Type State

CS242

Lecture 9

APIs

Consider a typical use of a file API:

```
File f = open("MyFile");
Char c = f.read();
f.close()
```

APIs

File operations mutate state and can change what operations are legal

```
File f = open("MyFile"); // opening the file enables reading

Char c = f.read();

f.close() // but once the file is closed, reading is no

// legal even though f is still available
```

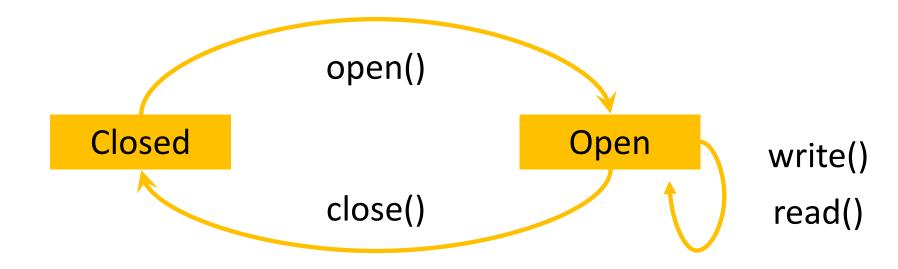
Stateful APIs

Many APIs are stateful in this way

- Correct usage depends on the state of the system
 - Which changes over time as API calls are made
- States and state changes are commonly modeled using finite state machines

Files

- Files can be open or closed
- A closed file can only be opened
- An open file can be read, written or closed



Discussion

Many similar examples of this two state machine in real APIs:

- File open/close
- Lock acquire/release
- Message encrypted/decrypted
- User logged in/logged out
- Resource initialized/unitialized

More Complex Machines

 File example can be refined, revealing more states specific to reading/writing

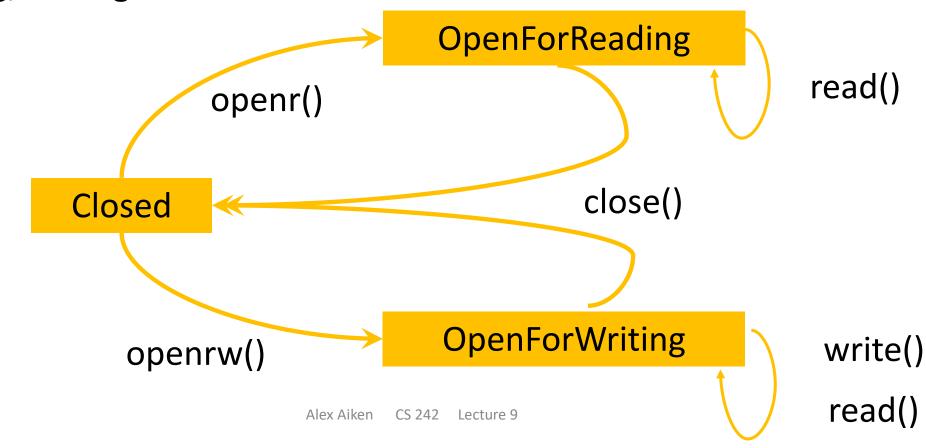
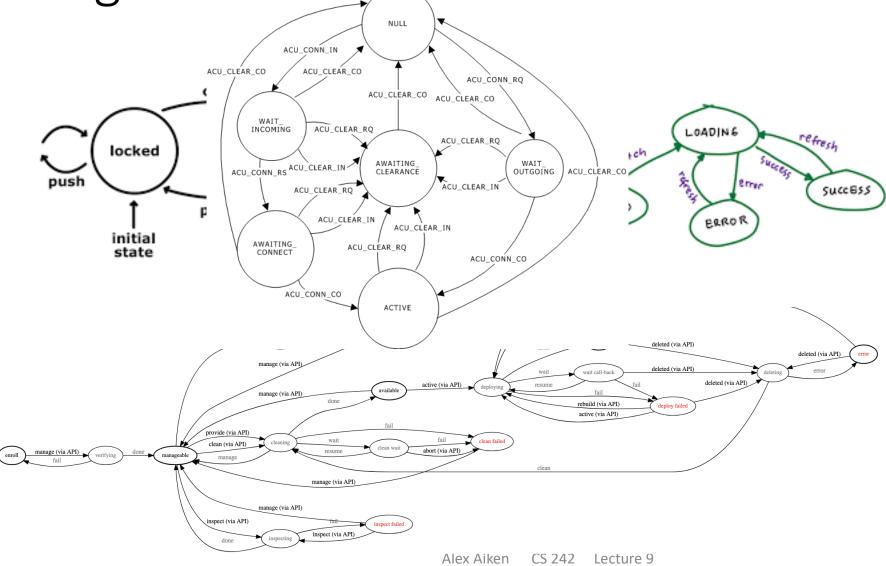


Image Search for "State Machine API"



Enforcing APIs

 We would like to enforce that API usage is correct according to a state machine specification

Guarantees

- We start in a correct state
- The only calls that are made are valid in the current state
- Can also enforce that we end in a correct state
 - e.g., Don't leave a resource locked forever

Type State

- Idea: Use the type system
- Each state machine state corresponds to a distinct type

