

# RAFTLIB

*Presented by: Jonathan Beard  
To: C++ Now 2016*

# RAFTLIB

Alternate Titles:

*This thing I started on a plane*

# RAFTLIB

Alternate Titles:

*This thing I started on a plane  
What's this RaftLib Thingy?*

# RAFTLIB

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*This thing I started on a plane*

*What's this RaftLib Thingy?*

*OMG, Another Threading Library...*

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*What's this RaftLib Thingy?*

*OMG, Another Threading Library...*

*Why I hate parallel programming*

# RAFTLIB

Alternate Titles:

*This thing I started on a plane  
What's this RaftLib Thingy?*

*OMG, Another Threading Library...  
Why I hate parallel programming  
A self help guide for pthread anxiety*



All thoughts, opinions are my own.  
RaftLib is not a product of ARM Inc.  
Please don't ask about ARM products or  
strategy. I will scowl and not answer.

Thank you 😎

# ABOUT ME

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my website

<http://www.jonathanbeard.io>



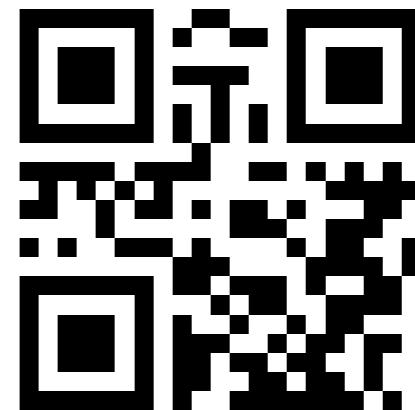
slides at

<http://goo.gl/cwT5UB>



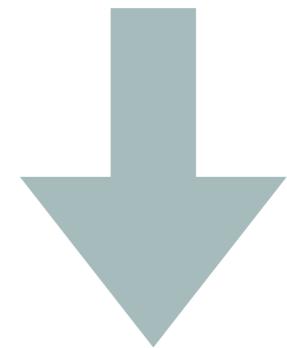
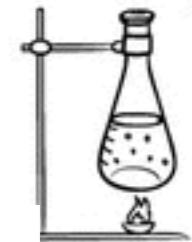
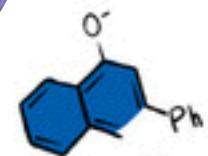
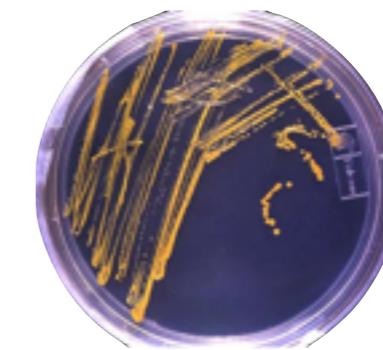
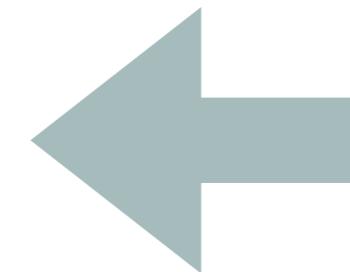
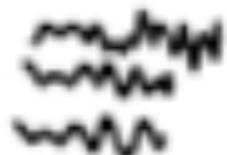
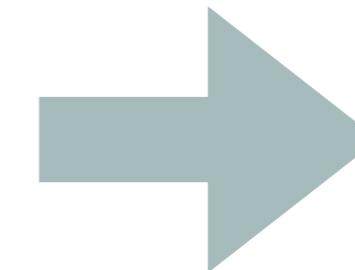
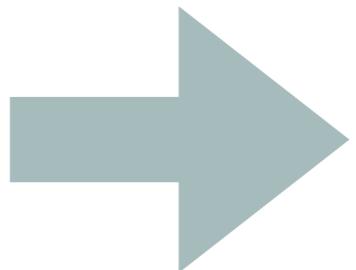
project page

[raftlib.io](http://raftlib.io)



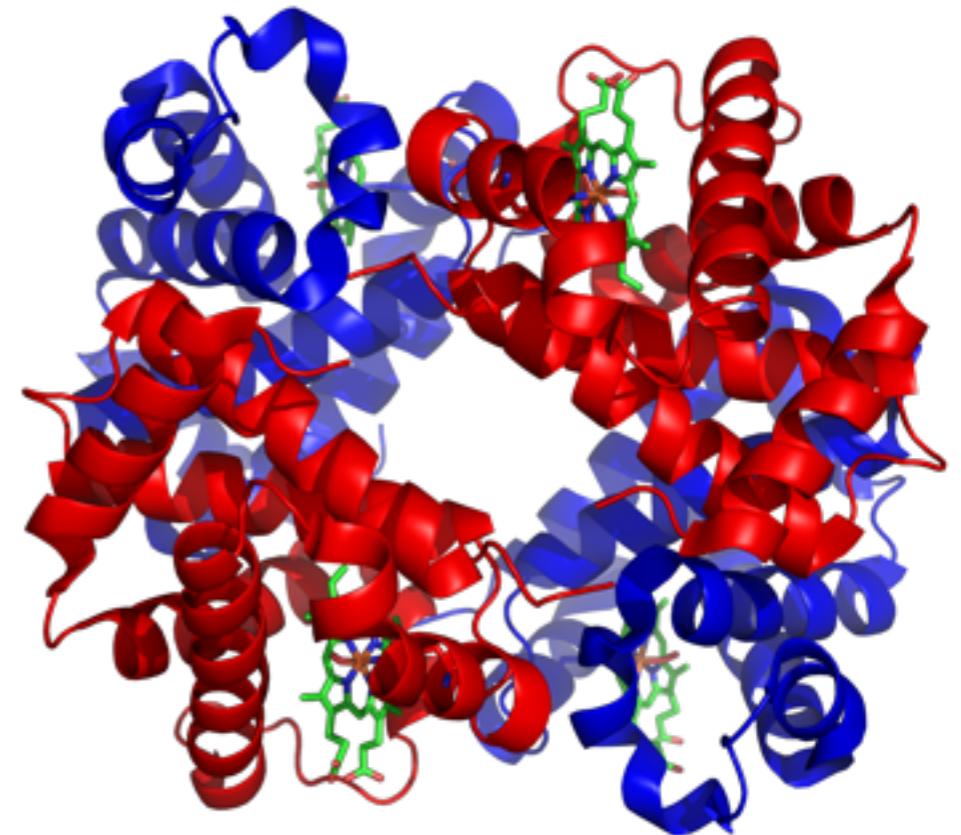
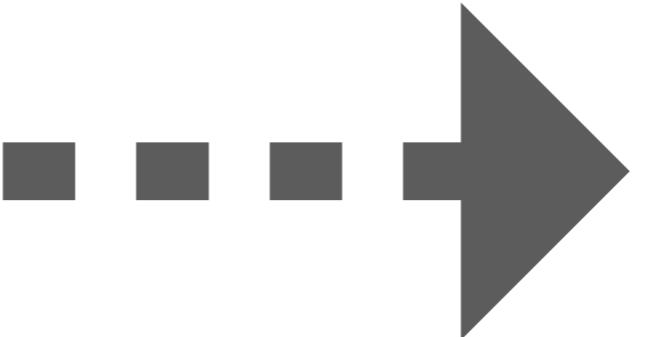
# WHERE I STARTED

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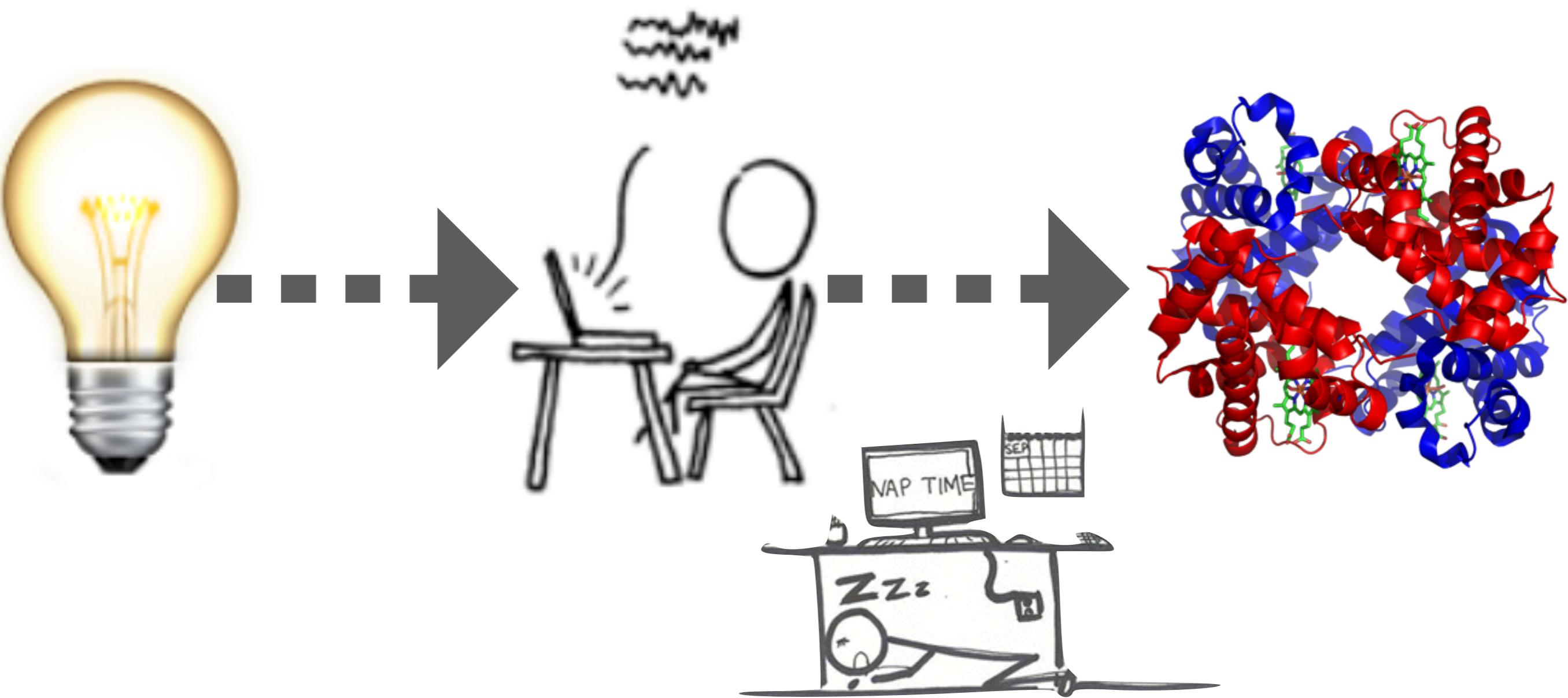
# WHERE I STARTED

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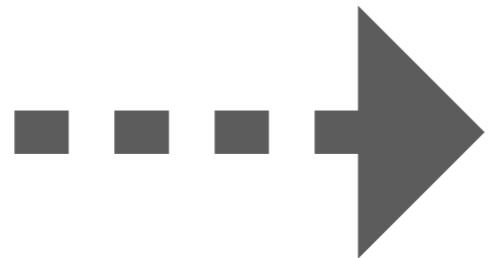
# WHERE I STARTED

---



# WHERE I STARTED

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```
const uint8_t fsm_table[STATES][CHARS] =  
{  
    /* 0 - @ */ {[9]}=0x21},  
    /* 1 - NAME */ {[0 ... 8]}=0x11,[11]=0x32},  
    /* 2 - SEQ */ {[1 ... 2]}=0x42,[5]=0x42,[8]=0x42,[11]=0x13},  
    /* 3 - \n */ {[1 ... 2]}=0x42,[5]=0x42,[8]=0x42,[10]=0x14},  
    /* 4 - NAME2 */ {[0 ... 8]}=0x54,[11]=0x15},  
    /* 5 - SCORE */ {[0 ... 10]}=0x65,[11]=0x16},  
    /* 6 - SEND/SO */ {[9]} = 0x71,[11] = 0x16}  
};
```

# NECESSITY DRIVES IDEAS

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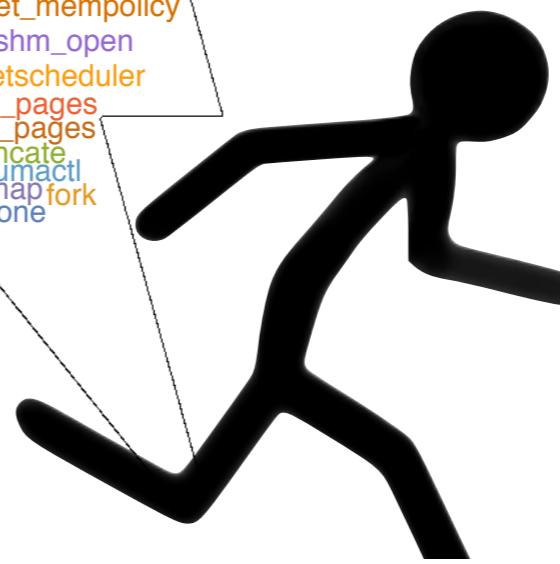
Java  
Storm Obj-C R  
Swift OpenMP Go  
X10 Fortran Perl  
C++ Javascript Ruby  
Python Chapel Rust  
CPascal Spark MPI  
Ada

# NECESSITY DRIVES IDEAS

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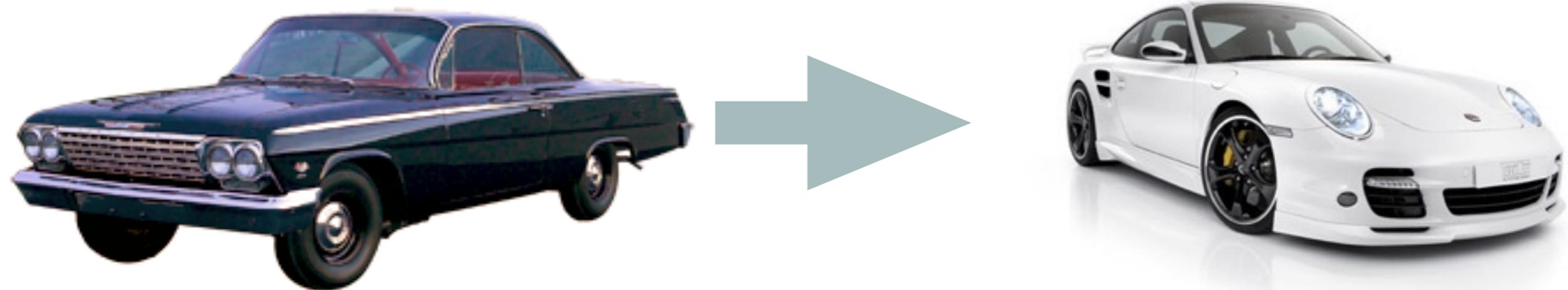
Java  
Obj-C R  
Storm Go  
Swift OpenMP Scala  
X10 Fortran Perl  
C++ Javascript Ruby  
Python Chapel Rust  
CPascal Spark MPI  
Ada

sched\_setscheduler  
pthread\_setaffinity  
pthread\_getaffinity  
posix\_memadvice  
sched\_setaffinity  
sched\_getaffinity  
unlink  
pthread\_create  
set\_mempolicy  
shm\_open  
sched\_getscheduler  
migrate\_pages  
move\_pages  
truncate  
numactl  
mmap  
fork  
clone



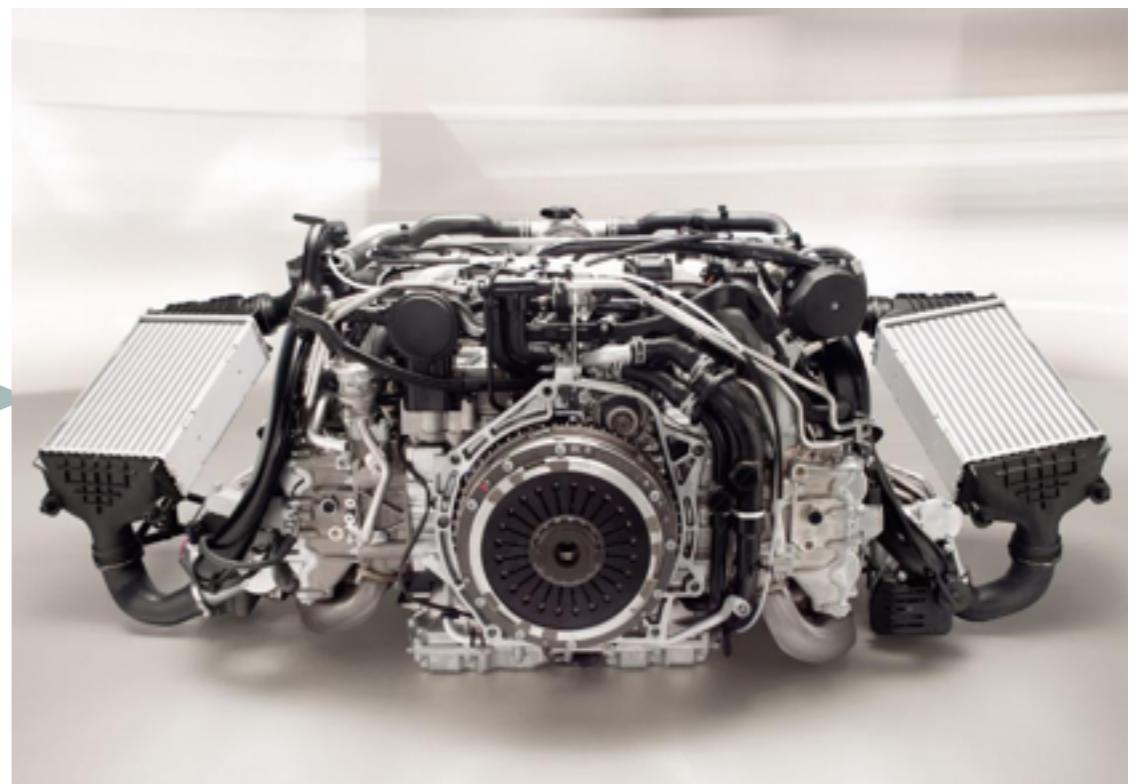
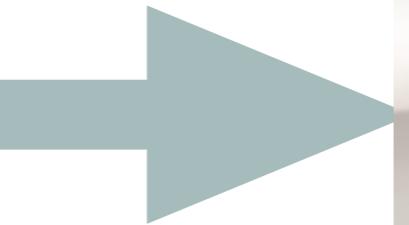
# AN (PERHAPS BAD) ANALOGY

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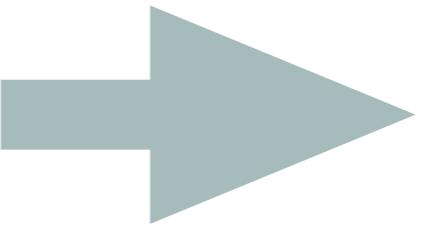
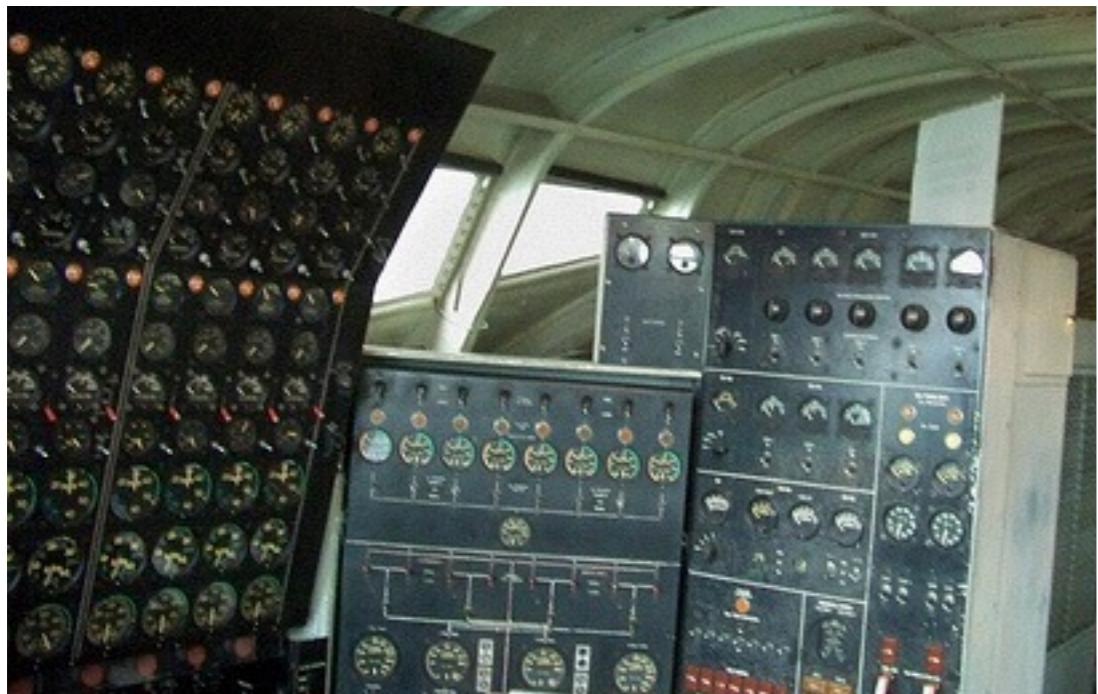
# AN (PERHAPS BAD) ANALOGY

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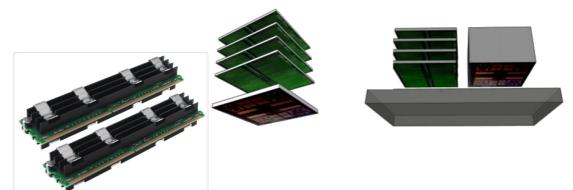
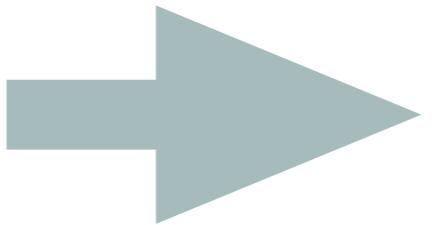
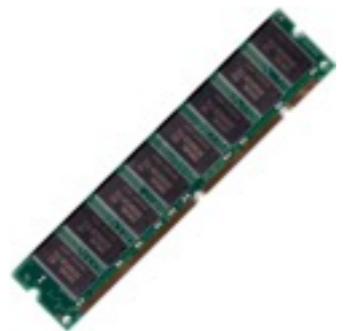
# AN (PERHAPS BAD) ANALOGY

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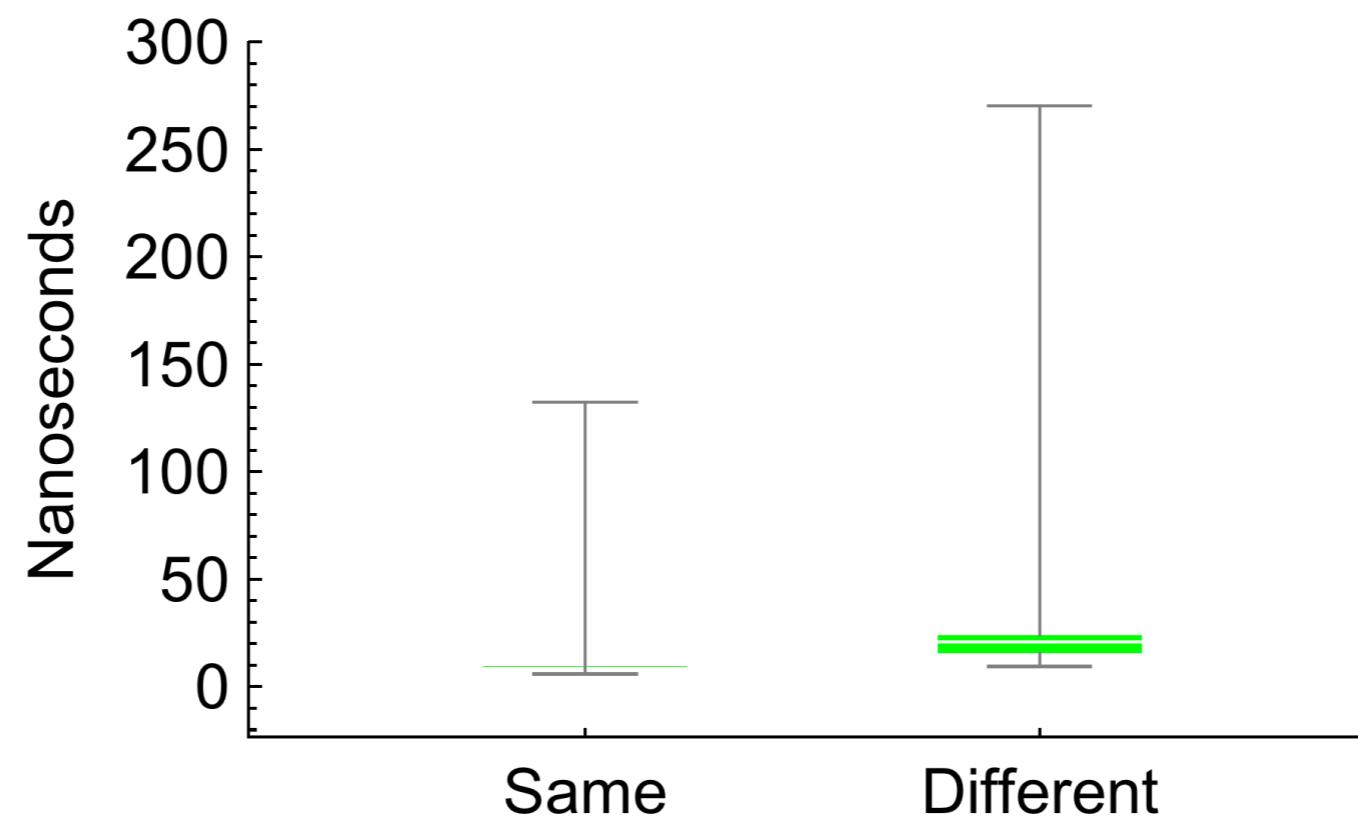
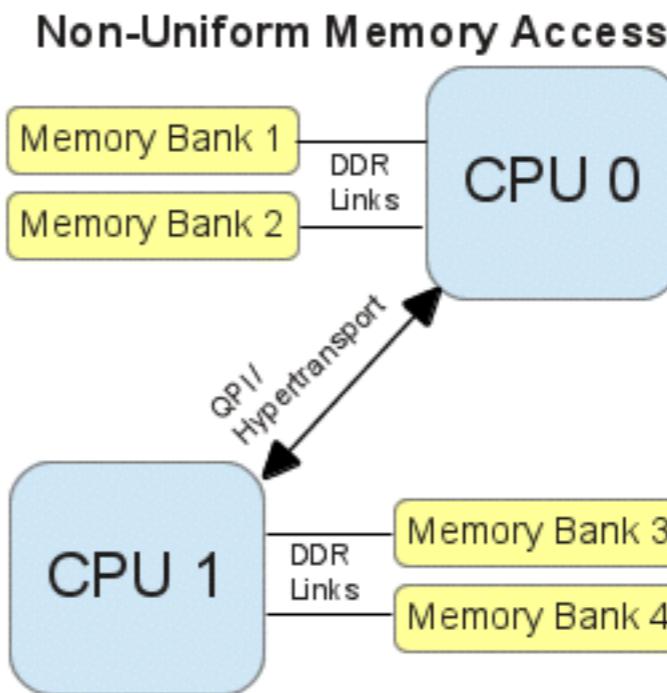
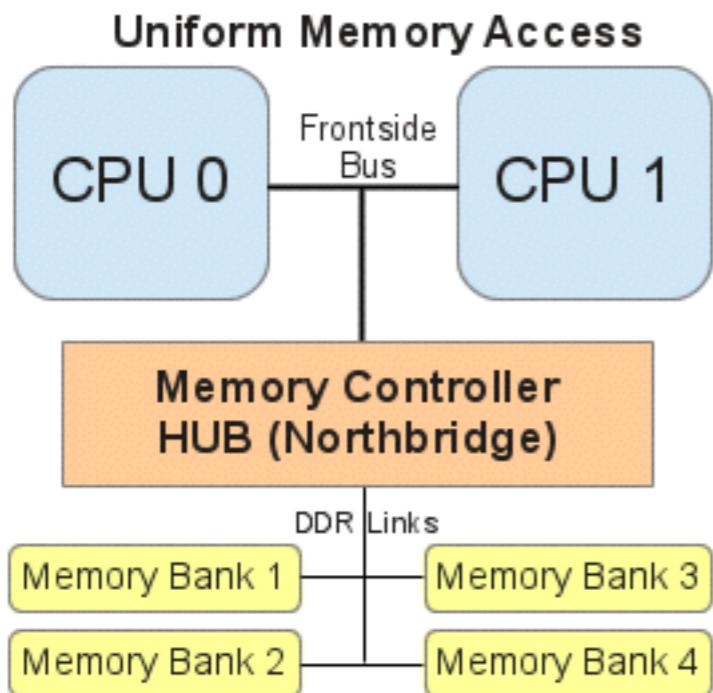
# ANALOGY PART TWO

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# TOPOLOGY

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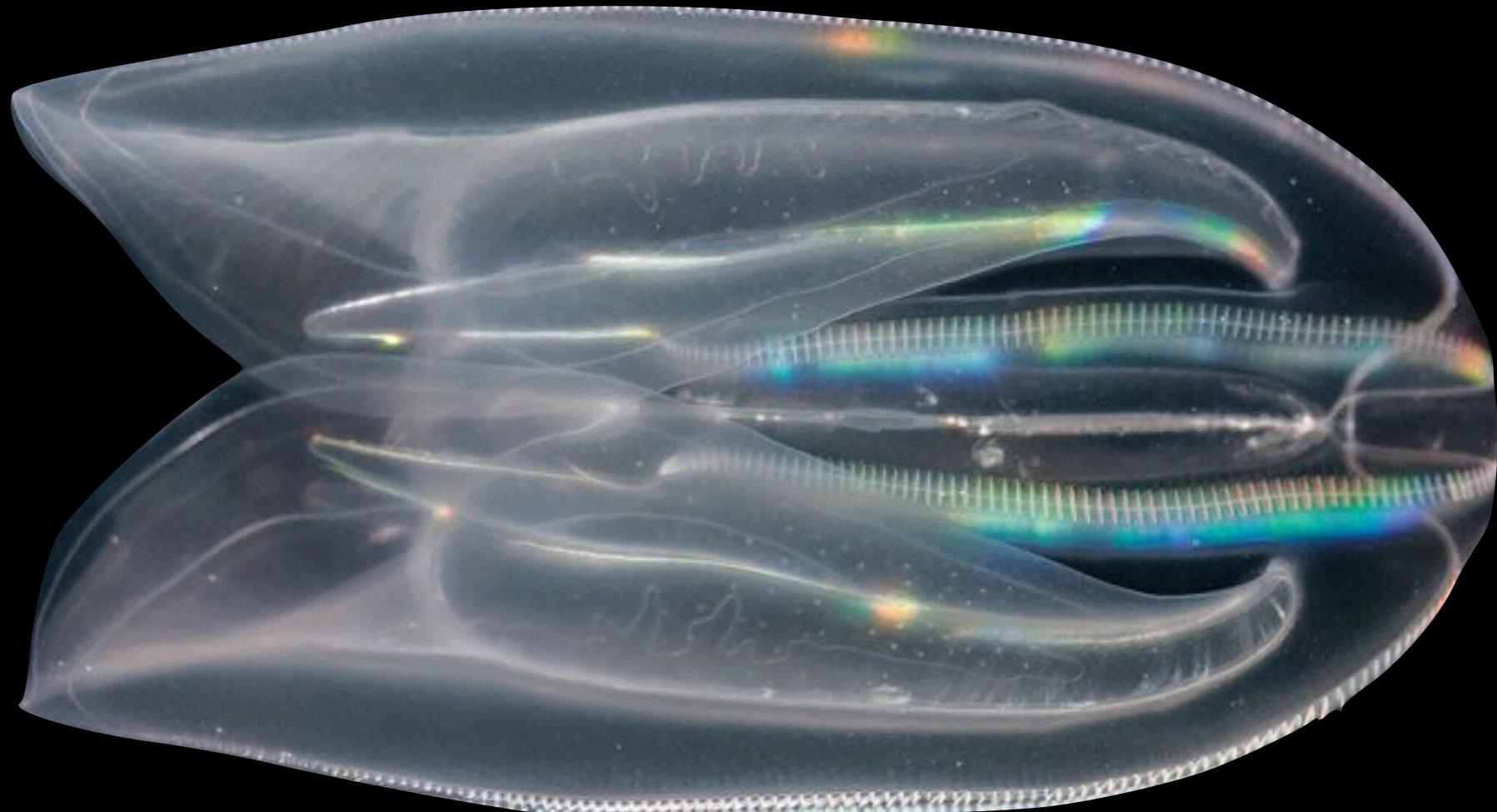
# THE JELLYFISH

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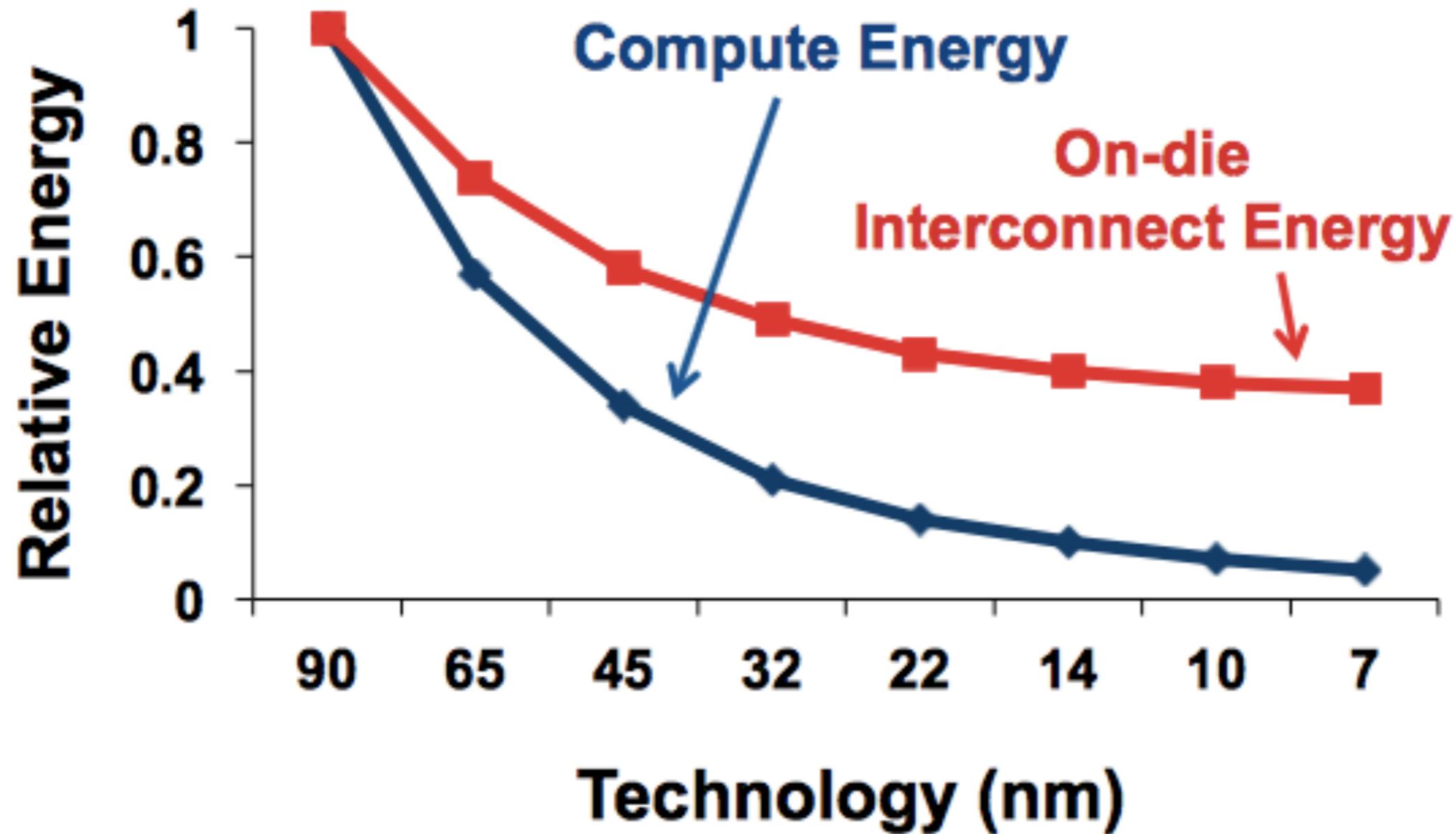
# THE FIRST ORGANISM TO OVERLAP ACCESS AND EXECUTION

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# DATA MOVEMENT DOMINATES

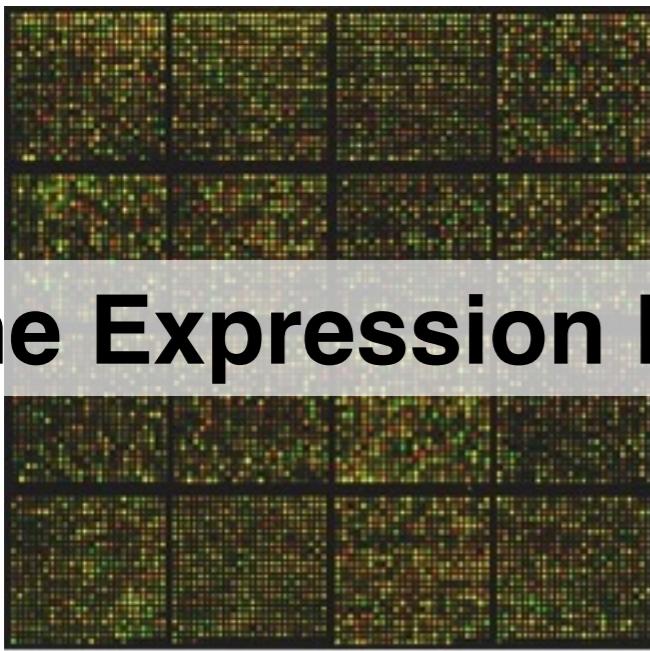
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Source: Shekhar Borkar, Journal of Lightwave Technology, 2013

# I SHOULDN'T HAVE TO CARE

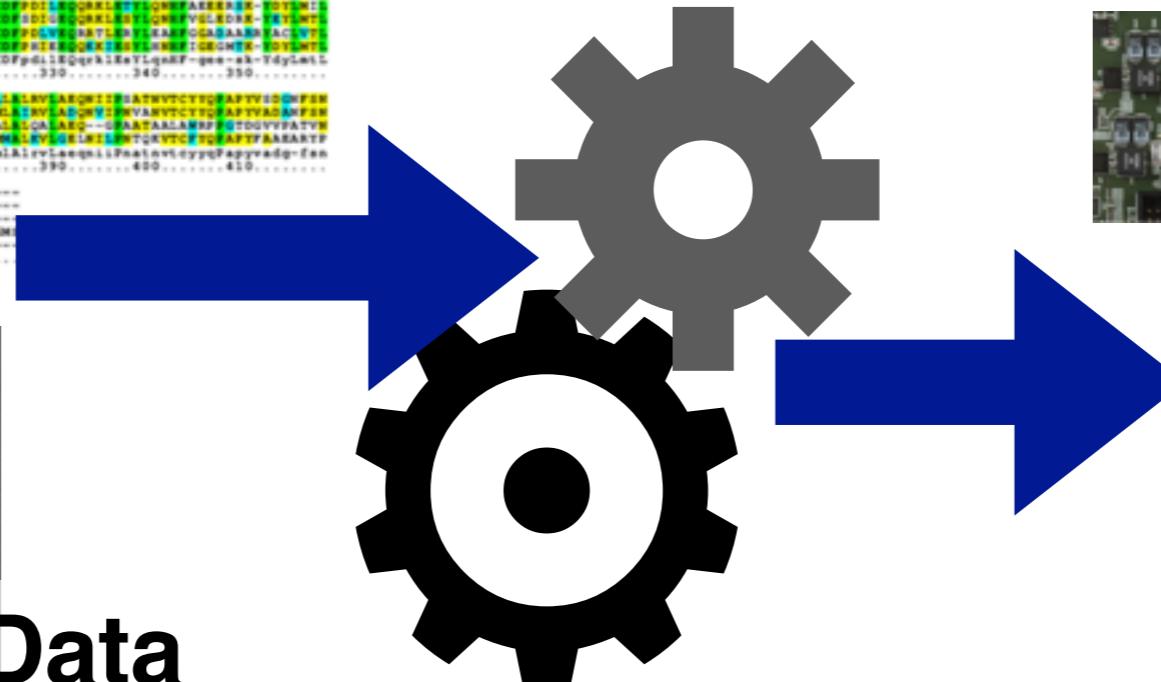
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## Multiple Sequence Alignment

Multiple Sequence Alignment of FAM46B and its Paralogs

