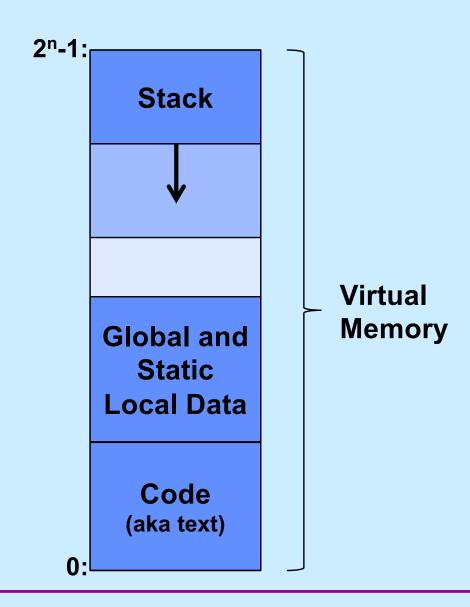
CS 33

Machine Programming (4)

Digression (Again): Where Stuff Is (Roughly)



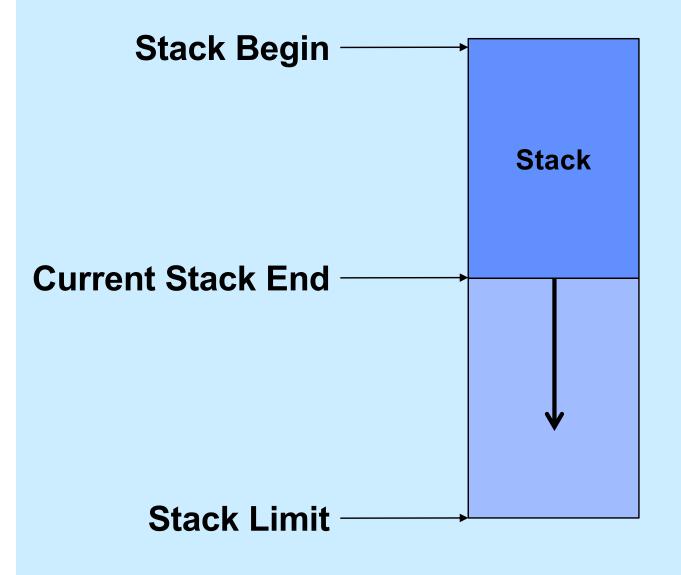
Function Call and Return

- Function A calls function B
- Function B calls function C

... several million instructions later

- C returns
 - how does it know to return to B?
- B returns
 - how does it know to return to A?

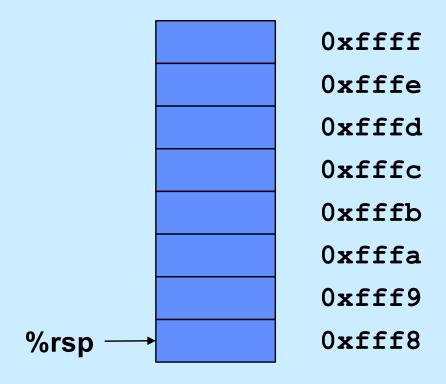
The Runtime Stack



Higher memory addresses

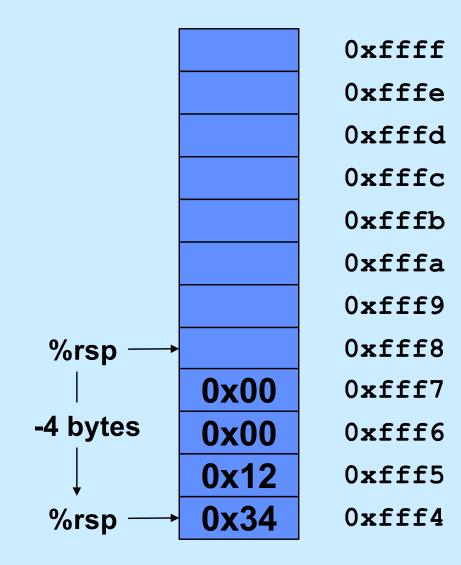
Lower memory addresses

Stack Operations

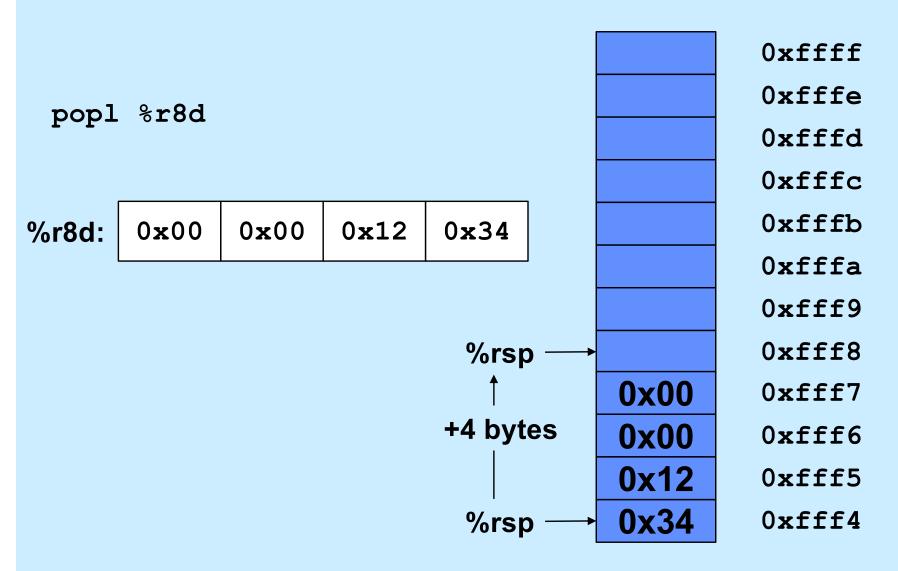


Push

pushl \$0x1234



Pop



Call and Return

0x1000: call func

0x1004: addq \$3, %rax

0x2000: func:

• • • • •

0x2200: movq \$6, %rax

0x2203: ret

0x2000: func:

