Dependently Typed Heaps

https://github.com/brunjlar/heap

About Me

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Outline

- Motivation
- Leftist Heaps
- Proving Theorems in Haskell
- Dependently Typed Heaps
- Reflection on Results
- Questions & Comments

Motivation

Motivation

- Types help catching errors at compile time.
- Some invariants cannot be expressed by “simple” types.
- Haskell steadily moves towards dependent types.
- I wanted to see whether it is possible to prove theorems in Haskell...
- ...and use this to encode some non-trivial invariants.
- Heaps seem to be a good example.

Leftist Heaps

Heap

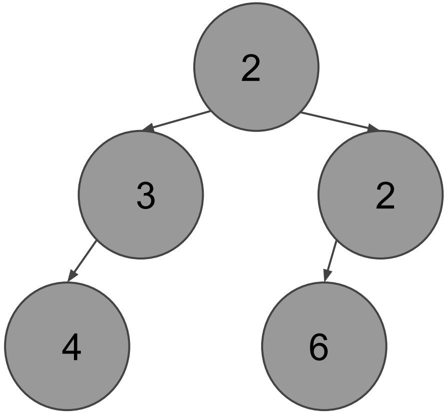
We consider binary trees whose nodes carry some payload and a **priority** (a natural number):

```
data Tree a = Empty | Node Natural a (Tree a) (Tree a)
```

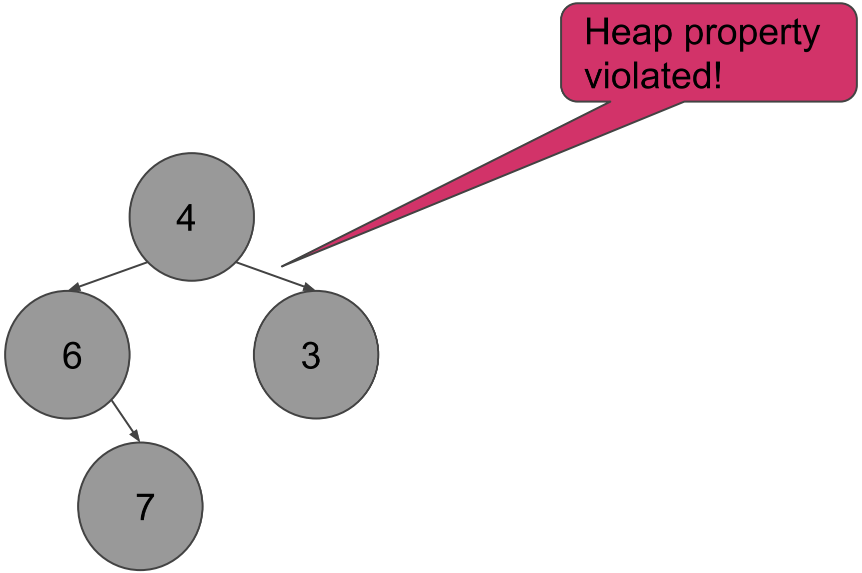
Such a tree is a **heap** if it satisfies the **heap property**: The priority of a node is not bigger than the priority of any of its children.

(So in a non-empty heap, the root has minimal priority.)

Heap



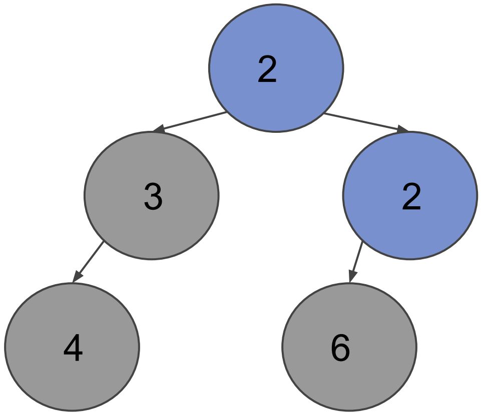
This is a heap.



