

Lecture 6L1 Logic Coverage eg. 2

- detecting type of Δ . 1 \rightarrow scalene, 2 \rightarrow isosceles, 3 \rightarrow eq. 4 \rightarrow not a Δ .
- If there is an int. variable, solve it in terms of inputs.

L2 Logic Coverage Specification

- Specifications can be in formal lang. or Eng. They have several logical conditions. If given in Eng., we need to derive logical predicates from the specification & then apply logic coverage as relevant.

Eg. name must be empty, state must be valid, pin 6 digits
 name != "" \wedge state in list \wedge (pin $\geq 000000 \wedge$ pin ≤ 999999)

- A CNF (conjunctive normal form) if it consists of clauses or disjuncts connected by conjunct operator (and). For a predicate of $a \wedge b \wedge c$, a major clause is made active by making all other clauses False.
- A DNF \rightarrow consists of clauses of conjuncts connected by disjunct op. Eg. $a \vee b \vee c$. A major clause is made active by making all other clauses false.

L3FSMs

- Transitions in FSMs have guards which are logical expressions.
- Train eg.

left door open

trainSpeed = 0 \wedge

pedfeetm = left \wedge

(locn = instn \vee

(currenstop \wedge driveridOpen \wedge locn = intunnel))

All doors close

- predicate has 6 clauses
- FSM model has 2 locaⁿ - in Staⁿ & in Tunnel. They both are false in ITR, which is not possible \Rightarrow error

L4 Logic Coverage Summary

- Predicate transformer deals with re-writing the predicate in the prog. so that they have only 1 clause each. It is mandatory for security testing where multiple clauses can lead to issues.
- ACC criteria is expensive for testers and can be overcome by re-writing the predicates.
- PT involves doing the control structure of the prog. using nested decision statements that represent how the various clauses in the predicate are connected to each other. The resulting prog. is difficult to read & maintain & hence not recommended.
- ACC criteria on the re-written prog. are not always eq \geq to the criteria on original prog.

e.g. if(a && b) : \rightarrow {
 s₁
 use:
 s₂

if (a)
 {
 if (b)
 {
 s₁
 use s₂
 close:
 s₂

ter : {t,t}, {t,f}, {f,t}

ter : { (t,t), (t,f), (f,t) }

L5 SMT Solvers

- The SAT problem determines if there exists an assignment of T/F values that satisfies a given propositional logic formula. It is NP-complete.
- But we deal with predicate logic, the SAT problem is undecidable. But some instances are decidable.
- Satisfiability modulo theories (SMT) is the problem of determining whether a given predicate is satisfiable.
- SMT solver uses SAT solver as underlying engine. They are used extensively in formal verification and program analysis, for proving the correctness of prog.