```
Esa_FAM46C -----KAE-----ESCTRCQHS---  
Esa_FAM46A MARGGTYFANLRLSCSPYTFPLG-----FGQ-----SQQGDFGGGEGGGCLYKCFPT  
Esa_FAM46D MPSE-----GARPA-----RAASVQGTAAATAVATAPAGSS-----  
consensus -----K-----A-----S-----P-----G-----S-----  
1 10 20 30 40 50 60 70 80 90 100 110 120 130 140 150 160 170 180 190 200 210 220 230 240 250 260 270 280 290 300 310 320 330 340 350 360 370 380 390 400 410 420 430 440  
Esa_FAM46C PFTV-----CNP-----T-----V-----P-----Q-----W-----R-----S-----  
Esa_FAM46A CATTC-----CNP-----T-----V-----P-----Q-----W-----R-----S-----  
Esa_FAM46D CAGQ-----CNP-----T-----V-----P-----Q-----W-----R-----S-----  
consensus LIP-----CNP-----T-----V-----P-----Q-----W-----R-----S-----  
1 10 20 30 40 50 60 70 80 90 100 110 120 130 140 150 160 170 180 190 200 210 220 230 240 250 260 270 280 290 300 310 320 330 340 350 360 370 380 390 400 410 420 430 440  
Esa_FAM46C D-----TQERK-----R-----MCS-----F-----D-----Q-----A-----E-----K-----X-----  
Esa_FAM46A D-----TQERK-----R-----MCS-----F-----D-----Q-----A-----E-----K-----X-----  
Esa_FAM46D D-----TQERK-----R-----MCS-----F-----D-----Q-----A-----E-----K-----X-----  
consensus V-----TQERK-----R-----MCS-----F-----D-----Q-----A-----E-----K-----X-----  
1 10 20 30 40 50 60 70 80 90 100 110 120 130 140 150 160 170 180 190 200 210 220 230 240 250 260 270 280 290 300 310 320 330 340 350 360 370 380 390 400 410 420 430 440  
Esa_FAM46C R-----S-----I-----L-----M-----A-----F-----D-----Y-----C-----T-----Q-----A-----F-----R-----  
Esa_FAM46A R-----S-----I-----L-----M-----A-----F-----D-----Y-----C-----T-----Q-----A-----F-----R-----  
Esa_FAM46D R-----S-----I-----L-----M-----A-----F-----D-----Y-----C-----T-----Q-----A-----F-----R-----  
consensus S-----S-----I-----L-----M-----A-----F-----D-----Y-----C-----T-----Q-----A-----F-----R-----  
1 10 20 30 40 50 60 70 80 90 100 110 120 130 140 150 160 170 180 190 200 210 220 230 240 250 260 270 280 290 300 310 320 330 340 350 360 370 380 390 400 410 420 430 440  
Esa_FAM46C T-----V-----A-----P-----P-----Y-----S-----A-----P-----F-----P-----F-----  
Esa_FAM46A T-----V-----A-----P-----P-----Y-----S-----A-----P-----F-----P-----F-----  
Esa_FAM46D T-----V-----A-----P-----P-----Y-----S-----A-----P-----F-----P-----F-----  
consensus T-----V-----A-----P-----P-----Y-----S-----A-----P-----F-----P-----F-----
```



# I SHOULDN'T HAV

.....

Multiple  
Align

Multiple Sequence Align

```
Esa_FAM6C -----KAE-----
Esa_FAM6A MARGGTYFANLRLSASP1
Esa_FAM6B HPSSEGAGRRAA3QVG1
Esa_FAM6D -----MSTIA-----
consensus -----K-----A-----1
                               10

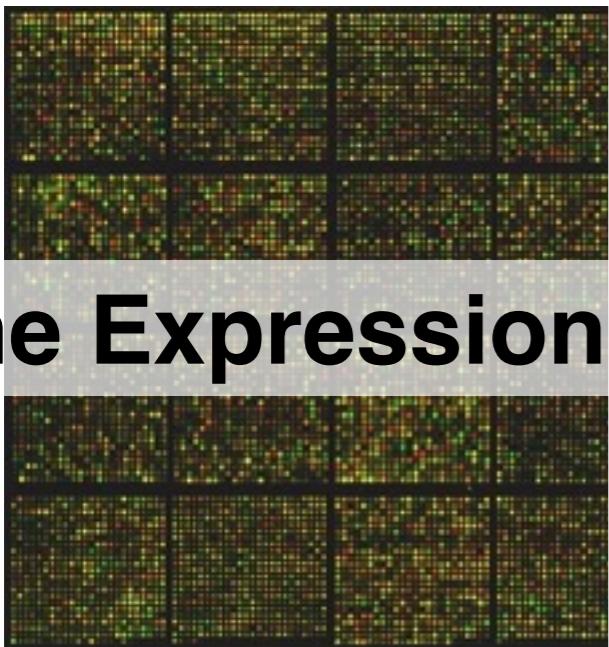
Esa_FAM6C -----FESTVYVETV1
Esa_FAM6A M-----FESTVYVETV1
Esa_FAM6B M-----FESTVYVETV1
Esa_FAM6D -----FESTVYVETV1
consensus M-----FESTVYVETV1
                               100

Esa_FAM6C FPTVCFDPF1TFPFTV1
Esa_FAM6A CPTVCFDPF1TFPFTV1
Esa_FAM6B CPTVCFDPF1TFPFTV1
Esa_FAM6D DPTVCFDPF1TFPFTV1
consensus DPTVCFDPF1TFPFTV1
                               240....250

Esa_FAM6C RQ-----TQERIKK1R
Esa_FAM6A RQ-----AS-DKIKK1R
Esa_FAM6B RQ-----AS-DKIKK1R
Esa_FAM6D RQ-----AS-DKIKK1R
consensus VqgFrP-----AS-DKIKK1R
                               300....310

Esa_FAM6C R-----R-----R-----R
Esa_FAM6A R-----R-----R-----R
Esa_FAM6B R-----R-----R-----R
Esa_FAM6D R-----R-----R-----R
consensus R-----R-----R-----R
                               360....370

