Smart Contracts

Smart Contracts & Bitcoin-Script

Lars Brünjes



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Reminder: The Simple UTxO-Model

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- The state of the blockchain is determined by the set of all UTxOs.

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- No output value is negative.

Example: A Simple Transaction

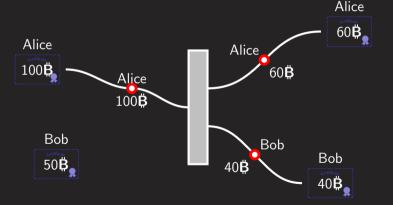
Alice holds 100 **B** and wants to send 40 **B** to Bob, who has 50 **B**.

Alice

Bob 50**B**

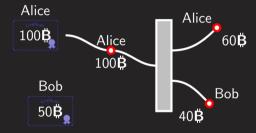
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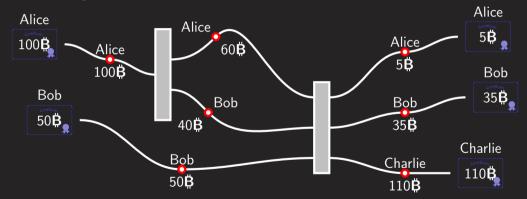
Example: A More Complex Transaction

After sending 40 \$\beta\$ to Bob, Alice and Bob want to send 55 \$\beta\$ each to Charlie.



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- In this sense, smart contracts are *only* "words": Intent does not matter. All that matters is the actual code.

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- While verifying the inputs of a Bitcoin transaction, input- and output-scripts are combined and executed. The result from this execution decides whether the spending transaction is entitled to spend the output.
- Details depend on the specific cryptocurrency, but the principle stays the same: Programs decide under which circumstances money may be spent.

Flavours of Smart Contracts

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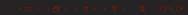
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- The creators of Bitcoin have chosen the first option, Ethereum and Cardano (for Plutus) use the second.



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- Many more exotic systems (like λ-calculus) are Turing-complete as well.



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Interlude: The Halting-Problem

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- To be more precise: Is there a program that will take an arbitrary program as input and then decide whether that program will halt?
- The answer is no! Assume there was a Python-function halt solving the Halting-Problem. Then consider the following Python-function:

```
def paradox():
    if halt(paradox):
        while True:
        pass
```

If halt returns True, paradox gets stuck in an infinite loop, and if halt returns False, paradox will halt. Both are contradictions.

Consequences for Smart Contracts

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- So if one chooses a Turing-complete smart-contract language, it is impossible to guarantee in advance whether a script will stop. Nor is it possible to know for how long the script will run, even if it eventually halts.
- As a consequence, one either has to decide against using a Turing-complete language or be prepared to interrupt a running script after some finite time has passed.

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- Bitcoin Script does not allow any loops, so a program written in Bitcoin Script can never get stuck in an infinite loop.
- In spite of its simplicity, Bitcoin Script is quite powerful and flexible and allows for a plethora of different kinds of transaction verification.
- On the other hand, Bitcoin Script is too limited to allow for real smart contracts implementing complex financial transactions.

Stacks

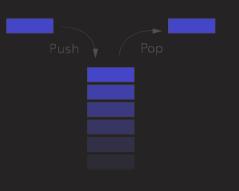
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- Forth is a (relatively) popular higher programming language that is stack based as well. Other examples are the Java Virtual Machine (JVM) and Microsoft's Common Language Runtime (CLR).
- There are no variables in Bitcoin Script. Instead, data is put onto the stack and processed there.



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- The transaction is valid if all inputs are valid in this sense (and if all other conditions are satisfied, so the sum of all input-values is greater than the sum of all output-values etc.).
- This script-mechanism extends the simple UTxO-model by allowing for more complex input-validation, going beyond digital signature verification.

- The following program written in Bitcoin Script calculates $(2+3) \cdot 4$:
- 2 3 op_add 4 op_mul

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3
op_add
4
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2 3 op_a<u>dd</u>

• 4 op_mul



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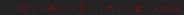
4

• op_mul



Exercises

- Compute $(10-3) \cdot (4+7)$ using Bitcoin Script! (*Hint*: Use op_sub!)
- Write a Bitcoin Script program which squares the number on top of the stack.
 (Hint: Use op_dup!)
- Write a Bitcoin Script programm which computes $x^2 + y^2$, where x and y are the two top-most numbers on the stack. (*Hint*: Use op_swap!)
- Write a Bitcoin Script Programm which computes $x \cdot y$ if y < x and x + y if y >= x, where x and y are the two top-most numbers on the stack (x on top, y below). (*Hint*: Use op_2dup, op_lessthan, op_if, op_else, and op_endif!)
- You can find a list of all Bitcoin-Script commands on https://en.bitcoin.it/wiki/Script.
- There is a nice online simulator on https://siminchen.github.io/bitcoinIDE/build/editor.html.



Pay to Public Key Hash in Bitcoin Script

- The vast majority of all Bitcoin transaction uses ordinary "hash of public key"-addresses.
- What do input- and output-scripts look like in this case?
 - <sig> <pubKey>
 - op_dup op_hash160 <pubKeyHash> op_equalverify op_checksig
 - The input-script puts the digital signature and the public key onto the stack. The
 output-script checks whether the hash of this public key has the right value and
 whether the signature is correct.

