Intro to Software Testing chapter 8.1

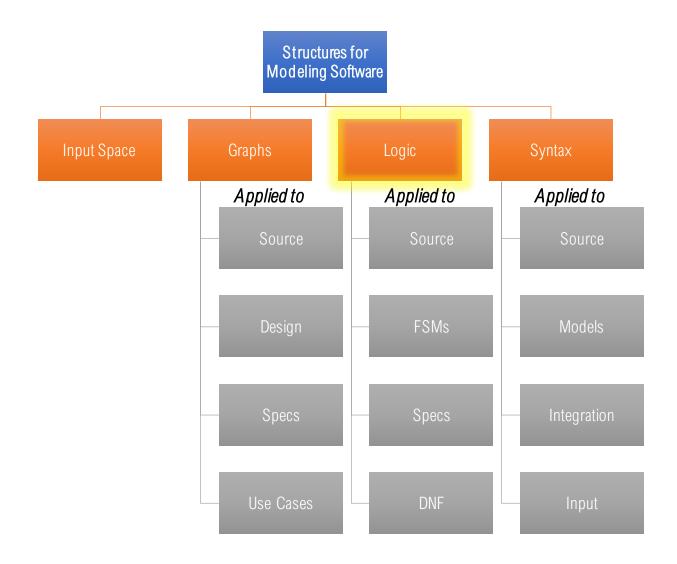
Semantic Logic Coverage

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https://go.gmu.edu/SWE637

Adapted from slides by Jeff Offutt and Bob Kurtz

LOGIC COVERAGE



Semantic Logic Criteria

Logic expressions show up in many situations

Logic expressions can come from many sources

Decisions in programs

FSMs and state charts

Requirements

Covering logic expressions is *required* by the US Federal Aviation Administration (FAA) for safety-critical software (regulation DO-178B/C)

MCDC – Modified Condition/Decision Coverage

Tests are intended to choose some subset of the total number of *truth* assignments (e.g. true/false) to the expressions

Logic Predicates and Clauses

Predicate: an expression that evaluates to a boolean value

Predicates can contain

Boolean variables or function calls

Non-Boolean variables that contain relational operators (<, >, ==, <=, >=, !=)

Logical operators

- negation operator (Java: !)
- ^ and operator (Java: &&)
- ∨ *or* operator (Java: ||)
- ⊕ *exclusive or* operator (Java: ^)
- implication operator
- equivalence operator

Clause: a predicate without logical operators

Semantic Logic Criteria

$$(a < b) \lor f(z) \land d \lor (m >= n*o)$$

This predicate has four clauses:

(a < b): a relational expression

f(z): a boolean function

d: a boolean variable

 $(m \ge n^*o)$: a relational expression

Predicate and Clause statistics

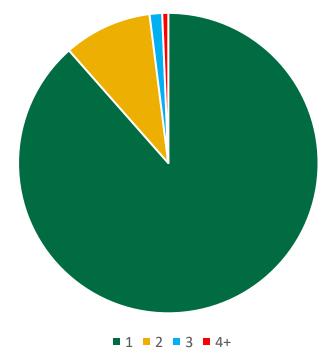
From a study of 63 open-source programs which contained > 400,000 predicates:

88.5% had only 1 clause

9.5% had 2 clauses

1.35% had 3 clauses

Only **0.65%** had 4 or more clauses



Predicate and Clause statistics

Predicates can be derived from:

- Decisions in program source code
- Guards on finite state machine transitions
- Decisions in UML activity graphs
- Requirements, both formal and informal
- SQL queries
- ...and more

Logic Coverage Criteria

We use predicates in testing to:

Develop a model of the software as one or more predicates Require tests to satisfy some combination of clauses

Abbreviations:

P is a set of predicates

p is a predicate in P

C is the set of clauses in P

Cp is the set of clauses in predicate p

c is a clause in C

Predicate Coverage

The first (and simplest) criterion requires that each predicate be evaluated to both true and false

DEFINITION

Predicate Coverage (PC) – For each **p** in **P**, **TR** contains two requirements: **p** evaluates to true, and **p** evaluates to false

When predicates come from conditions on graph edges, this is equivalent to *edge coverage*

Predicate Coverage Example

$$(a < b) \lor d \land (m >= n*o)$$

Predicate must evaluate to true

Predicate must evaluate to false

Predicate Coverage Example

$$(a < b) \lor d \land (m >= n*o)$$

Predicate must evaluate to true

= true

```
Example test case:

a=5, b=10, d=true, m=1, n=1, o=1

= ((5 < 10) \lor true) \land (1 >= 1*1)

= (true \lor true) \land (true)
```

Predicate Coverage Example

$$(a < b) \lor d \land (m >= n*o)$$

Predicate must evaluate to true

```
Example test case:
```

$$= ((5 < 10) \lor true) \land (1 >= 1*1)$$

- = (true \vee true) \wedge (true)
- = true

Predicate must evaluate to false

Example test case:

$$= ((10 < 5) \lor false) \land (1 >= 1*1)$$

- = (false \vee false) \wedge (true)
- = false

Clause Coverage

Predicate coverage does not require evaluation of all clauses, so a more complete criterion is clause coverage

Clause Coverage (CC) — For each c in C, TR contains two requirements: c evaluates to true, and c evaluates to false

Does clause coverage subsume predicate coverage?

Clause Coverage

Predicate coverage does not require evaluation of all clauses, so a more complete criterion is clause coverage

DEFINITION

Clause Coverage (CC) — For each c in C, TR contains two requirements: c evaluates to true, and c evaluates to false

Does clause coverage subsume predicate coverage?

