

Diskrete Mathematik

Solution 5

5.1 Computing Representations of Relations

a) We have $\rho^3 = \{(1, 1), (1, 3), (2, 2), (4, 4)\}$ and

$$M^{\rho^*} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 \end{bmatrix}$$

5.2 Operations on Relations

	Relation	reflexive	symmetric	transitive
a)	$< \circ $	✗	✗	✓
b)	$ \cup \equiv_2$	✓	✗	✗
c)	$ \cup ^{-1}$	✓	✓	✗

a) Two numbers (a, b) are in the relation whenever there exists an x such that $a < x$ and $x | b$. This relation is not reflexive, since $(1, 1) \notin < \circ |$. Moreover, it is not symmetric, because $(1, 2) \in < \circ |$, but $(2, 1) \notin < \circ |$. This relation is transitive. For any (a, b, c) , assume that there exist some x and y , such that $a < x$, $x | b$, $b < y$ and $y | c$. From $x | b$ it follows that $x \leq b$, hence, $a < x \leq b < y$. Therefore, $a < y$ and $y | c$.

b) Two numbers (a, b) are in the relation whenever $a | b$ or $a \equiv_2 b$. This relation is reflexive, since for any a , we have $a \equiv_2 a$ (alternatively, one could use the fact that $a | a$). It is, however, not symmetric, because $(1, 2) \in | \cup \equiv_2$, but $(2, 1) \notin | \cup \equiv_2$. It is also not transitive, since $(3, 1) \in | \cup \equiv_2$ and $(1, 2) \in | \cup \equiv_2$, but $(3, 2) \notin | \cup \equiv_2$.

c) Two numbers (a, b) are in the relation whenever $a | b$ or $b | a$. This relation is reflexive, since for any a , we have $a | a$. It is also symmetric, because for any (a, b) , we trivially have $a | b$ or $b | a$ if and only if $b | a$ or $a | b$. The relation is, however, not transitive, since $(3, 1) \in | \cup |^{-1}$ and $(1, 2) \in | \cup |^{-1}$ but $(3, 2) \notin | \cup |^{-1}$.

5.3 A False Proof

a) For an arbitrary $x \in A$, there does not always exist a $y \in A$ such that $x \rho y$.

b) Consider the following counterexample: $A = \{1, 2\}$ and $\rho = \{(1, 1)\}$. The relation ρ is symmetric and transitive. However, it is not reflexive, since $2 \rho 2$ does not hold.

5.4 An Equivalence Relation

a) We prove that \sim satisfies all properties of an equivalence relation.

Reflexivity: For any point $(x, y) \in \mathbb{R}^2 \setminus \{(0, 0)\}$, we have $(x, y) \sim (x, y)$, because one can choose $\lambda = 1$ in the definition of \sim .

Symmetry: Let $x_1, y_1, x_2, y_2 \in \mathbb{R} \setminus \{0\}$ and assume that $(x_1, y_1) \sim (x_2, y_2)$. It follows that $x_1 = \lambda x_2$ and $y_1 = \lambda y_2$ for some $\lambda > 0$. Hence, $x_2 = \frac{1}{\lambda} x_1$ and $y_2 = \frac{1}{\lambda} y_1$, where $\frac{1}{\lambda} > 0$. Therefore, $(x_2, y_2) \sim (x_1, y_1)$.

Transitivity: Let $x_1, y_1, x_2, y_2, x_3, y_3 \in \mathbb{R} \setminus \{0\}$ and assume that $(x_1, y_1) \sim (x_2, y_2)$ and $(x_2, y_2) \sim (x_3, y_3)$. This means that $(x_1, y_1) = (\lambda_1 x_2, \lambda_1 y_2)$ and $(x_2, y_2) = (\lambda_2 x_3, \lambda_2 y_3)$ for some $\lambda_1, \lambda_2 > 0$. It follows that $(x_1, y_1) = (\lambda_1 \lambda_2 x_3, \lambda_1 \lambda_2 y_3)$, where $\lambda > 0$ is defined as $\lambda_1 \lambda_2$. Hence, $(x_1, y_1) \sim (x_3, y_3)$.

b) An equivalence class $[(x, y)]_\sim$ contains all points on the ray through the origin $(0, 0)$ and the point (x, y) (excluding the origin). Note that no equivalence class can contain the origin $(0, 0)$ (\sim is only defined on $\mathbb{R}^2 \setminus \{(0, 0)\}$).

