Humuhumunukunukuapua'a UFMG

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15 de Março de 2023

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1 Matematica

1.1 2-SAT

```
// solve() retorna um par. o first fala se eh possivel
// atribuir, o second fala se cada variavel eh verdadeira
// O(|V|+|E|) = O(\#variaveis + \#restricoes)
138 struct sat {
e6c
       int n, tot;
       vector < vector < int >> g;
789
       vector < int > vis, comp, id, ans;
        stack<int> s;
4ce
141
        sat() {}
        sat(int n_{-}) : n(n_{-}), tot(n), g(2*n) {}
172
        int dfs(int i, int& t) {
f32
            int lo = id[i] = t++:
cf0
efc
            s.push(i), vis[i] = 2;
48e
            for (int j : g[i]) {
740
                if (!vis[j]) lo = min(lo, dfs(j, t));
994
                else if (vis[j] == 2) lo = min(lo, id[j]);
d64
            }
            if (lo == id[i]) while (1) {
3de
                int u = s.top(); s.pop();
3c3
                vis[u] = 1, comp[u] = i;
9c5
91d
                if ((u>1) < n \text{ and } ans[u>1] == -1) ans[u>1] = \sim u\&1:
2ef
                if (u == i) break;
60d
            }
253
            return lo;
dec
        }
74a
        void add_impl(int x, int y) { // x -> y = !x ou y
26a
            x = x >= 0 ? 2*x : -2*x-1;
2b8
            y = y >= 0 ? 2*y : -2*y-1;
            g[x].push_back(v);
a1e
            g[y^1].push_back(x^1);
1e2
ef0
e85
        void add_cl(int x, int y) { // x ou y
0b5
            add_impl(\sim x, y);
254
487
        void add_xor(int x, int y) { // x xor y
0b7
            add_cl(x, y), add_cl(\sim x, \sim y);
9a1
        }
```

```
978
        void add_eq(int x, int y) { // x = y
c86
             add_xor(\simx, y);
b91
        }
b10
        void add_true(int x) { // x = T
18b
             add_impl(\sim x, x);
9e2
d14
        void at most one(vector<int> v) { // no max um verdadeiro
54d
             g.resize(2*(tot+v.size()));
f14
             for (int i = 0; i < v.size(); i++) {</pre>
869
                 add_impl(tot+i, ~v[i]);
a8f
                 if (i) {
b6a
                      add_impl(tot+i, tot+i-1);
3d3
                      add_impl(v[i], tot+i-1);
0f7
                 }
             }
084
258
             tot += v.size();
b00
        }
        pair < bool, vector < int >> solve() {
a8e
27b
             ans = vector < int > (n, -1):
6bb
             int t = 0:
0de
             vis = comp = id = vector\langle int \rangle (2*tot, 0);
53 c
             for (int i = 0; i < 2*tot; i++) if (!vis[i]) dfs(i, t);</pre>
f88
             for (int i = 0; i < tot; i++)</pre>
4 c 9
                 if (comp[2*i] == comp[2*i+1]) return {false, {}};
997
             return {true. ans}:
7b3
ef6 }:
```

1.2 Avaliacao de Interpolacao

```
// Dado 'n' pontos (i, y[i]), i \in [0, n),
// avalia o polinomio de grau n-1 que passa
// por esses pontos em 'x'
// Tudo modular, precisa do mint
//
// O(n)
ee8 mint evaluate_interpolation(int x, vector<mint> y) {
        int n = y.size();
80e
184
        vector < mint > sulf(n+1, 1), fat(n, 1), ifat(n):
6fa
        for (int i = n-1; i >= 0; i--) sulf[i] = sulf[i+1] * (x - i);
29b
        for (int i = 1; i < n; i++) fat[i] = fat[i-1] * i;</pre>
0da
        ifat[n-1] = 1/fat[n-1]:
        for (int i = n-2; i >= 0; i--) ifat[i] = ifat[i+1] * (i + 1);
3db
```

```
ca1
        mint pref = 1, ans = 0;
        for (int i = 0; i < n; pref *= (x - i++)) {
5ea
            mint num = pref * sulf[i+1];
42f
            mint den = ifat[i] * ifat[n-1 - i];
b4e
            if ((n-1 - i)\%2) den *= -1;
0bd
            ans += v[i] * num * den;
03f
ce6
ba7
        return ans;
4fe }
1.3 Berlekamp-Massey
// guess_kth(s, k) chuta o k-esimo (0-based) termo
// de uma recorrencia linear que gera s
// Para uma rec. lin. de ordem x, se passar 2x termos
// vai gerar a certa
// Usar aritmetica modular
// O(n^2 log k), em que n = |s|
        int n = c.size();
ff2
9ee
        assert(c.size() <= s.size());</pre>
d09
564
            vector <T> ret(a.size() + b.size() - 1);
```

999

ce8

x.resize(n);

T ret = 0;

```
b7c template < typename T> T evaluate (vector < T> c, vector < T> s, ll k) {
        auto mul = [\&] (const vector T &a. const vector T &b) {
d75
            for (int i = 0; i < a.size(); i++) for (int j = 0; j <
   b.size(); j++)
                ret[i+j] += a[i] * b[j];
cff
83d
            for (int i = ret.size()-1; i \ge n; i--) for (int j = n-1;
   i >= 0; i--)
112
                ret[i-j-1] += ret[i] * c[j];
16d
            ret.resize(min<int>(ret.size(), n));
            return ret;
edf
3b9
        };
        vector < T > a = n == 1 ? vector < T > ({c[0]}) : vector < T > ({0, 1}),
   x = \{1\}:
        while (k) {
95f
7f1
            if (k\&1) x = mul(x, a):
            a = mul(a, a), k >>= 1;
8ea
```

```
e72
        for (int i = 0; i < n; i++) ret += x[i] * s[i];
edf
        return ret:
7e2 }
192 template < typename T > vector < T > berlekamp_massey(vector < T > s) {
        int n = s.size(), 1 = 0, m = 1;
222
        vector < T > b(n), c(n):
        T 1d = b[0] = c[0] = 1;
46e
620
        for (int i = 0; i < n; i++, m++) {</pre>
793
            T d = s[i]:
ab6
            for (int j = 1; j \le 1; j++) d += c[j] * s[i-j];
5f0
            if (d == 0) continue;
8b4
            vector <T> temp = c:
369
            T coef = d / ld;
ba6
            for (int j = m; j < n; j++) c[j] -= coef * b[j-m];
88f
            if (2 * 1 \le i) 1 = i + 1 - 1, b = temp, 1d = d, m = 0;
76a
        }
90c
        c.resize(1 + 1);
844
        c.erase(c.begin());
        for (T\& x : c) x = -x;
0dc
807
        return c:
4d9 }
2cf template < typename T > T guess_kth(const vector < T > & s, ll k) {
cc3
        auto c = berlekamp_massey(s);
96a
        return evaluate(c, s, k):
697 }
```

1.4 Binomial Distribution

1.5 Convolucao de GCD / LCM

```
// O(n log(n))
// multiple_transform(a)[i] = \sum_d a[d * i]
bbe template < typename T > void multiple_transform (vector < T > & v, bool
   inv = false) {
64a
        vector < int > I(v.size()-1);
847
        iota(I.begin(), I.end(), 1);
        if (inv) reverse(I.begin(), I.end());
674
dad
        for (int i : I) for (int j = 2; i*j < v.size(); j++)</pre>
            v[i] += (inv ? -1 : 1) * v[i*j];
a8a
338 }
// gcd convolution(a, b)[k] = \sum {gcd(i, i) = k} a i * b i
fe2 template < typename T > vector < T > gcd_convolution(vector < T > a,
   vector <T> b) {
bdf
        multiple_transform(a), multiple_transform(b);
799
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
dea
        multiple_transform(a, true);
3f5
        return a;
984 }
// divisor_transform(a)[i] = \sum_{d|i} a[i/d]
be7 template < typename T> void divisor_transform(vector < T>& v, bool inv
   = false) {
        vector < int > I(v.size()-1);
64a
847
        iota(I.begin(), I.end(), 1);
5ea
        if (!inv) reverse(I.begin(), I.end());
        for (int i : I) for (int j = 2; i*j < v.size(); j++)
dad
14f
            v[i*i] += (inv ? -1 : 1) * v[i]:
aa7 }
// lcm_convolution(a, b)[k] = \sum_{lcm(i, j) = k} a_i * b_j
b1b template < typename T > vector < T > lcm_convolution (vector < T > a,
   vector <T> b) {
        divisor_transform(a), divisor_transform(b);
3af
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
799
        divisor_transform(a, true);
d8f
3f5
        return a:
f5a }
1.6 Coprime Basis
// Dado um conjunto de elementos A constroi uma base B
// de fatores coprimos tal que todo elemento A[i]
// pode ser fatorado como A[i] = \prod B[j]^p_ij
```

```
// Sendo n o numero de inserts, a complexidade esperada fica
// O(n*(n*loglog(MAX) + log(MAX)^2))
// No pior caso, podemos trocar n*loglog(MAX) por
// se MAX <= 1e6 fica 8*n
// se MAX <= 1e9 fica 10*n
// se MAX <= 1e18 fica 16*n
// se MAX <= 1e36 fica 26*n
ebc template <typename T> struct coprime_basis {
        vector <T> basis:
        coprime_basis() {}
60e
        coprime_basis(vector<T> v) { for (T i : v) insert(i); }
055
845
        void insert(T z) {
            int n = basis.size():
сЗс
efe
            basis.push_back(z);
43c
            for (int i = n; i < basis.size(); i++) {</pre>
21c
                 for (int j = (i != n) ? i+1 : 0; j < basis.size();</pre>
   j++) {
                     if (i == j) continue;
4ce
024
                     T &x = basis[i];
c91
                     if (x == 1) {
fac
                         j = INF;
5e2
                         continue;
6e0
                     }
544
                     T \& v = basis[i];
3c9
                     T g = gcd(x, y);
                     if (g == 1) continue;
e10
15b
                     v /= g, x /= g;
8c6
                     basis.push_back(g);
                }
069
422
fe8
            basis.erase(remove(basis.begin(), basis.end(), 1),
    basis.end());
        }
1a5
4ba
        vector < int > factor(T x) {
21d
            vector < int > fat(basis.size());
6fd
            for (int i = 0; i < basis.size(); i++) {</pre>
25 c
                 while (x \% basis[i] == 0) x /= basis[i], fat[i]++;
            }
8de
6a7
            return fat;
b5d
        }
671 }:
```

1.7 Crivo de Eratosthenes

```
// "O" crivo
// Encontra maior divisor primo
// Um numero eh primo sse divi[x] == x
// fact fatora um numero <= lim
// A fatoracao sai ordenada
// crivo - 0(n log(log(n)))
// fact - O(log(n))
f12 int divi[MAX];
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
        for (int i = 2; i <= lim; i++) if (divi[i] == 1)
018
            for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
349 }
470 void fact(vector<int>& v, int n) {
       if (n != divi[n]) fact(v, n/divi[n]);
        v.push_back(divi[n]);
ab4
1db }
// Crivo linear
// Mesma coisa que o de cima, mas tambem
// calcula a lista de primos
//
// O(n)
f12 int divi[MAX];
fd3 vector<int> primes;
fb9 void crivo(int lim) {
d5a
        divi[1] = 1;
        for (int i = 2; i <= lim; i++) {</pre>
f70
            if (divi[i] == 0) divi[i] = i, primes.push_back(i);
3eb
            for (int j : primes) {
3ba
522
                if (j > divi[i] or i*j > lim) break;
                divi[i*j] = j;
00b
491
            }
85a
        }
519 }
```

```
// Crivo de divisores
// Encontra numero de divisores
// ou soma dos divisores
// O(n log(n))
f12 int divi[MAX];
fb9 void crivo(int lim) {
f53
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
424
        for (int i = 2: i <= lim: i++)
594
            for (int j = i; j <= lim; j += i) {</pre>
                // para numero de divisores
                 divi[j]++;
9e0
                // para soma dos divisores
278
                 divi[j] += i;
c58
fc1 }
// Crivo de totiente
// Encontra o valor da funcao
// totiente de Euler
// O(n log(log(n)))
5f4 int tot[MAX];
fb9 void crivo(int lim) {
a27
        for (int i = 1; i <= lim; i++) {</pre>
bc9
            tot[i] += i:
feb
            for (int j = 2*i; j <= lim; j += i)</pre>
                 tot[i] -= tot[i];
837
678
        }
212 }
// Crivo de funcao de mobius
//
// O(n log(log(n)))
4e1 char meb[MAX];
fb9 void crivo(int lim) {
649
        for (int i = 2; i <= lim; i++) meb[i] = 2;</pre>
        meb[1] = 1:
ace
```

```
842
        for (int i = 2; i <= lim; i++) if (meb[i] == 2)
            for (int j = i; j <= lim; j += i) if (meb[i]) {</pre>
8d8
686
                if (meb[j] == 2) meb[j] = 1;
ae1
                meb[i] *= i/i\%i ? -1 : 0;
97f
            }
9bc }
// Crivo linear de funcao multiplicativa
// Computa f(i) para todo 1 <= i <= n, sendo f
// uma funcao multiplicativa (se gcd(a,b) = 1,
// entao f(a*b) = f(a)*f(b)
// f prime tem que computar f de um primo. e
// add_prime tem que computar f(p^{(k+1)}) dado f(p^k) e p
// Se quiser computar f(p^k) dado p e k, usar os comentarios
// O(n)
fd3 vector <int> primes;
623 int f[MAX], pot[MAX];
//int expo[MAX];
5c4 void sieve(int lim) {
        // Funcoes para soma dos divisores:
fc9
        auto f_prime = [](int p) { return p+1; };
31c
        auto add_prime = [](int fpak, int p) { return fpak*p+1; };
        //auto f_pak = [](int p, int k) {};
02d
        f[1] = 1;
f70
        for (int i = 2; i <= lim; i++) {</pre>
e6b
            if (!pot[i]) {
e74
                primes.push_back(i);
                f[i] = f_prime(i), pot[i] = i;
                //\expo[i] = 1;
b71
3b9
            for (int p : primes) {
                if (i*p > lim) break;
b9f
569
                if (i\%p == 0) {
b97
                    f[i*p] = f[i / pot[i]] * add_prime(f[pot[i]], p);
                    // se for descomentar, tirar a linha de cima também
                    //f[i*p] = f[i / pot[i]] * f_pak(p, expo[i]+1);
                    //\expo[i*p] = \expo[i]+1;
                    pot[i*p] = pot[i] * p;
51f
c2b
                    break;
643
                } else {
                    f[i*p] = f[i] * f[p];
9ef
                    pot[i*p] = p;
638
```

```
//expo[i*p] = 1;
6f7 }
f31 }
1bb }
350 }
```

1.8 Deteccao de ciclo - Tortoise and Hare

```
// Linear no tanto que tem que andar pra ciclar,
// O(1) de memoria
// Retorna um par com o tanto que tem que andar
// do f0 ate o inicio do ciclo e o tam do ciclo
58d pair<11, 11> find_cycle() {
273
        11 \text{ tort} = f(f0);
b2b
        ll hare = f(f(f0));
b1b
        11 t = 0;
683
        while (tort != hare) {
            tort = f(tort);
b4d
4b2
            hare = f(f(hare));
c82
            t++;
93d
        }
        11 st = 0;
0e8
909
        tort = f0:
683
        while (tort != hare) {
b4d
            tort = f(tort):
1a2
            hare = f(hare);
397
            st++;
c91
        }
73d
        11 len = 1;
3cd
        hare = f(tort):
683
        while (tort != hare) {
1a2
            hare = f(hare):
040
            len++;
f1a
ebd
        return {st, len};
899 }
```

1.9 Division Trick

```
5bf }
```

1.10 Equação Diofantina Linear

```
// Encontra o numero de solucoes de a*x + b*y = c,
// em que x \in [lx. rx] e v \in [lv. rv]
// Usar o comentario para recuperar as solucoes
// (note que o b ao final eh b/gcd(a, b))
// Cuidado com overflow! Tem que caber o quadrado dos valores
//
// O(log(min(a, b)))
c5e template < typename T > tuple < ll, T, T > ext_gcd(ll a, ll b) {
        if (!a) return {b, 0, 1};
        auto [g, x, y] = ext_gcd < T > (b%a, a);
c4b
        return \{g, y - b/a*x, x\};
c59
8a8 }
// numero de solucoes de a*[lx, rx] + b*[ly, ry] = c
14c template<typename T = 11 > // usar __int128 se for ate 1e18
2a4 ll diophantine(ll a, ll b, ll c, ll lx, ll rx, ll ly, ll ry) {
        if (lx > rx or ly > ry) return 0;
a98
       if (a == 0 \text{ and } b == 0) \text{ return } c ? 0 : (rx-lx+1)*(ry-ly+1);
        auto [g, x, y] = ext_gcd < T > (abs(a), abs(b));
8ce
9c3
       if (c % g != 0) return 0;
       if (a == 0) return (rx-lx+1)*(ly <= c/b and c/b <= ry);
249
       if (b == 0) return (rv-lv+1)*(lx <= c/a and c/a <= rx):
4ce
        x *= a/abs(a) * c/g, y *= b/abs(b) * c/g, a /= g, b /= g;
fb1
        auto shift = [\&](T qt) \{ x += qt*b, y -= qt*a; \};
b20
efa
        auto test = [&](T& k, ll mi, ll ma, ll coef, int t) {
866
            shift((mi - k)*t / coef):
79d
            if (k < mi) shift(coef > 0 ? t : -t);
74d
           if (k > ma) return pair <T, T > (rx+2, rx+1);
41f
           T x1 = x:
633
            shift((ma - k)*t / coef);
            if (k > ma) shift(coef > 0 ? -t : t);
c5b
4a9
            return pair<T, T>(x1, x);
8e1
       };
        auto [11, r1] = test(x, lx, rx, b, 1);
639
38e
        auto [12, r2] = test(y, 1y, ry, a, -1);
c43
        if (12 > r2) swap(12, r2);
       T l = max(11, 12), r = min(r1, r2);
50a
339
        if (1 > r) return 0:
42f
       11 k = (r-1) / abs(b) + 1;
839
        return k; // solucoes: x = 1 + [0, k)*|b|
```

```
98e }
```

1.11 Euclides estendido

```
// Acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
// Assume a, b >= 0
//
// O(log(min(a, b)))

2be tuple<11, 11, 11> ext_gcd(11 a, 11 b) {
3bd     if (!a) return {b, 0, 1};
550     auto [g, x, y] = ext_gcd(b%a, a);
c59     return {g, y - b/a*x, x};
354 }
```

1.12 Exponenciacao rapida

```
// (x^y mod m) em O(log(y))
03c ll pow(ll x, ll y, ll m) \{ // \text{ iterativo} \}
c85
        ll ret = 1:
1 b8
         while (y) {
895
             if (y & 1) ret = (ret * x) % m;
23b
             y >>= 1;
сс5
             x = (x * x) % m;
020
edf
         return ret;
12b }
03c ll pow(ll x, ll y, ll m) \{ // \text{ recursivo} \}
13a
        if (!y) return 1;
426
        11 ans = pow(x*x\%m, y/2, m);
88d
        return y%2 ? x*ans%m : ans;
7d4 }
```

1.13 Fast Walsh Hadamard Transform

```
// FWHT<'!'>(f) eh SOS DP
// FWHT<'&'>(f) eh soma de superset DP
// Se chamar com ^, usar tamanho potencia de 2!!
//
// O(n log(n))

382 template<char op, class T> vector<T> FWHT(vector<T> f, bool inv = false) {
b75    int n = f.size();
```

```
d78
        for (int k = 0; (n-1) >> k; k++) for (int i = 0; i < n; i++) if
                                                                             078
                                                                                     for (int i = 1; i < n/2; i++) roots[i] = roots[i-1]*r;</pre>
                                                                             63f }
   (i >> k&1) {
            int j = i^(1 << k);
29e
627
            if (op == '\^') f[i] += f[i], f[i] = f[i] - 2*f[i];
                                                                             8a2 template < typename T > void fft (vector < T > & a, bool f, int N,
a38
            if (op == ', ') f[i] += (inv ? -1 : 1) * f[j];
                                                                                vector < int > % rev) {
93c
            if (op == '&') f[j] += (inv ? -1 : 1) * f[i];
                                                                                     for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],</pre>
                                                                             bc7
1bb
        }
                                                                                a[rev[i]]):
578
        if (op == ', and inv) for (auto& i : f) i /= n;
                                                                             12b
                                                                                     int 1, r, m;
abe
        return f;
                                                                             cb4
                                                                                     vector <T> roots(N);
50e }
                                                                                     for (int n = 2; n <= N; n *= 2) {</pre>
                                                                             192
                                                                             0f4
                                                                                         get_roots(f, n, roots);
                                                                             5dc
                                                                                         for (int pos = 0; pos < N; pos += n) {
1.14 FFT
                                                                             432
                                                                                             1 = pos + 0, r = pos + n/2, m = 0;
                                                                             a88
                                                                                             while (m < n/2) {
// Chamar convolution com vector < complex < double >> para FFT
                                                                             297
                                                                                                 auto t = roots[m] * a[r];
// Precisa do mint para NTT
                                                                                                 a[r] = a[1] - t;
                                                                             254
//
                                                                             b8f
                                                                                                 a[1] = a[1] + t;
// O(n log(n))
                                                                             2c9
                                                                                                 1++, r++, m++;
                                                                             d89
// Para FFT
                                                                                         }
                                                                             1fd
488 void get_roots(bool f, int n, vector<complex<double>>& roots) {
                                                                             185
                                                                                     }
f26
        const static double PI = acosl(-1);
                                                                             235
                                                                                     if (f) {
        for (int i = 0; i < n/2; i++) {</pre>
71a
                                                                             1c5
                                                                                         auto invN = T(1) / T(N);
b1e
            double alpha = i*((2*PI)/n);
                                                                             557
                                                                                         for (int i = 0; i < N; i++) a[i] = a[i] * invN;</pre>
1a1
            if (f) alpha = -alpha;
                                                                             256
                                                                                     }
069
            roots[i] = {cos(alpha), sin(alpha)};
                                                                             1b1 }
804
        }
de5 }
                                                                             bf5 template < typename T > vector < T > convolution (vector < T > & a,
                                                                                vector<T>& b) {
// Para NTT
                                                                             87a
                                                                                     vector<T> l(a.begin(), a.end()), r(b.begin(), b.end());
9f7 template <int p>
                                                                                     int N = 1.size()+r.size()-1;
                                                                             e0a
97b void get_roots(bool f, int n, vector<mod_int<p>>& roots) {
                                                                             f03
                                                                                     int n = 1, log_n = 0;
1e6
        mod_int  r;
                                                                             0a4
                                                                                     while (n \le N) n *= 2, log_n++;
de9
        int ord;
                                                                             808
                                                                                     vector < int > rev(n):
57a
        if (p == 998244353) {
                                                                             603
                                                                                     for (int i = 0; i < n; i++) {</pre>
9b6
            r = 102292;
                                                                             434
                                                                                         rev[i] = 0;
81b
            ord = (1 << 23);
                                                                             f44
                                                                                         for (int j = 0; j < log_n; j++) if (i>>j&1)
121
        } else if (p == 754974721) {
                                                                             4ff
                                                                                             rev[i] = 1 << (log_n-1-i);
43a
            r = 739831874;
                                                                             256
f0a
            ord = (1 << 24);
                                                                             143
                                                                                     assert(N <= n);</pre>
d48
        } else if (p == 167772161) {
                                                                             fa4
                                                                                     l.resize(n):
a2a
            r = 243;
                                                                             7 e 4
                                                                                     r.resize(n);
033
            ord = (1 << 25):
                                                                                     fft(l, false, n, rev);
                                                                             56e
        } else assert(false);
5a4
                                                                             fcf
                                                                                     fft(r, false, n, rev);
                                                                             917
                                                                                     for (int i = 0; i < n; i++) l[i] *= r[i];
547
        if (f) r = r^(p - 1 - ord/n);
                                                                             88b
                                                                                     fft(1, true, n, rev);
ee2
        else r = r^(ord/n);
                                                                                     l.resize(N):
                                                                             5e1
be4
        roots[0] = 1;
```

```
792
        return 1;
bd6 }
// NTT
6c8 template < int p, typename T > vector < mod_int < p >> ntt (vector < T > & a,
    vector < T > & b) {
d52
        vector < mod_int < p >> A(a.begin(), a.end()), B(b.begin(),
    b.end()):
d29
        return convolution(A, B);
3bf }
// Convolucao de inteiro
// Precisa do CRT
// Tabela de valores:
// [0,1]
              - <int, 1>
// [-1e5, 1e5] - <11, 2>
// [-1e9, 1e9] - <__int128, 3>
b3c template < typename T, int mods >
eec vector<T> int_convolution(vector<int>& a, vector<int>& b) {
        static const int M1 = 998244353, M2 = 754974721, M3 =
    167772161:
bf5
        auto c1 = ntt < M1 > (a, b);
221
        auto c2 = (mods \ge 2 ? ntt \le M2 \ge (a, b) : vector \le mod_int \le M2 \ge ());
        auto c3 = (mods >= 3 ? ntt < M3 > (a, b) : vector < mod_int < M3 >> ());
f9b
2da
        vector <T> ans:
5c5
        for (int i = 0; i < c1.size(); i++) {</pre>
c09
             crt < T > at (c1[i].v, M1);
316
             if (mods \ge 2) at = at * crt<T>(c2[i].v, M2);
             if (mods >= 3) at = at * crt<T>(c3[i].v, M3);
987
b2b
             ans.push back(at.a):
             if (at.a > at.m/2) ans.back() -= at.m:
26d
b9f
        return ans;
ba7
5e8 }
1.15 Gauss
// Resolve sistema linear
// Retornar um par com o numero de solucoes
// e alguma solucao, caso exista
//
// O(n^2 * m)
```

67a template < typename T> 728 pair<int, vector<T>> gauss(vector<vector<T>> a, vector<T> b) { const double eps = 1e-6; f92 int n = a.size(), m = a[0].size(); for (int i = 0; i < n; i++) a[i].push_back(b[i]);</pre> 2f0 3cb vector<int> where(m. -1); 237 for (int col = 0, row = 0; col < m and row < n; col++) { f05 int sel = row; for (int i=row; i<n; ++i)</pre> b95 e55 if (abs(a[i][col]) > abs(a[sel][col])) sel = i; 2c4 if (abs(a[sel][col]) < eps) continue;</pre> 1ae for (int i = col: i <= m: i++)</pre> dd2 swap(a[sel][i], a[row][i]); 2c3 where [col] = row; 0 c 0 for (int i = 0; i < n; i++) if (i != row) {</pre> T c = a[i][col] / a[row][col]: 96c d5c for (int j = col; j <= m; j++)</pre> c8f a[i][j] -= a[row][j] * c; 490 } b70 row++; } 3d8 b1d vector <T> ans(m, 0); e1a for (int i = 0; i < m; i++) if (where[i] != -1)</pre> 12a ans[i] = a[where[i]][m] / a[where[i]][i]; 603 for (int i = 0; i < n; i++) { 501 T sum = 0;a75 for (int j = 0; j < m; j++)</pre> 5a9 sum += ans[j] * a[i][j]; b1f if (abs(sum - a[i][m]) > eps)6cd return pair(0, vector<T>()); ec9 } for (int i = 0; i < m; i++) if (where[i] == -1)</pre> 12e 018 return pair(INF, ans); 280 return pair(1, ans); 292 } 1.16 Gauss - **Z**2 // D eh dimensao do espaco vetorial // add(v) - adiciona o vetor v na base (retorna se ele jah pertencia ao span da base) // coord(v) - retorna as coordenadas (c) de v na base atual (basis^T.c

= v)

```
// recover(v) - retorna as coordenadas de v nos vetores na ordem em
   que foram inseridos
// coord(v).first e recover(v).first - se v pertence ao span
// Complexidade:
// add, coord, recover: O(D^2 / 64)
2a3 template <int D> struct Gauss_z2 {
3c1
        bitset <D> basis[D], keep[D];
b16
        int rk. in:
482
        vector < int > id;
37f
        Gauss_z2 () : rk(0), in(-1), id(D, -1) {};
04e
        bool add(bitset < D > v) {
42c
            in++:
fb0
            bitset <D> k;
            for (int i = D - 1; i \ge 0; i--) if (v[i]) {
659
                if (basis[i][i]) v ^= basis[i], k ^= keep[i];
189
4e6
                else {
ea6
                    k[i] = true, id[i] = in, keep[i] = k;
                    basis[i] = v, rk++;
6ce
8a6
                    return true:
b34
                }
09c
            }
d1f
            return false;
58b
0f6
        pair < bool, bitset < D >> coord(bitset < D > v) {
944
            bitset <D> c:
659
            for (int i = D - 1; i \ge 0; i--) if (v[i]) {
                if (basis[i][i]) v ^= basis[i], c[i] = true;
a39
8af
                else return {false, bitset <D>()};
            }
a08
5db
            return {true, c};
a08
330
        pair < bool, vector < int >> recover(bitset < D > v) {
22e
            auto [span, bc] = coord(v);
af8
            if (not span) return {false, {}};
f79
            bitset <D> aux:
5a0
            for (int i = D - 1; i >= 0; i--) if (bc[i]) aux ^= keep[i];
ea9
            vector<int> oc:
            for (int i = D - 1; i >= 0; i--) if (aux[i])
   oc.push_back(id[i]);
001
            return {true, oc};
b75
        }
d0a }:
```

1.17 Integracao Numerica

if (n <= 64) {

return;

int mid = n/2;

}

r[i+j] += a[i] * b[j];

memset(E, 0, sizeof(E[0])*n);

for (int i = 0; i < mid; i++) {</pre>

T * atmp = tmp, *btmp = tmp+mid, *E = tmp+n;

```
// Metodo de Simpson 3/8
// Integra f no intervalo [a, b], erro cresce proporcional a (b - a)^5
676 const int N = 3*100; // multiplo de 3
287 ld integrate(ld a, ld b, function < ld(ld) > f) {
        ld s = 0, h = (b - a)/N;
b4d
067
        for (int i = 1; i < N; i++) s += f(a + i*h)*(i%3 ? 3 : 2);
0da
        return (f(a) + s + f(b))*3*h/8;
c7e }
1.18 Inverso Modular
// Computa o inverso de a modulo b
// Se b eh primo, basta fazer
// a^{(b-2)}
f0a ll inv(ll a, ll b) {
        return a > 1 ? b - inv(b\%a, a)*b/a : 1;
ae1
cf9 }
// computa o inverso modular de 1..MAX-1 modulo um primo
a88 ll inv[MAX]:
Of 2 \operatorname{inv}[1] = 1:
Ofa for (int i = 2; i < MAX; i++) inv[i] = MOD - MOD/i*inv[MOD%i]%MOD;
1.19 Karatsuba
// Os pragmas podem ajudar
// Para n \sim 2e5, roda em < 1 s
//
// O(n^1.58)
//#pragma GCC optimize("Ofast")
//#pragma GCC target ("avx,avx2")
77a template < typename T > void kar(T* a, T* b, int n, T* r, T* tmp) {
```

for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)

d4c

510

212

505

bb8

194

2d7

4f1

c65

```
c72
             atmp[i] = a[i] + a[i+mid];
             btmp[i] = b[i] + b[i+mid];
4b9
a3f
38a
        kar(atmp, btmp, mid, E, tmp+2*n);
        kar(a, b, mid, r, tmp+2*n);
b1e
        kar(a+mid, b+mid, mid, r+n, tmp+2*n);
229
c65
        for (int i = 0: i < mid: i++) {</pre>
735
            T \text{ temp} = r[i+mid];
de7
            r[i+mid] += E[i] - r[i] - r[i+2*mid];
            r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
f1e
f72
        }
28f }
e38 template < typename T > vector < T > karatsuba (vector < T > a, vector < T > b)
   {
        int n = max(a.size(), b.size());
ba3
a84
        while (n&(n-1)) n++;
        a.resize(n), b.resize(n);
ca9
        vector\langle T \rangle ret(2*n), tmp(4*n);
ae0
644
        kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
edf
        return ret:
f87 }
1.20 Logaritmo Discreto
```

```
// Resolve logaritmo discreto com o algoritmo baby step giant step
// Encontra o menor x tal que a^x = b (mod m)
// Se nao tem, retorna -1
//
// O(sqrt(m) * log(sqrt(m))
d41
da8 int dlog(int b, int a, int m) {
        if (a == 0) return b ? -1 : 1; // caso nao definido
9f8
d41
        a \%= m, b \%= m;
       int k = 1, shift = 0;
a10
        while (1) {
31e
6e3
           int g = gcd(a, m);
d47
           if (g == 1) break;
d41
9bc
           if (b == k) return shift;
642
           if (b % g) return -1;
            b \neq g, m \neq g, shift++;
9ab
            k = (11) k * a / g % m;
515
d41
af7
        int sq = sqrt(m)+1, giant = 1;
```

```
975
        for (int i = 0; i < sq; i++) giant = (11) giant * a % m;
d41
0b5
        vector < pair < int , int >> baby;
33f
        for (int i = 0, cur = b; i <= sq; i++) {
496
            babv.emplace back(cur. i):
            cur = (11) cur * a % m;
16c
622
        }
        sort(baby.begin(), baby.end());
eb4
d41
        for (int j = 1, cur = k; j <= sq; j++) {</pre>
9 c 9
ace
            cur = (11) cur * giant % m;
78ъ
            auto it = lower_bound(baby.begin(), baby.end(), pair(cur,
INF)):
d26
            if (it != baby.begin() and (--it)->first == cur)
ac3
                return sq * j - it->second + shift;
b9d
        }
d41
daa
        return -1;
739 }
1.21 Miller-Rabin
```

```
// Testa se n eh primo, n \leq 3 * 10^18
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
d8b ll mul(ll a, ll b, ll m) {
        ll ret = a*b - ll((long double)1/m*a*b+0.5)*m;
e7a
074
        return ret < 0 ? ret+m : ret:</pre>
2f3 }
03c ll pow(ll x, ll y, ll m) {
13a
        if (!v) return 1;
        ll ans = pow(mul(x, x, m), y/2, m);
dbc
        return y%2 ? mul(x, ans, m) : ans;
7fa
539 }
1a2 bool prime(ll n) {
        if (n < 2) return 0:
1aa
237
        if (n <= 3) return 1;</pre>
9de
        if (n \% 2 == 0) return 0:
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
        // com esses primos, o teste funciona garantido para n <= 2^64
        // funciona para n <= 3*10^24 com os primos ate 41
771
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
```

```
795265022}) {
da0
             11 x = pow(a, d, n);
             if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue;
709
4a2
            for (int j = 0; j < r - 1; j++) {
                 x = mul(x, x, n);
10f
df0
                 if (x == n - 1) break:
1ff
e1b
             if (x != n - 1) return 0;
e74
6a5
        return 1;
5ba }
```

1.22 Pollard's Rho Alg

```
// Usa o algoritmo de deteccao de ciclo de Floyd
// com uma otimizacao na qual o gcd eh acumulado
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
d8b ll mul(ll a, ll b, ll m) {
e7a
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
        return ret < 0 ? ret+m : ret;</pre>
074
2f3 }
03c ll pow(ll x, ll y, ll m) {
13a
        if (!y) return 1;
dbc
        ll ans = pow(mul(x, x, m), y/2, m);
7fa
        return y%2 ? mul(x, ans, m) : ans;
539 }
1a2 bool prime(ll n) {
       if (n < 2) return 0;</pre>
1aa
237
       if (n <= 3) return 1;
       if (n % 2 == 0) return 0;
9de
f6a
       ll r = \_builtin\_ctzll(n - 1), d = n >> r;
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
   795265022}) {
            ll x = pow(a, d, n);
da0
709
           if (x == 1 or x == n - 1 or a % n == 0) continue;
```

```
4a2
            for (int j = 0; j < r - 1; j++) {
                x = mul(x, x, n);
10f
df0
                if (x == n - 1) break;
1ff
e1b
            if (x != n - 1) return 0;
e74
6a5
        return 1;
5ba }
9cf ll rho(ll n) {
0f9
        if (n == 1 or prime(n)) return n;
f7c
        auto f = [n](11 x) \{return mul(x, x, n) + 1;\};
8a5
        11 x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
533
        while (t \% 40 != 0 or gcd(prd, n) == 1) {
            if (x==y) x = ++x0, y = f(x);
8a0
e13
            q = mul(prd, abs(x-y), n);
21f
            if (q != 0) prd = q;
            x = f(x), y = f(f(y)), t++;
450
379
        }
002
        return gcd(prd, n);
523 }
5b7 vector<ll> fact(ll n) {
1 b 9
        if (n == 1) return {};
0ec
        if (prime(n)) return {n};
0ed
        11 d = rho(n);
1de
        vector < 11 > 1 = fact(d), r = fact(n / d);
        l.insert(l.end(), r.begin(), r.end());
3af
792
        return 1;
902 }
1.23 Produto de dois long long mod m
// 0(1)
d8b ll mul(ll a, ll b, ll m) { // a*b % m
        ll ret = a*b - ll((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
1.24 Simplex
// Maximiza c^T x s.t. Ax <= b, x >= 0
//
// O(2^n), porem executa em O(n^3) no caso medio
```

```
395 const double eps = 1e-7;
493 namespace Simplex {
        vector < vector < double >> T;
14e
        int n. m:
43e
        vector < int > X, Y;
c51
        void pivot(int x, int y) {
8e6
            swap(X[y], Y[x-1]);
            for (int i = 0; i <= m; i++) if (i != y) T[x][i] /=
d03
   T[x][v];
            T[x][y] = 1/T[x][y];
33c
38b
            for (int i = 0; i \le n; i++) if (i != x \text{ and } abs(T[i][y]) >
   eps) {
774
                 for (int j = 0; j <= m; j++) if (j != y) T[i][j] -=
   T[i][y] * T[x][j];
                T[i][y] = -T[i][y] * T[x][y];
a7d
        }
e05
        // Retorna o par (valor maximo, vetor solucao)
        pair < double , vector < double >> simplex(
6f8
                 vector < vector < double >> A, vector < double >> b,
e9d
   vector < double > c) {
            n = b.size(), m = c.size();
5bb
002
            T = vector(n + 1, vector < double > (m + 1));
2d9
            X = vector < int > (m);
0c2
            Y = vector < int > (n);
            for (int i = 0; i < m; i++) X[i] = i;</pre>
115
51f
            for (int i = 0; i < n; i++) Y[i] = i+m;</pre>
            for (int i = 0; i < m; i++) T[0][i] = -c[i];
5 b 5
603
            for (int i = 0; i < n; i++) {</pre>
                 for (int j = 0; j < m; j++) T[i+1][j] = A[i][j];
ba6
                 T[i+1][m] = b[i]:
eca
07 c
667
            while (true) {
714
                 int x = -1, y = -1;
2db
                 double mn = -eps;
                 for (int i = 1; i <= n; i++) if (T[i][m] < mn) mn =
c29
   T[i][m], x = i;
                if (x < 0) break;
af2
                for (int i = 0; i < m; i++) if (T[x][i] < -eps) { y = }
   i; break; }
                 if (y < 0) return {-1e18, {}}; // sem solucao para Ax</pre>
4a6
   \leq b
7fb
                 pivot(x, y);
```

```
472
            }
667
            while (true) {
714
                int x = -1, y = -1;
2db
                double mn = -eps;
                for (int i = 0; i < m; i++) if (T[0][i] < mn) mn =</pre>
562
   T[0][i], y = i;
9ъ0
                if (y < 0) break;
                mn = 1e200;
034
                for (int i = 1; i \le n; i++) if (T[i][y] > eps and
5af
   T[i][m] / T[i][y] < mn
48f
                    mn = T[i][m] / T[i][y], x = i;
53b
                if (x < 0) return {1e18, {}}; // c^T x eh ilimitado
7fb
                pivot(x, y);
            }
81e
290
            vector < double > r(m);
32f
            for(int i = 0; i < n; i++) if (Y[i] < m) r[Y[i]] =
   T[i+1][m];
e59
            return {T[0][m], r};
        }
7a4
a64 }
     Teorema Chines do Resto
// Combina equacoes modulares lineares: x = a (mod m)
// O m final eh o lcm dos m's, e a resposta eh unica mod o lcm
```

```
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
3bd
        if (!a) return {b, 0, 1};
550
        auto [g, x, y] = ext_gcd(b%a, a);
c59
        return \{g, y - b/a*x, x\};
537 }
bfe template < typename T = 11 > struct crt {
        Ta, m;
5f3
        crt() : a(0), m(1) {}
        crt(T a_, T m_) : a(a_), m(m_) {}
7eb
        crt operator * (crt C) {
911
238
            auto [g, x, y] = ext\_gcd(m, C.m);
dc0
            if ((a - C.a) \% g) a = -1;
4f9
            if (a == -1 or C.a == -1) return crt(-1, 0);
d09
            T lcm = m/g*C.m;
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
```

```
1f2 }
0d9 };

1.26 Totiente
```

```
// O(sqrt(n))
a7e int tot(int n){
0f6    int ret = n;

505    for (int i = 2; i*i <= n; i++) if (n % i == 0) {
        while (n % i == 0) n /= i;
125         ret -= ret / i;
34a    }
af4    if (n > 1) ret -= ret / n;

edf    return ret;
fae }
```

2 Grafos

2.1 AGM Direcionada

```
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
//
// O(m log(n))
3c9 struct node {
f31
        pair<ll, int> val;
4e4
       ll lazy;
b19
       node *1, *r;
       node() {}
f93
c53
        node(pair<int, int> v) : val(v), lazy(0), l(NULL), r(NULL) {}
a9c
        void prop() {
768
            val.first += lazy;
b87
            if (1) 1->lazy += lazy;
d3b
            if (r) r->lazy += lazy;
c60
            lazv = 0:
05Ъ
296 }:
de5 void merge(node*& a, node* b) {
c11
        if (!a) swap(a, b);
```

```
802
        if (!b) return;
626
        a->prop(), b->prop();
        if (a->val > b->val) swap(a, b);
d04
4b0
        merge(rand()%2 ? a->1 : a->r, b);
b82 }
d01 pair<ll, int> pop(node*& R) {
e8f
        R->prop();
22e
        auto ret = R->val;
        node* tmp = R;
af0
        merge(R->1, R->r);
3f3
6c9
        R = R - > 1;
        if (R) R->lazy -= ret.first;
3e4
7c3
        delete tmp;
edf
        return ret;
c4e }
6f6 void apaga(node* R) { if (R) apaga(R->1), apaga(R->r), delete R; }
f13 ll dmst(int n, int r, vector<pair<int, int>, int>>& ar) {
        vector<int> p(n); iota(p.begin(), p.end(), 0);
94e
        function < int(int) > find = [&](int k) { return
   p[k] == k?k:p[k] = find(p[k]); };
2d7
        vector < node *> h(n);
56f
        for (auto e : ar) merge(h[e.first.second], new node({e.second,
   e.first.first}));
fd1
        vector < int > pai(n, -1), path(n);
        pai[r] = r;
66e
04b
        11 \text{ ans} = 0;
603
        for (int i = 0; i < n; i++) { // vai conectando todo mundo
2a3
            int u = i, at = 0;
            while (pai[u] == -1) {
cae
daa
                if (!h[u]) { // nao tem
947
                    for (auto i : h) apaga(i);
77c
                    return LINF:
dd1
167
                path[at++] = u, pai[u] = i;
                auto [mi, v] = pop(h[u]);
55e
64c
                ans += mi;
5e2
                if (pai[u = find(v)] == i) { // ciclo
86f
                    while (find(v = path[--at]) != u)
621
                         merge(h[u], h[v]), h[v] = NULL, p[find(v)] = u;
57a
                    pai[u] = -1;
                }
850
            }
ce8
5df
        }
947
        for (auto i : h) apaga(i);
```

```
ba7     return ans;
e02 }
```

2.2 Articulation Points

```
// Computa os pontos de articulação (vertices criticos) de um grafo
// art[i] armazena o numero de novas componentes criadas ao deletar
   vertice i
// se art[i] >= 1, entao vertice i eh ponto de articulacao
// O(n+m)
1a8 int n;
789 vector < vector < int >> g;
4ce stack < int > s:
b66 vector < int > id, art;
3e1 int dfs_art(int i, int& t, int p = -1) {
cf0
        int lo = id[i] = t++;
18e
        s.push(i);
        for (int j : g[i]) if (j != p) {
cac
            if (id[j] == -1) {
9a3
206
               int val = dfs_art(j, t, i);
               lo = min(lo, val);
0 c 3
                if (val >= id[i]) {
588
66a
                    art[i]++;
bd9
                    while (s.top() != j) s.pop();
2eb
                    s.pop();
1f3
                // if (val > id[i]) aresta i-j eh ponte
238
328
            else lo = min(lo, id[j]);
762
3bd
        if (p == -1 and art[i]) art[i]--;
253
        return lo;
8e1 }
d79 void compute_art_points() {
597
        id = vector < int > (n, -1);
       art = vector < int > (n, 0);
6bb
       int t = 0;
d41
        for (int i = 0; i < n; i++) if (id[i] == -1)
625
            dfs_art(i, t, -1);
379 }
```

2.3 Bellman-Ford

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
14e int n, m;
248 int d[MAX]:
e93 vector <pair <int, int>> ar; // vetor de arestas
                               // peso das arestas
9e2 vector<int> w;
6be bool bellman ford(int a) {
8ec
        for (int i = 0; i < n; i++) d[i] = INF;</pre>
8a8
        d[a] = 0:
4e3
        for (int i = 0; i <= n; i++)
891
            for (int j = 0; j < m; j++) {</pre>
6e4
                if (d[ar[j].second] > d[ar[j].first] + w[j]) {
                    if (i == n) return 1;
705
e93
                    d[ar[j].second] = d[ar[j].first] + w[j];
84b
                }
a82
            }
bb3
        return 0;
6eb }
2.4 Block-Cut Tree
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulação
// Blocos sao componentes 2-vertice-conexos maximais
// Uma 2-coloração da arvore eh tal que uma cor são
// os blocos, e a outra cor sao os pontos de art.
// Funciona para grafo nao conexo
// art[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se art[i] >= 1, i eh ponto de articulação
// Para todo i <= blocks.size()</pre>
```

// blocks[i] eh uma componente 2-vertce-conexa maximal

```
// edgblocks[i] sao as arestas do bloco i
// tree[i] eh um vertice da arvore que corresponde ao bloco i
// pos[i] responde a qual vertice da arvore vertice i pertence
// Arvore tem no maximo 2n vertices
//
// O(n+m)
d10 struct block_cut_tree {
        vector < vector < int >> g, blocks, tree;
43b
        vector < vector < pair < int , int >>> edgblocks;
4ce
        stack<int> s:
6c0
        stack<pair<int, int>> s2;
2bb
        vector < int > id, art, pos;
763
        block_cut_tree(vector<vector<int>> g_) : g(g_) {
af1
            int n = g.size();
37a
            id.resize(n, -1), art.resize(n), pos.resize(n);
6f2
            build();
       }
6bd
        int dfs(int i, int& t, int p = -1) {
df6
            int lo = id[i] = t++;
cf0
            s.push(i);
18e
827
            if (p != -1) s2.emplace(i, p);
            for (int j : g[i]) if (j != p and id[j] != -1)
53f
   s2.emplace(i, j);
            for (int j : g[i]) if (j != p) {
cac
                if (id[j] == -1) {
9a3
121
                    int val = dfs(j, t, i);
                    lo = min(lo, val);
                    if (val >= id[i]) {
588
66a
                        art[i]++:
                        blocks.emplace_back(1, i);
483
                        while (blocks.back().back() != j)
110
138
                             blocks.back().push_back(s.top()), s.pop();
128
                        edgblocks.emplace_back(1, s2.top()), s2.pop();
                         while (edgblocks.back().back() != pair(j, i))
47e
                             edgblocks.back().push_back(s2.top()),
bce
   s2.pop();
870
                    // if (val > id[i]) aresta i-j eh ponte
85 c
```

```
328
                 else lo = min(lo, id[j]);
            }
344
3bd
            if (p == -1 and art[i]) art[i]--;
253
            return lo:
726
        }
0a8
        void build() {
6bb
            int t = 0;
            for (int i = 0; i < g.size(); i++) if (id[i] == -1) dfs(i,</pre>
abf
   t, -1);
56c
            tree.resize(blocks.size()):
            for (int i = 0; i < g.size(); i++) if (art[i])</pre>
f7d
                 pos[i] = tree.size(), tree.emplace_back();
965
973
            for (int i = 0; i < blocks.size(); i++) for (int j :</pre>
   blocks[i]) {
403
                 if (!art[i]) pos[i] = i;
                 else tree[i].push_back(pos[j]),
101
   tree[pos[j]].push_back(i);
3df
        }
c03
056 };
2.5 Blossom
// Matching maximo em grafo geral
//
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
042 vector < int > g[MAX];
128 int match[MAX]; // match[i] = com quem i esta matchzado ou -1
1f1 int n, pai[MAX], base[MAX], vis[MAX];
26a queue < int > q;
107 void contract(int u, int v, bool first = 1) {
        static vector < bool > bloss:
165
fbe
        static int 1;
418
        if (first) {
a47
            bloss = vector < bool > (n, 0);
042
            vector < bool > teve(n, 0);
ddf
            int k = u; l = v;
31e
            while (1) {
297
                 teve[k = base[k]] = 1;
```

