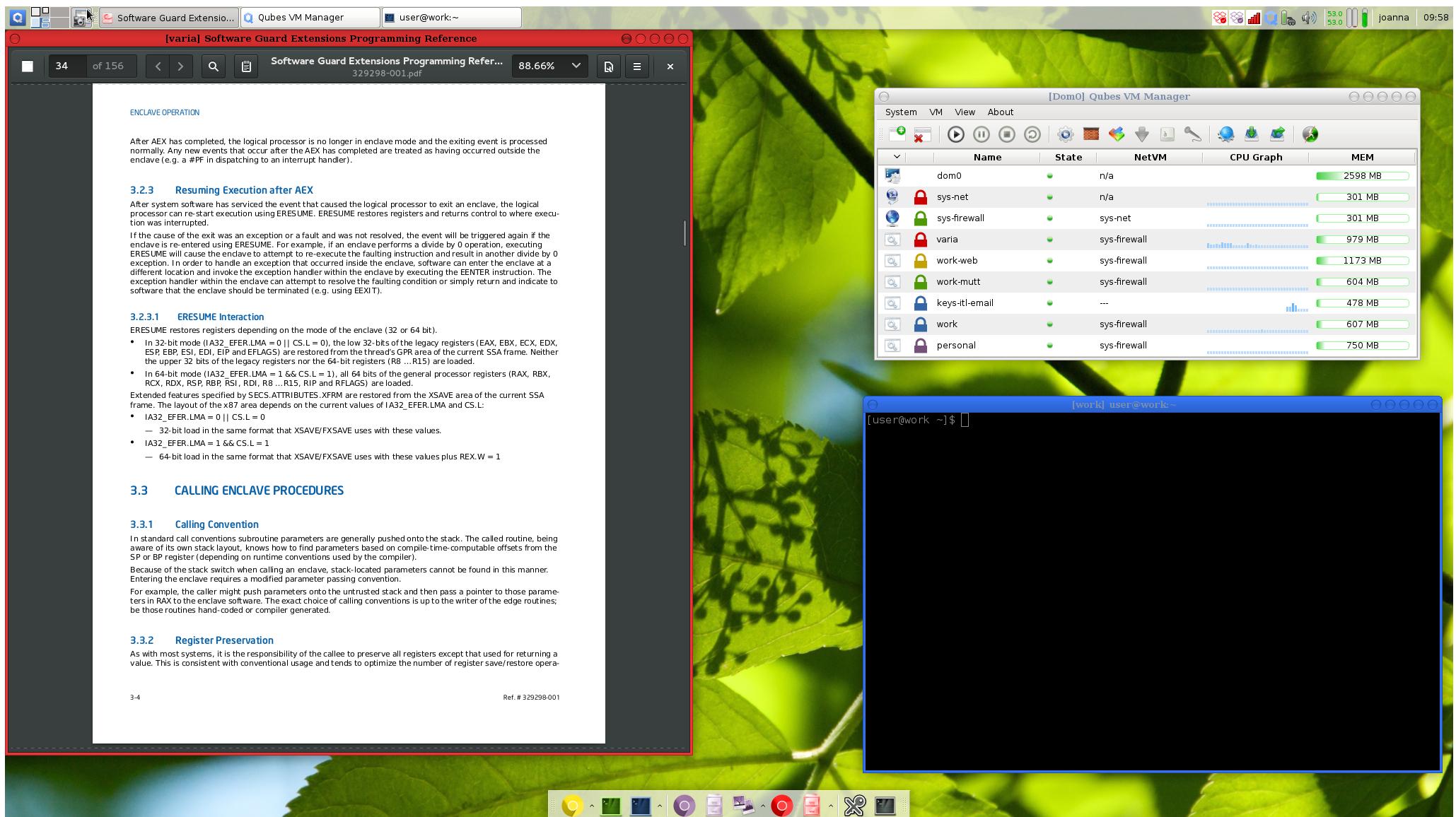
Root of trust
(admin)



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Isolation is just part of the story!





Software Guard Extension... Qubes VM Manager user@work:~

[varia] Software Guard Extensions Programming Reference

34 of 156

< > Q PDF Software Guard Extensions Programming Refer...

