# Principle and Interface Techniques of Microcontroller

## --8051 Microcontroller and Embedded Systems Using Assembly and C

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杭州 • 浙江大学 • 2022

# Chapter 5

# Addressing Modes

# Addressing Modes

The CPU can access data in various ways, which are called addressing modes

Immediate
Register
Direct
Register indirect
Indexed
Bit
Relative

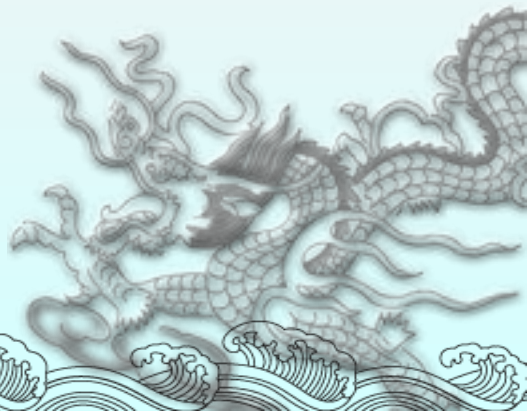
# Immediate Mode

The source operand is a constant

**The immediate data must be preceded by the pound sign, "#". Can load information into any registers, including 16-bit DPTR register.**

DPTR can also be accessed as two 8-bit registers, the high byte DPH and low byte DPL

```
MOV   A, #25H          ;load 25H into A
MOV   R4, #62          ;load 62 into R4
MOV   B, #40H          ;load 40H into B
MOV   DPTR, #4521H     ;DPTR=4512H
MOV   DPL, #21H        ;This is the same
MOV   DPH, #45H        ;as above
```



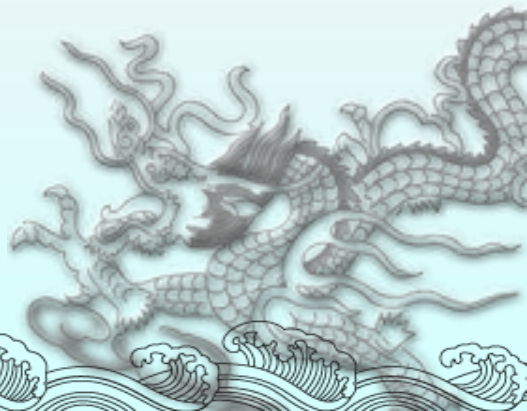
# Immediate Mode

We can use EQU directive to access immediate data

```
        Count    EQU  30
        ... ...
        MOV   R4,#COUNT ;R4=1EH
        MOV DPTR,#MYDATA ;DPTR=200H
        ORG  200H
MYDATA:  DB “America”
```

We can also use immediate addressing mode to send data to 8051 ports

```
        MOV   P1, #55H
```



# Register Addressing Mode

- Use registers to hold the data to be manipulated

```
MOV  A, R0        ;copy contents of R0 into A
MOV  R2, A        ;copy contents of A into R2
ADD  A, R5        ;add contents of R5 to A
ADD  A, R7        ;add contents of R7 to A
MOV  R6, A        ;save accumulator in R6
```

- The source and destination registers must match in size

```
MOV  DPTR, #25F5H
MOV  R7, DPL
MOV  R6, DPH
```

- The movement of data between Rn registers is not allowed

```
MOV R4,R7 is invalid
```

# Direct Addressing Mode

It is most often used the direct addressing mode to access RAM locations 30 – 7FH

The entire 128 bytes of RAM can be accessed.
The register bank locations are accessed by the register names.

```
MOV   A, 4        ;is same as
MOV   A, R4       ;which means copy R4 into A
```

Direct addressing mode

Register addressing mode

Contrast this with immediate addressing mode

There is no "#" sign in the operand

```
MOV  R0, 40H      ;save content of 40H in R0
MOV  56H, A       ;save content of A in 56H
```