```
if (match[k] == -1) break;
116
dfa
                k = pai[match[k]];
68b
d31
            while (!teve[l = base[l]]) l = pai[match[l]];
5d6
2e9
        while (base[u] != 1) {
            bloss[base[u]] = bloss[base[match[u]]] = 1:
e29
8fa
            pai[u] = v;
0b0
            v = match[u];
            u = pai[match[u]];
a51
58e
71c
        if (!first) return;
95e
        contract(v. u. 0):
        for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
6ee
594
            base[i] = 1;
            if (!vis[i]) q.push(i);
ca7
29a
            vis[i] = 1;
857
       }
e35 }
f10 int getpath(int s) {
88f
        for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1, vis[i] =
   0:
ded
        vis[s] = 1; q = queue < int > (); q.push(s);
402
        while (q.size()) {
be1
            int u = q.front(); q.pop();
            for (int i : g[u]) {
bdc
                if (base[i] == base[u] or match[u] == i) continue;
7a2
                if (i == s or (match[i] != -1 and pai[match[i]] != -1))
e35
4f2
                    contract(u, i);
                else if (pai[i] == -1) {
e2e
545
                    pai[i] = u;
                    if (match[i] == -1) return i;
f6a
818
                    i = match[i]:
                    vis[i] = 1; q.push(i);
29d
                }
90e
0b5
            }
634
daa
        return -1;
a16 }
83f int blossom() {
1a4
        int ans = 0:
        memset(match, -1, sizeof(match));
315
2e3
        for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
f76
            for (int j : g[i]) if (match[j] == -1) {
                match[i] = i:
1bc
```

```
f1d
                match[i] = i;
Odf
                ans++:
c2b
                break;
723
            }
        for (int i = 0; i < n; i++) if (match[i] == -1) {
da8
7e3
            int j = getpath(i);
5f2
            if (j == -1) continue;
Odf
            ans++;
3a0
            while (j != -1) {
                int p = pai[j], pp = match[p];
ef0
348
                match[p] = j;
fe9
                match[j] = p;
55d
                j = pp;
797
            }
        }
f70
ba7
        return ans;
fcd }
```

2.6 Centro de arvore

```
// Retorna o diametro e o(s) centro(s) da arvore
// Uma arvore tem sempre um ou dois centros e estes estao no meio do
    diametro
//
// O(n)
042 vector <int> g[MAX]:
df1 int d[MAX], par[MAX];
544 pair<int, vector<int>> center() {
a95
        int f, df;
36d
         function < void(int) > dfs = [&] (int v) {
             if (d[v] > df) f = v, df = d[v];
d47
             for (int u : g[v]) if (u != par[v])
e68
                 d[u] = d[v] + 1, par[u] = v, dfs(u);
1a5
90d
        };
        f = df = par[0] = -1, d[0] = 0:
1b0
41e
        dfs(0):
c2d
        int root = f:
0f6
        f = df = par[root] = -1, d[root] = 0;
14e
         dfs(root):
761
        vector<int> c;
87 e
         while (f != -1) {
999
             if (d[f] == df/2 \text{ or } d[f] == (df+1)/2) \text{ c.push_back}(f);
19c
            f = par[f];
```

```
3bf
        }
00f
        return {df, c};
9c7 }
2.7 Centroid
// Computa os 2 centroids da arvore
//
// O(n)
97a int n, subsize[MAX];
042 vector < int > g[MAX];
98f void dfs(int k, int p=-1) {
bd2
        subsize[k] = 1:
6e5
        for (int i : g[k]) if (i != p) {
801
            dfs(i, k);
2e3
            subsize[k] += subsize[i]:
1b2
5a5 }
2e8 int centroid(int k, int p=-1, int size=-1) {
       if (size == -1) size = subsize[k]:
e73
8df
        for (int i : g[k]) if (i != p) if (subsize[i] > size/2)
            return centroid(i, k, size);
bab
839
        return k;
b6a }
f20 pair <int, int > centroids(int k=0) {
051
        dfs(k):
        int i = centroid(k), i2 = i;
       for (int j : g[i]) if (2*subsize[j] == subsize[k]) i2 = j;
8dd
        return {i, i2}:
0cb
cf4 }
2.8 Centroid decomposition
// decomp(0, k) computa numero de caminhos com 'k' arestas
// Mudar depois do comentario
// O(n log(n))
042 vector < int > g[MAX];
ba8 int sz[MAX], rem[MAX];
747 void dfs(vector<int>& path, int i, int l=-1, int d=0) {
```

```
547
         path.push_back(d);
 75f
         for (int j : g[i]) if (j != 1 and !rem[j]) dfs(path, j, i,
    d+1);
 3e9 }
 071 int dfs_sz(int i, int l=-1) {
         sz[i] = 1:
         for (int j : g[i]) if (j != 1 and !rem[j]) sz[i] += dfs_sz(j,
    i);
 191
         return sz[i]:
 86b }
 85a int centroid(int i, int 1, int size) {
 994
         for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
 735
             return centroid(j, i, size);
 d9a
         return i:
 96e }
 d79 ll decomp(int i, int k) {
         int c = centroid(i, i, dfs sz(i));
 a67
         rem[c] = 1:
         // gasta O(n) aqui - dfs sem ir pros caras removidos
 04b
         11 \text{ ans} = 0;
 020
         vector<int> cnt(sz[i]);
 878
         cnt[0] = 1:
 0a8
         for (int j : g[c]) if (!rem[j]) {
 5b4
             vector < int > path;
 baf
             dfs(path, j);
 1a1
             for (int d: path) if (0 \le k-d-1 \text{ and } k-d-1 \le sz[i])
 285
                 ans += cnt[k-d-1]:
 e8b
             for (int d : path) cnt[d+1]++;
         }
 fa2
 1 c 1
         for (int j : g[c]) if (!rem[j]) ans += decomp(j, k);
 3f1
         rem[c] = 0;
 ba7
         return ans;
 193 }
 2.9 Centroid Tree
// Constroi a centroid tree
// p[i] eh o pai de i na centroid-tree
// dist[i][k] = distancia na arvore original entre i
// e o k-esimo ancestral na arvore da centroid
//
// O(n log(n)) de tempo e memoria
```

```
845 vector < int > g[MAX], dist[MAX];
c1e int sz[MAX], rem[MAX], p[MAX];
071 int dfs_sz(int i, int l=-1) {
        sz[i] = 1:
02c
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
e5c
   i);
191
        return sz[i];
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
735
            return centroid(j, i, size);
d9a
       return i;
96e }
324 void dfs_dist(int i, int 1, int d=0) {
        dist[i].push_back(d);
541
       for (int j : g[i]) if (j != l and !rem[j])
5a1
82a
            dfs_dist(j, i, d+1);
645 }
27e void decomp(int i, int l = -1) {
       int c = centroid(i, i, dfs_sz(i));
1b9
       rem[c] = 1, p[c] = 1;
534
       dfs_dist(c, c);
a2a
       for (int j : g[c]) if (!rem[j]) decomp(j, c);
ebd }
76c void build(int n) {
235
       for (int i = 0; i < n; i++) rem[i] = 0, dist[i].clear();
867
        for (int i = 0: i < n: i++) reverse(dist[i].begin().
   dist[i].end());
a78 }
2.10 Dijkstra
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = LINF.
// entao x nao alcanca i
// O(m log(n))
eff ll d[MAX];
```

```
c0d vector<pair<int, int>> g[MAX]; // {vizinho, peso}
1a8 int n;
abc void dijkstra(int v) {
        for (int i = 0; i < n; i++) d[i] = LINF;</pre>
        d[v] = 0;
a7f
88 c
        priority_queue < pair < ll, int >> pq;
b32
        pq.emplace(0, v);
265
        while (pq.size()) {
a25
            auto [ndist, u] = pq.top(); pq.pop();
953
            if (-ndist > d[u]) continue;
cda
            for (auto [idx, w] : g[u]) if (d[idx] > d[u] + w) {
331
                 d[idx] = d[u] + w;
a84
                pq.emplace(-d[idx], idx);
c56
        }
e5c
fec }
2.11 Dinitz
// O(min(m * max_flow, n^2 m))
// Grafo com capacidades 1: O(\min(m \text{ sqrt}(m), m * n^{2/3}))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
472 struct dinitz {
        const bool scaling = false; // com scaling -> 0(nm
   log(MAXCAP)),
        int lim;
                                     // com constante alta
206
670
        struct edge {
358
            int to, cap, rev, flow;
7f9
            bool res;
d36
            edge(int to_, int cap_, int rev_, bool res_)
                : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
a94
f70
        };
        vector<vector<edge>> g;
002
216
        vector < int > lev, beg;
a71
        11 F;
190
        dinitz(int n) : g(n), F(0) {}
087
        void add(int a, int b, int c) {
bae
            g[a].emplace_back(b, c, g[b].size(), false);
4c6
            g[b].emplace_back(a, 0, g[a].size()-1, true);
5c2
        }
```

```
123
        bool bfs(int s, int t) {
            lev = vector \langle int \rangle (g.size(), -1); lev[s] = 0;
90f
64c
            beg = vector < int > (g.size(), 0);
            queue < int > q; q.push(s);
8b2
402
            while (q.size()) {
                int u = q.front(); q.pop();
be1
bd9
                for (auto& i : g[u]) {
                     if (lev[i.to] != -1 or (i.flow == i.cap)) continue;
dbc
b4f
                     if (scaling and i.cap - i.flow < lim) continue;</pre>
                     lev[i.to] = lev[u] + 1;
185
8ca
                     q.push(i.to);
f97
                }
e87
            }
0de
            return lev[t] != -1;
742
dfb
        int dfs(int v, int s, int f = INF) {
50b
            if (!f or v == s) return f;
            for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
88f
027
                 auto& e = g[v][i];
                if (lev[e.to] != lev[v] + 1) continue;
206
                int foi = dfs(e.to, s, min(f, e.cap - e.flow));
ee0
                if (!foi) continue:
749
                 e.flow += foi, g[e.to][e.rev].flow -= foi;
3c5
45c
                return foi;
618
            }
bb3
            return 0;
4b1
ff6
        11 max flow(int s. int t) {
            for (lim = scaling ? (1 << 30) : 1; lim; lim /= 2)
a86
9d1
                 while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
4ff
            return F:
8b9
        }
86f }:
// Recupera as arestas do corte s-t
dbd vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
        g.max_flow(s, t);
f07
68 c
        vector < pair < int , int >> cut;
1b0
        vector < int > vis(g.g.size(), 0), st = {s};
321
        vis[s] = 1;
        while (st.size()) {
3c6
            int u = st.back(); st.pop_back();
b17
322
            for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)
c17
                 vis[e.to] = 1, st.push_back(e.to);
d14
        }
481
        for (int i = 0; i < g.g.size(); i++) for (auto e : g.g[i])</pre>
9d2
            if (vis[i] and !vis[e.to] and !e.res) cut.emplace_back(i,
```

```
e.to):
d1b
        return cut;
1e8 }
2.12 Dominator Tree
// Codigo do Kawakami. Se vira pra usar ai
//
// build - O(m log(n))
// dominates - 0(1)
1a8 int n;
bbf namespace d_tree {
        vector < int > g[MAX];
        // The dominator tree
        vector<int> tree[MAX];
b39
5af
        int dfs_1[MAX], dfs_r[MAX];
        // Auxiliary data
        vector < int > rg[MAX], bucket[MAX];
a2e
3ef
        int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
44b
        int ancestor[MAX], label[MAX];
563
        vector<int> preorder;
76a
        void dfs(int v) {
6a1
            static int t = 0;
db6
            pre[v] = ++t;
            sdom[v] = label[v] = v;
767
a3d
            preorder.push_back(v);
d08
            for (int nxt: g[v]) {
56c
                if (sdom[nxt] == -1) {
eed
                    prv[nxt] = v;
900
                    dfs(nxt);
f48
2b5
                rg[nxt].push_back(v);
            }
5a1
        }
d6a
        int eval(int v) {
62e
c93
            if (ancestor[v] == -1) return v;
a75
            if (ancestor[ancestor[v]] == -1) return label[v]:
f33
            int u = eval(ancestor[v]);
b49
            if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
66e
            ancestor[v] = ancestor[u]:
c24
            return label[v];
```

0b9

}

4b2 void dfs2(int v) { static int t = 0: 6a1 330 $dfs_1[v] = t++;$ 5e0 for (int nxt: tree[v]) dfs2(nxt); $dfs_r[v] = t++;$ 8e2 cfa c2c void build(int s) { for (int i = 0; i < n; i++) {</pre> 603 e6f sdom[i] = pre[i] = ancestor[i] = -1; 2e1 rg[i].clear(); 50a tree[i].clear(); 666 bucket[i].clear(); 3ba 772 preorder.clear(); c6c dfs(s); if (preorder.size() == 1) return; 12b 3c7 for (int i = int(preorder.size()) - 1; i >= 1; i--) { int w = preorder[i]; 6c6 for (int v: rg[w]) { a52 int u = eval(v): 5c1 if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] = sdom[u];</pre> a17 018 680 bucket[sdom[w]].push_back(w); ancestor[w] = prv[w]; ea7 b99 for (int v: bucket[prv[w]]) { 5c1 int u = eval(v): idom[v] = (u == v) ? sdom[v] : u;977 aff 2cc bucket[prv[w]].clear(); 0a3 for (int i = 1; i < preorder.size(); i++) {</pre> d0c 6c6 int w = preorder[i]; if (idom[w] != sdom[w]) idom[w] = idom[idom[w]]; 14b 32f tree[idom[w]].push back(w); c58 idom[s] = sdom[s] = -1;8ac 1b6 dfs2(s):d09 } // Whether every path from s to v passes through u bool dominates(int u. int v) { 490 c75 if (pre[v] == -1) return 1; // vacuously true 2ea return dfs_1[u] <= dfs_1[v] && dfs_r[v] <= dfs_r[u]; 332 } ce9 };

2.13 Euler Path / Euler Cycle

```
// Para declarar: 'euler < true > E(n); ' se quiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu.
// e um vector de {vertice, id da aresta para chegar no vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao 'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
// get_cycle() acha um ciclo euleriano se o grafo for euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
//
// O(n+m)
63f template <bool directed=false > struct euler {
1a8
        int n:
4c0
        vector < vector < pair < int , int >>> g;
        vector<int> used:
d63
30f
        euler(int n_) : n(n_), g(n) {}
50f
        void add(int a, int b) {
4cd
            int at = used.size();
c51
            used.push back(0):
74e
            g[a].emplace_back(b, at);
fab
            if (!directed) g[b].emplace_back(a, at);
        }
411
d41 #warning chamar para o src certo!
eed
        pair < bool, vector < pair < int, int >>> get_path(int src) {
baf
            if (!used.size()) return {true, {}};
b25
            vector < int > beg(n, 0);
4ec
            for (int& i : used) i = 0:
            // {{vertice, anterior}, label}
            vector<pair<int, int>, int>> ret, st = {{src, -1},
363
   -1}};
            while (st.size()) {
3c6
                int at = st.back().first.first:
8ff
002
                int& it = beg[at];
8 a 1
                while (it < g[at].size() and used[g[at][it].second])</pre>
   it++;
8e4
                if (it == g[at].size()) {
944
                    if (ret.size() and ret.back().first.second != at)
b82
                         return {false, {}};
420
                    ret.push_back(st.back()), st.pop_back();
```

```
2c0
                } else {
                     st.push_back({{g[at][it].first, at},
   g[at][it].second});
eb8
                    used[g[at][it].second] = 1;
                }
396
b3a
a19
            if (ret.size() != used.size()+1) return {false, {}};
f77
            vector < pair < int , int >> ans;
fdf
            for (auto i : ret) ans.emplace_back(i.first.first,
   i.second):
            reverse(ans.begin(), ans.end());
459
            return {true, ans};
997
844
9b6
        pair < bool, vector < pair < int, int >>> get_cycle() {
baf
            if (!used.size()) return {true, {}};
            int src = 0;
ad1
            while (!g[src].size()) src++;
34b
687
            auto ans = get_path(src);
            if (!ans.first or ans.second[0].first !=
33c
   ans.second.back().first)
                return {false, {}};
b82
            ans.second[0].second = ans.second.back().second;
350
868
            ans.second.pop_back();
ba7
            return ans;
48f
        }
711 }:
```

2.14 Euler Tour Tree

```
// Mantem uma floresta enraizada dinamicamente
// e permite queries/updates em sub-arvore
//
// Chamar ETT E(n, v), passando n = numero de vertices
// e v = vector com os valores de cada vertice (se for vazio,
// constroi tudo com 0
//
// link(v, u) cria uma aresta de v pra u, de forma que u se torna
// o pai de v (eh preciso que v seja raiz anteriormente)
// cut(v) corta a resta de v para o pai
// query(v) retorna a soma dos valores da sub-arvore de v
// update(v, val) soma val em todos os vertices da sub-arvore de v
// update_v(v, val) muda o valor do vertice v para val
// is_in_subtree(v, u) responde se o vertice u esta na sub-arvore de v
// Tudo O(log(n)) com alta probabilidade
878 mt19937 rng((int)
```

```
chrono::steady_clock::now().time_since_epoch().count());
9f9 template < typename T > struct ETT {
        // treap
3c9
        struct node {
ed1
            node *1, *r, *p;
fa4
            int pr, sz;
875
            T val, sub, lazy;
53e
            int id;
ffd
            bool f; // se eh o 'first'
5ef
            int qt_f; // numero de firsts na subarvore
            node(int id_, T v, bool f_ = 0) : l(NULL), r(NULL),
7a8
   p(NULL), pr(rng()),
62b
                sz(1), val(v), sub(v), lazy(), id(id_), f(f_),
   qt_f(f_) {}
a9c
            void prop() {
d09
                if (lazy != T()) {
021
                    if (f) val += lazy;
971
                    sub += lazy*sz;
b87
                    if (1) 1->lazy += lazy;
d3b
                    if (r) r->lazy += lazy;
30c
bfd
                lazy = T();
0fc
            }
0.1e
            void update() {
8da
                sz = 1, sub = val, qt_f = f;
171
                if (1) 1->prop(), sz += 1->sz, sub += 1->sub, qt_f +=
   1->qt_f;
117
                if (r) r - prop(), sz += r - sz, sub += r - sub, qt_f +=
   r->qt_f;
ccb
        };
bff
bb7
        node* root:
73c
        int size(node* x) { return x ? x->sz : 0; }
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
bcf
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
161
            1->prop(), r->prop();
            if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p = i = 1;
ff5
            else join(1, r->1, r->1), r->1->p = i = r;
982
bda
            i->update();
84d
a20
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
26a
            if (!i) return void(r = 1 = NULL);
c89
            i->prop();
            if (key + size(i->1) < v) {
d9e
```

```
448
                 split(i\rightarrow r, i\rightarrow r, r, v, key+size(i\rightarrow l)+1), l = i;
a21
                 if (r) r - p = NULL;
6e8
                if (i->r) i->r->p = i;
            } else {
396
                 split(i->1, 1, i->1, v, key), r = i;
98d
                 if (1) 1->p = NULL;
5a3
899
                if (i->1) i->1->p = i;
            }
18b
bda
            i->update();
134
ac7
        int get_idx(node* i) {
6cf
            int ret = size(i->1);
482
            for (: i->p: i = i->p) {
fbf
                 node* pai = i->p;
8a6
                if (i != pai->l) ret += size(pai->l) + 1;
            }
e22
edf
            return ret;
479
048
        node* get_min(node* i) {
433
            if (!i) return NULL:
f8e
            return i->1 ? get_min(i->1) : i;
0de
f03
        node* get_max(node* i) {
433
            if (!i) return NULL;
424
            return i->r ? get_max(i->r) : i;
e92
        // fim da treap
4fb
        vector < node *> first, last;
        ETT(int n, vector<T> v = {}) : root(NULL), first(n), last(n) {
f82
с5е
            if (!v.size()) v = vector<T>(n);
            for (int i = 0: i < n: i++) {</pre>
603
a00
                 first[i] = last[i] = new node(i, v[i], 1):
                 join(root, first[i], root);
469
            }
8ac
ec3
83f
        ETT(const ETT& t) { throw logic_error("Nao copiar a ETT!"); }
c09
        \simETT() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
                 node* x = q.back(); q.pop_back();
e5d
                if (!x) continue;
ee9
1c7
                q.push_back(x->1), q.push_back(x->r);
bf0
                 delete x;
            }
653
672
        }
```

```
153
        pair < int , int > get_range(int i) {
670
            return {get_idx(first[i]), get_idx(last[i])};
        }
ada
7af
        void link(int v, int u) { // 'v' tem que ser raiz
890
            auto [lv, rv] = get_range(v);
f13
            int ru = get_idx(last[u]);
4b4
            node* V;
df9
            node *L, *M, *R;
117
            split(root, M, R, rv+1), split(M, L, M, lv);
f1e
            V = M:
a28
            join(L, R, root);
e66
            split(root, L, R, ru+1);
367
            join(L, V, L);
7e8
            join(L, last[u] = new node(u, T() /* elemento neutro */),
   L);
a28
            join(L, R, root);
849
4e6
        void cut(int v) {
892
            auto [1, r] = get_range(v);
df9
            node *L, *M, *R;
dca
            split(root, M, R, r+1), split(M, L, M, 1);
de6
            node *LL = get_max(L), *RR = get_min(R);
710
            if (LL and RR and LL->id == RR->id) { // remove duplicata
e8b
                 if (last[RR->id] == RR) last[RR->id] = LL:
992
                 node *A, *B;
6b3
                 split(R, A, B, 1);
                 delete A:
10c
9d5
                 R = B;
7c0
            }
            join(L, R, root);
a28
a0d
            join(root, M, root);
6ff
        }
808
        T query(int v) {
892
            auto [1, r] = get_range(v);
df9
            node *L, *M, *R;
dca
            split(root, M, R, r+1), split(M, L, M, 1);
d43
            T ans = M->sub:
69d
            join(L, M, M), join(M, R, root);
            return ans;
ba7
ede
93ъ
        void update(int v, T val) { // soma val em todo mundo da
   subarvore
892
            auto [1, r] = get_range(v);
```

```
df9
            node *L, *M, *R;
                                                                           830
                                                                           753
dca
            split(root, M, R, r+1), split(M, L, M, 1);
409
            M->lazy += val;
69d
            join(L, M, M), join(M, R, root);
                                                                           bb3
                                                                           192 }
61c
129
        void update_v(int v, T val) { // muda o valor de v pra val
ac1
            int l = get_idx(first[v]);
df9
            node *L, *M, *R;
d0c
            split(root, M, R, l+1), split(M, L, M, 1);
            M->val = M->sub = val;
25 e
69d
            join(L, M, M), join(M, R, root);
630
934
        bool is_in_subtree(int v, int u) { // se u ta na subtree de v
890
            auto [lv, rv] = get_range(v);
6ec
            auto [lu, ru] = get_range(u);
732
            return lv <= lu and ru <= rv;</pre>
a21
        }
355
        void print(node* i) {
            if (!i) return;
eae
            print(i->1);
a1e
743
            cout << i->id+1 << " ";
f15
            print(i->r);
59f
        }
                                                                           1a8
065
        void print() { print(root); cout << endl; }</pre>
                                                                           ce2
045 }:
                                                                           f82
                                                                           ebd
2.15 Floyd-Warshall
                                                                           6a9
                                                                           405
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
                                                                           97 c
// returna 1 sse ha ciclo negativo
                                                                           27b
// d[i][i] deve ser 0
                                                                           9c9
// para i != j, d[i][j] deve ser w se ha uma aresta
                                                                           f09
// (i, j) de peso w, INF caso contrario
                                                                           e0a
//
                                                                           74c
// O(n^3)
                                                                           97b
                                                                           15c
                                                                           bfb
1a8 int n;
ae5 int d[MAX][MAX];
                                                                           a22
                                                                           c16
73c bool floyd_warshall() {
                                                                           8c0
```

for (int k = 0; k < n; k++)

for (int i = 0; i < n; i++)

for (int j = 0; j < n; j++)</pre>

d[i][j] = min(d[i][j], d[i][k] + d[k][j]);

e22 830

f90

0ab

```
830 for (int i = 0; i < n; i++)
753 if (d[i][i] < 0) return 1;

bb3 return 0;
192 }
```

2.16 Functional Graph

```
// rt[i] fala o ID da raiz associada ao vertice i
// d[i] fala a profundidade (0 sse ta no ciclo)
// pos[i] fala a posicao de i no array que eh a concat. dos ciclos
// build(f, val) recebe a funcao f e o custo de ir de
// i para f[i] (por default, val = f)
// f_k(i, k) fala onde i vai parar se seguir k arestas
// path(i, k) fala o custo (soma) seguir k arestas a partir de i
// Se quiser outra operacao, da pra alterar facil o codigo
// Codigo um pouco louco, tenho que admitir
// build - O(n)
// f_k - O(log(min(n, k)))
// path - O(\log(\min(n, k)))
6ef namespace func_graph {
        int n:
        int f[MAX], vis[MAX], d[MAX];
        int p[MAX], pp[MAX], rt[MAX], pos[MAX];
        int sz[MAX], comp:
        vector < vector < int >> ciclo;
        11 val[MAX], jmp[MAX], seg[2*MAX];
        11 op(11 a, 11 b) { return a+b; }; // mudar a operacao aqui
        void dfs(int i, int t = 2) {
            vis[i] = t;
            if (vis[f[i]] \ge 2) \{ // comeca ciclo - f[i] eh o rep.
                d[i] = 0, rt[i] = comp;
                sz[comp] = t - vis[f[i]] + 1;
                p[i] = pp[i] = i, jmp[i] = val[i];
                ciclo.emplace_back();
                ciclo.back().push_back(i);
            } else {
                if (!vis[f[i]]) dfs(f[i], t+1);
                rt[i] = rt[f[i]]:
195
                if (sz[comp]+1) { // to no ciclo
d0f
                    d[i] = 0;
97b
                    p[i] = pp[i] = i, jmp[i] = val[i];
bfb
                    ciclo.back().push_back(i);
c20
                } else { // nao to no ciclo
```

```
00d
                    d[i] = d[f[i]]+1, p[i] = f[i];
                    pp[i] = 2*d[pp[f[i]]] == d[pp[pp[f[i]]]]+d[f[i]]?
511
   pp[pp[f[i]]] : f[i];
                    jmp[i] = pp[i] == f[i] ? val[i] : op(val[i],
   op(jmp[f[i]], jmp[pp[f[i]]]));
                }
db8
003
            }
e4a
            if (f[ciclo[rt[i]][0]] == i) comp++; // fim do ciclo
29a
            vis[i] = 1;
       }
0ba
1da
        void build(vector<int> f_, vector<int> val_ = {}) {
bcb
            n = f_size(), comp = 0;
527
            if (!val .size()) val = f :
            for (int i = 0; i < n; i++)</pre>
830
                f[i] = f_{i}, val[i] = val_{i}, vis[i] = 0, sz[i] = -1;
998
e74
            ciclo.clear();
            for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
158
            int t = 0;
6bb
            for (auto& c : ciclo) {
daa
                reverse(c.begin(), c.end());
336
                for (int j : c) {
ea5
                    pos[i] = t;
85b
948
                    seg[n+t] = val[i];
c82
                    t++:
25e
                }
cbc
            for (int i = n-1; i; i--) seg[i] = op(seg[2*i],
   seg[2*i+1]);
       }
90b
283
        int f_k(int i, ll k) {
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
                if (big <= k) k -= big, i = pp[i];</pre>
ded
                else k--, i = p[i];
584
09c
            }
77e
            if (!k) return i;
            return ciclo[rt[i]][(pos[i] - pos[ciclo[rt[i]][0]] + k) %
   sz[rt[i]]];
f34
       }
        ll path(int i, ll k) {
047
            auto guery = [&](int 1, int r) {
3cf
3e4
                11 a = 0:
47a
                for (1 += n, r += n; 1 <= r; ++1/=2, --r/=2) {
                    if (1\%2 == 1) q = op(q, seg[1]);
27 e
                    if (r\%2 == 0) q = op(q, seg[r]);
1f2
```

```
598
                }
bef
                return q;
6e1
            };
b73
            11 \text{ ret} = 0;
            while (d[i] and k) {
1b1
77b
                int big = d[i] - d[pp[i]];
327
                if (big <= k) k -= big, ret = op(ret, jmp[i]), i =</pre>
   pp[i];
f9e
                else k--, ret = op(ret, val[i]), i = p[i];
7e3
            }
еЗс
            if (!k) return ret;
a9e
            int first = pos[ciclo[rt[i]][0]], last =
   pos[ciclo[rt[i]].back()]:
            // k/sz[rt[i]] voltas completas
            if (k/sz[rt[i]]) ret = op(ret, k/sz[rt[i]] * query(first,
430
   last));
            k %= sz[rt[i]];
9af
e3c
            if (!k) return ret:
8ea
            int l = pos[i], r = first + (pos[i] - first + k - 1) %
   sz[rt[i]]:
982
            if (1 <= r) return op(ret, query(1, r));</pre>
687
            return op(ret, op(query(1, last), query(first, r)));
380
        }
51f }
2.17 HLD - aresta
// SegTree de soma
// query / update de soma das arestas
//
// Complexidades:
// build - O(n)
// query_path - 0(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
826 namespace hld {
c0d
        vector<pair<int, int> > g[MAX];
e65
        int pos[MAX], sz[MAX];
7c0
        int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
```

```
Осе
        void build_hld(int k, int p = -1, int f = 1) {
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
180
418
            for (auto& i : g[k]) if (i.first != p) {
                auto [u, w] = i:
dd2
a76
                sobe[u] = w; pai[u] = k;
                h[u] = (i == g[k][0] ? h[k] : u);
0 c 1
                build hld(u, k, f): sz[k] += sz[u]:
da7
865
                if (sz[u] > sz[g[k][0].first] or g[k][0].first == p)
                    swap(i, g[k][0]);
9a3
804
            }
667
            if (p*f == -1) build hld(h\lceil k \rceil = k, -1, t = 0):
4dd
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build_hld(root);
c83
            seg::build(t, v);
ea2
3fc
        11 query_path(int a, int b) {
2d5
            if (a == b) return 0:
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
29b
            if (h[a] == h[b]) return seg::query(pos[b]+1, pos[a]);
            return seg::query(pos[h[a]], pos[a]) +
fca
   query_path(pai[h[a]], b);
87f
        void update_path(int a, int b, int x) {
920
d54
            if (a == b) return:
            if (pos[a] < pos[b]) swap(a, b);
aa1
            if (h[a] == h[b]) return (void)seg::update(pos[b]+1,
   pos[a], x);
701
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
   b, x);
dbf
d0a
        11 query_subtree(int a) {
b9f
            if (sz[a] == 1) return 0;
2f6
            return seg::query(pos[a]+1, pos[a]+sz[a]-1);
77f
acc
        void update_subtree(int a, int x) {
            if (sz[a] == 1) return:
a5a
            seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
9cd
a46
        int lca(int a, int b) {
7be
aa1
            if (pos[a] < pos[b]) swap(a, b);
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
219
        }
```

599 }

2.18 HLD - vertice

```
// SegTree de soma
// query / update de soma dos vertices
//
// Complexidades:
// build - \Omega(n)
// query_path - 0(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
826 namespace hld {
        vector < int > g[MAX];
042
e65
        int pos[MAX], sz[MAX];
bd4
        int peso[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
Осе
        void build_hld(int k, int p = -1, int f = 1) {
            v[pos[k] = t++] = peso[k]; sz[k] = 1;
b18
b94
            for (auto& i : g[k]) if (i != p) {
78d
                pai[i] = k;
                h[i] = (i == g[k][0] ? h[k] : i):
26e
                build_hld(i, k, f); sz[k] += sz[i];
193
                if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
cd1
    g[k][0]);
d94
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
a83
        void build(int root = 0) {
1f8
a34
            t = 0:
295
            build hld(root):
c83
            seg::build(t, v);
ea2
        }
3fc
        11 query_path(int a, int b) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
4bf
            if (h[a] == h[b]) return seg::query(pos[b], pos[a]);
            return seg::query(pos[h[a]], pos[a]) +
fca
    query_path(pai[h[a]], b);
c17
        }
920
        void update_path(int a, int b, int x) {
```

```
aa1
            if (pos[a] < pos[b]) swap(a, b);</pre>
            if (h[a] == h[b]) return (void) seg::update(pos[b], pos[a],
198
   x);
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
701
   b, x);
421
        }
        11 query_subtree(int a) {
d0a
            return seg::query(pos[a], pos[a]+sz[a]-1);
b3e
ba2
        void update_subtree(int a, int x) {
acc
a22
            seg::update(pos[a], pos[a]+sz[a]-1, x);
480
7be
        int lca(int a, int b) {
aa1
            if (pos[a] < pos[b]) swap(a, b);
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
219
        }
de3 }
```

2.19 HLD sem Update

```
// query de min do caminho
// Complexidades:
// build - O(n)
// query_path - O(log(n))
826 namespace hld {
c0d
        vector<pair<int, int> > g[MAX];
       int pos[MAX], sz[MAX];
e65
7c0
       int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
ea2
        int men[MAX], seg[2*MAX];
0ce
        void build_hld(int k, int p = -1, int f = 1) {
180
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
            for (auto& i : g[k]) if (i.first != p) {
418
                sobe[i.first] = i.second; pai[i.first] = k;
1f5
                h[i.first] = (i == g[k][0] ? h[k] : i.first);
6fa
                men[i.first] = (i == g[k][0] ? min(men[k], i.second) :
   i.second);
                build_hld(i.first, k, f); sz[k] += sz[i.first];
4b2
                if (sz[i.first] > sz[g[k][0].first] or g[k][0].first
bc3
   == p)
9a3
                    swap(i, g[k][0]);
ea4
            }
```

```
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
8ec
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build hld(root):
3ae
            for (int i = 0; i < t; i++) seg[i+t] = v[i];</pre>
            for (int i = t-1; i : i--) seg[i] = min(seg[2*i].
8db
    seg[2*i+1]);
ea5
        int query_path(int a, int b) {
f 04
490
            if (a == b) return INF;
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
98f
            if (h[a] != h[b]) return min(men[a], query_path(pai[h[a]],
   b));
46b
            int ans = INF, x = pos[b]+1+t, y = pos[a]+t;
            for (; x \le y; ++x/=2, --y/=2) ans = min({ans, seg[x],
646
   seg[y]});
ba7
            return ans;
3a9
        }
ee6 };
2.20 Isomorfismo de arvores
// thash() retorna o hash da arvore (usando centroids como vertices
   especiais).
// Duas arvores sao isomorfas sse seu hash eh o mesmo
// O(|V|.log(|V|))
91f map < vector < int >, int > mphash;
df6 struct tree {
1a8
        int n;
789
        vector < vector < int >> g;
        vector<int> sz, cs;
347
        tree(int n_) : n(n_), g(n_), sz(n_) {}
1b5
        void dfs_centroid(int v, int p) {
76b
588
            sz[v] = 1;
fa7
            bool cent = true:
```

for (int u : g[v]) if (u != p) {

if(sz[u] > n/2) cent = false;

dfs_centroid(u, v), sz[v] += sz[u];

if (cent and $n - sz[v] \le n/2$) cs.push_back(v);

18e

365

e90

есе

1f6

}

```
368
784
        int fhash(int v, int p) {
544
            vector < int > h;
332
            for (int u : g[v]) if (u != p) h.push_back(fhash(u, v));
1c9
            sort(h.begin(), h.end());
            if (!mphash.count(h)) mphash[h] = mphash.size();
3ac
bbc
            return mphash[h];
748
        }
38f
       11 thash() {
23a
            cs.clear():
3a5
            dfs_centroid(0, -1);
16d
            if (cs.size() == 1) return fhash(cs[0], -1);
772
            11 h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
            return (min(h1, h2) << 30) + max(h1, h2);</pre>
fae
       }
138
4dd };
```

2.21 Kosaraju

```
// O(n + m)
1a8 int n:
042 vector < int > g[MAX];
58d vector <int > gi[MAX]; // grafo invertido
c5a int vis[MAX];
ee6 stack<int> S:
a52 int comp[MAX]; // componente conexo de cada vertice
1ca void dfs(int k) {
        vis[k] = 1:
59a
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
8d5
            if (!vis[g[k][i]]) dfs(g[k][i]);
58f
        S.push(k);
89c }
436 void scc(int k, int c) {
59a
       vis[k] = 1;
        comp[k] = c;
52c
        for (int i = 0; i < (int) gi[k].size(); i++)</pre>
ff0
bf6
            if (!vis[gi[k][i]]) scc(gi[k][i], c);
088 }
db8 void kosaraju() {
991
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
158
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
```

```
991
        for (int i = 0; i < n; i++) vis[i] = 0;
d32
         while (S.size()) {
70b
            int u = S.top();
7de
            S.pop();
f43
            if (!vis[u]) scc(u, u);
 207
        }
e21 }
2.22 Kruskal
// Gera e retorna uma AGM e seu custo total a partir do vetor de
    arestas (edg)
// do grafo
//
// O(m log(m) + m a(m))
// 864875
1b9 vector<tuple<int, int, int>> edg; // {peso,[x,y]}
// DSU em O(a(n))
4a6 void dsu build():
d78 int find(int a);
369 void unite(int a, int b);
 c67 pair<ll, vector<tuple<int, int, int>>> kruskal(int n) {
8d2
         dsu build(n):
e31
         sort(edg.begin(), edg.end());
854
        11 cost = 0;
979
        vector<tuple<int, int, int>> mst;
fea
        for (auto [w,x,y]: edg) if (find(x) != find(y)) {
9de
            mst.emplace_back(w, x, y);
45f
            cost += w;
05a
            unite(x.v):
ca2
5df
        return {cost, mst};
b6a }
2.23 Kuhn
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// 'recover' recupera o min vertex cover como um par de
```

```
// {caras da particao A, caras da particao B}
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinic
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
6c6 struct kuhn {
14e
        int n. m:
789
        vector < vector < int >> g;
       vector < int > vis, ma, mb;
40e
        kuhn(int n_, int m_) : n(n_), m(m_), g(n),
8af
            vis(n+m), ma(n, -1), mb(m, -1) {}
        void add(int a, int b) { g[a].push_back(b); }
        bool dfs(int i) {
caf
            vis[i] = 1:
29a
29b
            for (int j : g[i]) if (!vis[n+j]) {
8c9
                vis[n+j] = 1;
                if (mb[j] == -1 or dfs(mb[j])) {
2cf
bfe
                    ma[i] = j, mb[j] = i;
8a6
                    return true;
b17
                }
82a
d1f
            return false:
4ef
bf7
        int matching() {
            int ret = 0, aum = 1;
1ae
5a8
            for (auto& i : g) shuffle(i.begin(), i.end(), rng);
            while (aum) {
392
618
                for (int j = 0; j < m; j++) vis[n+j] = 0;
                aum = 0:
c5d
830
                for (int i = 0; i < n; i++)
                    if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
01f
085
            }
edf
            return ret;
2ee
b0d }:
ebf pair < vector < int >, vector < int >> recover (kuhn & K) {
e80
       K.matching();
       int n = K.n, m = K.m;
50c
       for (int i = 0; i < n+m; i++) K.vis[i] = 0;</pre>
940
       for (int i = 0; i < n; i++) if (K.ma[i] == -1) K.dfs(i);</pre>
bde
```

```
8ad
        vector < int > ca, cb;
576
        for (int i = 0; i < n; i++) if (!K.vis[i]) ca.push_back(i);</pre>
f24
        for (int i = 0; i < m; i++) if (K.vis[n+i]) cb.push_back(i);</pre>
        return {ca, cb};
55f }
2.24 LCA com binary lifting
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
//
// Complexidades:
// build - O(n log(n))
// lca - O(log(n))
677 vector < vector < int > > g(MAX);
41c int n, p;
e75 int pai[MAX2][MAX];
999 int in[MAX], out[MAX];
1ca void dfs(int k) {
fdf
        in[k] = p++;
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
9b7
            if (in[g[k][i]] == -1) {
                pai[0][g[k][i]] = k;
ba6
c38
                dfs(g[k][i]);
            }
e2d
        out[k] = p++;
26f
691 }
c11 void build(int raiz) {
a67
        for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
c63
        p = 0, memset(in, -1, sizeof in);
        dfs(raiz);
ecb
        // pd dos pais
511
        for (int k = 1; k < MAX2; k++) for (int i = 0; i < n; i++)
d38
            pai[k][i] = pai[k - 1][pai[k - 1][i]];
530 }
00f bool anc(int a, int b) { // se a eh ancestral de b
        return in[a] <= in[b] and out[a] >= out[b]:
2d6 }
7be int lca(int a, int b) {
        if (anc(a, b)) return a;
```

```
e52
        if (anc(b, a)) return b;
        // sobe a
        for (int k = MAX2 - 1; k >= 0; k--)
f70
            if (!anc(pai[k][a], b)) a = pai[k][a];
acf
847
        return pai[0][a];
5c4 }
// Alternativamente:
// 'binary lifting' gastando O(n) de memoria
// Da pra add folhas e fazer queries online
// 3 vezes o tempo do binary lifting normal
//
// build - O(n)
// kth, lca, dist - O(log(n))
9c6 int d[MAX], p[MAX], pp[MAX];
d40 void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
e9d void add_leaf(int i, int u) {
       p[i] = u, d[i] = d[u]+1;
e0b
        pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]] : u;
b15
33f }
c37 int kth(int i, int k) {
       int dd = max(0, d[i]-k):
935
        while (d[i] > dd) i = d[pp[i]] >= dd ? pp[i] : p[i];
d9a
        return i;
f3c }
7be int lca(int a, int b) {
a69
       if (d[a] < d[b]) swap(a, b);
       while (d[a] > d[b]) a = d[pp[a]] >= d[b] ? pp[a] : p[a];
6cd
984
       while (a != b) {
932
            if (pp[a] != pp[b]) a = pp[a], b = pp[b];
e7c
            else a = p[a], b = p[b];
4ea
3f5
        return a;
21d }
4fe int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)]; }
042 vector < int > g[MAX];
3ab void build(int i, int pai=-1) {
```

```
5cf
        if (pai == -1) set_root(i);
15f
        for (int j : g[i]) if (j != pai) {
d31
            add_leaf(j, i);
b21
            build(j, i);
43b
        }
74a }
2.25 LCA com HLD
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b. lca(a, b) = a
// Para buildar pasta chamar build(root)
// anc(a, b) responde se 'a' eh ancestral de 'b'
//
// Complexidades:
// build - O(n)
// lca - O(log(n))
// anc - 0(1)
042 vector <int> g[MAX];
713 int pos[MAX], h[MAX], sz[MAX];
ff1 int pai[MAX], t;
8bf void build(int k, int p = -1, int f = 1) {
bce
        pos[k] = t++; sz[k] = 1;
e26
        for (int& i : g[k]) if (i != p) {
78d
            pai[i] = k;
26e
            h[i] = (i == g[k][0] ? h[k] : i);
            build(i, k, f); sz[k] += sz[i];
cb8
cd1
            if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i, g[k][0]);
917
        if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
3da
1b9 }
7be int lca(int a, int b) {
        if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
ca5
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
219 }
00f bool anc(int a, int b) {
db5
        return pos[a] \le pos[b] and pos[b] \le pos[a] + sz[a] - 1;
272 }
2.26 LCA com RMQ
```

// Assume que um vertice eh ancestral dele mesmo, ou seja,

```
// se a eh ancestral de b, lca(a, b) = a
// dist(a, b) retorna a distancia entre a e b
//
// Complexidades:
// build - O(n)
// lca - O(1)
// dist - 0(1)
1a5 template < typename T > struct rmq {
        vector <T> v;
517
fcc
        int n; static const int b = 30;
70e
        vector < int > mask, t;
18e
        int op(int x, int y) { return v[x] < v[y] ? x : y; }
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
6ad
        rmq() {}
43c
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) 
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {</pre>
2e5
                 at = (at << 1) &((1 << b) -1);
a61
                 while (at and op(i, i-msb(at&-at)) == i) at ^= at&-at;
76a
53 c
            }
            for (int i = 0; i < n/b; i++) t[i] =</pre>
243
   b*i+b-1-msb(mask[b*i+b-1]);
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
39d
   i+(1<<j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
   t[n/b*(j-1)+i+(1<<(j-1))]);
2d3
c92
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1)); }
        T query(int 1, int r) {
b7a
27b
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
7bf
            int ans = op(small(1+b-1), small(r));
e80
            int x = 1/b+1, y = r/b-1;
            if (x <= y) {
e25
                int j = msb(y-x+1);
a4e
                 ans = op(ans, op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
002
4b6
            }
ba7
            return ans;
6bf
        }
021 };
065 namespace lca {
042
        vector < int > g[MAX];
8ec
        int v[2*MAX], pos[MAX], dep[2*MAX];
8bd
        int t;
2de
        rmq < int > RMQ;
```

```
4cf
        void dfs(int i, int d = 0, int p = -1) {
c97
            v[t] = i, pos[i] = t, dep[t++] = d;
cac
            for (int j : g[i]) if (j != p) {
8ec
                dfs(j, d+1, i);
                v[t] = i, dep[t++] = d;
cf2
843
            }
        }
d6a
789
        void build(int n, int root) {
            t = 0:
a34
14e
            dfs(root);
3f4
            RMQ = rmq < int > (vector < int > (dep, dep + 2*n - 1));
657
        }
7be
        int lca(int a, int b) {
ab7
            a = pos[a], b = pos[b];
9c0
            return v[RMQ.query(min(a, b), max(a, b))];
5db
        }
b5d
        int dist(int a, int b) {
670
            return dep[pos[a]] + dep[pos[b]] - 2*dep[pos[lca(a, b)]];
5b7
        }
645 }
```

2.27 Line Tree

```
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
//
// build - O(n log(n))
// query - 0(log(n))
1a8 int n;
3ae namespace linetree {
        int id[MAX], seg[2*MAX], pos[MAX];
f37
43f
        vector<int> v[MAX], val[MAX];
        vector<pair<int, pair<int, int> > ar;
430
dc6
        void add(int a, int b, int p) { ar.push_back({p, {a, b}}); }
        void build() {
0a8
b09
            sort(ar.rbegin(), ar.rend());
            for (int i = 0; i < n; i++) id[i] = i, v[i] = \{i\},
    val[i].clear();
8bb
            for (auto i : ar) {
c91
                int a = id[i.second.first], b = id[i.second.second];
f6f
                if (a == b) continue;
```

```
c58
                if (v[a].size() < v[b].size()) swap(a, b);</pre>
fb8
                for (auto j : v[b]) id[j] = a, v[a].push_back(j);
482
                val[a].push_back(i.first);
                for (auto j : val[b]) val[a].push_back(j);
78b
                v[b].clear(), val[b].clear();
e39
012
8e8
            vector < int > vv:
            for (int i = 0; i < n; i++) for (int j = 0; j < 0
2ce
   v[i].size(); j++) {
                pos[v[i][j]] = vv.size();
e52
941
                if (j + 1 < v[i].size()) vv.push_back(val[i][j]);</pre>
                else vv.push_back(0);
1cb
475
            }
bb4
            for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];</pre>
            for (int i = n-1; i; i--) seg[i] = min(seg[2*i],
   seg[2*i+1]);
        }
9fe
4ea
        int query(int a, int b) {
            if (id[a] != id[b]) return 0; // nao estao conectados
596
ab7
            a = pos[a], b = pos[b];
d11
            if (a > b) swap(a, b);
199
            b--:
38a
            int ans = INF:
            for (a += n, b += n; a <= b; ++a/=2, --b/=2) ans =
   min({ans, seg[a], seg[b]});
            return ans:
ba7
952
        }
00f }:
2.28 Link-cut Tree
// Link-cut tree padrao
// Todas as operacoes sao O(log(n)) amortizado
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
062
            node() \{ p = ch[0] = ch[1] = -1; \}
f43
        };
5f3
        node t[MAX];
971
        bool is_root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
```

t[t[x].p].ch[1] != x);

}

cf1

