

## C - Week - 2

### L1 No. representation in a computer

- Humans use variety of symbols. Computers are built using only 0 & 1.
- Inputs of quad.  $10^n \rightarrow a, b, c$ . Are these arbitrary no.? How are they stored? How are they recognized as no.?

-  $5 \rightarrow 5 \times 10^0$ ,  $27 \rightarrow 2 \times 10^1 + 7 \times 10^0$

- Binary no.

$$5 \text{ (decimal)} = 4 + 0 + 1$$
$$2^2 \quad 2^1 \quad 2^0$$

$$(5)_{10} \rightarrow (101)_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$$
$$= 101 \text{ (binary)}$$

$$27 = 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$

- How many bits do we need in gen.?

$$8 \text{ bit no.} \rightarrow 0 \rightarrow 255 \quad \boxed{2^n - 1}$$
$$2^8 - 1 = 256 - 1 = 255$$

- Physical memory  $\rightarrow$  circuits, uniform size preferable. CPU registers. Engg. Tradeoffs.
- C  $\rightarrow$  basic int. data type  $\rightarrow$  32 bits  
 $\hookrightarrow$  primitive data type, portability.

## L2 Negative No. & Hexadecimal Representations

MSB  $\leftarrow$  00001111  $\rightarrow$  LSB  
Most Significant Bit  $\rightarrow$  least

- For negative integers, MSB  $\rightarrow$  1 (-ve)  
0 (+ve)

Signed  $\rightarrow$  0011 = +3

Magnitude 1011 = -3

Convention

1 bit for sign  $\rightarrow$  0 to  $2^{N-1}-1$   $\rightarrow$  undemanded  
 $\rightarrow$   $-(2^{N-1}-1)$  to 0  $\rightarrow$  -dant

### - Complement Convention

$x \geq 0 \Rightarrow$  store  $x$  as reg. binary no.

$x < 0 \Rightarrow$  store complement of  $x$

Complement  $\Rightarrow 2^{N-1} - x$  for  $N$ -bit no.

Eg. 8 bits  $17 \Rightarrow 00010001$

$$\begin{aligned} -17 &\Rightarrow 2^8 - 17 = 239 \\ &= 11101111 \end{aligned}$$

largest +value  $\Rightarrow +127 = 01111111$

" "  $\Rightarrow -128 = 10000000$

Adv.  $\rightarrow$  covers full range, easy to compute in hardware, arithmetic easy

- Binary is too difficult for humans to read  $\Rightarrow$  hexadecimal.

$16 = 2^4 \Rightarrow$  1 digit in base 16 = 4 bits

$$(0101)_2 = (5)_{10} = (5)_{16}$$

$$(1001)_2 = (9)_{10} = (9)_{16}$$

$$(1010)_2 = (10)_{10} = (A)_{16}$$

$$(1101)_2 = (13)_{10} = (D)_{16}$$

0 ... 9, A, B, C, D, E, F

- 8 bit values -

$$\left( \begin{array}{cc} 0001 & 0001 \\ \hline 1 & 1 \end{array} \right)_2 = (17)_{10} = (11)_{16}$$

$$\left( \begin{array}{cc} 0101 & 1011 \\ \hline 5 & B \end{array} \right)_2 = (91)_{10} = (5B)_{16}$$

8 bits  $\Rightarrow 0x00$  to  $0xFF$

- Integers -

1. fixed width - possibly CPU or lang. dep.
2. -ve nos.  $\rightarrow$  2's complement is preferred.
3. hexadecimal - useful intermediate.

### L3 Floating Pt. Representa<sup>n</sup>

$\pi = 3.14$ ,  $c = 3 \times 10^8$  m/s  $\rightarrow$  scientific nota<sup>n</sup>

- Fixed pt. nota<sup>n</sup> - integers with scaling factor

16 bit  $\rightarrow (-2^{15}, +2^{15}-1) \times 2^{-8}$  scale factor

- Floating pt. nota<sup>n</sup>

32 bits  $\rightarrow$  1 sign, 8 exponent, 23 bits value

- single precision (32 bits)  $\rightarrow 10^{-38} \rightarrow 10^{38}$
- double " (64 bits)  $\rightarrow 10^{-208}$  to  $10^{208}$  (scientific computing)
- Half " (16 bits)  $\rightarrow$  used in ML.

