Semi-Offline Attack

Android Full-Disk Encryption

Who / What / Why

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Agenda

- Topic Introduction
- Online Attack The Baseline
- Offline Attack The Old-fashioned Way
- Semi-Offline Attack
- Questions

Topic Introduction

Full-Disk Encryption on Mobiles

• Situation: Mobile is lost or stolen

Objective: Prevent access to data stored on mobile

Security Control: Encryption at rest (storage encryption)

• Limitation 1: Data is accessible in cleartext when device

is running

• Limitation 2: Everyone in possession of the mobile can...

... launch an attack on the ciphertext

... launch an attack on the encryption cipher

... launch an attack on the crypto system

Elements of Full-Disk Encryption

We need...

- Disk encryption key
- Disk encryption cipher
- A process to en- and decrypt data

[Master Key]

[AES-CBC]

[dm-crypt]

Wait a minute...

Do we have to remember an AES key?

We also need...

- Key Encryption Key (KEK)
- Key Derivation Function (KDF) to derive the KEK from user value

Encryption Algorithms for FDE

Advanced Encryption Standard (AES)

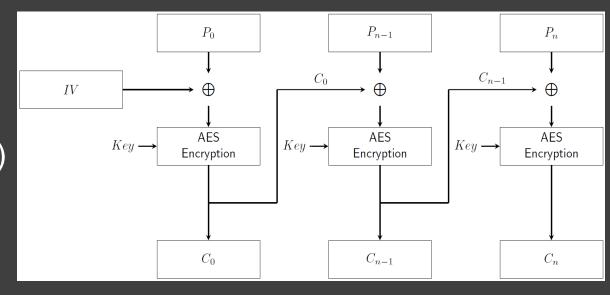
- Block Size: 128 bits
- Key Size: 128, 192, 256 bits

Block Cipher Mode

- Cipher Block Chaining (CBC)
- Alternative Mode XTS
 - XEX Encryption Mode with Tweak and Ciphertext Stealing (XTS)
 - Needs recompilation of Android source

AES-CBC

- Used for encryption and as building blocks
- Message dependency to prevent
 - Ciphertext manipulation (block swapping, insertion, or deletion)
 - Frequency analysis and dictionary attacks
- Initialisation Vector (IV) is required



Encrypted Salt-Sector Initialisation Vector

IV Properties

- Use only once with certain key
- Unpredictable value (same size as block size)
- Must be known to encrypting and decrypting party (can be public)

ESSIV

- Create a per-sector IV based on the encryption key and the sector number
- Allows to encrypt storage sector-by-sector
 - Else the whole storage needs re-encryption after an intermediate block changes

Master Key

- Created by create_encrypted_random_key() in cryptfs.c
- Read twice 16 Byte from /dev/urandom on first-time encryption
 - Master Key and Salt
 - Both values are static until smartphone is wiped!
- Protected by the KEK derived from user value
 - AES-CBC Encryption
- Stored in encrypted form in the crypto footer

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Key Derivation Functions

Password-Based Key Derivation Function 2 (PBKDF2)

- Used before Android 4.4
- Only computationally expensive processing power is cheap

scrypt

- Used since Android 4.4
- Computationally expensive AND memory-hard
- Uses PBKDF2 as building block

Screen Unlock value (PIN / Password)

Apply KDF > Result = KEK

Android – What Is Encrypted?

Compatibility Definition Document (CDD)

- Encrypt the /data partition and non-removable internal SD card partitions
- USERDATA partition contains all user stored information

Crypto Footer – stores all crypto information for FDE

- Located on the METADATA partition
- Encrypted master key
- Salt
- KDF and cipher information

Online Attack

The Benchmark Baseline

Pre-Conditions – An Attacker Needs...

- Physical access to the device
 - For as long as the brute-force attack takes
- A brute-force script
 - In the absence of any Android Debug Bridge (ADB) security controls
- A physical attack machine
 - Justin Engler & Paul Vines <u>Automated PIN Cracking</u> at DC21

Online Attack Process

- Used as baseline for the research
- Enabled ADB and trusted the Attacking Host
- Attack Duration (10'000 PIN)
 - Approx. 22 hrs 40 min
 - 2000 timeout triggered