• • • • • •

0x2200: movq \$6, %rax

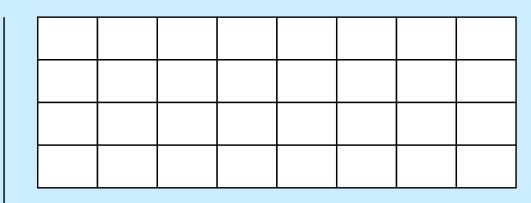
0x2203: ret

→ 0x1000: call func

Call and Return

0x1004: addq \$3, %rax

stack growth



00	00	0	00	0	00	10	00
00	00	00	0f	ff	f1	00	00

%rax

%rip

→ 0x2000: func:

• • • • •

0x2200: movq \$6, %rax

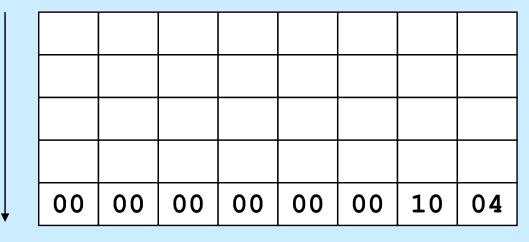
0x2203: ret

0x1000: call func

Call and Return

0x1004: addq \$3, %rax

stack growth



0xffff10018
0xffff10010
0xffff10008
0xffff10000
0xffff0fff8 <</pre>

00	00	00	00	00	00	20	00
00	00	00	0f	ff	f0	ff	f8

%rax

%rip

0x2000: func:

• • • • •

0x2200: movq \$6, %rax

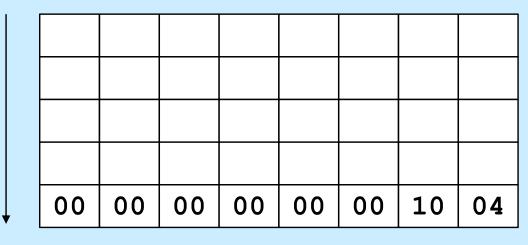
→ 0x2203: ret

0x1000: call func

Call and Return

0x1004: addq \$3, %rax

stack growth



0xffff10018
0xffff10010
0xffff10008
0xffff10000
0xffff0fff8 <</pre>

00	00	00	00	00	00	00	06
00	00	00	00	00	00	22	03
00	00	00	0f	ff	f0	ff	f8

%rax

%rip

0x2000: func:

• • • • •

0x2200: movq \$6, %rax

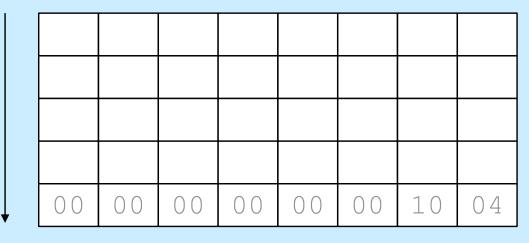
0x2203: ret

0x1000: call func

Call and Return

0x1004: addq \$3, %rax

stack growth



00	00	00	00	00	00	00	06
00	00	0	00	0	00	10	04
00	00	00	0f	ff	f1	00	00

%rax

%rip

Arguments and Local Variables

```
int mainfunc() {
  long array[3] =
          {2,117,-6};
  long sum =
          ASum(array, 3);
    ...
  return sum;
}
```

- Local variables usually allocated on stack
- Arguments to functions pushed onto stack

```
long ASum(long *a,
    unsigned long size) {
    long i, sum = 0;
    for (i=0; i<size; i++)
        sum += a[i];
    return sum;
}</pre>
```

 Local variables may be put in registers (and thus not on stack)

Arguments and Local Variables

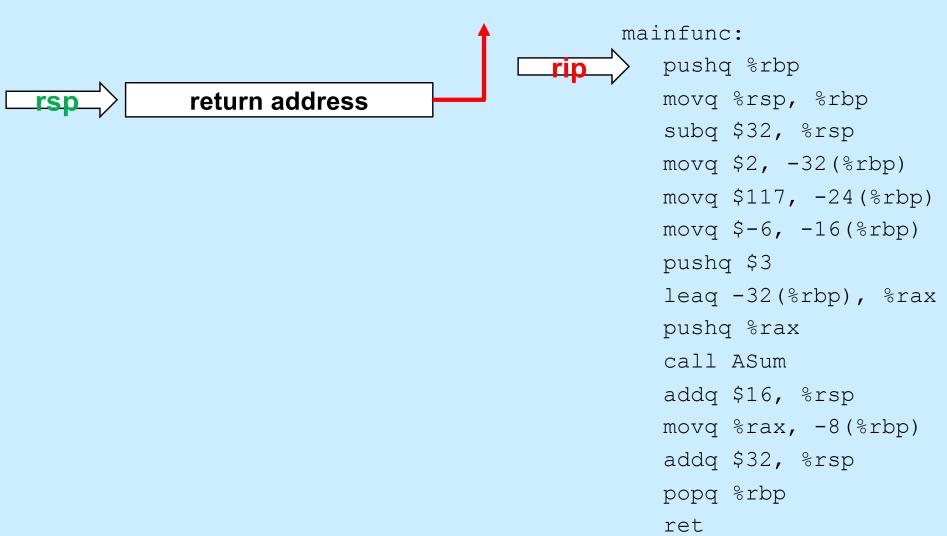
mainfunc:

```
pushq %rbp
                         # save old %rbp
                         # set %rbp to point to stack frame
movq %rsp, %rbp
subq $32, %rsp
                         # alloc. space for locals (array and sum)
movq \$2, -32(\$rbp) # initialize array[0]
movq $117, -24(%rbp) # initialize array[1]
movq \$-6, -16(\$rbp) # initialize array[2]
pusha $3
                         # push arg 2
leaq -32(%rbp), %rax
                    # array address is put in %rax
                         # push arg 1
pushq %rax
call ASum
addq $16, %rsp
               # pop args
movg %rax, -8(%rbp) # copy return value to sum
addq $32, %rsp
                         # pop locals
popq %rbp
                         # pop and restore old %rbp
ret
```

