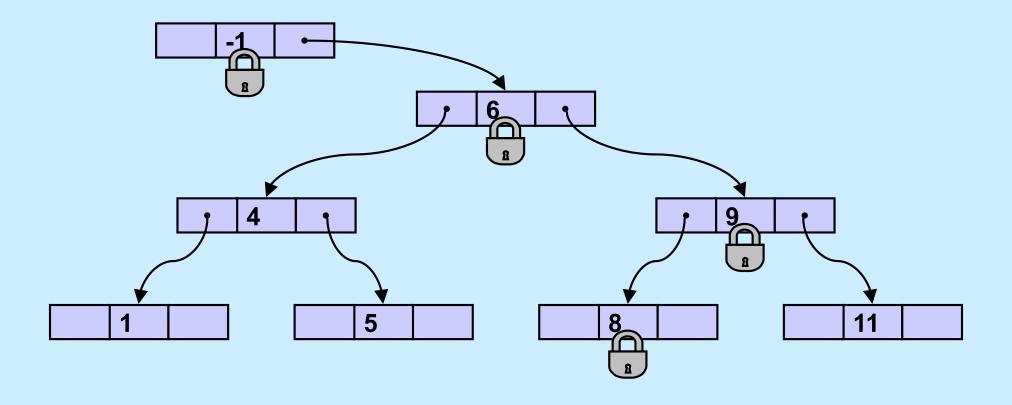
**CS 33** 

**Multithreaded Programming IV** 

# Doing It Right ...



#### C Code: Fine-Grained Search I

```
enum locktype {l read, l write};
                                         } else {
                                              lock(lt, &next->lock);
#define lock(lt, lk) ((lt) == l read)?
                                              if (key == next->key) {
      pthread rwlock rdlock(lk):
                                                result = next;
      pthread rwlock wrlock(lk)
                                              } else {
                                                pthread rwlock unlock (
Node *search(int key,
                                                     &parent->lock);
    Node *parent, Node **parentp,
                                                result = search(key,
    enum locktype lt) {
                                                    next, parentpp, lt);
   // parent is locked on entry
                                                 return result;
 Node *next;
 Node *result;
  if (key < parent->key) {
    if ((next = parent->lchild)
        == 0)
      result = 0;
```

#### C Code: Fine-Grained Search II

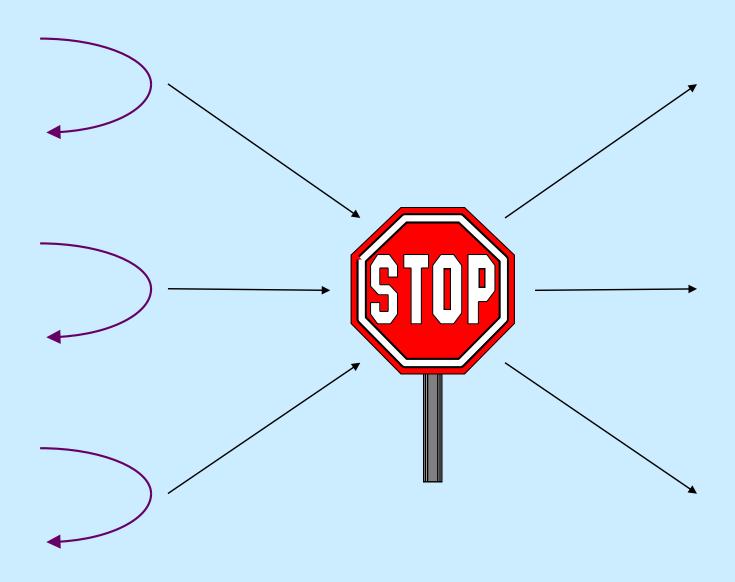
```
} else {
                                           } else {
  if ((next = parent->rchild)
                                            pthread rwlock unlock (
      == 0)
                                                 &parent->lock);
     result = 0;
                                             result = search(key,
   } else {
                                                 next, parentpp, lt);
                                            return result;
     lock(lt, &next->lock);
     if (key == next->key) {
       result = next;
                                      if (parentpp != 0) {
                                        // parent remains locked
                                        *parentpp = parent;
                                      } else
                                        pthread rwlock unlock (
                                             &parent->lock);
                                      return result;
```

