

DP Formulations

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July 26, 2017

Hopscotch

$$DP(i, j) := \text{Max number of pennies attainable by starting at square } (i, j) \\ := \max_{(i', j') \text{ within distance } k \text{ of } (i, j) \text{ s.t. } pennies(i', j') > pennies(i, j)} DP(i', j') + pennies(i, j)$$

Base case is if no squares within distance k have $pennies(i', j') > pennies(i, j)$, then value at state is $pennies(i, j)$.

Answer is in $DP(0, 0)$.

$$\begin{aligned} \text{Time Complexity} &= \text{Number of states} * \text{Time to compute each state} \\ &= O(n^2) \cdot O(k) \\ &= O(n^2 k). \end{aligned}$$

Forming Quiz Teams

$$\begin{aligned} DP(\text{subset}) &:= \text{Lowest cost pairing of people in subset} \\ &\quad (\text{let's assume people in subset are numbered } p_0, p_1, \dots, p_k) \\ &:= \min_{1 \leq j \leq k} \text{distance}(p_0, p_j) + DP(\text{subset} \setminus \{p_0, p_j\}) \\ DP(\emptyset) &:= 0 \end{aligned}$$

Answer is in $DP(\mathcal{U})$ - the full subset of all $2N$ people.

$$\begin{aligned} \text{Time Complexity} &= \text{Number of states} * \text{Time to compute each state} \\ &= O(2^{2N}) \cdot O(N) \\ &= O(N \cdot 2^{2N}). \end{aligned}$$