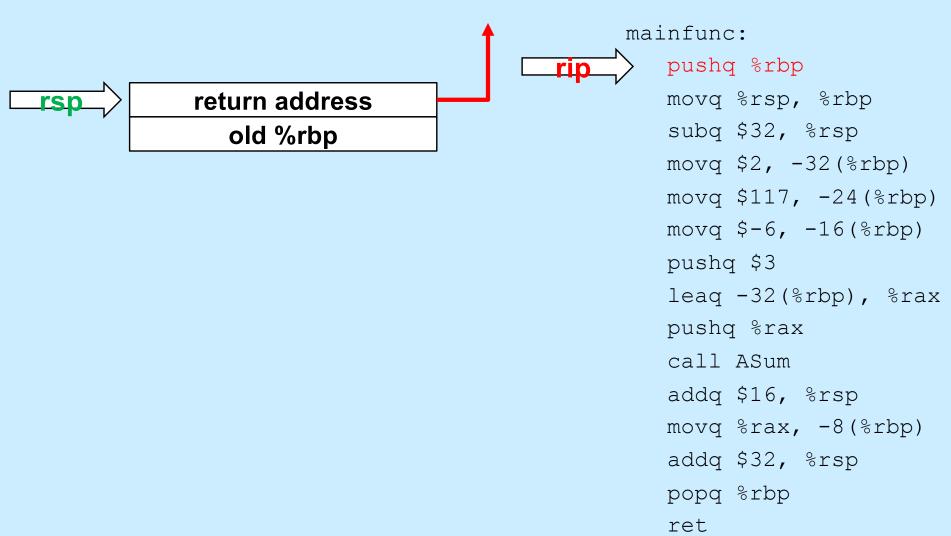
Arguments and Local Variables

```
ASum:
   pushq %rbp
                              # save old %rbp
   movq %rsp, %rbp
                              # set %rbp to point to stack frame
                              # i in %rcx
   movq $0, %rcx
   movq $0, %rax
                              # sum in %rax
   movq 16(%rbp), %rdx
                              # copy arg 1 (array) into %rdx
loop:
   cmpq 24(%rbp), %rcx # i < size?</pre>
   jge done
   addq (%rdx, %rcx, 8), %rax # sum += a[i]
   incq %rcx
                              # i++
   ja loop
done:
                              # pop and restore %rbp
   popq %rbp
   ret
```