# Stack and Direct Addressing Mode

Only direct addressing mode is allowed for pushing or popping the stack

PUSH A is invalid

**Pushing the accumulator onto the stack must be coded as**

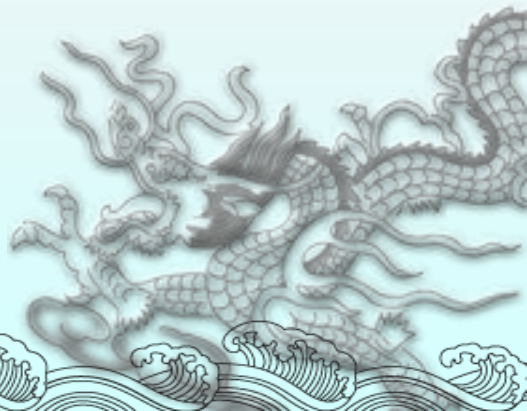
PUSH 0E0H or PUSH ACC

**Example 5-1**

Show the code to push R5 and A onto the stack and then pop them back them into R2 and B, where B = A and R2 = R5

**Solution:**

```
PUSH  05          ;push R5 onto stack
PUSH  0E0H        ;push register A onto stack
POP  0F0H         ;pop top of stack into B
                  ;now register B = register A
POP  02           ;pop top of stack into R2
                  ;now R2=R6
```



# Register Indirect Addressing Mode

A register is used as a pointer to the data

Only register R0 and R1 are used for this purpose

R2 – R7 cannot be used to hold the address of an operand located in RAM

When R0 and R1 hold the addresses of RAM locations, they must be preceded by the “@” sign

```
MOV  A, @R0        ;move contents of RAM whose
                   ;address is held by R0 into A
MOV  @R1, B        ;move contents of B into RAM
                   ;whose address is held by R1
```

# Register Indirect Addressing Mode

Uses registers R0 or R1 for 8-bit address:

```
MOV PSW, #0          ; use register bank 0
MOV R0, #0x3C
MOV @R0, #3
```

Uses DPTR register for 16-bit addresses:

```
MOV DPTR, #0x9000        ; DPTR ← 9000H
MOVX A, @DPTR            ; A ← [9000H]
```

Note that 9000 is an address in external memory

**Example 5-2**

Write a program to copy the value 55H into RAM memory locations 40H to 41H using (a) direct addressing mode, (b) register indirect addressing mode without a loop, and (c) with a loop

**Solution:**

```
(a)       MOV  A, #55H        ;load A with value 55H
          MOV 40H,A           ;copy A to RAM location 40H
          MOV 41H.A           ;copy A to RAM location 41H
(b)
          MOV A,#55H          ;load A with value 55H
          MOV R0,#40H         ;load the pointer. R0=40H
          MOV @R0,A           ;copy A to RAM R0 points to
          INC R0              ;increment pointer. Now R0=41h
          MOV @R0,A           ;copy A to RAM R0 points to
(c)
          MOV A,#55H          ;A=55H
          MOV R0,#40H         ;load pointer.R0=40H,
          MOV R2,#02          ;load counter, R2=3
AGAIN:    MOV @R0,A           ;copy 55 to RAM R0 points to
          INC R0              ;increment R0 pointer
          DJNZ  R2, AGAIN     ;loop until counter = zero
```

# Register Indirect Addressing Mode

- The advantage is that it makes accessing data dynamic rather than static as in direct addressing mode

  *Looping is not possible in direct addressing mode*

**Example 5-4**

Write a program to copy a block of 10 bytes of data from 35H to 60H

**Solution:**

```
        MOV  R0, #35H      ;source pointer
        MOV  R1, #60H      ;destination pointer
        MOV  R3, #10       ;counter
BACK:   MOV  A, @R0        ;get a byte from source
        MOV  @R1, A        ;copy it to destination
        INC  R0            ;increment source pointer
        INC  R1            ;increment destination pointer
        DJNZ  R3,BACK      ;keep doing for ten bytes
```

