

## Week - 1

L 1.1

30/6/21

CLASSTIME Pg. No. 1

Date / /

## 1. Introduction to datasets

- patterns of computation becomes easier
    - 1. School data - name, gender, DOB, city, MM MP MC MT Total etc
    - 2. Shopping data - shop name, cust name, item, category
    - 3. Word data - word, type of speech, no of letters
  - this data to illustrate CT

## 2. Introduction to course

- CT is easiest for computer to how to write program
  - every computer or program requires steps  $\Rightarrow$  details  $\Rightarrow$  depends on how much you have expertise
  - > programming  $\rightarrow$  lang. to talk to computer
  - pattern - sequence of steps . Eg. Books , clothes

$\Rightarrow$  Video 3 & 4 are repetitions of 1 & 2 respectively

~~20-5~~ L 1-2

- Concept of variables, iterators & filtering
  - to know how many students in a class?  $\text{HTT} \Rightarrow 5$
  - Iteration - giving them the sequence of objects. u have <sup>init.</sup> and for each obj. u keep a count (gen. from 0)
  - Variable - quantity is a no. (maybe) which can change eventually. (value keeps changing)
  - Eg. Avg. marks of maths = sum of all maths marks / total students
  - keeping track of <sup>see more</sup> 2 variables in single iteration
  - in the word data  $\Rightarrow$  want only words  $\Rightarrow$  concept of filtering  $\Rightarrow$  to catch some particular thing / from data

## Video - 6:

6. Tutorial feel Video - 5. 18 minutes

(a) no food item to purchase  $A_F$  count  $\Rightarrow \star \dots 10$

(b) only " "

$$NF_{\text{sum}} = 1 \dots 6$$

- Eg. Calculate the avg. word length of desired data  
 $\text{Word count} = \alpha \times 2.65$   
 $T_{\text{Total letters}} = \alpha \times 2.65 \cdot 329 \quad \text{Avg} = \frac{T_{\text{Total}}}{\text{WC.}} = \frac{329 \times 5}{65} \approx 25$

Video 7 b1.3

- Iterations using combination of filtering conditions
- some prob. we'll take up  $\Rightarrow$  then iterate
- class data  $\Rightarrow$  Eg. how many Girls from Chennai
  - $\hookrightarrow$  gender, city  $\Rightarrow$  2 items + 2 filters (5/30)
  - $\hookrightarrow$  now to do in 1 step  $\Rightarrow$  'and' condition  
both should be true  $\Rightarrow$  to select the card
- Eg. been b/w. the half year
  - $\hookrightarrow$  Male =  $\varnothing X \dots 8$
  - $\hookrightarrow$  Female =  $\varnothing Y \dots 7$
- Eg. Girls / Boys doing better in Maths
  - $\hookrightarrow$  not look at max.
  - $\hookrightarrow$  avg. of boys/girls and then compare  
boys/girls count, boys/girls TMaths marks  
in 1 iteration
- algorithm problem - systematic procedure

Video 8

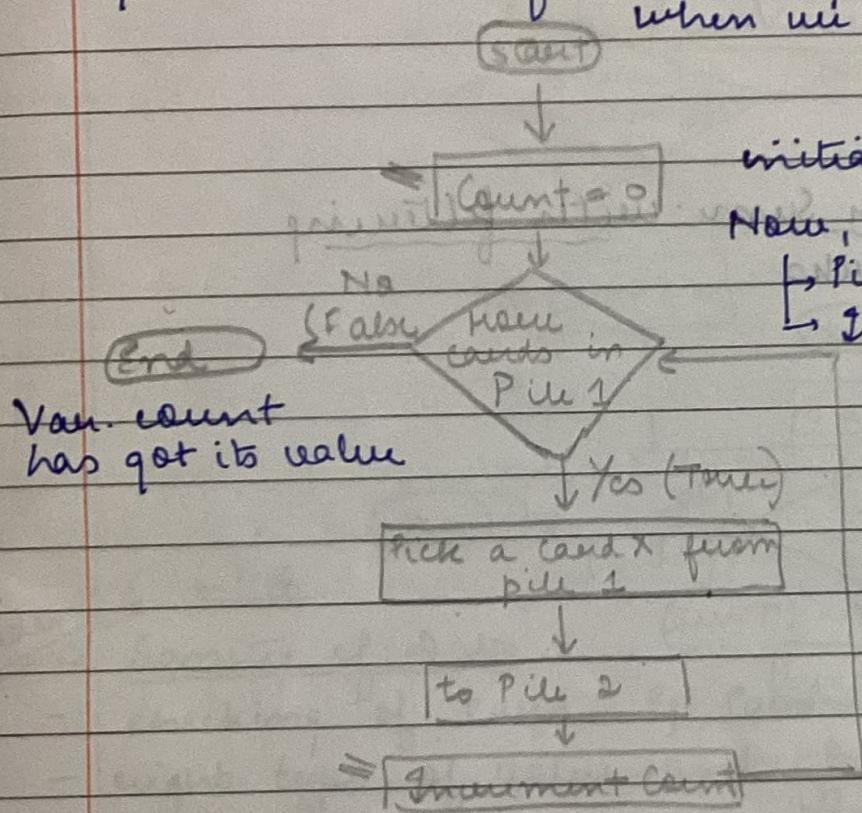
### 8. Tutorial for Video 7

- Eg. Word data - Eg. no. of adjectives,  $> 4$  letters
  - $\hookrightarrow$  LargeAdjectiveCount =  $\varnothing X \dots 7$
- Bill data - Eg. How much Sunriser spent -
  - $\hookrightarrow$  SVS Stores - ₹ 843 BigBazaar - ₹ 342694
- Word data - Eg. compare the size of avg. preposition  
and avg. pronoun (in one iteration)
  - $\hookrightarrow$  Prep.Count  $\varnothing X \dots 10$  ✓      PrepLength  $\varnothing X \dots 27$
  - $\hookrightarrow$  ProCount  $\varnothing Y \dots 7$  ✓      ProLength  $\varnothing X \dots 17$
  - $\hookrightarrow$  After that calculate avg. for comparing

Video-9 L 1.4

### 9. Introduce " to Flowcharts

- there is initializa " step, then we keep repeating
- step wise procedures in formal way
  - Flowcharts and Pseudocode
- symbols in flowcharts
  - Process / Activity  $\Rightarrow$  Set of op's that s the value of data.
  - Flowline / Arrow  $\Rightarrow$  Shows the order of exec' of program steps
  - Decision  $\Rightarrow$  determines which path the program will take
  - Terminal  $\Rightarrow$  indicates 'start' or 'end'
- Eg. Flowchart for counting cards

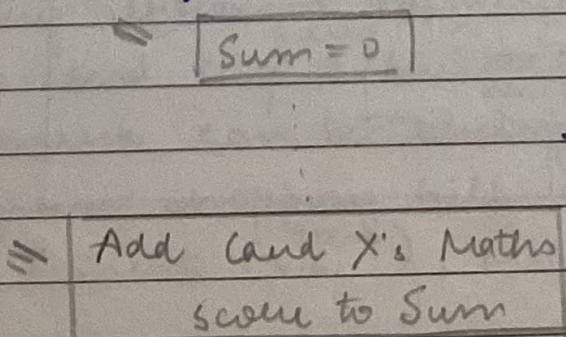


Vari. count has got its value

initializes the count to 0  
Now, iterator starts  
Pile 1  $\rightarrow$  Pile 2  
Inc. the value of var.

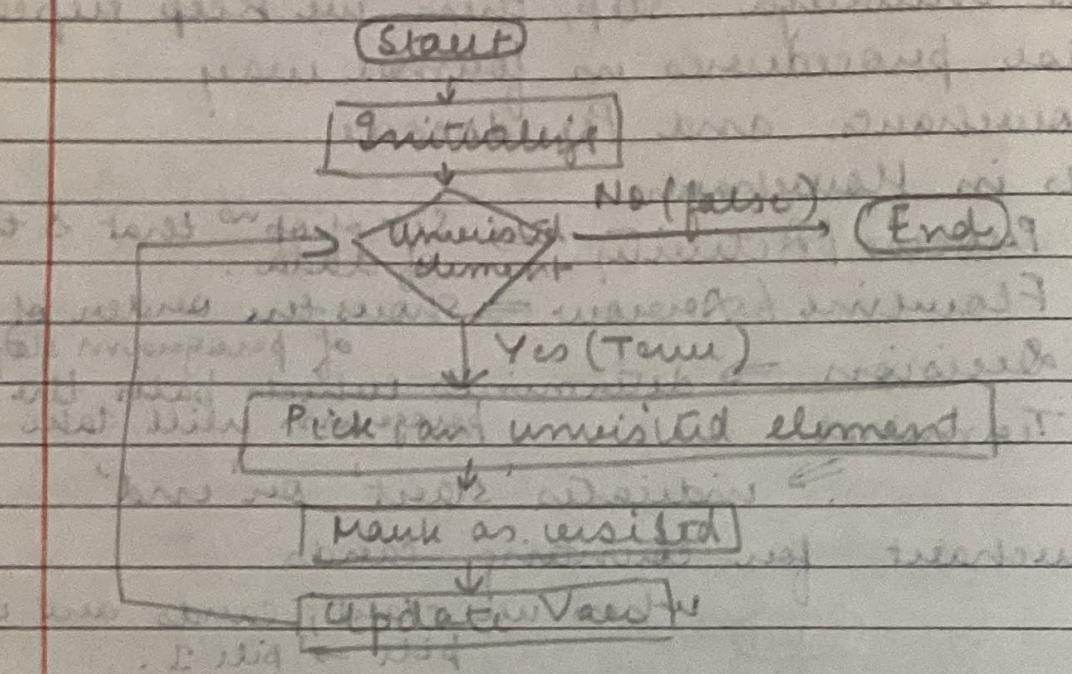
go back to "Start"

