

Huffman Coding: File Compression using Greedy Algorithm

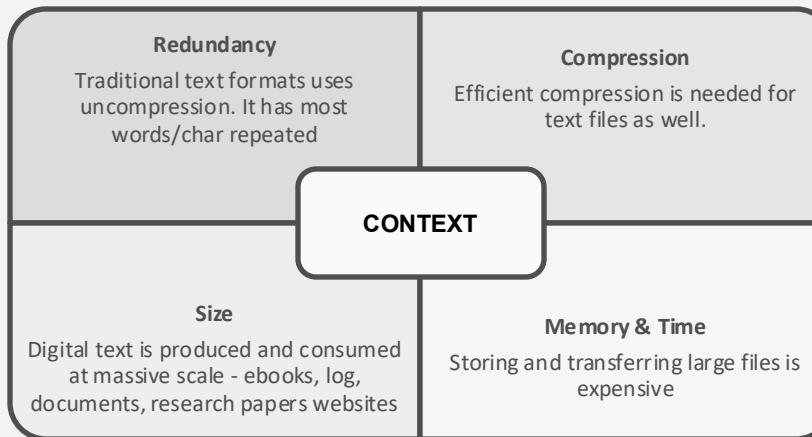
Northeastern University | CS5800 Algorithms | Prof. Aida Sharif Rohani

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Introduction

Key Question.

How can we design a lossless compression system
that significantly reduces text file size while enabling fast, real-time
decompression for reading?



Huffman tool

Reduce storage size without losing information.



Reader Application

Decode compressed data efficiently enough to display pages of a book on demand.



Application: A Fast, Chunk-Based Book Reader
The algorithm powers a GUI that decompresses and displays large books page-by-page for efficient reading.

Encodings



File Type	Extensions	Existing Compression / Encoding Method	Is File Already Compressed?	Expected Huffman Compression Ratio	Explanation
Plain Text	.txt, .log, .csv, .json, .xml	None	✗ No	40% – 80% smaller	Text has skewed character frequency; Huffman is ideal.
Documents (Office XML)	.docx, .xlsx, .pptx	ZIP (DEFLATE: LZ77 + Huffman)	✓ Already compressed	File becomes 20–200% larger	Office files are ZIP containers; content already compressed using Huffman + LZ.
PDF	.pdf	Flate/DEFLATE, LZW, JPEG, JP2	✓ Already compressed	+20% to +300% expansion	PDF streams already use entropy coding; random-like distribution.
Images (raw)	.bmp, .ppm, .tiff (uncompressed)	None (sometimes RLE for TIFF)	✗ No	30% – 70% smaller	Raw pixel data compresses fairly well.
Images (compressed)	.jpg, .jpeg	DCT + Quantization + Huffman	✓ Yes	Huge expansion: 300% – 800%	JPEG already uses Huffman coding inside. Compressing again makes it worse.
Images (compressed)	.png	DEFLATE (LZ77 + Huffman)	✓ Yes	Very large expansion: 200% – 600%	PNG uses entropy coding and filters; nearly incompressible.
Audio (raw)	.wav, .pcm, .aiff	None	✗ No	10% – 40% smaller	Raw amplitude distributions slightly skewed; small gains.
Audio (compressed)	.mp3, .aac, .flac	MP3: MDCT + Huffman / AAC: Huffman / FLAC: Rice/Huffman	✓ Yes	Massive expansion: 300% – 1000%	Audio codecs already use Huffman coding internally.
Video (raw)	.yuv	None	✗ No	10% – 30% smaller	Pixel values partly skewed; limited improvement.
Video (compressed)	.mp4, .mov, .mkv, .avi	H.264/HEVC/AV1 (CABAC, CAVLC, entropy coding)	✓ Yes	Very large expansion: 200% – 800%	These codecs use advanced entropy coding more efficient than Huffman.
Python/Source Code	.py, .java, .cpp, .html, .css, .js	None	✗ No	30% – 70% smaller	High redundancy and repeated keywords; good for Huffman.
Binary Executables	.exe, .dll, .bin	Often packed or randomized	⚠ Sometimes	Likely expansion: 20% – 500%	Binaries include many random bytes or pre-packed segments.
Archives	.zip, .7z, .rar, .gz, .whl	DEFLATE, LZMA, PPMD	✓ Fully compressed	Always expands	These formats already use Huffman, arithmetic coding, or LZ — cannot compress again.