# Indexed Addressing Mode and On-chip ROM Access

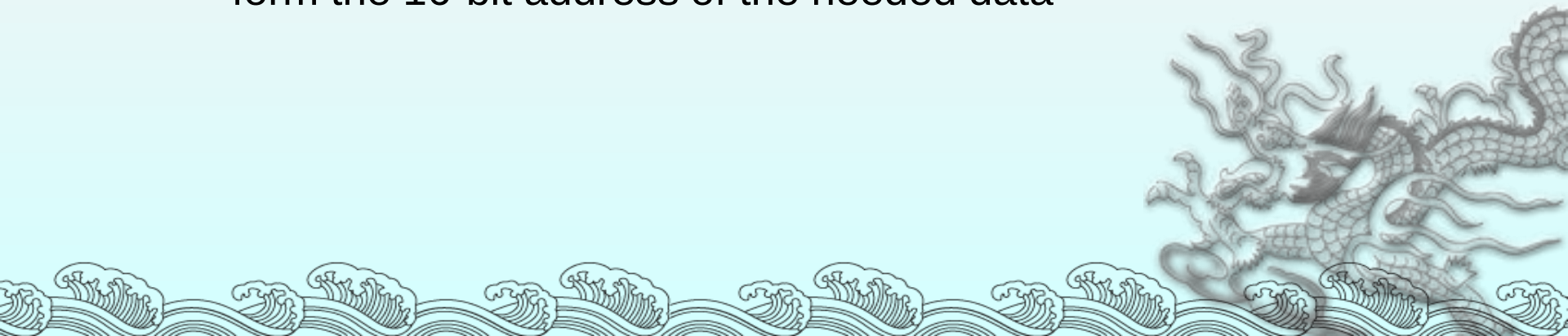
Indexed addressing mode is widely used in accessing data elements of look-up table entries located in the program ROM.

The instruction used for this purpose is

MOVC A,@A+DPTR

Use instruction MOVC, “C” means code

The contents of A are added to the 16-bit register DPTR to form the 16-bit address of the needed data



**Example 5-5**

In this program, assume that the word “USA” is burned into ROM locations starting at 200H. And that the program is burned into ROM locations starting at 0. Analyze how the program works and state where “USA” is stored after this program is run.

**Solution:**

```
          ORG  0000H                ;burn into ROM starting at 0
           MOV  DPTR,#200H           ;DPTR=200H look-up table addr
           CLR  A                    ;clear A(A=0)   DPTR=200H
           MOVC  A,@A+DPTR           ;get the char from code space
           MOV R0, A                 ;save it in R0
           INC DPTR                  ;DPTR=201 point to next char
           CLR   A                   ;clear A(A=0)
           MOVC  A, @A+DPTR          ;get the next char
           MOV  R1, A                ;save it in R1
           INC DPTR                  ;DPTR=202 point to next char
           CLR  A                    ;clear A(A=0)
           MOVC  A, @A+DPTR          ;get the next char
           MOV  R2, A                ;save it in R2
Here:      SJMP HERE                 ;stay here
;Data is burned into code space starting at 200H
            ORG 200H
MYDATA:  DB   “USA”
            END                      ;end of program
```

DPTR=200H, A=55H

# Look-up Table

The look-up table allows access to elements of a frequently used table with minimum operations.

**Example 5-6**

Write a program to get the x value from P1 and send $x^2$ to P2, continuously

**Solution:**

```
        ORG 0
        MOV  DPTR, #300H     ;LOAD TABLE ADDRESS
        MOV  A, #0FFH        ;A=FF
        MOV  P1, A           ;CONFIGURE P1 INPUT PORT
BACK:   MOV  A, P1           ;GET X
        MOV A,@A+DPTR        ;GET X SQAURE FROM TABLE
        MOV P2,A             ;ISSUE IT TO P2
        SJMP BACK            ;KEEP DOING IT
        ORG 300H
XSQR_TABLE:
        DB    0,1,4,9,16,25,36,49,64,81
        END
```

# Indexed Addressing Mode and MOVX

- In many applications, the size of program code does not leave any room to share the 64K-byte code space with data