```
ed6
         void rotate(int x) {
497
             int p = t[x].p, pp = t[p].p;
fc4
             if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
             bool d = t[p].ch[0] == x;
             t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
             if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a76
8fa
             t[x].p = pp, t[p].p = x;
49b
        }
07c
        void splay(int x) {
18c
             while (!is_root(x)) {
497
                int p = t[x].p, pp = t[p].p;
0 c 5
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
    == x) ? x : p):
64f
                rotate(x);
            }
d8d
        }
4fa
f16
        int access(int v) {
0eb
             int last = -1:
01a
             for (int w = v; w+1; last = w, splay(v), w = t[v].p)
024
                 splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
             return last;
0a4
e89
        int find_root(int v) {
5e3
             access(v);
3de
             while (t[v].ch[0]+1) v = t[v].ch[0];
f05
             return splay(v), v;
ee7
        }
142
        void link(int v. int w) { // v deve ser raiz
5e3
             access(v);
10d
            t[v].p = w;
c56
4e6
         void cut(int v) { // remove aresta de v pro pai
5e3
             access(v):
264
             t[v].ch[0] = t[t[v].ch[0]].p = -1;
5f5
bbb
        int lca(int v, int w) {
948
             return access(v), access(w);
b6d
        }
e4e }
     Link-cut Tree - aresta
// Valores nas arestas
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
//
```

```
// Todas as operacoes sao O(log(n)) amortizado
1ef namespace lct {
        struct node {
3c9
19f
            int p, ch[2];
810
            ll val, sub;
aa6
            bool rev:
04a
            int sz, ar;
4e4
            ll lazy;
            node() {}
f93
7a8
            node(int v, int ar_) :
546
            p(-1), val(v), sub(v), rev(0), sz(ar_{-}), ar(ar_{-}), lazy(0) {
b07
                ch[0] = ch[1] = -1:
53b
            }
6e0
        };
c53
        node t[2*MAX]; // MAXN + MAXQ
        map<pair<int, int>, int> aresta;
99e
e4d
        int sz;
95a
        void prop(int x) {
dc1
            if (t[x].lazy) {
                if (t[x].ar) t[x].val += t[x].lazy;
25e
2ab
                t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
edc
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
            }
1ba
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1:
c3d
50e
230
            t[x].lazy = 0, t[x].rev = 0;
f9d
        }
564
        void update(int x) {
            t[x].sz = t[x].ar, t[x].sub = t[x].val;
1a3
8ca
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
621
                prop(t[x].ch[i]);
c4f
                t[x].sz += t[t[x].ch[i]].sz;
269
                t[x].sub += t[t[x].ch[i]].sub;
400
            }
28b
        bool is_root(int x) {
971
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
        void rotate(int x) {
ed6
```

```
497
            int p = t[x].p, pp = t[p].p;
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
fc4
251
            bool d = t[p].ch[0] == x;
461
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
8fa
            t[x].p = pp, t[p].p = x;
444
            update(p), update(x);
f31
       }
238
        int splay(int x) {
            while (!is_root(x)) {
18c
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x):
0 c 5
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
                rotate(x);
64f
72c
            }
            return prop(x), x;
aab
08f
        }
        int access(int v) {
f16
0eb
            int last = -1:
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
            return last;
294
9f1
        void make_tree(int v, int w=0, int ar=0) { t[v] = node(w, ar);
 }
        int find_root(int v) {
e89
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
        }
16a
82f
        bool conn(int v. int w) {
            access(v), access(w);
2cf
b9b
            return v == w ? true : t[v].p != -1;
        }
ec0
277
        void rootify(int v) {
5e3
            access(v);
            t[v].rev ^= 1;
a02
        }
a05
971
        11 query(int v, int w) {
b54
            rootify(w), access(v);
249
            return t[v].sub;
652
3fa
        void update(int v, int w, int x) {
b54
            rootifv(w). access(v):
```

```
12c
            t[v].lazy += x;
74 f
204
        void link_(int v, int w) {
821
            rootify(w);
389
            t[w].p = v;
523
6b8
        void link(int v, int w, int x) { // v--w com peso x
379
            int id = MAX + sz++;
110
            aresta[make_pair(v, w)] = id;
            make_tree(id, x, 1);
a88
c88
            link_(v, id), link_(id, w);
58c
e63
        void cut (int v. int w) {
            rootify(w), access(v);
b54
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
264
7cd
       }
031
        void cut(int v, int w) {
b0f
            int id = aresta[make_pair(v, w)];
            cut_(v, id), cut_(id, w);
a4a
840
       }
        int lca(int v, int w) {
bbb
5e3
            access(v):
a8b
            return access(w);
524
       }
9ce }
```

2.30 Link-cut Tree - vertice

```
// Valores nos vertices
// make_tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
//
// Todas as operacoes sao O(log(n)) amortizado
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
            ll val, sub;
810
aa6
            bool rev;
e4d
            int sz;
4e4
            ll lazy;
f93
            node() {}
aa0
            node(int v) : p(-1), val(v), sub(v), rev(0), sz(1),
   lazy(0) {
```

```
b07
                ch[0] = ch[1] = -1;
            }
c4e
2b7
        };
        node t[MAX];
5f3
95a
        void prop(int x) {
            if (t[x].lazy) {
dc1
9f7
                t[x].val += t[x].lazy, t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
edc
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
e26
            }
aa2
            if (t[x].rev) {
                swap(t[x].ch[0], t[x].ch[1]);
f95
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
c3d
50e
230
            t[x].lazy = 0, t[x].rev = 0;
c62
        }
564
        void update(int x) {
            t[x].sz = 1, t[x].sub = t[x].val;
ec2
8ca
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
621
                prop(t[x].ch[i]);
c4f
                t[x].sz += t[t[x].ch[i]].sz;
269
                t[x].sub += t[t[x].ch[i]].sub;
400
            }
        }
da7
971
        bool is root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
       }
cf1
        void rotate(int x) {
ed6
497
            int p = t[x].p, pp = t[p].p;
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
fc4
            bool d = t[p].ch[0] == x;
251
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
8fa
            t[x].p = pp, t[p].p = x;
444
            update(p), update(x);
f31
       }
238
        int splay(int x) {
            while (!is_root(x)) {
18c
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
0c5
   == x) ? x : p);
```

```
64f
                rotate(x);
72c
aab
            return prop(x), x;
08f
f16
       int access(int v) {
0eb
            int last = -1:
d9f
            for (int w = v: w+1: update(last = w), splay(v), w =
   t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
            return last:
294
       }
f17
        void make_tree(int v, int w) { t[v] = node(w); }
e89
        int find root(int v) {
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
16a
       }
f94
        bool connected(int v, int w) {
2cf
            access(v), access(w);
            return v == w ? true : t[v].p != -1;
b9b
ec6
277
        void rootify(int v) {
5e3
            access(v):
a02
            t[v].rev ^= 1;
a05
971
       11 query(int v, int w) {
            rootify(w), access(v);
b54
249
            return t[v].sub;
652
3fa
        void update(int v, int w, int x) {
            rootify(w), access(v);
b54
12c
            t[v].lazv += x;
74f
142
        void link(int v. int w) {
821
            rootify(w);
389
            t[w].p = v;
8a8
       }
031
        void cut(int v, int w) {
b54
            rootify(w), access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
d9a
       int lca(int v, int w) {
bbb
5e3
            access(v):
a8b
            return access(w);
524
        }
f9f }
```

2.31 Max flow com lower bound

```
// add(a, b, l, r):
// adiciona aresta de a pra b, onde precisa passar f de fluxo, l <= f
    <= r
// add(a, b, c):
// adiciona aresta de a pra b com capacidade c
//
// Mesma complexidade do Dinitz
cd5 struct lb_max_flow : dinitz {
        vector < int > d;
5ce
d8c
        lb_max_flow(int n) : dinitz(n + 2), d(n, 0) {}
b12
        void add(int a, int b, int l, int r) {
c97
            d[a] -= 1;
            d[b] += 1:
f1b
            dinitz::add(a, b, r - 1);
4c0
ed4
        }
        void add(int a, int b, int c) {
087
0f3
             dinitz::add(a, b, c);
039
        }
7a1
        bool has_circulation() {
            int n = d.size();
50c
            11 cost = 0;
854
603
            for (int i = 0: i < n: i++) {
                if (d[i] > 0) {
c69
f56
                     cost += d[i];
57a
                     dinitz::add(n, i, d[i]);
c72
                } else if (d[i] < 0) {</pre>
b76
                     dinitz::add(i, n+1, -d[i]);
b42
                }
676
            }
067
            return (dinitz::max_flow(n, n+1) == cost);
110
7bd
        bool has_flow(int src, int snk) {
387
             dinitz::add(snk, src, INF);
            return has_circulation();
e40
cc1
4eb
        11 max_flow(int src, int snk) {
            if (!has_flow(src, snk)) return -1;
ee8
4ad
            dinitz::F = 0;
fe5
            return dinitz::max flow(src, snk);
619
        }
e8f };
```

2.32 MinCostMaxFlow

```
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)</pre>
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s pra t
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
// Se nao tiver aresta com custo negativo, nao precisa do SPFA
// O(nm + f * m log n)
123 template < typename T > struct mcmf {
670
        struct edge {
b75
            int to, rev, flow, cap; // para, id da reversa, fluxo,
   capacidade
7f9
            bool res; // se eh reversa
            T cost; // custo da unidade de fluxo
635
            edge(): to(0), rev(0), flow(0), cap(0), cost(0),
   res(false) {}
1d7
            edge(int to_, int rev_, int flow_, int cap_, T cost_, bool
   res )
                : to(to_), rev(rev_), flow(flow_), cap(cap_),
   res(res_), cost(cost_) {}
723
        };
002
        vector < vector < edge >> g;
168
        vector < int > par_idx, par;
f1e
        T inf;
a03
        vector <T> dist;
        mcmf(int n) : g(n), par_idx(n), par(n),
   inf(numeric_limits <T>::max()/3) {}
91c
        void add(int u, int v, int w, T cost) { // de u pra v com cap
   w e custo cost
            edge a = edge(v, g[v].size(), 0, w, cost, false);
2fc
234
            edge b = edge(u, g[u].size(), 0, 0, -cost, true);
b24
            g[u].push_back(a);
c12
            g[v].push_back(b);
0ed
        vector<T> spfa(int s) { // nao precisa se nao tiver custo
8bc
   negativo
            deque < int > q;
871
3d1
            vector < bool > is_inside(g.size(), 0);
577
            dist = vector <T>(g.size(), inf);
```

```
a93
             dist[s] = 0;
a30
             q.push_back(s);
ecb
             is_inside[s] = true;
14d
             while (!q.empty()) {
b1e
                 int v = q.front();
ced
                 q.pop_front();
48d
                 is_inside[v] = false;
76e
                 for (int i = 0; i < g[v].size(); i++) {</pre>
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
9d4
e61
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
943
                          dist[to] = dist[v] + cost;
ed6
                          if (is_inside[to]) continue;
020
                          if (!q.empty() and dist[to] > dist[q.front()])
   q.push_back(to);
b33
                          else q.push_front(to);
                          is_inside[to] = true;
b52
2d1
                     }
8cd
                 }
             }
f2c
8d7
             return dist;
96c
2a2
        bool dijkstra(int s, int t, vector<T>& pot) {
489
             priority_queue < pair < T, int > , vector < pair < T, int > > ,
   greater<>> q;
577
             dist = vector<T>(g.size(), inf);
             dist[s] = 0;
a93
             q.emplace(0, s);
115
402
             while (q.size()) {
                 auto [d, v] = q.top();
91b
833
                 q.pop();
                 if (dist[v] < d) continue;</pre>
68b
76e
                 for (int i = 0; i < g[v].size(); i++) {</pre>
9d4
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
                     cost += pot[v] - pot[to];
e8c
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
943
                          dist[to] = dist[v] + cost;
441
                          q.emplace(dist[to], to);
                          par_idx[to] = i, par[to] = v;
88b
873
                     }
                 }
de3
9d4
1d4
             return dist[t] < inf;</pre>
        }
c68
```

```
3d2
        pair < int , T > min_cost_flow(int s, int t, int flow = INF) {
3dd
            vector<T> pot(g.size(), 0);
9e4
            pot = spfa(s); // mudar algoritmo de caminho minimo aqui
d22
            int f = 0:
ce8
            T ret = 0:
4a0
            while (f < flow and dijkstra(s, t, pot)) {</pre>
                 for (int i = 0; i < g.size(); i++)</pre>
bda
                     if (dist[i] < inf) pot[i] += dist[i];</pre>
d2a
                 int mn_flow = flow - f, u = t;
71b
045
                 while (u != s){
90f
                     mn_flow = min(mn_flow,
07d
                         g[par[u]][par_idx[u]].cap -
   g[par[u]][par_idx[u]].flow);
                     u = par[u];
3d1
                 }
935
                ret += pot[t] * mn_flow;
1f2
476
                 u = t;
045
                 while (u != s) {
                     g[par[u]][par_idx[u]].flow += mn_flow;
e09
                     g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
498
3d1
                     u = par[u];
bcc
                 }
04d
                 f += mn_flow;
            }
36d
15b
            return make_pair(f, ret);
        }
cc3
        // Opcional: retorna as arestas originais por onde passa flow
182
        vector < pair < int , int >> recover() {
24a
            vector < pair < int , int >> used;
2a4
            for (int i = 0; i < g.size(); i++) for (edge e : g[i])</pre>
587
                 if(e.flow == e.cap && !e.res) used.push_back({i,
   e.to}):
f6b
            return used;
390
        }
697 };
```

2.33 Prufer code

```
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de 0 a n-1
// Todo array com n-2 posicoes e valores de
// O a n-1 sao prufer codes validos
//
// O(n)
47d vector <int > to_prufer(vector <pair <int, int >> tree) {
        int n = tree.size()+1;
2cf
        vector<int> d(n, 0);
        vector < vector < int >> g(n);
4aa
f87
        for (auto [a, b] : tree) d[a]++, d[b]++,
f60
             g[a].push_back(b), g[b].push_back(a);
c5a
        vector < int > pai(n, -1);
260
        queue < int > q; q.push(n-1);
402
        while (q.size()) {
             int u = q.front(); q.pop();
be1
             for (int v : g[u]) if (v != pai[u])
34 c
9c9
                 pai[v] = u, q.push(v);
        }
70d
399
        int idx, x;
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
4b8
        vector < int > ret;
b28
        for (int i = 0; i < n-2; i++) {
d4b
             int y = pai[x];
e81
             ret.push_back(y);
666
             if (-d[y] == 1 \text{ and } y < idx) x = y;
367
             else idx = x = find(d.begin()+idx+1, d.end(), 1) -
    d.begin();
5f9
edf
        return ret;
d3b }
4d8 vector < pair < int , int >> from prufer (vector < int > p) {
455
        int n = p.size()+2;
        vector < int > d(n, 1);
126
650
        for (int i : p) d[i]++;
85b
        p.push_back(n-1);
399
        int idx, x;
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
1df
        vector<pair<int, int>> ret;
        for (int v : p) {
b06
             ret.push_back({x, y});
dab
666
             if (--d[y] == 1 \text{ and } y < idx) x = y;
             else idx = x = find(d.begin()+idx+1, d.end(), 1) -
367
    d.begin();
```

```
c3b
edf
        return ret;
765 }
2.34 Sack (DSU em arvores)
// Responde queries de todas as sub-arvores
// offline
//
// O(n log(n))
6bf int sz[MAX], cor[MAX], cnt[MAX];
042 vector < int > g[MAX];
6df void build(int k, int d=0) {
e8f
        sz[k] = 1:
01a
        for (auto& i : g[k]) {
            build(i, d+1); sz[k] += sz[i];
30f
925
            if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
011
189 }
74f void compute(int k, int x, bool dont=1) {
        cnt[cor[k]] += x:
de9
828
        for (int i = dont; i < g[k].size(); i++)</pre>
            compute(g[k][i], x, 0);
b5c
896 }
dc4 void solve(int k, bool keep=0) {
        for (int i = int(g[k].size())-1; i >= 0; i--)
32a
b4c
            solve(g[k][i], !i);
4a0
        compute(k. 1):
        // agora cnt[i] tem quantas vezes a cor
        // i aparece na sub-arvore do k
830
        if (!keep) compute(k, -1, 0);
8bc }
     Stable Marriage
// Emparelha todos os elementos de A com elementos de B
// de forma que nao exista um par x \in A, y \in B
// e x nao pareado com y tal que x prefira parear com y
// e y prefira parear com x.
//
// a[i] contem os elementos de B ordenados por preferencia de i
```

```
// b[j] contem os elementos de A ordenados por preferencia de j
// |A| <= |B|
//
// Retorna um vetor v de tamanho |A| onde v[i] guarda o match de i.
// O(|A| * |B|)
 380 vector<int> stable_marriage(vector<vector<int>> &a,
    vector < vector < int >> &b) {
 652
         int n = a.size(), m = b.size();
 83e
         assert(a[0].size() == m and b[0].size() == n and n <= m);
 017
         vector < int > match(m, -1), it(n, 0);
 e6f
         vector inv b(m, vector<int>(n));
 a34
         for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
 9f2
             inv_b[i][b[i][j]] = j;
 26a
         queue < int > q;
         for (int i = 0; i < n; i++) q.push(i);</pre>
 5af
 402
         while (q.size()) {
 379
             int i = q.front(); q.pop();
             int j = a[i][it[i]];
 4b8
 57 c
             if (match[j] == -1) match[j] = i;
 02d
             else if (inv_b[j][i] < inv_b[j][match[j]]) {</pre>
 5d1
                 q.emplace(match[j]);
 e7d
                 it[match[j]]++;
 f1d
                 match[j] = i;
 bc4
             } else q.emplace(i), it[i]++;
 258
         }
 825
         vector < int > ret(n);
 d72
         for (int i = 0; i < m; i++) if (match[i] != -1) ret[match[i]]</pre>
    = i:
 edf
         return ret:
 Off }
2.36 Tarjan para SCC
// O(n + m)
042 vector < int > g[MAX];
 4ce stack<int> s:
 a42 int vis[MAX], comp[MAX];
3fd int id[MAX];
// se quiser comprimir ciclo ou achar ponte em grafo nao direcionado,
// colocar um if na dfs para nao voltar pro pai da DFS tree
```

```
f32 int dfs(int i, int& t) {
                                                                            57e
                                                                                    };
cf0
        int lo = id[i] = t++;
18e
        s.push(i);
                                                                                    for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
                                                                            158
0c2
        vis[i] = 2;
                                                                                    if (!dag) ret.clear();
                                                                            d8f
        for (int j : g[i]) {
                                                                                     return ret;
48e
                                                                            edf
740
            if (!vis[j]) lo = min(lo, dfs(j, t));
                                                                            d6b }
994
            else if (vis[j] == 2) lo = min(lo, id[j]);
d64
        }
                                                                            2.38 Vertex cover
        // aresta de i pro pai eh uma ponte (no caso nao direcionado)
                                                                            // Encontra o tamanho do vertex cover minimo
        if (lo == id[i]) while (1) {
3de
                                                                            // Da pra alterar facil pra achar os vertices
3c3
            int u = s.top(); s.pop();
                                                                            // Parece rodar com < 2 s pra N = 90
            vis[u] = 1, comp[u] = i;
9c5
            if (u == i) break;
2ef
                                                                            // O(n * 1.38^n)
266
        }
                                                                            76a namespace cover {
253
        return lo;
                                                                                     const int MAX = 96;
38a }
                                                                            042
                                                                                    vector < int > g[MAX];
                                                                            823
                                                                                     bitset < MAX > bs[MAX];
f93 void tarjan(int n) {
                                                                            1a8
                                                                                    int n:
        int t = 0:
6bb
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
991
                                                                            697
                                                                                     void add(int i, int j) {
                                                                                         if (i == j) return;
                                                                            bd0
3be
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i, t);</pre>
                                                                            78c
                                                                                         n = max({n, i+1, j+1});
ea1 }
                                                                            200
                                                                                         bs[i][j] = bs[j][i] = 1;
                                                                            203
                                                                                    }
      Topological Sort
2.37
                                                                            6c0
                                                                                     int rec(bitset < MAX > m) {
                                                                                         int ans = 0:
// Retorna uma ordenacaoo topologica de g
                                                                            1a4
// Se g nao for DAG retorna um vetor vazio
                                                                            25b
                                                                                         for (int x = 0; x < n; x++) if (m[x]) {
//
                                                                            002
                                                                                             bitset < MAX > comp;
// O(n + m)
                                                                            4bf
                                                                                             function < void(int) > dfs = [&](int i) {
                                                                            b96
                                                                                                 comp[i] = 1, m[i] = 0;
042 vector < int > g[MAX];
                                                                            0c3
                                                                                                 for (int j : g[i]) if (m[j]) dfs(j);
                                                                            815
                                                                                             };
b6a vector<int> topo_sort(int n) {
                                                                            963
                                                                                             dfs(x);
        vector < int > ret(n,-1), vis(n,0);
46e
                                                                            d34
                                                                                             int ma, deg = -1, cvc = 1;
f51
                                                                            417
                                                                                             for (int i = 0; i < n; i++) if (comp[i]) {</pre>
        int pos = n-1, dag = 1;
        function < void(int) > dfs = [&](int v) {
36d
                                                                            d0b
                                                                                                 int d = (bs[i]&comp).count();
cca
            vis[v] = 1:
                                                                            18a
440
            for (auto u : g[v]) {
                                                                            c1f
152
                if (vis[u] == 1) dag = 0;
                                                                            d8e
532
                else if (!vis[u]) dfs(u);
                                                                            269
e37
            }
```

d44

ret[pos--] = v, vis[v] = 2;

```
702
3f9
                 comp[ma] = 0;
                 // ou ta no cover, ou nao ta no cover
                 ans += min(1 + rec(comp), deg + rec(comp & ~bs[ma]));
1dd
            }
6e6
ba7
            return ans;
2ec
        }
f5c
        int solve() {
3c5
            bitset < MAX > m:
603
            for (int i = 0; i < n; i++) {</pre>
939
                m[i] = 1;
f90
                for (int j = 0; j < n; j++)
741
                     if (bs[i][j]) g[i].push_back(j);
13e
4f9
            return rec(m);
708
        }
9c5 }
```

2.39 Virtual Tree

```
// Comprime uma arvore dado um conjunto S de vertices, de forma que
// o conjunto de vertices da arvore comprimida contenha S e seja
// minimal e fechado sobre a operacao de LCA
// Se |S| = k, a arvore comprimida tem menos que 2k vertices
// As arestas de virt possuem a distancia do vertice ate o vizinho
// Retorna a raiz da virtual tree
// lca::pos deve ser a ordem de visitacao no dfs
// voce pode usar o LCAcomHLD, por exemplo
// O(k log(k))
b36 vector <pair <int, int>> virt[MAX];
d41 #warning lembrar de buildar o LCA antes
c14 int build_virt(vector<int> v) {
        auto cmp = [&](int i, int j) { return lca::pos[i] <</pre>
   lca::pos[i]; };
074
        sort(v.begin(), v.end(), cmp);
        for (int i = v.size()-1; i; i--) v.push_back(lca::lca(v[i],
   v[i-1])):
074
        sort(v.begin(), v.end(), cmp);
d76
        v.erase(unique(v.begin(), v.end()), v.end());
37 c
        for (int i = 0; i < v.size(); i++) virt[v[i]].clear();</pre>
        for (int i = 1; i < v.size(); i++) virt[lca::lca(v[i-1],</pre>
   v[i])].clear();
```

3 DP

3.1 Divide and Conquer DP

```
// Particiona o array em k subarrays
// minimizando o somatorio das queries
//
// O(k n log n), assumindo quer query(1, r) eh O(1)
547 ll dp[MAX][2];
94b void solve(int k, int l, int r, int lk, int rk) {
de6
        if (1 > r) return;
109
        int m = (1+r)/2, p = -1;
d2b
        auto& ans = dp[m][k&1] = LINF;
6e2
        for (int i = max(m, lk); i <= rk; i++) {</pre>
07b
            ll at = dp[i+1][\sim k\&1] + query(m, i);
57d
            if (at < ans) ans = at, p = i;</pre>
8f5
        }
1ee
        solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
d3e }
cf1 ll DC(int n. int k) {
321
        dp[n][0] = dp[n][1] = 0;
f27
        for (int i = 0; i < n; i++) dp[i][0] = LINF;</pre>
        for (int i = 1; i <= k; i++) solve(i, 0, n-i, 0, n-i);
b76
        return dp[0][k&1];
8e7
5e9 }
```

3.2 Longest Common Subsequence

```
// Computa a LCS entre dois arrays usando
// o algoritmo de Hirschberg para recuperar
//
// O(n*m), O(n+m) de memoria
```

```
eaf int lcs_s[MAX], lcs_t[MAX];
a6d int dp[2][MAX];
// dp[0][j] = max lcs(s[li...ri], t[lj, lj+j])
d12 void dp_top(int li, int ri, int lj, int rj) {
        memset(dp[0], 0, (rj-lj+1)*sizeof(dp[0][0]));
d13
753
        for (int i = li: i <= ri: i++) {</pre>
9aa
            for (int j = rj; j >= lj; j--)
83b
                dp[0][i - 1i] = max(dp[0][i - 1i],
                (lcs_s[i] == lcs_t[j]) + (j > lj ? dp[0][j-1 - lj] :
741
   0));
            for (int j = lj+1; j <= rj; j++)
04c
939
                dp[0][j-1j] = max(dp[0][j-1j], dp[0][j-1-1j]);
09f
        }
58f }
// dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
ca0 void dp_bottom(int li, int ri, int lj, int rj) {
        memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
0dd
        for (int i = ri; i >= li; i--) {
3a2
            for (int j = lj; j <= rj; j++)</pre>
49 c
dbb
                dp[1][j - 1j] = max(dp[1][j - 1j],
                (lcs_s[i] == lcs_t[j]) + (j < rj ? dp[1][j+1 - lj] :
4da
   0));
            for (int j = rj-1; j >= lj; j--)
6ca
769
                dp[1][j-1j] = max(dp[1][j-1j], dp[1][j+1-1j]);
        }
19b
e8a }
93c void solve(vector<int>& ans, int li, int ri, int lj, int rj) {
        if (li == ri){
2ad
49c
            for (int j = lj; j <= rj; j++)</pre>
                if (lcs_s[li] == lcs_t[j]){
a66
                    ans.push back(lcs t[i]):
c2b
                    break;
                }
840
505
            return;
126
534
        if (lj == rj){
753
            for (int i = li; i <= ri; i++){</pre>
                if (lcs_s[i] == lcs_t[lj]){
88f
531
                    ans.push_back(lcs_s[i]);
c2b
                    break:
                }
68a
a03
            }
505
            return;
76d
        }
```

```
dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
ade
        int i_{-} = 0, mx = -1;
d7a
aee
        for (int j = lj-1; j <= rj; j++) {
da8
            int val = 0:
            if (j >= lj) val += dp[0][j - lj];
2bb
b9e
            if (j < rj) val += dp[1][j+1 - lj];</pre>
ba8
            if (val >= mx) mx = val, i_ = i;
14e
        }
6f1
        if (mx == -1) return:
c2a
        solve(ans, li, mi, lj, j_), solve(ans, mi+1, ri, j_+1, rj);
dd5 }
058 vector<int> lcs(const vector<int>& s, const vector<int>& t) {
        for (int i = 0; i < s.size(); i++) lcs_s[i] = s[i];</pre>
953
577
        for (int i = 0; i < t.size(); i++) lcs_t[i] = t[i];</pre>
        vector < int > ans:
dab
599
        solve(ans, 0, s.size()-1, 0, t.size()-1);
ba7
        return ans:
17c }
3.3 Mochila
// Resolve mochila, recuperando a resposta
// O(n * cap), O(n + cap) de memoria
add int v[MAX], w[MAX]; // valor e peso
582 int dp[2][MAX_CAP];
// DP usando os itens [1, r], com capacidade = cap
0d6 void get_dp(int x, int 1, int r, int cap) {
f8f
        memset(dp[x], 0, (cap+1)*sizeof(dp[x][0]));
574
        for (int i = 1; i \le r; i++) for (int j = cap; j \ge 0; j--)
            if (i - w[i] >= 0) dp[x][i] = max(dp[x][i], v[i] + dp[x][i]
3a9
    - w[i]]);
b95 }
5ab void solve(vector<int>& ans, int 1, int r, int cap) {
        if (1 == r) {
893
9ff
            if (w[1] <= cap) ans.push_back(1);</pre>
505
            return:
13a
        }
ee4
        int m = (1+r)/2;
```

a57

int mi = (li+ri)/2;

```
283
        get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
056
        int left_cap = -1, opt = -INF;
c94
        for (int j = 0; j \le cap; j++)
2f2
            if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
91d
                opt = at, left_cap = j;
da3
        solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap - left_cap);
d75 }
0d7 vector<int> knapsack(int n, int cap) {
        vector < int > ans:
dab
1e0
        solve(ans, 0, n-1, cap);
ba7
        return ans:
e4d }
3.4 SOS DP
// O(n 2^n)
// soma de sub-conjunto
e03 vector<ll> sos_dp(vector<ll> f) {
       int N = __builtin_ctz(f.size());
        assert((1<<N) == f.size());
e59
       for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
   mask++)
            if (mask >> i&1) f [mask] += f [mask^{(1 << i)}]:
796
abe
        return f;
bec }
// soma de super-conjunto
e03 vector<ll> sos_dp(vector<ll> f) {
        int N = builtin ctz(f.size()):
6c0
        assert((1<<N) == f.size());
e59
        for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
5a5
   mask++)
            if (\sim mask >> i\&1) f[mask] += f[mask^(1<<ii)];
a3c
        return f;
abe
dbd }
3.5 Subset sum
// Retorna max(x <= t tal que existe subset de w que soma x)</pre>
// O(n * max(w))
// O(max(w)) de memoria
```

```
efd int subset_sum(vector<int> w, int t) {
        int pref = 0, k = 0;
bb5
417
        while (k < w.size() and pref + w[k] <= t) pref += w[k++];</pre>
        if (k == w.size()) return pref;
1e7
        int W = *max_element(w.begin(), w.end());
444
        vector < int > last, dp(2*W, -1);
44d
        dp[W - (t-pref)] = k;
d7b
54d
        for (int i = k; i < w.size(); i++) {</pre>
288
            last = dp;
15f
            for (int x = 0; x < W; x++) dp[x+w[i]] = max(dp[x+w[i]],
   last[x]);
17b
            for (int x = 2*W - 1; x > W; x--)
303
                for (int j = max(0, last[x]); j < dp[x]; j++)
595
                     dp[x-w[j]] = max(dp[x-w[j]], j);
867
        }
        int ans = t;
2fb
1 c 1
        while (dp[W - (t-ans)] < 0) ans --;
ba7
        return ans;
d88 }
```

4 Problemas

4.1 Angle Range Intersection

```
// Computa intersecao de angulos
// Os angulos (arcos) precisam ter comprimeiro < pi
// (caso contrario a intersecao eh estranha)
//
// Tudo 0(1)
32a struct angle_range {
75e
        static constexpr ld ALL = 1e9, NIL = -1e9;
395
        ld 1, r;
        angle_range() : 1(ALL), r(ALL) {}
c77
        angle_range(1d 1_, 1d r_) : 1(1_), r(r_) { fix(1), fix(r); }
894
        void fix(ld& theta) {
4ee
da7
            if (theta == ALL or theta == NIL) return;
323
            if (theta > 2*pi) theta -= 2*pi;
868
            if (theta < 0) theta += 2*pi;</pre>
625
        }
2ee
        bool empty() { return 1 == NIL; }
931
        bool contains(ld q) {
40f
            fix(q);
4d7
            if (1 == ALL) return true;
fec
            if (1 == NIL) return false;
```

```
if (1 < r) return 1 < q and q < r;
6a6
075
            return q > 1 or q < r;</pre>
800
9c7
        friend angle_range operator &(angle_range p, angle_range q) {
743
            if (p.l == ALL or q.l == NIL) return q;
            if (q.1 == ALL or p.1 == NIL) return p;
20f
7d5
            if (p.1 > p.r \text{ and } q.1 > q.r) \text{ return } \{\max(p.1, q.1),
   min(p.r, q.r)};
            if (q.1 > q.r) swap(p.1, q.1), swap(p.r, q.r);
aa6
            if (p.1 > p.r) {
868
249
                if (q.r > p.1) return \{max(q.1, p.1), q.r\};
6f7
                 else if (q.1 < p.r) return {q.1, min(q.r, p.r)};</pre>
270
                return {NIL, NIL};
337
            }
5a8
            if (max(p.1, q.1) > min(p.r, q.r)) return {NIL, NIL};
bcb
            return {max(p.1, q.1), min(p.r, q.r)};
142
        }
5e1 };
```

4.2 Area da Uniao de Retangulos

```
// O(n log(n))
// 5d8d2f
aa4 namespace seg {
6b3
        pair < int , 11 > seg[4*MAX];
b1b
        ll lazy[4*MAX], *v;
1a8
        int n;
e01
        pair < int , ll > merge(pair < int , ll > 1, pair < int , ll > r) {
719
            if (1.second == r.second) return {1.first+r.first,
   1.second:
            else if (1.second < r.second) return 1;</pre>
53b
aa0
            else return r;
d82
        }
6fc
        pair < int, ll > build(int p=1, int l=0, int r=n-1) {
3c7
            lazv[p] = 0;
bf8
            if (1 == r) return seg[p] = {1, v[1]};
            int m = (1+r)/2;
ee4
            return seg[p] = merge(build(2*p, 1, m), build(2*p+1, m+1,
432
   r));
f94
d9e
        void build(int n2, 11* v2) {
680
            n = n2, v = v2:
6f2
            build();
f8a
        }
```

```
ceb
        void prop(int p, int 1, int r) {
208
            seg[p].second += lazy[p];
2c9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
3c7
            lazv[p] = 0;
bf2
693
        pair < int, ll > query (int a, int b, int p=1, int l=0, int r=n-1)
{
6b9
            prop(p, 1, r);
527
            if (a <= 1 and r <= b) return seg[p];</pre>
            if (b < 1 or r < a) return {0, LINF};</pre>
9b7
ee4
            int m = (1+r)/2;
            return merge (query (a, b, 2*p, 1, m), query (a, b, 2*p+1,
eeb
   m+1. r)):
786
07c
        pair < int , 11 > update(int a, int b, int x, int p=1, int 1=0,
   int r=n-1) {
6b9
            prop(p, 1, r);
9a3
            if (a \le 1 \text{ and } r \le b) {
                lazy[p] += x;
b94
6b9
                prop(p, 1, r);
534
                return seg[p];
821
e9f
            if (b < 1 or r < a) return seg[p];</pre>
ee4
            int m = (1+r)/2;
086
            return seg[p] = merge(update(a, b, x, 2*p, 1, m),
579
                     update(a, b, x, 2*p+1, m+1, r));
c65
        }
043 }:
eb5 ll seg_vec[MAX];
8be 11 area_sq(vector<pair<int, int>, pair<int, int>>> &sq){
28 c
        vector<pair<int, int>, pair<int, int>>> up;
60a
        for (auto it : sq){
619
            int x1, y1, x2, y2;
ae0
            tie(x1, y1) = it.first;
68e
            tie(x2, y2) = it.second;
80f
            up.push_back(\{\{x1+1, 1\}, \{y1, y2\}\}\});
            up.push_back({{x2+1, -1}, {y1, y2}});
aee
6c3
        }
092
        sort(up.begin(), up.end());
049
        memset(seg_vec, 0, sizeof seg_vec);
6fe
        11 H_MAX = MAX;
156
        seg::build(H_MAX-1, seg_vec);
7ba
        auto it = up.begin();
04b
        11 \text{ ans} = 0;
f14
        while (it != up.end()){
```

```
07f
            11 L = (*it).first.first;
            while (it != up.end() && (*it).first.first == L){
718
127
                int x, inc, y1, y2;
d35
                tie(x, inc) = it->first;
d3d
               tie(y1, y2) = it->second;
5d1
                seg::update(y1+1, y2, inc);
40d
9b1
            }
852
            if (it == up.end()) break;
            11 R = (*it).first.first;
d8a
f59
           11 W = R-L:
efd
            auto jt = seg::query(0, H_MAX-1);
91a
           11 H = H_MAX - 1;
            if (jt.second == 0) H -= jt.first;
e8a
8df
            ans += W*H;
5c8
       }
ba7
        return ans;
385 }
```

4.3 Area Maxima de Histograma

```
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
//
// O(n)
15e ll area(vector<int> v) {
b73
       11 \text{ ret} = 0;
4ce
        stack<int> s:
       // valores iniciais pra dar tudo certo
       v.insert(v.begin(), -1);
447
d56
        v.insert(v.end(), -1);
1f8
        s.push(0);
0be
        for(int i = 0; i < (int) v.size(); i++) {</pre>
            while (v[s.top()] > v[i]) {
78e
                11 h = v[s.top()]; s.pop();
265
                ret = max(ret, h * (i - s.top() - 1));
de1
            }
40a
18e
            s.push(i);
020
edf
        return ret;
e43 }
```

4.4 Binomial modular

```
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)
97c ll divi[MAX]:
398 ll expo(ll a, ll b, ll m) {
        if (!b) return 1:
1 c 1
399
        ll ans = expo(a*a\%m, b/2, m);
        if (b\%2) ans *= a:
751
2e9
        return ans%m;
754 }
f0a ll inv(ll a, ll b){
        return 1 < a ? b - inv(b%a,a)*b/a : 1;
bca
041 }
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
3bd
        if (!a) return {b, 0, 1};
550
        auto [g, x, y] = ext_gcd(b\%a, a);
        return \{g, y - b/a*x, x\};
c59
537 }
bfe template < typename T = 11> struct crt {
627
        Ta, m;
        crt(): a(0), m(1) {}
5f3
        crt(T a_, T m_) : a(a_), m(m_) {}
7eb
911
        crt operator * (crt C) {
238
            auto [g, x, y] = ext_gcd(m, C.m);
dc0
            if ((a - C.a) \% g) a = -1;
            if (a == -1 or C.a == -1) return crt(-1, 0);
4f9
            T lcm = m/g*C.m;
d09
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
1f2
        }
0d9 };
6f2 pair<11, 11> divide_show(11 n, int p, int k, int pak) {
        if (n == 0) return {0, 1};
d02
        11 blocos = n/pak, falta = n%pak;
        ll periodo = divi[pak], resto = divi[falta];
2ce
616
        ll r = expo(periodo, blocos, pak)*resto%pak;
445
        auto rec = divide_show(n/p, p, k, pak);
a51
        11 y = n/p + rec.first;
```

```
bb9
        r = r*rec.second % pak;
90f
        return {y, r};
533 }
6e6 ll solve_pak(ll n, ll x, int p, int k, int pak) {
d34
        divi[0] = 1:
        for (int i = 1; i <= pak; i++) {</pre>
f2b
            divi[i] = divi[i-1];
901
            if (i%p) divi[i] = divi[i] * i % pak;
840
51a
        }
        auto dn = divide_show(n, p, k, pak), dx = divide_show(x, p, k,
   pak),
162
             dnx = divide\_show(n-x, p, k, pak);
        11 y = dn.first-dx.first-dnx.first, r =
768
            (dn.second*inv(dx.second, pak)%pak)*inv(dnx.second,
b64
   pak)%pak;
035
        return expo(p, y, pak) * r % pak;
d78 }
9dd ll solve(ll n, ll x, int mod) {
        vector < pair < int , int >> f;
490
        int mod2 = mod;
c3b
7b4
        for (int i = 2; i*i <= mod2; i++) if (mod2%i==0) {</pre>
aff
            int c = 0:
75b
            while (mod2\%i==0) mod2 /= i, c++;
2a1
            f.push_back({i, c});
fe7
Off
        if (mod2 > 1) f.push_back({mod2, 1});
e96
        crt ans(0, 1);
a13
        for (int i = 0; i < f.size(); i++) {</pre>
702
            int pak = 1;
7e4
            for (int j = 0; j < f[i].second; j++) pak *= f[i].first;</pre>
            ans = ans * crt(solve_pak(n, x, f[i].first, f[i].second,
   pak), pak);
7fd
5fb
        return ans.a;
689 }
4.5 Closest pair of points
// O(nlogn)
915 pair <pt, pt > closest_pair_of_points(vector <pt > v) {
3d2
        int n = v.size();
        sort(v.begin(), v.end());
```

fca

```
31 c
        for (int i = 1; i < n; i++) if (v[i] == v[i-1]) return</pre>
   {v[i-1], v[i]};
        auto cmp_y = [&](const pt &1, const pt &r) {
c20
b53
            if (1.y != r.y) return 1.y < r.y;
            return 1.x < r.x;</pre>
920
55a
        };
62e
        set < pt, decltype(cmp_y) > s(cmp_y);
3d9
        int 1 = 0, r = -1;
6a2
        11 d2_min = numeric_limits<11>::max();
4d5
        pt pl, pr;
bd1
        const int magic = 5;
a55
        while (r+1 < n) {
7f1
            auto it = s.insert(v[++r]).first:
c92
            int cnt = magic/2;
773
            while (cnt-- and it != s.begin()) it--;
a01
d68
            while (cnt++ < magic and it != s.end()) {</pre>
f19
                 if (!((*it) == v[r])) {
67e
                     11 d2 = dist2(*it, v[r]);
74e
                     if (d2_min > d2) {
                         d2_min = d2;
229
841
                         pl = *it;
4f2
                         pr = v[r];
7d9
                     }
10a
                 }
40d
                it++;
801
eb0
            while (1 < r \text{ and } sq(v[1].x-v[r].x) > d2_min)
   s.erase(v[1++]);
de6
c74
        return {pl, pr};
f90 }
4.6 Coloração de Grafo de Intervalo
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
// O(n log(n))
```

615 vector<int> coloring(vector<pair<int, int>>& v) {

vector<pair<int, pair<int, int>>> ev;

ev.push_back({v[i].first, {1, i}});

for (int i = 0; i < n; i++) {</pre>

int n = v.size();

3d2

c08

603

150

```
cda
            ev.push_back({v[i].second, {0, i}});
6a4
49e
        sort(ev.begin(), ev.end());
        vector < int > ans(n), avl(n);
360
265
        for (int i = 0; i < n; i++) avl.push_back(n-i);</pre>
4bf
        for (auto i : ev) {
cbe
            if (i.second.first == 1) {
021
                ans[i.second.second] = avl.back();
a00
                avl.pop_back();
            } else avl.push_back(ans[i.second.second]);
e98
3a6
ba7
        return ans:
83a }
```

4.7 Conectividade Dinamica DC

```
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
8f2 typedef pair <int, int > T;
1cd namespace data {
553
        int n, ans;
573
        int p[MAX], sz[MAX];
        stack<int> S;
ee6
        void build(int n2) {
e5b
1e3
            n = n2:
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
8a6
0b2
            ans = n;
        }
cba
1b1
        int find(int k) {
006
            while (p[k] != k) k = p[k];
839
            return k;
        }
c1e
072
        void add(T x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
           if (a == b) return S.push(-1);
843
e7d
            ans - -;
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
582
            sz[b] += sz[a];
84b
            p[a] = b;
e1a
        }
5eb
        int query() {
```

```
ba7
            return ans;
        }
35c
5cf
        void rollback() {
465
            int u = S.top(); S.pop();
61c
            if (u == -1) return:
270
            sz[p[u]] -= sz[u];
546
            p[u] = u:
Odf
            ans++;
456
        }
568 };
357 int ponta[MAX]; // outra ponta do intervalo ou -1 se for query
4f0 int ans[MAX], n, q;
487 T qu[MAX];
47b void solve(int l = 0, int r = q-1) {
0b1
        if (1 >= r) {
            ans[1] = data::query(); // agora a estrutura ta certa
8c0
505
            return;
        }
f77
962
        int m = (1+r)/2, qnt = 1;
fc7
        for (int i = m+1; i <= r; i++) if (ponta[i]+1 and ponta[i] < 1)</pre>
37d
            data::add(qu[i]), qnt++;
221
        solve(1, m);
593
        while (--qnt) data::rollback();
        for (int i = 1; i <= m; i++) if (ponta[i]+1 and ponta[i] > r)
a2c
37d
            data::add(qu[i]), qnt++;
37b
        solve(m+1, r):
        while (qnt--) data::rollback();
281
0d4 }
4.8 Conectividade Dinamica LCT
// Offline com link-cut trees
// O(n log(n))
1ef namespace lct {
3c9
        struct node {
            int p, ch[2];
19f
a2a
            int val, sub;
aa6
            bool rev;
f93
            node() {}
            node(int v) : p(-1), val(v), sub(v), rev(0) { ch[0] = }
   ch[1] = -1; }
        };
cac
c53
        node t[2*MAX]; // MAXN + MAXQ
```

```
99e
        map<pair<int, int>, int> aresta;
e4d
        int sz:
95a
        void prop(int x) {
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
50e
693
            t[x].rev = 0;
750
564
        void update(int x) {
e8d
            t[x].sub = t[x].val:
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
78d
                t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
3e4
            }
       }
9bf
971
        bool is_root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
        void rotate(int x) {
ed6
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
8fa
            t[x].p = pp, t[p].p = x;
444
            update(p), update(x);
       }
f31
238
        int splay(int x) {
            while (!is_root(x)) {
18c
497
                int p = t[x].p. pp = t[p].p:
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
0 c 5
   == x) ? x : p);
64f
                rotate(x):
72c
aab
            return prop(x), x;
08f
        int access(int v) {
f16
0eb
            int last = -1:
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
```

```
3d3
            return last;
        }
294
952
        void make_tree(int v, int w=INF) { t[v] = node(w); }
82f
        bool conn(int v, int w) {
2cf
            access(v). access(w):
b9b
            return v == w ? true : t[v].p != -1;
ec0
        }
277
        void rootify(int v) {
5e3
            access(v);
            t[v].rev ^= 1;
a02
a05
        }
        int query(int v, int w) {
a1d
b54
            rootifv(w). access(v):
249
            return t[v].sub;
c28
        }
204
        void link_(int v, int w) {
821
            rootify(w);
389
            t[w].p = v;
523
        void link(int v, int w, int x) { // v--w com peso x
6b8
379
            int id = MAX + sz++;
110
            aresta[make_pair(v, w)] = id;
ab6
            make_tree(id, x);
c88
            link_(v, id), link_(id, w);
984
        }
e63
        void cut_(int v, int w) {
            rootify(w), access(v);
b54
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
7cd
031
        void cut(int v, int w) {
            int id = aresta[make_pair(v, w)];
b0f
a4a
            cut_(v, id), cut_(id, w);
        }
840
0d7 }
893 void dyn_conn() {
        int n, q; cin >> n >> q;
c5f
d6e
        vector <int > p(2*q, -1); // outra ponta do intervalo
b4f
        for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
fbf
        vector<pair<int, int>> qu(q);
139
        map<pair<int, int>, int> m;
        for (int i = 0; i < q; i++) {</pre>
abf
3c2
            char c; cin >> c;
ef6
            if (c == '?') continue;
602
            int a, b; cin >> a >> b; a--, b--;
d11
            if (a > b) swap(a, b);
8a1
            qu[i] = \{a, b\};
```

```
8d7
            if (c == '+') {
                p[i] = i+q, p[i+q] = i;
94b
906
                m[make_pair(a, b)] = i;
8a0
            } else {
                int j = m[make_pair(a, b)];
412
                p[i] = j, p[j] = i;
ac2
            }
0da
9e5
        }
447
        int ans = n;
        for (int i = 0; i < q; i++) {</pre>
abf
87d
            if (p[i] == -1) {
886
                cout << ans << endl: // numero de comp conexos</pre>
5e2
            }
b35
            int a = qu[i].first, b = qu[i].second;
69d
            if (p[i] > i) { // +
c4d
ac5
                if (lct::conn(a, b)) {
                    int mi = lct::query(a, b);
18f
                    if (p[i] < mi) {</pre>
993
dd3
                         p[p[i]] = p[i];
5e2
                         continue;
474
                    lct::cut(qu[p[mi]].first, qu[p[mi]].second), ans++;
6f7
                    p[mi] = mi;
6ea
9a9
d1d
                lct::link(a, b, p[i]), ans--;
            } else if (p[i] != i) lct::cut(a, b), ans++; // -
9d0
c03
       }
56a }
```

4.9 Conj. Indep. Maximo com Peso em Grafo de Intervalo

```
// Retorna os indices ordenados dos intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento total
// O(n log(n))
31e vector<int> ind_set(vector<tuple<int, int, int>>& v) {
        vector<tuple<int, int, int>> w;
b27
f14
        for (int i = 0: i < v.size(): i++) {</pre>
            w.push_back(tuple(get<0>(v[i]), 0, i));
e85
6f0
            w.push_back(tuple(get<1>(v[i]), 1, i));
17f
d1d
        sort(w.begin(), w.end());
844
        vector < int > nxt(v.size());
c22
        vector<pair<ll, int>> dp(v.size());
```