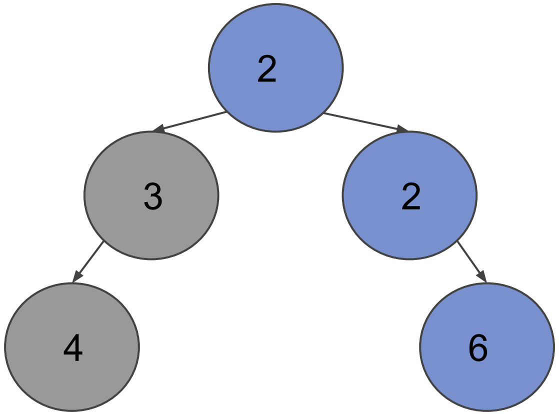
This is **not** a heap.

Rank

The **rank** of a heap is the length of its **right spine.**



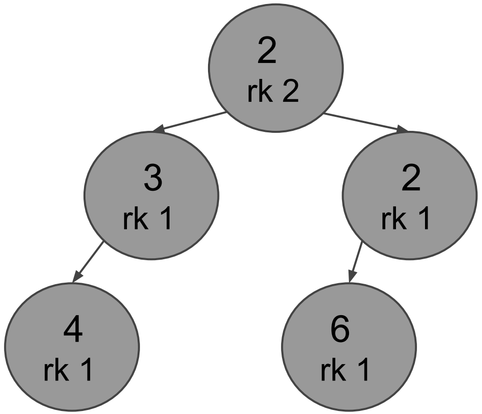
Rank 2



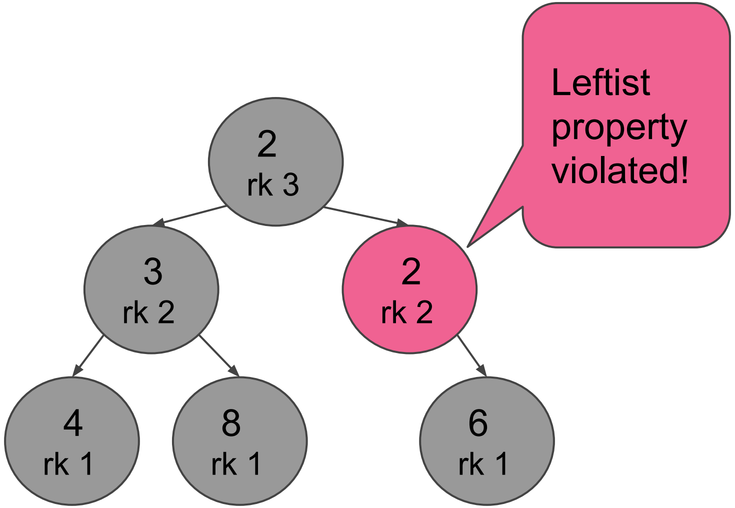
Rank 3

Leftist Heap

A heap is **leftist** if in each node, the rank of the left child is not smaller than the rank of the right child.



leftist



not leftist

Merging Leftist Heaps

```haskell
data Heap a = Empty | Node !Natural !Natural a (Heap a) (Heap a) deriving (Show, Functor)

rank :: Heap a -> Natural
rank Empty              = 0
rank (Node _ r _ _ _) = r

priority :: Heap a -> Maybe Natural
priority Empty              = Nothing
priority (Node p _ _ _ _) = Just p

singleton :: Natural -> a -> Heap a
singleton p x = Node p 1 x Empty Empty

merge :: Heap a -> Heap a -> Heap a
merge Empty                h'                   = h'
merge h                    Empty                = h
merge h@(Node p _ x ys zs) h'@(Node q _ _ _ _)
      | q < p                                   = merge h' h
      | otherwise                               =
          let h''@(Node _ r _ _ _) = merge zs h'
              r'                   = rank ys
          in if r <= r'
             then Node p (succ r ) x ys  h''
             else Node p (succ r') x h'' ys
```