# C Code: Add with Fine-Grained Synchronization I

# C Code: Add with Fine-Grained Synchronization II

```
newnode = malloc(sizeof(Node));
newnode->key = key;
newnode->lchild = newnode->rchild = 0;
pthread_rwlock_init(&newnode->lock, 0);
if (name < parent->name)
   parent->lchild = newnode;
else
   parent->rchild = newnode;
pthread_rwlock_unlock(&parent->lock);
return 1;
```

### **Barriers**



#### A Solution?

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
} else while (!(count == number)) {
   pthread_cond_wait(&cond_var, &m);
}
pthread_mutex_unlock(&m);
```

#### **How About This?**

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
   count = 0;
} else while (!(count == number)) {
   pthread_cond_wait(&cond_var, &m);
}
pthread_mutex_unlock(&m);
```

#### And This ...

```
pthread_mutex_lock(&m);
if (++count == number) {
   pthread_cond_broadcast(&cond_var);
   count = 0;
} else {
   pthread_cond_wait(&cond_var, &m);
}
pthread_mutex_unlock(&m);
```

Quiz 1

Does it work?

a) definitely

#### **Barrier in POSIX Threads**

```
pthread mutex lock(&m);
if (++count < number) {</pre>
  int my generation = generation;
  while (my generation == generation) {
    pthread cond wait(&waitQ, &m);
 else {
  count = 0;
  generation++;
  pthread cond broadcast (&waitQ);
pthread mutex unlock (&m);
```

#### **More From POSIX!**

## Why cond\_wait is Weird ...

```
pthread_cond_wait(pthread_cond_t *c, pthread_mutex_t *m) {
    pthread_mutex_unlock(m);
    sem_wait(c->sem);
    pthread_mutex_lock(m);
}

pthread_cond_signal(pthread_cond_t *c) {
    sem_post(c->sem);
}
```

#### **Deviations**

Signals



VS.



- Cancellation
  - tamed lightning

# **Signals**



- who gets them?
- who needs them?



– how do you respond to them?

# **Dealing with Signals**

- Per-thread signal masks
- Per-process signal vectors
- One delivery per signal

# Signals and Threads

```
int pthread_kill(pthread_t thread, int signo);
```

thread equivalent of kill

thread equivalent of sigprocmask

# **Asynchronous Signals (1)**

```
int main() {
  void handler(int);
   signal(SIGINT, handler);
void handler(int sig) {
```

# **Asynchronous Signals (2)**

```
int main() {
                                 void handler(int sig) {
 void handler(int);
                                   ... // deal with signal
  signal(SIGINT, handler);
                                   printf("equally important "
  ... // complicated program
                                      "message: %s\n", message);
 printf("important message:
     "%s\n", message);
  ... // more program
```

#### Quiz 2

```
int main() {
  void handler(int);
  signal(SIGINT, handler);
  ... // complicated program
 pthread mutex lock(&mut);
 printf("important message: "
     "%s\n", message);
 pthread mutex unlock (&mut);
  ... // more program
```

```
void handler(int sig) {
    ... // deal with signal

pthread_mutex_lock(&mut);
 printf("equally important "
    "message: %s\n", message);
 pthread_mutex_unlock(&mut);
}
```

# Does this work? a) yes b) no

# **Synchronizing Asynchrony**

```
computation state t state;
sigset t set;
int main() {
  pthread_t thread;
  sigemptyset(&set);
  sigaddset(&set, SIGINT);
  pthread sigmask (SIG BLOCK,
   &set, 0);
  pthread create (&thread, 0,
   monitor, 0);
  long running procedure();
```

```
void *monitor(void *dummy) {
  int sig;
  while (1) {
    sigwait(&set, &sig);
    display(&state);
  }
  return(0);
}
```

#### **Cancellation**



# **Sample Code**

```
void *thread code(void *arg) {
  node t *head = 0;
  while (1) {
    node t *nodep;
    nodep = (node t *) malloc(sizeof(node t));
    if (read(0, &node->value,
        sizeof(node->value)) == 0) {
      free (nodep);
      break;
                               pthread cancel(thread);
    nodep->next = head;
    head = nodep;
  return head;
```

#### **Cancellation Concerns**

- Getting cancelled at an inopportune moment
- Cleaning up

#### **Cancellation State**

#### Pending cancel

```
- pthread cancel (thread)
```

#### Cancels enabled or disabled

```
- int pthread_setcancelstate(
     {PTHREAD_CANCEL_DISABLE
     PTHREAD_CANCEL_ENABLE},
     &oldstate)
```