Rationale

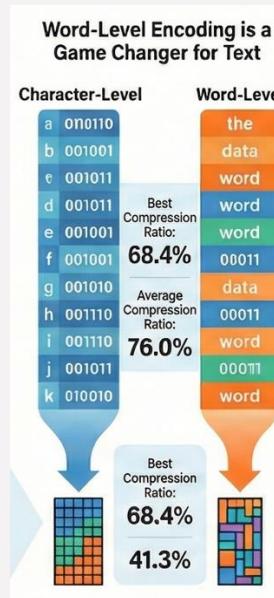
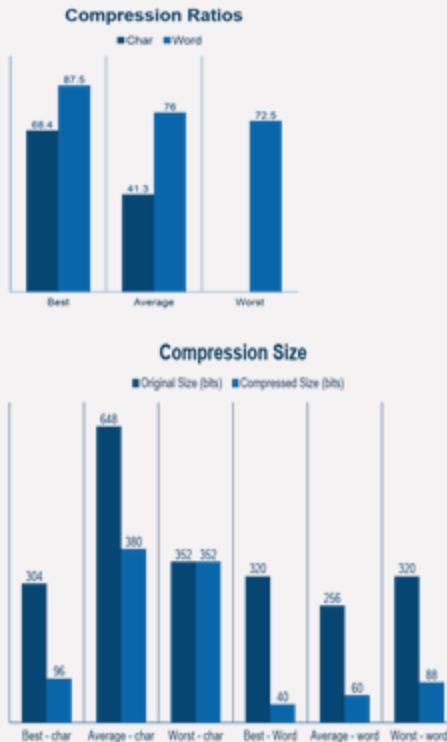


- Huffman Coding is a **classic Greedy Algorithm** that produces an **optimal prefix-free encoding** based on symbol frequencies.
- It is widely used in real systems (ZIP, JPEG, MP3), but classroom examples rarely show practical applications.
- This project extends the Huffman algorithm beyond theory, building a **complete working system**:
 - A compression tool
 - A decompression module
 - A chunk-based paging system
 - A Python Tkinter based book reader

Why this matters:

- Demonstrates how algorithmic theory can be used to build a real application.
- Highlights trade-offs in designing usable compression systems (speed, memory, chunking, file formats).

Analysis - Word level?



- Character-level Huffman coding compresses but leaves redundancy:
"the", "and", "to", "of" appear thousands of times.
- Word-level encoding:
 - Dramatically reduces redundancy
 - Produces shorter average codewords
 - Improves readability after decoding
 - Works perfectly for book-style text
- Compression improvement observed:
- Word-level coding achieved up to 70% space savings vs raw .txt.

Huffman.py

Q

```
compress_file(file_path, mode)

Steps:
    └── Read input file
    └── Instantiate HuffmanCoding() object
    └── Call → encode(text, code_output_path, mode)
    └── write encoded data → compress_<file>.huff
    └── write freq table → compress_<file>.huff.freq.json
    └── visualize_tree(root) ← optional
        |   └-- Graphviz Digraph OR Custom Matplotlib
        |   └-- save PNG image
    └── return success
```

```
encode()
    └── build_frequency_table()
    └── build_huffman_tree()
        |   └-- create Node for each symbol
        |   └-- heapq.heappush()
        |   └-- heapq.heappop()
        |   └-- merge two smallest nodes
        |   └-- return root
    └── generate_codes(root)
        |   └-- DFS traversal
        |   └-- prefix + '0' for left
        |   └-- prefix + '1' for right
    └── calculate_compressed_bits
    └── write_codebook_summary_file
    └── return (encoded_bitstring, freq_table)
```

Huffman tool



Char level encoding

```
: !python code/huffman_tool.py sample_char.txt compress char
[INFO] Huffman tree saved as /Users/bhalchandra/SEM1_NEU/huffman_tool/data/../output/sample_char.txt_tree.png.png
[INFO] Compression complete. Compressed file saved to: /Users/bhalchandra/SEM1_NEU/huffman_tool/data/../output/sample_char.txt.huff

: from collections import deque
print(''.join(deque(open('/Users/bhalchandra/SEM1_NEU/huffman_tool/data/sample_char.txt_codes.txt'), 2)))
Actual compression ratio (file sizes): 44.65%
Original size: 422448 bytes, Compressed size: 233833 bytes
```

Decoding Char level encoded .huff file

```
: !python ../code/huffman_tool.py sample_char.txt.huff decompress char
[INFO] Decompression complete using char-level Huffman decoding.
[INFO] Output: /Users/bhalchandra/SEM1_NEU/huffman_tool/output/uncompressed_sample_char.txt.txt
```

