

Efficient Object Search in Game Maps

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Monash University

- Introduction
- Problem definition
- Previous work
 - Interval Heuristic Polyanya
 - Incremental Euclidean Restriction (IER)
- Grid Tree
 - Euclidean Hub Labeling
 - Grid Tree
- Results

Object Search

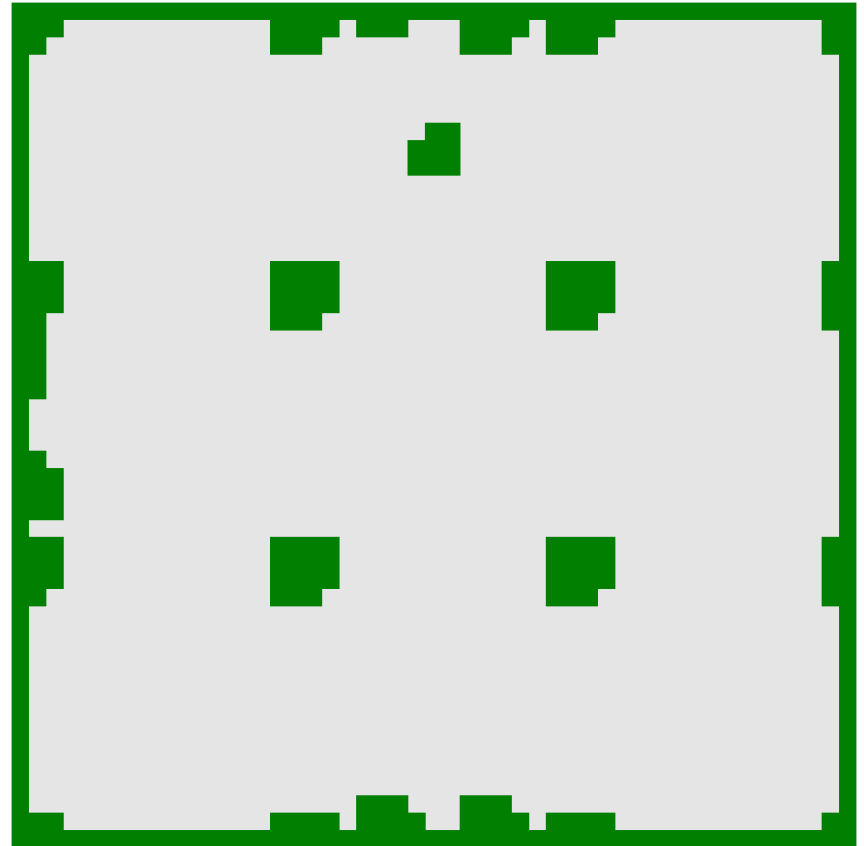


Object Search

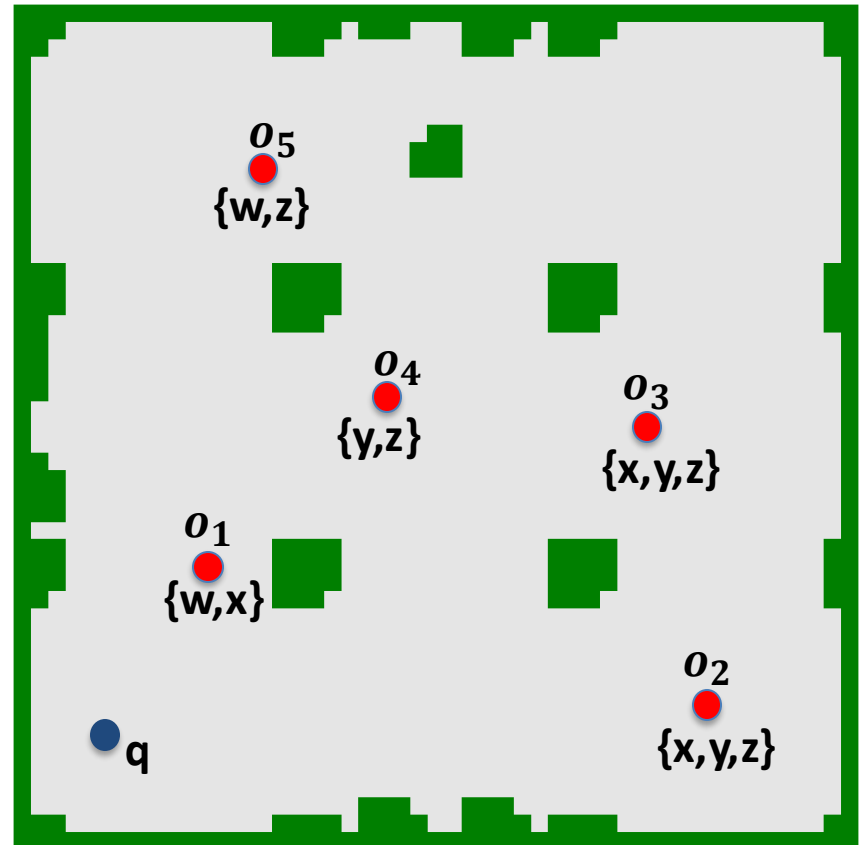


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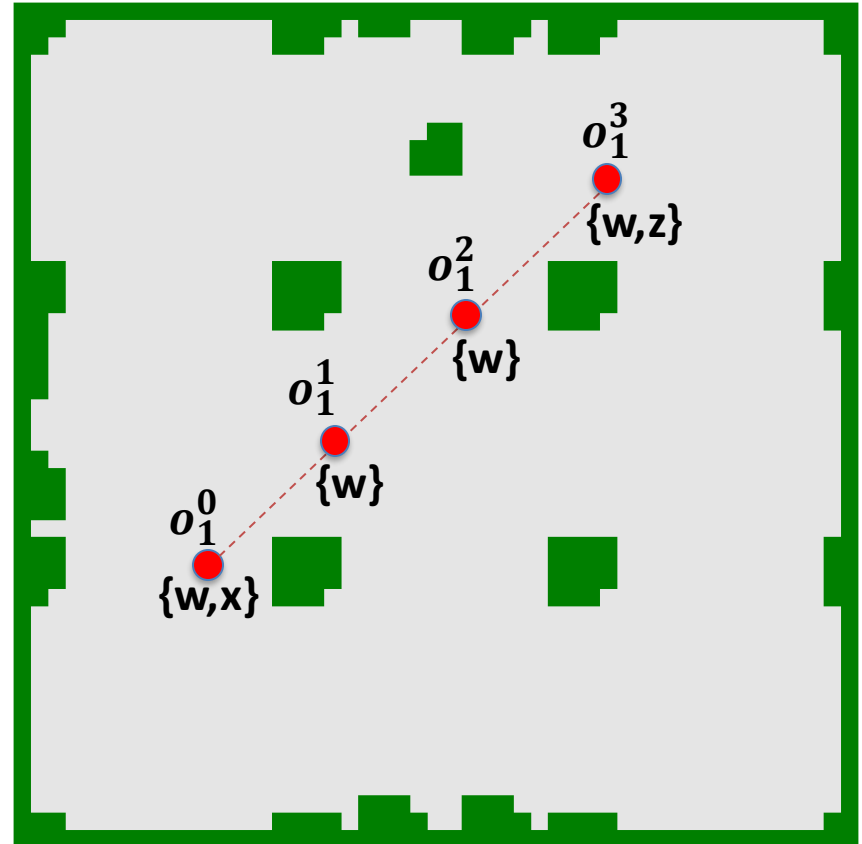
- Euclidean plane:
 - polygonal obstacles
 - a set of vertices
 - set of closed edges