```
0eb
        int last = -1;
723
        for (auto [fim, t, i] : w) {
25a
            if (t == 0) {
                nxt[i] = last;
4ca
5e2
                continue;
5fd
78b
            dp[i] = \{0, 0\}:
            if (last != -1) dp[i] = max(dp[i], dp[last]);
cb8
911
            pair<11, int> pega = {get<2>(v[i]), -(get<1>(v[i]) -
   get<0>(v[i]) + 1)};
5d3
            if (nxt[i] != -1) pega.first += dp[nxt[i]].first,
   pega.second += dp[nxt[i]].second;
b08
            if (pega > dp[i]) dp[i] = pega;
7cb
            else nxt[i] = last;
381
            last = i;
        }
b7c
977
        pair<11, int > ans = \{0, 0\};
919
        int idx = -1:
        for (int i = 0; i < v.size(); i++) if (dp[i] > ans) ans =
ceb
   dp[i], idx = i;
4b8
        vector<int> ret;
fdd
        while (idx != -1) {
d69
            if (get < 2 > (v[idx]) > 0 and
                (nxt[idx] == -1 or get<1>(v[nxt[idx]]) <</pre>
a05
   get<0>(v[idx]))) ret.push_back(idx);
e4f
            idx = nxt[idx]:
042
        }
0ea
        sort(ret.begin(), ret.end());
        return ret;
edf
c4d }
```

4.10 Convex Hull Dinamico

```
// insert - O(log n) amortizado
// is_inside - O(log n)
0b9 struct upper {
af8
        set <pt> se;
80b
        set <pt>::iterator it;
25 c
        int is_under(pt p) { // 1 -> inside ; 2 -> border
fe0
            it = se.lower_bound(p);
            if (it == se.end()) return 0:
633
a94
            if (it == se.begin()) return p == *it ? 2 : 0;
ca0
            if (ccw(p, *it, *prev(it))) return 1;
402
            return ccw(p, *prev(it), *it) ? 0 : 2;
dba
        }
```

```
eaa
        void insert(pt p) {
712
            if (is_under(p)) return;
            if (it != se.end()) while (next(it) != se.end() and
   !ccw(*next(it), *it, p))
                it = se.erase(it);
316
be3
            if (it != se.begin()) while (--it != se.begin() and
   !ccw(p, *it, *prev(it)))
                it = se.erase(it);
316
0c8
            se.insert(p);
5da
       }
750 }:
06f struct dyn_hull {
d93
        upper U, L;
333
        int is_inside(pt p) {
632
            int u = U.is\_under(p), l = L.is\_under({-p.x, -p.y});
4 c.c
            if (!u or !1) return 0;
fc0
            return max(u, 1);
478
        void insert(pt p) {
86c
            U.insert(p);
925
            L.insert({-p.x, -p.y});
64b
285
        int size() {
7c2
            int ans = U.se.size() + L.se.size();
1c9
            return ans <= 2 ? ans/2 : ans-2;</pre>
ad5
        }
65e }:
4.11 Distancia maxima entre dois pontos
```

```
// \max_{dist2(v)} - O(n \log(n))
// max_dist_manhattan - O(n)
// Quadrado da Distancia Euclidiana (precisa copiar convex_hull, ccw e
   pt)
859 ll max_dist2(vector<pt> v) {
        v = convex_hull(v);
a14
       if (v.size() \le 2) return dist2(v[0], v[1%v.size()]):
       11 ans = 0:
04b
323
       int n = v.size(), j = 0;
603
        for (int i = 0: i < n: i++) {</pre>
057
            while (!ccw(v[(i+1)%n]-v[i], pt(0, 0), v[(i+1)%n]-v[i])) j
   = (j+1) \%n;
```

```
e7a
            ans = \max(\{ans, dist2(v[i], v[j]), dist2(v[(i+1)%n],
   v[i])});
        }
1f6
ba7
        return ans;
bda }
// Distancia de Manhattan
c51 template < typename T> T max_dist_manhattan(vector < pair < T, T>> v) {
8eb
        T min_sum, max_sum, min_dif, max_dif;
4f5
        min sum = max sum = v[0].first + v[0].second:
271
        min_dif = max_dif = v[0].first - v[0].second;
c25
        for (auto [x, y] : v) {
1cb
            min sum = min(min sum, x+v):
683
            max_sum = max(max_sum, x+y);
782
            min_dif = min(min_dif, x-y);
            max_dif = max(max_dif, x-y);
af7
e3a
9f0
        return max(max_sum - min_sum, max_dif - min_dif);
4e9 }
     Distinct Range Query
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
789 namespace perseg { };
53d int qt[MAX];
edc void build(vector<int>& v) {
        int n = v.size();
16b
        perseg::build(n);
663
        map<int, int> last;
05e
        int at = 0;
603
        for (int i = 0; i < n; i++) {</pre>
817
            if (last.count(v[i])) {
a58
                perseg::update(last[v[i]], -1);
69a
                at++;
d1f
            }
4f2
            perseg::update(i, 1);
460
            qt[i] = ++at;
efe
            last[v[i]] = i:
d6f
        }
0f4 }
9e3 int query(int 1, int r) {
080
        return perseg::query(1, r, qt[r]);
```

4.13 Distinct Range Query com Update

```
// build - O(n log(n))
// query - O(log^2(n))
// update - O(log^2(n))
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
def
        tree_order_statistics_node_update>;
042 int v[MAX], n, nxt[MAX], prv[MAX];
f60 map<int, set<int> > ocor;
e04 namespace bit {
        ord_set<pair<int, int>> bit[MAX];
0a8
        void build() {
            for (int i = 1; i <= n; i++) bit[i].insert({nxt[i-1],</pre>
3e1
   i-1});
78a
            for (int i = 1; i <= n; i++) {</pre>
                int j = i + (i\&-i);
edf
                if (j <= n) for (auto x : bit[i]) bit[j].insert(x);</pre>
d03
            }
5cb
af6
        }
d3f
        int pref(int p, int x) {
7c9
            int ret = 0;
bbf
            for (; p; p -= p\&-p) ret += bit[p].order_of_key({x, -INF});
edf
            return ret;
0e1
        }
d50
        int query(int 1, int r, int x) {
e55
            return pref(r+1, x) - pref(1, x);
9b4
ff3
        void update(int p, int x) {
f17
            int p2 = p;
            for (p++; p \le n; p += p\&-p) {
5ed
                bit[p].erase({nxt[p2], p2});
ca8
f6b
                bit[p].insert({x, p2});
            }
3df
151
        }
c63 }
0a8 void build() {
```

```
383
        for (int i = 0; i < n; i++) nxt[i] = INF;</pre>
        for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
7b3
d07
        vector < pair < int , int >> t;
        for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
348
        sort(t.begin(), t.end());
3fd
        for (int i = 0; i < n; i++) {</pre>
603
b40
            if (i and t[i].first == t[i-1].first)
565
                 prv[t[i].second] = t[i-1].second;
            if (i+1 < n \text{ and } t[i].first == t[i+1].first)
a8b
                 nxt[t[i].second] = t[i+1].second;
12f
48d
        }
a23
        for (int i = 0: i < n: i++) ocor[v[i]].insert(i):
1d7
        bit::build();
d44 }
aae void muda(int p, int x) {
f92
        bit::update(p, x);
c3d
        nxt[p] = x:
97c }
4ea int query(int a, int b) {
a0a
        return b-a+1 - bit::query(a, b, b+1);
511 }
ff3 void update(int p, int x) { // mudar valor na pos. p para x
c0b
        if (prv[p] > -INF) muda(prv[p], nxt[p]);
        if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
4ae
5bf
        ocor[v[p]].erase(p);
4b4
        if (!ocor[x].size()) {
19d
            muda(p, INF);
8d4
            prv[p] = -INF:
        } else if (*ocor[x].rbegin() < p) {</pre>
f6c
5b5
            int i = *ocor[x].rbegin();
f64
            prv[p] = i;
19d
            muda(p, INF);
5f2
            muda(i, p);
f36
        } else {
            int i = *ocor[x].lower_bound(p);
d46
33f
            if (prv[i] > -INF) {
f17
                 muda(prv[i], p);
8f9
                 prv[p] = prv[i];
bc4
            } else prv[p] = -INF;
523
            prv[i] = p;
597
            muda(p, i);
```

```
2ac     }
c96     v[p] = x; ocor[x].insert(p);
38e }
```

4.14 Dominator Points

```
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
// o multiset 'quina' guarda informacao sobre os pontos
// nao dominados por um elemento do conjunto que nao dominam
// outro ponto nao dominado por um elemento do conjunto
// No caso, armazena os valores de x+y esses pontos
//
// Complexidades:
// is_dominated - O(log(n))
// insert - O(log(n)) amortizado
// query - O(1)
e2a struct dominator_points {
        set < pair < int , int >> se;
baf
        multiset < int > quina;
4dd
        bool is_dominated(pair<int, int> p) {
a85
            auto it = se.lower_bound(p);
80f
            if (it == se.end()) return 0;
633
            return it->second >= p.second;
ab4
28f
99Ъ
        void mid(pair<int, int> a, pair<int, int> b, bool rem) {
29a
            pair < int , int > m = {a.first+1, b.second+1};
b19
            int val = m.first + m.second;
638
            if (!rem) quina.insert(val);
731
            else quina.erase(quina.find(val));
241
7c4
        bool insert(pair<int, int> p) {
            if (is_dominated(p)) return 0;
fb4
            auto it = se.lower_bound(p);
80f
            if (it != se.begin() and it != se.end())
ca9
d4a
                mid(*prev(it), *it, 1);
            while (it != se.begin()) {
1fa
049
                it--;
23 c
                if (it->second > p.second) break;
b86
                if (it != se.begin()) mid(*prev(it), *it, 1);
316
                it = se.erase(it):
```

```
acd
433
            it = se.insert(p).first;
69e
            if (it != se.begin()) mid(*prev(it), *it, 0);
96d
            if (next(it) != se.end()) mid(*it, *next(it), 0);
6a5
            return 1:
688
        }
5eb
        int query() {
956
            if (!quina.size()) return INF;
add
            return *quina.begin();
h8b
        }
09f };
4.15 DP de Dominação 3D
// Computa para todo ponto i,
// dp[i] = 1 + max_{i} dominado por i dp[i]
// em que ser dominado eh ter as 3 coordenadas menores
// Da pra adaptar facil para outras dps
//
// O(n log^2 n), O(n) de memoria
c53 void lis2d(vector<vector<tuple<int, int, int>>>& v, vector<int>&
    dp, int 1, int r) {
893
        if (1 == r) {
56f
            for (int i = 0; i < v[1].size(); i++) {</pre>
                 int ii = get<2>(v[1][i]);
8b5
1ce
                 dp[ii] = max(dp[ii], 1);
4b0
505
            return;
3e4
        }
        int m = (1+r)/2;
ee4
62b
        lis2d(v, dp, 1, m);
325
        vector<tuple<int, int, int>> vv[2];
d44
        vector < int > Z;
871
        for (int i = 1; i <= r; i++) for (auto it : v[i]) {</pre>
2ef
            vv[i > m].push_back(it);
042
            Z.push_back(get<1>(it));
        }
0d1
        sort(vv[0].begin(), vv[0].end());
e9f
9b5
        sort(vv[1].begin(), vv[1].end());
0d1
        sort(Z.begin(), Z.end());
        auto get_z = [&](int z) { return lower_bound(Z.begin(),
    Z.end(), z) - Z.begin(); };
        vector < int > bit(Z.size());
c51
181
        int i = 0;
```

```
e9a
        for (auto [y, z, id] : vv[1]) {
             while (i < vv[0].size() and get<0>(vv[0][i]) < y) {</pre>
6bd
397
                 auto [y2, z2, id2] = vv[0][i++];
                 for (int p = get_z(z_2)+1; p \le Z.size(); p += p\&-p)
ea0
                     bit[p-1] = max(bit[p-1], dp[id2]);
300
            }
82c
d3b
            int q = 0;
            for (int p = get_z(z); p; p -= p&-p) q = max(q, bit[p-1]);
fd9
614
             dp[id] = max(dp[id], q + 1);
acc
c25
        lis2d(v, dp, m+1, r);
4d6 }
4de vector<int> solve(vector<tuple<int, int, int>> v) {
3d2
        int n = v.size();
        vector<tuple<int, int, int, int>> vv;
cd4
603
        for (int i = 0; i < n; i++) {</pre>
             auto [x, y, z] = v[i];
9be
             vv.emplace_back(x, y, z, i);
5bb
64c
        sort(vv.begin(), vv.end());
bd3
e11
        vector < vector < tuple < int , int , int >>> V;
603
        for (int i = 0; i < n; i++) {</pre>
a5b
            int j = i;
808
            V.emplace_back();
c01
             while (j < n and get <0>(vv[j]) == get <0>(vv[i])) {
ba6
                 auto [x, y, z, id] = vv[j++];
                 V.back().emplace_back(y, z, id);
cbb
8bd
            }
452
            i = j-1;
ac4
388
        vector < int > dp(n);
839
        lis2d(V. dp. 0. V.size()-1):
898
        return dp;
bOa }
4.16 Gray Code
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// O(2<sup>n</sup>)
df6 vector<int> gray_code(int n) {
73f
        vector < int > ret(1 << n);</pre>
```

for (int i = 0; i < (1<<n); i++) ret[i] = i^(i>>1);

f29

```
edf
                         return ret;
 840 }
 4.17 Half-plane intersection
 // Cada half-plane eh identificado por uma reta e a regiao ccw a ela
 // O(n log n)
  f4f vector <pt> hp_intersection(vector <line> &v) {
                         deque < pt > dq = \{\{INF, INF\}, \{-INF, INF\}, \{-INF, -INF\}, \{INF, -INF
            -INF}};
 d41 #warning considerar trocar por compare_angle
  de3
                         sort(v.begin(), v.end(), [&](line r, line s) { return
           angle(r.q-r.p) < angle(s.q-s.p); \});
                         for(int i = 0; i < v.size() and dq.size() > 1; i++) {
 5e9
 c69
                                    pt p1 = dq.front(), p2 = dq.back();
 6c6
                                    while (dq.size() and !ccw(v[i].p, v[i].q, dq.back()))
 47b
                                               p1 = dq.back(), dq.pop_back();
                                    while (dq.size() and !ccw(v[i].p, v[i].q, dq.front()))
 0a2
 7 cf
                                               p2 = dq.front(), dq.pop_front();
 4d9
                                    if (!dq.size()) break;
 606
                                    if (p1 == dq.front() and p2 == dq.back()) continue;
                                    dq.push_back(inter(v[i], line(dq.back(), p1)));
  c9b
 65 c
                                    dq.push_front(inter(v[i], line(dq.front(), p2)));
 fdd
                                    if (dq.size() > 1 and dq.back() == dq.front())
           dq.pop_back();
 4d8
 b2b
                         return vector < pt > (dq.begin(), dq.end());
 f56 }
 4.18 Heap Sort
 // O(n log n)
 f18 void down(vector<int>& v, int n, int i) {
                         while ((i = 2*i+1) < n) {
 583
                                    if (i+1 < n and v[i] < v[i+1]) i++;</pre>
 b27
                                    if (v[i] < v[(i-1)/2]) break;
 322
                                    swap(v[i], v[(i-1)/2]);
 170
                        }
 724 }
eb6 void heap_sort(vector<int>& v) {
```

```
3d2
        int n = v.size();
                                                                              4c5
        for (int i = n/2-1; i \ge 0; i--) down(v, n, i);
61d
                                                                              0d7
917
        for (int i = n-1; i > 0; i--)
                                                                              6d4
37f
             swap(v[0], v[i]), down(v, i, 0);
                                                                              886
b33 }
                                                                              38d
                                                                              306
                                                                              6db
4.19 Hungaro
                                                                              da3
                                                                              979
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)</pre>
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
//
// O(n^3)
                                                                              //
a6a template < typename T > struct hungarian {
1a8
        int n:
        vector < vector < T >> a;
a08
f36
        vector <T> u. v:
                                                                              bb6
5ff
        vector < int > p, way;
                                                                              796
f1e
        T inf:
                                                                              1bc
                                                                              dfb
c3f
        hungarian(int n_{-}): n(n_{-}), u(n+1), v(n+1), p(n+1), way(n+1) {
                                                                              874
b2f
            a = vector < vector < T >> (n, vector < T > (n));
                                                                              8c0
1f3
            inf = numeric_limits <T>::max();
                                                                              4e9
78f
        }
                                                                              61 c
d67
        pair <T. vector <int >> assignment() {
                                                                              d1d
             for (int i = 1; i <= n; i++) {</pre>
78a
                                                                              603
8c9
                 p[0] = i;
                                                                              bf3
625
                 int j0 = 0;
ce7
                 vector <T> minv(n+1, inf);
                                                                              1bf
241
                 vector<int> used(n+1, 0);
016
                 do {
                                                                              962
472
                     used[j0] = true;
                                                                              6c0
d24
                     int i0 = p[j0], j1 = -1;
                                                                              964
7e5
                     T delta = inf;
9ac
                     for (int j = 1; j <= n; j++) if (!used[j]) {</pre>
                                                                              04b
                         T cur = a[i0-1][j-1] - u[i0] - v[j];
7bf
                                                                              45b
                         if (cur < minv[j]) minv[j] = cur, wav[j] = j0;</pre>
9f2
                                                                              2d9
                         if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
821
                                                                              3a1
                     }
4d1
                                                                              ebe
f63
                     for (int j = 0; j <= n; j++)
                                                                              ba7
2c5
                         if (used[j]) u[p[j]] += delta, v[j] -= delta;
                         else minv[j] -= delta;
6ec
6d4
                     j0 = j1;
f4f
                 } while (p[i0] != 0);
                                                                              4.21 LIS - recupera
016
                 do {
```

```
int j1 = way[j0];
                    p[j0] = p[j1];
                     j0 = j1;
                } while (i0);
            vector < int > ans(n);
            for (int j = 1; j \le n; j++) ans [p[j]-1] = j-1;
            return make_pair(-v[0], ans);
        }
64c };
4.20 Inversion Count
// Computa o numero de inversoes para transformar
// l em r (se nao tem como, retorna -1)
// O(n log(n))
37b template < typename T > 11 inv_count(vector < T > 1, vector < T > r = {}) {
        if (!r.size()) {
            r = 1:
            sort(r.begin(), r.end());
        }
        int n = 1.size();
        vector < int > v(n), bit(n):
        vector < pair < T, int >> w;
        for (int i = 0; i < n; i++) w.push_back({r[i], i+1});</pre>
        sort(w.begin(), w.end());
        for (int i = 0; i < n; i++) {</pre>
            auto it = lower_bound(w.begin(), w.end(), make_pair(l[i],
   0));
            if (it == w.end() or it->first != l[i]) return -1; // nao
   da
            v[i] = it->second;
            it->second = -1;
        }
        11 \text{ ans} = 0:
        for (int i = n-1; i >= 0; i--) {
            for (int j = v[i]-1; j; j = j\&-j) ans += bit[j];
            for (int j = v[i]; j < n; j += j\&-j) bit[j]++;
        }
        return ans;
eef }
```

```
// Calcula e retorna uma LIS
// O(n.log(n))
121 template < typename T> vector < T> lis(vector < T>& v) {
        int n = v.size(), m = -1;
f0c
        vector <T> d(n+1. INF):
        vector < int > l(n);
aec
       d[0] = -INF;
007
603
        for (int i = 0; i < n; i++) {</pre>
            // Para non-decreasing use upper_bound()
4fd
            int t = lower_bound(d.begin(), d.end(), v[i]) - d.begin();
            d[t] = v[i], l[i] = t, m = max(m, t);
3ad
89 c
4ff
        int p = n;
5a9
        vector <T> ret:
        while (p--) if (1[p] == m) {
cdf
883
            ret.push_back(v[p]);
76b
            m - - ;
f83
        reverse(ret.begin(),ret.end());
969
edf
        return ret:
474 }
4.22 LIS - tamanho
// Calcula o tamanho da LIS
//
// O(n log(n))
84b template < typename T > int lis(vector < T > &v){
        vector <T> ans;
2da
5e0
        for (T t : v){
            // Para non-decreasing use upper_bound()
            auto it = lower_bound(ans.begin(), ans.end(), t);
fe6
d7f
            if (it == ans.end()) ans.push_back(t);
b94
            else *it = t;
1f5
1eb
        return ans.size();
402 }
      Minimum Enclosing Circle
```

// O(n) com alta probabilidade

```
254
        pt operator * (double c) const { return pt(x*c, y*c); }
701
        pt operator / (double c) const { return pt(x/c, y/c); }
54d }:
2f9 double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
dd5 double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
e7c double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
3f4 pt center(pt p, pt q, pt r) {
5d9
        pt a = p-r, b = q-r;
        pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
        return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x), c)) /
   cross(a, b):
fc8 }
aa8 struct circle {
f41
        pt cen;
c12
        double r:
898
        circle(pt cen_, double r_) : cen(cen_), r(r_) {}
83 c
        circle(pt a, pt b, pt c) {
            cen = center(a, b, c);
13d
1f1
            r = dist(cen, a);
bc1
        }
cd5
        bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
2a6 };
806 circle minCirc(vector<pt> v) {
f21
        shuffle(v.begin(), v.end(), rng);
ae0
        circle ret = circle(pt(0, 0), 0);
618
        for (int i = 0; i < v.size(); i++) if (!ret.inside(v[i])) {</pre>
            ret = circle(v[i], 0);
16a
f11
            for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
881
                ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
                for (int k = 0; k < j; k++) if (!ret.inside(v[k]))
b8c
43f
                    ret = circle(v[i], v[j], v[k]);
5f8
            }
        }
6a1
```

chrono::steady_clock::now().time_since_epoch().count());

 $pt(double x_{=} = 0, double y_{=} = 0) : x(x_{=}), y(y_{=}) {}$

pt operator + (const pt& p) const { return pt(x+p.x, y+p.y); }

pt operator - (const pt& p) const { return pt(x-p.x, y-p.y); }

22c const double EPS = 1e-12;

double x, y;

878 mt19937 rng((int)

b2a struct pt {

662

be7

7af

b23

```
edf
        return ret;
eba }
4.24 Minkowski Sum
// Computa A+B = \{a+b : a \setminus A, b \setminus B\}, em que
// A e B sao poligonos convexos
// A+B eh um poligono convexo com no max |A|+|B| pontos
//
// O(|A|+|B|)
539 vector<pt> minkowski(vector<pt> p, vector<pt> q) {
        auto fix = [](vector < pt > & P) {
            rotate(P.begin(), min_element(P.begin(), P.end()),
515
    P.end());
018
            P.push_back(P[0]), P.push_back(P[1]);
f24
        };
889
        fix(p), fix(q);
        vector < pt > ret;
8af
        int i = 0, j = 0;
692
2ee
        while (i < p.size()-2 \text{ or } j < q.size()-2) {
898
            ret.push_back(p[i] + q[j]);
            auto c = ((p[i+1] - p[i]) ^ (q[i+1] - q[i]));
732
            if (c >= 0) i = min<int>(i+1, p.size()-2);
ebc
81e
            if (c \le 0) j = min \le int > (j+1, q.size()-2);
9ff
edf
        return ret;
d7c }
c3e ld dist_convex(vector<pt> p, vector<pt> q) {
dc2
        for (pt& i : p) i = i * -1;
        auto s = minkowski(p, q);
44c
        if (inpol(s, pt(0, 0))) return 0;
95d
        return 1:
6a5
        ld ans = DINF;
921
073
        for (int i = 0; i < s.size(); i++) ans = min(ans,
f04
                 disttoseg(pt(0, 0), line(s[(i+1)%s.size()], s[i])));
ba7
        return ans;
2f5 }
4.25 MO
// Para ter o bound abaixo, escolher
// SQ = n / sqrt(q)
//
// O(n * sqrt(q))
```

```
0d2 const int MAX = 1e5+10;
6ff const int SQ = sqrt(MAX);
b69 int v[MAX];
b65 int ans, freq[MAX];
9da inline void insert(int p) {
        int o = v[p];
591
        freq[o]++;
        ans += (freq[o] == 1);
21d }
a25 inline void erase(int p) {
ae0
        int o = v[p];
        ans -= (freq[o] == 1);
7ee
        freq[o]--;
ba2
dc7 }
e51 inline ll hilbert(int x, int y) {
        static int N = 1 << (__builtin_clz(0) - __builtin_clz(MAX));</pre>
100
        int rx, ry, s;
b72
        11 d = 0:
        for (s = N/2; s > 0; s /= 2) {
43b
c95
            rx = (x \& s) > 0, ry = (y \& s) > 0;
e3e
            d += s * 11(s) * ((3 * rx) ^ ry);
d2e
            if (rv == 0) {
5aa
                if (rx == 1) x = N-1 - x, y = N-1 - y;
9dd
                swap(x, y);
e2d
            }
888
        }
        return d;
be2
7fa }
bac #define HILBERT true
617 vector<int> MO(vector<pair<int, int>> &q) {
c3b
        ans = 0:
        int m = q.size();
c23
3f8
        vector < int > ord(m);
be8
        iota(ord.begin(), ord.end(), 0);
6a6 #if HILBERT
        vector<ll> h(m):
        for (int i = 0; i < m; i++) h[i] = hilbert(q[i].first,</pre>
74c
   q[i].second):
        sort(ord.begin(), ord.end(), [&](int 1, int r) { return h[1] <</pre>
   h[r]; });
8c1 #else
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
```

```
9c9
            if (q[1].first / SQ != q[r].first / SQ) return q[1].first
   < q[r].first;
            if ((q[1].first / SQ) % 2) return q[1].second >
   q[r].second:
            return q[1].second < q[r].second;</pre>
a66
        }):
bec
f2e #endif
435
        vector < int > ret(m):
3d9
        int 1 = 0, r = -1;
8b0
        for (int i : ord) {
6c6
            int ql, qr;
4f5
            tie(al. ar) = a[i]:
            while (r < qr) insert(++r);</pre>
026
232
            while (1 > q1) insert(--1);
75e
            while (1 < q1) erase(1++);</pre>
fe8
            while (r > qr) erase(r--);
            ret[i] = ans;
381
116
edf
        return ret;
fb7 }
4.26 MO - DSU
// Dado uma lista de arestas de um grafo, responde
// para cada query(1, r), quantos componentes conexos
// o grafo tem se soh considerar as arestas l. l+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura rollbackavel
//
// O(m sqrt(q) log(n))
8d3 struct dsu {
553
        int n, ans;
2e3
        vector < int > p, sz;
ee6
        stack<int> S;
4b8
        dsu(int n_{-}) : n(n_{-}), ans(n), p(n), sz(n) {
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
aae
        }
        int find(int k) {
1b1
006
            while (p[k] != k) k = p[k];
839
            return k:
c1e
        }
553
        void add(pair<int, int> x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
```

```
e7d
             ans --;
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
582
            sz[b] += sz[a];
            p[a] = b;
84b
720
35 c
        int query() { return ans; }
5cf
        void rollback() {
465
            int u = S.top(); S.pop();
61 c
            if (u == -1) return:
270
            sz[p[u]] -= sz[u];
546
            p[u] = u;
Odf
            ans++:
456
        }
9c1 };
1a8 int n;
e93 vector<pair<int, int>> ar;
// 9d242b
617 vector<int> MO(vector<pair<int, int>> &q) {
547
        int SQ = sqrt(q.size()) + 1;
c23
        int m = q.size();
3f8
        vector<int> ord(m);
be8
        iota(ord.begin(), ord.end(), 0);
d01
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
9c9
                 if (q[1].first / SQ != q[r].first / SQ) return
   q[l].first < q[r].first;
                return q[1].second < q[r].second;</pre>
a66
b90
                });
        vector < int > ret(m);
435
3bd
        dsu small(n):
dd5
        for (int i = 0: i < m: i++) {
            auto [1, r] = q[ord[i]];
5ec
acc
            if (1 / SQ == r / SQ) {
00c
                for (int k = 1; k <= r; k++) small.add(ar[k]);</pre>
b99
                ret[ord[i]] = small.query();
64e
                for (int k = 1; k <= r; k++) small.rollback();</pre>
259
            }
        }
6b0
        for (int i = 0; i < m; i++) {</pre>
dd5
176
            dsu D(n):
ae9
            int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
e25
            int last_r = fim;
            int j = i-1;
ebc
```

```
00c
            while (j+1 < m and q[ord[j+1]].first / SQ ==</pre>
   q[ord[i]].first / SQ) {
                auto [1, r] = q[ord[++i]];
                if (1 / SQ == r / SQ) continue;
f58
59b
                while (last r < r) D.add(ar[++last r]);</pre>
                for (int k = 1; k <= fim; k++) D.add(ar[k]);</pre>
2cf
9b2
                ret[ord[j]] = D.query();
572
                for (int k = 1; k <= fim; k++) D.rollback();</pre>
9c8
            }
bdf
            i = j;
e99
        return ret;
edf
9d2 }
4.27 MO em Arvores
// Problema que resolve: https://www.spoj.com/problems/COT2/
//
// Complexidade sendo c = O(update) e SQ = sgrt(n):
// O((n + q) * sqrt(n) * c)
1bc const int MAX = 40010, SQ = 400;
042 vector < int > g[MAX];
c54 namespace LCA { ... }
249 int in[MAX], out[MAX], vtx[2 * MAX];
81b bool on [MAX];
4c3 int dif, freq[MAX];
9e2 vector < int > w;
d9a void dfs(int v, int p, int &t) {
        vtx[t] = v, in[v] = t++;
659
        for (int u : g[v]) if (u != p) {
18e
            dfs(u, v, t);
c53
e0f
        }
        vtx[t] = v, out[v] = t++;
217
42b }
```

e5f void update(int p) { // faca alteracoes aqui

int v = vtx[p];

bbc

```
0ec
        if (not on[v]) { // insere vtx v
31c
             dif += (freq[w[v]] == 0);
b20
             frea[w[v]]++;
        }
cf7
        else { // retira o vertice v
4e6
0a9
             dif \rightarrow (freq[w[v]] == 1);
fd3
             freq[w[v]]--;
2c8
        }
73e
        on[v] = not on[v];
ea9 }
a3a vector < tuple < int , int , int >> build_queries (const vector < pair < int ,
   int>>& a) {
ea6
        LCA::build(0);
        vector<tuple<int, int, int>> ret;
f77
        for (auto [1, r] : q){
aa9
d24
             if (in[r] < in[l]) swap(l, r);
6f9
             int p = LCA::lca(1, r);
826
             int init = (p == 1) ? in[1] : out[1];
07a
             ret.emplace_back(init, in[r], in[p]);
        }
b0e
edf
        return ret;
8e6 }
f31 vector < int > mo_tree(const vector < pair < int , int >> & vq) {
6bb
        int t = 0;
dab
        dfs(0, -1, t);
af1
        auto q = build_queries(vq);
f48
        vector<int> ord(q.size());
be8
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&] (int 1, int r) {
d01
             int bl = get<0>(q[1]) / SQ, br = <math>get<0>(q[r]) / SQ;
d8d
596
             if (bl != br) return bl < br;</pre>
158
             else if (b1 % 2 == 1) return get<1>(q[1]) < get<1>(q[r]);
f1d
             else return get<1>(q[1]) > get<1>(q[r]);
0a8
        });
80e
        memset(freq, 0, sizeof freq);
bf6
        dif = 0:
        vector<int> ret(q.size());
ff2
        int 1 = 0, r = -1;
3d9
8b0
        for (int i : ord) {
             auto [ql, qr, qp] = q[i];
3 c 7
             while (r < qr) update(++r);</pre>
af7
```

```
d6b
             while (1 > q1) update (--1);
             while (1 < q1) update(1++);</pre>
951
             while (r > qr) update(r--);
6a1
            if (qp < 1 \text{ or } qp > r)  { // se LCA estah entre as pontas
3d8
74b
                 update(qp);
2e1
                 ret[i] = dif;
74b
                 update(qp);
e83
0fe
             else ret[i] = dif;
Ofd
edf
        return ret;
48d }
```

4.28 Palindromic Factorization

```
// Precisa da eertree
// Computa o numero de formas de particionar cada
// prefixo da string em strings palindromicas
// O(n log n), considerando alfabeto O(1)
070 struct eertree { ... };
0e7 ll factorization(string s) {
        int n = s.size(), sz = 2;
b19
580
        eertree PT(n):
147
        vector \langle int \rangle diff (n+2), slink (n+2), sans (n+2), dp (n+1);
0ec
        dp[0] = 1;
        for (int i = 1; i <= n; i++) {</pre>
78a
c58
            PT.add(s[i-1]);
a7c
            if (PT.size()+2 > sz) {
6c4
                diff[sz] = PT.len[sz] - PT.len[PT.link[sz]];
241
                if (diff[sz] == diff[PT.link[sz]])
d6f
                     slink[sz] = slink[PT.link[sz]];
f53
                else slink[sz] = PT.link[sz];
eb9
                sz++;
            }
f6a
            for (int v = PT.last; PT.len[v] > 0; v = slink[v]) {
911
                sans[v] = dp[i - (PT.len[slink[v]] + diff[v])];
297
                if (diff[v] == diff[PT.link[v]])
85d
f20
                    sans[v] = (sans[v] + sans[PT.link[v]]) % MOD;
071
                dp[i] = (dp[i] + sans[v]) % MOD;
e5e
            }
fc0
5f0
        return dp[n];
3a7 }
```

4.29 Parsing de Expressao

```
// Operacoes associativas a esquerda por default
// Para mudar isso, colocar em r_assoc
// Operacoes com maior prioridade sao feitas primeiro
cc1 bool blank(char c) {
f34
        return c == ' ';
ec3 }
8e4 bool is_unary(char c) {
        return c == '+' or c == '-';
b6b }
76d bool is_op(char c) {
010
        if (is_unary(c)) return true;
        return c == '*' or c == '/' or c == '+' or c == '-':
31c
4e4 }
fa3 bool r_assoc(char op) {
        // operator unario - deve ser assoc. a direita
cf0
        return op < 0;</pre>
c5c }
79d int priority(char op) {
        // operator unario - deve ter precedencia maior
103
        if (op < 0) return INF;</pre>
        if (op == '*' or op == '/') return 2;
727
        if (op == '+' or op == '-') return 1;
439
daa
        return -1;
966 }
c15 void process_op(stack<int>& st, stack<int>& op) {
        char o = op.top(); op.pop();
88 c
        if (o < 0) {
91c
4e6
            o *= -1:
1e2
            int 1 = st.top(); st.pop();
Off
            if (o == '+') st.push(1);
            if (o == '-') st.push(-1);
7e9
320
        } else {
14c
            int r = st.top(); st.pop();
1e2
            int 1 = st.top(); st.pop();
1e4
            if (o == '*') st.push(1 * r);
f55
            if (o == '/') st.push(1 / r);
605
            if (o == '+') st.push(1 + r);
            if (o == '-') st.push(1 - r);
c40
```

```
0aa
        }
2b2 }
439 int eval(string& s) {
        stack<int> st, op;
212
        bool un = true;
d0c
1 c f
       for (int i = 0: i < s.size(): i++) {
            if (blank(s[i])) continue;
68d
            if (s[i] == '(') {
139
367
                op.push('(');
99d
                un = true:
b88
            } else if (s[i] == ')') {
709
                while (op.top() != '(') process_op(st, op);
75e
                op.pop();
                un = false;
ce2
003
            } else if (is_op(s[i])) {
                char o = s[i]:
4d0
                if (un and is_unary(o)) o *= -1;
37 c
                while (op.size() and (
ae3
                             (!r_assoc(o) and priority(op.top()) >=
   priority(o)) or
c41
                             (r_assoc(o) and priority(op.top()) >
   priority(o))))
                    process_op(st, op);
c47
c00
                op.push(o);
99d
                un = true;
196
            } else {
                int val = 0:
c2b
                while (i < s.size() and isalnum(s[i]))</pre>
                    val = val * 10 + s[i++] - '0':
169
                i--;
25d
                st.push(val);
ce2
                un = false:
            }
442
b19
        }
7f6
        while (op.size()) process_op(st, op);
123
        return st.top();
05c }
     RMQ com Divide and Conquer
// Responde todas as queries em
// O(n log(n))
f74 typedef pair <pair <int, int>, int> iii;
```

```
7c6 #define f first
Oab #define s second
87d int n, q, v[MAX];
e3f iii qu[MAX];
aeb int ans[MAX], pref[MAX], sulf[MAX];
0e3 void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
8a3
        if (1 > r or q1 > qr) return;
        int m = (1+r)/2:
ee4
        int qL = partition(qu+ql, qu+qr+1, [=](iii x){return x.f.s <</pre>
1b1
    m;}) - qu;
        int qR = partition(qu+qL, qu+qr+1, [=](iii x){return x.f.f
eb0
    <=m;}) - qu;
        pref[m] = sulf[m] = v[m]:
3cd
9f9
        for (int i = m-1; i >= 1; i--) pref[i] = min(v[i], pref[i+1]);
        for (int i = m+1; i <= r; i++) sulf[i] = min(v[i], sulf[i-1]);</pre>
ea8
b2a
        for (int i = qL; i < qR; i++)
             ans[qu[i].s] = min(pref[qu[i].f.f], sulf[qu[i].f.s]);
f3a
364
        solve(l, m-1, ql, qL-1), solve(m+1, r, qR, qr);
13e }
4.31 Segment Intersection
// Verifica, dado n segmentos, se existe algum par de segmentos
// que se intersecta
//
// O(n log n)
6e0 bool operator < (const line& a, const line& b) { // comparador pro
191
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
    b.p.x))
780
            return ccw(a.p, a.q, b.p);
        return ccw(a.p, b.q, b.p);
dc0
e36 }
8e2 bool has intersection(vector<line> v) {
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
576
a08
            return interseg(a.first, b.first);
3e6
        }:
e1b
        vector<pair<pt, pair<int, int>>> w;
f 14
        for (int i = 0; i < v.size(); i++) {</pre>
```

```
876
            if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
            w.push_back({v[i].p, {0, i}});
e1d
034
            w.push_back({v[i].q, {1, i}});
220
        sort(w.begin(), w.end());
d1d
7f2
        set < pair < line , int >> se;
e58
       for (auto i : w) {
bfd
            line at = v[i.second.second];
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
                if (nxt != se.end() and intersects(*nxt, {at,
   i.second.second})) return 1;
257
                if (nxt != se.begin() and intersects(*(--nxt), {at,
   i.second.second})) return 1;
78f
                se.insert({at, i.second.second});
08b
            } else {
                auto nxt = se.upper_bound({at, i.second.second}), cur
884
   = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
                    and intersects(*nxt, *(--prev))) return 1;
4fb
                se.erase(cur):
cca
            }
e27
a00
        }
bb3
        return 0;
196 }
```

4.32 Sequencia de de Brujin

```
// Se passar sem o terceiro parametro, gera um vetor com valores
// em [0, k) de tamanho k^n de forma que todos os subarrays ciclicos
// de tamanho n ocorrem exatamente uma vez
// Se passar com um limite lim, gera o menor vetor com valores
// em [0, k) que possui lim subarrays de tamanho n distintos
// (assume que lim <= k^n)</pre>
//
// Linear no tamanho da resposta
860 vector<int> de_brujin(int n, int k, int lim = INF) {
        if (k == 1) return vector < int > (lim == INF ? 1 : n, 0);
b55
        vector<int> 1 = {0}, ret; // 1 eh lyndon word
5f6
        while (true) {
667
c86
            if (1.size() == 0) {
1 b 9
                if (lim == INF) break;
daf
                1.push_back(0);
            }
bae
686
            if (n % 1.size() == 0) for (int i : 1) {
728
                ret.push_back(i);
```

```
c99
                if (ret.size() == n+lim-1) return ret;
            }
56e
630
            int p = 1.size();
905
            while (1.size() < n) 1.push_back(1[1.size()%p]);</pre>
            while (l.size() and l.back() == k-1) l.pop_back();
e7f
            if (1.size()) 1.back()++;
88a
2ef
        }
edf
        return ret;
197 }
```

4.33 Shortest Addition Chain

```
// Computa o menor numero de adicoes para construir
// cada valor, comecando com 1 (e podendo salvar variaveis)
// Retorna um par com a dp e o pai na arvore
// A arvore eh tao que o taminho da raiz (1) ate x
// contem os valores que devem ser criados para gerar x
// A profundidade de x na arvore eh dp[x]
// DP funciona para ateh 300, mas a arvore soh funciona
// para ateh 148
// recuperacao certa soh ateh 148 (erra para 149, 233, 298)
3de pair < vector < int > , vector < int >> addition_chain() {
16f
        int MAX = 301:
875
        vector < int > dp(MAX), p(MAX);
1ab
        for (int n = 2; n < MAX; n++) {
7c0
            pair < int , int > val = {INF , -1};
212
            for (int i = 1; i < n; i++) for (int j = i; j; j = p[j])
94a
                if (j == n-i) val = min(val, pair(dp[i]+1, i));
eb3
            tie(dp[n], p[n]) = val;
efe
            if (n == 9) p[n] = 8;
ba1
            if (n == 149 \text{ or } n == 233) dp[n]--;
bcd
717
        return {dp, p};
84f }
```

4.34 Simple Polygon

```
// Verifica se um poligono com n pontos eh simples
//
// O(n log n)

6e0 bool operator < (const line& a, const line& b) { // comparador pro
    sweepline
191    if (a.p == b.p) return ccw(a.p, a.q, b.q);
231    if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or a.p.x+eps <
    b.p.x))</pre>
```

```
780
            return ccw(a.p, a.q, b.p);
        return ccw(a.p, b.q, b.p);
dc0
e36 }
6f3 bool simple(vector<pt> v) {
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
576
e72
            if ((a.second+1)%v.size() == b.second or
80e
                (b.second+1)%v.size() == a.second) return false;
a08
            return interseg(a.first, b.first);
1 c.5
        }:
41a
        vector<line> seg;
e1b
        vector<pair<pt, pair<int, int>>> w;
f14
        for (int i = 0: i < v.size(): i++) {</pre>
            pt at = v[i], nxt = v[(i+1)%v.size()];
0a8
828
            if (nxt < at) swap(at, nxt);</pre>
937
            seg.push_back(line(at, nxt));
f7e
            w.push_back({at, {0, i}});
            w.push_back({nxt, {1, i}});
69 c
            // casos degenerados estranhos
            if (isinseg(v[(i+2)%v.size()], line(at, nxt))) return 0;
ae8
            if (isinseg(v[(i+v.size()-1)%v.size()], line(at, nxt)))
688
   return 0;
cba
       }
d1d
        sort(w.begin(), w.end());
7f2
        set < pair < line , int >> se;
e58
        for (auto i : w) {
ff8
            line at = seg[i.second.second];
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
7c4
                if (nxt != se.end() and intersects(*nxt, {at,
   i.second.second})) return 0;
                if (nxt != se.begin() and intersects(*(--nxt), {at,
b34
   i.second.second})) return 0:
78f
                se.insert({at. i.second.second}):
            } else {
537
884
                auto nxt = se.upper_bound({at, i.second.second}), cur
   = nxt, prev = --cur;
                if (nxt != se.end() and prev != se.begin()
b64
403
                    and intersects(*nxt, *(--prev))) return 0;
cca
                se.erase(cur);
            }
7be
        }
d17
6a5
        return 1;
af3 }
```

4.35 Sweep Direction

```
// Passa por todas as ordenacoes dos pontos definitas por "direcoes"
// Assume que nao existem pontos coincidentes
//
// O(n^2 \log n)
4b8 void sweep_direction(vector<pt> v) {
3d2
        int n = v.size():
163
        sort(v.begin(), v.end(), [](pt a, pt b) {
3a5
            if (a.x != b.x) return a.x < b.x;
572
            return a.y > b.y;
79a
        });
b89
        vector < int > at(n);
516
        iota(at.begin(), at.end(), 0);
b79
        vector<pair<int, int>> swapp;
        for (int i = 0; i < n; i++) for (int j = i+1; j < n; j++)
25e
95f
            swapp.push_back({i, j}), swapp.push_back({j, i});
        sort(swapp.begin(), swapp.end(), [&](auto a, auto b) {
269
            pt A = rotate90(v[a.first] - v[a.second]);
134
247
            pt B = rotate90(v[b.first] - v[b.second]);
            if (quad(A) == quad(B) and !sarea2(pt(0, 0), A, B)) return
615
   a < b;
224
            return compare_angle(A, B);
        });
5e7
        for (auto par : swapp) {
4e6
e24
            assert(abs(at[par.first] - at[par.second]) == 1);
            int 1 = min(at[par.first], at[par.second]),
a96
                r = n-1 - max(at[par.first], at[par.second]);
0.43
            // l e r sao quantos caras tem de cada lado do par de
                pontos
            // (cada par eh visitado duas vezes)
9cf
            swap(v[at[par.first]], v[at[par.second]]);
            swap(at[par.first], at[par.second]);
1c0
        }
241
6bb }
```

4.36 Triangulação de Delaunay

```
// Computa a triangulacao de Delaunay, o dual
// do diagrama de Voronoi (a menos de casos degenerados)
// Retorna um grafo indexado pelos indices dos pontos, e as arestas
// sao as arestas da triangulacao
// As arestas partindo de um vertice ja vem ordenadas por angulo,
// ou seja, se o vertice v nao esta no convex hull, (v, v_i, v_{i+1})
// eh um triangulo da triangulacao, em que v_i eh o i-esimo vizinho
// Usa o alg d&c, precisa representar MAX_COOR^4, por isso __int128
// pra aguentar valores ateh 1e9
```

```
//
// Propriedades:
// 1 - 0 grafo tem no max 3n-6 arestas
// 2 - Para todo triangulo, a circunf. que passa pelos 3 pontos
      nao contem estritamente nenhum ponto
// 3 - A MST euclidiana eh subgrafo desse grafo
// 4 - Cada ponto eh vizinho do ponto mais proximo dele
//
// O(n log n)
2ad typedef struct QuadEdge* Q;
ba5 struct QuadEdge {
53e
        int id:
114
        pt o;
41e
       Q rot, nxt;
       bool used;
3e5
3fc
        QuadEdge(int id_ = -1, pt o_ = pt(INF, INF)) :
            id(id_), o(o_), rot(nullptr), nxt(nullptr), used(false) {}
4ba
        Q rev() const { return rot->rot; }
00f
        Q next() const { return nxt; }
сЗс
        Q prev() const { return rot->next()->rot; }
188
0d4
        pt dest() const { return rev()->o; }
828 }:
91b Q edge(pt from, pt to, int id_from, int id_to) {
сбе
        Q e1 = new QuadEdge(id_from, from);
        Q e2 = new QuadEdge(id_to, to);
61b
8f6
        Q e3 = new QuadEdge;
       Q e4 = new QuadEdge;
5ca
e69
        tie(e1->rot, e2->rot, e3->rot, e4->rot) = \{e3, e4, e2, e1\};
        tie(e1->nxt, e2->nxt, e3->nxt, e4->nxt) = \{e1, e2, e4, e3\};
f22
1ad
        return e1:
c70 }
d8d void splice(Q a, Q b) {
a6f
        swap(a->nxt->rot->nxt, b->nxt->rot->nxt);
da4
        swap(a->nxt, b->nxt);
a58 }
167 void del_edge(Q& e, Q ne) { // delete e and assign e <- ne
        splice(e, e->prev());
cc0
        splice(e->rev(), e->rev()->prev());
eec
7ea
        delete e->rev()->rot, delete e->rev();
        delete e->rot; delete e;
524
6b2
        e = ne:
```

```
18b }
d08 Q conn(Q a, Q b) {
        Q = edge(a->dest(), b->o, a->rev()->id, b->id);
        splice(e, a->rev()->prev());
f2b
d37
        splice(e->rev(), b);
6bf
        return e:
f78 }
d64 bool in_c(pt a, pt b, pt c, pt p) { // p ta na circunf. (a, b, c) ?
268
        \_int128 p2 = p*p, A = a*a - p2, B = b*b - p2, C = c*c - p2;
        return sarea2(p, a, b) * C + sarea2(p, b, c) * A + sarea2(p,
   c. a) * B > 0:
b54 }
540 pair < Q, Q > build_tr(vector < pt > & p, int 1, int r) {
        if (r-1+1 \le 3) {
            Q = edge(p[1], p[1+1], 1, 1+1), b = edge(p[1+1], p[r],
2eb
   1+1, r);
            if (r-l+1 == 2) return {a, a->rev()};
912
0ec
            splice(a->rev(), b);
сЗс
            ll ar = sarea2(p[1], p[1+1], p[r]);
            Q c = ar ? conn(b, a) : 0;
1af
021
            if (ar >= 0) return \{a, b > rev()\};
9db
            return {c->rev(), c};
bce
        }
ee4
        int m = (1+r)/2;
328
        auto [la, ra] = build_tr(p, l, m);
        auto [lb, rb] = build_tr(p, m+1, r);
b93
667
        while (true) {
b99
            if (ccw(lb->o, ra->o, ra->dest())) ra = ra->rev()->prev();
458
            else if (ccw(lb->o, ra->o, lb->dest())) lb =
   lb->rev()->next();
f97
            else break:
24a
ca5
        Q b = conn(lb->rev(), ra);
        auto valid = [&](Q e) { return ccw(e->dest(), b->dest(),
713
   b->o); };
        if (ra->o == la->o) la = b->rev();
ee1
63f
        if (1b->o == rb->o) rb = b;
        while (true) {
667
            Q L = b - > rev() - > next();
71e
            if (valid(L)) while (in_c(b->dest(), b->o, L->dest(),
d11
   L->next()->dest()))
                del_edge(L, L->next());
1c0
            Q R = b - > prev();
c76
2b0
            if (valid(R)) while (in_c(b->dest(), b->o, R->dest(),
```