Esa_FAM6C T-----V-----P-----T
Esa_FAM6A T-----V-----P-----T
Esa_FAM6B T-----V-----P-----T
Esa_FAM6D T-----V-----P-----T
consensus T-----V-----P-----T
                               420....430
```



# HARDWARE

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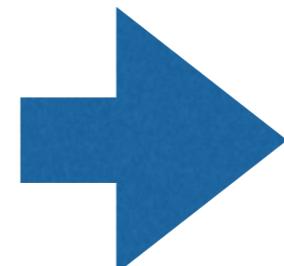
**1984**

**Cray X-MP/48**  
**\$19 million / GFLOP**

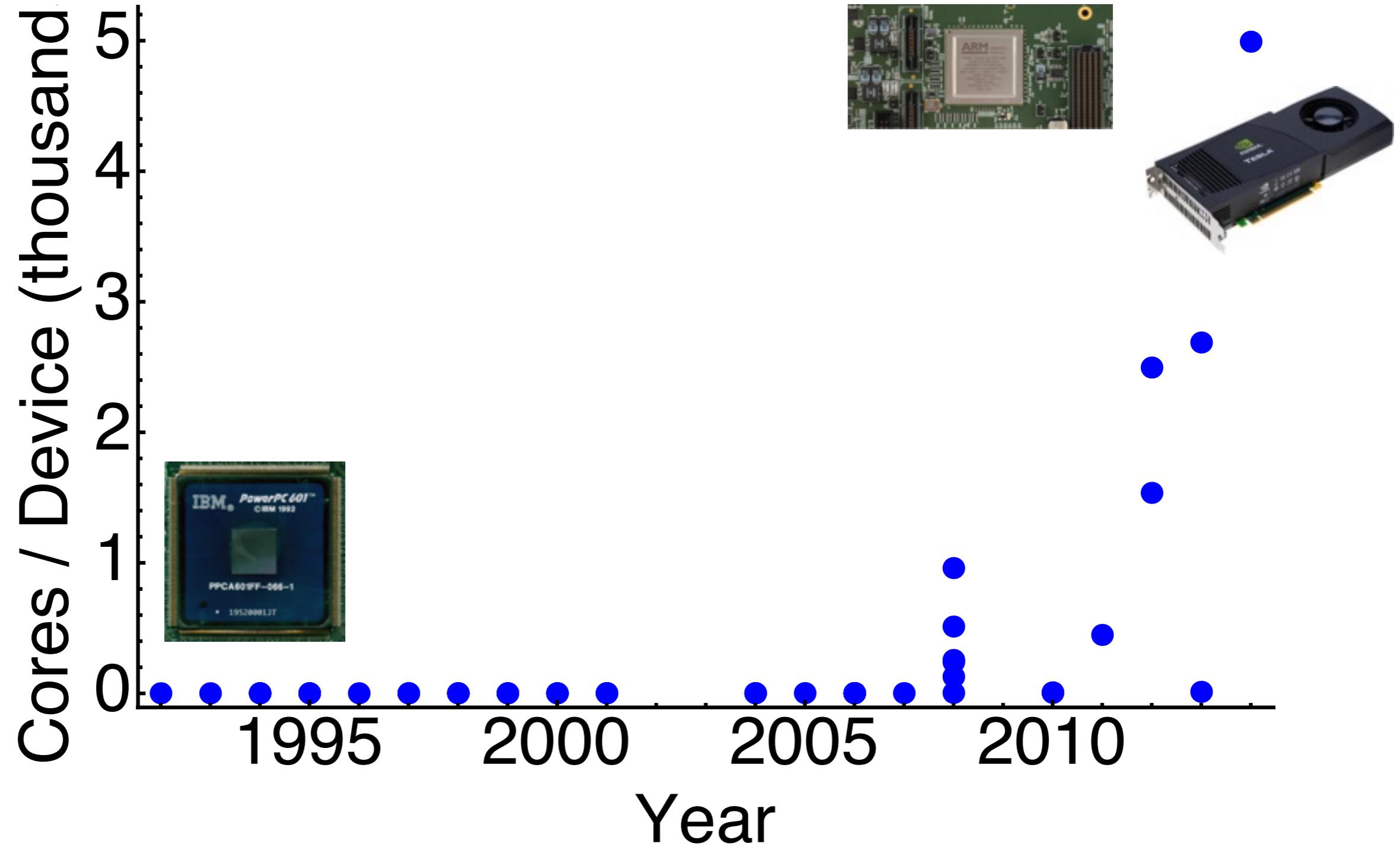


**2015**

**\$.08 / GFLOP**



# CORES PER DEVICE



# FINANCIAL INCENTIVE

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*Sequential JS:* \$4-7/line

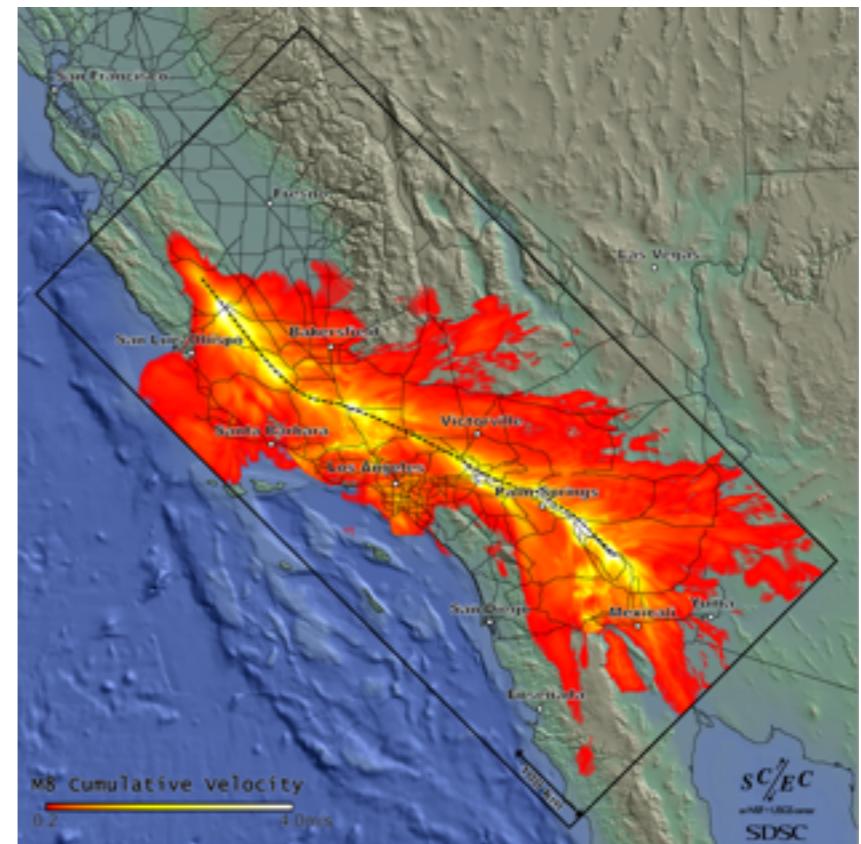
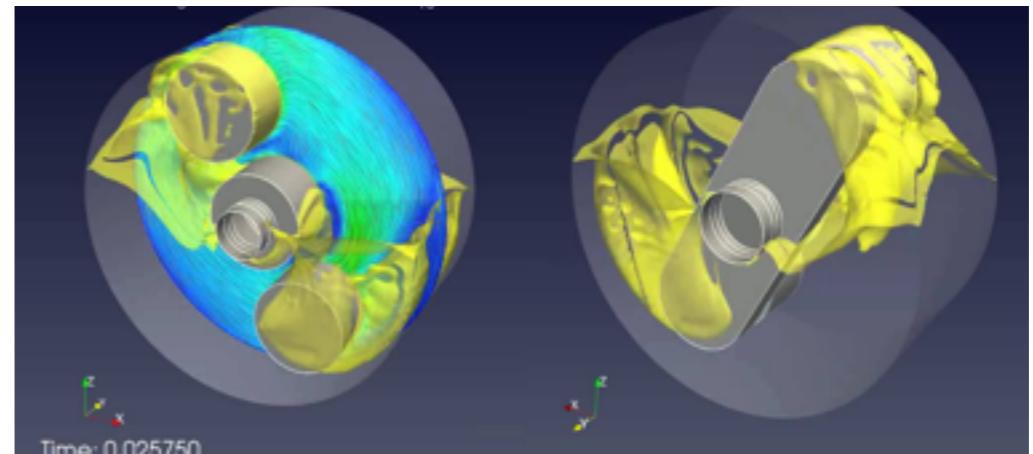
*Sequential Java:* \$5-10/line

*Embedded Code:* \$30-50/line

*HPC Code:* \$100/line

# WHY YOU SHOULD CARE

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# PRODUCTIVITY / EFFICIENCY AN EQUALIZER

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- Titan SC estimates for porting range from 5, to > 20 million USD
- Most code never really optimized for machine topology, wasting \$\$ (time/product) and energy
- **Getting the most out of what you have is an equalizer**



Google Cloud Platform



Java  
Storm Obj-C R  
Swift Go  
OpenMP Scala  
X10 Fortran Perl  
C++ Javascript Ruby  
Python Chapel Rust  
CPascal Spark MPI  
Ada

# HAVES AND HAVE NOTS

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## Gov't / Big Business

- *Lots of \$\$*
- *Can hire the best people*
- *Can acquire the rest*
- *Plenty of compute resources*

## Start-up / Small Government

- *Not a lot of \$\$*
- *Often can't hire the best people*
- *Left to the mercy of cloud providers*



# AN IDEA

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Let's make  
computers super  
fast, and easy to  
program



# WHERE TO START



## Brook for GPUs: Stream Computing on Graphics Hardware

Ian Buck Tim Foley Daniel Horn Jeremy Sugerman Kayvon Fatahalian Mike Houston Pat Hanrahan  
Stanford University

### Abstract

In this paper, we present Brook for GPUs, a system for general-purpose computation on programmable graphics hardware. Brook extends C to include simple data-parallel constructs, making use of the GPU as a streaming coprocessor. We present a compiler and runtime system that abstracts and virtualizes many aspects of graphics hardware. In addition, we present an analysis of the effectiveness of the GPU as a compute engine compared to the CPU, to determine the limits of the GPU as a streaming processor. Finally, we evaluate our system with five applications, the SAXPY and SGEMV BLAS operators, image segmentation, FFT, and ray tracing. For these applications, we demonstrate that our Brook implementations perform comparable to hand-coded GPU code and up to seven times faster than their CPU counterparts.

**CR Categories:** I.3.1 [Computer Graphics]: Hardware Architecture—Graphics processors; D.3.2 [Programming Languages]: Language Classifications—Parallel Languages

**Keywords:** Programmable Graphics Hardware, Data Parallel Computing, Stream Computing, GPU Computing, Brook

### 1 Introduction

In recent years, commodity graphics hardware has rapidly evolved from being a fixed-function pipeline into having programmable vertex and fragment shaders. While this programmability is intended for real-time shading, it has been observed that these processors feature instruction sets generic enough to perform computation beyond the domain of rendering. Applications such as linear algebra [Krueger et al. 2003], physics [Woo et al. 2003], [Huang et al. 2003], and a compiler and tracer [Purcell et al. 2003; Carr et al. 2002] have been demonstrated on GPUs.

Originally, GPUs could only be programmed using assembly languages. Microsoft's HLSL, NVIDIA's Cg, and OpenGL's GLSL are examples of high-level assembly level, C-like programming language [Microsoft 2003; Microsoft 2003; Kessenich et al. 2003]. However, these languages do not assist the programmer in controlling the aspects of the graphics hardware such as shading, texture assembly, rendering order programs, or constructing scene primitives. As a result, the implementation of applications requires extensive knowledge of the latest graphics APIs as well as an understanding of the features and limitations of

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*SIGMOD '08, June 9–12, 2008, Vancouver, BC, Canada.*  
Copyright 2008 ACM 978-1-60558-102-6/08/06 ...\$5.00

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### ABSTRACT

In this paper, we present SPADE – the System S declarative stream processing engine. System S is a large-scale, distributed data stream processing middleware under development at IBM T. J. Watson Research Center. As a front-end for rapid application development for System S, SPADE provides (1) an intermediate language for flexible composition of parallel and distributed data-flow graphs, (2) a toolkit of type-generic, built-in stream processing operators, that support scalar as well as vectorized processing and can seamlessly inter-operate with user-defined operators, and (3) a rich set of stream adapters to ingest/publish data from/to outside sources. More importantly, SPADE automatically brings performance optimization and scalability to System S applications. To that end, SPADE employs a code generation framework to create highly-optimized applications that run natively on the Stream Processing Core (SPC), the execution and communication substrate of System S, and take full advantage of other System S services. SPADE allows developers to construct their applications with fine granular stream operators without worrying about the performance implications that might exist, even in a distributed system. SPADE's optimizing compiler automatically maps applications into appropriately sized execution units in order to minimize communication overhead, while at the same time exploiting available parallelism. By virtue of the scalability of the System S runtime and SPADE's effective code generation and optimization, we can scale applications to a large number of nodes. Currently, we can run SPADE jobs on  $\approx 500$  processors within more than 100 physical nodes in a tightly connected cluster environment. SPADE has been in use at IBM Research to create real-world streaming applica-

tion

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*SIGMOD '08, June 9–12, 2008, Vancouver, BC, Canada.*  
Copyright 2008 ACM 978-1-60558-102-6/08/06 ...\$5.00

modern hardware. In addition, the user is forced to express parallel primitives in terms of graphics hardware primitives such as textures and triangles. As a result, general-purpose GPU computing is limited to only the most advanced graphics developers.

This paper presents *Brook*, a programming environment that provides developers with a view of the GPU as a streaming coprocessor. The main contributions of this paper are:

- The presentation of the Brook stream programming model.
- The demonstration that the GPU can be used as a streaming processor.
- The demonstration of how various GPU hardware limitations can be virtualized or extended using our compiler and runtime system; specifically, the GPU memory system, the number of supported streams, and support for user-defined functions.
- The presentation of a cost model for comparing GPU vs. CPU performance tradeoffs to better understand under what circumstances the GPU outperforms the CPU.

### TOOLS FOR SIMULATION AND ANALYSIS OF HETEROGENEOUS PIPELINED

### WASHINGTON UNIVERSITY

### SCHOOL OF ENGINEERING AND APPLIED SCIENCE

### DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

### X-SIM AND X-EVAL:

### by

Saurabh Gayen

Prepared under the direction of Professor Mark A. Franklin

### SPADE: The System S Declarative Stream Processing Engine

Washington University in partial fulfillment of the

requirements for the degree of

M. KUNDUNG WATSON  
IBM Thomas J. Watson Research Center, Hawthorne, NY 10532 USA  
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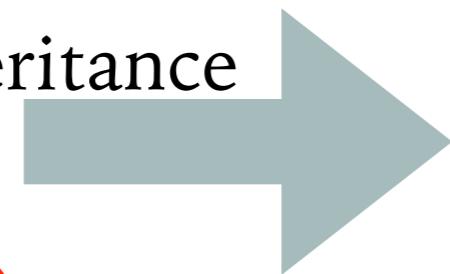
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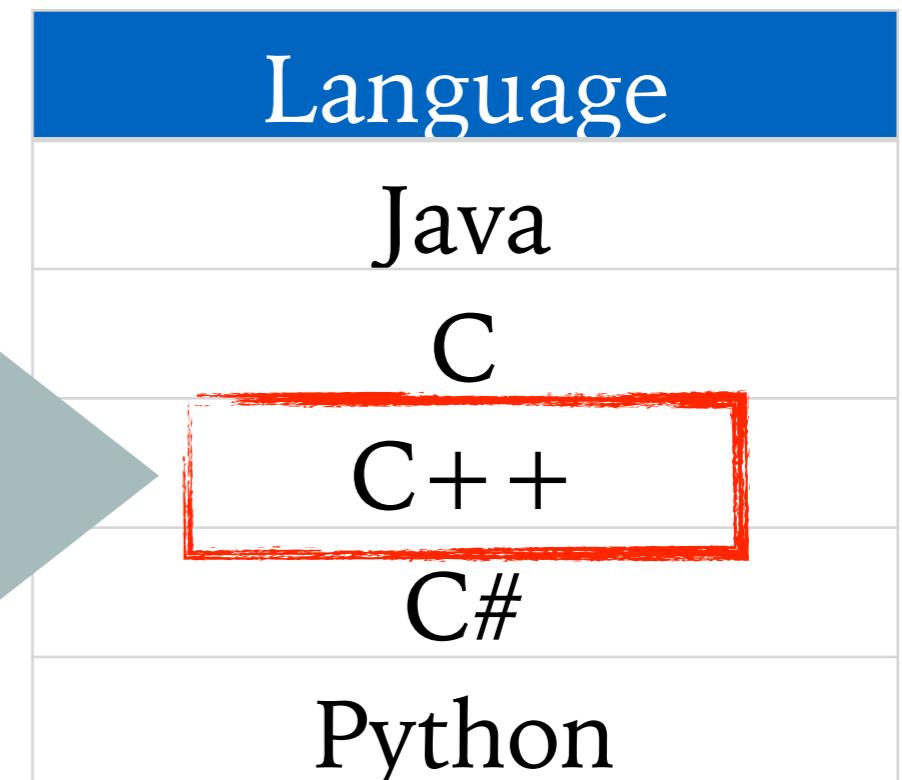
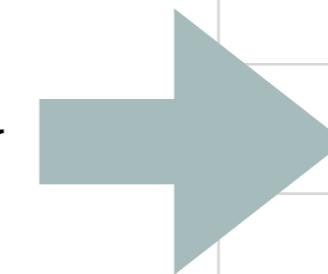
# WHERE TO START

---

- ~~Fast Languages~~
- Object oriented
  - Multiple inheritance
  - Simple types
  - .... lots more



Library



## C++ Streaming Template Library

Simplifies parallelization of code

Abstracts the details

Auto-manages allocation types, data movement.

software download: <http://raftlib.io>

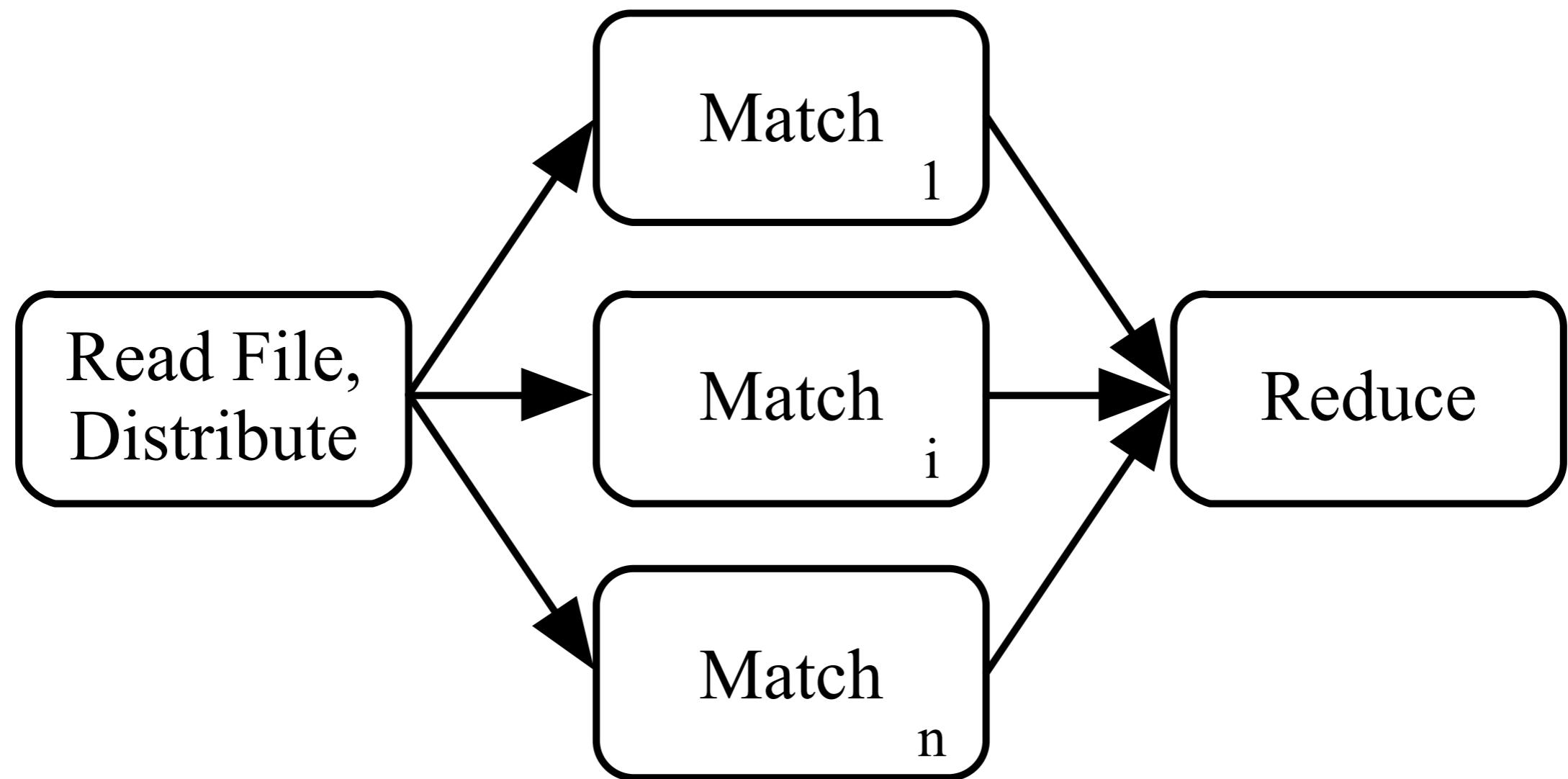
# ISSUES THAT NEED SOLUTIONS

---

- Usable Generic API
- Where to run the code
- Where to allocate memory
- How big of allocations to start off with
- How many user vs. kernel threads
- ...Gilligan had a better chance of having a 3 hour tour
- ...than we have of fitting all the issues on a single slide

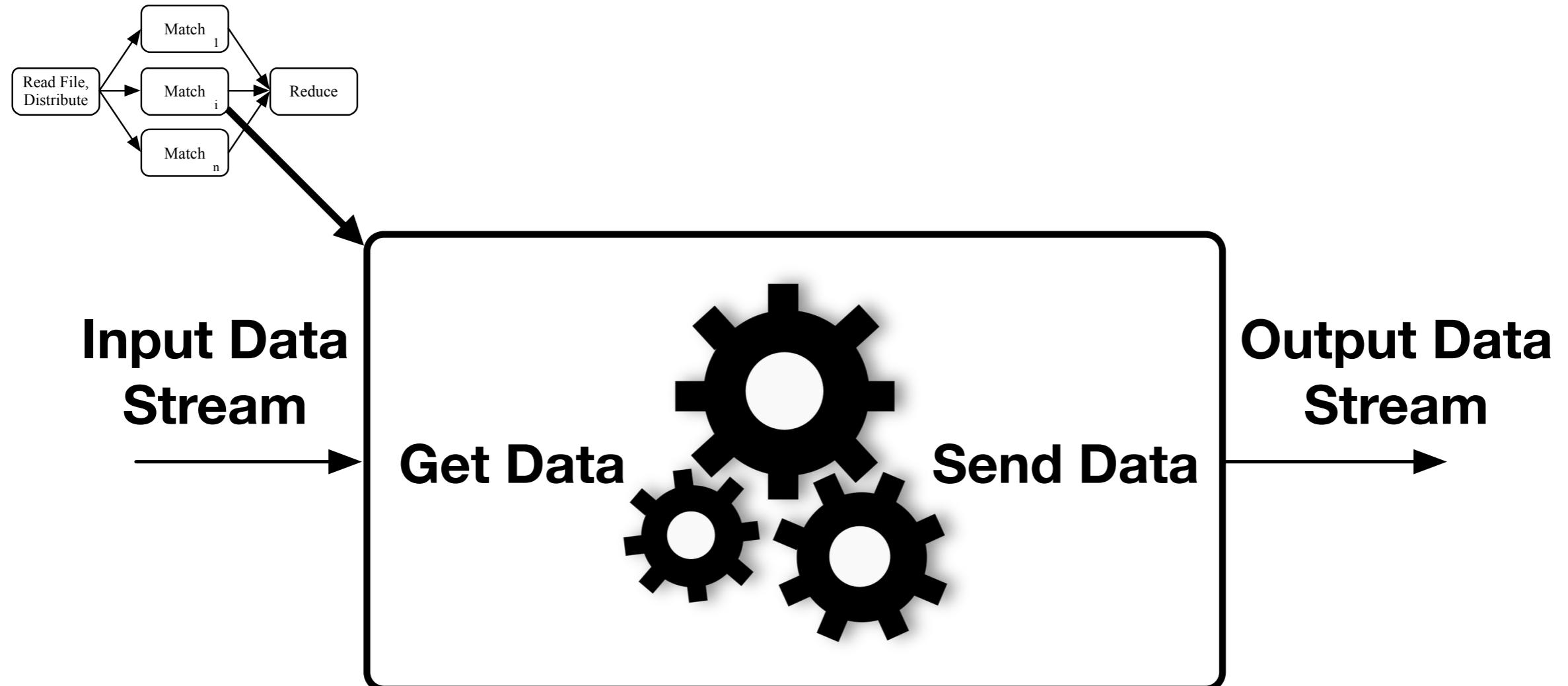
# STREAM PROCESSING (STRING SEARCH EXAMPLE)

---



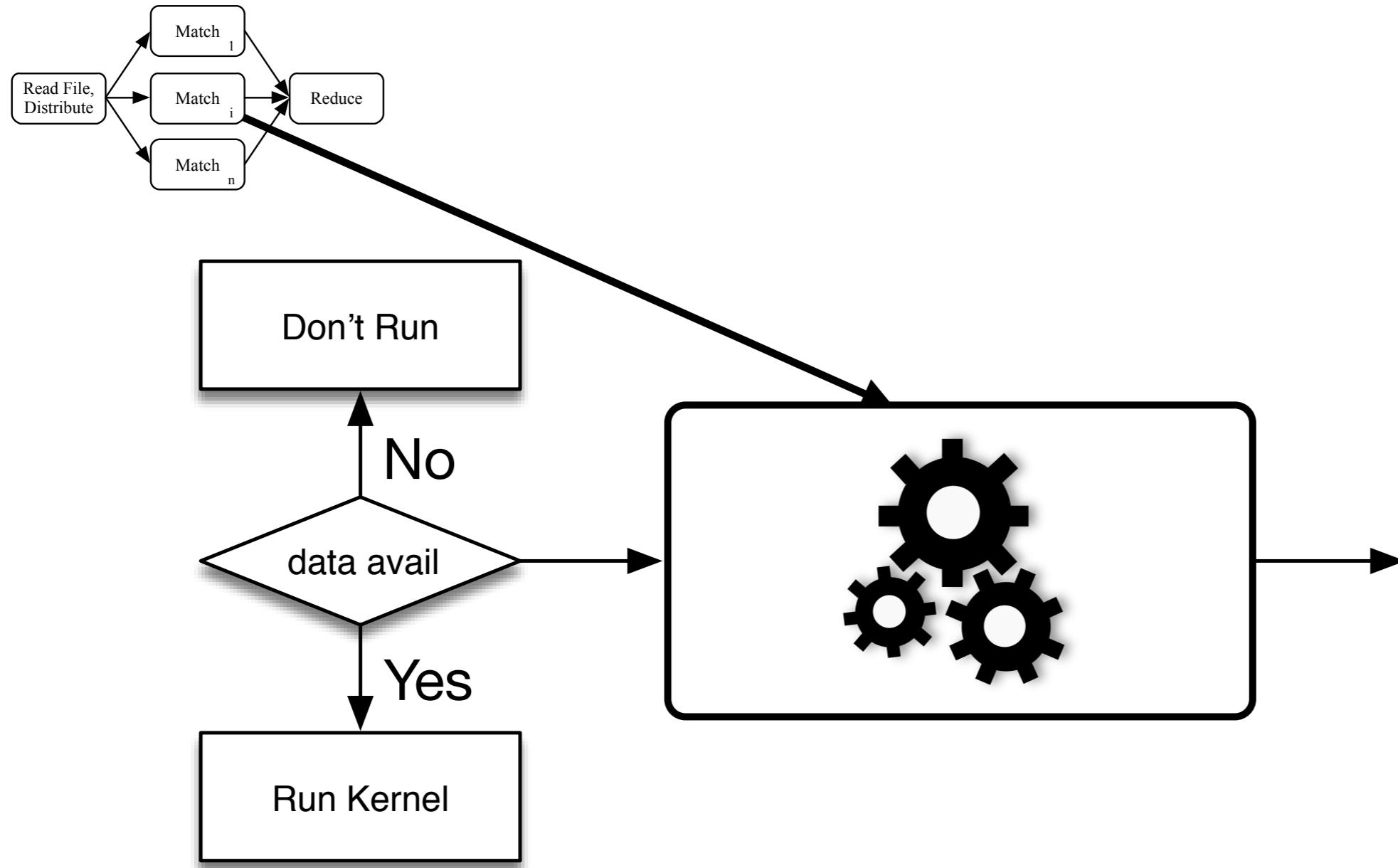
# STREAM PROCESSING (STRING SEARCH EXAMPLE)

---



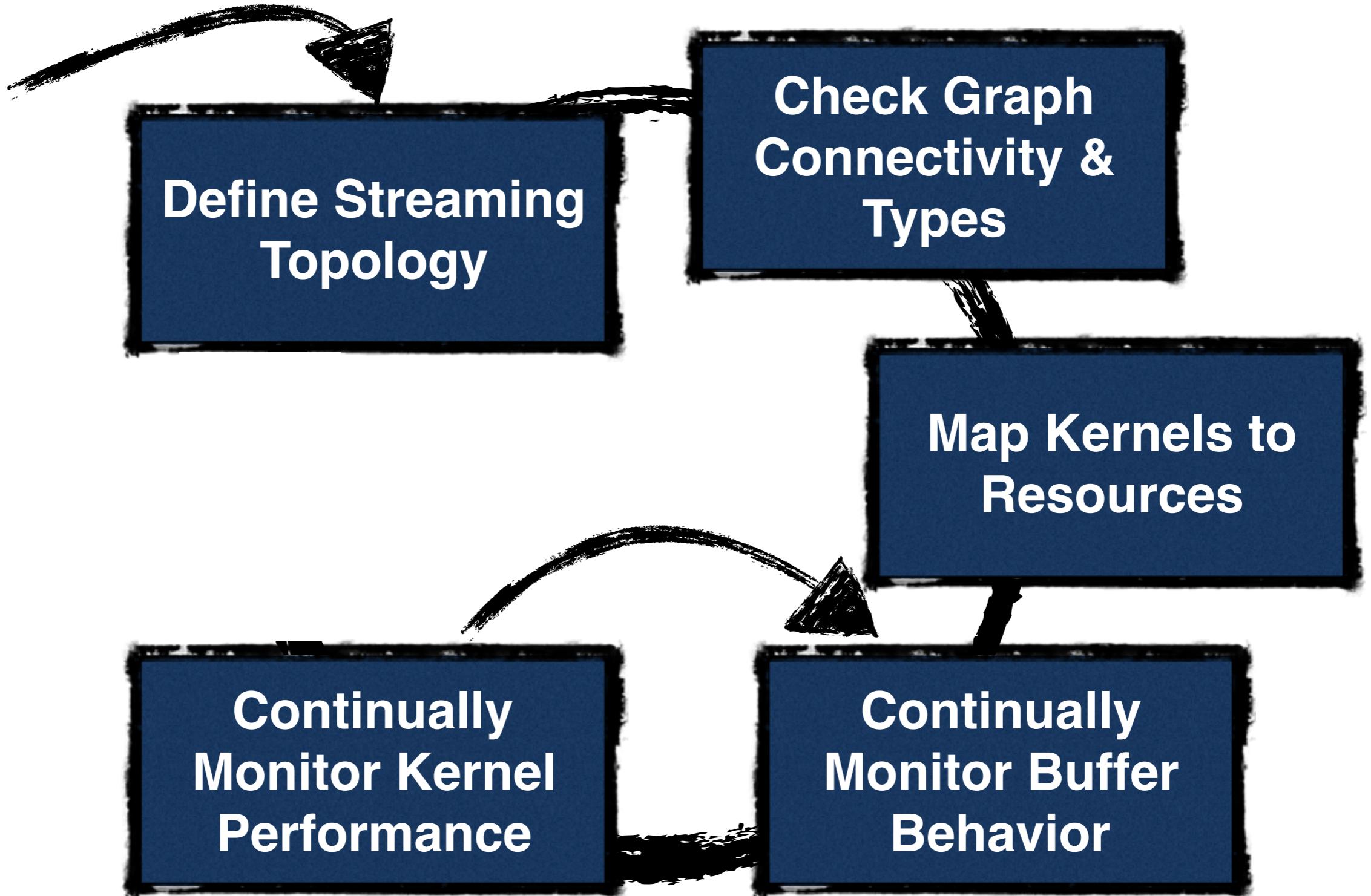
# STREAM PROCESSING (STRING SEARCH EXAMPLE)

---



# HOW IT WORKS (HIGH LEVEL)

---



# Stream Processing (Search Ex)