No! Consider $a \lor b$, the clauses can be satisfied with (a=true, b=false) and (a=false, b=true) but the predicate is always true

Clause Coverage Example

$$(a < b) \lor d \land (m >= n*o)$$

Clauses must evaluate to true

In this case we arbitrarily choose to set all clauses to true at the same time, but that is not required

Clause Coverage Example

$$(a < b) \lor d \land (m >= n*o)$$

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d: true

$$(m >= n*o) : m=1, n=1, o=1$$

Clause Coverage Example

$$(a < b) \lor d \land (m >= n*o)$$

Clauses must evaluate to true

In this case we arbitrarily choose to set all clauses to true at the same time, but that is not required

d: true

$$(m >= n*o) : m=1, n=1, o=1$$

Clauses must evaluate to false

(a < b): a = 10, b = 5

d: false

(m >= n*o) : m=1, n=2, o=2

Limitations of PC and CC

PC does not fully exercise all the clauses, especially with short-circuit evaluation

 $(a < b) \lor (c < d)$: if (a < b) is true, then (c < d) is not evaluated CC does not always ensure PC (and so CC *does not subsume* PC)

We can satisfy CC without causing the entire predicate to be both true and false

That doesn't seem like something we want to do...

The simplest solution is to just test all combinations

Combinatorial Coverage

DEFINITION

Combinatorial Coverage (CoC) – For each p in P, TR contains a test requirement for the clauses in C_p to evaluate to each possible combination of truth values

Combinatorial coverage is conceptually simple and complete, but very expensive

Results in 2^N tests, where N is the number of clauses

Combinatorial Coverage Example

	a <b< th=""><th>d</th><th>m>=n*0</th><th>$((a < b) \lor d) \land (m >= n*o)$</th></b<>	d	m>=n*0	$((a < b) \lor d) \land (m >= n*o)$
1	Т	T	T	Т
2	Т	T	F	F
3	Т	F	T	Т
4	Т	F	F	F
5	F	T	Т	Т
6	F	T	F	F
7	F	F	Т	F
8	F	F	F	F

Is there a better way?

Testing literature has lots of suggestions, some of them confusing or conflicting

The general idea is simple: *test each clause independently* from the other clauses

But what does "independently" mean?

Active Clauses

Clause coverage has a weakness: the values don't always make a difference

Consider the first test case for clause coverage where each clause was true

$$((5 < 10) \lor true) \land (1 >= 1*1)$$

If we change *only one clause*, then only changing the last clause can change the evaluation of the predicate!

Given these values, the third clause determines the predicate

Active Clauses

To really test the results of a clause, the clause should be the determining factor to the value of the predicate

Determination: a selected clause c_i in predicate p, called the *major clause*, determines p if and only if the values of the remaining *minor clauses* c_i are such that changing the value of c_i changes the value of p

Such a condition makes clause c_i active

Active Clauses

Consider $p = a \lor b$

If a=false, then b determines p If b=false, then a determines p

Consider $p = a \land b$

If a=true, then b determines p If b=true, then a determines p

The goal is to find test cases for each selected major clause when the clause determines the value of the predicate

This is formalized in a *family of criteria* that have subtle but important differences

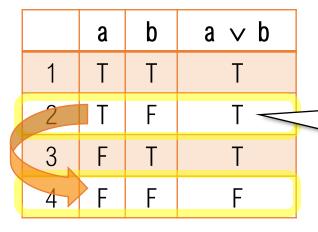
Active Clause Coverage Example

$$p = a \lor b$$

	a	b	a ∨ b
1	Т	Т	T
2	Т	F	T
3	F	Т	T
4	F	F	F

Active Clause Coverage Example

$$p = a \lor b$$



Select **a** as the major clause, then choose values for minor clause **b** such that changing **a** changes **p**, so **a** determines **p** and **a** is the active clause

Active Clause Coverage Example

$$p = a \vee b$$

	a	b	a ∨ b
1	Т	Т	T
2	I T	F	T <
3	F	T	T
4	F	F	F

Select **a** as the major clause, then choose values for minor clause **b** such that changing **a** changes **p**, so **a** determines **p** and **a** is the *active clause*

	a	b	a∨ b
1	Т	Т	T
2	Т	F	Т
3	F	Т	T <
4	F	F	F

Select **b** as the major clause, then choose values for minor clause **a** such that changing **b** changes **p**, so **b** determines **p** and **b** is the active clause

Calculating Determination

Finding values for minor clauses is easy for simple predicates, but hard for complex ones

Definitional approach:

 $p_{c=\mathrm{true}}$ is predicate p with every occurrence of clause c replaced by true

 $p_{c ext{=false}}$ is predicate p with every occurrence of clause c replaced by false

To find values for minor clauses, exclusive or $p_{c=\text{true}}$ and $p_{c=\text{false}}$

$$p_c = p_{c=true} \oplus p_{c=false}$$

After solving, p_c describes exactly the values needed for c to determine p

Consider $p = a \lor b$

Consider
$$p = a \lor b$$

 $p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$

Consider
$$p = a \lor b$$

 $p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$
 $p_a = (\text{true} \lor b) \oplus (\text{false} \lor b)$