- Eg. Flowchart for sum of maths marks



Rest remains the same  
Accumulate the sum

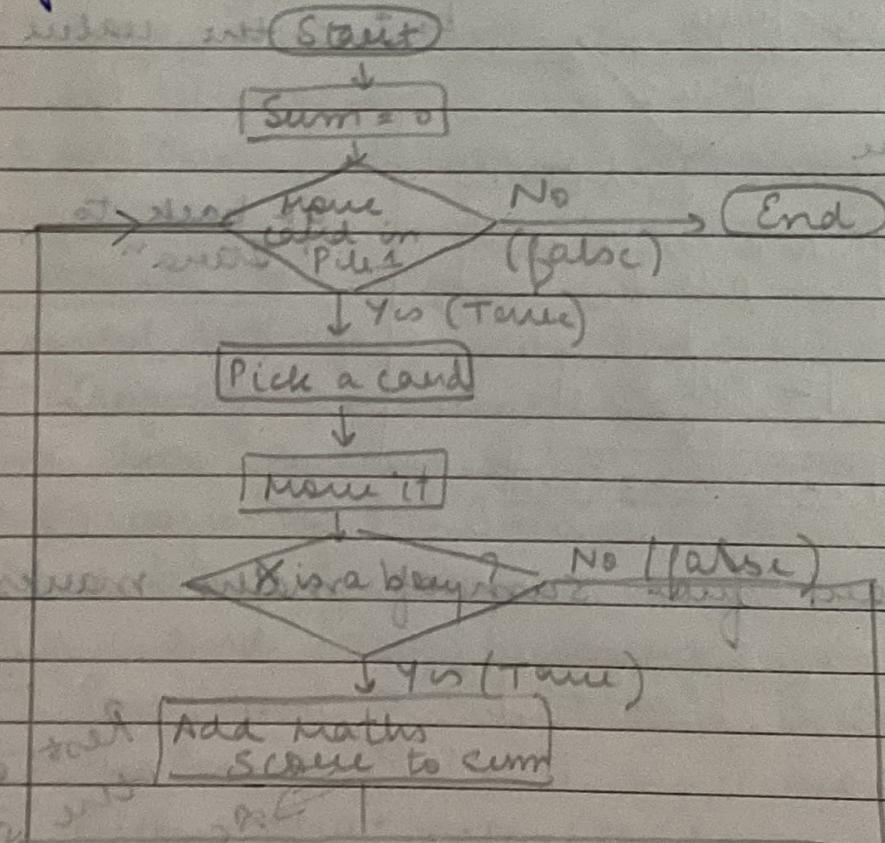
- Generic flowchart for iteration



~~Video-10~~ L1-5

- Flowchart for sum with filtering

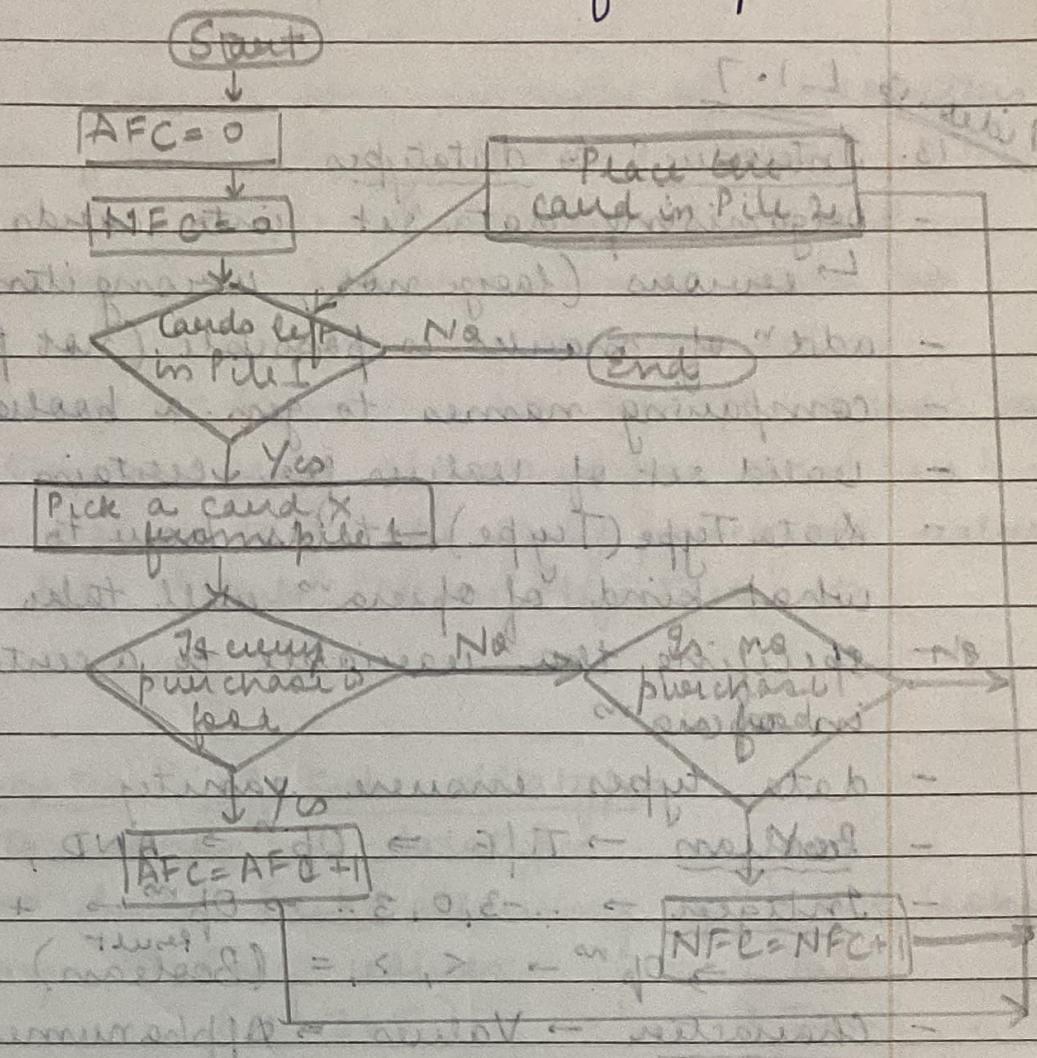
- By Maths marks



Video-11

11. Tutorial

- Eg. All Food Purchase and no food purchase.



Video-12

L1.612. Sanity of Data

- checking of card. Eg. Card no. → cannot be Q
- eight type of value and further they should be in range
- typing mistakes → sanity of numbers [8th. is money]
- data entry error → name pasted in city
- Eg. D.o.b. and city ⇒ swapped values
- Total mark can be wrong (added wrongly)
- Errors from shopping bill data - wrong item name, T-shirts (3.6<sup>term</sup>), Milk (Qty → 2.1.8), etc.

- Word data set - card no. (3F), considered  
Swami → Noun (x Adj.), open (letter count = 6x)

### Video 13 L 1.7

#### Introducing to datatypes

- Organising data set into cards  
↳ errors (large nos., wrong item, extra decimal)
- add<sup>n</sup> of marks is possible (not just multiplication)
- comparing names to gen. a boolean-type (T/F)
- valid set of values for certain elements
- Data Type (Type) → tells how to use that element, what kind of opera<sup>n</sup>s will take place
- specifying the variable to a certain constraint in opera<sup>n</sup>s
- data types ensure sanity
- Boolean → T/F  $\Rightarrow$  Op<sup>n</sup>  $\rightarrow$  AND, OR
- Integer → ... -3, 0, 3 ...  $\Rightarrow$  Op<sup>m</sup>  $\rightarrow$  +, -, ×, ÷  $\Rightarrow$  some constraints  
 $\Rightarrow$  Op<sup>n</sup>  $\rightarrow$  <, >, = (Boolean)  $\frac{21}{7} = 3 \Rightarrow \frac{22}{7}, 3$
- Character → Values  $\Rightarrow$  Alphanumeric A, a, 1, 0, etc.  
 $\Rightarrow$  Spec. charac. . ; , \$, #, @  
 $\Rightarrow$  Op<sup>n</sup>  $\rightarrow$  = (Result Boolean)

### Video 14 L 1.8

#### Subtypes of basic datatypes

- Sequence number  $\rightarrow$  0 - Range (Reasonable)  
↳ no integer of <sup>m</sup> doesn't make sense.
- Marks  $\rightarrow$  +ve, 0 - 100 (For one subject)  
↳ + (add), - but ×, ÷ no sense
- Count  $\rightarrow$  +ve, 0 - Max Range  
↳ +, -  $\Rightarrow$  Integer Result, <, >, =  $\Rightarrow$  Boolean

All these were integers

- Characters → Gender → Values M/F  
↳ = Boolean Result
- String - Values → any sequence of characters  
↳ char in string? ↳ Boolean Result  
↳ = ⇒ Boolean Result
- Name (Strings) → Values → strings with no spcl. characters
- City (Strings) - Strings with no spcl. charac. value
- Words - Strings with alphanumeric and . , ; ...
- Category - Values can take only precribed values like "Verb", "Noun", "Adjective", etc.

### L 1.9 Transformation of sub-datatypes

- we have to <sup>to</sup> work on sub-datatypes
- Integer → Date (like allowing op<sup>n</sup>) but looks like string ⇒ Date values 0 → 1 Jan  
(0-365) 1 → 2 Jan ...
- point(0) = "1 Jan"      op<sup>n</sup>      Result  
point(31) = "1 Feb"      point      String  
                        <, >, =      Boolean
- Int → Frac. marks 62.5 → what to do?
- we can use type for real nos. ⇒ Float
  - ↳ we might have 2 decimal places
  - ↳ multiply with 100 ⇒ it becomes int. ⇒ now perform integer operations

62.5 → 6250      op<sup>n</sup>      Result  
 point(6250) = "62.5"      point      String  
                         + -      Marks  
                         >, <, =      Boolean

- Int → Amount ⇒ same 27.50 ⇒ 2750 ...

- Int.  $\Rightarrow$  Quantity  $\Rightarrow$  could be frac. " also  
 $\hookrightarrow 1.25 \rightarrow 125$  stone, print (125) = "1.25"

L 1.10

### Introduc<sup>n</sup> to complex datatypes

- 2 ways of bundle  $\rightarrow$  record & list
- record  $\rightarrow$  data items with no. of elements
  - $\hookrightarrow$  data type with multiple fields  $\begin{cases} \text{name} \\ \text{value} \end{cases}$
  - $\hookrightarrow$  also cfd. struct & tuple
- sanity of field values is ensured by data-type specifications

Marks card  
(Record)

Name of field	Type
sig. no.	Seg No
name	Name
gender	Gender
dob	Date
city	City
marks1	
marks2	
c	
T	
	4] Marks

- Word in Para  
(Record)
- List - a seq. of data elements (a seq. of record)  
 $\hookrightarrow$  Marks lists  $\rightarrow$  is the data type for over data set of all marks cards  
 $\rightarrow$  Each element in the seq. is of marks card record data type
- Shopping bill  
(Record)
- Record types  $\rightarrow$  collec<sup>n</sup> with same or diff. type  
 $\hookrightarrow$  List type - seq. of items, with same data type.

Item  $\begin{cases} \text{Name} \\ \text{Category} \\ \text{Cost} \\ \text{Quantity} \end{cases}$

# Week - 2

Needs Rev.

CLASSTIME Pg. No. 9  
Date / /

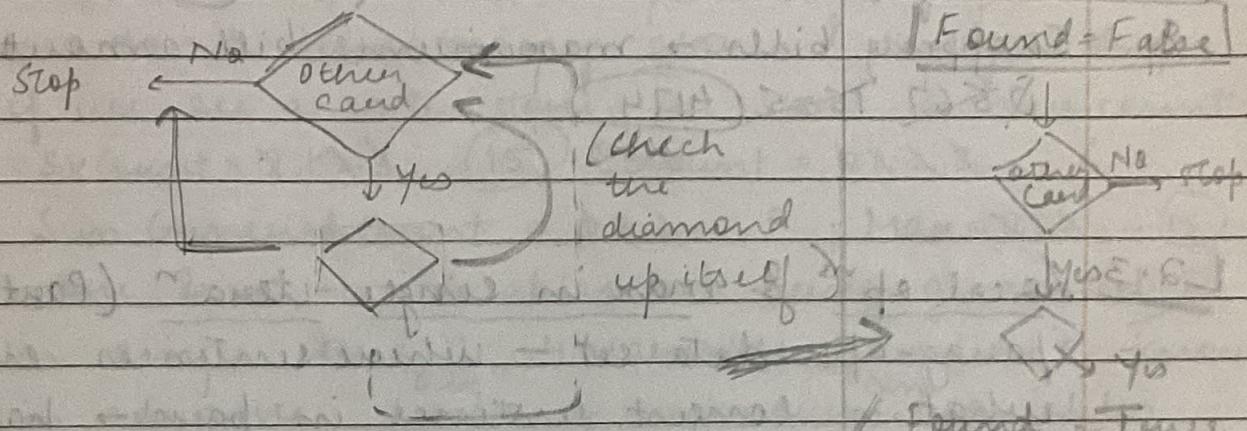
## L2.1 Conditional termina " in Itera"

- some search for a specific student  $\Rightarrow$  for criteria
- Eg. Student with  $> 90$  in all subj.  $\rightarrow$  stop or not?
- Eg. any customer  $\rightarrow$  food + apparel + lottery (3)  $\rightarrow$  1 stop
- Eg. First web in third sentence. Scount = 2  $\rightarrow$  now web
- $\hookrightarrow$  go on to second sentence find the web in third sentence  $\Rightarrow$  word No. 12  $\Rightarrow$  no need to go ahead

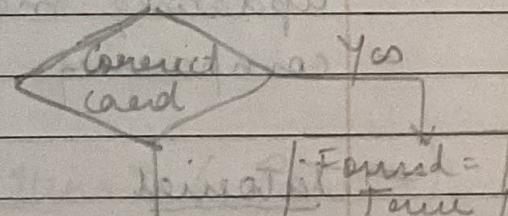
Bit unclear

Find / Search Pattern

Card Found?



- conditioning itself in doing
- this gives the way to stop
- "itera" with a "termina"
- if I stop, that means I found it  $\Rightarrow$  no student then found would still be false
- exist or not  $\Rightarrow$  kind of check



- L2.2 Local op<sup>ns</sup> & max in single item (Part 1)
- Ex. Best total marks - 270 252 276 (28)  
 → contains check & when marks beats it, replace it  
 → keep variable - replace the last mark.  
 → filtering + checking against the variable

### Tutorial

- Ex. Words data → find the word with maximum no. of letters → max = 22 X (11)  
 Ex. Shopping bills → maximum bill amount = 567 7525 (4174)

- L2.3 Local op<sup>ns</sup> & max in single item (Part 2)
- paragraph data set - when sentences ends, what is longest sentence in para → has most no. of words
  - hit full stop → keep count separate → better way
  - keep only max. no. of words till then, longest sentence so far, Max = 478 20 (25)  
 Count = 1, 2, 3, 4 ; 1, 2, 3, 4, ... so on
  - keeping track of longest sentence, count variable is looking changed (resetting) and updating max count, in 1 item
  - comparing 1 variable with other variable

### ⇒ Tutorial

- ex. of words until a fullstop is taken to be sentence
- Ex. Score data set & shopping bill is bit problematic to perform, because words are in a paragraph so they are in a ordered way

Eg. Scores dataset are arranged in a month way.  
To find out max. no. of students born in a month  
Max = 214                          Count = 21; 21; 21... so on

### L2.4 Local Ops & Max in single item^n (Part 3)

- which shop is doing well  $\rightarrow$  how many people go there?  $\times$   $\rightarrow$  how many bills they generate
- we don't have many stores? segregate shops if we see a new store, we start a new count
- SVCount = 2123... 15      BBCount = 2123... 6  
Sum General Count = 212... 9 ; Max = 212... 15
- more & more variables are being generated
- in single item^n, 4 variable, we didn't know in advance, 1<sup>single</sup> max & unknown<sup>single</sup> variables

### Tutorial

- Find the city with max. no. of students in Scores dataset; Max = 212345678
- Enode = 212      Salem = 21      Thenmai = 21234
- Madurai = 21234      Ambur = 21      Villupuram = 2123
- Bengaluru = 212345      Tiruchy = 2123      Teni = 21
- Nagpurcoil = 21

### L2.5 Local Ops and Max in Single Item^n (Part 4)

- Shopping bills  $\rightarrow$  SV stores had 15 bills, value of bills was less, but diff. from BB
- total value of bill the shop earned

$$SV = \phi 567 + 1341 = 1908 \dots 7120 \quad SG = \alpha 354 \dots 4009$$

$$\boxed{BB} \quad \alpha 1525 = 5699 \quad 12486 \quad Max = \phi 567 \quad 1525 \quad 1908 \quad 12486$$

- "itera" to add all bills and compare with the max and further update it.