```
template < class T > class search : public raft::kernel
{
public:
    search( const std::string && term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        output.addPort< std::size_t >( "0" );
    }

    search( const std::string & term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        output.addPort< std::size_t >( "0" );
    }

    virtual ~search() = default;

    virtual raft::kstatus run()
    {
        auto &chunk( input[ "0" ].template peek< T >() );
        auto it( chunk.begin() );
        do
        {
            it = std::search( it, chunk.end(),
                term.begin(), term.end() );
            if( it != chunk.end() )
            {
                output[ "0" ].push( it.location() );
                it += 1;
            }
            else
            {
                break;
            }
        }
        while( true );
        input[ "0" ].unpeek();
        input[ "0" ].recycle();
        return( raft::proceed );
    }
private:
    const std::size_t term_length;
    const std::string term;
};
```

# Stream Processing (Search Ex)

```
template < class T > class search : public raft::kernel
{
public:
    search( const std::string && term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        output.addPort< std::size_t >( "0" );
    }

    search( const std::string & term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        output.addPort< std::size_t >( "0" );
    }

    virtual ~search() = default;

    virtual raft::kstatus run()
    {
        auto &chunk( input[ "0" ].template peek< T >() );
        auto it( chunk.begin() );
        do
        {
            it = std::search( it, chunk.end(),
                term.begin(), term.end() );
            if( it != chunk.end() )
            {
                output[ "0" ].push( it.location() );
                it += 1;
            }
            else
            {
                break;
            }
        }
        while( true );
        input[ "0" ].unpeek();
        input[ "0" ].recycle();
        return( raft::proceed );
    }
private:
    const std::size_t term_length;
    const std::string term;
};
```

# Stream Processing (Search Ex)

```
template < class T > class search : public raft::kernel
{
public:
    search( const std::string && term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        outUBLIC
    }
    search( const std::string && term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        output.addPort< std::size_t >( "0" );
    }
    virtual
    {
        search( const std::string &term ) : raft::kernel(),
            term_length( term.length() ),
            term( term )
        {
            input.addPort< T >( "0" );
            output.addPort< std::size_t >( "0" );
        }
        {
            break;
        }
    }
    while( true );
    input[ "0" ].unpeek();
    input[ "0" ].recycle();
    return( raft::proceed );
}
private:
    const std::size_t term_length;
    const std::string term;
};
```

# Stream Processing (Search Ex)

```
template < class T > class search : public raft::kernel
{
public:
    search( const std::string && term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        output.addPort< std::size_t >( "0" );
    }

    search( const std::string & term ) : raft::kernel(),
        term_length( term.length() ),
        term( term )
    {
        input.addPort< T >( "0" );
        output.addPort< std::size_t >( "0" );
    }

    virtual ~search() = default;

    virtual raft::kstatus run()
    {
        auto &chunk( input[ "0" ].template peek< T >() );
        auto it( chunk.begin() );
        do
        {
            it = std::search( it, chunk.end(),
                term.begin(), term.end() );
            if( it != chunk.end() )
            {
                output[ "0" ].push( it.location() );
                it += 1;
            }
            else
            {
                break;
            }
        }
        while( true );
        input[ "0" ].unpeek();
        input[ "0" ].recycle();
        return( raft::proceed );
    }

private:
    const std::size_t term_length;
    const std::string term;
};
```

# Stream Processing (Search Ex)

```
template < class T > class search : public raft::kernel
{
public:
    search( const s
    {
        input.addPor
        output.addPo
    }

    search( const s
    {
        input.addPor
        output.addPo
    }

    virtual ~search
    {
        auto &chunk(
        auto it( chunk.begin() );
        do
        {
            it = std::search( it, chunk.end(),
                            term.begin(), term.end() );
            if( it != chunk.end() )
            {
                output[ "0" ].push( it.location() );
                it += 1;
            }
            else
            {
                break;
            }
        }
        while( true );
        input[ "0" ].unpeek();
        input[ "0" ].recycle( );
        return( raft::proceed );
    }

private:
    const std::size_t
    const std::string term;
};
```

# Stream Processing (Search Ex)

```
int
main( int argc, char **argv )
{
    using chunk = raft::filechunk< 256 >;
    using fr    = raft::filereader< chunk, false >;
    using search = search< chunk >;
    using print = raft::print< std::size_t, '\n'>;

    const std::string term( "Alice" );
    raft::map m;

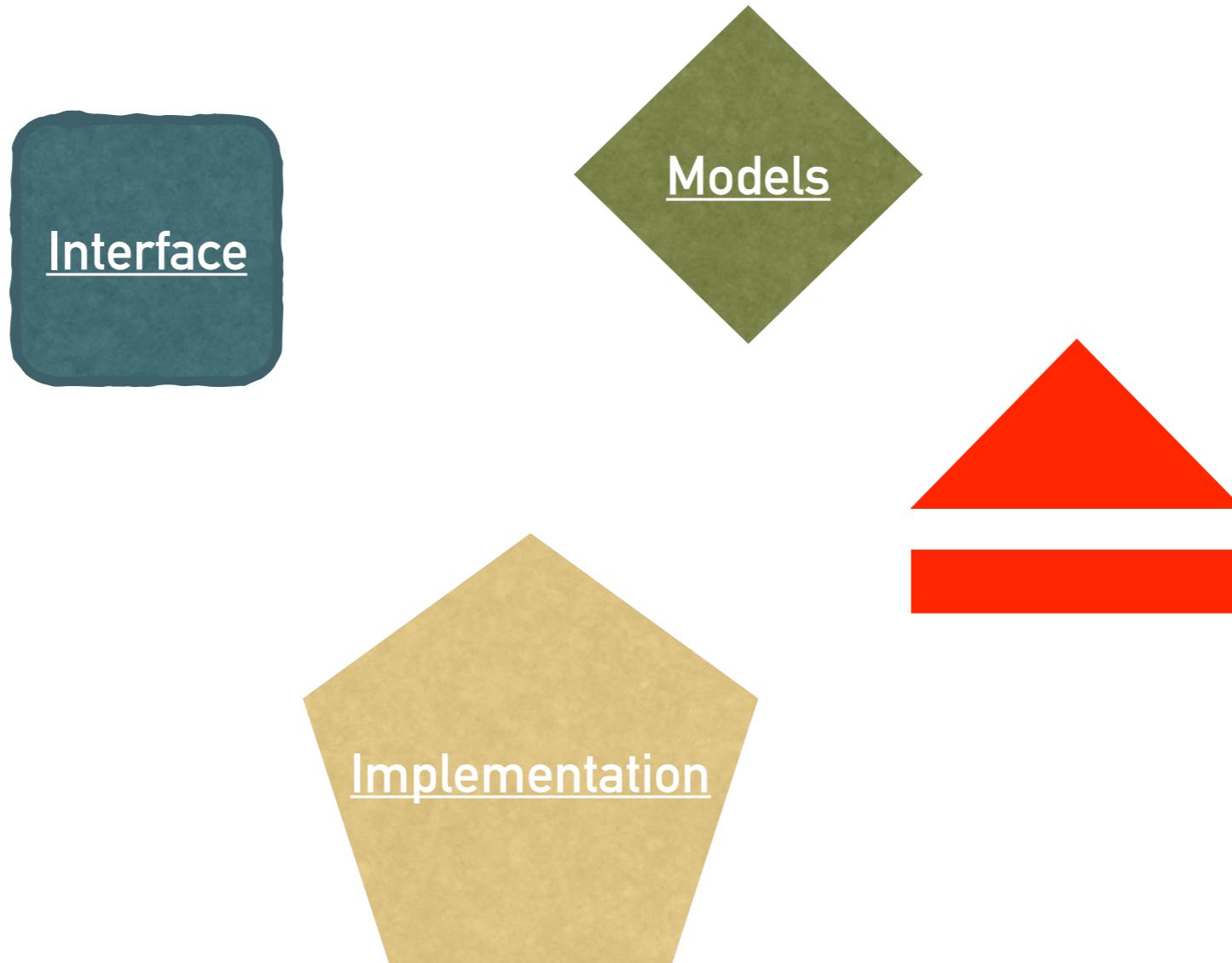
    fr.read( argv[ 1 ], 1, term.length() );
    search find( term );
    print p;

    m += read >> raft::order::out >> find >> raft::order::out >> p;
    m.exe();