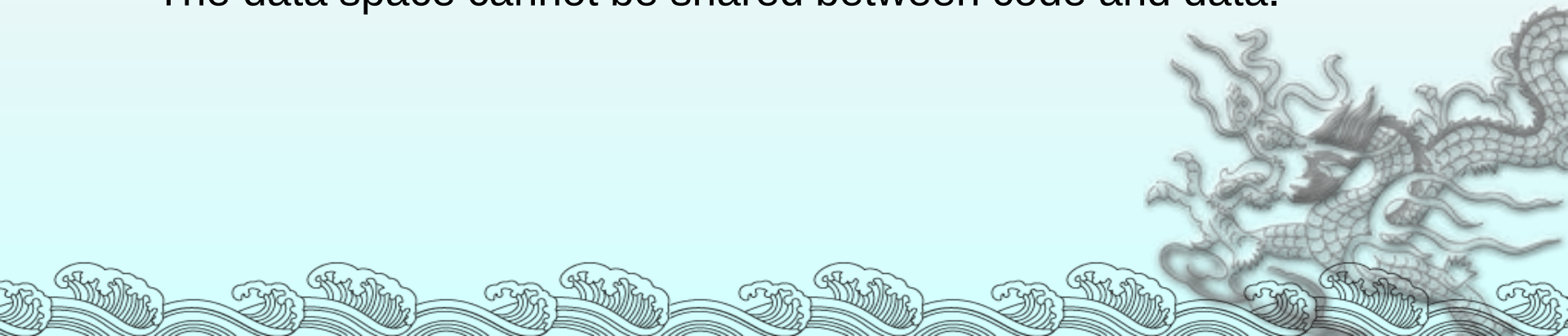
  The 8051 has another 64K bytes of memory space set aside exclusively for data storage

  *This data memory space is referred to as external memory and it is accessed only by the MOVX instruction*

- The 8051 has a total of 128K bytes of memory space

  64K bytes of code and 64K bytes of data.

  The data space cannot be shared between code and data.



# Bit Addressing

- Many microprocessors allow program to access registers and I/O ports in byte size only

  However, in many applications we need to check a single bit

- One unique and powerful feature of the 8051 is single-bit operation

  Single-bit instructions allow the programmer to set, clear, move, and complement individual bits of a port, memory, or register

  It is registers, RAM, and I/O ports that need to be bit-addressable

  *ROM, holding program code for execution, is not bit-addressable*

# Bit Addressing RAM

- The bit-addressable RAM location are 20H to 2FH

  These 16 bytes provide 128 bits of RAM bit-addressability, since

  16 × 8 = 128 ( 0 to 127 (in decimal) or 00 to 7FH )

