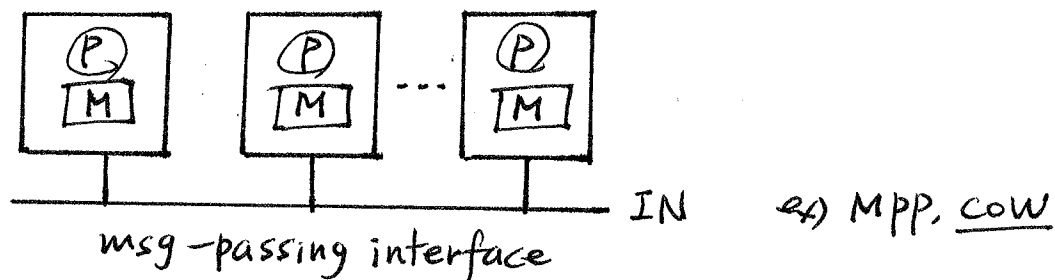
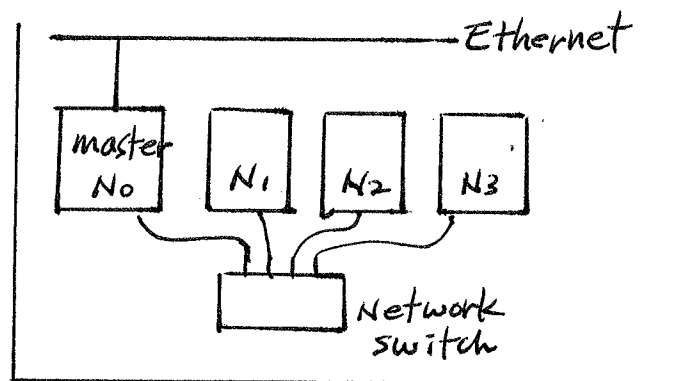


### Ch3 Distributed-mem programming with MPI



```
$> mpicc ...
$> mpic++ pl.cpp
$> mpiexec -n 4 ./a.out
    mpirun -np 4 ./a.out
```



```
MPI-Init (NULL, NULL); // initialize a mpi session
```

```
MPI-Finalize ();
```

```
— MPI-comm-size (MPI_COMM_WORLD, &comm-size);
```

Communicator  
Collection of user  
defined processes  
gets the number of processes from  
-n (4)

int comm-size

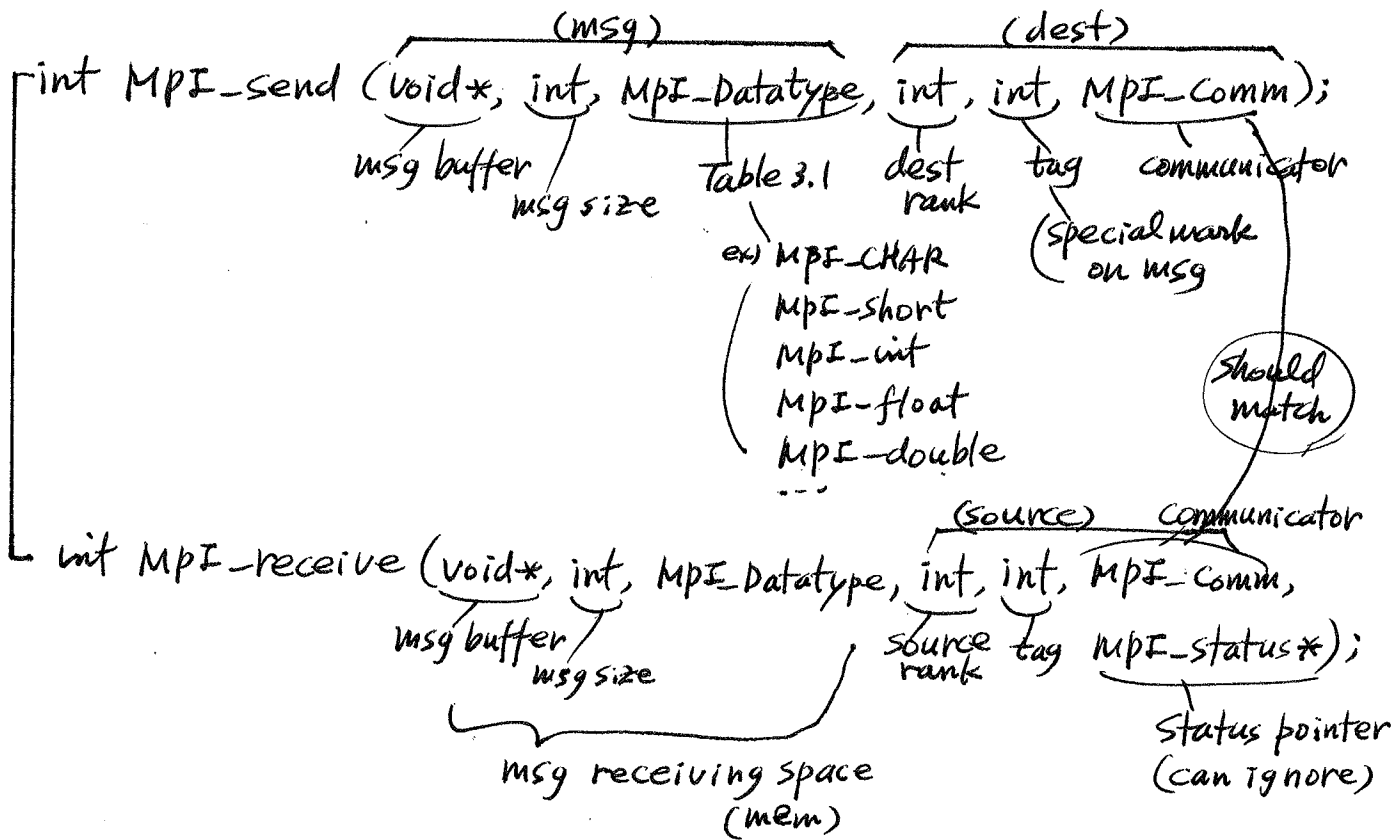
```
— MPI-comm-rank (MPI_COMM_WORLD, &my-rank);
```