Input-Script:

 $304402200cc8b0471a38edad2ff9f9799521b7d948054817793c980eaf3a6637ddfb939702201c1a801461d4c3cf4de4e7336454dba0dd70b89d71f221e991cb6a79df1a860d01\\02ce9f5972fe1473c9b6948949f676bbf7893a03c5b4420826711ef518ceefd8dc$

Output-Script:

```
op_dup op_hash160 1290b657a78e201967c22d8022b348bd5e23ce17 op_equalverify op_checksig

304402200cc8b0471a38ed...
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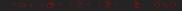
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 $304402200 co8b047 1a38 edad 2ff 9f 97995 21b7 d948054817793 c980 eaf 3a6637 ddf b939702201 c1a801461 d4c3 cf4 de4e733645 ddba0dd 70b89d71f221 e991 cb6a79df 1a860d01\\02 ce9f5972 fe1473 c9b6948949 f676 bbf7893 a03 c5b4420826711 ef518 ceefd 8dc$

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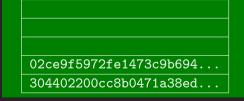
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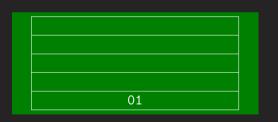
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- Bitcoin Script supports this via op_checkmultisig.

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- One possible application of this is to for example generate an output with value zero, then use a script which starts with op_return and contains arbitrary data afterwards.

Other kinds of Bitcoin Scripts — Riddles

 The output of transaction a4bfa8ab6435ae5f25dae9d89e4eb67dfa94283ca751f393c1ddc5a837bbc31b contained the following script:

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op_hash256
6fe28c0ab6f1b372c1a6a246ae63f74f931e8365e15a089c68d619000000000
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- To spend that output, one had to find a number with the given hash.
- This riddle was eventually solved: The given hash turned out to be the hash of the genesis-block-header.

• In the year 2013, Peter Todd created scripts whose outputs could be spend by anybody who found a hash collision for SHA-1:

op_2dup op_equal op_not op_verify op_sha1 op_swap op_sha1 op_equal

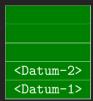
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- In February 2017, somebody claimed the reward of 2.48 ₿.

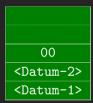


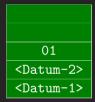


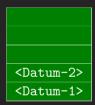


<Datum-1> <Datum-2> op_2dup op_equal op_not op_verify op_sha1 op_swap op_sha1 op_equal

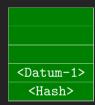
<Datum-2>
<Datum-1>
<Datum-2>
<Datum-1>

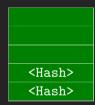


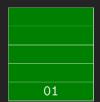












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- Both types of time locks are for example used in Bitcoin Lightning.

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- Instead they can use the following script:

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op_if <in three months> op_checklocktimeverify op_drop
<PubKey-Charlie> op_checksigverify 1 op_else 2 op_endif
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```

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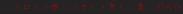
• Alice and Bob can access their money at any time using this script:

```
0 <Sig-Alice> <Sig-Bob> 0
```

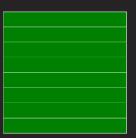
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- After three months, Charlie and either Alice or Bob can use the following script instead:
 - 0 <Sig-Alice/Bob> <Sig-Charlie> 1



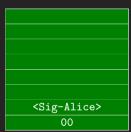
```
<Sig-Alice>
<Sig-Bob>
op_if
<in three months>
op_checklocktimeverify
op_drop
<PubKey-Charlie>
op_checksigverify
01
op_else
02
op_endif
<PubKey-Alice>
<PubKey-Bob>
02
op_checkmultisig
```



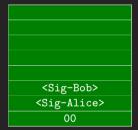
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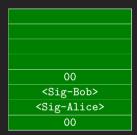
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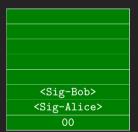
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<Sig-Alice>
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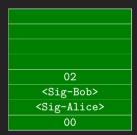
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op_endif
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op_checkmultisig
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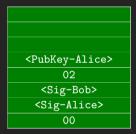
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<Sig-Alice>
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<PubKey-Bob>
<PubKey-Alice>
02
<Sig-Bob>
<Sig-Alice>
00

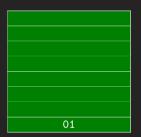
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op_if
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op_checklocktimeverify
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op_checkmultisig
```

02

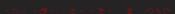
<PubKey-Bob>
<PubKey-Alice>
02

<Sig-Bob>
<Sig-Alice>
00

```
<Sig-Alice>
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op_if
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op_checklocktimeverify
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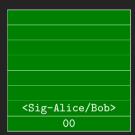
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<Sig-Alice/Bob>
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op_if
<in drei Monaten>
op_checklocktimeverify
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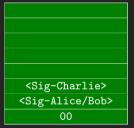
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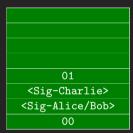
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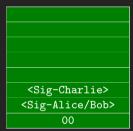
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<Sig-Alice/Bob>
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op_checklocktimeverify
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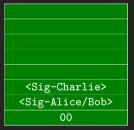
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op_if
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op_checklocktimeverify
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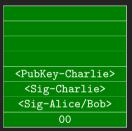
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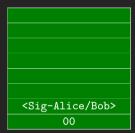
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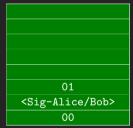
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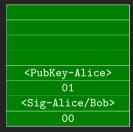
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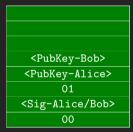
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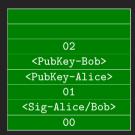
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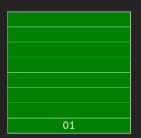
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INPUT OUTPUT