  The first byte of internal RAM location 20H has bit address 0 to 7H

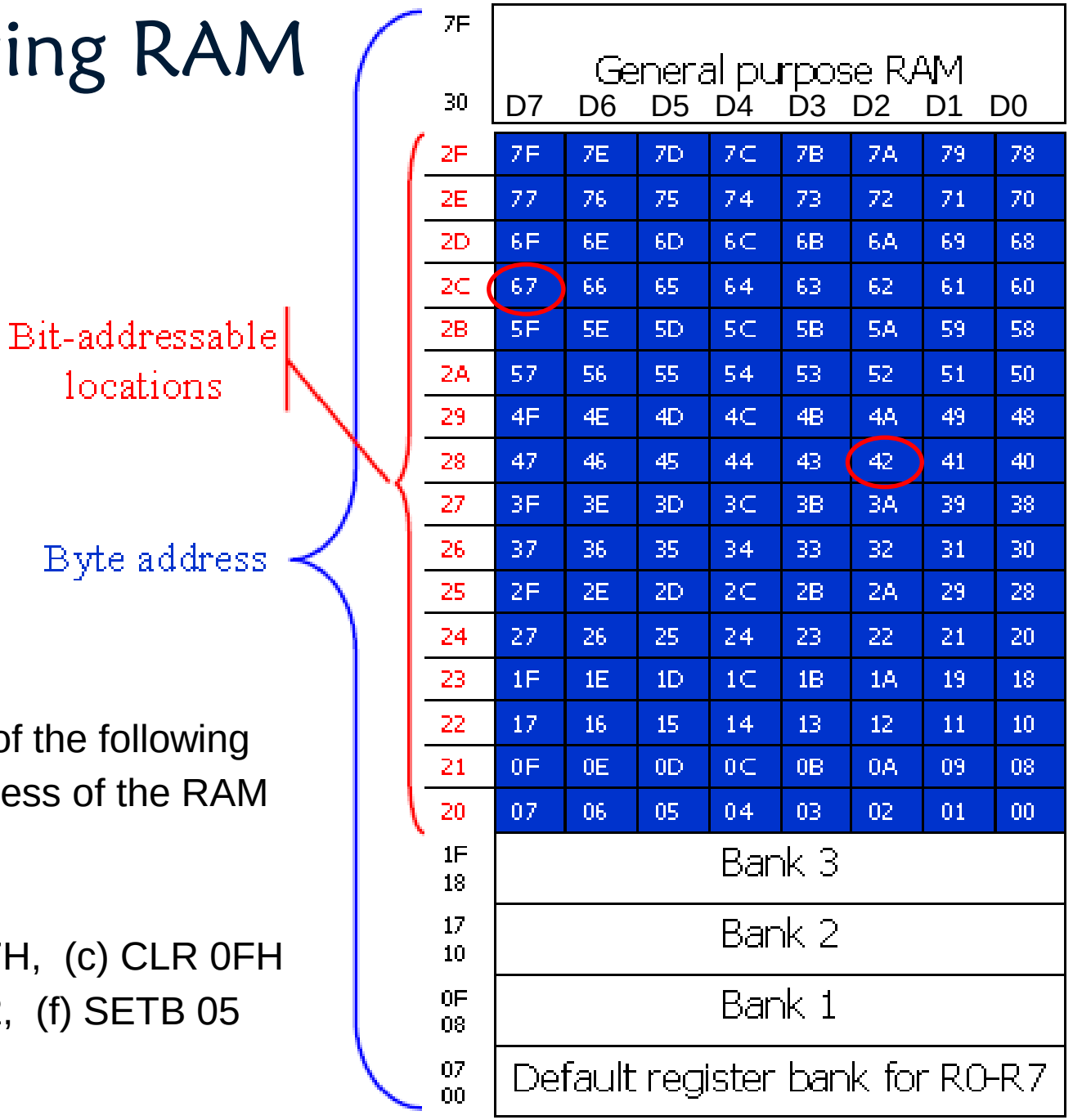
  The last byte of 2FH has bit address 78H to 7FH

- Internal RAM locations 20-2FH are both byte-addressable and bit addressable

  Bit address 00-7FH belong to RAM byte addresses 20-2FH

  Bit address 80-F7H belong to SFR P0, P1, …

# Bit Addressing RAM




**Example 5-7**

Find out to which by each of the following bits belongs. Give the address of the RAM byte in hex

(a) SETB 42H, (b) CLR 67H, (c) CLR 0FH
(d) SETB 28H, (e) CLR 12, (f) SETB 05

# Bit Addressing RAM

- To avoid confusion regarding the addresses 00 – 7FH

  The 128 bytes of RAM have the byte addresses of 00 – 7FH can be accessed in byte size using various addressing modes

  (Direct and register-indirect)

  The 16 bytes of RAM locations 20 – 2FH have bit address of 00 – 7FH

  We can use only the single-bit instructions and these instructions use only direct addressing mode

# Bit Addressing RAM

Instructions that are used for signal-bit operations are as following

| Instruction | Function |
|---|---|
| SETB bit | Set the bit (bit = 1) |
| CLR bit | Clear the bit (bit = 0) |
| CPL bit | Complement the bit (bit = NOT bit) |
| JB bit, target | Jump to target if bit = 1 (jump if bit) |
| JNB bit, target | Jump to target if bit = 0 (jump if no bit) |
| JBC bit, target | Jump to target if bit = 1, clear bit (jump if bit, then clear) |