### → Tutorial

- words data set → which part of speech contributes most charac. to the text? letter sound
- Pronoun =  $\alpha 2 \dots$       Max =  $\alpha 23 * 7 = 1324 \dots 145$
- Verb =  $\alpha 3 \dots 35$       Noun =  $\alpha 3124 \dots 145$
- Adjective =  $\alpha 9 \dots 65$       Preposin =  $\alpha 2 \dots 27$
- Adverb =  $\alpha 9 \dots 13$       Article =  $\alpha 3 \dots 18$
- Conjunc =  $\alpha 3 \dots 9$

### L2.6 Max in 1 & 2 iteras (non-nested)

- which card is most freq. in paragraph dataset; see each card and make new variables accordingly
- creating a variable for each word
- count is becoming large, very tricky to not make mistakes
- 2<sup>nd</sup> iteras, keep variable max at change it as and when req.
- so we can do both the iteras in 1 go
- Max = 6 (the)
- we don't know which word, but that can be done by separate variable

### ⇒ Tutorial

- Shopping bill → no. of bills / customers
- when we see new customer, we will make a new variable each time, if that person comes again, we'll add the count
- we can do the max. and inc. no. of variables in "item" Max = 4 ("item") → Ahmed

### L2.7 Max. in single "item" w/o losing info.

- so who got max marks? not only max. marks, but the person who is associated to it.
- if we update max, and change card no. (unique)
- Max Maths = 98 97      Max.Card No  $\Rightarrow$  11, 23
- 2 people who got max. maths marks, if we got 99 marks, then replace both people.
- SEO → applica<sup>n</sup> of max. frequency, keep relating unique words gives best results
- Google engines gives best relevant words
- we need to keep 2 items → 1<sup>st</sup> to find freq. count; 2<sup>nd</sup> to use as an applican<sup>n</sup>

### ⇒ Tutorial

- max. bill amt. & which customer has paid that max. amt.

$$\text{Max.bill} = \text{Rs } 567 \text{ 1525}$$

$$\underline{\underline{4174}}$$

Cus. - Naveen Suneetha

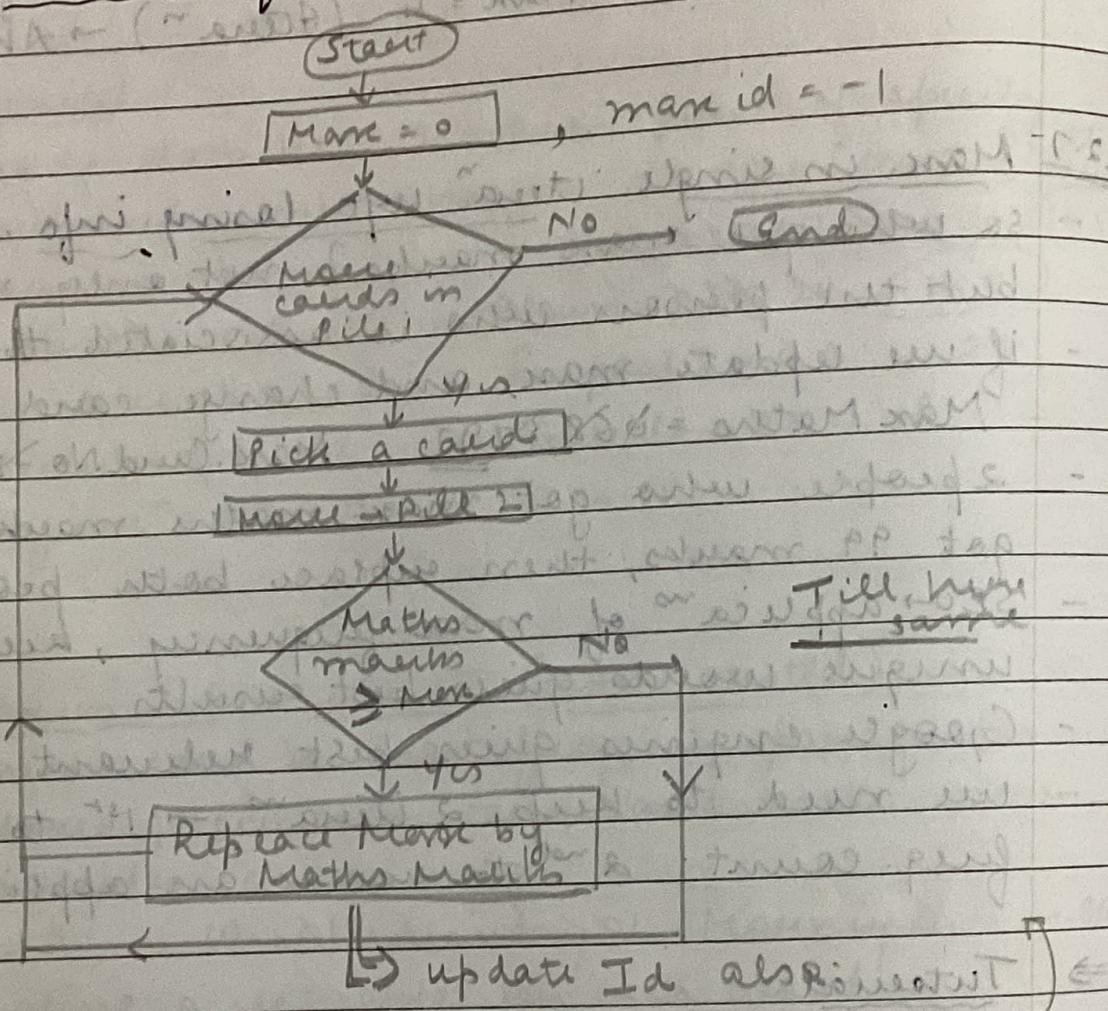
Sudipto Akshaya

### ⇒ Tutorial

- initialisa<sup>n</sup> of variables; count = 0; sum = 0
- Name = (String) Naveen
- Min - - - Max in chemistry scores

- Min (43) ----- Max (93)  
 = marking to 35x                          = marking to 94x  
 0. Range 100  
 Max marking  
 Min marking

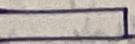
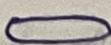
## L2.8 / concurrent for mark marks



- keep track of card which has the max marks
- we can keep diff. no. of people who had scored highest marks
- append of a & b is a,b (helps to create a list)

L 2.9

## Introduc<sup>n</sup> to Pseudocode

- pseudocode - textual representa<sup>n</sup> of procedures
- node types - process 
- decision 
- terminal 
- arrows indicate opera<sup>n</sup> flow
- adv → visual representa<sup>n</sup> of computa<sup>n</sup>, easy to understand
- disadvantages → size problem, collabora<sup>n</sup> restricted, version req. (modifica<sup>n</sup> needed)
- Step 0 Start
  - " 1 Initialize count = 0
  - " 2 check cards in Pile 1
  - " 3 If no more card, Step 8
  - " 4 Pick a card x
  - " 5 Move to Pile 2
  - " 6 Increment count
  - " 7 go to step 2
  - " 8 End
- succinct nota<sup>n</sup> for computational processes
- better textual representa<sup>n</sup> for conditional & repeated execution.

Start

Count = 0

while (Pile 1 has more cards) {

Pick a card x from pile 1

Move x to pile 2

Increment Count

}

End

## L2.10 Pseudocode with "data" & filtering

- will always <sup>not</sup> use slant & end
- sum of maths marks

this is a  
assignment  
statement  
eg.  $C = C + 1$

Sum = 0  
while (Pile 1 has cards) {  
    Pick a card X  
    Move X to Pile 2  
    Add X Maths mark to Sum  
} Sum = Sum + Maths  
Maths field of card X

- sum of boys math marks

== Equality  
= Assignment

Sum = 0  
while (Pile 1) {  
    Pick a card X  
    Move to Pile 2  
    if (x.Gender == M) {  
        Sum = Sum + X.Maths  
    }  
}

- sum of both girls & boys marks

alternative  
branch  
for  
conditional  
execution

BoySum, GirlSum = 0  
while (Pile 1) {  
    Pick a card X  
    Move to Pile 2  
    if (x.Gender == M) {  
        BoySum = BoySum + X.Maths  
    }  
    else {  
        GirlSum = GirlSum + X.Maths  
    }  
}

- find the max. maths marks

MaxM = 0  
while (Pile 1) {  
    Pick a card X  
    Move to Pile 2  
    if (x.Maths > MaxM) {  
        MaxM = x.Maths  
    }  
}

MaxCard = 0  
while (Pile 1) {  
    Pick a card X  
    Move to Pile 2  
    if (x.Gender == M) {  
        if (x.Maths > MaxCard) {  
            MaxCard = x.Maths  
        }  
    }  
}

## → Tutorial

- count = 0

while (Pile 1 has cards) {

Cards X → Pile 2

if (X.TownCity == "Chennai") {

if (X.Gender == "Female") {

count = count + 1;

}

}

Move X to Pile 2

}

- we could do 2 if statements in 1 way

if (X.TownCity == "Chennai" & X.Gender == "Female") {

L 3.1 Structure of datasets in table

- each card has is a unit of info, diff attributes or fields → organize as a table
- each card is a row now + 9 columns
- new card → new row
- we can do with same with word dataset  
9d word Type Length
- what're the attributes in shopping bills?
- some are 1 but item, category and so on
- which ill create a variable no. of columns  
→ no longer a neat table
- variable no. of rows? tag to unique Id
- some things get duplicated multiple times
- with this, problem ↑ as entries increase
- we can create multiple tables

9d/Shop/Customers Table    9d/Item Cat. Qty/Prcd

the 9d links the 2 tables

⇒ Tutorial

Eg. Long Num = 0 . Table1 (cards) ; Read the first sheet Num = 0 ; Table2 (empty) ; some of table1  
Num → > 7 ⇒ LN = LN + 1  
→ else ⇒ SN = SN + 1      LN = 012347  
SN = 0123456897

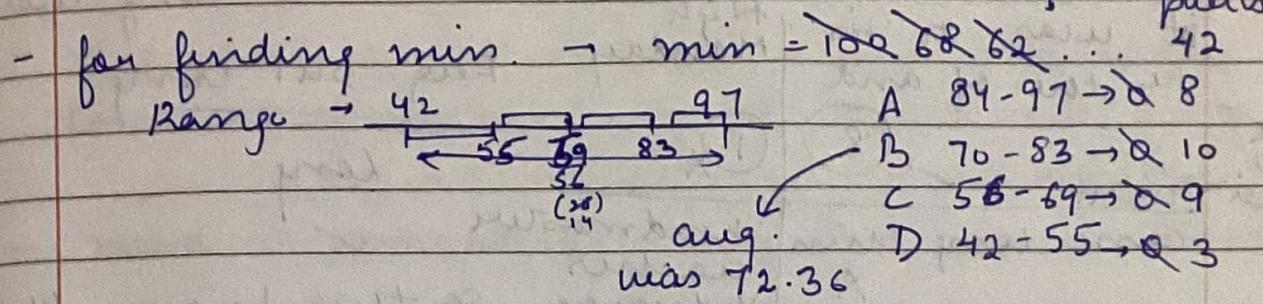
L 3.2 Below avg. students in 2 items (non-nested) & graded collacn

- we want to find avg, and then who are below avg. → most probably should be half
- count =  $\sum_{i=1}^{30}$  ... 30       $\frac{217.1}{30}$  = avg. = 72.36
- sum =  $\sum_{i=1}^{30}$  ... 2171       $\frac{30}{30}$
- so who're below avg. students?

b) count =  $\frac{1}{2} \times 2 \times 3 \dots \frac{15}{2}$  ( $30\%$ ) (anything  $> 72$  or  $< 72$ )

→ it's not necessarily always half

- Ques. the student, how do they have done? → grades  
Range → lowest to highest real marks  
(20) (97)  $78 \rightarrow$  4 equal parts



- this gives a better pic. → how class is doing

### ⇒ Tutorial

Q. all the bills above avg.

1. find avg:

count =  $27 \dots 30$

2.  $2^{nd}$  item who are

sum =  $2567 \dots 24413$

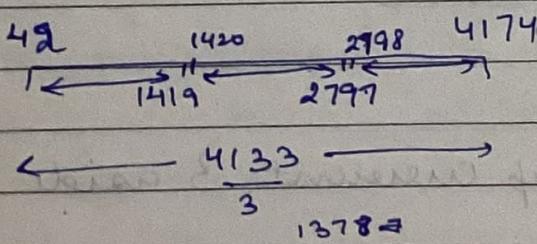
24413

32

→ so, new those bills who are  
above 813.76

Aug. count = 8

Ex. Low exp.	1	feel that
Med. exp.	26	Max. = <u>..... 4174</u>
High exp	3	Min = <u>5000 ... 42</u>



- Systematic Procedure of Hypothesis Verification
- L 3.3
- freq. words are gen. smaller and large words are less freq.
  - sum =  $\frac{329}{65} \approx 5$   $5 >$  is small word
  - count =  $\frac{18}{65} \approx 5$   $5 <$  is long word
  - it was to his  
be the of and  
this  
confirms short  
came long thin  
they  
refute
  - 15
  - 8
  - 19
  - 23
  - Monday
  - headmaster
  - teacher
  - wednesday
  - low
  - long
  - so this proves that freq. words are gen. smaller
  - this confirms
  - 47 distinct words  $\rightarrow$  31 confirm
  - 16 refute
  - there is correla<sup>n</sup> b/w. length of word & freq. of how do they occur

### L 3.4 Three Peiges Problem

- award prize to best student ; top 3 students of a class ; min. both  $\rightarrow$  1 girl & 1 boy

1 → }  
2 → }  
3 → }  
criticisms

- Top 3 marks  $\rightarrow$  keep current 3 aside

maths	phy	chem
1 2 3	1 2 3	1 2 3

- go on replacing & checking

281  $\rightarrow$  M

276  $\rightarrow$  F

261  $\rightarrow$  F

L 3.5Intro. to procedures & parameters

- compare the highest T with highest marks of each subj. added up.., if diff. is less, that means that highest topper is doing better in all subj.
- this is a procedure → repeatedly executing; some set of instances + changing the field name (parameters)
- max → procedure, P, C, M, Totals ⇒ parameters

e.g. Max (Maths)  $\Rightarrow$  97

Func<sup>n</sup>

L 3.6 Pseudocode for procedures & parameters (Part 1)

Procedure SumMaths (gen)

Sum = 0

(gen, fld)

gets assigned

while (Pile 1 has cards) {

gen(f)

Pick card X

Move X to pile 2

if (X.Gender == gen) {

Sum = Sum + X.Math

X.fld

}

}

return (Sum)

\* procedure to

sum up Maths marks

End SumMaths

⇒ 1 GirlChemSum = SumMaths (F, Chem.)