    return( EXIT_SUCCESS );
}
```

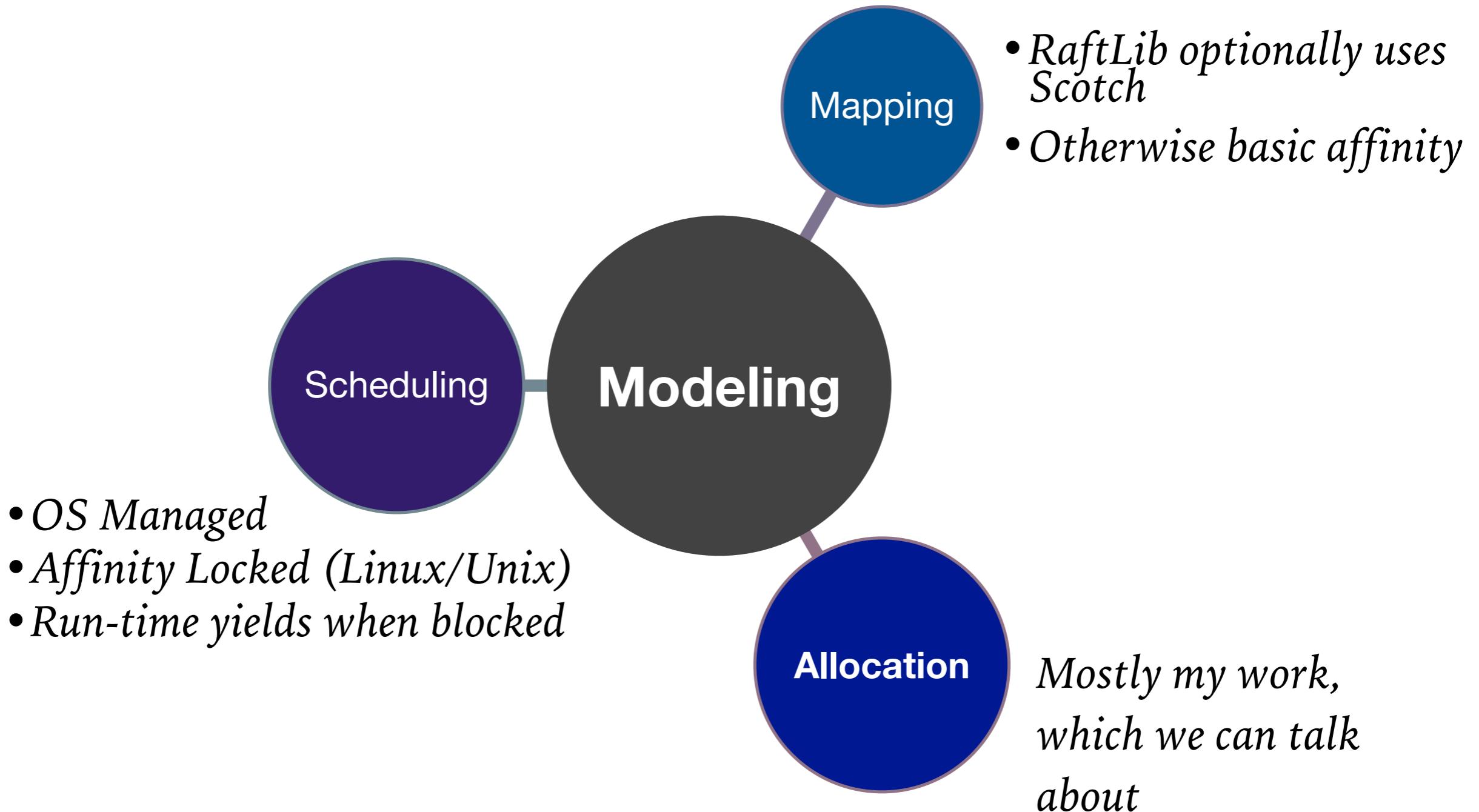
# CHOOSE YOUR ADVENTURE

---

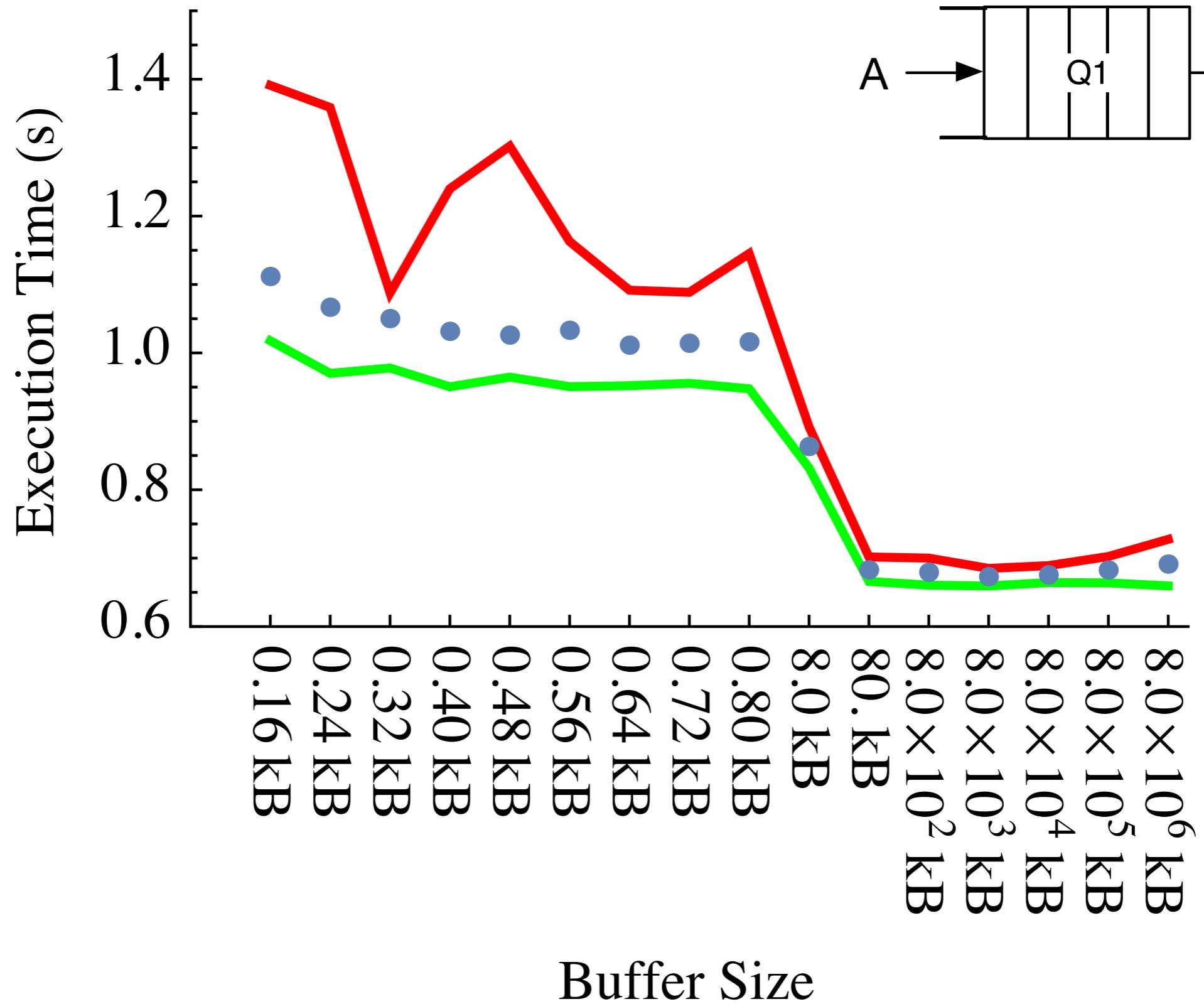


# MODELING

---



# MODELING ISSUES - SIZE OF Q1



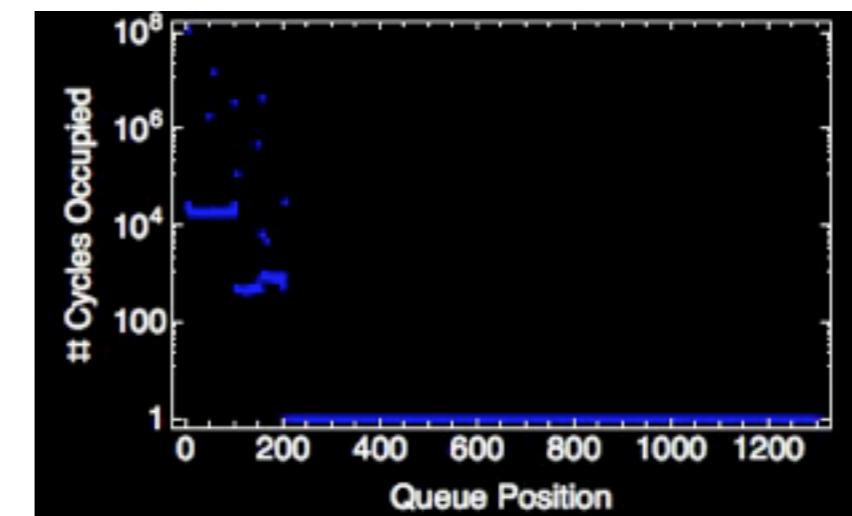
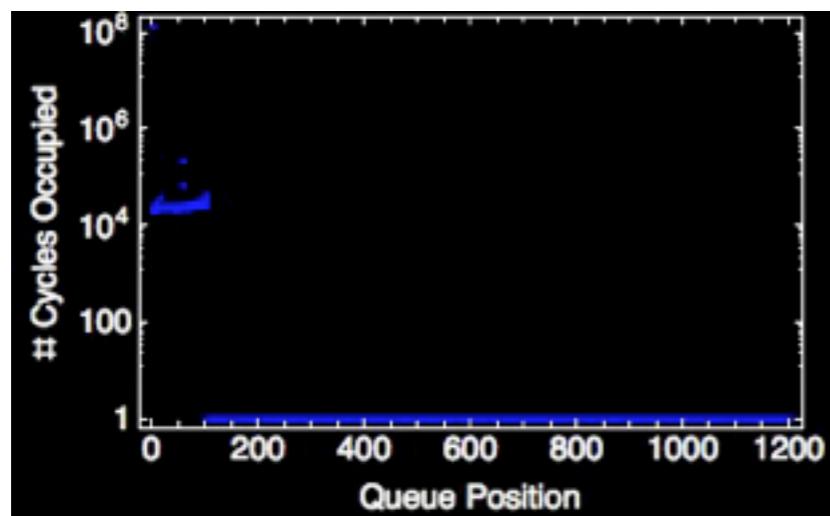
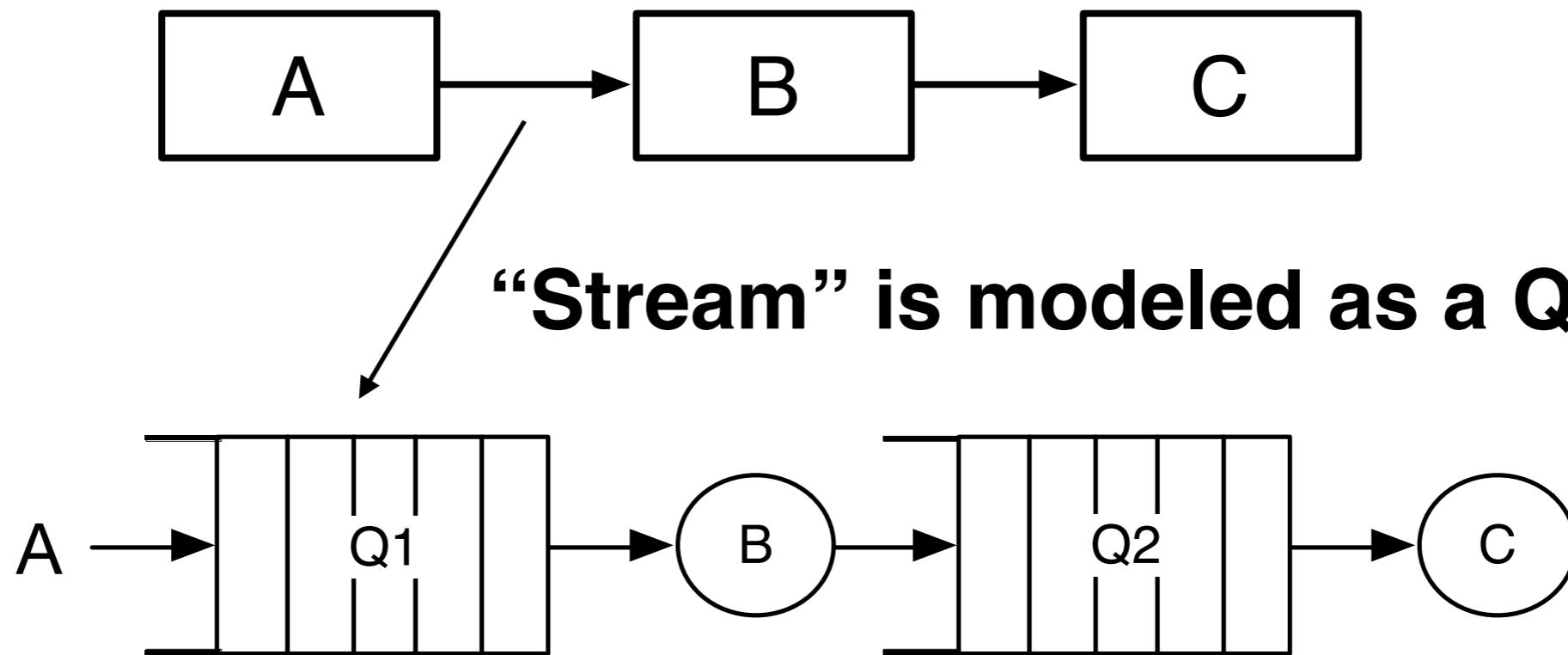
# WHY DO WE NEED TO SOLVE

---

- Buffer allocations take time and energy
- Programmers are horrible at deciding (too many parameters)
- Hardware specific locations matter (NUMA)
- Re-allocating with an accelerator takes even more time (bus latency, hand-shakes, etc.)
- Must be solved in conjunction with partitioning/scheduling/placement

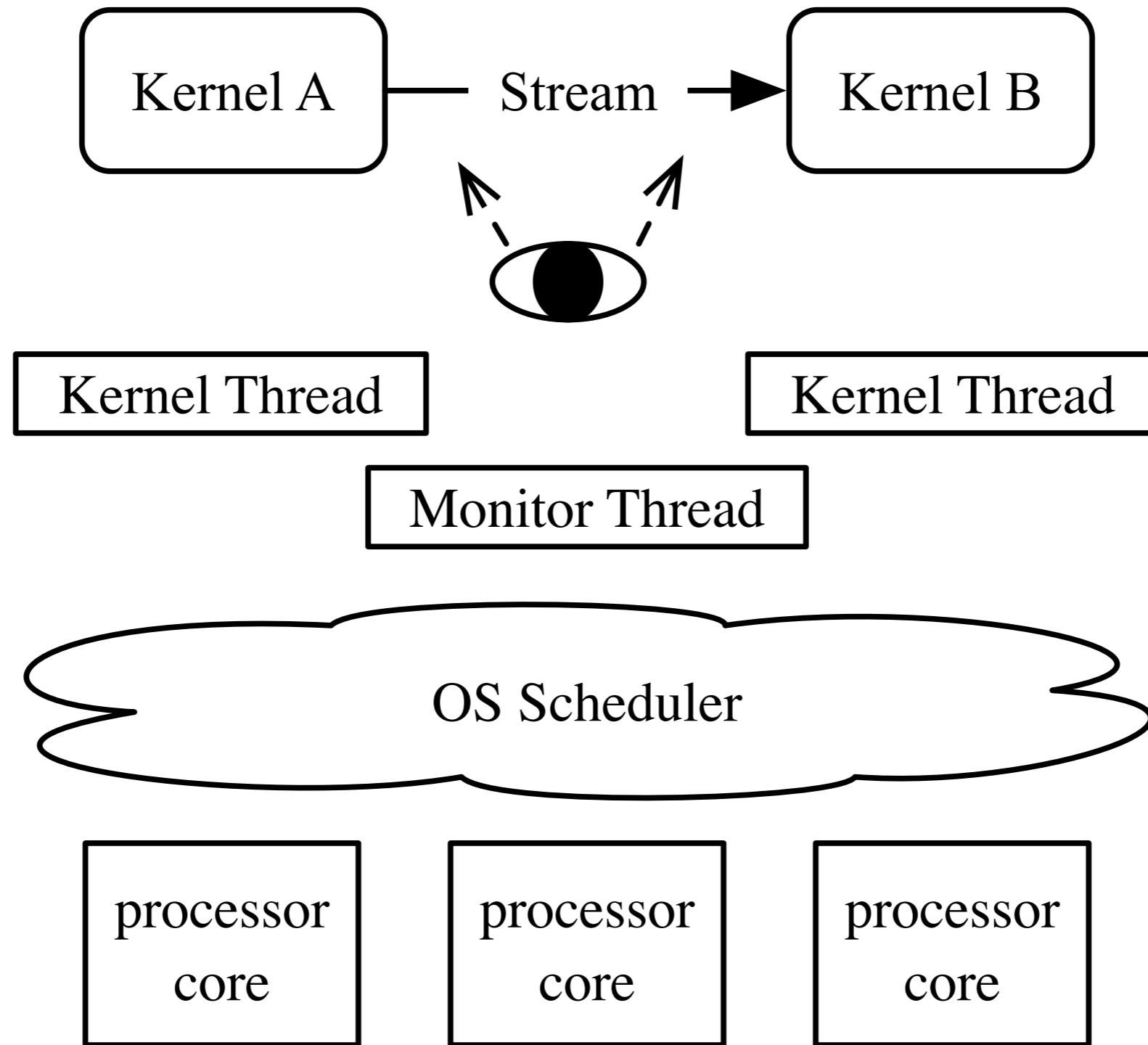
# MODELING

---



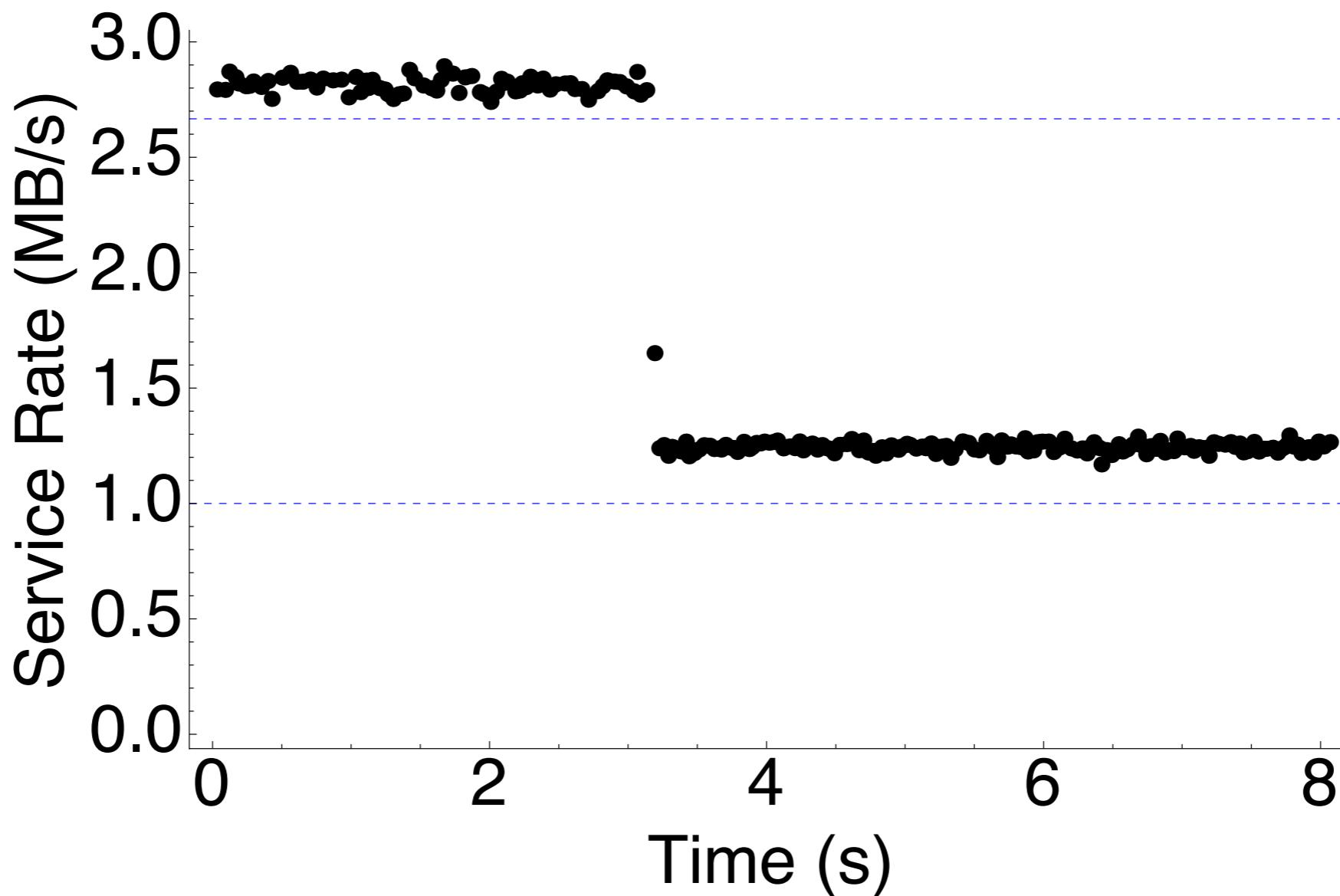
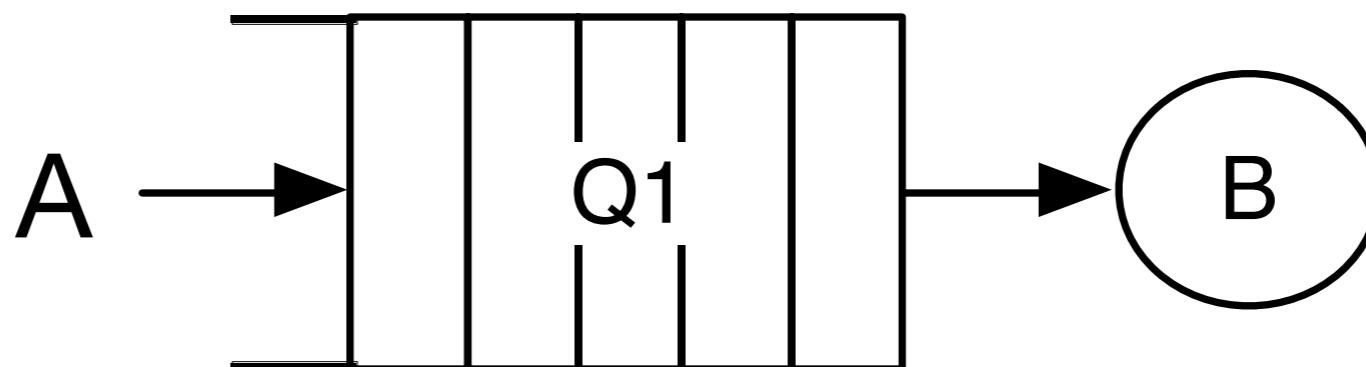
# MONITOR

---



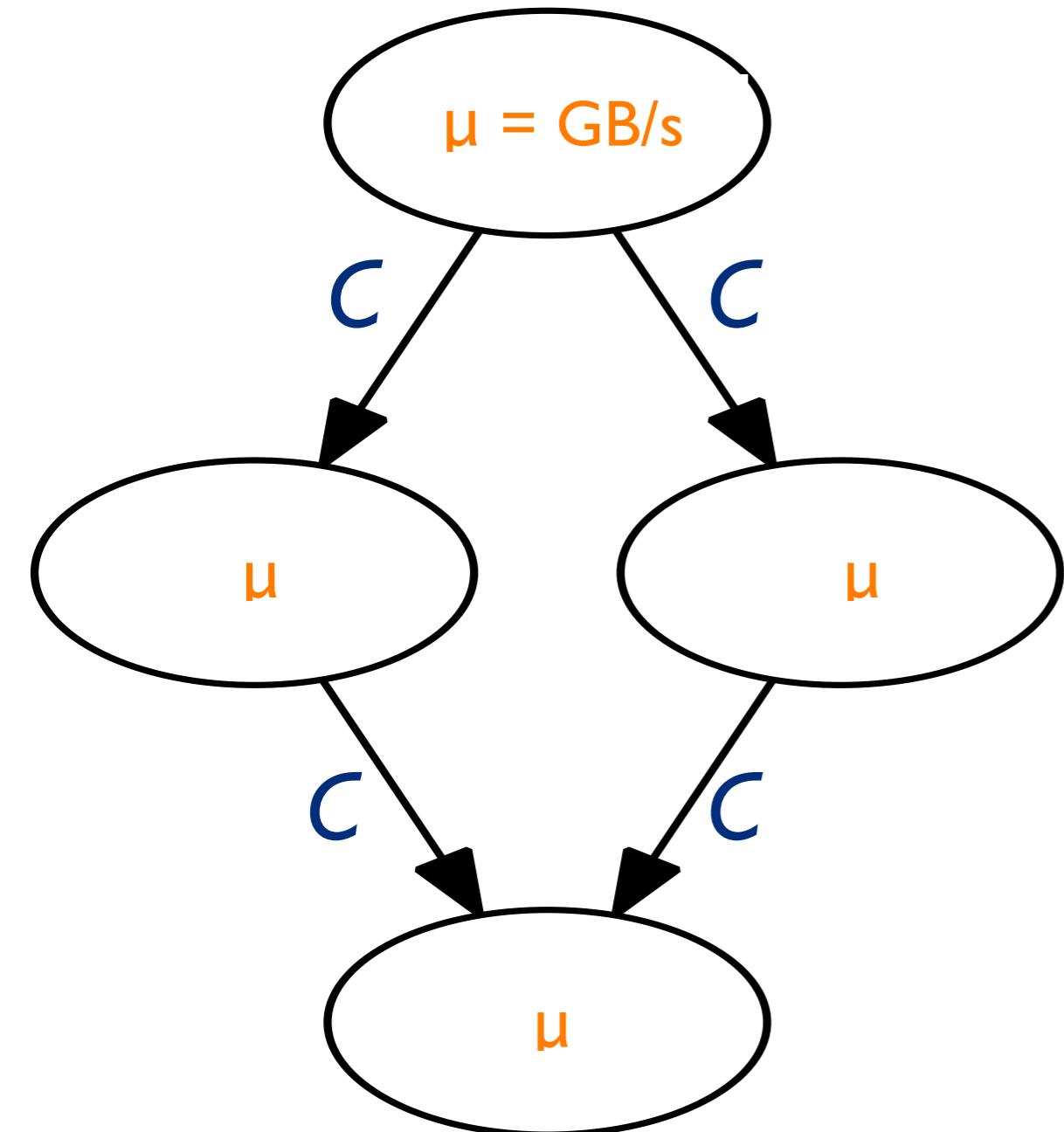
# APPROXIMATE INSTRUMENTATION

---



# SOLVE FOR THROUGHPUT QUICKLY

---



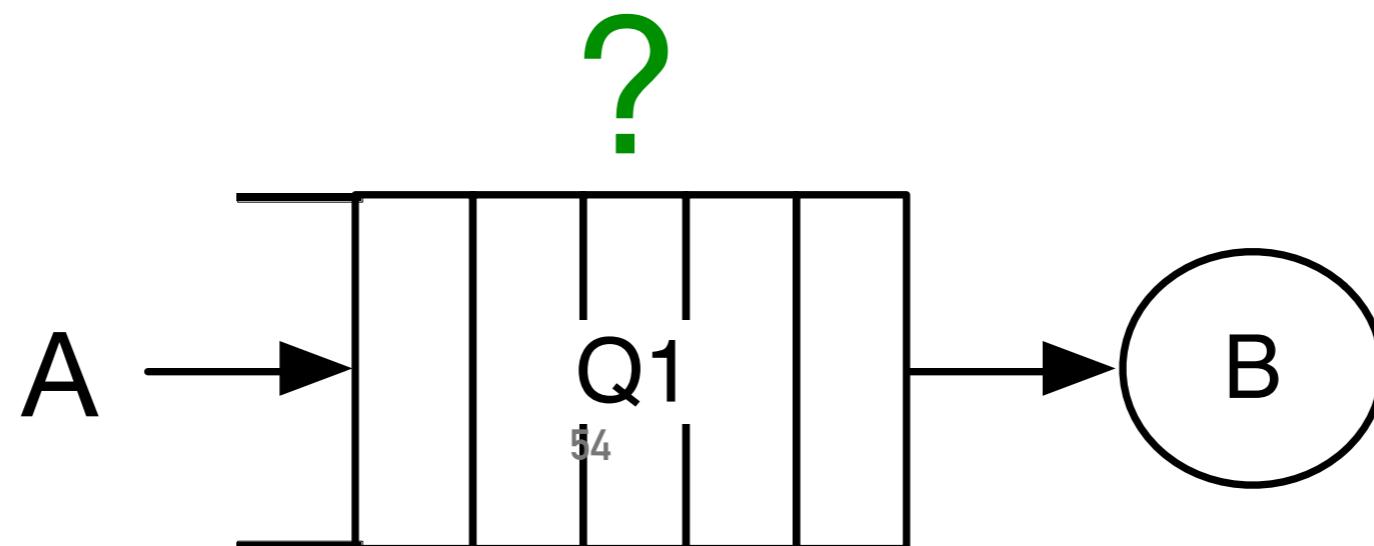
Use network flow model  
to quickly estimate flow  
within a streaming graph

Decompose queueing  
network and solve each  
queueing station  
independently

# HOW WOULD YOU GET THE RIGHT BUFFER SIZE?

---

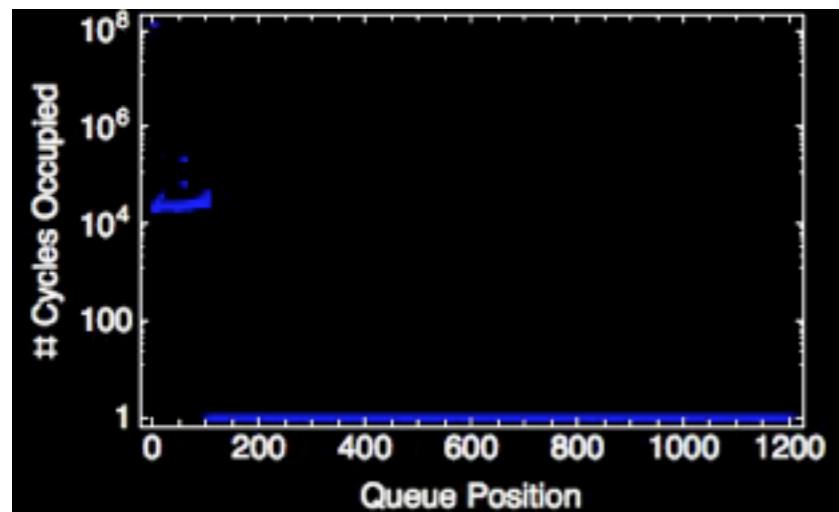
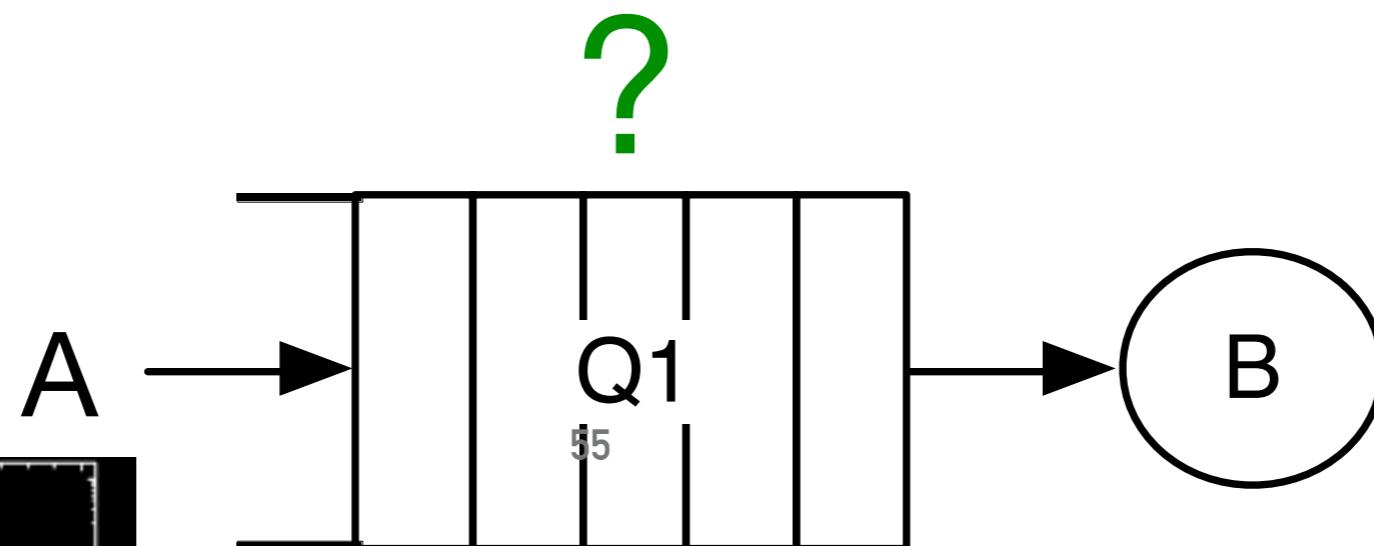
*Queueing Model to Solve for Buffer Size*



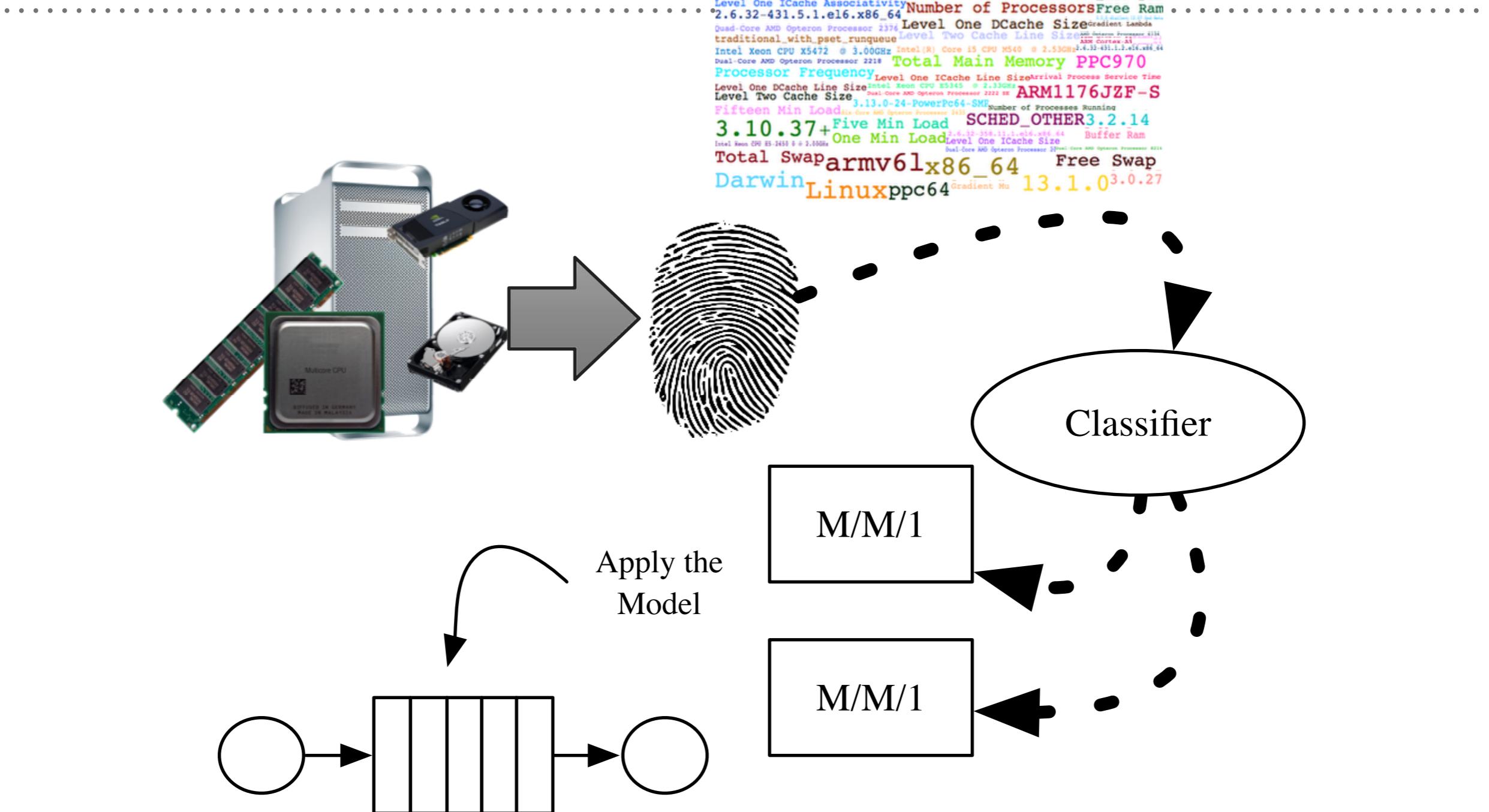
# HOW WOULD YOU GET THE RIGHT BUFFER SIZE?

---

*Queueing Model to Solve for Buffer Size*

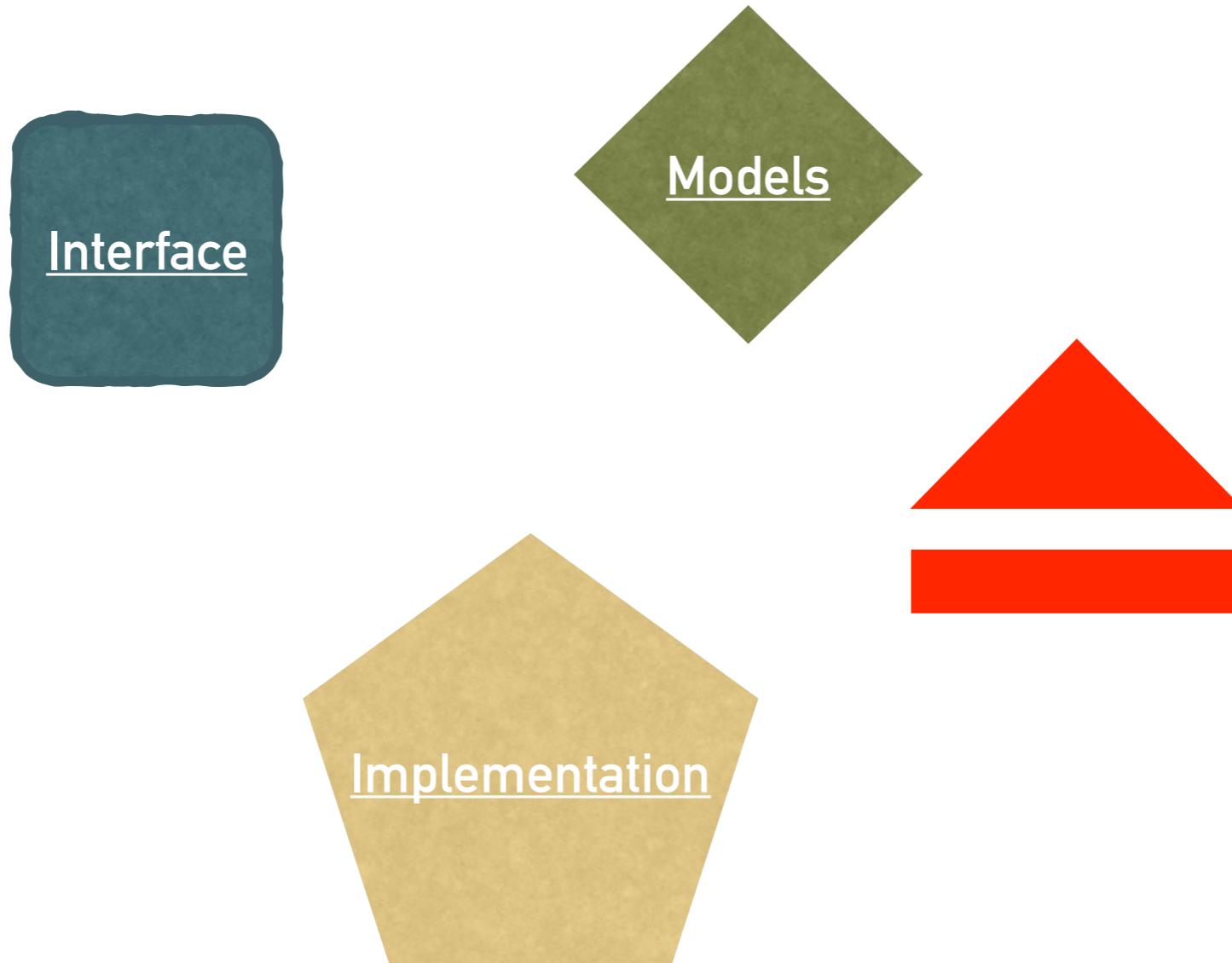


# ML BASED MODEL SELECTION



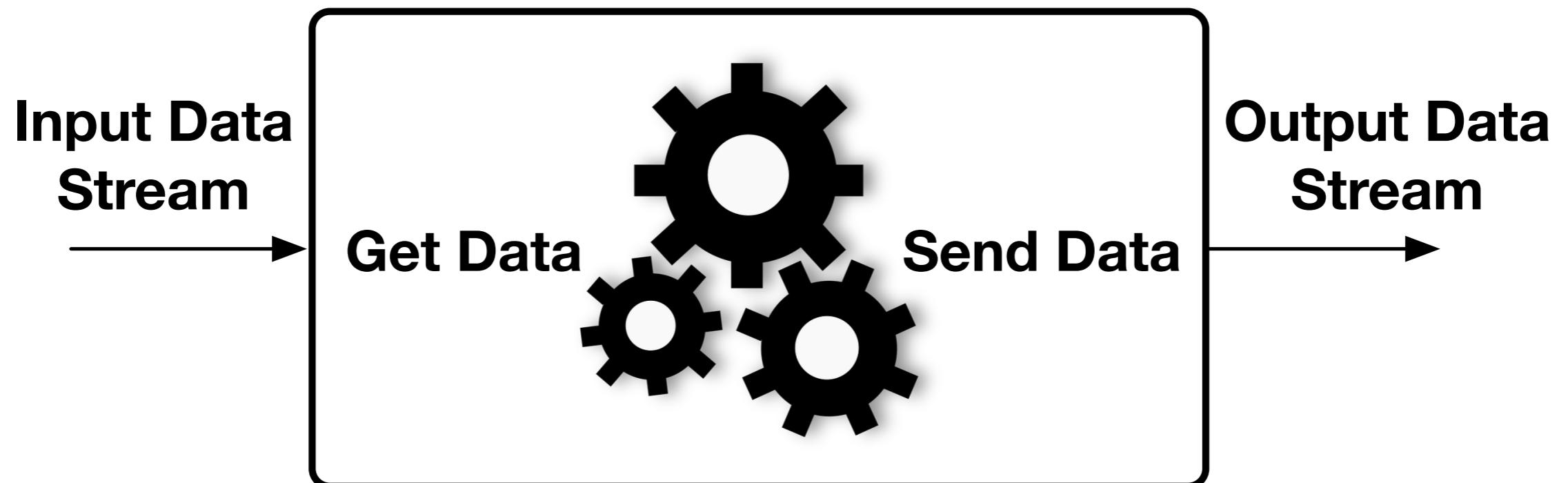
# CHOOSE YOUR ADVENTURE

---



# COMPUTE KERNEL

---



# COMPUTE KERNEL

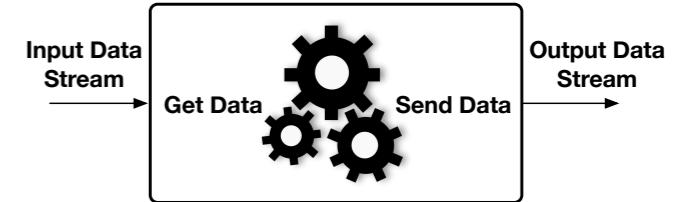
```
class akernel : public raft::kernel
{
public:
    akernel() : raft::kernel()
    {
        //add input ports
        input.addPort< /** type **>("x0", "x1", "x...");
        //add output ports
        output.addPort< /** type **>("y0", "y1", "y...");

