Maintenance Scheduling System

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Agenda FiraCode Nerd Font

- ► Introduction to Maintenance Scheduling
- Architecture of a Scheduling System
- Possible Contributions to Operation Research

Mathematical Notation: Sets

$$a \in A^m_{b,c}(t, x, y)$$

- a: set element
- A: set itself
- b: set element from set B
- c: set element from set C
- m: model formulation m
- t: time
- x: value of decision variable from a different model
- y: value of decision variable from a different model

Mathematical Notation: Parameters

 $name_of_parameter_{a,b}(t, x, y)$

- parameters are functions of set elements and input parameters
- ▶ a: set element from A
- b: set element from B
- t: time
- x: value of decision variable from another model
- y: value of decision variable from another model

Mathematical Notation: Variables

$x_{a,b}^m(t)$

- variables are functions of set elements, specified model, and time
- x: decision variable
- a: set element from A
- b: set element from B
- m: specifing the model
- t: time
- Notice: decision variables cannot depend on other decision variables as it would make them belong to the same model.

Strategic

 $\epsilon_{pr}(\tau) \in \mathbb{R}^+ \quad \forall p \in P(\tau) \quad \forall r \in R(\tau)$

Meta variables: $s \in S$ (1) $\beta(\tau)$ (2) $\tau \in [0, \infty]$ (3) Minimize: $\sum_{w \in W(\tau)} \sum_{p \in P(\tau)} strategic_value_{wp}(\tau) \cdot \alpha_{wp}(\tau)$ $+ \sum_{p \in P(\tau)} \sum_{r \in R(\tau)} strategic_penalty \cdot \epsilon_{pr}(\tau)$ $+ \sum_{p \in P(\tau)} \sum_{w1 \in W(\tau)} \sum_{w2 \in W(\tau)} \text{ clustering_value}_{w1,w2} \cdot \alpha_{w1p}(\tau) \cdot \alpha_{w2p}(\tau)$ (4) Subject to: $\sum_{w \in W(\tau)} work_order_work_{wr} \cdot \alpha_{wp}(\tau) \leq \ resource_{pr}(\tau,\beta(\tau)) + \epsilon_{pr}(\tau) \quad \forall p \in P(\tau) \quad \forall r \in R(\tau)$ (5) $\sum_{w \in W(\tau)} \alpha_{wp}(\tau) = 1 \quad \forall p \in P(\tau)$ (6) $\alpha_{wp}(\tau) = 0 \quad \forall (w, p) \in exclude(\tau)$ (7) $\alpha_{wp}(\tau) = 1 \quad \forall (w, p) \in include(\tau)$ (8) $\alpha_{wp}(\tau) \in \{0, 1\} \quad \forall w \in W(\tau) \quad \forall p \in P(\tau)$ (9)

(10)