gets my-rank  
among comm-size

int my-rank

— SPMD programming  
different task on  
different process

- 2 basic communication functions (p2p: point-to-point)



- for a successful communication,

$\left\{ \begin{array}{l} \text{send-comm} = \text{recv-comm} \\ \text{send-tag} = \text{recv-tag} \\ \text{dest} = \text{source} \end{array} \right.$   
 also,  $\left\{ \begin{array}{l} \text{send-type} = \text{recv-type} \\ \text{send-buffer-size} \leq \text{recv-buffer-size} \end{array} \right.$

-  $\left\{ \begin{array}{l} \text{sender} - \text{must specify receiver} \\ \text{receiver} - \text{can use wild card} \end{array} \right.$

"MPI-ANY-SOURCE" (for source rank)

- problem: receiver can receive a msg with unknown {size, source, tag}

- sol:

`MPI_Get_Count(..)` using status



↓  
ex)

```

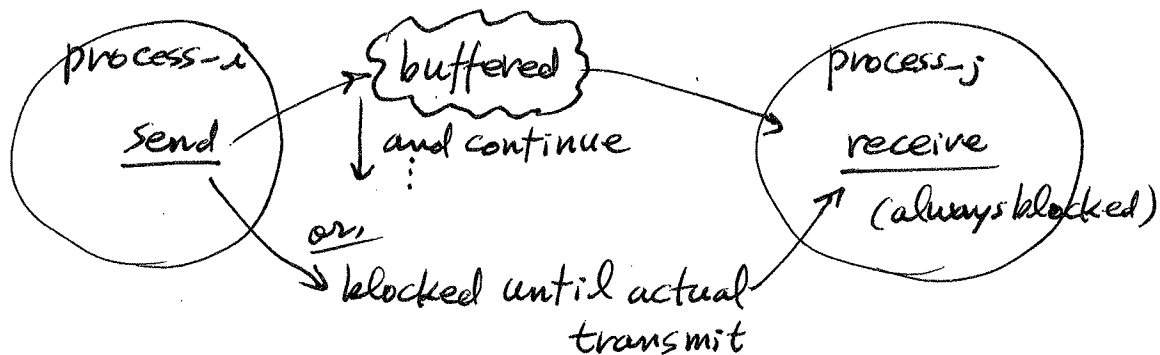
MPI_status status;
MPI_Recv(....., &status);
    (Status.MPI_Src → source
     Status.MPI_Tag → tag)
MPI_Get_count(&status, recv_type, &count);

```

size of msg

— should avoid hanging receive (#sender)

