

0.1 Smooth scrolling on EGA

Performing a full-screen redraw per frame would kill the CPU, as it would require updating every pixel on all four EGA planes. Achieving a 60Hz frame rate while refreshing the entire screen is impossible under these constraints. If we were to run the following code, which simply fills all memory banks, it would run at 5 frames per seconds.

```
# define SC_INDEX    0x3C4
# define SC_DATA     0x3C5
# define SC_MAPMASK  0x02

char far *EGA = (unsigned char far*)0xA0000000L;

void selectPlane (int plane) {
    outp ( SC_INDEX , SC_MAPMASK );
    outp ( SC_DATA , 1 << plane );
}

void WriteScreen(void){
    int i, bank_id;

    for (bank_id = 0 ; bank_id < 4 ; bank_id++) {
        selectPlane(bank_id);
        for (i = 0; i < 40 * 200; i++) {
            EGA[i] = 0x00;
        }
    }
}
```

“

Clearly, there is a whole world of awesome things there that we just couldn't do on the PC...

You can just redraw the whole screen, but then it turns out...

Well, you're going five frames a second.

John Carmack¹

”

¹Interview with Lex Fridman in 2022, this quote is around 1h:41m.

So, how did they create a smooth scrolling game with these limitations? The solution was twofold:

- Utilizing specific EGA hardware tricks
- Updating only the portion of the screen that has changed

Each method is described in detail in the following sections.

0.1.1 EGA Virtual Screen and Pel Panning

The EGA card adds a powerful approach to linear addressing: the logical width of the virtual screen in VRAM does not need to match the physical screen display width. A programmer can define a virtual screen width of up to 4096 pixels and use the physical screen as a window onto any part of the virtual screen. What's more, the virtual screen's logical height can extend up to the total VRAM capacity. The code below demonstrates how to set a custom logical width.

```
CRTC_INDEX = 03D4h
CRTC_OFFSET = 19

;=====
;

;  set wide virtual screen
;

;=====

mov dx,CRTC_INDEX
mov al,CRTC_OFFSET
mov ah,[BYTE PTR width] ;screen width in bytes
shr ah,1                ;register expresses width
                         ;in word instead of byte
out dx,ax
```

The displayed area of the virtual screen at any time is determined by setting the display memory address for fetching video data, configured via the CRTC Start Address register. The default address is A000:0000h, though the offset can be adjusted to any other address.

```
CRTC_INDEX    = 03D4h
CRTC_STARTHIGH = 12

;=====
;
;  VW_SetScreen
;
;=====

cli                      ; disable interrupts

    mov cx,[crtc]          ; [crtc] is start address
    mov dx,CRTC_INDEX      ; set CRTR register
    mov al,CRTC_STARTHIGH  ; start address high register
    out dx,al
    inc dx                 ; port 03D5h
    mov al,ch
    out dx,al              ; set address high
    dec dx                 ; set CRTR register
    mov al,0dh              ; start address low register
    out dx,al
    mov al,cl
    inc dx                 ; port 03D5h
    out dx,al              ; set address low

    sti                    ; enable interrupts

    ret
```

To pan down a scan line, it requires only to increase the CRTC start address by the logical width. Horizontal panning is achieved by incrementing the start address by one byte. In EGA's planar graphics modes, the eight bits in each video RAM byte correspond to eight consecutive on-screen pixels, meaning the screen can move horizontally in steps of 8 pixels. This is rather coarse and not really smooth horizontal scrolling. Fortunately, another EGA register addresses this limitation.

0.1. SMOOTH SCROLLING ON EGA

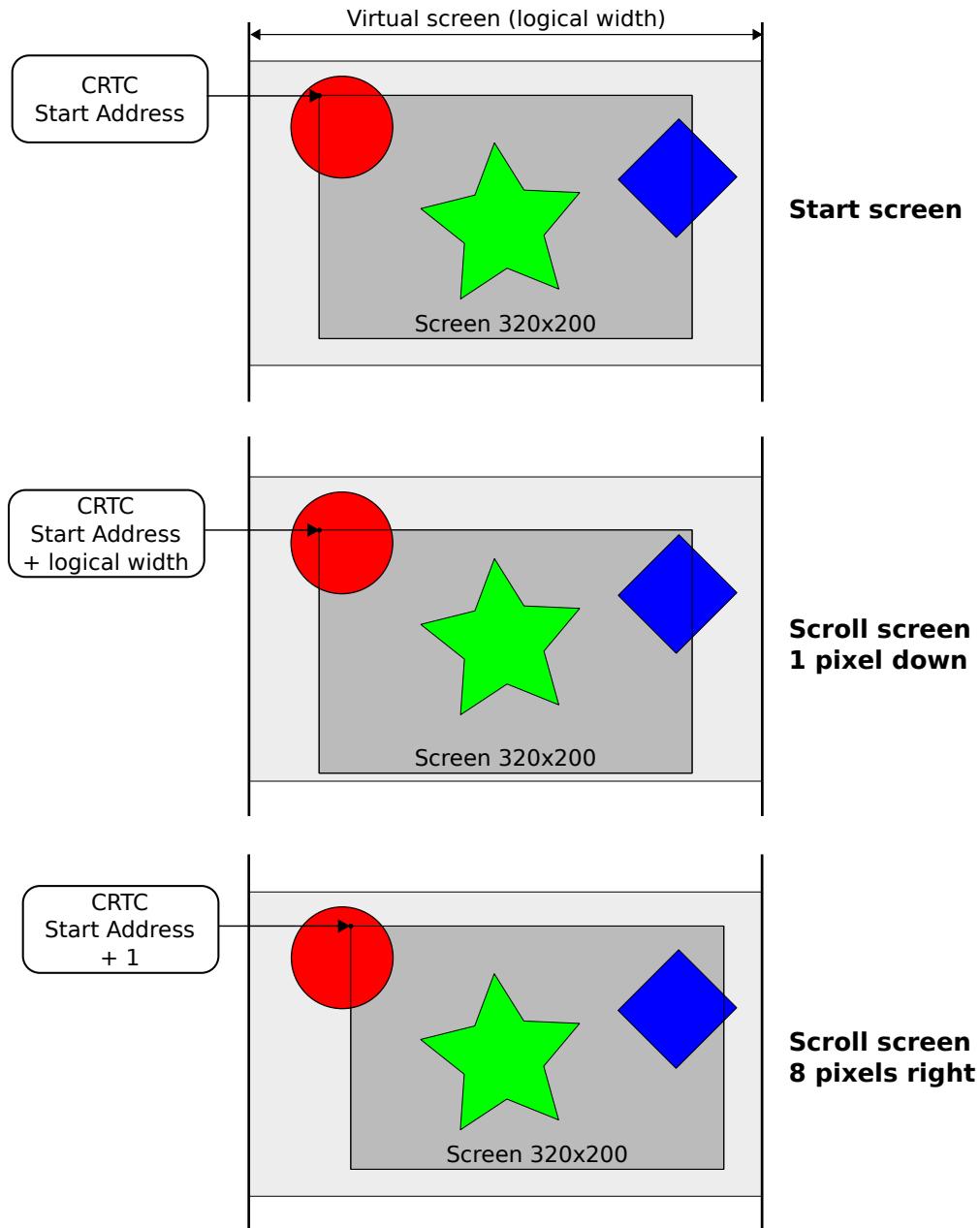


Figure 1: EGA screen scrolling by updating the CRTC Start Address.

0.1.2 Horizontal Pel Panning

Smooth horizontal pixel scrolling is managed by the "Horizontal Pel Panning" register in the Attribute Controller (ATC), allowing fine adjustment in 1-pixel increments up to 7 pixels to the left.

However, programming the ATC requires attention: the ATC Index register only uses the lower five bits (bits 0-4) as its internal index. The next most significant bit, bit 5, controls the source of the video data send to the monitor by the EGA card. When bit 5 is set to 1, the output of the color palette controls the displayed pixels; this is normal operation. When bit 5 is 0, video data doesn't come from the color palette, and the screen becomes a solid color. To maintain normal video operation, bit 5 must be set to 1 by writing 20h to the register.

```

ATTR_INDEX = 03C0h
ATTR_PELPAN = 19

;=====
;

;  set horizontal panning
;

;=====

    mov dx, ATTR_INDEX
    mov al, ATTR_PELPAN or 20h ;horizontal pel panning register
                                ;(bit 5 is high to keep palette
                                ;RAM addressing on)
    out dx,al
    mov al,[BYTE pel]          ;pel pan value [0 to 8]
    out dx,al

```

Smooth horizontal scrolling on EGA involves adjusting the CRTC Start Address alongside fine-tuning within an 8-pixel range using horizontal pel panning. The following steps accomplish this:

- Calculate the panning offset in pixels for both x- and y-direction.
- Smooth y-panning is achieved by adding the (logical width *) y to the CRTC start address.
- For coarse horizontal scrolling, increase the CRTC start address by $x / 8$ bytes.
- Fine horizontal adjustment is applied using $(x \% 8)$ to the horizontal pel panning register.

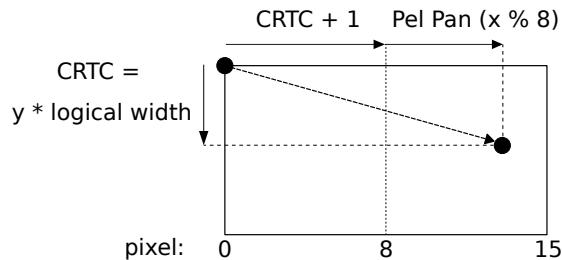


Figure 2: Smooth scrolling in EGA.

0.2 Adaptive Tile Refreshment

So far, we have built a virtual screen in VRAM allowing smooth one pixel moves in both axes using only EGA registers. However, once the screen edge is reached, a full-screen redraw would be too slow, causing a drop to 5 fps.

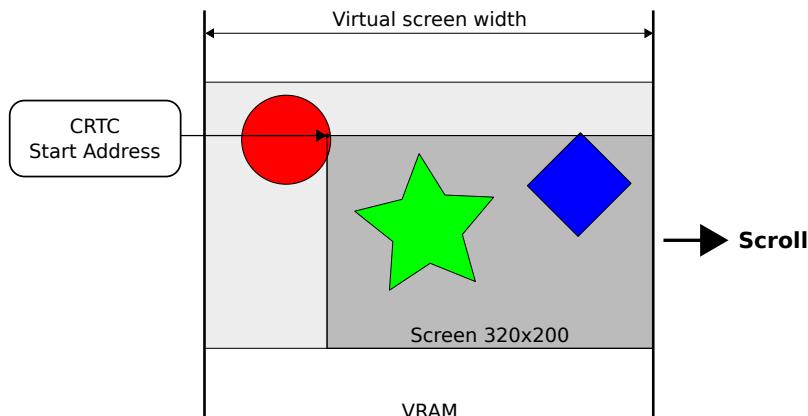


Figure 3: Screen reaching the edge of virtual screen.

One option for screen refresh is to load the entire level into VRAM, setting the logical width to match the level width. By adjusting the CRTC Start Address and Horizontal Pel Panning registers, smooth scrolling can be achieved. However, loading the entire level requires far more VRAM than is available. The first level, *Horsh Radish Hill*, is 136 tiles wide and 37 tiles high. Each tile is 128 bytes, which means a total of $136 \times 37 \times 128$ bytes = 644KB is required to store the entire level in VRAM. Since we only have 256KB available, this is not a feasible option.

Trivia : At the time of writing this book, most video cards contain more than 1GB VRAM, sufficient to store all Commander Keen levels at once in VRAM.

To address this, John Carmack invented "Adaptive Tile Refresh" (ATR). The core idea is to refresh only those areas of the screen that need to change.

Let's look at *Commander Keen 1: Marooned on Mars* in Figure 4. This is the first level of Marooned, immediately to the right of the crashed Bean-with-Bacon Megarocket. The first figure is the start of the level, the second figure is after Keen has scrolled one tile (16 pixels) to the right through the world. They look almost identical to the naked eye, don't they?

Now, if we perform a difference on both images, you can see which tiles need to be changed upon screen refresh. The trick behind the scrolling is to only redraw tiles that actually changed after panning 16 pixels (one tile). For matching tiles there is nothing to do and they are skipped entirely. In Figure 4.21, there are large swathes of constant background tiles. In total, only 69 tiles out of the 260 tiles need to be refreshed, which is 27% of the screen!

This is also the part where the game designers Tom Hall and Jogn Romero had to help the engine. Smooth scrolling is inversely proportional to the number of tiles to redraw. A checkerboard tile pattern basically means a full-screen redraw and would kill the CPU. So to avoid costly "jolts", the game designers built tile maps with a lot of reperating tiles.