Consider
$$p = a \lor b$$

$$p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_a = (\text{true} \lor b) \oplus (\text{false} \lor b)$$

$$p_a = \text{true} \oplus b$$

Consider
$$p = a \lor b$$

$$p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_a = (\text{true} \lor b) \oplus (\text{false} \lor b)$$

$$p_a = \text{true} \oplus b$$

$$p_a = \neg b$$

Consider
$$p = a \lor b$$

$$p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_a = (\text{true} \lor b) \oplus (\text{false} \lor b)$$

$$p_a = \text{true} \oplus b$$

$$p_a = \neg b$$

Thus a determines p when b=false

Consider $p = a \wedge b$

Consider
$$p = a \wedge b$$

 $p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$

Consider
$$p = a \land b$$

 $p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$
 $p_a = (\text{true} \land b) \oplus (\text{false} \land b)$

Consider
$$p = a \land b$$

$$p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_a = (\text{true} \land b) \oplus (\text{false} \land b)$$

$$p_a = b \oplus \text{false}$$

```
Consider p = a \land b
p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}
p_a = (\text{true} \land b) \oplus (\text{false} \land b)
p_a = b \oplus \text{false}
p_a = b
```

Consider
$$p = a \land b$$

$$p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_a = (\text{true} \land b) \oplus (\text{false} \land b)$$

$$p_a = b \oplus \text{false}$$

$$p_a = b$$

Thus a determines p when b = true

Consider $p = a \lor (b \land c)$

Consider
$$p = a \lor (b \land c)$$

 $p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$

Consider
$$p = a \lor (b \land c)$$

$$p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_a = (\text{true} \lor (b \land c)) \oplus (\text{false} \lor (b \land c))$$

```
Consider p = a \lor (b \land c)
p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}
p_a = (\text{true} \lor (b \land c)) \oplus (\text{false} \lor (b \land c))
p_a = \text{true} \oplus b \land c
```

```
Consider p = a \lor (b \land c)
p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}
p_a = (\text{true} \lor (b \land c)) \oplus (\text{false} \lor (b \land c))
p_a = \text{true} \oplus b \land c
p_a = \neg (b \land c)
```

```
Consider p = a \lor (b \land c)
p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}
p_a = (\text{true} \lor (b \land c)) \oplus (\text{false} \lor (b \land c))
p_a = \text{true} \oplus b \land c
p_a = \neg (b \land c)
p_a = \neg b \lor \neg c
```

Consider
$$p = a \lor (b \land c)$$

$$p_a = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_a = (\text{true} \lor (b \land c)) \oplus (\text{false} \lor (b \land c))$$

$$p_a = \text{true} \oplus b \land c$$

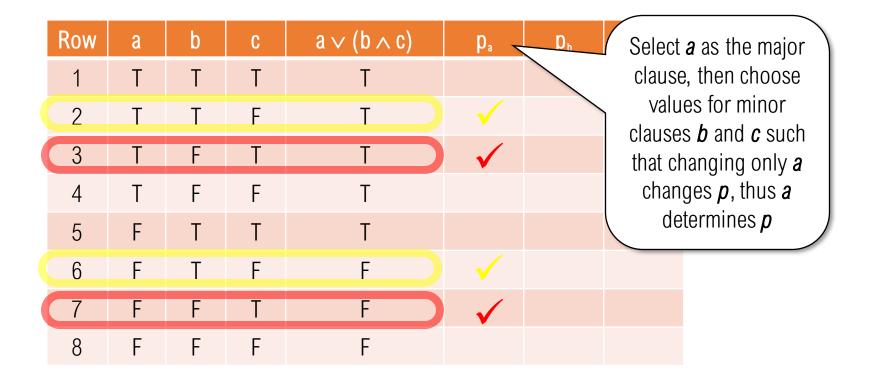
$$p_a = \neg (b \land c)$$

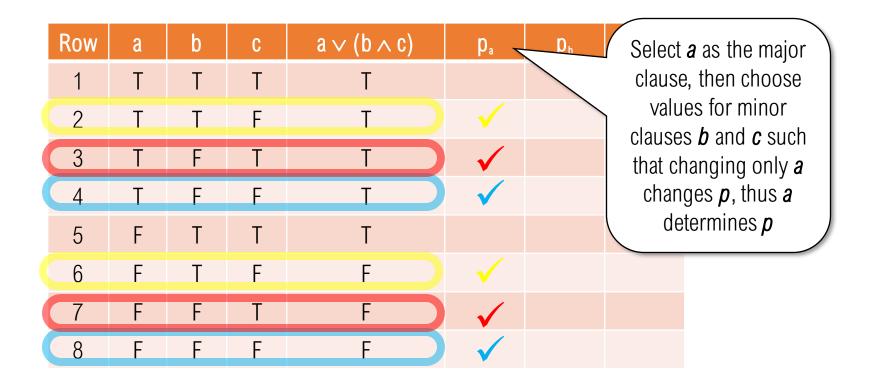
$$p_a = \neg b \lor \neg c$$

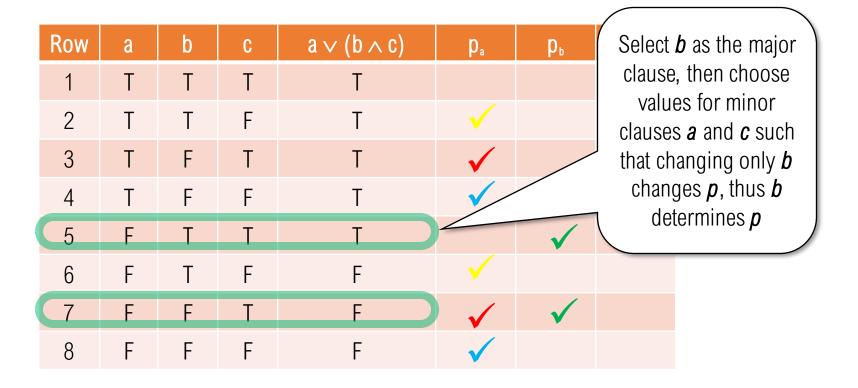
Thus a determines p when b=false or c=false