```
R->prev()->dest()))
541
                 del_edge(R, R->prev());
            if (!valid(L) and !valid(R)) break;
a3a
            if (!valid(L) or (valid(R) and in_c(L->dest(), L->o, R->o,
   R->dest())))
36c
                b = conn(R, b \rightarrow rev());
666
            else b = conn(b->rev(), L->rev());
94d
a2b
        return {la, rb};
689 }
b58 vector < vector < int >> delaunay (vector < pt > v) {
3d2
        int n = v.size():
397
        auto tmp = v;
135
        vector < int > idx(n);
        iota(idx.begin(), idx.end(), 0);
295
        sort(idx.begin(), idx.end(), [&](int 1, int r) { return v[1] <</pre>
   v[r]; });
        for (int i = 0; i < n; i++) v[i] = tmp[idx[i]];</pre>
5d8
        assert(unique(v.begin(), v.end()) == v.end());
780
4aa
        vector < vector < int >> g(n);
        bool col = true:
4ec
        for (int i = 2; i < n; i++) if (sarea2(v[i], v[i-1], v[i-2]))
   col = false;
        if (col) {
bf5
aa4
            for (int i = 1; i < n; i++)</pre>
                 g[idx[i-1]].push_back(idx[i]),
   g[idx[i]].push_back(idx[i-1]);
96b
            return g;
0ae
d36
        Q e = build_tr(v, 0, n-1).first;
        vector <Q> edg = {e};
113
5d1
        for (int i = 0; i < edg.size(); e = edg[i++]) {</pre>
3ed
            for (Q at = e: !at->used: at = at->next()) {
                 at->used = true;
60d
cf8
                 g[idx[at->id]].push_back(idx[at->rev()->id]);
                 edg.push_back(at->rev());
15d
9f2
            }
d19
        }
96b
        return g;
b43 }
      Triangulos em Grafos
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
```

```
// O(m sqrt(m) log(n)) se chamar para todos os vertices
c0d vector<pair<int, int>> g[MAX]; // {para, peso}
d41 #warning o 'g' deve estar ordenado
9a5 vector<pair<int, pair<int, int>>> get_triangles(int i) {
771
        vector<pair<int, pair<int, int>>> tri;
b23
        for (pair<int, int> j : g[i]) {
2b3
            int a = i, b = j.first;
6dd
            if (g[a].size() > g[b].size()) swap(a, b);
            for (pair<int, int> c : g[a]) if (c.first != b and c.first
eb0
   > i.first) {
525
                auto it = lower_bound(g[b].begin(), g[b].end(),
   make_pair(c.first, -INF));
                if (it == g[b].end() or it->first != c.first) continue;
f55
                tri.push_back({j.second+c.second+it->second, {a == i ?
0aa
   b : a, c.first}});
            }
b5e
7e1
f5e
        return tri;
036 }
```

5 Primitivas

5.1 Aritmetica Modular

```
// O mod tem q ser primo
429 template <int p> struct mod_int {
c68
         ll expo(ll b, ll e) {
c85
             ll ret = 1;
c87
             while (e) {
cad
                 if (e % 2) ret = ret * b % p;
                 e /= 2, b = b * b % p;
9d2
c42
edf
             return ret;
734
        }
        11 inv(11 b) { return expo(b, p-2); }
1f6
4d7
         using m = mod_int;
d93
         int v;
fe0
        mod_int() : v(0) {}
e12
         mod int(ll v ) {
019
            if (v_ >= p or v_ <= -p) v_ %= p;</pre>
bc6
             if (v_{-} < 0) v_{-} += p;
```

```
2e7
            v = v_{-};
7f3
74d
        m& operator +=(const m& a) {
2fd
            v += a.v:
ba5
            if (v >= p) v -= p;
357
            return *this;
c8b
       }
eff
        m& operator -=(const m& a) {
8b4
            v -= a.v;
            if (v < 0) v += p;
008
357
            return *this;
f8d
4c4
        m& operator *=(const m& a) {
            v = v * 11(a.v) \% p;
8a5
357
            return *this;
d4c
3f9
        m& operator /=(const m& a) {
5d6
            v = v * inv(a.v) % p;
357
            return *this;
62d
        m operator -(){ return m(-v); }
d65
b3e
        m& operator ^=(ll e) {
            if (e < 0) {
06d
6e2
                v = inv(v);
00c
                e = -e;
275
94a
            v = expo(v, e\%(p-1));
357
            return *this;
ba3
423
        bool operator ==(const m& a) { return v == a.v; }
        bool operator !=(const m& a) { return v != a.v; }
69f
1c6
        friend istream& operator >>(istream& in, m& a) {
d1c
            11 val: in >> val:
            a = m(val):
d48
091
            return in;
870
       }
44f
        friend ostream& operator <<(ostream& out, m a) {</pre>
5a0
            return out << a.v;</pre>
214
        }
399
        friend m operator +(m a, m b) { return a += b; }
        friend m operator -(m a, m b) { return a -= b; }
f9e
        friend m operator *(m a, m b) { return a *= b; }
9c1
        friend m operator /(m a, m b) { return a /= b; }
51b
08f
        friend m operator ^(m a, ll e) { return a ^= e; }
1af };
```

```
055 typedef mod_int<(int)1e9+7> mint;
5.2 Big Integer
// Complexidades: (para n digitos)
// Soma, subtracao, comparacao - O(n)
// Multiplicacao - O(n log(n))
// Divisao, resto - O(n^2)
864 struct bint {
        static const int BASE = 1e9;
669
990
        vector < int > v;
3bd
        bool neg;
609
        bint() : neg(0) {}
        bint(int val) : bint() { *this = val; }
d53
        bint(long long val) : bint() { *this = val; }
e8f
a0f
        void trim() {
f42
            while (v.size() and v.back() == 0) v.pop_back();
df8
            if (!v.size()) neg = 0;
        }
8e3
        // converter de/para string | cin/cout
294
        bint(const char* s) : bint() { from_string(string(s)); }
        bint(const string& s) : bint() { from_string(s); }
548
4ab
        void from_string(const string& s) {
            v.clear(), neg = 0;
0a6
d72
            int ini = 0;
            while (ini < s.size() and (s[ini] == '-' or s[ini] == '+'
8e2
   or s[ini] == '0'))
71d
                if (s[ini++] == '-') neg = 1;
            for (int i = s.size()-1; i >= ini; i -= 9) {
883
05e
                int at = 0:
5b1
                for (int j = max(ini, i - 8); j <= i; j++) at = 10*at</pre>
   + (s[j]-'0');
1fd
                v.push_back(at);
a5a
df8
            if (!v.size()) neg = 0;
e9a
2ff
        string to_string() const {
8be
            if (!v.size()) return "0";
793
            string ret;
73e
            if (neg) ret += '-';
            for (int i = v.size()-1; i >= 0; i--) {
3e9
582
                string at = ::to_string(v[i]);
ced
                int add = 9 - at.size();
```

```
75e
                if (i+1 < v.size()) for (int j = 0; j < add; j++) ret</pre>
   += '0';
f9f
                ret += at;
f64
            }
edf
            return ret:
770
d2f
        friend istream& operator >> (istream& in. bint& val) {
eb6
            string s; in >> s;
966
            val = s;
091
            return in:
328
99d
        friend ostream& operator << (ostream& out, const bint& val) {
8b9
            string s = val.to string():
396
            out << s:
fe8
            return out;
ce1
        }
        // operators
        friend bint abs(bint val) {
60a
c5f
            val.neg = 0;
d94
            return val;
44b
bee
        friend bint operator - (bint val) {
815
            if (val != 0) val.neg ^= 1;
d94
            return val:
326
        bint& operator=(const bint& val) { v = val.v, neg = val.neg;
   return *this: }
        bint& operator=(long long val) {
249
0a6
            v.clear(), neg = 0;
            if (val < 0) neg = 1, val *= -1;
3a6
fdc
            for (; val; val /= BASE) v.push_back(val % BASE);
            return *this:
357
220
        int cmp(const bint& r) const { // menor: -1 | igual: 0 |
3bd
   maior: 1
            if (neg != r.neg) return neg ? -1 : 1;
b14
0bb
            if (v.size() != r.v.size()) {
                int ret = v.size() < r.v.size() ? -1 : 1;</pre>
ff7
91b
                return neg ? -ret : ret;
1f6
478
            for (int i = int(v.size())-1; i >= 0; i--) {
405
                if (v[i] != r.v[i]) {
                    int ret = v[i] < r.v[i] ? -1 : 1;</pre>
2e5
91b
                    return neg ? -ret : ret;
                }
9a9
c32
            }
```

```
bb3
            return 0;
07d
152
        friend bool operator < (const bint& 1, const bint& r) { return
   1.cmp(r) == -1;}
        friend bool operator > (const bint& 1, const bint& r) { return
   1.cmp(r) == 1: }
        friend bool operator <= (const bint& l. const bint& r) { return
   1.cmp(r) <= 0;}
        friend bool operator >= (const bint& 1, const bint& r) { return
954
   1.cmp(r) >= 0: 
        friend bool operator == (const bint& 1, const bint& r) { return
   1.cmp(r) == 0: }
        friend bool operator!=(const bint& 1. const bint& r) { return
   1.cmp(r) != 0: }
        bint& operator +=(const bint& r) {
38e
6bf
            if (!r.v.size()) return *this;
            if (neg != r.neg) return *this -= -r;
a93
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {
                if (i == v.size()) v.push_back(0);
e28
                v[i] += c + (i < r.v.size() ? r.v[i] : 0);
08f
                if ((c = v[i] >= BASE)) v[i] -= BASE;
baa
            }
8bb
357
            return *this;
ab1
54c
        friend bint operator+(bint a, const bint& b) { return a += b; }
        bint& operator -=(const bint& r) {
9c8
6bf
            if (!r.v.size()) return *this;
524
            if (neg != r.neg) return *this += -r;
358
            if ((!neg and *this < r) or (neg and r < *this)) {</pre>
                *this = r - *this:
b10
a10
                neg ^= 1;
357
                return *this:
807
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {</pre>
256
                v[i] = c + (i < r.v.size() ? r.v[i] : 0);
9ef
                if ((c = v[i] < 0)) v[i] += BASE;
c8c
687
            }
0eb
            trim():
357
            return *this;
f72
        friend bint operator-(bint a, const bint& b) { return a -= b; }
f44
        // operators de * / %
6b0
        bint& operator *=(int val) {
            if (val < 0) val *= -1, neg ^= 1;</pre>
bca
566
            for (int i = 0, c = 0; i < v.size() or c; i++) {
```

```
e28
                if (i == v.size()) v.push_back(0);
                long long at = (long long) v[i] * val + c;
352
6a3
                v[i] = at % BASE;
                c = at / BASE;
b3d
            }
cb1
            trim();
0eb
357
            return *this:
a57
480
        friend bint operator *(bint a, int b) { return a *= b; }
        friend bint operator *(int a, bint b) { return b *= a; }
d5c
13b
        using cplx = complex <double >;
        void fft(vector<cplx>& a, bool f, int N, vector<int>& rev)
   const {
bc7
            for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],</pre>
   a[rev[i]]);
            vector < cplx > roots(N);
bad
192
            for (int n = 2; n <= N; n *= 2) {</pre>
4e9
                const static double PI = acos(-1);
                for (int i = 0; i < n/2; i++) {
71a
                     double alpha = (2*PI*i)/n;
40d
                     if (f) alpha = -alpha;
1a1
3f6
                     roots[i] = cplx(cos(alpha), sin(alpha));
                }
f16
3e9
                for (int pos = 0; pos < N; pos += n)
                     for (int 1 = pos, r = pos+n/2, m = 0; m < n/2;
   1++, r++, m++) {
297
                         auto t = roots[m]*a[r];
254
                         a[r] = a[1] - t:
b8f
                         a[1] = a[1] + t;
b0d
                     }
            }
e07
3f1
            if (!f) return;
            auto invN = cplx(1)/cplx(N);
08ъ
            for (int i = 0: i < N: i++) a[i] *= invN:</pre>
873
c75
0e0
        vector < long long > convolution(const vector < int > & a, const
   vector < int > & b) const {
ff9
            vector < cplx > 1(a.begin(), a.end()), r(b.begin(), b.end());
996
            int ln = l.size(), rn = r.size(), N = ln+rn+1, n = 1,
   log_n = 0;
            while (n \le N) n \le 1, \log_n + 1;
821
808
            vector < int > rev(n);
603
            for (int i = 0; i < n; i++) {</pre>
434
                rev[i] = 0:
f44
                for (int j = 0; j < log_n; j++) if (i>>j&1)
4ff
                     rev[i] = 1 << (log_n-1-j);
256
            }
```

```
230
            l.resize(n), r.resize(n);
            fft(1, false, n, rev), fft(r, false, n, rev);
a89
917
            for (int i = 0; i < n; i++) l[i] *= r[i];
88b
            fft(1, true, n, rev);
            vector < long long > ret;
7ae
c14
            for (auto& i : 1) ret.push_back(round(i.real()));
edf
            return ret:
917
633
        vector<int> convert_base(const vector<int>& a, int from, int
   to) const {
498
            static vector < long long > pot(10, 1);
671
            if (pot[1] == 1) for (int i = 1; i < 10; i++) pot[i] =</pre>
   10*pot[i-1]:
4b8
            vector < int > ret;
156
            long long at = 0;
608
            int digits = 0;
941
            for (int i : a) {
                at += i * pot[digits];
412
035
                digits += from;
                while (digits >= to) {
684
0c8
                    ret.push_back(at % pot[to]);
cf9
                    at /= pot[to];
fd4
                    digits -= to;
122
                }
87b
            }
944
            ret.push_back(at);
384
            while (ret.size() and ret.back() == 0) ret.pop_back();
edf
            return ret:
090
        bint operator*(const bint& r) const { // O(n log(n))
edb
2af
            bint ret:
968
            ret.neg = neg ^ r.neg;
            auto conv = convolution(convert_base(v, 9, 4),
d5d
   convert base(r.v. 9, 4)):
            long long c = 0;
a0e
a74
            for (auto i : conv) {
f6d
                long long at = i+c;
4cb
                ret.v.push_back(at % 10000);
a25
                c = at / 10000;
773
            }
            for (; c; c /= 10000) ret.v.push_back(c%10000);
3cb
0e2
            ret.v = convert_base(ret.v, 4, 9);
25 c
            if (!ret.v.size()) ret.neg = 0;
edf
            return ret;
c6b
359
        bint& operator*=(const bint& r) { return *this = *this * r; };
9a3
        bint& operator/=(int val) {
```

```
d9a
            if (val < 0) neg ^= 1, val *= -1;</pre>
            for (int i = int(v.size())-1, c = 0; i >= 0; i--) {
f18
2a7
                long long at = v[i] + c * (long long) BASE;
                v[i] = at / val;
e02
                c = at % val:
fb1
            }
fdb
0eb
            trim():
357
            return *this;
db6
        friend bint operator/(bint a, int b) { return a /= b; }
e74
4a9
        int operator %=(int val) {
23b
            if (val < 0) val *= -1;</pre>
156
            long long at = 0:
f31
            for (int i = int(v.size())-1; i >= 0; i--)
1b3
                at = (BASE * at + v[i]) % val;
            if (neg) at *= -1;
d22
            return at;
ce6
4b4
2fb
        friend int operator % (bint a, int b) { return a % = b; }
        friend pair < bint, bint > divmod(const bint& a_, const bint& b_)
   if (a_ == 0) return {0, 0};
611
            int norm = BASE / (b_.v.back() + 1);
d8a
            bint a = abs(a<sub>_</sub>) * norm;
b4e
027
            bint b = abs(b<sub>_</sub>) * norm;
14d
            bint a. r:
c91
            for (int i = a.v.size() - 1; i >= 0; i--) {
b71
                r *= BASE, r += a.v[i]:
                long long upper = b.v.size() < r.v.size() ?</pre>
   r.v[b.v.size()] : 0;
                int lower = b.v.size() - 1 < r.v.size() ?</pre>
   r.v[b.v.size() - 1] : 0;
                int d = (upper * BASE + lower) / b.v.back();
5d4
                r \rightarrow b*d:
                while (r < 0) r += b, d--; // roda O(1) vezes
30f
738
                q.v.push_back(d);
            }
c6a
a48
            reverse(q.v.begin(), q.v.end());
            q.neg = a_.neg ^ b_.neg;
ae2
88b
            r.neg = a_.neg;
8e5
            q.trim(), r.trim();
            return {q, r / norm};
0ef
4fd
        bint operator/(const bint& val) { return divmod(*this,
   val).first; }
        bint& operator/=(const bint& val) { return *this = *this /
   val: }
```

```
1f9
        bint operator%(const bint& val) { return divmod(*this,
   val).second: }
        bint& operator%=(const bint& val) { return *this = *this %
df5
   val; }
6c3 }:
5.3 Fracao
// Funciona com o Big Int
a4e template < typename T = int> struct frac {
a40
        T num, den;
e3f
        template < class U, class V>
        frac(U num_ = 0, V den_ = 1) : num(num_), den(den_) {
61d
bad
            assert(den != 0);
583
            if (den < 0) num *= -1, den *= -1;
a51
            T g = gcd(abs(num), den);
572
            num /= g, den /= g;
fbf
        }
51f
        friend bool operator < (const frac& 1, const frac& r) {</pre>
fa0
            return l.num * r.den < r.num * l.den;</pre>
a4e
4b5
        friend frac operator+(const frac& 1, const frac& r) {
b61
            return {1.num*r.den + 1.den*r.num, 1.den*r.den};
25f
74d
        friend frac operator-(const frac& 1, const frac& r) {
2cd
            return {1.num*r.den - 1.den*r.num, 1.den*r.den};
8a7
c80
        friend frac operator*(const frac& 1, const frac& r) {
510
            return {1.num*r.num, 1.den*r.den};
14b
        friend frac operator/(const frac& 1, const frac& r) {
a1b
8f3
            return {1.num*r.den, 1.den*r.num};
b2c
012
        friend ostream& operator << (ostream& out, frac f) {</pre>
            out << f.num << ',' << f.den;
37a
fe8
            return out;
b49
        }
cdb };
5.4 Geometria
c83 typedef double ld;
e3b const ld DINF = 1e18;
```

43a const ld pi = acos(-1.0);

107 const ld eps = 1e-9;

```
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
ba0
       return abs(a - b) <= eps;</pre>
bfc }
b2a struct pt { // ponto
       ld x, v;
3dd
        pt(1d x_{-} = 0, 1d y_{-} = 0) : x(x_{-}), y(y_{-}) {}
5bc
       bool operator < (const pt p) const {</pre>
059
            if (!eq(x, p.x)) return x < p.x;
f98
            if (!eq(y, p.y)) return y < p.y;
bb3
           return 0:
f61
a83
        bool operator == (const pt p) const {
ed0
            return eq(x, p.x) and eq(y, p.y);
589
cb9
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
       pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
4a8
       pt operator * (const ld c) const { return pt(x*c , y*c ); }
       pt operator / (const ld c) const { return pt(x/c , y/c ); }
a60
       1d operator * (const pt p) const { return x*p.x + y*p.y; }
3b6
6df
       ld operator ^ (const pt p) const { return x*p.y - y*p.x; }
5ed
       friend istream& operator >> (istream& in, pt& p) {
e37
            return in >> p.x >> p.y;
e45
       }
a8b }:
b3a struct line { // reta
       pt p, q;
0d6
       line() {}
4b8 line(pt p_, pt q_) : p(p_), q(q_) {}
       friend istream& operator >> (istream& in, line& r) {
8d7
            return in >> r.p >> r.q;
4cb
858
       }
7ab }:
// PONTO & VETOR
364 ld dist(pt p, pt q) { // distancia
       return hypot(p.y - q.y, p.x - q.x);
5f3
c68 }
9d7 ld dist2(pt p, pt q) { // quadrado da distancia
        return sq(p.x - q.x) + sq(p.y - q.y);
f24
80f }
```

```
483 ld norm(pt v) { // norma do vetor
        return dist(pt(0, 0), v);
cf7 }
589 ld angle(pt v) { // angulo do vetor com o eixo x
        ld ang = atan2(v.v.v.x):
        if (ang < 0) ang += 2*pi;</pre>
6f8
19c
        return ang;
404 }
298 ld sarea(pt p, pt q, pt r) { // area com sinal
        return ((a-p)^(r-a))/2:
1b1 }
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
        return eq(sarea(p, q, r), 0);
98c }
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
        return sarea(p, q, r) > eps;
fa7
85d }
1ef pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
e5c
        return pt(p.x * cos(th) - p.y * sin(th),
ff1
                p.x * sin(th) + p.y * cos(th));
41a }
ab1 pt rotate90(pt p) { // rotaciona 90 graus
a0d
        return pt(-p.v, p.x);
e4a }
// RETA
edc bool isvert(line r) { // se r eh vertical
       return eq(r.p.x, r.q.x);
87d
Ofb }
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
f65
        pt a = r.p - p, b = r.q - p;
b04
        return eq((a \hat{b}), 0) and (a * b) < eps;
726 }
98d ld get_t(pt v, line r) { // retorna t tal que t*v pertence a reta r
6ee
       return (r.p^r.q) / ((r.p-r.q)^v);
aOa }
```

```
256 pt proj(pt p, line r) { // projecao do ponto p na reta r
       if (r.p == r.q) return r.p;
bea
97a
       r.q = r.q - r.p; p = p - r.p;
        pt proj = r.q * ((p*r.q) / (r.q*r.q));
9f8
2cd
        return proj + r.p;
232 }
d5c pt inter(line r, line s) { // r inter s
146
        if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
205
       r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
543
        return r.q * get_t(r.q, s) + r.p;
111 }
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
       if (isinseg(r.p, s) or isinseg(r.q, s)
19b
c21
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
9fa
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
fcb ld disttoline(pt p, line r) { // distancia do ponto a reta
89a
        return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
1b7 }
bcc ld disttoseg(pt p, line r) { // distancia do ponto ao seg
73d
       if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
        if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
951
        return disttoline(p, r);
a19
367 }
11d ld distseg(line a, line b) { // distancia entre seg
        if (interseg(a, b)) return 0;
4df
       ld ret = DINF;
349
       ret = min(ret, disttoseg(a.p, b));
341
ceb
        ret = min(ret, disttoseg(a.q, b));
093
       ret = min(ret, disttoseg(b.p, a));
448
        ret = min(ret, disttoseg(b.q, a));
edf
        return ret;
222 }
// POLIGONO
// corta poligono com a reta r deixando os pontos p tal que
```

```
// ccw(r.p, r.q, p)
1a9 vector<pt> cut_polygon(vector<pt> v, line r) { // O(n)
8af
        vector<pt> ret;
        for (int j = 0; j < v.size(); j++) {</pre>
8a4
            if (ccw(r.p, r.q, v[j])) ret.push_back(v[j]);
dac
            if (v.size() == 1) continue;
dce
030
            line s(v[j], v[(j+1)\%v.size()]);
            pt p = inter(r, s);
ae3
a3d
            if (isinseg(p, s)) ret.push_back(p);
d44
8a1
        ret.erase(unique(ret.begin(), ret.end()), ret.end());
24d
        if (ret.size() > 1 and ret.back() == ret[0]) ret.pop_back();
edf
        return ret:
253 }
// distancia entre os retangulos a e b (lados paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
5f5 ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
        ld hor = 0, vert = 0;
080
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
34b
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
f5f
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
4fd
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
96f
        return dist(pt(0, 0), pt(hor, vert));
630 }
13d ld polarea(vector<pt> v) { // area do poligono
        ld ret = 0:
9c5
        for (int i = 0; i < v.size(); i++)</pre>
сбе
80f
            ret += sarea(pt(0, 0), v[i], v[(i + 1) \% v.size()]);
d03
        return abs(ret):
5df }
// se o ponto ta dentro do poligono: retorna O se ta fora.
// 1 se ta no interior e 2 se ta na borda
8e7 int inpol(vector\phit>& v, pt p) { // O(n)
8de
        int qt = 0;
f14
        for (int i = 0; i < v.size(); i++) {</pre>
bda
            if (p == v[i]) return 2;
            int j = (i+1)%v.size();
6af
e38
            if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
                if ((v[i]-p)*(v[j]-p) < eps) return 2;</pre>
97f
5e2
                continue:
48b
            }
388
            bool baixo = v[i].y+eps < p.y;</pre>
464
            if (baixo == (v[j].y+eps < p.y)) continue;</pre>
366
            auto t = (p-v[i])^(v[j]-v[i]);
```

```
1b4
            if (eq(t, 0)) return 2;
839
            if (baixo == (t > eps)) qt += baixo ? 1 : -1;
d13
        }
b84
        return qt != 0;
a64 }
6ff bool interpol(vector<pt> v1, vector<pt> v2) { // se dois poligonos
   se intersectam - O(n*m)
        int n = v1.size(), m = v2.size();
7d1
        for (int i = 0; i < n; i++) if (inpol(v2, v1[i])) return 1;</pre>
c36
ab8
        for (int i = 0; i < n; i++) if (inpol(v1, v2[i])) return 1;
        for (int i = 0; i < n; i++) for (int j = 0; j < m; j++)
523
0c8
            if (interseg(line(v1[i], v1[(i+1)%n]), line(v2[j],
   v2[(j+1)%m]))) return 1;
bb3
        return 0;
c58 }
494 ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre
        if (interpol(v1, v2)) return 0;
f6b
        ld ret = DINF;
349
        for (int i = 0; i < v1.size(); i++) for (int j = 0; j <</pre>
   v2.size(); j++)
6c2
            ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
   v1.size()]),
9d9
                         line(v2[j], v2[(j + 1) % v2.size()])));
edf
        return ret;
125 }
138 vector <pt> convex_hull(vector <pt> v) { // convex hull - O(n log(n))
        sort(v.begin(), v.end());
fca
d76
        v.erase(unique(v.begin(), v.end()), v.end());
        if (v.size() <= 1) return v;</pre>
52d
526
        vector<pt> 1, u;
        for (int i = 0; i < v.size(); i++) {</pre>
f14
fb2
            while (1.size() > 1 \text{ and } !ccw(1.end()[-2], 1.end()[-1],
   v[i]))
364
                1.pop_back();
c35
            l.push_back(v[i]);
58e
3e9
        for (int i = v.size() - 1; i >= 0; i--) {
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
f19
   v[i]))
                u.pop_back();
7a8
            u.push_back(v[i]);
a95
```

```
0ъ8
cfc
        1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
792
        return 1;
10d }
483 struct convex_pol {
        vector<pt> pol;
f50
        // nao pode ter ponto colinear no convex hull
d98
        convex_pol() {}
a04
        convex_pol(vector < pt > v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(\log(n))
8af
        bool is_inside(pt p) {
b6e
             if (pol.size() == 0) return false;
eae
             if (pol.size() == 1) return p == pol[0];
67f
             int 1 = 1, r = pol.size();
40c
             while (1 < r) {
                 int m = (1+r)/2:
ee4
                 if (ccw(p, pol[0], pol[m])) l = m+1;
48f
ef3
                 else r = m;
            }
91c
00a
             if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
9e7
             if (l == pol.size()) return false;
1c0
             return !ccw(p, pol[1], pol[1-1]);
6b0
        }
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
719
        int extreme(const function < bool(pt, pt) > & cmp) {
             int n = pol.size();
b1c
4a2
             auto extr = [&](int i, bool& cur_dir) {
22a
                 \operatorname{cur\_dir} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
61a
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
364
             };
63d
             bool last_dir, cur_dir;
             if (extr(0, last_dir)) return 0;
a0d
993
             int 1 = 0, r = n;
             while (1+1 < r) {
ead
ee4
                 int m = (1+r)/2;
f29
                 if (extr(m, cur_dir)) return m;
                 bool rel_dir = cmp(pol[m], pol[1]);
44a
b18
                 if ((!last_dir and cur_dir) or
261
                          (last_dir == cur_dir and rel_dir == cur_dir)) {
8a6
1 f 1
                     last_dir = cur_dir;
94a
                 else r = m:
```

```
606
792
            return 1;
56c
        }
316
        int max_dot(pt v) {
            return extreme([&](pt p, pt q) { return p*v > q*v; });
ec1
3b7
a54
        pair < int , int > tangents(pt p) {
            auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
ffb
8fd
            auto R = [k](pt q, pt r) \{ return ccw(p, q, r); \};
            return {extreme(L). extreme(R)};
fa8
736
        }
3ec }:
// CIRCUNFERENCIA
911 pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3 pontos
174
        b = (a + b) / 2;
       c = (a + c) / 2;
2ae
98b
        return inter(line(b, b + rotate90(a - b)),
3f8
                line(c, c + rotate90(a - c))):
a12 }
4b3 vector <pt> circ_line_inter(pt a, pt b, pt c, ld r) { // intersecao
   da circunf (c, r) e reta ab
        vector <pt> ret;
8af
f2b
        b = b-a, a = a-c:
4b1
       1d A = b*b;
       1d B = a*b;
20a
2e9
       1d C = a*a - r*r:
1fa
       1d D = B*B - A*C;
       if (D < -eps) return ret;</pre>
818
dc5
       ret.push_back(c+a+b*(-B+sqrt(D+eps))/A);
        if (D > eps) ret.push_back(c+a+b*(-B-sqrt(D))/A);
20e
edf
        return ret:
cd8 }
ad2 vector<pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao da
   circunf (a, r) e (b, R)
8af
        vector <pt> ret;
b7e
        1d d = dist(a, b);
       if (d > r+R \text{ or } d+min(r, R) < max(r, R)) return ret;
       1d x = (d*d-R*R+r*r)/(2*d);
398
183
       1d v = sqrt(r*r-x*x);
       pt v = (b-a)/d;
325
76e
       ret.push_back(a+v*x + rotate90(v)*y);
        if (y > 0) ret.push_back(a+v*x - rotate90(v)*y);
2cb
        return ret:
edf
```

```
fb1 }
6e0 bool operator <(const line& a, const line& b) { // comparador pra
   reta
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
f82
        if (!eq(angle(v1), angle(v2))) return angle(v1) < angle(v2);</pre>
        return ccw(a.p, a.q, b.p); // mesmo angulo
780
27e }
b14 bool operator ==(const line& a, const line& b) {
76c
        return !(a < b) and !(b < a);
449 }
// comparador pro set pra fazer sweep line com segmentos
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
d80
             // assume que os segmentos tem p < q
             if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
             if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
231
    b.p.x))
780
                 return ccw(a.p, a.q, b.p);
dc0
             return ccw(a.p, b.q, b.p);
243
367 };
// comparador pro set pra fazer sweep angle com segmentos
bef pt dir;
5b0 struct cmp_sweepangle {
        bool operator () (const line& a, const line& b) const {
522
             return get_t(dir, a) + eps < get_t(dir, b);</pre>
        }
653
97f };
5.5 Geometria - inteiro
2de #define sq(x) ((x)*(11)(x))
b2a struct pt { // ponto
e91
        int x, y;
        pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
df1
        bool operator < (const pt p) const {</pre>
5bc
95a
             if (x != p.x) return x < p.x;
89 c
             return y < p.y;</pre>
dcd
a83
        bool operator == (const pt p) const {
d74
             return x == p.x and y == p.y;
7b4
        }
```

```
cb9
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
a24
0ef
       pt operator * (const int c) const { return pt(x*c, y*c); }
        11 operator * (const pt p) const { return x*(11)p.x +
   y*(11)p.y; }
       11 operator ^ (const pt p) const { return x*(11)p.y -
   y*(11)p.x; }
        friend istream& operator >> (istream& in, pt& p) {
5ed
            return in >> p.x >> p.y;
e37
e45
840 };
b3a struct line { // reta
730
        pt p, q;
0d6
       line() {}
4b8
       line(pt p_, pt q_) : p(p_), q(q_) {}
       friend istream& operator >> (istream& in, line& r) {
8d7
            return in >> r.p >> r.q;
4cb
858
       }
7ab }:
// PONTO & VETOR
ea8 ll dist2(pt p, pt q) { // quadrado da distancia
f24
        return sq(p.x - q.x) + sq(p.y - q.y);
515 }
5a2 11 sarea2(pt p, pt q, pt r) { // 2 * area com sinal}
       return (q-p)^(r-q);
bf4 }
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
        return sarea2(p, q, r) == 0;
a08 }
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
276
        return sarea2(p, q, r) > 0;
42b }
c31 int quad(pt p) { // quadrante de um ponto
dbb
        return (p.x<0)^3*(p.y<0);
fcf }
2df bool compare_angle(pt p, pt q) { // retorna se ang(p) < ang(q)
        if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
9fc
ea1
        return ccw(q, pt(0, 0), p);
771 }
```

```
ab1 pt rotate90(pt p) { // rotaciona 90 graus
a0d
       return pt(-p.v, p.x);
e4a }
// RETA
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
f65
        pt a = r.p - p, b = r.q - p;
2ac
        return (a ^ b) == 0 and (a * b) <= 0;
c9f }
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
   de s
19b
        if (isinseg(r.p, s) or isinseg(r.q, s)
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
9fa
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
359 }
9e0 int segpoints(line r) { // numero de pontos inteiros no segmento
        return 1 + \_gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
dd8 }
88a double get_t(pt v, line r) { // retorna t tal que t*v pertence a
   reta r
1ad
       return (r.p^r.q) / (double) ((r.p-r.q)^v);
d27 }
// POI.TGONO
// quadrado da distancia entre os retangulos a e b (lados paralelos
   aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
485 ll dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
c59
       int hor = 0, vert = 0;
34b
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
f5f
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
4fd
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
80a
869
        return sq(hor) + sq(vert);
e13 }
9c3 ll polarea2(vector<pt> v) { // 2 * area do poligono
b73
        ll ret = 0:
        for (int i = 0: i < v.size(): i++)</pre>
сбе
```

```
532
            ret += sarea2(pt(0, 0), v[i], v[(i + 1) % v.size()]);
d03
        return abs(ret):
d5f }
// se o ponto ta dentro do poligono: retorna 0 se ta fora,
// 1 se ta no interior e 2 se ta na borda
8e7 int inpol(vector\phi) { // O(n)
        int qt = 0;
8de
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            if (p == v[i]) return 2;
bda
6af
            int j = (i+1)%v.size();
cc6
            if (p.y == v[i].y and p.y == v[j].y) {
547
                if ((v[i]-p)*(v[j]-p) <= 0) return 2;</pre>
5e2
                continue:
            }
b47
78c
            bool baixo = v[i].y < p.y;</pre>
057
            if (baixo == (v[j].v < p.v)) continue;</pre>
366
            auto t = (p-v[i])^(v[j]-v[i]);
            if (!t) return 2;
2ad
            if (baixo == (t > 0)) qt += baixo ? 1 : -1;
0bb
9cf
        }
b84
        return qt != 0;
afd }
138 vector<pt> convex_hull(vector<pt> v) { // convex hull - O(n log(n))
fca
        sort(v.begin(), v.end());
        v.erase(unique(v.begin(), v.end()), v.end());
d76
52d
       if (v.size() <= 1) return v;</pre>
526
        vector <pt> 1, u;
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            while (1.size() > 1 \text{ and } !ccw(1.end()[-2], 1.end()[-1],
fb2
   v[i]))
                l.pop_back();
364
c35
            l.push_back(v[i]);
58e
        for (int i = v.size() - 1; i >= 0; i--) {
3e9
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
f19
   v[i]))
7a8
                u.pop_back();
a95
            u.push_back(v[i]);
0b8
cfc
        1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
        return 1;
792
10d }
786 ll interior_points(vector<pt> v) { // pontos inteiros dentro de um
```

```
poligono simples
c4e
       11 b = 0:
        for (int i = 0; i < v.size(); i++)</pre>
c6e
            b += segpoints(line(v[i], v[(i+1)%v.size()])) - 1;
Осе
        return (polarea2(v) - b) / 2 + 1;
a1c
af2 }
483 struct convex_pol {
f50
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
d98
        convex_pol() {}
a04
        convex pol(vector<pt> v) : pol(convex hull(v)) {}
        // se o ponto ta dentro do hull - O(\log(n))
        bool is_inside(pt p) {
8af
b6e
            if (pol.size() == 0) return false;
            if (pol.size() == 1) return p == pol[0];
eae
67f
            int 1 = 1, r = pol.size();
            while (1 < r) {
40c
ee4
                int m = (1+r)/2;
48f
                if (ccw(p, pol[0], pol[m])) l = m+1;
ef3
                else r = m;
91c
            }
00a
            if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
9e7
            if (1 == pol.size()) return false;
            return !ccw(p, pol[1], pol[1-1]);
1 c 0
6b0
       }
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
719
        int extreme(const function < bool(pt, pt) > & cmp) {
b1c
            int n = pol.size();
4a2
            auto extr = [&](int i, bool& cur_dir) {
                cur dir = cmp(pol[(i+1)\%n], pol[i]);
22a
61a
                return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
364
            };
63d
            bool last_dir, cur_dir;
a0d
            if (extr(0, last_dir)) return 0;
993
            int 1 = 0, r = n;
ead
            while (1+1 < r) {
                int m = (1+r)/2:
ee4
                if (extr(m, cur_dir)) return m;
f29
                bool rel_dir = cmp(pol[m], pol[l]);
44a
                if ((!last_dir and cur_dir) or
b18
261
                        (last_dir == cur_dir and rel_dir == cur_dir)) {
8a6
                    1 = m:
1f1
                    last dir = cur dir:
```

```
94a
                } else r = m;
606
792
            return 1;
56c
        }
316
        int max dot(pt v) {
            return extreme([&](pt p, pt q) { return p*v > q*v; });
ec1
3b7
       }
a54
        pair < int , int > tangents(pt p) {
ffb
            auto L = [k](pt q, pt r) \{ return ccw(p, r, q); \};
8fd
            auto R = [&](pt q, pt r) { return ccw(p, q, r); };
fa8
            return {extreme(L), extreme(R)};
736
       }
3ec }:
6e0 bool operator <(const line& a, const line& b) { // comparador pra
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
        bool b1 = compare_angle(v1, v2), b2 = compare_angle(v2, v1);
036
73c
        if (b1 or b2) return b1:
        return ccw(a.p, a.q, b.p); // mesmo angulo
780
b61 }
b14 bool operator ==(const line& a, const line& b) {
76c
        return !(a < b) and !(b < a);</pre>
449 }
// comparador pro set pra fazer sweep line com segmentos
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
            // assume que os segmentos tem p < q
191
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
614
            if (a.p.x != a.q.x and (b.p.x == b.q.x or a.p.x < b.p.x))
780
                return ccw(a.p, a.q, b.p);
dc0
            return ccw(a.p. b.g. b.p):
       }
baf
677 };
// comparador pro set pra fazer sweep angle com segmentos
bef pt dir;
5b0 struct cmp_sweepangle {
       bool operator () (const line& a, const line& b) const {
            return get_t(dir, a) < get_t(dir, b);</pre>
261
dc5
       }
f6d };
```

5.6 Geometria 3D

```
c83 typedef double ld;
e3b const ld DINF = 1e18;
107 const ld eps = 1e-9;
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
            return abs(a - b) <= eps;</pre>
ba0
bfc }
b2a struct pt { // ponto
2eb
            ld x, y, z;
a50
            pt(1d x_{-} = 0, 1d y_{-} = 0, 1d z_{-} = 0) : x(x_{-}), y(y_{-}), z(z_{-})
  {}
5bc
            bool operator < (const pt p) const {</pre>
059
                    if (!eq(x, p.x)) return x < p.x;
f98
                    if (!eq(v, p.v)) return v < p.v;
44c
                    if (!eq(z, p.z)) return z < p.z;
bb3
                    return 0;
6cd
            }
a83
            bool operator == (const pt p) const {
41c
                    return eq(x, p.x) and eq(y, p.y) and eq(z, p.z);
fb5
            }
44b
            pt operator + (const pt p) const { return pt(x+p.x, y+p.y,
   z+p.z); }
392
            pt operator - (const pt p) const { return pt(x-p.x, y-p.y,
   z-p.z); }
fb7
            pt operator * (const ld c) const { return pt(x*c , y*c ,
   z*c ): }
            pt operator / (const ld c) const { return pt(x/c , y/c ,
7a1
   z/c ): }
a65
            ld operator * (const pt p) const { return x*p.x + y*p.y +
   z*p.z; }
7f6
            pt operator ^ (const pt p) const { return pt(y*p.z -
   z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }
5ed
            friend istream& operator >> (istream& in, pt& p) {
9bf
                    return in >> p.x >> p.y >> p.z;
5e8
            }
3ee };
b3a struct line { // reta
730
            pt p, q;
0d6
            line() {}
4 b 8
            line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
            friend istream& operator >> (istream& in, line& r) {
4cb
                    return in >> r.p >> r.q;
858
            }
```

```
7ab }:
79b struct plane { // plano
            array<pt, 3> p; // pontos que definem o plano
7 e 1
29b
            array < ld, 4 > eq; // equacao do plano
bb7
            plane() {}
            plane(pt p_, pt q_, pt r_) : p({p_, q_, r_}) { build(); }
fb0
            friend istream& operator >> (istream& in, plane& P) {
ca9
                    return in >> P.p[0] >> P.p[1] >> P.p[2];
2ab
70e
                    P.build();
544
            }
0a8
            void build() {
da2
                    pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
7d5
                    eq = \{dir.x, dir.y, dir.z, dir*p[0]*(-1)\};
41a
            }
d5d };
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de rotacao ao
   redor de z
2fb pt convert(ld rho, ld th, ld phi) {
           return pt(sin(phi) * cos(th), sin(phi) * sin(th),
   cos(phi)) * rho;
a4f }
// projecao do ponto p na reta r
256 pt proj(pt p, line r) {
bea
            if (r.p == r.q) return r.p;
            r.q = r.q - r.p; p = p - r.p;
97a
9f8
            pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
           return proj + r.p;
232 }
// projecao do ponto p no plano P
bla pt proj(pt p, plane P) {
            p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2]
   - P.p[0]:
b69
            pt norm = P.p[1] ^ P.p[2];
            pt proj = p - (norm * (norm * p) / (norm*norm));
6ab
            return proj + P.p[0];
467
4a0 }
// distancia
a45 ld dist(pt a, pt b) {
fd9
            return sqrt(sq(a.x-b.x) + sq(a.y-b.y) + sq(a.z-b.z));
```

```
2d0 }
// distancia ponto reta
137 ld distline(pt p, line r) {
            return dist(p, proj(p, r));
ce1
3c4 }
// distancia de ponto para segmento
d43 ld distseg(pt p, line r) {
            if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
73d
951
            if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
200
            return distline(p, r);
42c }
// distancia de ponto a plano com sinal
7cc ld sdist(pt p, plane P) {
150
            return P.eq[0]*p.x + P.eq[1]*p.y + P.eq[2]*p.z + P.eq[3];
d49 }
// distancia de ponto a plano
768 ld distplane(pt p, plane P) {
сЗе
            return abs(sdist(p, P));
33d }
// se ponto pertence a reta
099 bool isinseg(pt p, line r) {
a32
            return eq(distseg(p, r), 0);
31a }
// se ponto pertence ao triangulo definido por P.p
cd2 bool isinpol(pt p, vector<pt> v) {
fad
            assert(v.size() >= 3);
            pt norm = (v[1]-v[0]) ^ (v[2]-v[1]);
bf4
8a4
            bool inside = true:
            int sign = -1;
cec
            for (int i = 0; i < v.size(); i++) {</pre>
f14
834
                    line r(v[(i+1)\%3], v[i]);
2a9
                    if (isinseg(p, r)) return true;
4ef
                    pt ar = v[(i+1)\%3] - v[i];
320
                    if (sign == -1) sign = ((ar^(p-v[i]))*norm > 0);
                    else if (((ar^(p-v[i]))*norm > 0) != sign) inside
82b
   = false;
15e
            return inside;
aca
c81 }
```

```
// distancia de ponto ate poligono
361 ld distpol(pt p, vector<pt> v) {
3e7
            pt p2 = proj(p, plane(v[0], v[1], v[2]));
            if (isinpol(p2, v)) return dist(p, p2);
61a
349
            ld ret = DINF:
            for (int i = 0; i < v.size(); i++) {</pre>
f14
6af
                   int i = (i+1)\%v.size():
5ee
                    ret = min(ret, distseg(p, line(v[i], v[j])));
7b2
edf
            return ret:
a8d }
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
// CONCOR = segmento concorrente ao plano
e51 enum RETCODE {BOTH, ONE, PARAL, CONCOR};
26b pair < RETCODE, pt > intersect(plane P, line r) {
       1d d1 = sdist(r.p, P);
f8f
        1d d2 = sdist(r.q, P);
53a
        if (eq(d1, 0) and eq(d2, 0))
                    return pair(BOTH, r.p);
504
72c
        if (eq(d1, 0))
847
                    return pair(ONE, r.p);
485
        if (eq(d2, 0))
168
                    return pair(ONE, r.q);
3fb
        if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0)) 
463
            if (eq(d1-d2, 0)) return pair(PARAL, pt());
406
            return pair(CONCOR, pt());
91 c
c84
        1d frac = d1 / (d1 - d2);
        pt res = r.p + ((r.q - r.p) * frac);
394
        return pair(ONE, res);
b92 }
// rotaciona p ao redor do eixo u por um angulo a
787 pt rotate(pt p, pt u, ld a) {
           u = u / dist(u, pt());
            return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) *
   sin(a):
7f0 }
5.7 Matriz
945 #define MODULAR false
5ed template < typename T > struct matrix : vector < vector < T >> {
```

```
14e
        int n, m;
        void print() {
30f
            for (int i = 0; i < n; i++) {</pre>
603
                 for (int j = 0; j < m; j++) cout << (*this)[i][j] << "</pre>
70f
1fb
                cout << endl:</pre>
            }
d98
101
        }
        matrix(int n_, int m_, bool ident = false) :
aa3
b14
                 vector < vector < T > (n_, vector < T > (m_, 0)), n(n_), m(m_)  {
94e
            if (ident) {
df7
                assert(n == m):
a89
                for (int i = 0; i < n; i++) (*this)[i][i] = 1;
359
527
        }
b83
        matrix(const vector<vector<T>>& c) : vector<vector<T>>(c),
            n(c.size()), m(c[0].size()) {}
a3d
        matrix(const initializer list<initializer list<T>>& c) {
efc
f7e
            vector < vector < T >> val:
212
            for (auto& i : c) val.push_back(i);
            *this = matrix(val):
303
c50
        }
388
        matrix<T> operator*(matrix<T>& r) {
            assert(m == r.n);
1e2
82c
            matrix <T> M(n. r.m):
d69
            for (int i = 0; i < n; i++) for (int k = 0; k < m; k++)
df4
                 for (int j = 0; j < r.m; j++) {
                    T \text{ add} = (*this)[i][k] * r[k][j];
e34
f98 #if MODULAR
d41 #warning Usar matrix<11> e soh colocar valores em [0, MOD) na
   matriz!
8b6
                     M[i][i] += add%MOD;
                     if (M[i][j] >= MOD) M[i][j] -= MOD;
983
8c1 #else
7bb
                     M[i][i] += add;
f2e #endif
620
                }
474
            return M:
394
528
        matrix<T> operator^(ll e){
            matrix<T> M(n, n, true), at = *this;
f10
c87
            while (e) {
2e2
                if (e\&1) M = M*at;
cc2
                e >>= 1:
```

```
c80
                 at = at*at;
            }
eb6
474
            return M;
ca3
        }
        void apply_transform(matrix M, ll e){
582
            auto& v = *this;
1 c 3
c87
            while (e) {
9ba
                 if (e\&1) v = M*v;
                 e >>= 1;
cc2
419
                 M = M * M:
d86
            }
4e5
        }
70d }:
```

5.8 Matroid

```
// Matroids de Grafo e Particao
// De modo geral, toda Matroid contem um build() linear
// e uma funcao constante oracle()
// oracle(i) responde se o conjunto continua independente
// apos adicao do elemento i
// oracle(i, j) responde se o conjunto continua indepente
// apos trocar o elemento i pelo elemento j
// Intersecao sem peso O(r^2 n)
// em que n eh o tamanho do conjunto e r eh o tamanho da resposta
// Matroid Grafica
// Matroid das florestas de um grafo
// Um conjunto de arestas eh independente se formam uma floresta
// build() : O(n)
// oracle() : 0(1)
fda struct graphic_matroid {
5da
        int n, m, t;
32c
        vector < array < int , 2>> edges;
789
        vector < vector < int >> g;
62e
        vector < int > comp, in, out;
513
        graphic_matroid(int n_, vector<array<int, 2>> edges_)
            : n(n_), m(edges_.size()), edges(edges_), g(n), comp(n),
   in(n), out(n) {}
315
        void dfs(int u) {
ab8
            in[u] = t++;
17d
            for (auto v : g[u]) if (in[v] == -1)
863
                comp[v] = comp[u], dfs(v);
677
            out[u] = t;
```

