

Multi-agent Maintenance Scheduling: The Making of a Science

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: excess capacity for each day

usually quite similar.

This figure shows

Chapter 1

Introduction

Maintenance scheduling is in its nature a multi actor process. Many stakeholders have to coordinate in both time and space to allow for an efficient and effective execution. This thesis will propose a generalized multi-agent scheduling system and it will argue that for the field of maintenance scheduling to move forward similar approaches will have to be adopted. Other approaches may be very different but they will share many of the aspects.

This Ph.D. will present a generalized dynamic multi-model approach to maintenance scheduling which will be model after a practical maintenance handbook Palmer [2019]. This book written by the experienced practitioner Richard D. Palmer will be a guiding light throughout the thesis, so it serves as the main source and validation, or maybe invalidation is a better word, as we explore the academic maintenance scheduling literature and also, and more importantly, it will also be the source which above all else will us us through the perilous process of create a generalize model setup for maintenance scheduling.

1.1 The General Maintenance Scheduling Process

This section will provide an overview of the maintenance scheduling process in the most abstracted way possible. It will be important to understand this setup thoroughly as most industries that perform maintenance of a considerable scale follow this process. Many industries are of course unique and deviate from general framework in specific work but the fundamentals are

Chapter 2

Modelling the Generalized Setup

To model the maintenance process in its entirety we will need mathematical modelling tools that are powerful enough to describe the maintenance scheduling process. As the process requires multiple different actors we need a system that enables us to model each of them and specific

To effectively model a system that contains multiple actors also The system will be described in accordance with the ?? Palmer [2019].

The maintenance scheduling problem is NP-hard and real-time optimal solutions will never be a feasible approach unless we use a multi-model setup where each model enriches the overall solution in the way that it is most capable of.

2.1 Sets

2.1.1 Strategic

* W : WorkOrders * P : Periods * R : Resources

2.1.2 Tactical

* O : Operations * D : Days * R : Resources

2.1.3 Supervisor

* A : Activities * C : Crew

2.1.4 Operational

* K : Segments * A : Activities * I : Instances
* E : Events

2.2 Parameters

2.2.1 Strategic

* $strategic_value_{wp}$: value of WorkOrder in period
* $strategic_penalty$: penalty of exceeding capacity
* $clustering_value_{o1,o2}$: Gain from clustering two work orders
* $strategic_resource_{pr}$: Capacity by period and resource

2.2.2 Tactical

* $tactical_value_{do}$: value of tactical operation (RENAME)
* $tactical_penalty$: penalty of exceeding daily trait capacity
* $work_o$: hours of work in an operation (RENAME)
* $tactical_resource_{dr}$: hours of available work for r
* $start_start_{o1,o2}$: relations between operations that have start-start
* $finish_start_{o1,o2}$: relations between operations that have finish-start
* $number_o$: number of people required to complete the operation
* $operating_time_o$: allowed hours to be worked on the operation per day
* $duration_lower_o$: shortest allowed duration
* $duration_upper_o$: longest allowed duration

2.2.3 Supervisor

* $supervisor_value_{oc}$: value of assigning the operation to crew member c
* $number_o$: Number of workers that the operation needs to be assigned to
* $feasible_{ow}$: parameter specifying whether worker w can execute operation o

2.2.4 Operational

* $work_i$: total hours of work in an activity (COORDINATE NAMING)
* $preparation_{ij}$: preparation between job i and j
* $event_{te}$: start time of event e (LOGIC ERROR HERE)
* $duration_{te}$: duration of event e
* $constraint_{limit}$: Big M to model if statement (Metaheuristic does not care)
* $time_window_start_i$: time window start of job i
* $time_window_finish_i$: time window finish of job i
* $availability_start$: start of availability of operational model
* $availability_end$: end of availability of operational model

2.3 Variables

2.3.1 Strategic

- * α_{wp} : Assignment of WorkOrder w to period p
- * ϵ_{ptau} : penalty for exceeding capacity in period p for trait τ

2.3.2 Tactical

- * β_{do} : work for each day
- * μ_{cd} : penalty for exceeding capacity

2.3.3 Supervisor

- * γ_{at} : Determining the assignment of each activity to each technician.

2.3.4 Operational

- * δ_{ik} : processing time of job i 's k th segment
- * π_{ik} : binary variable specifying whether the k th segment for job i is active
- * λ_{ik} : start time of segment k of job i
- * Λ_{ik} : finish time of segment k of job i
- * ω_{ikte} : binary variable determining the position of event e in the ordering of segment k for job i (REWRITE LOGIC ERROR ON "t", MISSIN: constrain y not all can be one)
- * θ_i : binary variable for each job determining whether or not it is in the solution

2.4 Metavariables

- * τ : absolute time (RENAME TO TAU)

2.5 The Strategic Model

The Strategic Model have multiple different purposes.