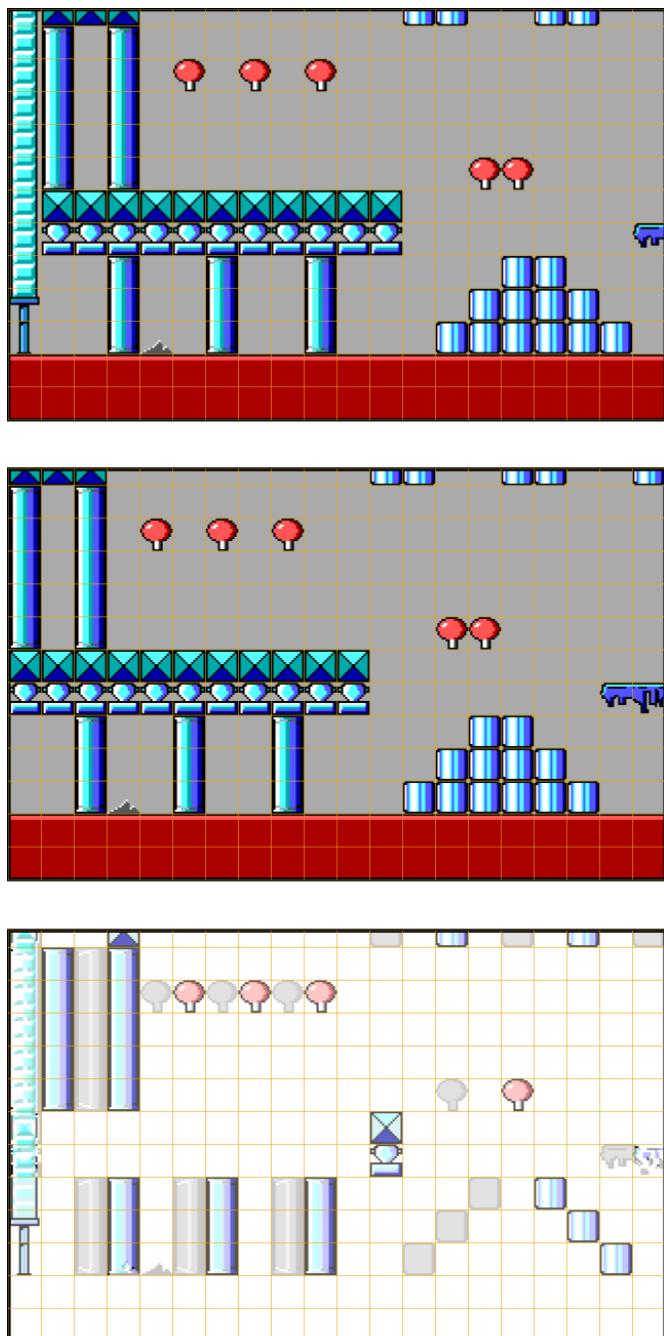
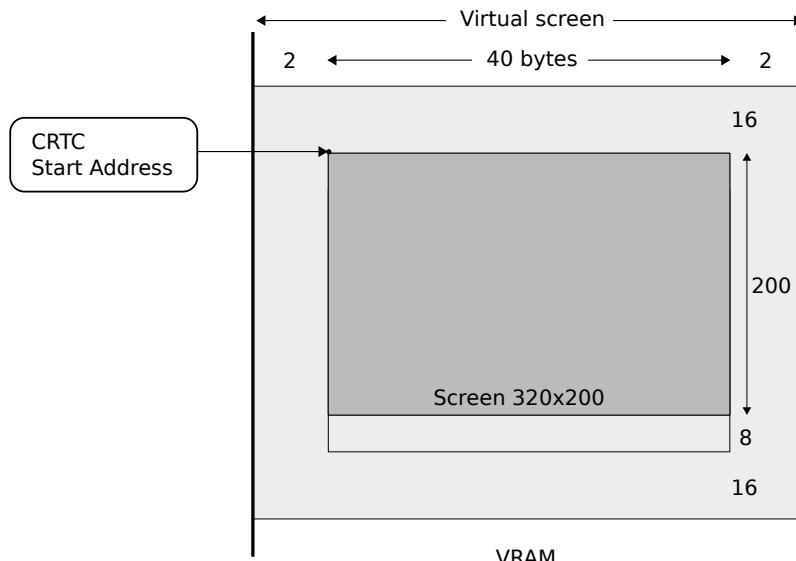


Figure 4: Start of the world, moved one tile to the right and difference.

0.2.1 Adaptive tile refreshment in Commander Keen 1-3

This section explains how ATR is working for the first three episodes of Commander Keen. Since there is no source code released for these episodes, ATR will be explained without code examples. The later versions of Commander Keen, including Keen Dreams, used a different, improved engine which will be explained in the next section².

The EGA screen in mode 0xD has a resolution of 320x200 pixels, or 40x200 bytes. Let's extend the height by 8 bytes to have a height of 208 pixels, so the screen fits nicely in 20x13 tiles. By making the virtual screen one tile higher and one wider on each side of the screen, the engine can scroll up to 16 pixels to any direction of the screen without any tile refresh, by simply adjusting the CRTC Start Address and Pel Pan registers.



Now, let's have a closer look at the EGA VRAM setup. The video memory is organized into three virtual screens:

- Page 0 and 1, which are used to switch between buffer and visible screen. The idea between two pages (double buffer) is that the code can draw in the second buffer while the first buffer is being shown on screen, which is then switched out during screen refresh. This ensures that no frame is ever displayed mid-drawing, which yields smooth, flicker-free animation.

²<https://retrocomputing.stackexchange.com/questions/22175/what-is-adaptive-tile-refresh-in-the-context-of-commander-keen>

0.2. ADAPTIVE TILE REFRESHMENT

- A master page containing a static page, which is copied to the buffer screen when performing the screen refresh.

Each virtual screen has a size of $44 \times 240 \times 4 = 42,420$ bytes. So within a 256KB EGA card there is enough VRAM available to keep all three virtual screens in memory. The page that is displayed on screen is selected by setting the CRTC Start Address register at which to begin fetching video data.

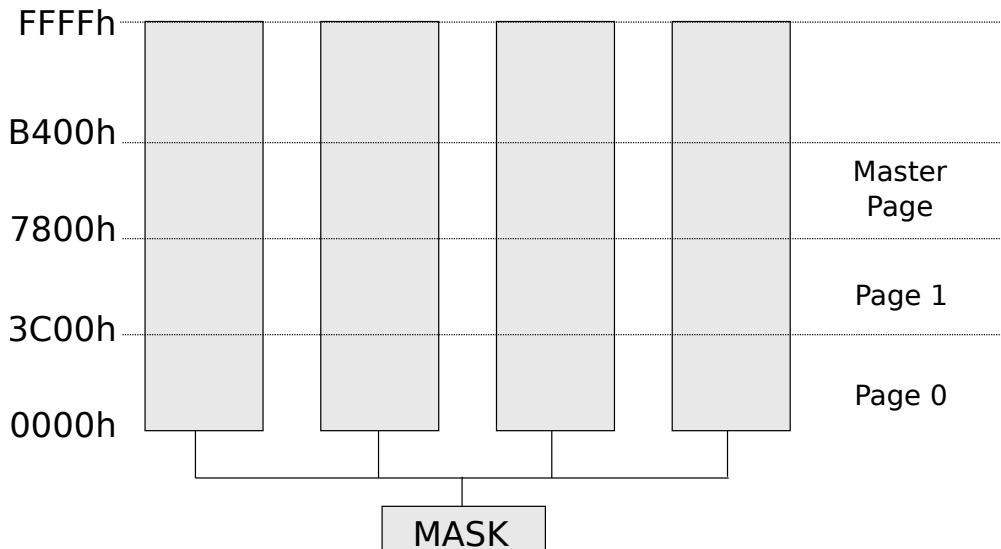


Figure 5: Virtual screen layout on EGA card.

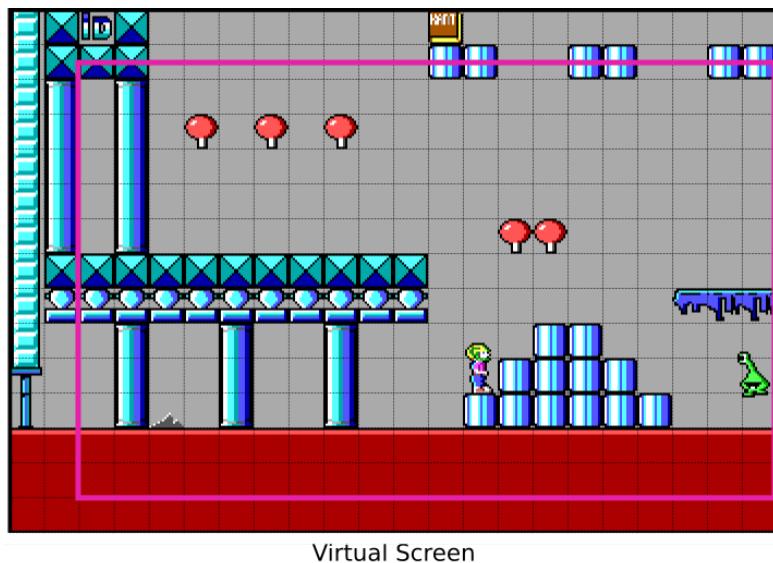
For both the buffer and visible screen a tile array is created to maintain which tiles are changed since last refresh.

```
byte update[2][UPDATESIZE];
```

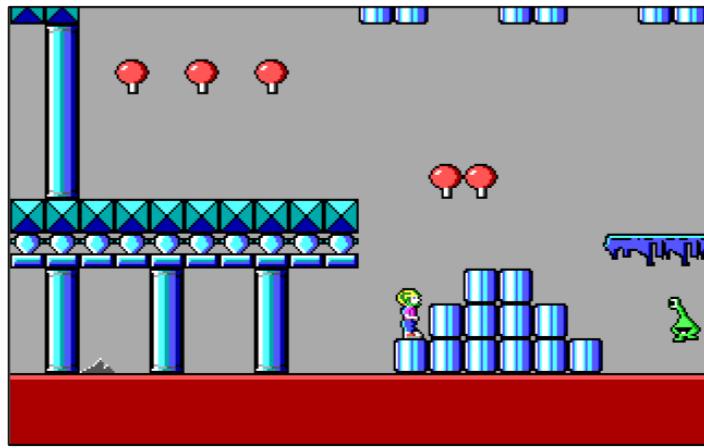
Steps to refresh the screen are as follows:

1. Check if the player has moved one tile in any direction.
2. Identify changed tiles and copy these to the master page, marking them in both tile arrays.
3. Refresh the buffer page by copying marked tiles from the master page.
4. Switch the view and buffer pages by adjusting the CRTC Start Address and Pel Panning registers.

Ignoring sprites for now, the following steps illustrate these stages. In the next four screenshots, we take you step-by-step through each of the stages. The player has reached the edge of the virtual screen and moves further to the right.



Virtual Screen



Display (320x200 pixels)

Figure 6: Start: Reach the end of the virtual screen.

0.2. ADAPTIVE TILE REFRESHMENT

The engine keeps track of which tile numbers are part of the virtual screen. Since the engine only refreshes the screen at tile size granularity, it can determine extremely fast what has changed on the screen by comparing tile numbers. If the tile number has changed, the tile is updated by copying tiles from RAM into the VRAM master page. The changed tiles are marked with a '1' in both tile arrays, meaning it needs to be updated upon next refresh.

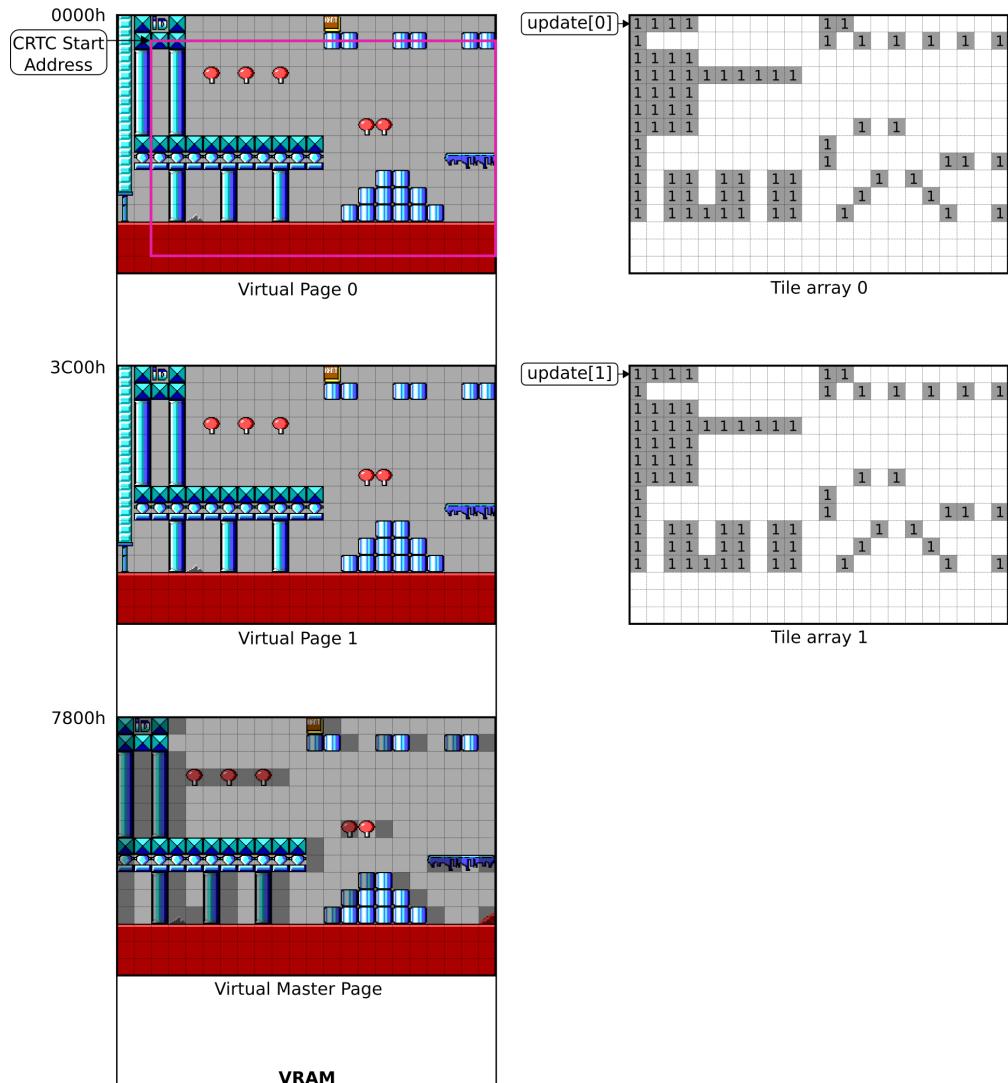


Figure 7: Update tiles in master page and update both tile arrays.

