

23. a.

$$\forall j \in \mathbb{N}, S_j \leq 1$$

$$S_j = \sum_{i=(2^{j-1})+1}^{2^j} \left( \frac{1}{i} \right) \quad \left\{ \begin{array}{l} \xrightarrow{2n} \\ \xrightarrow{n} \end{array} \right\} \quad \underbrace{\frac{1}{n+1} + \frac{1}{n+2} + \dots + \frac{1}{2n}}_{n \text{ terms}}$$

$$= \sum_{i=n+1}^{2n} \left( \frac{1}{i} \right)$$

$$\leq n \cdot \left( \frac{1}{n+1} \right) = \frac{n}{n+1}$$

$$\leq \underline{\underline{1}}$$

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2.3 b)

To prove: For any  $j \in \mathbb{N}$ ,  $\sum_{i=2^{j-1}+1}^{2^j} \frac{1}{i} \geq \frac{1}{2}$

$$n = 2^{j-1}$$

$$\sum_{i=n+1}^{2n} \frac{1}{i} \geq \frac{1}{n+1} \geq \frac{1}{2^{j-1}+1} \geq \frac{1}{2}$$

2.3 c)

To prove: For any  $k \in \mathbb{N}_0$ :

$$\sum_{i=1}^{2^k} \frac{1}{i} \leq k+1$$

Proof:

$$1 + \sum_{j=1}^k S_j \leq k+1$$

$$\sum_{j=1}^k S_j \leq k$$

$$\sum_{j=1}^k S_j \leq \sum_{j=1}^k 1 \leq k$$

To prove: For any  $k \in \mathbb{N}_0$ :

$$\sum_{i=1}^{2^k} \frac{1}{i} \geq \frac{k+1}{2}$$

Proof:

$$1 + \sum_{j=1}^k S_j \geq \frac{k+1}{2}$$

$$2 + 2 \sum_{j=1}^k S_j \geq k+1$$

$$2 \sum_{j=1}^k S_j \geq k-1$$

$$2 \sum_{j=1}^k S_j \geq 2 \sum_{j=1}^k \frac{1}{2} \geq 2 \cdot k \cdot \frac{1}{2} \geq k \geq k-1$$

2.5.a

```
for(a in range(1,n))  
  for(b in range(1,n))  
    for(c in range(1,n))  
      for(d in range(1,n))  
        if( f(a)+f(b)+f(c) = f(d))  
          output("YES")  
output("NO")
```



$$n \cdot n \cdot n \cdot n = n^4$$

$$\therefore O(n^4)$$

2.5.b

```
array: possible_sums = [false] * n3
for(a in range(1, n))
  for(b in range(1, n))
    for(c in range(1, n))
      sum_abc = f(a) + f(b) + f(c)
      if(sum_abc in range(1, n3))
        possible_sums(sum_abc) = true
for(d in range(1, n))
  if(possible_sums(f(d)))
    output("YES")
  output("NO")
```

$(n \cdot n \cdot n) + n = O(n^3)$