# I/O Port Bit Addresses

- While all of the SFR registers are byte addressable, some of them are also bit addressable

  The P0 – P3 are bit addressable

- We can access either the entire 8 bits or any single bit of I/O ports P0, P1, P2, and P3 without altering the rest
- When accessing a port in a single-bit manner, we use the syntax SETB X.Y

 X is the port number P0, P1, P2, or P3

Y is the desired bit number from 0 to 7 for data bits D0 to D7

ex. SETB P1.5 sets bit 5 of port 1 high

# I/O Port Bit Addresses

Notice that when code such as SETB P1.0 is assembled, it becomes SETB 90H

- ✓ The bit address for I/O ports

  P0 are 80H to 87H

  P1 are 90H to 97H

  P2 are A0H to A7H

  P3 are B0H to B7H

| P0 | P1 | P2 | P3 | Port Bit |
|---|---|---|---|---|
| P0.0 (80) | P1.0 (90) | P2.0 (A0) | P3.0 (B0) | D0 |
| P0.1 | P1.1 | P2.1 | P3.1 | D1 |
| P0.2 | P1.2 | P2.2 | P3.2 | D2 |
| P0.3 | P1.3 | P2.3 | P3.3 | D3 |
| P0.4 | P1.4 | P2.4 | P3.4 | D4 |
| P0.5 | P1.5 | P2.5 | P3.5 | D5 |
| P0.6 | P1.6 | P2.6 | P3.6 | D6 |
| P0.7 (87) | P1.7 (97) | P2.7 (A7) | P3.7 (B7) | D7 |

# Registers Bit-Addressability

- **Bit addresses 80 – F7H belong to SFR of P0, TCON, P1,SCON, P2, etc**
- **Only registers A, B, PSW, IP, IE, ACC, SCON, and TCON are bit-addressable**

  While all I/O ports are bit-addressable
- **In PSW register, two bits are set aside for the selection of the register banks**

  Upon RESET, bank 0 is selected

  We can select any other banks using the bit-addressability of the PSW

**Example 5-8**

Write a program to save the accumulator in R7 of bank 2.

**Solution:**

| CY | AC | -- | RS1 | RS0 | OV | -- | P |
|---|---|---|---|---|---|---|---|

```
CLR  PSW.3
SETB  PSW.4
MOV  R7, A
```

**Example 5-9**

While there are instructions such as JNC and JC to check the carry flag bit (CY), there are no such instructions for the overflow flag bit (OV). How would you write code to check OV?

**Solution:**

```
JB  PSW.2, TARGET          ;jump if OV=1
```

**Example 5-10**

While a program to save the status of bit P1.7 on RAM address bit 05.

**Solution:**

```
MOV C,P1.7
MOV 05,C
```

**Example 5-11**

Write a program to see if the RAM location 37H contains an even value. If so, send it to P2. If not, make it even and then send it to P2.

**Solution:**

```
        MOV A,37H          ;load RAM 37H into ACC
        JNB ACC.0,YES      ;if D0 of ACC 0? If so jump
        INC A              ;it's odd, make it even
YES:    MOV P2,A           ;send it to P2
```

**Example 5-12**

The status of bits P1.2 and P1.3 of I/O port P1 must be saved before they are changed. Write a program to save the status of P1.2 in bit location 06 and the status of P1.3 in bit location 07

**Solution:**

```
        CLR 06 ;clear bit addr. 06
        CLR 07 ;clear bit addr. 07
         JNB P1.2,OVER ;check P1.2, if 0 then jump
        SETB 06 ;if P1.2=1,set bit 06 to 1
OVER:   JNB P1.3,NEXT ;check P1.3, if 0 then jump
        SETB 07 ;if P1.3=1,set bit 07 to 1
NEXT: ...
```

# Using BIT

The BIT directive is a widely used directive to assign the bit-addressable I/O and RAM locations

**Example 5-14**

Assume that bit P2.3 is an input and represents the condition of an oven. If it goes high, it means that the oven is hot. Monitor the bit continuously. Whenever it goes high, send a high-to-low pulse to port P1.5 to turn on a buzzer.