The next step is to scan all tiles in buffer tile array (array 1) and for each tile marked '1', copy the corresponding tile from master to virtual page 1 page.

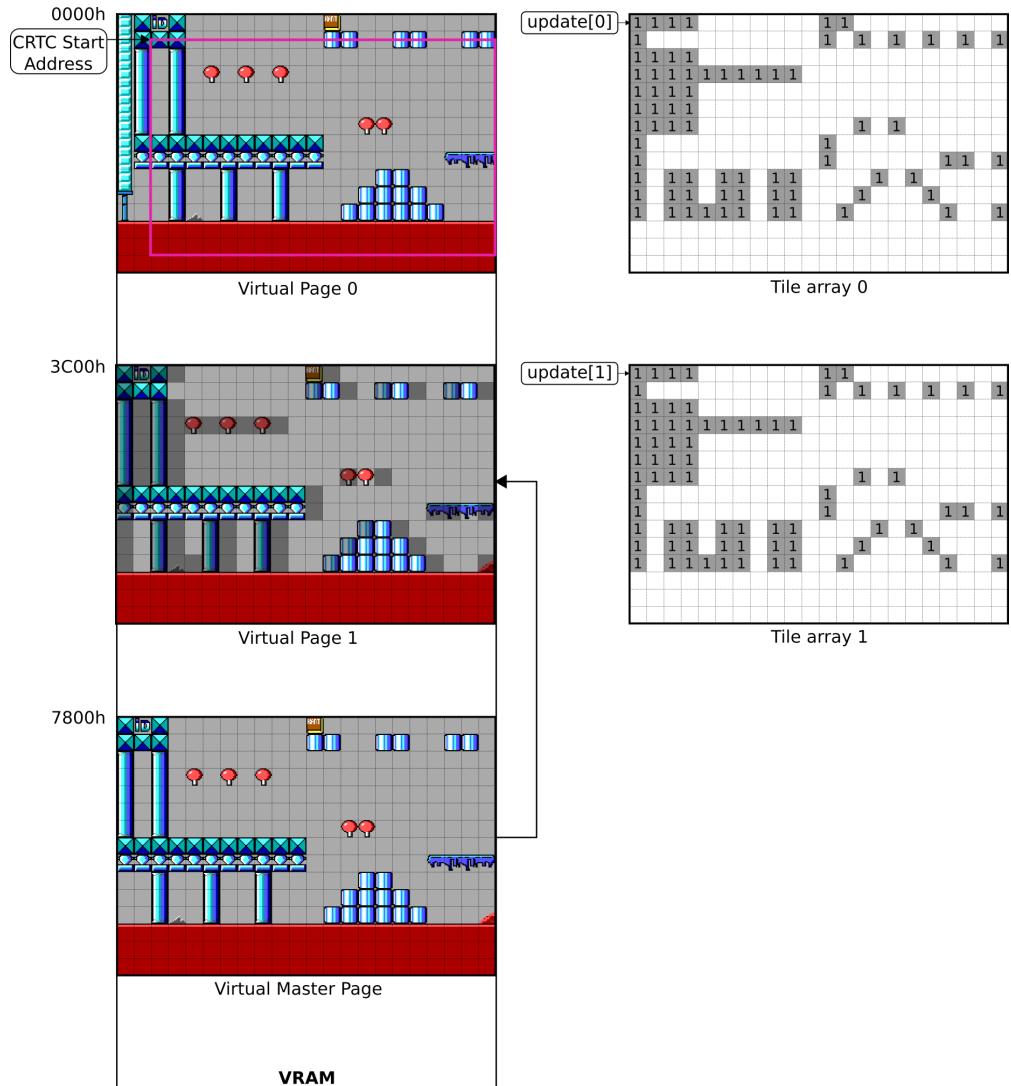


Figure 8: Copy changes tiles from virtual master page to page 1.

0.2. ADAPTIVE TILE REFRESHMENT

In the final step, point the visible screen to virtual page 1 by updating the CRTC start address. Finally, tile array 1 is cleared to '0'. Now the first step is repeated, but this time virtual page 0 acts as the buffer screen. Note that after swapping, tile array 0 keeps marked tiles from last update. This makes sense, as the current buffer page is not yet refreshed since it was displayed in the previous refresh cycle.

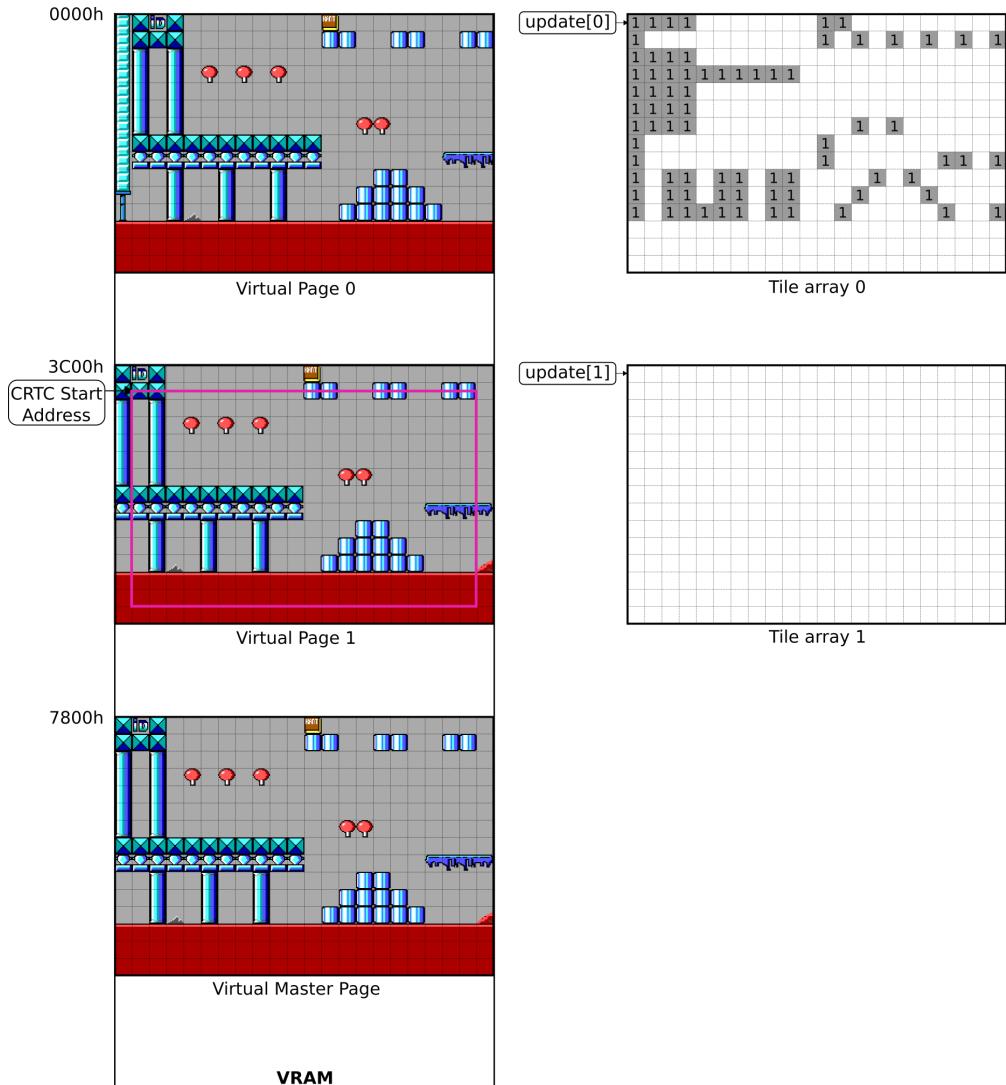


Figure 9: Update CRTC Start Address to virtual page 1 and empty tile array 1.

Now the buffer is refreshed and the CRTC address is updated, the final step is fine adjustment using the Pel Panning register.

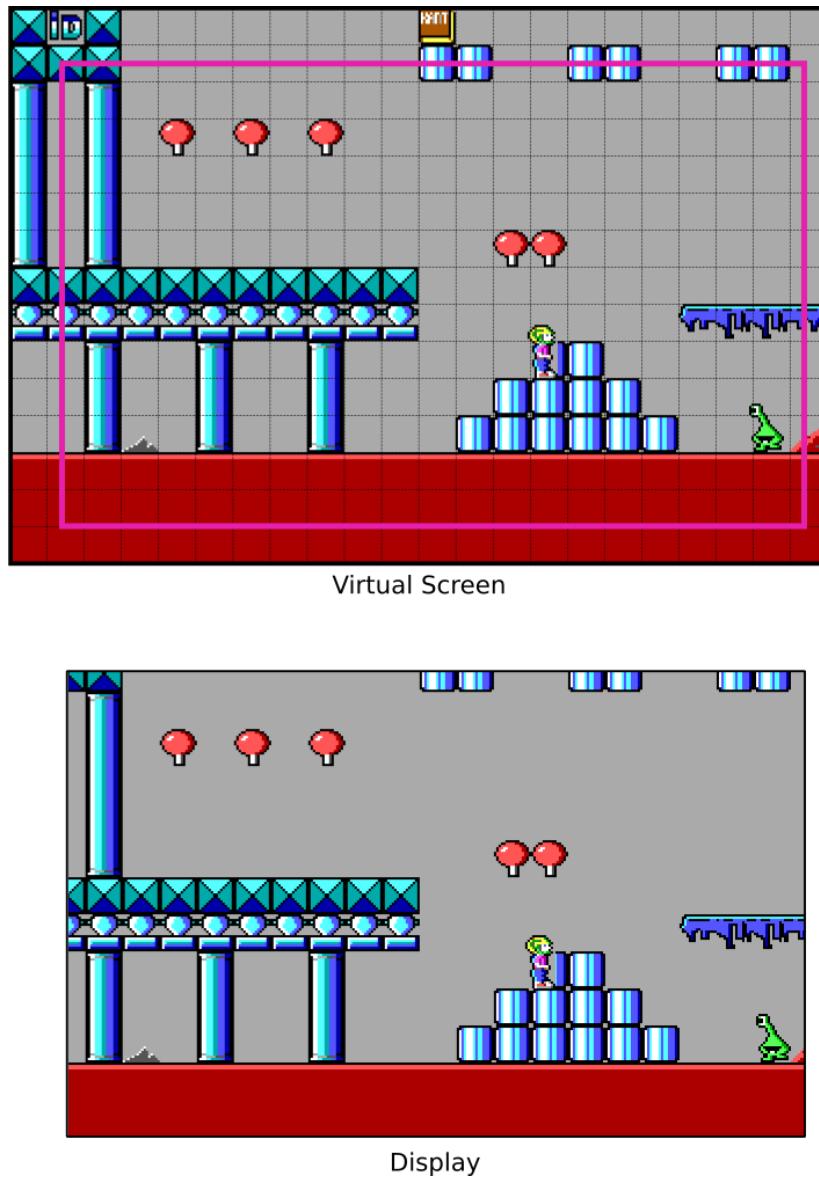


Figure 10: Screen is refreshed and scrolled to the right.

0.2.2 Wrap around the EGA Memory

John Carmack explored what would happen if you push the virtual screen over the 64kB border (address 0xFFFF) in video memory. It turned out that the EGA continues the virtual screen at 0x0000. This means you could wrap the virtual screen around the EGA memory and only need to add a stroke of tiles on one of the edges when Commander Keen moves more than 16 pixels.