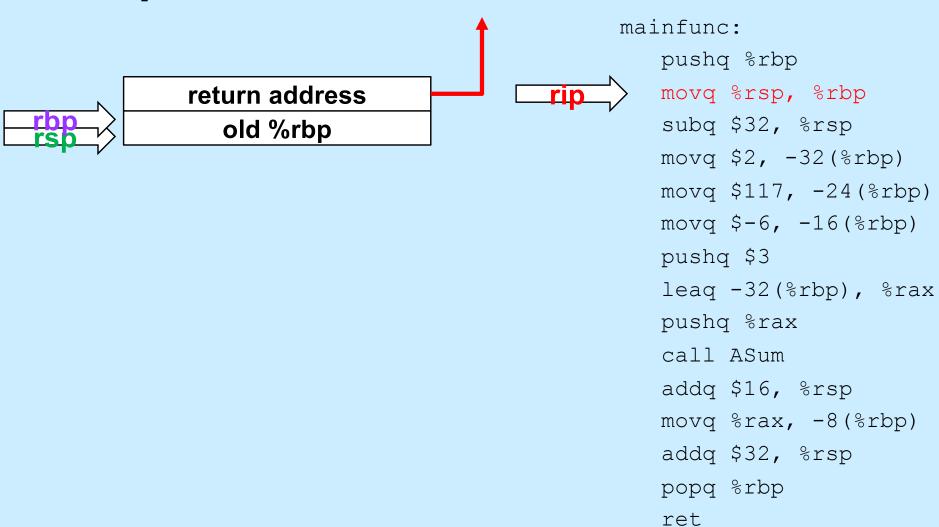
Enter mainfunc



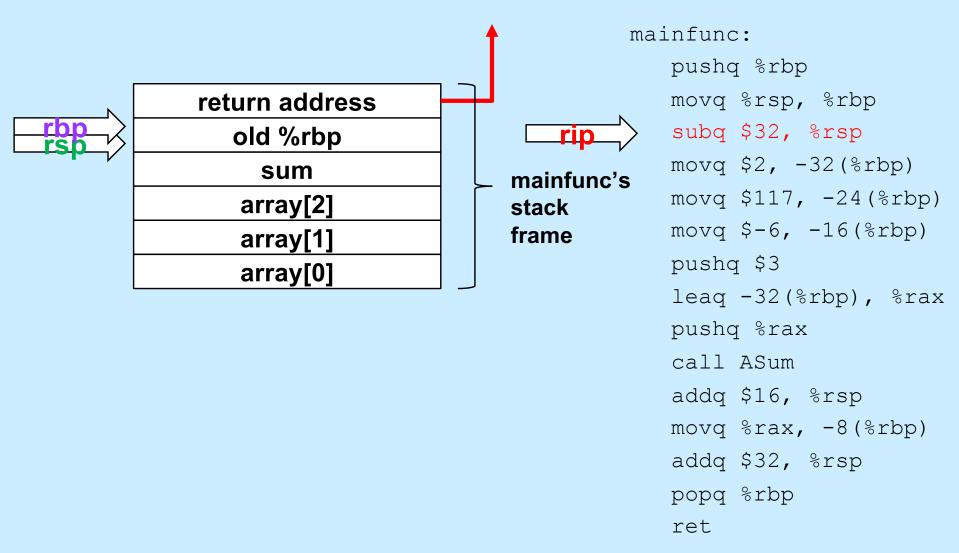
Enter mainfunc



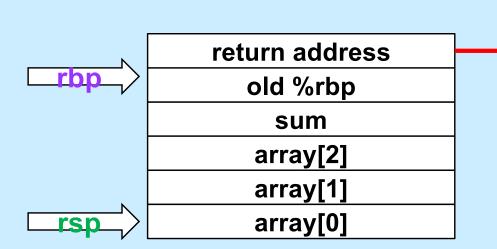
Setup Frame

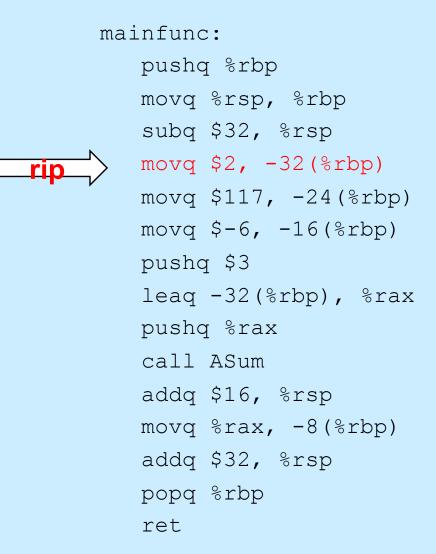


Allocate Local Variables

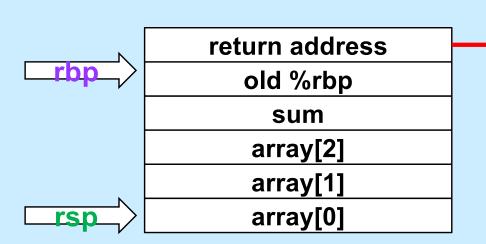


Initialize Local Array



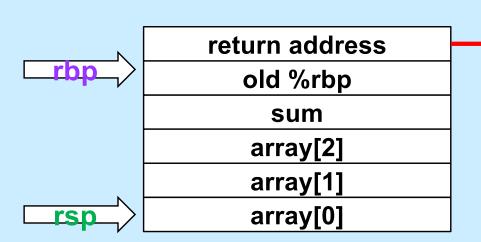


Initialize Local Array



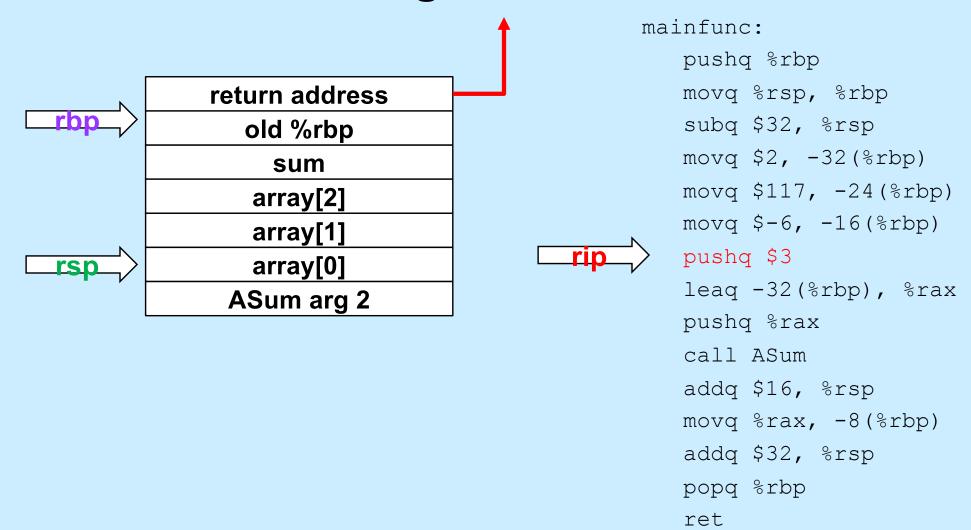
```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movq $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pushq $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Initialize Local Array

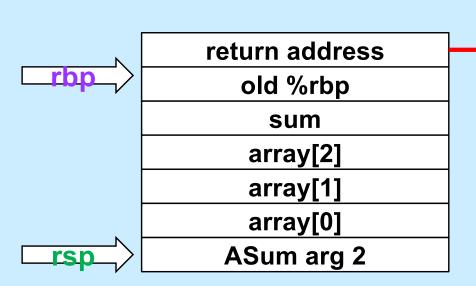


```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movq $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pushq $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Push Second Argument