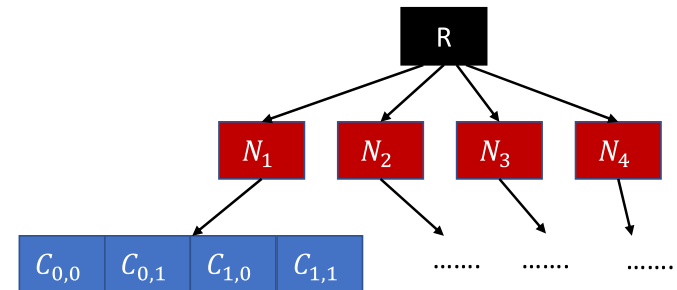
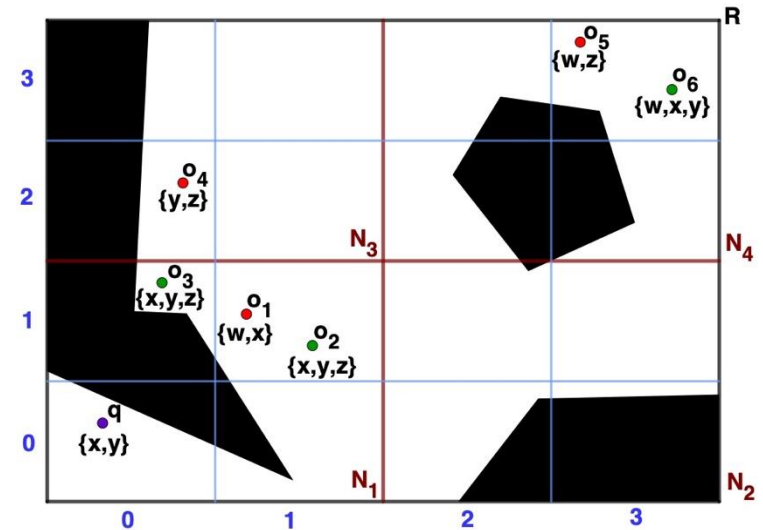
- Euclidean plane
- Objects
 - Each $o_i \in O$ is represented as a tuple $(o_i.\rho, o_i.\tau)$
 - $o_i.\rho$ is a two-dimensional point representing location
 - $o_i.\tau$ is the textual description represented as a set of keywords



- Euclidean plane
- Objects
- Timestamp
 - Discretised into a set of timestamps T
 - O^t denote the set of objects at a timestamp $t \in T$
 - $o_i^t = (o_i^t \cdot \rho, o_i^t \cdot \tau)$



- Euclidean plane
- Objects
- Timestamp
- Problem definition
 - Given a query q , issued at timestamp t and the set of objects O^t , find up to k objects closest from the query location among the objects that contain all query keywords



Object Lists

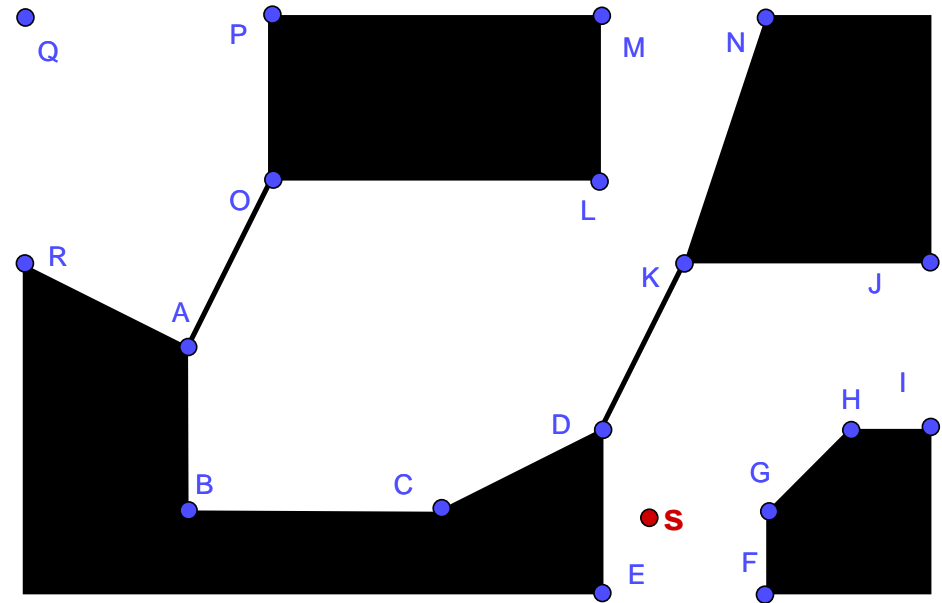
Cell	Objects
$C_{1,1}$	o_1, o_2
$C_{0,1}$	o_3
$C_{0,2}$	o_4
$C_{3,3}$	o_5, o_6

Keyword Lists of some nodes

Node	Keywords
R	w: 3, x: 4, y: 4, z: 4
N_1	w: 1, x: 3, y: 2, z: 2
$C_{1,1}$	w: 1, x: 2, y: 1, z: 1
$C_{0,1}$	x: 1, y: 1, z: 1

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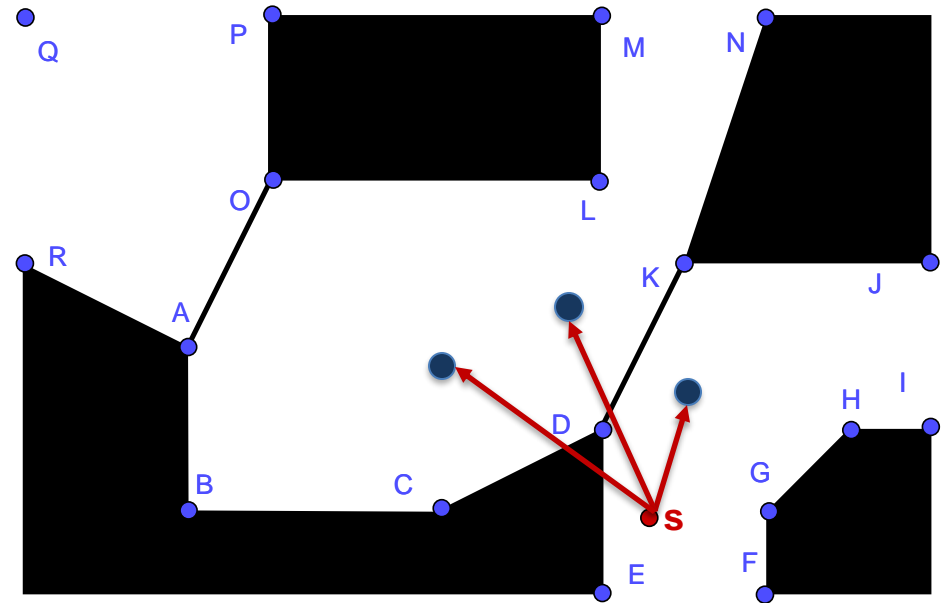
- Runs an A*-like search starting from s on navigation mesh with a set of convex polygons



Previous work

Interval Heuristic Polyanya

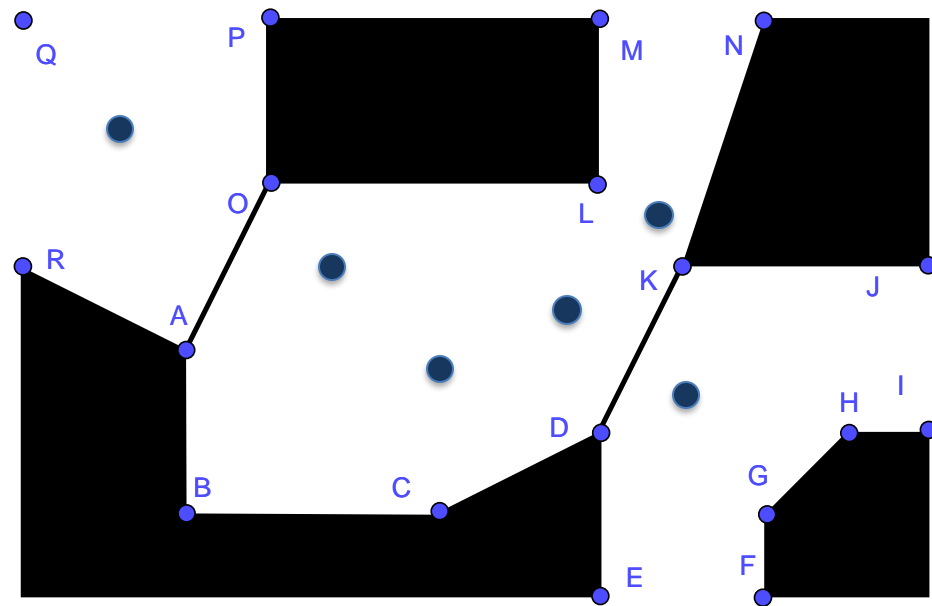
- Runs an A*-like search starting from s on navigation mesh with a set of convex polygons
- When the search reaches a polygon that contains an object, the object matching the keyword is added to the queue