Word level encoding

```
: !python code/huffman_tool.py sample_word.txt compress word
[INFO] Huffman tree saved as /Users/bhalchandra/SEM1_NEU/huffman_tool/data/../output/sample_word.txt_tree.png
[INFO] Compression complete. Compressed file saved to: /Users/bhalchandra/SEM1_NEU/huffman_tool/data/../output/sample_word.txt.huff

: from collections import deque
print(''.join(deque(open('/Users/bhalchandra/SEM1_NEU/huffman_tool/data/sample_word.txt_codes.txt'), 2)))
Actual compression ratio (file sizes): 84.67%
Original size: 422448 bytes, Compressed size: 64753 bytes
```

Decoding word level encoded .huff file

```
: !python ../code/huffman_tool.py sample_word.txt.huff decompress word
[INFO] Decompression complete using word-level Huffman decoding.
[INFO] Output: /Users/bhalchandra/SEM1_NEU/huffman_tool/output/uncompressed_sample_word.txt.txt
```

sample_char.txt	422 KB	Plain Text
sample_char.txt_codes.txt	3 KB	Plain Text

sample_char.txt.huff	234 KB	Document
sample_char.txt.huff.freq.json	18 bytes	Plain Text
uncompressed_sample_char.txt.txt	422 KB	Plain Text

sample_word.txt	422 KB	Plain Text
sample_word.txt_codes.txt	14 KB	Plain Text

sample_word.txt.huff	65 KB	Document
sample_word.txt.huff.freq.json	5 KB	Plain Text
uncompressed_sample_word.txt.txt	421 KB	Plain Text

Reader - pseudocode



```
CHUNKED-HUFFMAN-COMPRESS(Text, PAGE_SIZE)
    words = split Text into tokens
    chunks = group words into PAGE_SIZE blocks
    outputBinaries = empty
    metadata = empty list

    for each chunk:
        freq = frequency count(chunk)
        tree = build Huffman tree(freq)
        codes = generate codes(tree)

        encodedBits = concat(codes[word] for word in chunk)
        packedBytes, padding = pack_into_bytes(encodedBits)

        append packedBytes to output file
        store in metadata: offset, length, padding, freq

    write metadata.json
```

```
CHUNKED-HUFFMAN-DECODE(file.huff, metadata)
    for each chunkInfo in metadata:
        bytes = read chunk from file
        bitstring = unpack_bytes(bytes, chunkInfo.padding)

        tree = rebuild Huffman tree(chunkInfo.freq)
        decodedWords += decode_bits_using_tree(bitstring, tree)

    return decodedWords
```

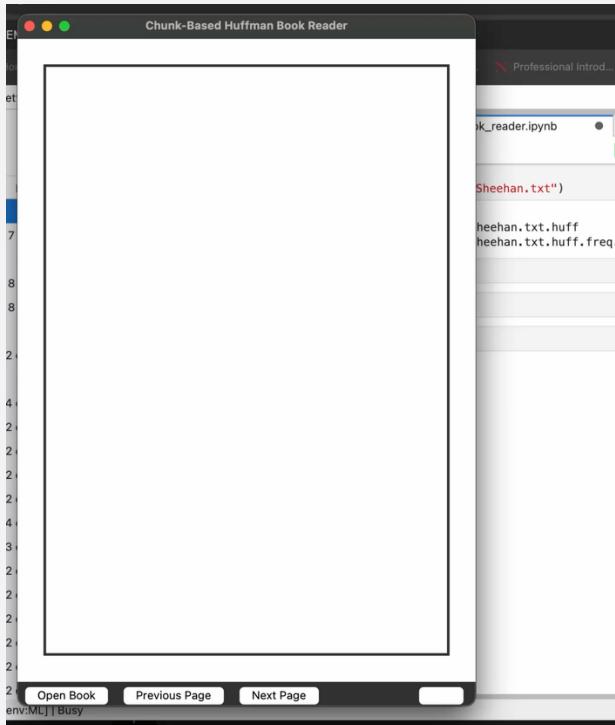
Inputs

- book.txt — raw book
- book.txt.huff — compressed binary
- book.txt.huff.freq.json — chunk metadata
- Chunk sizes, padding values, frequency tables

Outputs

- Decoded individual pages
- Page count
- Compression ratio
- Memory-efficient reading of entire text

Reader application



▶ Python Tkinter GUI

▶ Paper like reading experience

▶ Chunks of 250 words decoded per page

▶ Page Navigation controls

▶ Load any book with .huff and .json metadata file

	Hints on news reporting by Murray Sheehan		86 KB	Plain Text
	Hints on news reporting by Murray Sheehan.txt.huff		15 KB	Document