- Method invocations cause transitions in the state diagram
 - open(): ClosedFile -> OpenFile
 - close(): OpenFile -> ClosedFile
 - read(): OpenFile -> OpenFile

A Possible Implementation

```
Class FileRecord {
                                    Class ClosedFile {
  File f;
                                           FileRecord r;
                                           new ClosedFile(r: FileRecord);
  Str name; }
                                           open(): OpenFile; }
Class OpenFile {
      FileRecord r;
      new OpenFile(name: Str);
      read(): Char;
      close(): ClosedFile }
```

Example

```
f = new OpenFile("foo")  // f is of type OpenFile
c = f.read()
x = f.close()  // x is of type ClosedFile
// but ...
d = f.read()  // f has been closed!
```

Discussion

- The OpenFile and ClosedFile types allow us to express the transition from an open file to a closed file in the type system
 - And even back again
- The types only permit valid operations for the state of the file

- But ...
 - Nothing gets rid of the OpenFile objects when the file is closed
 - We could null out the FileRecord, but the OpenFile object would still exist

What Do We Want?

```
f = new OpenFile("foo")
c = f.read()
x = f.close()
// ideally at this point the f object would be deallocated
// and it would also be impossible to invoke operations on f
// So...
d = f.read() // this should be a type error!
```

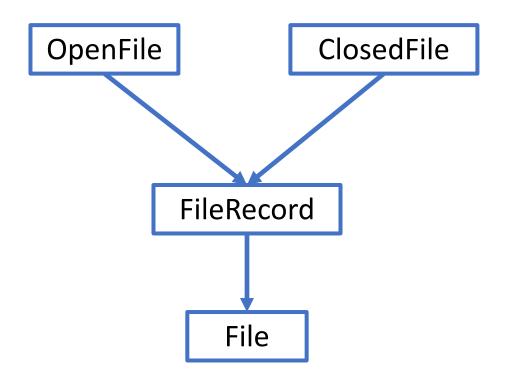
The Real Problem

- Types are not the problem
 - We can express the state transitions as functions that map one type to another
- Keeping track of object references is the problem
 - In particular, tracking *aliasing* between names

Aliasing

 OpenFile and ClosedFile methods both modify the state of the File object in the FileRecord

- An example of aliasing
 - Having multiple names (ways to access) the same value
- Either object can change the state of the file without the knowledge of the other



Aliasing Control

A Classic Example

```
copy(char *x, char *y) {
    ...
}
```

But what about copy(a,a)?

A Classic Example

```
copy(restrict char *x, restrict char *y) {
    ...
}
```

Semantics: In C, a restricted pointer cannot be aliased to any other pointer in scope.

A Point of View

- Aliasing is bad
- State can be modified through one name and those changes are visible through a different name
 - Leads to subtle and difficult bugs
- But aliasing is very common in real programs
 - Impossible to avoid
 - E.g., references passed as arguments to functions
 - Object-oriented code is particularly prone to generating aliasing

Idea #1

Maybe aliasing is not the problem ...

Problems arise only when aliasing is combined with mutation

- So, disallow mutation!
 - The pure functional programming viewpoint

Could This Really Work?

- People have studied pure functional languages for decades
 - No mutation, whenever a data structure is changed a copy is made
- A surprising number of standard algorithms are just as efficient without mutation of state
 - Sometimes just amortized bounds, but that is still quite good!
- But there are some operations that seem to fundamentally require mutation to be efficient
 - Update in place of an array is O(1)
 - The best known functional update is O(log N) in the size of the array

A Practical Approach

Split the world into mutable and immutable values

Rust

```
let x = 5  // immutable
let mut x = 5  //mutable
x = 3  // only allowed if x is mutable
```

• ML

```
    let x = 5  // immutable
    let x = ref 5  // mutable
    x := 3
```

Separating Mutable & Immutable

- Not entirely a new idea
 - E.g., const in C
- Gaining in popularity
 - More languages are making this distinction
 - With immutability being the default
- Now accepted as a good idea
 - Limit the possibility of mutation to places it is really needed
 - Make these points obvious in the syntax & types

Idea #2

- Control aliasing in the type system
 - Track it, restrict it, or even disallow it
- Ownership types
 - Track aliases using types
- There is a large literature on ownership types
 - Some quite elaborate ...

Ownership in Rust

- Rust has a simple ownership model
- There is always a single *owner* variable of every object
 - Owning = responsible for the resources of the object
- Implications
 - An object with no owner is deallocated
 - When an owner goes out of scope, the owned object is deallocated
 - Copies transfer ownership
 - x = y removes ownership from y and transfers it to x
 - y can no longer be used after the assignment

Ownership Example

```
fn main(){
  let v = vec[1,2,3]; // v owns the vector
  let v2 = v; // moves ownership to v2
  // println(v) compile-time error!
  display(v2); // ownership is moved to display
  // println(v2); compile-time error!
fn display(v:Vec<i32>){
  println(v);
  // v goes out of scope here and the vector is deallocated
```