Previous work

Interval Heuristic Polyanya

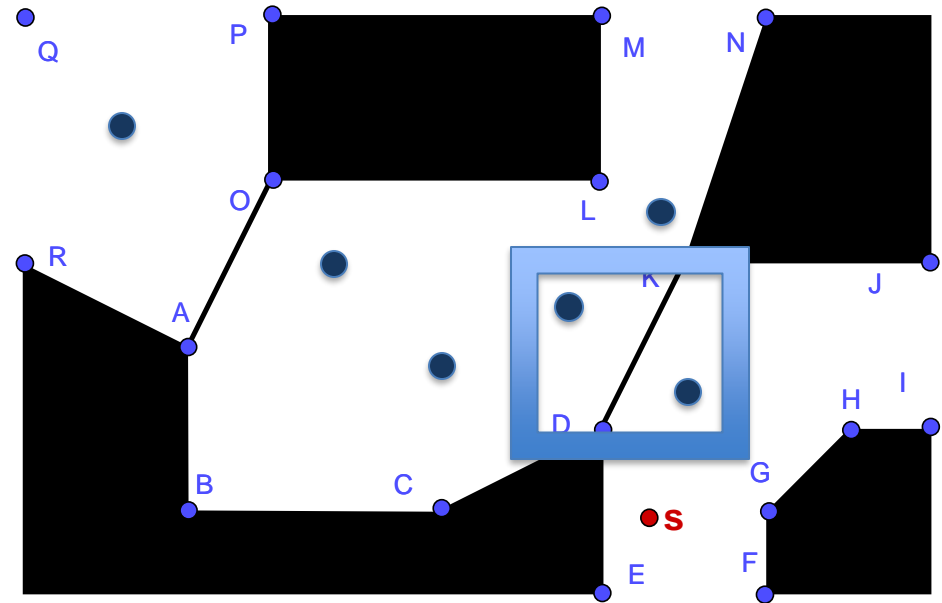
- Runs an A*-like search starting from s on navigation mesh with a set of convex polygons
- When the search reaches a polygon that contains an object, the object matching the keyword is added to the queue
- Objects are stored in navigation mesh, handling object updates is quite efficient



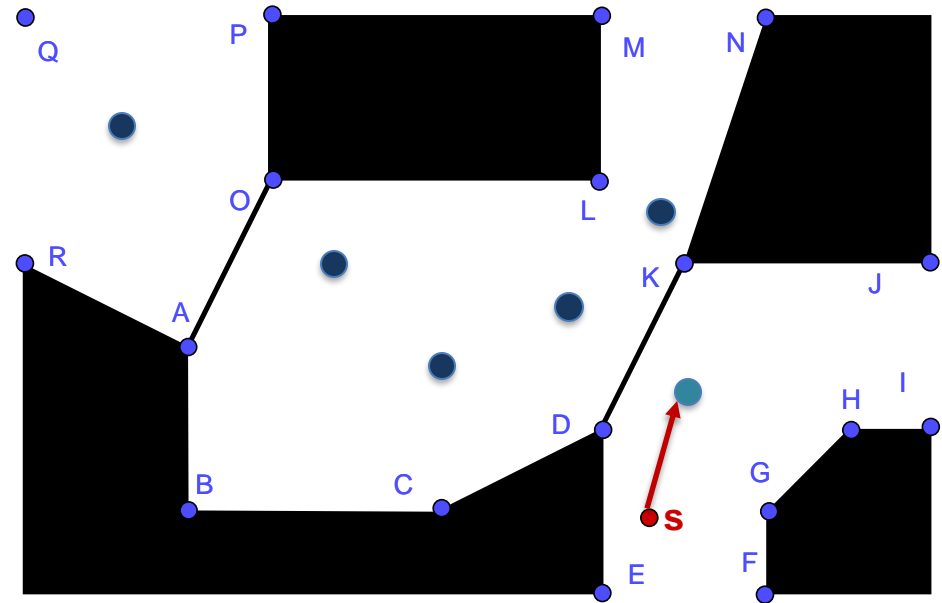
Previous work

IER [2]

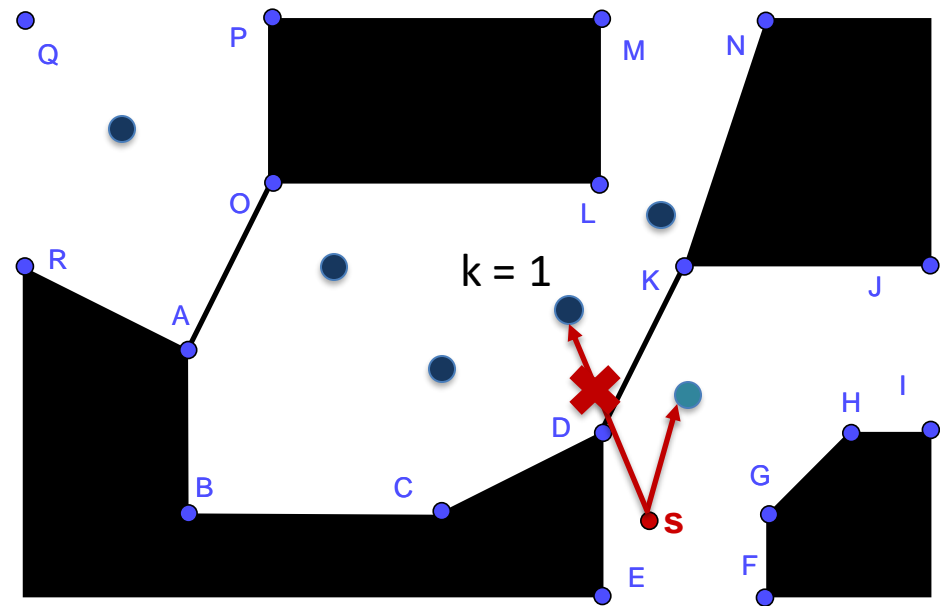
- Employs IR-tree to store the objects and incrementally retrieves nearest objects to the query location



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- Employs IR-tree to store the objects and incrementally retrieves nearest objects to the query location
- For each retrieved object, it calls Polyanya to compute its actual distance from the query
- Terminates when the Euclidean distance of next retrieved object is no smaller than the actual distances of kNNs



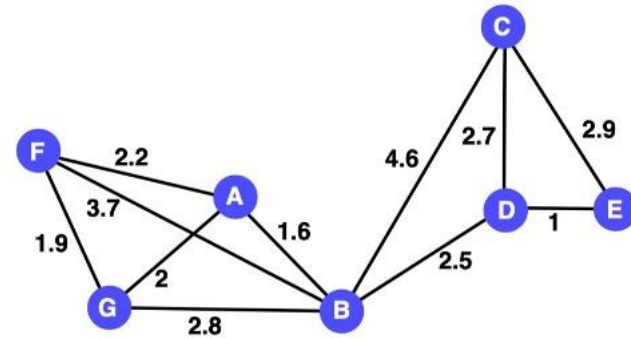
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- **Grid Tree**
 - Euclidean Hub Labeling
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- State-of-the-art Euclidean shortest pathfinding algorithm