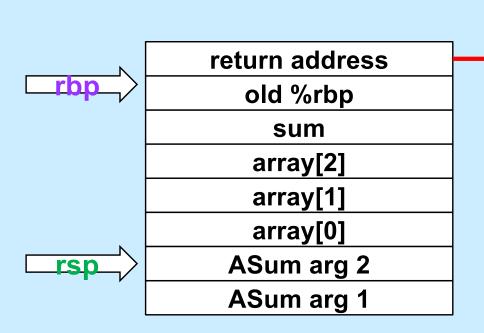


Get Array Address



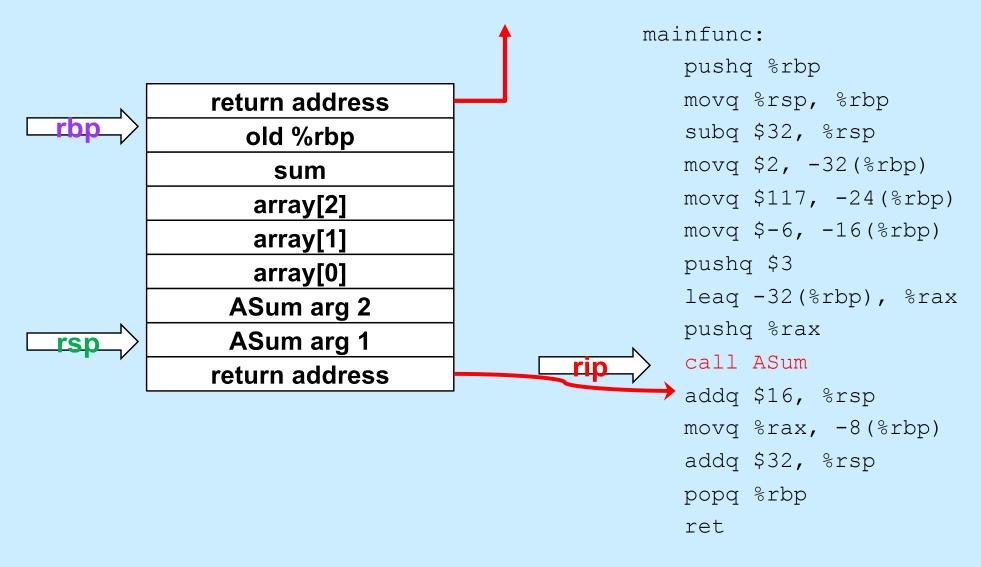
```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq \$-6, -16(%rbp)
   pushq $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Push First Argument

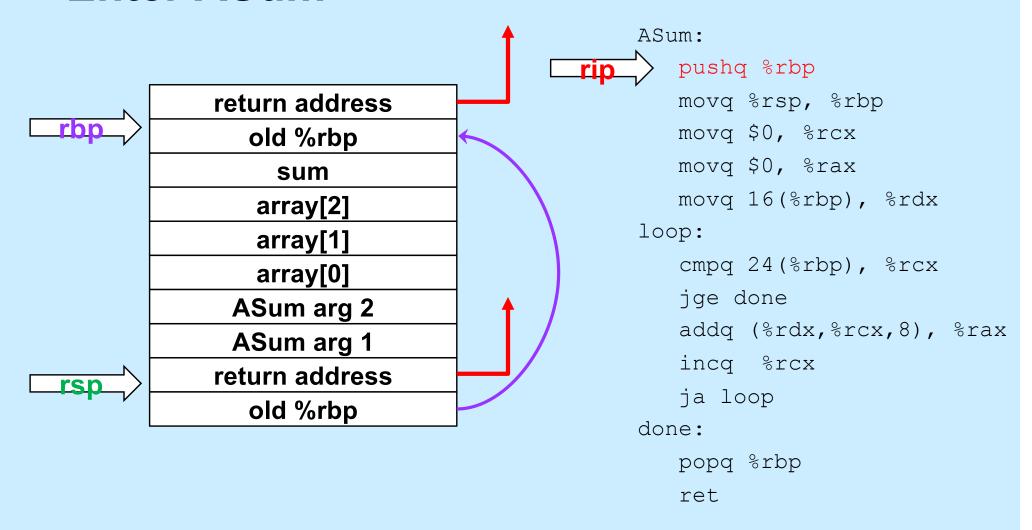


mainfunc: pushq %rbp movq %rsp, %rbp subq \$32, %rsp movq \$2, -32(%rbp) movq \$117, -24(%rbp) movg \$-6, -16(%rbp)pusha \$3 leaq -32(%rbp), %rax pushq %rax call ASum addq \$16, %rsp movq %rax, -8(%rbp) addq \$32, %rsp popq %rbp ret

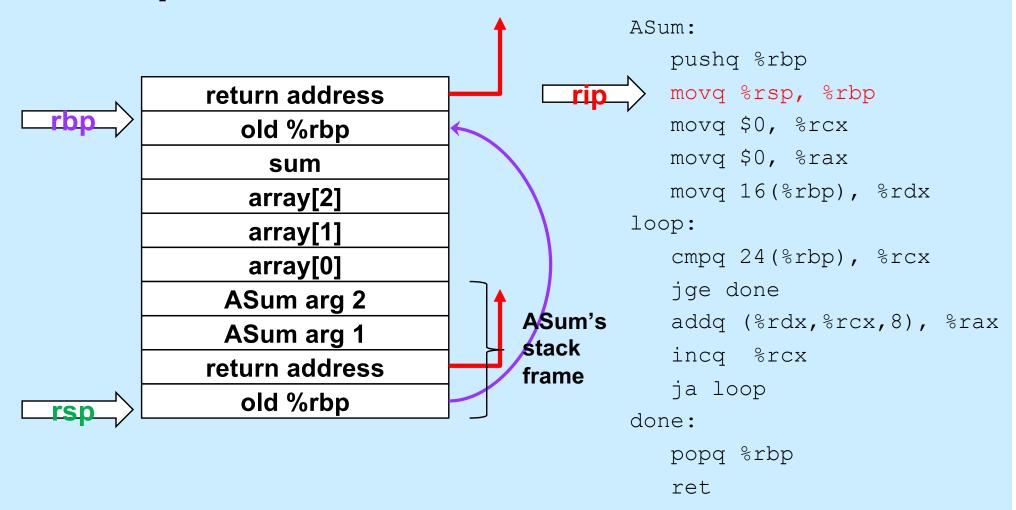
Call ASum



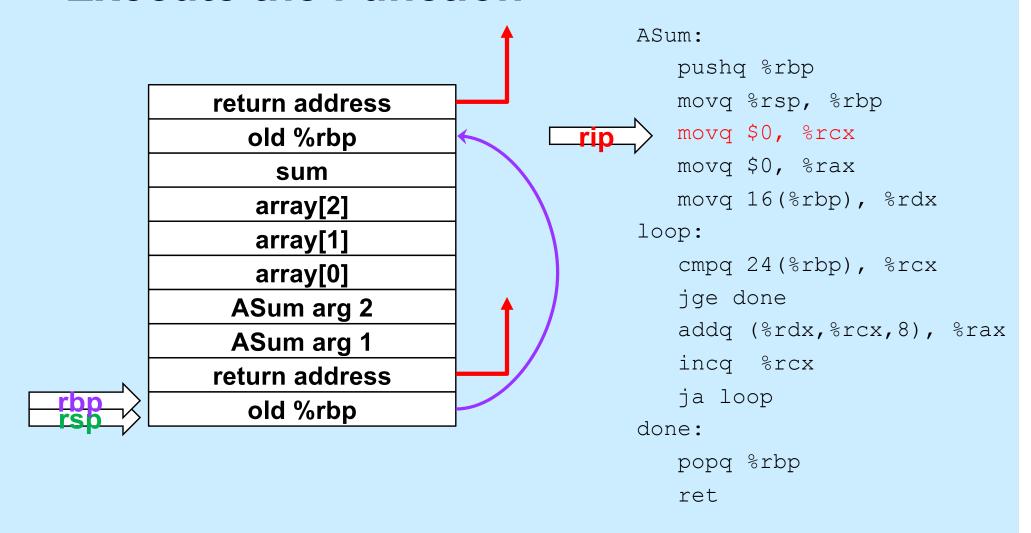
Enter ASum



Setup Frame



Execute the Function



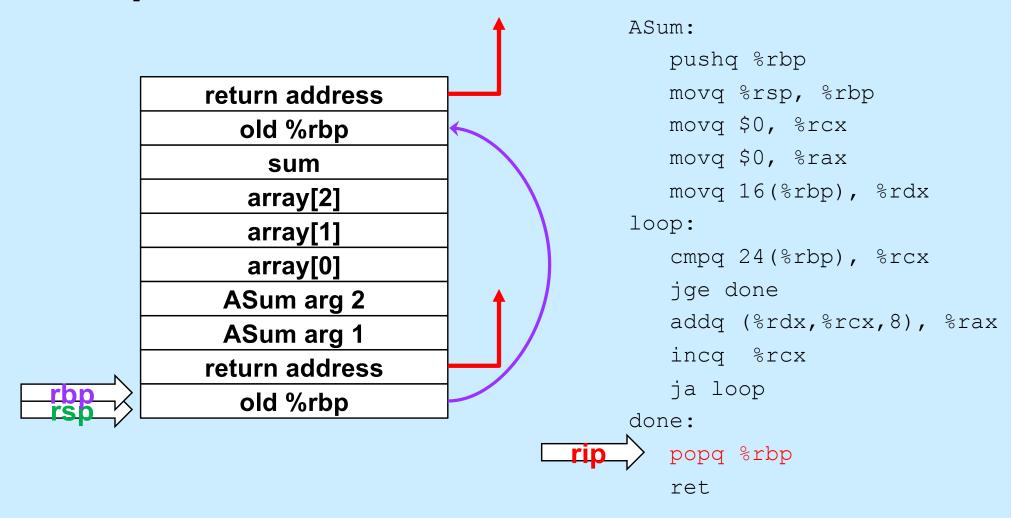
Quiz 1

What's at 24(%rbp)?

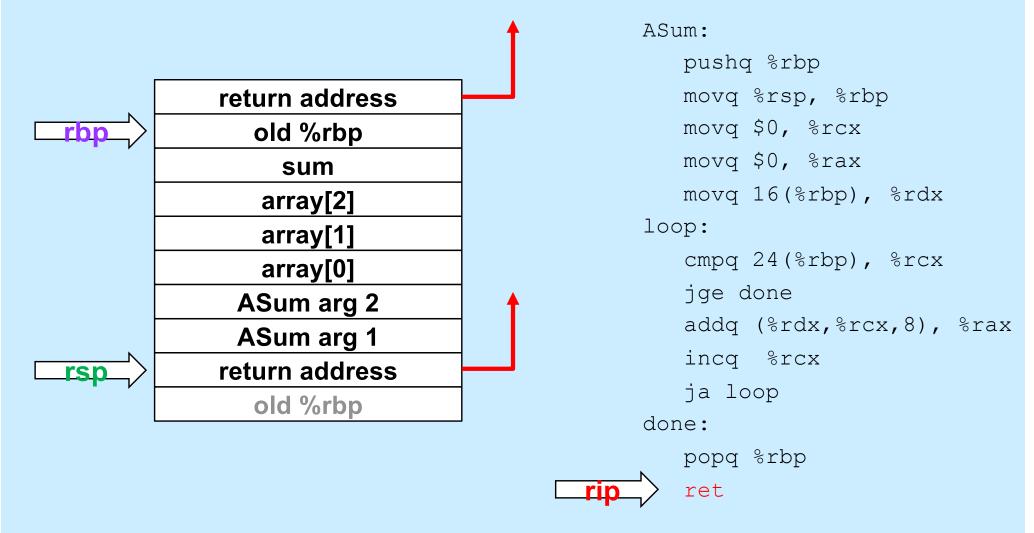
- a) a local variable
- b) the first argument to ASum
- c) the second argument to ASum
- d) something else