2 BoyChemSum = " " (B, Chem.)

if ( $1 > 2$ ) {

"compa. to girls")

}

else if

} "compa. to boys"

- procedure may not return a value, like updating marks → it is a separate statement
- procedures are pseudo-code-templates that work in diff. situa", they help to modularize pseudocodes

### L3.1 Pseudocode for Parameters (Part 2)

Eg. is there a single student who is best performer across subj.

- same computer " with a parameters to deliver the field

Procedure Maxmarks (fld)

Markst = 0

{ if (x.fld > markst) is true  
markst = x.fld } else skip

} if (MarkTotal == SubjectTotal \* x) is true  
return (MarkTotal) = markst  
end

if MarkTotal == SubjectTotal then true

Single Topper = True otherwise false

else {

} Single Topper = False

→ Tutorial

- 3 common errors

$\text{SumBB}, \text{CountBB} = \text{AddIfBB}(x, \text{SumBB}, \text{CountBB})$

Ex.

Procedure AddIfBB( . . . , . . . , . . . )

- #1 wrong Procedure Name; name used should be same in both places
- #2 wrong no. of parameters; if we have 3 parameters in procedure, take same in the code also
- #3 incorrect order of parameters; order is essential at both the places; assignments will be wrong.

L3.8 Pseudocode for 3 Peunje Problem

- Top 3 Peunje
  - basic → is total marks
  - at least 1 suely. topper
  - intop 3 → min 1 boy & min 1 girl

Procedure Top 3 Marks (sulij)

mark, 2<sup>nd</sup> mark, 3<sup>rd</sup> mark = 0

while (R != 1) and S contains numbers

Pick a candidate → switching him

if (x.sulij. > mark) {

3<sup>rd</sup> mark = 2<sup>nd</sup> mark; 2<sup>nd</sup> mark = mark; mark = x.sulij

}

if (mark > x.sulij. > 2<sup>nd</sup> mark) {

3<sup>rd</sup> mark = 2<sup>nd</sup> mark; 2<sup>nd</sup> mark = mark; mark = x.sulij

Ends Top 3 Marks code block numbered

Suppose present which was earlier numbered

- but, we record only 3rd mark ; between (3rd mark)
  - we should keep track of card no. that has the highest marks
- Pseudocode

~~Mark = 0 (2nd mark = 3rd)~~ mark b/w numbers

~~max = maxid = 1; 2nd maxid = 2; 3rd maxid = 3~~ maxid = 1-11

~~maths 3 = Top 3 marks (Maths) were 1st~~

~~phy 3 = Top 3 marks (Phy) were 1st~~

~~Chem 3 = Top 3 marks (Chem) were 1st~~

~~marks abt~~

~~update. Id. & marksmarkabt 2nd~~

~~third maxid = 3rd id used to determine~~

Procedure Subj-Topper procedure

(card, Mmark, Pmark, Cmark)

if (Card. Maths ≥ Maths ~~or~~ Phy ~~or~~ Chem) {

return (true); } else everything is false

else { true; need to ~~return~~ } I should

knock off boy (last) - card - held good

also { true; need to ~~return~~ }

knock off girl (first) - card - held good

End Subj-Topper

- gender problem? 3 boys, knock off boy (last) and replace it till we find the girl

L3.9

### Side effects of procedures

- procedures to sum up any type of marks
- what abt. the set of cards? Pile 1 is fixed.
- 3rd Parameter deck → Scan Deck, the procedure should also restore deck
- procedure modifies some data during computation

- seq. of cards may be disturbed, it might matter when some dec./circ. order is imp't.
- each procedure comes with a contract
  - Functionality → parameters will be passed ↳ expected in return
  - Side effects → is it predictable
- contract specifies interface → can a procedure implementa"(code) provide interface is unaffected
- there are places where side-effects are undesirable
  - └ no guarantee needed
  - └ undesirable
  - └ side effects are goal

# Week 4

Date / /

- Tutorial
- OR statement      Part of speech == Noun or Verb
- OR is true if even 1 statement is true
- The cond<sup>n</sup> can both be true simultaneously
- They are mutually exclusive.
- Not mutually exclusive --- don't split that
- AND is true iff. both the statements are true

## L4.1 Concept of nested items (Näive)

- birthday paradox → 25 people → more 50% probability → same month & date
- "1<sup>st</sup> item" - 1<sup>st</sup> card → inner item<sup>n</sup>  
  ↑  
  what class    Aksnay<sup>a</sup>    whole class    it's a item<sup>n</sup>  
  ↓  
  this pile    Abusim    within item<sup>n</sup>  
  more + big.
- 2 customers who shop in a same shop

## L4.2 Binning concept of nested item<sup>n</sup>

- group students by months ; separated out into bins by months → get 12 bins ; each bin do same double item<sup>n</sup>
- cards per bin decrease ; the double item<sup>n</sup> becomes easier
- but if use binning wisely so that it eases task ; take clue from the data

Tutorial

Q. Find score cards same maths with same place

~~Euode = 68, 72, 0, 18~~

~~Salem = 1~~

~~Ambur = 6~~

~~Banglore = 74, 63, 97, 78, 11, 21, 23~~

~~Tenn = 15~~

~~Chennai = 57, 42, 71, 64, 51, 62, 27~~

~~Madurai = 4, 13, 24, 27, 28~~

~~Vellore = 7, 20, 29~~

~~Trichy = 14, 16, 19~~

~~Nagercoil = 22~~

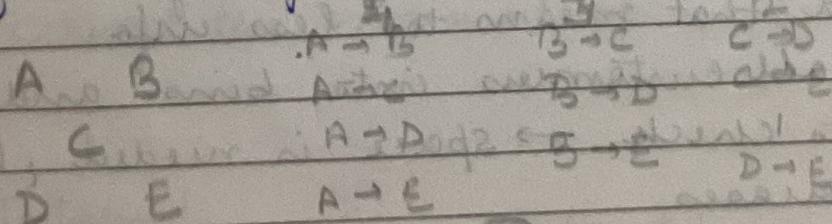
Naive methods  $\rightarrow$  435 comparisons

$$\boxed{\frac{n \times (n-1)}{2}} \Rightarrow n \text{ no. of rows in that table}$$

Binning methods  $\rightarrow$  only 55 comparisons  
 drastically red., if data is reasonably distributed

L4.3 Emp. of binning

- red. no. of comparisons
- some computer - e.g. comparisons of each card with all other cards
- obs. that if we can organize cards into bins formal way of determining the reduc^n in comp.



$$\boxed{\frac{N(N-1)}{2}}$$

for N obj.,

no. of comparison

- this is same as no. of ways of choosing 2 obj. from N obj.  $N_C^2 = \frac{N(N-1)}{2}$
- the no. of comparisons grows exponentially

- computer  $\approx 10^8$  comparisons / sec

A		$9 \times 8 = 36$
B	C	
D	E	$G_1$
F		$H_2$

reduced  
comp. by a 4

so for each bin

$3 \times 2 = 6$

2

$3 + 3 + 3 = 9$

~~for N items~~  $\left\lceil \frac{N \times (N-1)}{2} \right\rceil$

- if we use  $k$  bins of equal size, no. of items in each bin =  $\frac{N}{k}$

~~Now, Total comparison~~

$$\left\lceil \frac{N \times \left( \frac{N-1}{k+1} \right)}{2} \right\rceil$$

~~Reduc" by a factor =~~

$$\left( \frac{(N-1)}{\left( \frac{N-1}{k+1} \right)} \right)$$

- ### L 4.4 Concept of binning to avoid complexity
- loyal customer - visited many times to a same shop  $\rightarrow$  bin as we go along
  - if 1 card that means too less info.
  - separate all customers into bins and iterate to check  $\rightarrow$  space is needed, but time decreases

### L 4.5 Concept of Fair Teams

- help each other in studies  $\rightarrow$  group study
- pair of students X. Maths > Y. Maths  
 $\Delta X \cdot Phy < Y \cdot Phy$

- pairs which have same totals  $150 + 25$  bins
- binning is a generic method, comparisons reduced; inside the bin we again compare to make the pairs

#### L4.6 Procedure to find same d.o.b

- birthday  $\rightarrow$  parameter; check along the whole deck
- q. 30<sup>th</sup> April  $\rightarrow$  return (False)  
 13<sup>th</sup> May  $\rightarrow$  but 30<sup>th</sup> April, we have already checked  $\rightarrow$  return in such a way that specific  $\Rightarrow$  are marked

#### L4.7 Procedure to find person with their name

- when down to get person, then went backward to find name  $\Rightarrow$  now directly move further to find new person
- keep a user name
- nested loop  $\rightarrow$  iterating over the card, then started the inner loop

#### Tutorial

 $i = 1$  $v = 0$  $c = 0$ while( $i \leq x.\text{letterCount}$ )if  $i^{\text{th}}$  letter  $x.\text{word}$  is vowel     $v = v + 1$   
    else  
         $c = c + 1$  $i = i + 1$

- Summary
- iterator → most commonly used in CT, it represents the procedure of some task repeated
- initialize → count → stop
- variables → S the values, at each item "it", it is initialized & update of vars. are done thru assignment statements
- terminals  $\square$ , process  $\square$ , decision  $\square$
- use of flowcharts & pseudocode
- items to systematically go thru a set of items
- sanity of diff. data fields of the items
- basic data types → boolean, int., char.
- list → all data items typically have the same datatype
- record → multiple named fields, each can be of diff. datatype
- filtering - a decision at each repeated step
- sometimes we have to write the same pair of code again & again with small diff.  
↳ use procedure (small diff. → parameter)  
Eg. find max for each subject
- most common use of an iterator is to create an accumulators
- doing 2 items - one nested within other
- we can reduce the number of comparisons by using binning methods

# Week - 5

CLASSTIME	Pg. No.
Date	/ /

## L5.1 Intro. to collec<sup>n</sup> & list data stucture.

- collec<sup>n</sup> is sets ; relationships are relations
- comp. has memory RAM → Random Access Memory → u can't go to that place directly & change it
- no side-effect , with the help of bookmark
- list → indexing the positions , we can make several lists

(outer layer) (inner)  
 Chennai list ( ) may list = students born  
 in May from Chennai (nested layer)

## L5.2 Pseudocode for lists

- variables keep track of intermediate values
- simplest collec<sup>n</sup> is a list → seq. of values .  
 nota<sup>n</sup> for lists
- [1, 13, 2] ["cane", "school"] [] → empty set
- append & lists . . . l1 ++ (2) → <sup>→ sy. impt.</sup> combine
- l1 is [1, 3] l2 [2, 17, 1] l1 + l2 → [1, 3, 2, 17, 1]
- extend l with item x l = l ++ [x], list of length 1
- list of student in may , list of students from Chennai

Maylist = []

while (Table 1 has entries) {

Read first

```
if (x. MOB == "May") {
```

```
  Maylist = Maylist ++ [x. Seq.]
```

- we need to iterate over a list

foreach x in l {

do sth. with x

may chmnalist = []  
 foreach x in MayList {  
 foreach y in chmnalist {  
 if ( $x == y$ ) {  
 mclist = mclist + + [x]  
 }
 }
 }

- L 5.3 Operations on data collected in three periods
- Top 3 in atleast 1 subject
  - maths, phy, chem. set  $\Rightarrow$  student who managed to come in all
  - we do not disturb the original card set - that is key thing -

- L 5.4 Pseudocode for 3 Period problem
- Procedure Top 3 Marks (Subj.)
- ↳ max, 2<sup>nd</sup> max, 3<sup>rd</sup> max
  - ↳ we get return of third max (cutoff)
  - $\Rightarrow$  cutoffs M/P/C, now mpC list = [] tail - while (table 1)
    - if ( $x.\text{maths} \geq \text{cutoff M}$ )  
 $mclist = mclist + + [x.\text{Seq. No.}]$
  - $\Rightarrow$  we got three firsts

$\Rightarrow$  mp list [] foreach x in mp list {  
 foreach y in plist {  
 if  $x == y$  {  
 mp list = mp list + + [x]
 }
 }
}

- now we'll match mp list & c list and append to form mpc list
- lists are useful to collect items that share some property

### L5.5 Insertion Sort & Ordered List

- arranging in some order  $\rightarrow$  find a place to insert  $\rightarrow$  building up a new list
- insertion sort  $\rightarrow$  we are inserting it along with sorting it out the things  $\Rightarrow$  (nested iteratn)
- building up a ordered list.
- we sort the bookmark
  - $\rightarrow (0, 210) \Rightarrow$  Ordered list
  - $\rightarrow (1, 198), (0, 210)$
  - $\rightarrow (3, 173), (2, 188), (1, 198), (0, 210) \dots$
- each ordered list keeps a track of set of card in inc. / dec. order.