Grid Tree

Euclidean Hub Labeling

- State-of-the-art Euclidean shortest pathfinding algorithm
- Employs the popular hub labeling technique [4]



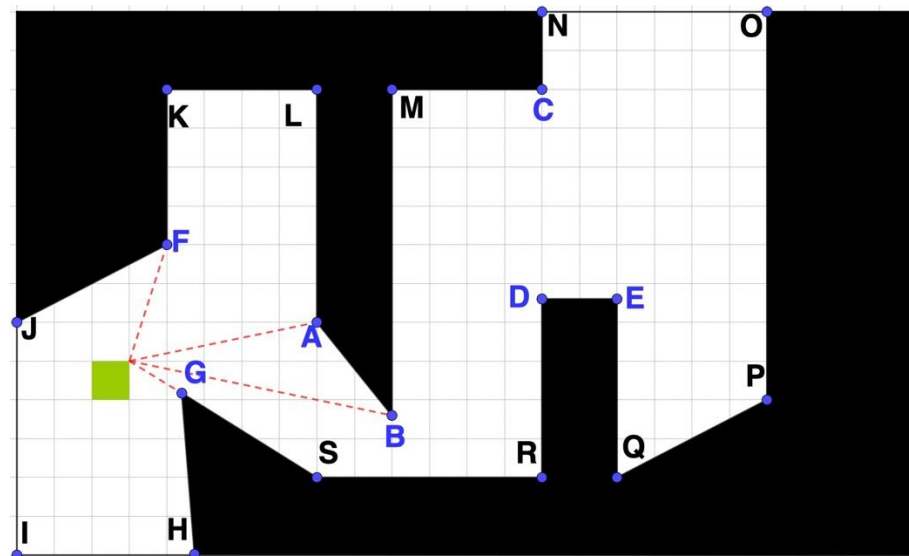
Vertex	Hub labels
A	(A, 0)
B	(A, 1.6), (B, 0)
C	(A, 6.2), (B, 4.6), (C, 0)
D	(A, 4.1), (B, 2.5), (C, 2.7), (D, 0)
E	(A, 5.1), (B, 3.5), (C, 2.9), (D, 1), (E, 0)
F	(A, 2.2), (B, 3.7), (F, 0)
G	(A, 2), (B, 2.8), (F, 1.9), (G, 0)

Table 1: Hub labels constructed for graph above

Grid Tree

Euclidean Hub Labeling

- State-of-the-art Euclidean shortest pathfinding algorithm
- Employs the popular hub labeling technique
- Superimpose uniform grids across a given map, and store via labels in the grids to allow fast lookup and retrieval during query phase



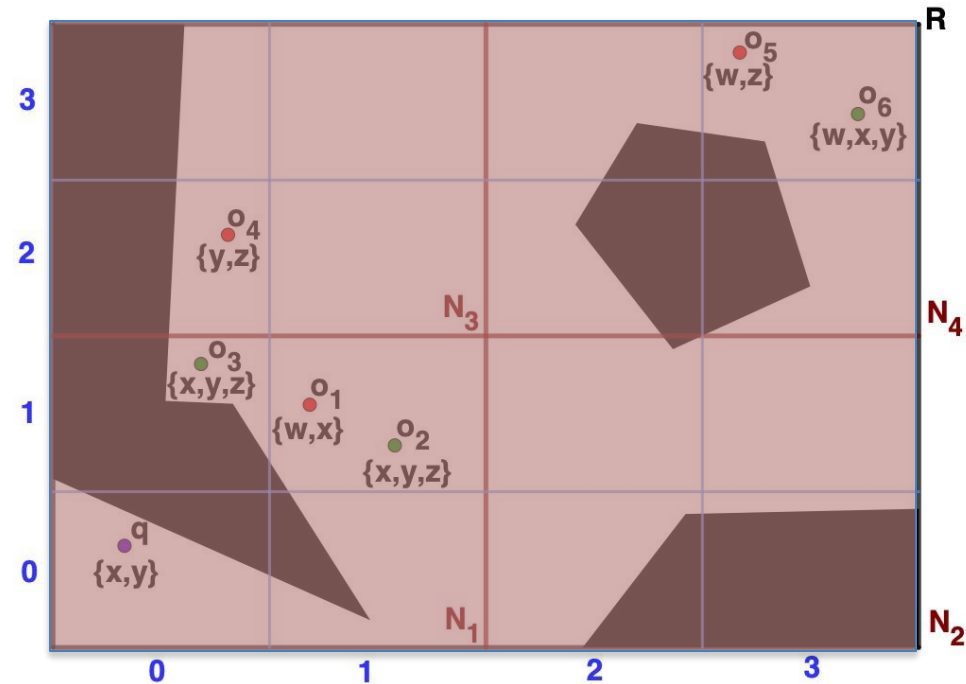
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E	(A, 5.1), (B, 3.5), (C, 2.9), (D, 1), (E, 0)
F	(A, 2.2), (B, 3.7), (F, 0)
G	(A, 2), (B, 2.8), (F, 1.9), (G, 0)

Hub nodes	Via labels
A	(A, 0), (B, 1.6), (G, 2), (F, 2.2)
B	(B, 0), (G, 2.8), (F, 3.7)
F	(F, 0), (G, 1.9)
G	(G, 0)

Table 2: Hub labels and via labels for the figure above

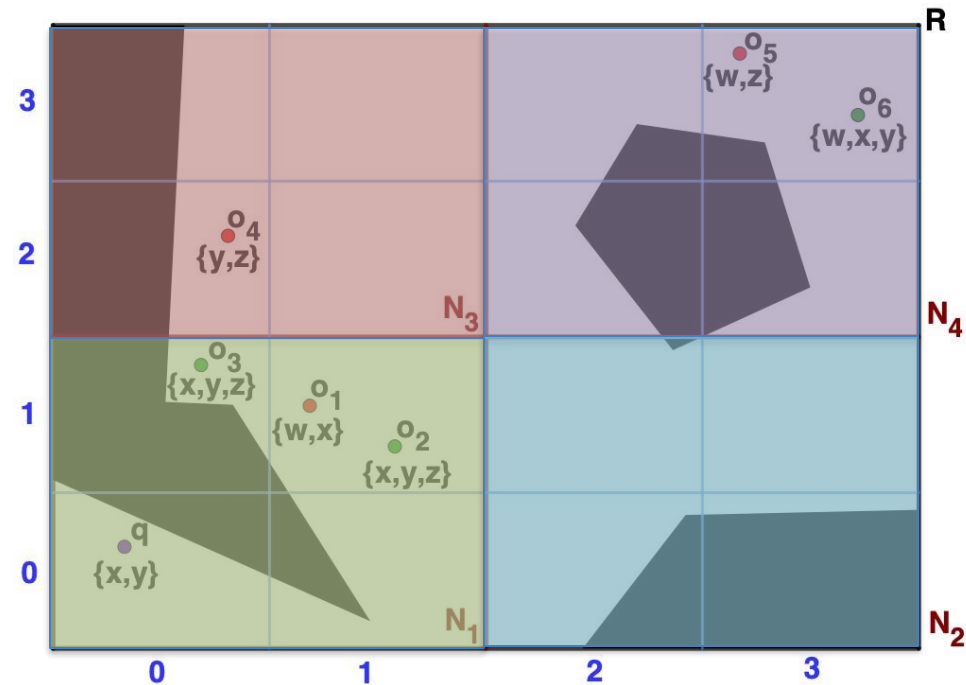
Grid Tree Structure

- Starting from root node, recursively divides each node into 4 equal sized children until size of each child node is smaller than a threshold