```
ASum:
   pushq %rbp
   movq %rsp, %rbp
   movq $0, %rcx
   movq $0, %rax
   movq 16(%rbp), %rdx
loop:
   cmpq 24(%rbp), %rcx
   jge done
   addq (%rdx,%rcx,8), %rax
   incq %rcx
   ja loop
done:
   popq %rbp
   ret
```

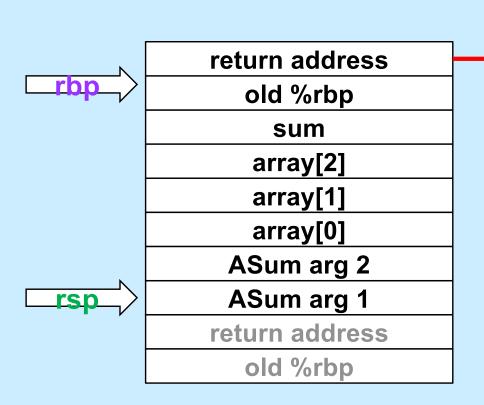
Prepare to Return



Return

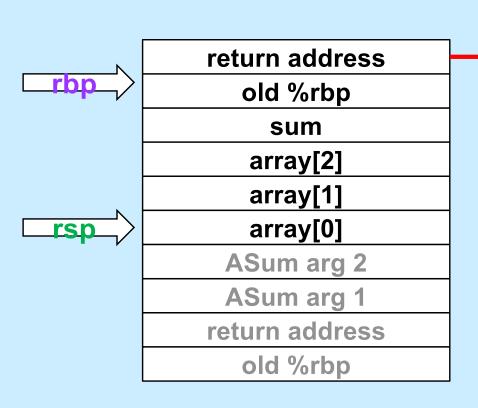


Pop Arguments



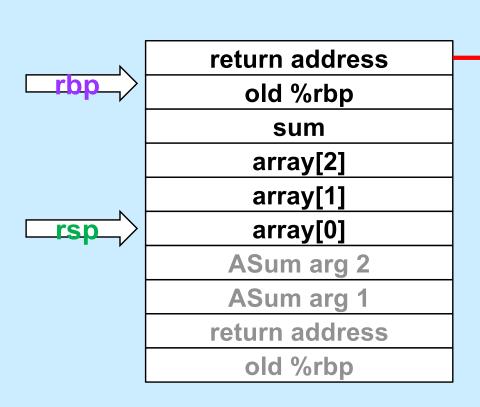
```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Save Return Value



```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq \$-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

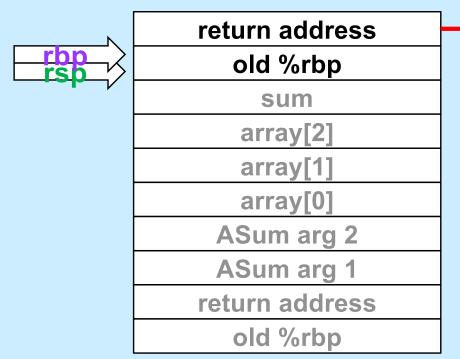
Pop Local Variables