Row	a	b	С	$a \vee (b \wedge c)$	p _a	p _b	p _c
1	T	T	T	Т			
2	T	T	F	Т			
3	T	F	T	Т			
4	T	F	F	Т			
5	F	T	T	Т			
6	F	T	F	F			
7	F	F	T	F			
8	F	F	F	F			

Row	a	b	С	$a \vee (b \wedge c)$	p _a	Select a as the major
1	T	T	T	T		clause, then choose
2	T	T	F	T		values for minor clauses b and c such
3	T	F	T	T		that changing only a
4	Т	F	F	T		changes \boldsymbol{p} , thus \boldsymbol{a}
5	F	T	T	Т		determines <i>p</i>
6	F	T	F	F		
7	F	F	T	F		
8	F	F	F	F		







Row	a	b	С	$a \vee (b \wedge c)$	p _a	Рь	p _c
1	T	T	T	T			
2	T	T	F	Т			
3	T	F	T	Т	\checkmark		
4	T	F	F	Т	\checkmark		
5	F	T	T	T		\checkmark	
6	F	T	F	F			
7	F	F	T	F	\checkmark	\checkmark	
8	F	F	F	F	\checkmark		

						Select	\boldsymbol{c} as the r	najor \
Row	a	b	С	$a \vee (b \wedge c)$	p _a	clause, th	en choose	e values
1	T	T	T	Т			r clauses a	
2	T	Т	F	T	√		changing ges p , thu	
3	Т	F	Т	Т			termines ,	
4	Τ	F	F	Т	/			
5	F	T	T	T		√	\checkmark	
6	F	T	F	F			\checkmark	
7	F	F	Т	F	\checkmark			
8	F	F	F	F	√			

	Row	а	b	С	a □ (b □ c)	p _a	p _b		of color- ve can tag
Ī	1	T	Т	Т	Т				ning pairs
	2	T	T	F	Т	√(1) ~		with ID i	numbers
	3	Т	F	T	Т	√ (2)			
	4	T	F	F	Т	√ (3)			
	5	F	T	T	Т		√ (4)	√ (5)	
	6	F	T	F	F	√ (1)		√ (5)	
	7	F	F	T	F	√ (2)	√ (4)		
	8	F	F	F	F	√ (3)			

Active Clause Coverage

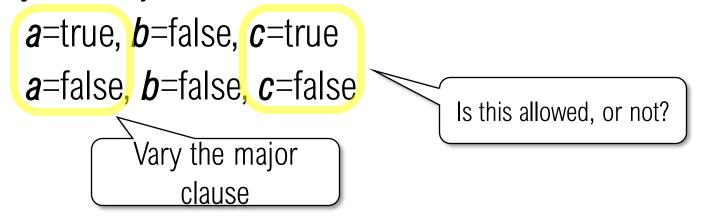
DEFINITION

Active Clause Coverage (ACC) – For each p in P and each major clause c_i in C_p , choose minor clauses c_j (j!=i) such that c_i determines p. TR has two requirements for c_i : c_i evaluates to true and c_i evaluates to false

MCDC Ambiguity

Do the minor clauses have to *retain the same values* while the major clause changes between true and false?

Example: given $p = a \lor (b \land c)$, if a is the major clause then when we vary the major clause:



This question has caused confusion among safety-critical testers for years

Resolving the Ambiguity

Three possible answers (which leads to three different coverage criteria)

Minor clauses do not need to be the same

Minor clauses can *force the predicate* to become both true and false

Minor clauses do need to be the same

General Active Clause Coverage

Minor clauses do not need to be the same

DEFINITION

General Active Clause Coverage (GACC) — For each p in P and each major clause c_i in C_p , choose minor clauses c_j (j!=i) such that c_i determines p. TR has two requirements for c_i : c_i evaluates to true and c_i evaluates to false. The values chosen for minor clauses c_j do not need to be the same when c_i is true as when c_i is false, and the value of p does not need to change.

It is possible to satisfy GACC without satisfying predicate coverage

Correlated Active Clause Coverage

Minor clauses can force the predicate

JEFINITION

Correlated Active Clause Coverage (CACC) — For each p in P and each major clause c_i in C_p , choose minor clauses c_j (j!=i) such that c_i determines p. TR has two requirements for c_i : c_i evaluates to true and c_i evaluates to false. The values chosen for minor clauses c_j must cause p to be true for one value of major clause c_i and false for the other value of c_i .

Subsumes predicate coverage This is "masking MCDC"*

Restrictive Active Clause Coverage

Minor clauses do need to be the same

DEFINITION

Restricted Active Clause Coverage (RACC) — For each p in P and each major clause c_i in C_p , choose minor clauses c_j (j!=i) such that c_i determines p. TR has two requirements for c_i : c_i evaluates to true and c_i evaluates to false. The values chosen for minor clauses c_j must must be the same when c_i is true as when c_i is false.

This is "unique-cause MCDC", the common interpretation of MCDC* Often leads to *infeasible test requirements*

Active Clause Comparison

To satisfy an active clause coverage criterion, pick a pair of tests where...

