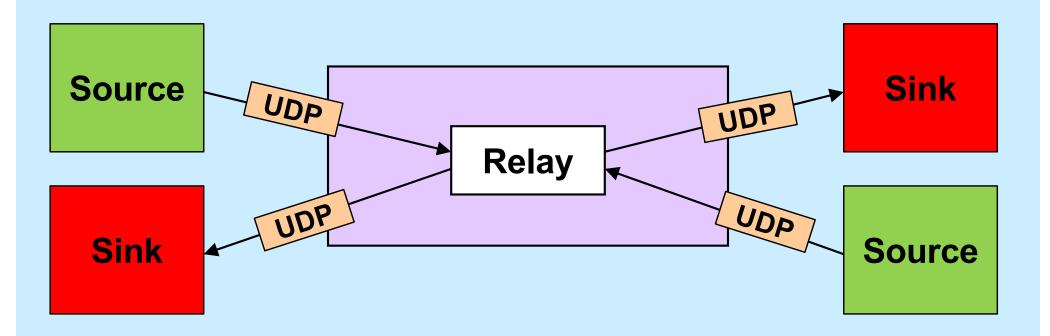
**CS 33** 

**Event-Based Programming** 

# **Stream Relay**



### Solution?

```
while(...) {
    size = read(left, buf, sizeof(buf));
    write(right, buf, size);
    size = read(right, buf, sizeof(buf));
    write(left, buf, size);
}
```

# **Select System Call**

# Relay Sketch

```
void relay(int left, int right) {
   fd set rd, wr;
   int maxFD = max(left, right) + 1;
   FD ZERO(&rd); FD SET(left, &rd); FD SET(right, &rd);
   FD ZERO(&wr); FD SET(left, &wr); FD SET(right, &wr);
   while (1) {
      select(maxFD, &rd, &wr, 0, 0);
      if (FD ISSET(left, &rd))
         read(left, bufLR, sizeof(message t));
      if (FD ISSET(right, &rd))
         read(right, bufRL, sizeof(message t));
      if (FD ISSET(right, &wr))
         write(right, bufLR, sizeof(message t));
      if (FD ISSET(left, &rd))
         write(left, bufRL, sizeof(message t));
```

# Relay (1)

```
void relay(int left, int right) {
  fd_set rd, wr;
  int left_read = 1, right_write = 0;
  int right_read = 1, left_write = 0;
  message_t bufLR;
  message_t bufRL;
  int maxFD = max(left, right) + 1;
```

# Relay (2)

```
while(1) {
  FD ZERO (&rd);
 FD ZERO(&wr);
  if (left read)
    FD SET(left, &rd);
  if (right read)
    FD SET (right, &rd);
  if (left write)
    FD SET(left, &wr);
  if (right write)
    FD SET(right, &wr);
  select(maxFD, &rd, &wr, 0, 0);
```

# Relay (3)

```
if (FD_ISSET(left, &rd)) {
    read(left, bufLR, sizeof(message_t));
    left_read = 0;
    right_write = 1;
}
if (FD_ISSET(right, &rd)) {
    read(right, bufRL, sizeof(message_t));
    right_read = 0;
    left_write = 1;
}
```

# Relay (4)

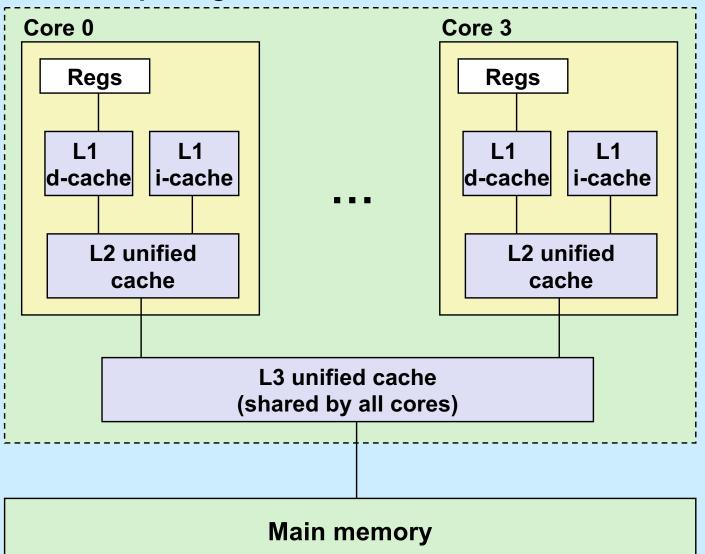
```
if (FD_ISSET(right, &wr)) {
    write(right, bufLR, sizeof(message_t));
    left_read = 1;
    right_write = 0;
}
if (FD_ISSET(left, &wr)) {
    write(left, bufRL, sizeof(message_t));
    right_read = 1;
    left_write = 0;
}
return 0;
}
```

