

# EC 504 – Spring 2023 – Homework 1

Due Wednesday Feb 8, 2023 at 8 PM Boston Time

**NOTE: This HW is to be turned in to Gradescope as a PDF.**

Start reading Chapters 1, 2, 3 and 4 in CLRS . They give a very readable introduction to Algorithms.. Also glance at Appendix A for math tricks which we will use from time to time.

1. (40 pts) In CLRS do Exercise 3.1-4 on page 53, Problem 3.2-7 on page 60, Problem 3-2 on page 61 , Problem 4.3-9 on page 88.
2. (30 pts) Place the following functions in order from asymptotically smallest to largest. As a convenience you may use  $f(n) < O(g(n))$  to mean  $f(n) \in O(g(n))$  and  $f(n) = g(n)$  to mean  $f(n) \in \Theta(g(n))$ . Please use  $=$  when you are sure that it is  $\Theta(g(n))$ . )  
 $n^2 + 3n \log(n) + 5$  ,  $n^2 + n^{-2}$  ,  $n^{n^2} + n!$  ,  $n^{\frac{1}{n}}$  ,  $n^{n^2-1}$  ,  $\ln n$  ,  $\ln(\ln n)$  ,  $3^{\ln n}$  ,  $2^n$  ,  
 $(1+n)^n$  ,  $n^{1+\cos n}$  ,  $\sum_{k=1}^{\log n} \frac{n^2}{2^k}$  ,  $1$  ,  $n^2 + 3n + 5$  ,  $\log(n!)$  ,  $\sum_{k=1}^n \frac{1}{k}$  ,  $\prod_{k=1}^n (1 - \frac{1}{k^2})$  ,  $(1 - 1/n)^n$

Giving the the algebra and explanation for the tricky cases can get some extra credit (even if you get it wrong!). Don't have to be perfect to get a good score.

1. (40 pts)
  - (a) Given the equation,  $T(n) = 2T(n/2) + n$ , guess a solution of the form:

$$T(n) = c_1 n + c_2 n \log_2(n) .$$

Find the coefficients  $c_1, c_2$  to determine the exact solution assuming a value  $T(1)$  at the bottom the recursion.

- (b) Generalize this to the case to the equation  $T(n) = aT(n/b) + n^k$  and guess the solution of the form:

$$T(n) = c_1 n^\gamma + c_2 n^k$$

using  $b^\gamma = a$  and assuming  $\gamma \neq k$ . First show if you drop the  $n^k$  using the **homogeneous** equation  $T(n) = aT(n/b)$  the form  $c_1 n^\gamma$  is a solution! (What is  $c_1$ ?) Second drop  $aT(n/b)$  and show  $c_2 n^k$  is a solution (What is  $c_2$ ?) With both terms (and  $\gamma \neq k$ ) the full solution is just the sum of the two terms but only one or the other dominates!

- (c) What happens when the two solution collide (i.e have the same power, i.e  $\gamma = \log(a)/\log(b) = k$ .) Now show that the leading solution is as  $n$  goes to infinity is  $T(n) = \Theta(n^k \log n)$  <sup>1</sup>

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<sup>1</sup>If you are ambitious you find exact solutions for part b and c above are of the form  $T(n) = c_1 n^\gamma + c_2 n^k$  and  $T(n) = c_1 n^k + c_2 n \log(n) n^k$  explicitly determine the  $c$ 's for each case respectively. We did that in part for example. BUT note we already know the leading terms for the Master Equation without this extra effort! **Neat trick!**