329298-001.pdf

88.66%

Q

☰

X

ENCLAVE OPERATION

After AEX has completed, the logical processor is no longer in enclave mode and the exiting event is processed normally. Any new events that occur after the AEX has completed are treated as having occurred outside the enclave (e.g. a #PF in dispatching to an interrupt handler).

3.2.3 Resuming Execution after AEX

After system software has serviced the event that caused the logical processor to exit an enclave, the logical processor can re-start execution using ERESUME. ERESUME restores registers and returns control to where execution was interrupted.

If the cause of the exit was an exception or a fault and was not resolved, the event will be triggered again if the enclave is re-entered using ERESUME. For example, if an enclave performs a divide by 0 operation, executing ERESUME will cause the enclave to attempt to re-execute the faulting instruction and result in another divide by 0 exception. In order to handle an exception that occurred inside the enclave, software can enter the enclave at a different location and invoke the exception handler within the enclave by executing the EENTER instruction. The exception handler within the enclave can attempt to resolve the faulting condition or simply return and indicate to software that the enclave should be terminated (e.g. using EXIT T).

3.2.3.1 ERESUME Interaction

ERESUME restores registers depending on the mode of the enclave (32 or 64 bits).

- In 32-bit mode (IA32_EFER.LMA = 0 || CS.L = 0), the low 32-bits of the legacy registers (EAX, EBX, ECX, EDX, ESP, EBP, ESI, EDI, EIP and EFLAGS) are restored from the thread's CPR area of the current SSA frame. Neither the upper 32-bits of the legacy registers nor the 64-bit registers (R8 ... R15) are loaded.
- In 64-bit mode (IA32_EFER.LMA = 1 && CS.L = 1), all 64 bits of the general processor registers (RAX, RBX, RCX, RDX, RSP, RBP, RSI, RDI, R8 ... R15, RIP and RFLAGS) are loaded.

Extended features specified by SECS.ATTRIBUTES_XFRM are restored from the XSAVE area of the current SSA frame. The layout of the x87 area depends on the current values of IA32_EFER.LMA and CS.L:

- IA32_EFER.LMA = 0 || CS.L = 0
 - 32-bit load in the same format that XSAVE/FXSAVE uses with these values.
- IA32_EFER.LMA = 1 && CS.L = 1
 - 64-bit load in the same format that XSAVE/FXSAVE uses with these values plus REX.W = 1

3.3 CALLING ENCLAVE PROCEDURES

3.3.1 Calling Convention

In standard call conventions subroutine parameters are generally pushed onto the stack. The called routine, being aware of its own stack layout, knows how to find parameters based on compile-time-computable offsets from the SP or BP register (depending on runtime conventions used by the compiler).

Because of the stack switch when calling an enclave, stack-located parameters cannot be found in this manner. Entering the enclave requires a modified parameter passing convention.

For example, the caller might push parameters onto the untrusted stack and then pass a pointer to those parameters in RAX to the enclave software. The exact choice of calling conventions is up to the writer of the edge routines; be those routines hand-coded or compiler generated.

3.3.2 Register Preservation

As with most systems, it is the responsibility of the callee to preserve all registers except that used for returning a value. This is consistent with conventional usage and tends to optimize the number of register save/restore operations.

Biggest challenge for Qubes OS is
how to do *desktop integration* (seamless UX)
without compromising isolation!



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Example #7: Almighty admins?