**Solution:**

```
OVEN_HOT BIT   P2.3
BUZZER     BIT   P1.5
HERE:      JNB  OVEN_HOT,HERE ;keep monitoring
           ACALL  DELAY
           CPL  BUZZER              ;sound the buzzer
           ACALL  DELAY
           SJMP  HERE
```

# Using EQU

Use the EQU to assign addresses

- Defined by names, like P1.7 or P2
- Defined by addresses, like 97H or 0A0H

**Example 5-24**

A switch is connected to pin P1.7. Write a program to check the status of the switch and make the following decision.

(a) If SW = 0, send "0" to P2

(b) If SW = 1, send "1" to P2

**Solution:**

```
SW          EQU    P1.7
MYDATA      EQU    P2
HERE:       MOV C,SW
            JC OVER
            MOV MYDATA,#'0'
            SJMP HERE
OVER:       MOV MYDATA,#'1'
            SJMP HERE
            END
```

```
SW          EQU  97H
MYDATA      EQU  0A0H
```

# Extra 128 Byte On-chip RAM in 8052

The 8052 has another 128 bytes of on-chip RAM with addresses 80 – FFH

- **It is often called upper memory**

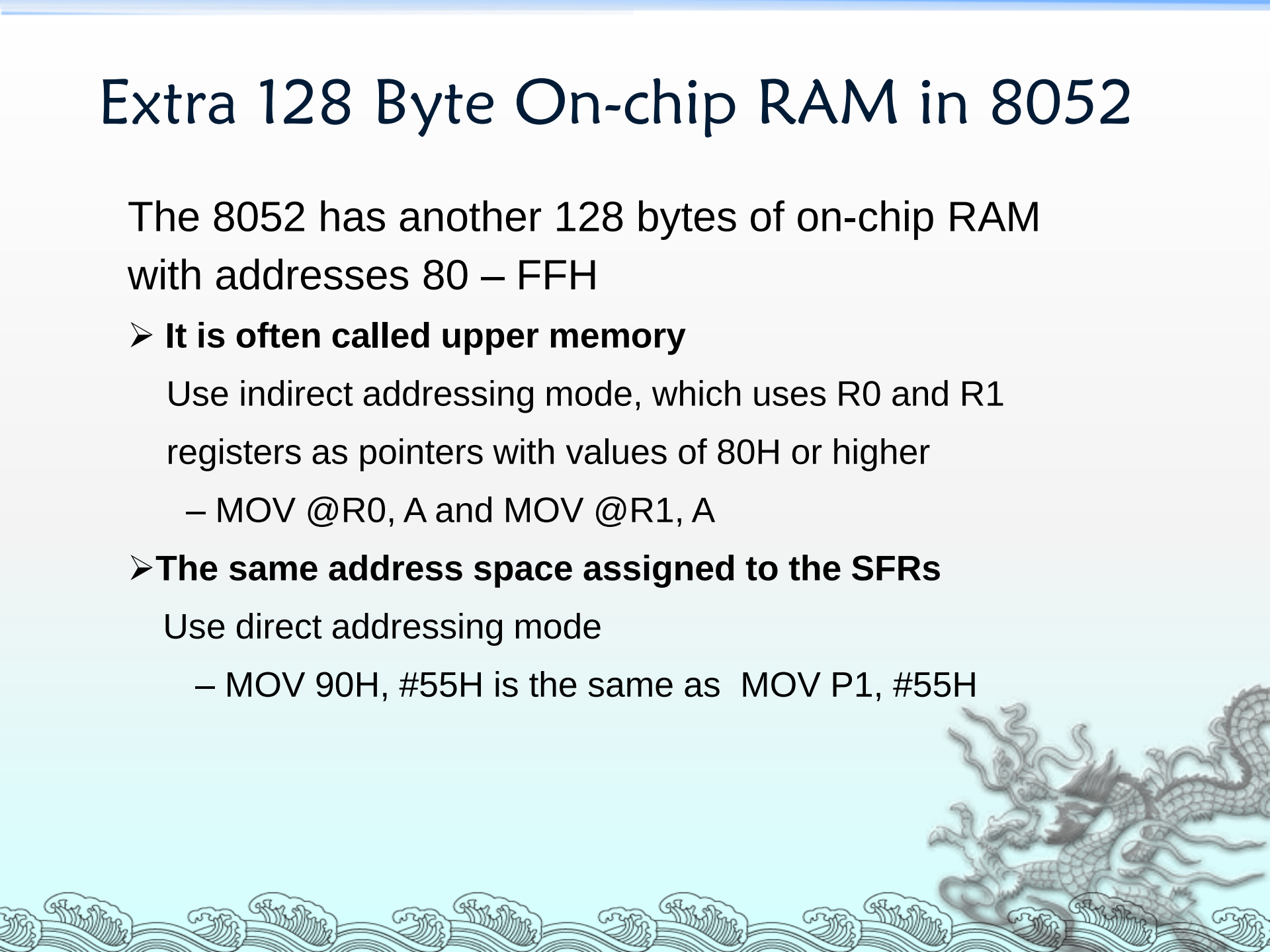
  Use indirect addressing mode, which uses R0 and R1 registers as pointers with values of 80H or higher