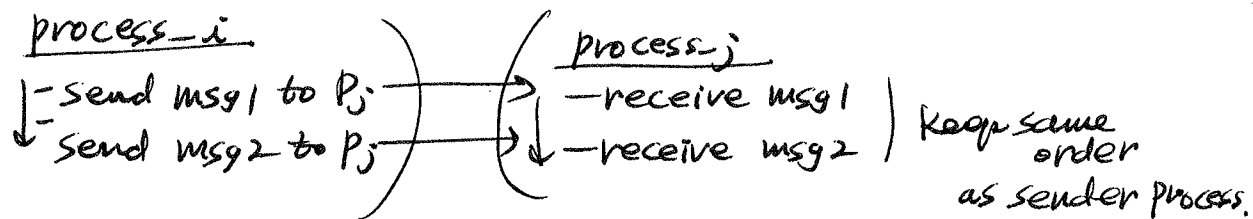
— Semantics of MPI-Send / MPI-Recv



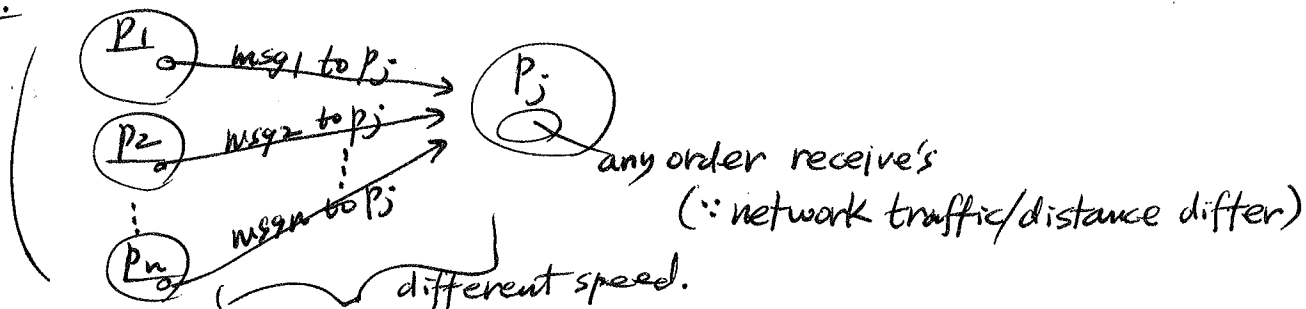
— hybrid scheme:

using cut-off,  $\begin{cases} \text{msg\_size} < \text{cut-off} \\ \Rightarrow \text{buffered} \\ \text{else} \\ \Rightarrow \text{blocked} \end{cases}$

— Non-overtaking (used in MPI) — between 2 processes

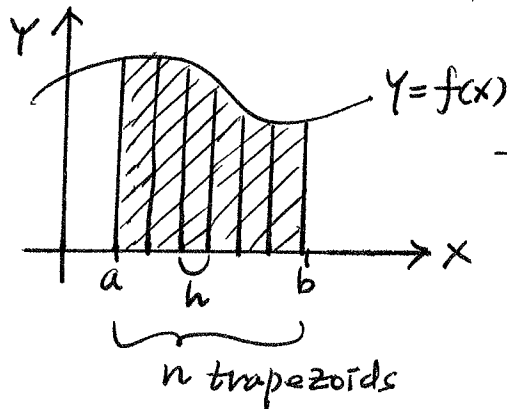


vs.

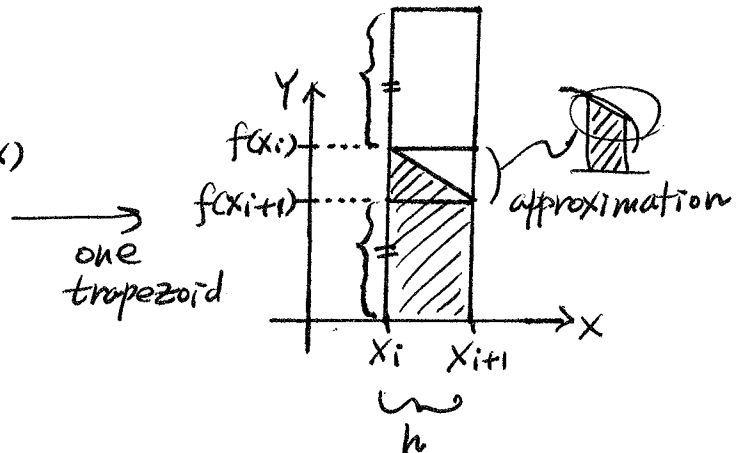


# - Trapezoid computation in MPI

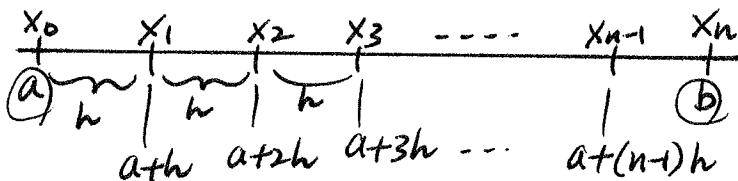
review



$$(h = \frac{b-a}{n})$$



$$\begin{aligned} \text{local integral} \\ = \left( \frac{f(x_i) + f(x_{i+1})}{2} \right) * h \end{aligned}$$



