

# CS532100 Numerical Optimization Homework 2

109062710 Bijon Setyawan Raya

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1. Consider the linear least square problem:

$$\min_{\vec{x} \in \mathbb{R}^2} \|A\vec{x} - \vec{b}\|^2,$$

where

$$A = \begin{bmatrix} 4 & 8 \\ 2 & 4 \\ 1 & 2 \end{bmatrix}, \vec{b} = \begin{pmatrix} 21/4 \\ 0 \\ 0 \end{pmatrix}$$

- (a) (10%) Write its normal equation.

$$A^T A \vec{x} = A^T \vec{b} \quad (1)$$

$$\begin{bmatrix} 4 & 2 & 1 \\ 8 & 4 & 2 \end{bmatrix} \begin{bmatrix} 4 & 8 \\ 2 & 4 \\ 1 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 4 & 2 & 1 \\ 8 & 4 & 2 \end{bmatrix} \begin{bmatrix} 21/4 \\ 0 \\ 0 \end{bmatrix} \quad (2)$$

- (b) (10%) Express  $\vec{b} = \vec{b}_1 + \vec{b}_2$  such that  $\vec{b}_1$  is in the subspace spanned by  $A$ 's column vectors, and  $\vec{b}_2$  is orthogonal to  $A$ 's column vectors. From matrix  $A$ , we know that its basis is its first column which is  $\begin{bmatrix} 4 & 2 & 1 \end{bmatrix}^T$ . From its basis column, we can express  $\vec{b}$  in  $\vec{b}_1 + \vec{b}_2$ . Let,

$$\vec{b}_1 = \begin{bmatrix} 4s \\ 2s \\ s \end{bmatrix}, \vec{b}_2 = \begin{bmatrix} t_1 \\ t_2 \\ t_3 \end{bmatrix} \quad (3)$$

Then we can express  $\vec{b}$  such as

$$\vec{b} = \vec{b}_1 + \vec{b}_2 \quad (4)$$

$$\begin{bmatrix} 21/4 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 4s + t_1 \\ 2s + t_2 \\ s + t_3 \end{bmatrix} \quad (5)$$

Since  $\vec{b}_1$  is a column space of  $A$ , and  $\vec{b}_2$  is orthogonal to  $\vec{b}_1$ , then

$$\begin{bmatrix} 4s \\ 2s \\ s \end{bmatrix} \begin{bmatrix} t_1 & t_2 & t_3 \end{bmatrix} = 0 \quad (6)$$

$$4s \cdot t_1 + 2s \cdot t_2 + s \cdot t_3 = 0 \quad (7)$$

We can add  $4s \cdot t_1 + 2s \cdot t_2 + s \cdot t_3 = 0$  at the last index of our matrix equality. Then

$$\begin{bmatrix} 4s + t_1 \\ 2s + t_2 \\ s + t_3 \\ 4s \cdot t_1 + 2s \cdot t_2 + s \cdot t_3 \end{bmatrix} = \begin{bmatrix} 21/4 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

Solving the equality above, we then have  $s = 1$ ,  $t_1 = 5/4$ ,  $t_2 = -2$ , and  $t_3 = -1$ . Thus,

$$\vec{b}_1 = \begin{bmatrix} 4s \\ 2s \\ s \end{bmatrix}, \vec{b}_2 = \begin{bmatrix} 5/4 \\ -2 \\ -1 \end{bmatrix} \quad (8)$$

$$\begin{bmatrix} 21/4 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 4 \\ 2 \\ 1 \end{bmatrix} + \begin{bmatrix} 5/4 \\ -2 \\ -1 \end{bmatrix} \quad (9)$$

- (c) (10%) Show that  $\vec{z} \in \mathbb{R}^2$  is a least square solution for  $A\vec{x} = \vec{b}$  if and only if  $\vec{z}$  is part of a solution to the larger linear system:

$$\begin{bmatrix} 0 & A^T \\ A & I \end{bmatrix} \begin{bmatrix} \vec{z} \\ \vec{y} \end{bmatrix} = \begin{bmatrix} 0 \\ \vec{b} \end{bmatrix}$$

2. In Note05 (Page 16), memoryless BFGS iteration matrix  $H_{k+1}$  can be derived from considering the Hestenes–Stiefel form of the nonlinear conjugate gradient method. Recalling that  $\vec{s}_k = \alpha_k \vec{p}_k$ , we have that the search direction for this method is given by

$$\begin{aligned} \vec{p}_{k+1} &= -\nabla f_{k+1} + \frac{\nabla f_{k+1}^T \vec{y}_k}{\vec{y}_k^T \vec{p}_k} \vec{p}_k \\ &= -\nabla f_{k+1} + \frac{\nabla f_{k+1}^T \vec{y}_k}{\vec{y}_k^T \vec{s}_k} \vec{s}_k \\ &= -(I - \frac{\vec{s}_k \vec{y}_k^T}{\vec{y}_k^T \vec{s}_k}) \nabla f_{k+1} \\ &= -\hat{H}_{k+1} \nabla f_{k+1} \end{aligned}$$

However, the matrix  $\hat{H}_{k+1}$  is neither symmetric nor positive definite.