Method	Nexus 4	Nexus 6
input tap	5.1 sec	3.3 sec
input text	2.8 sec	1.8 sec

Countermeasures

Two different prompts with two different countermeasures

	Nexus 4 (Android 4.4)		Nexus 6 (Android 5.0)	
	Start-Up	Screen-Unlock	Start-Up	Screen-Unlock
Timeout (Threshold/Duration)	10 / 30 sec	5 / 30 sec	10 / 30 sec	5 / 30 sec
Action (Threshold/Activity)	30 / Wipe	>100 / None	30 / Wipe	>100 / None

- Potential Improvements
 - Increasing penalty timeouts
 - Enforce device wipe after Xth failed attempt

Online Attack Conclusions

- Google has secured ADB sufficiently enough
 - Disabled by default / Authentication on connect
- Complexity of physical attack devices
 - Needs to support more than just PIN
 - Needs to recognize successful attempt
- Time Costs
 - 10'000 PIN took almost one day!
- Missing Security Controls
 - Too many failed attempts with no final penalty!









Offline Attack

The Old-fashioned Way

Pre-Conditions – An Attacker Needs...

- Physical access to device
 - Only for as long as the next two steps take
 - Mostly this will be due to device theft or loss
- A means to image a partition
 - USERDATA (encrypted data)
 - METADATA (crypto footer)

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Partition Imaging

Best case: Unlocked Bootloader

- Connect the smartphone with the attack host
- fastboot boot <path-to-recovery.img>
- Start the imaging process with dd and send to the host using nc

Worst case: Locked bootloader, ADB disabled

- Unlocking the bootloader wipes USERDATA
 (L. Simon and R. Anderson, University of Cambridge,
 <u>Security Analysis of Android Factory Resets</u>)
- Image via JTAG (M. F. Breeuwsma
 <u>Forensic imaging of embedded systems using JTAG</u>)

Offline Attack Script

- Initially written by Thomas Cannon and Seyton Bradford from viaForensics (ships with Santoku Linux) / Nikolay Elenkov
- PoC script handling PIN numbers only
- Procedure:
 - Apply KDF to candidate PIN
 - → Decrypt master key
 - → Decrypt Header file
 - → Check at offset 1080-1082 if ext4 magic signature (0xEF53) found

Offline Attack Analysis

- bruteforce stdcrypto.py [header] [footer]
 - Header file: min. first 1088 bytes of userdata
 - Footer file: complete crypto footer
- Attack Duration (10'000 PIN)
 - Approx. 51 min

Countermeasures

- On Device Prevent partition imaging
 - Locked Bootloader
 - Disabled ADB and ADB authentication
- Off Device Make brute-force attack expensive
 - Use of a strong KDF
 - Is an arms race with CPU power and memory increase

Offline Attack Conclusions

- Time Efficiency: 51 min compared to 22 hrs!
- Missing Security Control
 - Something to prevent offline brute-forcing
- Google responded with improvements for Android 5.0!

x

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Good News, Everyone!

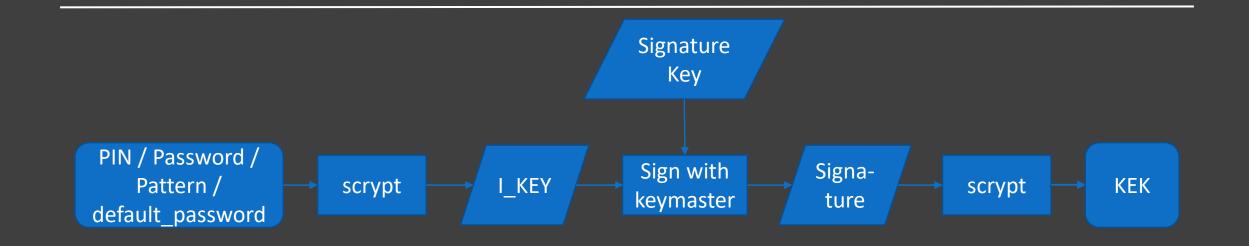
Improvements in Android 5.0

- Google addressed the Offline Attack
 - Hardware-binding of the encryption key material to the device

To be clear, the master key is not bound to the device. It is the Key Encryption Key which is cryptographically bound to the device

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New KEK Generation Process



Semi-Offline Attack

"What is wanted is not the will to believe, but the will to find out, which is the exact opposite." (Bertrand Russel)

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Pre-Conditions – An Attacker Needs...

