

Problem 1.1:

From the question we have, for every $a \in A$, $b \in B$ and $c \in C$

$$a) \ a \subseteq \gamma_1(b) \text{ iff } \alpha_1(a) \leq b$$

$$b) \ b \subseteq \gamma_2(c) \text{ iff } \alpha_2(b) \leq c$$

We have to prove, $a \subseteq \gamma_1\gamma_2(c)$ iff $\alpha_2\alpha_1(a) \leq c$

From the contractive property we get,

$$\alpha_2\gamma_2(c) \leq c$$

$$\alpha_2(b) \leq c \quad [\text{As } b \subseteq \gamma_2(c)]$$

$$\alpha_2\alpha_1(a) \leq c \quad [\text{As } \alpha_1(a) \leq b]$$

And from the expansive property we get,

$$a \subseteq \gamma_1\alpha_1(a)$$

$$a \subseteq \gamma_1(b) \quad [\text{As } \alpha_1(a) \leq b]$$

$$a \subseteq \gamma_1\gamma_2(c) \quad [\text{As } b \subseteq \gamma_2(c)]$$

So, $\alpha_2\alpha_1$ and $\gamma_1\gamma_2$ form a Galois connection

Problem 1.2:

From Galois connection property we know, for every $a \in A$ and every $c \in C$

$$c \subseteq \gamma(a) \text{ iff } \alpha(c) \leq a$$

To prove α is monotone, we need to prove that for any c_1, c_2 if $c_1 \subseteq c_2$ then $\alpha(c_1) \leq \alpha(c_2)$.

From the expansive property we get,

$$c_2 \subseteq \gamma(\alpha(c_2))$$

$$\text{So, } c_1 \subseteq \gamma(\alpha(c_2)) \quad [\text{As } c_1 \subseteq c_2]$$

$$\text{So, } \alpha(c_1) \leq \alpha(c_2) \quad [\text{From Galois connection property}]$$

Problem 1.3:

From Galois connection property we know, for every $a \in A$ and every $c \in C$

$c \subseteq \gamma(a)$ iff $\alpha(c) \leq a$

To prove γ is monotone, we need to prove that for any a_1, a_2 if $a_1 \leq a_2$ then $\gamma(a_1) \subseteq \gamma(a_2)$.

From the contractive property we get,

$$\alpha(\gamma(a_1)) \leq a_1$$

So, $\alpha(\gamma(a_1)) \leq a_2$ [As $a_1 \leq a_2$]

So, $\gamma(a_1) \subseteq \gamma(a_2)$ [From Galois connection property]

Problem 2.2:

$(\lambda x. \lambda y. x)(\lambda z. (\lambda x. \lambda y. x) z ((\lambda x. z x)(\lambda x. z x)))$

Normal Order:

$\rightarrow \lambda y. (\lambda z. (\lambda x. \lambda y. x) z ((\lambda x. z x)(\lambda x. z x)))$

$\rightarrow \lambda y. ((\lambda x. \lambda y. x) ((\lambda x. z x)(\lambda x. z x)))$

Applicative Order:

$\rightarrow (\lambda x. \lambda y. x)((\lambda x. \lambda y. x) ((\lambda x. z x)(\lambda x. z x)))$

$\rightarrow (\lambda x. \lambda y. x)((\lambda x. \lambda y. x) (z (\lambda x. z x)))$

Problem 2.3:

$\text{interpret}(x) = x$

$\text{interpret}(\lambda x. E1) = \text{let } E2 = \text{interpret}(E1)$

In case E2 of

$x = \lambda x. E2$

$_ = \text{interpret}(E2)$

$\text{interpret}(E1 E2) = \text{let } E3 = \text{interpret}(E1)$

In case E3 of

$\lambda x. E4 = \text{interpret}(E4[E2/x])$

$x = E3 \text{ interpret}(E2)$

$_ = \text{interpret}(E3 E2)$