### L5.6 Pseudocode for insertion list

- sorting a list make computations easier
- finding top k values, finding duplicates, grouping by percentiles (quarters)
- insert sort - create a 2<sup>nd</sup> sorted list, start with an empty list, repeatedly insert next value
- list l in arranged order  $\rightarrow$  insert a new element x at a appropriate place

Procedure to list (l, n)

newl = []

or not inserted = False

for each z in l {

if not(inserted) {

if (z <= z) {

newl = newl + [z]

not inserted = True

newl = newl + [z]

}

if (not(inserted)) {

newl = newl + [z]

- ] - - (018, 0), (8P1, 1) -

End

$\Rightarrow$  sorted list = []

foreach (z) in before (n) do

sortedList = Insert (sortedList, z)

return (sortedList)

invariant - second list is always sorted

$\Rightarrow$  [] is sorted, since it is empty

"insert" sort is a natural sorting

algorithm

- L5.7 Systematic procedure for hypothesis verification
- the person who does good in maths is good in phy. also
  - distribute in grades → we had a hypothesis
  - either there is confirm / refute
  - this hypothesis isn't much associa"

- L5.8 Pseudocode for systematic procedure
- list eg. → calculating student's performance
  - algorithm → assign grades
    - count A of Maths in A of Phy
    - count of B of Maths in A/B of Phy
    - use count to confirm / refute

⇒ Tutorial

- $\text{length}(l) \Rightarrow \text{length}([20, 30, 40, 50, 10]) \rightarrow 5$
- $\text{first}(l) \Rightarrow \text{first}([20, 30, 40, 50, 10]) \rightarrow 20$
- $\text{last}(l)$ 
  - ↳ if empty → ND
- $\text{rest}(l) \Rightarrow \text{rest}([20, 30, 40, 50, 10]) \rightarrow [30, 40, 50, 10]$
- $\text{init}(l) \Rightarrow \text{init}([20, 30, 40, 50, 10]) \rightarrow [20, 30, 40, 50]$
- $\text{member}(l, e) \Rightarrow \text{rest}(l) = l - \text{first}(l)$ 
  - $\text{init}([20, 30, 40, 50, 10]) = [20, 30, 40, 50]$
  - $\text{init}(l) = l - \text{last}(e)$
  - $\text{member}(([20, 30, 40], 30) \Rightarrow \text{True}$
  - $\text{member}(([20, 30, 40], 10) \Rightarrow \text{False}$

→ Tutorial

- X → books insertion

$sBooks = []$

updated version

somebooks = []

" " = Insert [2, somebooks]

L5.9 Intro. to Train Dataset

- train & station dataset
- pair of terrain city  $\Rightarrow$  city (paired train)
- MWF  $\rightarrow$  days that train moves [022017  
022021]
- next data  $\rightarrow$  list of cities and all the terrain that pass thru it

Equivalent  $\Leftarrow$

(1) appear -

(1) tried -

(1) test -

(1) take -

(1) time -

(1, 1) remove -

# WEEK - 6

CLASSTIME Pg. No.  
Date / /

## L6.1 Rela<sup>n</sup> of customers - Part 1

- Segments - segment of customers for ad
- Abhinav -Appealed food Toiletries (Revs of items)
- Neerja 270
- Akshaya 476 5 . . . so on
- we got a table, rows are customers, columns are categories
- Radha, Akhil, Julia, George. → Food only
- 1 grp. with more than 10 & 1 grp. with 10 or below 10
- atleast everybody buys food across all bills irrespective of the grp.

## L6.2 Rela<sup>n</sup> of customers - Part 2

- 6 customers - each compared with everyone. 15 comparisons will be there
- compute the distance & the pair which has the least difference becomes a similar customer
- Rectilinear / Manhattan distance
- tables are good way of organizing things
- grouping in some 1 column

## L6.3 Intro. to dictionary data structure

- which shop gives the max. no. of food items
- iterate them card and in inner dict "in" count all the food items in that list ↳ SV stores

- which customer buys most food? there are many customers — how to handle?
- so the customer who buys the max food item is from the shop that sells max food item — thus, hypothesis ✓

#### L6.4 Concept of dictionary to solve birthday problem

- dictionary — indexing, index is full b.o.d.
- build dictionary in 1 go

DOB      Card No.

3 June      1

5 May      3

any element  
which has  
2 entries

so on → only 1 b.o.d. is common and also

dictionary told that there aren't any more duplicates

- also, keep track of diff. birthdays that come up
- each year in dictionary store. can be accessed using key

#### L6.5 Resolve pronoun with its matching noun

- group together by some category
- pronoun & noun - bracket

Pronoun - o, g, - . . . 61

Noun - 2, 3, 4, 64

g → his ; go before 2, 3, 4 and it should sensible also

4 7 9 11

- dictionary by category - and we are storing the seq. no. as key
- partitioning each word into any of the categories  $\rightarrow$  binning

### L6.6 Pseudocode for dictionaries

- a list has seq. of values, can iterate them list
- dictionary stores key-value pairs
  - $\mapsto$  chemmarks (value) for each student (key)
  - $\mapsto$  source sta<sup>n</sup> (value) of train route (key)
- present the key to extract the value  
 $m = \text{chemmarks}["Rahul"] \Rightarrow$  random access allowed
- seq. of 'key:value' pairs  $\{ "K": 92, "L": 86 \}$
- empty dictionary {}
  - $\mapsto$  chemmarks ["Rahul"] = 92 (either replaced or added)
- the dictionary must exist to create a new entry

chem = {}

while (more names) {

    Read first name into list  
    name = X.Name

    marks = X.chem

    chem [name] = marks

} More X to name 2

- keys(d) is the list of keys of dictionary  
for each k in keys(d) {
  - do something

total = count = 0

for each k in keys (chem) {

total = total + chem[k]

count = count + 1

}

chemavg = total / count

- use of a dictionary is to accumulate values

sum ["Kohli"] = sum ["Kohli"] + score

but, in beginning

sum ["Kohli"] = score

is key (d, k) - returns True if k is a key in d

{ if iskey (sum, "Kohli") {

else {

} first assign

- 1 Problem  $\Rightarrow$  Take time proportional to size of the dictionary  $\Rightarrow$  random access

### L6.7

Pseudocode for real-time egs. using dictionary  
- find the customer who buys highest no. of food items

food D { }

while (Table 1 has rows) { }

read the first row

Customer = X. Cust Name

items = X. Items

```

for each row {
    if (row.category == "Food") {
        if (is key (foodD, cust)) {
            foodD[cust] = foodD[cust] + 1
        }
    } else {
        foodD[cust] = 1
    }
}
more rows → table 2
}

```

### L6.8 Dic. comparison to find common elements

- train dataset  $\begin{cases} 1 & \text{dic. - train is the key} \\ 2 & \text{dic. - station is the key} \end{cases}$
  - Merges directly to  $\rightarrow$  Pune
  - 02283      } Stations list  $\times$       New which all  
12224      }                                  trains reach Pune
- Best should  
be Mumbai

### L6.9 Procedure to find rela<sup>n</sup> bus. customers

- 1 contractor keeps the comp. list of all values
  - ↳ sub-contractor works for 1 person and iterate them 1 card
- sub-contractor doesn't have to do anything abt. the name
- this becomes modular & easy, the decision-making is now divided
- parameter is the card, return value is high
- side-effect  $\rightarrow$  call for str., he does sth - else

- L6.10 Side effects for list & dictionary
- compare 2 lists for duplicate items → nested loop
  - what if lists are sorted out? use `first()` and `rest()` to cut down list to be scanned
  - now, 2<sup>nd</sup> list has been modified inside the procedure → side-effect
  - for this, we can copy the main list to my own copy and change it accordingly
  - in some cases, side effects are intended, like in places to update the pure collection

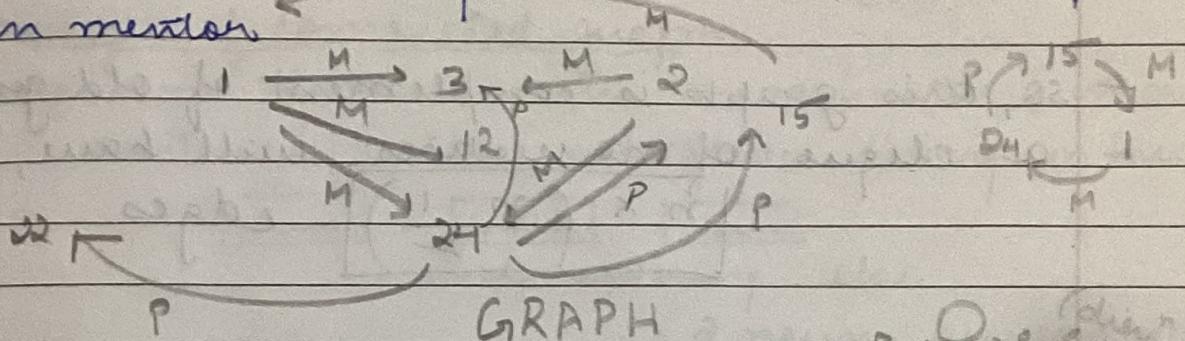
# Week - 7

CLASSTIME Pg. No.

Date / /

## L7.1 Intro. to Graph Data Struct.

- $A \xrightarrow{S} B$  A's Mark in S - B's Mark in S  
student helps student in subj.s btw. [10, 20]
- all pairs, hence, nested items  
 $1 \xrightarrow{M} [7, 16, 6, 27, 21, 4]$  and so on find for each student, how many other students that person can mention



- cycles of length 2/3 with list of adjacent / neighbours

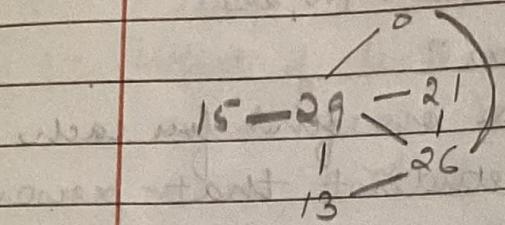
## L7.2 Intro. to matrices & implementation of matrices

- list keeps a seq. of values, no random access, for value at i, scan ( $i-1$ ) elements
- dictionary stores key-value pairs, supports random access
- Matrix — 2d table, m rows, n columns
  - ↳ rows, columns are numbered from 0
  - ↳  $0 \leq i \leq m-1, 0 \leq j \leq n-1 \Rightarrow \text{matrix}[i][j]$
- outer key  $\rightarrow$  rows, inner key  $\rightarrow$  columns
  - ↳ creates a nested dictionary
- create  $\rightarrow$  mymat = Create matrix (30, 45)
- sort (key(d)) — ascending order

```
foreach a in sort (keys (matrix)) {
    " " sort (keys (matrix[a])) {
        do sth
    }
}
```

### L7.3 Undirected graphs & cliques

- similar  $[0, 9] \rightarrow$  thus, undirect, only similar
- Eg.  $22 \sim 13$   $\rightarrow$  undirected edge



Clique  $\rightarrow$  congregate together and exclude other people

- = so, this graph is complement of old graph
- a clique of  $n$  vertices will have,  $\boxed{n \times \left( \frac{n-1}{2} \right)}$  edges

### L7.4 Concept of popular students using graph

- more popular, in middle, they get help from many people and they'll help many people

- no. of edges in nodes is degree  $\rightarrow$  in degree

$21, 26 \xrightarrow{M} 10 \xrightarrow{M} 13, 29, 26$  and so on  
complete the list

### L7.5 Pseudocodes for eval-time egs using graphs

- A can mentor B, diff. btw. 10 - 20 marks
- create a dictionary for marks in subj.
- create a mentoring graph, represent as a matrix

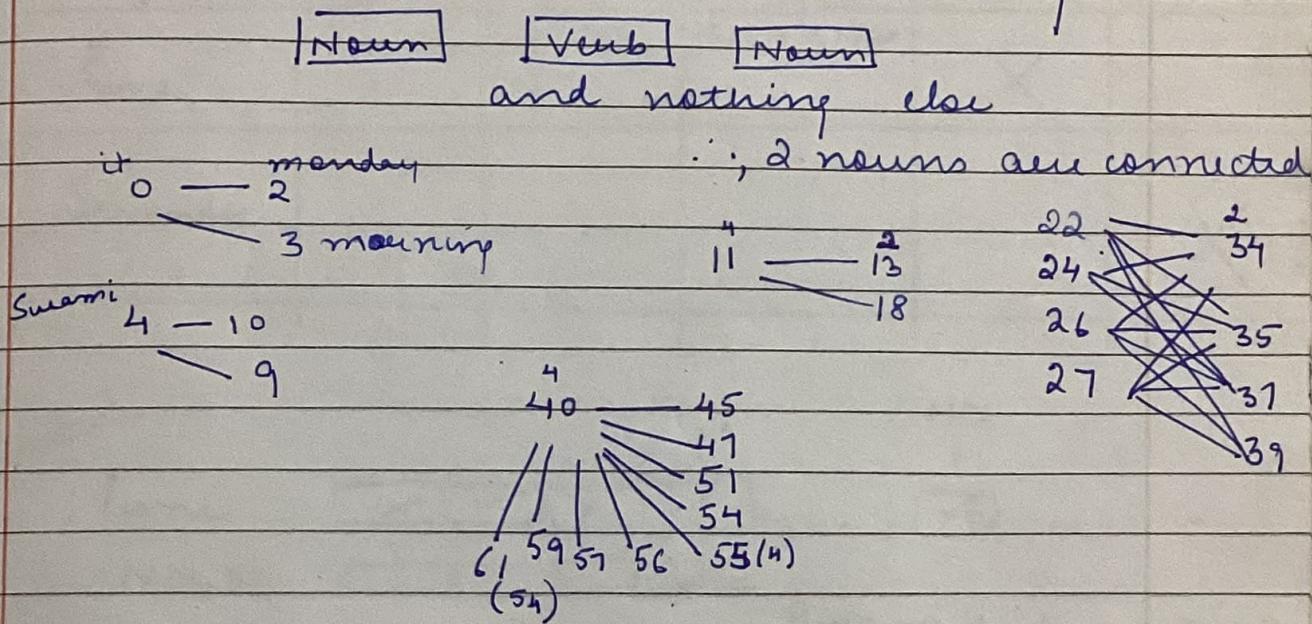
$M[i][j] = 1$  - edge from  $i \rightarrow j$

$M[i][j] = 0$  - no edge

- dictionary helps to keep the list of things explicitly w/o any repetition
- iterate them matrix to aggregate info. from the graph

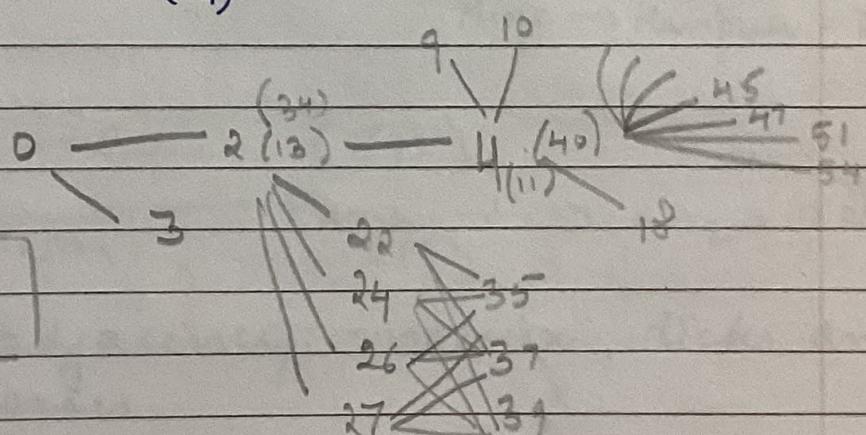
### L7.6 Concept of connected graph

- Rela<sup>n</sup> → 2 noun with a verb b/w., in same sentence, with not other thing in b/w.