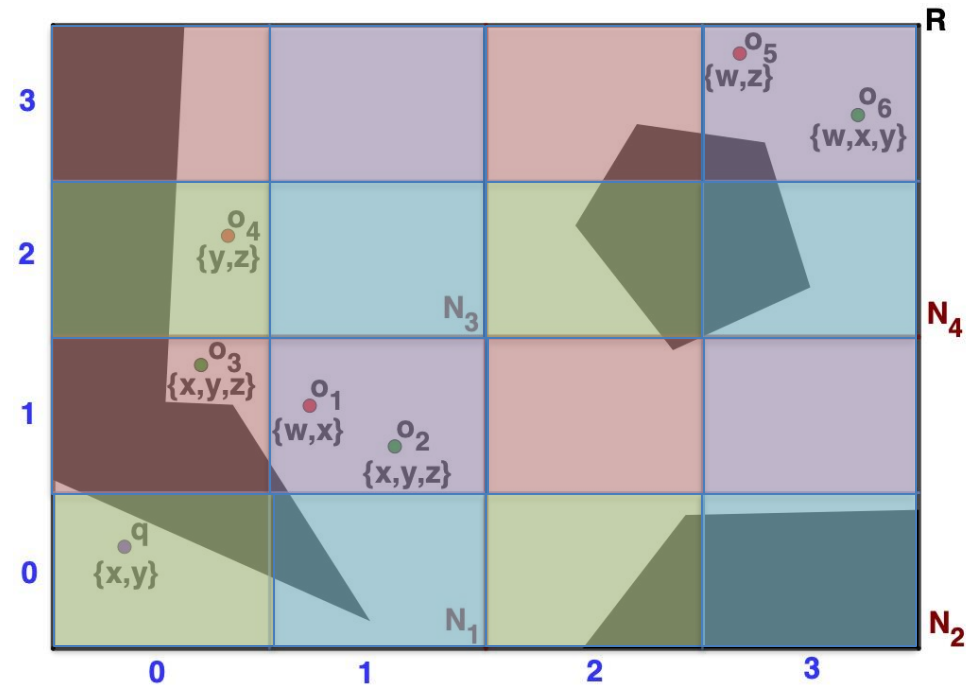
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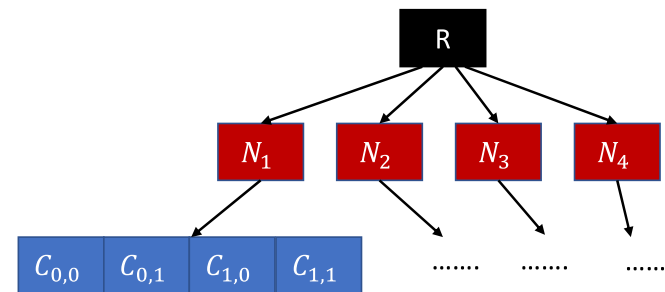
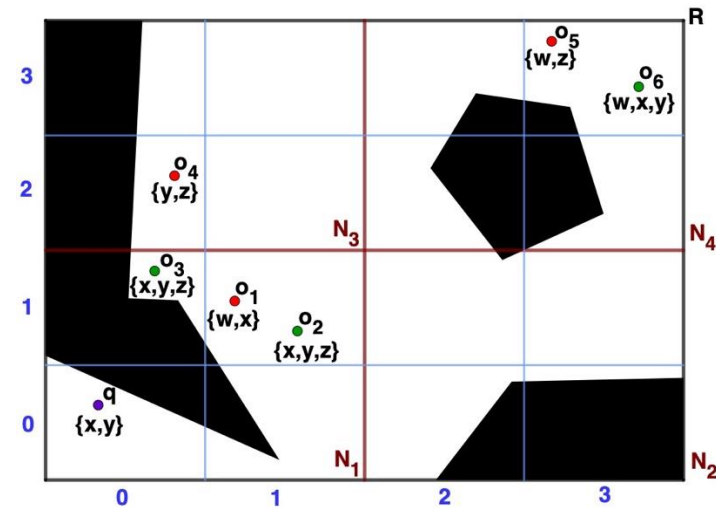
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Grid Tree Structure

- Starting from root node, recursively divides each node into 4 equal sized children until size of each child node is smaller than a threshold
- Each node stores an object list containing the IDs of the objects that are located inside it. It also stores a keyword list containing unique keywords with its frequency

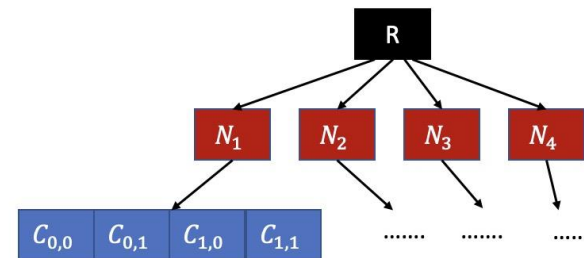
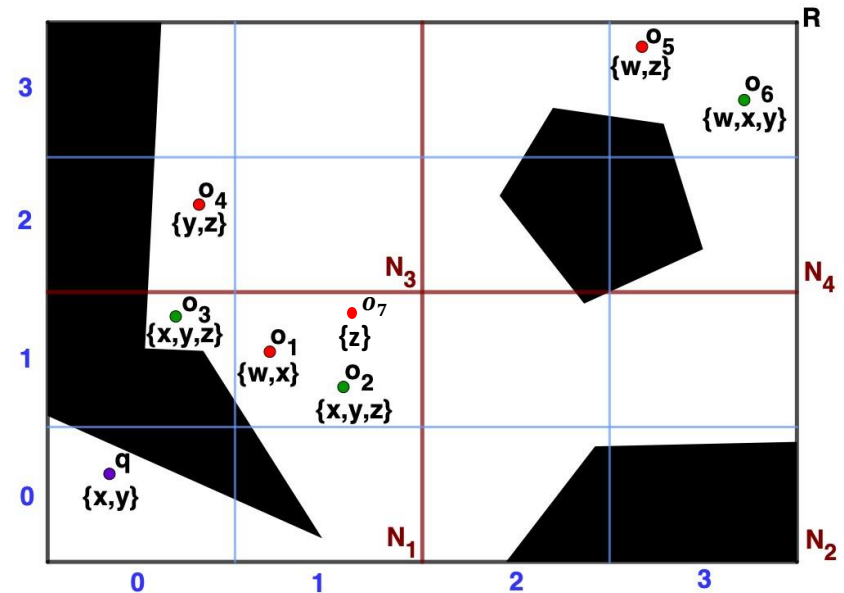


Object Lists	
Cell	Objects
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$C_{0,1}$	o_3
$C_{0,2}$	o_4
$C_{3,3}$	o_5, o_6

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Grid Tree Update

■ Insert



Object Lists

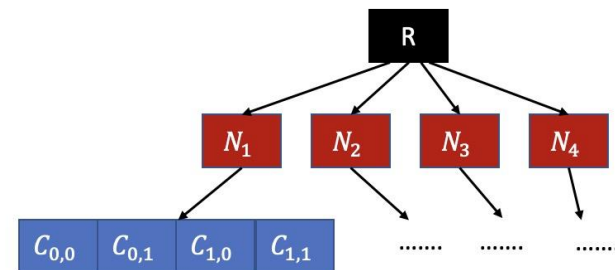
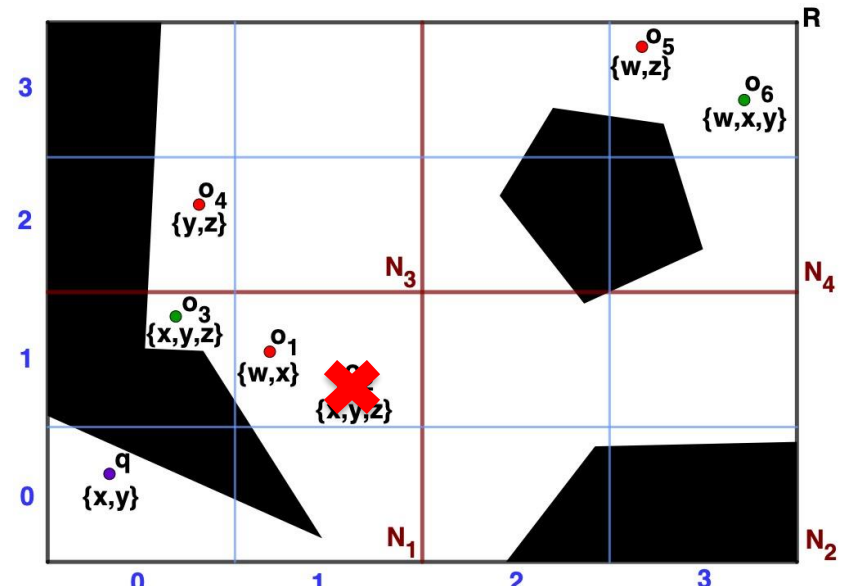
Cell	Objects
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Grid Tree Update