Another Ownership Example

```
A() {
X = \text{new Foo()}; // X \text{ is the owner}
           // ownership is transferred to the argument of Bar
Y = Bar(X)
                // and then back to Y
// Y goes out of scope and the vector is deallocated
Bar(Z) {
  return Z // ownership is transferred back to the caller
```

Lifetimes

• Rust reasons about aliasing/ownership by using *lifetimes*

- The lifetime of a variable is the span between
 - The definition (first use)
 - The last use
- Roughly: Lifetimes of aliases cannot overlap
 - Thus it is important to minimize lifetimes
 - We will refine this rule shortly

Lifetimes

```
fn main(){
  let v = vec[1,2,3]; // vector v owns the object
  let v2 = v; // moves ownership to v2
  // println(v) compile-time error!
  display(v2); // ownership is moved to display
  // println(v2); compile-time error!
fn display(v:Vec<i32>){
   println(v);
  // v goes out of scope here and the vector is deallocated
```

Lifetimes: A Compile Time Error

```
fn main(){
  let v = vec[1,2,3]; // vector v owns the object
  let v2 = v;
                     // moves ownership to v2
  println(v)
                     // compile-time error!
                     // ownership is moved to display
  display(v2);
  // println(v2); compile-time error!
fn display(v:Vec<i32>){
   println(v);
  // v goes out of scope here and the vector is deallocated
```

Lifetimes: A Fix

```
fn main(){
  let v = vec[1,2,3]; // vector v owns the object 2
  println(v)
               // compile-time error!
  let v2 = v; // moves ownership to v2
  display(v2); // ownership is moved to display
  // println(v2); compile-time error!
fn display(v:Vec<i32>){
  println(v);
  // v goes out of scope here and the vector is deallocated
```

Aliasing Control in Rust

- Having only a single owner is painful in many situations
 - Can never have another name for an object or even a piece of an object
 - E.g., makes it impossible to name a subarray
 - And we often don't need to take ownership anyway
- Rust allows the creation of explicit aliases
 - called references or borrows
- There are two kinds of references:
 - mutable
 - immutable

Example: Immutable Reference

```
A() {
X = new Foo(); // X is the owner
Y = &X
             // An immutable reference to X; X is still the owner.
             // pass an immutable reference to Bar
Bar(Y)
Bar(&Z) {
 ... = .. Z ... // can read from Z in Bar as many times as we like
     Global.x = Z but if we try to store Z somewhere that outlives Bar we get an error
```

Example: Immutable Reference

```
A() {
X = new Foo(); // X is the owner
Y = &X
               // An immutable reference to X; X is still the owner.
Bar(Y, Y) // pass two immutable references to Bar
Bar(&A, &B) {
 ... = .. A ... // can read from A and B in Bar as many times as we like
 ... = ... B ...
```

Example: Mutable Reference

```
A() {
X = new Foo(); // X is the owner
Y = &mut X; // Y is a mutable reference to X
Bar(Y) // pass a mutable reference to Bar
Bar(&mut Z) {
 Z.f = ... // can mutate Z
```

Example: Mutable Reference

```
A() {
X = new Foo(); // X is the owner
Y = &mut X; // Y is a mutable reference to X
Bar(Y, Y) // Error: Cannot have two mutable references to X
Bar(&mut A, &mut B) { // since A and B are mutable, they cannot alias
 A.f = ... // can mutate A
 B.f = ... // can mutate B
```

Reference Rules

- A reference cannot outlive its referent
 - No dangling references ...
- A mutable reference cannot be aliased

- Meaning:
 - There can be one mutable reference to an object in scope
 - There can be any number of immutable references
 - The owner's lifetime must contain the lifetime of all references

Example: Immutable Reference

Back to Type State

 Recall type state allowed us to express state changes in the type system:

```
// f: OpenFile
x = f.close() // x is of type ClosedFile
```

 Ownership types can be used to make objects corresponding to the previous API state inaccessible

The Type State Example

```
f = new OpenFile("foo") // f is the owner
c = f.read()
x = f.close() // x is of type ClosedFile, transfers file ownership f -> x
// then ...
d = f.read() // generates a type error: f is not the owner
```

Type State Abstraction

 Objects of a given state support only the methods available in that state

- Methods that change the internal API state
 - Return a new object of the appropriate type
 - Transfer ownership to the new object so that the references to the previous state become inaccessible
- Thus only objects of the correct type are ever accessible by the program

Discussion

- Ownership rules are very restrictive
 - Program must be linear in owned objects
 - Exactly one owner at all times
 - Implies no copies can be made
- Three techniques help:
 - Using immutable data wherever possible
 - Deep copies are OK (cloning)
 - Borrowing creates a reference that can be used
 - Does not transfer ownership
 - Implies a borrowed reference cannot deallocate an object
 - The owner cannot deallocate an object until all borrowed references are returned
 - Borrowed references have a different syntax and type

Ownership in Practice

- Ownership has been studied for > 20 years
- Rust is the first full language to support ownership types
 - The major new feature
- Experience is that Rust's ownership system helps
 - Enables manually managed memory without the bugs
 - Makes it possible to write efficient and correct code
 - Ownership types are the key
 - Which is not to say ownership is always easy to use ... it's not!