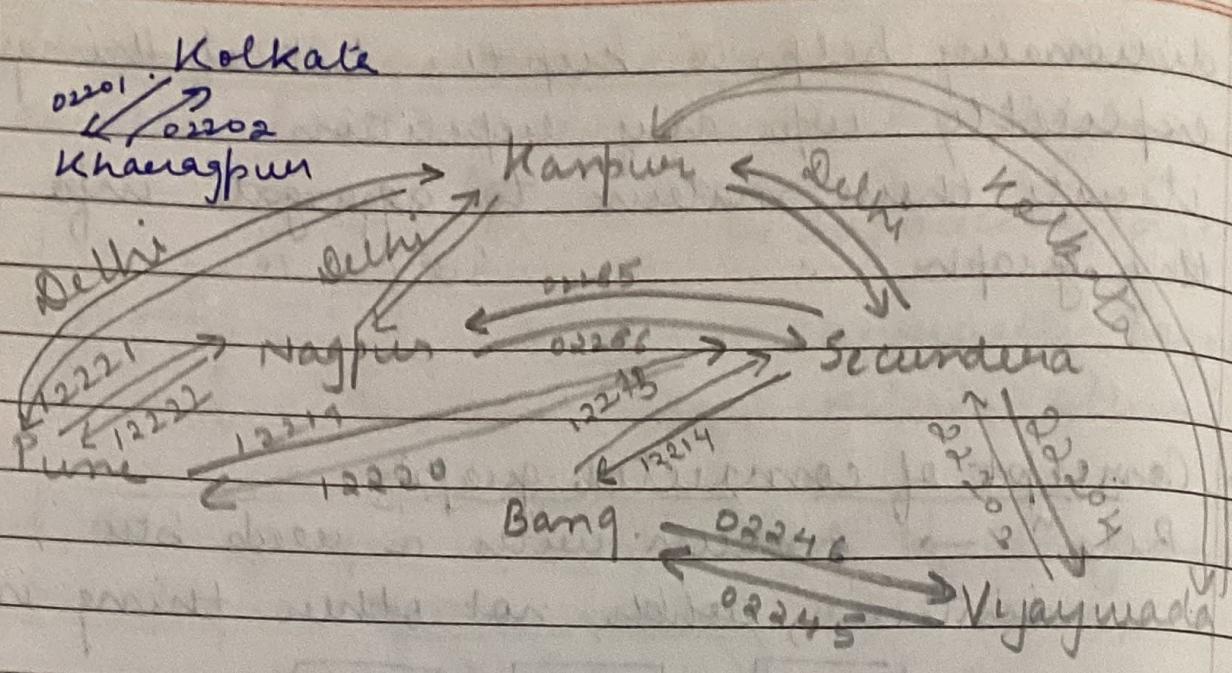
Coherent  
Para

Connected  
Graph



### L7.7 Represent direct trains using graph

- Rela<sup>n</sup> b/w. stations with a edge



# Week - 8

CLASSTIME Pg. No.  
Date / /

## L8.1 Represent graph as adjacency matrix

T0	K	Na	NDSL	P	S
From ↓	Kolkata	X	$10200 = 1131 \text{ km}$ $10262 = 1131$	$10273 = 1529(2)$ $10269 = 1457(1)$	return -
	Nagpur	X			$02286 - 577$
	NDSL			X	
	Pune		$10263 - 1520$	X	
Secunder.					X

		Kanpur	ND
Pune	<input checked="" type="checkbox"/> <del>go to Kolkata</del> <input checked="" type="checkbox"/> <del>Mumbai - go to ND</del>		
NDSL	<input type="checkbox"/> <input type="checkbox"/> <input type="checkbox"/> <input checked="" type="checkbox"/>		

$$\text{Pune} \rightarrow \text{Kanpur} = D_1 + D_2$$

Secunder. → Pune (597)

→ p.d → Nagpur (579) → ~~Pune~~ ~~go to Mumbai~~ ~~ND~~ (1520)  
~~(889)~~ ~~most~~ ~~2nd~~

- Tables → adjacency matrix, ticks and crosses

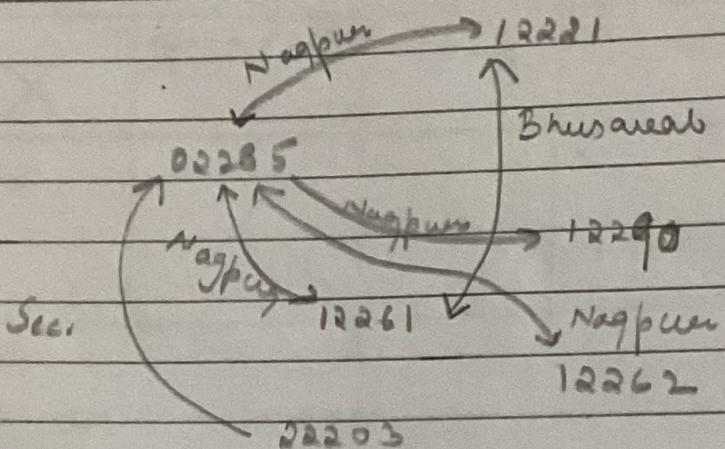
	Sec	Vija	NDSL	Kan	Kol
B	<input type="checkbox"/> 688	<input type="checkbox"/> 720	<input type="checkbox"/> 1686 1067		
S				<input type="checkbox"/> 439 440	
ND				<input type="checkbox"/> 1014 1090	
Kol					<input type="checkbox"/> 1218
WJ.					

$$B \rightarrow V \rightarrow Ked \rightarrow Kan = 728 + 1218 + 1014 = 2960$$

~~$$B \rightarrow Sec \rightarrow ND \rightarrow Kan = 688 + 1667 + 439 = 2794$$~~

- Matrix multiplication - adds values

L8.2 Construction of graph where trains are nodes and stations are edges



- concept of delay propagation

L8.3 Pseudocode for finding direct b/w 2 sta<sup>n</sup>

- compute pair of sta<sup>n</sup> connected by a direct train

```

for each t in keys(trains) {
    stations { } = stations
    for each s in stations { } {
        stations {[trains[t][start]]} = Time
        stations {[trains[t][end]]} = Time
    }
}
  
```

$n = \text{length}(\text{keys}(\text{station}))$

direct = create matrix ( $n, n$ )

start index [i]

initial print - print every station present -  
man is forming a route in steps (stations) also of  
strides [s] = distance

new) + fifth one - print listing of  
present steps also print a label -

1st int formant via steps (train) also not  
i = strides [train (+), [start]] +

initially takeable with configurations -

transitively [i][j] implies transitively along -  
(transitive)

meet transitive (direct)

- 1 hop route - stations A & B such that u can reach B from A by changing 1 train
- we can modify to direct to 1 hop (at most)
- 2 hops -  $i \xrightarrow{o} k \xrightarrow{r} j$
- we can extend 'n hops' to allow 1 more hop
- Path - seq. of edges from  $A \rightarrow B$ . A direct edge is a path of length 1
- N nodes  $\rightarrow$  shortest paths from  $A \rightarrow B$  has at most  $(N-1)$  edges
- Transitive closure - pair of nodes connected by a path

### L8-4 Pseudocode for edge labelled graph

- keep a track connecting stations
  - ↳ each entry in the matrix is now a dictionary
  - ↳ initial empty - no direct connec<sup>n</sup>
  - ↳ add a key for each train
- for each direct train record the dist it travels trains [t] [distance]
- compute the shortest distance
- graph directlist [i][j] = 1 (is the shortest distance)
  - = 0 (no direct train)

str index

```

mst = null
for i in range(n):
    for j in range(i+1, n):
        if (train[i][j] != None) & (train[i][j] != 0):
            mst.append((i, j))
            print(f"Train {i} connects station {i} to {j} via {train[i][j]}")
            print(f"Distance between {i} and {j} is {train[i][j]} km")
            print(f"Train {i} connects station {j} to {i} via {train[j][i]}")
            print(f"Distance between {j} and {i} is {train[j][i]} km")
            print("-----")

```

L9.1 DFS (Depth First Search) & recursive procedure  
call - Part 1

Vijaywada → Monday

Vijaywada (sun) 00:55, 22/03 Sec. 00:15, 22/04 Vij × (Sun) (X)

6:15 (sun) → 01:05 Tues X  
 4:25, 00:45 RNGT 17:00, 02/04 Vij × (Sun) (X)

10:00 (sun) → 00:00

10:15, BLR (sun) (X)

02/04

16:00

(sun) p 2681

- depth first search - goes deeper and deeper
  - ↳ u have starting time → start seen recursive procedure it calls itself

explore (Vijaywada, mon, 00:00) Vij  
 explore (Sec, MON, 6:15) ex (RNGT, MON, 10:00) SEC

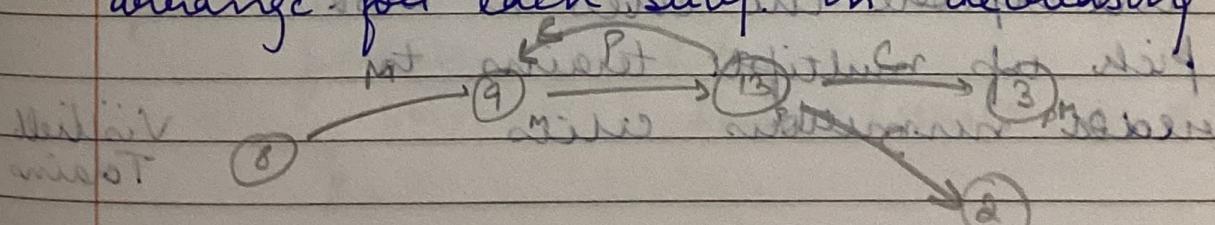
RNG

→ recursion → procedures call itself

- DFS creates a tree with edges of graph, gives the order of journey
- keeping track of seen

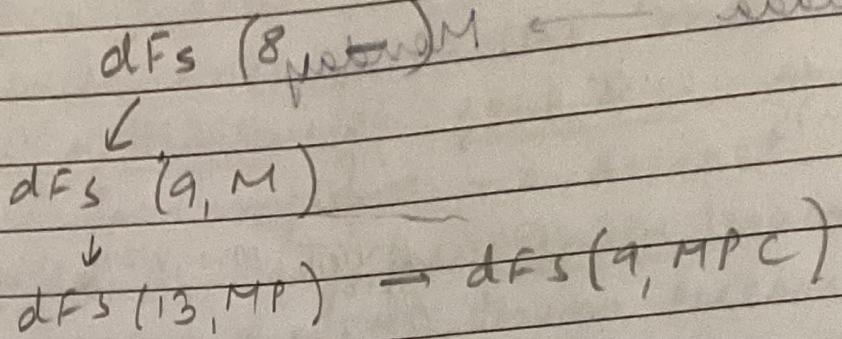
L9.2 Part 2 (Kruskal's algorithm) ← and HC  
 mentioning graph minst and terminals

- arrange for each subject in decreasing order



- finding the graph as we go along

- last 2 nodes



## L9.3 Part - 3

12269 → Monday → Exploring connecting train

~~vergelijk ons verleden oop - neem daar~~  
~~meewerken~~ ~~met een~~ ~~stem~~  
~~verleden~~ ~~met een~~ ~~stem~~  
dit is vijf Ex (22246, M, 22:31) Ex (22004, Tu, 1:45) 02246/45  
Nettij Ex (22246, M, 22:31) Ex (22004, Tu, 1:45) 02202/01  
↓ Bnub.

Ex(02202-Tu, 20:15)

Heidi has a husband - Michael

- spanning tree → reaches all the stations  
the geography is included in this

## L9.4 Pant - 4

- 24 hrs → Explore (Bhusaval) Visited
  - alternate b/w train ~~stop~~ no intermediate cities and explore these cities & also generate - pick up multiple trains to reach numerous cities Visithell

Vishnu  
Toraino

grate off sea as afterglow and graining -

### L 9.5 Pseudocode for recursion

- many computations — inductively
  - ↳ directly return the value → base
  - ↳ compute value in terms of smaller arguments — inductive

$$\text{factorial}(0) = 1$$

$$\text{for } n > 0, \text{ factorial}(n) = n * \text{factorial}(n-1)$$

Procedure fac(n)

```
if (n == 0){  
    return(1)  
}
```

```
else {  
    return (n * fac(n-1))  
}
```

End

factorial(n)  
is suspended till  
factorial(n-1)  
returns the  
value.