- global integral

$$= \left( \frac{f(x_0) + f(x_1)}{2} \right) * h + \left( \frac{f(x_1) + f(x_2)}{2} \right) * h + \dots + \left( \frac{f(x_{n-1}) + f(x_n)}{2} \right) * h$$

$$= \frac{h}{2} \left[ (f(x_0) + f(x_1)) + (f(x_1) + f(x_2)) + (f(x_2) + f(x_3)) + \dots + (f(x_{n-1}) + f(x_n)) \right]$$

$$= h \left[ \left( \frac{f(x_0)}{2} + f(x_1) + f(x_2) + \dots + f(x_{n-1}) \right) + \frac{f(x_n)}{2} \right]$$

Serial code

```

h = (b-a) / n;
approx = (f(a) + f(b)) / 2;

for (i = 1 ~ n-1)
{
    x_i = a + i * h;
    approx += f(x_i);
}

approx = approx * h;
  
```

— Trapezoid — parallel code (MPI) — msg passing

concept {

- divide  $n$  into  $\frac{n}{p}$
- each process computes  $\frac{n}{p}$  trapezoids
- process  $\phi$  combines local integrals

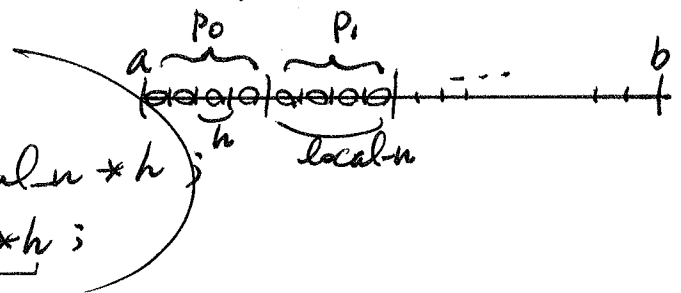
get  $a, b, n$ ; { using only 1 process, get inputs and  
send them to all other processes

$h = (b-a) \times n$ ;

$\text{local-}n = n/p$ ;

$\text{local-}a = a + \text{my-rank} \times \text{local-}n \times h$ ;

$\text{local-}b = \text{local-}a + \text{local-}n \times h$ ;



$\text{local-integral} = \text{Trap}(\text{local-}a, \text{local-}b, \text{local-}n, h)$ ;

if ( $\text{my-rank} \neq \phi$ )

send  $\text{local-integral}$  to process  $\phi$ ;

else

$\text{tot-integral} \leftarrow \text{local-integral (of } P_0)$ ;

for ( $i = 1 \sim p-1$ )

(Receive  $\text{local-integral}$  from Process- $i$ ;

$\text{tot-integral} += \text{local-integral}$ ;

display  $\text{tot-integral}$ ;

— I/O for MPI program (suggestion)

(— any process accesses std-out

(— only process  $\phi$  accesses std-in

- Collective Comm ( $\neq$  P2P — send/receive)

Communication function involves all the processes in a communicator.

ex) global sum computation

- MPI\_Reduce ( $\&$  local-integral, — input data } must be  
 $\&$  tot-integral, — output data } different  
 1, — count  
 MPI\_Double, — data type  
 MPI\_Sum, — see Table 3.2  
 $\emptyset$ , — dest rank (MPI\_MAX,  
 MPI\_Comm\_World); MPI\_MIN  
 MPI\_Proc  
 MPI\_Sum  
 :)

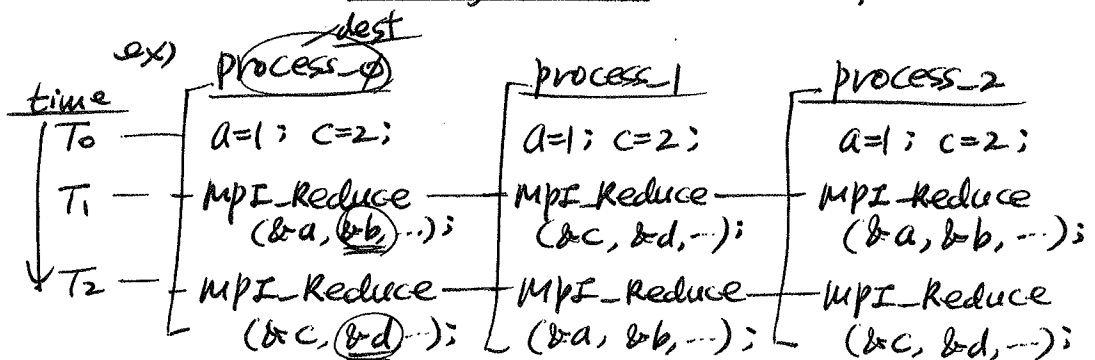
1. all processes must call the same collective func.

ex) one process calls MPI\_Recv(-)  $\rightarrow$  Crash  
 others call MPI\_Reduce(-)

2. arguments to collective func. must be compatible from processes

3. output data arg. is only used in dest-process; others should use NULL.