  – MOV @R0, A and MOV @R1, A

- **The same address space assigned to the SFRs**

  Use direct addressing mode

  – MOV 90H, #55H is the same as MOV P1, #55H



# MOV 指令集

**注意：字节数、周期**

MOV A, Rn

MOV A, direct

MOV direct, #data

MOV @Ri, #data

MOV Rn, direct

MOV DPTR, #data16

MOVX A, @Ri

MOVX A, @DPTR

MOVC A, @A+DPTR

msc 51指令速查表

| 类别 | 指令格式 | 功能简述 | 字节数 | 周期 |
|---|---|---|---|---|
| 数据传送类指令期 | MOV A, Rn | 寄存器送累加器 | 1 | 1 |
| | MOV Rn A | 累加器送寄存器 | 1 | 1 |
| | MOV A,@Ri | 内部 RAM 单元送累加器 | 1 | 1 |
| | MOV @Ri, A | 累加器送内部 RAM 单元 | 1 | 1 |
| | MOV A, #data | 立即数送累加器 | 2 | 1 |
| | MOV A, direct | 直接寻址单元送累加器 | 2 | 1 |
| | MOV direct, A | 累加器送直接寻址单元 | 2 | 1 |
| | MOV Rn #data | 立即数送寄存器 | 2 | 1 |
| | MOV direct, #data | 立即数送直接寻址单元 | 3 | 2 |
| | MOV @Ri, #data | 立即数送内部 RAM 单元 | 2 | 1 |
| | MOV direct, Rn | 寄存器送直接寻址单元 | 2 | 2 |
| | MOV Rn, direct | 直接寻址单元送寄存器 | 2 | 2 |
| | MOV direct,@Ri | 内部 RAM 单元送直接寻址单元 | 2 | 2 |
| | MOV @Ri, direct | 直接寻址单元送内部 RAM 单元 | 2 | 2 |
| | MOV direct2，direct1 | 直接寻址单元送直接寻址单元 | 3 | 2 |
| | MOV DPTR, #data16 | 16 位立即数送数据指针 | 3 | 2 |
| | MOVX A,@Ri | 外部 RAM 单元送累加器(8 位地址) | 1 | 2 |
| | MOVX @Ri, A | 累加器送外部 RAM 单元(8 位地址) | 1 | 2 |
| | MOVX A,@DPTR | 外部 RAM 单元送累加器(16 位地址) | 1 | 2 |
| | MOVX @DPTR, A | 累加器送外部 RAM 单元(16 位地址) | 1 | 2 |
| | MOVC A,@A+DPTR | 查表数据送累加器(DPTR 为基址) | 1 | 2 |
| | MOVC A,@A+PC | 查表数据送累加器(PC 为基址) | 1 | 2 |

# THANK YOU!!