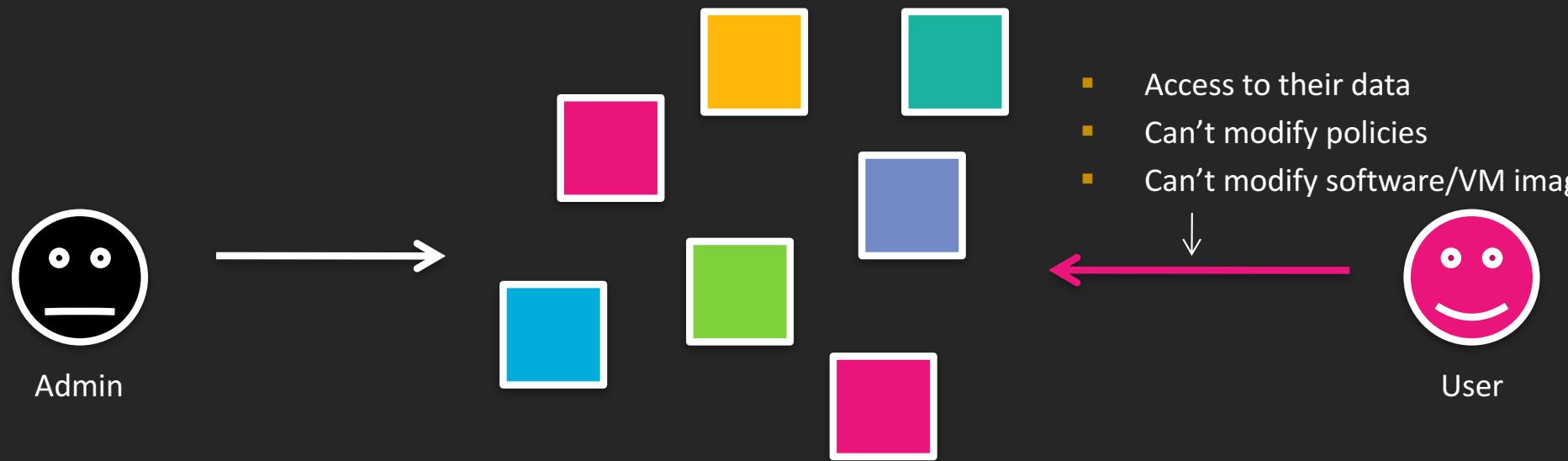


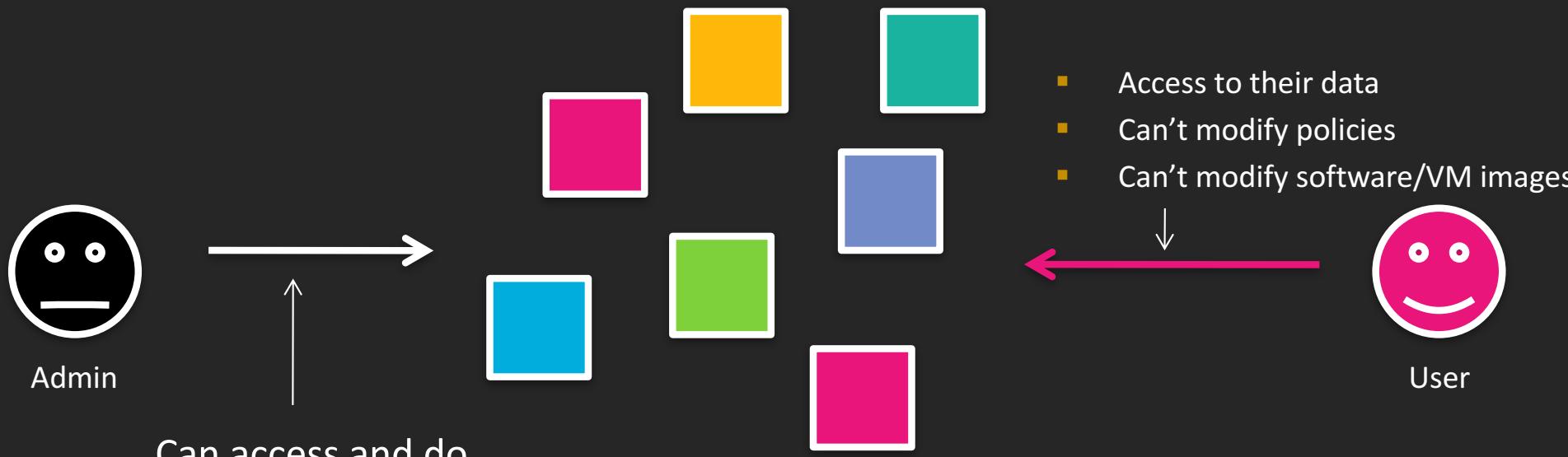
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Admins can steal all our data :(