- base → empty list

inductively → compute value in terms first  
element & rest

if  $l == []$ , sum is 0

or  $\text{first}(l) + \text{rest}(l)$

or  $\text{last}(l) + \text{init}(l)$

if ( $l == []$ ) {

return(0)

return ( $\text{first}(l) +$   
 $\text{listsum}(\text{rest}(l))$ )

- build up a sorted prefix, extend it by inserting the next element in the correct pos

sortedList = []

```
foreach z in l {
```

sortedList = Insert(sortedList, z)

}

- insert sort, inductively  $\Rightarrow$  list of length 1 or less is sorted, for longer, inserit  $\text{first}(l)$  into rest( $l$ )

```

    if (length(L) <= 1) {
        return(1)
    }
    else {
        return (insert (insert (rest(L))
    }
}

```

- w/o properly defined base cases, recursive procedures will never terminate  $\Rightarrow$  suspend the current computation till the recursion computation terminates

Lq.6 Pseudocode for DFS & recursive procedure

- what are vertices from node  $i$ ?
- start from  $i$ , visit a neighbor  $j \rightarrow$  bfs
- $dfs \rightarrow i \rightarrow j \rightarrow k \rightarrow \dots$  and so on
- maintain info about visited nodes in a dictionary visited  $\rightarrow$  not to explore again

visited = {}

visited =  $dfs(\text{graph}, \text{visited}, i)$

keys (visited) is set of nodes that can be reached from  $i$

Procedure  $dfs(\text{graph}, \text{visited}, i)$

visited[i] = True

foreach  $j$  in columns (graph) {

if ( $\text{graph}[i][j] == 1$  & not iskey (visited))

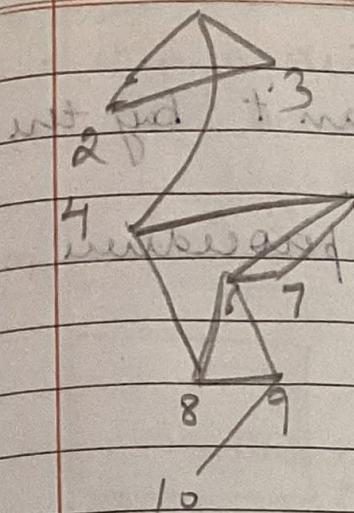
visited =  $dfs(\text{graph}, \text{visited}, j)$

}

return (visited)

End

$\Rightarrow$  keys (visited) includes all nodes, the graph is connected



visited = {4 : True, 1 : True}

2: Train, 3: Train, 5: Train, 6: Train  
7: Train, 8: Train, 9: Train

principles for Teamwork -

DFS (geophin, assisted, 4) w/

~~DESC "Gardens and sites")~~  
~~DESC "Gardens and sites")~~

~~DESC("autoform - 3")~~ int

DFS( " 400 je wege ti  
DFS( " )

~~DE 5 "Deutsche Zeitung"~~

DE's project's approach to triple  
DE's project's approach to triple

~~DESC~~ 8) ~~image ab' has 2 minif~~

sets up a field - tested area

• jels, os grader - "Mid en oppdrag

plus! and sometimes!

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10. Mr. John Wright, of the U.S.

10. *Leucosia* *lutea* (Fabricius) *lutea*

*pers* | *o*

*100*

But ~ and with a taste

planar file

I don't care if it's terrible

animal

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## Week - 10

CLASSTIME Pg. No.

Date / /

### L10.1 Concept of Encapsulation & Object

- Manual comes with TV, u can't buy the manual separately
- other aspects of organising procedures
- all the cards are inside the box → object
- capsule consist all the data and procedures in the box → encapsula<sup>n</sup>
- it relates of OOP (object-oriented programming)
- object → some static . of things and do some opera<sup>n</sup>s on it.

Score dataset - obj. is whole stacks  
Shopping Bill " " - 1 shop is 1 obj.

- 1 customer be 1 obj.

Shop Obj.

SV, BB, GG

Aust. Obj.

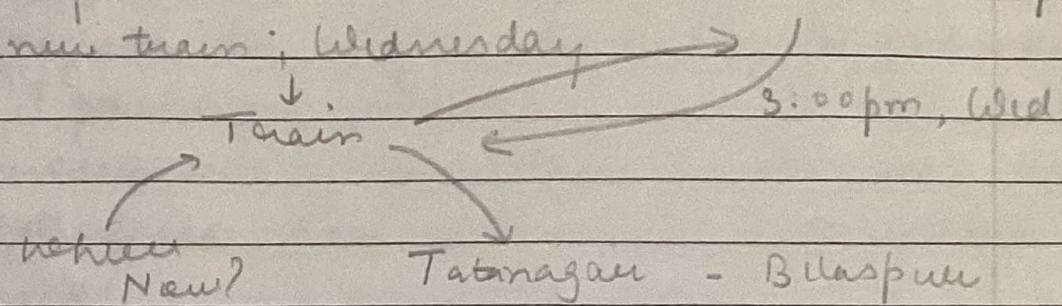
Ahmed ..

Bill Obj.

### L10.2 Concept of class & its instance c/d obj.

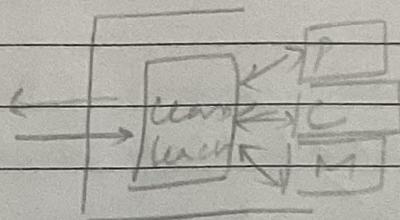
- procedures are inside the obj.  
<sup>train</sup> → trains → 1 obj.  
<sup>data set</sup> → stations → 1 obj.
- depending on what is the ques<sup>n</sup> we have to choose obj wisely
- 1 card of train dataset is a template of generating trains

- 1 train card in train dataset
  - it is a design of an obj. → class → not a obj. itself but it represents an obj.
- 12262 → class      clock → obj



#### L10.3 Concept of abstraction

- from where, this obj. comes from?
- we create 1 obj. that obj. creates other obj. if needed.



#### L10.4 Application of class, obj., encapsulation & abstraction - Part 1

Jhansi ← → Kanpur (same day)

Tuesday

12214 (Mon)  
12269 (Mon)  
10:22 12270 (Tues) DDL  
12260 (Mon)  
12273 (Mon)  
12259 (Mon) 06:30  
12274 (Tues)

Tamai, TDM, JPN9, T) - after -

next train journey ext priceformat -

L10.5 Part 2

file note required with  
trainee file → T11 (Visti) file SEE NASTN

Expense STN

KGP

KGP 01:15

KOL 4:15

TATA 11:35

BLSP 17:35

SEEN TRAIN

02202 (Tu)

12262 (W)

12261 (W)

Expense TRAIN

02202 (Tu)

12262

12261 (Th)

? many curves file want number 12262 (Thu)  
other file sent file 1 stens new  
beber fi file.

ExS (ACGP, W, 01:10)

ExT (02202, Tu)

ExS (KOL, W, 04:15)

ExT (12262, W)

ExS (TATA, W, 11:35)

X

ExS (BLSP, W, 17:35)

X

L10.6 "decomposes" of study group problem

- 4 obj. - CT, PhyT, MathT, ChemT
- decomposing the scenario into 4 obj

- L10.7 Formalized nota" and summary of encapsula"
- It seems to be more natural to attach procedures to the data element
  - Encapsulate the data elements & procedure into 1 package, which is cld an obj., i.e., oop
  - The procedures encapsulated in the obj. act as an interface thru which the ext. would interact with obj.
  - The obj. hides details of implementation from ext. world, separates "what" from "how"
  - procedures was also doing the same things, taking parameters & return values, hiding was also there, but procedures have undesirable side-effects, and we can't store the value
  - datatypes —

- ↳ basic — int., char., boolean
- ↳ subtype of datatype restricts values & opera"
- ↳ compound — record, list, string  
collection of cards is also a datatype

### Procedure

### Encapsulation

- Aumarks takes a field -datatype ClassAvg has as parameter
- e.g.  $\text{AvgTotal} = \text{Aumarks}$  (total) — we call avg() for each

$\text{AvgMaths} = \text{Aumarks}$  (Maths)

obj. of Class Avg datatype

- e.g.  $\text{AvgTotal} = \text{CT. Avg}()$

$\text{AvgMaths} = \text{Mat. Avg}()$

- if the dataset is not static, means a student gets added to the class

- Class Ave should be able to restrict access to its fields, done by private/public field encapsulation procedures due to modularity
- when compared to procedures

does not have access, if I change one file it is different compiler and not -  
other changes that will affect my application  
file will

and "encapsulation" is called when file don't  
want many "extern" declarations between them  
exist other external code will be considered  
public, unless another I mentioned public  
declaration must be numbered two, went into each  
other file won't make any sense, so if this  
— applies -

instead want this - said of  
methods & variables declared as private to application or  
private, tail bracket - boundaries of  
application so code is used for  $\sim$  like a  
"blurred" number

and return application - like a class where we  
intervene (use number) is mentioned as  
this  $\rightarrow$   $\text{class} = \text{class} . p$  -  
the top () for the file - (later)

part and for file environment -  $\text{class} . e$   
application (extern)

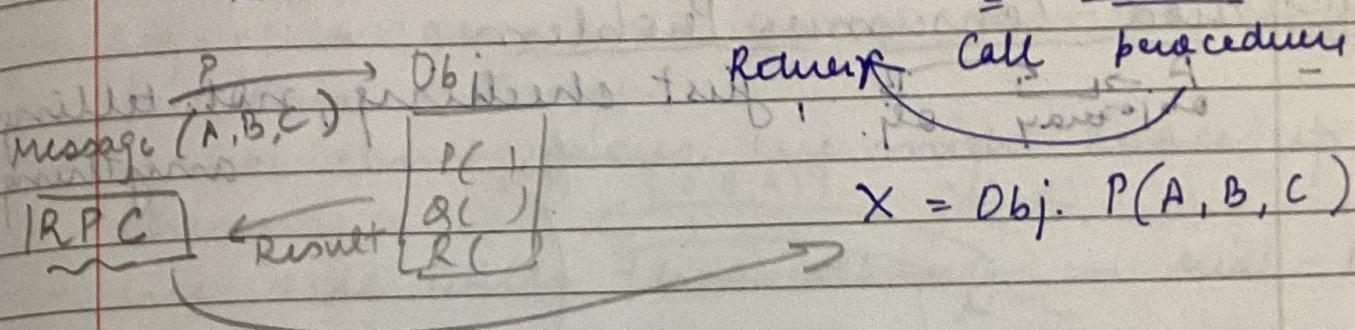
( $\text{class} . T$ ) =  $\text{class} . e$  . p -

part EM = environment

— answer private tail is treated as if it -  
code will at file and declared

### L11.1 Concept of msg. passing using Remote Procedure Call (RPC)

- we wait to get the ans. → u r stuck
- we're telling the obj., the procedure is inside the object
- control from two-way communication
- flight / hotel → may thing, u put your thing and then u get it later
- the sub-contractor can be in a diff. city
- Remote Procedure Call  $x = P(A, B, C)$

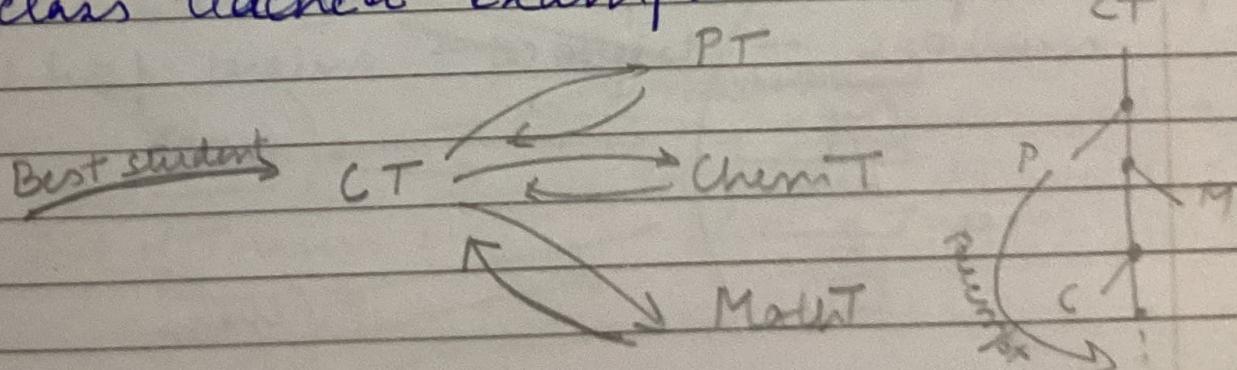


Eg. Ordering Food from a restaurant,

- phenomenal way of abstraction

### L11.2 Concept of concurrent exec" using polling and pumping

- class teacher example

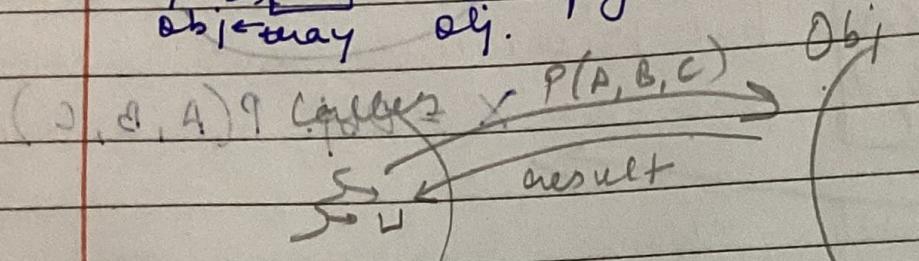


- things are going in kind of parallel

- concurrent → diff. things happening at the same time
    - ↳ multiple parts are active and each thing is doing sth. else
  - preemption - preempts sth. and does sth. more imp. first, kind of poking multiple threads — thing started and simultaneously u started sth.

