## GDR and ASF Winter School on Distributed Systems and Networks

# Hardware accelerators for virtualization

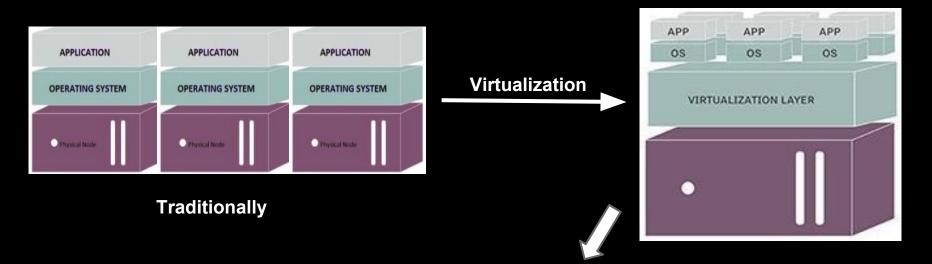
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## Agenda

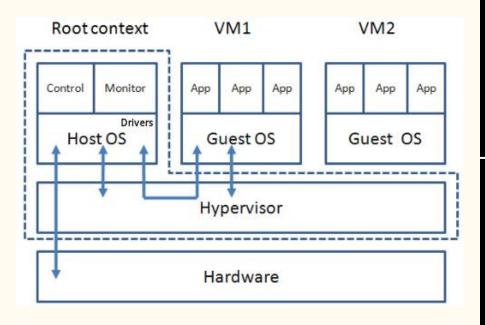
- Context
- State of the art
- Problem Statement
- Contribution
- Evaluation
- Conclusion

## Context



- Resources management and energy saving (vm consolidation)
- Scalability and easy deployment of applications
- Maintenance (vm migration) and fault tolerance
- etc.

### Classic virtualization architecture



### **Root context Role**

- VM management
- I/O drivers
- Monitoring tools
- etc.

### **Hypervisor Role**

- Scheduler
- Memory allocator
- Interrupt manager
- etc.

## Virtualization side effects

### Limits of classical architecture

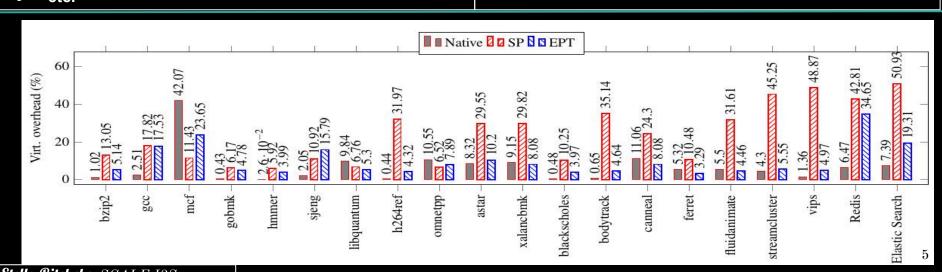
- Hypervisor intervention (cpu time consumed by context switches)
- Interference between VMs (pollution, e.g. cache pollution)
- Hardware abstraction from the VM (VMs are black boxes)
- etc.



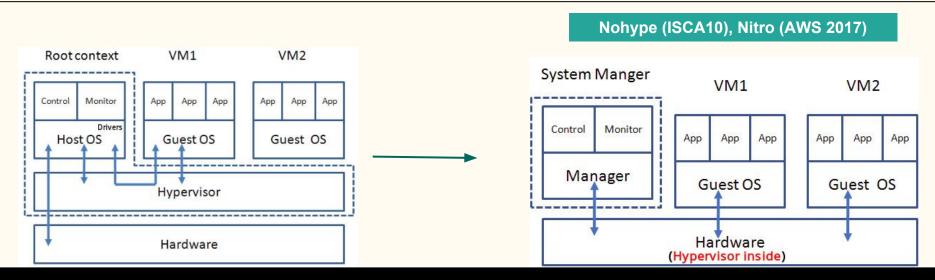
VT-d / VT-x

**Promising approaches: Hardware Assisted Virtualization** 

- **SRIOV** CAT
- **APICV**
- etc.