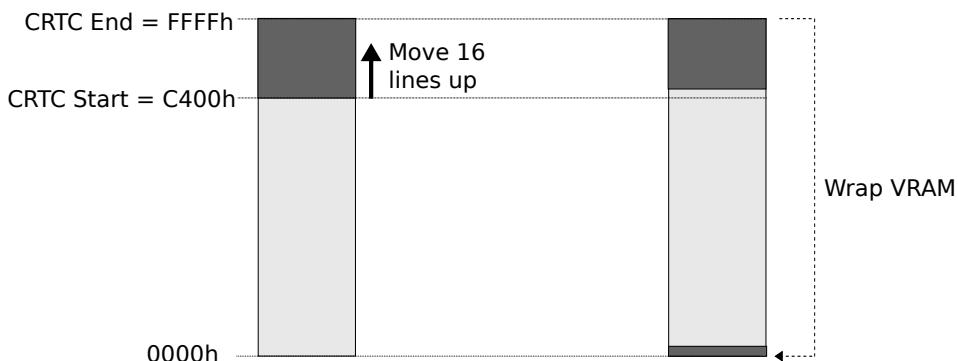


Figure 11: Wrap virtual screen around the EGA memory

“

I finally asked what actually happens if you just go off the edge [OF THE VRAM]?

If you take your [CRTC] start and you say OK, I can move over and I get to what should be the bottom of the memory window. [...] What happens if I start at 0xFFFF at the very end of the 64k block? It turns out it just wraps back around to the top of the block.

I'm like oh well this makes everything easy. You can just scroll the screen everywhere and all you have to draw is just one new line of tiles.

It just works. We no longer had the problem of having fields of similar colors. It doesn't matter what you're doing, you could be having a completely unique world and you're just drawing the new strip.

John Carmack³

”

³An explanation further elaborated during the same interview with Lex Fridman in 2022.

There was however an issue with the introduction of Super VGA cards, which had typically more than 256kB RAM⁴. This resulted in crippled backwards compatibility and the wrapping around 0xFFFF did not work anymore on these cards.

There is an easy solution to resolve this issue. As you can see in Figure 5 on page 10, the space between 0xB400 and 0xFFFF is not used and contains enough space for another virtual screen. In case the start address is between 0xC400 and 0xFFFF the corresponding screen is copied to the opposite end of the buffer, as illustrated in Figure 12.

“

I was in a tough position. Do I have to track every single one of these [SUPER VGA CARDS] and it was a madhouse back then with 20 different video card vendors with all slightly different implementations of their non-standard functionality. Either I needed to natively program all of the cards or I kind of punt. I took the easy solution of when you finally did run to the edge of the screen I accepted a hitch and just copied the whole screen up.

John Carmack⁵

”

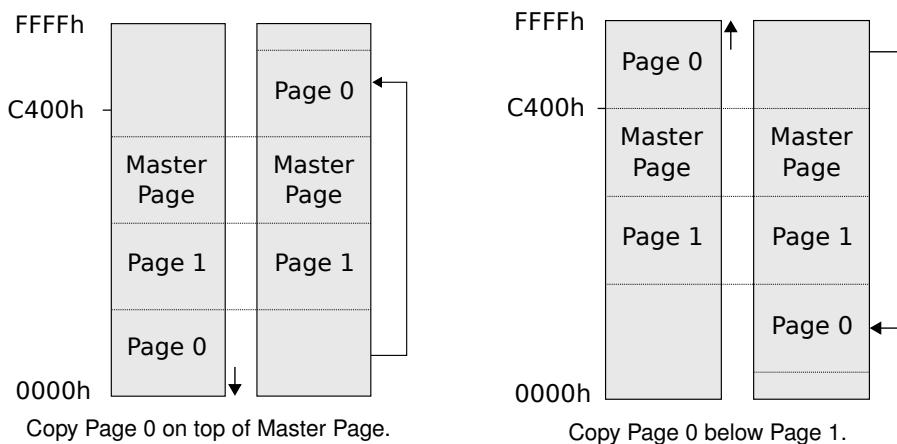


Figure 12: Move screen to opposite end of VRAM buffer

⁴In 1989 the VESA consortium standardized an API to use Super VGA modes in a generic way. One of the first modes was 640x480 at 256 colors requiring at least 256kB RAM, which from a hardware constraint resulted in 512kB.

⁵And again the same interview with Lex Fridman in 2022.

```
#define SCREENSPACE      (SCREENWIDTH*240)
#define FREEEGAMEM        (0x100001-31*SCREENSPACE)

screenmove = deltay*16*SCREENWIDTH + deltax*TILEWIDTH;
for (i=0;i<3;i++)
{
    screenstart[i] += screenmove;
    if (compatability && screenstart[i] > (0x100001-
SCREENSPACE) )
    {
        // move the screen to the opposite end of the buffer
        screencopy = screenmove>0 ? FREEEGAMEM : -FREEEGAMEM;
        oldscreen = screenstart[i] - screenmove;
        newscreen = oldscreen + screencopy;
        screenstart[i] = newscreen + screenmove;
        // Copy the screen to new location
        VW_ScreenToScreen (oldscreen,newscreen ,
                           PORTTILESWIDE*2,PORTTILESHIGH*16);

        if (i==screenpage)
            VW_SetScreen (newscreen+oldpanadjust ,oldpanx &
xpanmask);
    }
}
```

As explained in Section ?? each pixel is encoded by four bits, which are spread across the four EGA banks. So how can we copy four VRAM planes fast enough, without noticing any performance hit? To copy eight pixels, byte-aligned, one would have to do four read and four writes.

But how can we copy four VRAM planes fast enough, without noticing any performance hit? As explained in Section 3.3.7 each pixel is encoded by four bits, which are spread across the four EGA banks. To copy eight pixels, byte-aligned, one would have to do four read and four writes.

Having a closer look at the EGA card one notice there is a latch placed in front of the ALU, which can be re-purposed for the greater good. With a little creativity the ALU in front of each bank can be setup to use only the latch for writing⁶. With such a setup, upon doing one read, four latches are populated at once and four bytes in the bank are written with only one write to the RAM. Now it is possible to copy the entire buffer fast enough, without notifying any performance impact.

⁶The mask trick is discussed in *Game Engine Black Book: Wolfenstein 3D* for the VGA card. The trick is the same for the EGA card.

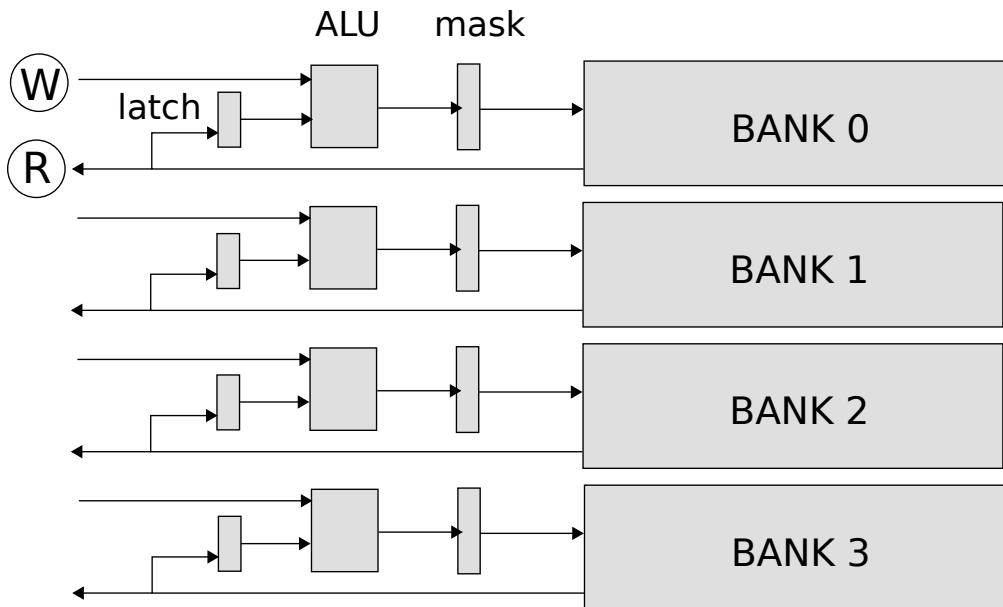


Figure 13: Latches memorize read operations from each bank. The memorized value can be used for later writes.

```

GC_INDEX      = 0x3CE      ;Graphics Controller register
GC_MODE       = 5          ;mode register
SC_INDEX      = 0x3C4      ;Sequence register
SC_MAPMASK    = 2          ;map mask register

;=====
; Set EGA mode to read/write from latch
;=====
cli           ;interrupts disabled
mov dx,GC_INDEX          ;mode 1, each memory plane is
mov ax,GC_MODE+256*1      ;written with the content of
out dx,ax                 ;the latches only

mov dx,SC_INDEX           ;enable writing to all 4 planes
mov ax,SC_MAPMASK+15*256  ;at once
out dx,ax
sti                      ;interrupts enabled

```

0.2.3 Adaptive tile refreshment in Keen Dreams

The EGA memory wrapping results in an improved, more simplified Adaptive Tile Refreshment algorithm. First, a virtual screen and tile array is defined by making the display one tile taller and wider, creating what is known as the viewport. This allows the engine to scroll the display up to 16 pixels to the right and bottom without refreshing tiles, by simply adjusting the CRTC Start Address and Pel Pan register.

An additional column is then added to the viewport to support horizontal scrolling. Finally, the tile array is further expanded to allow the entire viewport to move up to two tiles in any direction.

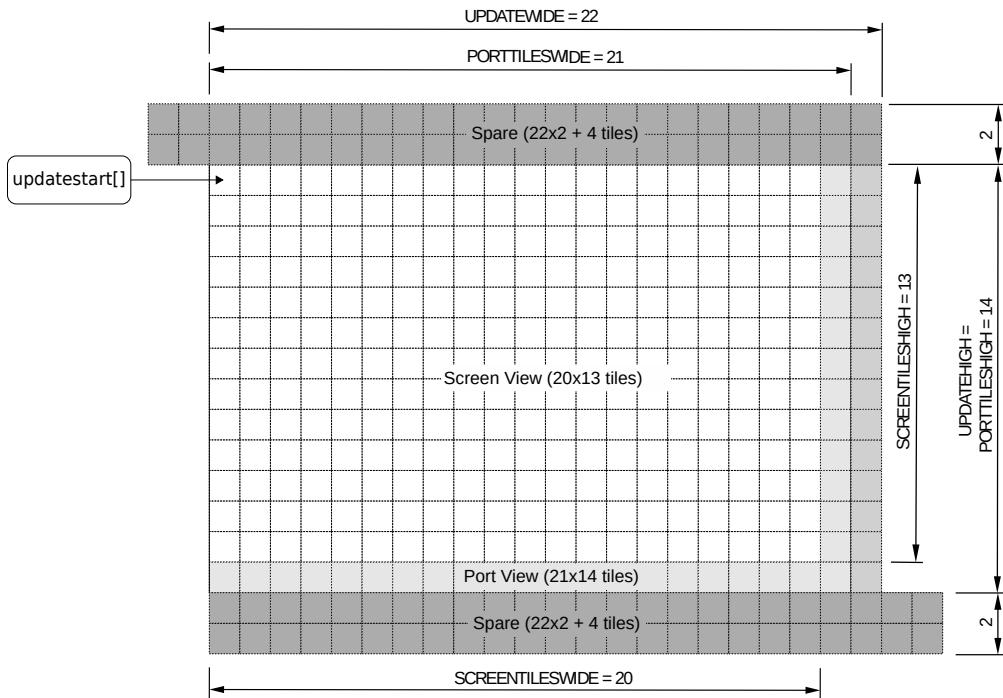


Figure 14: Tile array layout.

The full refresh cycle, including sprite updates, proceeds as follows:

1. Verify if the player has moved one tile in any direction.
2. Copy the new column or row of tiles to the master page, update both the tile array pointers, and mark the tiles in both arrays.

3. Refresh the buffer page by scanning the tile index array. If a tile is marked, copy it from the master page to the buffer page.
4. Iterate through the sprite removal list, copying the corresponding image block from the master page to the buffer page to clear the sprite.
5. Iterate through the sprite list, copying each sprite image block to the buffer page.
6. Switch the view and buffer pages by adjusting the CRTC Start Address and Pel Panning registers.

```
void RF_Refresh (void)
{
    updateptr = updatestart[otherpage];

    RFL_AnimateTiles ()// update animated tiles

    // copy newly scrolled and animated tiles
    // from the master to buffer screen
    EGAWRITEVIDEO(1);
    EGAMAPMASK(15);      // write 4 bytes of VRAM at once
    RFL_UpdateTiles (); // copy from master to buffer page
    RFL_EraseBlocks (); // remove sprites

    // update sprites
    EGAWRITEVIDEO(0);
    RFL_UpdateSprites ();

    // display the changed screen (swap view and buffer)
    VW_SetScreen(bufferofs+panadjust,panx & xpanmask);
}
```

0.2. ADAPTIVE TILE REFRESHMENT

Let's go over the refresh cycle step by step. At the start, all three virtual pages display the same viewport, with virtual page 0 currently shown on-screen. Both tile arrays are empty, meaning no tile updates are required in the next refresh cycle. The player has reached the edge of the virtual screen and moves further to the left.