4.  $\neq$  tags and calling order is used for matching



Assume, MPI\_Sum.

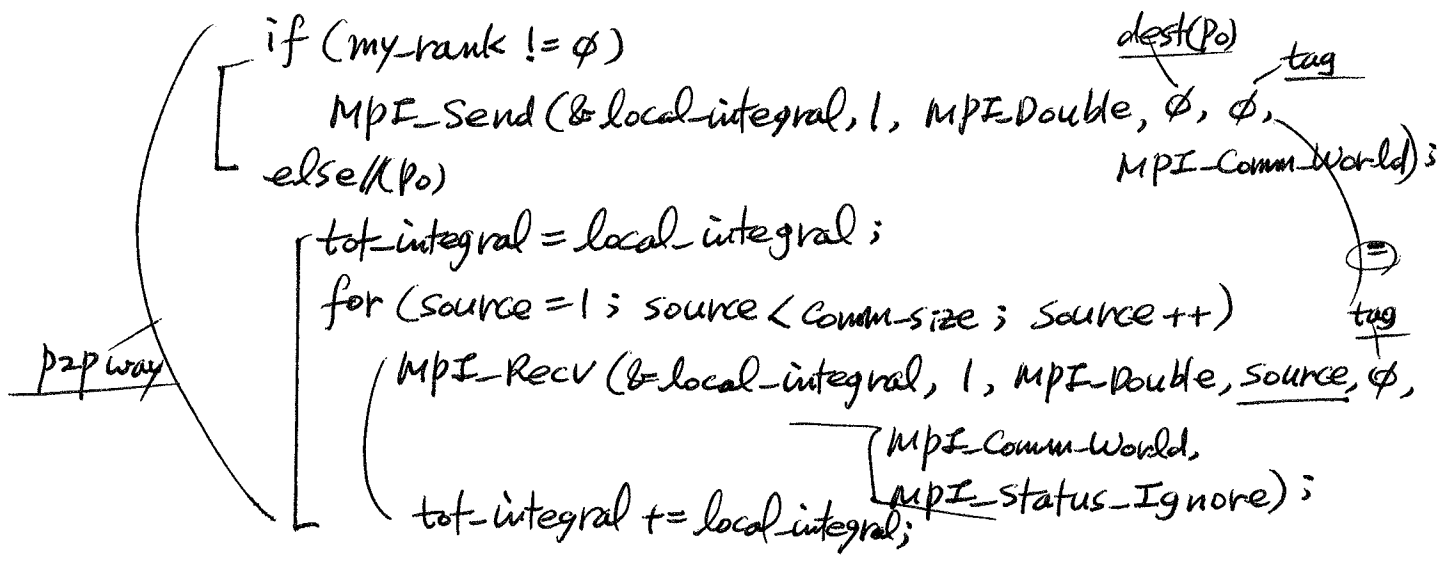
[ expected output:  $b=3, d=6$

[ actual output:  $b=1+2+1=4, d=2+1+2=5$

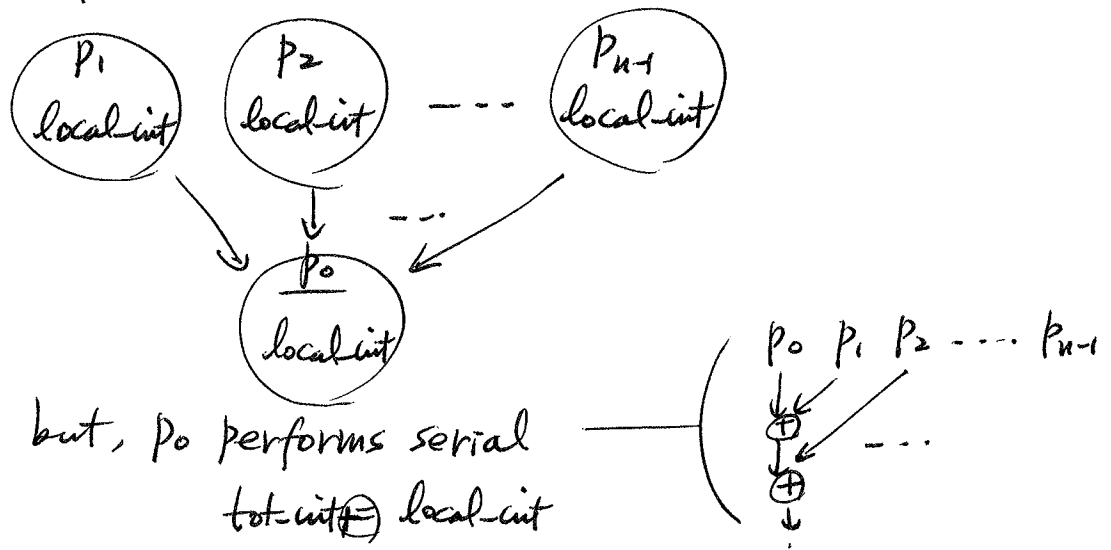
# Collective Communication

## MPI-Reduce

ex) trapezoid computing