```
d83
945
        void build(vector<int> I) {
            t = 0:
a34
741
            for (int u = 0; u < n; u++) g[u].clear(), in[u] = -1;
667
            for (int e : I) {
                auto [u, v] = edges[e];
d00
125
                g[u].push_back(v), g[v].push_back(u);
a8a
809
            for (int u = 0; u < n; u++) if (in[u] == -1)
                comp[u] = u, dfs(u);
a7d
207
        }
f31
        bool is_ancestor(int u, int v) {
a68
            return in[u] <= in[v] and in[v] < out[u];</pre>
0c2
        }
e6b
        bool oracle(int e) {
453
            return comp[edges[e][0]] != comp[edges[e][1]];
687
        }
f75
        bool oracle(int e, int f) {
574
            if (oracle(f)) return true;
            int u = edges[e][in[edges[e][0]] < in[edges[e][1]]];</pre>
622
ff2
            return is_ancestor(u, edges[f][0]) != is_ancestor(u,
   edges[f][1]);
       }
8a9
691 };
// Matroid de particao ou cores
// Um conjunto eh independente se a quantidade de elementos
// de cada cor nao excede a capacidade da cor
// Quando todas as capacidades sao 1, um conjunto eh independente
// se todas as suas cores sao distintas
//
// build() : O(n)
// oracle() : 0(1)
994 struct partition_matroid {
501
        vector < int > cap, color, d;
        partition_matroid(vector<int> cap_, vector<int> color_)
608
04d
            : cap(cap_), color(color_), d(cap.size()) {}
945
        void build(vector<int> I) {
def
            fill(d.begin(), d.end(), 0);
e9d
            for (int u : I) d[color[u]]++;
c58
514
        bool oracle(int u) {
0a1
            return d[color[u]] < cap[color[u]];</pre>
703
        }
f7f
        bool oracle(int u, int v) {
2f7
            return color[u] == color[v] or oracle(v);
```

```
4b4
        }
caa }:
// Intersecao de matroid sem pesos
// Dadas duas matroids M1 e M2 definidas sobre o mesmo
// conjunto I, retorna o maior subconjunto de I
// que eh independente tanto para M1 quanto para M2
//
// O(r^2*n)
// Matroid "pesada" deve ser a M2
132 template < typename Matroid1, typename Matroid2 >
801 vector <int > matroid intersection(int n. Matroid1 M1. Matroid2 M2) {
f5b
        vector < bool > b(n):
a64
        vector < int > I[2];
a8b
        bool converged = false;
0 c 1
        while (!converged) {
742
            I[0].clear(), I[1].clear();
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
99d
09d
            M1.build(I[1]), M2.build(I[1]);
            vector < bool > target(n), pushed(n);
289
26a
            queue < int > q;
5c5
            for (int u : I[0]) {
2h2
                target[u] = M2.oracle(u);
c1b
                if (M1.oracle(u)) pushed[u] = true, q.push(u);
0e6
            }
3fe
            vector < int > p(n, -1);
07a
            converged = true;
402
            while (q.size()) {
                int u = q.front(); q.pop();
be1
5c6
                if (target[u]) {
                     converged = false;
101
c32
                     for (int v = u; v != -1; v = p[v]) b[v] = !b[v];
c2b
                     break:
                }
a80
                for (int v : I[!b[u]]) if (!pushed[v]) {
e78
                     if ((b[u] and M1.oracle(u, v)) or (b[v] and
   M2.oracle(v. u)))
                         p[v] = u, pushed[v] = true, q.push(v);
bae
533
                }
1d9
5e7
b68
        return I[1];
381 }
// Intersecao de matroid com pesos
```

```
// Dadas duas matroids M1 e M2 e uma funcao de pesos w, todas
   definidas sobre
// um conjunto I retorna o maior subconjunto de I (desempatado pelo
// que eh independente tanto para M1 quanto para M2
// A resposta eh construida incrementando o tamanho conjunto I de 1 em
// Se nao tiver custo negativo, nao precisa de SPFA
// O(r^3*n) com SPFA
// O(r^2*n*log(n)) com Dijkstra e potencial
42a template < typename T, typename Matroid1, typename Matroid2>
2b5 vector < int > weighted_matroid_intersection(int n, vector < T > w,
   Matroid1 M1, Matroid2 M2) {
6c9
        vector < bool > b(n), target(n), is_inside(n);
        vector<int> I[2], from(n);
563
        vector < pair < T, int >> d(n);
e35
        auto check_edge = [&](int u, int v) {
169
            return (b[u] and M1.oracle(u, v)) or (b[v] and
249
   M2.oracle(v, u));
253
        }:
        while (true) {
667
742
            I[0].clear(), I[1].clear();
994
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
            // I[1] contem o conjunto de tamanho I[1].size() de menor
                peso
09d
            M1.build(I[1]), M2.build(I[1]);
            for (int u = 0; u < n; u++) {
687
                target[u] = false, is_inside[u] = false, from[u] = -1;
ea5
                d[u] = {numeric_limits <T>::max(), INF};
961
392
8d3
            deque <T> q;
476
            sort(I[0].begin(), I[0].end(), [&](int i, int j){ return
   w[i] < w[i]; });
            for (int u : I[0]) {
5c5
                target[u] = M2.oracle(u);
2b2
5a7
                if (M1.oracle(u)) {
4ef
                    if (is_inside[u]) continue;
7cc
                    d[u] = \{w[u], 0\};
427
                     if (!q.empty() and d[u] > d[q.front()])
    q.push_back(u);
655
                    else q.push_front(u);
                    is_inside[u] = true;
4ae
764
                }
            }
add
402
            while (q.size()) {
```

```
97a
                int u = q.front(); q.pop_front();
6f3
                is_inside[u] = false;
57a
                for (int v : I[!b[u]]) if (check_edge(u, v)) {
                    pair <T, int > nd(d[u].first + w[v], d[u].second +
   1):
61b
                    if (nd < d[v]) {
                        from[v] = u, d[v] = nd;
6ac
bd7
                        if (is_inside[v]) continue;
                        if (q.size() and d[v] > d[q.front()])
   q.push_back(v);
275
                         else q.push_front(v);
                         is_inside[v] = true;
b3f
                    }
                }
a3b
            }
563
cc8
            pair < T, int > mini = pair (numeric_limits < T >:: max(), INF);
489
            int targ = -1;
259
            for (int u : I[0]) if (target[u] and d[u] < mini)</pre>
                mini = d[u], targ = u;
2b9
            if (targ != -1) for (int u = targ; u != -1; u = from[u])
e14
d89
                b[u] = !b[u], w[u] *= -1;
f97
            else break;
c7d
b68
        return I[1];
8e7 }
```

6 Estruturas

6.1 BIT

```
// BIT de soma 1-based, v 0-based
// Para mudar o valor da posicao p para x,
// faca: poe(x - query(p, p), p)
// l_bound(x) retorna o menor p tal que
// query(1, p+1) > x (0 based!)
//
// Complexidades:
// build - O(n)
// poe - O(log(n))
// query - O(log(n))
// l_bound - O(log(n))

1a8 int n;
7f4 int bit[MAX];
b69 int v[MAX];
```

```
0a8 void build() {
b91
        bit[0] = 0:
        for (int i = 1; i <= n; i++) bit[i] = v[i - 1];</pre>
33c
        for (int i = 1: i <= n: i++) {
78a
            int j = i + (i & -i);
edf
b8a
            if (j <= n) bit[j] += bit[i];</pre>
b39
        }
5dc }
// soma x na posicao p
235 void poe(int x, int p) {
        for (; p <= n; p += p & -p) bit[p] += x;
ec3 }
// soma [1, p]
Obf int pref(int p) {
7c9
        int ret = 0;
        for (; p; p -= p & -p) ret += bit[p];
805
        return ret:
edf
601 }
// soma [a, b]
4ea int query(int a, int b) {
70c
        return pref(b) - pref(a - 1);
3dd }
e4a int l_bound(ll x) {
        int p = 0;
676
        for (int i = MAX2; i+1; i--) if (p + (1 << i) <= n
729
             and bit[p + (1 << i)] <= x) x -= bit[p += (1 << i)];
74e
        return p;
c89 }
6.2 BIT 2D
// BIT de soma, update incrementa posicao
// Tem que construir com um vetor com todos os pontos
// que vc quer um dia atualizar (os pontos q vc vai chamar update)
//
// Complexidades:
// construir - O(n log(n))
// update e query - O(log^2(n))
a6b template < class T = int > struct bit2d {
acf
        vector <T> X;
a84
        vector < vector < T >> Y, t;
```

```
709
        int ub(vector<T>& v, T x) {
dde
            return upper_bound(v.begin(), v.end(), x) - v.begin();
9cc
        }
5cb
        bit2d(vector<pair<T, T>> v) {
            for (auto [x, y] : v) X.push_back(x);
2e1
fd4
            sort(X.begin(), X.end());
1ee
            X.erase(unique(X.begin(), X.end()), X.end());
d56
            t.resize(X.size() + 1);
d12
            Y.resize(t.size());
3d0
            sort(v.begin(), v.end(), [](auto a, auto b) {
e8f
                return a.second < b.second; });</pre>
            for (auto [x, y] : v) for (int i = ub(X, x); i < t.size();
   i += i\&-i
b75
                if (!Y[i].size() or Y[i].back() != y)
   Y[i].push_back(y);
            for (int i = 0; i < t.size(); i++) t[i].resize(Y[i].size()</pre>
   + 1);
8cc
        void update(T x, T y, T v) {
e78
2a9
            for (int i = ub(X, x); i < t.size(); i += i\&-i)
                for (int j = ub(Y[i], y); j < t[i].size(); j += j&-j)</pre>
   t[i][j] += v;
533
        }
5d2
        T query(T x, T y) {
966
            T ans = 0;
c54
            for (int i = ub(X, x); i; i = i\&-i)
4fb
                for (int j = ub(Y[i], y); j; j -= j&-j) ans += t[i][j];
ba7
            return ans;
62d
        }
46d
        T query (T x1, T y1, T x2, T y2) {
fcf
            return query (x2, y2) -query (x2, y1-1) -query (x1-1, y2)
   y2)+query(x1-1, y1-1);
232
6a7 };
6.3 BIT com update em range
// Operacoes 0-based
// query(l, r) retorna a soma de v[l..r]
// update(l, r, x) soma x em v[l..r]
//
// Complexidades:
```

```
// build - O(n)
// query - O(log(n))
// update - O(log(n))
e04 namespace bit {
        11 bit[2][MAX+2];
3ba
1a8
        int n:
        void build(int n2, int* v) {
61c
            n = n2:
1 e 3
535
            for (int i = 1; i <= n; i++)</pre>
                 bit [1] [min(n+1, i+(i\&-i))] += bit[1][i] += v[i-1];
edd
db0
        }
637
        ll get(int x, int i) {
b73
            11 \text{ ret} = 0;
360
            for (; i; i -= i&-i) ret += bit[x][i];
edf
            return ret;
99c
        }
20 c
        void add(int x, int i, ll val) {
503
            for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
bf6
        }
162
        11 get2(int p) {
            return get(0, p) * p + get(1, p);
с7с
153
        }
02a
        11 query(int 1, int r) {
ff5
            return get2(r+1) - get2(1);
633
        }
089
        void update(int 1, int r, ll x) {
e5f
            add(0, 1+1, x), add(0, r+2, -x);
f58
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
        }
e5f
f91 };
6.4 BIT-Sort Tree
// Tipo uma MergeSort Tree usando Bit
// Apesar da complexidade ser pior, fica melhor na pratica.
// query(1, r, k) retorna o numero de elementos menores que k
// no intervalo [1, r]
// Usa O(n log(n)) de memoria
//
```

// Complexidades:

 $// \text{ query - O(log^2(n))}$

// construir - O(n log^2(n))

```
6fa template < typename T > struct ms_bit {
1a8
        int n:
b2f
        vector < vector < T >> bit;
        ms_bit(vector < T > \& v) : n(v.size()), bit(n+1) {
899
            for (int i = 0; i < n; i++)</pre>
830
d51
                for (int j = i+1; j \le n; j += j\&-j)
dad
                     bit[j].push_back(v[i]);
535
            for (int i = 1; i <= n; i++)
                sort(bit[i].begin(), bit[i].end());
eec
        }
b4d
257
        int p_query(int i, T k) {
7c9
            int ret = 0:
be8
            for (i++; i; i -= i&-i)
                ret += lower_bound(bit[i].begin(), bit[i].end(), k) -
1bd
   bit[i].begin();
            return ret;
edf
6f9
690
        int query(int 1, int r, T k) {
83d
            return p_query(r, k) - p_query(l-1, k);
bcc
8d0 }:
```

6.5 Convex Hull Trick Dinamico

```
// para double, use LINF = 1/.0, div(a, b) = a/b
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
//
// O(log(n)) amortizado por insercao
// O(log(n)) por query
72c struct Line {
073
        mutable ll a, b, p;
        bool operator<(const Line& o) const { return a < o.a; }</pre>
8e3
        bool operator<(ll x) const { return p < x; }</pre>
abf
469 }:
326 struct dynamic_hull : multiset <Line, less <>> {
        11 div(ll a, ll b) {
a20
            return a / b - ((a ^ b) < 0 and a % b);
        }
a8a
bbb
        void update(iterator x) {
```

```
b2a
             if (next(x) == end()) x -> p = LINF;
772
             else if (x->a == next(x)->a) x->p = x->b >= next(x)->b ?
   LINF : -LINF;
             else x \rightarrow p = div(next(x) \rightarrow b - x \rightarrow b, x \rightarrow a - next(x) \rightarrow a);
424
0c4
        }
71c
         bool overlap(iterator x) {
f18
             update(x);
cfa
             if (next(x) == end()) return 0;
a4a
             if (x->a == next(x)->a) return x->b >= next(x)->b:
d40
             return x - p >= next(x) - p;
901
        }
176
         void add(ll a, ll b) {
1c7
             auto x = insert({a, b, 0});
             while (overlap(x)) erase(next(x)), update(x);
4ab
dbc
             if (x != begin() and !overlap(prev(x))) x = prev(x),
    update(x);
0fc
             while (x != begin() and overlap(prev(x)))
                 x = prev(x), erase(next(x)), update(x);
4d2
        }
48f
4ad
        11 query(ll x) {
229
             assert(!empty());
7d1
             auto 1 = *lower_bound(x);
d41 #warning cuidado com overflow!
aba
             return 1.a * x + 1.b;
3f5
        }
8f2 };
6.6 Convex Hull Trick Estatico
// adds tem que serem feitos em ordem de slope
// queries tem que ser feitas em ordem de x
//
// linear
4b5 struct CHT {
942
         int it;
ac1
        vector<ll> a. b:
45e
        CHT():it(0){}
0bb
        ll eval(int i. ll x){
93d
             return a[i]*x + b[i];
b2a
63a
        bool useless(){
a20
             int sz = a.size();
35f
             int r = sz-1, m = sz-2, 1 = sz-3;
```

```
d41 #warning cuidado com overflow!
d71
            return (b[1] - b[r])*(a[m] - a[1]) <
413
                (b[1] - b[m])*(a[r] - a[1]);
a0c
        }
bf4
        void add(ll A, ll B){
7f5
            a.push_back(A); b.push_back(B);
565
            while (!a.empty()){
233
                if ((a.size() < 3) || !useless()) break;</pre>
ecb
                a.erase(a.end() - 2);
                b.erase(b.end() - 2);
568
b21
            }
165
        }
81b
        ll get(ll x){
            it = min(it, int(a.size()) - 1);
d27
46a
            while (it+1 < a.size()){</pre>
                if (eval(it+1, x) > eval(it, x)) it++;
3c4
f97
                else break;
            }
fe9
420
            return eval(it, x);
88a
       }
303 };
6.7 DSU
// Une dois conjuntos e acha a qual conjunto um elemento pertence por
   seu id
// find e unite: O(a(n)) \sim = O(1) amortizado
8d3 struct dsu {
825
        vector < int > id, sz;
b33
        dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), id.end(), 0); }
0cf
        int find(int a) { return a == id[a] ? a : id[a] = find(id[a]);
  }
        void unite(int a, int b) {
440
            a = find(a), b = find(b);
605
            if (a == b) return;
d54
            if (sz[a] < sz[b]) swap(a, b);</pre>
956
6d0
            sz[a] += sz[b], id[b] = a;
ea7
       }
8e1 };
// DSU de bipartido
//
```

```
// Une dois vertices e acha a qual componente um vertice pertence
// Informa se a componente de um vertice e bipartida
//
// find e unite: O(log(n))
8d3 struct dsu {
6f7
        vector<int> id, sz, bip, c;
5b4
        dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {
db8
            iota(id.begin(), id.end(), 0);
f25
        }
ef0
        int find(int a) { return a == id[a] ? a : find(id[a]); }
f30
        int color(int a) { return a == id[a] ? c[a] : c[a] ^
   color(id[a]); }
440
        void unite(int a, int b) {
            bool change = color(a) == color(b);
263
            a = find(a), b = find(b);
605
            if (a == b) {
a89
4ed
                if (change) bip[a] = 0;
505
                return;
            }
32d
956
            if (sz[a] < sz[b]) swap(a, b);
efe
            if (change) c[b] = 1;
2cd
            sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
22b
        }
118 };
// DSU Persistente
//
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
//
// find e unite: O(log(n))
8d3 struct dsu {
        vector < int > id, sz, ti;
733
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) {
db8
            iota(id.begin(), id.end(), 0);
        }
aad
5e6
        int find(int a, int t) {
            if (id[a] == a or ti[a] > t) return a;
6ba
```

```
ea5
            return find(id[a], t);
6cb
        void unite(int a, int b, int t) {
84f
            a = find(a, t), b = find(b, t);
            if (a == b) return;
d54
956
            if (sz[a] < sz[b]) swap(a, b);
            sz[a] += sz[b], id[b] = a, ti[b] = t;
35d
513
       }
6c6 };
// DSU com rollback
// checkpoint(): salva o estado atual de todas as variaveis
// rollback(): retorna para o valor das variaveis para
// o ultimo checkpoint
//
// Sempre que uma variavel muda de valor, adiciona na stack
// find e unite: O(log(n))
// checkpoint: 0(1)
// rollback: O(m) em que m e o numero de vezes que alguma
// variavel mudou de valor desde o ultimo checkpoint
8d3 struct dsu {
825
        vector < int > id. sz:
27 c
        stack<stack<pair<int&, int>>> st;
98d
        dsu(int n) : id(n), sz(n, 1) {
1cc
            iota(id.begin(), id.end(), 0), st.emplace();
8cd
        void save(int &x) { st.top().emplace(x, x); }
bdf
        void checkpoint() { st.emplace(); }
30d
5cf
        void rollback() {
ba9
            while(st.top().size()) {
6bf
                auto [end, val] = st.top().top(); st.top().pop();
                end = val;
149
            }
f9a
25a
            st.pop();
       }
3c6
ef0
        int find(int a) { return a == id[a] ? a : find(id[a]); }
440
        void unite(int a, int b) {
```

```
605
            a = find(a), b = find(b);
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);
803
            save(sz[a]), save(id[b]);
            sz[a] += sz[b], id[b] = a;
6d0
1b9
c6e };
6.8 Li-Chao Tree
// Adiciona retas (ax+b), e computa o minimo entre as retas
// em um dado 'x'
// Cuidado com overflow!
// Se tiver overflow, tenta comprimir o 'x' ou usar
// convex hull trick
//
// O(log(MA-MI)), O(n) de memoria
5b0 template<11 MI = 11(-1e9), 11 MA = 11(1e9) > struct lichao {
        struct line {
12d
            ll a. b:
cef
            array<int, 2> ch;
fdf
            line(ll a_{-} = 0, ll b_{-} = LINF):
423
                 a(a_{-}), b(b_{-}), ch(\{-1, -1\}) \{\}
888
            11 operator ()(11 x) { return a*x + b; }
d1d
        };
17b
        vector<line> ln;
df8
        int ch(int p, int d) {
            if (ln[p].ch[d] == -1) {
e85
9af
                ln[p].ch[d] = ln.size();
cdc
                ln.emplace_back();
bc1
ef2
            return ln[p].ch[d];
86a
        }
        lichao() { ln.emplace_back(); }
021
        void add(line s, ll l=MI, ll r=MA, int p=0) {
c33
3e3
            11 m = (1+r)/2;
911
            bool L = s(1) < ln[p](1);
d37
            bool M = s(m) < ln[p](m);
03b
            bool R = s(r) < ln[p](r);
825
            if (M) swap(ln[p], s), swap(ln[p].ch, s.ch);
            if (s.b == LINF) return;
cac
            if (L != M) add(s, 1, m-1, ch(p, 0));
f6d
898
            else if (R != M) add(s, m+1, r, ch(p, 1));
76e
        }
```

```
092
        11 query(int x, 11 1=MI, 11 r=MA, int p=0) {
                                                                            0d7
            11 m = (1+r)/2, ret = ln[p](x);
                                                                            fa8
11b
9db
            if (ret == LINF) return ret;
                                                                            77f
529
            if (x < m) return min(ret, query(x, 1, m-1, ch(p, 0)));
                                                                            01e
            return min(ret, query(x, m+1, r, ch(p, 1)));
                                                                            89ъ
                                                                                    }
81a
fba
59b }:
                                                                            c06
                                                                            6b9
                                                                                        prop(p, 1, r);
                                                                            6f3
6.9 Li-Chao Tree - Lazy
                                                                            33b
                                                                            90d
// Sendo N = MA-MI:
                                                                            da9
// insert({a, b}) minimiza tudo com ax+b - O(log N)
                                                                            c55
// insert(\{a, b\}, 1, r) minimiza com ax+b no range [1, r] - 0(\log^2 N)
                                                                            953
                                                                                    }
// shift({a, b}) soma ax+b em tudo - O(1)
// shift({a, b}, l, r) soma ax+b no range [l, r] - O(log^2 N)
                                                                            5df
// query(x) retorna o valor da posicao x - O(\log N)
                                                                            6b9
                                                                                        prop(p, 1, r);
                                                                            90d
// No inicio eh tudo LINF, se inserir {0, 0} fica tudo 0
                                                                            911
                                                                            d37
// O(n log N) de memoria ; O(n) de memoria se nao usar as operacoes de
                                                                            03b
   range
                                                                            c3f
                                                                            cac
41c template <int MI = int(-1e9), int MA = int(1e9) > struct lichao {
                                                                            c49
        struct line {
b3a
                                                                            29e
12d
            ll a, b;
                                                                            ceb
            11 la, lb; // lazy
158
                                                                            a8e
cef
            arrav<int, 2> ch:
                                                                               r=MA) {
            line(ll a_{-} = 0, ll b_{-} = LINF):
fdf
                                                                            6b9
                                                                                        prop(p, 1, r);
b09
                a(a_{-}), b(b_{-}), la(0), lb(0), ch(\{-1, -1\})  {}
                                                                            2d3
            11 operator ()(11 x) { return a*x + b; }
888
                                                                            1dd
92e
        };
                                                                            90d
17b
        vector < line > ln;
                                                                            f1e
                                                                            952
df8
        int ch(int p, int d) {
                                                                            375
                                                                                    }
e85
            if (ln[p].ch[d] == -1) {
9af
                ln[p].ch[d] = ln.size();
                                                                            97a
cdc
                ln.emplace_back();
                                                                               r=MA) {
bc1
            }
                                                                            6b9
                                                                                        prop(p, 1, r);
ef2
            return ln[p].ch[d];
                                                                            90d
86a
                                                                            9a3
        lichao() { ln.emplace_back(); }
021
                                                                            ada
                                                                            505
                                                                                            return:
ceb
        void prop(int p, int 1, int r) {
                                                                            570
                                                                                        }
ff8
            if (ln[p].la == 0 and ln[p].lb == 0) return;
                                                                           1dd
1d3
            ln[p].a += ln[p].la, ln[p].b += ln[p].lb;
                                                                            fdd
579
            if (1 != r) {
                                                                            751
b9e
                int pl = ch(p, 0), pr = ch(p, 1);
```

```
ln[pl].la += ln[p].la, ln[pl].lb += ln[p].lb;
        ln[pr].la += ln[p].la, ln[pr].lb += ln[p].lb;
    ln[p].la = ln[p].lb = 0;
11 query(int x, int p=0, int l=MI, int r=MA) {
    ll ret = ln[p](x);
    if (ln[p].ch[0] == -1 and ln[p].ch[1] == -1) return ret;
    int m = 1 + (r-1)/2;
    if (x \le m) return min(ret, query(x, ch(p, 0), 1, m));
    return min(ret, querv(x, ch(p, 1), m+1, r));
void push(line s, int p, int l, int r) {
    int m = 1 + (r-1)/2;
    bool L = s(1) < ln[p](1);
    bool M = s(m) < ln[p](m);
    bool R = s(r) < ln[p](r);
    if (M) swap(ln[p].a, s.a), swap(ln[p].b, s.b);
    if (s.b == LINF) return;
    if (L != M) push(s, ch(p, 0), 1, m);
    else if (R != M) push(s, ch(p, 1), m+1, r);
void insert(line s, int a=MI, int b=MA, int p=0, int l=MI, int
    if (a \le 1 \text{ and } r \le b) \text{ return push}(s, p, l, r);
    if (b < 1 or r < a) return;</pre>
    int m = 1 + (r-1)/2;
    insert(s, a, b, ch(p, 0), 1, m);
    insert(s, a, b, ch(p, 1), m+1, r);
void shift(line s, int a=MI, int b=MA, int p=0, int l=MI, int
    int m = 1 + (r-1)/2;
    if (a \le 1 \text{ and } r \le b)
        ln[p].la += s.a, ln[p].lb += s.b;
    if (b < 1 or r < a) return;
    if (ln[p].b != LINF) {
        push(ln[p], ch(p, 0), 1, m);
```

6.10 MergeSort Tree

```
// Se for construida sobre um array:
//
        count(i, j, a, b) retorna quantos
        elementos de v[i..j] pertencem a [a, b]
//
        report(i, j, a, b) retorna os indices dos
//
//
        elementos de v[i..j] que pertencem a [a, b]
//
        retorna o vetor ordenado
// Se for construida sobre pontos (x, y):
//
        count(x1, x2, y1, y2) retorna quantos pontos
//
        pertencem ao retangulo (x1, y1), (x2, y2)
//
        report(x1, x2, y1, y2) retorna os indices dos pontos que
//
        pertencem ao retangulo (x1, y1), (x2, y2)
//
        retorna os pontos ordenados lexicograficamente
//
        (assume x1 \le x2, y1 \le y2)
//
// kth(y1, y2, k) retorna o indice do ponto com k-esimo menor
// x dentre os pontos que possuem y em [y1, y2] (0 based)
// Se guiser usar para achar k-esimo valor em range, construir
// com ms_tree t(v, true), e chamar kth(1, r, k)
//
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log(n))
// count - O(log(n))
// report - O(\log(n) + k) para k indices retornados
// kth - O(log(n))
c6c template <typename T = int> struct ms_tree {
        vector<tuple<T, T, int>> v;
6f7
1a8
        vector<vector<tuple<T, T, int>>> t; // {y, idx, left}
5ee
6ae
        vector <T> vy;
        ms_tree(vector < pair < T, T >> & vv) : n(vv.size()), t(4*n), vy(n) {
78c
e80
            for (int i = 0; i < n; i++) v.push_back({vv[i].first,</pre>
   vv[i].second, i});
fca
            sort(v.begin(), v.end());
```

```
224
             build(1, 0, n-1);
01a
             for (int i = 0; i < n; i++) vy[i] = get < 0 > (t[1][i+1]);
45e
dac
        ms_tree(vector<T>& vv, bool inv = false) { // inv: inverte
   indice e valor
             vector < pair < T, T >> v2;
8e8
e1e
             for (int i = 0; i < vv.size(); i++)</pre>
                 inv ? v2.push_back({vv[i], i}) : v2.push_back({i,
196
   vv[i]});
             *this = ms_tree(v2);
cca
f23
2c6
        void build(int p, int l, int r) {
1d2
             t[p].push_back({get<0>(v[1]), get<0>(v[r]), 0}); //
   \{\min_{x, \max_{x}} 0\}
5c8
             if (1 == r) return t[p].push_back({get<1>(v[1]),
   get <2>(v[1]), 0});
             int m = (1+r)/2;
ee4
bd9
             build (2*p, 1, m), build (2*p+1, m+1, r);
             int L = 0, R = 0;
32d
             while (t[p].size() <= r-l+1) {</pre>
a 0.3
68e
                 int left = get<2>(t[p].back());
                 if (L > m-1 \text{ or } (R+m+1 \le r \text{ and } t[2*p+1][1+R] \le
4aa
   t[2*p][1+L])) {
8cf
                     t[p].push_back(t[2*p+1][1 + R++]);
da0
                     get<2>(t[p].back()) = left;
5e2
                     continue;
ce0
                 }
249
                 t[p].push_back(t[2*p][1 + L++]);
339
                 get < 2 > (t[p].back()) = left + 1;
208
            }
2eb
        }
        int get_1(T y) { return lower_bound(vy.begin(), vy.end(), y) -
   vv.begin(); }
ebb
        int get_r(T y) { return upper_bound(vy.begin(), vy.end(), y) -
   vy.begin(); }
f62
        int count(T x1, T x2, T y1, T y2) {
902
             function < int (int, int, int) > dfs = [&] (int p, int 1, int
   r) {
                 if (1 == r \text{ or } x2 < get < 0 > (t[p][0]) \text{ or } get < 1 > (t[p][0])
7c6
   < x1) return 0;
2bb
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2)
   return r-1;
784
                 int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
                 return dfs(2*p, nl, nr) + dfs(2*p+1, l-nl, r-nr);
eb6
```

```
122
             };
                                                                              75f
                                                                              987
7cb
             return dfs(1, get_l(y1), get_r(y2));
f65
                                                                              e8d
002
        vector<int> report(T x1, T x2, T y1, T y2) {
                                                                              42d
4b8
             vector < int > ret:
                                                                              aa2
             function < void(int, int, int) > dfs = [&](int p, int l, int
                                                                               c51
85e
    r) {
                                                                              5fd
                                                                                       }
                 if (1 == r \text{ or } x2 < get<0>(t[p][0]) \text{ or } get<1>(t[p][0])
                                                                              ea6
882
                                                                               c13 };
    < x1) return;
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2) {
8da
                     for (int i = 1; i < r; i++)</pre>
e00
    ret.push_back(get<1>(t[p][i+1]));
505
                     return:
067
784
                 int nl = get < 2 > (t[p][1]), nr = get < 2 > (t[p][r]);
                 dfs(2*p, nl, nr), dfs(2*p+1, l-nl, r-nr);
194
                                                                              81f
12b
             };
8ad
             dfs(1, get_l(y1), get_r(y2));
                                                                              3fc
edf
             return ret;
                                                                              12b
668
        }
                                                                              9d9
985
        int kth(T y1, T y2, int k) {
                                                                              f8d
902
             function<int(int, int, int)> dfs = [&](int p, int 1, int
                                                                              4f0
   r) {
                                                                              94a
                 if (k >= r-1) {
150
                                                                              1f2
                     k \rightarrow r-1:
941
                                                                              2eb
daa
                     return -1:
                                                                              ba7
b8d
                 }
                                                                              013
8da
                 if (r-l == 1) return get<1>(t[p][l+1]);
                 int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
                                                                              614
784
                                                                              13b
072
                 int left = dfs(2*p, nl, nr);
                                                                              4c0 };
                 if (left != -1) return left;
3b6
04d
                 return dfs(2*p+1, l-nl, r-nr);
            };
a1b
                                                                              cdc
7cb
             return dfs(1, get_l(y1), get_r(y2));
635
                                                                              7cd
1ce };
                                                                               c96
                                                                              d4d
6.11 Min queue - deque
                                                                              d92
                                                                              7ae
// Tudo O(1) amortizado
                                                                              489
                                                                              656
1dc template < class T > struct minqueue {
                                                                              ef1
                                                                                       }
                                                                              787
248
        deque<pair<T, int>> q;
                                                                              23a
3fc
        void push(T x) {
                                                                              7f3
56e
             int ct = 1;
                                                                              19c
953
```

while (q.size() and x < q.front().first)</pre>

```
ct += q.front().second, q.pop_front();
            q.emplace_front(x, ct);
        void pop() {
            if (q.back().second > 1) q.back().second--;
            else q.pop_back();
        T min() { return q.back().first; }
6.12 Min queue - stack
// Tudo O(1) amortizado
557 template < class T > struct minstack {
        stack<pair<T, T>> s;
        void push(T x) {
            if (!s.size()) s.push({x, x});
            else s.emplace(x, std::min(s.top().second, x));
        T top() { return s.top().first; }
        T pop() {
            T ans = s.top().first;
            s.pop();
            return ans;
        int size() { return s.size(); }
        T min() { return s.top().second; }
1dc template < class T> struct minqueue {
        minstack <T> s1, s2;
        void push(T x) { s1.push(x); }
        void move() {
            if (s2.size()) return;
            while (s1.size()) {
                T x = s1.pop();
                s2.push(x);
            }
        T front() { return move(), s2.top(); }
        T pop() { return move(), s2.pop(); }
        int size() { return s1.size()+s2.size(); }
        T min() {
cd6
            if (!s1.size()) return s2.min();
```

```
// Funciona do C++11 pra cima
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
        tree_order_statistics_node_update>;
3a1
// para declarar:
// ord set < int > s:
// coisas do set normal funcionam:
// for (auto i : s) cout << i << endl;
// cout << s.size() << endl:
// k-esimo maior elemento O(log|s|):
// k=0: menor elemento
// cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
// cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int>> com o
// segundo parametro sendo algo para diferenciar
// os ementos iguais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k
```

6.14 Priority Queue DS

```
// Mantem updates aplicados em uma estrutura de dados
// que permita rollback e nao seja amortizada.
// Cada update possui uma prioridade,
// sendo possivel remover o update com maior prioridade.
// Os updates devem ser comutativos, ou seja, o estado
// da estrutura deve ser o mesmo independente da ordem
// que eles sejam aplicados.
//
// Complexidades:
// update - O(log(n) + T(n))
// query - T(n)
```

```
// pop - O(log(n) * T(n)) amortizado
// onde T(n) eh a complexidade do update
// assumes all priorities are distinct
945 template < typename DS, typename UPD > struct priority_queue_ds {
df4
a7e
        vector<tuple<UPD, int, int>> upd; // {u, p, idx_in_pos}
866
        set < pair < int , int >> st;
927
        vector < int > pos;
cf0
        priority_queue_ds(int n) : D(n) {}
6af
        void update(UPD u, int p) {
9ab
            D.update(u);
d07
            st.emplace(p, pos.size());
6ca
            upd.emplace_back(u, p, pos.size());
e3d
            pos.push_back(upd.size() - 1);
        }
6af
427
        int query(int a) {
            return D.find(a);
aa3
        }
2d3
42d
        void pop() {
25f
            int k = 1, min_p; // k = number of pops we will do
43e
            vector<tuple<UPD, int, int>> small, big;
639
            auto it = st.end():
231
            for (int qt = 0; qt++ < (k+1)/2;) {
049
                it--;
3ab
                 min_p = it->first;
80f
                int i = pos[it->second];
e82
                 if (qt > 1) big.push_back(upd[i]);
84b
                 k = max < int > (k, upd.size() - i);
b9a
            }
b3d
            for (int i = 0; i < k; i++) {</pre>
a62
                 D.rollback();
868
                 auto [u, p, idx] = upd.rbegin()[i];
86d
                 if (p < min_p) small.emplace_back(u, p, idx);</pre>
            }
588
23e
            st.erase(prev(st.end()));
623
            upd.erase(upd.end() - k, upd.end());
a 25
            small.insert(small.end(), big.rbegin(), big.rend());
            for (auto [u, p, idx] : small) {
06f
```

```
9ab
                D.update(u);
c8e
                upd.emplace_back(u, p, idx);
a7d
                pos[idx] = upd.size() - 1;
            }
ec7
bd1
        }
54a };
6.15 Range color
```

```
// update(1, r, c) colore o range [1, r] com a cor c,
// e retorna os ranges que foram coloridos {1, r, cor}
// query(i) returna a cor da posicao i
//
// Complexidades (para q operacoes):
// update - O(log(q)) amortizado
// query - O(log(q))
df6 template < typename T> struct color {
f0c
        set < tuple < int , int , T >> se;
071
        vector<tuple<int, int, T>> update(int 1, int r, T val) {
            auto it = se.upper_bound({r, INF, val});
9 c 4
753
            if (it != se.begin() and get<1>(*prev(it)) > r) {
e91
                auto [L, R, V] = *--it;
3f0
                se.erase(it);
bfd
                se.emplace(L, r, V), se.emplace(r+1, R, V);
683
            }
            it = se.lower_bound({1, -INF, val});
d9e
516
            if (it != se.begin() and get<1>(*prev(it)) >= 1) {
                auto [L, R, V] = *--it;
e91
3f0
                se.erase(it);
75a
                se.emplace(L, 1-1, V), it = se.emplace(1, R, V).first;
b65
            }
d7b
            vector < tuple < int , int , T >> ret;
            for (; it != se.end() and get<0>(*it) <= r; it =</pre>
   se.erase(it))
8c0
                ret.push_back(*it);
            se.emplace(1, r, val);
b4a
edf
            return ret;
b6c
        T query(int i) {
ff9
c31
            auto it = se.upper_bound({i, INF, T()});
            if (it == se.begin() or get<1>(*--it) < i) return -1; //</pre>
   nao tem
            return get <2>(*it);
53d
daf
9e9 };
```

6.16 RMQ < O(n), O(1) > - min queue

```
// O(n) pra buildar, query O(1)
// Se tiver varios minimos, retorna
// o de menor indice
1a5 template < typename T > struct rmq {
517
        vector <T> v;
fcc
        int n; static const int b = 30;
        vector<int> mask, t;
70e
183
        int op(int x, int y) { return v[x] \le v[y] ? x : y; }
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1)); }
6ad
        rmq() {}
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) {
43c
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {</pre>
2e5
                at = (at << 1) & ((1 << b) -1);
a61
c00
                 while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
c2f
ea4
            for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
39d
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
   i+(1<< j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
    t[n/b*(j-1)+i+(1<<(j-1))]);
41a
        }
e34
        int index_query(int 1, int r) {
27b
            if (r-l+1 \le b) return small(r, r-l+1);
            int x = 1/b+1, y = r/b-1;
e80
fd3
            if (x > y) return op(small(1+b-1), small(r));
            int j = msb(y-x+1);
a4e
ea3
            int ans = op(small(1+b-1), op(t[n/b*j+x],
    t[n/b*j+y-(1<<j)+1]));
be6
            return op(ans, small(r));
62a
093
        T query(int 1, int r) { return v[index_query(1, r)]; }
bab }:
6.17 SegTreap
// Muda uma posicao do plano, e faz query de operacao
// associativa e comutativa em retangulo
// Mudar ZERO e op
// Esparso nas duas coordenadas, inicialmente eh tudo ZERO
//
```

// Para query com distancia de manhattan <= d, faca

```
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Valores no X tem que ser de O ateh NX
// Para q operacoes, usa O(q log(NX)) de memoria, e as
// operacoes custa O(log(q) log(NX))
55b const int ZERO = INF;
560 const int op(int 1, int r) { return min(1, r); }
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T > struct treap {
3c9
        struct node {
b19
            node *1, *r;
ee1
            int p;
            pair<11, 11> idx; // {y, x}
850
36d
            T val, mi;
            node(ll x, ll y, T val_) : l(NULL), r(NULL), p(rng()),
bc2
                idx(pair(y, x)), val(val_), mi(val) {}
1b5
            void update() {
01e
                mi = val:
d6e
182
                if (1) mi = op(mi, 1->mi);
b68
                if (r) mi = op(mi, r \rightarrow mi);
282
            }
6e1
        };
bb7
        node* root;
        treap() { root = NULL; }
84b
        \simtreap() {
cec
609
            vector < node *> q = {root};
402
            while (a.size()) {
                node* x = q.back(); q.pop_back();
e5d
ee9
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                delete x;
653
            }
50e
        treap(treap&& t) : treap() { swap(root, t.root); }
225
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
bcf
            if (!l or !r) return void(i = 1 ? 1 : r);
986
80e
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
            else join(1, r->1, r->1), i = r;
fa0
            i->update();
bda
```

```
671
        void split(node* i, node*& 1, node*& r, pair<11, 11> idx) {
c82
26a
             if (!i) return void(r = 1 = NULL);
13c
             if (i->idx < idx) split(i->r, i->r, r, idx), l = i;
d26
             else split(i \rightarrow 1, l, i \rightarrow 1, idx), r = i;
             i->update();
bda
143
        }
d3b
        void update(ll x, ll y, T v) {
df9
             node *L, *M, *R;
             split(root, M, R, pair(y, x+1)), split(M, L, M, pair(y,
8b2
   x));
             if (M) M \rightarrow val = M \rightarrow mi = v:
1e4
9e5
             else M = new node(x, v, v):
69d
             join(L, M, M), join(M, R, root);
58e
        }
91b
        T query(ll ly, ll ry) {
df9
             node *L, *M, *R;
             split(root, M, R, pair(ry, LINF)), split(M, L, M, pair(ly,
1c0
   0));
             T ret = M ? M->mi : ZERO:
0f7
69d
             join(L, M, M), join(M, R, root);
             return ret:
edf
        }
1ae
bdf };
46a template < typename T> struct segtreap {
c4f
        vector < treap < T >> seg;
6e7
        vector < int > ch[2]:
e4e
        ll NX;
        segtreap(11 NX_{-}) : seg(1), NX(NX_{-}) \{ ch[0].push_back(-1), \}
    ch[1].push_back(-1); }
a71
        int get ch(int i, int d){
             if (ch[d][i] == -1) {
e51
2d6
                 ch[d][i] = seg.size();
                 seg.emplace_back();
23e
842
                 ch[0].push_back(-1), ch[1].push_back(-1);
3e1
             return ch[d][i];
968
        }
bb6
        T query(ll lx, ll rx, ll ly, ll ry, int p, ll l, ll r) {
10c
003
             if (rx < 1 or r < 1x) return ZERO;</pre>
fOf
             if (lx <= l and r <= rx) return seg[p].query(ly, ry);</pre>
             11 m = 1 + (r-1)/2:
e6a
```

```
354
            return op(query(lx, rx, ly, ry, get_ch(p, 0), l, m),
                                                                            c10
                    query(lx, rx, ly, ry, get_ch(p, 1), m+1, r));
                                                                            2c3
                                                                                    ll query(int a, int b, int p=1, int l=0, int r=n-1) {
060
a5e
                                                                            6b9
                                                                                        prop(p, 1, r);
f48
        T query(11 lx, 11 rx, 11 ly, 11 ry) { return query(1x, rx, ly,
                                                                            527
                                                                                        if (a <= 1 and r <= b) return seg[p];</pre>
   ry, 0, 0, NX); }
                                                                            786
                                                                                        if (b < 1 \text{ or } r < a) \text{ return } 0:
                                                                                        int m = (1+r)/2;
                                                                            ee4
249
        void update(ll x, ll y, T val, int p, ll l, ll r) {
                                                                            b1f
                                                                                        return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
            if (1 == r) return seg[p].update(x, y, val);
73c
                                                                            4c5
                                                                                    }
e6a
            11 m = 1 + (r-1)/2;
                                                                            cfb
                                                                                    11 update(int a, int b, int x, int p=1, int l=0, int r=n-1) {
            if (x <= m) update(x, y, val, get_ch(p, 0), 1, m);</pre>
                                                                            6b9
                                                                                        prop(p, 1, r);
cc5
5a2
            else update(x, y, val, get_ch(p, 1), m+1, r);
                                                                            9a3
                                                                                        if (a <= 1 and r <= b) {</pre>
980
            seg[p].update(x, y, val);
                                                                            b94
                                                                                            lazy[p] += x;
cc2
                                                                            6b9
                                                                                            prop(p, 1, r);
517
        void update(11 x, 11 y, T val) { update(x, y, val, 0, 0, NX); }
                                                                            534
                                                                                            return seg[p];
                                                                                        }
40a };
                                                                            821
                                                                            e9f
                                                                                        if (b < l or r < a) return seg[p];</pre>
                                                                            ee4
                                                                                        int m = (1+r)/2;
6.18 SegTree
                                                                                        return seg[p] = update(a, b, x, 2*p, 1, m) +
                                                                            fdb
                                                                            7fd
                                                                                            update(a, b, x, 2*p+1, m+1, r);
// Recursiva com Lazy Propagation
                                                                            75 c
                                                                                    }
// Query: soma do range [a, b]
                                                                            0af };
// Update: soma x em cada elemento do range [a, b]
// Pode usar a seguinte funcao para indexar os nohs:
// f(1, r) = (1+r) | (1!=r), usando 2N de memoria
                                                                            // Se tiver uma seg de max, da pra descobrir em O(log(n))
//
                                                                            // o primeiro e ultimo elemento >= val numa range:
// Complexidades:
// build - O(n)
                                                                            // primeira posicao >= val em [a, b] (ou -1 se nao tem)
// query - O(log(n))
                                                                            119 int get_left(int a, int b, int val, int p=1, int l=0, int r=n-1) {
// update - 0(log(n))
                                                                            6b9
                                                                                    prop(p, 1, r);
                                                                            f38
                                                                                    if (b < l or r < a or seg[p] < val) return -1;
aa4 namespace seg {
                                                                            205
                                                                                    if (r == 1) return 1:
005
        11 \text{ seg}[4*MAX], lazy[4*MAX];
                                                                            ee4
                                                                                    int m = (1+r)/2;
052
        int n. *v:
                                                                            753
                                                                                    int x = get_left(a, b, val, 2*p, l, m);
                                                                                    if (x != -1) return x:
                                                                            50e
d22
        11 build(int p=1, int l=0, int r=n-1) {
                                                                                    return get_left(a, b, val, 2*p+1, m+1, r);
                                                                            сЗс
3c7
            lazy[p] = 0;
                                                                            68c }
6cd
            if (1 == r) return seg[p] = v[1];
ee4
            int m = (1+r)/2:
                                                                            // ultima posicao >= val em [a, b] (ou -1 se nao tem)
193
            return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1, r);
                                                                            992 int get_right(int a, int b, int val, int p=1, int l=0, int r=n-1) {
c71
        }
                                                                            6b9
                                                                                    prop(p, 1, r);
0d8
        void build(int n2. int* v2) {
                                                                            f38
                                                                                    if (b < l or r < a or seg[p] < val) return -1;</pre>
680
            n = n2, v = v2;
                                                                            205
                                                                                    if (r == 1) return 1;
6f2
            build():
                                                                                    int m = (1+r)/2:
                                                                            ee4
acb
        }
                                                                                    int x = get_right(a, b, val, 2*p+1, m+1, r);
                                                                            1 b 1
ceb
        void prop(int p, int l, int r) {
                                                                                    if (x != -1) return x;
                                                                            50e
cdf
            seg[p] += lazy[p]*(r-l+1);
                                                                            6a7
                                                                                    return get_right(a, b, val, 2*p, 1, m);
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
2c9
                                                                            1b7 }
```

3c7

lazv[p] = 0;

```
// Se tiver uma seg de soma sobre um array nao negativo v, da pra
// descobrir em O(\log(n)) o maior j tal que v[i]+v[i+1]+...+v[j-1] <
6a9 int lower_bound(int i, ll& val, int p, int l, int r) {
        prop(p, 1, r);
6b9
6e8
        if (r < i) return n:
        if (i <= 1 and seg[p] < val) {</pre>
b5d
bff
            val -= seg[p];
041
            return n;
634
        }
Зсе
        if (1 == r) return 1;
ee4
       int m = (1+r)/2:
514
       int x = lower_bound(i, val, 2*p, 1, m);
ee0
        if (x != n) return x;
8b9
        return lower_bound(i, val, 2*p+1, m+1, r);
2b8 }
```

6.19 SegTree 2D Iterativa

```
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
//
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
//
// Complexidades:
// build - O(n^2)
// \text{ query - } O(\log^2(n))
// update - O(log^2(n))
731 int seg[2*MAX][2*MAX], n;
0a8 void build() {
919
        for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
            if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
c81
fe9
            if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
d51
        }
499 }
```