        virtual raft::kstatus run()
        {
            /** get data from input ports */
            auto &valFromX( input[ "x..." ].peek< /** type of "x..." **>() );
            /** do something with data */

            const auto ret_val( do_something( valFromX ) );

            output[ "y..." ].push( ret_val );

            input[ "x..." ].unpeek();
            input[ "x..." ].recycle();
            return( raft::proceed /** or stop **/ );
        }
    };
}
```

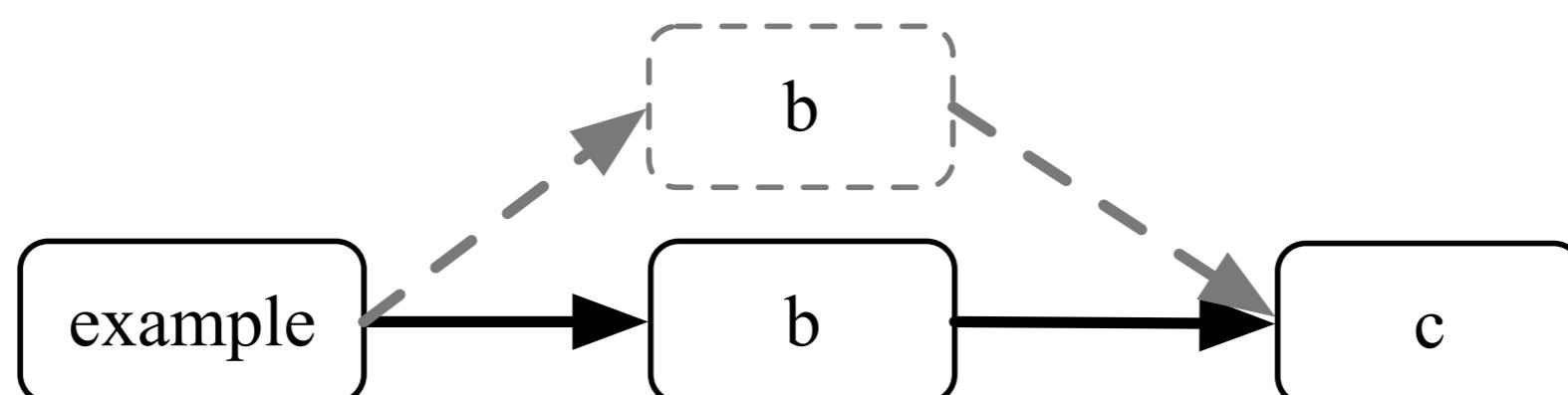
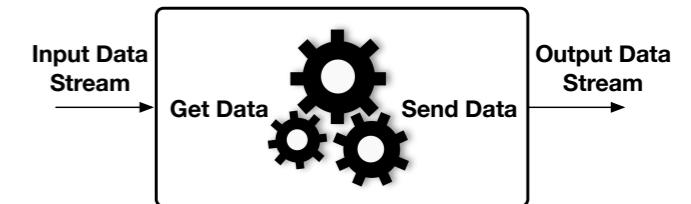


# COMPUTE KERNEL

```
class example : public raft::parallel_k
{
public:
    example() : parallel_k()
    {
        input.addPort< /* some type */ >( "0" );
        /* add a starter output port */
        addPort();
    }

    /** implement virtual function */
    virtual std::size_t addPort()
    {
        return( (this)->addPortTo< /* type */ >( output /* direction */ ) );
    }
}
```

"..."

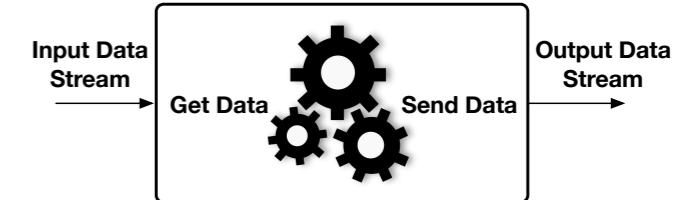


# RECEIVING DATA

```
/**  
 * return reference to memory on  
 * in-bound stream  
 */  
template< class T >  
T& peek( raft::signal *signal = nullptr )
```

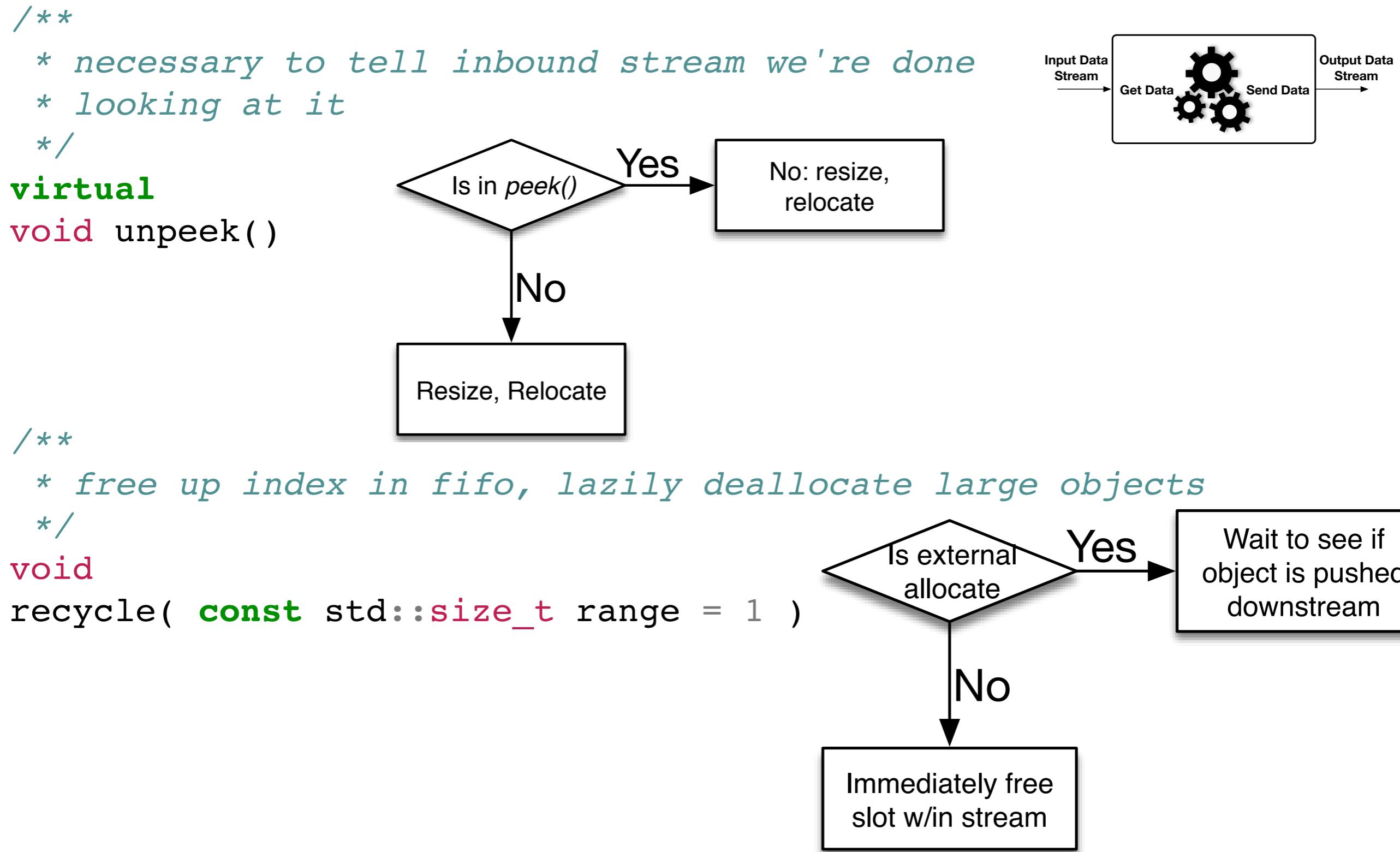
```
template< class T >  
autorelease< T, peekrange > peek_range( const std::size_t n )
```

- Returns object with “special access to stream”
- Operator [] overload returns auto pair
- Direct reference as in peek() for each element



```
template < class T > struct autopair  
{  
    T &ele;  
    Buffer::Signal &sig;  
};
```

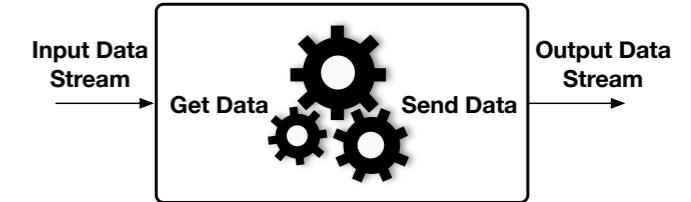
# RECEIVING DATA



# RECEIVING DATA

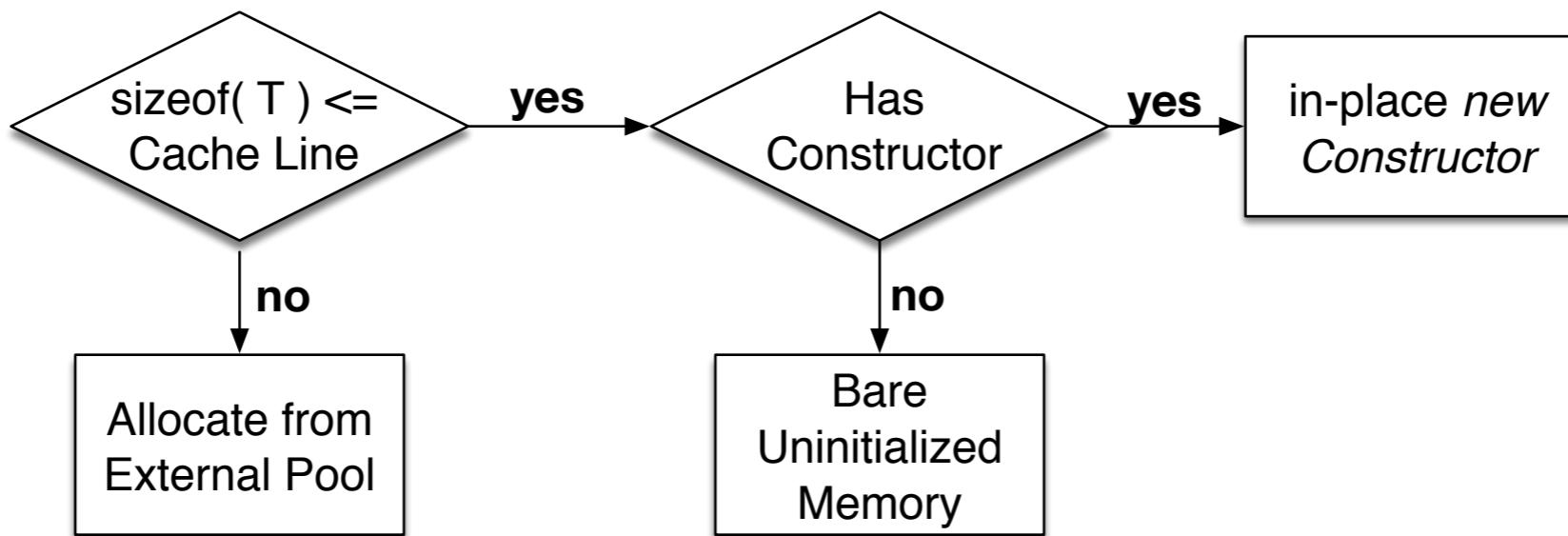
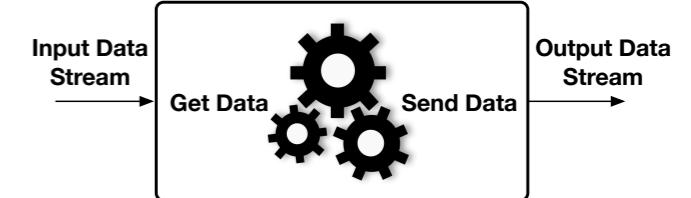
---

```
/**  
 * these pops produce a copy  
 */  
template< class T >  
void pop( T &item, raft::signal *signal = nullptr )  
  
template< class T > using pop_range_t =  
std::vector< std::pair< T , raft::signal > >;  
  
template< class T >  
void pop_range( pop_range_t< T > &items,  
                const std::size_t n_items )  
  
/**  
 * no copy, slightly higher overhead, "smart object"  
 * implements peek, unpeek, recycle  
 */  
template< class T >  
autorelease< T, poptype > pop_s()
```



# SENDING DATA

```
/** in-place allocation */
template < class T,
           class ... Args >
T& allocate( Args&&... params )
```

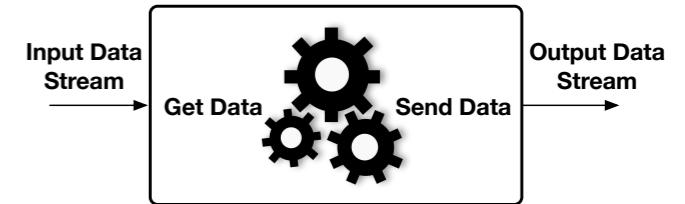


```
/** in-place alloc of range for fundamental types */
template < class T >
auto allocate_range( const std::size_t n ) ->
    std::vector< std::reference_wrapper< T > >
```

# SENDING DATA

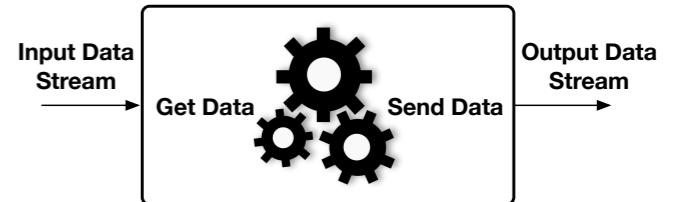
---

```
/** release data to stream **/  
virtual  
void send( const raft::signal = raft::none )  
  
/** release data to stream **/  
virtual  
void send_range( const raft::signal = raft::none )  
  
/** oops, don't need this memory **/  
virtual void deallocate()
```



# SENDING DATA

---



```
/** multiple forms **/
template < class T >
void push( const T &item, const raft::signal signal = raft::none )

/** insert from container within run() function to stream */
template< class iterator_type >
void insert( iterator_type begin,
            iterator_type end,
            const raft::signal signal = raft::none )
```

# INCLUDED KERNELS

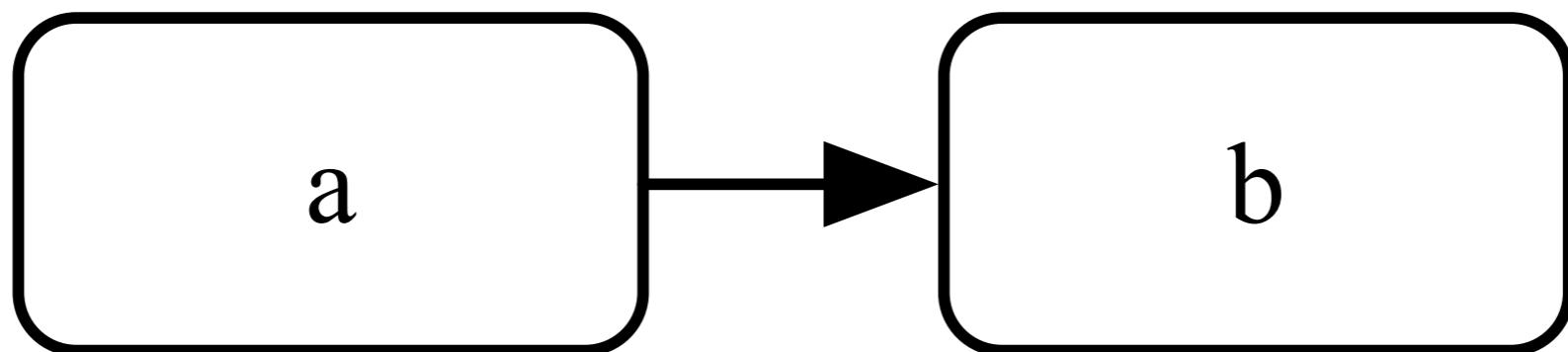
---

```
/**  
 * thread safe print, specialization for '\n' vs. '\0'  
 */  
template< typename T, char delim = '\0' > class print  
  
/** read from iterator to streams **/  
static  
raft::readeach< T, Iterator >  
read_each( Iterator &&begin,  
           Iterator &&end )  
  
/** write from iterator to streams **/  
template < class T, class BackInsert >  
static  
writeeach< T, BackInsert >  
write_each( BackInsert &&bi )  
{  
    return( writeeach< T, BackInsert >( bi ) );  
}
```

# CONNECTING COMPUTE KERNELS

---

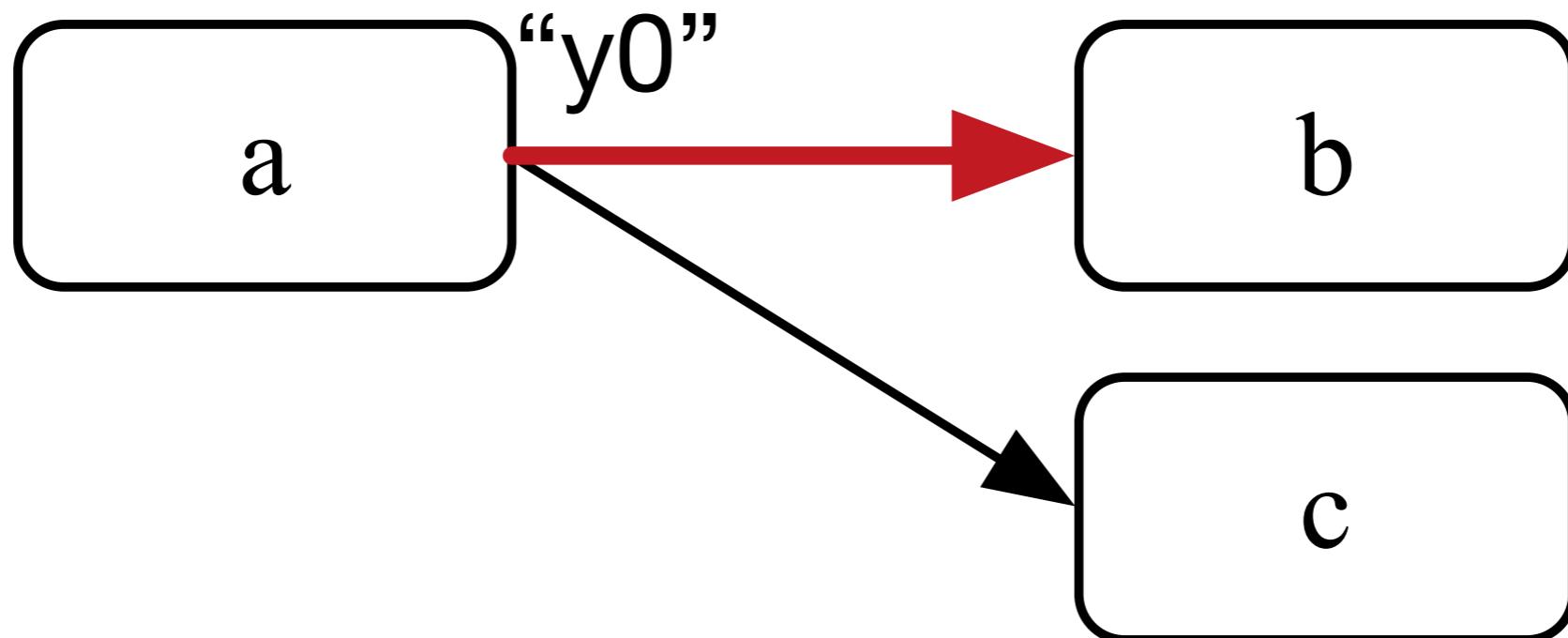
```
raft::map m;  
/* example only */  
raft::kernel a, b;  
m += a >> b;
```



# CONNECTING COMPUTE KERNELS

---

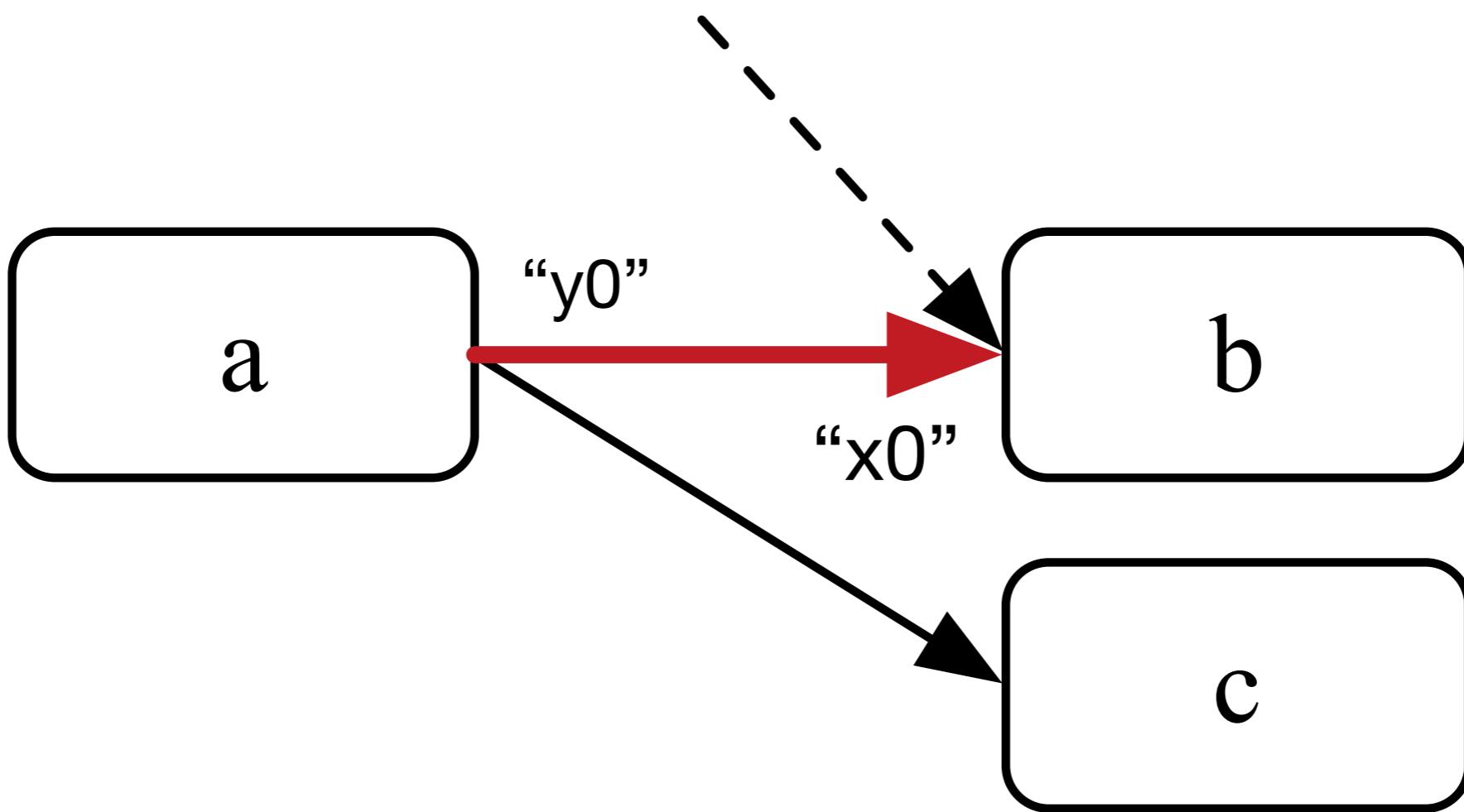
```
raft::map m;  
/** example only **/  
raft::kernel a, b;  
m += a[ "y0" ] >> b;
```



# CONNECTING COMPUTE KERNELS

---