- Insert
- Delete



Object Lists

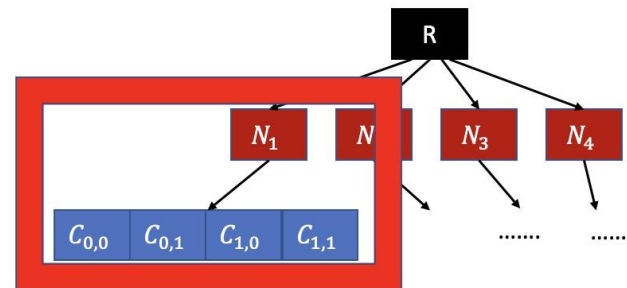
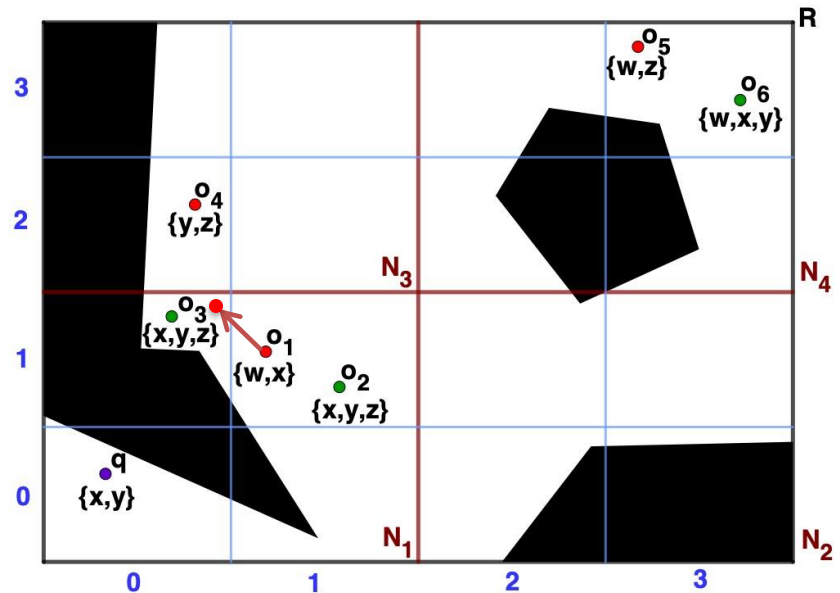
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N_1	w: 1, x: 3, y: 2, z: 2
$C_{1,1}$	w: 1, x: 1
$C_{0,1}$	x: 1, y: 1, z: 1

Grid Tree Update

- Insert
- Delete
- Location change of an object



Object Lists

Cell	Objects
$C_{1,1}$	o_2
$C_{0,1}$	o_3, o_1
$C_{0,2}$	o_4
$C_{3,3}$	o_5, o_6

Keyword Lists of some nodes

Node	Keywords
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- Query is similar to the standard best-first search algorithm on trees

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- The grid tree is traversed based on the minimum distances of the nodes and the nodes that do not contain query keywords are removed

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- The grid tree is traversed based on the minimum distances of the nodes and the nodes that do not contain query keywords are removed
- Euclidean Hub Labeling (EHL) is used to compute shortest distance from the query to the objects containing query keywords

Grid Tree

Query algorithm

Algorithm 1: Boolean k NN query processing

Input: $q.\rho, q.\tau, k$: query location, query keywords and k

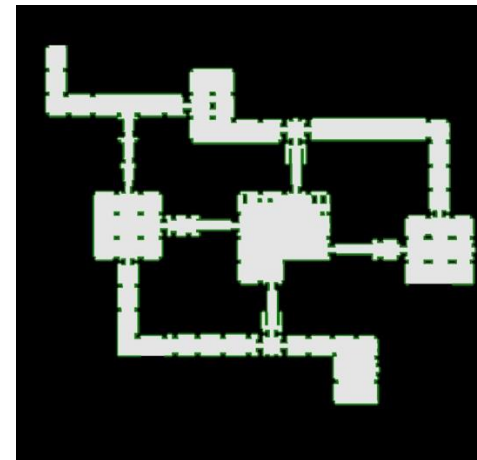
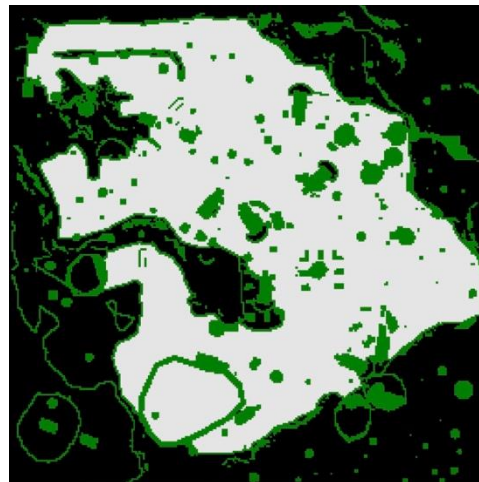
Output: R : query results

```
1  $R = \phi; d^k = \infty;$ 
2 Initialise a min-heap  $H$  with the root node of Grid Tree ;
3 while  $H \neq \phi$  do
4   | deheap an entry  $e$  from  $H$  ;
5   | if  $e.key \geq d^k$  then
6   |   | return  $R$ ;
7   | if  $e$  is an object then
8   |   | compute  $d(q.\rho, e.\rho)$ ;
9   |   | if  $d(q.\rho, e.\rho) < d^k$  then
10  |   |   | update  $R$  and  $d^k$  by object  $e$ ;
11  | else if  $e$  is a leaf node then
12  |   | for each object  $o_i^t$  in the object list of  $e$  do
13  |   |   | if  $q.\tau \subseteq o_i^t.\tau$  then
14  |   |   |   | insert  $o_i^t$  in  $H$  with key
15  |   |   |   |    $mindist(q.\rho, o_i^t.\rho)$ ;
16  | else
17  |   | for each child node  $c$  of  $e$  do
18  |   |   | if  $c$  contains all query keywords  $q.\tau$  then
19  |   |   |   | insert  $c$  in  $H$  with key  $mindist(q.\rho, c)$ ;
20 return  $R$ ;
```

- Generalisation of boolean k NN query
- Top- k spatial keyword query
- Keyword range query

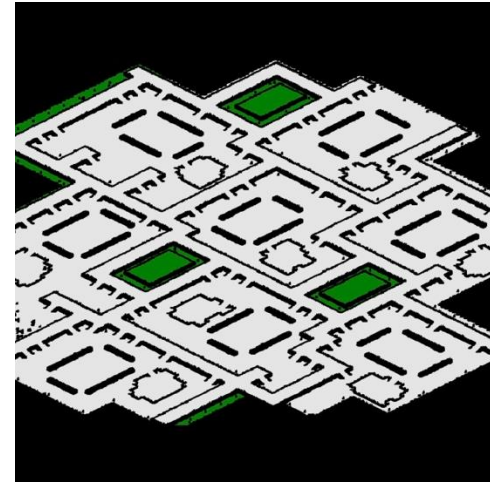
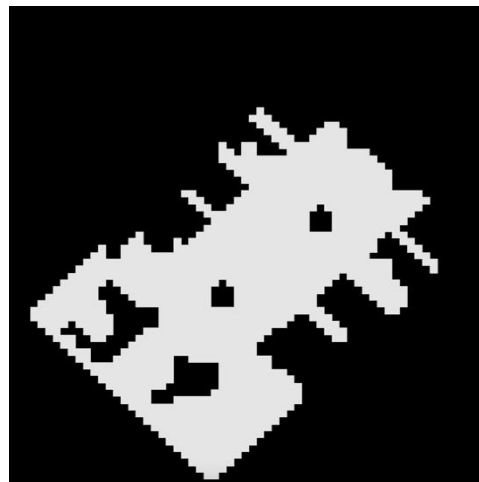
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- Benchmark maps: widely used game map benchmarks [5]