```
251 int query(int x1, int y1, int x2, int y2) {
         int ret = 0, y3 = y1 + n, y4 = y2 + n;
827
83e
         for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
             for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
0f2
                  if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
554
6b0
                  if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
c01
                  if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
5d4
                  if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
2d0
             }
edf
         return ret;
ff1 }
767 void update(int x, int y, int val) {
        int y2 = y += n;
66a
192
         for (x += n; x; x /= 2, y = y2) {
970
             if (x >= n) seg[x][y] = val;
ba9
             else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
             while (y /= 2) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
3b1
        }
68b
62e }
6.20 SegTree Beats
// \text{ query(a, b)} - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
// updatemin(a, b, x) faz com que v[i] \leftarrow min(v[i], x),
// para i em [a, b]
// updatemax faz o mesmo com max, e updatesum soma x
// em todo mundo do intervalo [a. b]
//
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - O(log^2 (n)) amortizado
// (se nao usar updatesum, fica log(n) amortizado)
7c6 #define f first
Oab #define s second
f39 namespace beats {
3c9
         struct node {
526
             int tam:
125
             ll sum, lazy; // lazy pra soma
4f3
             ll mi1, mi2, mi; // mi = #mi1
c61
             ll ma1, ma2, ma; // ma = #ma1
```

```
426
            node(11 x = 0) {
ba6
                sum = mi1 = ma1 = x:
b29
                mi2 = LINF, ma2 = -LINF;
                mi = ma = tam = 1:
62c
c60
                lazv = 0:
b00
770
            node(const node& 1. const node& r) {
a95
                sum = 1.sum + r.sum, tam = 1.tam + r.tam;
c60
                lazv = 0;
                if (1.mi1 > r.mi1) {
797
230
                     mi1 = r.mi1, mi = r.mi;
ea2
                     mi2 = min(1.mi1, r.mi2):
f1e
                } else if (1.mi1 < r.mi1) {</pre>
                     mi1 = 1.mi1. mi = 1.mi:
e34
4b3
                     mi2 = min(r.mi1, l.mi2);
ef2
                } else {
a39
                     mi1 = 1.mi1, mi = 1.mi+r.mi;
                     mi2 = min(1.mi2, r.mi2):
83d
a92
                if (1.ma1 < r.ma1) {</pre>
cd0
                     ma1 = r.ma1, ma = r.ma;
6a0
96d
                     ma2 = max(1.ma1, r.ma2);
                } else if (1.ma1 > r.ma1) {
3c0
                     ma1 = l.ma1, ma = l.ma;
ae0
                     ma2 = max(r.ma1, 1.ma2);
2ca
da8
                } else {
                     ma1 = l.ma1, ma = l.ma+r.ma;
db2
c05
                     ma2 = max(1.ma2. r.ma2):
                }
11c
1ba
            }
4b4
            void setmin(11 x) {
55e
                if (x >= ma1) return;
                sum += (x - ma1)*ma:
463
be5
                if (mi1 == ma1) mi1 = x:
                if (mi2 == ma1) mi2 = x:
0a0
b81
                ma1 = x:
0 c 3
            }
6cb
            void setmax(ll x) {
e25
                if (x <= mi1) return;</pre>
7e8
                sum += (x - mi1)*mi;
                if (ma1 == mi1) ma1 = x:
0bb
                if (ma2 == mi1) ma2 = x;
c32
1ff
                mi1 = x;
            }
a86
4cf
            void setsum(ll x) {
fe8
                mi1 += x, mi2 += x, ma1 += x, ma2 += x;
620
                sum += x*tam:
```

```
c46
                 lazv += x;
b53
47f
        };
        node seg[4*MAX];
62b
        int n, *v;
052
93ъ
        node build(int p=1, int l=0, int r=n-1) {
d84
             if (1 == r) return seg[p] = {v[1]};
             int m = (1+r)/2:
ee4
3d6
             return seg[p] = \{build(2*p, 1, m), build(2*p+1, m+1, r)\};
444
        }
0d8
        void build(int n2. int* v2) {
680
             n = n2. v = v2:
6f2
             build();
        }
acb
        void prop(int p, int 1, int r) {
ceb
8ce
             if (1 == r) return;
             for (int k = 0; k < 2; k++) {
abd
                 if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
d07
                 seg[2*p+k].setmin(seg[p].ma1);
843
f79
                 seg[2*p+k].setmax(seg[p].mi1);
585
431
             seg[p].lazy = 0;
7ee
055
        pair < pair < ll. 11 > . 11 > guery (int a. int b. int p=1. int l=0.
   int r=n-1) {
e07
             if (b < 1 \text{ or } r < a) \text{ return } \{\{LINF, -LINF\}, 0\}:
             if (a \le 1 \text{ and } r \le b) \text{ return } \{\{seg[p].mi1, seg[p].ma1\},
9be
   seg[p].sum};
6b9
             prop(p, 1, r);
             int m = (1+r)/2;
ee4
             auto L = query(a, b, 2*p, 1, m), R = query(a, b, 2*p+1.
e6f
   m+1. r):
             return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)}, L.s+R.s};
96d
e9d
2c8
        node updatemin(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
744
             if (b < l or r < a or seg[p].ma1 <= x) return seg[p];</pre>
309
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
                 seg[p].setmin(x);
ccd
534
                 return seg[p];
bbf
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2;
96a
             return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
4db
                              updatemin(a, b, x, 2*p+1, m+1, r)};
```

```
aad
044
        node updatemax(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
b59
             if (b < l \text{ or } r < a \text{ or } seg[p].mi1 >= x) return seg[p];
a9e
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
                 seg[p].setmax(x);
e8a
534
                 return seg[p];
             }
e9b
6b9
             prop(p, 1, r);
             int m = (1+r)/2:
ee4
             return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
ee3
98b
                              updatemax(a, b, x, 2*p+1, m+1, r)};
323
        }
        node updatesum(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
e9f
             if (b < 1 or r < a) return seg[p];
             if (a \le 1 \text{ and } r \le b) \{
9a3
8f4
                 seg[p].setsum(x);
534
                 return seg[p];
             }
596
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2;
7b6
             return seg[p] = \{updatesum(a, b, x, 2*p, 1, m),
483
                              updatesum(a, b, x, 2*p+1, m+1, r)};
111
        }
0d2 }:
```

6.21 SegTree Colorida

```
// Cada posicao tem um valor e uma cor
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
// query(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
// [a, b] que tem cor c
// paint(c1, c2, a, b) faz com que todo mundo
// em [a, b] que tem cor c1 passe a ter cor c2
//
// Complexidades:
// construir - O(n log(n)) espaco e tempo
// query - O(log(n))
// update - O(log(n))
// paint - O(log(n)) amortizado
04f struct seg_color {
3c9
        struct node {
```

```
b19
            node *1, *r;
0f9
            int cnt;
9ca
            ll val, lazv;
277
            node() : 1(NULL), r(NULL), cnt(0), val(0), lazy(0) {}
01e
            void update() {
d0a
                cnt = 0, val = 0;
bc4
                for (auto i : {1, r}) if (i) {
c89
                     i->prop();
281
                     cnt += i->cnt, val += i->val;
                }
684
554
            }
a9c
            void prop() {
2dd
                if (!lazv) return:
3f7
                val += lazy*(ll)cnt;
b64
                for (auto i : {1, r}) if (i) i->lazy += lazy;
c60
                lazv = 0;
e24
            }
514
        };
1a8
        int n:
9ъ0
        vector < node *> seg;
        seg_color(vector<pair<int, int>>& v, int c) : n(v.size()),
   seg(c, NULL) {
830
            for (int i = 0; i < n; i++)</pre>
9b7
                seg[v[i].second] = insert(seg[v[i].second], i,
   v[i].first, 0, n-1);
94a
        }
3c7
        \simseg_color() {
dde
            queue < node *> q;
3a6
            for (auto i : seg) q.push(i);
402
            while (q.size()) {
20b
                auto i = q.front(); q.pop();
dab
                if (!i) continue:
7 c 7
                q.push(i->1), q.push(i->r);
5ce
                delete i;
            }
c60
139
        }
40b
        node* insert(node* at, int idx, int val, int l, int r) {
1a4
            if (!at) at = new node();
            if (1 == r) return at->cnt = 1, at->val = val, at;
232
            int m = (1+r)/2;
ee4
            if (idx \le m) at->1 = insert(at->1, idx, val, 1, m);
137
3e6
            else at->r = insert(at->r, idx, val, m+1, r);
cff
            return at->update(), at;
        }
d6e
```

```
870
        11 query(node* at, int a, int b, int l, int r) {
            if (!at or b < 1 or r < a) return 0;</pre>
61b
d9f
            at->prop();
cb2
            if (a <= l and r <= b) return at->val;
            int m = (1+r)/2:
ee4
            return query(at->1, a, b, 1, m) + query(at->r, a, b, m+1,
4c4
   r);
8c3
        11 query(int c, int a, int b) { return query(seg[c], a, b, 0,
e54
        void update(node* at, int a, int b, int x, int l, int r) {
91c
fba
            if (!at or b < l or r < a) return;</pre>
d9f
            at->prop():
            if (a <= 1 and r <= b) {</pre>
9a3
e9a
                at -> lazy += x;
                return void(at->prop());
cb2
051
            }
ee4
            int m = (1+r)/2:
0b0
            update(at->1, a, b, x, 1, m), update(at->r, a, b, x, m+1,
   r);
7b4
            at->update();
9fd
a40
        void update(int c, int a, int b, int x) { update(seg[c], a, b,
   x, 0, n-1); }
        void paint(node*& from, node*& to, int a, int b, int l, int r)
70c
   {
10f
            if (to == from or !from or b < l or r < a) return;</pre>
e85
            from ->prop();
            if (to) to->prop();
889
9a3
            if (a \le 1 \text{ and } r \le b) \{
                if (!to) {
24d
38f
                     to = from;
140
                     from = NULL:
505
                     return:
e5f
ee4
                int m = (1+r)/2;
                paint(from->1, to->1, a, b, 1, m), paint(from->r,
1cb
   to->r, a, b, m+1, r);
72d
                to->update();
270
                delete from;
140
                from = NULL:
505
                return:
a0e
019
            if (!to) to = new node();
            int m = (1+r)/2;
ee4
            paint(from->1, to->1, a, b, 1, m), paint(from->r, to->r,
1 cb
   a. b. m+1. r):
```

```
45a
            from ->update(), to ->update();
4aa
471
        void paint(int c1, int c2, int a, int b) { paint(seg[c1],
   seg[c2], a, b, 0, n-1); }
293 }:
     SegTree Esparsa - Lazy
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
//
// Complexidades:
// build - 0(1)
// query - O(log(n))
// update - O(log(n))
aa4 namespace seg {
6de
        int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
e9a
        int get_l(int i){
3db
            if (L[i] == 0) L[i] = ptr++;
a96
            return L[i]:
b6e
        }
943
        int get_r(int i){
71b
            if (R[i] == 0) R[i] = ptr++;
283
            return R[i];
43a
        }
e71
        void build() { ptr = 2; }
        void prop(int p, int 1, int r) {
ceb
b77
            if (!lazy[p]) return;
            seg[p] = r-l+1 - seg[p];
76c
213
            if (1 != r) lazy[get_l(p)]^=lazy[p],
   lazy[get_r(p)]^=lazy[p];
            lazy[p] = 0;
3c7
        }
20b
158
        int query(int a, int b, int p=1, int l=0, int r=N-1) {
6b9
            prop(p, 1, r);
786
            if (b < 1 \text{ or } r < a) \text{ return } 0;
527
            if (a <= 1 and r <= b) return seg[p];</pre>
```

return query(a, b, get_1(p), 1, m)+query(a, b, get_r(p),

int m = (1+r)/2;

ee4

818

0d9

m+1, r);

}

```
51f
        int update(int a, int b, int p=1, int l=0, int r=N-1) {
6b9
            prop(p, 1, r);
e9f
            if (b < l or r < a) return seg[p];
            if (a <= 1 and r <= b) {</pre>
9a3
                lazy[p] ^= 1;
ab6
6b9
                prop(p, 1, r);
534
                return seg[p];
8e4
            int m = (1+r)/2;
ee4
            return seg[p] = update(a, b, get_l(p), l, m)+update(a, b,
43a
   get_r(p), m+1, r);
1dc
dc3 };
```

6.23 SegTree Esparsa - O(q) memoria

```
// Query: min do range [a, b]
// Update: troca o valor de uma posicao
// Usa O(q) de memoria para q updates
//
// Complexidades:
// query - 0(log(n))
// update - O(log(n))
13d template < typename T > struct seg {
3c9
        struct node {
            node* ch[2];
d53
970
            char d;
ca0
            T v;
            T mi:
c4e
d4e
            node(int d_, T v_, T val) : d(d_), v(v_) {
e71
                ch[0] = ch[1] = NULL;
d6e
                mi = val;
065
            node(node* x) : d(x->d), v(x->v), mi(x->mi) {
b32
                ch[0] = x -> ch[0], ch[1] = x -> ch[1];
c99
            }
cb7
01e
            void update() {
909
                mi = numeric_limits <T>::max();
                for (int i = 0; i < 2; i++) if (ch[i])
151
b5a
                     mi = min(mi, ch[i]->mi);
fe3
            }
530
        };
```

```
bb7
        node* root;
9c5
        char n;
ba7
        seg() : root(NULL), n(0) {}
512
        \simseg() {
4c0
            std::vector<node*> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
                if (!x) continue;
ee9
73f
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                delete x;
d3e
            }
d8c
        }
1a6
        char msb(T v, char l, char r) { // msb in range (l, r]
            for (char i = r; i > 1; i--) if (v>>i&1) return i;
8e4
            return -1;
daa
688
        }
        void cut(node* at, T v, char i) {
430
            char d = msb(v ^a at -> v, at -> d, i);
677
23b
            if (d == -1) return; // no need to split
            node* nxt = new node(at);
ebf
d43
            at -> ch[v>> d&1] = NULL:
34f
            at -> ch[!(v>>d&1)] = nxt;
150
            at -> d = d:
0ъ3
        }
6e5
        node* update(node* at, T idx, T val, char i) {
            if (!at) return new node(-1, idx, val);
c8c
d67
            cut(at, idx, i);
1a2
            if (at->d == -1) { // leaf
792
                at->mi = val;
ce6
                return at;
a6f
            }
b29
            bool dir = idx>>at->d&1;
c8f
            at->ch[dir] = update(at->ch[dir], idx, val, at->d-1);
7b4
            at->update();
            return at;
ce6
76d
85 c
        void update(T idx, T val) {
8f4
            while (idx >> n) n++:
61e
            root = update(root, idx, val, n-1);
        }
79d
9d8
        T query(node* at, T a, T b, T l, T r, char i) {
            if (!at or b < l or r < a) return numeric_limits<T>::max();
df0
            if (a <= 1 and r <= b) return at->mi;
fd3
```

```
841
            T m = 1 + (r-1)/2;
c85
            if (at->d < i) {</pre>
                if ((at->v>>i&1) == 0) return query(at, a, b, 1, m,
c59
   i-1);
                else return query(at, a, b, m+1, r, i-1);
ca4
934
373
            return min(query(at->ch[0], a, b, 1, m, i-1),
   query(at->ch[1], a, b, m+1, r, i-1));
2db
        T query(T 1, T r) { return query(root, 1, r, 0, (T(1) \le n) - 1,
034
   n-1); }
d7f }:
```

6.24 SegTree Iterativa

```
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - O(log(n))
6a4 int seg[2 * MAX];
1a8 int n;
0a8 void build() {
d15
        for (int i = n - 1; i; i--) seg[i] = seg[2*i] + seg[2*i+1];
9a8 }
4ea int query(int a, int b) {
       int ret = 0;
7c9
728
       for (a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
           if (a % 2 == 1) ret += seg[a];
4ea
244
           if (b \% 2 == 0) ret += seg[b];
ac0
edf
        return ret;
24a }
ff3 void update(int p, int x) {
37d
        seg[p += n] = x;
c8c
        while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
02d }
```

6.25 SegTree Iterativa com Lazy Propagation

```
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
aa4 namespace seg {
6db
        11 seg[2*MAX], lazy[2*MAX];
1a8
9ъ3
        ll junta(ll a, ll b) {
534
            return a+b;
e26
        }
        // soma x na posicao p de tamanho tam
1b4
        void poe(int p, ll x, int tam, bool prop=1) {
517
            seg[p] += x*tam;
6ae
            if (prop and p < n) lazy[p] += x;</pre>
8bc
        }
        // atualiza todos os pais da folha p
b1e
        void sobe(int p) {
d5a
            for (int tam = 2; p /= 2; tam *= 2) {
4ca
                seg[p] = junta(seg[2*p], seg[2*p+1]);
388
                poe(p, lazy[p], tam, 0);
acd
            }
        }
b76
        // propaga o caminho da raiz ate a folha p
        void prop(int p) {
a0a
076
            int tam = 1 << (LOG-1);</pre>
            for (int s = LOG; s; s--, tam /= 2) {
0a8
4b1
                int i = p >> s;
27 c
                if (lazy[i]) {
                    poe(2*i, lazy[i], tam);
860
e38
                    poe(2*i+1, lazy[i], tam);
b97
                    lazy[i] = 0;
de8
                }
3ed
            }
e29
        }
```

```
61 c
        void build(int n2, int* v) {
                                                                            16a
                                                                                    vector < Data > seg;
1e3
            n = n2:
                                                                            1a8
                                                                                    int n;
95f
            for (int i = 0; i < n; i++) seg[n+i] = v[i];
            for (int i = n-1; i; i--) seg[i] = junta(seg[2*i],
                                                                            d45
                                                                                    seg_pa(int n_) {
   seg[2*i+1]);
                                                                            e95
                                                                                        n = n:
                                                                            fc3
f4c
            for (int i = 0; i < 2*n; i++) lazy[i] = 0;
                                                                                        seg = vector < Data > (4*n);
8<sub>bb</sub>
                                                                            ce0
                                                                                    }
4f3
        11 query(int a, int b) {
                                                                                    void prop(int p, int 1, int r) {
                                                                            ceb
b73
            11 \text{ ret} = 0:
                                                                            d5a
                                                                                        int tam = r-1+1:
b48
            for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2) {
                                                                            c3f
                                                                                        11 &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r =
a8e
                if (a%2 == 1) ret = junta(ret, seg[a]);
                                                                               seg[p].set_r,
c58
                if (b%2 == 0) ret = junta(ret, seg[b]);
                                                                                            &add_a = seg[p].add_a, &add_r = seg[p].add_r;
                                                                            a1b
510
            }
edf
            return ret;
                                                                            c02
                                                                                        if (set_a != LINF) {
38b
                                                                            660
        }
                                                                                            set_a += add_a, set_r += add_r;
                                                                            06e
                                                                                            sum = set_a*tam + set_r*tam*(tam+1)/2;
a28
        void update(int a, int b, int x) {
                                                                            579
                                                                                            if (1 != r) {
            int a2 = a += n, b2 = b += n, tam = 1;
c2d
                                                                            ee4
                                                                                                int m = (1+r)/2;
Off
            for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
32a
                if (a\%2 == 1) poe(a, x, tam);
                                                                            886
                                                                                                 seg[2*p].set_a = set_a;
9da
                if (b\%2 == 0) poe(b, x, tam);
                                                                            358
                                                                                                 seg[2*p].set_r = set_r;
            }
9bc
                                                                            ed6
                                                                                                 seg[2*p].add_a = seg[2*p].add_r = 0;
0f7
            sobe(a2), sobe(b2);
adc
        }
                                                                            f0c
                                                                                                 seg[2*p+1].set_a = set_a + set_r * (m-l+1);
6dc }:
                                                                            471
                                                                                                 seg[2*p+1].set_r = set_r;
                                                                            d48
                                                                                                 seg[2*p+1].add_a = seg[2*p+1].add_r = 0;
                                                                            a97
                                                                                            }
6.26 SegTree PA
                                                                            823
                                                                                            set_a = LINF, set_r = 0;
                                                                            953
                                                                                            add_a = add_r = 0;
// Segtree de PA
                                                                            da7
                                                                                        } else if (add_a or add_r) {
// update_set(1, r, A, R) seta [1, r] para PA(A, R),
                                                                            18b
                                                                                            sum += add_a*tam + add_r*tam*(tam+1)/2;
// update_add soma PA(A, R) em [1, r]
                                                                            579
                                                                                            if (1 != r) {
// query(l, r) retorna a soma de [l, r]
                                                                            ee4
                                                                                                int m = (1+r)/2:
// PA(A, R) eh a PA: [A+R, A+2R, A+3R, ...]
                                                                            ff0
                                                                                                 seg[2*p].add_a += add_a;
//
                                                                            ec0
                                                                                                 seg[2*p].add_r += add_r;
// Complexidades:
// construir - O(n)
                                                                            06c
                                                                                                 seg[2*p+1].add_a += add_a + add_r * (m-l+1);
// update_set, update_add, query - O(log(n))
                                                                            a6d
                                                                                                 seg[2*p+1].add_r += add_r;
                                                                            8af
dc7 struct seg_pa {
                                                                            953
                                                                                            add_a = add_r = 0;
350
        struct Data {
                                                                                        }
                                                                            ab7
8f5
            ll sum:
                                                                                    }
                                                                            07f
662
            11 set_a, set_r, add_a, add_r;
            Data(): sum(0), set_a(LINF), set_r(0), add_a(0), add_r(0)
9b7
                                                                            0b7
                                                                                    int inter(pair<int, int> a, pair<int, int> b) {
```

98c

if (a.first > b.first) swap(a, b);

{}

};

eb6

```
eef
            return max(0, min(a.second, b.second) - b.first + 1);
628
be1
        11 set(int a, int b, ll aa, ll rr, int p, int l, int r) {
6b9
            prop(p, 1, r);
457
            if (b < l or r < a) return seg[p].sum;</pre>
            if (a <= 1 and r <= b) {</pre>
9a3
91c
                seg[p].set a = aa:
                seg[p].set_r = rr;
774
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
8ee
            }
ee4
            int m = (1+r)/2;
963
            int tam 1 = inter({1, m}, {a, b});
            return seg[p].sum = set(a, b, aa, rr, 2*p, 1, m) +
c34
                set(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
365
8e2
        }
f55
        void update_set(int 1, int r, 11 aa, 11 rr) {
6f7
            set(1, r, aa, rr, 1, 0, n-1);
913
5f6
        11 add(int a, int b, ll aa, ll rr, int p, int l, int r) {
6b9
            prop(p, 1, r);
457
            if (b < 1 or r < a) return seg[p].sum;</pre>
            if (a <= 1 and r <= b) {</pre>
9a3
359
                seg[p].add_a += aa;
1ee
                seg[p].add_r += rr;
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
            }
d19
ee4
            int m = (1+r)/2;
963
            int tam_1 = inter({1, m}, {a, b});
586
            return seg[p].sum = add(a, b, aa, rr, 2*p, 1, m) +
695
                add(a, b, aa + rr * tam_1, rr, 2*p+1, m+1, r);
904
        }
848
        void update add(int 1. int r. ll aa. ll rr) {
            add(1, r, aa, rr, 1, 0, n-1);
afa
81e
        }
f45
        11 query(int a, int b, int p, int l, int r) {
6b9
            prop(p, 1, r);
786
            if (b < 1 or r < a) return 0;</pre>
e9a
            if (a <= 1 and r <= b) return seg[p].sum;</pre>
            int m = (1+r)/2:
ee4
b1f
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
f6e
        11 query(int 1, int r) { return query(1, r, 1, 0, n-1); }
bfc.
bc4 };
```

6.27 SegTree Persistente

```
// SegTree de soma, update de somar numa posicao
//
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - O(log(n))
// update - O(log(n))
54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
6de const int MAXS = 2*MAX+UPD*LOG;
f6e namespace perseg {
bd6
        11 seg[MAXS];
f4e
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
        int n, *v;
3 c 4
        ll build(int p, int l, int r) {
6cd
            if (1 == r) return seg[p] = v[1];
855
            L[p] = cnt++, R[p] = cnt++;
            int m = (1+r)/2;
ee4
275
            return seg[p] = build(L[p], 1, m) + build(R[p], m+1, r);
39d
        void build(int n2, int* v2) {
0d8
680
            n = n2, v = v2;
856
            rt[0] = cnt++:
c50
            build(0, 0, n-1);
a2e
        }
f45
        11 query(int a, int b, int p, int l, int r) {
786
            if (b < 1 \text{ or } r < a) \text{ return } 0;
527
            if (a <= l and r <= b) return seg[p];</pre>
            int m = (1+r)/2;
ee4
            return query(a, b, L[p], 1, m) + query(a, b, R[p], m+1, r);
1ed
4d2
182
        11 query(int a, int b, int tt) {
            return query(a, b, rt[tt], 0, n-1);
c13
726
bb3
        11 update(int a, int x, int lp, int p, int l, int r) {
747
            if (1 == r) return seg[p] = seg[lp]+x;
ee4
            int m = (1+r)/2;
ab8
            if (a \le m)
b48
                return seg[p] = update(a, x, L[lp], L[p]=cnt++, l, m)
   + seg[R[p]=R[lp]];
```

```
8a9
            return seg[p] = seg[L[p]=L[lp]] + update(a, x, R[lp],
   R[p] = cnt ++ , m+1, r);
788
        int update(int a, int x, int tt=t) {
6f6
            update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
ab3
e0d
            return t;
d63
       }
26f };
     SegTree Persistente com Lazy
// Nao propaga, meio estranho de mexer, mas da
//
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, b, x, t) faz um update v[a..b]+=x a partir da
```

```
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
// build - O(n)
// query - 0(log(n))
// update - O(log(n))
54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
ab3 const int MAXS = 2*MAX + 4*UPD*LOG:
f6e namespace perseg {
9eb
        int seg[MAXS];
f4e
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
        int n, *v;
        int build(int p, int 1, int r) {
adf
6cd
            if (1 == r) return seg[p] = v[1];
855
            L[p] = cnt++, R[p] = cnt++;
ee4
            int m = (1+r)/2;
01d
            return seg[p] = max(build(L[p], 1, m), build(R[p], m+1,
   r));
ffd
       }
        void build(int n2, int *v2) {
0d8
            n = n2, v = v2;
680
            rt[0] = cnt++:
856
            build(0, 0, n-1);
c50
a2e
        }
976
        int query(int a, int b, int p, int l, int r) {
27b
            if (b < 1 or r < a) return -INF;</pre>
793
            if (a <= l and r <= b) return lazy[p] + seg[p];</pre>
ee4
            int m = (1+r)/2;
```

int ret = lazy[p] + max(query(a, b, L[p], 1, m), query(a,

7a2

```
b, R[p], m+1, r));
            return ret;
edf
9a7
        }
442
        int query(int a, int b, int tt) {
            return query(a, b, rt[tt], 0, n-1);
c13
a05
bc1
        int update(int a, int b, int x, int lp, int p, int l, int r) {
3f6
            tie(seg[p], lazy[p], L[p], R[p]) = {seg[lp], lazy[lp],
   L[lp], R[lp]};
847
            if (b < l or r < a) return seg[p] + lazy[p];</pre>
32a
            if (a \le 1 \text{ and } r \le b) \text{ return } seg[p] + (lazy[p] += x);
            int m = (1+r)/2:
ee4
24a
            seg[p] = max(update(a, b, x, L[lp], L[p] = cnt++, l, m),
bdb
                          update(a, b, x, R[lp], R[p] = cnt++, m+1, r);
1ed
            lazy[p] = lazy[lp];
1b7
            return seg[p] + lazy[p];
877
        }
        int update(int a, int b, int x, int tt=t) {
cbf
            assert(tt <= t):
aa8
661
            update(a, b, x, rt[tt], rt[++t]=cnt++, 0, n-1);
e0d
            return t:
        }
aad
f27 };
      Sparse Table
// Resolve RMQ
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - 0(1)
cca namespace sparse {
710
        int m[MAX2][MAX], n;
        void build(int n2, int* v) {
61 c
1e3
            n = n2;
            for (int i = 0; i < n; i++) m[0][i] = v[i];</pre>
78e
            for (int j = 1; (1<<j) <= n; j++) for (int i = 0; i+(1<<<math>j)
a1c
    <= n; i++)
5d5
                 m[j][i] = min(m[j-1][i], m[j-1][i+(1<<(j-1))]);
cae
        }
        int query(int a, int b) {
4ea
            int j = __builtin_clz(1) - __builtin_clz(b-a+1);
ee5
```

return min(m[i][a], m[i][b-(1<<i)+1]);</pre>

dc3

fba

}

7aa } // Resolve qualquer operacao associativa //

6.30 Sparse Table Disjunta

```
// MAX2 = log(MAX)
// Complexidades:
// build - O(n log(n))
// query - 0(1)
cca namespace sparse {
        int m[MAX2][2*MAX], n, v[2*MAX];
5f7
        int op(int a, int b) { return min(a, b); }
860
        void build(int n2, int* v2) {
1e3
            n = n2:
df4
            for (int i = 0; i < n; i++) v[i] = v2[i];
a84
            while (n&(n-1)) n++;
3d2
            for (int j = 0; (1<<j) < n; j++) {
1c0
                int len = 1<<j;</pre>
d9b
                for (int c = len; c < n; c += 2*len) {
332
                    m[j][c] = v[c], m[j][c-1] = v[c-1];
                    for (int i = c+1; i < c+len; i++) m[j][i] =</pre>
   op(m[j][i-1], v[i]);
                    for (int i = c-2; i >= c-len; i--) m[j][i] =
432
   op(v[i], m[j][i+1]);
eda
                }
            }
f4d
ce3
9e3
        int query(int 1, int r) {
f13
            if (1 == r) return v[1];
            int j = __builtin_clz(1) - __builtin_clz(l^r);
e6d
d67
            return op(m[j][1], m[j][r]);
        }
a7b
fd8 }
```

6.31 Splay Tree

```
// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO
// NODE MAIS PROFUNDO VISITADO
// Todas as operacoes sao O(log(n)) amortizado
// Se quiser colocar mais informacao no node,
// mudar em 'update'
538 template < typename T > struct splaytree {
3c9
        struct node {
183
            node *ch[2], *p;
```

```
e4d
            int sz:
f48
            T val;
            node(T v) {
da0
696
                 ch[0] = ch[1] = p = NULL;
a26
                 sz = 1:
250
                 val = v;
2d0
            }
01e
            void update() {
a26
                 sz = 1;
c7c
                 for (int i = 0; i < 2; i++) if (ch[i]) {
d5f
                     sz += ch[i] -> sz;
486
                 }
f45
            }
aa3
        };
bb7
        node* root;
fbc
        splaytree() { root = NULL; }
214
        splaytree(const splaytree& t) {
cbf
             throw logic_error("Nao copiar a splaytree!");
1f1
        }
891
        \simsplaytree() {
609
            vector < node *> q = {root};
402
             while (q.size()) {
                 node* x = q.back(); q.pop_back();
e5d
ee9
                 if (!x) continue;
73f
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                 delete x;
d3e
            }
837
        }
94f
        void rotate(node* x) { // x vai ficar em cima
d9b
             node *p = x->p, *pp = p->p;
ecf
            if (pp) pp - ch[pp - ch[1] == p] = x;
            bool d = p \rightarrow ch[0] == x;
286
d63
            p - ch[!d] = x - ch[d], x - ch[d] = p;
bad
            if (p->ch[!d]) p->ch[!d]->p = p;
fc2
            x->p = pp, p->p = x;
            p->update(), x->update();
1ea
007
3fa
        node* splay(node* x) {
a39
            if (!x) return x;
4ea
            root = x:
3cf
             while (x->p) {
d9b
                 node *p = x->p, *pp = p->p;
359
                 if (!pp) return rotate(x), x; // zig
еЗс
                 if ((pp->ch[0] == p)^(p->ch[0] == x))
```

```
a2b
                    rotate(x), rotate(x); // zigzag
4b2
                else rotate(p), rotate(x); // zigzig
028
            }
ea5
            return x;
379
        }
313
        node* insert(T v, bool lb=0) {
b64
            if (!root) return lb ? NULL : root = new node(v);
002
            node *x = root, *last = NULL;;
31e
            while (1) {
5d7
                bool d = x -> val < v;
Ofd
                if (!d) last = x;
c2e
                if (x->val == v) break;
c16
                if (x->ch[d]) x = x->ch[d];
4e6
                else {
dea
                    if (lb) break;
055
                    x \rightarrow ch[d] = new node(v);
99c
                    x - ch[d] - p = x;
30e
                    x = x - ch[d];
c2b
                    break:
                }
68a
            }
1ab
0b6
            splay(x);
61c
            return lb ? splay(last) : x;
622
        }
сОс
        int size() { return root ? root->sz : 0; }
2ca
        int count(T v) { return insert(v, 1) and root->val == v; }
        node* lower_bound(T v) { return insert(v, 1); }
111
26b
        void erase(T v) {
446
            if (!count(v)) return;
bce
            node *x = root, *1 = x -> ch[0];
            if (!1) {
268
8b1
                root = x -> ch[1];
32e
                if (root) root->p = NULL;
8f3
                return delete x:
a86
5e7
            root = 1, 1->p = NULL;
902
            while (1->ch[1]) 1 = 1->ch[1];
bab
            splay(1);
f0e
            1 - ch[1] = x - ch[1];
7d9
            if (1->ch[1]) 1->ch[1]->p = 1;
bf0
            delete x:
62a
            1->update();
007
24a
        int order_of_key(T v) {
62b
            if (!lower_bound(v)) return root ? root->sz : 0;
1cc
            return root -> ch [0] ? root -> ch [0] -> sz : 0;
b00
        }
```

```
db6
        node* find_by_order(int k) {
084
            if (k >= size()) return NULL;
52f
            node* x = root;
31e
             while (1) {
20f
                 if (x->ch[0] \text{ and } x->ch[0]->sz >= k+1) x = x->ch[0]:
4e6
a1c
                     if (x->ch[0]) k -= x->ch[0]->sz;
                     if (!k) return splay(x);
1dc
eb8
                     k--, x = x->ch[1];
                }
aca
e01
            }
0de
        }
19c
        T min() {
52f
             node* x = root;
6f6
             while (x->ch[0]) x = x->ch[0]; // max -> ch[1]
3e9
            return splay(x)->val;
70e
        }
4ff };
      Splay Tree Implicita
// vector da NASA
// Um pouco mais rapido q a treap
// O construtor a partir do vector
// eh linear, todas as outras operacoes
// custam O(log(n)) amortizado
081 template < typename T > struct splay {
3c9
        struct node {
183
            node *ch[2], *p;
e4d
            int sz;
875
            T val, sub, lazy;
aa6
            bool rev;
da0
            node(T v) {
696
                 ch[0] = ch[1] = p = NULL;
                 sz = 1;
a26
1e4
                 sub = val = v;
c60
                lazv = 0;
b67
                 rev = false;
48f
a9c
            void prop() {
0ec
                 if (lazy) {
924
                     val += lazy, sub += lazy*sz;
091
                     if (ch[0]) ch[0]->lazy += lazy;
1a8
                     if (ch[1]) ch[1]->lazy += lazy;
a98
                }
1bb
                if (rev) {
```

```
80a
                     swap(ch[0], ch[1]);
                     if (ch[0]) ch[0]->rev ^= 1;
628
adc
                    if (ch[1]) ch[1]->rev ^= 1;
                }
30a
a32
                lazy = 0, rev = 0;
            }
6bf
            void update() {
01e
0 c 3
                sz = 1, sub = val;
c7c
                for (int i = 0; i < 2; i++) if (ch[i]) {
0.5f
                     ch[i]->prop();
d5f
                     sz += ch[i] -> sz;
4a1
                     sub += ch[i]->sub:
6c1
                }
e98
b4a
        };
bb7
        node* root;
        splay() { root = NULL; }
5d9
        splay(node* x) {
9b1
4ea
            root = x:
32e
            if (root) root->p = NULL;
371
1b7
        splay(vector < T > v) { // O(n)}
950
            root = NULL:
806
            for (T i : v) {
                node* x = new node(i);
2a0
bd1
                x - > ch[0] = root:
37a
                if (root) root->p = x;
4ea
                root = x;
                root ->update();
a0a
17c
            }
c6b
        }
        splay(const splay& t) {
a9e
e62
            throw logic_error("Nao copiar a splay!");
d4d
5ab
        \simsplav() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue:
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
73f
bf0
                delete x;
            }
d3e
        }
a1c
        int size(node* x) { return x ? x->sz : 0; }
73c
```

```
94f
         void rotate(node* x) { // x vai ficar em cima
d9b
             node *p = x \rightarrow p, *pp = p \rightarrow p;
ecf
             if (pp) pp - ch[pp - ch[1] == p] = x;
286
             bool d = p \rightarrow ch[0] == x;
d63
             p - ch[!d] = x - ch[d], x - ch[d] = p;
             if (p->ch[!d]) p->ch[!d]->p = p;
bad
fc2
             x->p = pp, p->p = x;
             p->update(), x->update();
1ea
007
6a0
         node* splaya(node* x) {
a39
             if (!x) return x;
be6
             root = x, x->update();
3cf
             while (x->p) {
d9b
                  node *p = x->p, *pp = p->p;
                  if (!pp) return rotate(x), x; // zig
359
                  if ((pp->ch[0] == p)^(p->ch[0] == x))
еЗс
a2b
                      rotate(x), rotate(x); // zigzag
4b2
                  else rotate(p), rotate(x); // zigzig
028
ea5
             return x;
         }
21a
a7f
         node* find(int v) {
a2e
             if (!root) return NULL;
52f
             node *x = root;
6cd
             int key = 0;
31e
             while (1) {
857
                  x->prop();
ba1
                  bool d = key + size(x->ch[0]) < v;
877
                  if (\text{kev} + \text{size}(x->\text{ch}[0]) != v \text{ and } x->\text{ch}[d]) {
15e
                      if (d) key += size(x->ch[0])+1;
                      x = x - > ch[d]:
30e
a30
                  } else break;
             }
3c3
152
             return splaya(x);
f19
сОс
         int size() { return root ? root->sz : 0; }
c26
         void join(splay<T>& 1) { // assume que l < *this</pre>
690
             if (!size()) swap(root, l.root);
579
             if (!size() or !l.size()) return:
bee
             node* x = 1.root;
             while (1) {
31e
857
                  x->prop();
34d
                  if (!x->ch[1]) break;
bd8
                  x = x -> ch[1];
fa3
147
             1.splaya(x), root->prop(), root->update();
42b
             x \rightarrow ch[1] = root, x \rightarrow ch[1] \rightarrow p = x:
```

```
0aa
            root = 1.root, 1.root = NULL;
a0a
            root ->update();
7e6
5ed
        node* split(int v) { // retorna os elementos < v</pre>
398
            if (v <= 0) return NULL:</pre>
060
            if (v >= size()) {
f87
                node* ret = root:
                root = NULL;
950
8c9
                ret->update();
edf
                return ret;
d0f
            }
adc
            find(v);
a59
            node* 1 = root -> ch[0]:
4df
            root -> ch[0] = NULL;
5a3
            if (1) 1->p = NULL;
            root ->update();
a0a
792
            return 1;
        }
826
        T& operator [](int i) {
511
9d4
            find(i):
ae0
            return root -> val;
829
231
        void push_back(T v) { // 0(1)
a01
            node* r = new node(v);
0de
            r \rightarrow ch[0] = root;
            if (root) root->p = r;
b11
b13
            root = r, root->update();
315
        }
b7a
        T query(int 1, int r) {
95f
            splay <T > M(split(r+1));
5ff
            splay <T> L(M.split(1));
d1c
            T ans = M.root->sub;
49c
            M.join(L), join(M);
ba7
            return ans:
ca3
41f
        void update(int 1, int r, T s) {
95f
            splay <T > M(split(r+1));
5ff
            splay <T > L(M.split(1));
996
            M.root->lazy += s;
49c
            M. join(L), join(M);
9e9
        }
8c1
        void reverse(int 1, int r) {
95f
            splay <T > M(split(r+1));
5ff
            splay <T> L(M.split(1));
945
            M.root->rev ^= 1;
49c
            M. join(L), join(M);
        }
c1a
```

```
2fb
        void erase(int 1, int r) {
95f
            splay <T > M(split(r+1));
5ff
            splay <T> L(M.split(1));
dcc
            join(L);
        }
68e
a35 };
6.33 Split-Merge Set
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge, que custa O(log(N)) amortizado
// Usa O(min(N, n log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
2dc template < typename T, bool MULTI = false, typename SIZE_T = int > struct
   sms {
3c9
        struct node {
b19
            node *1, *r;
15f
            SIZE T cnt:
658
            node() : 1(NULL), r(NULL), cnt(0) {}
01e
            void update() {
a 0 1
                cnt = 0:
d8a
                if (1) cnt += 1->cnt;
e49
                if (r) cnt += r->cnt;
74d
            }
        };
84f
bb7
        node* root;
fd0
        T N;
f34
        sms() : root(NULL), N(0) {}
83b
        sms(T v) : sms() { while (v >= N) N = 2*N+1; }
        sms(const sms& t) : root(NULL), N(t.N) {
5e1
3af
            for (SIZE_T i = 0; i < t.size(); i++) {</pre>
a0f
                T at = t[i];
e6d
                SIZE_T qt = t.count(at);
                insert(at, qt);
a43
f42
                i += qt-1;
1e9
ea8
        }
a 96
        sms(initializer_list<T> v) : sms() { for (T i : v) insert(i); }
2dd
        \simsms() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
```

```
ee9
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                delete x;
653
            }
f0d
        }
fdc
        friend void swap(sms& a, sms& b) {
49e
            swap(a.root, b.root), swap(a.N, b.N);
984
83e
        sms& operator =(const sms& v) {
768
            sms tmp = v;
420
            swap(tmp, *this);
357
            return *this:
e9b
        SIZE_T size() const { return root ? root->cnt : 0; }
d06
17f
        SIZE_T count(node* x) const { return x ? x->cnt : 0; }
75a
        void clear() {
0a0
            sms tmp;
4ac
            swap(*this, tmp);
fcb
       }
a06
        void expand(T v) {
bc3
            for (; N < v; N = 2*N+1) if (root) {
                node* nroot = new node();
63 c
956
                nroot ->1 = root;
897
                root = nroot:
a0a
                root ->update();
            }
dd9
9f0
       }
b14
        node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
            if (!at) at = new node();
1a4
893
            if (1 == r) {
435
                at->cnt += qt;
beb
                if (!MULTI) at->cnt = 1:
ce6
                return at;
            }
a53
841
            T m = 1 + (r-1)/2;
a02
            if (idx <= m) at->1 = insert(at->1, idx, qt, 1, m);
8d9
            else at->r = insert(at->r, idx, qt, m+1, r);
cff
            return at->update(), at;
83b
        void insert(T v, SIZE_T qt=1) { // insere 'qt' ocorrencias de
cf7
   , 17 ,
882
            if (qt <= 0) return erase(v, -qt);</pre>
72b
            assert(v >= 0);
f52
            expand(v);
            root = insert(root, v, qt, 0, N);
5e9
```

```
f62
        }
        node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
f06
28c
            if (!at) return at;
54b
            if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt - qt;
4e6
841
                T m = 1 + (r-1)/2:
281
                if (idx \le m) at->1 = erase(at->1, idx, qt, 1, m);
ba1
                else at->r = erase(at->r, idx, qt, m+1, r);
7b4
                at->update();
d3d
135
            if (!at->cnt) delete at, at = NULL;
ce6
            return at:
e1f
43d
        void erase(T v, SIZE_T qt=1) { // remove 'qt' ocorrencias de
9c3
            if (v < 0 \text{ or } v > N \text{ or } !qt) \text{ return};
            if (qt < 0) insert(v, -qt);</pre>
9dc
b1d
            root = erase(root, v, qt, 0, N);
        }
b32
        void erase_all(T v) { // remove todos os 'v'
868
347
            if (v < 0 \text{ or } v > N) return:
            root = erase(root, v, numeric_limits < SIZE_T >:: max(), 0, N);
9f2
569
        }
0fe
        SIZE_T count(node* at, T a, T b, T 1, T r) const {
61b
            if (!at or b < 1 or r < a) return 0:
            if (a <= 1 and r <= b) return at->cnt:
0fe
841
            T m = 1 + (r-1)/2:
84a
            return count(at->1, a, b, 1, m) + count(at->r, a, b, m+1,
   r);
4e6
        }
        SIZE_T count(T v) const { return count(root, v, v, 0, N); }
0a9
        SIZE T order of kev(T v) { return count(root, 0, v-1, 0, N): }
ffc
        SIZE_T lower_bound(T v) { return order_of_key(v); }
df2
        const T operator [](SIZE_T i) const { // i-esimo menor elemento
e68
809
            assert(i >= 0 and i < size());
c43
            node* at = root:
4a5
            T 1 = 0, r = N;
            while (1 < r) {
40 c
841
                T m = 1 + (r-1)/2;
5c2
                if (count(at->1) > i) at = at->1, r = m;
4e6
                else {
b4a
                    i -= count(at->1);
                    at = at - r; 1 = m+1;
ded
fa6
                }
```

```
41a
792
            return 1;
67f
78c
        node* merge(node* 1. node* r) {
            if (!1 or !r) return 1 ? 1 : r;
347
504
            if (!1->1 \text{ and } !1->r) { // folha}
599
                if (MULTI) 1->cnt += r->cnt;
55d
                delete r;
792
                return 1:
92c
            }
f58
            1->1 = merge(1->1, r->1), 1->r = merge(1->r, r->r);
f4f
            1->update(), delete r:
792
            return 1;
06a
f59
        void merge(sms& s) { // mergeia dois sets
068
            if (N > s.N) swap(*this, s);
785
            expand(s.N);
938
            root = merge(root, s.root);
            s.root = NULL:
ee2
2f6
        }
        node* split(node*& x, SIZE_T k) {
dc6
7ca
            if (k <= 0 or !x) return NULL;</pre>
640
            node* ret = new node():
            if (!x->l \text{ and } !x->r) x->cnt -= k, ret->cnt += k;
386
4e6
            else {
                if (k \le count(x->1)) ret->1 = split(x->1, k);
85e
4e6
06f
                     ret->r = split(x->r, k - count(x->1));
                     swap(x->1, ret->1);
cfd
63b
674
                ret ->update(), x ->update();
379
            }
            if (!x->cnt) delete x, x = NULL;
d5b
edf
            return ret;
f18
        }
02b
        void split(SIZE_T k, sms& s) { // pega os 'k' menores
e63
            s.clear():
6e5
            s.root = split(root, min(k, size()));
            s.N = N:
e3c
9a6
        // pega os menores que 'k'
        void split_val(T k, sms& s) { split(order_of_key(k), s); }
131
2d2 };
```

6.34 SQRT Tree

```
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entag a guery eh praticamente O(1))
//
// build - O(n log log n)
// query - O(log log n)
97a namespace sqrtTree {
        int n, *v;
052
        int pref[4][MAX], sulf[4][MAX], getl[4][MAX], entre[4][MAX],
ec7
   sz[4]:
5f7
        int op(int a, int b) { return min(a, b); }
        inline int getblk(int p, int i) { return (i-getl[p][i])/sz[p];
c72
 }
2c6
        void build(int p, int l, int r) {
            if (1+1 >= r) return;
bc8
368
            for (int i = 1: i \le r: i++) getl[p][i] = 1:
f16
            for (int L = 1; L <= r; L += sz[p]) {</pre>
191
                int R = min(L+sz[p]-1, r);
89 c
                pref[p][L] = v[L], sulf[p][R] = v[R];
59f
                for (int i = L+1; i <= R; i++) pref[p][i] =</pre>
    op(pref[p][i-1], v[i]);
d9a
                for (int i = R-1; i >= L; i--) sulf[p][i] = op(v[i],
    sulf[p][i+1]);
221
                build(p+1, L, R);
c7b
695
            for (int i = 0; i <= sz[p]; i++) {</pre>
ca5
                int at = entre[p][l+i*sz[p]+i] = sulf[p][l+i*sz[p]];
759
                for (int j = i+1; j <= sz[p]; j++)</pre>
    entre[p][1+i*sz[p]+j] = at =
23a
                         op(at, sulf[p][l+j*sz[p]]);
            }
c51
        }
861
0d8
        void build(int n2, int* v2) {
680
            n = n2, v = v2;
            for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
44 c
            build(0, 0, n-1);
c50
940
        }
9e3
        int query(int 1, int r) {
792
            if (1+1 \ge r) return 1 = r ? v[1] : op(v[1], v[r]);
1 ba
            int p = 0:
4ba
            while (getblk(p, 1) == getblk(p, r)) p++;
            int ans = sulf[p][1], a = getblk(p, 1)+1, b = getblk(p,
9e4
```

```
r)-1;
8bf
            if (a \le b) ans = op(ans, entre[p][get1[p][1]+a*sz[p]+b]);
            return op(ans, pref[p][r]);
dea
589
        }
8ff }
6.35 Treap
// Todas as operacoes custam
// O(log(n)) com alta probabilidade, exceto meld
// meld custa O(log^2 n) amortizado com alta prob.,
// e permite unir duas treaps sem restricao adicional
// Na pratica, esse meld tem constante muito boa e
// o pior caso eh meio estranho de acontecer
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T > struct treap {
3c9
        struct node {
b19
            node *1, *r;
284
            int p, sz;
36d
            T val, mi;
            node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
   mi(v) {}
01e
            void update() {
a26
                sz = 1:
d6e
                mi = val:
bd7
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
a54
                if (r) sz += r->sz, mi = min(mi, r->mi);
660
            }
c1b
        };
bb7
        node* root;
        treap() { root = NULL; }
84b
2d8
        treap(const treap& t) {
465
            throw logic_error("Nao copiar a treap!");
1e9
        }
        \simtreap() {
cec
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
1 c 7
bf0
                delete x;
653
            }
```