```
raft::map m;  
/** example only */  
raft::kernel a, b;  
m += a[ "y0" ] >> b[ "x0" ];
```

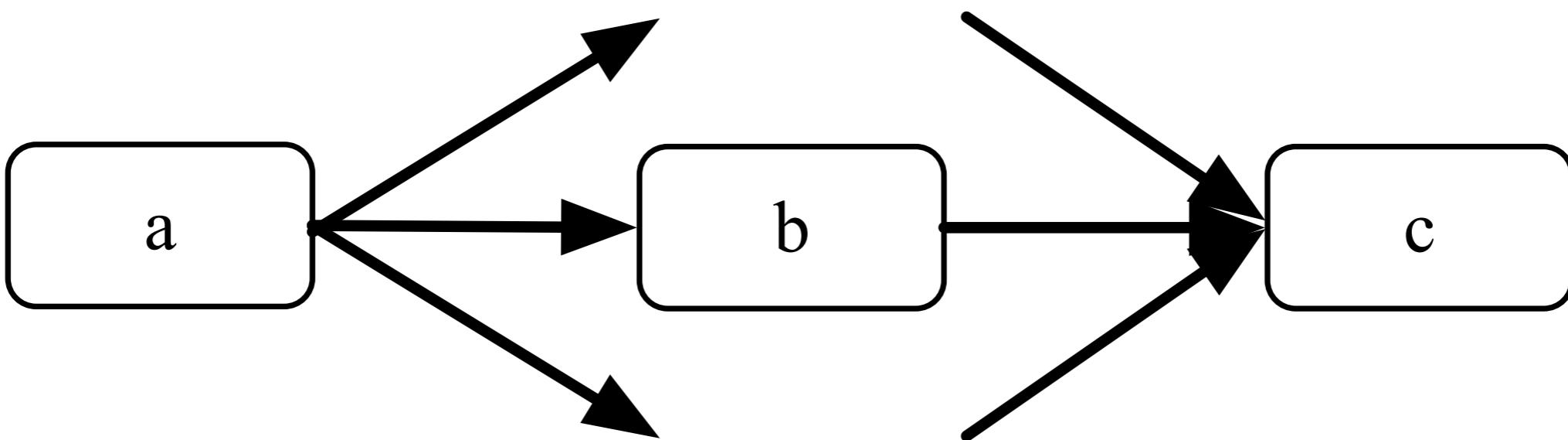


# CONNECTING COMPUTE KERNELS

---

```
raft::map m;  
/** example only **/  
raft::kernel a, b, c;  
m += a <= b >= c;
```

*Topology user specifies*

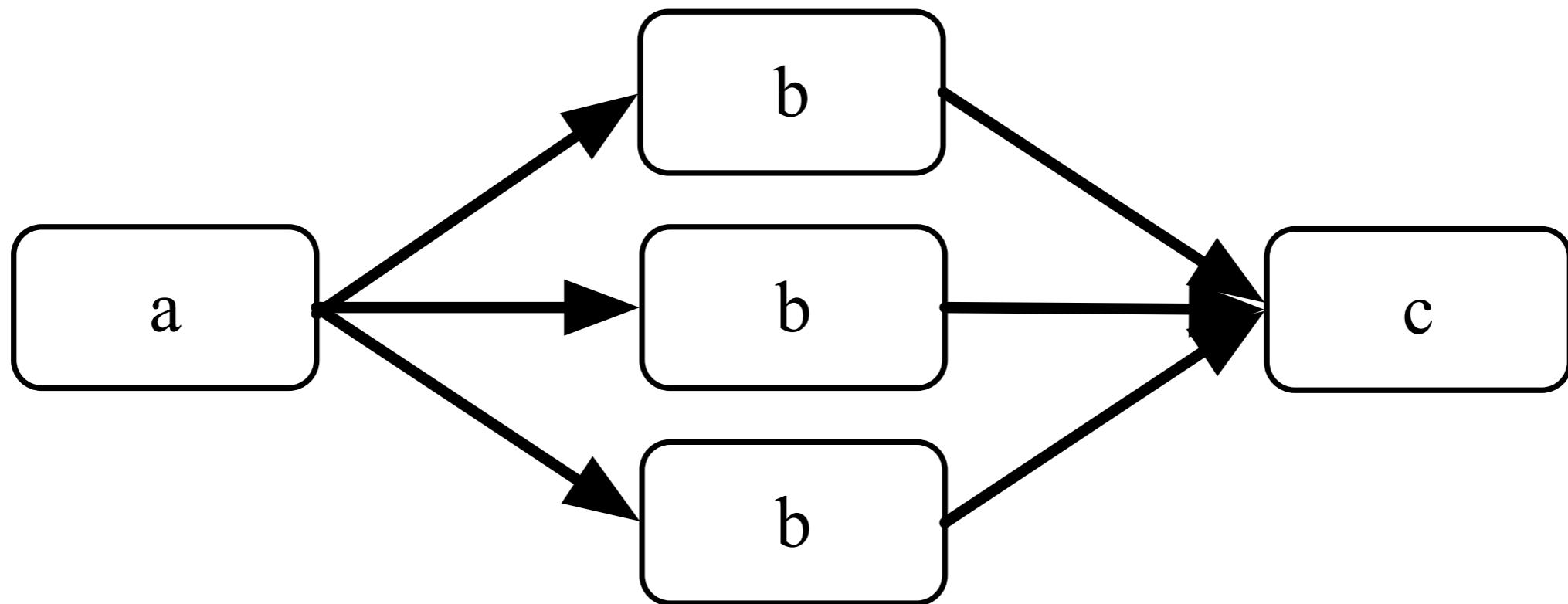


# CONNECTING COMPUTE KERNELS

---

```
raft::map m;  
/** example only **/  
raft::kernel a, b, c;  
m += a <= b >= c;
```

*RaftLib Turns Into*

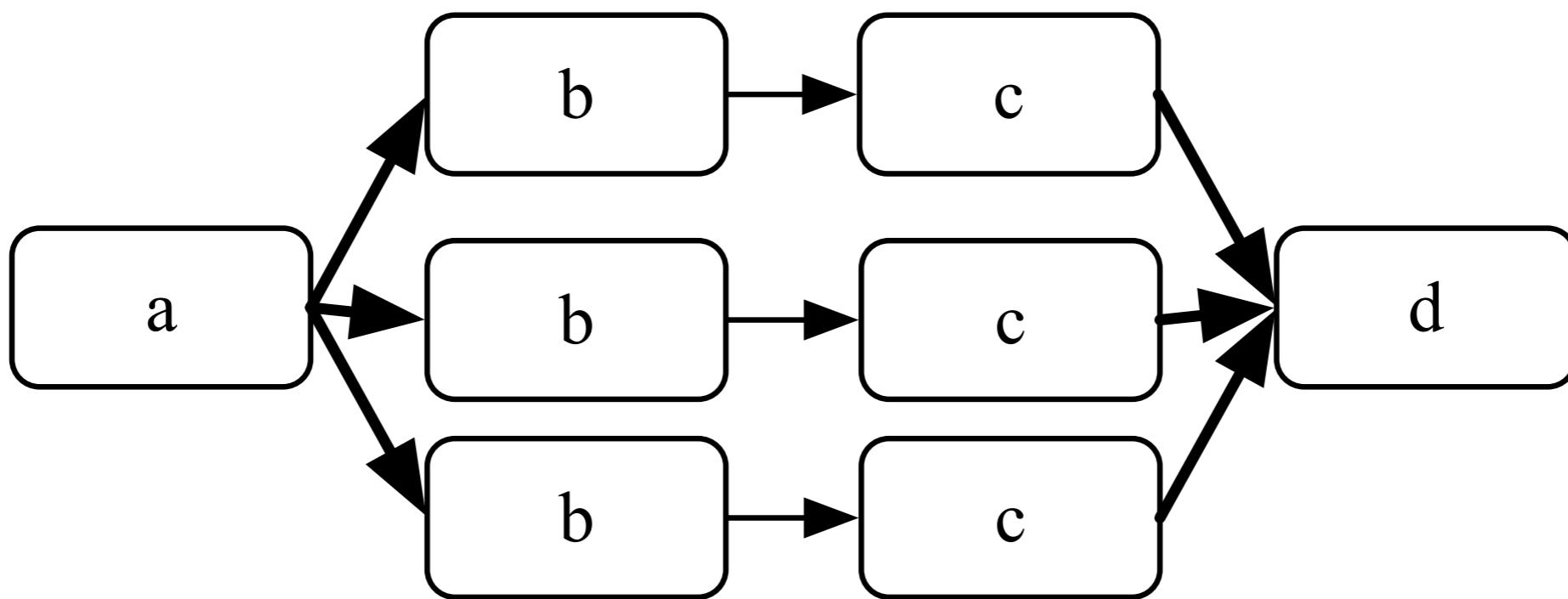


# CONNECTING COMPUTE KERNELS

---

```
raft::map m;  
/** example only **/  
raft::kernel a, b, c, d;  
m += a <= b >> c >= d;
```

*RaftLib Turns Into*

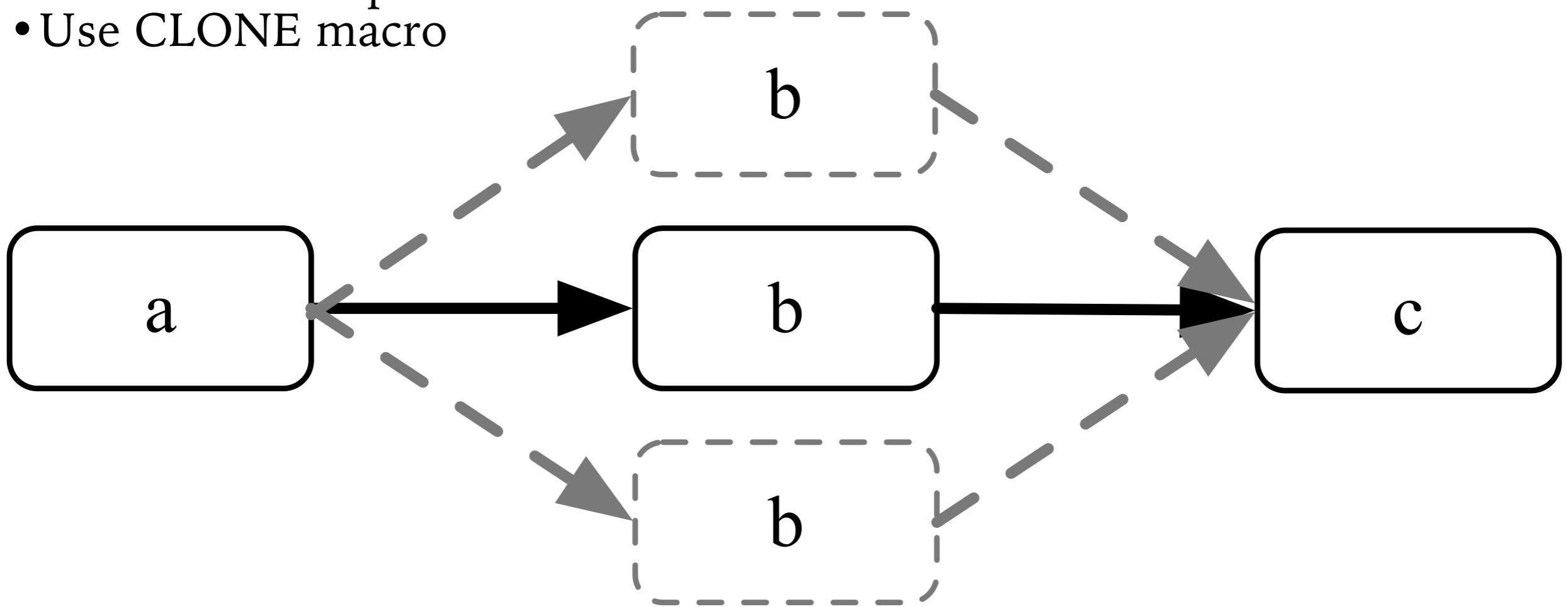


# CONNECTING COMPUTE KERNELS

---

```
raft::map m;  
/* example only */  
raft::kernel a, b, c;  
m += a >> raft::order::out >> b >> raft::order::out >> c;
```

- Clone at **SESE** points
- Use CLONE macro

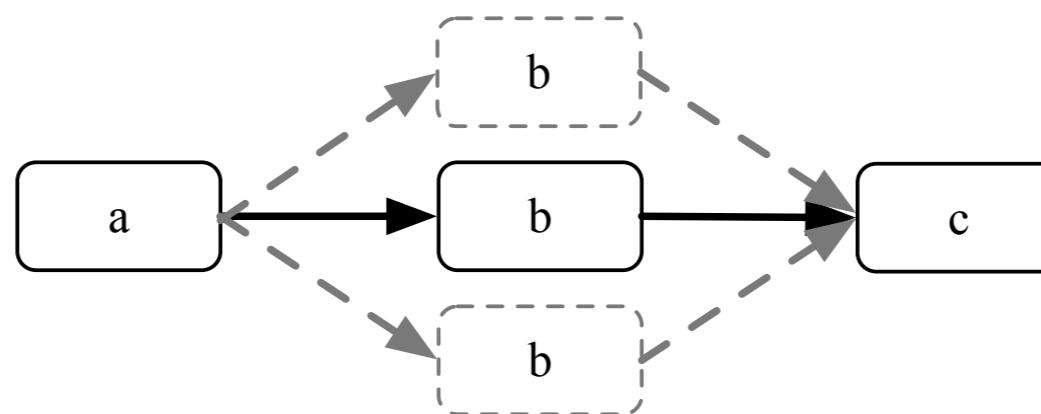


# AUTO-PARALLELIZATION

---

- 1) RaftLib figures out methodology
- 2) Runtime calls CLONE() macro of kernel “b”

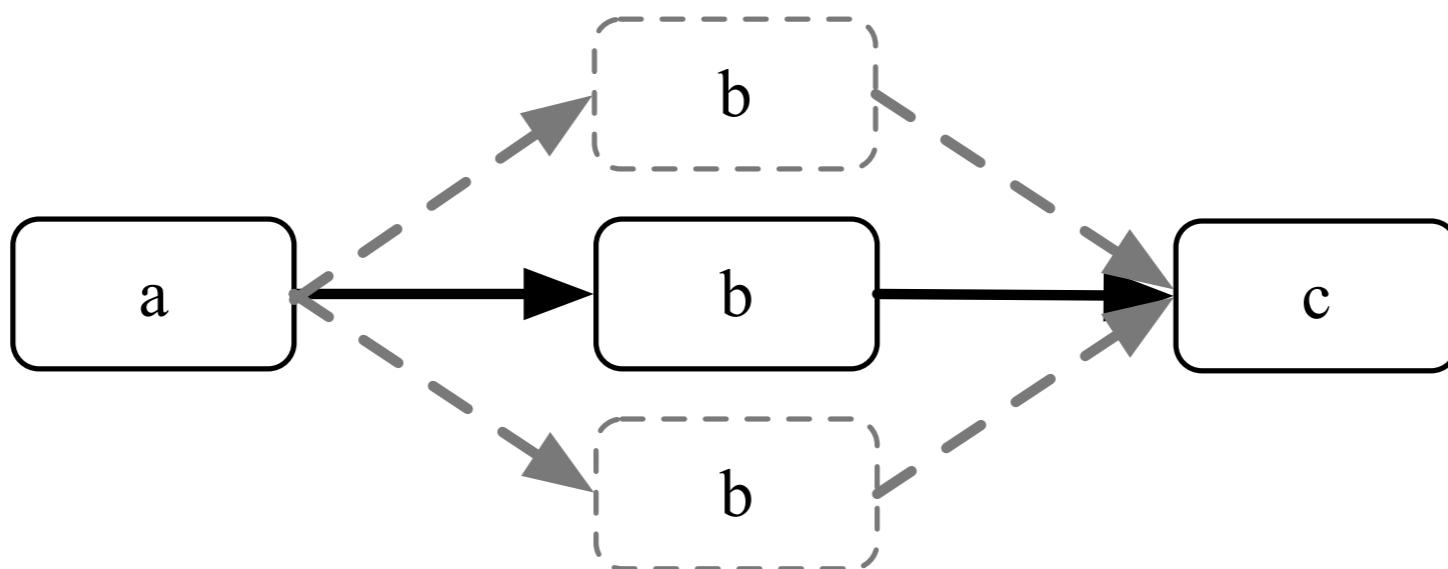
```
#define CLONE()\nvirtual raft::kernel* clone()\n{\n    auto *ptr( \n        new typename std::remove_reference< decltype( *this ) >::type( ( *( \n( typename std::decay< decltype( *this ) >::type * ) \n    this ) ) ) );\n    /** RL needs to dealloc this one **/\n    ptr->internal_alloc = true;\n    return( ptr );\n}
```



# AUTO-PARALLELIZATION

---

- 3) Lock output port container of “a”
- 4) Register new port
- 5) Decide where to run it
- 6) Allocate memory for stream
- 7) Unlock output container of kernel “a”  
//do same on output side of “b”



# CONNECTION TODO ITEMS

---

- Better SESE implementation
- Decide on syntax for set of kernels

```
raft::map m;  
/** example only **/  
raft::kernel a, b, c, d;  
m += src <= raft::kset( a, b, c, d ) >= dst;
```

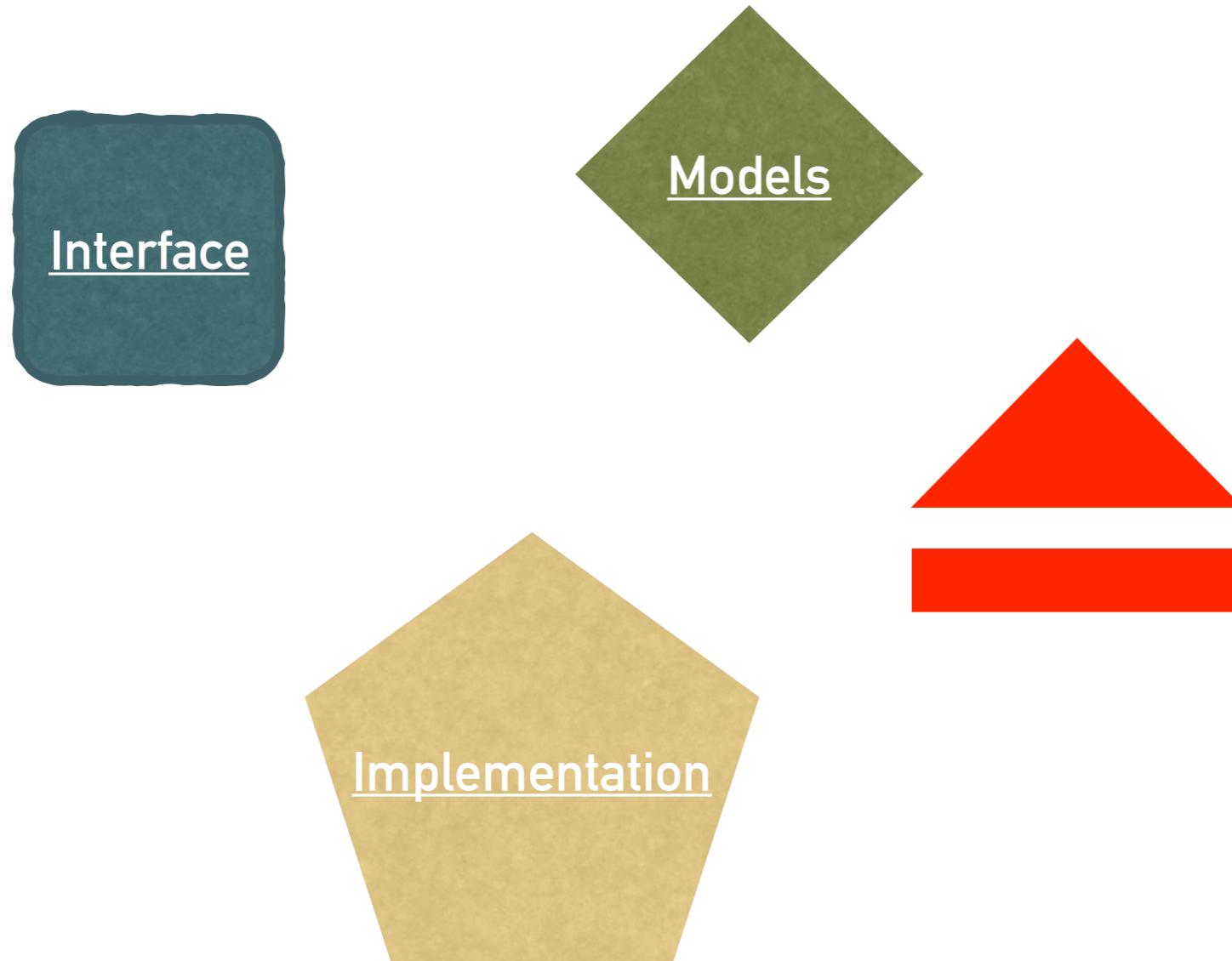
- Address space stream modifier (new VM space)

