

## L4: TSP Problem L5: Vehicle Routing Problem Routing Vehicle Routing Problem (VRP) – Introduction • When it comes to logistics, an important question is "How should a given set of locations be visited?" This is referred to as the routing problem In the TSP, a single salesman needs to visit all the nodes and return back e all the locations are visited (and the demand is satisfied) by The aim of any routing problem is to ens The TSP does not consider attaining Minimal travel distance Minimal use of resources Or both > Any capacity restrictions Multiple salesmen (or vehicles) Does not consider a specific (set) of starting points Thus, the Vehicle Routing Problem (VRP) is an extension of the TSP In this lecture, we will explore two problems 1. Travelling Salesman Problem (TSP) The simplest form of the VRP (Single depot VRP) is one where 2. Vehicle Routing Problem (VRP) There is a single depot It is noted that though these problems are NP-Hard, a learner needs to understand the mathematical construct to solve it via any Heuristic technique There are "n" nodes (customers) Each node requires a certain quantity that needs to be de There are a fixed number of vehicles available, and the capacity of the vehicles is known The distance (between the nodes, as well as from the depot to the nodes) is known The vehicles must start from the depot and return to the depot after visiting the customer nodes Hence, in this lecture, the focus will be on formulating these problems mathematically Travelling Salesman Problem (TSP) · There are also Multiple depot VRPs, where the customers can be served by any vehicle from any depot In this problem, a salesman needs to visi "n" nodes and the only return to the starting point The salesman can start from any node (from the "n" possible nodes) · In this lecture, the discussion is only on the Single Depot VRP were use this In the TSP, it is assumed that the salesman can go from any node to another node Single Depot VRP – Types lan now The salesman must not perform any sub-tours (form closed paths that does not include all nodes) Then, the question is "What route should the salesman take to visit all the cities, such that the total distance travelled is minimal" In the Single Depot VRP there are three types of problems which are possible TSP – Verbal Formulation All the vehicles have to be used (it is assumed that every vehicle will have at least 1 node attached to it) 2. Minimum number of vehicles must be used 3. Using enough vehicles such that the total distance travelled is minimum i, j → set of nodes {1,2,3...n} $D_{ij} \rightarrow distance$ between node-i and node-j (Note: $D_{ij}$ is set to a large value) Decision variable: Single Depot VRP - Verbal Formulation $\mathbf{X}_{ij}$ = 1 if the salesman visits node-j immediately after visiting node-i There are several ways to formulate the VRP X<sub>ii</sub> = 0 otherwise The following formulation is based on Achuthan and Caccetta (1991), 'Integer linear programming formulation for a vehicle routing problem', European Journal of Operational Research, Vol. 52, pp. 86-89. U. → Auxiliary variable to eliminate sub-tours Objective Function: In this formulation, the number of vehicles is known and all vehicles must be used The total distance travelled by the salesman must be minimized Sets & Parameters: i, j → set of nodes {1,2,3...n} D. → distance between node-i and node-i (Note: D. is set to a lor) 1. The salesman must visit all the nodes exactly once → set of vehicles {1,2,...k} 2. The salesman must not do any sub-tours (sub-tour elimination constraints) Q, → Demand at node-j parameters → Capacity for a Vehicle $X_{i,j}$ =1 if vehicle travels from node-i to node-j, 0 otherwise TSP - Mathematical Formulation U → Auxiliary variable to eliminate sub-tour Objective Function: Objective function: Minimize the total travel distance The total distance travelled by the sales $Minimize \sum_{i=1}^{K} \sum_{j=1}^{K} D_{i,j} * X_{i,j}$ Every node must be reached either from the depot (the source) or another node Every node must end either in another customer node or at the depot (the destination) There are exactly V-routes that start from the depot and that terminate at the depot (ensuring vehicles are used) Constraints: 4. The load carried in each route does not exceed the vehicle capacity 1. The salesman must visit all the nodes exactly once 5. Non-Negativity (Binary variable) Constraints $\frac{1. \, \text{Every node must be reached either from the depot (the source) or another node}}{\sum_{i=1,l \neq j}^{n} X_{i,j} + X_{0,j} = 1 \, \forall \, j \in \{1,...,n\} \, \text{where node $^*0$}^* \, \text{representes the depot } in the property of the property of$ $X_{i,j}=1 \ \forall \ j \ \in \{1,\dots,n\}$ The salesman must not do any sub-tours $U_i-U_j+n*X_{i,j}\leq n-1\ \forall\ i\in\{1,\dots,n-1\}\ j\in\{2,\dots,n\}\ and\ i\neq j$ $\sum_{i=1}^{n} X_{i,j} + X_{i,0} = 1 \,\forall \, i \in \{1,...,n\} \, \text{where node "0" representes the depot}$ 3. Non-Negativity $U_j \ge 0, X_{i,j} \in \{0,1\} \forall i, j$ 3. There are exactly V-routes that start from the depot and that terminate at the depot (ensuring all vehicles are 4. The load carried in each route does not exceed the vehicle capacity \( \mathcal{V} \) $U_i - U_j + W * X_{i,j} \le W - Q_j \; \forall \; i,j \; \in \{(0,i),(i,j),(i,0) : 1 \le i \; \neq j \le 0\}$ $U_{depot \ is \ destination} \ - \ U_{depot \ is \ source} = W$ In the above "1" represents the first node in N and "0" represents the depot 5. Non-Negativity (Binary Decision Variable): $X_{i,j} \in \{0,1\} \ \forall i,j$