**CS 33** 

# Caching and Program Optimization

# Intel Core i5 and i7 Cache Hierarchy

#### **Processor package**



#### L1 i-cache and d-cache:

32 KB, 8-way, Access: 4 cycles

#### L2 unified cache:

256 KB, 8-way, Access: 11 cycles

#### L3 unified cache:

8 MB, 16-way, Access: 30-40 cycles

7.00033. 30 10 cycles

Block size: 64 bytes for

all caches

# **Accessing Memory**

- Program references memory (load)
  - if not in cache (cache miss), data is requested from RAM
    - » fetched in units of 64 bytes
      - aligned to 64-byte boundaries (low-order 6 bits of address are zeroes)
    - » if memory accessed sequentially, data is pre-fetched
    - » data stored in cache (in 64-byte cache lines)
      - stays there until space must be re-used (least recently used is kicked out first)
  - if in cache (cache hit) no access to RAM needed
- Program modifies memory (store)
  - data modified in cache
  - eventually written to RAM in 64-byte units

### Quiz 1

The previous slide said that 64 bytes of memory from contiguous locations are transferred at a time. Suppose we have memory that transfers 128 contiguous bytes in the same amount of time. If we have a program that reads memory one byte at a time from random (but valid) memory locations, how much faster will it run with the new memory system than with the old?

- a) half as fast
- b) roughly the same speed
- c) twice as fast
- d) four times as fast

# **Layout of C Matrices in Memory**

- C matrices allocated in row-major order
  - each row in contiguous memory locations
- Stepping through columns in one row:

```
- for (i = 0; i < n; i++)
sum += a[0][i];
```

- accesses successive elements
- data fetched from RAM in 64-byte units
- Stepping through rows in one column:

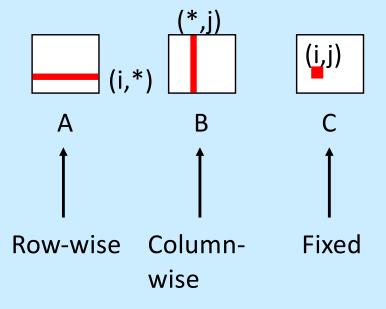
```
- for (i = 0; i < n; i++)
sum += a[i][0];
```

- accesses distant elements
- if array element is 8 bytes, 56 bytes (out of 64) are not used
  - » effective throughput reduced by factor of 8

# **Matrix Multiplication (ijk)**

```
/* ijk */
for (i=0; i<n; i++) {
  for (j=0; j<n; j++) {
    sum = 0.0;
    for (k=0; k<n; k++)
        sum += a[i][k] * b[k][j];
    c[i][j] = sum;
    }
}</pre>
```

#### Inner loop:

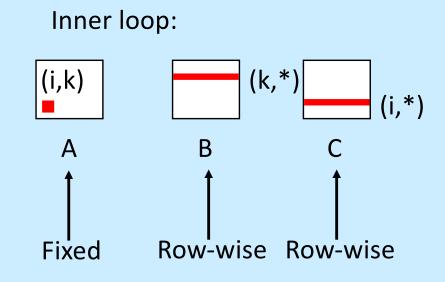


### Misses per inner loop iteration:

<u>A</u> <u>B</u> <u>C</u> 0.125 1.0 0.0

# **Matrix Multiplication (kij)**

```
/* kij */
for (k=0; k<n; k++) {
  for (i=0; i<n; i++) {
    r = a[i][k];
    for (j=0; j<n; j++)
        c[i][j] += r * b[k][j];
  }
}</pre>
```