Tactical

Meta variables: $s \in S$ (11) $\alpha(\tau)$ (12) (13) $\tau \in [0, \infty]$ Minimize: $\sum_{\mathbf{o} \in O(\tau, \mathbf{o}(\tau))} \sum_{\mathbf{d} \in D(\tau)} \textit{tactical_value}_{d\mathbf{o}}(\tau) \cdot \beta_{d\mathbf{o}}(\tau) + \sum_{r \in R(\tau)} \sum_{\mathbf{d} \in D(\tau)} \textit{tactical_penalty} \cdot \mu_{rd}(\tau)$ (14)Subject to: $\sum \quad work_o(\tau) \cdot \beta_{do}(\tau) \leq tactical_resource_{dr}(\tau) + \mu_{rd}(\tau) \forall d \in D(\tau) \quad \forall r \in R(\tau)$ (15) $o \in O(\tau, \alpha(\tau))$ latest finish_o(τ) $\qquad \qquad \sigma_{do}(\tau) = \mathit{duration}_o(\tau) \quad \forall o \in \mathit{O}(\tau, \alpha(\tau))$ (16) $\sigma_{d^* \in D_{\mathit{dereshop}(\tau)}(\tau)} \sigma_{d^* \circ}(\tau) = \mathit{duration}_{o}(\tau) \cdot \eta_{do}(\tau) \quad \forall o \in \mathit{O}(\tau, \alpha(\tau)) \quad \forall d \in \mathit{D}(\tau)$ (17) $\sum_{o \in O(\tau, \alpha(\tau))} \eta_{do}(\tau) = 1, \quad \forall d \in D(\tau)$ $\sum_{d \in D(\tau)} d \cdot \sigma_{do1}(\tau) + \Delta_o(\tau) = \sum_{d \in D(\tau)} d \cdot \sigma_{do2}(\tau) \quad \forall (o1, o2) \in \mathit{finish_start}_{o1, o2}$ (18) $\sum_{d \in D(\tau)} d \cdot \sigma_{do1}(\tau) = \sum_{d \in D(\tau)} d \cdot \sigma_{do2}(\tau) \quad \forall (o1, o2) \in \textit{start_start}_{o1, o2}$ (19) $\beta_{do}(\tau) \le number_o(\tau) \cdot operating time_o \quad \forall d \in D(\tau) \quad \forall o \in O(\tau, \alpha(\tau))$ (20) $\beta_{do}(\tau) \in \mathbb{R}$ $\forall d \in D(\tau) \ \forall o \in O(\tau, \alpha(\tau))$ $\mu_{rd}(\tau) \in \mathbb{R}$ $\forall r \in R(\tau) \forall d \in D(\tau)$ (22) $\sigma_{do}(\tau) \in \{0, 1\}$ $\forall d \in D(\tau) \forall o \in O(\tau, \alpha(\tau))$ (23) $n_{do}(\tau) \in \{0, 1\}$ $\forall d \in D(\tau) \ \forall o \in O(\tau, \alpha(\tau))$ (24) $\Delta_o(\tau) \in \{0, 1\} \quad \forall o \in O(\tau, \alpha(\tau))$ (25)

Supervisor

 $\rho_a(\tau) \in [lower \ activity \ work_a(\tau), work_a(\tau)] \quad \forall a \in A(\tau, \alpha(\tau))$

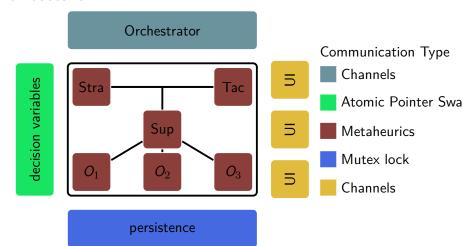
Meta variables: $z \in Z$ (26) $\alpha(\tau)$ (27) $\theta(\tau)$ (28)(29) $\tau \in [0, \infty]$ Maximize: $\sum \quad \sum \ \, \textit{supervisor_value}_{\mathsf{at}}(\tau, \lambda_{\mathsf{t}}(\tau), \Lambda_{\mathsf{t}}(\tau)) \cdot \gamma_{\mathsf{at}}(\tau)$ (30)Subject to: $\sum_{\mathbf{a} \in A_{\mathbf{o}}(\tau,\alpha(\tau))} \rho_{\mathbf{a}}(\tau) = \mathit{work}_{\mathbf{o}}(\tau) \quad \forall \mathbf{o} \in \mathit{O}(\tau,\alpha(\tau))$ (31) $\sum_{\mathbf{t} \in T(\tau)} \sum_{\mathbf{a} \in A_0(\tau, \alpha(\tau))} \gamma_{\mathbf{a}\mathbf{t}}(\tau) = \phi_{\mathbf{o}}(\tau) \cdot \mathit{number}_{\mathbf{o}}(\tau) \quad \forall \mathbf{o} \in O(\tau, \alpha(\tau))$ (32) $\sum_{\mathbf{o} \in O_{\mathbf{w}}(\tau, \alpha(\tau))} \phi_{\mathbf{o}}(\tau) = |O_{\mathbf{w}}(\tau, \alpha(\tau))| \quad \forall \mathbf{w} \in W(\tau, \alpha(\tau))$ (33) $\sum_{\mathbf{a} \in A_{\sigma}(\tau, \alpha(\tau))} \gamma_{\mathbf{a}\mathbf{t}}(\tau) \leq 1 \quad \forall \mathbf{o} \in O(\tau, \alpha(\tau)) \quad \forall \mathbf{t} \in T(\tau)$ (34) $\gamma_{st}(\tau) \le feasible_{st}(\theta(\tau)) \quad \forall o \in O(\tau, \alpha(\tau)) \quad \forall t \in T(\tau)$ (35) $\gamma_{st}(\tau) \in \{0, 1\} \quad \forall o \in O(\tau, \alpha(\tau)) \quad \forall t \in T(\tau)$ (36)