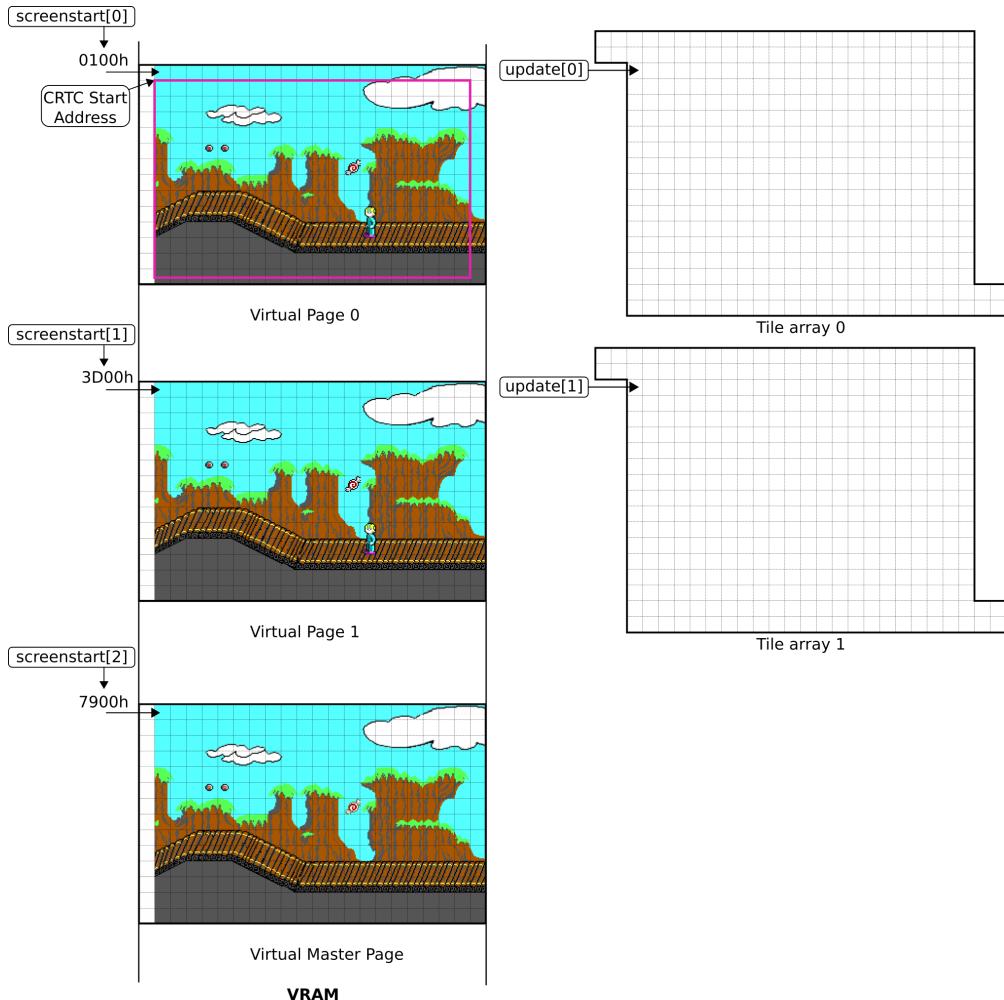


Figure 15: VRAM and tile arrays before scrolling to the left.

Both VRAM and tile array pointers shift left to introduce a new column of tiles. The VRAM pointer is lowered by 2 bytes (1 tile), and the tile array pointer is decreased by 1 byte. A new column of tiles is copied from RAM to the master page, while the leftmost column in both tile arrays is marked with '1' to ensure it updates in the next refresh cycle. Animated tiles are then copied to the master page, and the corresponding tiles in the array are also marked with '1'.

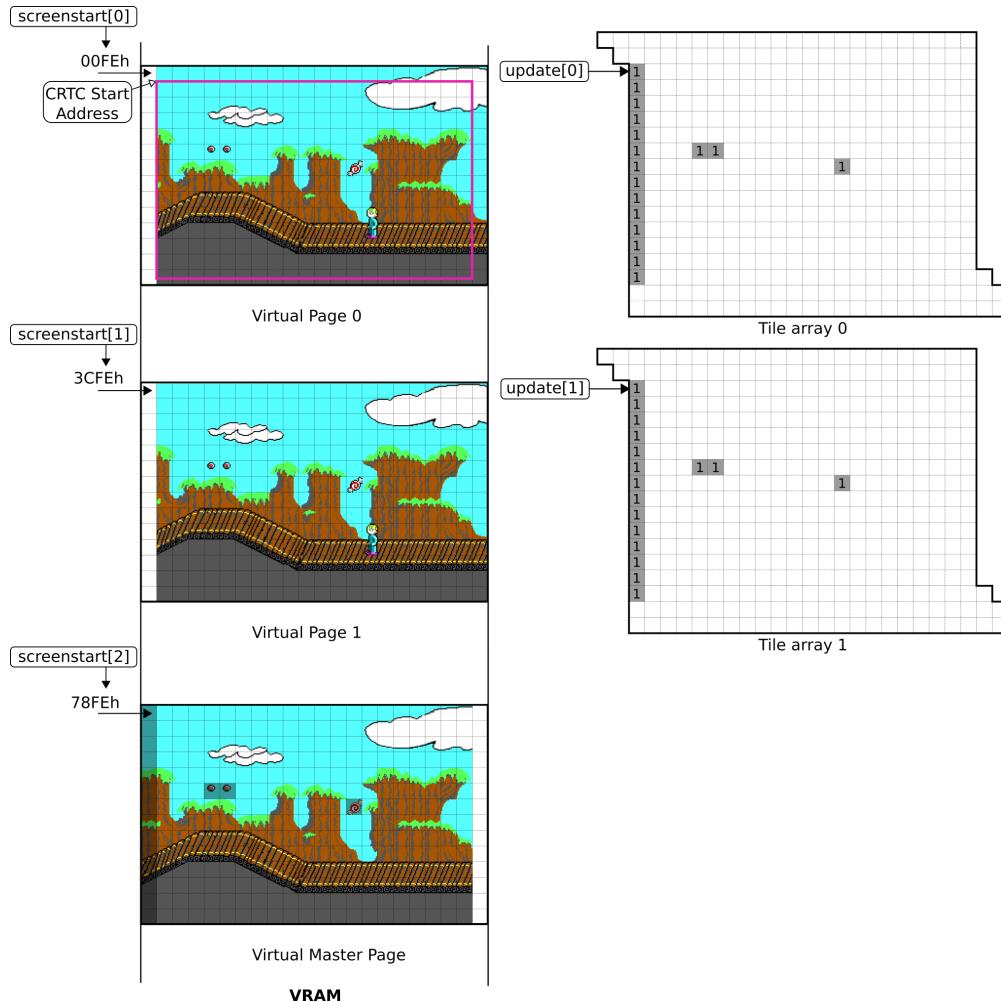


Figure 16: Update Virtual master with new left column and animated tiles.

0.2. ADAPTIVE TILE REFRESHMENT

Next, the engine scans all '1's and copies the corresponding tiles from the master to the buffer page.

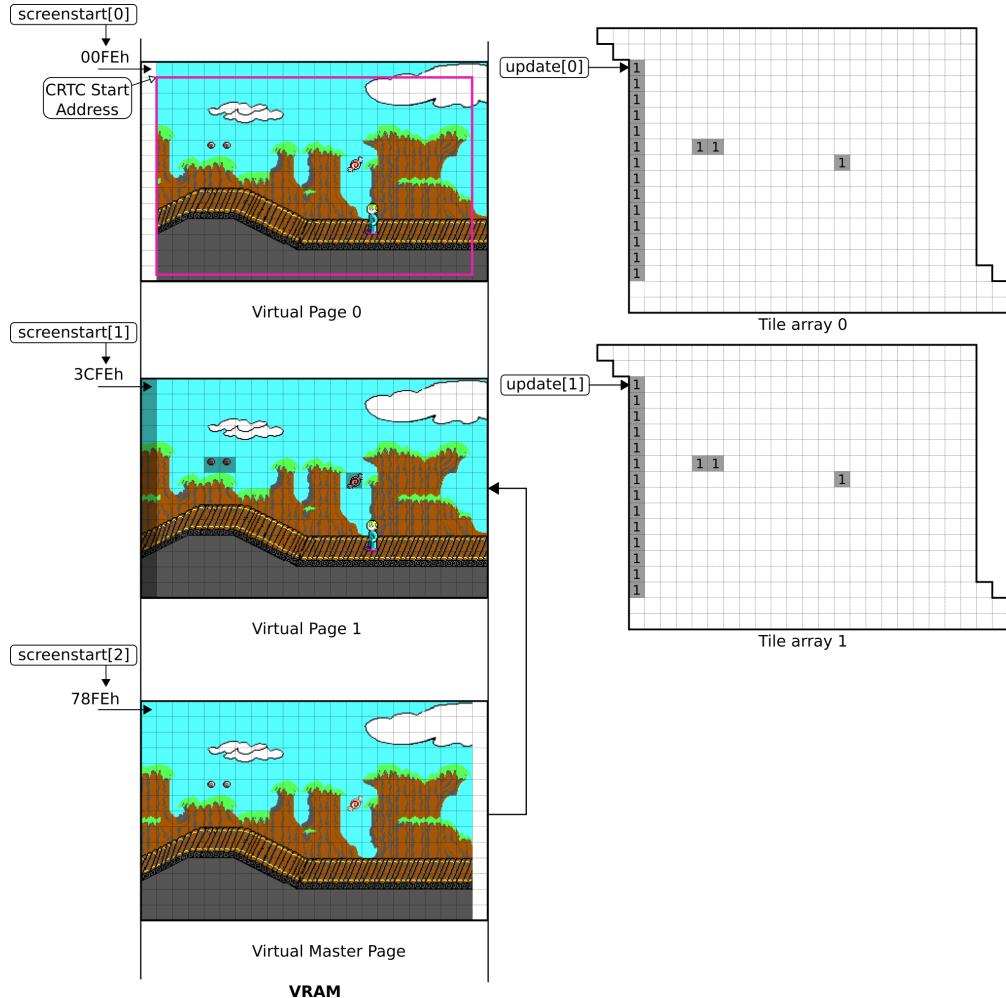


Figure 17: Copy changed tiles from master to buffer page.

Each sprite removed from the screen has its location and size stored in the sprite removal list, which removes the sprite by copying a specific section from the master screen to the virtual page. Tiles overlapping with the removal block are marked with a '2'.

```
typedef struct
{
    int      screenx, screeny;
    int      width, height;
} eraseblocktype;

//sprite removal list for Page 0 and Page 1
eraseblocktype  eraselist[2][MAXSPRITES], *eraselistptr[2];
```

```
void RFL_EraseBlocks (void)
{
    eraseblocktype *block,*done;
    unsigned pos;

    block = &eraselist[0][0];
    done = eraselistptr[0];

    while (block != done)
    {
        [...]

        //
        // erase the block by copying from the master screen
        //
        pos = ylookup[block->screeny]+block->screenx;
        block->width = (block->width + (pos&1) + 1)& ~1;
        pos &= ~1;           // make sure a word copy gets used
        VW_ScreenToScreen (masterofs+pos,bufferofs+pos,
                           block->width,block->height);

        [...]
        block++;
    }
}
```

Trivia : The '2' marking is nowhere used in the engine, most likely it is used in the original ATR algorithm.

0.2. ADAPTIVE TILE REFRESHMENT

The final step in updating the buffer is copying sprites to their new locations, with each corresponding tile marked with a '3' in the tile array.

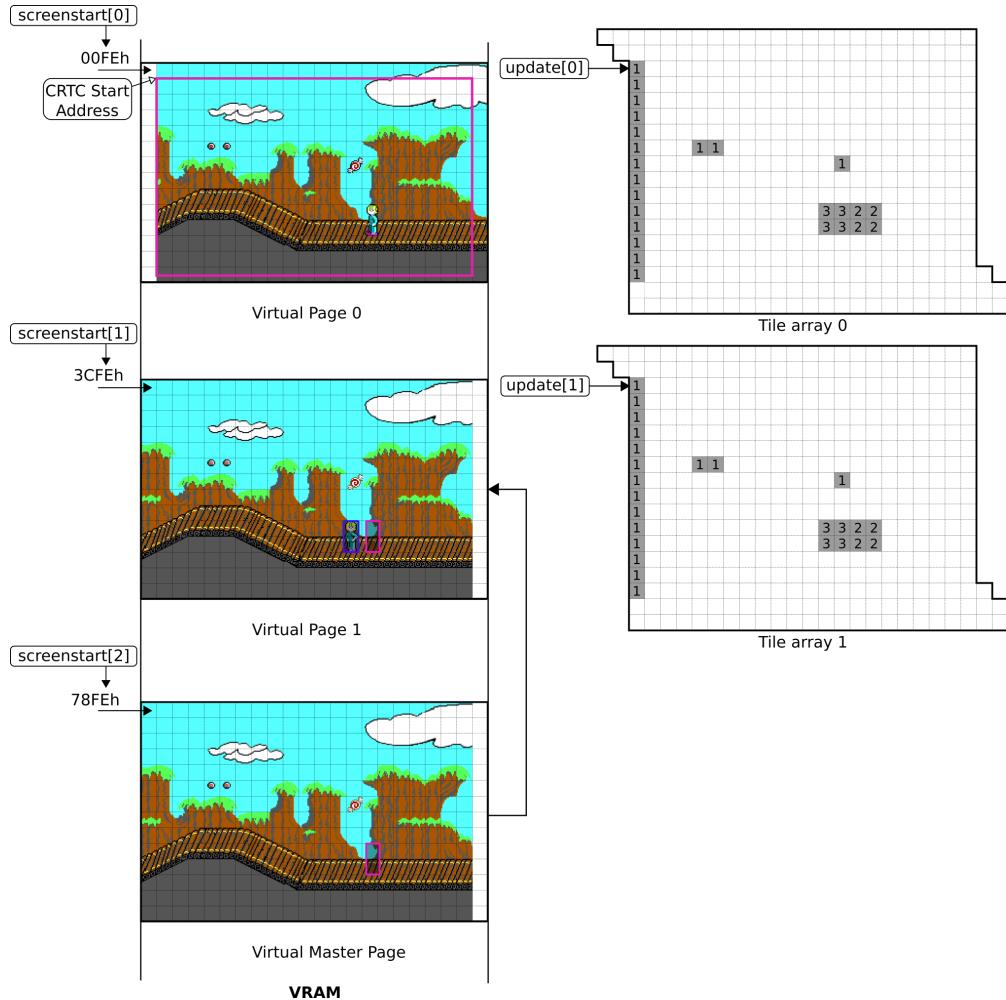


Figure 18: Removing and updating sprites to the buffer screen.