Criterion	Major clause determines p	Major clause changes value	Changing major clause changes p	Minor clauses are held the same
GACC	✓	✓		
CACC	✓	✓	✓	
RACC	✓	✓	✓	✓

Evaluation process

- 1. Select a major clause
- 2. Determine the conditions for the minor clauses where the major clause determines the predicate
- 3. For GACC, select a pair of conditions where the value of the major clause changes (the value of the minor clauses and *p* may or may not change)
- 4. For CACC, select a pair of conditions where the value of the major clause changes and the value of *p* changes (the minor clauses may or may not change)
- 5. For RACC, select a pair of conditions where the value of the major clause changes, the values of the minor clauses are held constant, and the value of p changes

Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

1. Select a major clause -- a

	a	b	С	p=(a∧b)∨(¬a∧¬b∧c)
1	Т	T	Τ	T
2	Т	T	F	T
3	T	F	T	F
4	Т	F	F	F
5	F	T	T	F
6	F	T	F	F
7	F	F	T	T
8	F	F	F	F

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$

2. Determine the conditions for the minor clauses where the major clause determines the predicate

Using the XOR approach:

$$p_{a} = p_{a=\text{true}} \oplus p_{a=\text{false}}$$

$$p_{a} = (T \land b) \lor (\neg T \land \neg b \land c) \oplus (F \land b) \lor (\neg F \land \neg b \land c)$$

$$p_{a} = b \lor (F \land \neg b \land c) \oplus F \lor (T \land \neg b \land c)$$

$$p_{a} = b \oplus \neg b \land c$$

$$p_{a} = b \lor c$$

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$

2. Determine the conditions for the minor clauses where the major clause determines the predicate

			_		
	a	b	С	$p=(a \land b) \lor (\neg a \land \neg b \land c)$	p _a
1	Τ	Τ	Τ	Т	√ 1
2	Τ	Τ	F	T	
3	Т	F	T	F	
4	T	F	F	F	
5	F	Τ	Τ	F	√ 1
6	F	T	F	F	
7	F	F	Τ	Т	
8	F	F	F	F	

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$

2. Determine the conditions for the minor clauses where the major clause determines the predicate

	a	b	С	p=(a∧b)∨(¬a∧¬b∧c)	p _a
1	T	T	T	Т	√ 1
2	T	T	F	Т	√ 2
3	T	F	Τ	F	
4	Τ	F	F	F	
5	F	T	Т	F	√ 1
6	F	Τ	F	F	√2
7	F	F	Τ	Т	
8	F	F	F	F	

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$

2. Determine the conditions for the minor clauses where the major clause determines the predicate

	a	b	С	p=(a∧b)∨(¬a∧¬b∧c)	p _a
1	T	T	T	Т	√ 1
2	T	T	F	Т	√ 2
3	Τ	F	Τ	F	√3
4	T	F	F	F	
5	F	T	Τ	F	√ 1
6	F	T	F	F	√ 2
7	F	F	Τ	T	√3
8	F	F	F	F	

Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

2. Determine the conditions for the minor clauses where the major clause determines the predicate

Using the truth table approach:

		a	b	C	p=(a∧b)∨(¬a∧¬b∧c)	P _a
	1	Τ	Τ	Τ	T	√ 1
	2	Τ	Т	F	Т	√ 2
	3	Τ	F	Τ	F	√ 3
-	4	Τ	F	F	F	
	5	F	Τ	Т	F	√ 1
	6	F	Τ	F	F	√ 2
	7	F	F	Τ	T	√ 3
	8	F	F	F	F	

Select inputs such that \boldsymbol{a} changes, \boldsymbol{b} and \boldsymbol{c} do not change... but \boldsymbol{p} DOES NOT change, thus \boldsymbol{a} DOES NOT determines \boldsymbol{p} when $\neg \boldsymbol{b} \wedge \neg \boldsymbol{c}$

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$

2. Determine the conditions for the minor clauses where the major clause determines the predicate

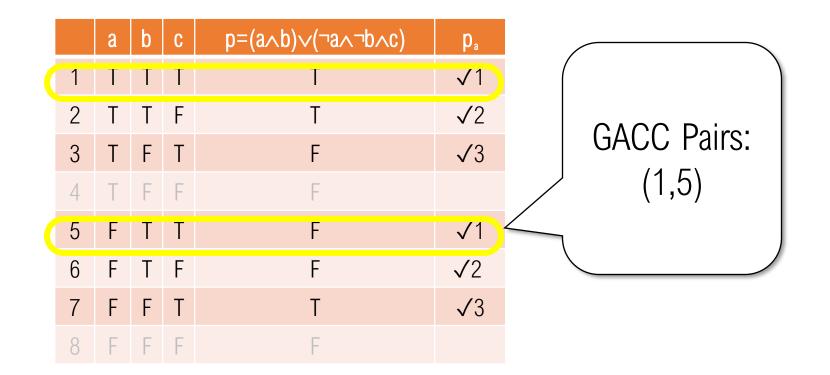
	a	b	С	p=(a∧b)√(¬a∧¬b∧c)	p _a
	a	U	U	$p-(a\wedge b) \vee (-a\wedge -b\wedge c)$	Pa
1	Т	T	T	Т	√ 1
2	Т	T	F	Т	√ 2
3	T	F	Τ	F	√ 3
4	T	F	F	F	
5	F	T	Τ	F	√ 1
6	F	T	F	F	√ 2
7	F	F	Τ	Т	√ 3
8	F	F	F	F	

To satisfy an active clause coverage criterion, pick a pair of tests where...