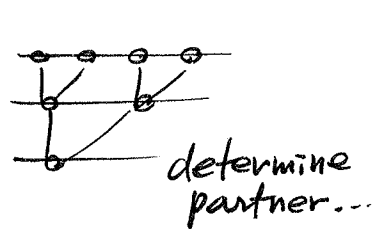


→ it looks like:



⇒ better way is using full-tree reduction.

MPI-Reduce (&local-integral, &tot-integral, 1, MPI\_Double,

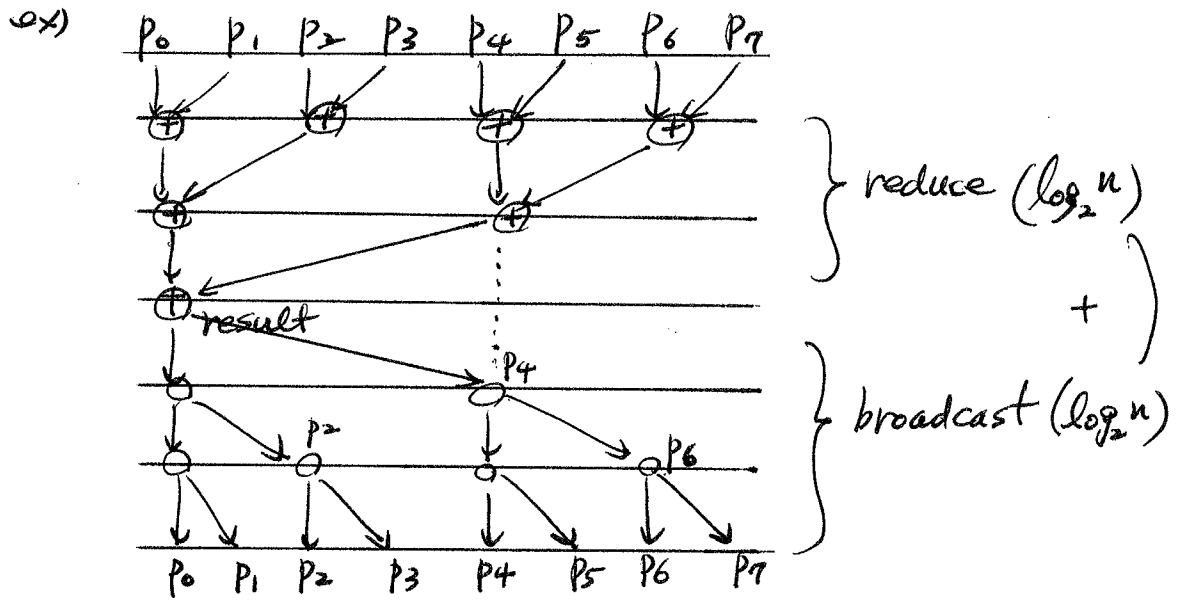


MPI\_Sum, 0, MPI\_Comm\_World);

