

Forces-Directed Graph Layout

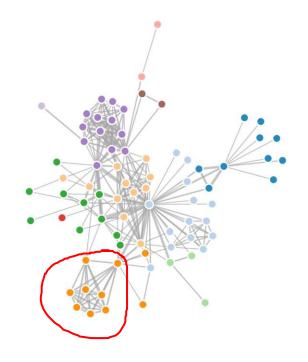
Scientific Visualization Professor Eric Shaffer



Force-Directed Graph Layout

Idiom: force-directed placement

- visual encoding
 - -link connection marks, node point marks
- considerations
 - -spatial position: no meaning directly encoded
 - left free to minimize crossings
 - -proximity semantics?
 - · sometimes meaningful
 - sometimes arbitrary, artifact of layout algorithm
 - tension with length
 - -long edges more visually salient than short
- tasks
 - –explore topology; locate paths, clusters
- scalability
 - -node/edge density E < 4N





Force-directed Layout Intuition

Goals

- Vertices well-distributed on the display
- Edges cross each other as little as possible.

Approach

- Simulate the graph as a physical system.
- Nodes are electrically charged particles and repulse each other
- Edges act as springs that attract connected nodes

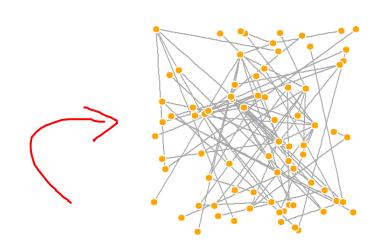
Result

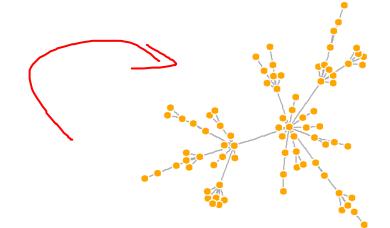
- Nodes are evenly distributed through the chart area
- Nodes which share more connections are closer to each other

"Fruchterman-Reingold is one of the most used force-directed layout algorithms out there." - The Internet

Graph Drawing by Force-Directed Placement Fruchterman, Thomas M. J.; Reingold, Edward M. (1991)

Developed at University of Illinois







Fruchterman-Reingold Details

$$F_{a}(n_{i}, n_{j}) = \frac{|p_{i} - p_{j}|^{2}}{k}$$

$$F_{r}(n_{i}, n_{j}) = -\frac{k^{2}}{|p_{i} - p_{j}|}$$

$$k = C \sqrt{\frac{area}{number\ of\ vertices}}$$

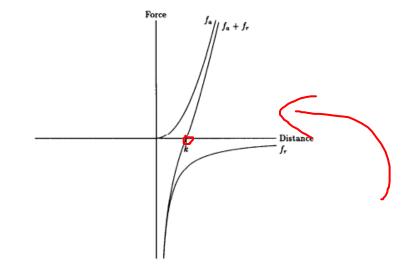


 p_i is the position of that vertex

 ${\it C}$ is a constant found experimentally to yield good result

k if vertices are uniformly distributed, would be radius of empty circle around vertex

the distance at which the forces will balance



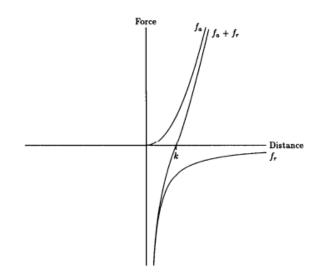


Fruchterman-Reingold Details

$$F_a(n_i, n_j) = \frac{|p_i - p_j|^2}{k}$$

$$F_r(n_i, n_j) = \frac{k^2}{|p_i - p_j|}$$

$$k = C \sqrt{\frac{area}{number\ of\ vertices}}$$



For each vertex n_i calculate a sum of forces:

- $F_a(n_i, n_j)$ between n_i and all neighbors n_j
- $\overline{F_r(n_i, n_j)}$ between n_i and all vertices n_j
- Each force has a magnitude and direction $\overleftarrow{p_j-p_i}$
- For repulsion, direction is negated
- Move p_i accordingly

Do this until change in positions below a threshold ...or you hit your limit on the number of iterations



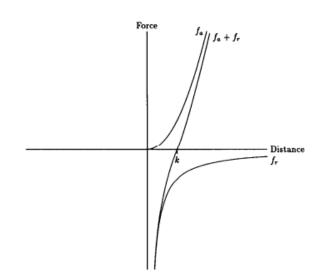
Fruchterman-Reingold...More Details

$$F_a(n_i, n_j) = \frac{|p_i - p_j|^2}{k}$$
 $F_r(n_i, n_j) = -\frac{k^2}{|p_i - p_j|}$

$$F_r(n_i, n_j) = \frac{k^2}{|p_i - p_j|}$$

$$k = C \sqrt{\frac{area}{number\ of\ vertices}}$$

- Direction $p_i p_i$ should be unit length vector
- Each iteration, stop vertices from moving outside display
- Apply some scale factor to the amount of movement of each vertex
 - Reduce this scale factor each iteration

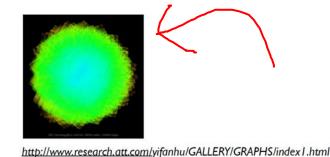


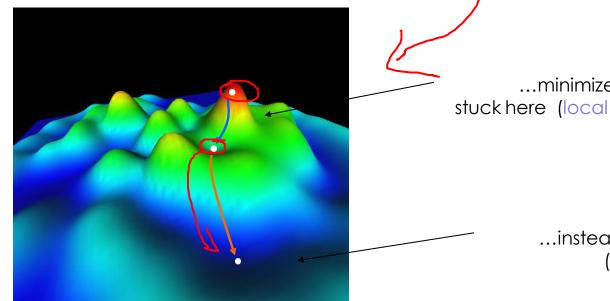


Disadvantages

- Computationally slow... $O(n^3)$
 - Can be sped up using spatial partitioning and calculating repulsion for only near-by nodes
 - Also can employ multi-level computation like Barnes-Hut and handle 1M node graph in seconds
- Visual limit is probably around 10K vertices/edges...hairball problem

Layout can get trapped in locally optimal rather than globally optimal state





...minimizer may get stuck here (local minimum)

> ...instead of arriving here (global minimum)



Force-Directed Layout Advantages

- Simple to implement
- Can work on any graph (e.g. not just trees)
- Can change force functions to use other data for specific applications
 - e.g. directed edges, edge weights, different classes of nodes, etc.
- Can be interactive

Works pretty well