### L4 Character Representa<sup>n</sup> & Encoding

- comp. process text, encode text into binary, max. 200 charac.  $\rightarrow 2^8$  bits are enough.
- ASCII  $\rightarrow$  maps charac. to no./codes
- unicode  $\rightarrow$  consistent encoding, representa<sup>n</sup> specifies code pts., diff. encodings possible.
- ASCII  $\rightarrow$  8 bits / charac., but Ltd.  
Unicode + UTF-8  $\rightarrow$  8  $\rightarrow$  32 bits / charac.  
 $\hookrightarrow$  complex, needs careful handling
- use of std. libraries recommended.
- comp. store & process bits. Memory (array of bits). Registers (grp. of bits)

## L5 Instructions Encoding

- Basic ops<sup>m</sup> defined for CPU.
  - ↳ ALU, Load/Store (Memory/Storage), Branching (Control flow)
- not all instructions are available on all CPUs. But, can replicate intended behaviour thru other means.
- Eg. of RISC-V - 32bit-ISA architecture.
  - ↳ other and even smaller encodings are possible.



Eg. 32-bit memory

1M loca<sup>n</sup>  $\Rightarrow$  4 MB

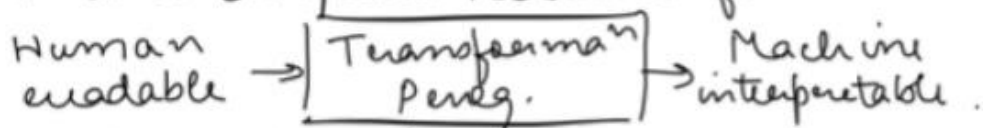
Code: 200K<sup>n</sup>  $\Rightarrow$  800 KB

Data: 300K<sup>n</sup>  $\Rightarrow$  1200 KB

no physical distinc<sup>n</sup> btw.  
code & data.

## L6 Program Compila<sup>n</sup>

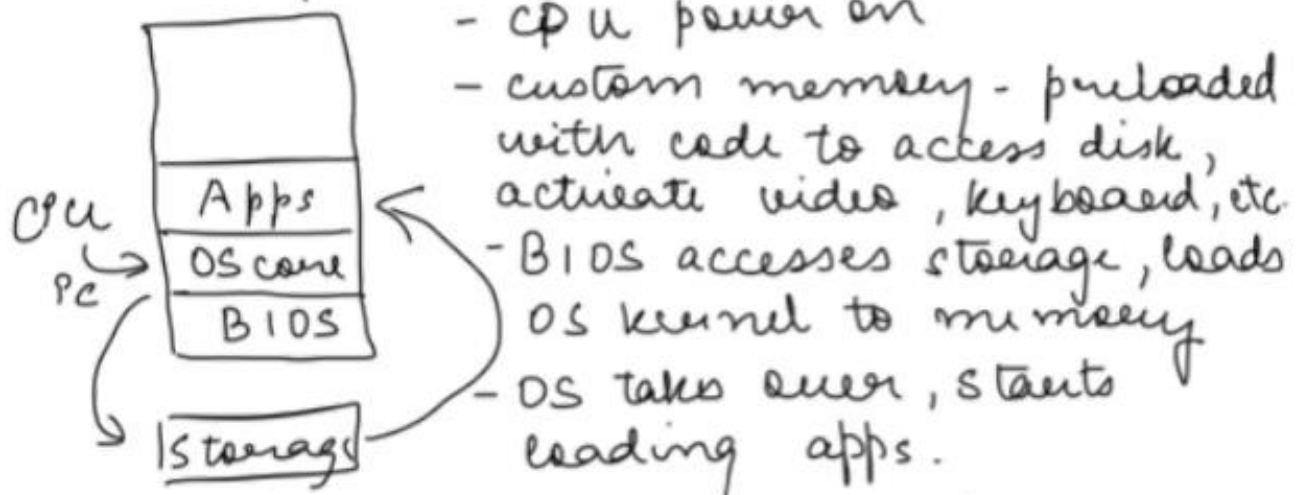
- programs as text  $\rightarrow$  human readable, colours, background, etc. (not part of prog. but IDE), indenta<sup>n</sup> & formatting may be relevant, files help to organize text into computer readable form.



- Compiler is previously created prog.  
Input prog. is data for compiler.  
Output of a prog. is a new executable prog.
- many steps are involved in converting human-readable prog. to machine interpretable prog.  $\rightarrow$  syntax check, map statements to equivalent codes, generate actual machine code.
- Optimiza<sup>m</sup>  $\rightarrow$  speed, prog. size, etc.
- can be CPU dependent.

## L7 Role of OS in a Computer

- write the prog. to enter into comp. I/O using peripherals, save data as encoded text using file sys.
- compile the prog., it is multiprocessing (can execute several prog.)
- Key things → compiler app, business/web based, memory mgmt.
- execute the prog. → load into memory after flow of control.
- In all these activities, OS acts as the coordinator.
- OS → is a prog., 1<sup>st</sup> prog. executed on startup, coordinates loading & executing other prog.
- Kernel - OS core - almost negligible interac<sup>n</sup> with humans.
- Booting / Bootstrapping



"Fly by pulling yourself up with your bootstraps."