### State of the art



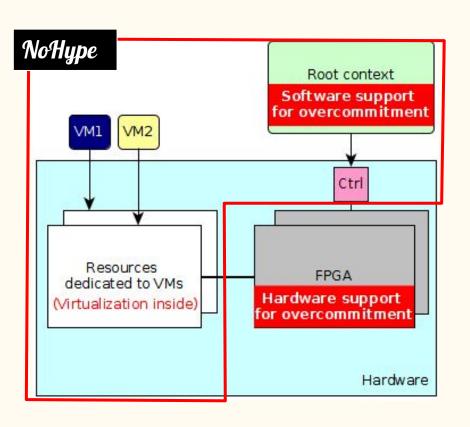
Limits of virtualization with hypervisor inside the hardware : rigidity of hypervisor

- In case of security issues (Meltdown & Spectre)
- In case of hypervisor updates (requires to change the hardware)
- Does not support Over Provisioning → PhD project
- etc.

# Importance of over provisioning

- Workloads are not static
- Cloud users overestimate resources
- New workloads: Faas
- Cloud provider competition for costs reducing
- New hardware feature to facilitate overcommitment (e.g. PML)

## Contribution



The goal of my thesis:

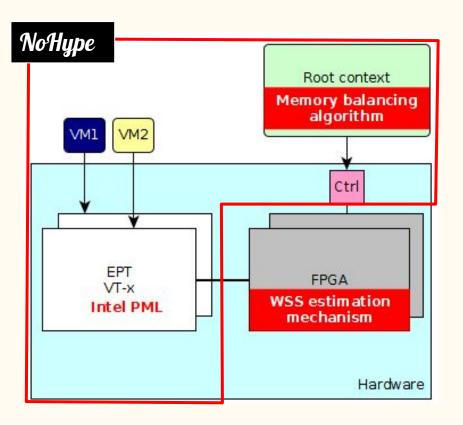
(Re)design hardware features and software supports (when needed) to improve overcommitment of:

**CPU** 

**Memory (first step)** 

Without impacting VMs perf.

## Memory overcommitment



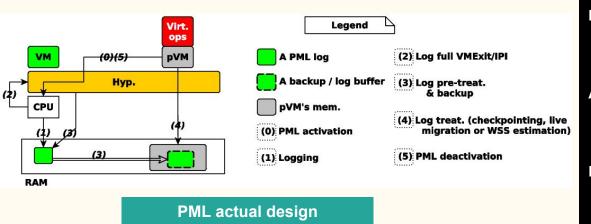
- 1. Working Set Size (WSS) estimation: determine the effective needs of the VMs
- 2. VM memory balancing

Existing methods: all software based  $\rightarrow$  induce overheads

- Geiger [xxx'10]
- Exclusive cache [xxx'10]
- VMWare [xxx'10]
- etc.

Our solution: WSS estimation method based on Intel PML (hardware feature introduced in 2016 by Intel in collaboration with VMWare)

## Intel Page Modification Logging (PML)



### **Description:**

 Allows hypervisor to monitor the guest-physical pages modified by the VM

#### Aim:

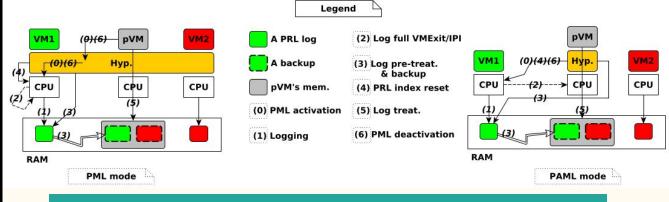
Facilitate working set statistics during VM operations

#### Limits:

- Impact on VMs (VMExits)
- Overhead (CPU time used)

### PRL (Page Reference Logging)

- a new PML design that we propose
- Adresses PML limits



Page Reference Logging (PRL), the new design that we propose.

## Implementation and Evaluation methodology

- 1. Implementation and Evaluation frameworks
  - a. Gem5, a computer architecture simulator
  - b. FPGA boards [may be]
  - c. Xen and KVM virtualization systems
- 2. Benchmarks
  - a. SPEC benchmarks, CloudSuite, BigBench
- 3. Metrics
  - a. Performance impact on VMs (#VMEXITs)
  - b. Resource gain
  - c. Energy gain (number of active servers, electricity consumption)
  - d. Resource consumption (by our solution): #CPU and #memory
- 4. Comparison with state-of-the-art solutions
  - a. see the previous slide

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Questions?