- Schedule Work Order out across the weekly periods
- Prioritize all the different released work orders
- Respect the available weekly hours available for each trait

The Strategic model is responsible for grouping work orders into weekly or biweekly periods depending on which kind of maintenance setup that one is running. This kind of model closely resembles a variant of the multi-compartment multi-knapsack problem.

Minimize:

$$\sum_{w \in W(\tau, \alpha(\tau))} \sum_{p \in P(\tau)} strategic_value_{wp}(\tau) \cdot \alpha_{wp}(\tau) \quad (2.1)$$

$$+ \sum_{p \in P(\tau)} \sum_{r \in R(\tau)} strategic_penalty \cdot \epsilon_{pr}(\tau) \quad (2.2)$$

$$+ \sum_{p \in P(\tau)} \sum_{w1 \in W(\tau, \alpha(\tau))} \sum_{w2 \in W(\tau, \alpha(\tau))} clustering_value_{w1, w2} \cdot \alpha_{w1p}(\tau) \cdot \alpha_{w2p}(\tau) \quad (2.3)$$

Subject to:

$$\begin{aligned} \sum_{w \in W(\tau, \alpha(\tau))} work_order_work_{wr} \cdot \alpha_{wp}(\tau) \\ \leq resource_{pr}(\tau) + \epsilon_{pr}(\tau) \\ \forall p \in P(\tau), \forall r \in R(\tau) \end{aligned} \quad (2.4)$$

$$\begin{aligned} \sum_{w \in W(\tau, \alpha(\tau))} \alpha_{wp}(\tau) = 1 \\ \forall p \in P(\tau) \end{aligned} \quad (2.5)$$

$$\begin{aligned} \alpha_{wp}(\tau) = 0 \\ \forall (w, p) \in exclude(\tau) \end{aligned} \quad (2.6)$$

$$\begin{aligned} \alpha_{wp}(\tau) = 1 \\ \forall (w, p) \in include(\tau) \end{aligned} \quad (2.7)$$

$$\begin{aligned} \alpha_{wp}(\tau) \in \{0, 1\} \\ \forall w \in W(\tau, \alpha(\tau)), \forall p \in P(\tau) \end{aligned} \quad (2.8)$$

$$\begin{aligned} \epsilon_{pr}(\tau) \in \mathbb{R}^+ \\ \forall p \in P(\tau), \forall r \in R(\tau) \end{aligned} \quad (2.9)$$

$$\tau \in [0, \infty] \quad (2.10)$$

2.6 The Tactical Model

- Respect precedence constraints
- Respect daily resource requirements for each trait
- Penalize exceeded daily capacity

After the strategic model has optimized its schedule the tactical agent will continue scheduling the output at a more detailed level. This means that now the tactical agent will schedule out on each of the days of the work orders scheduled by the strategic agent.

below we see the model for the tactical agent.

Minimize:

$$\begin{aligned} & \sum_{o \in O(\tau, \alpha(\tau))} \sum_{d \in D(\tau)} tactical_value_{do}(\tau) \cdot \beta_{do}(\tau) \\ & + \sum_{r \in R(\tau)} \sum_{d \in D(\tau)} tactical_penalty \cdot \mu_{rd}(\tau) \end{aligned} \quad (2.11)$$

Subject to:

$$\begin{aligned} & \sum_{o \in O(\tau, \alpha(\tau))} work_o(\tau) \cdot \beta_{do}(\tau) \\ & \leq tactical_resource_{dr}(\tau) + \mu_{rd}(\tau) \\ & \forall d \in D(\tau), \forall r \in R(\tau) \end{aligned} \quad (2.12)$$

$$\begin{aligned} & \sum_{d^* \in D_{duration_upper_o}(\tau)} \sigma_{d^*o}(\tau) \\ & \geq duration_lower_o(\tau) \cdot \eta_{do}(\tau) \\ & \forall o \in O(\tau, \alpha(\tau)) \quad \forall d \in D(\tau) \end{aligned} \quad (2.13)$$

$$\begin{aligned} & \sum_{d^* \in D_{duration_upper_o}(\tau)} \sigma_{d^*o}(\tau) \\ & \leq duration_upper_o(\tau) \cdot \eta_{do}(\tau) \\ & \forall o \in O(\tau, \alpha(\tau)) \forall d \in D(\tau) \end{aligned} \quad (2.14)$$