Game	#Maps	# Cells	# Trav. Cells	# Vertices
DA	67	151,420	15,911	1182.9
DAO	156	134,258	21,322	1727.6
BG	75	262,144	73,930	1294.4
SC	75	446,737	263,782	11487.5

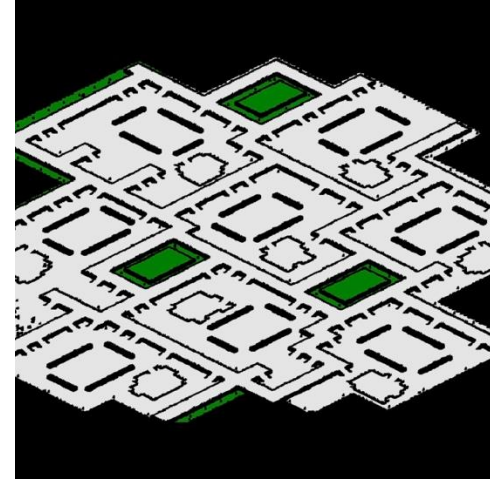
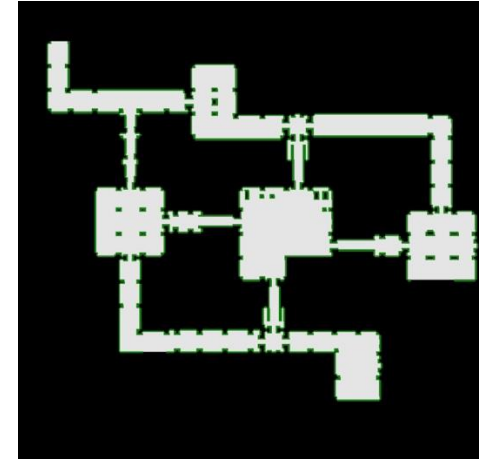
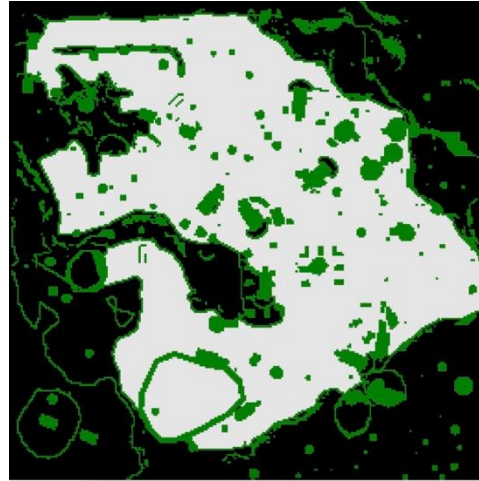
Table 1: Total number of maps, and average number of total cells, traversable cells and vertices in each benchmark.



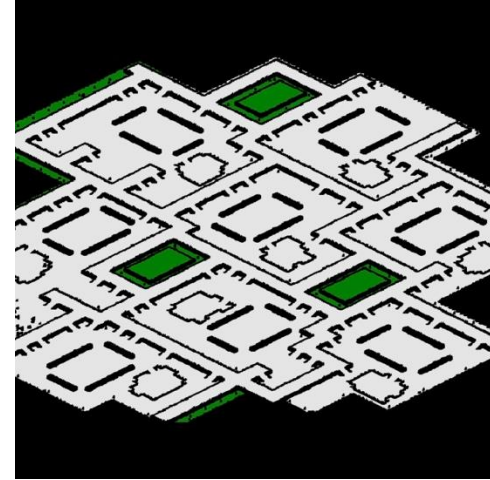
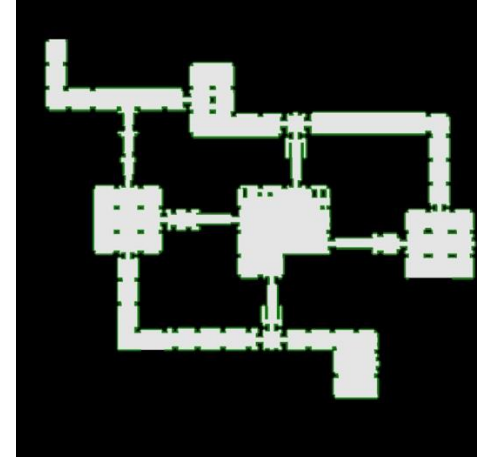
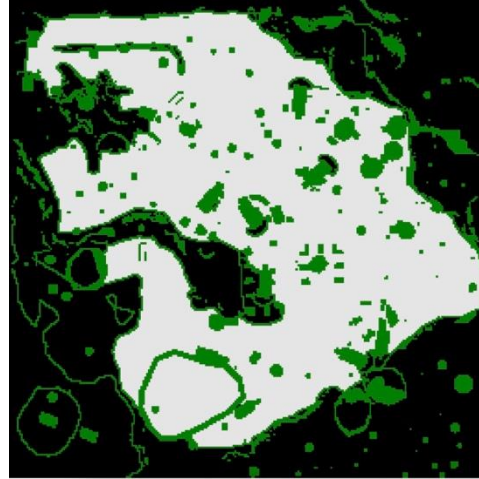
Results

Experimental setup

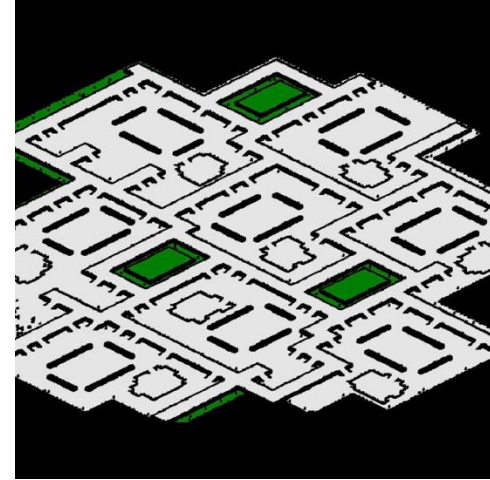
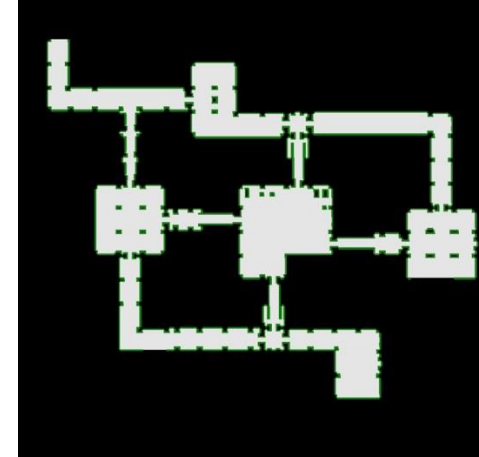
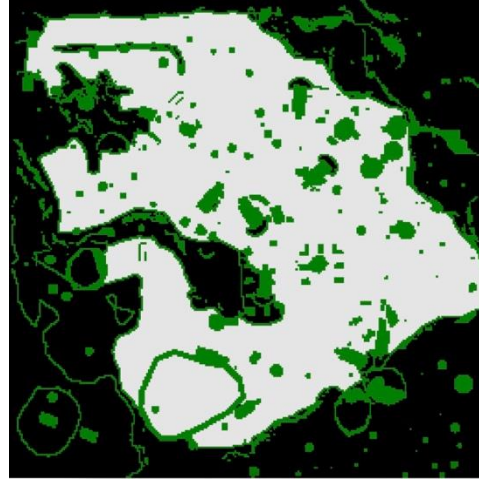
- Benchmark maps: widely used game map benchmarks [5]
- Object density: 0.1%, 1% and 10%