Weakly typed heaps

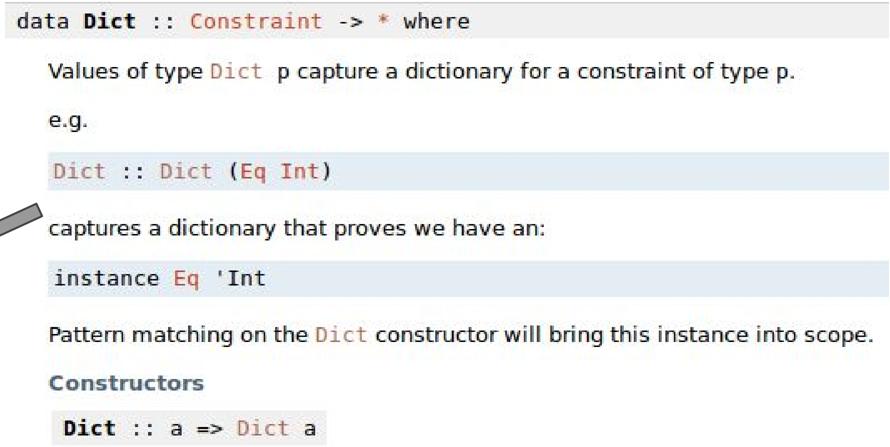
- Neither heap property nor leftist property are enforced by the compiler.
- “Classical solution”: smart constructors, but those only catch errors at runtime.
- Algorithms like “merge” can easily be done wrong.
- Can we define leftist heaps in such a way that the **compiler** prevents us from constructing “illegal” heaps?

Proving Theorems in Haskell

Technical Tools

- Constraints to express statements

- reified statements (dictionaries) (see “constraints” library by Kmett on Hackage)



```
data Dict :: Constraint -> * where
```

Values of type `Dict p` capture a dictionary for a constraint of type p.

e.g.

```
Dict :: Dict (Eq Int)
```

captures a dictionary that proves we have an:

```
instance Eq 'Int
```

Pattern matching on the `Dict` constructor will bring this instance into scope.

Constructors

```
Dict :: a => Dict a
```

- Singleton Types (see “singletons” library by Eisenberg & Stolarek on Hackage)

A Simple Example

```
data Peano = Z | S Peano deriving (Show, Read, Eq)

infix 4 ??

type family (m :: Peano) ?? (n :: Peano) :: Ordering where
    'Z   ?? 'Z   = 'EQ
    'Z   ?? _    = 'LT
    'S _ ?? 'Z   = 'GT
    'S m ?? 'S n = m ?? n

type (m :: Peano) < (n :: Peano) = (m ?? n) ~ 'LT

data SingPeano :: Peano -> * where

    SZ :: SingPeano 'Z
    SS :: SingPeano n -> SingPeano ('S n)

ltS :: SingPeano n -> Dict (n < 'S n)
ltS SZ = Dict
ltS (SS n) = ltS n
```