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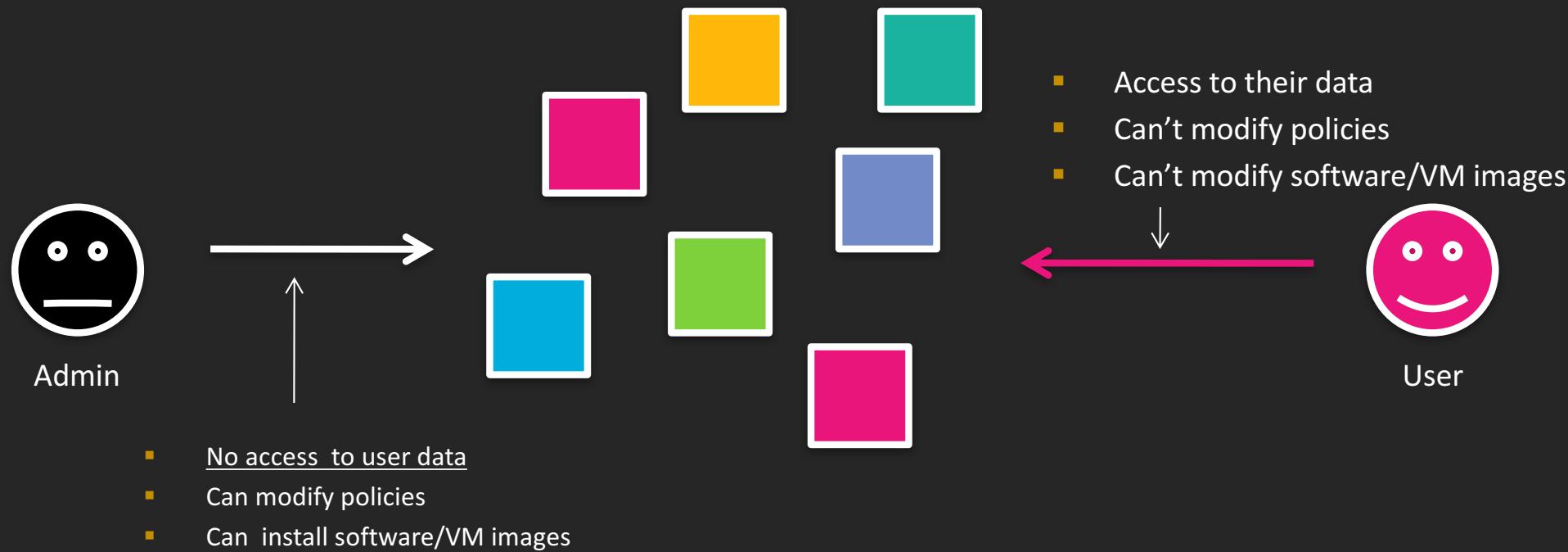


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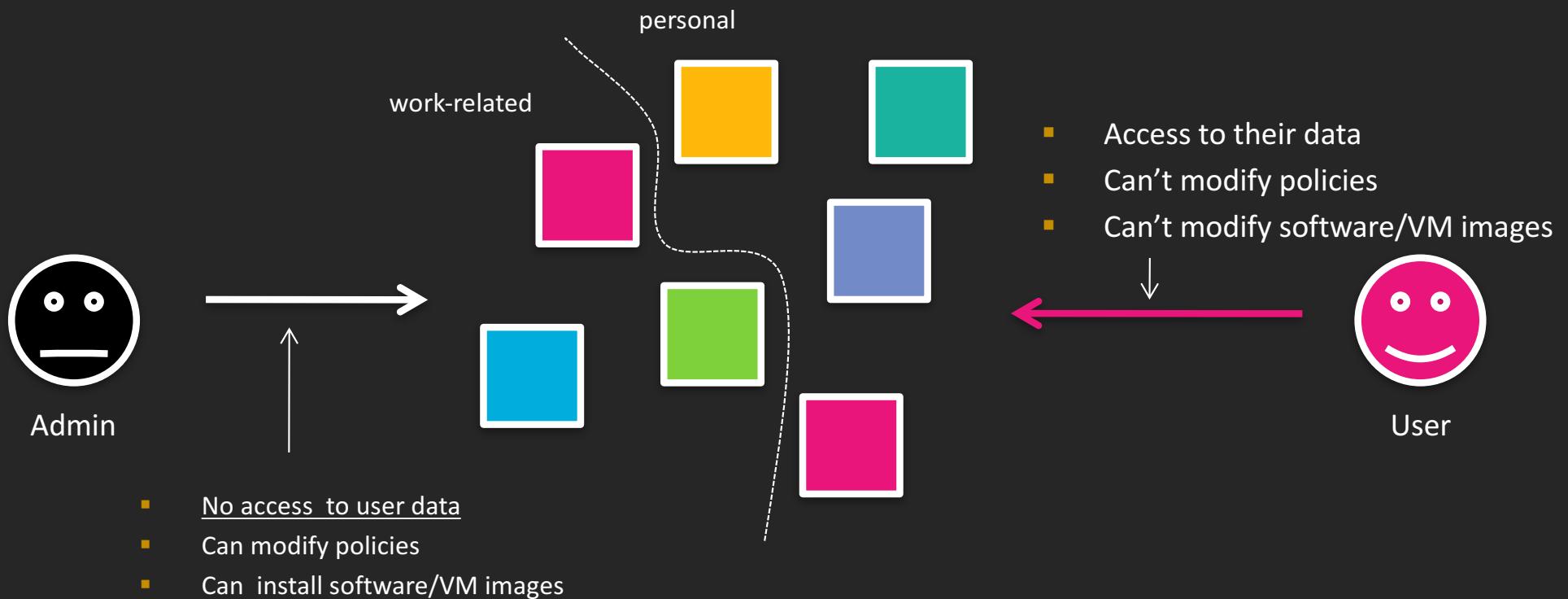
Hmm...