reduction operation      dest P0

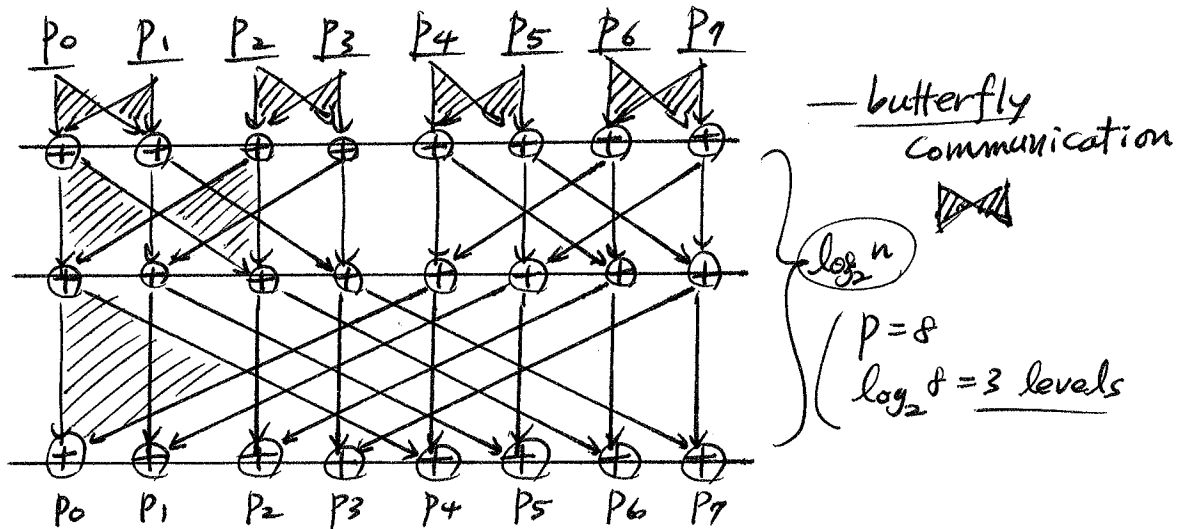
- MPI\_Allreduce { useful for the situation that  
 (all processes need the reduced result.  
 (for further operation)

- naive method: reduce to  $P_0$  and broadcast to all



ex) trapezoid Computing

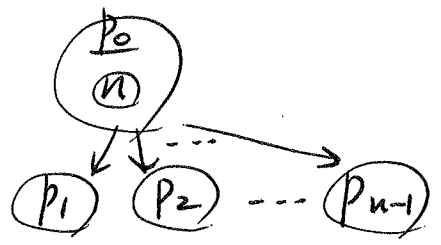
MPI\_Allreduce (local-integral, tot-integral, 1, MPI\_Double,  
 MPI\_Sum, MPI\_Comm\_World);  
 # dest-process (all processes are destination)





# - MPI-Bcast

ex)  $p_0$  accesses input  $n$  and distributes it to all others.



- in fact, this is serial in p-2-p comm. with for-loop way, i.e.,

send to  $p_1$ ;  
send to  $p_2$ ;  
⋮  
send to  $p_{n-1}$ ;

p2p way

```
if (my_rank == 0) // p0
{
    n = atoi(argv[1]);
    for (dest = 1; dest < Comm_sz; dest++)
        MPI_Send(&n, 1, MPI_Int, dest, 0, MPI_Comm_world);
}
else // other processes
    MPI_Recv(&n, 1, MPI_Int, 0, 0, MPI_Comm_world);
```

Annotations: In the `MPI_Send` call, `0` is the tag. In the `MPI_Recv` call, `0` is the source ( $p_0$ ) and `0` is the tag.

vs.

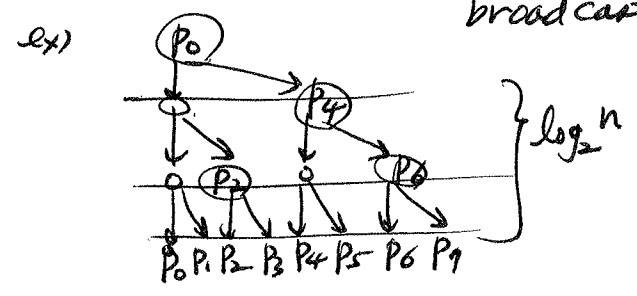
broadcast way

```
if (my_rank == 0)
{
    n = atoi(argv[1]);
    MPI_Bcast(&n, 1, MPI_Int, 0, MPI_Comm_world);
}
```

(in para - for source process  $p_0$ )  
(out para - for other processes)

\* Should call this out of  $p_0$   
∴  $p_0$  - sends  
others - receive

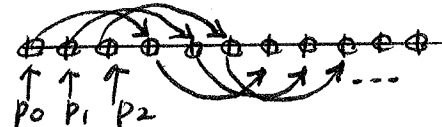
- MPI-Bcast performs full tree broadcasting.

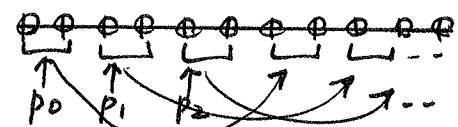


# Data distribution

ex)  $p=3$

block partition —  $0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11$   
 $\underbrace{\hspace{1.5cm}}_{p_0} \quad \underbrace{\hspace{1.5cm}}_{p_1} \quad \underbrace{\hspace{1.5cm}}_{p_2}$

cyclic partition:  $\begin{cases} p_0: 0, 3, 6, 9 \\ p_1: 1, 4, 7, 10 \\ p_2: 2, 5, 8, 11 \end{cases}$  

block-cyclic partition  $\begin{cases} p_0: 0, 1, 6, 7 \\ p_1: 2, 3, 8, 9 \\ p_2: 4, 5, 10, 11 \end{cases}$  

## MPI-Scatter

Assume:  $n$  is evenly divisible by  $p$ , and block partition.