Criterion	Major clause determines p	Major clause changes value	Changing major clause changes p	Minor clauses are held the same
GACC	√	✓		
CACC	√	✓	✓	
RACC	✓	✓	✓	✓

Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

3. For GACC, select a pair of conditions where the value of the major clause changes (the value of the minor clauses and *p* may or may not change)



Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

	a	b	С	p=(a∧b)∨(¬a∧¬b∧c)	p _a
1	Т	Τ	Τ	T	√ 1
2	Т	Τ	F	Т	√2
3	Т	F	T	F	√ 3
4	Т	F	F	F	
5	F	Т	Τ	F	√ 1
6	F	Τ	F	F	√ 2
7	F	F	T	Т	√3
8	F	F	F	F	

Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

	a	b	C	p=(a∧b)∨(¬a∧¬b∧c)	p _a	GACC Pairs:
1	Τ	Τ	Τ	T	√ 1	(1,5) or (1,6)
2	T	Τ	F	Т	√ 2	or (1,7) okay
3	T	F	T	F	√ 3	
4	Τ	F	F	F		because p
5	F	Τ	Τ	F	√ 1	does not
6	F	Τ	F	F	√ 2	need to
7	F	F	Τ	T	√3	change
8	F	F	F	F		

Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

3. For GACC, select a pair of conditions where the value of the major clause changes (the value of the minor clauses and *p* may or may not change)

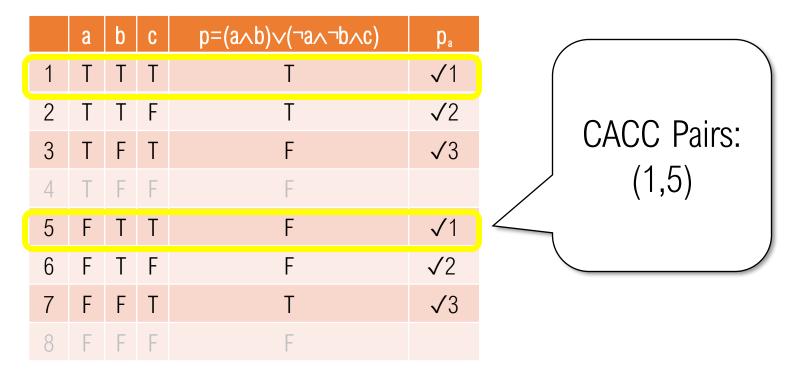
	a	b	С	p=(a∧b)∨(¬a∧¬b∧c)	p _a
1	Τ	Τ	Τ	T	√ 1
2	T	Τ	F	Т	√ 2
3	T	F	Τ	F	√3
4	Τ	F	F	F	
5	F	Τ	Τ	F	√ 1
6	F	T	F	F	√ 2
7	F	F	Τ	Т	√3
8	F	F	F	F	

GACC Pairs: any one of (1,5), (1,6), (1,7), (2,5), (2,6), (2,7), (3,5), (3,6), (3,7)

To satisfy an active clause coverage criterion, pick a pair of tests where...

Criterion	Major clause determines p	Major clause changes value	Changing major clause changes p	Minor clauses are held the same
GACC	\checkmark	√		
CACC	✓	✓	✓	
RACC	✓	✓	✓	✓

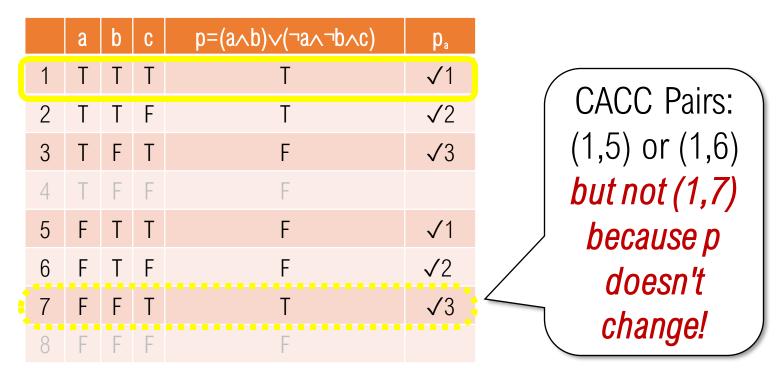
Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$



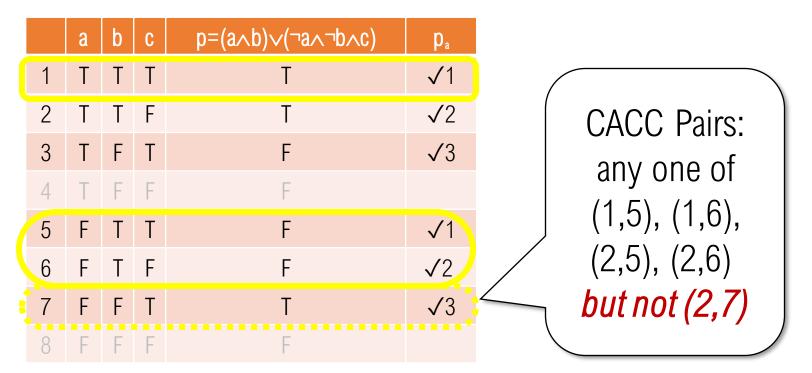
Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

	a	b	С	p=(a∧b)∨(¬a∧¬b∧c)	p _a
1	Τ	T	Τ	T	√ 1
2	Τ	T	F	T	√2
3	Т	F	Т	F	√ 3
4	Т	F	F	F	
5	F	T	Τ	F	√ 1
6	F	T	F	F	√2
7	F	F	Τ	T	√3
8	F	F	F	F	