What we want instead:



What we want instead:



Check our Qubes OS new Admin API for implementation details



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Occasionally mishaps happen still...

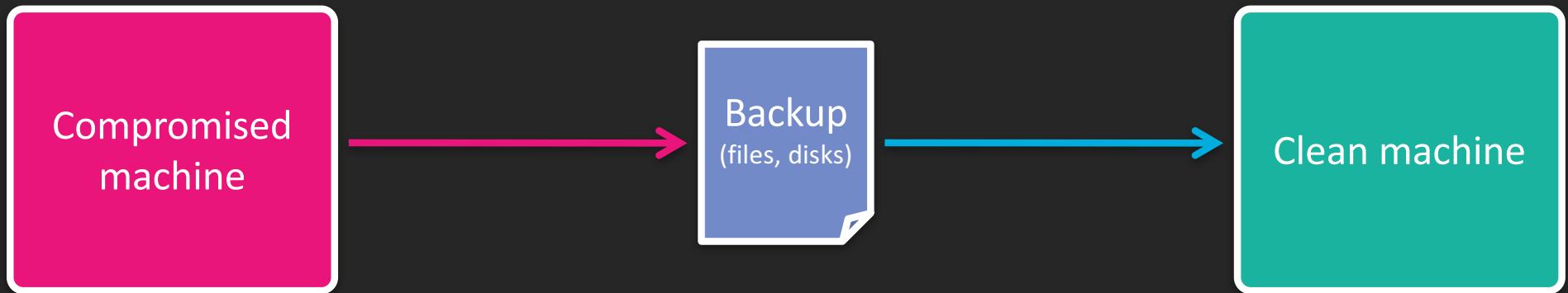


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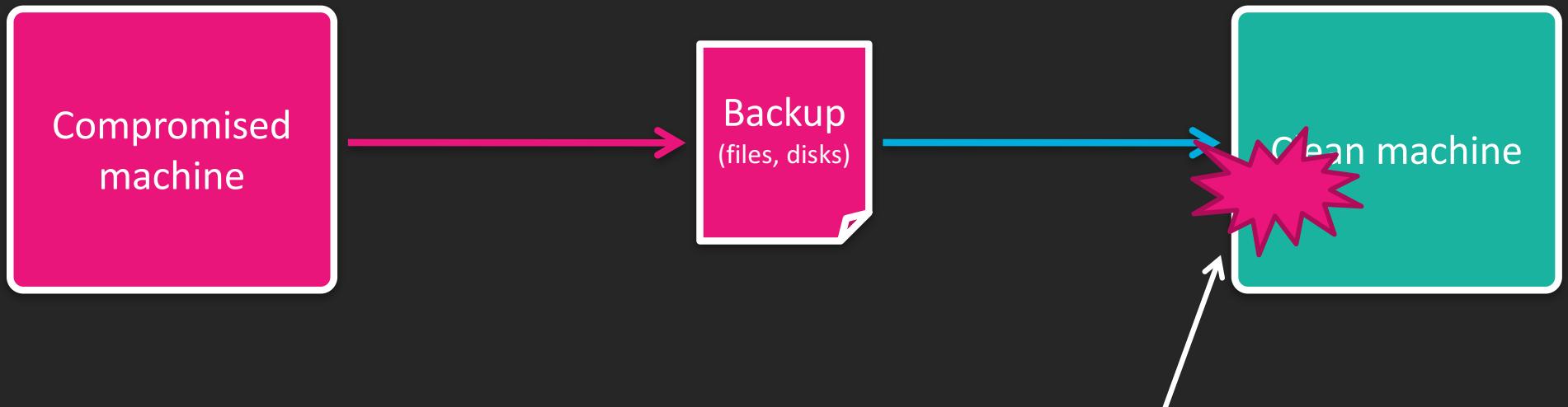
Example #8: Plan B