Observations

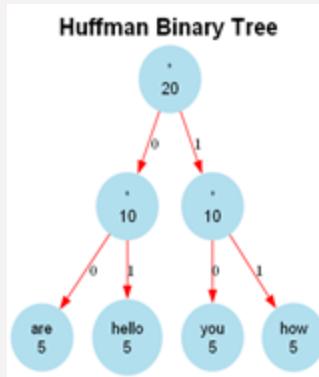


- More repetition → better compression
- Huffman coding efficiency depends on frequency distribution
- Worst-case for char: no gain, overhead may slightly increase size
- Visual tree helps understand code assignment
- Huffman Coding is not just theoretical, it powers real systems.
- This project demonstrates complete integration from **algorithm** → **compression** → **metadata** → **decoding** → **UI**.
- Shows how algorithm design, data structures and UX can combine to produce a working, practical application.

Best case Example: "hello how are you
hello how are you hello how are you hello
how are you hello how are you"

Best case Example: "Huffman coding is a data compression algorithm.
Huffman coding assigns shorter codes to more frequent data."

Worst case Example: "this is a worst case example for huffman coding algorithm where no any word being repeated, its used only ones"



Huffman Encoding Report (Word-Level)

Original Text Length: 20 tokens

Symbol Table:

Symbol	Frequency	Huffman Code	Bits Used
--------	-----------	--------------	-----------

Symbol	Frequency	Huffman Code	Bits Used
<hr/>			
are	5	00	10
hello	5	01	10
you	5	10	10
how	5	11	10
<hr/>			

Original size (bits): 320

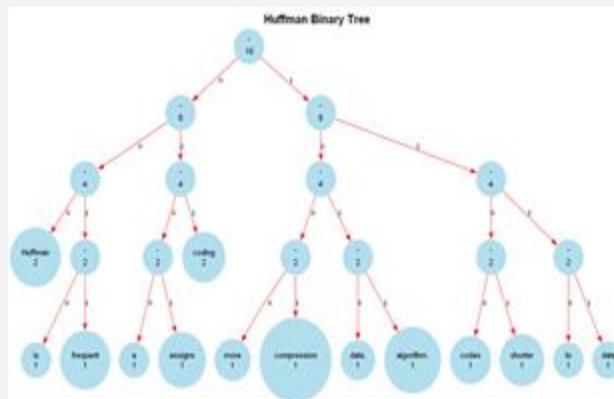
Compressed size (bits): 40

Compression ratio: 87.50%

Unique symbols: 4

Most frequent symbol: 'hello' (5 times)

Least frequent symbol: 'hello' (5 times)



Huffman Encoding Report (Word-Level)

Original Text Length: 16 tokens

Original size (bits): 256

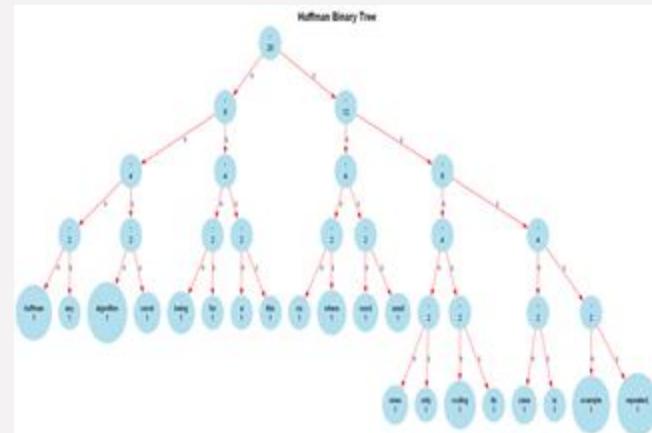
Compressed size (bits): 60

Compression ratio: 76.56%

Unique symbols: 14

Most frequent symbol: 'Huffman' (2 times)

Least frequent symbol: 'is' (1 times)



Huffman Encoding Report (Word-Level)

Original Text Length: 20 tokens

Original size (bits): 320

Compressed size (bits): 88

Compression ratio: 72.50%

Unique symbols: 20

Most frequent symbol: 'this' (1 times)

Least frequent symbol: 'this' (1 times)

Conclusion



Implemented system



- Reduces storage usage substantially
- Correctly reconstructs text
- Ensures fast page loading
- Works well across multiple books



Huffman Coding is not just theoretical,
it powers real systems



This project demonstrates complete
integration from **algorithm** → **compression**
→ **metadata** → **decoding** → **UI**.



Shows how algorithm design, data
structures and UX can combine to produce
a working, practical application.

THANK YOU