Once the buffer refresh is completed, the engine needs only to update the CRTC start address to virtual page 1. At this point, the visible tile array is emptied, and the pointer resets to its starting location. The new buffer array (virtual page 0) retains the marked tiles since this screen has not yet refreshed. At this point, the refresh cycle restarts, with virtual page 0 now functioning as the buffer.

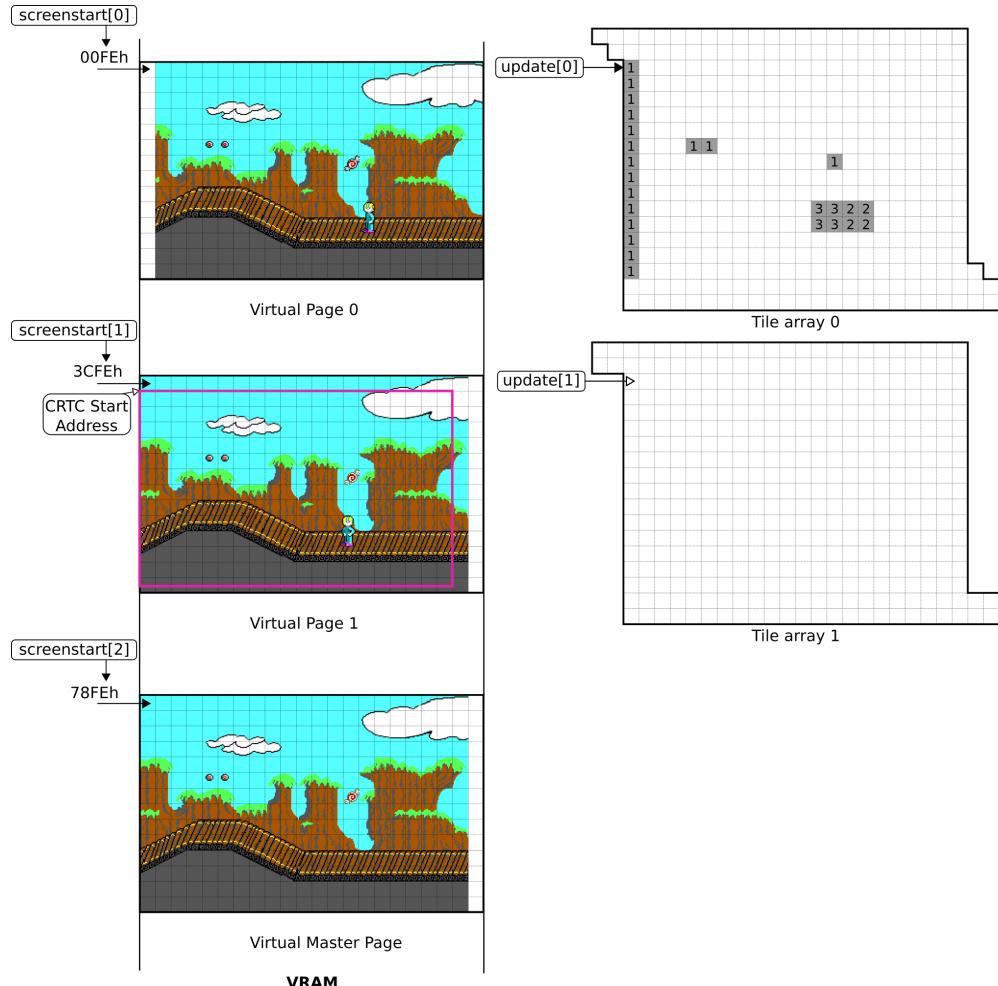


Figure 19: Swap buffer and visible screen by updating CRTC start address.

The final step in scrolling is to adjust the horizontal fine-pixel alignment by setting the Pel Pan register. *Et voilà*, the screen scrolls smoothly to the left. This scrolling process is significantly more efficient than the original ATR engine, resulting in only 8% of tiles to be refreshed!

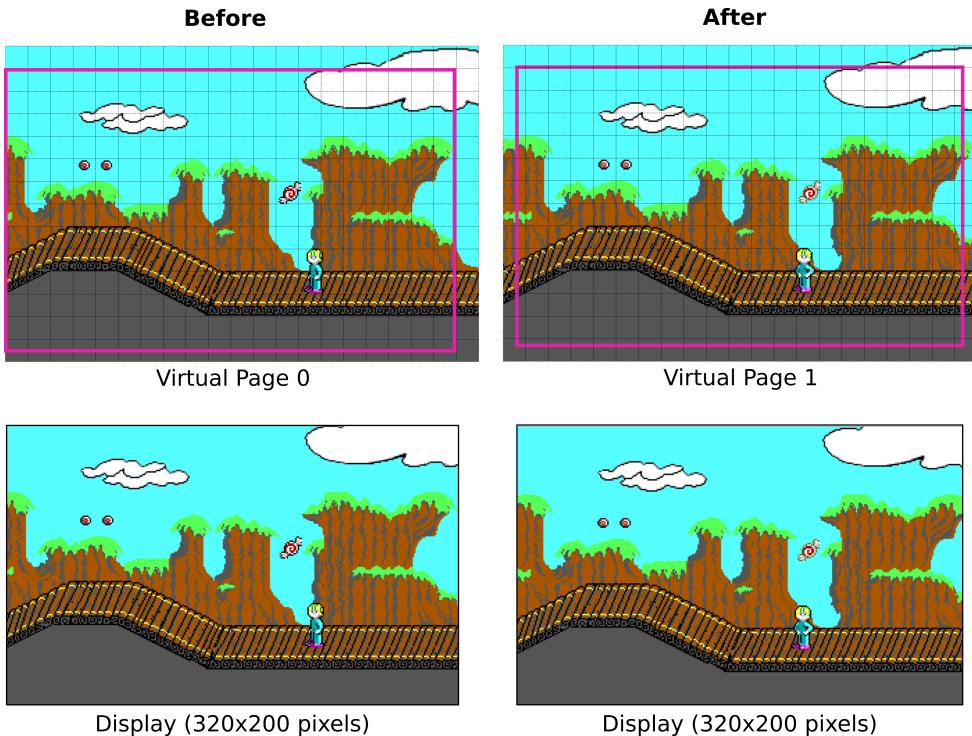


Figure 20: Left screen is before scrolling, right screen after scrolling.

If the player continues to move to the left, the tile array 0 pointer lowers a second time, underscoring why additional space is needed to allow the entire viewport to float up to two tiles in all directions.

0.3 Drawing layer for layer

Each tile on the screen can contain up to three layers: the background layer, the foreground layer and a sprite layer. Rendering the final screen requires multiple redraws on the same tile to achieve the correct layering:

1. Draw the background layer or a combined background and foreground tile.
2. Draw the sprites.
3. Re-draw the foreground layer if a sprite should not appear on top.

In the best case, the engine must read and write 128 bytes ($2 * 16 \text{ bytes} * 4 \text{ memory banks}$) to VRAM for a background tile only. In the worst case, reading and writing up to 608 bytes are needed (background, foreground, sprite, and foreground re-drawn), with several bitwise operations involved. It must run quickly and therefore most of the code is written in assembly.

0.3.1 Draw background and foreground tiles

Drawing the background layer is straightforward, as it only requires copying 128 bytes to VRAM. Drawing foreground tiles on top of the background layer requires an additional mask, which defines which pixels are overwritten by the foreground tile.

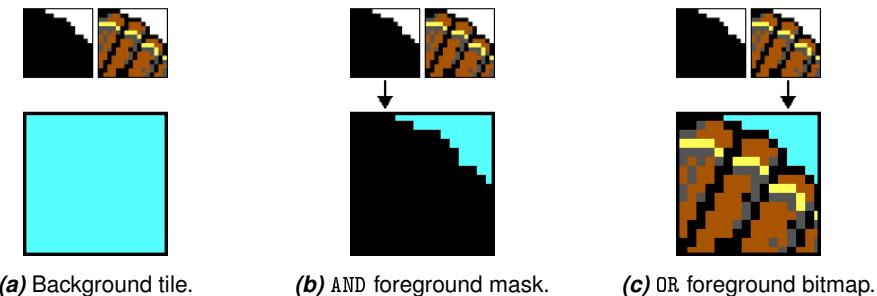


Figure 21: Draw masked foreground tile.

Combining the background and foreground layers involves an AND-bitwise operator to clear the background and the OR-bitwise operator to write the foreground layer.

```
PUBLIC  RFL_NewTile

[...]
es,[mapsegs+2]           ;foreground plane
mov bx,[es:si]
mov es,[mapsegs]          ;background plane
mov si,[es:si]

[...]
or  bx,bx                 ;do we have foreground tile?
jz  @@singletile          ;draw background tile only
jmp @@maskeddraw          ;draw both together

[...]
@@maskeddraw:
shl bx,1
mov ss,[grsegs+STARTTILE16M*2+bx]
shl si,1
mov ds,[grsegs+STARTTILE16*2+si]

xor si,si                 ;first word of tile data

mov ax,SC_MAPMASK+0001b*256 ;map mask for plane 0

mov di,[cs:screenstartcs]

@@planeloopm:
WORDOUT
tileofs = 0
lineoffset = 0
REPT 16
    mov bx,[si+tileofs]    ;background tile
    and bx,[ss:tileofs]    ;mask
    or  bx,[ss:si+tileofs+32] ;masked data
    mov [es:di+lineoffset],bx
tileofs = tileofs + 2
lineoffset = lineoffset + SCREENWIDTH
ENDM
```

0.3.2 Drawing sprites

The next step is to render sprites on the screen. Most home computers of that era had built-in sprite functionality on the video card. For example, on a MSX computer, one could simply enter

```
PUT SPRITE <SpriteNumber>, <X>, <Y>, Color
```

to display a sprite on the screen. Updating the (X, Y) coordinates would move the sprite, with the display adapter handling everything else. Unfortunately, the concept of sprites did not exist on EGA cards, so game developers had to implement their own solution.

A challenge arises from the fact that sprites can move freely across the screen and are not byte-aligned. To address this, the bit-shifting technique described in section ?? is used. When caching a sprite into memory, each sprite is copied four times, with each copy shifted by two or more pixels using this technique. The property `*spr->shifts` determines the number of bit shifts applied to each of the four copies.

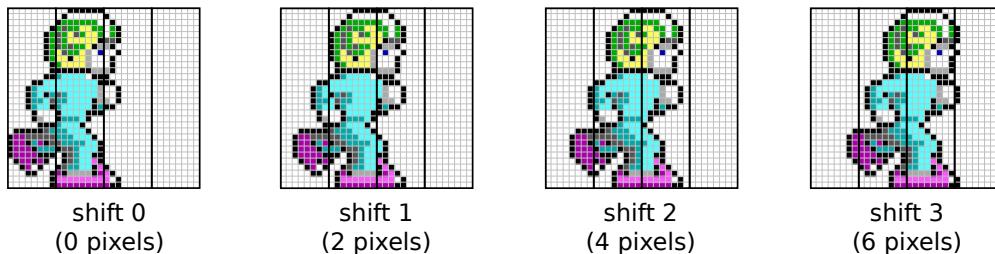


Figure 22: Sprite shifted in 4 steps.