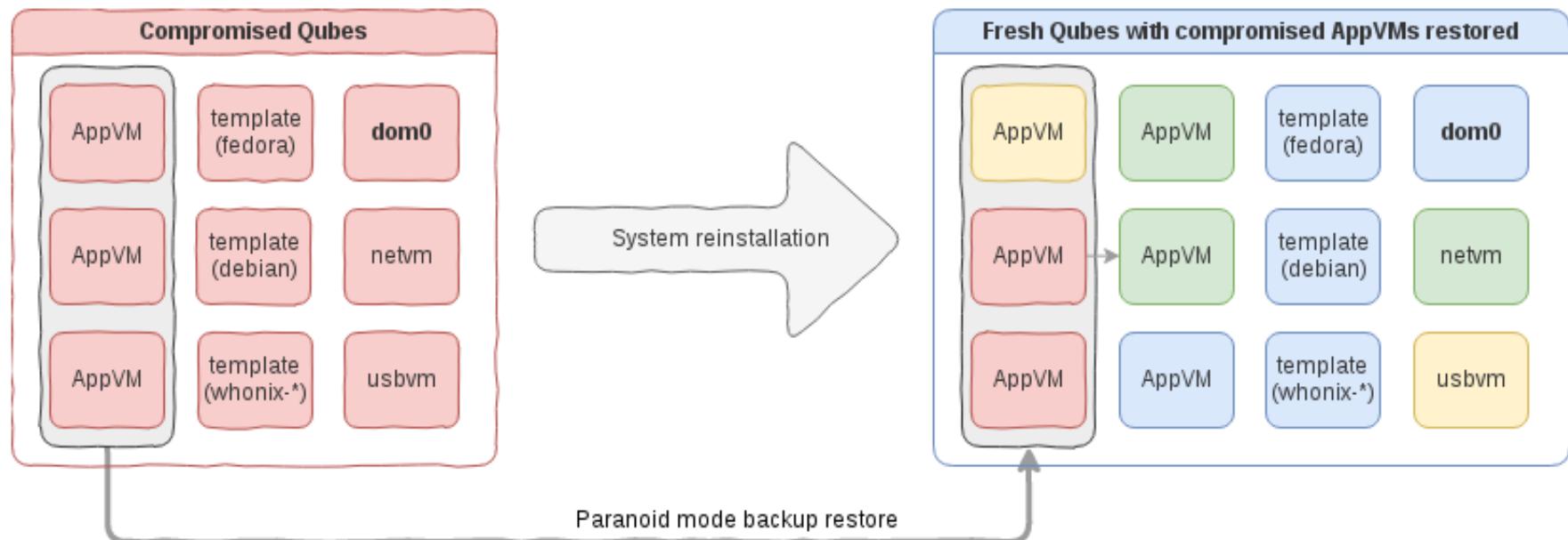
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Qubes (Paranoid) Backup Restore



Security through Distrusting



Division of Duty

- Mics (#1)
- Stateless laptop (#2)
- Multi signatures (#3/4)
- Tunneling (#5)
- Qubes (#6/7)
- Tunneling (#5)
- Qubes Backup Restore (#8)