## L11.3 Prod.-Consumer Problem

- 1.3 Prod. - Consumer Problem



(Consumer)

(iii) Collinear  $\Leftrightarrow Q(A, B, C), P(A, B, C)$   
Prove  $\leftarrow$

~~privated~~ Ansys  
~~Toronto~~

Class Teacher

PMS  
Math

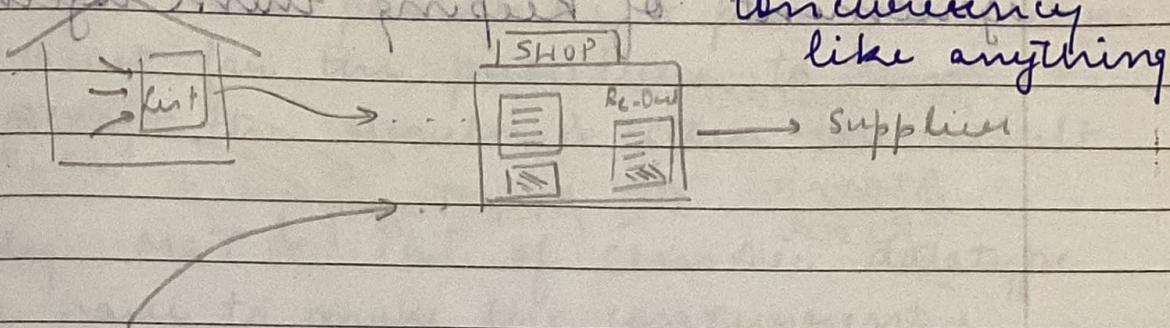
~~Phylogenetic trees~~  
~~about rules~~

3 (math)

→ (chem)

- atomic - reading & writing should't happen simultaneously  
(no concurrency)

- L11.4 Real time applica<sup>n</sup> of concurrency
- shopping bill data - customers are not under each other shopping with printer board
  - ~~customer~~ most popular fo! **Concurrencey like anything**



- no. of possibilities increases
- seq. of consumer things can be arranged but this is tough
- if we don't focus, we might loose up things
- we deal with these things in real-life always

L11.5 Applica<sup>n</sup> of concurrency to speedup the exec

- Train dataset
- all trains are running (printed) p3
- at any time, multiple trains are running
- On to a sta<sup>n</sup>, how many trains you have.
- resources like track / station is pre-occupied
- 1 hop trains → 2 object → 2 stations

write to db to insert, update, delete  
Stn & Train  
List of Trains  
List of Stations

traversing list (1) & (2) along with them &  
at point (1) (2) numbered & permanent

June ← Junc  
Borrower

earliest person and (A)9 = 8 gets into

## L11.6 Concept of Msg. Broadcasting

- send a msg. to everyone together
- broadcasting the question — to everyone
- agreement of keeping them together

principles w.r.t  
Broadcast

Atomic  
Read

R Am. Today

Atomic  
Write

## L11.7 Concurrency principles in input / output operations

- between → search query → submit ans.
- I/O (Input - Output) - interaction with user
- Forms (univite str.) - with restrictions
  - Submit button (with sanity)
- And result is shown, then computer gives the output
- This I/O work is also concurrency with a atomic model.
- Eg. Chatting also

## L11.8 Formalized Notes & Summary of Concurrency

- till now, only 1 step was executed at a time
- multiple lines occurring simultaneously → concurrency
- 2 procedure calls  $P()$  &  $Q()$  is concurrent
- Concurrency → procedure calls has to be unbundled
- The step  $B = P(A)$  has many actions

- For obj.  $X \& Y : X.P(A) \& Y.P(A)$ , may be executed concurrently  $\Rightarrow$  related to OOP
- $B = X.P(A)$  has actions
  - 1. start the procedure & pass parameters
  - 2. wait for the procedure to complete
  - 3. Accept the result & store the result

Ex. 2 obj. MAT & PHT of class Ave datatype  
(we have to make them concurrent)

- we need start ready & result (average)
  - start: start of average in ready
  - PHT: start of individual slot when
  - wait: wait for MAT & ready, average instead of HT & ready (avg)
- - Ave Maths = MAT.result (average)
  - Ave Physics = PHT.result (average)
- if, now want to add new students to both MAT & PHT

MAT. start (news, new M1)  
PHT. start (news, new M2)

$\Rightarrow$  MAT. start (average)  
MAT. start (adds, new M1)  
wait(MAT.ready (average)) & MAT.ready (adds)  
Ave Maths = MAT.result (average)

→ we can add aready & add ready fields

- to prevent corrupt<sup>n</sup> of the list, we have to make sure that we do only 1 list op<sup>n</sup> at a time. We can say that list field is atomic.
- for different coordination<sup>n</sup>, the obj. MAT should be shared.
- Polling - The caller started the procedure and then just waited for them to finish.
- Pre-emption - It waits for user request & stop it.
- Prod. - Consumer Model - I keep the things in tray (buffer) & consumer will take it from there.
- this prevents unwanted race condition b/w. read & write of the buffer.

## Week-12

CLASSTIME Pg. No.

Date / /

L12.1 Tutorial -1 - (Pseudocodes) for assignment -

Q. Find the no. of students based on total marks.

→ If  $A \geq 0$  & if  $X$  Total  $\geq 200$  then  $A = 1$  else (ii)

→  $B = 0$ , else,  $B = B + 1$  ans

→  $C = 0$ , else { If  $X$  Total  $\geq 200$  then  $C = C + 1$  (iii)}

marks calculated if number less than 200 then  $B = B + 1$  (ii)

If - Else

else { else }

$D = D + \text{minutes}$

else with alternation

with alternation

Q. Find the no. of students based on marks &

city of topper

→ Nested loop "if (name) {

"if (no.) {

do stn.

}

Q. Find the no. of students above avg. marks

→ multiple non-nested loops

2nd loop "ans." is dependent on the  
ans of 1st loop"

Q. Find the no. of pair of students who score  
same marks in maths

→ Nested loops "with else (multiple loops) " -

with 2nd loop to find marks & rank of student  
marks for subjects and average points = pair marks

3

"star [ ]

1. when  $i = j$  then  $i = j$  -

Q. Find subj. topper of 3 subjects in

→ same code for 3 subjs. → insert one loop into a nested loop with subj. as parameter

- Advantages of procedure -

- (i) avoid repetition of code
- (ii) divide complex problem into smaller one
- (iii) makes it easy to read the code
- (iv) modify of procedure becomes easier

- return

Terminates the procedure & flow of code goes back to the procedure call

- exitloop

Terminates the only loop from which the statement has been called

A = isEven (num)

Procedure isEven (n)

if ( $n \% 2 == 0$ ) {

return (true)

$i = 1, j = 1, n = 4, sum = 0$

while ( $i < n$ ) {

if ( $i == j$ ) {

exitloop

$j = j + 1$

$i = i + 1$

}

End isEven

where return repeats一直到 for next print  
sum = sum + j

## L12.2 Tutorial - 2 (List and Dictionaries)

- list is collection of values in seq. order, can store values of diff. datatypes, duplicates allowed, it preserves the order of elements  
[ ] nota"

- $l = []$  or  $length(l) == 0$ , denotes l is empty

- $list = list ++ list$   
↳ append & extend

- append → merges 2 lists ~ star [ ] -  
 l1 = [ ] l2 = ["a"] l1.append(l2) = [ "a" ]
- extend → adds element to list list -  
 l1 = [ ] append x l1.extend([x]) = [1, 2, 3, 4]  
x = 4 www  
empt
- eg. nest = [ ] q1. p = first; 2:4, 3:9 } nest -  
 d = [{"a": "A", "b": "B"}]  
 l = l + d,  $\Rightarrow$  l = [[1], 2:4, 3:9]]
- then, l = l + d,  $\Rightarrow$  l = [{""}, {"a": "A", "b": "B"}]]
- nest and set - nest and set -
- [ ] length = 0 inf st (l), is undefined = not inf st (l)
  - [1, 2, 3, 4] length = len first (l) last (l) = 4  
nest (l) = [2, 3, 4] init (l) = [1, 2, 3]

returns  
/ output  
types

length (l)  $\Rightarrow$  always integer  
 nest (l) = init (l)  $\Rightarrow$  always list  
 first (l) = nest (l)  $\Rightarrow$  depends on element  
next item of first of l - opt -

### - Insertion Sort

input list = [5, 2, 1, 3, 4], sorted list = []  
 foreach x in input list {  
 sorted list = insert an element (x, n)  
 }

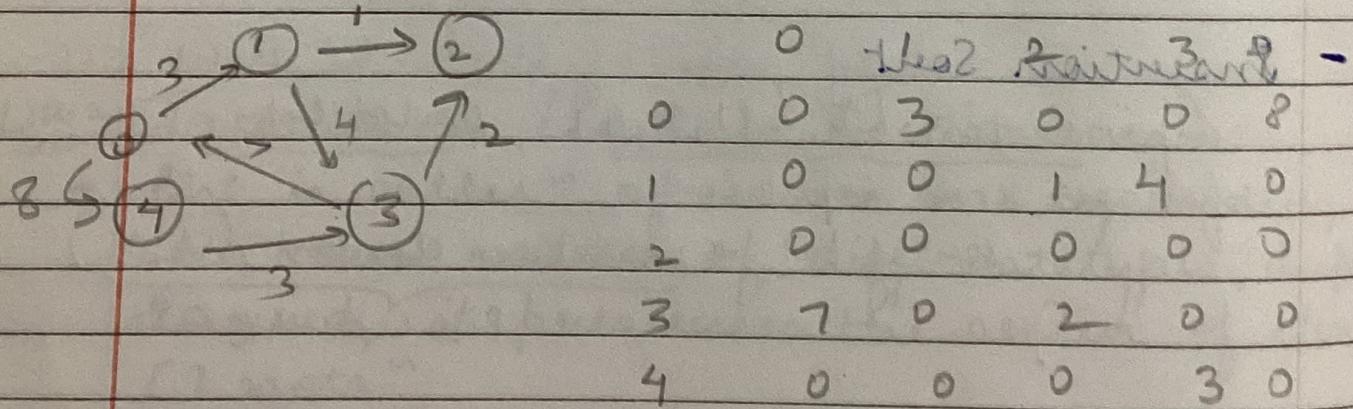
Procedures

- Dictionary - it is a collec" of element stored as key: value pair.
- It can store duplicate values, but keys are unique, no order preserved.

- `{ nota }`, value is accessed by key [key "nota"]
  - `keys(d)` → return a list of keys of dictionary  
 $d = \{ "a": 1, "b": 2 \}$   
 $\text{keys}(d) = ["a", "b"]$
  - `iskey(d, k)` → True if  $k$  is a key of a dictionary  
 $\text{iskey}(d, "a") \rightarrow \text{return (True)}$
  - **Index** **Key**
    - list has indices - dict has keys
    - it indicates the pos of element - unique identity to identify the key
    - it always integer - any datatype!
- $l = [10, 20, 30]$        $d = \{ 1: 10, 2: 20 \}$
- index  $l[0] = 10$        $d[2] = 20$

### L12.3 Tutorial - 3 (Graph & Adjacency Matrix)

- edge - labelled directed graph



The 2 Row Part -						
0	0	3	0	0	8	
1	0	0	1	4	0	
2	0	0	0	0	0	
3	7	0	2	0	0	
4	0	0	0	3	0	

elements for "edges"  $\Rightarrow a_{ij}$  - preexisting -  
 need edges : put as zeros  
 zero speed link, center Steinbeck units may fit -  
 decreased values are, in sum

matrix = { 0: {0:0, 1:3, 2:0, 3:0, 4:8},  
 1: {0:0, 1:0, 2:1, 3:4, 4:0},  
 2: ... similarity  
 3: ... for other ts  
 4: ... }

matrix  $[0][1] = 3$   
 "       $[0][4] = 3$       } list of edges

→ Procedure: triangulation (row by col.) - 2.7D -  
↳ produces a matrix of row x col.

matrix = {}

二

```
while (i < events) {
```

matrix [:] = [3]

J-20

while ( $j < \text{cols}$ ) {

$$mature [i][j] = 0$$

二十一

卷之三

efte fere as oratulas intendo amed -  
metum (metum

maturing matured

was and will remain my best

trash we consider as dirty

~~spine transverse tubercles at primary twist~~

L12.4 Tutorial - 4 (Recursion & DFS)

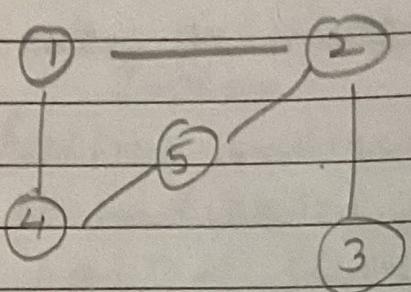
- Recursion = is a process in which a procedure calls itself. It simplifies the pseudo code.
- Problems with base problem only can be changed to recursive call.

```

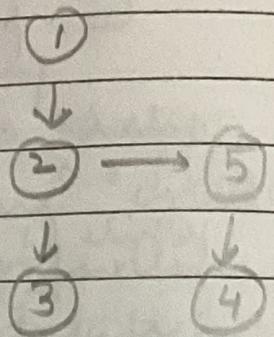
Procedure fac(n) → if (n == 0) {
    fact = 1
    return (1)
}
while (n > 0) {
    fact = fact * n
    n = n - 1
}
else {
    return (n * fact)
}
return (fact)

```

- DFS → recursive algorithm



Input: graph



Output : True

## L12.5 Final Summary

- Express problem solutions as seq. of steps
- find common patterns btw sol'n
- data sc. problems → on dataset
- first scanning to collect relevant info.
- processing the info.
- finally, organising the relationships - in a form
- Iterator - that scans the dataset (most powerful construct)
- Variables - storage units at higher level
- Assignment statements & accumulators
- Pseudocode & Flowcharts

- basic datatypes - boolean, int, char, string
- subtypes for more constraint
- list → items are of same datatype
- dictionary → Record → multiple datatype fields  
↳ we can keep adding fields
- concept of filtering
- procedure - to separate out the code
- parameters - user used for procedure
- this makes the code compact & easy to manage
- sequential iteration - to find above avg. students
- nested iteration - beginning to end. time  
↳ and store in list, dictionary, graph
- foreach iterates thru values in a list
- insertion sort - is a natural sorting algorithm

foreach k in keys(d) [ ]  
 {  
 } do sth with d[k] ] dictionary provides random access

- graphs represent relationships
- matrices are 2-d tables → m[i][j]
- Recursion - base case & inductive step  
↳ suspend the current computation till the recursive computation terminates
- Encapsulation - packs procedures with data elements  
↳ provides more modularity  
↳ disadj. → obj. need to created
- concurrency → allows several tasks to be executed simultaneously  
↳ producer-consumer model  
↳ atomic  
↳ Input/Output type