$$\begin{aligned} & \sum_{o \in O(\tau, \alpha(\tau))} \eta_{do}(\tau) = 1, \\ & \forall d \in D(\tau) \\ & \sum_{d \in D(\tau)} d \cdot \beta_{do1}(\tau) + \Delta_o(\tau) = \sum_{d \in D(\tau)} d \cdot \beta_{do2}(\tau) \\ & \forall (o1, o2) \in finish_start_{o1, o2} \end{aligned} \quad (2.15)$$

$$\begin{aligned} & \sum_{d \in D(\tau)} d \cdot \beta_{do1}(\tau) = \sum_{d \in D} d \cdot \beta_{do2}(\tau) \\ & \forall (o1, o2) \in start_start_{o1, o2} \end{aligned} \quad (2.16)$$

$$\begin{aligned} & \beta_{do}(\tau) \leq number_o(\tau) \cdot operating_time_o \\ & \forall d \in D(\tau) \forall o \in O(\tau, \alpha(\tau)) \end{aligned} \quad (2.17)$$

$$\begin{aligned} & \beta_{do}(\tau) \in \{0, 1\} \\ & \forall d \in D(\tau) \forall o \in O(\tau, \alpha(\tau)) \end{aligned} \quad (2.18)$$

$$\begin{aligned} & \mu_{rd}(\tau) \in \mathbb{R} \\ & \forall r \in R(\tau) \forall d \in D(\tau) \end{aligned} \quad (2.19)$$

$$\begin{aligned} & \sigma_{do}(\tau) \in \{0, 1\} \\ & \forall d \in D(\tau) \forall o \in O(\tau, \alpha(\tau)) \end{aligned} \quad (2.20)$$

$$\begin{aligned} & \eta_{do}(\tau) \in \{0, 1\} \\ & \forall d \in D(\tau) \forall o \in O(\tau, \alpha(\tau)) \end{aligned} \quad (2.21)$$

$$\begin{aligned} & \Delta_o(\tau) \in \{duration_lower_o(\tau), \\ & \quad duration_upper_o(\tau)\} \\ & \forall o \in O(\tau, \alpha(\tau)) \end{aligned} \quad (2.22)$$

$$\tau \in [0, \infty] \quad (2.23)$$

The tactical model is responsible for providing an initial suggestion for a weekly schedule,

2.7 The Supervisor Model

The maintenance supervisor is considered the most central person in a maintenance scheduling system. All the work of the planner and scheduler should be considered a service for the supervisor.

The supervisor has multiple different responsibilities among them are:

- Assigning work orders
- Creating a daily schedule
- Keeping the schedule up-to-date

Maximize:

$$\sum_{a \in A(\tau, \beta(\tau))} \sum_{t \in T(\tau)} supervisor_value_{at}(\tau) \cdot \gamma_{at}(\tau) \quad (2.24)$$

Subject to:

$$\sum_{a \in A_o(\tau, \beta(\tau))} \rho_a(\tau) = work_o(\tau) \quad (2.25)$$

$$\forall o \in O(\tau, \sigma(\tau))$$

$$\sum_{t \in T(\tau)} \sum_{a \in A_o(\tau, \beta(\tau))} \gamma_{at}(\tau) = \phi_o(\tau) \cdot number_o(\tau) \quad (2.26)$$

$$\forall o \in O(\tau, \sigma(\tau))$$

$$\sum_{o \in O_w(\tau, \sigma(\tau))} \phi_o(\tau) = operations_in_work_order_w \quad (2.27)$$

$$\forall w \in W(\tau, \alpha(\tau))$$

$$\sum_{a \in A_o(\tau, \beta(\tau))} \gamma_{at}(\tau) \leq 1 \quad (2.28)$$

$$\forall o \in O(\tau, \sigma(\tau)) \forall t \in T(\tau)$$

$$\gamma_{at}(\tau) \leq feasible_{at}(\theta(\tau)) \quad (2.29)$$

$$\forall o \in O(\tau, \sigma(\tau)) \forall t \in T(\tau)$$

$$\gamma_{at}(\tau) \in \{0, 1\} \quad (2.30)$$

$$\forall o \in O(\tau, \sigma(\tau)) \forall t \in T(\tau)$$

$$\rho_a(\tau) \in [lower_activity_work_a(\tau), work_a(\tau)] \quad (2.31)$$

$$\forall a \in A(\tau, \beta(\tau))$$

$$\tau \in [0, \infty] \quad (2.32)$$

the operation can be assigned to a specific operational model comes from the operational model itself and is captured in the.

Can this be done? What should the Supervisor have here? He should have what is necessary to handle the.