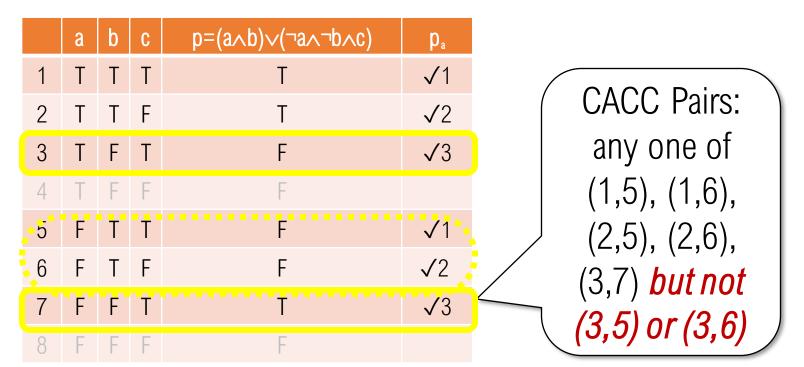
Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$



Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$



Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

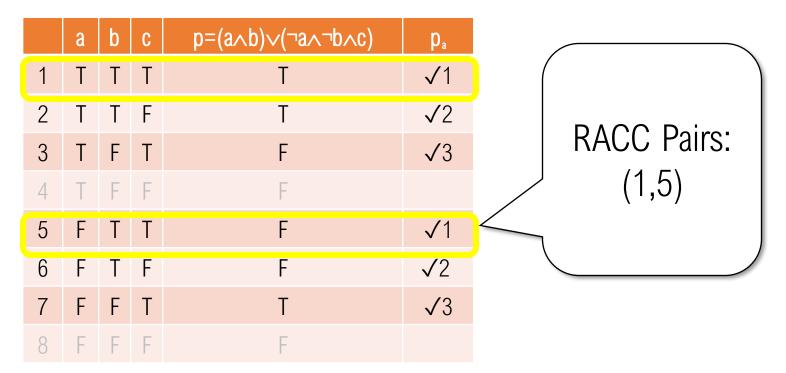


To satisfy an active clause coverage criterion, pick a pair of tests where...

Criterion	Major clause determines p	Major clause changes value	Changing major clause changes p	Minor clauses are held the same
GACC	✓	✓		
CACC	\checkmark	✓	✓	
RACC	\checkmark	\checkmark	✓	\checkmark

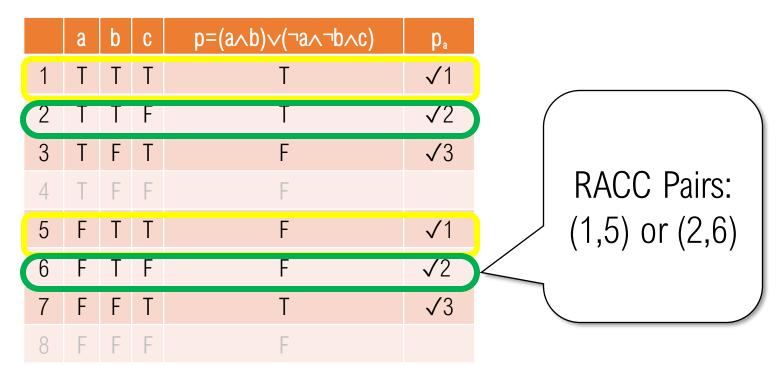
Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

5. For RACC, select a pair of conditions where the value of the major clause changes, the values of the minor clauses are held constant, and the value of p changes



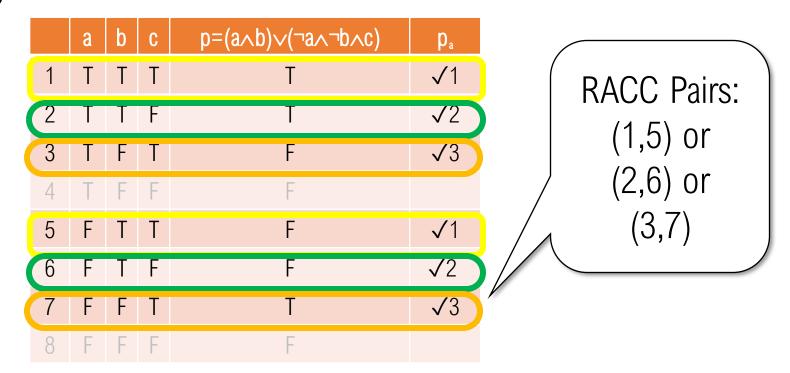
Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

5. For RACC, select a pair of conditions where the value of the major clause changes, the values of the minor clauses are held constant, and the value of p changes



Consider $p = (a \land b) \lor (\neg a \land \neg b \land c)$

5. For RACC, select a pair of conditions where the value of the major clause changes, the values of the minor clauses are held constant, and the value of p changes



Inactive Clause Coverage

Taking the opposite approach — major clauses *do not affect* the predicates

EFINITION

Inactive Clause Coverage (ICC) – For each p in P and each major clause c_i in C_p , choose minor clauses c_j (j!=i) such that c_i does not determine p. TR has four requirements for c_i : (1) c_i evaluates to true with p true, (2) c_i evaluates to false with p true, (3) c_i evaluates to true with p false, and (4) c_i evaluates to false with p false.