(37)

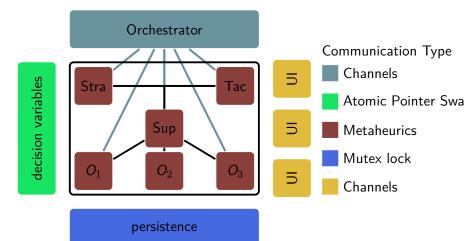
Operational

Meta variables:	
$t \in T(\tau)$	(38)
$\alpha(\tau)$	(39)
$\gamma(\tau)$ $\tau \in [0, \infty]$	(40)
	(41)
Maximize:	
\sum $\delta_{ak}(au)$	(42)
$a \in A(\tau, \gamma_t(\tau)) \ k \in K(\gamma(\tau))$	
Subject to:	
$\sum_{k \in K(\gamma(\tau))} \delta_{ak}(\tau) \cdot \pi_{ak}(\tau) = activity_work_{a}(\tau, \rho(\tau)) \cdot \theta (\tau) \forall a \in A(\tau, \gamma_{t}(\tau))$	(43)
$\lambda_{a21}(\tau) \ge \Lambda_{a1/ast(a1)}(\tau) + preparation_{a1,a2} \forall a1 \in A(\tau, \gamma_t(\tau)) \forall a2 \in A(\tau, \gamma_t(\tau))$	(44)
$\lambda_{ak}(\tau) \ge \Lambda_{ak-1}(\tau) - constraint_limit \cdot (2 - \pi_{ak}(\tau) + \pi_{ak-1}(\tau))$	
$\forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau))$	(45)
$\delta_{ak}(\tau) = \Lambda_{ak}(\tau) - \lambda_{ak}(\tau) \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau))$	(46)
$\lambda_{ak}(\tau) \ge event_{ie} + duration_{ie} - constraint_limit \cdot (1 - \omega_{akie}(\tau))$	
$\forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \forall i \in I(\tau) \forall e \in E(\tau)$	(47)
$\Lambda_{ak}(\tau) \le event_{ie} + constraint_limit \cdot \omega_{akie}(\tau)$	
$\forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \forall i \in I(\tau) \forall e \in E(\tau)$	(48)
$\lambda_{a1}(\tau) \ge time_window_start_a(\beta(\tau)) \forall a \in A(\tau, \gamma_t(\tau))$	(49)
$\Lambda_{alast(a)}(\tau) \le time_window_finish_a(\beta(\tau)) \forall a \in A(\tau, \gamma_t(\tau))$	(50)
$\pi_{ak}(\tau) \in \{0,1\}$ $\forall a \in A(\tau, \gamma_t(\tau))$ $\forall k \in K(\gamma(\tau))$	(51)
$\lambda_{ak}(\tau) \in [availability_start(\tau), availability_finish(\tau)]$	
$\forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau))$	(52)
$\Lambda_{ak}(\tau) \in [availability_start(\tau), availability_finish(\tau)]$	
$\forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau))$	(53)
$\delta_{ak}(\tau) \in [0, work_{a_to_{-o}(a)}(\tau)] \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau))$	(54)
$\omega_{akie}(\tau) \in \{0,1\} \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \forall i \in I(\tau) \forall e \in E(\tau)$	(55)
$\theta_a(\tau) \in \{0, 1\} \forall a \in A(\tau, \gamma_t(\tau))$	(56)

Architecture

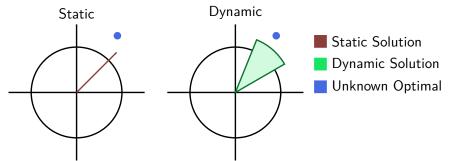


Architecture



- 1. Reactive Versus Static Constraints
- 2. Dynamic
- 3. Business
- 4. Technical
- 5. Academic

Dynamic versus Static Models



- Mathematical models guide direction but does not provide direct solutions.
- ► Static solutions are rarely fully executable.
- Dynamic models are less constrained and ensure a contained optimal solution.

Dynamic

Business