Dependently Typed Heaps

Dependently Typed Heaps

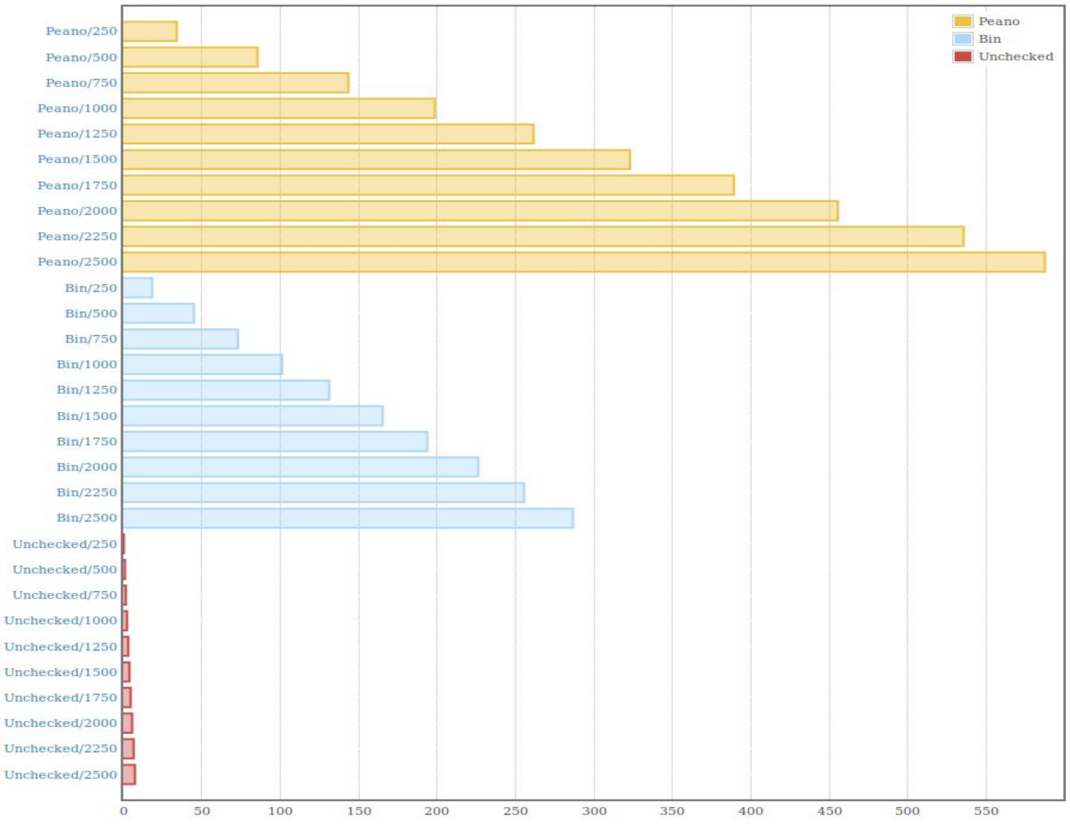
```
data Heap' nat (p :: Maybe nat) (r :: nat) a where

    Empty :: Heap' nat 'Nothing (Zero nat) a

    Node :: ( (p   <=. p')
            , (p   <=. p'')
            , (r'' <=  r')
            )
            => !(Sing nat p)
            -> !(Sing nat (Succ nat r''))
            -> !a
            -> !(Heap' nat p' r' a)
            -> !(Heap' nat p'' r'' a)
            -> Heap' nat ('Just p) (Succ nat r'') a
```

Type-Safe Merging

```haskell
merge :: Nat nat => Heap'' nat p a -> Heap'' nat q a -> Heap'' nat (Min' p q) a
merge (Heap'' Empty)                         h'                                = h'
merge h                                      (Heap'' Empty)                    = h
merge h@(Heap'' (Node p _ x ys zs)) h'@(Heap'' (Node q _ _ _ _)) =
    alternative (ltGeqDec q p)
        (using (minSymm p q) $ merge h' h) $
        let h'' = merge (Heap'' zs) h'
        in  case h'' of
            Heap'' Empty                   -> error "impossible branch"
            Heap'' h'''@(Node _ r _ _ _) ->
                using (minProd' p (priority zs) (Just' q)) $
                    alternative (leqGtDec r $ rank ys)
                        (Heap'' $ Node p (succ' r) x ys h''')
                        (Heap'' $ Node p (succ' $ rank ys) x h''' ys)
```

Reflection on Results

Benchmark



Peano
Bin
Unchecked
Peano/250
Peano/500
Peano/750
Peano/1000
Peano/1250
Peano/1500
Peano/1750
Peano/2000
Peano/2250
Peano/2500
Bin/250
Bin/500
Bin/750
Bin/1000
Bin/1250
Bin/1500
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Unchecked/250
Unchecked/500
Unchecked/750
Unchecked/1000
Unchecked/1250
Unchecked/1500
Unchecked/1750
Unchecked/2000
Unchecked/2250
Unchecked/2500
0
50
100
150
200
250
300
350
400
450
500
550

Type Safety

- Haskell's type system is powerful enough to encode non-trivial invariants on the type-level.

- Haskell functions can serve as "proofs" for statements about types.

- **Caution:** Compiler gives no termination guarantee, so would accept a non-terminating proof.

Questions & Comments