Displaying the correct shifted sprite is as simple as

```
#define G_P_SHIFT    4 // global >> ?? = pixels

//Set x,y to top-left corner of sprite
y+=spr->orgy>>G_P_SHIFT;
x+=spr->orgx>>G_P_SHIFT;

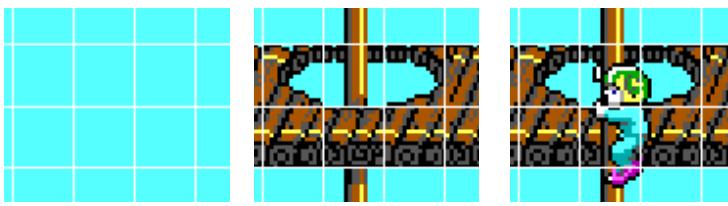
shift = (x&7)/2; // Set sprite shift
```

```
void CAL_CacheSprite (int chunk, char far *compressed)
{
[...]
//
// make the shifts!
//
switch (spr->shifts)
{
case 1: // no shifts
[...]

case 2: // one shift of 4 pixels
for (i=0;i<2;i++)
{
    dest->sourceoffset[i] = shiftstarts[0];
    dest->planesize[i] = smallplane;
    dest->width[i] = spr->width;
}
for (i=2;i<4;i++)
{
    dest->sourceoffset[i] = shiftstarts[1];
    dest->planesize[i] = bigplane;
    dest->width[i] = spr->width+1;
}
CAL_ShiftSprite ((unsigned)grsegs[chunk],dest->
sourceoffset[0],
    dest->sourceoffset[2],spr->width,spr->height,4);
break;

case 4: // four shifts of 2 pixels
[...]
default:
    Quit ("CAL_CacheSprite: Bad shifts number!");
}
}
```

If multiple sprites are displayed on the same tile, each sprite is assigned a priority from 0 to 3 to determine the drawing order. A sprite with a higher priority number is always drawn on top of sprites with a lower priority. Since sprites are always drawn on top of tiles, this can create unnatural situations, such as when Commander Keen is climbing through a hole, as illustrated in Figure 23.

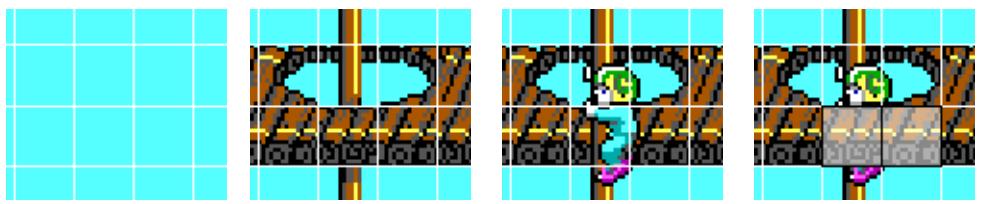


(a) Background tile. (b) Foreground tile. (c) Sprite on top.

Figure 23: Unnatural situation where Commander Keen is in front of a hole.

To draw sprites 'inside' a foreground tile, a trick is used by introducing a priority foreground tile in the `tinf` info table. The attribute is named INTILE ("in-front"). If a foreground tile has its INTILE attribute high bit set (80h), sprites with priority 0-2 will not be drawn over it. Only sprites with priority 3 are drawn over the foreground tile. This results in the following drawing order:

1. Draw the background tile and the masked foreground tile.
2. Draw sprites with priority 0, 1, and 2 (in that order), and mark the corresponding tile in the tile buffer array with a '3', as illustrated in Figure ?? on page ??.
3. Scan the tile buffer array for tiles marked with '3'. If the corresponding foreground tile's INTILE attribute high bit (80h) is set, re-draw the foreground tile.
4. Finally, draw sprites with priority 3. These sprites are always drawn on top of everything.



(a) Background tile. (b) Foreground tile. (c) Sprite on top. (d) Redraw foreground tile.

Figure 24: Draw sprite inside a tile, by redrawing foreground tile.

```
PROC  RFL_MaskForegroundTiles
PUBLIC RFL_MaskForegroundTiles

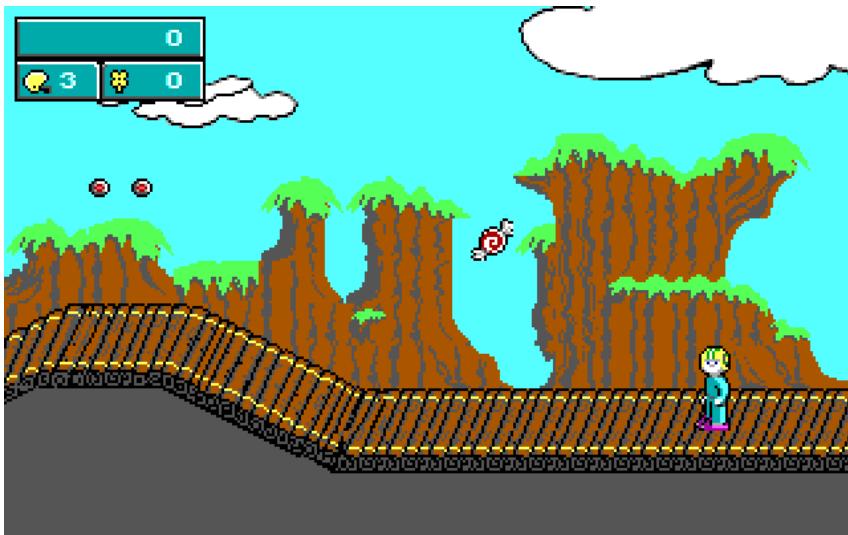
[...]

@@realstart:
    mov di,[updateptr]
    mov bp,(TILESWIDE+1)*TILESHIGH+2
    add bp,di           ; when di = bx,
    push di             ; all tiles have been scanned
    mov cx,-1           ; definately scan the entire thing
;=====
; scan for a 3 in the update list
;=====
@@findtile:
    mov ax,ss
    mov es,ax           ; scan in the data segment
    mov al,3             ; check for tiles marked as '3's
    pop di              ; place to continue scanning from
    repne scasb
    cmp di,bp
    je  @@done
;=====
; found a tile, see if it needs to be masked on
;=====
    push di
    sub di,[updateptr]
    shl di,1
    mov si,[updatemapofs-2+di] ; offset from originmap
    add si,[originmap]
    mov es,[mapsegs+2]         ; foreground map plane segment
    mov si,[es:si]              ; foreground tile number
    or si,si
    jz  @@findtile            ; 0 = no foreground tile
    mov bx,si
    add bx,INTILE            ; INTILE tile info table
    mov es,[tinf]
    test [BYTE PTR es:bx],80h ; high bit = masked tile
    jz  @@findtile

; mask the tile
```

0.3.3 Screen refresh

Flipping between the pages is as simple as setting the CRTC start address registers to page 0 or page 1 starting point. However, there is one issue to solve. If you were to run it, every once in a while the expected screen shown below...



...would instead appear distorted, showing both misalignment and parts of two pages:



This problem results from the timing between updating the CRTC start address and the screen refresh. The start address is latched by the EGA's internal circuitry exactly once per frame, usually at the beginning of the vertical retrace. Although the CRTC start address is a 16-bit value, the `out` instruction can only write 8 bits at a time.

This issue can be illustrated with the following setup: the current CRTC start address (Page 0) points to 0000h, while the buffer (Page 1) points to 3C00h. After moving one tile to the left, Page 0 now points to FFFEh in VRAM, and Page 1 points to 3BFEh. Since Page 1 is the updated buffer, it will be displayed in the next refresh cycle. However, due to poor timing in updating the vertical retrace and start address, the CRTC only picks up the first byte of the address, 3Bh, setting the start address to 3B00h instead of 3BFEh.

```
CRTC_INDEX = 03D4h
CRTC_STARTHIGH = 12

cli                      ; disable interrupts
mov cx,[crtc]             ; [crtc] is start address
mov dx,CRTC_INDEX         ; set CRTR register
mov al,CRTC_STARTHIGH    ; start address high register
out dx,al
inc dx                   ; port 03D5h
mov al,ch
out dx,al                ; set address high

;***** VERTICAL RETRACE STARTS HERE !!!!!!! ****
;***** AND SHOWS 2 PARTIAL FRAMEBUFFERS *****

dec dx                   ; set CRTR register
mov al,0dh                ; start address low register
out dx,al
mov al,cl
inc dx                   ; port 03D5h
out dx,al                ; set address low
sti                      ; enable interrupts

ret
```

One solution to this issue is to update the start address when the vertical retrace signal is detected through the Input Status 1 Register, specifically bit 3 of 3DAh. Unfortunately, by the time this status is observed by the program, the start address for the next frame has already been latched, occurring immediately when the vertical retrace pulse begins.

To address this, the start address must be updated at a point sufficiently distant from the start of the vertical retrace. This requires identifying a signal that confirms the completion of a horizontal or vertical retrace and the beginning of a new scan line, far enough from the vertical retrace to ensure the new start address is latched during the next vertical sync. The Display Enable Status signal, accessible via the Input Status 1 Register, provides this information; a value of 1 indicates that the display is within a horizontal or vertical retrace⁷.

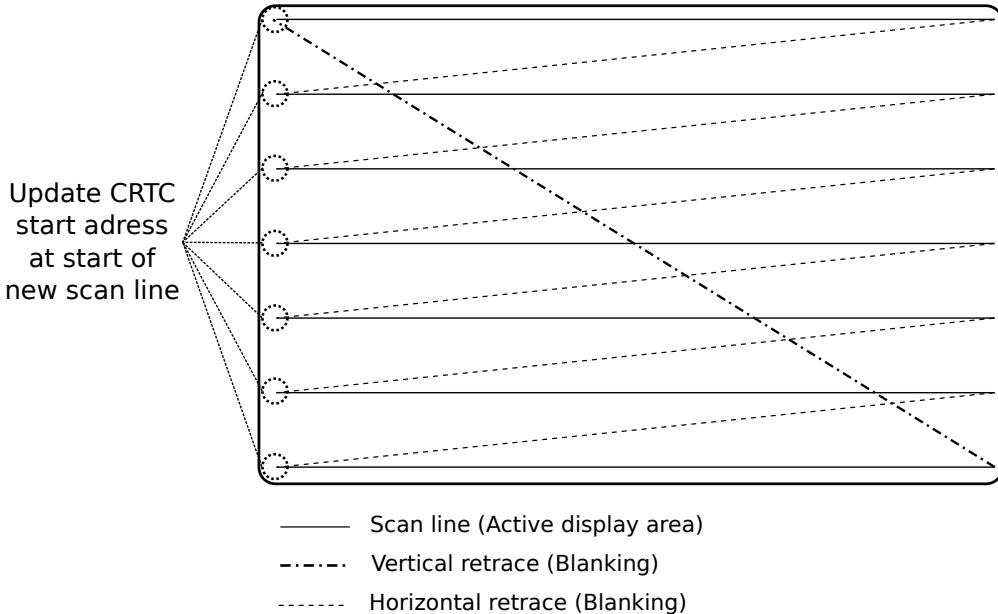


Figure 25: Update CRTC start address at beginning of new scan line.

To guarantee that the start address updates at the very beginning of a new scan line, the process first waits until the current scan line completes. Then, it waits for a full retrace, checking that the Display Enable Status returns to '0'. At this point, the CRTC address is updated.

Trivia : Regardless of CPU speed, the game's speed is limited by the CRTC sync to the monitor's refresh rate, which is 60Hz for EGA and 70Hz for VGA.

⁷Documentation is a bit unclear here. The IBM technical documentation for VGA explains retrace takes place when bit 0 of the Input Status Register 1 is set to high ('1'). The IBM technical EGA documentation explains the opposite, saying when bit 0 is set low ('0') a retrace is taking place. For now, we assume source code and VGA documentation is correct, retrace takes place on a '1'.

```
; =====
;
;  VW_SetScreen
;
; =====

    mov dx,03DAh          ; Status Register 1
;
;  wait until the CRTC just starts scaning a displayed line
;  to set the CRTC start
;
    cli

@@waitnodisplay:        ;wait until scan line is finished
    in al,dx
    test al,01b
    jz @@waitnodisplay

@@waitdisplay:           ;wait until retrace is finished
    in al,dx
    test al,01b
    jnz @@waitdisplay

endif

;
;  set CRTC start
;
    mov cx,[crtc]
    mov dx,CRTC_INDEX
    mov al,0ch      ;start address high register
    out dx,al
    inc dx
    mov al,ch
    out dx,al
    dec dx
    mov al,0dh      ;start address low register
    out dx,al
    mov al,cl
    inc dx
    out dx,al
```

0.3.4 Manage Refresh Timing

After each screen refresh, a certain amount of time, referred to as "ticks", has passed. The tick count depends on several factors, such as the number of tiles refreshed and the waiting period for a screen's vertical retrace. Since all game actions and reactions rely on the tick interval between two refreshes, it is important to keep this interval consistent.