#### Asynchronous vs. deferred cancels

```
- int pthread_setcanceltype(
     {PTHREAD_CANCEL_ASYNCHRONOUS,
     PTHREAD_CANCEL_DEFERRED),
     &oldtype)
```

#### **Cancellation Points**

- aio\_suspend
- close
- creat
- fcntl (when F\_SETLCKW is the command)
- fsync
- mq\_receive
- mq\_send
- msync
- nanosleep
- open
- pause
- pthread\_cond\_wait
- pthread\_cond\_timedwait
- pthread\_join

- pthread\_testcancel
- read
- sem\_wait
- sigwait
- sigwaitinfo
- sigsuspend
- sigtimedwait
- sleep
- system
- tcdrain
- wait
- waitpid
- write

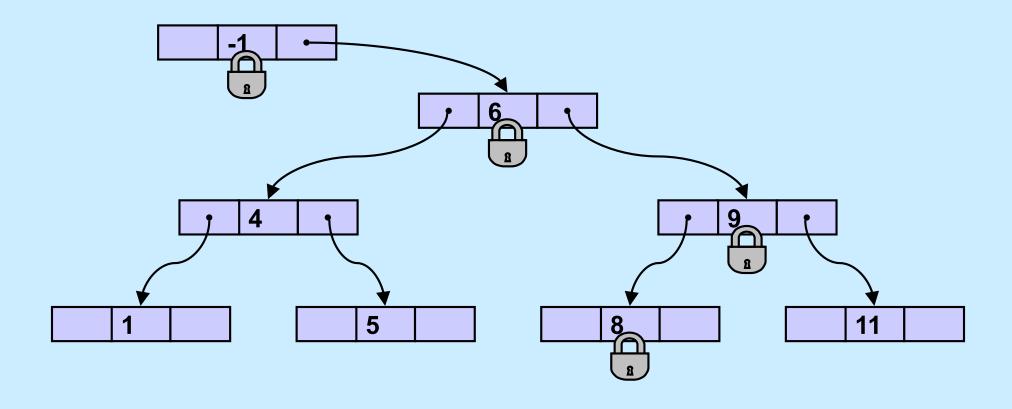
# Cleaning Up

- void pthread\_cleanup\_push((void)(\*routine)(void \*),void \*arg)
- void pthread cleanup pop(int execute)

# Sample Code, Revisited

```
void *thread code(void *arg) {
                                          void cleanup(void *arg) {
  node t *head = 0;
                                            node t **headp = arg;
  pthread cleanup push (
                                            while(*headp) {
      cleanup, &head);
                                              node t *nodep = head->next;
    while (1) {
                                              free(*headp);
      node t *nodep;
                                              *headp = nodep;
      nodep = (node t *)
      malloc(sizeof(node t));
      if (read(0, &node->value,
          sizeof(node->value)) == 0) {
        free (nodep);
        break;
      nodep->next = head;
      head = nodep;
  pthread cleanup pop(0);
  return head;
```

# A More Complicated Situation ...



# Start/Stop



#### Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  while (s->state == stopped)
    pthread cond wait(&s->queue, &s->mutex);
  pthread mutex unlock(&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast(&s->queue);
  pthread mutex unlock(&s->mutex);
```

# Start/Stop

#### Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  while (s->state == stopped)
    pthread cond wait (&s->queue,
      &s->mutex);
  pthread mutex unlock (&s->mutex);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast(&s->queue);
  pthread mutex unlock(&s->mutex);
```



#### Quiz 3

You're in charge of designing POSIX threads. Should *pthread\_cond\_wait* be a cancellation point?

- a) no
- b) yes; cancelled threads must acquire mutex before invoking cleanup handler
- c) yes; but they don't acquire mutex

# Start/Stop





#### Start/Stop interface

```
void wait for start(state t *s) {
  pthread mutex lock(&s->mutex);
  pthread cleanup push (
    pthread mutex unlock, &s);
  while(s->state == stopped)
    pthread cond wait (&s->queue, &s->mutex);
  pthread cleanup pop(1);
void start(state t *s) {
  pthread mutex lock(&s->mutex);
  s->state = started;
  pthread cond broadcast (&s->queue);
  pthread mutex unlock(&s->mutex);
```

#### **Cancellation and Conditions**

```
pthread_mutex_lock(&m);
pthread_cleanup_push(pthread_mutex_unlock, &m);
while(should_wait)
   pthread_cond_wait(&cv, &m);

// ... (code perhaps containing other cancellation points)
pthread_cleanup_pop(1);
```