Compartmentalization



Plan B having

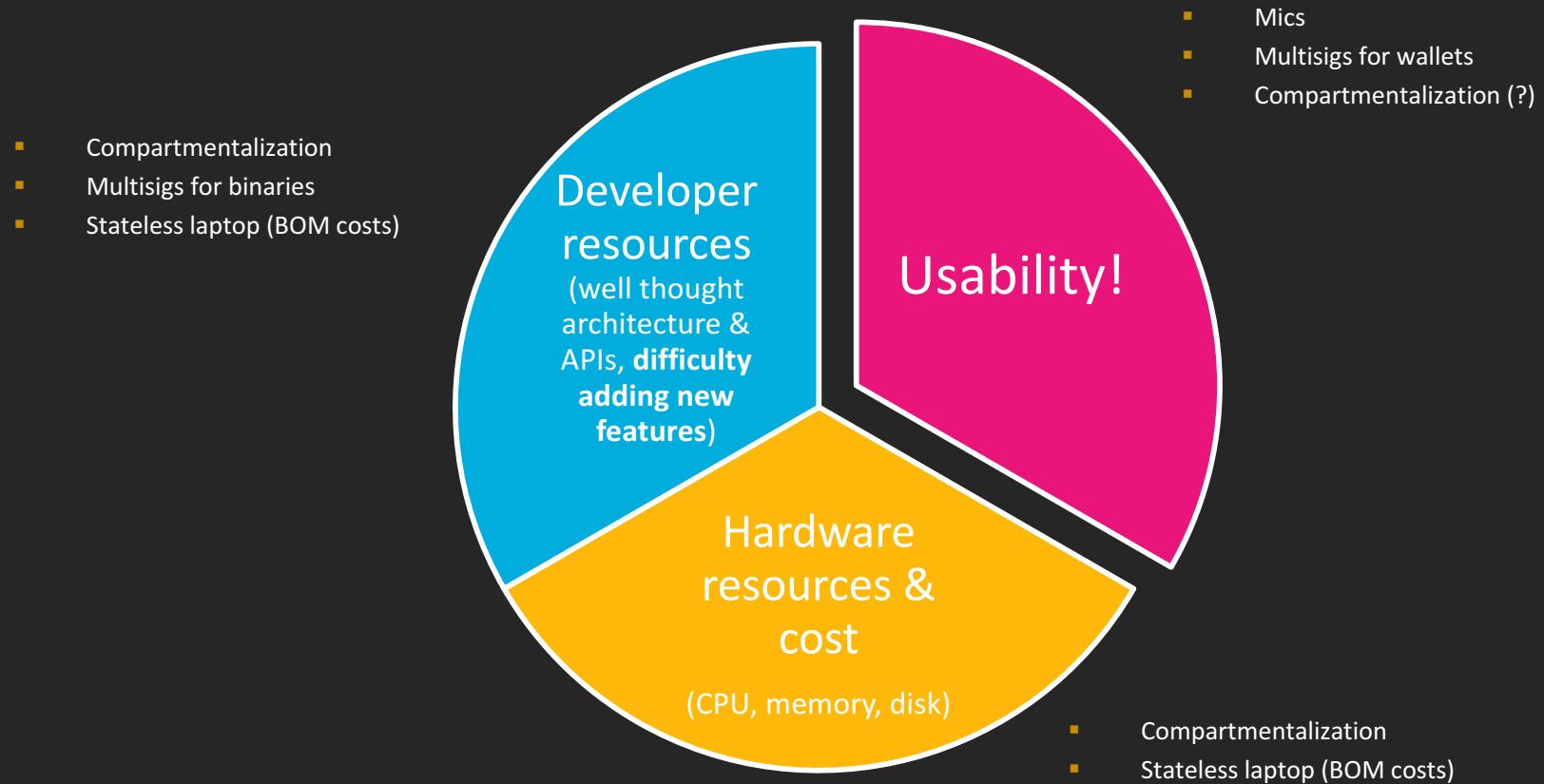
- Qubes Backup Restore (#8)



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Tradeoffs?





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Thanks!

<https://qubes-os.org>

<https://invisiblethingslab.com>

<https://blog.invisiblethings.org>

<https://github.com/rootkovska>

@QubesOS // Twitter for Qubes OS

@rootkovska // Personal Twitter

427F11FD0FAA4B080123F01CDDFA1A3E36879494 // Qubes Master Key

ED727C306E766BC85E621AA65FA6C3E4D9AFBB99 // Personal Master Key

This presentation made in MS Office on Qubes OS 3.2