Without controlling the tick interval, the state and speed of actors can become unpredictable, potentially causing them to move too quickly or even "warp" to an unexpected location. To manage refresh intervals effectively, a minimum and maximum number of ticks are defined within the refresh loop.

```
#define MINTICS      2
#define MAXTICS      6

void RF_Refresh (void)
{
    [...]

// 
// calculate tics since last refresh for adaptive timing
//
    do
    {
        newtime = TimeCount;
        tics = newtime - lasttimecount;
    } while (tics < MINTICS);
    lasttimecount = newtime;

    if (tics > MAXTICS)
    {
        TimeCount -= (tics - MAXTICS);
        tics = MAXTICS;
    }
}
```

0.4 Actors and sprites

0.4.1 A.I.

To simulate enemies, some objects are allowed to "think" and take actions like walking, shooting or emitting sounds. These thinking objects are called "actors". Actors are pro-

grammed via a state machine. They can be aggressive (chase you), just running in any direction, or dump (throwing things at you). To model their behavior, all enemies have an associated state:

- Chase Keen
- Smash Keen
- Shoot projectile
- Climb and slide from pole
- Walking around
- Turn into flower
- Special Boss (Boobus)

Each state has associated think, reaction and contact method pointers. There is also a next pointer to indicate which state the actor should transition to when the current state is completed.

```
typedef struct
{
    int      leftshapenum,rightshapenum; // Sprite to render
                                         // on screen
    enum     {step,slide,think,stepphink,slidethink} progress;
    boolean skipable;
    boolean pushtofloor;   // Make sure sprites stays
                           // connected with ground
    int  tictime;          // How long stay in that state
    int  xmove;
    int  ymove;
    void (*think) ();
    void (*contact) ();
    void (*react) ();
    void *nextstate;
} statetype;
```

All actors have a state chain, as example the tater trooper.

```

statetype s_taterwalk1 = {TATERTROOPWALK1SPR,TATERTROOPWALKR1SPR,step ,
    false , true ,10, 128,0, TaterThink , NULL , WalkReact, &s_taterwalk2};
statetype s_taterwalk2 = {TATERTROOPWALK2SPR,TATERTROOPWALKR2SPR,step ,
    false , true ,10, 128,0, TaterThink , NULL , WalkReact, &s_taterwalk3};
statetype s_taterwalk3 = {TATERTROOPWALK3SPR,TATERTROOPWALKR3SPR,step ,
    false , true ,10, 128,0, TaterThink , NULL , WalkReact, &s_taterwalk4};
statetype s_taterwalk4 = {TATERTROOPWALK4SPR,TATERTROOPWALKR4SPR,step ,
    false , true ,10, 128,0, TaterThink , NULL , WalkReact, &s_taterwalk1};

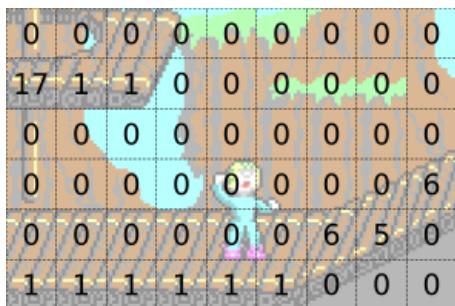
statetype s_taterattack1 = {TATERTROOPLUNGE1SPR,TATERTROOPLUNGER1SPR,
    step ,false , false ,12, 0,0, NULL , NULL , BackupReact, &s_taterattack2 };
statetype s_taterattack2 = {TATERTROOPLUNGE1SPR,TATERTROOPLUNGER2SPR,
    step ,false , false ,20, 0,0, NULL , NULL , DrawReact, &s_taterattack3};
statetype s_taterattack3 = {TATERTROOPLUNGE1SPR,TATERTROOPLUNGER1SPR,
    step ,false , false ,8, 0,0, NULL , NULL , DrawReact, &s_taterwalk1};

```

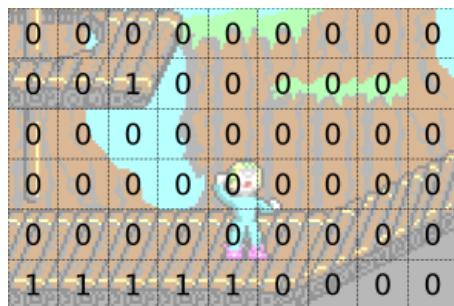
All types of enemies (including Boobus) have their own state machine. They often share the same reactions (e.g. WalkReact and ProjectileReact), but often have their own thinking state.

0.4.2 Clipping

Before drawing a sprite on the screen, the engine determines if the boundaries of a sprite are hitting a wall or floor. This is called clipping and ensures an actor doesn't fall through a floor or walks through a vertical wall. To define whether a tile is a wall or floor, a tile is enriched with tile information, as explained in section ?? on page ???. Each foreground tile contains a NORTHWALL, SOUTHWALL, EASTWALL and WESTWALL, as explained in section ???. A number greater than 0 means the tile is a wall or floor when approaching from a given direction.



(a) Wall type map NORTHWALL



(b) Wall type map EASTWALL

Figure 26: Foreground tile clipping information.

As example, when a sprite, moving from right to left, is hitting a wall on the left side, it will

update the sprite movement to ensure the sprite clips to the east wall of the left tile as illustrated in Figure 27. The east/west wall clipping logic is covered by `ClipToEastWalls()` and `ClipToWestWalls()` functions.

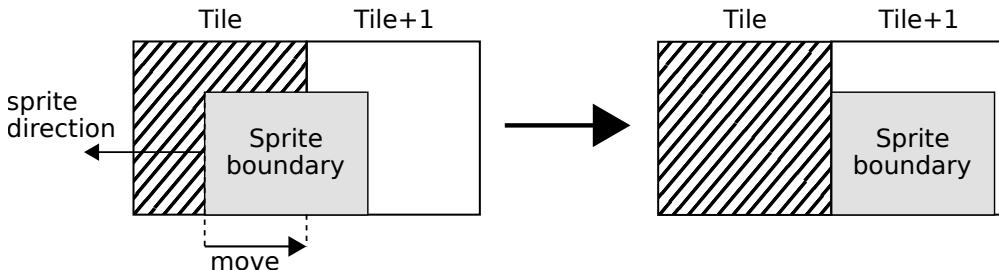


Figure 27: Clipping to east wall when moving from the west.

```
void ClipToEastWalls (objtype *ob)
{
    ...
    for (y=top;y<=bottom;y++)
    {
        map = (unsigned far *)mapsegs[1] +
            mapwidthtable[y]/2 + ob->tileleft;

        //Check if we hit EAST wall
        if (ob->hiteast = tinf[EASTWALL+*map])
        {
            //Clip left side actor to left side
            //of next right tile
            move = ( (ob->tileleft+1)<<G_T_SHIFT ) - ob->left;
            MoveObjHoriz (ob,move);
            return;
        }
    }
}
```

For clipping top and bottom the engine also needs to take clipping to slopes into account. After the sprite is clipped to the top or bottom of the wall tile, an offset can be applied to move a sprite up or down a slope. The offset is defined by a lookup table, where the midpoint of the sprite and the wall type from the foreground tile defines the offset.

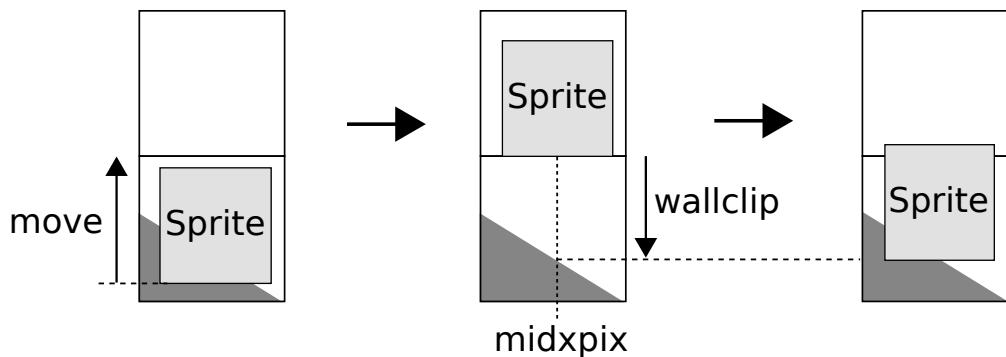


Figure 28: Clipping north wall with slope.

```
// walltype / x coordinate (0–15)

int wallclip[8][16] = {      // the height of a given point in a tile
{ 256, 256, 256, 256, 256, 256, 256, 256, 256, 256, 256, 256, 256, 256, 256, 256, 256},
{ 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0},
{ 0,0x08,0x10,0x18,0x20,0x28,0x30,0x38,0x40,0x48,0x50,0x58,0x60,0x68,0x70,0x78},
{0x80,0x88,0x90,0x98,0xa0,0xa8,0xb0,0xb8,0xc0,0xc8,0xd0,0xd8,0xe0,0xe8,0xf0,0xf8},
{ 0,0x10,0x20,0x30,0x40,0x50,0x60,0x70,0x80,0x90,0xa0,0xb0,0xc0,0xd0,0xe0,0xf0},
{0x78,0x70,0x68,0x60,0x58,0x50,0x48,0x40,0x38,0x30,0x28,0x20,0x18,0x10,0x08, 0},
{0xf8,0xf0,0xe8,0xe0,0xd8,0xd0,0xc8,0xc0,0xb8,0xb0,0xa8,0xa0,0x98,0x90,0x88,0x80},
{0xff,0xe0,0xd0,0xc0,0xb0,0xa0,0x90,0x80,0x70,0x60,0x50,0x40,0x30,0x20,0x10, 0}
};
```

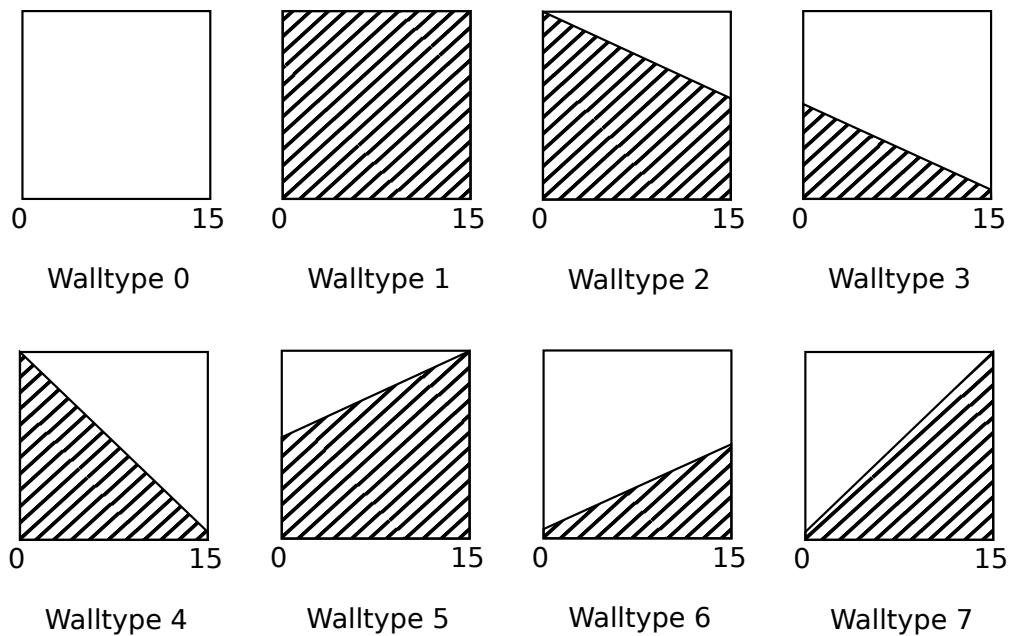


Figure 29: Eight different walltypes (slopes) defined.

```
void ClipToEnds (objtype *ob)
{
    ...

    //Get midpoint of sprite [0-15]
    midxpix = (ob->midx&0xf0) >> 4;

    map = (unsigned far *)mapsegs[1] +
        mapbwidhtable[oldtilebottom-1]/2 + ob->tilemidx;
    for (y=oldtilebottom-1 ; y<=ob->tilebottom ; y++, map+=
        mapwidth)
    {
        //Do we hit a NORTH wall
        if (wall = tinf[NORTHWALL+*map])
        {
            //offset from tile border clip
            clip = wallclip[wall&7][midxpix];
            //Clip bottom side actor to top side tile + offset-1
            move = ( (y<<G_T_SHIFT)+clip - 1) - ob->bottom;
            if (move<0 && move>=maxmove)
            {
                ob->hitnorth = wall;
                MoveObjVert (ob,move);
                return;
            }
        }
    }
}
```

The game and all actors are defined in a global coordinate system, which is scaled to 16 times a pixel. The higher resolution enables more precision of movements and better simulation of movement acceleration. Conversion between global and pixel coordinates can be easily performed by bit shift operations.

0.5 use later

```
#define G_P_SHIFT 4 // global >> ?? = pixels

void RFL_CalcOriginStuff (long x, long y)
{
    originxglobal = x;
    originyglobal = y;
    [...]

    //panning 0-15 pixels
    panx = (originxglobal>>G_P_SHIFT) & 15;
    //move CRTC + 1 when move >7 pixels
    pansx = panx & 8;
    pany = pansy = (originyglobal>>G_P_SHIFT) & 15;
    //Start location in VRAM
    panadjust = panx/8 + ylookup[pany];
}
```

To take full advantage of this optimization, the refresh algorithm maintains a list of tiles that are already copied on the master page via `tilecache` variable. If a tile is already on the master page the algorithm copies the tile from that location to its destination instead of the RAM location in memory, saving the four separated writes to each memory plane.

0.5. USE LATER