Why bother? It's useful for testing safety interlock systems to ensure that during certain circumstances a control variable does not have any effect on operation

General Inactive Clause Coverage

DEFINITION

General Inactive Clause Coverage (GICC) — For each p in P and each major clause c_i in C_p , choose minor clauses c_j (j!=i) such that c_i does not determine p. The values chosen for minor clauses c_j do not need to be the same when c_i is true as when c_i is false.

DEFINITION

Restricted Inactive Clause Coverage (RICC) – For each p in P and each major clause c_i in C_p , choose minor clauses c_j (j!=i) such that c_i does not determine p. The values chosen for minor clauses c_j must be the same when c_i is true as when c_i is false.

Inactive Clause Comparison

To satisfy an inactive clause coverage criterion, pick a pair of tests where...

Criterion	Major clause determines p	Major clause changes	Changing major clause changes p	Minor clauses are held the same
GICC)		
RICC				

By definition, if the major clause does not determine p, then changing the major clause will not change p

GICC Example

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$
Selecting clause c as the major clause

For GICC, select a pair of conditions where the major clause does not determine p, the value of the major clause changes, and the value of p does

not change

	a	b	С	$p=(a \land b) \lor (\neg a \land \neg b \land c)$	p _c	
1	T	Τ	T	T		
2	Τ	T	F	T		CICC Daire
3	Τ	F	T	F		GICC Pairs:
4	T	F	F	F		(1,2)
5	F	Т	T	F		
6	F	Т	F	F		
7	F	F	Т	Т	√ 1	c determines p for
8	F	F	F	F	√ 1	rows 7 and 8

GICC Example

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$

For GICC, select a pair of conditions where the major clause does not determine p, the value of the major clause changes, and the value of p does

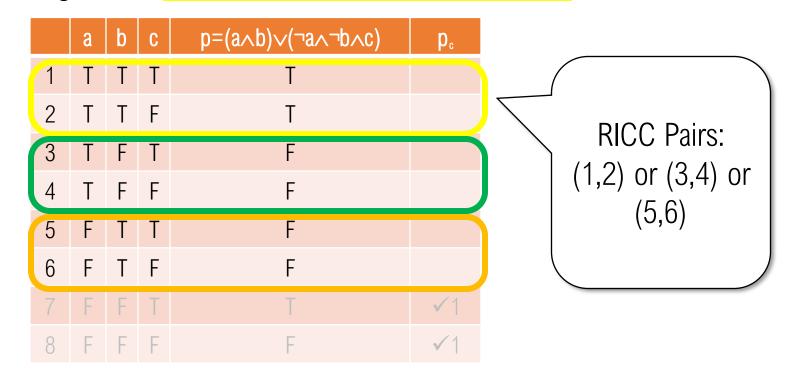
not change

	a	b	С	p=(a^b)~(¬a^¬b^c)	p _c
1	Т	T	T	Т	
2	Т	T	F	Т	
3	Т	F	Τ	F	
4	Т	F	F	F	
5	F	Τ	Τ	F	
6	F	T	F	F	
7	F	F	Τ	Т	√1
8	F	F	F	F	√ 1

RICC Example

Consider
$$p = (a \land b) \lor (\neg a \land \neg b \land c)$$

For RICC, select a pair of conditions where the major clause does not determine p, the value of the major clause changes, and the value of p does not change, and the minor clauses are the same



Infeasibility

Infeasible test requirements are common

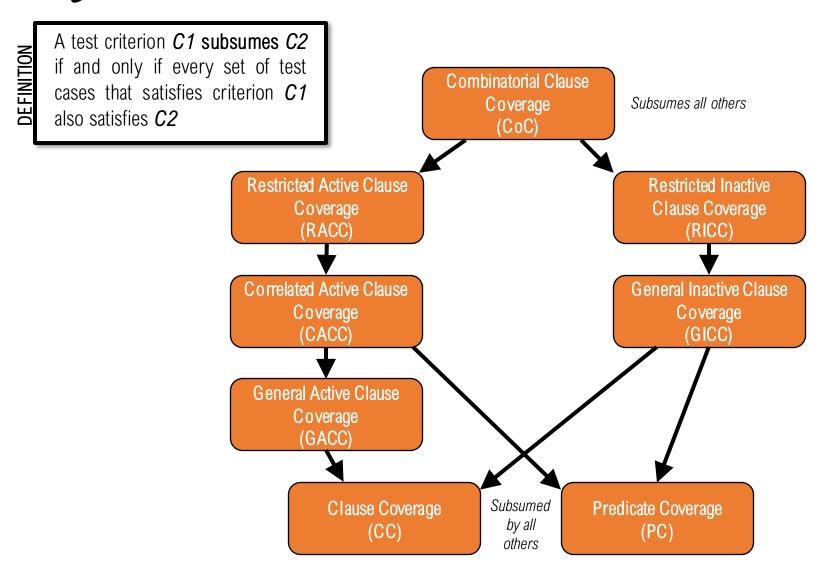
Given
$$p = (a > b \land b > c) \lor (c > a)$$

If $(a > b)$ =true and $(b > c)$ =true, then $(c > a)$ =true is infeasible

As with ISP and graph criteria, infeasible test requirements must be *recognized and ignored*

Recognizing infeasible test requirements is difficult, and in general *undecidable*

Logic Criteria Subsumption



Logic Coverage Summary

Predicates are often very simple – in practice, most have fewer than 3 clauses

That's good news, because fewer clauses significantly simplifies testing

With only one clause, predicate coverage is sufficient

With 2 or 3 clauses, combinatorial coverage may be practical

With more complex clauses, ACC and ICC criteria are practical

Testing safety-critical software often requires MCDC (or RACC or CACC)