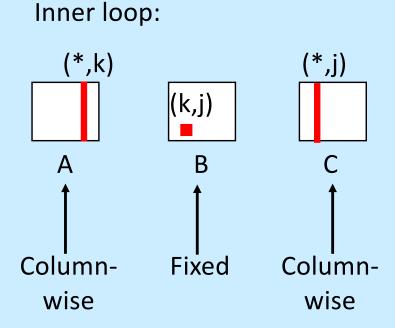


#### Misses per inner loop iteration:

<u>A</u> <u>B</u> <u>C</u> 0.0 0.125 0.125

# **Matrix Multiplication (jki)**

```
/* jki */
for (j=0; j<n; j++) {
  for (k=0; k<n; k++) {
    r = b[k][j];
    for (i=0; i<n; i++)
        c[i][j] += a[i][k] * r;
  }
}</pre>
```



#### Misses per inner loop iteration:

<u>A</u> <u>B</u> <u>C</u> 1.0 0.0 1.0

## **Summary of Matrix Multiplication**

```
for (i=0; i<n; i++)
  for (j=0; j<n; j++) {
    sum = 0.0;
  for (k=0; k<n; k++)
    sum += a[i][k] * b[k][j];
  c[i][j] = sum;
}</pre>
```

```
for (k=0; k<n; k++)
for (i=0; i<n; i++) {
  r = a[i][k];
  for (j=0; j<n; j++)
    c[i][j] += r * b[k][j];
}</pre>
```

```
for (j=0; j<n; j++)
  for (k=0; k<n; k++) {
    r = b[k][j];
    for (i=0; i<n; i++)
      c[i][j] += a[i][k] * r;
}</pre>
```

#### ijk (& jik):

- 2 loads, 0 stores
- misses/iter = **1.125**

#### kij (& ikj):

- 2 loads, 1 store
- misses/iter = **0.25**

#### jki (& kji):

- 2 loads, 1 store
- misses/iter = **2.0**

### In Real Life ...

 Multiply two 1024x1024 matrices of doubles on sunlab machines

```
ijk» 4.185 seconds
```

```
kij» 0.798 seconds
```

```
jki» 11.488 seconds
```

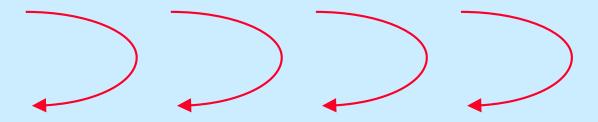
**CS 33** 

Multithreaded Programming I

# **Multithreaded Programming**

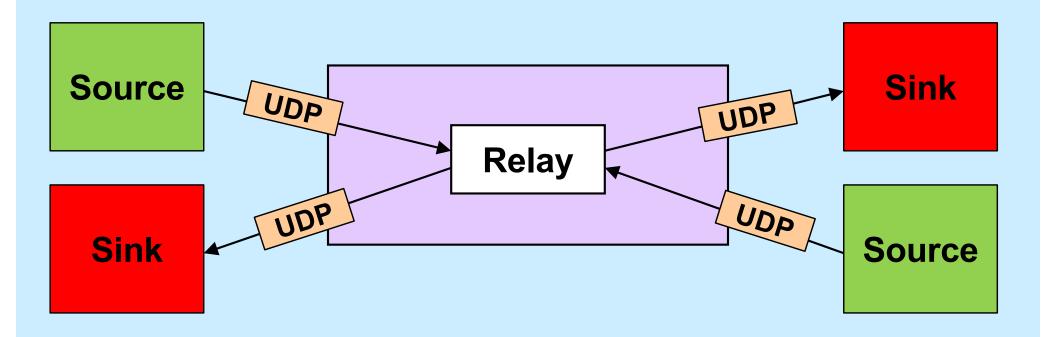
- A thread is a virtual processor
  - an independent agent executing instructions
- Multiple threads
  - multiple independent agents executing instructions