2.8 The Operational Model

Here the o is a single operation and $o2$ is another operation. It is crucial to understand here that the main decision variable, x defines an ordering of the operations that a single operational agent will do the operations in.

The $\lambda_{ak}(\tau)$ is the start time of job i in segment k and $\Lambda_{ak}(\tau)$ is the finish time of job i in segment k . $\delta_{ak}(\tau)$ is the processing time of

In the supervisor model shown in ?? the set O and W comes from the tactical algorithm and value v and the information of whether or not

each segment.

Maximize:

$$\sum_{a \in A(\tau, \gamma_t(\tau))} \sum_{k \in K(\gamma(\tau))} \delta_{ak}(\tau) \quad (2.33)$$

Subject to:

$$\begin{aligned} & \sum_{k \in K(\gamma(\tau))} \delta_{ak}(\tau) \cdot \pi_{ak}(\tau) \\ &= \text{activity_work}_a(\tau, \rho(\tau)) \cdot \theta(\tau) \\ & \forall a \in A(\tau, \gamma_t(\tau)) \end{aligned} \quad (2.34)$$

$$\begin{aligned} \lambda_{a21}(\tau) &\geq \Lambda_{a1\text{last}(a1)}(\tau) + \text{preparation}_{a1,a2} \\ & \forall a1 \in A(\tau, \gamma_t(\tau)) \forall a2 \in A(\tau, \gamma_t(\tau)) \end{aligned} \quad (2.35)$$

$$\begin{aligned} \lambda_{ak}(\tau) &\geq \Lambda_{ak-1}(\tau) \\ &- \text{constraint_limit} \cdot (2 - \pi_{ak}(\tau) + \pi_{ak-1}(\tau)) \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \end{aligned} \quad (2.36)$$

$$\begin{aligned} \delta_{ak}(\tau) &= \Lambda_{ak}(\tau) - \lambda_{ak}(\tau) \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \end{aligned} \quad (2.37)$$

$$\begin{aligned} \lambda_{ak}(\tau) &\geq \text{event}_{ie} + \text{duration}_{ie} \\ &- \text{constraint_limit} \cdot (1 - \omega_{akie}(\tau)) \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \\ & \forall i \in I(\tau) \forall e \in E(\tau) \end{aligned} \quad (2.38)$$

$$\begin{aligned} \Lambda_{ak}(\tau) &\leq \text{event}_{ie} + \text{constraint_limit} \cdot \omega_{akie}(\tau) \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \\ & \forall i \in I(\tau) \forall e \in E(\tau) \end{aligned} \quad (2.39)$$

$$\begin{aligned} \lambda_{a1}(\tau) &\geq \text{time_window_start}_a(\beta(\tau)) \\ & \forall a \in A(\tau, \gamma_t(\tau)) \end{aligned} \quad (2.40)$$

$$\begin{aligned} \Lambda_{a\text{last}(a)}(\tau) &\leq \text{time_window_finish}_a(\beta(\tau)) \\ & \forall a \in A(\tau, \gamma_t(\tau)) \end{aligned} \quad (2.41)$$

$$\begin{aligned} \pi_{ak}(\tau) &\in \{0, 1\} \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \end{aligned} \quad (2.42)$$

$$\begin{aligned} \lambda_{ak}(\tau) &\in [\text{availability_start}(\tau), \\ & \text{availability_finish}(\tau)] \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \end{aligned} \quad (2.43)$$

$$\begin{aligned} \Lambda_{ak}(\tau) &\in [\text{availability_start}(\tau), \\ & \text{availability_finish}(\tau)] \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \end{aligned} \quad (2.44)$$

$$\begin{aligned} \delta_{ak}(\tau) &\in [0, \text{work}_{a_to_o(a)}(\tau)] \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \end{aligned} \quad (2.45)$$

$$\begin{aligned} \omega_{akie}(\tau) &\in \{0, 1\} \\ & \forall a \in A(\tau, \gamma_t(\tau)) \forall k \in K(\gamma(\tau)) \\ & \forall i \in I(\tau) \forall e \in E(\tau) \end{aligned} \quad (2.46)$$

$$\begin{aligned} \theta_a(\tau) &\in \{0, 1\} \\ & \forall a \in A(\tau, \gamma_t(\tau)) \end{aligned} \quad (2.47)$$

possible sections * All Stochasticity will be handled by user interaction. * In maintenance scheduling, data is assumed to be correct. Yes that must be the prevailing idea here.

Bibliography

Richard D. Palmer. *Maintenance Planning and Scheduling Handbook, 4th Edition*. McGraw Hill, 4th edition edition, September 2019.