```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
```

ret

Prepare to Return



```
mainfunc:
   pushq %rbp
   movq %rsp, %rbp
   subq $32, %rsp
   movq $2, -32(%rbp)
   movg $117, -24(%rbp)
   movq $-6, -16(%rbp)
   pusha $3
   leaq -32(%rbp), %rax
   pushq %rax
   call ASum
   addq $16, %rsp
   movq %rax, -8(%rbp)
   addq $32, %rsp
   popq %rbp
   ret
```

Return

return address
old %rbp
sum
array[2]
array[1]
array[0]
ASum arg 2
ASum arg 1
return address
old %rbp

mainfunc: pushq %rbp movq %rsp, %rbp subq \$32, %rsp movq \$2, -32(%rbp) movq \$117, -24(%rbp) movg \$-6, -16(%rbp)pushq \$3 leaq -32(%rbp), %rax pushq %rax call ASum addq \$16, %rsp movq %rax, -8(%rbp) addq \$32, %rsp popq %rbp

ret

Using Registers

- ASum modifies registers:
 - %rsp
 - %rbp
 - %rcx
 - %rax
 - %rdx
- Suppose its caller uses these registers

```
movq $33, %rcx
movq $167, %rdx
pushq $6
pushq array
call ASum
    # assumes unmodified %rcx and %rdx
addq $16, %rsp
addq %rax,%rcx  # %rcx was modified!
addq %rdx, %rcx  # %rdx was modified!
```

```
ASum:
   pushq %rbp
   movq %rsp, %rbp
   movq $0, %rcx
   movq $0, %rax
   movq 16(%rbp), %rdx
loop:
   cmpq 24(%rbp), %rcx
   jge done
   addq (%rdx,%rcx,8), %rax
   incq %rcx
   ja loop
done:
   popq %rbp
   ret
```

Register Values Across Function Calls

- ASum modifies registers:
 - %rsp
 - %rbp
 - %rcx
 - %rax
 - %rdx
- May the caller of ASum depend on its registers being the same on return?
 - ASum saves and restores %rbp and makes no net changes to %rsp
 - » their values are unmodified on return to its caller
 - %rax, %rcx, and %rdx are not saved and restored
 - » their values might be different on return

```
ASum:
   pushq %rbp
   movq %rsp, %rbp
   movq $0, %rcx
   movq $0, %rax
   movq 16(%rbp), %rdx
loop:
   cmpq 24(%rbp), %rcx
   jge done
   addq (%rdx,%rcx,8), %rax
   incq %rcx
   ja loop
done:
   popq %rbp
   ret
```

Register-Saving Conventions

Caller-save registers

 if the caller wants their values to be the same on return from function calls, it must save and restore them

```
pushq %rcx
call func
popq %rcx
```

Callee-save registers

 if the callee wants to use these registers, it must first save them, then restore their values before returning

```
func:
    pushq %rbx
    movq $6, %rbx
    ...
    popq %rbx
```

x86-64 General-Purpose Registers: Usage Conventions

%rax	Return value
%rbx	Callee saved
%rcx	Caller saved
%rdx	Caller saved
%rsi	Caller saved
%rdi	Caller saved
%rsp	Stack pointer
%rbp	Base pointer

%r8	Caller saved
%r9	Caller saved
%r10	Caller saved
%r11	Caller Saved
%r12	Callee saved
%r13	Callee saved
%r14	Callee saved
%r15	Callee saved

Recursive Function

```
/* Recursive popcount */
long pcount_r(unsigned long x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

Registers

- %rax, %rdx used without first saving
- %rbx used, but saved at beginning & restored at end

```
pcount r:
   pushq %rbp
   movq %rsp, %rbp
   pushq %rbx
   movq 16(%rbp), %rbx
   movq $0, %rax
   testq %rbx, %rbx
   je .L3
   movq %rbx, %rax
   shrq $1, %rax
   pushq %rax
   call pcount r
   addq $8, %rsp
   movq %rbx, %rdx
   andq $1, %rdx
   leaq (%rdx,%rax), %rax
.L3:
         %rbx
   popq
         %rbp
   popq
   ret
```

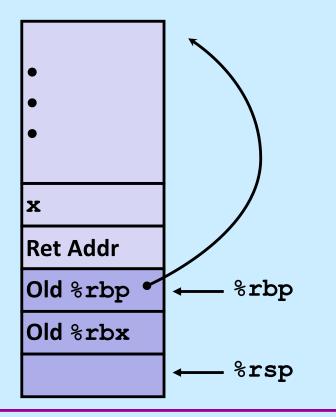
```
/* Recursive popcount */
long pcount_r(unsigned long x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

Actions

- save old value of %rbx on stack
- store x in %rbx

```
%rbx x
```

```
pcount_r:
    pushq %rbp
    movq %rsp, %rbp
    pushq %rbx
    movq 16(%rbp), %rbx
    • • •
```



```
/* Recursive popcount */
long pcount_r(unsigned long x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
movq $0, %rax
testq %rbx, %rbx
je .L3
.L3:
.ret
```

Actions

%rbx x

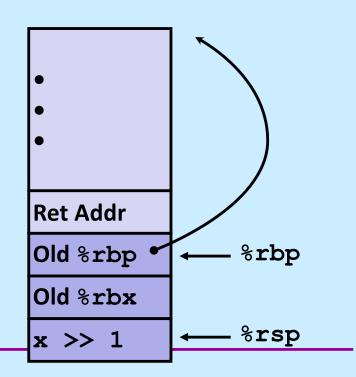
```
/* Recursive popcount */
long pcount_r(unsigned long x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

movq %rbx, %rax shrq \$1, %rax pushq %rax call pcount_r

Actions

- push x >> 1 on stack as arg
- make recursive call
- Effect
 - %rax set to function result
 - %rbx still has value of x

%rbx x



```
/* Recursive popcount */
long pcount_r(unsigned long x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
movq %rbx, %rdx
addq $8, %rsp
andq $1, %rdx
leaq (%rdx, %rax), %rax
```