```
raft::map m;  
/** example only **/  
raft::kernel a, b, c;  
m += a >> raft::vm::part >> b >> raft::vm::part >> c;
```

- Anything else?

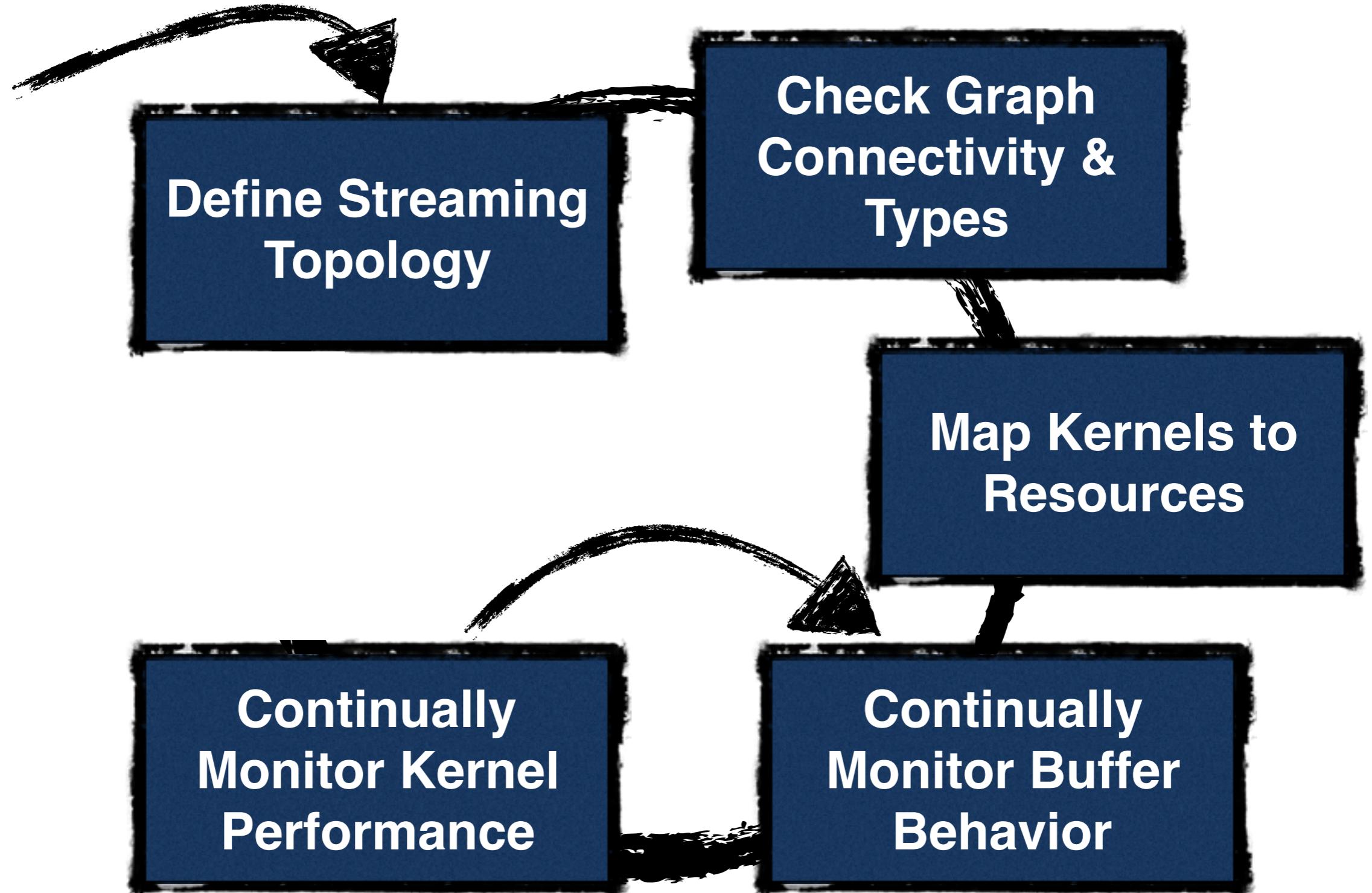
# CHOOSE YOUR ADVENTURE

---



# IMPLEMENTATION DETAILS

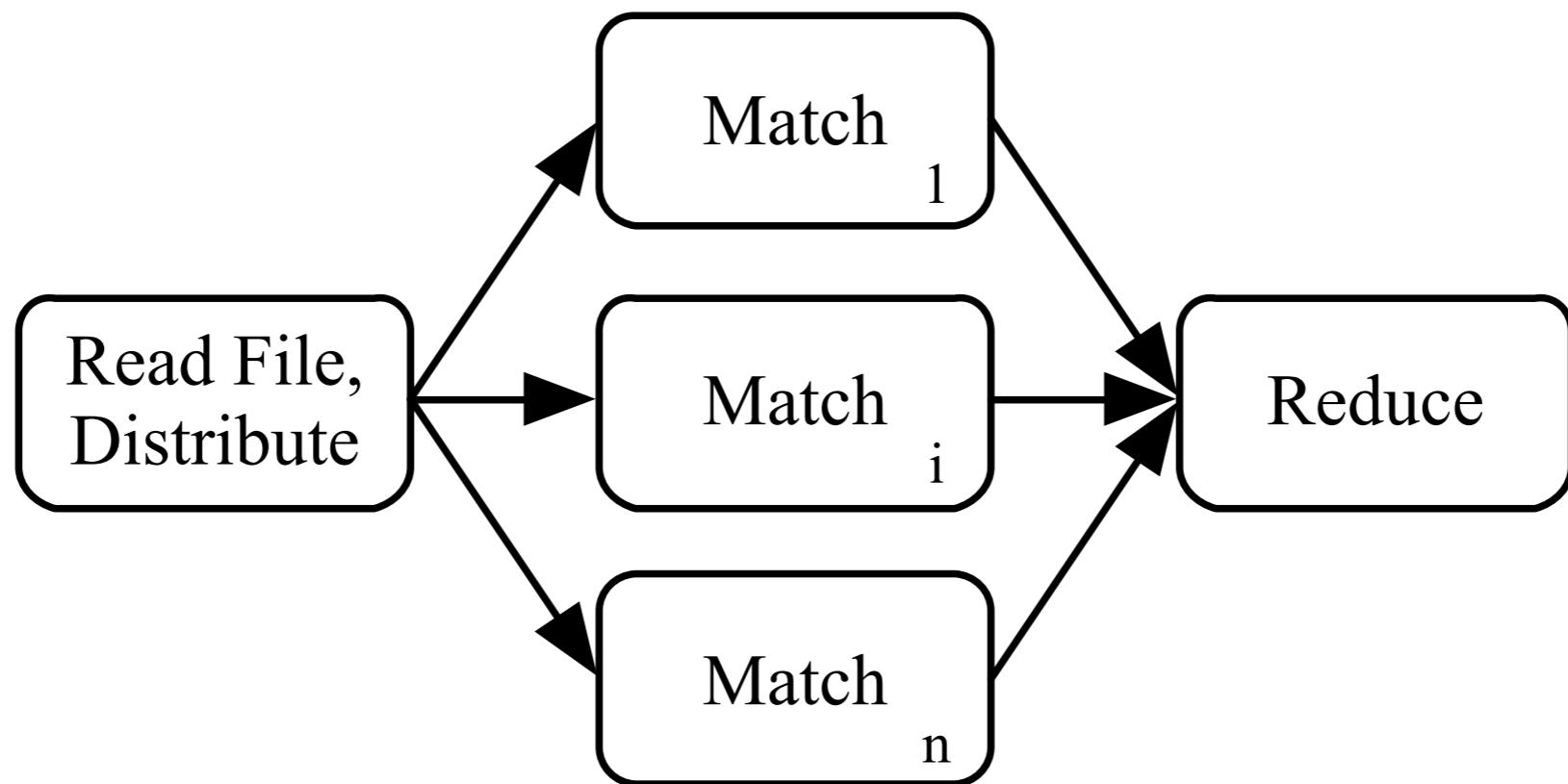
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# TOPOLOGY CHECK

---

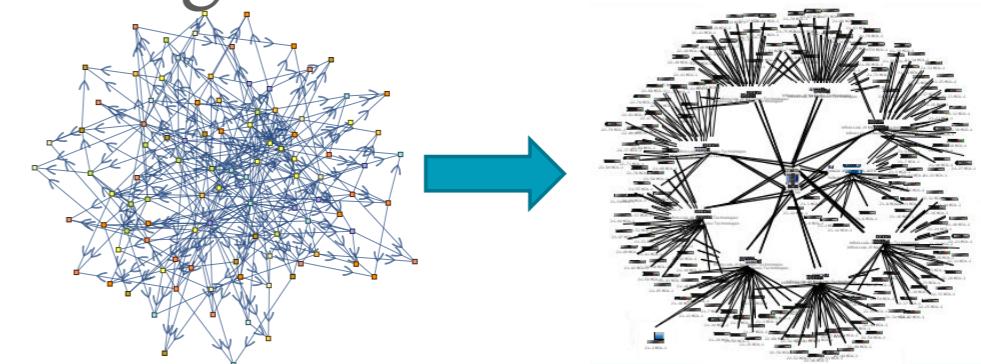
- Add kernels to map
- Check type of each link (potential for type optimization)
- Handle static split/joins (produce any new kernels)
- DFS to ensure no unconnected edges



# PARTITION (LINUX / UNIX)

---

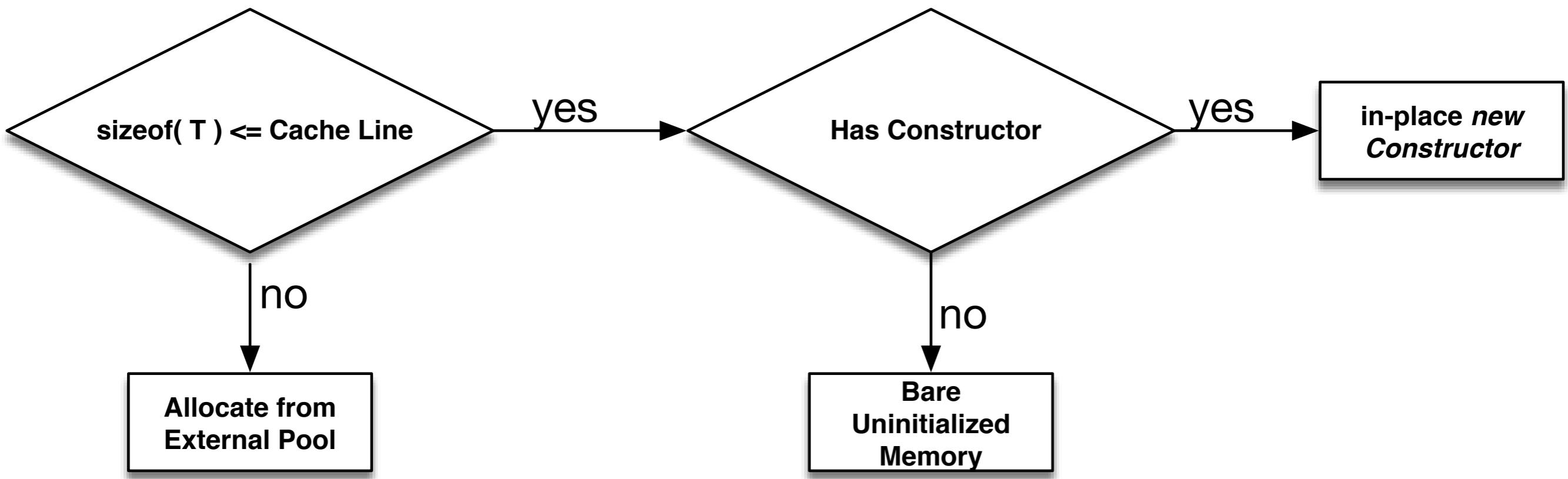
- Take RaftLib representation, convert to Scotch format
- Use fixed communications cost at each edge for initial partition
  - 2 main dimensions:
    - Flow between each edge in the application graph
    - Bandwidth available between compute resource
- Set affinity to partitioned compute cores
- Repartition as needed @ run-time
- TODO: incorporate OpenMPI utilities *hwloc* and *netloc* to get cross-platform hardware topology information



# CHOOSE ALLOCATIONS

---

- Alignment
  - SIMD ops often require memory alignment, RaftLib takes an alignment by default approach for in-stream allocations
- In-stream vs. External Pool Allocate (template allocators)



# MONITOR BEHAVIOR - ALLOCATORS

---

- Two options for figuring out optimal buffer size while running
  - model based (*discussed on modeling adventure path*)
  - branch & bound search
- separate thread, exits when app done
- pseudocode:

```
while( not done )  
{  
    if( queue_utilization > .5 )  
    {  
        queue->resize();  
    }  
    sleep( ALLOC_INTERVAL );  
}
```

# RUN-TIME LOCK FREE FIFO RESIZING

```
.....  
enum access_key : key_t { allocate      = 0,  
                           allocate_range = 1,  
                           push          = 3,  
                           recycle        = 4,  
                           pop           = 5,  
                           peek          = 6,  
                           size          = 7,  
                           N };  
}  
struct ThreadAccess  
{  
    union  
    {  
        std::uint64_t whole = 0; /* just in case, default zero */  
        dm::key_t     flag[ 8 ];  
    };  
    std::uint8_t padding[ L1D_CACHE_LINE_SIZE - 8 /* padding */ ];  
}  
#if defined __APPLE__ || defined __linux  
    __attribute__((aligned( L1D_CACHE_LINE_SIZE )))  
#endif  
    volatile thread_access[ 2 ];  
  
    /* std::memory_order_relaxed */  
    std::atomic< std::uint64_t > checking_size = { 0 };
```

# RUN-TIME LOCK FREE FIFO RESIZING

---

- Optimization...wait for the right conditions

```
if( rpt < wpt )
{
    //perfect to copy w/std::memcpy
}
```

- Factory allocators

```
/** allocator factory map */
std::map< Type::RingBufferType , instr_map_t* > const_map;

/** initialize some factories */
const_map.insert( std::make_pair( Type::Heap , new instr_map_t() ) );

const_map[ Type::Heap ]->insert(
    std::make_pair( false /* no instrumentation */ ,
                    RingBuffer< T, Type::Heap, false >::make_new_fifo ) );
const_map[ Type::Heap ]->insert(
    std::make_pair( true /* yes instrumentation */ ,
                    RingBuffer< T, Type::Heap, true >::make_new_fifo ) );
...many more
```

# MONITOR BEHAVIOR - PARALLELIZATION

---

- Mechanics covered in interface, simple model here
- Run in separate thread, term on exit

```
/** apply criteria */
if( in_utilization > .5 && out_utilization < .5 )
{
    //tag kernel
    auto &tag( ag[ reinterpret_cast< std::uintptr_t >( kernel ) ] );
    tag += 1;
    if( tag == 3 )
    {
        dup_list.emplace_back( kernel );
    }
}

//after checking all kernels, handle duplication
```

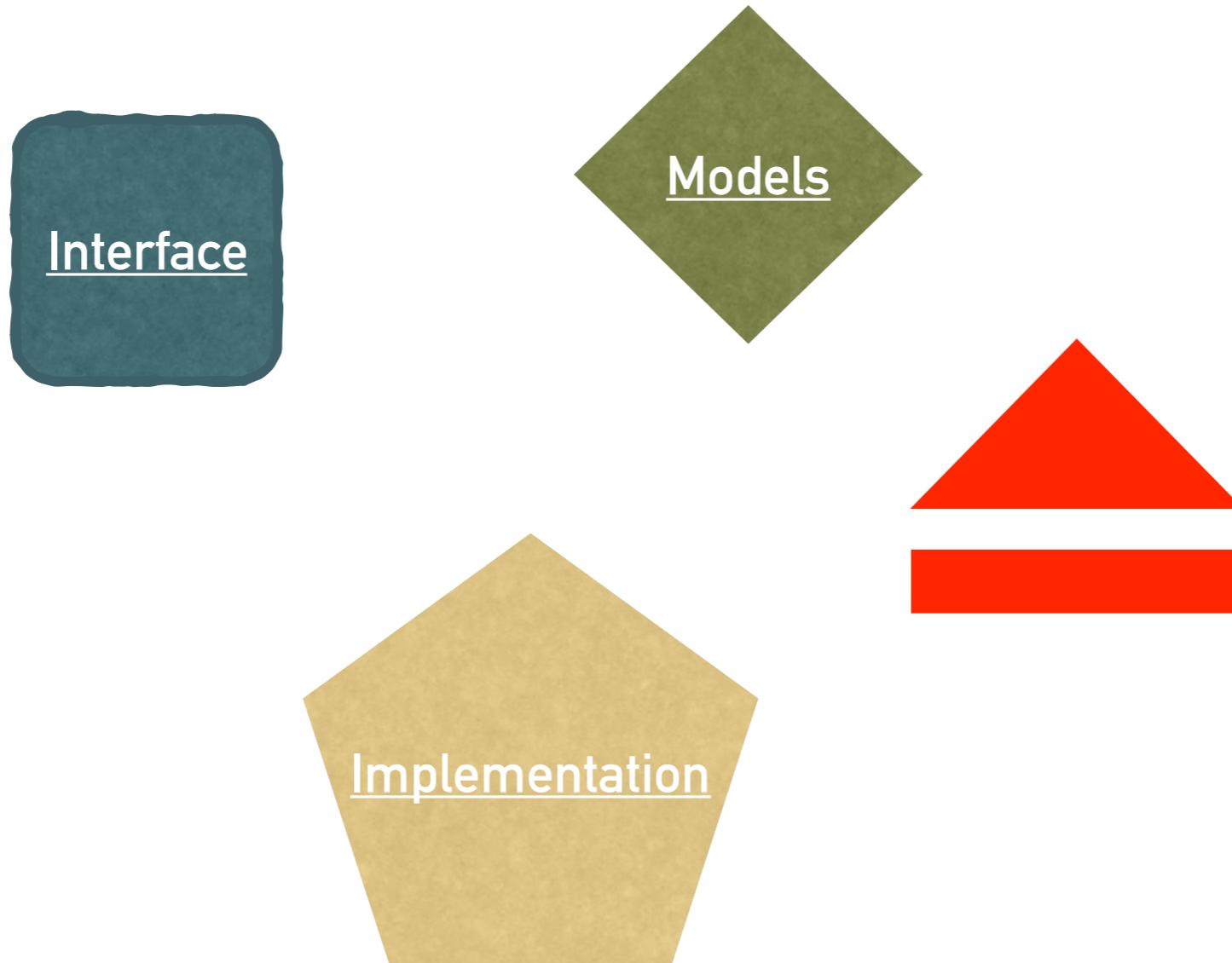
# IMPLEMENTATION TODO ITEMS

---

- Find fast SVM library to integrate (research code used LibSVM) for buffer model selection
- Integrate more production-capable network flow model for run-time re-partitioning choices
- Performant TCP links....
- RDMA on wish list
- QThreads Integration (see pool scheduler)
- *hwloc* and *netloc* integration (see `partition_scotch` )
- Perf data caching (useful for initial partition)

# CHOOSE YOUR ADVENTURE

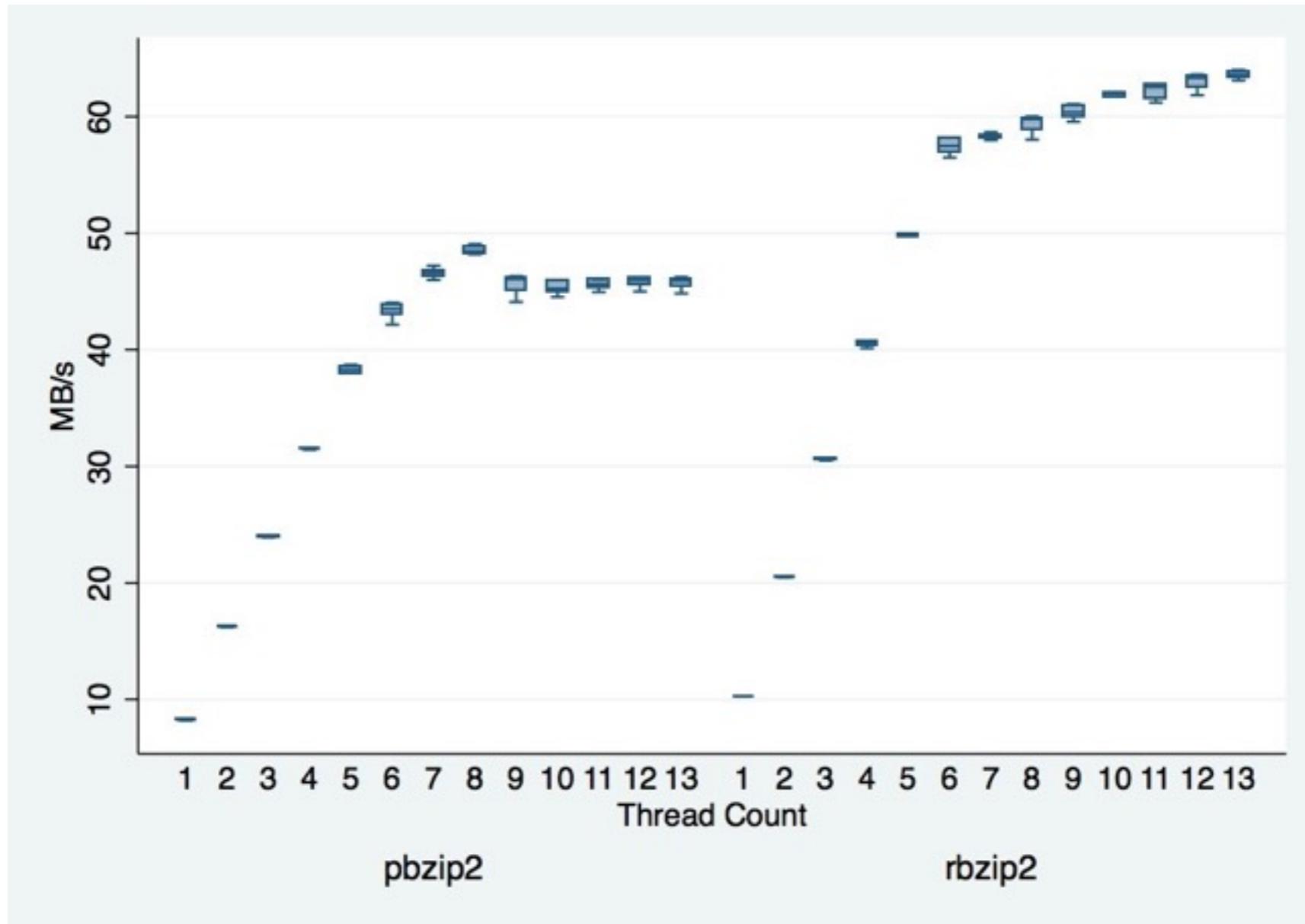
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# PERFORMANCE

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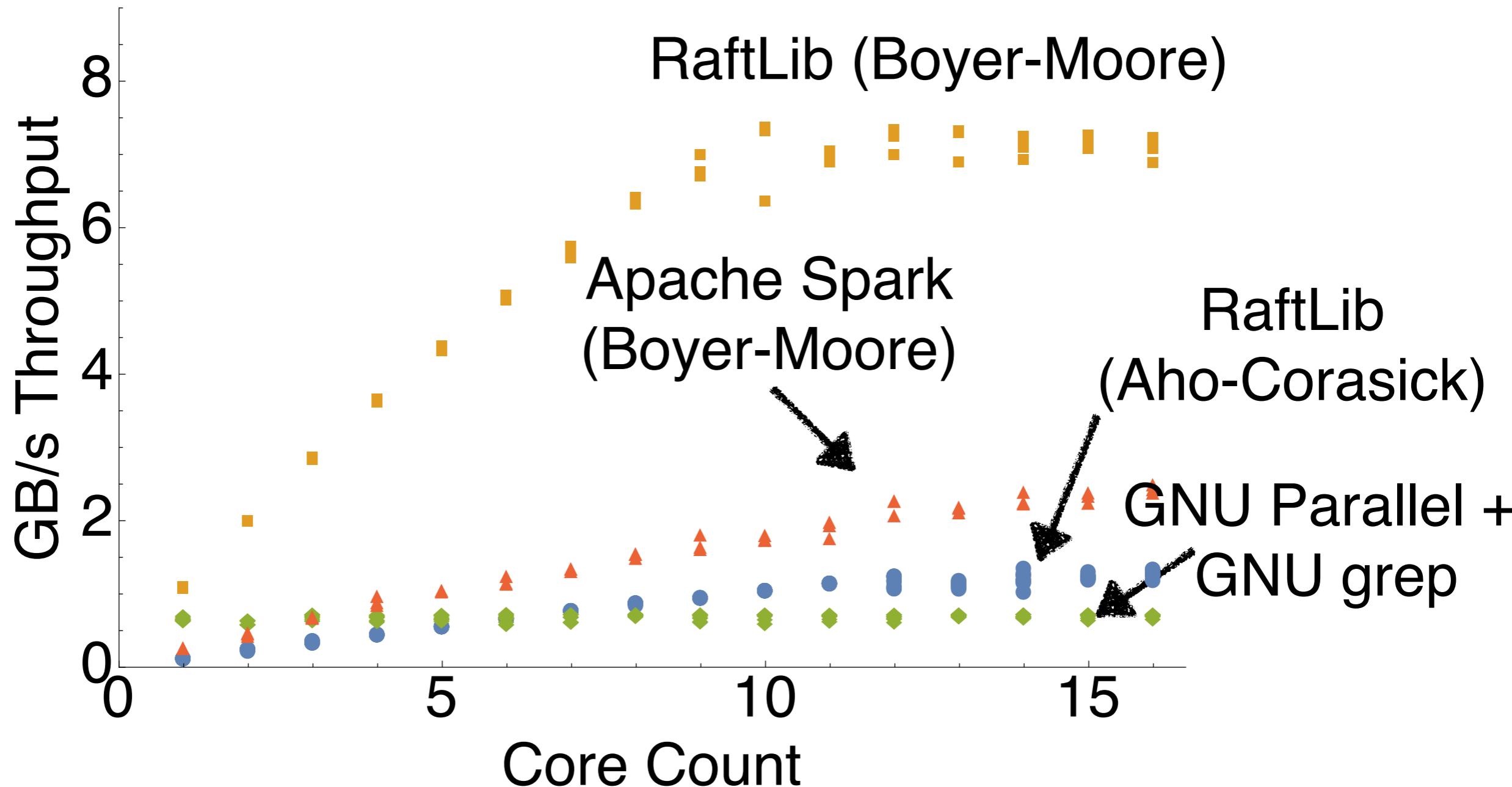
- Decent compared to pthread stock implementation of pbzip2
- Parallel Bzip2 Example: <https://goo.gl/xyQAhm>



# PERFORMANCE

---

- Fixed string search compared to Apache Spark, GNU Parallel + GNU Grep





# ABOUT ME

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my website

<http://www.jonathanbeard.io>



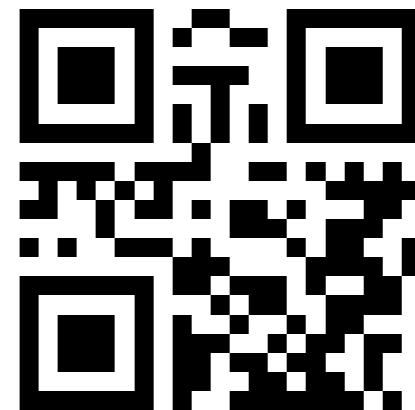
slides at

<http://goo.gl/cwT5UB>



project page

[raftlib.io](http://raftlib.io)



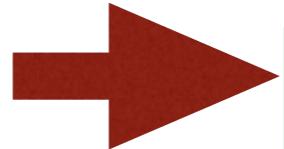


**video of talk given at #CppNow2016**

<http://goo.gl/mbxAwK>

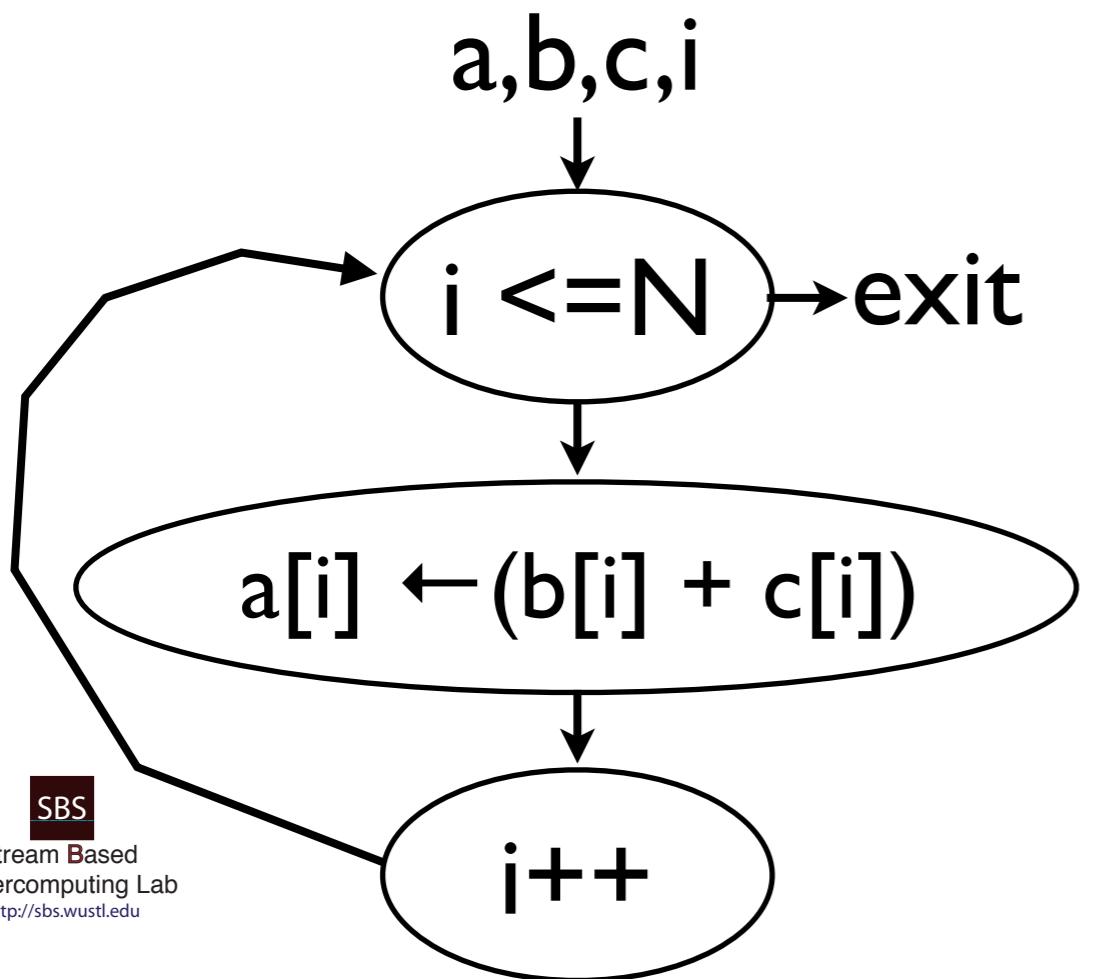


# Stream Processing



```
for i←0 through N do  
  a[i] ←(b[i] + c[i])  
  i++  
end do
```

## Traditional Control Flow



## Streaming