— process  $\phi$  receives (accesses) size  $n$  vector, and  
 send the  $\begin{cases} 1st \frac{n}{p} \text{ elements to } p_0 \\ 2nd \frac{n}{p} \text{ elements to } p_1 \\ 3rd \frac{n}{p} \text{ elements to } p_2, \text{ and so on} \\ \dots \end{cases}$

ex)  $\text{if}(\text{my\_rank} == \phi) // p_0$  (ref: not evenly divisible case, use MPI\_Scatterv)

$\{$   $\text{int}^* a = \text{new int}[n];$   
 get input values for  $a[n]$  array (vector); (receiving buffer size)  
MPI\_Scatter ( $a$ ,  $\text{local\_n}$ , MPI\_Int,  $\text{local\_a}$ ,  $\text{local\_n}$ ,  
 (sending buffer\* (sending buffer size (type (sender) (receiving buffer\* (int local a[i]  
 — local size — to each process  
 $\text{MPI\_Int}, \phi, \text{comm})$ ;  
 type (receiver) source  $p_0$   
 $\}$   
 delete[]  $a$ ;  
 else  
 others MPI\_Scatter ( $a$ ,  $\text{local\_n}$ , MPI\_Int,  $\text{local\_a}$ ,  $\text{local\_p}$ ,  
 sending buffer info (MPI\_Int,  $\phi$ ,  $\text{comm}$ );  
 receiving buffer

## — Performance/time checking

— time checking for only computation part (not for elapsed time)

```
#include "timer.h" → macro def
double start, finish;
[
  Get-Time(start);
  //
  Get-Time(finish);
]
display(finish-start); — in μ seconds
```

(posix lib. func) get time of day (&t, NULL);  
now = t.tv-sec + t.tv-usec/1000000.0; }

needs <sys/time.h>

— MPI supports MPI\_Wtime()

```
double start, finish;
start = MPI_Wtime();
[
  //
  finish = MPI_Wtime();
] — each process reports exec time.
display my-rank, (finish-start); — in μ seconds
```

— parallel time checking — slowest process's exec. time

```
MPI_Barrier(MPI_Comm_World);
local_start = MPI_Wtime();
[
  //
  local_finish = MPI_Wtime();
]
local_elapsed = local_finish - local_start;
MPI_Reduce(&local_elapsed, &elapsed, 1, MPI_Double,
           MPI_Max, 0, MPI_Comm_World);
If(my_rank == 0)
  display elapsed;
```

parallel elapsed time (longest)

reduction operation

dest = p<sub>0</sub>

— Running MPI program on a hybrid system, each node is multicored.

— by default, only 1 process works on each node.

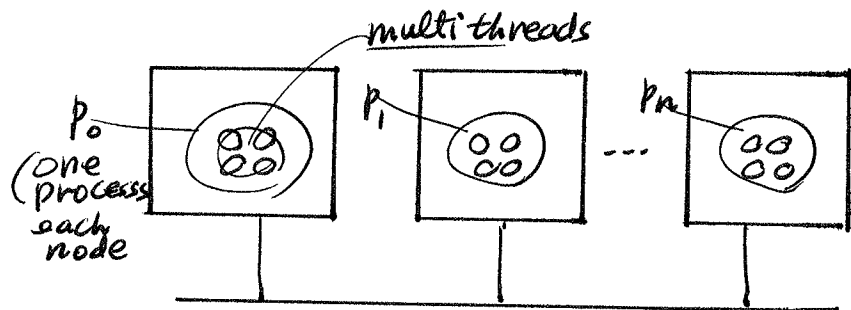
⇒ MPI + OpenMP model

(in each node (multicored)  
1 process

ex) pragma omp parallel for ---

— How to compile/run ?

```
$> mpicc -fopenmp p1.cpp ↓
$> mpi run -n 4 ./a.out ↓
```



ex) #include <mpi.h>  
#include <omp.h>

---

```
...
MPI_Init(NULL, NULL); // MPI starts
MPI_Comm_size(MPI_Comm_World, &comm_sz); // get #processes
MPI_Comm_rank(MPI_Comm_World, &my_rank); // get my rank
// mpi operations
```

#pragma omp parallel for num\_threads(4) ---

[ // omp operations

[ // mpi operations.

...