# Why Threads?



- Many things are easier to do with threads
- Many things run faster with threads

# A Simple Example

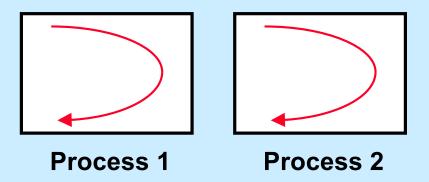


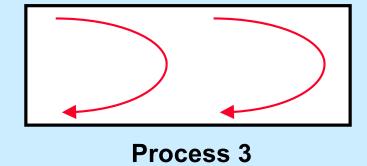
### **Life With Threads**

```
void copy(int source, int destination) {
   struct args *targs = args;
   char buf[BSIZE];

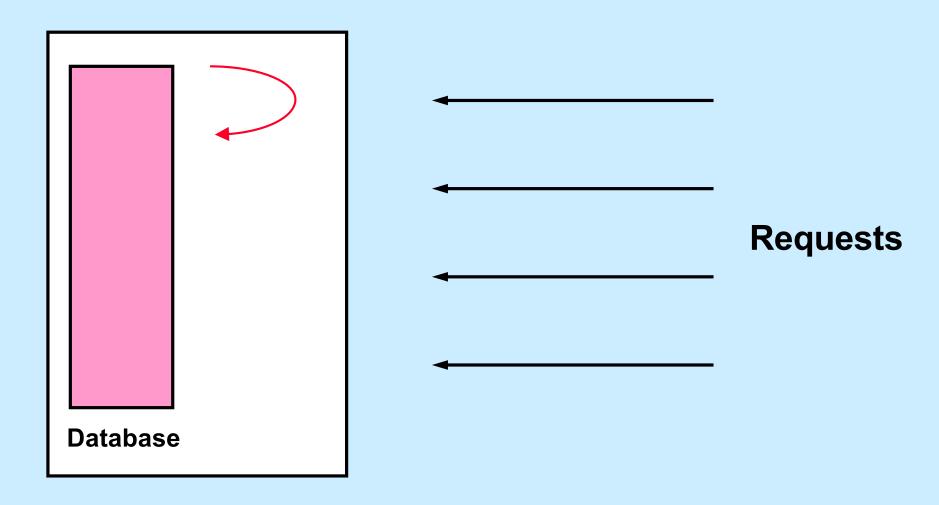
while(1) {
    int len = read(source, buf, BSIZE);
    write(destination, buf, len);
  }
}
```

