

## More KMP

Brief:

- \* we have skeleton machine  $M_p$
- \* pattern  $p$ , text  $t$
- \* looking for first occurrence of  $p$  in  $t$

Failure function  $\sim$  computed based on  $p$

e.g.  $p = aabbbaab \rightarrow \text{length}(p) = 7$

$t = abaabaaatbaaab$

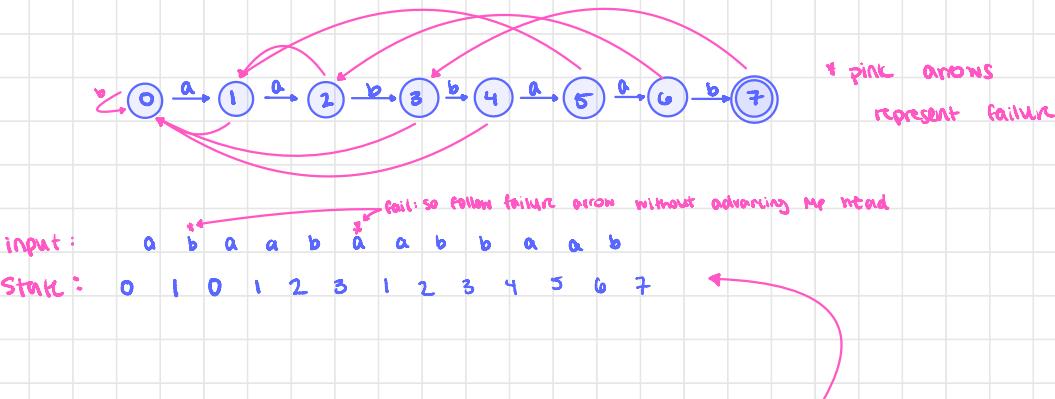
Skeleton machine  $M_p$ :



failure function:

$p_i$	a	a	b	b	a	a	b
$i$	1	2	3	4	5	6	7
$f(i)$	0	1	0	0	1	2	3

Add edges to  $M_p$  based on failure function:



Here are the first few steps we took to determine the sequence of states:

1. Initially,  $M_p$  is in state 0
2. upon reading first letter of  $t$ ,  $M_p$  enters state 1 because  $t_1 = p_1 = "a"$
3. next, the machine reads that  $t_2 \neq p_2$ , so  $M_p$  enters state 0 (follow the failure arrow for state 1) without moving the input head
4. since  $t_2 \neq p_2$ , case 2\* prevails, and  $M_p$  remains in state 0 and advances its head to  $t_3$
- ...

\* see previous class notes

## Failure Function alg.:

**input:** pattern of length  $\geq 1$   
**output:** failure function  $f(p)$

```

function FAILURE FUNCTION( $p = p_1 \dots p_l$ )
L1 {  $f(1) \leftarrow 0$ 
       $i \leftarrow 0$ 
L2 for  $j \in \{2, \dots, l\}$  do
L3    $i \leftarrow f(j-1)$ 
L4   while  $p_j \neq p_{i+1}$  and  $i > 0$  do
         $i \leftarrow f(i)$ 
    end while
L5   if  $p_j \neq p_{i+1}$  and  $i = 0$  then
         $f(j) \leftarrow 0$ 
    else
L6      $f(j) \leftarrow i + 1$ 
    end if
end for
end function

```

for  $j=4$ :

$i = f(3) = 0$   
 $\text{(i) } \text{(ii) } \checkmark$   
while  $p_4 \neq p_1$  and  $i > 0$   
if  $p_4 \neq p_{i+1}$  and  $i = 0$   
 $f(4) = 0$

for  $j=5$ :

$i = f(4) = 0$   
 $\text{(i) } \text{(ii) } \checkmark$   
while  $p_5 \neq p_1$  and  $i > 0$   
if  $p_5 \neq p_1$  and  $i = 0$   
else:  
 $f(5) = i + 1 = 0 + 1 = 1$

for  $j=6$ :

$i = f(5) = 1$   
 $\text{(i) } \text{(ii) } \checkmark$   
while  $p_6 \neq p_2$  and  $i > 0$   
if  $p_6 \neq p_2$  and  $i = 0$   
else:  $f(6) = i + 1 = 1 + 1 = 2$

for  $j=7$ :

$i = f(6) = 2$   
 $\text{(i) } \text{(ii) } \checkmark$   
while  $p_7 \neq p_3$  and  $i > 0$   
if  $p_7 \neq p_3$  and  $i = 0$   
else:  $f(7) = i + 1 = 2 + 1 = 3$

Let's walk through the algorithm on  $p = ambabab$

$$f(1) = 0$$

$$i = 0$$

for  $j=2$

$$i = f(j-1) = f(1) = 0$$

$$p_2 = p_2 = a \quad \checkmark$$

$$p_{1+1} = p_1 = a \quad \checkmark$$

while  $p_2 \neq p_{i+1}$  and  $i > 0$   
 $\times$

do not enter while loop

$$\text{if } p_2 \neq p_{i+1} \text{ and } i = 0 \quad \checkmark$$

skip to "else"

$$f(2) = i + 1 \quad \text{so } f(2) = 1$$

for  $j=3$

$$i = f(j-1) = f(2) = 1$$

$\text{(i) } \text{(ii) } \checkmark$   
while  $p_3 \neq p_2$  and  $i > 0$   
 $\checkmark$

$$i = f(i) = f(1) = 0$$

$\text{(i) } \text{(ii) } \checkmark$   
if  $p_3 \neq p_1$  and  $i = 0$   
 $\checkmark$

$$f(3) = f(3) = 0$$

i	1	2	3	4	5	6	7
$f(i)$	0	1	0	0	1	2	3

Yippee it works!