- Physical access to device
 - For as long as the brute-force attack takes
 - Mostly this will be due to device theft or loss
- A means to image a partition
 - USERDATA (encrypted data)
 - METADATA (crypto footer)

Attack Application

Client

- Runs on attack host
- Based on the Offline Attack script

Server

- Runs on target device
- Based on cryptfs.c

Tasks

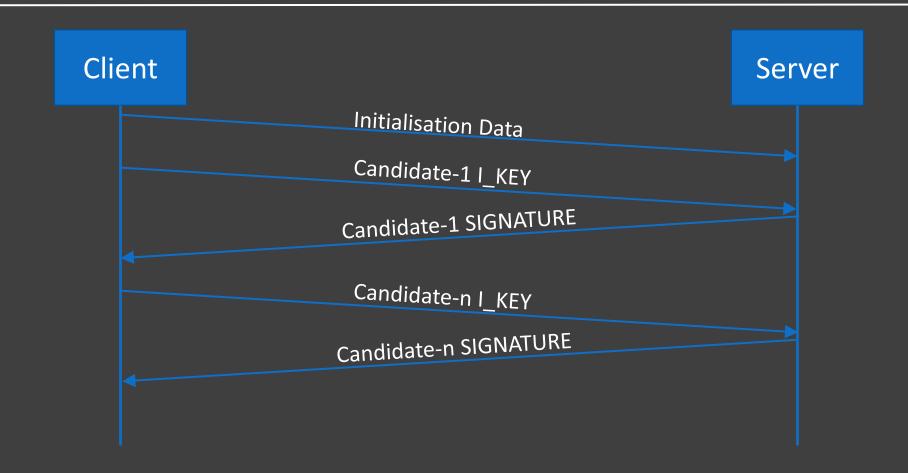
Client (attack host)

- 1. Process the crypto footer
- 2. Brute-force strategy
- 3. Perform KDF operations
- 4. Decrypt *master key*
- 5. Decrypt *userdata*

Server (target device)

- 1. Initialize *keymaster* trustlet with received key blob
- 2. Request RSA signature on intermediate key candidate

Communication



The Semi-Offline Attack (1/2)

- We have the device
- We have the attack application
- We have the two partitions
- We wipe the smartphone
 - After the wipe, we have a new master key and RSA key pair!

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The Semi-Offline Attack (2/2)

- 1. adb push poc-server /data/local/tmp
- 2. adb forward tcp:8888 tcp:9999
- 3. adb shell su -c '/data/local/tmp/poc-server 9999'
- 4. poc-client.py METADATA HEADER -a 127.0.0.1 -p 8888 -s --bruteforce
- Attack Duration (10'000 PIN)
 - Approx. 2 hrs 8 min

Demo Semi-Offline Attack

```
Received candidate from host
 Sign candidate and send result back to host
>> Test New Candidate <<
 Received candidate from host
 Sign candidate and send result back to host
>> Test New Candidate <<
  Received candidate from host
 Sign candidate and send result back to host
>> Test New Candidate <<
 Received candidate from host
 Sign candidate and send result back to host
>> Test New Candidate <<
 Received candidate from host
 Sign candidate and send result back to host
>> Test New Candidate <<
 Received candidate from host
>>> Disconnect <<<
 Close Socket: 127.0.0.1:34896
```

```
Candidate: 0800
        Decrypted Master Key:
                0xFCF79B53BF9371DE2EBDBA08937823E0
Bruteforce Duration
                        937.375999998
```

Bits, Please! 30.06.2016

- Gal Beniamini
 - <u>@laginimaineb</u>
 - http://bits-please.blogspot.com
- Read out the Keymaster key!
- Consequences for Semi-Offline Attack
 - Only one client-server interaction necessary
 - After that everything could run offline

Countermeasures

- All countermeasures of the Offline Attack still exists
- Enforce locked bootloader by all manufacturers
- Remove JTAG Interface > still gives the option to unsolder chip
- Key change mechanism for keymaster key
 - Attacker has to solve the RSA problem
 - Or directly brute-force the 16 byte master key?
 - But when to destroy the keymaster key?

Semi-Offline Attack Conclusions

- After the Offline Attack
 - Google said it had fixed the attack
 - Technically YES
 - Actually NO
- Time Efficiency: 2 hrs 8 min compared to 51 min / 22 hrs!
- The issues
 - Wrong attack conditions: How long does an attacker have access?
 - After imaging, no need to keep the mobile in it's original state
 - The internal keymaster key NEVER changes
- Google was already working on improvements!

Comparison of the Attacks

Attack Method	Attack Duration (10'000 PINs)			
Online Attack	81'549 sec	22 hrs 40 min		
Offline Attack	3'068 sec	51 min		
Semi-Offline Attack	7'622 sec	2 hrs 8 min		

Questions