### **Processes vs. Threads**

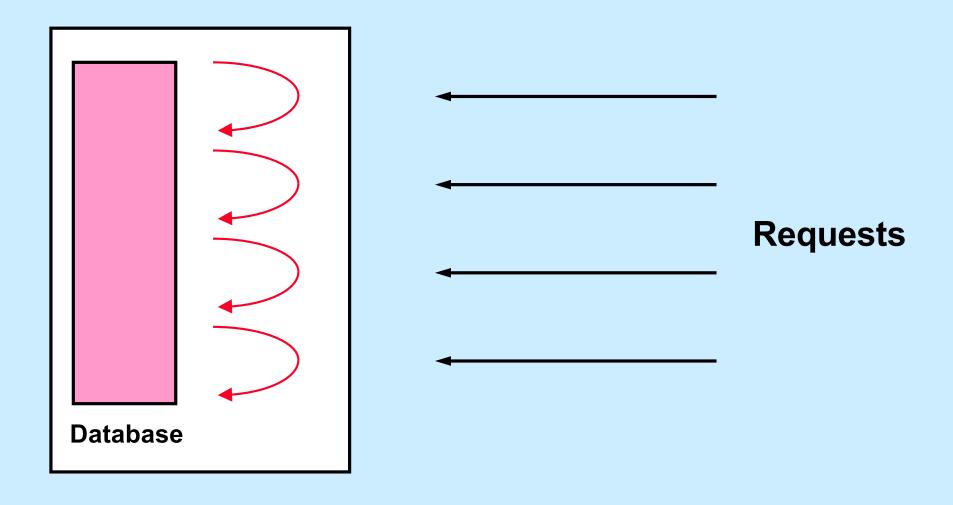




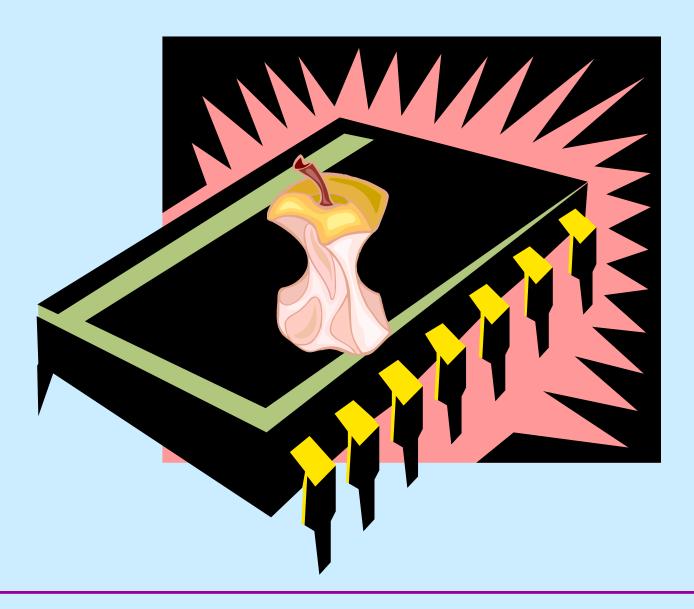
# Single-Threaded Database Server



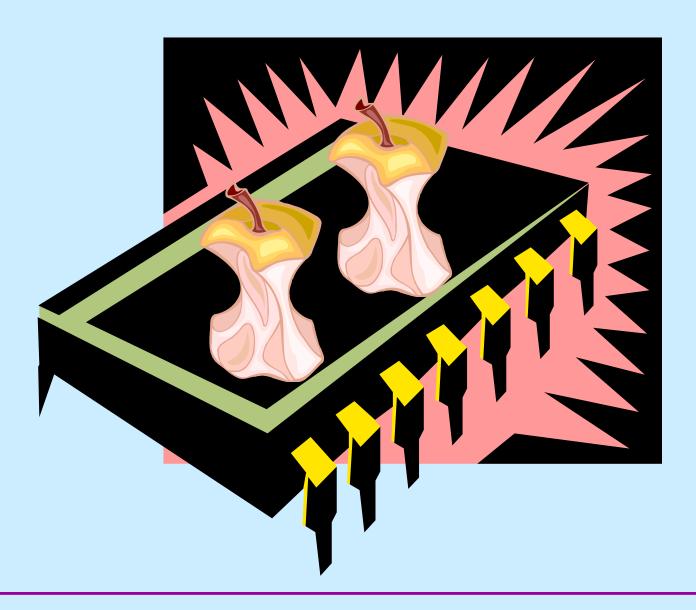
### **Multithreaded Database Server**



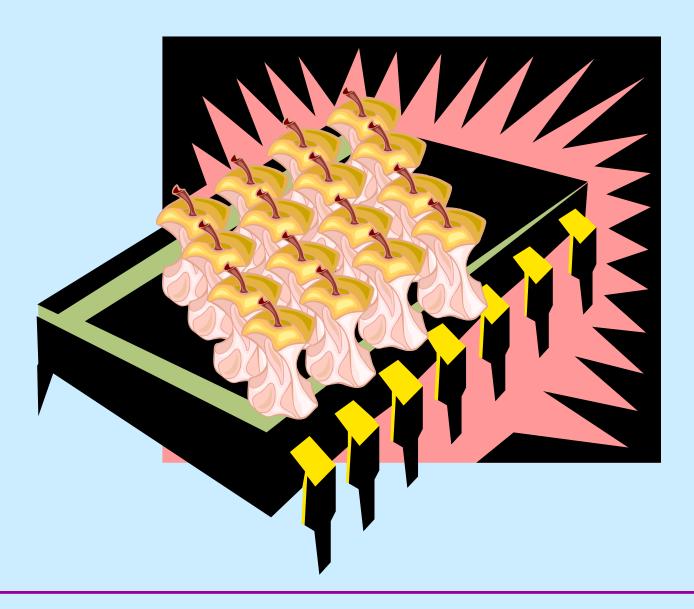
# **Single-Core Chips**



# **Dual-Core Chips**



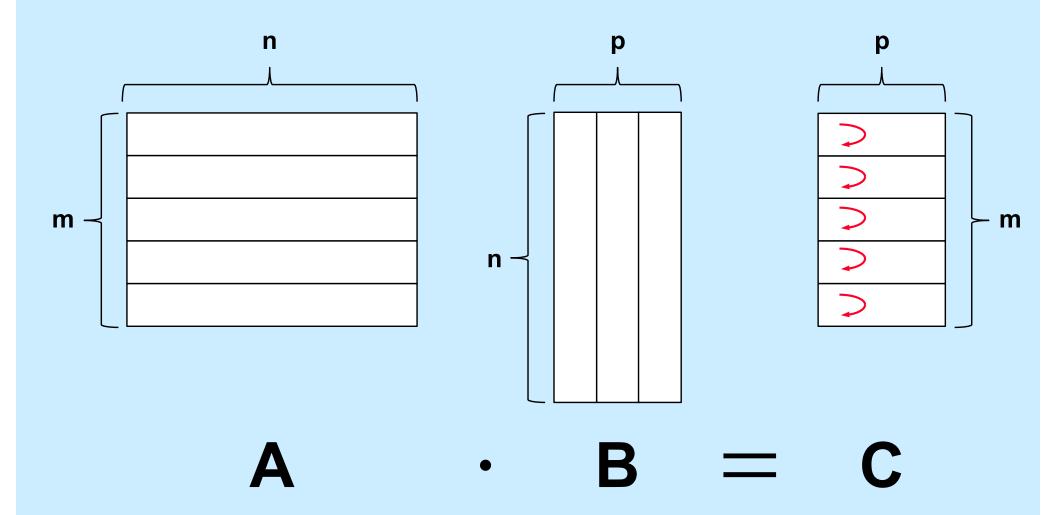
# **Multi-Core Chips**



### **Good News/Bad News**

- © Good news
  - multi-threaded programs can take advantage of multi-core chips (single-threaded programs cannot)
- **Bad news** 
  - it's not easy
    - » must have parallel algorithm
      - employing at least as many threads as processors
      - threads must keep processors busy
        - doing useful work

# **Matrix Multiplication Revisited**



### **Standards**

• POSIX  $1003.4a \rightarrow 1003.1c \rightarrow 1003.1j$ 

- Microsoft
  - Win32/64

# **Creating Threads**

```
long A[M][N], B[N][P], C[M][P];
 for (i=0; i<M; i++) // create worker threads</pre>
   pthread create(&thr[i], 0, matmult, i);
void *matmult(void *arg) {
  long i = (long) arg;
  // compute row i of the product C of A and B
```

### When Is It Finished?

# Example (1)

```
#include <stdio.h>
#include <pthread.h>
#include <string.h>
#define M 3
#define N 4
#define P 5
long A[M][N];
long B[N][P];
long C[M][P];
void *matmult(void *);
```

```
main() {
  long i;
  pthread_t thr[M];
  int error;

// initialize the matrices
...
```

# Example (2)

```
for (i=0; i<M; i++) { // create worker threads
 if (error = pthread create(
    &thr[i],
    0,
    matmult,
    (void *)i)) {
   fprintf(stderr, "pthread create: %s", strerror(error));
   exit(1);
for (i=0; i<M; i++) // wait for workers to finish their jobs
 pthread join(thr[i], 0)
/* print the results ... */
```

# Example (3)

```
void *matmult(void *arg) {
  long row = (long) arg;
  long col;
  long i;
  long t;
  for (col=0; col < P; col++) {</pre>
   t = 0;
   for (i=0; i<N; i++)
     t += A[row][i] * B[i][col];
   C[row][col] = t;
  return(0);
```