- (a) (10%) Please show that the matrix  $\hat{H}_{k+1}$  is singular.  
(You can only prove it for the case  $\nabla f_k, \vec{p}_k, \vec{y}_k, \vec{s}_k \in \mathbb{R}^2$  for all  $k \in \mathbb{N}$ .)  
asd
- (b) (0%) Please read the reference book (Page 180) to understand the derivation of the inverse hessian formula in Note05 (Page 16).  
(you don't need to write anything in this subproblem.)

$$H_{k+1} = (I - \frac{\vec{s}_k \vec{y}_k^T}{\vec{y}_k^T \vec{s}_k})(I - \frac{\vec{y}_k \vec{s}_k^T}{\vec{y}_k^T \vec{s}_k}) + \frac{\vec{s}_k \vec{s}_k^T}{\vec{y}_k^T \vec{s}_k}$$

3. (10%) The total least square problem is to solve the following problem

$$\min_{\vec{x}, \|\vec{x}\|=1} \vec{x}^T A^T A \vec{x}$$

where  $A$  is an  $m \times n$  matrix. Here we assume  $m > n$ . Let  $A = U\Sigma V^T$  be the SVD of  $A$ , where  $U$  is the matrix of left singular vectors,  $V$  is the matrix of right singular vectors, and  $\Sigma$  is a diagonal matrix with diagonal elements  $\sigma_1, \sigma_2, \dots, \sigma_n$ . Moreover,  $U$  and  $V$  are orthogonal matrices, and  $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_n$ . Show the solution of the above problem is the  $\sigma^2$ .

4. Consider the following linear programming problem:

$$\begin{array}{ll} \max_{x_1, x_2} & z = x_1 + x_2 \\ \text{s.t.} & x_1 + 2x_2 \leq 4 \\ & 4x_1 + 2x_2 \leq 12 \\ & -x_1 + x_2 \leq 1 \\ & x_1, x_2 \geq 0 \end{array}$$

- (a) (10%) Please refer Note08 (Page 2) to draw the figure of the constraints by any means, and use that to solve the problem.
- (b) (10%) Derive its dual problem and solve the dual problem by any means. Compare the solutions of the primal and the dual problems.
- (c) (10%) Verify the complementarity slackness condition.
- (d) (10%) Transform the problem to the standard form.
- (e) (10%) Solve it by the simplex method, as provided in Figure 1, using  $\vec{x}_0 = (0, 0)$ . Indicate  $B_k, N_k, \vec{s}_k, \vec{d}_k, p_k, q_k, \gamma_k$  in each step.

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- (1) Given a basic feasible point  $\vec{x}_0$  and the corresponding index set  $\mathcal{B}_0$  and  $\mathcal{N}_0$ .
  - (2) For  $k = 0, 1, \dots$
  - (3) Let  $B_k = A(:, \mathcal{B}_k)$ ,  $N_k = A(:, \mathcal{N}_k)$ ,  $\vec{x}_B = \vec{x}_k(\mathcal{B}_k)$ ,  $\vec{x}_N = \vec{x}_k(\mathcal{N}_k)$ ,  
and  $\vec{c}_B = \vec{c}_k(\mathcal{B}_k)$ ,  $\vec{c}_N = \vec{c}_k(\mathcal{N}_k)$ .
  - (4) Compute  $\vec{s}_k = \vec{c}_N - N_k^T (B_k^{-1})^T \vec{c}_B$  (pricing)
  - (5) If  $\vec{s}_k \geq 0$ , return the solution  $\vec{x}_k$ . (found optimal solution)
  - (6) Select  $q_k \in \mathcal{N}_k$  such that  $\vec{s}_k(i_{q_k}) < 0$ ,  
where  $i_{q_k}$  is the index of  $q_k$  in  $\mathcal{N}_k$
  - (7) Compute  $\vec{d}_k = B_k^{-1} A_k(:, q_k)$ . (search direction)
  - (8) If  $\vec{d}_k \leq 0$ , return **unbounded**. (unbounded case)
  - (9) Compute  $[\gamma_k, i_p] = \min_{i, \vec{d}_k(i) > 0} \frac{\vec{x}_B(i)}{\vec{d}_k(i)}$  (ratio test)  
(The first return value is the minimum ratio;  
the second return value is the index of the minimum ratio.)
  - (10)  $x_{k+1} \begin{pmatrix} \mathcal{B} \\ \mathcal{N} \end{pmatrix} = \begin{pmatrix} \vec{x}_B \\ \vec{x}_N \end{pmatrix} + \gamma_k \begin{pmatrix} -\vec{d}_k \\ \vec{e}_{i_{q_k}} \end{pmatrix}$   
( $\vec{e}_{i_{q_k}} = (0, \dots, 1, \dots, 0)^T$  is a unit vector with  $i_{q_k}$ th element 1.)
  - (11) Let the  $i_p$ th element in  $\mathcal{B}$  be  $p_k$ . (pivoting)  
 $\mathcal{B}_{k+1} = (\mathcal{B}_k - \{p_k\}) \cup \{q_k\}$ ,  $\mathcal{N}_{k+1} = (\mathcal{N}_k - \{q_k\}) \cup \{p_k\}$
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Figure 1: The simplex method for solving (minimization) linear programming