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- Keyword: 100 item descriptions obtained from ChatGPT for each benchmark and extracted keywords from it

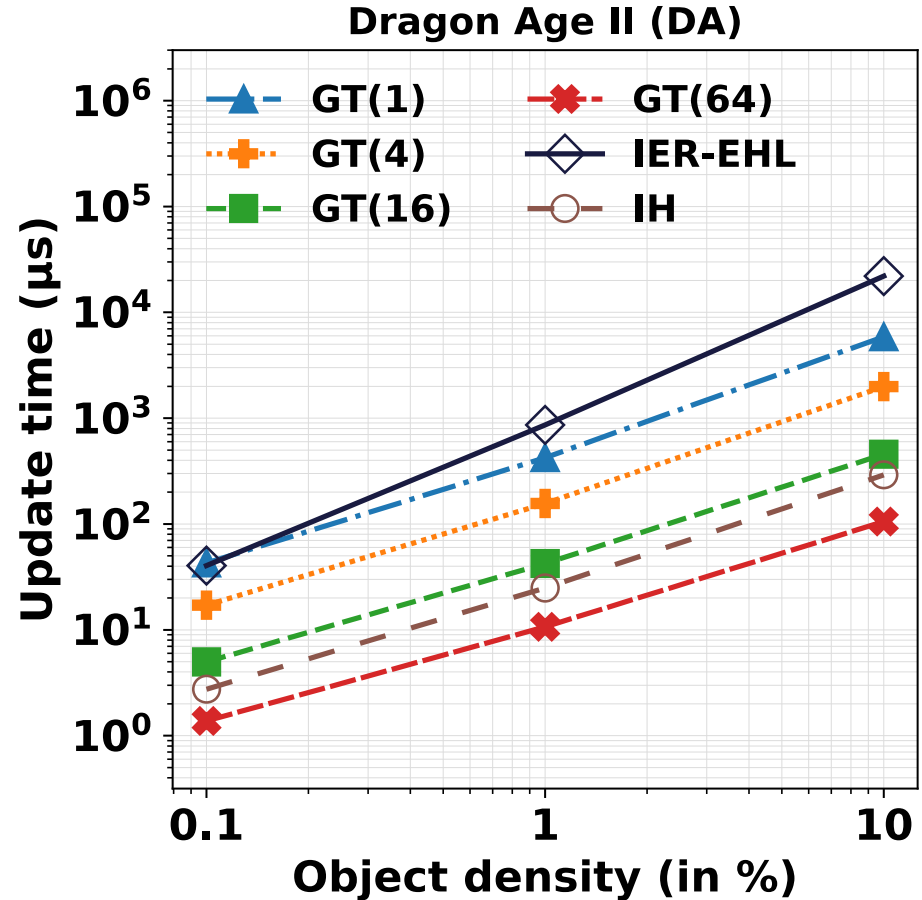
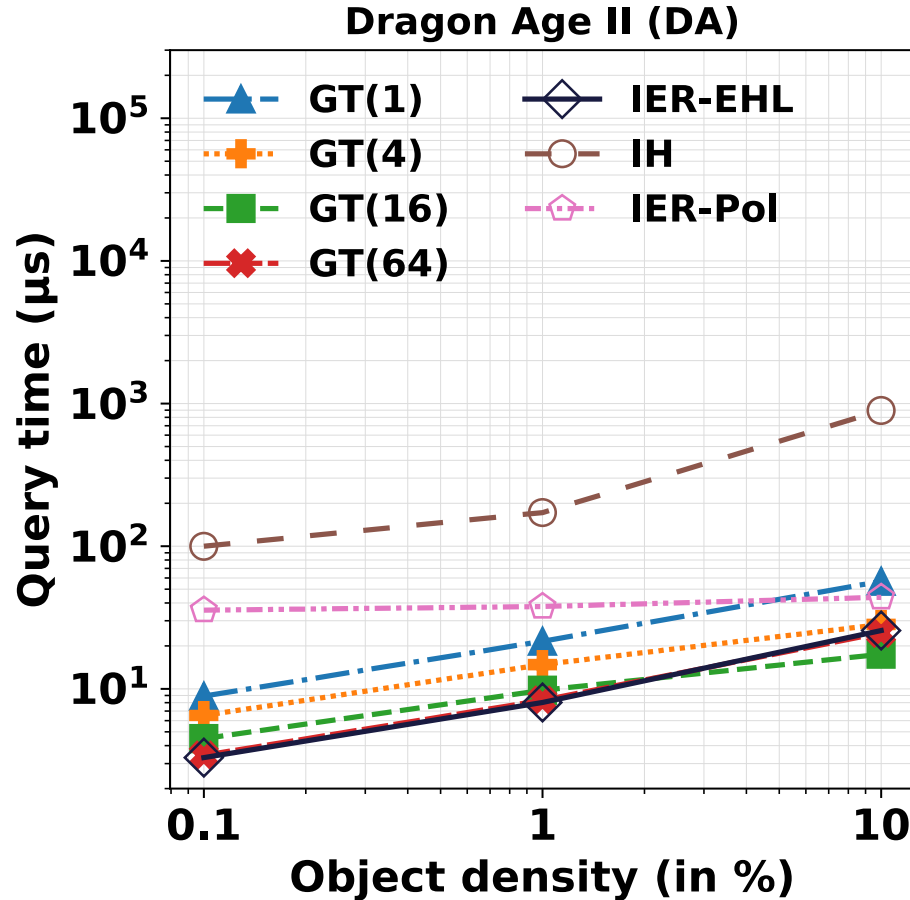


- Benchmark maps: widely used game map benchmarks [5]
- Object density: 0.1%, 1% and 10%
- Keyword: 100 item descriptions obtained from ChatGPT for each benchmark and extracted keywords from it
- Query: randomly generate 100 queries per timestamp for 50 timestamps



Results

Effect of object density



Effect of object density on query time and update time for DA on default settings ($k=3$, mobility = 70%, # of query keywords = 2)

- Mobility of objects
- k
- Number of query keywords
- Object distribution

Thank you for listening 😊

- [1] Zhao, S., Taniar, D., & Harabor, D. (2018). Fast k-nearest neighbor on a navigation mesh. In *Proceedings of the International Symposium on Combinatorial Search* (Vol. 9, No. 1, pp. 124-131).
- [2] Zhao, S., Harabor, D. D., & Taniar, D. (2018). Faster and more robust mesh-based algorithms for obstacle k-nearest neighbour. *arXiv preprint arXiv:1808.04043*.
- [3] Du, J., Shen, B., & Cheema, M. A. (2023). Ultrafast Euclidean Shortest Path Computation Using HubLabeling. *Proceedings of the AAAI Conference on Artificial Intelligence*, 37(10), 12417-12426. <https://doi.org/10.1609/aaai.v37i10.26463>
- [4] Abraham, I., Delling, D., Goldberg, A. V., & Wernke, R. F. (2011). A hub-based labeling algorithm for shortest paths in road networks. In *Experimental Algorithms: 10th International Symposium, SEA 2011, Kolimpari, Chania, Crete, Greece, May 5-7, 2011. Proceedings 10* (pp. 230-241). Springer Berlin Heidelberg.

Image sources:

<https://www.geospatialworld.net/news/sensewheres-indoor-location-volume-hits-50-billion-requests/>

<https://www.mecalux.com/blog/fully-automated-warehouse>

<https://www.theverge.com/2022/2/22/22945814/amazon-astro-home-robot-photo-video>