Assume

- %rax holds value from recursive call
- %rbx holds x

Actions

- pop argument from stack
- compute (x & 1) + computed value

Effect

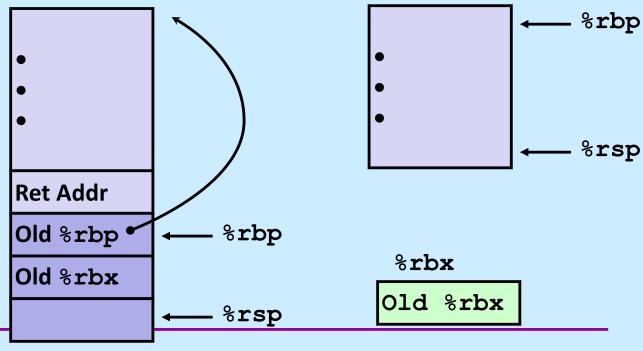
- %rax set to function result

```
/* Recursive popcount */
long pcount_r(unsigned long x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
popq %rbx
popq %rbp
ret
```

Actions

- restore values of %rbx and %rbp
- return (and pop return address)



CS33 Intro to Computer Systems

XII-47

Observations About Recursion

- Handled without special consideration
 - stack frames mean that each function call has private storage
 - » saved registers & local variables
 - » saved return pointer
 - register-saving conventions prevent one function call from corrupting another's data
 - stack discipline follows call / return pattern
 - » if P calls Q, then Q returns before P
 - » last-in, first-out
- Also works for mutual recursion
 - P calls Q; Q calls P

Passing Arguments in Registers

Observations

- accessing registers is much faster than accessing primary memory
 - » if arguments were in registers rather than on the stack, speed would increase
- most functions have just a few arguments

Actions

- change calling conventions so that the first six arguments are passed in registers
 - » in caller-save registers
- any additional arguments are pushed on the stack

Why Bother with a Base Pointer?

- It (%rbp) points to the beginning of the stack frame
 - making it easy for people to figure out where things are in the frame
 - but people don't execute the code ...
- The stack pointer always points somewhere within the stack frame
 - it moves about, but the compiler knows where it is pointing
 - » a local variable might be at 8(%rsp) for one instruction, but at 16(%rsp) for a subsequent one
 - » tough for people, but easy for the compiler
- Thus the base pointer is superfluous
 - it can be used as a general-purpose register

x86-64 General-Purpose Registers: Updated Usage Conventions

%rax	Return value
%rbx	Callee saved
%rcx	Argument #4
%rdx	Argument #3
%rsi	Argument #2
%rdi	Argument #1
%rsp	Stack pointer
%rbp	Callee saved

%r8	Argument #5
%r9	Argument #6
%r10	Caller saved
%r11	Caller Saved
%r12	Callee saved
%r13	Callee saved
%r14	Callee saved
%r15	Callee saved

Recursive Function (Improved)

```
/* Recursive popcount */
long pcount_r(unsigned long x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

Registers

- the single argument (x) is passed in %rdi
- %rbx is a callee-saved register and thus is saved and restored
- %rax is caller-saved
- %rbp isn't used

```
pcount r:
         $0, %rax
  movq
  testq %rdi, %rdi
     . L8
  jе
  pushq %rbx
  movq %rdi, %rbx
  shrq $1, %rdi
  call
         pcount r
  andq
         $1, %rbx
  addq %rbx, %rax
  popq %rbx
.L8:
  ret
```

Summary

- What's pushed on the stack
 - return address
 - saved registers
 - » caller-saved by the caller
 - » callee-saved by the callee
 - local variables
 - function parameters
 - » those too large to be in registers (structs)
 - » those beyond the six that we have registers for
 - large return values (structs)
 - » caller allocates space on stack
 - » callee copies return value to that space

Quiz 2

Suppose function A is compiled using the convention that %rbp is used as the base pointer, pointing to the beginning of the stack frame. Function B is compiled using the convention that there's no need for a base pointer. Will there be any problems if A calls B or if B calls A?

- a) Neither case will work
- b) A calling B works, but B calling A doesn't
- c) B calling A works, but A calling B doesn't
- d) Both work

Tail Recursion

```
int factorial(int x) {
                               int factorial(int x) {
  if (x == 1)
                                 return f2(x, 1);
    return x;
 else
                               int f2(int a1, int a2) {
    return
                                 if (a1 == 1)
      x*factorial(x-1);
                                    return a2;
                                 else
                                    return
                                      f2(a1-1, a1*a2);
```

No Tail Recursion (1)

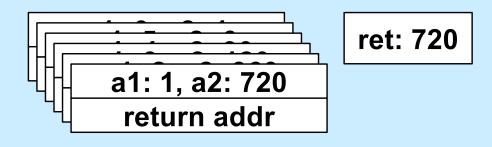
x: 6
return addr
x: 5
return addr
x: 4
return addr
x: 3
return addr
x: 2
return addr
x: 1
return addr

No Tail Recursion (2)

x: 6
return addr
x: 5
return addr
x: 4
return addr
x: 3
return addr
x: 2
return addr
x: 1
return addr

ret: 720
ret: 120
ret: 24
ret: 6
ret: 2

Tail Recursion



Code: gcc -O1

```
f2:
      movl %esi, %eax
      cmpl $1, %edi
      je .L5
      subq $8, %rsp
      movl %edi, %esi
      imull %eax, %esi
      subl
             $1, %edi
      call f2 # recursive call!
      addq $8, %rsp
.L5:
      ret
```

Not Using the Stack ...

```
f2:
    movl %esi, %eax
    cmpl $1, %edi
    je .L5
    movl %edi, %esi
    imull %eax, %esi
    subl $1, %edi
    call f2 # recursive call!
.L5:
    ret
```

Not Recursive!

```
movl %esi, %eax
cmpl $1, %edi
je .L5
movl %edi, %esi
imull %eax, %esi
subl $1, %edi
ja f2 # goto!
.L5:
ret
```

Code: gcc -O2

```
f2:
       cmpl $1, %edi
       movl %esi, %eax
             .L8
       je
.L12:
       imull %edi, %eax
       subl $1, %edi
                             loop!
       cmpl $1, %edi
       jne .L12
.L8:
       ret
```