```
50e
        }
        int size(node* x) { return x ? x->sz : 0; }
73c
        int size() { return size(root); }
b2b
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
bcf
986
             if (!1 or !r) return void(i = 1 ? 1 : r);
80e
             if (1->p > r->p) join(1->r, r, 1->r), i = 1;
fa0
             else join(1, r->1, r->1), i = r;
             i->update();
bda
671
        }
        void split(node* i, node*& 1, node*& r, T v) {
ece
26a
             if (!i) return void(r = l = NULL);
f05
             if (i\rightarrow val < v) split(i\rightarrow r, i\rightarrow r, r, v), l = i:
             else split(i->1, 1, i->1, v), r = i;
807
bda
             i->update();
        }
2cd
3fc
        void split_leg(node* i, node*& 1, node*& r, T v) {
             if (!i) return void(r = 1 = NULL);
26a
181
             if (i->val <= v) split_leq(i->r, i->r, r, v), l = i;
58f
             else split_leq(i \rightarrow 1, l, i \rightarrow 1, v), r = i;
bda
             i->update();
70f
e13
        int count(node* i, T v) {
6b4
             if (!i) return 0;
352
             if (i->val == v) return 1;
8d0
             if (v < i->val) return count(i->1, v);
4d0
             return count(i->r, v);
5e6
26d
        void index_split(node* i, node*& 1, node*& r, int v, int key =
   0) {
26a
             if (!i) return void(r = 1 = NULL);
c10
             if (\text{key} + \text{size}(i->1) < v) index_split(i->r, i->r, r, v,
   key+size(i->1)+1), l = i;
e5a
             else index_split(i->1, 1, i->1, v, key), r = i;
bda
             i->update();
ccf
        }
a1f
        int count(T v) {
e06
             return count(root, v);
980
c27
        void insert(T v) {
980
             if (count(v)) return;
031
             node *L, *R;
             split(root, L, R, v);
d42
585
             node* at = new node(v);
59f
             join(L, at, L);
a 28
             join(L, R, root);
37 c
        }
```

```
26b
        void erase(T v) {
                                                                              d3b
                                                                                                   if (r) r->lazy += lazy;
                                                                                              }
df9
            node *L, *M, *R;
                                                                              cea
b6b
            split_leq(root, M, R, v), split(M, L, M, v);
                                                                              1bb
                                                                                              if (rev) {
f17
            if (M) delete M;
                                                                              e4f
                                                                                                   swap(1, r);
f38
            M = NULL:
                                                                              dc8
                                                                                                   if (1) 1->rev ^= 1:
                                                                              f2f
                                                                                                   if (r) r->rev ^= 1;
a28
            join(L, R, root);
b92
        }
                                                                              3e5
                                                                                              }
e77
        void meld(treap& t) { // segmented merge
                                                                              a32
                                                                                               lazy = 0, rev = 0;
4a6
            node *L = root, *R = t.root;
                                                                                          }
                                                                              ca6
950
            root = NULL:
                                                                                          void update() {
                                                                              0.1e
6b1
            while (L or R) {
                                                                              0c3
                                                                                               sz = 1, sub = val;
fe2
                 if (!L or (L and R and L->mi > R->mi)) std::swap(L, R);
                                                                              a09
                                                                                               if (1) 1->prop(), sz += 1->sz, sub += 1->sub;
5e1
                 if (!R) ioin(root, L, root), L = NULL;
                                                                              095
                                                                                               if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub;
3c9
                 else if (L->mi == R->mi) {
                                                                              360
                                                                                          }
a76
                     node* LL;
                                                                              d37
                                                                                      };
439
                     split(L, LL, L, R->mi+1);
359
                     delete LL;
                                                                              bb7
                                                                                      node* root;
2a3
                 } else {
a76
                     node* LL;
                                                                              84b
                                                                                      treap() { root = NULL; }
537
                     split(L, LL, L, R->mi);
                                                                              2d8
                                                                                      treap(const treap& t) {
dbb
                     join(root, LL, root);
                                                                              465
                                                                                          throw logic_error("Nao copiar a treap!");
                 }
                                                                              1e9
                                                                                      }
f4f
            }
576
                                                                              cec
                                                                                      \simtreap() {
689
                                                                              609
                                                                                          vector < node *> q = {root};
            t.root = NULL;
8e7
        }
                                                                              402
                                                                                          while (q.size()) {
651 }:
                                                                              e5d
                                                                                               node* x = q.back(); q.pop_back();
                                                                              ee9
                                                                                               if (!x) continue;
                                                                              1c7
                                                                                               q.push_back(x->1), q.push_back(x->r);
      Treap Implicita
6.36
                                                                              bf0
                                                                                               delete x;
                                                                              653
                                                                                          }
// Todas as operacoes custam
                                                                                      }
                                                                              50e
// O(log(n)) com alta probabilidade
                                                                              73c
                                                                                      int size(node* x) { return x ? x->sz : 0; }
878 mt19937 rng((int)
                                                                              b2b
                                                                                      int size() { return size(root): }
   chrono::steady_clock::now().time_since_epoch().count());
                                                                                      void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
                                                                              bcf
                                                                              986
                                                                                          if (!1 or !r) return void(i = 1 ? 1 : r);
aa1 template < typename T > struct treap {
                                                                              161
                                                                                          1->prop(), r->prop();
3c9
        struct node {
                                                                              80e
                                                                                          if (1->p > r->p) join(1->r, r, 1->r), i = 1;
b19
            node *1, *r;
                                                                              fa0
                                                                                          else join(1, r->1, r->1), i = r;
284
            int p, sz;
                                                                              bda
                                                                                          i->update();
875
            T val, sub, lazy;
                                                                              b57
                                                                                      }
aa6
            bool rev;
                                                                                      void split(node* i, node*& 1, node*& r, int v, int key = 0) {
                                                                              a20
8dc
            node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
                                                                              26a
                                                                                          if (!i) return void(r = 1 = NULL);
   sub(v), lazy(0), rev(0) {}
                                                                              c89
                                                                                          i->prop();
            void prop() {
a9c
                                                                              5bd
                                                                                          if (key + size(i->1) < v) split(i->r, i->r, r, v,
0ec
                 if (lazy) {
                                                                                 key+size(i->1)+1), l = i;
924
                     val += lazv, sub += lazv*sz;
                                                                              219
                                                                                          else split(i \rightarrow 1, l, i \rightarrow 1, v, key), r = i;
```

b87

if (1) 1->lazy += lazy;

```
bda
            i->update();
d37
231
        void push_back(T v) {
2e0
            node* i = new node(v);
7ab
            join(root, i, root);
46b
b7a
        T query(int 1, int r) {
df9
            node *L, *M, *R;
            split(root, M, R, r+1), split(M, L, M, 1);
dca
d43
            T ans = M->sub:
69d
            join(L, M, M), join(M, R, root);
ba7
            return ans;
1f7
41f
        void update(int 1, int r, T s) {
df9
            node *L, *M, *R;
            split(root, M, R, r+1), split(M, L, M, 1);
dca
8f6
            M \rightarrow lazy += s;
69d
            join(L, M, M), join(M, R, root);
29f
        void reverse(int 1, int r) {
8c1
df9
            node *L, *M, *R;
            split(root, M, R, r+1), split(M, L, M, 1);
dca
66a
            M \rightarrow rev ^= 1:
69d
            join(L, M, M), join(M, R, root);
ea8
        }
139 }:
```

6.37 Treap Persistent Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
6cf mt19937_64 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
3c9 struct node {
b19
        node *1, *r;
f14
        ll sz, val, sub;
304
        node(11 \ v) : 1(NULL), r(NULL), sz(1), val(v), sub(v) {}
        node(node* x) : l(x->l), r(x->r), sz(x->sz), val(x->val),
   sub(x->sub) {}
01e
        void update() {
0c3
            sz = 1, sub = val;
77e
           if (1) sz += 1->sz, sub += 1->sub;
d6e
           if (r) sz += r->sz, sub += r->sub:
124
            sub %= MOD;
472
       }
```

```
95f };
bc9 ll size(node* x) { return x ? x \rightarrow sz : 0; }
 761 void update(node* x) { if (x) x->update(); }
828 node* copy(node* x) { return x ? new node(x) : NULL; }
b02 node* join(node* 1, node* r) {
         if (!1 or !r) return 1 ? copy(1) : copy(r);
e1f
48b
         node* ret;
49f
         if (rng() % (size(l) + size(r)) < size(l)) {</pre>
7eb
             ret = copv(1);
             ret -> r = join(ret -> r, r);
cc1
784
         } else {
             ret = copy(r);
4c5
             ret ->1 = join(1, ret ->1);
551
7a0
74f
         return update(ret), ret;
2cc }
723 void split(node* x, node*& 1, node*& r, 11 \text{ v}, 11 \text{ key} = 0) {
         if (!x) return void(l = r = NULL);
421
b4b
         if (kev + size(x->1) < v) {
72f
             1 = copy(x);
d70
             split(1->r, 1->r, r, v, key+size(1->1)+1);
710
        } else {
303
             r = copy(x);
417
             split(r->1, 1, r->1, v, key);
3d1
da2
         update(1), update(r);
666 }
f9e vector < node *> treap;
139 void init(const vector<ll>& v) {
         treap = {NULL};
 bbd
 969
         for (auto i : v) treap[0] = join(treap[0], new node(i));
 286 }
6.38 Wavelet Tree
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
```

```
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - O(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
597 int n. v[MAX]:
578 vector < int > esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
f8d void build(int b = 0, int e = n, int p = 1, int l = MINN, int r =
   MAXN) {
        int m = (1+r)/2; esq[p].push_back(0); pref[p].push_back(0);
58f
        for (int i = b; i < e; i++) {</pre>
f2f
            esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
6b9
            pref[p].push_back(pref[p].back()+v[i]);
26f
206
8ce
        if (1 == r) return;
        int m2 = stable_partition(v+b, v+e, [=](int i){return i <=</pre>
3a7
347
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
Ofb }
540 int count(int i, int j, int x, int y, int p = 1, int l = MINN, int
   r = MAXN)  {
       if (y < 1 \text{ or } r < x) \text{ return } 0;
2ad
4db
       if (x <= l and r <= y) return j-i;</pre>
       int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        return count(ei, ej, x, y, 2*p, l, m)+count(i-ei, j-ej, x, y,
   2*p+1, m+1, r);
3cf }
f62 int kth(int i, int j, int k, int p=1, int l = MINN, int r = MAXN) {
Зсе
        if (1 == r) return 1:
ddc
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
585
        return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
28b
8b6 }
f2c int sum(int i, int j, int x, int y, int p = 1, int l = MINN, int r
   = MAXN)
        if (y < 1 \text{ or } r < x) \text{ return } 0;
2ad
```

```
2a9
        if (x <= 1 and r <= y) return pref[p][j]-pref[p][i];</pre>
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
43b
        return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x, y,
   2*p+1, m+1, r);
b6d }
b84 int sumk(int i, int j, int k, int p = 1, int l = MINN, int r =
   MAXN) {
8a1
       if (1 == r) return 1*k;
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
50c
       if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
        return pref[2*p][ei]-pref[2*p][ei]+sumk(i-ei, j-ej, k-(ej-ei),
   2*p+1. m+1. r):
940 }
```

7 Strings

7.1 Aho-corasick

```
// query retorna o somatorio do numero de matches de
// todas as stringuinhas na stringona
//
// insert - O(|s| log(SIGMA))
// build - O(N), onde N = somatorio dos tamanhos das strings
// query - O(|s|)
eal namespace aho {
807
        map < char , int > to[MAX];
        int link[MAX], idx, term[MAX], exit[MAX], sobe[MAX];
c87
bfc
        void insert(string& s) {
05e
            int at = 0;
            for (char c : s) {
b4f
                auto it = to[at].find(c);
b68
1c9
                if (it == to[at].end()) at = to[at][c] = ++idx;
361
                else at = it->second;
ff4
142
            term[at]++, sobe[at]++;
6eb
d41 #warning nao esquece de chamar build() depois de inserir
0a8
        void build() {
26a
            queue < int > q;
537
            q.push(0);
dff
            link[0] = exit[0] = -1;
402
            while (q.size()) {
379
                int i = q.front(); q.pop();
```

```
3 c 4
                for (auto [c, j] : to[i]) {
                                                                            34f
                                                                                         while (s[n-len[last]-2] != c) last = link[last];
                    int 1 = link[i];
                                                                            289
5da
                                                                                        if (!t[last][c]) {
102
                    while (1 != -1 and !to[1].count(c)) 1 = link[1];
                                                                            dab
                                                                                             int prev = link[last];
                    link[j] = 1 == -1 ? 0 : to[1][c];
                                                                            553
                                                                                             while (s[n-len[prev]-2] != c) prev = link[prev];
7a5
                    exit[j] = term[link[j]] ? link[j] : exit[link[j]];
                                                                            fb2
                                                                                            link[sz] = t[prev][c];
3ab
                    if (exit[j]+1) sobe[j] += sobe[exit[j]];
                                                                                            len[sz] = len[last]+2;
                                                                            3f5
6f2
113
                    q.push(j);
                                                                            1f8
                                                                                            t[last][c] = sz++:
                }
                                                                            f8b
f1d
367
            }
                                                                            344
                                                                                        gt[last = t[last][c]]++;
        }
                                                                                    }
768
                                                                            b1d
bc0
        int query(string& s) {
                                                                            f17
                                                                                    int size() { return sz-2; }
86d
            int at = 0. ans = 0:
                                                                            2af
                                                                                    11 propagate() {
b4f
            for (char c : s){
                                                                            b73
                                                                                        11 \text{ ret} = 0:
                while (at != -1 and !to[at].count(c)) at = link[at];
1ca
                                                                            ebb
                                                                                        for (int i = n; i > 1; i--) {
                at = at == -1 ? 0 : to[at][c];
5b9
                                                                            fd3
                                                                                             qt[link[i]] += qt[i];
                ans += sobe[at]:
                                                                            db5
2b1
                                                                                             ret += qt[i];
b85
            }
                                                                            074
                                                                                        }
ba7
                                                                            edf
                                                                                        return ret;
            return ans;
                                                                                    }
038
                                                                            ef6
a30 }
                                                                            a2e };
                                                                            7.3 KMP
7.2 eertree
// Constroi a eertree, caractere a caractere
                                                                            // matching(s, t) retorna os indices das ocorrencias
                                                                            // de s em t
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
                                                                            // autKMP constroi o automato do KMP
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
                                                                            // Complexidades:
// substrings pal. com repeticao
                                                                            // pi - O(n)
                                                                            // match - 0(n + m)
// O(n) amortizado, considerando alfabeto O(1)
                                                                            // construir o automato - O(|sigma|*n)
                                                                            // n = |padrao| e m = |texto|
8eb struct eertree {
7cc
       vector < vector < int >> t;
                                                                            ea8 template < typename T > vector < int > pi(T s) {
                                                                                    vector < int > p(s.size());
42e
       int n, last, sz;
                                                                            019
                                                                                    for (int i = 1, j = 0; i < s.size(); i++) {</pre>
        vector < int > s, len, link, qt;
                                                                            725
745
                                                                                         while (j \text{ and } s[j] != s[i]) j = p[j-1];
                                                                            a51
d36
        eertree(int N) {
                                                                            973
                                                                                        if (s[j] == s[i]) j++;
            t = vector(N+2, vector(26, int()));
                                                                            f8c
ec8
                                                                                        p[i] = j;
            s = len = link = qt = vector < int > (N+2);
                                                                                    }
                                                                            e0a
cee
            s[0] = -1;
                                                                            74e
                                                                                    return p;
cd1
288
            link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
                                                                            f50 }
            sz = 2, last = 0, n = 1;
688
        }
                                                                            c10 template < typename T> vector <int> matching(T& s, T& t) {
8dc
                                                                            658
                                                                                    vector<int> p = pi(s), match;
244
        void add(char c) {
                                                                            a1b
                                                                                    for (int i = 0, j = 0; i < t.size(); i++) {</pre>
```

6be

while (j and s[j] != t[i]) j = p[j-1];

692

s[n++] = c -= 'a';

```
c4d
            if (s[i] == t[i]) j++;
            if (i == s.size()) match.push_back(i-j+1), j = p[j-1];
310
028
ed8
        return match;
c82 }
a2d struct KMPaut : vector < vector < int >> {
47 c
        KMPaut(){}
6c7
        KMPaut (string& s) : vector < vector < int >> (26,
   vector < int > (s.size()+1)) {
503
            vector < int > p = pi(s);
04b
            auto& aut = *this;
4fa
            aut[s[0]-'a'][0] = 1:
            for (char c = 0; c < 26; c++)
19a
5d3
                for (int i = 1; i <= s.size(); i++)</pre>
                     aut[c][i] = s[i] - a' == c ? i+1 : aut[c][p[i-1]];
42b
4bb
        }
79b };
7.4 Manacher
// manacher recebe um vetor de T e retorna o vetor com tamanho dos
   palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// ret[2*i+1] = tamanho maior palindromo centrado em i e i+1
//
// Complexidades:
// manacher - O(n)
// palindrome - <0(n), 0(1)>
// pal_end - 0(n)
28a template < typename T > vector < int > manacher (const T& s) {
18f
        int 1 = 0, r = -1, n = s.size();
fc9
        vector < int > d1(n), d2(n);
603
        for (int i = 0; i < n; i++) {</pre>
821
            int k = i > r ? 1 : min(d1[l+r-i], r-i);
            while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k]) k++;
61a
61e
            d1[i] = k--;
9f6
            if (i+k > r) l = i-k, r = i+k;
950
        1 = 0, r = -1;
e03
603
        for (int i = 0: i < n: i++) {
            int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
a 64
2c6
            while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k]) k++;
            d2[i] = --k:
eaa
26d
            if (i+k-1 > r) l = i-k, r = i+k-1;
```

4fe

}

```
c41
        vector<int> ret(2*n-1);
e6b
        for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;
        for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
e1d
edf
        return ret:
ebb }
// verifica se a string s[i..j] eh palindromo
cac template < typename T > struct palindrome {
f97
        vector < int > man;
b2d
        palindrome(const T& s) : man(manacher(s)) {}
9d7
        bool query(int i, int j) {
bad
            return man[i+i] >= i-i+1:
1e7
        }
60c };
// tamanho do maior palindromo que termina em cada posicao
7cb template < typename T > vector < int > pal_end(const T& s) {
        vector<int> ret(s.size());
e57
fde
        palindrome <T> p(s);
d51
        ret[0] = 1;
88e
        for (int i = 1; i < s.size(); i++) {</pre>
a32
            ret[i] = min(ret[i-1]+2, i+1);
6ea
             while (!p.query(i-ret[i]+1, i)) ret[i]--;
78e
        }
edf
        return ret;
8bd }
7.5 Min/max suffix/cyclic shift
// Computa o indice do menor/maior sufixo/cyclic shift
// da string, lexicograficamente
//
// O(n)
016 template < typename T > int max_suffix(T s, bool mi = false) {
476
        s.push_back(*min_element(s.begin(), s.end())-1);
        int ans = 0;
1a4
88e
        for (int i = 1; i < s.size(); i++) {</pre>
            int i = 0:
eec
708
            while (ans+j < i and s[i+j] == s[ans+j]) j++;
7a2
            if (s[i+j] > s[ans+j]) {
b52
                 if (!mi or i != s.size()-2) ans = i;
e51
            } else if (j) i += j-1;
69 c
ba7
        return ans;
f2a }
```

```
a1a template < typename T > int min_suffix(T s) {
        for (auto& i : s) i *= -1;
76b
        s.push_back(*max_element(s.begin(), s.end())+1);
09d
925
        return max_suffix(s, true);
ec0 }
97c template < typename T > int max_cyclic_shift(T s) {
163
        int n = s.size();
        for (int i = 0; i < n; i++) s.push_back(s[i]);</pre>
1ad
20a
        return max_suffix(s);
d34 }
08a template < typename T> int min_cyclic_shift(T s) {
        for (auto& i : s) i *= -1;
76b
7be
        return max_cyclic_shift(s);
c7a }
7.6 String Hashing
// Complexidades:
// construtor - O(|s|)
// operator() - 0(1)
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
463 int uniform(int 1, int r) {
        uniform_int_distribution < int > uid(1, r);
a7f
        return uid(rng);
f54
d9e }
9e0 template <int MOD> struct str_hash { // 116fcb
c63
        static int P;
dcf
        vector<ll> h, p;
ea8
        str_hash(string s) : h(s.size()), p(s.size()) {
            p[0] = 1, h[0] = s[0];
7a2
ad7
            for (int i = 1; i < s.size(); i++)</pre>
                p[i] = p[i - 1] * P % MOD, h[i] = (h[i - 1] * P + s[i]) % MOD;
84 c
1ef
        ll operator()(int 1, int r) { // retorna hash s[l...r]
af7
749
            11 hash = h[r] - (1 ? h[1 - 1]*p[r - 1 + 1]%MOD : 0);
dfd
            return hash < 0 ? hash + MOD : hash;</pre>
3ba
116 }:
217 template <int MOD > int str_hash < MOD >:: P = uniform (256, MOD - 1); //
```

1 > |sigma|

7.7 String Hashing - modulo 2⁶¹ - 1

```
// Quase duas vezes mais lento
// Complexidades:
// build - O(|s|)
// operator() - 0(1)
9d0 const ll MOD = (111<<61) - 1;
e38 ll mulmod(ll a. ll b) {
ff3
        const static 11 LOWER = (111<<30) - 1, GET31 = (111<<31) - 1;
        11 11 = a\&LOWER, h1 = a>>30, 12 = b\&LOWER, h2 = b>>30;
410
d54
        11 m = 11*h2 + 12*h1, h = h1*h2;
        ll ans = 11*12 + (h>>1) + ((h&1)<<60) + (m>>31) +
784
    ((m\&GET31) << 30) + 1;
1dd
        ans = (ans\&MOD) + (ans>>61), ans = (ans\&MOD) + (ans>>61);
        return ans - 1:
c0f
f98 }
798 mt19937 64
    rng(chrono::steady_clock::now().time_since_epoch().count());
f89 ll uniform(ll l, ll r) {
        uniform_int_distribution < ll > uid(1, r);
969
        return uid(rng):
f54
cac }
d7d struct str_hash {
        static 11 P;
c20
dcf
        vector<ll> h, p;
ea8
        str_hash(string s) : h(s.size()), p(s.size()) {
7a2
            p[0] = 1, h[0] = s[0];
ad7
            for (int i = 1; i < s.size(); i++)</pre>
632
                p[i] = mulmod(p[i-1], P), h[i] = (mulmod(h[i-1], P), h[i])
   P) + s[i])%MOD;
507
        }
        11 operator()(int 1, int r) { // retorna hash s[1...r]
af7
538
            ll hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0);
            return hash < 0 ? hash + MOD : hash;</pre>
dfd
544
        }
148 }:
6c5 ll str_hash::P = uniform(256, MOD - 1); // l > |sigma|
7.8 Suffix Array - O(n log n)
// kasai recebe o suffix array e calcula lcp[i],
```

// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],..,n-1]

```
//
// Complexidades:
// suffix_array - O(n log(n))
// kasai - O(n)
733 vector <int > suffix_array(string s) {
b38
        s += "$";
        int n = s.size(), N = max(n, 260);
043
2f3
       vector < int > sa(n), ra(n);
       for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];</pre>
29b
0a2
        for (int k = 0; k < n; k ? k *= 2 : k++) {
5ce
            vector < int > nsa(sa), nra(n), cnt(N);
            for(int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,
   cnt[ra[i]]++;
            for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
4 c 4
            for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]] = nsa[i];
368
            for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=</pre>
   ra[sa[i]] !=
                 ra[sa[i-1]] or ra[(sa[i]+k)\%n] != ra[(sa[i-1]+k)\%n];
f86
26b
d5e
            if (ra[sa[n-1]] == n-1) break;
11e
057
        return vector < int > (sa.begin()+1, sa.end());
ff3 }
481 vector<int> kasai(string s, vector<int> sa) {
232
        int n = s.size(), k = 0;
        vector < int > ra(n), lcp(n);
408
676
        for (int i = 0; i < n; i++) ra[sa[i]] = i;</pre>
740
        for (int i = 0: i < n: i++, k -= !!k) {
            if (ra[i] == n-1) { k = 0; continue; }
199
1de
            int j = sa[ra[i]+1];
891
            while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
d98
            lcp[ra[i]] = k;
a07
5ed
        return lcp;
fbe }
    Suffix Array - O(n)
// Rapidao
// Computa o suffix array em 'sa', o rank em 'rnk'
// e o lcp em 'lcp'
```

```
// query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
// Complexidades
// O(n) para construir
// query - 0(1)
1a5 template < typename T > struct rmq {
517
        vector <T> v;
fcc
        int n; static const int b = 30;
        vector < int > mask, t;
70e
183
        int op(int x, int y) { return v[x] \leftarrow v[y] ? x : y; }
ee1
        int msb(int x) { return builtin clz(1) - builtin clz(x): }
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1));}
        rmq() {}
6ad
43c
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) {
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {</pre>
2e5
a61
                 at = (at << 1) & ((1 << b) -1);
c00
                 while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
c2f
            }
            for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
ea4
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
39d
   i+(1<< i) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
   t[n/b*(j-1)+i+(1<<(j-1))]);
41a
e34
        int index_query(int 1, int r) {
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
27b
e80
            int x = 1/b+1, y = r/b-1;
            if (x > y) return op(small(l+b-1), small(r));
fd3
a4e
            int j = msb(y-x+1);
            int ans = op(small(1+b-1), op(t[n/b*j+x],
   t[n/b*j+y-(1<<j)+1]));
            return op(ans, small(r));
be6
62a
093
        T query(int 1, int r) { return v[index_query(1, r)]; }
bab };
9d7 struct suffix_array {
ac0
        string s;
1a8
        int n;
        vector<int> sa, cnt, rnk, lcp;
5b4
2de
        rmq<int> RMQ;
        bool cmp(int a1, int b1, int a2, int b2, int a3=0, int b3=0) {
d6e
91d
            return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3 < b3):
```

```
82d
        template < typename T > void radix(int* fr, int* to, T* r, int N,
4a4
   int k) {
            cnt = vector < int > (k+1, 0);
c17
            for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;
bac
            for (int i = 1; i <= k; i++) cnt[i] += cnt[i-1];</pre>
703
            for (int i = N-1: i+1: i--) to [--cnt[r[fr[i]]]] = fr[i]:
000
6f3
        }
d66
        void rec(vector<int>& v, int k) {
a76
            auto &tmp = rnk, &m0 = lcp;
3a9
            int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
7f8
            vector < int > R(sz2+3);
74f
            for (int i = 1, i = 0; i < sz2; i += i/(3) R[i++] = i;
b30
            radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
207
            radix(&tmp[0], &R[0], &v[0]+1, sz2, k);
5f1
            radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
af5
            int dif = 0;
            int 10 = -1, 11 = -1, 12 = -1;
ed9
            for (int i = 0; i < sz2; i++) {</pre>
d81
                if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or
   v[tmp[i]+2] != 12)
                     10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 =
b43
   v[tmp[i]+2], dif++;
199
                if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
1f5
                else R[tmp[i]/3+sz] = dif;
            }
d18
47f
            if (dif < sz2) {</pre>
146
                rec(R. dif):
746
                for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;</pre>
            } else for (int i = 0; i < sz2; i++) sa[R[i]-1] = i;</pre>
105
            for (int i = 0, j = 0; j < sz2; i++) if (sa[i] < sz)
   tmp[j++] = 3*sa[i];
            radix(&tmp[0], &m0[0], &v[0], sz, k);
7ce
74d
            for (int i = 0; i < sz2; i++)</pre>
                sa[i] = sa[i] < sz ? 3*sa[i]+1 : 3*(sa[i]-sz)+2;
с9е
            int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
332
            while (p \ge 0 \text{ and } p2 \ge 0) {
1c9
                if ((sa[p2]%3==1 and cmp(v[m0[p]], v[sa[p2]],
3b3
   R[m0[p]/3],
                     R[sa[p2]/3+sz])) or (sa[p2]%3==2 and cmp(v[m0[p]],
0ce
   v[sa[p2]],
                     v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz],
af6
```

```
R[sa[p2]/3+1]))
300
                     sa[at--] = sa[p2--];
cb0
                 else sa[at--] = m0[p--];
214
            }
f2b
             while (p >= 0) sa[at--] = m0[p--];
             if (N\%3==1) for (int i = 0; i < N; i++) sa[i] = sa[i+1];
eb6
ee6
        }
938
        suffix_array(const string& s_) : s(s_), n(s.size()), sa(n+3),
                 cnt(n+1), rnk(n), lcp(n-1) {
e62
9fe
             vector < int > v(n+3);
f9b
             for (int i = 0; i < n; i++) v[i] = i;</pre>
eba
             radix(&v[0], &rnk[0], &s[0], n, 256):
e6d
             int dif = 1:
             for (int i = 0; i < n; i++)</pre>
830
419
                 v[rnk[i]] = dif += (i and s[rnk[i]] != s[rnk[i-1]]);
7cf
             if (n \ge 2) rec(v, dif);
             sa.resize(n):
fb9
76f
             for (int i = 0: i < n: i++) rnk[sa[i]] = i:</pre>
892
             for (int i = 0, k = 0; i < n; i++, k -= !!k) {
668
                 if (rnk[i] == n-1) {
5a4
                     k = 0:
5e2
                     continue;
9df
                 }
39a
                 int j = sa[rnk[i]+1];
891
                 while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
825
                 lcp[rnk[i]] = k;
a3e
9ff
             RMQ = rmq < int > (lcp);
        }
9a8
588
        int query(int i, int j) {
d97
             if (i == i) return n-i:
             i = rnk[i], j = rnk[j];
223
c3a
             return RMQ.query(min(i, j), max(i, j)-1);
940
71c
        pair<int, int> next(int L, int R, int i, char c) {
024
             int l = L, r = R+1;
40c
             while (1 < r) {
                 int m = (1+r)/2:
ee4
                 if (i+sa[m] >= n \text{ or } s[i+sa[m]] < c) 1 = m+1;
e7e
                 else r = m:
ef3
ebe
575
             if (1 == R+1 \text{ or } s[i+sa[1]] > c) \text{ return } \{-1, -1\};
eb7
             L = 1:
```

```
9e2
            1 = L, r = R+1;
            while (1 < r) {
40 c
                int m = (1+r)/2;
ee4
                if (i+sa[m] >= n or s[i+sa[m]] <= c) l = m+1;
1a1
ef3
                else r = m:
            }
b5b
56a
            R = 1-1:
e13
            return {L, R};
71b
        // quantas vezes 't' ocorre em 's' - O(|t| log n)
66d
        int count_substr(string& t) {
b2b
            int L = 0, R = n-1;
c9d
            for (int i = 0: i < t.size(): i++) {</pre>
de0
                tie(L, R) = next(L, R, i, t[i]);
4fc
                if (L == -1) return 0;
            }
cff
fbf
            return R-L+1;
aaa
        // exemplo de f que resolve o problema
           https://codeforces.com/edu/course/2/lesson/2/5/practice/contes/t/2699656/%pidesmi/D
57e
        ll f(ll k) { return k*(k+1)/2; }
        11 dfs(int L, int R, int p) { // dfs na suffix tree chamado em
   pre ordem
c54
            int ext = L != R ? RMQ.query(L, R-1) : n - sa[L];
            // Tem 'ext - p' substrings diferentes que ocorrem 'R-L+1'
               vezes
            // O LCP de todas elas eh 'ext'
            ll ans = (ext-p)*f(R-L+1);
f80
            // L eh terminal, e folha sse L == R
            if (sa[L]+ext == n) L++;
63 c
            // se for um SA de varias strings separadas como s#t$u&,
               usar no lugar do if de cima
            // (separadores < 'a', diferentes e inclusive no final)</pre>
            // while (L <= R && (sa[L]+ext == n || s[sa[L]+ext] <
               'a')) {
            // L++;
            // }
            while (L <= R) {</pre>
add
                int idx = L != R ? RMQ.index_query(L, R-1) : -1;
5a8
                if (idx == -1 or lcp[idx] != ext) idx = R;
5ef
```

```
478
                ans += dfs(L, idx, ext);
28d
                L = idx+1;
590
ba7
            return ans;
e21
        }
        // sum over substrings: computa, para toda substring t
            distinta de s,
        // \sum f(# ocorrencias de t em s) - 0 (n)
        11 sos() { return dfs(0, n-1, 0); }
ca8
6fa }:
```

7.10 Suffix Array Dinamico

```
// Mantem o suffix array, lcp e rank de uma string,
// premitindo push_front e pop_front
// O operador [i] return um par com sa[i] e lcp[i]
// lcp[i] tem o lcp entre sa[i] e sa[i-1] (lcp[0] = 0)
//
// Construir sobre uma string de tamanho n: O(n log n)
// push_front e pop_front: O(log n) amortizado
2fe struct dyn_sa {
3c9
        struct node {
1d4
            int sa. lcp:
ed1
            node *1, *r, *p;
f0d
            int sz, mi;
17b
            node(int sa_, int lcp_, node* p_) : sa(sa_), lcp(lcp_),
543
                l(NULL), r(NULL), p(p_{-}), sz(1), mi(lcp) {}
0.1e
            void update() {
58f
                sz = 1, mi = lcp;
bd7
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
a54
                if (r) sz += r->sz, mi = min(mi, r->mi);
27 c
            }
574
        };
bb7
        node* root;
295
        vector<ll> tag; // tag of a suffix (reversed id)
        string s; // reversed
ac0
cf4
        dyn_sa() : root(NULL) {}
e45
        dyn_sa(string s_) : dyn_sa() {
            reverse(s_.begin(), s_.end());
ae4
519
            for (char c : s_) push_front(c);
2a7
        }
```

```
a86
        \sim dvn_sa() {
            vector < node *> q = {root};
609
402
            while (q.size()) {
                node* x = q.back(); q.pop_back();
e5d
                if (!x) continue:
ee9
                q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                delete x:
653
            }
8c1
        }
73c
        int size(node* x) { return x ? x->sz : 0; }
08e
        int mirror(int i) { return s.size()-1 - i; }
580
        bool cmp(int i, int i) {
            if (s[i] != s[j]) return s[i] < s[j];</pre>
a29
            if (i == 0 or j == 0) return i < j;</pre>
5b4
988
            return tag[i-1] < tag[j-1];</pre>
9fd
        }
919
        void fix_path(node* x) { while (x) x->update(), x = x->p; }
245
        void flatten(vector < node * > & v, node * x) {
8c8
            if (!x) return:
            flatten(v, x->1);
e96
2a2
            v.push_back(x);
            flatten(v, x->r);
42d
01f
        void build(vector<node*>& v, node*& x, node* p, int L, int R,
   11 1, 11 r) {
04c
            if (L > R) return void(x = NULL);
331
            int M = (L+R)/2:
            11 m = (1+r)/2;
3e3
7e5
            x = v[M];
63e
            x->p = p;
bb3
            tag[x->sa] = m;
            build(v, x->1, x, L, M-1, 1, m-1), build(v, x->r, x, M+1,
   R. m+1. r):
            x->update();
ca8
a3a
        void fix(node*& x, node* p, ll l, ll r) {
82f
            if (3*max(size(x->1), size(x->r)) \le 2*size(x)) return
   x->update();
3d1
            vector < node *> v;
            flatten(v. x):
Осс
            build(v, x, p, 0, v.size()-1, 1, r);
ea9
b86
b19
        node* next(node* x) {
728
            if (x->r) {
a91
                x = x - > r:
347
                while (x->1) x = x->1:
```

```
ea5
                 return x;
e7d
402
            while (x->p \text{ and } x->p->r == x) x = x->p;
137
            return x->p;
48b
        }
b68
        node* prev(node* x) {
e41
            if (x->1) {
a26
                x = x -> 1;
93c
                 while (x->r) x = x->r;
                return x;
ea5
9be
6a1
            while (x->p \text{ and } x->p->1 == x) x = x->p;
137
            return x->p:
73e
        }
4f7
        int get_lcp(node* x, node* y) {
75a
            if (!x or !y) return 0; // change defaut value here
            if (s[x->sa] != s[y->sa]) return 0;
e51
843
            if (x->sa == 0 \text{ or } y->sa == 0) \text{ return } 1;
            return 1 + query(mirror(x->sa-1), mirror(y->sa-1));
4d0
        }
8d6
        void add_suf(node*& x, node* p, int id, ll l, ll r) {
ad6
91e
            if (!x) {
8e3
                x = new node(id, 0, p);
8e2
                 node *prv = prev(x), *nxt = next(x);
65d
                 int lcp_cur = get_lcp(prv, x), lcp_nxt = get_lcp(x,
   nxt);
ca3
                 if (nxt) nxt->lcp = lcp_nxt, fix_path(nxt);
71f
                x->lcp = lcp_cur;
7b4
                 tag[id] = (1+r)/2;
ca8
                x->update();
505
                 return;
            }
d0e
4a3
            if (cmp(id, x->sa)) add_suf(x->1, x, id, 1, tag[x->sa]-1);
            else add_suf(x->r, x, id, tag[x->sa]+1, r);
сЗа
3db
            fix(x, p, 1, r);
c98
        }
ec2
        void push_front(char c) {
cc7
            s += c:
493
            tag.push_back(-1);
            add_suf(root, NULL, s.size() - 1, 0, 1e18);
05e
        }
1f2
7f3
        void rem_suf(node*& x, int id) {
6cf
            if (x->sa != id) {
864
                 if (tag[id] < tag[x->sa]) return rem_suf(x->1, id);
e6f
                return rem suf(x->r, id);
```

```
2ae
2cf
            node* nxt = next(x):
09b
            if (nxt) nxt->lcp = min(nxt->lcp, x->lcp), fix_path(nxt);
b20
            node *p = x - p, *tmp = x;
            if (!x->1 \text{ or } !x->r) {
f3f
2fd
                x = x->1 ? x->1 : x->r:
                if (x) x - p = p;
753
696
            } else {
7 f 7
                for (tmp = x->1, p = x; tmp->r; tmp = tmp->r) p = tmp;
f2a
                x->sa = tmp->sa, x->lcp = tmp->lcp;
482
                if (tmp->1) tmp->1->p = p;
14c
                if (p->1 == tmp) p->1 = tmp->1:
a94
                else p->r = tmp->1;
            }
97c
b5e
            fix_path(p);
7c3
            delete tmp;
510
        }
        void pop_front() {
15b
            if (!s.size()) return;
abe
342
            s.pop_back();
436
            rem_suf(root, s.size());
сбе
            tag.pop_back();
987
        }
530
        int guerv(node* x, 11 1, 11 r, 11 a, 11 b) {
            if (!x \text{ or } tag[x->sa] == -1 \text{ or } r < a \text{ or } b < 1) \text{ return}
e51
   s.size():
            if (a <= 1 and r <= b) return x->mi;
ef5
8eb
            int ans = s.size();
            if (a \le tag[x->sa] \text{ and } tag[x->sa] \le b) ans = min(ans,
   x->1cp);
            ans = min(ans, query(x->1, 1, tag[x->sa]-1, a, b));
d99
            ans = min(ans, querv(x->r, tag[x->sa]+1, r, a, b)):
261
ba7
            return ans;
4c8
        }
588
        int query(int i, int j) { // lcp(s[i..], s[j..])
209
            if (i == j) return s.size() - i;
29e
            11 a = tag[mirror(i)], b = tag[mirror(j)];
710
            int ret = query(root, 0, 1e18, min(a, b)+1, max(a, b));
            return ret;
edf
84e
        // optional: get rank[i], sa[i] and lcp[i]
        int rank(int i) {
044
396
            i = mirror(i);
52f
            node* x = root;
7 c.9
            int ret = 0:
```

```
f4c
             while (x) {
33e
                  if (tag[x->sa] < tag[i]) {</pre>
f9d
                      ret += size(x->1)+1;
a91
                      x = x - > r:
6dc
                  } else x = x - > 1:
a19
edf
             return ret;
        }
153
649
         pair<int, int> operator[](int i) {
             node* x = root:
52f
31e
             while (1) {
d4d
                  if (i < size(x->1)) x = x->1;
4e6
                  else {
85f
                      i \rightarrow size(x\rightarrow 1);
e03
                      if (!i) return {mirror(x->sa), x->lcp};
040
                      i--, x = x->r;
b9b
                 }
7a2
             }
90c
         }
4c2 }:
```

7.11 Suffix Automaton

```
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares
16e namespace sam {
        int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
c1a
0b8
        int nxt[2*MAX][26];
        void add(int c) {
e6a
17a
            int at = cur:
9a6
            len[sz] = len[cur]+1, cur = sz++;
500
            while (at != -1 and !nxt[at][c]) nxt[at][c] = cur, at =
   link[at]:
            if (at == -1) { link[cur] = 0; return; }
7ea
654
            int q = nxt[at][c];
fd9
            if (len[q] == len[at]+1) { link[cur] = q; return; }
            int qq = sz++;
31f
2c3
            len[qq] = len[at]+1, link[qq] = link[q];
9a9
            for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];
e76
            while (at !=-1 and nxt[at][c] == q) nxt[at][c] = qq, at =
   link[at]:
8b8
            link[cur] = link[q] = qq;
61a
94e
        void build(string& s) {
889
            cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
```

```
9fe
            for (auto i : s) add(i-'a');
17a
            int at = cur:
121
            while (at) acc[at] = 1, at = link[at];
0e7
       }
        // coisas que da pra fazer:
28c
        11 distinct_substrings() {
            11 \text{ ans} = 0;
04b
            for (int i = 1; i < sz; i++) ans += len[i] - len[link[i]];
a1e
ba7
0d7
a6c
        string longest_common_substring(string& S, string& T) {
419
            build(S):
            int at = 0, 1 = 0, ans = 0, pos = -1;
111
            for (int i = 0; i < T.size(); i++) {</pre>
d59
                while (at and !nxt[at][T[i]-'a']) at = link[at], 1 =
f2c
   len[at];
                if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'], 1++;
efa
                else at = 0, 1 = 0;
749
                if (1 > ans) ans = 1, pos = i;
a1a
2b3
            }
20f
            return T.substr(pos-ans+1, ans);
       }
930
46e
        11 dp[2*MAX];
455
        11 paths(int i) {
2a8
            auto& x = dp[i];
dee
            if (x) return x;
483
            x = 1:
            for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=
   paths(nxt[i][i]);
            return x;
ea5
d88
        void kth_substring(int k, int at=0) { // k=1 : menor substring
   lexicog.
            for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {</pre>
9d2
d58
                if (paths(nxt[at][i]) >= k) {
d02
                    cout << char('a'+i);</pre>
c43
                    kth_substring(k-1, nxt[at][i]);
505
                    return;
69a
                k -= paths(nxt[at][i]);
5f4
ef6
            }
a13
        }
c37 };
7.12 Trie
```

```
// trie T() constroi uma trie para o alfabeto das letras minusculas
// trie T(tamanho do alfabeto, menor caracter) tambem pode ser usado
//
// T.insert(s) - O(|s|*sigma)
// T.erase(s) - O(|s|)
// T.find(s) retorna a posicao, 0 se nao achar - 0(|s|)
// T.count_pref(s) numero de strings que possuem s como prefixo -
   O(|s|)
//
// Nao funciona para string vazia
ab5 struct trie {
e1a
        vector < vector < int >> to:
        vector < int > end , pref;
450
af0
        int sigma; char norm;
        trie(int sigma_=26, char norm_='a') : sigma(sigma_),
bb1
    norm(norm_) {
            to = {vector < int > (sigma)};
58a
            end = \{0\}, pref = \{0\};
86e
        }
fe1
64e
        void insert(string s) {
c67
            int x = 0:
7e7
            for(auto c : s) {
800
                 int &nxt = to[x][c-norm];
dd7
                 if(!nxt) {
0aa
                     nxt = to.size();
526
                     to.push_back(vector<int>(sigma));
770
                     end.push_back(0), pref.push_back(0);
933
                }
827
                 x = nxt, pref[x]++;
            }
34 c
e4e
            end[x]++;
        }
e6b
6b2
        void erase(string s) {
            int x = 0:
c67
b4f
            for(char c : s) {
800
                 int &nxt = to[x][c-norm];
10c
                x = nxt, pref[x]--;
d8e
                if(!pref[x]) nxt = 0;
885
            }
bf0
            end[x]--;
ddd
        }
        int find(string s) {
aee
c67
            int x = 0:
7e7
            for(auto c : s) {
                 x = to[x][c-norm]:
2ec
                if(!x) return 0;
a66
```

```
e12
            }
ea5
            return x;
e77
839
        int count_pref(string s) {
e2f
             return pref[find(s)];
f40
979 };
7.13 Z
// z[i] = lcp(s, s[i..n))
//
// Complexidades:
// z - O(|s|)
// \text{ match - } O(|s| + |p|)
a19 vector <int> get_z(string s) {
163
        int n = s.size();
2b1
        vector < int > z(n, 0);
fae
        int 1 = 0, r = 0;
6f5
        for (int i = 1; i < n; i++) {</pre>
0af
            if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
457
            while (i + z[i] < n and s[z[i]] == s[i + z[i]]) z[i]++:
            if (i + z[i] - 1 > r) 1 = i, r = i + z[i] - 1;
65e
5cd
        }
070
        return z;
74a }
```

```
Extra
8.1 debug.cpp
void debug_out(string s, int line) { cerr << endl; }</pre>
template < typename H, typename ... T>
void debug_out(string s, int line, H h, T... t) {
    if (s[0] != ',') cerr << "Line(" << line << ") ";</pre>
    do { cerr << s[0]; s = s.substr(1);</pre>
    } while (s.size() and s[0] != ',');
    cerr << " = " << h;
    debug_out(s, line, t...);
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__, __VA_ARGS__)
#define debug(...) 42
#endif
8.2 hash.sh
# Para usar (hash das linhas [11, 12]):
# bash hash.sh arquivo.cpp 11 12
sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P -fpreprocessed | tr
    -d '[:space:]' | md5sum | cut -c-6
8.3 stress.sh
P=a
make ${P} ${P}2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./${P} < in > out
    ./${P}2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    fі
    echo $i
done
```

```
8.4 makefile
CXX = g++
CXXFLAGS = -fsanitize=address, undefined -fno-omit-frame-pointer -g
   -Wall -Wshadow -std=c++17 -Wno-unused-result -Wno-sign-compare
   -Wno-char-subscripts #-fuse-ld=gold
8.5 fastIO.cpp
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
        if (ch == ',-') break;
        if (ch >= '0' && ch <= '9') break;
        ch = getchar();
    }
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0'):
    if (minus) return -result;
    else return result;
}
    pragma.cpp
// Otimizacoes agressivas, pode deixar mais rapido ou mais devagar
#pragma GCC optimize("Ofast")
// Auto explicativo
#pragma GCC optimize("unroll-loops")
// Vetorizacao
#pragma GCC target("avx2")
```

8.7 timer.cpp

// Para operacoes com bits

#pragma GCC target("bmi,bmi2,popcnt,lzcnt")

```
// timer T; T() -> retorna o tempo em ms desde que declarou
using namespace chrono;
struct timer : high_resolution_clock {
```

```
const time_point start;
    timer(): start(now()) {}
    int operator()() {
        return duration_cast < milliseconds > (now() - start).count();
};
8.8 template.cpp
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
typedef long long 11;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f3f11;
int main() { _
    exit(0);
8.9 rand.cpp
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution < int > uid(1, r);
    return uid(rng);
8.10 vimrc
d79 set ts=4 sw=4 mouse=a nu ai si undofile
3b3 function Hash(1)
        return system("sed '/^\#w/d' \| cpp -dD -P -fpreprocessed \|
   tr -d '[:space:]' \| md5sum \| cut -c-6", a:1)
Obe endfunction
db5 function PrintHash() range
c1b
        let 1 = getline(a:firstline, a:lastline)
800
        let stk = []
```

```
for i in range(len(1))
66e
            let ini = i
f41
           for c in str2list(l[i])
c68
dc7
                if c == char2nr('{'})
                    call add(stk, i)
500
                endif
e96
                if c == char2nr(')')
a13
08f
                    let ini = stk[-1]
                    call remove(stk, -1)
860
e96
                endif
bf9
            endfor
            echo Hash(join(l[ini:i], "\n"))[0:2] 1[i]
ae2
bf9
        endfor
Obe endfunction
6b0 vmap <C-H> :call PrintHash() <CR>
```