5.5 Order Relations on Quotient Sets

To show that $(A/\theta, \rho)$ is a partial order we need to show that the relation ρ is reflexive, antisymmetric, and transitive.

Notice that for all $a \in A$ the element $\text{least}([a]_\theta)$ is well-defined because $[a]_\theta \neq \emptyset$. Indeed, since θ is an equivalence relation, it is reflexive, so that $a \theta a$ and $a \in [a]_\theta$, by definition of $[a]_\theta$.

- **Reflexivity:** Let $[a]_\theta \in A/\theta$. One has

$$\begin{aligned} & [a]_\theta \rho [a]_\theta \\ \iff & \text{least}([a]_\theta) \leq \text{least}([a]_\theta) \quad (\text{Definition of } \rho) \\ \iff & \text{true.} \quad (\text{Reflexivity of } \leq) \end{aligned}$$

- **Antisymmetry:** Let $[a]_\theta, [b]_\theta \in A/\theta$. We have

$$\begin{aligned} & ([a]_\theta \rho [b]_\theta) \wedge ([b]_\theta \rho [a]_\theta) \\ \iff & (\text{least}([a]_\theta) \leq \text{least}([b]_\theta)) \wedge (\text{least}([b]_\theta) \leq \text{least}([a]_\theta)) \quad (\text{Definition of } \rho, \text{ twice}) \\ \implies & \text{least}([a]_\theta) = \text{least}([b]_\theta) \quad (\text{Antisymmetry of } \leq) \\ \implies & [a]_\theta \cap [b]_\theta \neq \emptyset \quad (\text{Definition of } \cap) \\ \implies & [a]_\theta = [b]_\theta. \quad (\text{Theorem 3.11}) \end{aligned}$$

- **Transitivity:** Let $[a]_\theta, [b]_\theta, [c]_\theta \in A/\theta$. We have

$$\begin{aligned} & ([a]_\theta \rho [b]_\theta) \wedge ([b]_\theta \rho [c]_\theta) \\ \iff & (\text{least}([a]_\theta) \leq \text{least}([b]_\theta)) \wedge (\text{least}([b]_\theta) \leq \text{least}([c]_\theta)) \quad (\text{Definition of } \rho, \text{ twice}) \\ \implies & \text{least}([a]_\theta) \leq \text{least}([c]_\theta) \quad (\text{Transitivity of } \leq) \\ \implies & [a]_\theta \rho [c]_\theta. \quad (\text{Definition of } \rho) \end{aligned}$$

5.6 Lifting an Operation to Equivalence Classes

a) We define the function $\text{sum} : A^2 \rightarrow A$ by

$$\text{sum}((a, b), (c, d)) \stackrel{\text{def}}{=} (ad + bc, bd).$$

Observe that $bd \neq 0$ since $b \neq 0$ and $d \neq 0$.

b) f is θ -consistent if and only if

$$(b_1 \theta b'_1 \text{ and } b_2 \theta b'_2) \implies f(b_1, b_2) \theta f(b'_1, b'_2)$$

is true for all $b_1, b_2, b'_1, b'_2 \in B$. Alternatively (and equivalently) we could say that f is θ -consistent if and only if

$$([b_1]_\theta = [b'_1]_\theta \text{ and } [b_2]_\theta = [b'_2]_\theta) \implies [f(b_1, b_2)]_\theta = [f(b'_1, b'_2)]_\theta$$

is true for all $b_1, b_2, b'_1, b'_2 \in B$.

c) Let $(a, b), (a', b'), (c, d), (c', d') \in A$ be arbitrary. We have

$$\begin{aligned} (a, b) \sim (a', b') \text{ and } (c, d) \sim (c', d') & \\ \iff ab' = ba' \text{ and } cd' = dc' & \quad (\text{def. } \sim) \\ \implies ab' \cdot dd' + cd' \cdot bb' = ba' \cdot dd' + dc' \cdot bb' & \\ \iff ad \cdot b'd' + bc \cdot b'd' = bd \cdot a'd' + bd \cdot b'c' & \quad (\text{comm.}) \\ \iff (ad + bc) \cdot b'd' = bd \cdot (a'd' + b'c') & \quad (\text{distr.}) \\ \iff (ad + bc, bd) \sim (a'd' + b'c', b'd') & \quad (\text{def. } \sim) \\ \iff \text{sum}((a, b), (c, d)) \sim \text{sum}((a', b'), (c', d')). & \quad (\text{def. sum}) \end{aligned}$$

Hence, sum is \sim -consistent.